
Adversarially Robust Learning with Tolerance

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Abstract

1 We initiate the study of tolerant adversarial PAC learning with respect to metric
2 perturbation sets. In adversarial PAC learning, an adversary is allowed to replace a
3 test point x with an arbitrary point in a closed ball of radius r centered at x . In the
4 tolerant version, the error of the learner is compared with the best achievable error
5 with respect to a slightly larger perturbation radius $(1 + \gamma)r$. This simple tweak
6 helps us bridge the gap between theory and practice and obtain the first PAC-type
7 guarantees for algorithmic techniques that are popular in practice. Furthermore, our
8 sample complexity bounds improve exponentially over best known (non-tolerant)
9 bounds in terms of the VC dimension of the hypothesis class. In particular, for
10 perturbation sets with doubling dimension d , we show that a variant of the “perturb-
11 and-smooth” algorithm PAC learns any hypothesis class \mathcal{H} with VC dimension v in
12 the γ -tolerant adversarial setting with $O\left(\frac{v(1+\gamma)^{O(d)}}{\epsilon}\right)$ samples. This guarantee
13 holds in the tolerant robust realizable setting. We extend this to the agnostic case
14 by designing a novel sample compression scheme based on the perturb-and-smooth
15 approach. This compression-based algorithm has a linear dependence on the
16 doubling dimension as well as the VC-dimension.

17 1 Introduction

18 Several empirical studies (Szegedy et al., 2014; Goodfellow et al., 2018) have demonstrated that
19 models trained to have a low accuracy on a data set often have the undesirable property that a small
20 perturbation to an input instance can change the label outputted by the model. For most domains this
21 does not align with human intuition and thus indicates that the learned models are not representing
22 the ground truth despite obtaining good accuracy on test sets.

23 The theory of PAC-learning characterizes the conditions under which learning is possible. For binary
24 classification, the following conditions are sufficient: a) unseen data should arrive from the same
25 distribution as training data, and b) the class of models should have a low capacity (as measured, for
26 example, by its VC dimension). If these conditions are met, an *Empirical Risk Minimizer* (ERM) that
27 simply optimizes model parameters to maximize accuracy on the training set learns successfully.

28 Recent work has studied test-time adversarial perturbations under the PAC-learning framework. If
29 an adversary is allowed to perturb data during test time then the conditions above do not hold, and
30 we cannot hope for the model to learn to be robust just by running ERM. Thus, the goal here is to
31 bias the learning process towards finding models where label-changing perturbations are rare. This is
32 achieved by defining a loss function that combines both classification error and the probability of
33 seeing label-changing perturbations, and learning models that minimize this loss on unseen data. It
34 has been shown that even though (robust) ERM can fail in this setting, PAC learning is still possible
35 as long as we know during training the kinds of perturbations we want to guard against at test
36 time (Montasser et al., 2019). This result holds for all perturbation sets. However, the learning

37 algorithm is significantly more complex than robust ERM and requires a large number of samples
38 (with the best known sample complexity bounds potentially being exponential in the VC-dimension).

39 We study a *tolerant* version of the adversarially robust learning framework and restrict the perturba-
40 tions to balls in a general metric space with finite doubling dimension. We show this slight shift in
41 the learning objective yields significantly improved sample complexity bounds through a simpler
42 learning paradigm than what was previously known. In fact, we show that a version of the common
43 “perturb-and-smooth” paradigm successfully PAC-learns any class of bounded VC dimension in this
44 setting.

45 **Learning in general metric spaces.** What kinds of perturbations should a learning algorithm guard
46 against? Any transformation of the input that we believe should not change its label could be a viable
47 perturbation for the adversary to use. The early works in this area considered perturbations contained
48 within a small ℓ_p -ball of the input. More recent work has considered other transformations such as
49 a small rotation, or translation of an input image (Engstrom et al., 2019; Fawzi & Frossard, 2015;
50 Kanbak et al., 2018; Xiao et al., 2018), or even adding small amounts of fog or snow (Kang et al.,
51 2019). It has also been argued that small perturbations in some *feature space* should be allowed as
52 opposed to the input space (Inkawhich et al., 2019; Sabour et al., 2016; Xu et al., 2020; Song et al.,
53 2018; Hosseini & Poovendran, 2018). This motivates the study of more general perturbations.

54 We consider a setting where the input comes from a domain that is equipped with a distance metric
55 and allows perturbations to be within a small metric ball around the input. Earlier work on general
56 perturbation sets (for example, (Montasser et al., 2019)) considered arbitrary perturbations. In this
57 setting one does not quantify the magnitude of a perturbation and thus cannot talk about small versus
58 large perturbations. Modeling perturbations using a metric space enables us to do that while also
59 keeping the setup general enough to be able to encode a large variety of perturbation sets by choosing
60 appropriate distance functions.

61 **Learning with tolerance.** In practice, we often believe that small perturbations of the input should
62 not change its label but we do not know *precisely* what small means. However, in the PAC-learning
63 framework for adversarially robust classification, we are required to define a precise perturbation set
64 and learn a model that has error arbitrarily close to the smallest error that can be achieved with respect
65 to that perturbation set. In other words, we aim to be arbitrarily close to a target that was picked
66 somewhat arbitrarily to begin with. Due to the uncertainty about the correct perturbation size, it is
67 more meaningful to allow for a wider range of error values. To achieve this, we introduce the concept
68 of tolerance. In the tolerant setting, in addition to specifying a perturbation size r , we introduce a
69 tolerance parameter γ that encodes our uncertainty about the size of allowed perturbations. Then, for
70 any given $\epsilon > 0$, we aim to learn a model whose error with respect to perturbations of size r is at
71 most ϵ more than the smallest error achievable with respect to perturbations of size $r(1 + \gamma)$.

72 2 Our results

73 In this paper we formalize and initiate the study of the problem of adversarially robust learning in
74 the tolerant setting for general metric spaces and provide two algorithms for the task. Both of our
75 algorithms rely on: 1) modifying the training data by randomly sampling points from the perturbation
76 sets around each data point, and 2) smoothing the output of the model by taking a majority over the
77 labels returned by the model for nearby points.

78 Our first algorithm starts by modifying the training set by randomly perturbing each training point
79 using a certain distribution (see Section 5 for details). It then trains a (non-robust) PAC learner (such
80 as ERM) on the perturbed training set to find a hypothesis h . Finally, it outputs a smooth version
81 of h . The smoothing step replaces $h(x)$ at each point x with the a majority label outputted by h on
82 the points around x . We show that for metric spaces of a fixed doubling dimension, this algorithm
83 successfully learns in the (robustly realizable) tolerant setting.

84 **Theorem 1** (Informal version of Theorem 10). *Let (X, dist) be a metric space with doubling*
85 *dimension d and \mathcal{H} a hypothesis class. Assuming robust realizability, \mathcal{H} can be learned tolerantly in*
86 *the adversarially robust setting using $O\left(\frac{(1+1/\gamma)^{O(d)} \text{VC}(\mathcal{H})}{\epsilon}\right)$ samples, where γ encodes the amount*
87 *of allowed tolerance, and ϵ is the desired accuracy.*

88 An interesting feature of the above result is the linear dependence of the sample complexity with
 89 respect to $VC(\mathcal{H})$. This is in contrast to the best known upper bound for non-tolerant adversarial
 90 setting (Montasser et al., 2019) which depends on the *dual VC dimension* of the hypothesis class
 91 and in general is exponential in $VC(\mathcal{H})$. Moreover, this is the first PAC type guarantee for the
 92 general perturb-and-smooth paradigm, indicating that the tolerant adversarial learning is the “right”
 93 learning model for studying these approaches. While the above method enjoys simplicity and can
 94 be computationally efficient, one downside is that its sample complexity grows exponentially with
 95 the doubling dimension. For instance, such algorithm cannot be used on high-dimensional data in
 96 the Euclidean space. Another limitation is that the guarantee holds only in the (robustly) realizable
 97 setting. We propose another algorithm that improves the dependence on doubling dimension, and
 98 works in the general agnostic setting.

99 **Theorem 2** (Informal version of Corollary 16). *Let (X, dist) be a metric space with doubling*
 100 *dimension d and \mathcal{H} a hypothesis class. Then \mathcal{H} can be learned tolerantly in the adversarially robust*
 101 *setting using $\tilde{O}\left(\frac{O(d)VC(\mathcal{H})\log(1+1/\gamma)}{\epsilon^2}\right)$ samples, where \tilde{O} hides logarithmic factors, γ encodes the*
 102 *amount of allowed tolerance, and ϵ is the desired accuracy.*

103 This algorithm exploits the connection between sample compression and adversarially robust learn-
 104 ing Montasser et al. (2019). However, unlike Montasser et al. (2019), our new compression scheme
 105 sidesteps the dependence on the dual VC dimension. As a result, we get an exponential improvement
 106 over the best known (nontolerant) sample complexity in terms of dependence on VC dimension.

