

521 **Appendices for the Submission: Learned Index with Dynamic  $\epsilon$**

522 Our appendices include the following content:

- 523 • Sec.A: the full **proof** of Theorem 1.
- 524 • Sec.B: the analysis about the **learned slopes** of other  $\epsilon$ -bounded methods.
- 525 • Sec.C: the details of the binary search and exponential search, and the connections between
- 526 **prediction error** and these specific **searching strategies**.
- 527 • Sec.D: the summarized **algorithm** of the proposed method.
- 528 • Sec.E: the discussion about how the proposed framework **inherits the good abilities** of
- 529 existing learned index methods.
- 530 • Sec.F: the **implementation details** of experiments.
- 531 • Sec.G: the detailed introduction and visualization of the adopted **datasets**.
- 532 • Sec.H: more **experimental results** including the overall index performance, ablation study
- 533 and theoretical validation on other datasets and methods. Besides, we explore an indica-
- 534 tive quantity (the *CV* value) to provide further insight into the rational of the proposed
- 535 framework.

536 **A Proof of Theorem 1**

537 Given a learned segment  $S_i : y = a_i x + b_i$ , denote  $c_i$  as the stored position of the last covered data  
 538 for the  $(i - 1)$ -th segment ( $c_1 = 0$  for the first segment). We can write the expectation of  $SegErr_i$   
 539 for the segment  $S_i$  as the following form:

$$\mathbb{E}[SegErr_i] = \mathbb{E} \left[ \sum_{j=0}^{\lceil j^* - 1 \rceil} |a_i X_j + b_i - (j + c_i + 1)| \right],$$

540 where  $j^*$  indicates the length of the segment, and  $X_j$  indicates the  $j$ -th key covered by the segment  $S_i$ .  
 541 As studied in [10], the linear-approximation problem with  $\epsilon$  guarantee can be modeled as random walk  
 542 processes. Specifically,  $X_j = X_0 + \sum_{k=0}^j G_k$  (for  $j \in \mathbb{Z}_{>0}$ ) where  $G_k$  is the key increment variable  
 543 whose mean and variance is  $\mu$  and  $\sigma^2$  respectively. Denote the  $Z_j = X_j - j/a_i + (b_i - c_i - 1)/a_i$  as the  
 544  $j$ -th position of the transformed random walk  $\{Z_j\}_{j \in \mathbb{N}}$ , and  $j^* = \max\{j \in \mathbb{N} | -\epsilon/a_i \leq Z_j \leq \epsilon/a_i\}$   
 545 as the random variable indicating the maximal position when the random walk is within the strip of  
 546 boundary  $\pm\epsilon/a_i$ . The expectation can be rewritten as

$$\begin{aligned} \mathbb{E} \left[ \sum_{j=0}^{\lceil j^* - 1 \rceil} |a_i X_j - j + (b_i - c_i - 1)| \right] &= a_i \mathbb{E} \left[ \sum_{j=0}^{\lceil j^* - 1 \rceil} |Z_j| \right] \\ &= a_i \sum_{n=1}^{\infty} \mathbb{E} \left[ \sum_{j=0}^{n-1} |Z_j| \right] \Pr(j^* = n). \end{aligned} \quad (3)$$

547 The last equality in Eq. (3) is due to the definition of expectation. Following the MET algorithm that  
 548 the  $S_i$  goes through the point  $(X_0, Y_0 = c_i + 1)$ , we get  $b_i = -a_i X_0 + c_i + 1$  and we can rewrite  
 549  $Z_j$  as the following form:

$$\begin{aligned} Z_0 &= 0, \quad Z_j \stackrel{j \geq 0}{=} X_j - X_0 - j/a_i = \sum_{k=1}^j G_k - j/a_i \\ &= \sum_{k=1}^j (G_k - 1/a_i) = \sum_{k=1}^j (W_k), \end{aligned}$$

550 where  $W_k$  is the walk increment variable of  $Z_j$ ,  $\mathbb{E}[W_k] = \mu - 1/a_i$  and  $Var[W_k] = \sigma^2$ . Under  
 551 the MET algorithm setting where  $a_i = 1/\mu$  and  $\epsilon \gg \sigma/\mu$ , the transformed random walk  $\{Z_j\}$  has  
 552 increments with zero mean and variance  $\sigma^2$ , and many steps are necessary to reach the random walk

553 boundary. With the Central Limit Theorem, we can assume that  $Z_j$  follows the normal distribution  
 554 with mean  $\mu_{zj}$  and variance  $\sigma_{zj}^2$ , and thus  $|Z_j|$  follows the folded normal distribution:

$$Z_j \sim \mathcal{N}((\mu - 1/a_i)j, j\sigma^2),$$

$$\mathbb{E}(|Z_j|) = \mu_{zj}[1 - 2\Phi(-\mu_{zj}/\sigma_{zj})] + \sigma_{zj}\sqrt{2/\pi}\exp(-\mu_{zj}^2/2\sigma_{zj}^2),$$

555 where  $\Phi$  is the normal cumulative distribution function. For the MET algorithm,  $a_i = 1/\mu$  and thus  
 556 the  $\mu_{zj} = 0$ ,  $\sigma_{zj} = \sigma\sqrt{j}$ , and  $\mathbb{E}(|Z_j|) = \sqrt{2/\pi}\sigma\sqrt{j}$ . Then the Eq. (3) can be written as

$$\begin{aligned} \frac{1}{\mu} \sum_{n=1}^{\infty} \mathbb{E} \left[ \sum_{j=0}^{n-1} |Z_j| \right] \Pr(j^* = n) &< \frac{1}{\mu} \sum_{n=1}^{\infty} \sum_{j=0}^{n-1} \mathbb{E}[|Z_j|] \Pr(j^* = n) \\ &= \frac{\sigma}{\mu} \sqrt{\frac{2}{\pi}} \sum_{n=1}^{\infty} \sum_{j=0}^{n-1} \sqrt{j} \Pr(j^* = n). \end{aligned} \quad (4)$$

