SCALEFUSION: SCALABLE INFERENCE OF SPATIAL-TEMPORAL DIFFUSION TRANSFORMERS FOR HIGH-RESOLUTION LONG VIDEO GENERATION

Anonymous Authors¹

ABSTRACT

Recent advancements in training diffusion models have made generating high-quality videos possible. Particularly, the spatial-temporal diffusion transformers (ST-DiTs) emerge as a promising diffusion model architecture for generating videos of high-resolution (1080p) and long duration (20 seconds). However, the quadratic scaling of compute cost with respect to resolution and duration, primarily due to spatial-temporal attention layers processing longer sequences, results in high inference latency of ST-DiTs. This hinders their applicability in time-sensitive scenarios. Existing sequence parallelism techniques, such as DeepSpeed-Ulysses and RingAttention, are not optimally scalable for ST-DiT inference across multiple GPU machines due to cross-machine communication overheads. To address this challenge, we introduce ScaleFusion, a scalable inference engine designed to optimize ST-DiT inference for high-resolution, long video generation. By leveraging the inherent structure of spatial-temporal attention layers, ScaleFusion effectively hides cross-machine communication overhead through novel intra-layer and inter-layer communication scheduling algorithms. This enables strong scaling of $3.60 \times$ on 4 Amazon EC2 p4d.24xlarge machines (32 A100 GPUs) against 1 machine (8 A100 GPUs). Our experiments demonstrate that ScaleFusion outperforms state-of-the-art techniques, achieving an average speedup of $1.36 \times$ (up to $1.58 \times$).

1 INTRODUCTION

High-quality AI-generated videos are becoming increasingly accessible thanks to the recent advancements in videorelated foundation models such as spatial-temporal diffusion transformers (ST-DiTs) (Zheng et al., 2024; Singer et al., 2022; Ma et al., 2024; Gupta et al., 2023). For example, OpenAI's Sora (Brooks et al., 2024) was the first to demonstrate the capability of generating 1080p 1-minutelong photo-realistic videos from text prompts. As ST-DiT models advance in capability and quality, generating highresolution long videos is becoming progressively feasible for industrial production.

However, serving diffusion models on modern GPUs is computationally expensive, especially when generating highresolution long videos. For example, generating a 1080p 4second-long video using the OpenSora ST-DiT model (1.1B
parameters) (Zheng et al., 2024) on a single A100 GPU
requires more than 5 minutes. This inference latency has
significantly limited applying generating high-resolution
long videos to time-sensitive use cases. While recent re-

search has explored innovative modeling paradigms to reduce inference latency (Song et al., 2022; Lu et al., 2022; Song et al., 2023; Liu et al., 2022), we propose to address this challenge through parallel computing, distributing the ST-DiT's workloads across multiple GPUs and parallelizing executions.

Generating high-resolution long videos essentially imposes long sequence workloads on the spatial and temporal attention layers of ST-DiTs. Existing optimizations for serving ST-DiTs employ sequence parallelism (SP) techniques (Liu et al., 2023; Jacobs et al., 2023) to distribute the workloads along sequence dimensions across multiple GPUs for parallel processing (Fang & Zhao, 2024; Zhao et al., 2024). However, these methods suffer from high communication overhead between GPU machines due to lack of computation-communication overlap. As a result, scaling these methods to further reduce ST-DiT inference latency with more compute resources is challenging.

To analyze the communication-overlap challenges of ST-DiTs, we first examine their computation-communication pattern. As shown in Figure 1, a typical ST-DiT architecture consists of alternating spatial and temporal attention layers. Distributed inference of this model on multiple GPUs involves sharding the input tensor along the temporal dimension for parallelizing spatial attention, followed by sharding along the spatial dimension for parallelizing temporal at-

054

 ¹Anonymous Institution, Anonymous City, Anonymous Region, Anonymous Country. Correspondence to: Anonymous Author
 <anon.email@domain.com>.

Preliminary work. Under review by the Machine Learning andSystems (MLSys) Conference. Do not distribute.



Figure 1. The distributed execution of ST-DiT introduce back-toback dependencies.

tention. All-to-All communication collective is inserted between these two layers to re-shard the tensor (Zhao et al., 2024). In this distributed computation-communication pattern, each operation is serially depdendent on the completion of its predecessor. As measured in our experiments, cross-GPU communication overhead accounts for 30-50% of the ST-DiT's end-to-end inference latency when multiple GPU machines are used, hindering scalability to larger GPU clusters.

To address this challenge, prior research has explored lossy techniques to hide communication overhead in the diffusion process for image generation (Li et al., 2024; Wang et al., 2024). These methods leverage the iterative nature of diffusion models, utilizing activation values from previous iteration to compute activations for the current iteration. While similar approaches could be applied to video generation, they inevitably compromise visual quality due to the reliance on stale information. To our knowledge, no existing method effectively hides communication overhead without sacrificing video generation quality.

In this work, we identify a fundamental property underlying the ST-DiT architecture: spatial-temporal independence. We observe that the spatial attention can be executed independently of the temporal dimension, i.e., treating the temporal dimension as a batch dimension. This property enables the input tensor to the spatial attention to be partitioned along the temporal dimension into multiple slices. As a result, the execution of the spatial attention and subsequent all-to-all operation can be pipelined across these slices, leading to the overlap of computation (spatial attention) and communication (all-to-all). A similar pipeline can be established for the temporal attention and its associated all-to-all operation.

Specifically, we devise intra- and inter-layer communication 097 scheduling algorithms to minimize communication over-098 head. Intra-layer communication schedules all-to-all oper-099 ation for each slice to overlap with its associated compu-100 tation layer (spatial or temporal attention) processing the previous slice. Inter-layer communication further schedules partial all-to-all operation for each slice to overlap with the previous computation layer processing the last slice. By 104 combining these techniques, we significantly reduce the 105 communication latency when scaling the model to multiple 106 GPU machines for the inference of ST-DiTs. 107

We summarize ScaleFusion's contributions as follows:



Figure 2. Illustration of video diffusion processes with ST-DiT.

- We analyze and quantify the high communication overhead of existing works utilizing SP techniques for accelerating the inference of ST-DiTs on multiple GPU machines, and identify the key property of ST-DiTs: spatial-temporal independence, which can be leveraged to reduce the communication overhead.
- We are the first to address the high communication overhead in the inference of ST-DiTs on multiple GPU machines and develop ScaleFusion consisting of novel intra-layer and inter-layer communication scheduling algorithms to achieve optimized computationcommunication overlap.
- We conduct comprehensive benchmarks for ScaleFusion and existing techniques on the inference of Open-Sora's ST-DiT model, and demonstrate that ScaleFusion can achieve speedup of 1.40× on average (up to 1.58×) compared to the state-of-the-art work. In addition, ScaleFusion achieves an average strong scaling of 3.60× on 4 p4d.24xlarge machines (32 A100 GPUs) and 1.93× on 2 machines (16 A100 GPUs) against 1 machine (8 A100 GPUs).

