

A Proof of Proposition 3.1

Proof. Let consider the subgraph $\mathcal{G}^{(i)}$ containing all the nodes that have been assigned to V_1 or V_2 at the end of iteration i of Algorithm 2. Let us denote $m^{(i)}$ the number of edges in the graph $\mathcal{G}^{(i)}$.

At the first iteration, the algorithm chooses the node 1, computes $n_1 = 0$ and $n_2 = 0$, and then assigns node 1 to V_1 . With only one node in $\mathcal{G}^{(1)}$, we have $m^{(1)} = 0$. By denoting $c^{(i)}$ the number of additional cut edges induces by the assignment of node i at iteration i , we have

$$\sum_{i=1}^1 c^{(i)} = c^{(1)} = 0 \geq \frac{m^{(1)}}{2} \quad (5)$$

Indeed, at the end of iteration 1, there is only one node assigned, hence the number of cut edges induced by this assignment is $c^{(1)} = 0$.

Suppose that $\sum_{i=1}^p c^{(i)} \geq \frac{m^{(p)}}{2}$ for a certain $p \in \{1, \dots, n-1\}$, let us prove that $\sum_{i=1}^{p+1} c^{(i)} \geq \frac{m^{(p+1)}}{2}$.

Indeed, at the iteration $p+1$, the algorithm chooses the node $(p+1)$ and computes n_1 and n_2 . Since n_1 represents the number of neighbors of the node $(p+1)$ in V_1 , if the node $p+1$ is added to V_2 , then $2 \times n_1$ edges would be cut (the factor 2 comes from the fact that between two nodes i and j , there are the edges (i, j) and (j, i)). Similarly, since n_2 represents the number of neighbors of the node $(p+1)$ in V_2 , if the node $(p+1)$ is added to V_1 , then $2 \times n_2$ edges would be cut. Notice also that there is a total of $2 \times n_1 + 2 \times n_2$ edges between the node $(p+1)$ and the nodes in $\mathcal{G}^{(p)}$. In the algorithm, the node $(p+1)$ is added to V_1 or V_2 such that we cut the most edges, indeed one has

$$c^{(p+1)} = \max(2n_1, 2n_2) \geq \frac{2n_1 + 2n_2}{2} = n_1 + n_2 .$$

Hence,

$$\sum_{i=1}^{p+1} c^{(i)} = \sum_{i=1}^p c^{(i)} + c^{(p+1)} \geq \frac{m^{(p)}}{2} + c^{(p+1)} \geq \frac{m^{(p)}}{2} + n_1 + n_2 \quad (6)$$

The number of edges that is added to the subgraph $\mathcal{G}^{(p)}$ when adding the node $(p+1)$ is equal to $2n_1 + 2n_2 = m^{(p+1)} - m^{(p)}$, hence,

$$\frac{m^{(p)}}{2} + n_1 + n_2 = \frac{m^{(p)}}{2} + \frac{m^{(p+1)} - m^{(p)}}{2} = \frac{m^{(p+1)}}{2} \quad (7)$$

We have shown that $\sum_{i=1}^1 c^{(i)} \geq \frac{m^{(1)}}{2}$ and that if $\sum_{i=1}^p c^{(i)} \geq \frac{m^{(p)}}{2}$ for a certain $p \in \{1, \dots, n-1\}$, then $\sum_{i=1}^{p+1} c^{(i)} \geq \frac{m^{(p+1)}}{2}$. Thus, $\sum_{i=1}^p c^{(i)} \geq \frac{m^{(p)}}{2}$ for any $p \in \{1, \dots, n\}$, especially for $p = n$ where $\mathcal{G}^{(n)} = \mathcal{G}$. By definition $\sum_{i=1}^n c^{(i)}$ is the total number of edges that are cut which also means that

$$\sum_{i=1}^n c^{(i)} = \text{Card} \{(i, j) \in E \mid (i \in V_1 \wedge j \in V_2) \vee (i \in V_2 \wedge j \in V_1)\} .$$

□

B Proof of Theorem 3.2 and Theorem 4.1

To properly derive the regret bounds, we will have to make connections between our setting and a standard linear bandit that chooses sequentially Tm arms. For that matter, let us consider an arbitrary

order on the set of edges E and denote $E[i]$ the i -th edge according to this order with $i \in \{1, \dots, m\}$. We define for all $t \in \{1, \dots, T\}$ and $p \in \{1, \dots, m\}$ the OLS estimator

$$\hat{\theta}_{t,p} = \mathbf{A}_{t,p}^{-1} b_{t,p} ,$$

where

$$\mathbf{A}_{t,p} = \lambda \mathbf{I}_{d^2} + \sum_{s=1}^{t-1} \sum_{b=1}^m z_s^{E[b]} z_s^{E[b]\top} + \sum_{k=1}^p z_t^{E[k]} z_t^{E[k]\top} ,$$

with λ a regularization parameter and

$$b_{t,p} = \sum_{s=1}^{t-1} \sum_{b=1}^m z_s^{E[b]} y_s^{E[b]} + \sum_{k=1}^p z_t^{E[k]} y_t^{E[k]} . \quad (8)$$

We define also the confidence set

$$C_{t,p}(\delta) = \left\{ \theta : \|\theta - \hat{\theta}_{t,p}\|_{\mathbf{A}_{t,p}^{-1}} \leq \sigma \sqrt{d^2 \log \left(\frac{1 + tmL^2/\lambda}{\delta} \right)} + \sqrt{\lambda} S \right\} , \quad (9)$$

where with probability $1 - \delta$, we have that $\theta_\star \in C_{t,p}(\delta)$ for all $t \in \{1, \dots, T\}$, $p \in \{1, \dots, m\}$ and $\delta \in (0, 1]$.