557 For the inner sum term in Eq. (4), we have  $(\sum_{j=0}^{n-1} \sqrt{j}) < \frac{2}{3}n\sqrt{n}$  since

$$\sum_{j=0}^{n-1} \sqrt{j} < \sum_{j=0}^{n-1} \sqrt{j} + \frac{\sqrt{n}}{2} < \int_0^n \sqrt{x} dx = \frac{2}{3}n\sqrt{n},$$

558 then the result in Eq. (4) becomes

$$\begin{aligned} \mathbb{E}[SegErr_i] &< \frac{2}{3} \sqrt{\frac{2}{\pi}} \frac{\sigma}{\mu} \sum_{n=1}^{\infty} n\sqrt{n} \Pr(j^* = n) \\ &= \frac{2}{3} \sqrt{\frac{2}{\pi}} \frac{\sigma}{\mu} \mathbb{E}[(j^*)^{\frac{3}{2}}] = \frac{2}{3} \sqrt{\frac{2}{\pi}} \frac{\sigma}{\mu} \mathbb{E} \left[ ((j^*)^2)^{\frac{3}{4}} \right] \\ &\leq \frac{2}{3} \sqrt{\frac{2}{\pi}} \frac{\sigma}{\mu} \left( \mathbb{E}[(j^*)^2] \right)^{\frac{3}{4}}, \end{aligned}$$

559 where the last inequality holds due to the Jensen inequality  $\mathbb{E}[X^{\frac{3}{4}}] \leq (\mathbb{E}[X])^{\frac{3}{4}}$ . Using  $\mathbb{E}[j^*] = \frac{\mu^2}{\sigma^2}\epsilon^2$   
 560 and  $Var[j^*] = \frac{2}{3} \frac{\mu^4}{\sigma^4}\epsilon^4$  derived in MET algorithm [10], we get  $\mathbb{E}[(j^*)^2] = \frac{5}{3} \frac{\mu^4}{\sigma^4}\epsilon^4$ , which yields the  
 561 following upper bound:

$$\mathbb{E}[SegErr_i] < \frac{2}{3} \sqrt{\frac{2}{\pi}} \left(\frac{5}{3}\right)^{\frac{3}{4}} \left(\frac{\mu}{\sigma}\right)^2 \epsilon^3.$$

562 For the lower bound, applying the triangle inequality into the Eq. (3), we have

$$\begin{aligned} &\frac{1}{\mu} \sum_{n=1}^{\infty} \mathbb{E} \left[ \sum_{j=0}^{n-1} |Z_j| \right] \Pr(j^* = n) \\ &> \frac{1}{\mu} \sum_{n=1}^{\infty} \mathbb{E} \left[ \left| \sum_{j=0}^{n-1} Z_j \right| \right] \Pr(j^* = n) \\ &= \frac{1}{\mu} \sum_{n=1}^{\infty} \mathbb{E}[|Z|] \Pr(j^* = n), \end{aligned} \quad (5)$$

563 where  $Z = \sum_{j=0}^{n-1} Z_j$ . Since  $Z_j \sim \mathcal{N}(0, \sigma_{zj}^2)$ , the  $Z$  follows the normal distribution:

$$Z \sim \mathbb{N}\left(\mu_Z = 0, \sigma_Z^2 = \sum_{j=0}^{n-1} \sigma_{zj}^2 + \sum_{j=0}^{n-1} \sum_{k=0, k \neq j}^{n-1} r_{jk} \sigma_{zj} \sigma_{zk}\right),$$

564 where  $r_{jk}$  is the correlation between  $Z_j$  and  $Z_k$ . Since  $\mu_Z = 0$ , the  $|Z|$  follows the folded normal  
 565 distribution with  $\mathbb{E}[|Z|] = \sigma_Z \sqrt{2/\pi}$ . Since the random walk  $\{Z_j\}$  is a process with i.i.d. increments,  
 566 the correlation  $r_{jk} \geq 0$ . With  $\sigma_{zj} = \sigma\sqrt{j} > 0$  and  $r_{jk} \geq 0$ , we have

$$\mathbb{E}[|Z|] > \sqrt{\frac{2}{\pi}} \sum_{j=0}^{n-1} \sigma_{zj} > \sigma \sqrt{n(n-1)/\pi} > \frac{\sigma(n-1)}{\sqrt{\pi}},$$

567 and the result in Eq. (5) becomes:

$$\begin{aligned}
\mathbb{E}[SegErr_i] &> \frac{1}{\mu} \sum_{n=1}^{\infty} \mathbb{E} \left[ \left| \sum_{j=0}^{n-1} Z_j \right| \right] \Pr(j^* = n) \\
&> \frac{\sigma}{\mu} \sqrt{\frac{1}{\pi}} \sum_{n=1}^{\infty} (n-1) \Pr(j^* = n) \\
&= \frac{\sigma}{\mu} \sqrt{\frac{1}{\pi}} \mathbb{E}[j^* - 1] = \sqrt{\frac{1}{\pi}} \left( \frac{\mu}{\sigma} \epsilon^2 - \frac{\sigma}{\mu} \right).
\end{aligned}$$

568 Since  $\epsilon \gg \frac{\sigma}{\mu}$ , we can omit the right term  $\sqrt{\frac{1}{\pi}} \frac{\sigma}{\mu}$  and finish the proof.

