

482 Appendix

483 A Optimal Monteiro-Svaiter Acceleration Framework

484 In this section, we present some general results that hold for the Monteiro-Svaiter Acceleration
485 framework. In particular, in the first part of this section (Section A.1), we present the proof of
486 Proposition 1.

487 A.1 Proof of Proposition 1

488 To begin with, we establish a potential function for Algorithm 1, as shown in Proposition 2. The
489 result is similar to Proposition 1 in [31], but for completeness we present its proof loosely following
490 the strategy in [44 Theorem 5.3]. To simplify the notations, we use f^* to denote the optimal $f(\mathbf{x}^*)$.

491 **Proposition 2.** *Consider the iterates generated by Algorithm 1. If f is convex, then*

$$A_{k+1}(f(\mathbf{x}_{k+1}) - f^*) + \frac{1}{2}\|\mathbf{z}_{k+1} - \mathbf{x}^*\|^2 \leq A_k(f(\mathbf{x}_k) - f^*) + \frac{1}{2}\|\mathbf{z}_k - \mathbf{x}^*\|^2. \quad (21)$$

492 Moreover, let $\sigma = \alpha_1 + \alpha_2$ and we have

$$\sum_{k=0}^{N-1} \frac{a_k^2}{\eta_k^2} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2 \leq \frac{1}{1 - \sigma^2} \|\mathbf{z}_0 - \mathbf{x}^*\|^2. \quad (22)$$

493 *Proof.* Since f is convex, it holds that

$$\begin{aligned} f(\mathbf{x}_k) - f(\hat{\mathbf{x}}_{k+1}) - \langle \nabla f(\hat{\mathbf{x}}_{k+1}), \mathbf{x}_k - \hat{\mathbf{x}}_{k+1} \rangle &\geq 0, \\ f(\mathbf{x}^*) - f(\hat{\mathbf{x}}_{k+1}) - \langle \nabla f(\hat{\mathbf{x}}_{k+1}), \mathbf{x}^* - \hat{\mathbf{x}}_{k+1} \rangle &\geq 0. \end{aligned}$$

494 By summing up the two inequalities with weights a_k and A_k respectively, we get

$$A_k(f(\mathbf{x}_k) - f^*) - (A_k + a_k)(f(\hat{\mathbf{x}}_{k+1}) - f^*) - a_k \langle \nabla f(\hat{\mathbf{x}}_{k+1}), \mathbf{x}^* - \hat{\mathbf{x}}_{k+1} - \frac{A_k}{a_k}(\hat{\mathbf{x}}_{k+1} - \mathbf{x}_k) \rangle \geq 0. \quad (23)$$

495 Let $\tilde{\mathbf{z}}_{k+1} = \hat{\mathbf{x}}_{k+1} + \frac{A_k}{a_k}(\hat{\mathbf{x}}_{k+1} - \mathbf{x}_k)$. By rearranging the terms, (23) can be rewritten as

$$(A_k + a_k)(f(\hat{\mathbf{x}}_{k+1}) - f^*) - A_k(f(\mathbf{x}_k) - f^*) \leq a_k \langle \nabla f(\hat{\mathbf{x}}_{k+1}), \tilde{\mathbf{z}}_{k+1} - \mathbf{x}^* \rangle. \quad (24)$$

496 Moreover, note that the update rule for \mathbf{z}_{k+1} in both (9) and (10) can be written as

$$\mathbf{z}_{k+1} - \mathbf{z}_k = -\frac{\hat{\eta}_k}{\eta_k} a_k \nabla f(\hat{\mathbf{x}}_{k+1}). \quad (25)$$

497 Also, since we also have $\mathbf{z}_k = \mathbf{y}_k + \frac{A_k}{a_k}(\mathbf{y}_k - \mathbf{x}_k)$ from (2), we can write

$$\begin{aligned} \tilde{\mathbf{z}}_{k+1} - \mathbf{z}_k &= \left[\hat{\mathbf{x}}_{k+1} + \frac{A_k}{a_k}(\hat{\mathbf{x}}_{k+1} - \mathbf{x}_k) \right] - \left[\mathbf{y}_k + \frac{A_k}{a_k}(\mathbf{y}_k - \mathbf{x}_k) \right] \\ &= \frac{A_k + a_k}{a_k}(\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k) = \frac{a_k}{\eta_k}(\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k), \end{aligned} \quad (26)$$

498 where we used the fact that $(A_k + a_k)\eta_k = a_k^2$ in the last equality (cf. (2)). Hence, combining (25)
499 and (26) leads to

$$\|\tilde{\mathbf{z}}_{k+1} - \mathbf{z}_{k+1}\| = \|\tilde{\mathbf{z}}_{k+1} - \mathbf{z}_k - (\mathbf{z}_{k+1} - \mathbf{z}_k)\| = \frac{a_k}{\eta_k} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k + \hat{\eta}_k \nabla f(\hat{\mathbf{x}}_{k+1})\| \leq \sigma \frac{a_k}{\eta_k} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|. \quad (27)$$

500 where we used (8) in the last inequality. In the following, we distinguish two cases depending on
501 $\hat{\eta}_k = \eta_k$ or $\hat{\eta}_k < \eta_k$. In both cases, we shall prove that

$$A_{k+1}(f(\mathbf{x}_{k+1}) - f^*) + \frac{1}{2}\|\mathbf{z}_{k+1} - \mathbf{x}^*\|^2 \leq A_k(f(\mathbf{x}_k) - f^*) + \frac{1}{2}\|\mathbf{z}_k - \mathbf{x}^*\|^2 - \frac{(1 - \sigma^2)a_k^2}{2\eta_k^2} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2. \quad (28)$$

502 If this is true, then Proposition 2 immediately follows. Indeed, since $\sigma < 1$, the last term in the
 503 right-hand side of (28) is negative, which implies (21). Moreover, (22) follows from summing the
 504 inequality in (28) from $k = 0$ to $N - 1$.