2 BACKGROUND

In this section, we briefly introduce the video diffusion processes, its neural network architectures, and existing works that run them for distributed inference on multiple GPU machines.

Video Diffusion Process with Spatial-Temporal Diffusion Transformers Figure 2 illustrates a typical diffusion process (Rombach et al., 2022; Ho et al., 2022) for video generation with spatial-temporal diffusion transformers (ST-DiT). It iteratively denoises a random noise sample through a sequence of sampling steps, gradually transforming it into a clean video. The number of sampling steps required varies depending on the specific methodology for training or inference, ranging from tens to hundreds. In each sampling step, a neural network predicts the denoising operation. This neural network evaluation, which is the major computational bottleneck, is run for every sampling step. To address this computational challenge, prior research has primarily focused on reducing the number of sampling steps through algorithmic innovations such as training-free methods (Song



Figure 3. Distributed execution of ST-DiT. For simplicity, the auxiliary layers such as layer normalizations are not displayed in the figure.

et al., 2022; Lu et al., 2022) and novel training paradigms (Song et al., 2023; Liu et al., 2022). In contrast, our work aims to accelerate the individual sampling steps by optimizing the inference system of the neural network through distributed computing.

Among all the neural network architectures for the diffusion process, ST-DiTs have emerged as a promising architecture for high-quality video generation in recent years (Brooks et al., 2024; Ma et al., 2024; Gupta et al., 2023; Yang et al., 2024; Zheng et al., 2024; Singer et al., 2022). The backbone of this architecture was initially proposed by ViViT (Arnab et al., 2021) for video classification tasks. However, the quadratic compute cost scaling of the attention computation in spatial and temporal layers with respect to video resolution and duration poses significant challenges for generating high-resolution long videos. In this work, we optimize the distributed inference for the ST-DiT architecture employing a ViViT backbone with repeated alternating spatial-temporal attention layers as in Figure 1.

Distributed Inference of ST-DiTs on Multiple GPUs Figure 3 demonstrates the distributed inference of a typical ST-DiT containing repeated alternating spatial and temporal layers on multiple GPUs. As the sizes of the spatial and temporal and dimensions scale linearly with the resolution and duration of the videos to be generated, existing systems 152 utilize sequence parallelism (SP) techniques for accelerating 153 the generation of high-resolution long videos (Zheng et al., 154 2024; Zhao et al., 2024). These systems partition the in-155 put tensor along spatial and temporal sequence dimensions, 156 distributing the resulting sharded tensors across multiple 157 GPUs for parallel processing. To compute spatial or tempo-158 ral attention, which requires global information along the 159 sequence dimension, communication operations are used to 160 gather the sharded tensors, enabling the execution of atten-161 tion layers. This section reviews the existing state-of-the-art 162 distributed inference technique proposed by (Zhao et al., 163 2024) for ST-DiTs. The discussion of model parallelism 164



Figure 4. The performance breakdown and strong scaling speedup when executing ST-DiT with the state-of-the-art work (Zhao et al., 2024) on single or multiple GPUs, where N represents the number machines and G represent the total number of GPUs.

techniques for other models is deferred to Section 6.

Consider an input tensor to an ST-DiT model has a shape of [B, T, S, C], where B denotes batch size, T and S denote the sizes of temporal and spatial sequence dimension, and C denotes the size of hidden dimension. Note that the attention computation of spatial and temporal layers occurs along the dimension of S and T, respectively. To run the model inference with P GPUs in parallel, the input tensor is split along the spatial dimension into P shards, each of which is distributed to a GPU for processing. Note that each shard has a shape of [B, T, S/P, C]. We name this layout spatial-distributed. As the execution of the spatial attention requires a complete spatial sequence dimension, (Zhao et al., 2024) inserts an all-to-all operation to gather shards from GPUs along the spatial sequence dimension and scatter along the temporal sequence dimension, resulting in a shard of shape of [B, T/P, S, C]. We name this layout temporal-distributed. Similarly, when executing the following temporal layer, a complete temporal sequence dimension is required. As the input tensor is temporal-distributed after the spatial layer, another all-to-all operation is inserted to transform temporal-distributed layout back to the spatialdistributed layout, by gathering shards from GPUs along the temporal dimension and scattering along the spatial dimension. After the execution of a pair of a spatial layer and a temporal layer, the output tensor transforms to the original spatial-distributed layout with the same shape as the input tensor, i.e., [B, T, S/P, C]. The aforementioned process repeats for the remaining layers.

3 CHALLENGES & OPPORTUNITIES

While there are works utilizing SP techniques to shard attention layers across multiple GPUs (Fang & Zhao, 2024; Zhao et al., 2024), none of them has addressed high communication overheads when scaling ST-DiT inference across multiple GPU machines to the best of our knowledge. Figure 4 165 illustrates the computation-communication breakdown of an existing state-of-the-art work (Zhao et al., 2024) generating 167 a 1080p 8 second video. While existing SP techniques has 168 negligible communication overhead (i.e. by only up to 3%) 169 in single-machine settings, they incur huge communication 170 overhead when scaling up to multiple machines. Specifi-171 cally, these techniques incur 34% time and 44% executing 172 communication operations in 2-machines and 4-machines 173 experiments, leading to 1.53× and 1.83× slowdowns com-174 pared to the theoretical strong scaling speedups. This is 175 because GPUs on the same machine are connected with 176 high-bandwidth interconnects such as NVLink or NVSwitch (NVIDIA, b) and their bandwidth is much larger than the 178 cross-machine networks. Thus, even with the state-of-the-179 art SP techniques, generating long high-resolution ST-DiT 180 does not scale efficiently with multiple GPU machines.

