

## A MenuGap( $X$ ) = 1 when $k = 1$

In this brief section we prove that when  $k = 1$ , for any sequence of  $x_i \in \mathbb{R}_{\geq 0}^+$ ,  $\text{MenuGap}(X) = 1$ .

**Claim 28.** When  $k = 1$ , for any  $X = \{x_i\}_{i=1}^N$ ,  $x_i \in \mathbb{R}_{\geq 0}^+$ ,  $\text{MenuGap}(X) = 1$ .

*Proof.* Note that when  $k = 1$ ,  $\|x_i\|_1 = x_i$ . Therefore,  $\text{MenuGap}(X, Q) = \sum_i \min_{j < i} (q_i - q_j)$ . We make the following observation which allows us to look at structured optimal solutions.

**Observation 29.** Any optimal solution  $Q$  to  $\text{MenuGap}(X)$  is monotone non-decreasing.

*Proof.* For the sake of contradiction, suppose we are given an optimal solution  $Q$  that is not monotone non-decreasing. Let  $i$  be the smallest index for which  $q_i < q_{i-1}$ . Then  $\text{gap}_i^{X, Q} = (q_i - q_{i-1})x_i < 0$ . Consider instead a solution  $Q'$  where  $q'_j = q_j$  for all  $j \neq i$  and  $q'_i = q_{i-1}$ . Now,  $\text{gap}_i^{X, Q'} = 0$ . Since  $Q_{\leq i-1} = Q'_{\leq i-1}$ ,  $\text{gap}_j^{X, Q} = \text{gap}_j^{X, Q'}$  for all  $j < i$ . Since  $q_{i-1} > q_i$ , for any  $j > i$  it holds that  $(q_j - q_{i-1}) < (q_j - q_i)$ . Therefore,  $q_i$  is not “setting the gap” for any point after it. Hence it also holds that  $\text{gap}_j^{X, Q} = \text{gap}_j^{X, Q'}$  for all  $j > i$ . Putting everything together we get that  $\text{MenuGap}(X, Q') - \text{MenuGap}(X, Q) = \text{gap}_i^{X, Q'} - \text{gap}_i^{X, Q} > 0$  contradicting the optimality of  $Q$ .  $\square$

With this observation in hand, since the  $q_i$  are monotone non-decreasing, without loss of generality it holds that  $\text{gap}_i^{X, Q} = \min_{j < i} q_i - q_j = q_i - q_{i-1}$  ( $q_{i-1} \geq q_j$  for all  $j < i$ ). Therefore, we get  $\text{MenuGap}(X, Q) = \sum_i q_i - q_{i-1} = q_N - q_0$ . Since  $q_0 = 0$  and  $0 \leq q_N \leq 1$ , we get that  $\text{MenuGap}(X, Q) \leq 1$ .

Finally, note that for any  $X$ , we can set  $q_N = 1$  and  $q_i = 0$  for all other  $i$ , proving that  $\text{MenuGap}(X) \geq 1$ .  $\square$

## B Omitted Proofs

*Proof of Lemma 8.* We prove that for all  $X, C$ ,  $\text{AlignGap}(X, C) \leq \text{MenuGap}(X)$ , which implies the lemma. For a given  $X, C$ , define:

- $\vec{q}_i := c_i \cdot \vec{x}_i$ , if  $\text{sgap}_i^{X, C} > 0$ .
- $\vec{q}_i := \arg \max_{j < i} \{c_j \cdot \vec{x}_j\}$ , if  $\text{sgap}_i^{X, C} \leq 0$ .