505 **Case I:** $\hat{\eta}_k = \eta_k$. Since by (9) we have $\mathbf{x}_{k+1} = \hat{\mathbf{x}}_{k+1}$ and $A_{k+1} = A_k + a_k$, (24) becomes

$$A_{k+1}(f(\mathbf{x}_{k+1}) - f^*) - A_k(f(\mathbf{x}_k) - f^*) \leq a_k \langle \nabla f(\mathbf{x}_{k+1}), \tilde{\mathbf{z}}_{k+1} - \mathbf{x}^* \rangle.$$

506 Using $\mathbf{z}_{k+1} = \mathbf{z}_k - a_k \nabla f(\mathbf{x}_{k+1})$ in (9), we have

$$\begin{aligned} & A_{k+1}(f(\mathbf{x}_{k+1}) - f^*) - A_k(f(\mathbf{x}_k) - f^*) \\ & \leq \langle \mathbf{z}_k - \mathbf{z}_{k+1}, \tilde{\mathbf{z}}_{k+1} - \mathbf{x}^* \rangle \\ & = \langle \mathbf{z}_k - \mathbf{z}_{k+1}, \tilde{\mathbf{z}}_{k+1} - \mathbf{z}_{k+1} \rangle + \langle \mathbf{z}_k - \mathbf{z}_{k+1}, \mathbf{z}_{k+1} - \mathbf{x}^* \rangle \\ & = \frac{1}{2} \|\mathbf{z}_k - \mathbf{z}_{k+1}\|^2 + \frac{1}{2} \|\tilde{\mathbf{z}}_{k+1} - \mathbf{z}_{k+1}\|^2 - \frac{1}{2} \|\tilde{\mathbf{z}}_{k+1} - \mathbf{z}_k\|^2 \\ & \quad + \frac{1}{2} \|\mathbf{z}_k - \mathbf{x}^*\|^2 - \frac{1}{2} \|\mathbf{z}_{k+1} - \mathbf{x}^*\|^2 - \frac{1}{2} \|\mathbf{z}_k - \mathbf{z}_{k+1}\|^2 \\ & \leq \frac{1}{2} \|\mathbf{z}_k - \mathbf{x}^*\|^2 - \frac{1}{2} \|\mathbf{z}_{k+1} - \mathbf{x}^*\|^2 - \frac{(1 - \sigma^2)a_k^2}{2\eta_k^2} \|\mathbf{x}_{k+1} - \mathbf{y}_k\|^2, \end{aligned} \quad (29)$$

507 where we used (26) and (27) in the last inequality. This immediately leads to (28) after rearranging
 508 the terms.

509 **Case II:** $\hat{\eta}_k < \eta_k$. Since $0 < \gamma_k < 1$ and $\mathbf{x}_{k+1} = \frac{(1-\gamma_k)A_k}{A_k + \gamma_k a_k} \mathbf{x}_k + \frac{\gamma_k(A_k + a_k)}{A_k + \gamma_k a_k} \hat{\mathbf{x}}_{k+1}$ according to (10),
 510 by Jensen's inequality we have $(A_k + \gamma_k a_k)f(\mathbf{x}_{k+1}) \leq \gamma_k(A_k + a_k)f(\hat{\mathbf{x}}_{k+1}) + (1 - \gamma_k)A_k f(\mathbf{x}_k)$,
 511 which further implies that

$$(A_k + \gamma_k a_k)(f(\mathbf{x}_{k+1}) - f^*) - A_k(f(\mathbf{x}_k) - f^*) \leq \gamma_k(A_k + a_k)(f(\hat{\mathbf{x}}_{k+1}) - f^*) - \gamma_k A_k(f(\mathbf{x}_k) - f^*).$$

512 Moreover, since $A_{k+1} = A_k + \gamma_k a_k$ by (10), together with (24) we obtain

$$A_{k+1}(f(\mathbf{x}_{k+1}) - f^*) - A_k(f(\mathbf{x}_k) - f^*) \leq \gamma_k a_k \langle \nabla f(\hat{\mathbf{x}}_{k+1}), \tilde{\mathbf{z}}_{k+1} - \mathbf{x}^* \rangle.$$

513 Using $\mathbf{z}_{k+1} = \mathbf{z}_k - \gamma_k a_k \nabla f(\hat{\mathbf{x}}_{k+1})$ in (10), we follow the same reasoning as in (29) to get:

$$\begin{aligned} & A_{k+1}(f(\mathbf{x}_{k+1}) - f^*) - A_k(f(\mathbf{x}_k) - f^*) \\ & \leq \langle \mathbf{z}_k - \mathbf{z}_{k+1}, \tilde{\mathbf{z}}_{k+1} - \mathbf{x}^* \rangle \\ & = \langle \mathbf{z}_k - \mathbf{z}_{k+1}, \tilde{\mathbf{z}}_{k+1} - \mathbf{z}_{k+1} \rangle + \langle \mathbf{z}_k - \mathbf{z}_{k+1}, \mathbf{z}_{k+1} - \mathbf{x}^* \rangle \\ & = \frac{1}{2} \|\mathbf{z}_k - \mathbf{z}_{k+1}\|^2 + \frac{1}{2} \|\tilde{\mathbf{z}}_{k+1} - \mathbf{z}_{k+1}\|^2 - \frac{1}{2} \|\tilde{\mathbf{z}}_{k+1} - \mathbf{z}_k\|^2 \\ & \quad + \frac{1}{2} \|\mathbf{z}_k - \mathbf{x}^*\|^2 - \frac{1}{2} \|\mathbf{z}_{k+1} - \mathbf{x}^*\|^2 - \frac{1}{2} \|\mathbf{z}_k - \mathbf{z}_{k+1}\|^2 \\ & \leq \frac{1}{2} \|\mathbf{z}_k - \mathbf{x}^*\|^2 - \frac{1}{2} \|\mathbf{z}_{k+1} - \mathbf{x}^*\|^2 - \frac{(1 - \sigma^2)a_k^2}{2\eta_k^2} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2, \end{aligned}$$

514 which also leads to (28). \square

515 Next, we prove a lower bound on A_N . Recall that \mathcal{B} denotes the set of iteration indices where the
 516 line search scheme backtracks, i.e., $\mathcal{B} \triangleq \{k : \hat{\eta}_k < \eta_k\}$.

517 **Lemma 3.** For any $N \geq 0$, it holds that

$$A_N \geq \frac{1}{4} \left(\sqrt{\hat{\eta}_0} + \sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \sqrt{\hat{\eta}_k} \right)^2. \quad (30)$$

518 *Proof.* To begin with, according to the update rule of A_{k+1} in (9) and (10) and the expression of a_k
 519 in (2), the sequence $\{A_k\}$ follows the dynamic:

$$A_{k+1} = \begin{cases} A_k + a_k, & \text{if } \hat{\eta}_k = \eta_k \ (k \notin \mathcal{B}); \\ A_k + \gamma_k a_k, & \text{if } \hat{\eta}_k < \eta_k \ (k \in \mathcal{B}), \end{cases} \quad \text{where } \gamma_k = \frac{\hat{\eta}_k}{\eta_k} \text{ and } a_k = \frac{\eta_k + \sqrt{\eta_k^2 + 4\eta_k A_k}}{2}.$$

520 Since we initialize $A_0 = 0$, we have $a_0 = \eta_0$. We further have $A_1 = \hat{\eta}_0$, since we get $A_1 =$
 521 $A_0 + a_0 = \hat{\eta}_0$ if $0 \notin \mathcal{B}$, while we get $A_1 = A_0 + \gamma_0 a_0 = \frac{\hat{\eta}_0}{\eta_0} \eta_0 = \hat{\eta}_0$ if $0 \in \mathcal{B}$. Moreover:

522 • In **Case I** where $k \notin \mathcal{B}$, we have

$$A_{k+1} = A_k + a_k = A_k + \frac{\eta_k + \sqrt{\eta_k^2 + 4\eta_k A_k}}{2} \geq A_k + \frac{\eta_k}{2} + \sqrt{\eta_k A_k} \geq \left(\sqrt{A_k} + \frac{\sqrt{\eta_k}}{2} \right)^2,$$

523 which further implies that $\sqrt{A_{k+1}} \geq \sqrt{A_k} + \frac{\sqrt{\eta_k}}{2} = \sqrt{A_k} + \frac{\sqrt{\hat{\eta}_k}}{2}$.