107 3 Related work

108 PAC-learning for adversarially robust classification has been studied extensively in recent years (Cul-
 109 lina et al., 2018; Awasthi et al., 2019; Montasser et al., 2019; Feige et al., 2015; Attias et al., 2019;
 110 Montasser et al., 2020a; Ashtiani et al., 2020). These works provide learning algorithms that guaran-
 111 tee low generalization error in the presence of adversarial perturbations in various settings. The most
 112 general result is due to (Montasser et al., 2019) which is proved for general hypothesis classes and
 113 perturbation sets. All of the above results assume that the learner knows the kinds of perturbations
 114 allowed for the adversary. Some more recent papers have considered scenarios where the learner
 115 does not even need to know that. Goldwasser et al. (2020) allow the adversary to perturb test data in
 116 unrestricted ways and are still able to provide learning guarantees. The catch is that it only works
 117 in the transductive setting and only if the learner is allowed to abstain from making a prediction on
 118 some test points. Montasser et al. (2021a) consider the case where the learner needs to infer the set of
 119 allowed perturbations by observing the actions of the adversary.

120 Tolerance was introduced by Ashtiani et al. (2020) but in the context of certification. They provide
 121 examples where certification is not possible unless we allow some tolerance. Montasser et al. (2021b)
 122 study transductive adversarial learning and provide a “tolerant” guarantee. Note that unlike our work,
 123 the main focus of this paper is on the transductive setting. Moreover, they do not specifically study
 124 tolerance with respect to metric perturbation sets. Without a metric, it is not meaningful to expand
 125 perturbation sets by a factor $(1 + \gamma)$ (as we do in the our definition of tolerance). Instead, they expand
 126 their perturbation sets by applying two perturbations in succession, which is akin to setting $\gamma = 1$. In
 127 contrast, our results hold in the more common inductive setting, and capture a more realistic setting
 128 where γ is any (small) real number larger than zero.

129 Like many recent adversarially robust learning algorithms (Feige et al., 2015; Attias et al., 2019),
 130 our first algorithm relies on calls to a non-robust PAC-learner. Montasser et al. (2020b) formalize the
 131 question of reducing adversarially robust learning to non-robust learning and study finite perturbation
 132 sets of size k . They show a reduction that makes $O(\log^2 k)$ calls to the non-robust learner and also
 133 prove a lower bound of $\Omega(\log k)$. It will be interesting to see if our algorithms can be used to obtain
 134 better bounds for the tolerant setting. Our first algorithm makes one call to the non-robust PAC-learner
 135 at training time, but needs to perform potentially expensive smoothing for making actual predictions
 136 (see Theorem 10).

137 The techniques of randomly perturbing the training data and smoothing the output classifier has been
 138 extensively used in practice and has shown good empirical success. Augmenting the training data with
 139 some randomly perturbed samples was used for handwriting recognition as early as in (Yaeger et al.,
 140 1996). More recently, “stability training” was introduced in (Zheng et al., 2016) for state of the art

141 image classifiers where training data is augmented with Gaussian perturbations. Empirical evidence
 142 was provided that the technique improved the accuracy against naturally occurring perturbations.
 143 Augmentations with non-Gaussian perturbations of a large variety were considered in (Hendrycks
 144 et al., 2019).

145 Smoothing the output classifier using random samples around the test point is a popular technique
 146 for producing *certifiably* robust classifiers. A certification, in this context, is a guarantee that given a
 147 test point x , all points within a certain radius of x receive the same label as x . Several papers have
 148 provided theoretical analyses to show that smoothing produces certifiably robust classifiers (Cao &
 149 Gong, 2017; Cohen et al., 2019; Lecuyer et al., 2019; Li et al., 2019; Liu et al., 2018; Salman et al.,
 150 2019; Levine & Feizi, 2020).

151 However, to the best of our knowledge, a PAC-like guarantee has not been shown for any algorithm
 152 that employs training data perturbations or output classifier smoothing, and our paper provides the
 153 first such analysis.

154 4 Notations and setup

155 We denote by X the input domain and by $Y = \{0, 1\}$ the binary label space. We assume that
 156 X is equipped with a metric dist . A hypothesis $h : X \rightarrow Y$ is a function that assigns a label
 157 to each point in the domain. A hypothesis class \mathcal{H} is a set of such hypotheses. For a sample
 158 $S = ((x_1, y_1), \dots, (x_n, y_n)) \in (X \times Y)^n$, we use the notation $S^X = \{x_1, x_2, \dots, x_n\}$ to denote
 159 the collection of domain points x_i occurring in S . The binary (also called 0-1) loss of h on data point
 160 $(x, y) \in X \times Y$ is defined by

$$\ell^{0/1}(h, x, y) = \mathbb{1}[h(x) \neq y],$$

161 where $\mathbb{1}[\cdot]$ is the indicator function. Let P be a probability distribution over $X \times Y$. Then the
 162 *expected binary loss* of h with respect to P is defined by

$$\mathcal{L}_P^{0/1}(h) = \mathbb{E}_{(x,y) \sim P}[\ell^{0/1}(h, x, y)]$$

163 Similarly, the *empirical binary loss* of h on sample $S = ((x_1, y_1), \dots, (x_n, y_n)) \in (X \times Y)^n$ is
 164 defined as $\mathcal{L}_S^{0/1}(h) = \frac{1}{n} \sum_{i=1}^n \ell^{0/1}(h, x_i, y_i)$. We also define the *approximation error* of \mathcal{H} with
 165 respect to P as $\mathcal{L}_P^{0/1}(\mathcal{H}) = \inf_{h \in \mathcal{H}} \mathcal{L}_P^{0/1}(h)$.

166 A *learner* \mathcal{A} is a function that takes in a finite sequence of labeled instances $S =$
 167 $((x_1, y_1), \dots, (x_n, y_n))$ and outputs a hypothesis $h = \mathcal{A}(S)$. The following definition abstracts
 168 the notion of PAC learning Vapnik & Chervonenkis (1971); Valiant (1984).

169 **Definition 3** (PAC Learner). *Let \mathcal{P} be a set of distributions over $X \times Y$ and \mathcal{H} a hypothesis class.*
 170 *We say \mathcal{A} PAC learns $(\mathcal{H}, \mathcal{P})$ with $m_{\mathcal{A}} : (0, 1)^2 \rightarrow \mathbb{N}$ samples if the following holds: for every*
 171 *distribution $P \in \mathcal{P}$ over $X \times Y$, and every $\epsilon, \delta \in (0, 1)$, if S is an i.i.d. sample of size at least*
 172 *$m_{\mathcal{A}}(\epsilon, \delta)$ from P , then with probability at least $1 - \delta$ (over the randomness of S) we have*

$$\mathcal{L}_P(\mathcal{A}(S)) \leq \mathcal{L}_P(\mathcal{H}) + \epsilon.$$

173 \mathcal{A} is called an *agnostic learner* if \mathcal{P} is the set of all distributions on $X \times Y$, and a *realizable learner* if
 174 $\mathcal{P} = \{P : \mathcal{L}_P(\mathcal{H}) = 0\}$.

175 The smallest function $m : (0, 1)^2 \rightarrow \mathbb{N}$ for which there exists a learner \mathcal{A} that satisfies the above
 176 definition with $m_{\mathcal{A}} = m$ is referred to as the (realizable or agnostic) *sample complexity* of learning
 177 \mathcal{H} .

178 The existence of sample-efficient PAC learners for VC classes is a standard result Vapnik & Chervo-
 179 nenkis (1971). We state the results formally in Appendix A.

180 4.1 Tolerant adversarial PAC learning

181 Let $\mathcal{U} : X \rightarrow 2^X$ be a function that maps each point in the domain to the set of its “admissible”
 182 perturbations. We call this function the *perturbation type*. The adversarial loss of h with respect to \mathcal{U}
 183 on $(x, y) \in X \times Y$ is defined by

$$\ell^{\mathcal{U}}(h, x, y) = \max_{z \in \mathcal{U}(x)} \{\ell^{0/1}(h, z, y)\}$$

184 The *expected adversarial loss* with respect to P is defined by $\mathcal{L}_P^{\mathcal{U}}(h) = \mathbb{E}_{(x,y) \sim P} \ell^{\mathcal{U}}(h, x, y)$. The
185 *empirical adversarial loss* of h on sample $S = ((x_1, y_1), \dots, (x_n, y_n)) \in (X \times Y)^n$ is defined by
186 $\mathcal{L}_S^{\mathcal{U}}(h) = \frac{1}{n} \sum_{i=1}^n \ell^{\mathcal{U}}(h, x_i, y_i)$. Finally, the *adversarial approximation error* of \mathcal{H} with respect to \mathcal{U}
187 and P is defined by $\mathcal{L}_P^{\mathcal{U}}(\mathcal{H}) = \inf_{h \in \mathcal{H}} \mathcal{L}_P^{\mathcal{U}}(h)$.