181 To reduce communication overhead, a common wisdom is 182 to execute the communication operations concurrently with 183 the computation operations and hide the communication 184 latency as much as possible (Aminabadi et al., 2022; Huang 185 et al., 2019; Harlap et al., 2018; Jiang et al., 2024; Jayarajan 186 et al., 2019). However, hiding the communication overhead 187 for ST-DiTs is non-trivial. We elucidate the challenges and 188 opportunities as follows. 189

190 Challenge 1: Back-to-back latency between communi-191 cation and computation operations. As discussed in Section 2, ST-DiT contains alternating spatial and tempo-193 ral layers and requires all-to-all operations that transform back and forth between the spatial-distributed layout and 195 temporal-distributed layout. As shown in Figure 1, since 196 the input tensor of each all-to-all operation is the output tensor of its previous spatial/temporal layer, and its output 198 tensor becomes the input tensor of the next spatial/temporal 199 layer, executing ST-DiT on multiple GPUs naturally intro-200 duce back-to-back dependencies between the communication operations and computation operations. Therefore, it is 202 non-trivial to execute communication and computation operations concurrently to hide the communication overhead.

204 **Opportunity 1: Intra-layer communication-computation overlap.** To solve the challenge 1, we observe that the 206 aforementioned back-to-back dependencies only exist along the spatial dimension in spatial layers and along the tem-208 poral dimension in temporal layers. Thus, the layer com-209 putation operations is totally independent along the other 210 dimension, i.e. temporal dimension in spatial layers and 211 spatial dimension in temporal layers. As a result, we can 212 divide the layer execution into slices where the back-to-213 back communication-computation dependencies only exist 214 within each slice. By doing so, we create new opportu-215 nities for overlapping the communication operations with 216 computation operations between different slices. 217

218 Challenge 2: Diminishing speedups with an increased

219



Figure 5. Computation and communication time of sharding the layer executions with the-state-of-the-art work (Zhao et al., 2024).



Figure 6. The breakdown of the execution time of spatial layers and temporal layers when generating videos with different resolutions and duration with the state-of-the-art work (Zhao et al., 2024).

number of slices. Figure 5 shows how communication and communication time vary when sharding the layer execution into different number of slices. While increasing the number of slices allows more communication operations to be overlapped (Figure 5a), it also increases the computation overhead (Figure 5b). Since the model is computation bounded as indicated by Figure 4, the computation overhead will cause diminished speedups when increasing the number of slices (Figure 5c). Thus, it is not always beneficial to overlap more communication operations by increasing the number of slices.

Opportunity 2: Inter-layer communication-computation overlap. Figure 6 illustrates the breakdown of the execution time with the state-of-the-art work on the computation and communication operations. We observe that the state-of-the-art implementation of typical video generation workloads are bounded by computation operations. Thus, instead of partitioning both the computation and communication operations by simply increasing the number of slices, we could avoid computation overhead by only further partitioning the communication operations. Even though the overall communication time could be slightly increased, it should be beneficial as long as it can be overlapped with computation operations. Specifically, we consider partition the communication operations and move part of them to overlap with the communication operations in the previous spatial/temporal layer. In this way, we could overlap more communication operations at no cost of computation overhead, thus further reducing the communication overhead.



Figure 7. Overview of ScaleFusion's key ideas. For simplicity, the figure demonstrates only the execution of one of the L pairs of spatial and temporal layers in a ST-DiT.

4 SCALEFUSION

In this work, we propose ScaleFusion, which composes of intra- and inter-layer communication scheduling algorithms.

4.1 Overview

We delinate the overview of the execution of ST-DiTs on multiple GPUs with the state-of-the-art existing work (Zhao et al., 2024) and ScaleFusion in Figure 7. We observe that the existing work suffer from high communication overhead, since they do not overlap the all-to-all communication operations with computations (Figure 7a). Instead, in this paper we propose the following key principle named *spatialtemporal independence*, which allows us to break each temporal and each spatial layer into multiple slices that can be executed independently. Based on this key principle, we propose ScaleFusion that composes of intra- and inter-layer communication scheduling algorithms to achieve efficient computation and communication overlap when executing ST-DiTs on multiple GPUs.

61 **Intra-Layer Communication Scheduling Algorithm** As 62 shown in Figure 7b, we schedule communication operations 63 of each slice concurrently with the computation operations 64 of the next slice in both spatial and temporal layers and 65 execute each communication and computation slice in a 66 pipelined fashion. Specifically, when executing the com-67 putation of one slice, we schedule the communication of 68 the next slice concurrently by leveraging the asynchronous 69 communication operation. In this way, ScaleFusion is able 70 to overlap communication with computation between differ-71 ence slices.

Inter-Layer Communication Scheduling Algorithm While intra-layer scheduling algorithm can overlap a large portion of communication with computation operations, the communication of the first slice and the computation on the last slice cannot be overlapped. Note that this issue cannot be simply solved by increasing the number of slices as it incurs computation overhead (see Section 5.3 details). We propose inter-layer communication scheduling algorithm, which further divide the first communication slice into multiple partitions, and schedule communication of a few partitions when executing the computation of the last slice in the previous layer, without blocking the critical execution path.

Through the intra- and inter-layer communication scheduling algorithms, we can achieve almost fully overlapped communication and thus significantly reduce the communication overhead.

4.2 Key Principle: Spatial-Temporal Independence

In the execution of a ST-DiT, we identify that the tokens in temporal dimension are totally independent when executing attention in spatial layers and similarly the tokens in spatial dimension are totally independent when executing attention in temporal layers.

Specifically, assume the input tensor has a shape of [B, T, S, C], where B, T, S, and C represents the size of batch, temporal, spatial, and hidden dimensions respectively. Suppose we have divided the input tensor x along the temporal dimension into three slices, and their corresponding input tensors are denoted as $x_{:,t_1,:}, x_{:,t_2,:}$, and $x_{:,t_3,:}$ respectively. Then we can compute the attention in spatial layers independently and then concatenated together, namely,

$$SpatialLayer(A2A_{T \to S}(cat_{T}(x_{:,t_{1},:}, x_{:,t_{2},:}, x_{:,t_{3},:})))$$

$$= cat_{T} \begin{pmatrix} A2A_{T \to S}(SpatialLayer(x_{:,t_{1},:})), \\ A2A_{T \to S}(SpatialLayer(x_{:,t_{2},:})), \\ A2A_{T \to S}(SpatialLayer(x_{:,t_{3},:})) \end{pmatrix}$$

$$(1)$$

where cat_T denotes tensor concatenate along the temporal dimension and $\operatorname{A2A}_{T \to S}$ denotes the all-to-all operation that scatters the temporal dimension and gathers the spatial dimension.

