

# Appendix

## A Missing Proofs in Section 3

**Proposition 1 ( $\beta$ -Pareto Efficiency, Upper Bound).** *There is no neutral rule  $f: \mathcal{L}(A)^n \rightarrow \mathcal{R}(A)$  that satisfies  $\epsilon$ -DP and  $\beta$ -Pareto efficiency with  $\beta > e^{\frac{n\epsilon}{m-1}}$ .*

*Proof.* Let  $f: \mathcal{L}(A)^n \rightarrow A$  be a voting rule satisfying  $\epsilon$ -DP and  $\beta$ -Pareto efficiency. Let  $P$  be a profile, where  $a_1 \succ a_2 \succ \dots \succ a_m$ , for all  $i \in N$ . Since  $a_1$  Pareto dominates  $a_2$ ,  $a_2$  Pareto dominates  $a_3$ , etc., we have

$$\begin{aligned} \mathbb{P}[f(P) = a_1] &\geq \beta \cdot \mathbb{P}[f(P) = a_2] \\ &\geq \beta^2 \cdot \mathbb{P}[f(P) = a_3] \\ &\dots \\ &\geq \beta^{m-1} \cdot \mathbb{P}[f(P) = a_m]. \end{aligned}$$

Then we claim that for profile  $P$ ,

$$\mathbb{P}[f(P) = a_1] \leq e^{n\epsilon} \cdot \mathbb{P}[f(P) = a_m].$$

Theorefore, we have  $\beta^{m-1} \leq e^{n\epsilon}$ , i.e.,  $\beta \leq e^{\frac{n\epsilon}{m-1}}$ , as desired.

Finally, we prove the claim above. In fact, for any voting rule  $f: \mathcal{L}(A)^n \rightarrow \mathcal{R}(A)$  satisfying  $\epsilon$ -DP and neutrality, we have

$$\mathbb{P}[f(P) = a] \leq e^{n\epsilon} \cdot \mathbb{P}[f(P) = b], \quad \text{for all } a, b \in A. \quad (1)$$

Now, we prove Equation (1). For any profile  $P, P' \in \mathcal{L}(A)^n$ , let  $\ell_0(\cdot, \cdot)$  represents the  $\ell_0$ -distance between them, i.e.,  $\ell_0(P, P') = \{j \in N : \succ_j \neq \succ'_j\}$ . Then, by considering the following operation **Op**, we can see that  $P$  can be transferred to  $P'$  through  $k = \ell_0(P, P')$  times of operations.

- **Op**: Choose a voter  $i \in N$  that  $\succ_i \neq \succ'_i$ , let  $\succ_i = \succ'_i$ .

Letting  $P_0, P_1, \dots, P_k$  denote all of the profiles, we have the following diagram.

$$P = P_0 \xrightarrow{\text{Op}} P_1 \xrightarrow{\text{Op}} P_2 \xrightarrow{\text{Op}} \dots \xrightarrow{\text{Op}} P_k = P'.$$

Notice that in each step, only one voter's preference is changed. Consequently, for each  $i$ ,  $P_i$  and  $P_{i+1}$  are neighboring profiles. Since  $f$  satisfies  $\epsilon$ -DP, we have

$$\mathbb{P}[f(P) = a] \leq e^\epsilon \cdot \mathbb{P}[f(P_1) = a] \leq e^{2\epsilon} \cdot \mathbb{P}[f(P_2) = a] \leq \dots \leq e^{k\epsilon} \cdot \mathbb{P}[f(P') = a].$$

Besides, for any given  $P, P' \in \mathcal{L}(A)^n$ , there are at most  $n$  distinct voters  $j \in N$  that  $\succ_j \neq \succ'_j$ . Therefore, for any profile  $P, P' \in \mathcal{L}(A)^n$  and any  $a \in A$ , we have

$$\mathbb{P}[f(P) = a] \leq e^{n\epsilon} \cdot \mathbb{P}[f(P') = a].$$

Now, for any profile  $P \in \mathcal{L}(A)^n$  and an arbitrarily chosen pair of alternatives  $a, b \in A$ , let  $P'$  be the profile transferred from  $P$  by swapping  $a$  and  $b$  in each voter's preference. By the neutrality of  $f$ , we have

$$\mathbb{P}[f(P) = a] = \mathbb{P}[f(P') = b] \text{ and } \mathbb{P}[f(P') = a] = \mathbb{P}[f(P) = b].$$

Then it follows that

$$\begin{aligned}\mathbb{P}[f(P) = a] &\leq e^{n\epsilon} \cdot \mathbb{P}[f(P') = a] \\ &= e^{n\epsilon} \cdot \mathbb{P}[f(P) = b],\end{aligned}$$

which completes the proof.  $\square$

**Proposition 2 ( $\beta$ -Pareto Efficiency, Lower Bound).** *Given  $\epsilon \in \mathbb{R}_+$ , Mechanism 1 satisfies  $\epsilon$ -DP and  $e^{\frac{n\epsilon}{2m-2}}$ -Pareto efficiency.*

*Proof.* Let  $\mathfrak{E}_{\text{Borda}}: \mathcal{L}(A)^* \rightarrow \mathcal{R}(A)$  denote the mapping introduced by BordaEXP. Then for any profile  $P \in \mathcal{L}(A)^*$  and alternative  $a \in A$ , we have

$$\mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = a] = \frac{e^{\text{Borda}_P(a)\epsilon/(2m-2)}}{\sum_{c \in A} e^{\text{Borda}_P(c)\epsilon/(2m-2)}}.$$

First, we establish the bound for Pareto efficiency. Given profile  $P \in \mathcal{L}(A)^n$ , suppose  $a, b \in A$  are a pair of alternatives that  $a \succ_j b$  for all  $j \in N$ . It follows that, for each voter  $j \in N$ , the number of alternatives that are considered worse than  $b$  according to her preference order  $\succ_j$  is strictly less than the number of alternatives considered worse than  $a$ . Formally, we have

$$|\{c \in A : a \succ_j c\}| - |\{c \in A : b \succ_j c\}| \geq 1, \quad \text{for all } j \in N.$$

By the definition of Borda score, we have

$$\text{Borda}_P(a) - \text{Borda}_P(b) \geq n.$$

Then it follows that

$$\begin{aligned}\mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = a] &= \frac{e^{\text{Borda}_P(a)\epsilon/(2m-2)}}{\sum_{c \in A} e^{\text{Borda}_P(c)\epsilon/(2m-2)}} \\ &\geq \frac{e^{\text{Borda}_P(b)\epsilon/(2m-2)} \cdot e^{n\epsilon/(2m-2)}}{\sum_{c \in A} e^{\text{Borda}_P(c)\epsilon/(2m-2)}} \\ &= e^{n\epsilon/(2m-2)} \cdot \mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = b],\end{aligned}$$

which indicates that  $\mathfrak{E}_{\text{Borda}}$  satisfies  $e^{\frac{n\epsilon}{2m-2}}$ -Pareto efficiency.

Then we prove the DP-bound. For all neighboring profiles  $P, P' \in \mathcal{L}(A)^n$ ,

$$\begin{aligned}\frac{\mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = a]}{\mathbb{P}[\mathfrak{E}_{\text{Borda}}(P') = a]} &= \frac{e^{\text{Borda}_P(a)\epsilon/(2m-2)}}{e^{\text{Borda}_{P'}(a)\epsilon/(2m-2)}} \cdot \frac{\sum_{c \in A} e^{\text{Borda}_{P'}(c)\epsilon/(2m-2)}}{\sum_{c \in A} e^{\text{Borda}_P(c)\epsilon/(2m-2)}} \\ &\leq e^{\epsilon/2} \cdot \sup_{P \in \mathcal{L}(A)^n} \frac{\sum_{c \in A} e^{\text{Borda}_{P'}(c)\epsilon/(2m-2)}}{\sum_{c \in A} e^{\text{Borda}_P(c)\epsilon/(2m-2)}} \\ &\leq e^{\epsilon},\end{aligned}$$

which completes the proof.  $\square$

**Lemma 1.** *Given  $\gamma > 0$ , a voting rule  $f$  satisfies  $\gamma$ -SD-efficiency if and only if*

$$\frac{1}{\gamma} \geq \sup_{P \in \mathcal{L}(A)^n, \xi \in \mathcal{R}(A)} \inf_{j \in N, y \in A} \frac{\sum_{x: x \succ_j y} \mathbb{P}[\xi = x]}{\sum_{x: x \succ_j y} \mathbb{P}[f(P) = x]}.$$