188 The following definition generalizes the setting of PAC adversarial learning to what we call the
189 *tolerant setting*, where we consider two perturbation types \mathcal{U} and \mathcal{V} . We say \mathcal{U} is *contained in* \mathcal{V} and
190 and write it as $\mathcal{U} \prec \mathcal{V}$ if $\mathcal{U}(x) \subset \mathcal{V}(x)$ for all $x \in X$.

191 **Definition 4** (Tolerant Adversarial PAC Learner). *Let \mathcal{P} be a set of distributions over $X \times Y$, \mathcal{H} a
192 hypothesis class, and $\mathcal{U} \prec \mathcal{V}$ two perturbation types. We say \mathcal{A} tolerantly PAC learns $(\mathcal{H}, \mathcal{P}, \mathcal{U}, \mathcal{V})$
193 with $m_{\mathcal{A}} : (0, 1)^2 \rightarrow \mathbb{N}$ samples if the following holds: for every distribution $P \in \mathcal{P}$ and every
194 $\epsilon, \delta \in (0, 1)$, if S is an i.i.d. sample of size at least $m_{\mathcal{A}}(\epsilon, \delta)$ from P , then with probability at least
195 $1 - \delta$ (over the randomness of S) we have*

$$\mathcal{L}_P^{\mathcal{U}}(\mathcal{A}(S)) \leq \mathcal{L}_P^{\mathcal{V}}(\mathcal{H}) + \epsilon.$$

196 We say \mathcal{A} is a *tolerant PAC learner* in the agnostic setting if \mathcal{P} is the set of all distributions over
197 $X \times Y$, and in the tolerantly realizable setting if $\mathcal{P} = \{P : \mathcal{L}_P^{\mathcal{V}}(\mathcal{H}) = 0\}$.

198 In the above context, we refer to \mathcal{U} as the *actual perturbation type* and to \mathcal{V} as the *reference*
199 *perturbation type*. The case where $\mathcal{U}(x) = \mathcal{V}(x)$ for all $x \in X$ corresponds to the usual adversarial
200 learning scenario (with no tolerance).

201 4.2 Tolerant adversarial PAC learning in metric spaces

202 If X is equipped with a metric $\text{dist}(\cdot, \cdot)$, then $\mathcal{U}(x)$ can be naturally defined by a ball of radius r
203 around x , i.e., $\mathcal{U}(x) = \mathcal{B}_r(x) = \{z \in X \mid \text{dist}(x, z) \leq r\}$. To simplify the notation, we sometimes
204 use $\ell^r(h, x, y)$ instead of $\ell^{\mathcal{B}_r}(h, x, y)$ to denote the adversarial loss with respect to \mathcal{B}_r .

205 In the tolerant setting, we consider the perturbation sets $\mathcal{U}(x) = \mathcal{B}_r(x)$ and $\mathcal{V}(x) = \mathcal{B}_{(1+\gamma)r}(x)$,
206 where $\gamma > 0$ is called the *tolerance parameter*. Note that $\mathcal{U} \prec \mathcal{V}$. We now define PAC learning with
207 respect to the metric space.

208 **Definition 5** (Tolerant Adversarial Learning in metric spaces). *Let (X, dist) be a metric space, \mathcal{H}
209 a hypothesis class, and \mathcal{P} a set of distributions of $X \times Y$. We say $(\mathcal{H}, \mathcal{P}, \text{dist})$ is *tolerantly PAC*
210 *learnable* with $m : (0, 1)^3 \rightarrow \mathbb{N}$ samples when for every $r, \gamma > 0$ there exist a PAC learner $\mathcal{A}_{r, \gamma}$ for
211 $(\mathcal{H}, \mathcal{P}, \mathcal{B}_r, \mathcal{B}_{(1+\gamma)r})$ that uses $m(\epsilon, \delta, \gamma)$ samples.*

212 **Remark 6.** *In this definition the learner receives γ and r as input but its sample complexity does
213 not depend on r (but can depend on γ). Also, as in Definition 4, the tolerantly realizable setting
214 corresponds to $\mathcal{P} = \{P : \mathcal{L}_P^{r(1+\gamma)}(\mathcal{H}) = 0\}$ while in the agnostic setting \mathcal{P} is the set of all
215 distributions over $X \times Y$.*

216 The doubling dimension and the doubling measure of the metric space will play important roles in
217 our analysis. We refer the reader to Appendix B for their definitions.

218 We will use the following lemma in our analysis, whose proof can be found in Appendix B:

219 **Lemma 7.** *For any family \mathcal{M} of complete, doubling metric spaces, there exist constants $c_1, c_2 > 0$
220 such that for any metric space $(X, \text{dist}) \in \mathcal{M}$ with doubling dimension d , there exists a measure μ
221 such that if a ball \mathcal{B}_r of radius $r > 0$ is completely contained inside a ball $\mathcal{B}_{\alpha r}$ of radius αr (with
222 potentially a different center) for any $\alpha > 1$, then $0 < \mu(\mathcal{B}_{\alpha r}) \leq (c_1 \alpha)^{c_2 d} \mu(\mathcal{B}_r)$. Furthermore, if
223 we have a constant $\alpha_0 > 1$ such that we know that $\alpha \geq \alpha_0$ then the bound can be simplified to
224 $0 < \mu(\mathcal{B}_{\alpha r}) \leq \alpha^{\zeta d} \mu(\mathcal{B}_r)$, where ζ depends on \mathcal{M} and α_0 .*

225 Later, we will set $\alpha = 1 + 1/\gamma$ where γ is the tolerance parameter. Since we are mostly interested in
226 small values of γ , suppose we decide on some loose upper bound $\Gamma \gg \gamma$. This corresponds to saying
227 that there exists some $\alpha_0 > 1$ such that $\alpha \geq \alpha_0$.

228 It is worth noting that in the special case of Euclidean metric spaces, we can set both c_1 and c_2 to be
229 1. In the rest of the paper, we will assume we have a loose upper bound $\Gamma \gg \gamma$ and use the simpler
230 bound from Lemma 24 extensively.

231 Given a metric space (X, d) and a measure μ defined over it, for any subset $Z \subseteq X$ for which $\mu(Z)$
232 is non-zero and finite, μ induces a *probability* measure P_Z^μ over Z as follows. For any set $Z' \subseteq Z$ in

233 the σ -algebra over Z , we define $P_Z^\mu(Z') = \mu(Z')/\mu(Z)$. With a slight abuse of notation, we write
 234 $z \sim Z$ to mean $z \sim P_Z^\mu$ whenever we know μ from the context.

235 Our learners rely on being able to sample from P_Z^μ . Thus we define the following oracle, which can
 236 be implemented efficiently for ℓ_p spaces.

237 **Definition 8** (Sampling Oracle). *Given a metric space (X, dist) equipped with a doubling measure
 238 μ , a sampling oracle is an algorithm that when queried with a $Z \subseteq X$ such that $\mu(Z)$ is finite, returns
 239 a sample drawn from P_Z^μ . We will use the notation $z \sim Z$ for queries to this oracle.*

240 5 The perturb-and-smooth approach for tolerant adversarial learning

241 In this section we focus on tolerant adversarial PAC learning in metric spaces (Definition 5), and
 242 show that VC classes are tolerantly PAC learnable in the tolerantly realizable setting. Interestingly,
 243 we prove this result using an approach that resembles the ‘‘perturb-and-smooth’’ paradigm which is
 244 used in practice (for example, (Cohen et al., 2019)). The overall idea is to ‘‘perturb’’ each training
 245 point x , train a classifier on the ‘‘perturbed’’ points, and ‘‘smooth out’’ the final hypothesis using a
 246 certain majority rule.

247 For this, we employ three perturbation types: \mathcal{U} and \mathcal{V} play the role of the *actual* and the *reference*
 248 perturbation type respectively. Additionally, we consider a perturbation type $\mathcal{W} : X \rightarrow 2^X$, which is
 249 used for smoothing. We assume $\mathcal{U} \prec \mathcal{V}$ and $\mathcal{W} \prec \mathcal{V}$. For this section, we will use metric balls for the
 250 three types. Specifically, if \mathcal{U} consists of balls of radius r for some $r > 0$, then \mathcal{W} will consist of
 251 balls of radius γr and \mathcal{V} will consist of balls of radius $(1 + \gamma)r$.