Similarly, for temporal layers, assume we have divided the input tensors along the spatial dimension into three partitions and denote the corresponding input tensors by $x_{:,:,s_1}$, $x_{:,:,s_2}$, and $x_{:,:,s_3}$, then we have,

$$TemporalLayer(cat_{S}(A2A_{S \to T}(x_{:,:,s_{1}}, x_{:,:,s_{2}}, x_{:,:,s_{3}})))$$

$$= cat_{S} \begin{pmatrix} A2A_{S \to T}(TemporalLayer(x_{:,:,s_{1}})), \\ A2A_{S \to T}(TemporalLayer(x_{:,:,s_{2}})), \\ A2A_{S \to T}(TemporalLayer(x_{:,:,s_{3}})) \end{pmatrix}$$

$$(2)$$

where cat_S denotes tensor concatenate along the spatial dimension and $\operatorname{A2A}_{S \to T}$ denotes the all-to-all operation that scatters the spatial dimension and gathers the temporal dimension.

Thus, instead of execute the spatial or temporal layer on the
whole input tensors, we could divide the input tensors into
multiple slices and execute each slice separately and finally
concatenate the results together. This key principle presents
potential opportunity for overlapping the communication
overhead with computation, for which we propose intra- and
inter-layer scheduling algorithms.

4.3 Intra-Layer Communication Scheduling Algorithm

282

329

285 As mentioned in Section 2, when executed with multiple 286 GPUs, each spatial and each temporal layer in ST-DiT incur 287 an all-to-all operation before the execution of the layer itself. 288 To hide the latency of all-to-all operations, we leverage 289 the key principle to divide the layer execution into several 290 slices. Since the execution of each slice is independent, we 291 can execute the all-to-all operation and the layer execution 292 separately of each slice and organize the execution in a way 293 such that the all-to-all operation execution is overlapped 294 with the layer computation from the subsequent slice. 295

296 Specifically, as shown in Figure 8a, we 1 introduce two 297 hyper-parameters N_T ($N_T \leq T$) and N_S ($N_S \leq S$), mean-298 ing the number of slices in the temporal and the spatial di-299 mension respectively. We 2 first rearrange the input tensor 300 x on each GPU to a shape of $[B, N_T, N_S, T/N_T, S/N_S/P]$, 301 where B, T, S, and C denote the size of batch, temporal, 302 spatial, and hidden dimension respectively and P denotes 303 the number of GPUs. Note that T and S do not need to 304 be divisible by N_T or N_S . If N_T does not divide T or 305 N_S does not divide S, we just evenly split T and S into 306 N_T and N_S parts respectively, and simply employ a jagged 307 tensor to represent x where each $x_{:,i,j}$ could have different 308 dimensions among all $i \in [0, N_T)$ and $j \in [0, N_S)$. When 309 executing the spatial layers, for communication operations, we **3** start by only executing the all-to-all operation in the 311 first temporal slice of a spatial layer. Then, immediately 312 after the all-to-all operation in the current temporal slice is 313 finished, we 4 start executing the all-to-all operation in the 314 following temporal slice. For computation operations, we 315 simply execute the computation operations in the order of 316 the temporal slices. To meet the data dependency, we block 317 the execution of the computation operations until the depen-318 dent all-to-all operation has finished. Similarly, for temporal 319 layers, we **S** execute all all-to-all operations in a back-to-320 back manner and we block the execution of the computation operations until the dependent all-to-all operation is finished. By dividing the layer execution into different slices, we can 323 overlap part of the all-to-all operations with the computation 324 operations and thus reducing the communication overhead. 325

Ideally, we only have the all-to-all operation in the first slice and the layer computation operations in the last slice that are not overlapped. Suppose the communication overhead is $Comm_T$ and $Comm_S$ for temporal layers and spatial layers respectively. By using the intra-layer communication scheduling algorithm, if the all-to-all operations in all slices except the last slice are overlapped completely, we could reduce the communication overhead to $Comm_T/N_T + Comm_S/N_S$.

According to the formula above, it is tempting to increase the number of spatial/temporal slices, since it could further reduce the communication overhead. However, in practice we found that using a large number of temporal/spatial slices leads to many small CUDA kernel launches for the computation operations, which causes high computation overhead and in turn slows down the overall execution time (see Section 5.3 for details). As a result, simply applying intra-layer communication scheduling algorithm can only overlap a certain amount of communication operations without incurring computation overhead. To solve this issue, we propose the following inter-layer communication scheduling algorithm to conduct inter-layer communication-computation overlap.

4.4 Inter-Layer Communication Scheduling Algorithm

While intra-layer communication scheduling algorithm can overlap the all-to-all communication operations and the layer computation operations, the all-to-all operation in the first slice and the layer computation in the last slice is still not overlapped. To further reduce the communication overhead without incurring computation overhead, we introduce the inter-layer communication scheduling algorithm, which overlaps part of the all-to-all communication in the first slice with the computation operations in the last slice.

The key insight behind inter-layer communication scheduling algorithm is that we could lift part of the all-to-all operations to overlap with the computation operations in the last slice of the previous spatial or temporal layer. Specifically, in the execution of a temporal layer, we find that each all-toall operation can be further decomposed along the temporal dimension. For example, assume we have partitioned both the spatial and temporal dimension into three slices, which we denote as s_1 , s_2 , s_3 , and t_1 , t_2 , t_3 respectively. Then, for *each* spatial partition s_i , we can **6** decompose the all-to-all operation into three partitions,

$$A2A_{T \to S}(cat_{T}(x_{:,t_{1},s_{i}}, x_{:,t_{2},s_{i}}, x_{:,t_{3},s_{i}})))$$

$$= cat_{T}\begin{pmatrix}A2A_{T \to S}(x_{:,t_{1},s_{i}}), \\ A2A_{T \to S}(x_{:,t_{2},s_{i}}), \\ A2A_{T \to S}(x_{:,t_{3},s_{i}})\end{pmatrix}$$
(3)

By doing so, we can \bigcirc lift the first two communication partitions to overlap with the computation operations in the last slice of the previous spatial layer. Similarly, in the execution of a spatial layer, we could also B divide the all-to-all operation in the first temporal slice into three



Figure 8. Illustration of (a) the intra-layer communication scheduling algorithm and (b) the inter-layer communication scheduling algorithm within one spatial layer and a (subsequent) temporal layer of a ST-DiT, when configuring $N_T = N_S = 3$.

partitions along the spatial dimension. Namely, for *each* temporal partition t_i , we have,

$$A2A_{S \to T}(cat_{S}(x_{:,t_{i},s_{1}}, x_{:,t_{i},s_{2}}, x_{:,t_{i},s_{3}})))$$

$$= cat_{S}\begin{pmatrix}A2A_{S \to T}(x_{:,t_{i},s_{1}}), \\ A2A_{S \to T}(x_{:,t_{i},s_{2}}), \\ A2A_{S \to T}(x_{:,t_{i},s_{3}}) \end{pmatrix}$$
(4)

And we can also **()** lift two communication partitions to overlap with the computation operations in the last slice of the previous temporal layer. Note that this process is only not applicable at the first (spatial) layer of ST-DiTs, simply because there is no previous temporal layer to overlap with. However, its communication overhead is negligible compared to the end-to-end execution time of the whole model.

