

Figure 5: Cables for Tracing Unseen Cables Experiment

Table 6: Tracing on Unseen Cables Results

Cable Reference	TR	1	2	3	4	5	Avg.
Tracing Success Rate	6/8	7/8	8/8	7/8	6/8	6/8	<b>40/48=83%</b>
Failures	(I) 2	(I) 1		(II) 1	(I) 1, (III) 1	(II) 1, (III) 1	

## 431 7 Appendix

### 432 7.1 Experiments Failure Mode Analysis

#### 433 7.1.1 Using LTODO for Tracing Cables Unseen During Training

- 434 (1) Retraces previously traced cable (went in a loop).
- 435 (2) Missteps onto a parallel cable.
- 436 (3) Skips a loop.

437 Figure 5 shows the cables tested on. The most common failure mode is (I), retracing previously  
 438 traced cable. This is commonly observed in cases with near parallel segments or in dense loop areas  
 439 within a knot.

#### 440 7.1.2 Using LTODO for Cable Inspection in Multi-Cable Settings

- 441 (I) Misstep in the trace, i.e. the trace did not reach any adapter.
- 442 (II) The trace reaches the wrong adapter.
- 443 (III) The trace reaches the correct adapter but is an incorrect trace.

444 The most common failure mode for the learned tracer, especially in Tier A3, is (I). One reason for  
 445 such failures is the presence of multiple twists along the cable path (particularly in Tier A3 setups,  
 446 which contain more complex inter-cable knot configurations). The tracer is also prone to deviating  
 447 from the correct path on encountering parallel cable segments. In Tier A2, we observe two instances  
 448 of failure mode (III), where the trace was almost entirely correct in that it reached the correct adapter  
 449 but skipped a section of the cable.

450 The most common failure modes across all tiers for the analytic tracer are (II) and (III). The analytic  
 451 tracer particularly struggles in regions of close parallel cable segments and twists. As a result of the  
 452 scoring metric, 87 of the 90 paths that we test reach an adapter; however, 45/90 paths did not reach  
 453 the correct adapter. Even for traces that reach the correct adapter, the trace is incorrect, jumping to  
 454 other cables and skipping sections of the true cable path.

#### 455 7.1.3 Using LTODO for Physical Robot Knot Tying from Demonstrations

- 456 (1) Trace missteps onto a parallel cable.
- 457 (2) Cable shifted during manipulation, not resulting in a knot at the end.

Table 7: Multi-Cable Tracing Results

	Analytic	Learned
Tier A1	3/30	<b>27/30</b>
Tier A2	2/30	<b>23/30</b>
Tier A3	1/30	<b>23/30</b>
Failures	(I) 3, (II) 45, (III) 36	(I) 14, (II) 1, (III) 2

Table 8: Learning From Demos

	Succ. Rate	Failures
Tier B1	5/5	-
Tier B2	4/5	(1) 1
Tier B3	4/5	(1) 1, (2) 1

Table 9: LTODO Experiments

	SGTM 2.0	LTODO (-LT)	LTODO (-CC)	LTODO
Tier C1	2/30	14/30	20/30	<b>24/30</b>
Tier C2	<b>28/30</b>	8/30	21/30	26/30
Tier C3	12/30	14/30	0/30	<b>19/30</b>
Failures	(A) 30, (B) 18	(D) 11, (F) 7 (G) 24, (H) 11	(B) 38, (C) 5, (E) 6	(B) 11, (D) 8 (F) 1

Table 10: LTODO and Physical Robot Experiments (90 total trials)

	Tier D1		Tier D2		Tier D3	
	SGTM 2.0	LTODO	SGTM 2.0	LTODO	SGTM 2.0	LTODO
Knot 1 Succ.	11/15	<b>12/15</b>	6/15	<b>11/15</b>	9/15	<b>14/15</b>
Knot 2 Succ.	-	-	-	-	2/15	<b>6/15</b>
Verif. Rate	<b>11/11</b>	8/12	<b>6/6</b>	6/11	<b>1/2</b>	2/6
Knot 1 Time (min)	1.1±0.1	2.1±0.3	3.5±0.7	3.9±1.1	1.8±0.4	2.0±0.4
Knot 2 Time (min)	-	-	-	-	3.1±1.2	7.5±1.6
Verif. Time (min)	5.7±0.9	6.1±1.4	6.4±1.8	10.1±0.7	5.4	9.6±1.5
Failures	(7) 4	(1) 2, (2) 1 (1) 2, (2) 1	(1) 3, (5) 6 (1) 3, (5) 6	(2) 2, (4) 1 (5) 1	(1) 3, (2) 3, (5) 3 (6) 2, (7) 2	(1) 2, (2) 3 (3) 1, (6) 3