524 • In **Case II** where $k \in \mathcal{B}$, we have $A_{k+1} = A_k + \gamma_k a_k \geq A_k$, which implies that $\sqrt{A_{k+1}} \geq \sqrt{A_k}$.

525 Considering the above, we obtain $\sqrt{A_N} \geq \sqrt{A_1} + \sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \frac{\sqrt{\hat{\eta}_k}}{2}$, which leads to (30). \square

526 Lemma 3 provides a lower bound on A_N in terms of the step sizes $\hat{\eta}_k$ in those iterations where the
 527 line search scheme does not backtrack, i.e., $k \notin \mathcal{B}$. The following lemma shows how we can further
 528 prove a lower bound in terms of all the step sizes $\{\hat{\eta}_k\}_{k=0}^{N-1}$.

529 **Lemma 4.** *We have*

$$\sum_{1 \leq k \leq N-1, k \in \mathcal{B}} \sqrt{\hat{\eta}_k} \leq \frac{1}{1 - \sqrt{\beta}} \left(\sqrt{\hat{\eta}_0} + \sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \sqrt{\hat{\eta}_k} \right). \quad (31)$$

530 As a corollary, we have

$$\sqrt{\hat{\eta}_0} + \sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \sqrt{\hat{\eta}_k} \geq \frac{1 - \sqrt{\beta}}{2 - \sqrt{\beta}} \sum_{k=0}^{N-1} \sqrt{\hat{\eta}_k}. \quad (32)$$

531 *Proof.* When the line search scheme backtracks, i.e., $k \in \mathcal{B}$, we have $\hat{\eta}_k \leq \beta \eta_k$. Therefore,

$$\sum_{1 \leq k \leq N-1, k \in \mathcal{B}} \sqrt{\hat{\eta}_k} \leq \sum_{1 \leq k \leq N-1, k \in \mathcal{B}} \sqrt{\beta \eta_k} \leq \sum_{k=1}^{N-1} \sqrt{\beta \eta_k} = \sqrt{\beta \eta_1} + \sum_{k=1}^{N-2} \sqrt{\beta \eta_{k+1}}. \quad (33)$$

532 Moreover, in the update of Algorithm 1, we have $\eta_{k+1} = \hat{\eta}_k / \beta$ if $k \notin \mathcal{B}$ (cf. Line 8) and $\eta_{k+1} = \hat{\eta}_k$
 533 otherwise (cf. Line 13). This implies that $\eta_1 \leq \hat{\eta}_0 / \beta$ and we further have

$$\begin{aligned} \sqrt{\beta \eta_1} + \sum_{k=1}^{N-2} \sqrt{\beta \eta_{k+1}} &= \sqrt{\beta \eta_1} + \sum_{1 \leq k \leq N-2, k \notin \mathcal{B}} \sqrt{\beta \eta_{k+1}} + \sum_{1 \leq k \leq N-2, k \in \mathcal{B}} \sqrt{\beta \eta_{k+1}} \\ &\leq \sqrt{\hat{\eta}_0} + \sum_{1 \leq k \leq N-2, k \notin \mathcal{B}} \sqrt{\hat{\eta}_k} + \sum_{1 \leq k \leq N-2, k \in \mathcal{B}} \sqrt{\beta \hat{\eta}_k} \\ &\leq \sqrt{\hat{\eta}_0} + \sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \sqrt{\hat{\eta}_k} + \sum_{1 \leq k \leq N-1, k \in \mathcal{B}} \sqrt{\beta \hat{\eta}_k}. \end{aligned} \quad (34)$$

534 We combine (33) and (34) to get

$$\sum_{1 \leq k \leq N-1, k \in \mathcal{B}} \sqrt{\hat{\eta}_k} \leq \sqrt{\hat{\eta}_0} + \sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \sqrt{\hat{\eta}_k} + \sum_{1 \leq k \leq N-1, k \in \mathcal{B}} \sqrt{\beta \hat{\eta}_k}.$$

535 By rearranging the terms and simple algebraic manipulation, we obtain (31) as desired. Finally, (32)
 536 follows by adding $\sqrt{\hat{\eta}_0} + \sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \sqrt{\hat{\eta}_k}$ to both sides of (31). \square

537 Now we are ready to prove Proposition 1

538 *Proof of Proposition 1.* By Proposition 2, the potential function $\phi_k \triangleq A_k(f(\mathbf{x}_k) - f^*) + \frac{1}{2} \|\mathbf{z}_k - \mathbf{x}^*\|^2$
 539 is non-increasing in each iteration. Hence, via a recursive augment we have $A_N(f(\mathbf{x}_N) - f^*) \leq$
 540 $\phi_N \leq \dots \leq \phi_0 = \frac{1}{2} \|\mathbf{z}_0 - \mathbf{x}^*\|^2$, which yields $f(\mathbf{x}_N) - f^* \leq \frac{\|\mathbf{z}_0 - \mathbf{x}^*\|^2}{2A_N}$. Moreover, combining
 541 Lemma 3 and (32) in Lemma 4 leads to the second inequality in Proposition 1. \square

542 A.2 Additional Supporting Lemmas

543 A crucial part of our analysis is to bound the path length of the sequence $\{\mathbf{y}_k\}_{k=0}^N$. This is done in
 544 Lemma 8. To achieve this goal we first present the results in Lemmas 5 and 7 which provide the required
 545 ingredients for proving the claim in Lemma 8. In our first intermediate result, we establish uniform
 546 upper bounds for the error terms $\|\mathbf{z}_k - \mathbf{x}^*\|$ and $\|\mathbf{x}_k - \mathbf{x}^*\|$.