In ScaleFusion, we design the inter-layer communication algorithm that generalizes the above example by defining another two hyperparameters, L_T and L_S , representing the numbers of the all-to-all operation partitions in the first slice of each spatial/temporal layer lifted to overlap with the computation operations in the last slice of the previous temporal/spatial layer. The example above is a special case of inter-layer communication scheduling algorithm by choosing $L_T = L_S = 2$. Intuitively, L_T and L_S controls the amount of communication operations in the first slice 378 of the spatial/temporal layer that are lifted to the previous 379 layer. With a larger value of L_T and L_S , we have more 380 communication operations in the spatial/temporal moved to 381 the previous temporal/spatial layer. Besides, by choosing 382 a different L_T and L_S , ScaleFusion is adaptable to differ-383 ent video generation workloads. If the temporal layer has 384

longer computation time than the spatial layer, we choose $L_T \ge L_S$, which creates more opportunities to overlap communication operations with computation operations in the temporal layer. On the other hand, if the spatial layer has longer computation time, we simply choose $L_T < L_S$.

Note that the inter-layer communication scheduling algorithm can not only further overlap the communication operations that are not able to be overlapped by the intra-layer communication scheduling algorithm, it will not increase the computation overhead when applied on top of the intralayer communication scheduling algorithm since we do not increase the number of partitions of computation operations. Thus, applying the inter-layer communication scheduling algorithm on top of the intra-layer computation scheduling algorithm can solve both challenge 1 and challenge 2 as described in Section 3.

Specifically, if the computation operations in each slice takes longer time than the communication operations, by using inter-layer communication scheduling algorithm, we can reduce the communication overhead to $(Comm_T + Comm_S)/N_SN_T$, since the only communication operations that are not overlapped with computations are the first spatial/temporal communication partition in the first slice of each temporal/spatial layer.

Additionally, by slicing the all-to-all operations, we also reduce the peak memory usage (see Section 5.2 for details). This is because each all-to-all operation requires a temporary buffer to receive the tensor values from other GPUs. Therefore, by dividing the all-to-all operations in different slices, we can release the corresponding temporary buffer after each slice of the all-to-all operation is finished, resultingin lower peak memory usages.

Overall, ScaleFusion maximally overlaps the communication operations with the computation operations in all evaluated video generation workloads, reduce the communication cost and peak memory footprint, and thus enable efficient high resolution long video generation on multiple GPUs.

5 EVALUATION

387

388

389

390

394

3953965.1 Experiment Setups

Baselines We compare ScaleFusion with two existing
implementations for ST-DiT: (1) OpenSora uses DeepSpeedUlysses (Jacobs et al., 2023) utilizing all-to-all operation to
transpose between the head and the spatial dimensions; (2)
DSP (Zhao et al., 2024) proposes a new SP technique that
reduces communication volume by transposing between the
temporal and spatial dimensions.

404 Hardware & Software We conduct all experiments on 405 Amazon p4d.24xlarge machines that are commonly 406 used in practice for large model inference. Each machine 407 is equipped with 8 NVIDIA A100 GPUs (40 GiB) with 408 NVIDIA Driver 535.183.01, CUDA 12.2, NCCL v2.19.3 409 and PyTorch v2.2.2 installed. The GPUs within the same 410 machine are connected with NVSwitch (NVIDIA, b) and 411 the machines are connected by AWS Elastic Fabric Adapter 412 (AWS) with 400 Gbps cross-machine network bandwidth. 413

414 Neural Networks We adopt the open-source implemen415 tation of ST-DiT, OpenSora v1.2 (Zheng et al., 2024) with
416 memory efficient attention enabled (Dao et al., 2022; Dao,
417 2023), as our codebase for all experiments.

418 Measurements For runtime latency measurement, we 419 gather the traces produced by NVIDIA Nsight System 420 (NVIDIA, a) when running each experiments. Then cal-421 culate the runtime and detailed performance breakdown 422 through analyzing the traces. Note that the end-to-end per-423 formance includes all communication and computation op-424 erations in the ST-DiT. Particularly, it includes the commu-425 nication operation of the first partition of the first (spatial) 426 layer when using ScaleFusion. For the theoretical lower 427 bound time (if presented), we simply use the minimum 428 time among all methods in single-machine experiments. In 429 multi-machine experiments, the theoretical lower bound is 430 the single-machine scaling time divided by the number of 431 machines. For peak memory usage measurement, we use 432 torch.cuda.max_memory_reserved. 433

434 **Hyperparameter Choices** We use $N_T = N_S = 4$ and 435 $L_T = 1$ and $L_S = 3$ by default since we found these hyper-436 parameters can already hide the communication overhead 437 pretty well in all our experiments with a variety of video gen-438 eration workloads, as discussed in Section 5.3. However, we 439 also make ScaleFusion configurable by providing interfaces that allow users to set these hyperparameters flexibly.

5.2 End-to-End Performance Evaluation

Figure 9 shows the end-to-end performance of prior works and ScaleFusion. We show normalized speedups of Scale-Fusion over OpenSora (OS), DSP (DSP), and ScaleFusion with only intra-layer communication scheduling algorithm (SO), and we show the exact latency numbers for Scale-Fusion (SF) on their corresponding the bar labels. Note that the OpenSora default implementation (OS) cannot scale up to 4 GPU machines. This is because OpenSora uses DeepSpeed-Ulysses (Jacobs et al., 2023) whose parallelism are limited by the number of heads. Since the benchmarked ST-DiT has a small number of heads (i.e., 16), OpenSora's ST-DiT cannot scale up to more than 16 GPUs and thus is not feasible to run on 4 machines of 32 GPUs. We make the following conclusions.

Comparison to Prior Works We found that both Open-Sora and DSP have high communication overhead, which is not hided by the computation. In contrast, ScaleFusion can almost hide all communication costs and achieves a superior speedup by $1.32 \times$ on average (up to $1.52 \times$) in 2-machines experiments, and by $1.40 \times$ on average (up to $1.58 \times$) in 4-machines experiments over the state-of-the-art work (Zhao et al., 2024). Besides, we also found that as the videos become longer and longer, ScaleFusion can achieve more speedup over prior works, demonstrating its effectiveness for speeding up long high-resolution video generation.