Notice that the confidence set $C_t(\delta)$ defined in Section 3 is exactly the confidence set $C_{t,m}(\delta)$ defined here. The definitions of the matrix $A_{t,m}$ and the vector $b_{t,m}$ follow the same reasoning.

B.1 Proof of Theorem 3.2

Proof. Recall that $(x_\star^{(1)}, \dots, x_\star^{(n)}) = \arg \max_{(x^{(1)}, \dots, x^{(n)})} \sum_{(i,j) \in E} x^{(i)\top} \mathbf{M}_\star x^{(j)}$ is the optimal joint arm, and we define for each edge $(i, j) \in E$ the optimal edge arm $z_\star^{(i,j)} = \text{vec}(x_\star^{(i)} x_\star^{(j)\top})$.

We recall that the α -pseudo-regret is

$$R_\alpha(T) \triangleq \sum_{t=1}^T \sum_{(i,j) \in E} \alpha \langle z_\star^{(i,j)}, \theta_\star \rangle - \langle z_t^{(i,j)}, \theta_\star \rangle \quad (10)$$

$$= R(T) - \sum_{t=1}^T \sum_{(i,j) \in E} (1 - \alpha) \langle z_\star^{(i,j)}, \theta_\star \rangle , \quad (11)$$

where the pseudo-regret $R(T)$ is defined by

$$R(T) = \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_\star^{(i,j)}, \theta_\star \rangle - \langle z_t^{(i,j)}, \theta_\star \rangle .$$

Let us borrow the notion of *Critical Covariance Inequality* introduced in [Chan et al., 2021], that is for a given round $t \in \{1, \dots, T\}$ and $p \in \{1, \dots, m\}$, the expected covariance matrix $\mathbf{A}_{t,p}$ satisfies the critical covariance inequality if

$$\mathbf{A}_{t-1,m} \preceq \mathbf{A}_{t,p} \preceq 2\mathbf{A}_{t-1,m} . \quad (12)$$

Let us now define the event D_t as the event where at a given round t , for all $p \in \{1, \dots, m\}$, $\mathbf{A}_{t,p}$ satisfies the critical covariance inequality (CCI).

We can write the pseudo-regret as follows:

$$\begin{aligned}
R(T) &= \sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \langle z_\star^{(i,j)}, \theta_\star \rangle - \langle z_t^{(i,j)}, \theta_\star \rangle + \mathbb{1}[D_t^c] \sum_{(i,j) \in E} \langle z_\star^{(i,j)}, \theta_\star \rangle - \langle z_t^{(i,j)}, \theta_\star \rangle \\
&\leq \underbrace{\sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \langle z_\star^{(i,j)}, \theta_\star \rangle - \langle z_t^{(i,j)}, \theta_\star \rangle}_{(a)} + \underbrace{LSm \sum_{t=1}^T \mathbb{1}[D_t^c]}_{(b)} .
\end{aligned}$$

We know that the approximation Max-CUT algorithm returns two subsets of nodes V_1 and V_2 such that there are at least $m/2$ edges between V_1 and V_2 , and to be more precise: at least $m/4$ edges from V_1 to V_2 and at least $m/4$ edges from V_2 to V_1 . Hence at each time t , if all the nodes of V_1 pull the node-arm x_t and all the nodes of V_2 pull the node-arm x'_t , we can derive the term (a) as follows:

$$\begin{aligned}
(a) &= \sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \langle z_\star^{(i,j)}, \theta_\star \rangle - \langle z_t^{(i,j)}, \theta_\star \rangle \\
&= \sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \langle z_\star^{(i,j)}, \theta_\star \rangle - \mathbb{1}[i \in V_1 \wedge j \in V_2] \langle z_t^{(i,j)}, \theta_\star \rangle \\
&\quad - \mathbb{1}[i \in V_2 \wedge j \in V_1] \langle z_t^{(i,j)}, \theta_\star \rangle \\
&\quad - \mathbb{1}[i \in V_1 \wedge j \in V_1] \langle z_t^{(i,j)}, \theta_\star \rangle \\
&\quad - \mathbb{1}[i \in V_2 \wedge j \in V_2] \langle z_t^{(i,j)}, \theta_\star \rangle .
\end{aligned}$$