## 569 B Learned Slopes of Other $\epsilon$ -Bounded Methods

570 As shown in Theorem 1, we have known how  $\epsilon$  impacts the  $SegErr_i$  of each segment learned by  
571 the MET algorithm, where the theoretical derivations largely rely on the slope condition  $a_i = 1/\mu$ .  
572 Here we prove that for other  $\epsilon$ -bounded methods, the learned slope of each segment (*i.e.*,  $a_i$  of  $S_i$ )  
573 concentrates on the reciprocal of the expected key interval as shown in the following Theorem.

574 **Theorem 2.** *Given an  $\epsilon \in \mathbb{Z}_{>1}$  and an  $\epsilon$ -bounded learned index algorithm  $\mathcal{A}$ . For a linear segment*  
575  *$S_i : y = a_i x + b_i$  learned by  $\mathcal{A}$ , denote its covered data and the number of covered keys as  $\mathcal{D}_i$*   
576 *and  $Len(\mathcal{D}_i)$  respectively. Assuming the expected key interval of  $\mathcal{D}_i$  is  $\mu_i$ , the learned slope  $a_i$*   
577 *concentrates on  $\tilde{a} = 1/\mu_i$  with bounded relative difference:*

$$\left(1 - \frac{2\epsilon}{\mathbb{E}[Len(\mathcal{D}_i)] - 1}\right)\tilde{a} \leq E[a_i] \leq \left(1 + \frac{2\epsilon}{\mathbb{E}[Len(\mathcal{D}_i)] - 1}\right)\tilde{a}.$$

578 *Proof.* For the learned linear segment  $S_i$ , denote its first predicted position and last predicted position  
579 as  $y'_0$  and  $y'_n$  respectively, we have its slope  $a_i = \frac{y'_n - y'_0}{x_n - x_0}$ . Notice that  $y_0 - \epsilon \leq y'_0 \leq y_0 + \epsilon$  and  
580  $y_n - \epsilon \leq y'_n \leq y_n + \epsilon$  due to the  $\epsilon$  guarantee, we have  $y_n - y_0 - 2\epsilon \leq y'_n - y'_0 \leq y_n - y_0 + 2\epsilon$  and  
581 the expectation of  $a_i$  can be written as

$$\mathbb{E}\left[\frac{y_n - y_0 + 2\epsilon}{x_n - x_0}\right] \leq E[a_i] = \frac{y'_n - y'_0}{x_n - x_0} \leq \mathbb{E}\left[\frac{y_n - y_0 - 2\epsilon}{x_n - x_0}\right].$$

582 Note that for any learned segment  $S_i$  whose first covered data is  $(x_0, y_0)$  and last covered data is  
583  $(x_n, y_n)$ , we have  $\mathbb{E}\left[\frac{x_n - x_0}{y_n - y_0}\right] = \mu_i$  and thus the inequalities become

$$\frac{1}{\mu} - \mathbb{E}\left[\frac{2\epsilon}{x_n - x_0}\right] \leq E[a_i] \leq \frac{1}{\mu} + \mathbb{E}\left[\frac{2\epsilon}{x_n - x_0}\right].$$

584 Since  $\tilde{a} = 1/\mu_i$  and  $\mathbb{E}[x_n - x_0] = (\mathbb{E}[Len(\mathcal{D}_i)] - 1)\mu_i$ , we finish the proof.  $\square$

585 The Theorem 2 shows that the relative deviations between learned slope  $a_i$  and  $\tilde{a}$  are within  
586  $2\epsilon/(\mathbb{E}[Len(\mathcal{D}_i)] - 1)$ . For the MET and PGM learned index methods, we have the following  
587 corollary that depicts preciser deviations without the expectation term  $\mathbb{E}[Len(\mathcal{D}_i)]$ .

588 **Corollary 2.1.** *For the MET method [10] and the optimal  $\epsilon$ -bounded linear approximation method*  
589 *that learns the largest segment length used in PGM [12], the slope relative differences are at  $O(1/\epsilon)$ .*

590 *Proof.* We note that the segment length of a learned segment is at  $O(\epsilon^2)$  for the MET algorithm,  
591 which is proved in the Theorem 1 of [10]. Since PGM achieves the largest learned segment length  
592 that is larger than the one of the MET algorithm, we finish the proof.  $\square$

## 593 C Connecting Prediction Error with Searching Strategy

594 As we mentioned in Section 3.1, we can find the true position of the queried data point in  $O(\log(N) +$   
595  $\log(|\hat{y} - y|))$  where  $N$  is the number of learned segments and  $|\hat{y} - y|$  is the absolute prediction error.  
596 A binary search or exponential search finds the stored true position  $y$  based on  $\hat{y}$ . It is worth pointing

597 out that the searching cost in terms of searching range  $|\hat{y} - y|$  of binary search strategy corresponds  
 598 to the maximum absolute prediction error  $\epsilon$ , whereas the one of exponential search corresponds to  
 599 the mean absolute prediction error (*MAE*). In this paper, we decouple the quantity  $SegErr_i$  as the  
 600 product of  $Len(\mathcal{D}_i)$  and  $MAE(\mathcal{D}_i|S_i)$  in the derivation of Theorem 1. Built upon the theoretical  
 601 analysis, we adopt exponential search in experiments to better leverage the predictive models.

602 To clarify, let’s consider a learned segment  $S_i$  with its covered data  $\mathcal{D}_i$ . Let  $|\hat{y}_k - y_k|$  be the absolute  
 603 prediction error of  $k$ -th data point covered by this segment, and  $\epsilon_i$  be the maximum absolute prediction  
 604 error of  $S_i$ , i.e.,  $|\hat{y}_k - y_k| \leq \epsilon_i$  for all  $k \in [len(\mathcal{D}_i)]$ .