*Proof.* Suppose that  $f$  is not  $\gamma$ -SD-efficient. Then there must be some profile  $P \in \mathcal{L}(A)^n$  that  $f(P)$  is  $\gamma$ -SD-dominated by some  $\xi \in \mathcal{R}(A)$ , i.e.,

$$\sum_{x: x \succ_j y} \mathbb{P}[\xi = x] \geq \frac{1}{\gamma} \cdot \sum_{x: x \succ_j y} \mathbb{P}[f(P) = x], \quad \text{for all } y \in A \text{ and } \succ_j \in P,$$

which is equivalent to

$$\frac{1}{\gamma} \leq \inf_{j \in N, y \in A} \frac{\sum_{x: x \succ_j y} \mathbb{P}[\xi = x]}{\sum_{x: x \succ_j y} \mathbb{P}[f(P) = x]}.$$

Therefore,  $f$  is  $\gamma$ -SD-efficient if and only if for each  $P \in \mathcal{L}(A)^n$ , there does not exist such  $\xi$ , i.e.,

$$\frac{1}{\gamma} \geq \inf_{j \in N, y \in A} \frac{\sum_{x: x \succ_j y} \mathbb{P}[\xi = x]}{\sum_{x: x \succ_j y} \mathbb{P}[f(P) = x]}, \quad \text{for all } y \in A \text{ and } P \in \mathcal{L}(A)^n,$$

which is equivalent to

$$\frac{1}{\gamma} \geq \sup_{P \in \mathcal{L}(A)^n, \xi \in \mathcal{R}(A)} \inf_{j \in A, y \in A} \frac{\sum_{x: x \succ_j y} \mathbb{P}[\xi = x]}{\sum_{x: x \succ_j y} \mathbb{P}[f(P) = x]}.$$

That completes the proof.  $\square$

**Proposition 3 ( $\gamma$ -SD-Efficiency, Upper Bound).** *Given  $\gamma \in \mathbb{R}_+$ , there is no neutral voting rule  $f: \mathcal{L}(A)^n \rightarrow \mathcal{R}(A)$  satisfying  $\epsilon$ -DP and  $\gamma$ -SD-efficiency with  $\gamma > \frac{(m-1)e^{n\epsilon}}{(m-1)e^{n\epsilon}+1}$ .*

*Proof.* Consider two profiles,  $P_1$  and  $P_2$ , where all voters in  $P_1$  share the same preference order  $a_1 \succ a_2 \succ \dots \succ a_m$ . In contrast, in  $P_2$ , the voters' preferences are  $a_m \succ' a_2 \succ' a_3 \succ' \dots \succ' a_{m-1} \succ' a_1$ . Then the unique SD-efficient lottery for  $P_1$  and  $P_2$  should be  $\mathbb{1}_{a_1}$  and  $\mathbb{1}_{a_m}$ , respectively. Here,  $\mathbb{1}_{a_1}$  and  $\mathbb{1}_{a_m}$  represent indicator functions defined as follows.

$$\mathbb{P}[\mathbb{1}_{a_1} = a] = \begin{cases} 1 & a = a_1 \\ 0 & \text{otherwise} \end{cases}, \quad \mathbb{P}[\mathbb{1}_{a_m} = a] = \begin{cases} 1 & a = a_m \\ 0 & \text{otherwise} \end{cases}.$$

Let  $f$  be any neutral voting rule satisfying  $\epsilon$ -DP. By Equation (1), we have

$$\begin{aligned} \mathbb{P}[f(P_1) = a_1] &\leq e^{n\epsilon} \cdot \mathbb{P}[f(P_2) = a_1] && \text{(by } \epsilon\text{-DP)} \\ &= e^{n\epsilon} \cdot \mathbb{P}[f(P_1) = a_m]. && \text{(by neutrality)} \end{aligned}$$

By symmetry, for any  $a \neq a_m$ ,  $\mathbb{P}[f(P_1) = a] \leq e^{n\epsilon} \cdot \mathbb{P}[f(P_1) = a_m]$ . Therefore,

$$\sum_{a \in A} \mathbb{P}[f(P_1) = a] \leq ((m-1)e^{n\epsilon} + 1) \cdot \mathbb{P}[f(P_1) = a_m] \leq 1,$$

i.e.,  $\mathbb{P}[f(P_1) = a_m] \leq \frac{1}{(m-1)e^{n\epsilon} + 1}$ . If there exists some  $\gamma$  that  $f$  satisfies  $\gamma$ -SD-efficiency, there does not exist any  $\xi \in \mathcal{R}(A)$ , such that

$$\sum_{x:x \succ y} \mathbb{P}[\xi = x] \geq \frac{1}{\gamma} \cdot \sum_{x:x \succ y} \mathbb{P}[f(P_1) = x], \quad \text{for all } y \in A.$$

Therefore, we have

$$\begin{aligned} \frac{1}{\gamma} &\geq \sup_{P \in \mathcal{L}(A)^n} \sup_{\xi \in \mathcal{R}(A)} \inf_{y \in A} \frac{\sum_{x:x \succ y} \mathbb{P}[\xi = x]}{\sum_{x:x \succ y} \mathbb{P}[f(P) = x]} \\ &\geq \sup_{\xi \in \mathcal{R}(A)} \inf_{y \in A} \frac{\sum_{x:x \succ y} \mathbb{P}[\xi = x]}{\sum_{x:x \succ y} \mathbb{P}[f(P_1) = x]} \\ &\geq \inf_{y \in A} \frac{\sum_{x:x \succ y} \mathbb{P}[\mathbb{1}_{a_1} = x]}{\sum_{x:x \succ y} \mathbb{P}[f(P_1) = x]} \\ &= \frac{1}{\max_{y \in A} \sum_{x:x \succ y} \mathbb{P}[f(P_1) = x]} \\ &= \frac{1}{1 - \mathbb{P}[f(P_1) = a_n]} \\ &\geq \frac{(m-1)e^{n\epsilon} + 1}{(m-1)e^{n\epsilon}}. \end{aligned}$$

In other words, we have  $\gamma \leq \frac{(m-1)e^{n\epsilon}}{(m-1)e^{n\epsilon} + 1}$ , as desired.  $\square$

**Proposition 4 ( $\gamma$ -SD-Efficiency, Lower Bound).** *Mechanism 2 satisfies  $\epsilon$ -DP and  $\frac{(m-1)e^\epsilon}{(m-1)e^\epsilon + 1}$ -SD-efficiency.*

*Proof.* Letting  $\mathfrak{E}_{\text{Anti}}: \mathcal{L}(A)^n \rightarrow \mathcal{R}(A)$  denote the mapping introduced by Mechanism 2.

For any neighboring profiles  $P, P' \in \mathcal{L}(A)^n$  that  $P_{-j} = P'_{-j}$  and  $\succ_j \neq \succ'_j$ , suppose that the chosen ballot in the mechanism is  $\succ_i$ . Then

$$\mathbb{P}[\mathfrak{E}_{\text{Anti}}(P) = a \mid i \neq j] = \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P') = a \mid i \neq j] \quad (\text{use } C \text{ to denote them})$$