458 Failure mode (1) occurs when the distractor cable creates near parallel sections to the cable of  
 459 interest for knot tying, causing the trace to misstep. Failure mode (2) occurs when the manipulation  
 460 sometimes slightly perturbs the rest of the cable’s position while moving one point of the cable,  
 461 causing the end configuration to not be a knot, as intended.

#### 462 7.1.4 Using LTODO for Knot Detection

- 463 (A) The system fails to detect a knot that is present—a false negative.  
 464 (B) The system detects a knot where there is no knot present—a false positive.  
 465 (C) The tracer retraces previously traced regions of cable.  
 466 (D) The crossing classification and correction schemes fail to infer the correct cable topology.  
 467 (E) The knot detection algorithm does not fully isolate the knot, also getting surrounding trivial  
 468 loops.  
 469 (F) The trace skips a section of the true cable path.  
 470 (G) The trace is incorrect in regions containing a series of close parallel crossings.  
 471 (H) The tracer takes an incorrect turn, jumping to another cable segment.

472 For SGTM 2.0, the most common failure modes are (A) and (B), where it misses knots or incorrectly  
 473 identifies knots when they are out of distribution. For LTODO (-LT), the most common failure  
 474 modes are (F), (G), and (H). All 3 failures are trace-related and result in knots going undetected or  
 475 being incorrectly detected. For LTODO (-CC), the most common failure modes are (B) and (E).  
 476 This is because LTODO (-CC) is unable to distinguish between trivial loops and knots without the  
 477 crossing cancellation scheme. By the same token, LTODO (-CC) is also unable to fully isolate a  
 478 knot from surrounding trivial loops. For LTODO, the most common failure mode is (B). However,  
 479 this is a derivative of failure mode (D), which is present in LTODO (-LT), LTODO (-CC), and  
 480 LTODO. Crossing classification is a common failure mode across all systems and is a bottleneck for  
 481 accurate knot detection. In line with this observation, we hope to dig deeper into accurate crossing  
 482 classification in future work.

#### 483 7.1.5 Using LTODO for Physical Robot Untangling

- 484 (1) Incorrect actions create a complex knot.

- 485 (2) The system misses a grasp on tight knots.
- 486 (3) The cable falls off the workspace.
- 487 (4) The cable drapes on the robot, creating an irrecoverable configuration.
- 488 (5) False termination.
- 489 (6) Manipulation failure.
- 490 (7) Timeout.

491 The main failure modes in LTODO are (1), (2), and (6). Due to incorrect cable topology estimates,  
 492 failure mode (1) occurs: a bad action causes the cable to fall into complex, irrecoverable states.  
 493 Additionally, due to the limitations of the cage-pinch dilation and endpoint separation moves, knots  
 494 sometimes get tighter during the process of untying. While the perception system is still able to  
 495 perceive the knot and select correct grasp points, the robot grippers bump the tight knot, moving  
 496 the entire knot and causing missed grasps (2). Lastly, we experience manipulation failures while  
 497 attempting some grasps as the YuMi has a conservative controller (6). We hope to resolve these  
 498 hardware issues in future work.

499 The main failure modes in SGTM 2.0 are (5) and (7). Perception experiments indicate that SGTM  
 500 2.0 has both false positives and false negatives for cable configurations that are out of distribution.  
 501 (5) occurs when out-of-distribution knots go undetected. (7) occurs when trivial loops are identified  
 502 as knots, preventing the algorithm from terminating.