- 605 • The binary search is conducted within the searching range  $[\hat{y}_k \pm \epsilon_i]$  for each data point <sup>2</sup>,  
 606 thus the mean search range is  $\frac{1}{len(\mathcal{D}_i)} \sum_{k=1}^{len(\mathcal{D}_i)} 2\epsilon_i = O(\epsilon_i)$ , which is independent of the  
 607 preciseness of the learned segment and an upper bound of  $MAE(\mathcal{D}_i|S_i)$ .
- 608 • The exponential search first finds the searching range where the queried data may exist by  
 609 centering around the  $\hat{y}$ , repeatedly doubling the range  $[\hat{y} \pm 2^q]$  where the integer  $q$  grows  
 610 from 0, and comparing the queried data with the data points at positions  $\hat{y} \pm 2^q$ . After finding  
 611 the specific range such that a  $q_k$  satisfies  $2^{\log(q_k)-1} \leq |\hat{y}_k - y_k| \leq 2^{\lceil \log(q_k) \rceil}$  for the  $k$ -th  
 612 data, a binary search is conducted to find the exact location. In this way, the mean search  
 613 range is  $\frac{1}{len(\mathcal{D}_i)} \sum_{k=1}^{len(\mathcal{D}_i)} (2^{\lceil \log(q_k) \rceil + 1}) = O(MAE(\mathcal{D}_i|S_i))$ , which can be much smaller  
 614 than  $O(\epsilon_i)$  especially for strong predictive models and the datasets having clear linearity.

## 615 D The Algorithm of Dynamic $\epsilon$ Adjustment

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**Algorithm** Dynamic  $\epsilon$  Adjustment with Pluggable  $\epsilon$  Learner

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**Input:**  $\mathcal{D}$ : Data to be indexed,  $\mathcal{A}$ : Learned index algorithm,  $\bar{\epsilon}$ : Expected  $\epsilon$ ,  $\rho$ : Length percentage  
 for look-ahead data

**Output:**  $\mathbf{S}$ : Learned segments with varied  $\epsilon$ s

- 1: initial parameters  $w_{1,2,3}$  of the learned function:  $f(\epsilon, \mu, \sigma) = w_1 \left(\frac{\mu}{\sigma}\right)^{w_2} \bar{\epsilon}^{w_3}$
  - 2: initial mean length of learned segments so far:  $Len(\mathcal{D}_S) \leftarrow 404$
  - 3:  $\mathbf{S} \leftarrow \emptyset$ ,  $(\hat{\mu}/\hat{\sigma}) \leftarrow 0$
  - 4: **repeat**
  - 5:   Get data statistic:
  - 6:    $(\mu/\sigma) \leftarrow lookahead(\mathcal{D}, Len(\mathcal{D}_S) \cdot \rho)$
  - 7:   Adjust  $\epsilon$  based on the learner:
  - 8:    $\epsilon^* \leftarrow \left(\widetilde{SegErr}/w_1 \left(\frac{\mu}{\sigma}\right)^{w_2}\right)^{1/w_3}$
  - 9:   Learn new segment  $S_i$  using adjusted  $\epsilon^*$ :
  - 10:    $[S_i, \mathcal{D}_i] \leftarrow \mathcal{A}(\mathcal{D}, \epsilon^*)$
  - 11:    $\mathbf{S} \leftarrow \mathbf{S} \cup S_i$
  - 12:    $\mathcal{D} \leftarrow \mathcal{D} \setminus \mathcal{D}_i$ ,  $\mathcal{D}_S \leftarrow \mathcal{D}_S \cup \mathcal{D}_i$
  - 13:   Online update  $Len(\mathcal{D}_S)$ :
  - 14:    $Len(\mathcal{D}_S) \leftarrow \text{running-mean}(Len(\mathcal{D}_S), Len(\mathcal{D}_i))$
  - 15:    $(\hat{\mu}/\hat{\sigma}) \leftarrow \text{running-mean}((\hat{\mu}/\hat{\sigma}), (\mu/\sigma))$
  - 16:   Train the learner with ground-truth:
  - 17:    $w_{1,2,3} \leftarrow optimize(f, S_i, SegErr_i)$
  - 18:    $\widetilde{SegErr} \leftarrow w_1 (\hat{\mu}/\hat{\sigma})^{w_2} \bar{\epsilon}^{w_3}$
  - 19: **until**  $\mathcal{D} = \emptyset$
- 

616 In Section 3.4, we provide detailed description about the initialization and adjustment sub-procedures.  
 617 The  $lookahead()$  and  $optimize()$  are in the “**Look-ahead Data**” and “**SegErr and Optimization**”  
 618 paragraph respectively.

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<sup>2</sup>The lower bound and upper bounds of searching ranges should be constricted to 0 and  $len(\mathcal{D}_i)$  respectively. For brevity, we omit the corner cases when comparing these two searching strategies as they both need to handle the out-of-bounds scenario.

## 619 E Inheriting the Abilities of Existing Works

620 In this Section, we discuss the benefits of our proposed framework brought by its pluggable property  
621 with two example scenarios, the dynamic data update and hard limitation on user-required index size.