Further, for any  $a \in A$ ,

$$\begin{aligned} \frac{\mathbb{P}[\mathfrak{E}_{\text{Anti}}(P) = a]}{\mathbb{P}[\mathfrak{E}_{\text{Anti}}(P') = a]} &= \frac{\mathbb{P}[i = j \wedge \mathfrak{E}_{\text{Anti}}(P) = a] + \mathbb{P}[i \neq j \wedge \mathfrak{E}_{\text{Anti}}(P) = a]}{\mathbb{P}[i = j \wedge \mathfrak{E}_{\text{Anti}}(P') = a] + \mathbb{P}[i \neq j \wedge \mathfrak{E}_{\text{Anti}}(P') = a]} \\ &= \frac{\frac{1}{n} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P) = a \mid i = j] + \frac{n-1}{n} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P) = a \mid i \neq j]}{\frac{1}{n} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P') = a \mid i = j] + \frac{n-1}{n} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P') = a \mid i \neq j]} \\ &= \frac{\frac{1}{n} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P) = a \mid i = j] + \frac{n-1}{n} \cdot C}{\frac{1}{n} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P') = a \mid i = j] + \frac{n-1}{n} \cdot C} \quad (P_{-j} = P'_{-j}) \\ &\leq \frac{e^\epsilon \cdot \frac{1}{n} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P') = a \mid i = j] + \frac{n-1}{n} \cdot C}{\frac{1}{n} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P') = a \mid i = j] + \frac{n-1}{n} \cdot C} \\ &\leq e^\epsilon, \end{aligned}$$

which indicates that  $\mathfrak{E}_{\text{Anti}}$  satisfies  $\epsilon$ -DP. On the other hand, given profile  $P$ , suppose the top-ranked and the last-ranked alternative of  $\succsim_i$  are  $a_{\text{top}}$  and  $a_{\text{last}}$ , respectively. Then, for any  $\xi \in \mathcal{R}(A)$ , we have

$$\frac{\sum_{x:x \succsim_i y} \mathbb{P}[\xi = x]}{\sum_{x:x \succsim_i y} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P) = x]} \leq \frac{1}{\sum_{x:x \succsim_i y} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P) = x]} = \frac{\sum_{x:x \succsim_i y} \mathbb{P}[\mathbb{1}_{a_{\text{top}}} = x]}{\sum_{x:x \succsim_i y} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P) = x]}.$$

Theorefore,

$$\begin{aligned} \sup_{\xi \in \mathcal{R}(A)} \inf_{y \in A} \frac{\sum_{x:x \succsim_i y} \mathbb{P}[\xi = x]}{\sum_{x:x \succsim_i y} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P) = x]} &= \inf_{y \in A} \frac{1}{\sum_{x:x \succsim_i y} \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P) = x]} \\ &= \frac{1}{1 - \mathbb{P}[\mathfrak{E}_{\text{Anti}}(P) = a_{\text{last}}]} \\ &= \frac{(m-1)e^\epsilon}{(m-1)e^\epsilon + 1}. \end{aligned}$$

By Lemma 1,  $\mathfrak{E}_{\text{Anti}}$  satisfies  $\frac{(m-1)e^\epsilon}{(m-1)e^\epsilon + 1}$ -SD-efficiency, which completes the proof.  $\square$

**Lemma 2.** *Given  $\gamma \leq 1$ ,  $\gamma$ -PC-efficiency implies  $\gamma$ -SD-efficiency.*

*Proof.* In fact, we only need to prove that for any  $\xi, \zeta \in \mathcal{R}(A)$ ,  $\xi \succeq^{\gamma-SD} \zeta$  implies  $\xi \succeq^{\gamma-PC} \zeta$ . Let  $\xi$  and  $\zeta$  be two lotteries satisfying  $\xi \succeq^{\gamma-SD} \zeta$ , i.e.,

$$\sum_{x \succ y} \mathbb{P}[\xi = x] \geq \frac{1}{\gamma} \cdot \sum_{x \succ y} \mathbb{P}[\zeta = y], \quad \text{for any } y \in A.$$

Then, on the one hand, we have

$$\begin{aligned} \sum_{x,y \in A \wedge x \succ y} \mathbb{P}[\xi = x] \cdot \mathbb{P}[\zeta = y] &= \sum_{y \in A} \mathbb{P}[\zeta = y] \cdot \sum_{x \succ y} \mathbb{P}[\xi = x] \\ &\geq \sum_{y \in A} \mathbb{P}[\zeta = y] \cdot \frac{1}{\gamma} \sum_{x \succ y} \mathbb{P}[\zeta = y] \\ &= \frac{1}{\gamma} \cdot \sum_{x,y \in A \wedge x \succ y} \mathbb{P}[\zeta = y] \cdot \mathbb{P}[\zeta = y]. \end{aligned}$$

On the other hand, we have

$$\begin{aligned} \frac{1}{\gamma} \cdot \sum_{x,y \in A \wedge x \succ y} \mathbb{P}[\zeta = y] \mathbb{P}[\xi = y] &= \sum_{x \in A} \mathbb{P}[\zeta = y] \cdot \sum_{y \prec x} \mathbb{P}[\xi = y] \\ &= \frac{1}{\gamma} \cdot \sum_{x \in A} \mathbb{P}[\zeta = y] \cdot \sum_{y \prec x} \left( 1 - \sum_{x \succeq y} \mathbb{P}[\xi = x] \right) \\ &\leq \frac{1}{\gamma} \cdot \sum_{x \in A} \mathbb{P}[\zeta = y] \cdot \sum_{y \prec x} \left( 1 - \frac{1}{\gamma} \sum_{x \succeq y} \mathbb{P}[\zeta = y] \right) \\ &\leq \frac{1}{\gamma} \cdot \sum_{x \in A} \mathbb{P}[\zeta = y] \cdot \sum_{y \prec x} \left( 1 - \sum_{x \succeq y} \mathbb{P}[\zeta = y] \right) \\ &= \frac{1}{\gamma} \cdot \sum_{x,y \in A \wedge x \succ y} \mathbb{P}[\zeta = y] \cdot \mathbb{P}[\zeta = y]. \end{aligned}$$

Then it follows that

$$\sum_{x,y \in A \wedge x \succ y} \mathbb{P}[\xi = x] \cdot \mathbb{P}[\zeta = y] \geq \frac{1}{\gamma} \cdot \sum_{x,y \in A \wedge x \succ y} \mathbb{P}[\zeta = x] \cdot \mathbb{P}[\xi = y],$$

which completes the proof.  $\square$

**Proposition 5 ( $\kappa$ -PC-Efficiency, Upper Bound).** *Given any  $\kappa, \epsilon \in \mathbb{R}_+$ , there is no voting rule  $f: \mathcal{L}(A)^n \rightarrow \mathcal{R}(A)$  satisfying  $\epsilon$ -DP and  $\kappa$ -PC-efficiency.*

*Proof.* Consider the profile  $P$ , where all voters share the same preference

$$a_1 \succ a_2 \succ \cdots \succ a_m.$$

Then the unique PC-efficient distribution on  $A$  is  $\mathbb{1}_{a_1}$ . Further, we have

$$\sum_{x,y: x \succ_i y} \mathbb{P}[\mathbb{1}_{a_1} = x] \cdot \mathbb{P}[f(P) = y] = \sum_{y: a_1 \succ_i y} \mathbb{P}[f(P) = y] = 1 - \mathbb{P}[f(P) = a_1].$$

However,

$$\sum_{x,y: x \succ_i y} \mathbb{P}[f(P) = x] \cdot \mathbb{P}[\mathbb{1}_{a_1} = y] = \sum_{x \in A} \mathbb{P}[f(P) = x] \cdot \sum_{y: x \succ_i y} \mathbb{P}[\mathbb{1}_{a_1} = y] = 0.$$

In other words, for all  $\kappa \in \mathbb{R}_+$ , the lottery  $\mathbb{1}_{a_1}$  can  $\kappa$ -PC-dominate any  $f(P)$ , which completes the proof.  $\square$

**Proposition 6 ( $\alpha$ -Condorcet Criterion, Lower Bound).** *Mechanism 3 satisfies  $e^\epsilon$ -Condorcet criterion and  $\epsilon$ -DP.*

*Proof.* Let  $\mathfrak{R}_{\text{CW}}: \mathcal{L}(A)^* \rightarrow \mathcal{R}(A)$  denote the mapping introduced by CWRR. Then for any profile  $P \in \mathcal{L}(A)^*$  and alternative  $a \in A$ , we have

$$\mathbb{P}[\mathfrak{R}_{\text{CW}}(P) = a] = \begin{cases} \frac{e^\epsilon}{e^\epsilon + m - 1}, & a = \text{CW}(P) \\ \frac{1}{e^\epsilon + m - 1}, & \text{otherwise} \end{cases}, \quad \text{for all } P \text{ that } \text{CW}(P) \text{ exists.}$$