## 503 7.2 Details on LTODO Methods

### 504 7.2.1 Over/Undercrossing Predictor

505 **Model Architecture and Inference:** The binary classification threshold of 0.275 is determined  
 506 by testing accuracy on a held-out validation set of 75 images on threshold values in the range  
 507  $[0.05, 0.95]$  at intervals of 0.05. Scores  $< 0.275$  indicate undercrossing predictions and scores  
 508  $\geq 0.275$  indicate overcrossing predictions. We output the raw prediction score and a scaled confi-  
 509 dence value (0.5 to 1) indicating the classifier’s probability.

## 510 7.3 Details on Robot Untangling using LTODO

### 511 7.3.1 Knot Definition

512 Consider a pair of points  $p_1$  and  $p_2$  on the cable path at time  $t$  with  $(p_1, p_2 \in \mathcal{C}_t)$ . Knot theory strictly  
 513 operates with closed loops, so to form a loop with the current setup, we construct an imaginary  
 514 cable segment with no crossings joining  $p_1$  to  $p_2$  [44]. This imaginary cable segment passes above  
 515 the manipulation surface to complete the loop between  $p_1$  and  $p_2$  (“ $p_1 \rightarrow p_2$  loop”). A knot exists  
 516 between  $p_1$  and  $p_2$  at time  $t$  if no combination of Reidemeister moves I, II (both shown in Figure 6),  
 517 and III can simplify the  $p_1 \rightarrow p_2$  loop to an unknot, i.e. a crossing-free loop. In this paper, we aim to  
 518 untangle semi-planar knots. For convenience, we define an indicator function  $k(s) : [0, 1] \rightarrow \{0, 1\}$   
 519 which is 1 if the point  $\theta(s)$  lies between any such points  $p_1$  and  $p_2$ , and 0 otherwise.

520 Based on the above knot definition, this objective is to remove all knots, such that  $\int k(s)_0^1 = 0$ .  
 521 In other words, the cable, if treated as a closed loop from the endpoints, can be deformed into an  
 522 unknot. We measure the success rate of the system at removing knots, as well as the time taken to  
 523 remove these knots.

### 524 7.3.2 State Definition

525 We construct line segments between consecutive points on the trace outputted by the learned cable  
 526 tracer (Section 4.1). Crossings are located at the points of intersection of these line segments. We  
 527 use the crossing classifier (Section 4.2) to estimate whether these crossings are over/undercrossings.  
 528 We also implement probabilistic crossing correction with the aim of rectifying classification errors,  
 529 as we describe in Section 4.2.2.

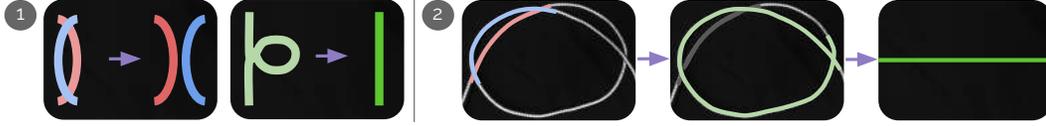


Figure 6: **Reidemeister Moves and Crossing Cancellation:** Left of part 1 depicts Reidemeister Move II. Right of part 1 depicts Reidemeister Move I. Part 2 shows that by algorithmically applying Reidemeister Moves II and I, we can cancel trivial loops, even if they visually appear as knots.

530 We denote the sequence of corrected crossings, in the order that they are encountered in the trace,  
 531 by  $\mathcal{X} = (c_1, \dots, c_n)$ , where  $n$  is the total number of crossings and  $c_1, \dots, c_n$  represent the crossings  
 532 along the trace.

### 533 7.3.3 Crossing Cancellation

534 Crossing cancellation allows for the simplification of cable structure by removing non-essential  
 535 crossings, shown in Figure 6. It allows the system to filter out some trivial configurations as Rei-  
 536 demeister moves maintain knot equivalence [44]. We cancel all pairs of consecutive crossings ( $c_i$ ,  
 537  $c_{i+1}$ ) in  $\mathcal{X}$  for some  $j$ ) that meet any of the following conditions:

- 538 • *Reidemeister I:*  $c_i$  and  $c_{i+1}$  are at the same location, or
- 539 • *Reidemeister II:*  $c_i$  and  $c_{i+1}$  are at the same set of locations as  $c_j$  and  $c_{j+1}$  ( $c_j, c_{j+1} \in \mathcal{X}$ ).  
 540 Additionally,  $c_i$  and  $c_{i+1}$  are either both overcrossings or both undercrossings. We also  
 541 cancel ( $c_j, c_{j+1}$ ) in this case.