Observe first that each  $\vec{q}_i \in [0, 1]^k$ , as each  $c_i \vec{x}_i \in [0, 1]^k$  (this follows because each component of  $\vec{x}_i$  is at most  $\|\vec{x}_i\|_\infty$ , and each  $c_i$  is at most  $1/\|\vec{x}_i\|_\infty$ ). Next, observe that if  $\text{sgap}_i^{X, C} \leq 0$ , then  $\text{gap}_i^{X, Q} = 0$ . This is by definition in bullet two above. Finally, observe that if  $\text{sgap}_i^{X, C} > 0$ , then  $\text{gap}_i^{X, Q} \geq \text{sgap}_i^{X, C}$ . This is because the set of  $\{\vec{q}_j\}_{j < i}$  is a subset of  $\{c_j \vec{x}_j\}_{j < i}$ , and because  $\vec{q}_i := c_i \cdot \vec{x}_i$  by bullet one. Therefore,  $\text{gap}_i^{X, Q} \geq \max\{0, \text{sgap}_i^{X, C}\}$  for all  $i$  and the lemma follows.  $\square$

*Proof of Claim 13.* Take  $M'$  to be exactly the same as  $M$ , except having removed all entries with price  $< c$ . For every value in the support of  $\mathcal{D}$  with  $p^M(\vec{v}) \geq c$  in  $M$ , we still have  $p^{M'}(\vec{v}) \geq c$ . This is simply  $\vec{v}$ 's favorite option in  $M$  is still available in  $M'$ , and all options in  $M'$  were also available in  $M$ . For any value with  $p^M(\vec{v}) < c$ , we clearly have  $p^{M'}(\vec{v}) \geq 0$ . So for all  $\vec{v}$ , we have  $p^{M'}(\vec{v}) \geq p^M(\vec{v}) - c$ , and the claim follows by taking an expectation with respect to  $\vec{v}$ .  $\square$

*Proof of Claim 15.* Simply let  $M_1$  denote the set of menu options from  $M$  whose price lies in  $[c \cdot 2^i, c \cdot 2^{i+1})$  for an odd integer  $i$ , and  $M_2$  denote the remaining menu options (which lie in  $[c \cdot 2^i, c \cdot 2^{i+1})$  for an even power of  $i$ ). It is easy to see that  $M_1$  is oddly-priced and  $M_2$  is evenly-priced. Then for all  $\vec{v}$ , we must have  $p^{M_1}(\vec{v}) + p^{M_2}(\vec{v}) \geq p^M(\vec{v})$ . This is because  $\vec{v}$ 's favorite menu

option from  $M$  appears in one of the two menus, and is necessarily  $\vec{v}$ 's favorite option on that menu (and they pay non-zero from the other menu). Taking an expectation with respect to  $\vec{v}$  yields that  $\text{Rev}(\mathcal{D}, M_1) + \text{Rev}(\mathcal{D}, M_2) \geq \text{Rev}(\mathcal{D}, M)$ , completing the proof.  $\square$

*Proof of Claim 18.* Recall that  $(1 + \varepsilon) \cdot \|\vec{v}\|_1 \geq \|\vec{x}_i\|_1$  for all  $\vec{v} \in B_i$ . Therefore, if we set a price of  $\|\vec{x}_i\|_1 / (1 + \varepsilon)$  for the grand bundle, every  $\vec{v} \in B_i$  would choose to purchase the grand bundle. This immediately implies the claim, as:  $\text{BRev}(\mathcal{D}) \geq \frac{\|\vec{x}_i\|_1}{1 + \varepsilon} \cdot \Pr_{\vec{v} \sim \mathcal{D}} \left[ \|\vec{v}\|_1 \geq \frac{\|\vec{x}_i\|_1}{1 + \varepsilon} \right] \geq \frac{\|\vec{x}_i\|_1}{1 + \varepsilon} \cdot \Pr_{\vec{v} \sim \mathcal{D}}[\vec{v} \in B_i]$ .  $\square$

*Proof of Claim 19.* Recall that  $\text{gap}_i^{X,Q} := \min_{j < i} \{\vec{x}_i \cdot (\vec{q}_i - \vec{q}_j)\}$ , and that  $\vec{q}_i := \vec{q}^M(\vec{x}_i)$ . For any fixed  $j < i$ , recall that because  $M$  was a truthful mechanism, we must have:

$$\begin{aligned} \vec{x}_i \cdot \vec{q}^M(\vec{x}_i) - p^M(\vec{x}_i) &\geq \vec{x}_i \cdot \vec{q}^M(\vec{x}_j) - p^M(\vec{x}_j) \\ \Rightarrow \vec{x}_i \cdot (\vec{q}_i - \vec{q}_j) &\geq p^M(\vec{x}_i) - p^M(\vec{x}_j) \\ \Rightarrow \vec{x}_i \cdot (\vec{q}_i - \vec{q}_j) &\geq p^M(\vec{x}_i)/2. \end{aligned}$$

