

## A Supplementary Material

### A.1 Proof of Lemma 1

Before proceeding with a proof of this statement, we recall the famous Hoeffding inequality [17] which can be applied to bounded independent random variables, regardless of their distribution.

**Proposition 1** (Hoeffding's inequality for bounded variables). *Let  $\xi_1, \dots, \xi_n$  be independent random variables such that  $a \leq \xi_i \leq b$  for all  $i \in [n]$ . Set  $\xi = \sum \xi_i$  and  $\mu = \mathbf{E}[\xi]$ . Then, for all  $\delta > 0$*

$$\Pr \{|\xi - \mu| \leq n\delta\} \geq 1 - 2 \exp(-2\delta^2 n / (b - a)^2).$$

*Proof.* Let  $Q \subseteq [n]$ . Let  $\mathbf{x} = (x_1, \dots, x_n)$  be a vector such that  $x_i$  has distribution  $U(0, 1)$  for  $i \in Q$  and  $x_i = 0$  otherwise. Define the random variable  $\theta := \text{sign}(\mathbf{x}, \beta_1) - \text{sign}(-\mathbf{x}, \beta_2)$ , where  $\beta_1$  and  $\beta_2$  are i.i.d. random variables taken uniformly at random from the set of unknown vectors  $\mathcal{B}$ . Clearly,  $\mathbf{E}[\theta] = \frac{2}{\ell} \text{pir}(Q, \mathcal{B})$ . From the lemma condition, it follows that there exists a set  $I \subseteq [m]$  of size at least  $(1 - \frac{1}{16\ell})m$  such that  $y_i = \text{sign}(\langle \mathbf{x}^{(i)}, \beta \rangle)$  for all  $i \in I$ . Moreover, one can find a subset  $I' \subseteq I$  such that  $|\{i \in I' : i \leq m/2\}| = |\{i \in I' : i > m/2\}| = \frac{8\ell-1}{16\ell}m$ . By applying Proposition 1 for a set  $\{y_i : i \in I', i \leq m/2\} \cup \{y_i : i \in I', i > m/2\}$  with  $\xi := \sum_{i \in I'} y_i$ ,  $\mu = \mathbf{E}[\xi] = \frac{|I'|}{2} \frac{2}{\ell} \text{pir}(Q, \mathcal{B}) = \frac{(8\ell-1)m}{8\ell^2} \text{pir}(Q, \mathcal{B})$ ,  $a = -1$ ,  $b = 1$  and  $\delta = 1/(2\ell)$ , we get

$$\Pr \{|\xi - \mu| \leq n\delta\} \geq 1 - 2 \exp(-\Omega(m\ell^{-2})).$$

Thus,

$$\begin{aligned} & \Pr \left\{ \widetilde{\text{pir}}(Q, m) = \text{pir}(Q, \mathcal{B}) \right\} = \\ & = \Pr \left\{ \text{pir}(Q, \mathcal{B}) - \frac{1}{2} < \frac{\ell}{m} \left( \sum_{i=1}^{m/2} y_i - \sum_{i=m/2+1}^m y_i \right) < \text{pir}(Q, \mathcal{B}) + \frac{1}{2} \right\} \\ & \geq \Pr \{|\xi - \mu| \leq n\delta\} \\ & \geq 1 - 2 \exp(-\Omega(m\ell^{-2})). \end{aligned}$$

This yields that  $\widetilde{\text{pir}}(Q, m)$  is a correct estimate of  $\text{pir}(Q, \mathcal{B})$  with probability  $1 - \exp(-\Omega(m\ell^{-2}))$ .  $\square$

### A.2 Proof of Lemma 2

*Proof.* By Lemma 1, one value  $z(Q_i, H_{\mathcal{B}})$  can be correctly estimated with probability  $1 - 1/(nm)$  by asking  $q'$ ,  $q' = O(\ell^2 \log(nm))$ , queries provided that a fraction of the erroneous responses to these queries is at most  $1/(16\ell)$ . We ask  $q = mq'$  queries and try to estimate all values  $\{z(Q_i, H_{\mathcal{B}}), i \in [m]\}$ . Recall that the total number of erroneous responses is bounded by  $\tau q$ . By the union bound, with probability  $1 - 1/n$ , the number of incorrect estimates is at most  $\frac{\tau q}{q'/(16\ell)} = 16\ell\tau m$ . This completes the proof.  $\square$

### A.3 Proof of Lemma 6

*Proof.* We shall make use of the random coding with expurgation technique. Let  $r := \alpha m$ . Consider a random binary matrix  $X$  of size  $m \times n$  whose entries are i.i.d. Bernoulli random variables which take the value 1 with probability  $p = \frac{c}{c+f}$  and the value 0 with probability  $1 - p$ . Let  $\mathcal{E}_t$  denote the event that the  $t$ -th column of  $X$  is *bad*, i.e., there exist some disjoint subsets  $U \subseteq [n]$ ,  $t \in U$ ,  $|U| = c$ , and  $W \subseteq [n]$ ,  $|W| = f$ , and less than  $2r + 1$  rows  $i \in [m]$  such that  $X_{i,j} = 1$  for all  $j \in U$  and  $X_{i,j} = 0$  for all  $j \in W$ . Let  $\nu$  be a Binomial random variable with parameters  $m$  and  $\bar{p} := p^c(1-p)^f$ . It is clear that  $\nu$  describes the number of rows in  $X$  that separate two sets of columns