Notice that $\sum_{(i,j) \in E} z_\star^{(i,j)} = \sum_{(i,j) \in E} \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}$, so one has

$$\begin{aligned}
(a) &= \underbrace{\sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \mathbb{1}[i \in V_1 \wedge j \in V_2] \left(\left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle - \langle z_t^{(i,j)}, \theta_\star \rangle \right)}_{(a_1)} \\
&\quad + \underbrace{\sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \mathbb{1}[i \in V_2 \wedge j \in V_1] \left(\left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle - \langle z_t^{(i,j)}, \theta_\star \rangle \right)}_{(a_2)} \\
&\quad + \underbrace{\sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \mathbb{1}[i \in V_1 \wedge j \in V_1] \left(\left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle - \langle z_t^{(i,j)}, \theta_\star \rangle \right)}_{(a_3)} \\
&\quad + \underbrace{\sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \mathbb{1}[i \in V_2 \wedge j \in V_2] \left(\left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle - \langle z_t^{(i,j)}, \theta_\star \rangle \right)}_{(a_4)} .
\end{aligned}$$

Let us analyse the first term:

$$(a_1) = \sum_{t=1}^T \mathbb{1}[D_t] \sum_{i=1}^n \sum_{\substack{j \in N_i \\ j > i}} \mathbb{1}[i \in V_1 \wedge j \in V_2] \left\langle \frac{2}{m} \sum_{(k,l) \in E} z_\star^{(k,l)} - \left(z_t^{(i,j)} + z_t^{(j,i)} \right), \theta_\star \right\rangle . \quad (13)$$

By defining $(x_*, x'_*) = \arg \max_{(x, x') \in \mathcal{X}^2} \langle z_{xx'} + z_{x'x}, \theta_* \rangle$, and noticing that in the case where a node i is in V_1 and a neighbouring node j in is V_2 , then $z_t^{(i,j)} = z_{x_t x'_t}$, we have,

$$\begin{aligned}
\frac{2}{m} \sum_{(k,l) \in E} \langle z_*^{(k,l)}, \theta_* \rangle &= \frac{2}{m} \sum_{k=1}^n \sum_{\substack{j \in \mathcal{N}_k \\ j > k}} \langle z_*^{(k,l)} + z_*^{(l,k)}, \theta_* \rangle \\
&\leq \frac{2}{m} \sum_{k=1}^n \sum_{\substack{j \in \mathcal{N}_k \\ j > k}} \langle z_{x_* x'_*} + z_{x'_* x_*}, \theta_* \rangle \\
&= \langle z_{x_* x'_*} + z_{x'_* x_*}, \theta_* \rangle \\
&\leq \langle z_{x_t x'_t} + z_{x'_t x_t}, \tilde{\theta}_{t-1, m} \rangle \quad \text{w.p } 1 - \delta \\
&= \langle z_t^{(i,j)} + z_t^{(j,i)}, \tilde{\theta}_{t-1, m} \rangle .
\end{aligned}$$

Plugging this last inequality in [\(I3\)](#) yields, with probability $1 - \delta$,

$$\begin{aligned}
(a_1) &\leq \sum_{t=1}^T \mathbb{1}[D_t] \sum_{i=1}^n \sum_{\substack{j \in \mathcal{N}_i \\ j > i}} \mathbb{1}[i \in V_1 \wedge j \in V_2] \langle z_t^{(i,j)} + z_t^{(j,i)}, \tilde{\theta}_{t-1, m} - \theta_* \rangle \\
&= \sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \mathbb{1}[i \in V_1 \wedge j \in V_2] \langle z_t^{(i,j)}, \tilde{\theta}_{t-1, m} - \theta_* \rangle .
\end{aligned}$$

We define, as in [Algorithm 1](#), $\mathbb{1} \left[z_t^{(i,j)} = z_{x_t x'_t} \right] \triangleq \mathbb{1}[i \in V_1 \wedge j \in V_2]$. Then, one has, with probability $1 - \delta$,