By definition, it is not hard to see that CWRR satisfies  $e^\epsilon$ -Condorcet criterion. Thus, we only need to prove that CWRR satisfies  $\epsilon$ -DP. In fact, for any neighboring profiles  $P, P' \in \mathcal{L}(A)^n$  and  $a \in A$ ,

$$\begin{aligned} \frac{\mathbb{P}[\mathfrak{R}_{\text{CW}}(P) = a]}{\mathbb{P}[\mathfrak{R}_{\text{CW}}(P') = a]} &\leq \frac{\max_{a \in A} \mathbb{P}[\mathfrak{R}_{\text{CW}}(P) = a]}{\max_{a \in A} \mathbb{P}[\mathfrak{R}_{\text{CW}}(P') = a]} \\ &\leq \frac{e^\epsilon}{e^\epsilon + m - 1} / \frac{1}{e^\epsilon + m - 1} \\ &= e^\epsilon, \end{aligned}$$

which completes the proof.  $\square$

**Proposition 7 ( $\eta$ -Condorcet Loser Criterion, Upper Bound).** *There is no voting rule satisfying  $\epsilon$ -DP and  $\eta$ -Condorcet loser criterion with  $\eta > e^\epsilon$ .*

*Proof.* Suppose  $f: \mathcal{L}(A)^n \rightarrow A$  be a voting rule satisfying  $\epsilon$ -DP and  $\eta$ -Condorcet loser criterion. Consider the profile  $P$  ( $n = 2k + 1$ ):

$$- k + 1 \text{ voters: } a_1 \succ a_2 \succ \cdots \succ a_m,$$

–  $k$  voters:  $a_m \succ a_{m-1} \succ \cdots \succ a_1$ .

By definition, we have  $w_P[a_m, a_i] = -1$ , for all  $a_i \in A \setminus \{a_m\}$ , i.e.,  $a_m$  is the Condorcet loser. Now, letting one voter change her preference from  $a_1 \succ a_2 \succ \cdots \succ a_m$  to  $a_m \succ a_{m-1} \succ \cdots \succ a_1$ , we can obtain another profile  $P'$ :

–  $k$  voters:  $a_1 \succ' a_2 \succ' \cdots \succ' a_m$ ,  
 –  $k+1$  voters:  $a_m \succ' a_{m-1} \succ' \cdots \succ' a_1$ .

Now we have  $w_{P'}[a_1, a_i] = -1$ , for all  $a_i \in A \setminus \{a_1\}$ , i.e.,  $a_1$  is the Condorcet loser for  $P'$ . Then

$$\begin{aligned} \mathbb{P}[f(P) = a_1] &\geq \eta \cdot \mathbb{P}[f(P) = a_m] && \text{(By } \eta\text{-Condorcet loser)} \\ &\geq e^{-\epsilon} \cdot \eta \cdot \mathbb{P}[f(P') = a_m] && \text{(By } \epsilon\text{-DP)} \\ &\geq e^{-\epsilon} \cdot \eta^2 \cdot \mathbb{P}[f(P') = a_1] && \text{(By } \eta\text{-Condorcet loser)} \\ &\geq e^{-2\epsilon} \cdot \eta^2 \cdot \mathbb{P}[f(P) = a_1], && (\epsilon\text{-DP}) \end{aligned}$$

which indicates that  $e^{-2\epsilon} \cdot \eta^2 \leq 1$ , i.e.,  $\eta \leq e^\epsilon$ . That completes the proof.  $\square$

**Proposition 8 ( $\eta$ -Condorcet Loser Criterion, Lower Bound).** *Mechanism 4 satisfies  $e^\epsilon$ -Condorcet loser criterion and  $\epsilon$ -DP.*

*Proof.* Let  $\mathfrak{R}_{\text{CL}}: \mathcal{L}(A)^* \rightarrow \mathcal{R}(A)$  denote the mapping introduced by CLRR, we have

$$\mathbb{P}[\mathfrak{R}_{\text{CL}}(P) = a] = \begin{cases} \frac{1}{(m-1)e^\epsilon + 1}, & a = \text{CL}(P) \\ \frac{e^\epsilon}{(m-1)e^\epsilon + 1}, & \text{otherwise} \end{cases}, \quad \text{for all } P \text{ that } \text{CL}(P) \text{ exists.}$$

By definition, it is not hard to see that CLRR satisfies  $e^\epsilon$ -Condorcet criterion. Thus, we only need to prove that CLRR satisfies  $\epsilon$ -DP. In fact, for any neighboring profiles  $P, P' \in \mathcal{L}(A)^n$  and  $a \in A$ ,

$$\begin{aligned} \frac{\mathbb{P}[\mathfrak{R}_{\text{CL}}(P) = a]}{\mathbb{P}[\mathfrak{R}_{\text{CL}}(P') = a]} &\leq \frac{\max_{a \in A} \mathbb{P}[\mathfrak{R}_{\text{CL}}(P) = a]}{\max_{a \in A} \mathbb{P}[\mathfrak{R}_{\text{CL}}(P') = a]} \\ &\leq \frac{e^\epsilon}{(m-1)e^\epsilon + 1} / \frac{1}{(m-1)e^\epsilon + 1} \\ &= e^\epsilon, \end{aligned}$$

which completes the proof.  $\square$

## B Missing Proofs in Section 4

### B.1 Results in Table 3 and their proofs

**Proposition 9.** *Given  $\epsilon \in \mathbb{R}_+$ , BordaEXP satisfies*

- (1)  $\frac{e^{\frac{n}{2} + (m-2) \cdot e^{\frac{n(m-2)}{4m-4}}}}{e^{\frac{n}{2} + (m-1) \cdot e^{\frac{n(m-2)}{4m-4}}}} - SD\text{-efficiency},$
- (2)  $e^{(\lfloor \frac{n}{2} \rfloor + 1) \cdot \frac{m}{2m-2} - \frac{n}{2}} - \text{Condorcet criterion},$
- (3)  $e^{\frac{n}{2m-2} - (\lceil \frac{n}{2} \rceil - 1) \frac{m}{2m-2}} - \text{Condorcet loser criterion}.$

*Proof.* Let  $\mathfrak{E}_{\text{Borda}}$  denote the voting rule introduced by BordaEXP. First, we prove (1). In fact,

$$\begin{aligned} \sup_{P, \xi} \inf_{j, y} \frac{\sum_{x \succ_j y} \mathbb{P}[\xi = x]}{\sum_{x \succ_j y} \mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = x]} &\leq \sup_P \inf_{j, y} \frac{1}{\sum_{x \succ_j y} \mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = x]} \\ &= \sup_P \inf_j \frac{1}{1 - \mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = a_{\text{last}}^j]} \\ &\leq \frac{1}{1 - \sup_P \inf_j \mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = a_{\text{last}}^j]}. \end{aligned}$$

where  $a_{\text{last}}^j$  denote the last-ranked alternative in  $\succ_j$ . By symmetry, we have

$$\sup_P \inf_j \mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = a_{\text{last}}^j] = \frac{e^{\frac{n(m-2)}{4m-4}}}{e^{\frac{n}{2} + (m-1) \cdot e^{\frac{n(m-2)}{4m-4}}}.$$

Then BordaEXP satisfies  $\frac{e^{\frac{n}{2} + (m-2) \cdot e^{\frac{n(m-2)}{4m-4}}}}{e^{\frac{n}{2} + (m-1) \cdot e^{\frac{n(m-2)}{4m-4}}}}$ . Second, we prove (2). By definition, for any profile  $P$  that  $\text{CW}(P)$  exists,  $\text{CW}(P)$  must defeat each alternative  $a \neq \{\text{CW}(P)\}$  in at least half of the votes, i.e.,  $\text{Borda}_P(\text{CW}(P)) \geq (m-1)(\lfloor \frac{n}{2} \rfloor + 1)$ . And for each  $a \neq \text{CW}(P)$ ,  $\text{Borda}_P(a) \leq (m-1)n - (\lfloor \frac{n}{2} \rfloor + 1)$ . Therefore,

$$\begin{aligned} \frac{\mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = \text{CW}(P)]}{\mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = a]} &\geq e^{\frac{(m-1)(\lfloor \frac{n}{2} \rfloor + 1)}{2m-2} - \frac{(m-1)n - (\lfloor \frac{n}{2} \rfloor + 1)}{2m-2}} \\ &= e^{(\lfloor \frac{n}{2} \rfloor + 1) \cdot \frac{m}{2m-2} - \frac{n}{2}}, \end{aligned}$$