252 **Definition 9** (Smoothed classifier). *For a hypothesis $h : X \rightarrow \{0, 1\}$, we let $\bar{h}_{\mathcal{W}}$ denote the classifier
 253 resulting from replacing the label $h(x)$ with the average label over $\mathcal{W}(x)$, that is*

$$\bar{h}_{\mathcal{W}}(x) = \mathbb{1} \left[\mathbb{E}_{x' \sim \mathcal{W}(x)} h(x') \geq 1/2 \right]$$

254 *For metric perturbation types, where \mathcal{W} is a ball of some radius r , we also use the notation \bar{h}_r and
 255 when the type \mathcal{W} is clear from context, we may omit the subscript altogether and simply write \bar{h} for
 256 the smoothed classifier.*

257 **The tolerant perturb-and-smooth algorithm** We propose the following learning algorithm, TPaS,
 258 for tolerant learning in metric spaces. Let the perturbation radius be $r > 0$ for the actual type $\mathcal{U} = \mathcal{B}_r$,
 259 and let $S = ((x_1, y_1), \dots, (x_m, y_m))$ be the training sample. For each $x_i \in S^X$, the learner samples
 260 a point $x'_i \sim \mathcal{B}_{r \cdot (1+\gamma)}(x_i)$ (using the sampling oracle) from the expanded reference perturbation
 261 set $\mathcal{V}(x_i) = \mathcal{B}_{(1+\gamma)r}(x_i)$. Let $S' = ((x'_1, y_1), \dots, (x'_m, y_m))$. TPaS then invokes a (standard, non-
 262 robust) PAC learner $\mathcal{A}_{\mathcal{H}}$ for the hypothesis class \mathcal{H} on the perturbed data S' . We let $\hat{h} = \mathcal{A}_{\mathcal{H}}(S')$
 263 denote the output of this PAC learner. Finally, TPaS outputs the \mathcal{W} -smoothed version of $\bar{h}_{\gamma r}$ for
 264 $\mathcal{W} = \mathcal{B}_{\gamma r}$. That is, $\bar{h}_{\gamma r}(x)$ is simply the majority label in a ball of radius γr around x with respect to
 265 the distribution defined by μ , see also Definition 9. We will prove below that this $\bar{h}_{\gamma r}$ has a small
 266 \mathcal{U} -adversarial loss. Algorithm 1 below summarizes our learning procedure.

Algorithm 1 Tolerant Perturb and Smooth (TPaS)

Input: Radius r , tolerance parameter γ , data $S = ((x_1, y_1), \dots, (x_m, y_m))$, access to sampling
 oracle \mathcal{O} for μ and PAC learner $\mathcal{A}_{\mathcal{H}}$.

Initialize $S' = \emptyset$

for $i = 1$ to m **do**

Sample $x'_i \sim \mathcal{B}_{(1+\gamma)r}(x_i)$

Add (x'_i, y_i) to S'

end for

Set $\hat{h} = \mathcal{A}_{\mathcal{H}}(S')$

Output: $\bar{h}_{\gamma r}$ defined by

$$\bar{h}_{\gamma r}(x) = \mathbb{1} \left[\mathbb{E}_{x' \sim \mathcal{B}_{\gamma r}(x)} \hat{h}(x') \geq 1/2 \right]$$

267 The following is the main result of this section.

268 **Theorem 10.** Let (X, dist) be an any metric space with doubling dimension d and doubling measure
 269 μ . Let \mathcal{O} be a sampling oracle for μ . Let \mathcal{H} be a hypothesis class and \mathcal{P} a set of distributions over
 270 $X \times Y$. Assume $\mathcal{A}_{\mathcal{H}}$ PAC learns \mathcal{H} with $m_{\mathcal{H}}(\epsilon, \delta)$ samples in the realizable setting. Then there exists
 271 a learner \mathcal{A} , namely TPAS, that

- 272 • Tolerantly PAC learns $(\mathcal{H}, \mathcal{P}, \text{dist})$ in the tolerantly realizable setting with sample complexity
- 273 bounded by $m(\epsilon, \delta, \gamma) = O(m_{\mathcal{H}}(\epsilon, \delta) \cdot (1 + 1/\gamma)^{\zeta d}) = O\left(\frac{\text{VC}(\mathcal{H}) + \log 1/\delta}{\epsilon} \cdot (1 + 1/\gamma)^{\zeta d}\right)$,
- 274 where γ is the tolerance parameter and d is the doubling dimension.
- 275 • Makes only one query to $\mathcal{A}_{\mathcal{H}}$
- 276 • Makes $m(\epsilon, \delta, \gamma)$ queries to sampling oracle \mathcal{O}

277 The proof of this theorem uses the following key technical lemma (proof can be found in Appendix C):
 278

Lemma 11. Let $r > 0$ be a perturbation radius, $\gamma > 0$ a tolerance parameter, and $g : X \rightarrow Y$ a classifier. For $x \in X$ and $y \in Y = \{0, 1\}$, we define

$$\Sigma_{g,y}(x) = \mathbb{E}_{z \sim \mathcal{B}_{r(1+\gamma)}(x)} \mathbb{1}[g(z) \neq y] \quad \text{and} \quad \sigma_{g,y}(x) = \mathbb{E}_{z \sim \mathcal{B}_{r\gamma}(x)} \mathbb{1}[g(z) \neq y].$$

279 Then $\Sigma_{g,y}(x) \leq \frac{1}{3} \cdot \left(\frac{1+\gamma}{\gamma}\right)^{-\zeta d}$ implies that $\sigma_{g,y}(z) \leq 1/3$ for all $z \in \mathcal{B}_r(x)$.

280 *Proof of Theorem 10.* Consider some $\epsilon_0 > 0$ and $0 < \delta < 1$ to be given (we will pick a suitable
 281 value of ϵ_0 later), and assume the PAC learner $\mathcal{A}_{\mathcal{H}}$ was invoked on the perturbed sample S' of size
 282 at least $m_A(\epsilon_0, \delta)$. According to definition 3, this implies that with probability $1 - \delta$, the output
 283 $\hat{h} = \mathcal{A}_{\mathcal{H}}(S)$ has (binary) loss at most ϵ_0 with respect to the data-generating distribution. Note that the
 284 relevant distribution here is the two-stage process of the original data generating distribution P and
 285 the perturbation sampling according to $\mathcal{V} = \mathcal{B}_{(1+\gamma)r}$. Since P is \mathcal{V} -robustly realizable, the two-stage
 286 process yields a realizable distribution with respect to the standard 0/1-loss. Thus, we have

$$\mathbb{E}_{(x,y) \sim P} \mathbb{E}_{z \sim \mathcal{B}_{r(1+\gamma)}(x)} \mathbb{1}[\hat{h}(z) \neq y] \leq \epsilon_0.$$

287 With Lemma 11, this becomes $\mathbb{E}_{(x,y) \sim P} \Sigma_{\hat{h},y}(x) \leq \epsilon_0$. For $\lambda > 0$, Markov's inequality then yields :

$$\mathbb{E}_{(x,y) \sim P} \mathbb{1}[\Sigma_{\hat{h},y}(x) \leq \lambda] > 1 - \epsilon_0/\lambda \quad (1)$$

Thus setting $\lambda = \frac{1}{3} \cdot \left(\frac{1+\gamma}{\gamma}\right)^{-\zeta d}$ and plugging in the result of the Lemma 11 to equation (1), we get

$$\mathbb{E}_{(x,y) \sim P} \mathbb{1}[\forall z \in \mathcal{B}_r(x), \sigma_{\hat{h},y}(z) \leq 1/3] > 1 - \epsilon_0/\lambda.$$

288 Since $\sigma_{\hat{h},y}(z) \leq 1/3$ implies that $\mathbb{1}[\mathbb{E}_{z' \sim \mathcal{B}_{\gamma r}(z)} \hat{h}(z') \geq 1/2] = y$, using the definition of the
 289 smoothed classifier $\bar{h}_{\gamma r}$ we get

$$\mathbb{E}_{(x,y) \sim P} \mathbb{1}[\exists z \in \mathcal{B}_r(x), \bar{h}_{\gamma r}(z) \neq y] \leq \epsilon_0/\lambda, \quad (2)$$

290 which implies $\mathcal{L}_P^r(\bar{h}_{\gamma r}) \leq \epsilon_0/\lambda$. Thus, for the robust learning problem, if we are given a desired
 291 accuracy ϵ and we want $\mathcal{L}_P^r(\bar{h}_{\gamma r}) \leq \epsilon$, we can pick $\epsilon_0 = \lambda\epsilon$. Putting it all together, we get
 292 sample complexity $m \leq O\left(\frac{\text{VC}(\mathcal{H}) + \log 1/\delta}{\epsilon_0}\right)$ where $\epsilon_0 = \lambda\epsilon$, and $\lambda = \frac{1}{3} \cdot \left(\frac{1+\gamma}{\gamma}\right)^{-\zeta d}$. Therefore,
 293 $m \leq O\left(\frac{\text{VC}(\mathcal{H}) + \log 1/\delta}{\epsilon} \cdot (1 + 1/\gamma)^{\zeta d}\right)$. \square

294 Since the dependence on d is exponential, the algorithm becomes impractical for high dimensions if γ
 295 is very small. However, since γ represents our uncertainty in the value of the true perturbation radius,
 296 it is natural to assume that it is a small but positive number. We can therefore ask whether there
 297 exists a threshold for each dimension such that if γ is above the threshold we can learn efficiently. In

298 particular, for any constant $c > 0$, we can ensure that $(1 + 1/\gamma)^{\zeta^d} \leq 1 + c$ if we set $\gamma \geq \frac{\zeta^d}{c}$. Thus the
 299 sample complexity of our learner does not depend on the dimension as long as $\gamma \geq \frac{\zeta^d}{c}$. For example,
 300 for a Euclidean space with the ℓ_∞ metric, we have $\zeta = 1$. Therefore setting $c = 1000$ would let us
 301 use a small γ for dimensions up to 1000.