542 We algorithmically perform alternating Reidemeister moves I and II as described. We iteratively  
 543 apply this step on the subsequence obtained until there are no such pairs left. We denote the final  
 544 subsequence, where no more crossings can be canceled, by  $\mathcal{X}'$ .

### 545 7.3.4 Knot Detection

546 We say that a subsequence of  $\mathcal{X}'$ ,  $\mathcal{K}_{ij} = (c_i, \dots, c_j)$ , defines a potential knot if:

- 547 •  $c_i$  is an undercrossing, and
- 548 •  $c_j$  is an overcrossing at the same location, and
- 549 • at least one intermediate crossing, i.e. crossing in  $\mathcal{X}'$  that is not  $c_i$  or  $c_j$ , is an overcrossing.

550 The first invariant is a result of the fact that all overcrossings preceding the first undercrossing (as  
 551 seen from an endpoint) are removable. We can derive this by connecting both endpoints from above  
 552 via an imaginary cable (as in Section 7.3.1): all such overcrossings can be removed by manipulating  
 553 the loop formed. The second invariant results from the fact that a cable cannot be knotted without a  
 554 closed loop of crossings. The third and final invariant can be obtained by noting that a configuration  
 555 where all intermediate crossings are undercrossings reduces to the unknot via the application of  
 556 the 3 Reidemeister moves. Therefore, for a knot to exist, it must have at least one intermediate  
 557 overcrossing.

558 Notably, these conditions are necessary, but not sufficient, to identify knots. However, they improve  
 559 the likelihood of bypassing trivial configurations and detecting knots. This increases the system's  
 560 efficiency by enabling it to focus its actions on potential knots.

### 561 7.3.5 Algorithmic Cage-Pinch Point Detection

562 As per the definition introduced in Section 7.3.4, given knot  $\mathcal{K}_{ij} = (c_i, \dots, c_j)$ ,  $c_i$  and  $c_j$  define the  
 563 segments that encompass the knot where  $c_i$  is an undercrossing and  $c_j$  is an overcrossing for the  
 564 same crossing. The pinch point is located on the overcrossing cable segment, intended to increase  
 565 space for the section of cable and endpoint being pulled through. The cage point is located on the

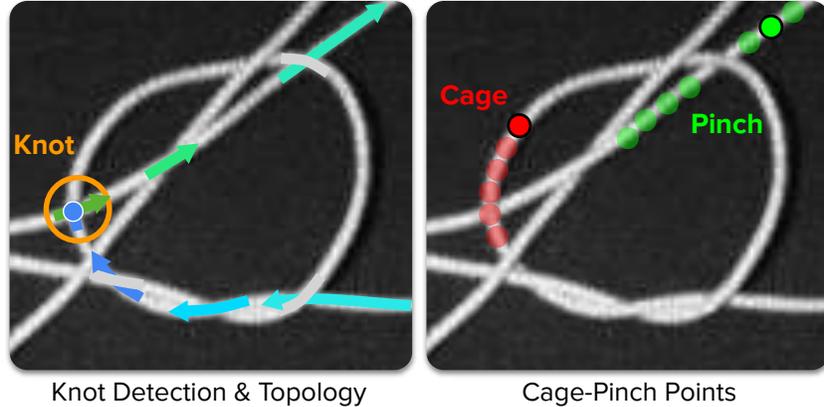


Figure 7: **Knot Detection and Cage Pinch Point Selection:** The left image shows using crossing cancellation rules from knot theory, the knot detection algorithm analytically determines where the knot begins in the cable. The right image shows the survey process for selecting the cap pinch points.