which indicates that BordaEXP satisfies  $e^{(\lfloor \frac{n}{2} \rfloor + 1) \cdot \frac{m}{2m-2} - \frac{n}{2}}$ -Condorcet criterion. Finally, we prove (3). By definition, for any profile  $P$  that  $\text{CL}(P)$  exists,  $a \neq \text{CL}(P)$  must be ranked than  $\text{CL}(P)$  in at least a half of votes, i.e.,  $\text{Borda}_P(\text{CL}(P)) \leq (m-1)(\lfloor \frac{n}{2} \rfloor - 1)$ . And for each  $a \neq \text{CL}(P)$ ,  $\text{Borda}_P(a) \geq n - \lceil \frac{n}{2} \rceil + 1$ . Therefore,

$$\begin{aligned} \frac{\mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = a]}{\mathbb{P}[\mathfrak{E}_{\text{Borda}}(P) = \text{CL}(P)]} &\geq e^{\frac{n - \lceil \frac{n}{2} \rceil + 1}{2m-2} - \frac{(m-1)(\lfloor \frac{n}{2} \rfloor - 1)}{2m-2}} \\ &= e^{\frac{n}{2m-2} - (\lceil \frac{n}{2} \rceil - 1) \frac{m}{2m-2}}, \end{aligned}$$

which indicates that BordaEXP satisfies  $e^{\frac{n}{2m-2} - (\lceil \frac{n}{2} \rceil - 1) \frac{m}{2m-2}}$ -Condorcet loser criterion.  $\square$

**Proposition 10.** *Given  $\epsilon \in \mathbb{R}_+$ , RD-Anti satisfies*

- (1) 1-Pareto efficiency,  
 (2)  $\frac{(\lfloor \frac{n}{2} \rfloor - 1)e^\epsilon + \lceil \frac{n}{2} \rceil + 1}{ne^\epsilon}$ -Condorcet criterion,  
 (3)  $\frac{(\lfloor \frac{n}{2} \rfloor - 1)e^\epsilon + \lceil \frac{n}{2} \rceil + 1}{ne^\epsilon}$ -Condorcet loser criterion.

*Proof.* First, given profile  $P$ , for any  $a, b \in A$ ,  $a$  Pareto dominates  $b$  means that  $a \succ_j b$  for all  $j \in N$ . Then  $a$ , the Pareto dominator, is never ranked last in any  $\succ_j$ . Therefore,  $\mathbb{P}[\mathfrak{C}_{\text{Anti}}(P) = a] \geq \mathbb{P}[\mathfrak{C}_{\text{Anti}}(P) = b]$ , which completes the proof of (1). Second, we prove (2). For any profile  $P \in \mathcal{L}(A)^n$ ,

$$|\{j \in N : a_{last}^j = \text{CW}(P)\}| \leq \lceil \frac{n}{2} \rceil - 1,$$

otherwise  $\text{CW}(P)$  will be the Condorcet loser. Therefore,

$$\mathbb{P}[\mathfrak{C}_{\text{Anti}}(P) = \text{CW}(P)] \geq \frac{\lceil \frac{n}{2} \rceil - 1}{n} \cdot \frac{1}{(m-1)e^\epsilon + 1} + \frac{\lfloor \frac{n}{2} \rfloor + 1}{n} \cdot \frac{e^\epsilon}{(m-1)e^\epsilon + 1}.$$

For any  $a \neq \text{CW}(P)$ ,

$$\mathbb{P}[\mathfrak{C}_{\text{Anti}}(P) = a] \leq \frac{e^\epsilon}{(m-1)e^\epsilon + 1}.$$

Hence, we have

$$\frac{\mathbb{P}[\mathfrak{C}_{\text{Anti}}(P) = \text{CW}(P)]}{\mathbb{P}[\mathfrak{C}_{\text{Anti}}(P) = a]} \geq \frac{(\lfloor \frac{n}{2} \rfloor - 1)e^\epsilon + \lceil \frac{n}{2} \rceil + 1}{ne^\epsilon},$$

which completes the proof. Finally, we prove (3). Given a profile  $P$  that  $\text{CL}(P)$  exists,

$$|\{j \in N : a_{last}^j = a\}| \leq \lceil \frac{n}{2} \rceil - 1,$$

otherwise  $a$  will be the Condorcet loser. Therefore,

$$\mathbb{P}[\mathfrak{C}_{\text{Anti}}(P) = a] \geq \frac{\lceil \frac{n}{2} \rceil - 1}{n} \cdot \frac{1}{(m-1)e^\epsilon + 1} + \frac{\lfloor \frac{n}{2} \rfloor + 1}{n} \cdot \frac{e^\epsilon}{(m-1)e^\epsilon + 1}.$$

Hoewver,

$$\mathbb{P}[\mathfrak{C}_{\text{Anti}}(P) = \text{CL}(P)] \leq \frac{e^\epsilon}{(m-1)e^\epsilon + 1}.$$

Hence we have

$$\frac{\mathbb{P}[\mathfrak{C}_{\text{Anti}}(P) = a]}{\mathbb{P}[\mathfrak{C}_{\text{Anti}}(P) = \text{CL}(P)]} \geq \frac{(\lfloor \frac{n}{2} \rfloor - 1)e^\epsilon + \lceil \frac{n}{2} \rceil + 1}{ne^\epsilon},$$

which completes the proof of (3).  $\square$

**Proposition 11.** Given  $\epsilon \in \mathbb{R}_+$ ,  $\text{CWRR}$  satisfies 1-Pareto efficiency,  $\frac{m-1}{m}$ -SD-efficiency, and 1-Condorcet loser criterion.

*Proof.* The bounds of Pareto efficiency and Condorcet loser criterion are evident, since for any profile  $P$ , neither a Pareto dominated alternative nor the Condorcet loser can be the Condorcet winner. Then we only need to prove the bound of SD-efficiency. Given profile  $P$ , we have

$$\sup_{P, \xi} \inf_{j, y} \frac{\sum_{x \succ_j y} \mathbb{P}[\xi = x]}{\sum_{x \succ_j y} \mathbb{P}[\mathfrak{R}_{\text{CW}}(P) = x]} \leq \frac{1}{1 - \sup_P \inf_j \mathbb{P}[\mathfrak{R}_{\text{CW}}(P) = a_{last}^j]}.$$

Then there are two possible cases for the profile, discussed as follows

1. If  $\text{CW}(P)$  exists, then there must exist some  $j$  that  $a_{last}^j \neq \text{CW}(P)$ . Therefore

$$\inf_j \mathbb{P}[\mathfrak{R}_{\text{CW}}(P) = a_{last}^j] = \frac{1}{e^\epsilon + m - 1}.$$

2. If  $\text{CW}(P)$  does not exist, then

$$\inf_j \mathbb{P}[\mathfrak{R}_{\text{CW}}(P) = a_{last}^j] = \frac{1}{m} \geq \frac{1}{e^\epsilon + m - 1}.$$

In other words, we have

$$\sup_P \inf_j \mathbb{P}[\mathfrak{R}_{\text{CW}}(P) = a_{last}^j] = \frac{1}{m},$$

which indicates that CWRR satisfies  $\frac{m-1}{m}$ -SD-efficiency. That completes the proof.  $\square$

**Proposition 12.** *Given  $\epsilon \in \mathbb{R}_+$ , CLRR satisfies 1-Pareto efficiency,  $\frac{(m-2)e^\epsilon + 1}{(m-1)e^\epsilon + 1}$ -SD-efficiency, and 1-Condorcet criterion.*

*Proof.* The bounds of Pareto efficiency and Condorcet criterion are evident, since for any profile  $P$ , neither a Pareto dominator nor the Condorcet winner can be the Condorcet loser. Then we only need to prove the bound of SD-efficiency. Given profile  $P$ , we have

$$\sup_{P, \xi} \inf_{j, y} \frac{\sum_{x \succ_j y} \mathbb{P}[\xi = x]}{\sum_{x \succ_j y} \mathbb{P}[\mathfrak{R}_{\text{CL}}(P) = x]} \leq \frac{1}{1 - \sup_P \inf_j \mathbb{P}[\mathfrak{R}_{\text{CL}}(P) = a_{last}^j]}.$$