302 **Computational complexity of the learner.** Assuming we have access to \mathcal{O} and an efficient algorithm
 303 for non-robust PAC-learning in the realizable setting, we can compute \hat{h} efficiently. Therefore, the
 304 learning can be done efficiently in this case. However, at the prediction time, we need to compute
 305 $\bar{h}(x)$ on new test points which requires us to compute an expectation. We can instead *estimate* the
 306 expectations using random samples from the sampling oracle. For a single test point x , if the number
 307 of samples we draw is $\Omega(\log 1/\delta)$ then with probability at least $1 - \delta$ we get the same result as that
 308 of the optimal $\bar{h}(x)$. Using more samples we can boost this probability to guarantee a similar output
 309 to that of \bar{h} on a larger set of test points.

310 6 Improved tolerant learning guarantees through sample compression

311 The perturb-and-smooth approach discussed in the previous section offers a general method for
 312 tolerant robust learning. However, one shortcoming of this approach is the exponential dependence
 313 of its sample complexity with respect to the doubling dimension of the metric space. Furthermore,
 314 the tolerant robust guarantee relied on the data generating distribution being tolerantly realizable.
 315 In this section, we propose another approach that addresses both of these issues. The idea is to
 316 adopt the perturb-and-smooth approach within a sample compression argument. We introduce the
 317 notion of a $(\mathcal{U}, \mathcal{V})$ -tolerant sample compression scheme and present a learning bound based on such
 318 a compression scheme, starting with the realizable case. We then show that this implies learnability
 319 in the agnostic case as well. Remarkably, this tolerant compression based analysis will yield bounds
 320 on the sample complexity that avoid the exponential dependence on the doubling dimension.

321 For a compact representation, we will use the general notation \mathcal{U} , \mathcal{V} , and \mathcal{W} for the three perturbation
 322 types (actual, reference and smoothing type) in this section and will assume that they satisfy the
 323 Property 1 below for some parameter $\beta > 0$. Lemma 11 implies that, in the metric setting, for any
 324 radius r and tolerance parameter γ the perturbation types $\mathcal{U} = \mathcal{B}_r$, $\mathcal{V} = \mathcal{B}_{(1+\gamma)r}$, and $\mathcal{W} = \mathcal{B}_{\gamma r}$ have

325 this property for $\beta = \frac{1}{3} \left(\frac{1+\gamma}{\gamma} \right)^{-\zeta^d}$.

326 **Property 1.** For a fixed $0 < \beta < 1/2$, we assume that the perturbation types \mathcal{V}, \mathcal{U} and \mathcal{W} are so
 327 that for any classifier h and any $x \in X$, any $y \in \{0, 1\}$ if

$$\mathbb{E}_{z \sim \mathcal{V}(x)}[h(z) = y] \geq 1 - \beta$$

328 then \mathcal{W} -smoothed class classifier $\bar{h}_{\mathcal{W}}$ satisfies $\bar{h}_{\mathcal{W}}(z) = y$ for all $z \in \mathcal{U}(x)$.

329 A compression scheme of size k is a pair of functions (κ, ρ) , where the compression function
 330 $\kappa : \bigcup_{i=1}^{\infty} (X \times Y)^i \rightarrow \bigcup_{i=1}^k (X \times Y)^i$ maps samples $S = ((x_1, y_1), (x_2, y_2), \dots, (x_m, y_m))$ of
 331 arbitrary size to sub-samples of S of size at most k , and $\rho : \bigcup_{i=1}^k (X \times Y)^i \rightarrow Y^X$ is a decompression
 332 function that maps samples to classifiers. The pair (κ, ρ) is a sample compression scheme for loss ℓ
 333 and class \mathcal{H} , if for any samples S realizable by \mathcal{H} , we recover the correct labels for all $(x, y) \in S$,
 334 that is, $\mathcal{L}_S(H) = 0$ implies that $\mathcal{L}_S(\kappa \circ \rho(S)) = 0$.

335 For tolerant learning, we introduce the following generalization of compression schemes:

336 **Definition 12** (Tolerant sample compression scheme). A sample compression scheme (κ, ρ) is a
 337 \mathcal{U}, \mathcal{V} -tolerant sample compression scheme for class \mathcal{H} , if for any samples S that are $\ell^{\mathcal{V}}$ realizable by
 338 \mathcal{H} , that is $\mathcal{L}_S^{\mathcal{V}}(\mathcal{H}) = 0$, we have $\mathcal{L}_S^{\mathcal{U}}(\kappa \circ \rho(S)) = 0$.

339 The next lemma establishes that the existence of a sufficiently small tolerant compression scheme
 340 for the class \mathcal{H} yields bounds on the sample complexity of tolerantly learning \mathcal{H} . The proof of the
 341 lemma is based on a modification of a standard compression based generalization bound. Appendix
 342 Section D provides more details.

343 **Lemma 13.** Let \mathcal{H} be a hypothesis class and \mathcal{U} and \mathcal{V} be perturbation types with \mathcal{U} included in \mathcal{V} . If
 344 the class \mathcal{H} admits a $(\mathcal{U}, \mathcal{V})$ -tolerant compression scheme of size bounded by $k \ln(m)$ for sample of
 345 size m , then the class is $(\mathcal{U}, \mathcal{V})$ -tolerantly learnable in the realizable case with sample complexity
 346 bounded by $m(\epsilon, \delta) = \tilde{O} \left(\frac{k + \ln(1/\delta)}{\epsilon} \right)$.

347 We next establish a bound on the tolerant compression size for general VC-classes, which will then
 348 immediately yield the improved sample complexity bounds for tolerant learning in the realizable case.
 349 The proof is sketched here; its full version has been moved to the Appendix for lack of space.

350 **Lemma 14.** *Let $\mathcal{H} \subseteq Y^X$ be some hypothesis class with finite VC-dimension $\text{VC}(\mathcal{H}) < \infty$, and
 351 let $\mathcal{U}, \mathcal{V}, \mathcal{W}$ satisfy the conditions in Property 1 for some $\beta > 0$. Then there exists a $(\mathcal{U}, \mathcal{V})$ -tolerant
 352 sample compression scheme for \mathcal{H} of size $\tilde{O}\left(\text{VC}(\mathcal{H}) \ln\left(\frac{m}{\beta}\right)\right)$.*

353 *Proof Sketch.* We will employ a boosting-based approach to establish the claimed compression sizes.
 354 Let $S = ((x_1, y_1), (x_2, y_2), \dots, (x_m, y_m))$ be a data-set that is $\ell^\mathcal{V}$ -realizable with respect to \mathcal{H} . We
 355 let $S_\mathcal{V}$ denote an ‘‘inflated data-set’’ that contains all domain points in the \mathcal{V} -perturbation sets of
 356 the $x_i \in S^X$, that is $S_\mathcal{V}^X := \bigcup_{i=1}^m \mathcal{V}(x_i)$. Every point $z \in S_\mathcal{V}^X$ is assigned the label $y = y_i$ of the
 357 minimally-indexed $(x_i, y_i) \in S$ with $z \in \mathcal{V}(x_i)$, and we set $S_\mathcal{V}$ to be the resulting collection of
 358 labeled data-points.

359 We then use the boost-by-majority method to encode a classifier g that (roughly speaking) has error
 360 bounded by β/m over (a suitable measure over) $S_\mathcal{V}$. This boosting method outputs a T -majority
 361 vote $g(x) = \mathbb{1}[\sum_{i=1}^T h_i(x)] \geq 1/2$ over weak learners h_i , which in our case will be hypotheses from
 362 \mathcal{H} . We prove that this error can be achieved with $T = 18 \ln\left(\frac{2m}{\beta}\right)$ rounds of boosting. We prove that
 363 each weak learner that is used in the boosting procedure can be encoded with $n = \tilde{O}(\text{VC}(\mathcal{H}))$ many
 364 sample points from S . The resulting compression size is thus $n \cdot T = \tilde{O}\left(\text{VC}(\mathcal{H}) \ln\left(\frac{m}{\beta}\right)\right)$.