$$\begin{aligned}
(a_1) &\leq \sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \mathbb{1} \left[z_t^{(i,j)} = z_{x_t x'_t} \right] \langle z_t^{(i,j)}, \tilde{\theta}_{t-1, m} - \theta_* \rangle \\
&= \sum_{t=1}^T \mathbb{1}[D_t] \sum_{k=1}^m \mathbb{1} \left[z_t^{E[k]} = z_{x_t x'_t} \right] \langle z_t^{E[k]}, \tilde{\theta}_{t-1, m} - \theta_* \rangle \\
&= \sum_{t=1}^T \mathbb{1}[D_t] \sum_{k=1}^m \mathbb{1} \left[z_t^{E[k]} = z_{x_t x'_t} \right] \langle z_t^{E[k]}, \tilde{\theta}_{t-1, m} - \hat{\theta}_{t-1, m} \rangle + \langle z_t^{E[k]}, \hat{\theta}_{t-1, m} - \theta_* \rangle \\
&\leq \sum_{t=1}^T \mathbb{1}[D_t] \sum_{k=1}^m \mathbb{1} \left[z_t^{E[k]} = z_{x_t x'_t} \right] \|z_t^{E[k]}\|_{\mathbf{A}_{t, k-1}^{-1}} \|\tilde{\theta}_{t-1, m} - \hat{\theta}_{t-1, m}\|_{\mathbf{A}_{t, k-1}} \\
&\quad + \mathbb{1} \left[z_t^{E[k]} = z_{x_t x'_t} \right] \|z_t^{E[k]}\|_{\mathbf{A}_{t, k-1}^{-1}} \|\hat{\theta}_{t-1, m} - \theta_*\|_{\mathbf{A}_{t, k-1}} \\
&\leq \sum_{t=1}^T \mathbb{1}[D_t] \sum_{k=1}^m \mathbb{1} \left[z_t^{E[k]} = z_{x_t x'_t} \right] \|z_t^{E[k]}\|_{\mathbf{A}_{t, k-1}^{-1}} \sqrt{2} \|\tilde{\theta}_{t-1, m} - \hat{\theta}_{t-1, m}\|_{\mathbf{A}_{t-1, m}} \quad (14) \\
&\quad + \mathbb{1} \left[z_t^{E[k]} = z_{x_t x'_t} \right] \|z_t^{E[k]}\|_{\mathbf{A}_{t, k-1}^{-1}} \sqrt{2} \|\hat{\theta}_{t-1, m} - \theta_*\|_{\mathbf{A}_{t-1, m}}
\end{aligned}$$

$$\leq \sum_{t=1}^T \sum_{k=1}^m \mathbb{1} \left[z_t^{E[k]} = z_{x_t x'_t} \right] 2\sqrt{2\beta_t(\delta)} \|z_t^{E[k]}\|_{\mathbf{A}_{t, k-1}^{-1}} \quad (15)$$

$$\leq \sum_{t=1}^T \sum_{k=1}^m 2\sqrt{2\beta_t(\delta)} \|z_t^{E[k]}\|_{\mathbf{A}_{t, k-1}^{-1}} , \quad (16)$$

with $\sqrt{\beta_t(\delta)} \leq \sigma \sqrt{d^2 \log \left(\frac{1+tmL^2/\lambda}{\delta} \right)} + \sqrt{\lambda}S$ and where (14) uses the critical covariance inequality (12), (15) comes from the definition of the confidence set $C_{t-1,m}(\delta)$ (9) and (16) upper bounds the indicator functions by 1.

Using a similar reasoning, we obtain the same bound for (a_2) :

$$(a_2) \leq \sum_{t=1}^T \sum_{k=1}^m 2\sqrt{2\beta_t(\delta)} \|z_t^{E[k]}\|_{\mathbf{A}_{t,k-1}^{-1}}. \quad (17)$$

Let us bound the terms (a_3) and (a_4) .

$$(a_3) = \sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \mathbb{1} \left[z_t^{(i,j)} = z_{x_t x_t} \right] \left(\left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle - \langle z_t^{(i,j)}, \theta_\star \rangle \right) \quad (18)$$

For all $x \in \mathcal{X}$, let γ_x be the following ratio

$$\gamma_x = \frac{\langle z_{xx}, \theta_\star \rangle}{\left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle}, \quad (19)$$

and let γ be the worst ratio

$$\gamma = \min_{x \in \mathcal{X}} \frac{\langle z_{xx}, \theta_\star \rangle}{\left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle}. \quad (20)$$

We have

$$\begin{aligned} (a_3) &= \sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \mathbb{1} \left[z_t^{(i,j)} = z_{x_t x_t} \right] \left(\left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle - \gamma_{x_t} \left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle \right) \\ &\leq \sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \mathbb{1} \left[z_t^{(i,j)} = z_{x_t x_t} \right] \left(\left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle - \gamma \left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle \right) \\ &= \sum_{t=1}^T \mathbb{1}[D_t] \sum_{(i,j) \in E} \mathbb{1} \left[z_t^{(i,j)} = z_{x_t x_t} \right] (1 - \gamma) \left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle \\ &\leq T \frac{m}{4} (1 - \gamma) \left\langle \frac{1}{m} \sum_{(k,l) \in E} z_\star^{(k,l)}, \theta_\star \right\rangle \\ &= \sum_{t=1}^T \sum_{(i,j) \in E} \frac{1}{4} (1 - \gamma) \langle z_\star^{(i,j)}, \theta_\star \rangle, \end{aligned} \quad (21)$$

where (21) comes from the fact that there is at most $m/4$ edges that goes from node in V_1 to other nodes in V_1 and that $\mathbb{1}[D_t] \leq 1$ for all t .