547 **Lemma 5.** Recall that $\sigma = \alpha_1 + \alpha_2$. For all $k \geq 0$, we have $\|\mathbf{z}_k - \mathbf{x}^*\| \leq \|\mathbf{z}_0 - \mathbf{x}^*\|$ and
 548 $\|\mathbf{x}_k - \mathbf{x}^*\| \leq \sqrt{\frac{2}{1-\sigma^2}} \|\mathbf{z}_0 - \mathbf{x}^*\|$.

549 *Proof.* To begin with, it follows from (21) in Proposition 2 that

$$\frac{1}{2} \|\mathbf{z}_k - \mathbf{x}^*\|^2 \leq A_k(f(\mathbf{x}_k) - f^*) + \frac{1}{2} \|\mathbf{z}_k - \mathbf{x}^*\|^2 \leq A_0(f(\mathbf{x}_0) - f^*) + \frac{1}{2} \|\mathbf{z}_0 - \mathbf{x}^*\|^2 = \frac{1}{2} \|\mathbf{z}_0 - \mathbf{x}^*\|^2.$$

550 Hence, we get $\|\mathbf{z}_k - \mathbf{x}^*\| \leq \|\mathbf{z}_0 - \mathbf{x}^*\|$ for any $k \geq 0$. To show the second inequality, we distinguish
 551 two cases and in both cases we will prove that

$$A_{k+1} \|\mathbf{x}_{k+1} - \mathbf{x}^*\|^2 \leq A_k \|\mathbf{x}_k - \mathbf{x}^*\|^2 + (A_{k+1} - A_k) \frac{2\sigma^2 a_k^2}{\eta_k^2} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2 + 2(A_{k+1} - A_k) \|\mathbf{z}_{k+1} - \mathbf{x}^*\|^2. \quad (35)$$

552 **Case I:** $\hat{\eta}_k = \eta_k$. Recall that in the proof of Proposition 2 we defined $\tilde{\mathbf{z}}_{k+1} = \hat{\mathbf{x}}_{k+1} + \frac{A_k}{a_k} (\hat{\mathbf{x}}_{k+1} - \mathbf{x}^*)$.

553 Since $\mathbf{x}_{k+1} = \hat{\mathbf{x}}_{k+1}$, we have $\mathbf{x}_{k+1} = \frac{A_k}{A_k + a_k} \mathbf{x}_k + \frac{a_k}{A_k + a_k} \tilde{\mathbf{z}}_{k+1}$ and by Jensen's inequality

$$\|\mathbf{x}_{k+1} - \mathbf{x}^*\|^2 \leq \frac{A_k}{A_k + a_k} \|\mathbf{x}_k - \mathbf{x}^*\|^2 + \frac{a_k}{A_k + a_k} \|\tilde{\mathbf{z}}_{k+1} - \mathbf{x}^*\|^2.$$

554 Furthermore, we have

$$\|\tilde{\mathbf{z}}_{k+1} - \mathbf{x}^*\|^2 \leq 2\|\tilde{\mathbf{z}}_{k+1} - \mathbf{z}_{k+1}\|^2 + 2\|\mathbf{z}_{k+1} - \mathbf{x}^*\|^2 \leq \frac{2\sigma^2 a_k^2}{\eta_k^2} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2 + 2\|\mathbf{z}_{k+1} - \mathbf{x}^*\|^2, \quad (36)$$

555 where we used (27) in the last inequality. By combining the above two inequalities, we obtain

$$(A_k + a_k) \|\mathbf{x}_{k+1} - \mathbf{x}^*\|^2 \leq A_k \|\mathbf{x}_k - \mathbf{x}^*\|^2 + a_k \frac{2\sigma^2 a_k^2}{\eta_k^2} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2 + 2a_k \|\mathbf{z}_{k+1} - \mathbf{x}^*\|^2, \quad (37)$$

556 which leads to (35) (note that $A_{k+1} = A_k + a_k$ in Case I).

557 **Case II:** Since $\mathbf{x}_{k+1} = \frac{(1-\gamma_k)A_k}{A_k + \gamma_k a_k} \mathbf{x}_k + \frac{\gamma_k(A_k + a_k)}{A_k + \gamma_k a_k} \hat{\mathbf{x}}_{k+1}$ and $\hat{\mathbf{x}}_{k+1} = \frac{A_k}{A_k + a_k} \mathbf{x}_k + \frac{a_k}{A_k + a_k} \tilde{\mathbf{z}}_{k+1}$, we
 558 have

$$\mathbf{x}_{k+1} = \frac{A_k}{A_k + \gamma_k a_k} \mathbf{x}_k + \frac{\gamma_k a_k}{A_k + \gamma_k a_k} \tilde{\mathbf{z}}_{k+1}.$$

559 Similarly, by Jensen's inequality we have

$$(A_k + \gamma_k a_k) \|\mathbf{x}_{k+1} - \mathbf{x}^*\|^2 \leq A_k \|\mathbf{x}_k - \mathbf{x}^*\|^2 + \gamma_k a_k \|\tilde{\mathbf{z}}_{k+1} - \mathbf{x}^*\|^2.$$

560 Combining this inequality with (36), we obtain

$$(A_k + \gamma_k a_k) \|\mathbf{x}_{k+1} - \mathbf{x}^*\|^2 \leq A_k \|\mathbf{x}_k - \mathbf{x}^*\|^2 + \gamma_k a_k \frac{2\sigma^2 a_k^2}{\eta_k^2} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2 + 2\gamma_k a_k \|\mathbf{z}_{k+1} - \mathbf{x}^*\|^2. \quad (38)$$

561 which leads to (35) (note that $A_{k+1} = A_k + \gamma_k a_k$ in Case II).

562 Now by summing (35) over $k = 0, \dots, N-1$, we get

$$A_N \|\mathbf{x}_N - \mathbf{x}^*\|^2 \leq \sum_{k=0}^{N-1} (A_{k+1} - A_k) \frac{2\sigma^2 a_k^2}{\eta_k^2} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2 + \sum_{k=0}^{N-1} 2(A_{k+1} - A_k) \|\mathbf{z}_{k+1} - \mathbf{x}^*\|^2 \quad (39)$$

$$\leq 2\sigma^2 \sum_{k=0}^{N-1} (A_{k+1} - A_k) \sum_{k=0}^{N-1} \frac{a_k^2}{\eta_k^2} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2 + 2\|\mathbf{z}_0 - \mathbf{x}^*\|^2 \sum_{k=0}^{N-1} (A_{k+1} - A_k) \quad (40)$$

$$\leq \frac{2\sigma^2}{1-\sigma^2} A_N \|\mathbf{z}_0 - \mathbf{x}^*\|^2 + 2A_N \|\mathbf{z}_0 - \mathbf{x}^*\|^2 \quad (41)$$

$$= \frac{2A_N}{1-\sigma^2} \|\mathbf{z}_0 - \mathbf{x}^*\|^2. \quad (42)$$

Hence, this implies that $\|\mathbf{x}_k - \mathbf{x}^*\|^2 \leq \frac{2}{1-\sigma^2} \|\mathbf{z}_0 - \mathbf{x}^*\|^2$ for any $k \geq 0$. \square

A key term appearing in several of our bounds is $\frac{a_{k+1}}{A_{k+1} + a_{k+1}}$. In the next lemma, we establish an upper bound for this ratio based on a factor of its previous value, for both cases of our algorithm.