$U$  and  $W$ . By the union bound, we obtain

$$\begin{aligned} \Pr\{\mathcal{E}_t\} &\leq \sum_{\substack{U, W \subseteq [n], t \in U \\ |U|=c, |W|=f, U \cap W = \emptyset}} \Pr\{\nu \leq 2r\} \\ &\leq \binom{n-1}{c+f-1} \binom{c+f-1}{f} \sum_{i=0}^{2r} \binom{m}{i} \bar{p}^i (1-\bar{p})^{m-i}. \end{aligned}$$

Note that the property on  $\alpha$  in the statement of this lemma implies that  $2r \leq \bar{p}m$ . Applying the monotonicity property of the Binomial distribution, we obtain

$$\Pr\{\mathcal{E}_t\} \leq (2r+1)n^{c+f-1} \binom{m}{2r} (1-\bar{p})^{m-2r} \bar{p}^{2r}.$$

If the above probability is bounded above by  $1/2$ , then  $X$  contains on average at most  $n/2$  bad columns. This would imply the existence of an  $(c, f, \alpha)$ -robust-cover-free matrix of size  $m \times n/2$ . In order to have  $\Pr\{\mathcal{E}_t\} \leq 1/2$ , it suffices to take  $m$  satisfying

$$(c+f-1) \log n + mH_2(2\alpha) + m(1-2\alpha) \log(1-\bar{p}) + 2\alpha m \log \bar{p} \leq o(m).$$

Thus, a required matrix exists if it holds that

$$m \geq \frac{(c+f-1) \log n}{(2\alpha-1) \log(1-\bar{p}) - 2\alpha \log \bar{p} - H_2(2\alpha)} (1 + o(1)).$$

□

#### A.4 Proof of Lemma 9

*Proof.* To make everything work, we indeed need to use some assumptions regarding the support of vectors and the magnitude of all the entries: (a) the support of any vector from  $\mathcal{B}$  is not fully included to the support of any other vector from  $\mathcal{B}$ , (b) the absolute value of each non-zero entry of  $\beta^{(i)}$  is bounded below by  $c_l$  and above by  $c_u$ .

Under Assumption 1, for  $j \neq i$  and large enough  $\text{Inf}$ , it holds that

$$s_{ij} := \text{sign}(\langle \text{Inf}^{(i)}(\mathbf{g}), \beta^{(j)} \rangle) = \text{sign}(\langle \text{Inf}^{(i)}(\mathbf{0}), \beta^{(j)} \rangle),$$

since the contribution of entries indexed by elements of  $\text{supp}(\beta^{(i)})$  to both inner products is negligible as  $\text{Inf} \rightarrow \infty$ . We also note that

$$\begin{aligned} \text{sign}(\langle \text{Inf}^{(i)}(\mathbf{0}), \beta^{(i)} \rangle) &= \text{sign}(\langle \mathbf{0}, \beta^{(i)} \rangle) = 1, \\ \text{sign}(\langle \text{Inf}^{(i)}(\mathbf{g}), \beta^{(i)} \rangle) &= \text{sign}(\langle \mathbf{g}, \beta^{(i)} \rangle), \end{aligned}$$

Set  $\xi = \sum_{j=1}^{m/2} y_j - \sum_{j=m/2+1}^m y_j$ . Then  $\mu = \mathbf{E}[\xi] = \frac{m}{2\ell} (\text{sign}(\langle \mathbf{g}, \beta^{(i)} \rangle) - 1)$ . By applying Proposition 1, we obtain

$$\begin{aligned} &\Pr \left\{ \text{sign}(\langle \mathbf{g}, \beta^{(i)} \rangle) - \frac{1}{2} < \frac{2\ell}{m} \xi + 1 < \text{sign}(\langle \mathbf{g}, \beta^{(i)} \rangle) + \frac{1}{2} \right\} \\ &= \Pr \left\{ \frac{m}{2\ell} \left( \text{sign}(\langle \mathbf{g}, \beta^{(i)} \rangle) - \frac{1}{2} \right) < \xi + \frac{m}{2\ell} < \frac{m}{2\ell} \left( \text{sign}(\langle \mathbf{g}, \beta^{(i)} \rangle) + \frac{1}{2} \right) \right\} \\ &= \Pr \left\{ \left| \xi - \mu \right| < \frac{m}{4\ell} \right\} \\ &\geq 1 - 2 \exp(-\Omega(\ell^{-2}m)). \end{aligned}$$

This yields the required statement. □