Then there are two possible cases for the profile, discussed as follows

1. If  $\text{CL}(P)$  exists, considering the profile  $P$ , where each  $a_{last}^j \neq \text{CL}(P)$  for each  $j \in N$ , we have

$$\inf_j \mathbb{P}[\mathfrak{R}_{\text{CL}}(P) = a_{last}^j] = \frac{e^\epsilon}{(m-1)e^\epsilon + 1}.$$

2. If  $\text{CL}(P)$  does not exist, then

$$\inf_j \mathbb{P}[\mathfrak{R}_{\text{CL}}(P) = a_{last}^j] = \frac{1}{m} \leq \frac{e^\epsilon}{(m-1)e^\epsilon + 1}.$$

In other words, we have

$$\sup_P \inf_j \mathbb{P}[\mathfrak{R}_{\text{CL}}(P) = a_{last}^j] = \frac{1}{m},$$

which indicates that CWRR satisfies  $\frac{(m-2)e^\epsilon + 1}{(m-1)e^\epsilon + 1}$ -SD-efficiency. That completes the proof.  $\square$

## B.2 Proofs of Theorems 1-6

**Theorem 1.** *There is no voting rule satisfying  $\epsilon$ -DP,  $\alpha$ -Condorcet criterion and  $\eta$ -Condorcet loser criterion with  $\alpha \cdot \eta > e^\epsilon$ .*

*Proof.* Consider the profile  $P$  ( $n = 2k + 1$ ):

$$- k + 1 \text{ voters: } a_1 \succ a_2 \succ \cdots \succ a_m,$$

–  $k$  voters:  $a_m \succ a_{m-1} \succ \dots \succ a_1$ .

By definition, we have  $\text{CW}(P) = a_1$ , since  $w_P[a_1, a_i] = 1$ , for all  $a_i \neq a_1$ . Now consider another profile  $P'$  with the same number of voters:

–  $k$  voters:  $a_1 \succ' a_2 \succ' \dots \succ' a_m$ ,

–  $k+1$  voters:  $a_m \succ a_{m-1} \succ \dots \succ a_1$ .

Then  $w_P[a_1, a_i] = -1$ , for all  $a_i \in A \setminus \{a\}$ , i.e.,  $a_1$  is a Condorcet loser. Since there is only one voter changes her preference from  $P$  to  $P'$ , we have

$$\begin{aligned}
\mathbb{P}[f(P) = a_1] &\geq \alpha \cdot \mathbb{P}[f(P) = a_2] && (\alpha\text{-Condorcet criterion}) \\
&\geq \alpha\eta \cdot \mathbb{P}[f(P) = a_m] && (\eta\text{-Condorcet loser criterion}) \\
&\geq e^{-\epsilon} \cdot \alpha\eta \cdot \mathbb{P}[f(P') = a_m] && (\epsilon\text{-DP}) \\
&\geq e^{-\epsilon} \cdot \alpha^2 \cdot \eta \cdot \mathbb{P}[f(P') = a_2] && (\alpha\text{-Condorcet criterion}) \\
&\geq e^{-\epsilon} \cdot \alpha^2 \cdot \eta^2 \cdot \mathbb{P}[f(P') = a_1] && (\eta\text{-Condorcet loser criterion}) \\
&\geq e^{-2\epsilon} \cdot \alpha^2 \cdot \eta^2 \cdot \mathbb{P}[f(P') = a_1], && (\epsilon\text{-DP})
\end{aligned}$$

which indicates that  $e^{-2\epsilon}\alpha^2\eta^2 \leq 1$ , i.e.,  $\alpha\eta \leq e^\epsilon$ . That completes the proof.  $\square$

**Theorem 2.** *If a neutral voting rule  $f: \mathcal{L}(A)^n \rightarrow A$  satisfies  $\epsilon$ -DP,  $\beta$ -Pareto efficiency, and  $\alpha$ -Condorcet criterion, then  $\alpha\beta^{m-2} \leq e^{n\epsilon}$ .*

*Proof.* Consider the following profile  $P$  ( $n = 2k + 1$ ):

–  $k+1$  voters:  $a_1 \succ a_2 \succ \dots \succ a_m$ ;

–  $k$  voters:  $a_2 \succ \dots \succ a_m \succ a_1$ .

By definition, we have  $w_P[a_1, a_i] = 1$ , for all  $a_i \in A \setminus \{a_1\}$ , which indicates that  $\text{CW}(P) = a_1$ . Also notice that  $w_P[a_i, a_j] = n$  for all  $i < j$ . Thus,  $a_i$  Pareto dominates  $a_j$  for all  $i < j$ . The relations among all alternatives are shown in the following graph.

$$a_1 \xrightarrow{\text{Condorcet Winner}} a_2 \xrightarrow{\text{Pareto}} a_3 \xrightarrow{\text{Pareto}} \dots \xrightarrow{\text{Pareto}} a_m. \quad (2)$$

Since  $f$  satisfies  $\alpha$ -Condorcet criterion and  $\beta$ -Pareto efficiency, we have

$$\begin{aligned}
\mathbb{P}[f(P) = a_1] &\geq \alpha \cdot \mathbb{P}[f(P) = a_2] \\
&\geq \alpha\beta \cdot \mathbb{P}[f(P) = a_3] \\
&\geq \dots \\
&\geq \alpha\beta^{m-2} \cdot \mathbb{P}[f(P) = a_m].
\end{aligned}$$

Now, consider another profile  $P'$ , where all voters' preferences are exactly the same:

$$a_m \succ a_{m-1} \succ \dots \succ a_1.$$

Then we have the following graph.

$$a_m \xrightarrow{\text{Condorcet Winner}} a_{m-1} \xrightarrow{\text{Pareto}} a_{m-2} \xrightarrow{\text{Pareto}} \dots \xrightarrow{\text{Pareto}} a_1.$$

Similarly, we have

$$\mathbb{P}[f(P') = a_m] \geq \alpha\beta^{m-2} \cdot \mathbb{P}[f(P') = a_1].$$

Notice that  $|\{j \in N : \succ_j \neq \succ'_j\}| = n$ . Therefore,

$$\begin{aligned} \mathbb{P}[f(P) = a_1] &\geq \alpha \beta^{m-2} \cdot \mathbb{P}[f(P) = a_m] \\ &\geq e^{-n\epsilon} \cdot \alpha \beta^{m-2} \cdot \mathbb{P}[f(P') = a_m] & (\epsilon\text{-DP}) \\ &\geq e^{-n\epsilon} \cdot \alpha^2 \beta^{2m-4} \cdot \mathbb{P}[f(P') = a_1] \\ &\geq e^{-2n\epsilon} \cdot \alpha^2 \beta^{2m-4} \cdot \mathbb{P}[f(P) = a_1]. & (\epsilon\text{-DP}) \end{aligned}$$

Then  $e^{-2n\epsilon} \alpha^2 \beta^{2m-4} \leq 1$ , i.e.,  $\alpha \beta^{m-2} \leq e^{n\epsilon}$ , which completes the proof.  $\square$

**Theorem 3.** *If a neutral voting rule  $f: \mathcal{L}(A)^n \rightarrow \mathcal{R}(A)$  satisfies  $\epsilon$ -DP,  $\beta$ -Pareto efficiency, and  $\alpha$ -Condorcet loser criterion, then  $\alpha \beta^{m-2} \leq e^{n\epsilon}$ .*

*Proof.* Consider the following profile  $P$  ( $n = 2k + 1$ ):

- $k + 1$  voters:  $a_1 \succ a_2 \succ \dots \succ a_m$ ;
- $k$  voters:  $a_2 \succ \dots \succ a_m \succ a_1$ .