622 We note that the data insert operation has been discussed in the adopted baseline methods, FITing-Tree  
623 [14] and PGM [12]. More importantly, neither of these two methods altered the notion of  $\epsilon$  when  
624 dealing with the data insertion, and they still relied on their  $\epsilon$ -bounded piece-wise segmentation  
625 algorithms. The proposed framework is still valid when using their respective solutions to handle  
626 the data insertion. Specifically, FITing-Tree proposes to introduce a buffer for each learned segment,  
627 which is used to store the inserted keys, and when the buffer is full, the data covered by the segment  
628 will be re-segmented (see Section 5 in [13]). PGM adopts a logarithmic method [28, 27] that maintains  
629 a series of sorted sets  $\{S_0, S_1, \dots, S_b\}$  where  $b = \theta(\log(|\mathcal{D}|))$ , and builds multiple PGM-INDEX  
630 models over the sets. When a key  $x$  is inserted, a new PGM-INDEX will be built over the merged  
631 sets (see Section 3 in [12]). In general, these solutions proposed by existing methods for inserting  
632 keys are based on *re-indexing for a piece of data along with the inserted data*, and the *re-indexing*  
633 *processes are the same as the original piece-wise linear segmentation processes* but for different  
634 data, therefore, we can still apply the proposed dynamic- $\epsilon$  framework for these methods in insertion  
635 scenarios just like we adjust  $\epsilon$  and learn index according to the new data to be re-indexed.

636 For the hard size limitation case, we observe that the existing work PGM introduced a multi-criteria  
637 variant that auto-tunes itself with pre-defined hard size requirement from users. Our proposed  
638 framework is pluggable and still valid when using the PGM variant to handle the size requirement.  
639 Specifically, given a space constraint, the multi-criteria PGM propose to iteratively estimate the  
640 relationship between  $\epsilon$  and *size* with a learnable function  $size(\epsilon) = a\epsilon^{-b}$ , and automatically outputs  
641 the index that minimizes its query time via different estimated  $\epsilon$ s. Given a size requirement, we can  
642 just do the same thing in dynamic  $\epsilon$  scene by setting our  $\tilde{\epsilon}$  as  $\epsilon$  estimated by the original PGM method.

## 643 F Implementation Details

644 All the experiments are conducted on a Linux server with an Intel Xeon Platinum 8163 2.50GHz  
645 CPU. We first introduce more details and the implementation of baseline learned index methods.  
646 *MET* [10] fixes the segment slope as the reciprocal of the expected key interval, and goes through  
647 the first available data point for each segment. *FITing-Tree* [14] adopts a greedy shrinking cone  
648 algorithm and the learned segments are organized with a  $B^+$ -tree. Here we use the `stx::btree` (v0.9)  
649 implementation [2] and set the filling factors of inner nodes and leaf nodes as 100%, *i.e.*, we adopt  
650 the full-paged filling manner. *Radix-Spline* [18] adopts a greedy spline interpolating algorithm to  
651 learn spline points, and the learned spline segments are organized with a flat radix table. We set  
652 the number of radix bits as  $r = 16$  for the Radix-Spline method, which means that the leveraged  
653 radix table contains  $2^{16}$  entries. *PGM* [12] adopts a convex hull based algorithm to achieve the  
654 minimum number of learned segments, and the segments can be organized with the help of binary  
655 search, *CSS-Tree* [29] and recursive structure. Here we implement the recursive version since it beats  
656 the other two variants in terms of indexing performance. For all the baselines and our method, we  
657 adopt exponential search to better leverage the predictive models since the query complexity using  
658 exponential search corresponds the preciseness of models (*MAE*) as we analyzed in Appendix C.

659 We then describe a few additional details of the proposed framework in terms of the  $\epsilon$ -learner  
660 initialization and the hyper-parameter setting. For the  $w_{1,2,3}$  of the  $\epsilon$ -learner shown in the Eq. (2), at  
661 the beginning, we learn the first five segments with the  $\epsilon$  sequence  $[\frac{1}{4}\tilde{\epsilon}, \frac{1}{2}\tilde{\epsilon}, \tilde{\epsilon}, 2\tilde{\epsilon}, 4\tilde{\epsilon}]$ , then track their  
662 rewarded  $SegErr_i$  and update the parameters  $w_{1,2,3}$  using least square regression. We empirically  
663 found that this light-weight initialization leads to better index performance compared to the versions  
664 with random parameter initialization, and it benefits the exploration of diverse  $\epsilon^*$ , *i.e.*, leading to  
665 the larger variance of the dynamic  $\epsilon$  sequence  $[\epsilon_1, \dots, \epsilon_i, \dots, \epsilon_N]$ . As for the hyper-parameter  $\rho$   
666 (described in the Section 3.4), we conduct grid search over  $\rho \in [0.1, 0.4, 0.7, 1.0]$  on Map an IoT  
667 datasets. We found that all the  $\rho$ s achieve better *N-MAE* trade-off (*i.e.*, smaller AUNEC results) than  
668 the fixed  $\epsilon$  versions. Since the setting  $\rho = 0.4$  achieves averagely best results on the two datasets, we  
669 set  $\rho$  to be 0.4 for the other datasets.

670 **G Dataset Details**

671 Our framework is verified on several widely adopted datasets having different data scales and  
 672 distributions. *Weblogs* [19, 14, 12] contains about 715M log entries for the requests to a university web  
 673 server and the keys are log timestamps. *IoT* [14, 12] contains about 26M event entries from different  
 674 IoT sensors in a building and the keys are recording timestamps. *Map* dataset [19, 14, 9, 12, 21]  
 675 contains location coordinates of 200M places that are collected around the world from the Open  
 676 Street Map [25], and the keys are the longitudes of these places. *Lognormal* [12] is a synthetic dataset  
 677 whose key intervals follow the lognormal distribution:  $\ln(G_i) \sim \mathcal{N}(\mu_{lg}, \sigma_{lg}^2)$ . To simulate the varied  
 678 data characteristics among different localities. We generate 20M keys with 40 partitions by setting  
 679  $\mu_{lg} = 1$  and setting  $\sigma_{lg}$  with a random number within  $[0.1, 1]$  for each partition.