The first line is simply restating incentive compatibility. The second line is basic algebra, and substituting  $\vec{q}_i := \vec{q}^M(\vec{x}_i)$ . The third line invokes the fact that  $p^M(\vec{x}_i) \geq 2^{2(i-1)+a}$ , while  $p^M(\vec{x}_j) < 2^{2(j-1)+a+1} \leq 2^{2(i-1)+a-1}$ .  $\square$

*Proof of Observation 21.* This follows immediately from weak Lagrangian duality. For a quick refresher on weak Lagrangian duality, observe that for any feasible solution to the LP defining  $\text{AlignGap}'(X)$  we must have  $\vec{x}_i \cdot (c_i \vec{x}_i - c_{i-1} \vec{x}_{i-1}) - \text{sgap}_i \geq 0$ . Therefore, for any feasible solution to the original LP, that solution is also feasible for  $\text{LagRel}_1(X)$ , and the objective is only larger. Therefore, the optimal solution to  $\text{LagRel}_1(X)$  must be at least as large as  $\text{AlignGap}'(X)$ .  $\square$

*Proof of Observation 22.* For all  $i$ ,  $\max\{0, \text{sgap}_i\} - \text{sgap}_i \leq 0$ . When  $\text{sgap}_i = 0$ , the maximum is achieved (and  $\text{sgap}_i := 0$  is feasible). Substituting  $\max\{0, \text{sgap}_i\} - \text{sgap}_i = 0$  for all  $i$  concludes the proof.  $\square$

*Proof of Proposition 27.* To ease notation throughout the proof, we'll use the notation  $\text{gap}_{\ell,j}^{X,Q} := \text{gap}_i^{X,Q}$ , where  $\vec{x}_i := \vec{x}_{\ell,j}$  ( $\vec{x}_i$  is the  $j^{\text{th}}$  point on layer  $\ell$ ). We will also use the notation  $(\ell', j') < (\ell, j)$  if  $\ell' < \ell$ , or  $\ell' = \ell$  and  $j' < j$  (that is, if the  $j^{\text{th}}$  point in the  $\ell^{\text{th}}$  layer comes before the  $j^{\text{th}}$  point in the  $\ell^{\text{th}}$  layer). To understand  $\text{gap}_{\ell,j}^{X,Q}$ , we need to understand which point ‘‘sets the gap’’ for  $\vec{x}_{\ell,j}$ , that is, which  $(\ell', j') := \arg \min_{(\ell', j') < (\ell, j)} \{(\vec{q}_{\ell',j'} - \vec{q}_{\ell,j}) \cdot \vec{x}_i\}$ .

We first analyze which point sets the gap for  $\vec{x}_{\ell,j}$  (for even  $\ell$ ; for odd  $\ell$  the gap is zero and we don't care which point sets it), and observe that it must either be  $\vec{q}_{\ell,j-1}$  or  $\vec{q}_{\ell-2, n_{\ell-2}-1}$  (that is, it must be the previous point in the same layer, or the final point in the previous even layer).

**Claim 30.** For all  $j$ , and all even  $\ell$ ,  $\text{gap}_{\ell,j}^{X,Q} = \vec{x}_{\ell,j} \cdot \vec{q}_{\ell,j} - \max\{\vec{x}_{\ell,j} \cdot \vec{q}_{\ell-2, n_{\ell-2}-1}, \vec{x}_{\ell,j} \cdot \vec{q}_{\ell,j-1}\}$ .<sup>11</sup>