By definition, we have  $w_P[a_1, a_i] = 1$ , for all  $a_i \in A \setminus \{a_1\}$ , which indicates that  $\text{CL}(P) = a_1$ . Also notice that  $w_P[a_i, a_j] = n$  for all  $i < j$ . Thus,  $a_i$  Pareto dominates  $a_j$  for all  $i < j$ . The relations among all alternatives are shown in the following graph.

$$a_1 \xrightarrow{\text{Pareto}} a_2 \xrightarrow{\text{Pareto}} \dots \xrightarrow{\text{Pareto}} a_{m-1} \xrightarrow{\text{Condorcet Loser}} a_m.$$

Since  $f$  satisfies  $\alpha$ -Condorcet loser criterion and  $\beta$ -Pareto efficiency, we have

$$\begin{aligned} \mathbb{P}[f(P) = a_1] &\geq \beta \cdot \mathbb{P}[f(P) = a_2] \\ &\geq \dots \\ &\geq \beta^{m-2} \cdot \mathbb{P}[f(P) = a_{m-1}] \\ &\geq \alpha \beta^{m-2} \cdot \mathbb{P}[f(P) = a_m]. \end{aligned}$$

Now, consider another profile  $P'$ , where all voters' preferences are exactly the same:

$$a_m \succ a_{m-1} \succ \dots \succ a_1.$$

Then we have the following graph.

$$a_m \xrightarrow{\text{Pareto}} a_{m-1} \xrightarrow{\text{Pareto}} \dots \xrightarrow{\text{Pareto}} a_2 \xrightarrow{\text{Condorcet Loser}} a_1.$$

Similarly, we have

$$\mathbb{P}[f(P') = a_m] \geq \alpha \beta^{m-2} \cdot \mathbb{P}[f(P') = a_1].$$

Notice that  $|\{j \in N : \succ_j \neq \succ'_j\}| = n$ . Therefore,

$$\begin{aligned} \mathbb{P}[f(P) = a_1] &\geq \alpha \beta^{m-2} \cdot \mathbb{P}[f(P) = a_m] \\ &\geq e^{-n\epsilon} \cdot \alpha \beta^{m-2} \cdot \mathbb{P}[f(P') = a_m] & (\epsilon\text{-DP}) \\ &\geq e^{-n\epsilon} \cdot \alpha^2 \beta^{2m-4} \cdot \mathbb{P}[f(P') = a_1] \\ &\geq e^{-2n\epsilon} \cdot \alpha^2 \beta^{2m-4} \cdot \mathbb{P}[f(P) = a_1]. & (\epsilon\text{-DP}) \end{aligned}$$

Then  $e^{-2n\epsilon} \alpha^2 \beta^{2m-4} \leq 1$ , i.e.,  $\alpha \beta^{m-2} \leq e^{n\epsilon}$ , which completes the proof.  $\square$

**Proposition 13.** *Condorcet method satisfies SD-efficiency on  $\mathcal{D}_C$ .*

*Proof.* Let  $P$  be an arbitrarily chosen profile in  $\mathcal{D}_C$ . Then we only need to proof that there does not exist  $\xi \in \mathcal{R}(A)$  that SD-dominates  $\mathbf{CM}(P)$ .

In fact, if there exists such a  $\xi$ , we can obtain by definition that

$$\sum_{b \succ_j a} \mathbb{P}[\xi = b] \geq \sum_{b \succ_j a} \mathbb{P}[\mathbf{CM}(P) = b], \quad \text{for all } j \in N \text{ and } a \in A.$$

Since for any  $a \in A$  that  $\mathbf{CW}(P) \succ_j a$ , we have

$$\sum_{b \succ_j a} \mathbb{P}[\mathbf{CM}(P) = b] = \mathbb{P}[\mathbf{CM}(P) = \mathbf{CW}(P)] = 1,$$

which indicates that

$$\sum_{b \succ_j a} \mathbb{P}[\xi = b] \geq 1.$$

Therefore, for any  $a \in A$  that  $\mathbf{CW}(P) \succ_j a$ ,  $\mathbb{P}[\xi = a] = 0$ . However, according to the definition of  $\mathbf{CW}(P)$ , each  $a \in A$  must be ranked behind  $\mathbf{CW}(P)$  in some  $\succ_j$ . Hence we have  $\xi = \mathbf{CM}(P)$ , a contradiction.  $\square$

**Theorem 4.** *There is no neutral voting rule  $f: \mathcal{L}(A)^n \rightarrow \mathcal{R}(A)$  satisfying  $\epsilon$ -DP,  $\alpha$ -Condorcet criterion, and  $\gamma$ -SD efficiency with  $\gamma > \frac{\alpha+m-1-\alpha e^{-n\epsilon}}{\alpha+m-1}$ .*

*Proof.* Consider the profile  $P$ , where all voters' vote are exactly the same, i.e.,

$$a_1 \succ_j a_2 \succ_j \cdots \succ_j a_m, \quad \text{for all } j \in N.$$

It is not hard to see that  $\mathbf{CW}(P) = a_1$ . Since  $f$  satisfies  $\alpha$ -Condorcet criterion, we have  $\mathbb{P}[f(P) = a] \leq \mathbb{P}[f(P) = a_1]/\alpha$ , for all  $a \in A \setminus \{a_1\}$ . Therefore,

$$1 = \mathbb{P}[f(P) = a_1] + \sum_{a \in A \setminus \{a_1\}} \mathbb{P}[f(P) = a] \leq \left(1 + \frac{m-1}{\alpha}\right) \mathbb{P}[f(P) = a_1],$$

i.e.,  $\mathbb{P}[f(P) = a_1] \geq \frac{\alpha}{\alpha+m-1}$ . Further, by Equation (1), we have

$$\mathbb{P}[f(P) = a_m] \geq e^{-n\epsilon} \cdot \mathbb{P}[f(P) = a_1] \geq \frac{\alpha e^{-n\epsilon}}{\alpha+m-1}.$$

However, for profile  $P$ , the unique SD-efficient lottery is  $\mathbb{1}_{a_1}$ . In other words, all lotteries  $\xi \in \mathcal{R}(A)$  that  $\xi \neq \mathbb{1}_{a_1}$  are  $\gamma$ -SD-dominated by  $\mathbb{1}_{a_1}$  with  $\gamma > 1$ . Further,

$$\begin{aligned} \frac{\sum_{x \succ y} \mathbb{P}[\mathbb{1}_{a_1} = x]}{\sum_{x \succ y} \mathbb{P}[f(P) = x]} &\geq \frac{\inf_{y \in A} \sum_{x \succ y} \mathbb{P}[\mathbb{1}_{a_1} = x]}{\sup_{y \in A} \sum_{x \succ y} \mathbb{P}[f(P) = x]} \\ &= \frac{1}{1 - \mathbb{P}[f(P) = a_m]} \\ &\geq \frac{1}{1 - \frac{\alpha e^{-n\epsilon}}{\alpha+m-1}} \\ &= \frac{\alpha+m-1-\alpha e^{-n\epsilon}}{\alpha+m-1}, \end{aligned}$$

i.e.,  $\mathbb{1}_{a_1}$  can  $\frac{\alpha+m-1-\alpha e^{-n\epsilon}}{\alpha+m-1}$ -dominates  $f(P)$ , which completes the proof.  $\square$

**Theorem 5.** *There is no neutral voting rule  $f: \mathcal{L}(A)^n \rightarrow \mathcal{R}(A)$  satisfying  $\epsilon$ -DP,  $\eta$ -Condorcet loser criterion, and  $\gamma$ -SD-efficiency with  $\gamma > \frac{e^{n\epsilon} - \eta}{e^{n\epsilon}}$ .*

*Proof.* Let  $m > 2$ . Consider the following profile  $P$  with  $k(m-2)$  voters ( $k \geq 1$ ).

- $k$  voters:  $y \succ a_1 \succ a_2 \succ \dots \succ x \succ a_{m-2}$ ,
- $k$  voters:  $y \succ a_2 \succ a_3 \succ \dots \succ x \succ a_1$ ,
- $k$  voters:  $y \succ a_3 \succ a_4 \succ \dots \succ x \succ a_2$ ,
- $k$  voters:  $\dots$ ,
- $k$  voters:  $y \succ a_{m-2} \succ a_1 \succ \dots \succ x \succ a_{m-3}$ .