680 We normalize the positions of stored data into the range  $[0, 1]$ , and thus the key-position distribution  
 681 can be modeled as Cumulative Distribution Function (CDF). We plot the CDFs and zoomed-in CDFs  
 682 of experimental datasets in Figure 7 and Figure 8 respectively, which intuitively illustrate the diversity  
 683 of the adopted datasets.

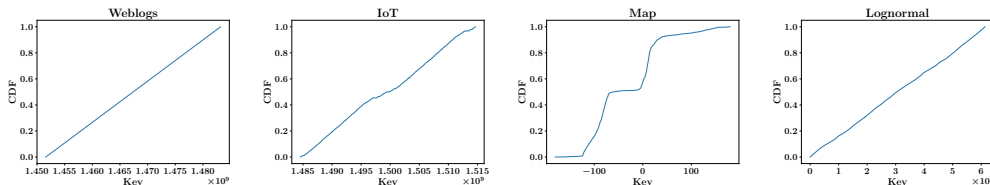


Figure 7: CDFs of adopted datasets.

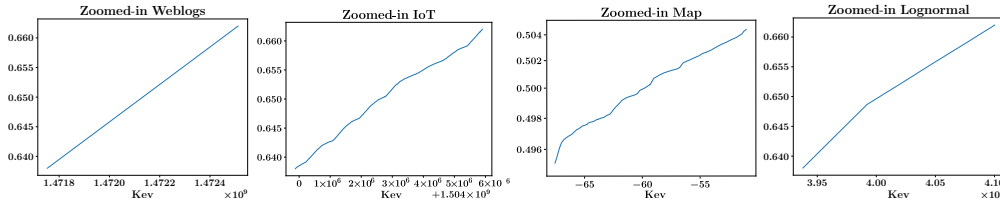


Figure 8: Zoomed-in CDFs of adopted datasets.

684 **H Additional Experimental Results**

685 **Overall Index Performance.** For the  $N$ -MAE trade-off improvements and the actual querying  
 686 efficiency improvements brought by the proposed framework, we illustrate more  $N$ -MAE trade-off  
 687 curves in Figure 9 and querying time results in Figure 10. We also mark the 99th percentile (P99)  
 688 latency as the right bar, which is a useful metric in industrial-scale practical systems. Recall that the  
 689  $N$ -MAE trade-off curve adequately reflects the *index size* and *querying time*: (1) the *segment size*  
 690 *in bytes* and  $N$  are only different by a constant factor, e.g., the size of a segment can be 128bit if it  
 691 consists of two double-precision float parameters (slope and intercept); (2) the querying operation can  
 692 be done in  $O(\log(N) + \log(|y - \hat{y}|))$  as we mentioned in Section 3.1, thus a better  $N$ -MAE trade-off  
 693 indicates a better querying efficiency. From these figures, we can see that the dynamic  $\epsilon$  versions of  
 694 all the baseline methods achieve better  $N$ -MAE trade-off and better querying efficiency, verifying the  
 695 effectiveness and the wide applicability of the proposed framework. Regards the p99 metrics, we can  
 696 see that the dynamic version achieves comparable or even better P99 results than the static version,  
 697 showing that the proposed method not only improves the average lookup time, but also has a good  
 698 robustness. This is because of that our method can effectively adjust  $\epsilon$  based on the expected  $\bar{\epsilon}$  and  
 699 data characteristic, making the  $\{\epsilon_i\}$  fluctuated within a moderate range.

700 **CV as an Indicative Quantity.** The *coefficient of variation* (CV) value, *i.e.*,  $CV = \sigma / \mu$ , plays an  
 701 important factor in our bounds to reflect the linearity degree of the data. We have seen that  $CV$  is  
 702 effective to help dynamically adjust  $\epsilon$  in our framework as shown in our experiments. Here we explore  
 703 that *whether the CV value can be an indicative quantity to shed light on what types of data will*