*Proof of Claim 30.* First, note that  $\text{gap}_{\ell,j}^{X,Q} := \min_{(\ell', j') < (\ell, j)} \{\vec{x}_{\ell,j} \cdot (\vec{q}_{\ell',j'} - \vec{q}_{\ell,j})\} = \vec{x}_{\ell,j} \cdot \vec{q}_{\ell,j} - \max_{(\ell', j') < (\ell, j)} \{\vec{x}_{\ell,j} \cdot \vec{q}_{\ell',j'}\}$ . To conclude the proof, simply observe that the first component of  $\vec{q}_{\ell',j'}$  is monotone increasing in  $\ell'$  (for fixed  $j'$ ), and the second component is monotone increasing in  $j'$  (for fixed  $\ell'$ ). Moreover, the second component of  $\vec{q}_{\ell', n_{\ell'}-1}$  is 1, and this is the maximum possible. Also, both components of  $\vec{x}_{\ell,j}$  are non-negative, and therefore we conclude that  $\vec{x}_{\ell,j} \cdot \vec{q}_{\ell-2, n_{\ell-2}-1} \geq \vec{x}_{\ell,j} \cdot \vec{q}_{\ell',j'}$  whenever  $(\ell', j') \leq (\ell-2, n_{\ell-2}-1)$  (in fact, this extends even to  $(\ell', j') \leq (\ell-1, n_{\ell-1}-1)$  as no new  $\vec{q}$  are introduced in layer  $\ell-1$ ). Also,  $\vec{x}_{\ell,j} \cdot \vec{q}_{\ell,j-1} \geq \vec{x}_{\ell,j} \cdot \vec{q}_{\ell',j'}$  whenever  $j' \leq j-1$ .  $\square$

Now that we know that the gap is set either by the last point in the previous layer, or the previous point in the current layer, we can nail down  $\text{gap}_{\ell,j}^{X,Q}$  exactly.

**Lemma 31.** For all even  $\ell > 2$ , and all  $j \in [0, n_\ell - 1]$ :  $\text{gap}_{\ell,j}^{X,Q} \geq \delta_\ell \frac{\sin(\theta_\ell)}{\sin((j+1)\theta_\ell)}$ .

<sup>11</sup>For simplicity of notation, define  $\vec{q}_{0,j} = \vec{0} = \vec{q}_{\ell,-1}$  for all  $\ell, j$ .

*Proof of Lemma 31.* To prove the lemma, we simply compute the inner product of  $\vec{x}_{\ell,j}$  with the three relevant vectors  $\vec{q}_{\ell,j}, \vec{q}_{\ell-2, n_{\ell-2}-1}, \vec{q}_{\ell,j-1}$ . To this end, recall that:

$$\begin{aligned}\vec{q}_{\ell,j} &= (z_{\ell}, 1 - \delta_{\ell} \cot((j+1)\theta_{\ell})), \\ \vec{q}_{\ell,j-1} &= (z_{\ell}, 1 - \delta_{\ell} \cot(j\theta_{\ell})), \\ \vec{q}_{\ell-2, n_{\ell-2}-1} &= (z_{\ell-2}, 1).\end{aligned}$$

Therefore, observe that

$$\begin{aligned}\vec{x}_{\ell,j} \cdot (\vec{q}_{\ell,j} - \vec{q}_{\ell,j-1}) &= \sin(j\theta_{\ell}) \cdot \delta_{\ell} \cdot (\cot(j\theta_{\ell}) - \cot((j+1)\theta_{\ell})) \\ &= \sin(j\theta_{\ell}) \cdot \delta_{\ell} \cdot \left( \frac{\cos(j\theta_{\ell})}{\sin(j\theta_{\ell})} - \frac{\cos((j+1)\theta_{\ell})}{\sin((j+1)\theta_{\ell})} \right) \\ &= \delta_{\ell} \cdot \frac{\cos(j\theta_{\ell}) \sin((j+1)\theta_{\ell}) - \sin(j\theta_{\ell}) \cos((j+1)\theta_{\ell})}{\sin((j+1)\theta_{\ell})} \\ &= \delta_{\ell} \cdot \frac{\sin(\theta_{\ell})}{\sin((j+1)\theta_{\ell})}.\end{aligned}$$