Lemma 6. *Without loss of generality assume $\beta > 1/5$. In Case I we have $\frac{a_{k+1}}{A_{k+1} + a_{k+1}} \leq \frac{1}{\sqrt{\beta}} \frac{a_k}{A_k + a_k}$. Otherwise, in Case II we have $\frac{a_{k+1}}{A_{k+1} + a_{k+1}} \leq \frac{2\sqrt{\beta}}{\sqrt{\beta} + 1} \frac{a_k}{A_k + a_k}$.*

Proof. By the choice of a_k in (2) we have $\eta_k(A_k + a_k) = a_k^2$ for all $k \geq 0$. As a result, we have

$$\frac{a_k}{A_k + a_k} = \frac{\eta_k}{a_k} = \frac{2\eta_k}{\eta_k + \sqrt{\eta_k^2 + 4\eta_k A_k}} = \frac{2}{1 + \sqrt{1 + 4\frac{A_k}{\eta_k}}}, \quad (43)$$

and similarly

$$\frac{a_{k+1}}{A_{k+1} + a_{k+1}} = \frac{2}{1 + \sqrt{1 + 4\frac{A_{k+1}}{\eta_{k+1}}}}. \quad (44)$$

In Case I, we have $\eta_{k+1} = \eta_k/\beta$ and $A_{k+1} \geq A_k$. Hence, it implies that $A_{k+1}/\eta_{k+1} \geq \beta A_k/\eta_k$, which leads to

$$\frac{a_{k+1}}{A_{k+1} + a_{k+1}} \leq \frac{2}{1 + \sqrt{1 + \frac{4\beta A_k}{\eta_k}}} \leq \frac{2}{\sqrt{\beta} + \sqrt{\beta + \frac{4\beta A_k}{\eta_k}}} = \frac{1}{\sqrt{\beta}} \frac{2}{1 + \sqrt{1 + 4\frac{A_k}{\eta_k}}} = \frac{1}{\sqrt{\beta}} \frac{a_k}{A_k + a_k}.$$

where the second inequality follows from the fact that $\beta \leq 1$.

In Case II, we have $\eta_{k+1} = \hat{\eta}_k = \gamma_k \eta_k$ and $A_{k+1} = A_k + \gamma_k a_k$. Since we also have $a_k \geq \eta_k$ and $\gamma_k \leq \beta$, we obtain $A_{k+1}/\eta_{k+1} \geq A_k/(\gamma_k \eta_k) + 1 \geq A_k/(\beta \eta_k) + 1$. Hence,

$$\frac{a_{k+1}}{A_{k+1} + a_{k+1}} \leq \frac{2}{1 + \sqrt{5 + \frac{4A_k}{\beta \eta_k}}} \leq \frac{2}{1 + \frac{1}{\sqrt{\beta}} \sqrt{1 + \frac{4A_k}{\eta_k}}} \leq \frac{2\sqrt{\beta}}{\sqrt{\beta} + 1} \frac{2}{1 + \sqrt{1 + \frac{4A_k}{\eta_k}}} = \frac{2\sqrt{\beta}}{\sqrt{\beta} + 1} \frac{a_k}{A_k + a_k},$$

where we used $\beta > 1/5$ in the second inequality and the fact that $1 + \frac{1}{\sqrt{\beta}}x \geq \frac{\sqrt{\beta}+1}{2\sqrt{\beta}}(1+x)$ for $x \geq 1$ in the last inequality. \square

Next, as a corollary of Lemma 6, we establish an upper bound on the series $\sum_{k=0}^{N-1} \frac{a_k}{A_k + a_k}$. Moreover, we use this result to establish an upper bound for $\sum_{k=0}^{N-1} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|$.

Lemma 7. *We have*

$$\sum_{k=0}^{N-1} \frac{a_k}{A_k + a_k} \leq \frac{1 + 2\sqrt{\beta} - \beta}{\sqrt{\beta} - \beta} \left(1 + \log \frac{A_N}{A_1}\right). \quad (45)$$

Moreover,

$$\sum_{k=0}^{N-1} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| \leq \sqrt{\frac{1}{1-\sigma^2} \frac{1 + 2\sqrt{\beta} - \beta}{\sqrt{\beta} - \beta} \left(1 + \log \frac{A_N}{A_1}\right)} \|\mathbf{z}_0 - \mathbf{x}^*\|. \quad (46)$$

Proof. Given the initial values of A_k and a_k we have

$$\sum_{k=0}^{N-1} \frac{a_k}{A_k + a_k} = 1 + \sum_{k=1}^{N-1} \frac{a_k}{A_k + a_k} = 1 + \sum_{k \in \mathcal{B}, k \geq 1} \frac{a_k}{A_k + a_k} + \sum_{k \notin \mathcal{B}, k \geq 1} \frac{a_k}{A_k + a_k} \quad (47)$$

582 Note that using the result in Lemma 6

$$\sum_{k \in \mathcal{B}, k \geq 1} \frac{a_k}{A_k + a_k} \leq \sum_{k=0}^{N-2} \frac{a_{k+1}}{A_{k+1} + a_{k+1}} \quad (48)$$

$$= \sum_{k \notin \mathcal{B}, k \geq 0} \frac{a_{k+1}}{A_{k+1} + a_{k+1}} + \sum_{k \in \mathcal{B}, k \geq 0} \frac{a_{k+1}}{A_{k+1} + a_{k+1}} \quad (49)$$

$$\leq \sum_{k \notin \mathcal{B}, k \geq 0} \frac{1}{\sqrt{\beta}} \frac{a_k}{A_k + a_k} + \sum_{k \in \mathcal{B}, k \geq 0} \frac{2\sqrt{\beta}}{\sqrt{\beta} + 1} \frac{a_k}{A_k + a_k} \quad (50)$$

$$\leq \frac{1}{\sqrt{\beta}} + \sum_{k \notin \mathcal{B}, k \geq 1} \frac{1}{\sqrt{\beta}} \frac{a_k}{A_k + a_k} + \sum_{k \in \mathcal{B}, k \geq 1} \frac{2\sqrt{\beta}}{\sqrt{\beta} + 1} \frac{a_k}{A_k + a_k}. \quad (51)$$