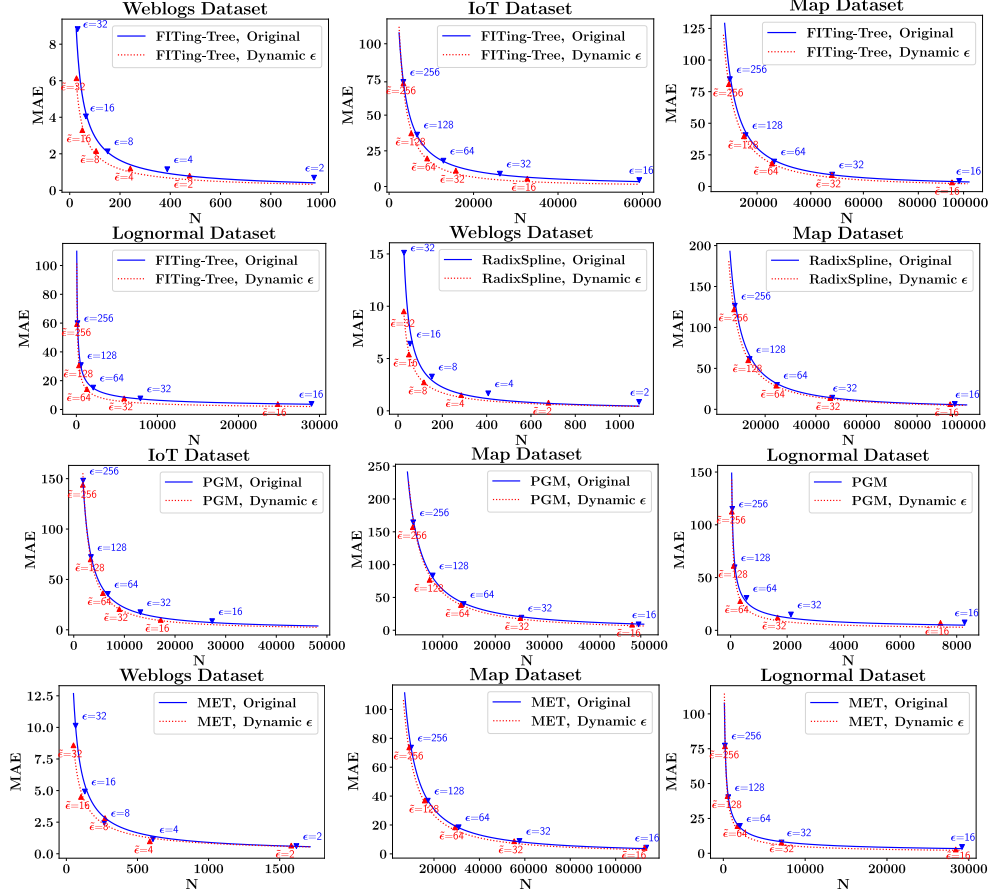


Figure 9: The additional  $N$ -MAE trade-off curves for learned index methods.

704 *benefit from our dynamic adjustment?* To be specific, we calculate the  $CV$  values of the experimental  
705 datasets and compare them with the trade-off improvements. The global  $CV$  values of IoT, Map,  
706 Lognormal and Weblogs are 65.24, 11.12, 0.85, and 0.013 respectively, while their AUNEC improved  
707 by 20.71%, 6.47%, 21.89% and 26.96% respectively. With the exception of IoT, the rest of the results  
708 show that *the smaller the  $CV$  value is, the greater the trade-off improvement of dynamic  $\epsilon$  brings.*  
709 We find that IoT is a locally linear but globally fluctuant dataset. We then divide the data into 5000  
710 segments and calculate their average  $CV$  values. The local  $CV$  values of IoT, Map, Lognormal  
711 and Weblogs are 0.95, 2.18, 0.63, and 0.005 respectively, which is consistent with the improvement  
712 trends. Intuitively, when the local  $CV$  value is small, the local data is hard-to-fit with a few linear  
713 segments if we adopt an improper  $\epsilon$ , and we need more fine-grained  $\epsilon$  adjustment rather than the  
714 fixed setting. Thus we can expect more performance improvements in this case. The calculation  
715 of actual  $CV$  values of real-world datasets helps to validate our  $\epsilon$  analysis based on the  $CV$  values,  
716 and provides further insight into the scenarios where the proposed method has strong potential to  
717 outperform existing methods.

718 **Ablation Study.** To examine the necessity and the effectiveness of the proposed framework, in  
719 Section 4.3, we compare the proposed framework with three dynamic  $\epsilon$  variants for the FITing-Tree  
720 method. Here we demonstrate the AUNEC relative changes for the Radix-Spline method with the  
721 same three variants in Table 4 and similar conclusions can be drawn.

722 **Theoretical Validation.** In Section 4.4, we show that all the learned index baseline methods learn  
723 similar segment slopes on the Map dataset. Here we illustrate the learned slope results on the IoT,  
724 Weblogs and Lognormal datasets in Figure 11, which supports the Theorem 2 that the learned segment  
725 slopes concentrate on the  $1/\mu_i$  with a bounded relative difference.