Similarly,

$$\begin{aligned}\vec{x}_{\ell,j} \cdot (\vec{q}_{\ell,j} - \vec{q}_{\ell-2, n_{\ell-2}-1}) &= (\delta_{\ell} + \delta_{\ell-1}) \cdot \cos(j\theta_{\ell}) - \delta_{\ell} \cot((j+1)\theta_{\ell}) \cdot \sin(j\theta_{\ell}) \\ &\geq \delta_{\ell} \cdot \cos(j\theta_{\ell}) - \delta_{\ell} \cot((j+1)\theta_{\ell}) \cdot \sin(j\theta_{\ell}) \\ &= \frac{\delta_{\ell}}{\sin((j+1)\theta_{\ell})} (\sin((j+1)\theta_{\ell}) \cos(j\theta_{\ell}) - \sin(j\theta_{\ell}) \cos((j+1)\theta_{\ell})) \\ &= \delta_{\ell} \frac{\sin(\theta_{\ell})}{\sin((j+1)\theta_{\ell})}.\end{aligned}$$

This means that no matter which point sets the gap (or if one of the points does not exist), the gap is at least  $\delta_{\ell} \frac{\sin(\theta_{\ell})}{\sin((j+1)\theta_{\ell})}$ .  $\square$

Finally, we need to sum over each even layer.

**Corollary 32.** For any even  $\ell > 2$ ,  $\sum_{j=0}^{n_{\ell}-1} \text{gap}_{\ell,j}^{X,Q} \geq \delta_{\ell} \cdot \ln(n_{\ell})/2$ .

*Proof of Corollary 32.* Consider the following sequence of calculations:

$$\begin{aligned}\sum_{j=0}^{n_{\ell}-1} \text{gap}_{\ell,j}^{X,Q} &\geq \sum_{j=0}^{n_{\ell}-1} \delta_{\ell} \frac{\sin(\theta_{\ell})}{\sin((j+1)\theta_{\ell})} \\ &\geq \delta_{\ell} \cdot (\theta_{\ell} - \theta_{\ell}^3/6) \cdot \sum_{j=0}^{n_{\ell}-1} \frac{1}{(j+1)\theta_{\ell}} \\ &\geq \delta_{\ell} \cdot (1 - \theta_{\ell}^2/6) \cdot \ln(n_{\ell}) \\ &\geq \delta_{\ell} \cdot \ln(n_{\ell})/2\end{aligned}$$

Above, the first line follows from Lemma 31. The second line uses the fact that  $\theta_{\ell} - \theta_{\ell}^3/6 \leq \sin(\theta_{\ell}) \leq \theta_{\ell}$ , because  $\theta_{\ell} \in [0, \pi/2]$ . The third line follows as the  $n^{\text{th}}$  harmonic sum is at least  $\ln(n)$ . The final line follows as  $\theta_{\ell}^2/6 = \pi^2/(24(n_{\ell}-1)^2) \leq 1/2$ .  $\square$

And finally, we can wrap up the proof of the proposition. Here, we just need to recall that  $\delta_{\ell} := \frac{1}{\alpha n_{\ell}} = \frac{1}{\alpha \ell \ln^2(\ell)}$ . Therefore, we conclude that:

$$\sum_{\ell \text{ even}} \sum_{j=0}^{n_{\ell}-1} \text{gap}_{\ell,j}^{X,Q} \geq \sum_{\ell \text{ even}} \delta_{\ell} \cdot \ln(n_{\ell})/2 = \sum_{\ell \text{ even}} \frac{1}{2\alpha \ell \ln(\ell)} = \infty. \quad \square$$

## C Proof of Corollary 11

We prove Corollary 11 by making use of Theorem 2 combined with the sequence  $X$  from Section 5. The only task is to confirm that  $\text{ARev}(\mathcal{D}) < \infty$  for the resulting  $\mathcal{D}$ , which essentially requires that we execute and analyze the construction fully. Let us quickly review the [HN19] construction, given as input a sequence  $X$ :

- Let  $B$  be a very large constant, to be defined later.
- Let  $\vec{v}_i := B^{2^i} \cdot \vec{x}_i / \|\vec{x}_i\|_1$  (for all  $i$ ).
- Let  $\mathcal{D}$  sample  $\vec{v}_i$  with probability  $1/B^{2^i}$  (for all  $i$ ).
- Let  $\mathcal{D}$  sample  $\vec{0}$  with probability  $1 - \sum_{i \geq 1} 1/B^{2^i}$ .