583 Hence, if we move the last term in the above upper bound to the left hand side and rescale both sides
584 of the resulted inequality we obtain

$$\sum_{k \in \mathcal{B}, k \geq 1} \frac{a_k}{A_k + a_k} \leq \frac{1 + \sqrt{\beta}}{\sqrt{\beta} - \beta} \left(1 + \sum_{k \notin \mathcal{B}, k \geq 1} \frac{a_k}{A_k + a_k} \right). \quad (52)$$

585 Now, if we replace the above upper bound into (47) we obtain

$$\sum_{k=0}^{N-1} \frac{a_k}{A_k + a_k} \leq \frac{1 + 2\sqrt{\beta} - \beta}{\sqrt{\beta} - \beta} \left(1 + \sum_{k \notin \mathcal{B}, k \geq 1} \frac{a_k}{A_k + a_k} \right). \quad (53)$$

586 Moreover, note that for $k \notin \mathcal{B}$, we have $A_{k+1} = A_k + a_k$. Hence,

$$\begin{aligned} \sum_{k \notin \mathcal{B}, k \geq 1} \frac{a_k}{A_k + a_k} &= \sum_{k \notin \mathcal{B}, k \geq 1} \left(1 - \frac{A_k}{A_{k+1}} \right) \leq \sum_{k \notin \mathcal{B}, k \geq 1} (\log(A_{k+1}) - \log(A_k)) \\ &\leq \sum_{k=1}^{N-1} (\log(A_{k+1}) - \log(A_k)) = \log \frac{A_N}{A_1}. \end{aligned}$$

587 Now if we replace the above upper bound, i.e., $\log \frac{A_N}{A_1}$ with $\sum_{k \notin \mathcal{B}, k \geq 1} \frac{a_k}{A_k + a_k}$ into the expression in
588 the right-hand side of (53) we obtain the result in (45).

589 Next, note that by Cauchy-Schwarz inequality, we have

$$\sum_{k=0}^{N-1} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| \leq \sqrt{\sum_{k=0}^{N-1} \frac{\eta_k^2}{a_k^2} \sum_{k=0}^{N-1} \frac{a_k^2}{\eta_k^2} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2} \leq \sqrt{\frac{1}{1 - \sigma^2} \sum_{k=0}^{N-1} \frac{\eta_k^2}{a_k^2} \|\mathbf{z}_0 - \mathbf{x}^*\|},$$

590 where the last inequality follows from (22). Moreover, based on the expression for a_k in (2) and the
591 result in (45) that we just proved, we have

$$\sum_{k=0}^{N-1} \frac{\eta_k^2}{a_k^2} = \sum_{k=0}^{N-1} \frac{a_k^2}{(A_k + a_k)^2} \leq \sum_{k=0}^{N-1} \frac{a_k}{A_k + a_k} \leq \frac{1 + 2\sqrt{\beta} - \beta}{\sqrt{\beta} - \beta} \left(1 + \log \frac{A_N}{A_1} \right).$$

592 Combining the two inequalities above leads to (46). \square

593 Now we are ready to present and prove Lemma 8 which characterizes a bound on the path length of
594 the sequence $\{\mathbf{y}_k\}_{k=0}^N$

595 **Lemma 8.** Consider the iterates generated by Algorithm 1. Then for any N ,

$$\sum_{k=0}^{N-1} \|\mathbf{y}_{k+1} - \mathbf{y}_k\| \leq C_2 \left(1 + \log \frac{A_N}{A_1} \right) \|\mathbf{z}_0 - \mathbf{x}^*\|. \quad (54)$$

596 where

$$C_2 = 2\sqrt{\frac{1}{1 - \sigma^2} \frac{1 + 2\sqrt{\beta} - \beta}{\sqrt{\beta} - \beta}} + \frac{1}{\sqrt{\beta}} \left(1 + \sqrt{\frac{2}{1 - \sigma^2}} \right) \frac{1 + 2\sqrt{\beta} - \beta}{\sqrt{\beta} - \beta} \quad (55)$$

597 *Proof.* By the triangle inequality, we have

$$\|\mathbf{y}_k - \mathbf{y}_{k+1}\| \leq \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| + \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_{k+1}\|. \quad (56)$$

598 We again distinguish two cases.

599 **Case I:** $\hat{\eta}_k = \eta_k$. In this case $\hat{\mathbf{x}}_{k+1} = \mathbf{x}_{k+1}$ and $\mathbf{y}_{k+1} = \frac{A_{k+1}}{A_{k+1}+a_{k+1}}\mathbf{x}_{k+1} + \frac{a_{k+1}}{A_{k+1}+a_{k+1}}\mathbf{z}_{k+1}$, hence

$$\|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_{k+1}\| = \|\mathbf{x}_{k+1} - \mathbf{y}_{k+1}\| = \frac{a_{k+1}\|\mathbf{z}_{k+1} - \mathbf{x}_{k+1}\|}{A_{k+1} + a_{k+1}} \leq \frac{1}{\sqrt{\beta}} \left(1 + \sqrt{\frac{2}{1-\sigma^2}}\right) \frac{a_k\|\mathbf{z}_0 - \mathbf{x}^*\|}{A_k + a_k}, \quad (57)$$

600 where we used Lemma 6 and the fact that $\|\mathbf{z}_{k+1} - \mathbf{x}_{k+1}\| \leq \|\mathbf{z}_{k+1} - \mathbf{x}^*\| + \|\mathbf{x}_{k+1} - \mathbf{x}^*\| \leq$
 601 $(1 + \sqrt{\frac{2}{1-\sigma^2}})\|\mathbf{z}_0 - \mathbf{x}^*\|$ in the last inequality. Therefore, using (56) and the above bound we have

$$\|\mathbf{y}_k - \mathbf{y}_{k+1}\| \leq \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| + \frac{1}{\sqrt{\beta}} \left(1 + \sqrt{\frac{2}{1-\sigma^2}}\right) \frac{a_k}{A_k + a_k} \|\mathbf{z}_0 - \mathbf{x}^*\|. \quad (58)$$

602 **Case II:** $\hat{\eta}_k < \eta_k$. Since $\mathbf{x}_{k+1} = \frac{A_k}{A_k + \gamma_k a_k} \mathbf{x}_k + \frac{\gamma_k a_k}{A_k + \gamma_k a_k} \tilde{\mathbf{z}}_{k+1}$ and $\hat{\mathbf{x}}_{k+1} = \frac{A_k}{A_k + a_k} \mathbf{x}_k + \frac{a_k}{A_k + a_k} \tilde{\mathbf{z}}_{k+1}$,
 603 we get

$$\hat{\mathbf{x}}_{k+1} = \frac{A_k}{A_k + a_k} \left(\mathbf{x}_{k+1} + \frac{\gamma_k a_k}{A_k} (\mathbf{x}_{k+1} - \tilde{\mathbf{z}}_{k+1}) \right) + \frac{a_k}{A_k + a_k} \tilde{\mathbf{z}}_{k+1} = \frac{A_k + \gamma_k a_k}{A_k + a_k} \mathbf{x}_{k+1} + \frac{(1 - \gamma_k) a_k}{A_k + a_k} \tilde{\mathbf{z}}_{k+1}.$$