The derivation of this bound for (a_3) gives the same one for (a_4)

$$(a_4) \leq \sum_{t=1}^T \sum_{(i,j) \in E} \frac{1}{4} (1 - \gamma) \langle z_\star^{(i,j)}, \theta_\star \rangle. \quad (22)$$

By rewriting (a), we have :

$$(a) \leq \sum_{t=1}^T \sum_{k=1}^m 4\sqrt{2\beta_t(\delta)} \|z_t^{E[k]}\|_{\mathbf{A}_{t,k-1}^{-1}} + \frac{1}{2}(1-\gamma)\langle z_\star^{(i,j)}, \theta_\star \rangle .$$

In [\[Chan et al., 2021\]](#), they bounded the term (b) as follows

$$LSm \sum_{t=1}^T \mathbf{1}[D_t^c] \leq LSm \left[d^2 \log_2 \left(\frac{TmL^2/\lambda}{\delta} \right) \right] . \quad (23)$$

We thus have the regret bounded by

$$R(T) \leq \sum_{t=1}^T \sum_{k=1}^m 4\sqrt{2\beta_t(\delta)} \|z_t^{E[k]}\|_{\mathbf{A}_{t,k-1}^{-1}} + \frac{1}{2}(1-\gamma)\langle z_\star^{(i,j)}, \theta_\star \rangle + LSm \left[d^2 \log_2 \left(\frac{TmL^2/\lambda}{\delta} \right) \right] ,$$

which gives us

$$R_{\frac{1+\gamma}{2}}(T) \leq \sum_{t=1}^T \sum_{k=1}^m 4\sqrt{2\beta_t(\delta)} \|z_t^{E[k]}\|_{\mathbf{A}_{t,k-1}^{-1}} + LSm \left[d^2 \log_2 \left(\frac{TmL^2/\lambda}{\delta} \right) \right] .$$

Let us bound the first term with the double sum as it is done in [\[Abbasi-Yadkori et al., 2011\]](#), [\[Chan et al., 2021\]](#):

$$\begin{aligned} & \sum_{t=1}^T \sum_{k=1}^m 4\sqrt{2\beta_t(\delta)} \|z_t^{E[k]}\|_{\mathbf{A}_{t,k-1}^{-1}} \\ & \leq \sum_{t=1}^T \sum_{k=1}^m \min \left(2LS, 4\sqrt{2\beta_t(\delta)} \|z_t^{E[k]}\|_{\mathbf{A}_{t,k-1}^{-1}} \right) \\ & \leq \sum_{t=1}^T \sum_{k=1}^m 4\sqrt{2\beta_t(\delta)} \min \left(LS, \|z_t^{E[k]}\|_{\mathbf{A}_{t,k-1}^{-1}} \right) \\ & \leq \sqrt{Tm \times 32\beta_T(\delta) \sum_{t=1}^T \sum_{k=1}^m \min \left((LS)^2, \|z_t^{E[k]}\|_{\mathbf{A}_{t,k-1}^{-1}}^2 \right)} \\ & \leq \sqrt{32Tm\beta_T(\delta) \sum_{t=1}^T \sum_{k=1}^m \max(2, (LS)^2) \log \left(1 + \|z_t^{E[k]}\|_{\mathbf{A}_{t,k-1}^{-1}}^2 \right)} \quad (24) \end{aligned}$$

$$\begin{aligned} & = \sqrt{32Tm\beta_T(\delta) \max(2, (LS)^2) \sum_{t=1}^T \sum_{k=1}^m \log \left(1 + \|z_t^{E[k]}\|_{\mathbf{A}_{t,k-1}^{-1}}^2 \right)} \\ & \leq \sqrt{32Tm\beta_T(\delta) \max(2, (LS)^2) d^2 \log \left(1 + \frac{TmL^2/\lambda}{d^2} \right)} \quad (25) \\ & \leq \sqrt{32Tmd^2 \max(2, (LS)^2) \log \left(1 + \frac{TmL^2/\lambda}{d^2} \right)} \left(\sigma \sqrt{d^2 \log \left(\frac{1 + TmL^2/\lambda}{\delta} \right)} + \sqrt{\lambda}S \right) \end{aligned}$$

where (24) uses the fact that for all $a, x \geq 0$, $\min(a, x) \leq \max(2, a) \log(1 + x)$, (25) uses the fact that $\sum_{t=1}^T \sum_{k=1}^m \log \left(1 + \|z_t^{E[k]}\|_{\mathbf{A}_{t,k-1}^{-1}}^2 \right) \leq d^2 \log \left(1 + \frac{TmL^2/\lambda}{d^2} \right)$ from Lemma 19.4 in Lattimore and Szepesvári [2018].

The final bound for the $\frac{1+\gamma}{2}$ -regret is

$$R_{\frac{1+\gamma}{2}}(T) \leq \sqrt{32Tmd^2 \max(2, (LS)^2) \log \left(1 + \frac{TmL^2/\lambda}{d^2} \right)} \left(\sigma \sqrt{d^2 \log \left(\frac{1 + TmL^2/\lambda}{\delta} \right)} + \sqrt{\lambda}S \right) + LSm \left[d^2 \log_2 \left(\frac{TmL^2/\lambda}{\delta} \right) \right]$$