[HN19] establishes that the above construction yields Theorem 2 (for sufficiently large  $B$ , as a function of  $\varepsilon$ ). To complete the proof of Corollary 11, we just need to relate  $\text{ARev}(\mathcal{D})$  for this construction to  $\text{AlignGap}(X)$ .

**Proposition 33.** *The construction above yields a  $\mathcal{D}$  satisfying  $\text{ARev}(\mathcal{D}) \leq \text{AlignGap}(X) + 1/B$ .*

*Proof.* Consider any mechanism  $M$ . We show that  $\text{AlignGap}(X) \geq \text{ARev}(\mathcal{D}, M) - 1/B$ . To see this, consider the following choice of  $C$ :

- If  $\vec{v}_i$  is parallel to  $\vec{q}^M(\vec{v}_i)$ , set  $c_i := \|\vec{q}^M(\vec{v}_i)\|_2 / \|\vec{x}_i\|_2$ .
- If  $\vec{v}_i$  is not parallel to  $\vec{q}^M(\vec{v}_i)$ , set  $c_i := 0$ .

We now need to lower bound  $\text{sgap}_i^{X,C}$ , when  $i$  satisfies the first bullet. Observe that, because  $M$  is truthful, we must have, for all  $j < i$ :

$$\begin{aligned} \vec{v}_i \cdot \vec{q}^M(\vec{v}_i) - p^M(\vec{v}_i) &\geq \vec{v}_i \cdot \vec{q}^M(\vec{v}_j) - p^M(\vec{v}_j) \\ &\Rightarrow p^M(\vec{v}_i) \leq p^M(\vec{v}_j) + B^{2^i} \vec{x}_i \cdot (c_i \vec{x}_i - c_j \vec{x}_j) / \|\vec{x}_i\|_1 \\ &\Rightarrow p^M(\vec{v}_i) \leq 2B^{2^{i-1}} + B^{2^i} \text{sgap}_i^{X,C} / \|\vec{x}_i\|_1 \end{aligned}$$

Above, the first line follows from incentive compatibility. The second line follows as  $\vec{q}^M(\vec{v}_i) = c_i \vec{x}_i$  for all  $i$  in the first bullet, and either  $\vec{q}^M(\vec{v}_j) = c_j \vec{x}_j$ , or  $c_j = 0$ . The final line follows by taking  $j := \arg \min_{j < i} \{\vec{v}_i \cdot (c_i \vec{v}_i - c_j \vec{v}_j)\}$ , and by observing that  $\vec{v}_j$  cannot possibly pay more than their value for the grand bundle.

We can then conclude that:

$$\begin{aligned} \text{ARev}(\mathcal{D}, M) &\leq \sum_i (2B^{2^{i-1}} + B^{2^i} \text{sgap}_i^{X,C} / \|\vec{x}_i\|_1) / B^{2^i} \\ &\leq \sum_i 2/B^{2^{i-1}} + \text{AlignGap}(X) \\ &\leq \text{AlignGap}(X) + 1/B. \end{aligned}$$

□

Because we can take  $B$  as large as we like, we can construct a  $\mathcal{D}$  such that  $\text{ARev}(\mathcal{D})$  is arbitrarily close to  $\text{AlignGap}(X)$ , while also maintaining that  $\text{Rev}(\mathcal{D})$  is arbitrarily close to  $\text{MenuGap}(X)$ . Because Theorem 9 provides a construction  $X$  such that  $\text{MenuGap}(X) / \text{AlignGap}(X) = \infty$ , the [HN19] construction, with sufficiently large  $B$ , yields a  $\mathcal{D}$  with  $\text{Rev}(\mathcal{D}) / \text{ARev}(\mathcal{D}) = \infty$ , completing the proof of Corollary 11.