365 Finally, the error bound β/m of g over $S_\mathcal{V}$ implies that the error in each perturbation set $\mathcal{V}(x_i)$ of a
 366 sample point $(x_i, y_i) \in S$ is at most β . Property 1 then implies $\mathcal{L}_S^\mathcal{U}(\bar{g}_\mathcal{W}) = 0$ for the \mathcal{W} -smoothed
 367 classifier $\bar{g}_\mathcal{W}$, establishing the $(\mathcal{U}, \mathcal{V})$ -tolerant correctness of the compression scheme. \square

368 This yields the following result

369 **Theorem 15.** *Let \mathcal{H} be a hypothesis class of finite VC-dimension and $\mathcal{V}, \mathcal{U}, \mathcal{W}$ be three perturbation
 370 types (actual, reference and smoothing) satisfying Property 1 for some $\beta > 0$. Then the sample
 371 complexity (omitting log-factors) of $(\mathcal{U}, \mathcal{V})$ -tolerantly learning \mathcal{H} is bounded by*

$$m(\epsilon, \delta) = \tilde{O}\left(\frac{\text{VC}(\mathcal{H}) \ln(1/\beta) + \ln(1/\delta)}{\epsilon}\right)$$

372 *in the realizable case, and in the agnostic case by*

$$m(\epsilon, \delta) = \tilde{O}\left(\frac{\text{VC}(\mathcal{H}) \ln(1/\beta) + \ln(1/\delta)}{\epsilon^2}\right)$$

373 *Proof.* The bound for the realizable case follows immediately from Lemma 14 and the subsequent
 374 discussion (in the Appendix). For the agnostic case, we employ a reduction from agnostic robust
 375 learnability to realizable robust learnability (Montasser et al., 2019; Moran & Yehudayoff, 2016).
 376 The reduction is analogous to the one presented in Appendix C of Montasser et al. (2019) for usual
 377 (non-tolerant) robust learnability with some minor modifications. Namely, for a sample S , we choose
 378 the largest subsample S' that is $\ell^\mathcal{V}$ -realizable (this will result in competitiveness with a $\ell^\mathcal{V}$ -optimal
 379 classifier), and we will use the boosting procedure described there for the $\ell^\mathcal{U}$ loss. For the sample sizes
 380 employed for the weak learners in that procedure, we can use the sample complexity for $\epsilon = \delta = 1/3$
 381 of an optimal $(\mathcal{U}, \mathcal{V})$ -tolerant learner in the realizable case (note that each learning problem during
 382 the boosting procedure is a realizable $(\mathcal{U}, \mathcal{V})$ -tolerant learning task). These modifications result in the
 383 stated sample complexity for agnostic tolerant learnability. \square

384 In particular, for the doubling measure scenario (as considered in the previous section), we obtain

385 **Corollary 16.** *For metric tolerant learning with tolerance parameter γ in doubling di-
 386 mension d the sample complexity of learning in the realizable case is bounded by*
 387 $m(\epsilon, \delta) = \tilde{O}\left(\frac{\text{VC}(\mathcal{H})\zeta d \ln(1+1/\gamma) + \ln(1/\delta)}{\epsilon}\right)$ *and in the agnostic case by* $m(\epsilon, \delta) =$
 388 $\tilde{O}\left(\frac{\text{VC}(\mathcal{H})\zeta d \ln(1+1/\gamma) + \ln(1/\delta)}{\epsilon^2}\right)$.

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488 Checklist

- 489 1. For all authors...
- 490 (a) Do the main claims made in the abstract and introduction accurately reflect the paper’s
491 contributions and scope? [Yes] Yes, the claims are supported by actual theorems and
492 proofs in the main text.
- 493 (b) Did you describe the limitations of your work? [Yes] We have discussed the limitations
494 of our main two results (theorems), including the dependency of the sample complexity
495 on each of the parameters and the computational costs.
- 496 (c) Did you discuss any potential negative societal impacts of your work? [N/A] This is a
497 theoretical paper and we do not foresee any immediate negative impacts.
- 498 (d) Have you read the ethics review guidelines and ensured that your paper conforms to
499 them? [Yes]
- 500 2. If you are including theoretical results...
- 501 (a) Did you state the full set of assumptions of all theoretical results? [Yes] Yes, the
502 assumptions are clearly stated. Also, wherever we had an informal theorem, we
503 have linked the full version of the theorem too (with all the necessary details and
504 assumptions).
- 505 (b) Did you include complete proofs of all theoretical results? [Yes] Yes. Some of the
506 proofs are deferred to the appendix for space constraints.
- 507 3. If you ran experiments...
- 508 (a) Did you include the code, data, and instructions needed to reproduce the main experi-
509 mental results (either in the supplemental material or as a URL)? [N/A]
- 510 (b) Did you specify all the training details (e.g., data splits, hyperparameters, how they
511 were chosen)? [N/A]
- 512 (c) Did you report error bars (e.g., with respect to the random seed after running experi-
513 ments multiple times)? [N/A]
- 514 (d) Did you include the total amount of compute and the type of resources used (e.g., type
515 of GPUs, internal cluster, or cloud provider)? [N/A]
- 516 4. If you are using existing assets (e.g., code, data, models) or curating/releasing new assets...
- 517 (a) If your work uses existing assets, did you cite the creators? [N/A]
- 518 (b) Did you mention the license of the assets? [N/A]
- 519 (c) Did you include any new assets either in the supplemental material or as a URL? [N/A]
- 520
- 521 (d) Did you discuss whether and how consent was obtained from people whose data you’re
522 using/curating? [N/A]
- 523 (e) Did you discuss whether the data you are using/curating contains personally identifiable
524 information or offensive content? [N/A]
- 525 5. If you used crowdsourcing or conducted research with human subjects...
- 526 (a) Did you include the full text of instructions given to participants and screenshots, if
527 applicable? [N/A]

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531

- (b) Did you describe any potential participant risks, with links to Institutional Review Board (IRB) approvals, if applicable? [N/A]
- (c) Did you include the estimated hourly wage paid to participants and the total amount spent on participant compensation? [N/A]

532 **A Standard results from VC theory**

533 Let X be a domain. For hypothesis h and $B \subseteq X$ let $h(B) = (h(b))_{b \in B}$.

534 **Definition 17** (VC-dimension). We say \mathcal{H} shatters $B \subseteq X$ if $|\{h(B) : h \in \mathcal{H}\}| = 2^{|B|}$. The
 535 VC-dimension of \mathcal{H} , denoted by $\text{VC}(\mathcal{H})$, is defined to be the supremum of the size of the sets that are
 536 shattered by \mathcal{H} .

537 **Theorem 18** (Existence of Realizable PAC Learners Hanneke (2016); Simon (2015); Blumer et al.
 538 (1989)). Let \mathcal{H} be a hypothesis class with bounded VC dimension. Then \mathcal{H} is PAC learnable in the
 539 realizable setting using $O\left(\frac{\text{VC}(\mathcal{H}) + \log(1/\delta)}{\epsilon}\right)$ samples.

540 **Theorem 19** (Existence of Agnostic PAC Learners Haussler (1992)). Let \mathcal{H} be a hypothesis class with
 541 bounded VC dimension. Then \mathcal{H} is PAC learnable in the agnostic setting using $O\left(\frac{\text{VC}(\mathcal{H}) + \log(1/\delta)}{\epsilon^2}\right)$
 542 samples.

543 **B Metric spaces**

544 **Definition 20.** A metric space (X, dist) is called a doubling metric if there exists a constant M such
 545 that every ball of radius r in it can be covered by at most M balls of radius $r/2$. The quantity $\log_2 M$
 546 is called the doubling dimension.

547 **Definition 21.** For a metric space (X, dist) , a measure μ defined on X is called a doubling measure
 548 if there exists a constant C , such that for all $x \in X$ and $r \in \mathbb{R}^+$, we have that $0 < \mu(\mathcal{B}_{2r}(x)) \leq$
 549 $C \cdot \mu(\mathcal{B}_r(x)) < \infty$. In this case, μ is called C -doubling.

550 It can be shown (Luukkainen & Saksman, 1998) that every complete metric with doubling dimension
 551 d has a C -doubling measure μ for some $C \leq 2^{cd}$ where c is a universal constant. For example,
 552 Euclidean spaces with an ℓ_p distance metric are complete and the Lesbesgue measure is a doubling
 553 measure.