□

B.2 Proof of Theorem 4.1

Proof. For the sake of completeness in the proof we recall that we defined the couples (x_*, x'_*) and $(\tilde{x}_*, \tilde{x}'_*)$ and the quantity Δ as follows:

$$(x_*, x'_*) = \arg \max_{(x, x') \in \mathcal{X}^2} \langle z_{xx'} + z_{x'x}, \theta_* \rangle$$

$$(\tilde{x}_*, \tilde{x}'_*) = \arg \max_{(x, x') \in \mathcal{X}} \langle m_{1 \rightarrow 2} \cdot z_{xx'} + m_{2 \rightarrow 1} \cdot z_{x'x} + m_1 \cdot z_{xx} + m_2 \cdot z_{x'x'}, \theta_* \rangle .$$

and

$$\Delta = \langle m_{1 \rightarrow 2} (z_{\tilde{x}_* \tilde{x}'_*} - z_{x_* x'_*}) + m_{2 \rightarrow 1} (z_{\tilde{x}'_* \tilde{x}_*} - z_{x'_* x_*}) + m_1 (z_{\tilde{x}_* \tilde{x}_*} - z_{x_* x_*}) + m_2 (z_{\tilde{x}'_* \tilde{x}'_*} - z_{x'_* x'_*}), \theta_* \rangle .$$

And we recall that in Algorithm 3, the tuple $(x_t, x'_t, \tilde{\theta}_{t-1, m})$ is obtained as follows:

$$(x_t, x'_t, \tilde{\theta}_{t-1, m}) = \arg \max_{(x, x', \theta) \in \mathcal{X}^2 \times C_{t-1}} \langle m_{1 \rightarrow 2} \cdot z_{xx'} + m_{2 \rightarrow 1} \cdot z_{x'x} + m_1 \cdot z_{xx} + m_2 \cdot z_{x'x'}, \theta \rangle$$

We can write the regret $R(T)$ as in the proof of Theorem 3.2:

$$\begin{aligned} R(T) &= \sum_{t=1}^T \mathbf{1}[D_t] \sum_{(i,j) \in E} \langle z_*^{(i,j)}, \theta_* \rangle - \langle z_t^{(i,j)}, \theta_* \rangle + \mathbf{1}[D_t^c] \sum_{(i,j) \in E} \langle z_*^{(i,j)}, \theta_* \rangle - \langle z_t^{(i,j)}, \theta_* \rangle \\ &\leq \underbrace{\sum_{t=1}^T \mathbf{1}[D_t] \sum_{(i,j) \in E} \langle z_*^{(i,j)}, \theta_* \rangle - \langle z_t^{(i,j)}, \theta_* \rangle}_{(a)} + \underbrace{LSm \sum_{t=1}^T \mathbf{1}[D_t^c]}_{(b)} \end{aligned}$$

Here, (b) doesn't change, we thus only focus on deriving (a).

$$\begin{aligned} (a) &= \sum_{t=1}^T \mathbf{1}[D_t] \sum_{(i,j) \in E} \langle z_*^{(i,j)}, \theta_* \rangle - \langle z_t^{(i,j)}, \theta_* \rangle \\ &\leq \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_*^{(i,j)}, \theta_* \rangle - \langle z_t^{(i,j)}, \theta_* \rangle \quad (\text{where } \mathbf{1}[D_t] \leq 1) \\ &= \underbrace{\sum_{t=1}^T \sum_{(i,j) \in E} \frac{m_{1 \rightarrow 2} + m_{2 \rightarrow 1}}{m} \langle z_*^{(i,j)}, \theta_* \rangle}_{(a_1)} + \sum_{t=1}^T \sum_{(i,j) \in E} \frac{m_1 + m_2}{m} \langle z_*^{(i,j)}, \theta_* \rangle - \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \theta_* \rangle \end{aligned}$$