Strong Scaling In our experiments, strong scaling measures the inference speedup as the number of GPUs increases while maintaining a constant workload. We found that prior works suffer from high communication overhead when scaling up to more than one GPU machines. In contrast, ScaleFusion has minimum communication overhead in all scenarios with an average strong scaling of $1.93 \times$ in the experiments of 2 machines and $3.60 \times$ of 4 machines against 1 machine.

Speedup Breakdown We observe that using only the intra-layer communication scheduling algorithm can provide $1.26 \times$ speedup on average (up to $1.41 \times$), and enabling inter-layer communication scheduling algorithms can provide another $1.08 \times$ speedup on average (up to $1.13 \times$) on top of the intra-layer communication scheduling algorithm. Thus, we conclude that both the intra- and inter-layer communication scheduling algorithms are needed to achieve the best performance in ScaleFusion.

Peak Memory Usage We conclude that while both Open-Sora and DSP suffers from out-of-memory (OOM) errors when generating high resolution or long video generation (e.g. when generating 4k videos for 16 seconds), ScaleFu-



Figure 9. The end-to-end performance of ScaleFusion compared to prior works, where " \times " denotes OOM errors, "N/A" means not applicable, and *N* represents the number of GPU machines. OS stands for OpenSora implementation with DeepSpeed-Ulysses (Jacobs et al., 2023), DSP stands for the dynamic sequence parallelism (Zhao et al., 2024), SO and SF stands for ScaleFusion (this work) with only intra-layer communication scheduling algorithm enabled and with both intra- and inter-layer communication scheduling algorithms enabled, respectively.

sion can generate such videos without hitting OOM errors. ScaleFusion achieves a $1.28 \times$ peak memory usage reduction on average (up to $1.43 \times$) compared to OpenSora and $1.87 \times$ reduction on average (up to $2.33 \times$) compared to DSP.

Overall, we conclude that ScaleFusion can achieve a significant speedup over all prior works, minimize the communication overhead, and thus enable efficient high-resolution long video generation.

5.3 Sensitivity Study on ScaleFusion's Hyperparameters

Intra-Laver Communication Scheduling Algorithm Figure 10 shows how the computation time and the nonoverlapped communication time varies across different number of slices with only intra-layer scheduling algorithm enabled. We make three conclusions from the figure. First, we conclude that sharding the layer execution into multiple slices incurs both computation and communication overhead (Figure 10a and 10b) that increases with a larger number of slices. Second, we find that non-overlapped communication time is decreasing with larger number of slices (Figure 10c). This is because only the first all-to-all operation in the first slice of each layer is not overlapped which decreases with an increasing number of slices. Third, we notice that while the total time is not always decreasing with a larger number of 489 slices (Figure 10d). Compared to the default hyperparame-490 ters we are using, i.e. $N_T = N_S = 4$, the best configuration 491 only brings an additional speedup of 2.7% on average. Thus, 492 we conclude that simply varying its hyperparameters cannot 493 achieve the same level of speedup compared to the case 494

where both intra- and inter-layer communication scheduling algorithms are enabled, i.e. by $1.08 \times$ speedup on average.

Inter-Layer Communication Scheduling Algorithm For inter-layer communication scheduling algorithm, we always set $L_T = 1$ and vary L_S since the computation operations in temporal layers are always shorter than the spatial layers in all experiments based on our measurements. Thus, it is always beneficial to lift communication operations from temporal layers to spatial layers but not the other way.

Figure 11 illustrates how the overall execution time varies with different hyperparameters in ScaleFusion. Compared to the default hyperparameters we are using, i.e. $N_T = N_S = 4$, $L_T = 1$, $L_S = 3$, the best configuration only brings 0.7% speedup on average. Thus, we conclude that while the inter-layer communication scheduling algorithm is sensitive to the specific hyperparameters, simply using the ScaleFusion's default hyperparameters can achieve close-tooptimal performance across a variety of scenarios.

6 RELATED WORKS

Sequence Parallelism To efficiently serving transformerbased large foundation models on multiple GPUs with long sequence length, different SP techniques have been studied to efficiently distribute the execution of attention layers on multiple GPUs.

DeepSpeed-Ulysses (Jacobs et al., 2023) leverages all-toall operation to gather and scatter tensors between sequence and attention head dimensions without computation-



Figure 10. Performance of ScaleFusion with different number of slices when only intra-layer communication scheduling algorithm is enabled. The X-axis ticks represent $N_T = N_S$. The star stands for the optimal configuration.

communication overlap. RingAttention (Liu et al., 2023) 522 divides the query, key, and value tensors along the sequence 523 524 dimension into multiple partitions and assigns each partition to a single GPU. RingAttention sends and receives each 525 526 key and value tensor partition from and to its neighboring GPU in a ring-like manner and executes the computation 527 concurrently with the current key and value tensor partition. 528 While RingAttention can overlap the computation with com-529 munication, the communication volume is much larger than 530 DeepSpeed Ulysses. This limits its scalability on multi-531 ple GPU machines where the inter-machine bandwidth is 532 constrained. DSP (Zhao et al., 2024) optimizes attention 533 operations with multiple sequence dimensions and proposes 534 to apply a similar idea to DeepSpeed-Ulysses. It uses all-to-535 all operation to gather and scatter tensors between different 536 sequence dimensions. However, unlike ScaleFusion, these 537 two techniques still incur high communication overhead in 538 end-to-end inference. 539

540 Pipeline Parallelism Pipeline parallelism (PP) (Huang 541 et al., 2019; Jayarajan et al., 2019; Harlap et al., 2018) 542 partition models into subsets of layers and assign the subsets 543 to multiple GPUs. In execution, PP shards the input tensor 544 along the batch dimension into micro-batches. Each GPU 545 processes one micro-batch and transfers the output to the 546 next GPU in a pipelined manner to keep all GPUs busy. PP 547 works orthogonal with ScaleFusion, since both the inter-548 and the intra-layer communication scheduling algorithm do 549



Figure 11. The end-to-end execution time of ScaleFusion with different number of slices when both intra- and inter-layer communication scheduling algorithms are enabled. The first level ticks of X-axis represent L_S and the second level ticks denotes the number of slices, i.e., $N_T = N_S$. The star stands for the optimal configuration.

not modify the batch dimension and thus can be applied in conjunction with PP.