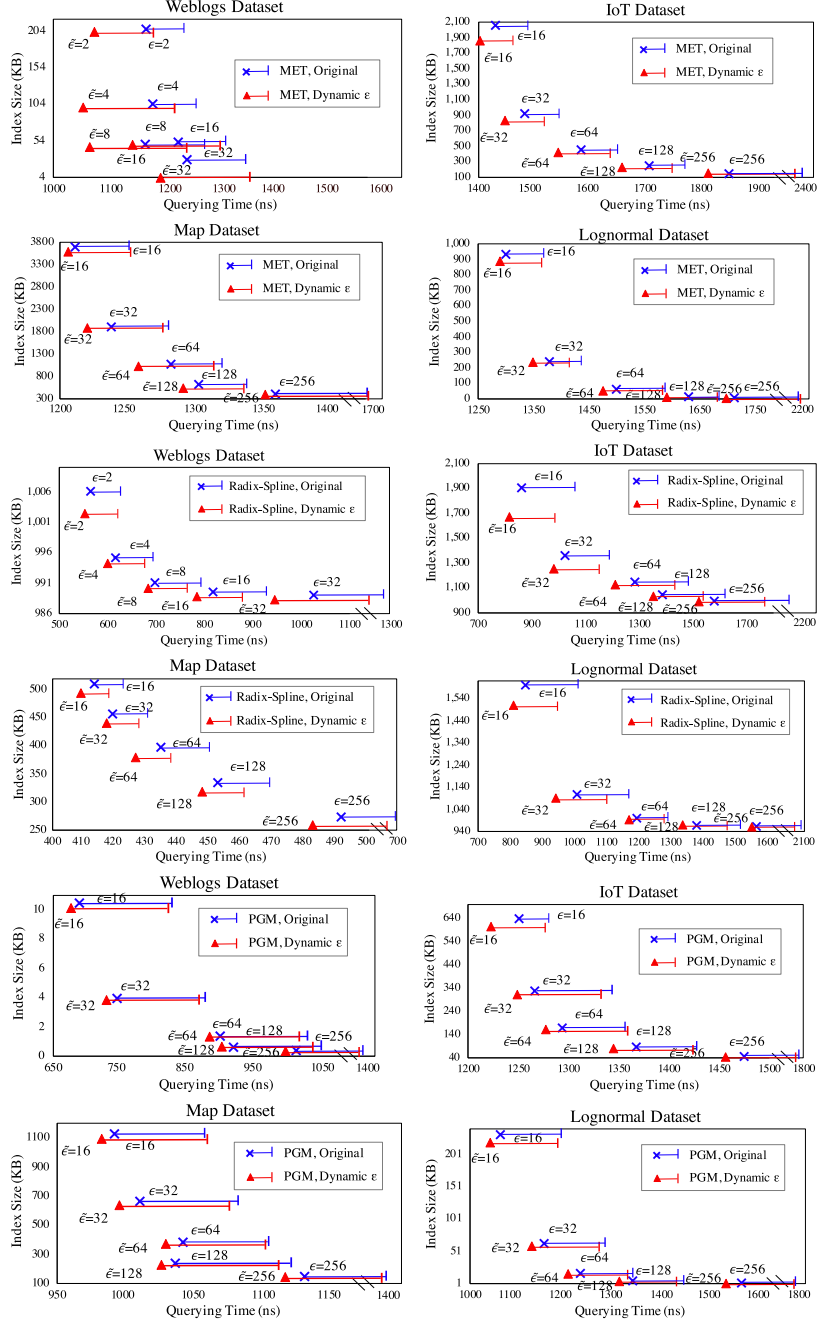


Figure 10: Improvements in terms of querying time for learned index methods with dynamic  $\epsilon$ .

726 Besides, for the comparison between the theoretical bounds and the actual  $SegErr_i$  of all the adopted  
727 learned index methods, we show more results on another two datasets *Gamma* and *Uniform* in Figure  
728 12, where the key intervals of the two datasets follow gamma distribution and uniform distribution  
729 respectively. These results show that the MET method gains actual  $SegErr_i$  within the bounds,  
730 verifying the correctness of the Theorem 1 again. Here all the learned index methods also achieve the  
731 same trends, showing that these methods have the same mathematical forms w.r.t. the  $SegErr_i$ ,  $\epsilon$   
732 and  $\mu/\sigma$ , and hence the  $\epsilon$ -learner can effectively learn the estimator and adaptively choose suitable  $\epsilon$ .



Table 4: The AUNEC relative changes of dynamic  $\epsilon$  variants compared to the Radix-Spline method with the proposed framework.

	Random $\epsilon$	Polynomial Learner	Least Square Learner
Weblogs	+81.23%	+56.20%	-9.56%
IoT	+74.78%	+53.28%	+9.81%
Map	+60.67%	+7.34%	+0.45%
Lognormal	+83.16%	+55.01%	-11.23%
Average	+74.96%	+42.96%	-2.63%

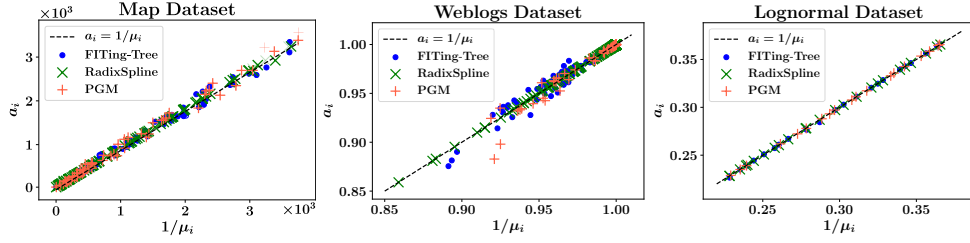


Figure 11: Learned slopes on the IoT, Weblogs and Lognormal datasets.

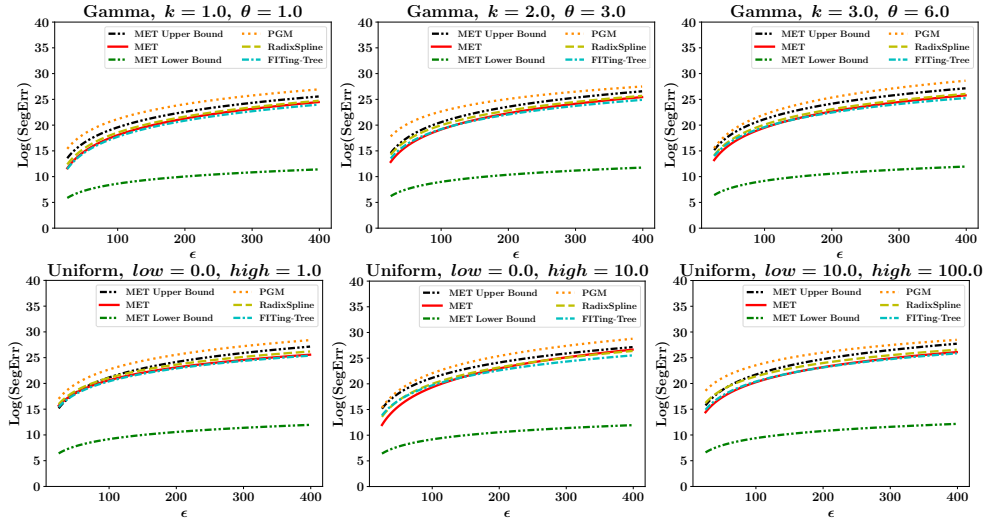


Figure 12: Illustrations of the derived bounds on *Gamma* and *Uniform* datasets.