We have

$$\begin{aligned}
(a_1) &= \sum_{t=1}^T \sum_{(i,j) \in E} \frac{2m_{1 \rightarrow 2}}{m} \langle z_{\star}^{(i,j)}, \theta_{\star} \rangle \\
&= \sum_{t=1}^T \sum_{i=1}^n \sum_{\substack{j \in \mathcal{N}_i \\ j > i}} \frac{2m_{1 \rightarrow 2}}{m} \langle z_{\star}^{(i,j)} + z_{\star}^{(j,i)}, \theta_{\star} \rangle \\
&\leq \sum_{t=1}^T \sum_{i=1}^n \sum_{\substack{j \in \mathcal{N}_i \\ j > i}} \frac{2m_{1 \rightarrow 2}}{m} \langle z_{x_{\star} x'_i} + z_{x'_i x_{\star}}, \theta_{\star} \rangle \\
&= \sum_{t=1}^T \sum_{i=1}^n \sum_{\substack{j \in \mathcal{N}_i \\ j > i}} \frac{2}{m} \langle m_{1 \rightarrow 2} \cdot z_{x_{\star} x'_i} + m_{2 \rightarrow 1} \cdot z_{x'_i x_{\star}}, \theta_{\star} \rangle \\
&= \sum_{t=1}^T \sum_{i=1}^n \sum_{\substack{j \in \mathcal{N}_i \\ j > i}} \frac{2}{m} \langle m_{1 \rightarrow 2} \cdot z_{x_{\star} x'_i} + m_{2 \rightarrow 1} \cdot z_{x'_i x_{\star}} + m_1 \cdot z_{x_{\star} x_{\star}} + m_2 \cdot z_{x'_i x'_i}, \theta_{\star} \rangle \\
&\quad - \sum_{t=1}^T \sum_{i=1}^n \sum_{\substack{j \in \mathcal{N}_i \\ j > i}} \frac{2}{m} \langle m_1 \cdot z_{x_{\star} x_{\star}} + m_2 \cdot z_{x'_i x'_i}, \theta_{\star} \rangle \\
&= \sum_{t=1}^T \sum_{i=1}^n \sum_{\substack{j \in \mathcal{N}_i \\ j > i}} \frac{2}{m} \langle m_{1 \rightarrow 2} \cdot z_{\tilde{x}_{\star} \tilde{x}'_i} + m_{2 \rightarrow 1} \cdot z_{\tilde{x}'_i \tilde{x}_{\star}} + m_1 \cdot z_{\tilde{x}_{\star} \tilde{x}_{\star}} + m_2 \cdot z_{\tilde{x}'_i \tilde{x}'_i}, \theta_{\star} \rangle - \frac{2}{m} \Delta \\
&\quad - \sum_{t=1}^T \sum_{i=1}^n \sum_{\substack{j \in \mathcal{N}_i \\ j > i}} \frac{2}{m} \langle m_1 \cdot z_{x_{\star} x_{\star}} + m_2 \cdot z_{x'_i x'_i}, \theta_{\star} \rangle \\
&= \sum_{t=1}^T \langle m_{1 \rightarrow 2} \cdot z_{\tilde{x}_{\star} \tilde{x}'_t} + m_{2 \rightarrow 1} \cdot z_{\tilde{x}'_t \tilde{x}_{\star}} + m_1 \cdot z_{\tilde{x}_{\star} \tilde{x}_{\star}} + m_2 \cdot z_{\tilde{x}'_t \tilde{x}'_t}, \theta_{\star} \rangle - \Delta \\
&\quad - \sum_{t=1}^T \langle m_1 \cdot z_{x_{\star} x_{\star}} + m_2 \cdot z_{x'_t x'_t}, \theta_{\star} \rangle \\
&\leq \sum_{t=1}^T \langle m_{1 \rightarrow 2} \cdot z_{x_t x'_t} + m_{2 \rightarrow 1} \cdot z_{x'_t x_t} + m_1 \cdot z_{x_t x_t} + m_2 \cdot z_{x'_t x'_t}, \tilde{\theta}_{t-1, m} \rangle - \Delta \quad \text{w.p } 1 - \delta \\
&\quad - \sum_{t=1}^T \langle m_1 \cdot z_{x_{\star} x_{\star}} + m_2 \cdot z_{x'_t x'_t}, \theta_{\star} \rangle \\
&= \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \tilde{\theta}_{t-1, m} \rangle - \sum_{t=1}^T \Delta - \sum_{t=1}^T \langle m_1 \cdot z_{x_{\star} x_{\star}} + m_2 \cdot z_{x'_t x'_t}, \theta_{\star} \rangle
\end{aligned}$$

By plugging the last upper bound in (a) and with probability $1 - \delta$, we have,

$$\begin{aligned}
(a) &\leq \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \tilde{\theta}_{t-1,m} \rangle - \sum_{t=1}^T \Delta - \sum_{t=1}^T \langle m_1 \cdot z_{x_* x_*} + m_2 \cdot z_{x'_* x'_*}, \theta_* \rangle \\
&\quad + \sum_{t=1}^T \sum_{(i,j) \in E} \frac{m_1 + m_2}{m} \langle z_*^{(i,j)}, \theta_* \rangle - \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \theta_* \rangle \\
&= \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \tilde{\theta}_{t-1,m} - \theta_* \rangle - \sum_{t=1}^T \Delta - \sum_{t=1}^T \langle m_1 \cdot z_{x_* x_*} + m_2 \cdot z_{x'_* x'_*}, \theta_* \rangle \\
&\quad + \sum_{t=1}^T \sum_{(i,j) \in E} \frac{m_1 + m_2}{m} \langle z_*^{(i,j)}, \theta_* \rangle \\
&= \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \tilde{\theta}_{t-1,m} - \theta_* \rangle - \sum_{t=1}^T \Delta - \sum_{t=1}^T \sum_{(i,j) \in E} \frac{m_1}{m} \gamma_{x_*} \langle z_*^{(i,j)}, \theta_* \rangle + \frac{m_2}{m} \gamma_{x'_*} \langle z_*^{(i,j)}, \theta_* \rangle \\
&\quad + \sum_{t=1}^T \sum_{(i,j) \in E} \frac{m_1 + m_2}{m} \langle z_*^{(i,j)}, \theta_* \rangle \\
&\leq \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \tilde{\theta}_{t-1,m} - \theta_* \rangle - \sum_{t=1}^T \Delta - \sum_{t=1}^T \sum_{(i,j) \in E} \frac{m_1 + m_2}{m} \gamma \langle z_*^{(i,j)}, \theta_* \rangle + \sum_{t=1}^T \sum_{(i,j) \in E} \frac{m_1 + m_2}{m} \langle z_*^{(i,j)}, \theta_* \rangle \\
&= \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \tilde{\theta}_{t-1,m} - \theta_* \rangle - \sum_{t=1}^T \Delta + \sum_{t=1}^T \sum_{(i,j) \in E} \frac{m_1 + m_2}{m} (1 - \gamma) \langle z_*^{(i,j)}, \theta_* \rangle \\
&= \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \tilde{\theta}_{t-1,m} - \theta_* \rangle - \sum_{t=1}^T \sum_{(i,j) \in E} \epsilon \langle z_*^{(i,j)}, \theta_* \rangle + \sum_{t=1}^T \sum_{(i,j) \in E} \frac{m_1 + m_2}{m} (1 - \gamma) \langle z_*^{(i,j)}, \theta_* \rangle \\
&= \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \tilde{\theta}_{t-1,m} - \theta_* \rangle + \sum_{t=1}^T \sum_{(i,j) \in E} \left[\frac{m_1 + m_2}{m} (1 - \gamma) - \epsilon \right] \langle z_*^{(i,j)}, \theta_* \rangle
\end{aligned}$$