604 Thus, given the above equality and the expression for \mathbf{y}_{k+1} , i.e., $\mathbf{y}_{k+1} = \frac{A_{k+1}}{A_{k+1}+a_{k+1}}\mathbf{x}_{k+1} +$
 605 $\frac{a_{k+1}}{A_{k+1}+a_{k+1}}\mathbf{z}_{k+1}$, we have

$$\|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_{k+1}\| \leq \frac{(1 - \gamma_k) a_k}{A_k + a_k} \|\tilde{\mathbf{z}}_{k+1} - \mathbf{z}_{k+1}\| + \left| \frac{(1 - \gamma_k) a_k}{A_k + a_k} - \frac{a_{k+1}}{A_{k+1} + a_{k+1}} \right| \|\mathbf{z}_{k+1} - \mathbf{x}_{k+1}\|. \quad (59)$$

606 Moreover, based on the result in (27), we can upper bound $\|\tilde{\mathbf{z}}_{k+1} - \mathbf{z}_{k+1}\|$ by $\sigma \frac{a_k}{\eta_k} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|$
 607 which implies that

$$\frac{(1 - \gamma_k) a_k}{A_k + a_k} \|\tilde{\mathbf{z}}_{k+1} - \mathbf{z}_{k+1}\| \leq \sigma \frac{(1 - \gamma_k) a_k^2}{\eta_k (A_k + a_k)} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| = \sigma (1 - \gamma_k) \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| \leq \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|$$

608 where the equality holds due to the definition of a_k , and the last inequality holds as both γ_k and σ are
 609 in $(0, 1)$. On the other hand, note that

$$\frac{(1 - \gamma_k) a_k}{A_k + a_k} - \frac{a_{k+1}}{A_{k+1} + a_{k+1}} \leq \frac{(1 - \gamma_k) a_k}{A_k + a_k} \leq \frac{a_k}{A_k + a_k}, \quad (60)$$

$$\frac{a_{k+1}}{A_{k+1} + a_{k+1}} - \frac{(1 - \gamma_k) a_k}{A_k + a_k} \leq \frac{2\sqrt{\beta} a_k}{\sqrt{\beta} + 1 (A_k + a_k)} - \frac{(1 - \gamma_k) a_k}{A_k + a_k} \leq \frac{a_k}{A_k + a_k}. \quad (61)$$

610 where in the second bound we used the result in Lemma 6 and the fact that $\frac{2\sqrt{\beta} a_k}{\sqrt{\beta} + 1} < 1$. Hence, we
 611 get

$$\|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_{k+1}\| \leq \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| + \frac{a_k \|\mathbf{z}_{k+1} - \mathbf{x}_{k+1}\|}{A_k + a_k} \leq \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| + \left(1 + \sqrt{\frac{2}{1-\sigma^2}}\right) \frac{a_k \|\mathbf{z}_0 - \mathbf{x}^*\|}{A_k + a_k}, \quad (62)$$

612 where the last inequality follows from the fact $\|\mathbf{z}_{k+1} - \mathbf{x}_{k+1}\| \leq \|\mathbf{z}_{k+1} - \mathbf{x}^*\| + \|\mathbf{x}_{k+1} - \mathbf{x}^*\|$ and
 613 the bounds in Lemma 5. Now by applying the above upper bound into (56) we obtain that

$$\|\mathbf{y}_k - \mathbf{y}_{k+1}\| \leq 2\|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| + \left(1 + \sqrt{\frac{2}{1-\sigma^2}}\right) \frac{a_k}{A_k + a_k} \|\mathbf{z}_0 - \mathbf{x}^*\|. \quad (63)$$

614 Considering the upper bounds established for $\|\mathbf{y}_k - \mathbf{y}_{k+1}\|$ in case I (equation (58)) and case II
 615 (equation (63)), we can conclude that

$$\|\mathbf{y}_k - \mathbf{y}_{k+1}\| \leq 2\|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| + \frac{1}{\sqrt{\beta}} \left(1 + \sqrt{\frac{2}{1-\sigma^2}}\right) \frac{a_k}{A_k + a_k} \|\mathbf{z}_0 - \mathbf{x}^*\|. \quad (64)$$

616 Finally, Lemma 8 follows from summing (64) over $k = 0$ to $N - 1$ and the result of Lemma 7. \square

Subroutine 1 Backtracking line search

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1: Input: iterate  $\mathbf{y} \in \mathbb{R}^d$ , gradient  $\mathbf{g} \in \mathbb{R}^d$ , Hessian approximation  $\mathbf{B} \in \mathbb{S}_+^d$ , initial trial step size  $\eta > 0$ 
2: Parameters: line search parameters  $\beta \in (0, 1)$ ,  $\alpha_1 \geq 0$  and  $\alpha_2 > 0$  such that  $\alpha_1 + \alpha_2 < 1$ 
3: Set  $\hat{\eta} \leftarrow \eta$ 
4: Compute  $\mathbf{s}_+ \leftarrow \text{LinearSolver}(\mathbf{I} + \hat{\eta}\mathbf{B}, -\hat{\eta}\mathbf{g}; \alpha_1)$  and  $\hat{\mathbf{x}}_+ \leftarrow \mathbf{y} + \mathbf{s}_+$ 
5: while  $\|\hat{\mathbf{x}}_+ - \mathbf{y} + \hat{\eta}\nabla f(\hat{\mathbf{x}}_+)\|_2 > (\alpha_1 + \alpha_2)\|\hat{\mathbf{x}}_+ - \mathbf{y}\|_2$  do
6:   Set  $\tilde{\mathbf{x}}_+ \leftarrow \hat{\mathbf{x}}_+$  and  $\hat{\eta} \leftarrow \beta\hat{\eta}$ 
7:   Compute  $\mathbf{s}_+ \leftarrow \text{LinearSolver}(\mathbf{I} + \hat{\eta}\mathbf{B}, -\hat{\eta}\mathbf{g}; \alpha_1)$  and  $\hat{\mathbf{x}}_+ \leftarrow \mathbf{y} + \mathbf{s}_+$ 
8: end while
9: if  $\hat{\eta} = \eta$  then
10:   Return  $\hat{\eta}$  and  $\hat{\mathbf{x}}_+$ 
11: else
12:   Return  $\hat{\eta}$ ,  $\hat{\mathbf{x}}_+$  and  $\tilde{\mathbf{x}}_+$ 
13: end if

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B Line Search Subroutine

In this section, we provide further details on our line search subroutine in Section 3.1. For completeness, the pseudocode of our line search scheme is shown in Subroutine 1. In Section B.1 we prove that Subroutine 1 will always terminate in a finite number of steps. In Section B.2 we provide the proof of Lemma 1.