566 undercrossing cable segment. To determine the pinch point, we search from crossing  $c_{u1}$  to crossing  
 567  $c_{u2}$ .  $c_{u1}$  is the previous undercrossing in the knot closest in the trace to  $j$ .  $u2 > j$  and  $c_{u2}$  is the  
 568 next undercrossing after the knot. We search in this region and select the most graspable region to  
 569 pinch at, where graspability ( $G$ ) is defined by the number of pixels that correspond to a cable within  
 570 a given crop and a requirement of sufficient distance from all crossings  $c_i$ . To determine the cage  
 571 point, we search from crossing  $c_i$  to  $c_k$  where  $i < k < j$  and  $c_k$  is the next undercrossing in the knot  
 572 closest in the trace to  $c_i$ . We similarly select the most graspable point. If no points in the search  
 573 space for either the cage or pinch point are graspable, meaning  $G < \mathcal{T}$  where  $\mathcal{T}$  is an experimentally  
 574 derived threshold value, we continue to step along the trace from  $c_{u2}$  for pinch and from  $c_k$  for cage  
 575 until  $G \geq \mathcal{T}$ . This search process is shown in Figure 7.

### 576 7.3.6 Manipulation Primitives

577 We use the same primitives as in SGTm 2.0 (Sliding and Grasping for Tangle Manipulation 2.0)  
 578 [6] to implement LTODo as shown in Figure 8 for untangling long cables. We add a *perturbation*  
 579 move.

580 **Cage-Pinch Dilation:** We use cage-pinch grippers as in Viswanath et al. [5]. We have one gripper  
 581 cage grasp the cable, allowing the cable to slide between the gripper fingers but not slip out. The  
 582 other gripper pinch grasps the cable, holding the cable firmly in place. This is crucial for preventing  
 583 knots in series from colliding and tightening during untangling. The *partial* version of this move  
 584 introduced by Shivakumar et al. [6] separates the grippers to a small, fixed distance of 5 cm.

585 **Reveal Moves:** First, we detect endpoints using a Mask R-CNN object detection model. If both  
 586 endpoints are visible, the robot performs an *Endpoint Separation Move* by grasping at the two end-  
 587 points and then pulling them apart and upwards, away from the workspace, allowing gravity to help  
 588 remove loops before placing the cable back on the workspace. If both endpoints are not visible, the  
 589 robot performs an *Exposure Move*. This is when it pulls in cable segments exiting the workspace.  
 590 Building on prior work, we add a focus on where this move is applied. While tracing, if we detect  
 591 the trace hits the edge, we perform an exposure move at the point where the trace exits the image.

592 **Perturbation Move:** If an endpoint or the cable segment near an endpoint has distracting cable  
 593 segments nearby, making it difficult for the analytic tracer to trace, we perturb it by grasping it and  
 594 translating in the x-y plane by uniformly random displacement in a  $10\text{cm} \times 10\text{cm}$  square in order to  
 595 separate it from slack.

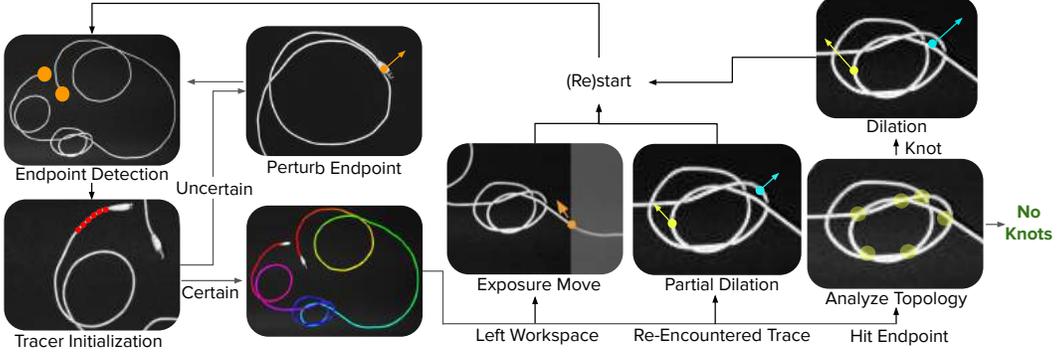


Figure 8: **Untangling Algorithm with LTOD**: We first detect the endpoints and initialize the tracer with start points. If we are not able to obtain start points, we perturb the endpoint and try again. Next, we trace. While tracing, if the cable exits the workspace, we pull the cable towards the center of the workspace. If the tracer gets confused and begins retracing a knot region, we perform a partial cage-pinch dilation that will loosen the knot, intended to make the configuration easier to trace on the next iteration. If the trace is able to successfully complete, we analyze the topology. If there are no knots, we are done. If there are knots, we perform a cage-pinch dilation and return to the first step.