Lossy Optimizations for Diffusion Models To reduce the communication overhead, prior works have found that diffusion models are resilient to stale inputs, meaning that even if using the outdated input tensors from previous sampling step, diffusion models can still generate images with good quality (Li et al., 2024; Wang et al., 2024). This property breaks the data dependency between consecutive layers, therefore allows concurrent executions of computation and communication to hide communication overhead. However, the accuracy of these methods decrease with an increasing number of GPUs. This is because using more devices implies a larger portion of input tensors are outdated. The quality of the generated contents could further degrade with more GPUs (Table 1 in Li et al. (2024)). Thus, directly applying these methods to ST-DiTs suffers from degraded quality in practice when scaling up to a large number of GPUs. In contrast, ScaleFusion offer lossless algorithm for the distributed inference of ST-DiTs while achieving strong scalability on multiple GPU machines.

7 CONCLUSION

In this paper, we propose ScaleFusion to optimize communication overhead in ST-DiTs. We identify the spatialtemporal independence that has not been utilized in prior works. We propose both intra-layer and inter-layer communication scheduling algorithms to miminize the communication overhead through computation-communication overlap. Our evaluations show that ScaleFusion can achieve an average speedup of $1.36 \times$ (up to $1.58 \times$) speedup compared to prior works. We conclude that ScaleFusion can significantly reduce the communication overhead and scale distributed inference ST-DiTs to multiple GPU machines for high-resolution long video generation.

550 REFERENCES

551

552

553

554

555

- Aminabadi, R. Y., Rajbhandari, S., Zhang, M., Awan, A. A., Li, C., Li, D., Zheng, E., Rasley, J., Smith, S., Ruwase, O., and He, Y. Deepspeed inference: Enabling efficient inference of transformer models at unprecedented scale, 2022. URL https://arxiv.org/abs/2207.0 556 0032. 557
- 558 Arnab, A., Dehghani, M., Heigold, G., Sun, C., Lučić, M., 559 and Schmid, C. Vivit: A video vision transformer, 2021. 560 URL https://arxiv.org/abs/2103.15691. 561
- AWS. Elastic Fabric Adapter. URL https://aws.am 562 azon.com/hpc/efa/. 563
- 564 Brooks, T., Peebles, B., Holmes, C., DePue, W., Guo, Y., 565 Jing, L., Schnurr, D., Taylor, J., Luhman, T., Luhman, 566 E., Ng, C., Wang, R., and Ramesh, A. Video generation 567 models as world simulators. 2024. URL https://op 568 enai.com/research/video-generation-m 569 odels-as-world-simulators. 570
- 571 Dao, T. Flashattention-2: Faster attention with better par-572 allelism and work partitioning, 2023. URL https: 573 //arxiv.org/abs/2307.08691.
- 574 Dao, T., Fu, D. Y., Ermon, S., Rudra, A., and Ré, C. Flashat-575 tention: Fast and memory-efficient exact attention with 576 io-awareness, 2022. URL https://arxiv.org/ab 577 s/2205.14135. 578
- 579 Fang, J. and Zhao, S. Usp: A unified sequence parallelism 580 approach for long context generative ai, 2024. URL 581 https://arxiv.org/abs/2405.07719. 582
- Gupta, A., Yu, L., Sohn, K., Gu, X., Hahn, M., Fei-Fei, L., 583 Essa, I., Jiang, L., and Lezama, J. Photorealistic video 584 generation with diffusion models, 2023. URL https: 585 //arxiv.org/abs/2312.06662. 586
- 587 Harlap, A., Narayanan, D., Phanishayee, A., Seshadri, V., 588 Devanur, N., Ganger, G., and Gibbons, P. Pipedream: 589 Fast and efficient pipeline parallel dnn training, 2018. 590 URL https://arxiv.org/abs/1806.03377. 591
- 592 Ho, J., Salimans, T., Gritsenko, A., Chan, W., Norouzi, M., 593 and Fleet, D. J. Video diffusion models. In Koyejo, S., 594 Mohamed, S., Agarwal, A., Belgrave, D., Cho, K., and 595 Oh, A. (eds.), Advances in Neural Information Processing 596 Systems, volume 35, pp. 8633–8646. Curran Associates, 597 Inc., 2022. URL https://proceedings.neurip 598 s.cc/paper_files/paper/2022/file/392 599 35c56aef13fb05a6adc95eb9d8d66-Paper-C 600 onference.pdf. 601
- Huang, Y., Cheng, Y., Bapna, A., Firat, O., Chen, M. X., 602 Chen, D., Lee, H., Ngiam, J., Le, Q. V., Wu, Y., and Chen, 603 604

Z. Gpipe: Efficient training of giant neural networks using pipeline parallelism, 2019. URL https://arxi v.org/abs/1811.06965.

- Jacobs, S. A., Tanaka, M., Zhang, C., Zhang, M., Song, S. L., Rajbhandari, S., and He, Y. Deepspeed ulysses: System optimizations for enabling training of extreme long sequence transformer models, 2023. URL https: //arxiv.org/abs/2309.14509.
- Jayarajan, A., Wei, J., Gibson, G., Fedorova, A., and Pekhimenko, G. Priority-based parameter propagation for distributed dnn training. In Talwalkar, A., Smith, V., and Zaharia, M. (eds.), Proceedings of Machine Learning and Systems, volume 1, pp. 132-145, 2019. URL https://proceedings.mlsys.org/paper files/paper/2019/file/3ed923f9f88108 cb066c6568d3df2666-Paper.pdf.
- Jiang, C., Tian, Y., Jia, Z., Zheng, S., Wu, C., and Wang, Y. Lancet: Accelerating mixture-of-experts training via whole graph computation-communication overlapping, 2024. URL https://arxiv.org/abs/2404.1 9429.
- Li, M., Cai, T., Cao, J., Zhang, Q., Cai, H., Bai, J., Jia, Y., Li, K., and Han, S. Distrifusion: Distributed parallel inference for high-resolution diffusion models. In Proceedings of the IEEE/CVF Conference on Computer Vision and Pattern Recognition (CVPR), pp. 7183–7193, June 2024.
- Liu, H., Zaharia, M., and Abbeel, P. Ring attention with blockwise transformers for near-infinite context, 2023. URL https://arxiv.org/abs/2310.01889.
- Liu, X., Gong, C., and Liu, Q. Flow straight and fast: Learning to generate and transfer data with rectified flow, 2022. URL https://arxiv.org/abs/2209.0 3003.
- Lu, C., Zhou, Y., Bao, F., Chen, J., Li, C., and Zhu, J. Dpm-solver: A fast ode solver for diffusion probabilistic model sampling in around 10 steps, 2022. URL https: //arxiv.org/abs/2206.00927.
- Ma, X., Wang, Y., Jia, G., Chen, X., Liu, Z., Li, Y.-F., Chen, C., and Qiao, Y. Latte: Latent diffusion transformer for video generation, 2024. URL https://arxiv.org/ abs/2401.03048.
- NVIDIA. NVIDIA NSight Systems, a. URL https: //developer.nvidia.com/nsight-systems.
- NVIDIA. NVLink and NVLink Switch, b. URL https: //www.nvidia.com/en-us/data-center/nv link/.