Then it is quite evident that  $\text{CL}(P) = x$ . Since  $f$  satisfies  $\eta$ -Condorcet loser criterion, we have  $\mathbb{P}[f(P) = a] \geq \eta \cdot \mathbb{P}[f(P) = x]$ , for all  $a \in A \setminus \{x\}$ . By Equation (1), we have  $\mathbb{P}[f(P) = x] \geq e^{-n\epsilon}$ . Since  $f$  satisfies  $\eta$ -Condorcet loser criterion,

$$\mathbb{P}[f(P) = a] \geq \eta \cdot \mathbb{P}[f(P) = x] \geq \alpha e^{-n\epsilon}, \quad \text{for all } a \in A \setminus \{x\}.$$

However, the unique SD-efficient lottery of  $P$  is  $\mathbb{1}_y$ , since  $\mathbb{1}_y$  can SD-dominate any other lotteries on  $A$ . Further,

$$\begin{aligned} \frac{\sum_{b \succ a} \mathbb{P}[\mathbb{1}_y = b]}{\sum_{b \succ a} \mathbb{P}[f(P) = b]} &\geq \frac{\inf_{a \in A} \sum_{b \succ a} \mathbb{P}[\mathbb{1}_y = b]}{\sup_{a \in A} \sum_{b \succ a} \mathbb{P}[f(P) = b]} \\ &= \frac{1}{1 - \inf_{1 \leq i \leq m-2} \min \mathbb{P}[f(P) = a_i]} \\ &\geq \frac{1}{1 - \alpha e^{-n\epsilon}}. \end{aligned}$$

In other words,  $\mathbb{1}_y$  can  $\frac{e^{n\epsilon} - \eta}{e^{n\epsilon}}$ -SD-dominate  $f(P)$ , which completes the proof.  $\square$

**Theorem 6.** *There is no neutral voting rule  $f: \mathcal{L}(A)^n \rightarrow \mathcal{R}(A)$  satisfying  $\epsilon$ -DP,  $\gamma$ -SD-efficiency, and  $\beta$ -Pareto efficiency with  $\gamma > \frac{e^{n\epsilon} - e^{n\epsilon} \beta^{2-m}}{e^{n\epsilon} - e^{n\epsilon} \beta^{2-m} + \beta - 1}$ .*

*Proof.* Consider the profile  $P$ , where all voters' preferences are the same, i.e.,

$$a_1 \succ_j a_2 \succ_j \dots \succ_j a_m, \quad \text{for all } j \in N.$$

By definition, for all  $i < j$ , any  $a_i$  Pareto dominates  $a_j$  in profile  $P$ . In other words, we have the following diagram

$$a_1 \xrightarrow{\text{Pareto}} a_2 \xrightarrow{\text{Pareto}} \dots \xrightarrow{\text{Pareto}} a_m.$$

Since  $f$  satisfies  $\beta$ -Pareto efficiency,  $\mathbb{P}[f(P) = a_{i+1}] \leq \beta \cdot \mathbb{P}[f(P) = a_i]$  holds for any  $i < m$ . By Equation (1),  $\mathbb{P}[f(P) = a_m] \geq e^{-n\epsilon} \cdot \mathbb{P}[f(P) = a_1]$ . Further,

$$\begin{aligned} \mathbb{P}[f(P) = a_1] &\leq e^{n\epsilon} \cdot \mathbb{P}[f(P) = a_m], \\ \mathbb{P}[f(P) = a_2] &\leq \frac{1}{\beta} \cdot \mathbb{P}[f(P) = a_1] \leq \frac{e^{n\epsilon}}{\beta} \cdot \mathbb{P}[f(P) = a_m], \\ &\dots \\ \mathbb{P}[f(P) = a_{m-1}] &\leq \frac{1}{\beta^{m-2}} \cdot \mathbb{P}[f(P) = a_1] \leq \frac{e^{n\epsilon}}{\beta^{m-2}} \cdot \mathbb{P}[f(P) = a_m]. \end{aligned}$$

By summing up the above inequalities, we have

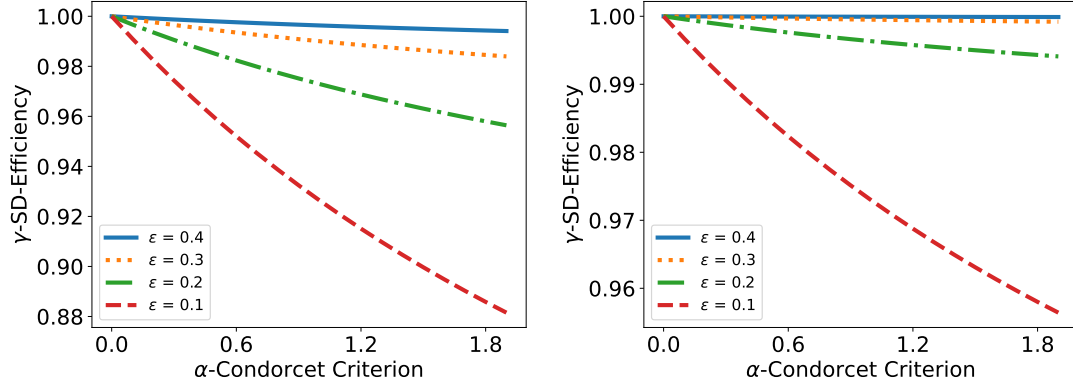
$$1 = \sum_{a \in A} \mathbb{P}[f(P) = a] \leq \left(1 + \left(1 + \frac{1}{\beta} + \cdots + \frac{1}{\beta^{m-2}}\right) e^{n\epsilon}\right) \cdot \mathbb{P}[f(P) = a_m],$$

i.e.,  $\mathbb{P}[f(P) = a_m] \geq \frac{\beta-1}{e^{n\epsilon} - e^{n\epsilon}\beta^{2-m} + \beta - 1}$ . However, the unique SD-efficient lottery of  $P$  is  $\mathbb{1}_{a_1}$ , since it can SD-dominate any other lottery. Further, we have

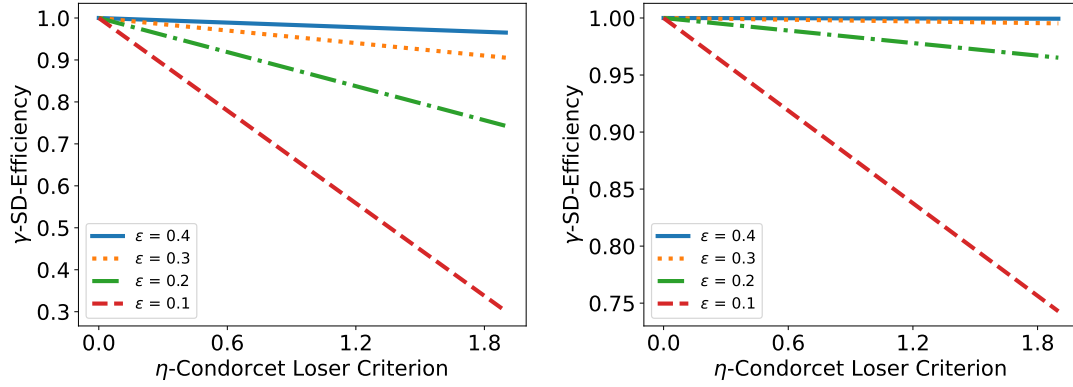
$$\begin{aligned} \frac{\sum_{b \succ a} \mathbb{P}[\mathbb{1}_y = b]}{\sum_{b \succ a} \mathbb{P}[f(P) = b]} &\geq \frac{\inf_{a \in A} \sum_{b \succ a} \mathbb{P}[\mathbb{1}_y = b]}{\sup_{a \in A} \sum_{b \succ a} \mathbb{P}[f(P) = b]} \\ &= \frac{1}{1 - \mathbb{P}[f(P) = a_m]} \\ &\geq \frac{1}{1 - \frac{\beta-1}{e^{n\epsilon} - e^{n\epsilon}\beta^{2-m} + \beta - 1}}. \end{aligned}$$

In other words,  $\mathbb{1}_{a_1}$  can  $\frac{e^{n\epsilon} - e^{n\epsilon}\beta^{2-m}}{e^{n\epsilon} - e^{n\epsilon}\beta^{2-m} + \beta - 1}$ -SD-dominates  $f(P)$ , which completes the proof.  $\square$

## C More Experimental Figures



**Fig. 1.** Tradeoff curves between  $\alpha$ -Condorcet criterion and  $\gamma$ -SD-efficiency under  $\epsilon$ -DP (upper bounds). Left:  $m = 5, n = 10$ . Right:  $m = 5, n = 20$ .



**Fig. 2.** Tradeoff curves between  $\eta$ -Condorcet loser criterion and  $\gamma$ -SD-efficiency under  $\epsilon$ -DP (upper bounds). Left:  $m = 5, n = 10$ . Right:  $m = 5, n = 20$ .