### 596 7.3.7 Cable Untangling System

597 Combining LTOD and the manipulation primitives from Section 7.3.6, the cable untangling al-  
 598 gorithm works as follows: First, detect endpoints and initialize the learned tracer with 6 steps of  
 599 the analytic tracer. If LTOD is unable to get these initialization points, perturb the endpoint from  
 600 which we are tracing and return to the endpoint detect step. Otherwise, during tracing, if the cable  
 601 leaves the workspace, perform an exposure move. If the trace fails and begins retracing itself, which  
 602 can happen in denser knots, perform a partial cage-pinch dilation as in [6]. If the trace completes  
 603 and reaches the other endpoint, analyze the topology. If knots are present, determine the cage-pinch  
 604 points for it, apply a cage-pinch dilation move to them, and repeat the pipeline. If no knots are  
 605 present, the cable is considered to be untangled. The entire system is depicted in Figure 8.

### 606 7.4 Details on Learning from Demos with LTOD

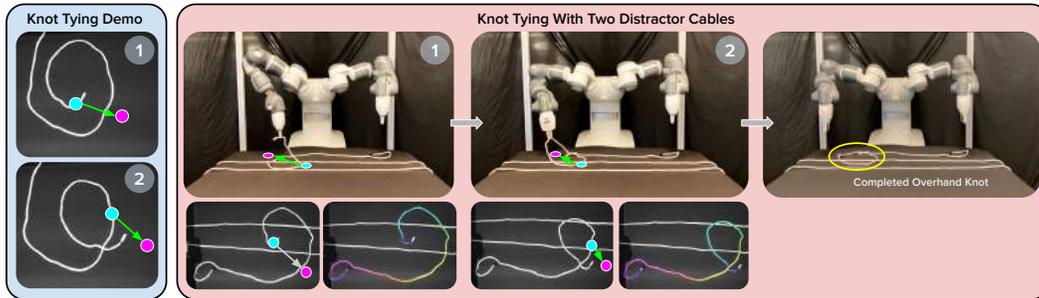


Figure 9: **Using LTOD for Learning from Demos**: The left (blue panel) displays the single human demonstration, indicating the pick and place points for tying an overhand knot. The right (pink panel) shows this demonstration successfully applied to the cable in a different configuration with 2 other distractor cables in the scene. The first step of the demonstration is achieved through an arc length relative action while the second step is achieved through a crossing relative action.

607 When performing state-based imitation, each of the pick and place points  $p_i$  from the demonstration  
 608 is parameterized in the following way: 1) find the point along the trace,  $T$ , closest to the chosen  
 609 point  $\hat{p}_i$  with index  $j$  in  $T$ , 2) find the displacement  $d_i = p_i - \hat{p}_i$  in the local trace-aligned coordinate  
 610 system of  $\hat{p}_i$ , 3) in memory, for point  $p_i$ , store  $d_i$ , arc length of  $\hat{p}_i$  ( $\sum_{x=1}^j T_x - T_{x-1}$ ), and the index  
 611 value of the crossing in the list of crossings just before  $\hat{p}_i$ .

612 When rolling out a policy using this demonstration, there are two ways to do so: 1) relative to the arc  
613 length along the cable, or 2) relative to the fraction of the arc length between the 2 crossing indices.  
614 The way to do so is to find the point on the cable with the same arc length as  $\hat{p}_i$  from the demo or the  
615 fractional arc length between the same 2 crossing indices, depending on the type of demonstration.  
616 Then, apply  $d_i$  in the correct trace-aligned coordinate system. An example demonstration is shown  
617 in Figure 9.

618 Additional conceivable ways, not explored in this work, include relative to the location of a knot  
619 along the cable or others.