- Rombach, R., Blattmann, A., Lorenz, D., Esser, P., and
 Ommer, B. High-resolution image synthesis with latent
 diffusion models, 2022. URL https://arxiv.org/
 abs/2112.10752.
- Singer, U., Polyak, A., Hayes, T., Yin, X., An, J., Zhang,
 S., Hu, Q., Yang, H., Ashual, O., Gafni, O., Parikh, D.,
 Gupta, S., and Taigman, Y. Make-a-video: Text-to-video
 generation without text-video data, 2022. URL https:
 //arxiv.org/abs/2209.14792.
- Song, J., Meng, C., and Ermon, S. Denoising diffusion implicit models, 2022. URL https://arxiv.org/ abs/2010.02502.
- Song, Y., Dhariwal, P., Chen, M., and Sutskever, I. Consistency models, 2023. URL https://arxiv.org/ab
 s/2303.01469.

- Wang, J., Fang, J., Li, A., and Yang, P. Pipefusion: Displaced patch pipeline parallelism for inference of diffusion transformer models, 2024. URL https: //arxiv.org/abs/2405.14430.
- Yang, Z., Teng, J., Zheng, W., Ding, M., Huang, S., Xu, J., Yang, Y., Hong, W., Zhang, X., Feng, G., et al. Cogvideox: Text-to-video diffusion models with an expert transformer. arXiv preprint arXiv:2408.06072, 2024.
- Zhao, X., Cheng, S., Chen, C., Zheng, Z., Liu, Z., Yang, Z., and You, Y. Dsp: Dynamic sequence parallelism for multi-dimensional transformers, 2024. URL https: //arxiv.org/abs/2403.10266.
- Zheng, Z., Peng, X., Yang, T., Shen, C., Li, S., Liu, H., Zhou, Y., Li, T., and You, Y. Open-sora: Democratizing efficient video production for all, March 2024. URL https: //github.com/hpcaitech/Open-Sora.

PSEUDO-CODE IMPLEMENTATION Α

Algorithm 1 ScaleFusion's implementation on each GPU.

Inputs: The input tensor x of shape [B, T, S/P, C], the video diffusion model \mathcal{M} , the number of temporal and spatial slices, N_T and N_S , and the number of partitions of the first temporal/spatial slice L_T and L_S . **Outputs:** The output tensor of shape [B, T, S/P, C]. 1: Create an empty tensor y of shape [B, T/P, S, C]2: Rearrange x to shape $[B, N_T, N_S, T/N_T, S/N_S/P, C]$ 3: Rearrange y to shape $[B, N_T, N_S, T/N_T/P, S/N_S, C]$ 4: for $i \in [0, N_S)$ do 5: for $j \in [0, L_T)$ do $E_{i,i}^S \leftarrow \text{ASYNCA2A}_{T \rightarrow S}(x_{:,i,i})$ 6: 7: for SpatialLayer, TemporalLayer in \mathcal{M} do for $i \in [L_T, N_T)$ do 8: for $j \in [L_S, N_S)$ do 9: $E_{i,j}^S \leftarrow \text{AsyncA2A}_{T \to S}(x_{:,i,j})$ 10: 11: for $i \in [0, N_T)$ do
$$\begin{split} tmp &\leftarrow \mathsf{cat}_S(\mathsf{Wait}(E^S_{i,0},\ldots,E^S_{i,S_N-1})) \\ y_{:,i,:} &\leftarrow \mathsf{SPATIALLAYER}(tmp) \end{split}$$
12: 13: for $j \in [0, L_S)$ do $E_{i,j}^T \leftarrow \operatorname{AsyncA2A}_{S \to T}(y_{:,i,j})$ 14: 15: 16: for $i \in [L_S, N_S)$ do for $j \in [L_T, N_T)$ do 17: $E_{i,i}^T \leftarrow \text{ASYNCA2A}_{S \to T}(x_{:,j,i})$ 18: for $i \in [0, N_S)$ do 19: $\begin{array}{l} tmp \leftarrow \mathbf{cat}_T(\mathbf{WAIT}(E_{0,i}^T, \dots E_{N_T-1,i}^T)) \\ x_{:,:,i} \leftarrow \mathbf{TEMPORALLAYER}(tmp) \end{array}$ 20: 21: if TemporalLayer is the last layer of \mathcal{M} then 22: 23: break for $j \in [0, L_T)$ do $E_{j,i}^S \leftarrow \operatorname{AsyncA2A}_{T \to S}(x_{:,j,i})$ 24: 694 25: 695 26: Rearrange x to shape [B, T, S/P, C]696 27: **return** *x* 697 698

Algorithm 1 describes the pseudo-code of ScaleFusion's im-699 plementation. To support asynchronous all-to-all operations, 700 we create one CUDA stream for computation operations and another CUDA stream for communication operations. For simplicity, we abstract out our implementation in a way where that calling an asynchronous all-to-all operation pro-704 705 duces a CUDA event by waiting on which we can block the default CUDA stream (i.e. the computation stream) and get 706 the return value of the all-to-all operation. Particularly, in the code each $E_{i,j}^S$ and each $E_{i,j}^T$ denotes a CUDA event 708 that represents the status of the (asynchronous) all-to-all 709 710 operation scheduled for the spatial or the temporal layer respectively. 711

712 Specifically, to keep the algorithm implementation and our 713 description clean, we initiate the first layer's communication 714

in line 4 to line 6 as if we had a previous temporal layer. For each spatial layer of the video diffusion model, we first initiate the communication that has not handled in the last temporal layer (line 8 to line 10). Then, we wait on all the events and concatenate the result tensors along the spatial dimension (line 12) and execute the spatial layer (line 13). We will then start conduct all-to-all operation for the next temporal layer on the output tensor y according to the inter-layer communication scheduling algorithm (line 14 to line 15). The execution of temporal layers would be similar except that if it is the last layer, we do not conduct communication for the next spatial layer (line 22).

Note that our implementation does not explicitly ensure the execution order of all-to-all operations. This is because by default each CUDA stream will execute each kernel in the kernel launch order and we only use one CUDA stream for all NCCL kernels (i.e. all-to-all operations). Thus, as long as we launch the communication operation in order, we do not have to synchronize between different communication operations to ensure the execution order.