By plugging (a) in the regret and with probability $1 - \delta$, we have,

$$R(T) \leq \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \tilde{\theta}_{t-1,m} - \theta_* \rangle + \sum_{t=1}^T \sum_{(i,j) \in E} \left[\frac{m_1 + m_2}{m} (1 - \gamma) - \epsilon \right] \langle z_*^{(i,j)}, \theta_* \rangle + LSm \sum_{t=1}^T \mathbb{1}[D_t^c]$$

which gives,

$$\begin{aligned}
R(T) - \sum_{t=1}^T \sum_{(i,j) \in E} \left[\frac{m_1 + m_2}{m} (1 - \gamma) - \epsilon \right] \langle z_*^{(i,j)}, \theta_* \rangle &\leq \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \tilde{\theta}_{t-1,m} - \theta_* \rangle + LSm \sum_{t=1}^T \mathbb{1}[D_t^c] \\
R_{1 - \left[\frac{m_1 + m_2}{m} (1 - \gamma) - \epsilon \right]}(T) &\leq \sum_{t=1}^T \sum_{(i,j) \in E} \langle z_t^{(i,j)}, \tilde{\theta}_{t-1,m} - \theta_* \rangle + LSm \sum_{t=1}^T \mathbb{1}[D_t^c]
\end{aligned}$$

The upper bound of the right hand term follows exactly what we have already done for Theorem [3.2](#) by applying the upper bounds [\(16\)](#) and [\(23\)](#) \square

C Additional information on the experiments

C.1 Table 1

The number of nodes in each graph is equal to 100. The random graph corresponds to a graph where for two nodes i and j in V , the probability that (i, j) and (j, i) is in E is equal to 0.6. The results for the random graph are averaged over 100 draws. The matching graph represents the graph where each node $i \in V$ has only one neighbour: $\text{Card}(\mathcal{N}_i) = 1$.

C.2 Figure 1

The graph used in this experiment is a complete graph of 10 nodes. The arm set $\mathcal{X} = \{e_1, \dots, e_d\}$ which gives $\mathcal{Z} = \{e_1, \dots, e_{d^2}\}$. The matrix \mathbf{M}_* is randomly initialized such that all elements of the matrix are drawn i.i.d. from a standard normal distribution, and then we take the absolute value of each of these elements to ensure that the matrix only contains positive numbers. We plotted the results by varying ζ from 0 to 1 with a step of 0.01. We conducted the experiment on 100 different matrices \mathbf{M}_* randomly initialized as explained above and plotted the average value of the obtained γ , ϵ , α_1 and α_2 .

C.3 Figure 2

For the last experiment, we used a complete graph of 5 nodes. The arm set $\mathcal{X} = \{e_1, \dots, e_d\}$ which gives $\mathcal{Z} = \{e_1, \dots, e_{d^2}\}$. The matrix \mathbf{M}_* is randomly initialized as explained in the previous experiment. We fixed $\zeta = 0$ and the horizon $T = 20000$. We ran the experiment 10 times and plotted the average values (shaded curve) and the moving average curve with a window of 100 steps for more clarity.

The Explore-Then-Commit algorithm has an exploration phase of $T/3$ rounds and then exploits by pulling the couple $(x_t, x'_t) = \arg \max_{(x, x')} \langle z_{xx'} + z_{x'x}, \hat{\theta}_t \rangle$. Note that we set the exploration phase to $T/3$ because most of the time, it was sufficient for the learner to have the estimated optimal pair (x_t, x'_t) equal to the real optimal pair (x_*, x'_*) .

Machine used for all the experiments : Macbook Pro, Apple M1 chip, 8-core CPU

The code is available [here](#).