554 The following lemmas follow straightforwardly from the definitions doubling metric and measures.

555 **Lemma 22.** Let (X, dist) be a doubling metric equipped with a C -doubling measure μ . Then for all
 556 $x \in X$, $r > 0$, and $\alpha > 1$, we have that $\mu(\mathcal{B}_{\alpha r}(x)) \leq C^{\lceil \log_2 \alpha \rceil} \cdot \mu(\mathcal{B}_r(x))$

557 *Proof.* Since μ is a measure, if $B, B' \subseteq X$ such that $B \subseteq B'$, then $\mu(B) \leq \mu(B')$. Let $R = 2^{\lceil \log_2 \alpha \rceil}$.
 558 It's clear that $R \geq \alpha$. Therefore $\mathcal{B}_{\alpha r}(x) \subseteq \mathcal{B}_R(x)$. Expanding $\mathcal{B}_r(x)$ by a factor of two $\lceil \log_2 \alpha \rceil$
 559 times, we get $\mathcal{B}_R(x)$, which means $\mu(\mathcal{B}_R(x)) \leq C^{\lceil \log_2 \alpha \rceil} \cdot \mu(\mathcal{B}_r(x))$. But since $\mathcal{B}_{\alpha r}(x) \subseteq \mathcal{B}_R(x)$,
 560 we get the desired result. \square

561 **Lemma 23.** Let (X, dist) be a doubling metric equipped with a C -doubling measure μ . Let $x, x' \in$
 562 X , $r > 0$, and $\alpha > 1$ be such that $\mathcal{B}_r(x') \subseteq \mathcal{B}_{\alpha r}(x)$. Then $\mu(\mathcal{B}_{\alpha r}(x)) \leq C^{\lceil \log_2(2\alpha) \rceil} \cdot \mu(\mathcal{B}_r(x'))$.

563 *Proof.* By Lemma 22, all we need to show is that $\mathcal{B}_{\alpha r}(x) \subseteq \mathcal{B}_{2\alpha r}(x')$. Indeed, let $y \in \mathcal{B}_{\alpha r}(x)$ be
 564 any point. Then, from triangle inequality, we have that

$$\begin{aligned} d(x', y) &\leq d(x, x') + d(x, y) \\ &\leq d(x, x') + \alpha r \end{aligned}$$

565 Moreover, since $x' \in \mathcal{B}_{\alpha r}(x)$, we have that $d(x, x') \leq \alpha r$. Substituting into the equation above, we
 566 get $d(x', y) \leq 2\alpha r$, which means $y \in \mathcal{B}_{2\alpha r}(x')$. \square

567 Finally, we also get:

568 **Lemma 24.** For any family \mathcal{M} of complete, doubling metric spaces, there exist constants $c_1, c_2 > 0$
 569 such that for any metric space $(X, \text{dist}) \in \mathcal{M}$ with doubling dimension d , there exists a measure μ
 570 such that if a ball \mathcal{B}_r of radius $r > 0$ is completely contained inside a ball $\mathcal{B}_{\alpha r}$ of radius αr (with
 571 potentially a different center) for any $\alpha > 1$, then $0 < \mu(\mathcal{B}_{\alpha r}) \leq (c_1 \alpha)^{c_2 d} \mu(\mathcal{B}_r)$.

572 *Proof.* We prove this when \mathcal{M} is the set of all complete, doubling metric spaces employing Lemmas
573 22 and 23, that can be found in Appendix, part B. We have that $C^{\lceil \log_2(2\alpha) \rceil} \leq (2\alpha)^{2 \log_2 C}$. Since
574 $\log_2 C \leq cd$, we get $(2\alpha)^{2 \log_2 C} \leq (2\alpha)^{cd}$. Thus $c_1 = 2$ and $c_2 = c$. \square

575 **Corollary 25.** *Suppose we have a constant $\alpha_0 > 1$ such that we know that $\alpha \geq \alpha_0$. Then the bound*
576 *in Lemma 24 can be further simplified to $0 < \mu(\mathcal{B}_{\alpha r}) \leq \alpha^{\zeta d} \mu(\mathcal{B}_r)$, where ζ depends on \mathcal{M} and α_0 .*
577 *Furthermore, if $c_1 = 1$ then we can set $\alpha_0 = 1$.*

578 *Proof.* $(c_1 \alpha)^{c_2 d} = \alpha^{c_2 d (1 + \log_{\alpha} c_1)} \leq \alpha^{c_2 d (1 + \log_{\alpha_0} c_1)} = \alpha^{\zeta d}$ for $\zeta = c_2 (1 + \log_{\alpha_0} c_1)$. If $c_1 = 1$,
579 then $\zeta = c_2$ for all α . \square

580 C Proof of Lemma 11

581 Let $X_{\text{err}} = \{z \in \mathcal{B}_{r(1+\gamma)}(x) \mid g(z) \neq y\}$. Then, we have that $\Sigma_{g,y}(x) =$
582 $\mathbb{E}_{z \sim \mathcal{B}_{r(1+\gamma)}(x)} \mathbb{1}[g(z) \neq y] = \frac{\mu(X_{\text{err}})}{\mu(\mathcal{B}_{r(1+\gamma)}(x))}$. Further, for all $z \in \mathcal{B}_r(x)$, we have
583 $\mathbb{E}_{z' \sim \mathcal{B}_{r\gamma}(z)} \mathbb{1}[g(z') \neq y] = \frac{\mu(X_{\text{err}} \cap \mathcal{B}_{r\gamma}(z))}{\mu(\mathcal{B}_{r\gamma}(z))}$.

584 Let $z \in \mathcal{B}_r(x)$. Since this implies that $\mathcal{B}_{r\gamma}(z) \subseteq \mathcal{B}_{r(1+\gamma)}(x)$, the worst case happens when
585 $X_{\text{err}} \subseteq \mathcal{B}_{r\gamma}(z)$. Therefore,

$$\begin{aligned} \sigma_{g,y}(x) &= \mathbb{E}_{z' \sim \mathcal{B}_{r\gamma}(z)} \mathbb{1}[g(z') \neq y] & (3) \\ &= \frac{\mu(X_{\text{err}} \cap \mathcal{B}_{r\gamma}(z))}{\mu(\mathcal{B}_{r\gamma}(z))} \\ &\leq \frac{\mu(X_{\text{err}})}{\mu(\mathcal{B}_{r\gamma}(z))} \\ &\leq \frac{\Sigma_{g,y}(x) \cdot \mu(\mathcal{B}_{r(1+\gamma)}(x))}{\mu(\mathcal{B}_{r\gamma}(z))} \\ &\leq \Sigma_{g,y}(x) \cdot \left(\frac{1+\gamma}{\gamma}\right)^{\zeta d}, \end{aligned}$$

586 where the last inequality is implied by Lemma 24. Thus, $\Sigma(x) \leq \frac{1}{3} \cdot \left(\frac{1+\gamma}{\gamma}\right)^{-\zeta d}$ implies that
587 $\sigma(z) \leq 1/3$ as claimed.

588 D Compression based bounds

589 D.1 Proof of Lemma 13

590 To prove the generalization bound for tolerant learning, we employ the following lemma that
591 establishes generalization for compression schemes for adversarial losses:

592 **Lemma 26** (Lemma 11, (Montasser et al., 2019)). *For any $k \in \mathbb{N}$ and fixed function $\rho : \bigcup_{i=1}^k (X \times$
593 $Y)^i \rightarrow Y^X$, for any distribution P over $X \times Y$ and any $m \in \mathbb{N}$, with probability at least $(1 - \delta)$
594 over an i.i.d. sample $S = ((x_1, y_1), (x_2, y_2), \dots, (x_m, y_m))$: if there exist indices i_1, i_2, \dots, i_k such
595 that*

$$\mathcal{L}_S^{\mathcal{U}}(\rho((x_{i_1}, y_{i_1}), (x_{i_2}, y_{i_2}), \dots, (x_{i_k}, y_{i_k}))) = 0$$

596 *then the robust loss of the decompression with respect to P is bounded by*

$$\begin{aligned} \mathcal{L}_P^{\mathcal{U}}(\rho((x_{i_1}, y_{i_1}), (x_{i_2}, y_{i_2}), \dots, (x_{i_k}, y_{i_k}))) \\ \leq \frac{1}{m-k} (k \ln(m) + \ln(1/\delta)) \end{aligned}$$

597 The above lemma implies that if (κ, ρ) is a compression scheme that compresses data-sets of size m
598 to at most $k \ln(m)$ data points, for class \mathcal{H} and robust loss $\ell^{\mathcal{U}}$, then the sample complexity (omitting
599 logarithmic factors) of robustly learning \mathcal{H} in the realizable case is bounded by

$$m(\epsilon, \delta) = \tilde{O}\left(\frac{k + \ln(1/\delta)}{\epsilon}\right)$$

600 For the tolerant setting, since every sample that is realizable with respect to $\ell^{\mathcal{V}}$ is also realizable with
601 respect to $\ell^{\mathcal{U}}$, if a $(\mathcal{U}, \mathcal{V})$ -tolerant compression scheme compresses to at most $k \ln(m)$ data-points
602 and decompresses all $\ell^{\mathcal{V}}$ -realizable samples S to functions that have $\ell^{\mathcal{U}}$ -loss 0 on S , then the lemma
603 implies the above bound for the $(\mathcal{U}, \mathcal{V})$ -tolerant sample complexity of learning \mathcal{H} .

604 D.2 Proof of Lemma 14

605 The proof of this Lemma employs the notions of a sample being ϵ -net or and ϵ -approximation for a
606 hypothesis class \mathcal{H} . A labeled data set $S = ((x_1, y_1), (x_2, y_2), \dots, (x_m, y_m))$ is an ϵ -net for class \mathcal{H}
607 with respect to distribution P over $X \times Y$ if for every hypothesis $h \in \mathcal{H}$ with $\mathcal{L}_P^{0/1}(h) \geq \epsilon$, there
608 exists an index j and $(x_j, y_j) \in S$ with $h(x_j) \neq y_j$. S is an ϵ -approximation for class \mathcal{H} with respect
609 to distribution P over $X \times Y$ if for every hypothesis $h \in \mathcal{H}$ we have $|\mathcal{L}_S^{0/1}(h) - \mathcal{L}_P^{0/1}(h)| \leq \epsilon$.
610 Standard VC-theory tells us that, for classes with bounded VC-dimension, sufficiently large samples
611 from P are ϵ -nets or ϵ -approximations with high probability (Haussler & Welzl, 1987).

612 *Proof.* We will employ a boosting-based approach to establish the claimed compression sizes. Let
613 $S = ((x_1, y_1), (x_2, y_2), \dots, (x_m, y_m))$ be a data-set that is $\ell^{\mathcal{V}}$ -realizable with respect to \mathcal{H} . We
614 let $S_{\mathcal{V}}$ denote an ‘‘inflated data-set’’ that contains all domain points in the perturbation sets of the
615 $x_i \in S^X$, that is

$$S_{\mathcal{V}}^X := \bigcup_{i=1}^m \mathcal{V}(x_i)$$

616 Every point $z \in S_{\mathcal{V}}^X$ is assigned the label $y = y_i$ of the minimally-indexed $(x_i, y_i) \in S$ with
617 $z \in \mathcal{V}(x_i)$, and we set $S_{\mathcal{V}}$ to be the resulting collection of labeled data-points. (Note that since the
618 sample S is assumed to be $\ell^{\mathcal{V}}$ -realizable, assigning it the label of some other corresponding data point
619 in case $z \in \mathcal{V}(x_i) \cap \mathcal{V}(x_j)$ for $x_i \neq x_j$, would not induce any inconsistencies). Now let D be the
620 probability measure over $S_{\mathcal{V}}^X$ defined by first sampling an index j uniformly from $[j] = \{1, 2, \dots, j\}$
621 and then sampling a domain point $z \sim \mathcal{V}(x_j)$ from the \mathcal{V} -perturbation set around the j -th sample
622 point in S . Note that this implies that if $D(B) \leq (\beta/m)$ for some subset $B \subseteq S_{\mathcal{V}}^X$, then

$$\mathbb{P}_{z \sim \mathcal{V}(x)}[z \in B] \leq \beta \quad (4)$$

623 for all $x \in S^X$.

624 We will now show that, by means of a compression scheme, we can encode a hypothesis g with
625 binary loss

$$\mathcal{L}_D^{0/1}(g) \leq \beta/m. \quad (5)$$

626 Property 1 together with Equation 4 then implies that the resulting \mathcal{W} -smoothed function \bar{g} has
627 \mathcal{U} -robust loss 0 on the sample S , $\mathcal{L}_S^{\mathcal{U}}(\bar{g}) = 0$. Since the smoothing is a deterministic operation once g
628 is fixed, this implies the existence of a $(\mathcal{U}, \mathcal{V})$ -tolerant compression scheme.

629 Standard VC-theory tells us that, for a class G of bounded VC-dimension, for any distribution over
630 $X \times Y$, and any $\epsilon, \delta > 0$, with probability at least $(1 - \delta)$ an i.i.d. sample of size $\Theta\left(\frac{\text{VC}(G) + \ln(1/\delta)}{\epsilon^2}\right)$
631 is an ϵ -approximation for the class G (Haussler & Welzl, 1987). This implies in particular, that there
632 exists a finite subset $S_{\mathcal{V}}^f \subset S_{\mathcal{V}}$ of size at most $\frac{4m^2 C \cdot \text{VC}(G)}{\beta^2}$ (for some constant C) with the property
633 that any classifier $g \in G$ with empirical (binary) loss at most $\beta/2m$ on $S_{\mathcal{V}}^f$ has loss $\mathcal{L}_D^{0/1}(g) \leq \beta/m$
634 with respect to the distribution D . We will choose such a set $S_{\mathcal{V}}^f$ for the class G of T -majority votes
635 over \mathcal{H} for $T = 18 \ln\left(\frac{2m}{\beta}\right)$. That is

$$G = \{g \in Y^X \mid \exists h_1, h_2, \dots, h_T \in \mathcal{H} : \\ g(x) = \mathbb{1}[\sum_{i=1}^T h_i(x) \geq 1/2]\}$$

636 The VC-dimension of G is bounded by $\text{VC}(G) = \mathcal{O}(T \cdot \text{VC}(\mathcal{H}) \log(T \text{VC}(\mathcal{H}))) =$
637 $\mathcal{O}(18 \ln\left(\frac{2m}{\beta}\right) \text{VC}(\mathcal{H}) \log(18 \ln\left(\frac{2m}{\beta}\right) \text{VC}(\mathcal{H})))$ (Shalev-Shwartz & Ben-David, 2014).

638 We will now show how to obtain the classifier g by means of a boosting approach on the finite data-set
639 $S_{\mathcal{V}}^f$. More specifically, we will use the boost-by-majority method. This method outputs a T -majority

640 vote $g(x) = \mathbb{1} [\sum_{i=1}^T h_i(x)] \geq 1/2$ over weak learners h_i , which in our case will be hypotheses from
 641 \mathcal{H} . After T iterations with γ -weak learners, the empirical loss over the sample S_V^f is bounded by
 642 $e^{-2\gamma^2 T}$ (see Section 13.1 in (Schapire & Freund, 2013)). Thus, with $\gamma = 1/6$, and $T = 18 \ln(\frac{2m}{\beta})$,
 643 we obtain

$$\mathcal{L}_{S_V^f}^{0/1}(g) \leq \frac{\beta}{2m}$$

644 which, by the choice of S_V^f implies

$$\mathcal{L}_D^{0/1}(g) \leq \beta/m$$

645 which is what we needed to show according to Equation 5.

646 It remains to argue that the weak learners to be employed in the boosting procedure can be encoded
 647 by a small number of sample points from the original sample S . For this part, we will employ a
 648 technique introduced earlier for robust compression (Montasser et al., 2019). Recall that the set S is
 649 \mathcal{V} -robustly realizable, which implies that the set S_V^f is (binary loss-) realizable by \mathcal{H} . By standard
 650 VC-theory, for every distribution D_i over S_V^f , there exists an ϵ -net of size $\mathcal{O}(\text{VC}(\mathcal{H})/\epsilon)$ (Haussler &
 651 Welzl, 1987). Thus, for every distribution D_i over S_V^f (that may occur during the boosting procedure),
 652 there exists a subsample S_i of S_V^f , of size at most $n = \mathcal{O}(3\text{VC}(\mathcal{H}))$ with the property that every
 653 hypothesis from \mathcal{H} that is consistent with S_i has binary loss at most $1/3$ with respect to D_i (thus
 654 can serve as a weak learner for margin $\gamma = 1/6$ in the above procedure). Now for every labeled
 655 point $(x, y) \in S_i$, there is a sample point $(x_j, y_j) \in S$ in the original sample S such that $x \in \mathcal{V}(x_j)$
 656 and $y = y_j$. Let S'_i be the collection of these corresponding original sample points. Note that any
 657 hypothesis $h \in \mathcal{H}$ that is \mathcal{V} -robustly consistent with S'_i is consistent with S_i . Therefore we can use
 658 the n original data-points in S'_i to encode the weak learner h_i (for the decoding any \mathcal{V} -robust ERM
 659 hypothesis can be chosen to obtain h_i).

660 To summarize, we will compress the sample S to the sequence S'_1, S'_2, \dots, S'_T of $n \cdot T =$
 661 $\mathcal{O}(\text{VC}(\mathcal{H}) \ln(\frac{m}{\beta}))$ sample points from S . To decode, we obtain the function g as a majority vote
 662 over the weak learner h_i and proceed to obtain the \mathcal{W} -smoothed function \bar{g} . This function \bar{g} satisfies
 663 $\mathcal{L}_S^{\mathcal{U}}(\bar{g}) = 0$ and by this we have established the existence of a \mathcal{U}, \mathcal{V} -tolerant compression scheme of
 664 size $\mathcal{O}(\text{VC}(\mathcal{H}) \ln(\frac{m}{\beta}))$ as claimed. \square