B.1 The line search subroutine terminates properly

Recall that in our line search scheme, we keep decreasing the step size $\hat{\eta}$ by a factor of β until we find a pair $(\hat{\eta}, \hat{\mathbf{x}}_+)$ satisfying (12) (also see Lines 5 and 6 in Subroutine 1). In the following lemma, we show that when the step size $\hat{\eta}$ is smaller than a certain threshold, then the pair $(\hat{\eta}, \hat{\mathbf{x}}_+)$ satisfies both conditions in (11) and (12), which further implies that Subroutine 1 will stop in a finite number of steps.

Lemma 9. Suppose Assumption 1 holds. If $\hat{\eta} < \frac{\alpha_2}{L_1 + \|\mathbf{B}\|_{\text{op}}}$ and $\hat{\mathbf{x}}_+$ is computed according to (13), then the pair $(\hat{\eta}, \hat{\mathbf{x}}_+)$ satisfies the conditions in (11) and (12).

Proof. By Definition 1, the pair $(\hat{\eta}, \hat{\mathbf{x}}_+)$ always satisfies the condition in (11) when $\hat{\mathbf{x}}_+$ is computed from (13). Hence, in the following we only need to prove that the condition in (12) also holds. Recall that $\mathbf{g} = \nabla f(\mathbf{y})$. By Assumption 1, the function f is L_1 -smooth and thus we have

$$\|\nabla f(\hat{\mathbf{x}}_+) - \mathbf{g}\| = \|\nabla f(\hat{\mathbf{x}}_+) - \nabla f(\mathbf{y})\| \leq L_1\|\hat{\mathbf{x}}_+ - \mathbf{y}\|.$$

Moreover, by using the triangle inequality, we get

$$\|\nabla f(\hat{\mathbf{x}}_+) - \mathbf{g} - \mathbf{B}(\hat{\mathbf{x}}_+ - \mathbf{y})\| \leq \|\nabla f(\hat{\mathbf{x}}_+) - \mathbf{g}\| + \|\mathbf{B}(\hat{\mathbf{x}}_+ - \mathbf{y})\| \leq (L_1 + \|\mathbf{B}\|_{\text{op}})\|\hat{\mathbf{x}}_+ - \mathbf{y}\|.$$

Hence, if $\hat{\eta} \leq \frac{\alpha_2}{L_1 + \|\mathbf{B}\|_{\text{op}}}$, we have

$$\hat{\eta}\|\nabla f(\hat{\mathbf{x}}_+) - \mathbf{g} - \mathbf{B}(\hat{\mathbf{x}}_+ - \mathbf{y})\| \leq \alpha_2\|\hat{\mathbf{x}}_+ - \mathbf{y}\|. \quad (65)$$

Finally, by using the triangle inequality, we can combine (11) and (65) to show that

$$\begin{aligned} \|\hat{\mathbf{x}}_+ - \mathbf{y} + \hat{\eta}\nabla f(\hat{\mathbf{x}}_+)\| &= \|\hat{\mathbf{x}}_+ - \mathbf{y} + \hat{\eta}(\mathbf{g} + \mathbf{B}(\hat{\mathbf{x}}_+ - \mathbf{y})) + \hat{\eta}(\nabla f(\hat{\mathbf{x}}_+) - \mathbf{g} - \mathbf{B}(\hat{\mathbf{x}}_+ - \mathbf{y}))\| \\ &\leq \|\hat{\mathbf{x}}_+ - \mathbf{y} + \hat{\eta}(\mathbf{g} + \mathbf{B}(\hat{\mathbf{x}}_+ - \mathbf{y}))\| + \|\hat{\eta}(\nabla f(\hat{\mathbf{x}}_+) - \mathbf{g} - \mathbf{B}(\hat{\mathbf{x}}_+ - \mathbf{y}))\| \\ &\leq \alpha_1\|\hat{\mathbf{x}}_+ - \mathbf{y}\| + \alpha_2\|\hat{\mathbf{x}}_+ - \mathbf{y}\| \\ &\leq (\alpha_1 + \alpha_2)\|\hat{\mathbf{x}}_+ - \mathbf{y}\|, \end{aligned}$$

which means the condition in (12) is satisfied. The proof is now complete. \square

B.2 Proof of Lemma 1

We follow a similar proof strategy as Lemma 3 in [34]. In the first case where $k \notin \mathcal{B}$, by definition, the line search subroutine accepts the initial step size η_k , i.e., $\hat{\eta}_k = \eta_k$. In the second case where $k \in \mathcal{B}$, the line search subroutine backtracks and returns the auxiliary iterate $\tilde{\mathbf{x}}_{k+1}$, which is computed from

641 (13) using the step size $\tilde{\eta}_k \triangleq \hat{\eta}_k / \beta$. Since the step size $\tilde{\eta}_k$ is rejected in our line search subroutine, it
 642 implies that the pair $(\tilde{\mathbf{x}}_{k+1}, \tilde{\eta}_k)$ does not satisfy (12), i.e.,

$$\|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k + \tilde{\eta}_k \nabla f(\tilde{\mathbf{x}}_{k+1})\| > (\alpha_1 + \alpha_2) \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\|. \quad (66)$$

643 Moreover, since we compute $\tilde{\mathbf{x}}_{k+1}$ from (13) using step size $\tilde{\eta}_k$, the pair $(\tilde{\eta}_k, \tilde{\mathbf{x}}_{k+1})$ also satisfies the
 644 condition in (11), which means

$$\|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k + \tilde{\eta}_k (\nabla f(\mathbf{y}_k) + \mathbf{B}_k(\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k))\| \leq \alpha_1 \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\|. \quad (67)$$

645 Hence, by using the triangle inequality, we can combine (66) and (67) to get

$$\begin{aligned} & \tilde{\eta}_k \|\nabla f(\tilde{\mathbf{x}}_{k+1}) - \nabla f(\mathbf{y}_k) - \mathbf{B}_k(\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k)\| \\ & \geq \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k + \tilde{\eta}_k \nabla f(\tilde{\mathbf{x}}_{k+1})\| - \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k + \tilde{\eta}_k (\nabla f(\mathbf{y}_k) + \mathbf{B}_k(\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k))\| \\ & > (\alpha_1 + \alpha_2) \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\| - \alpha_1 \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\| \\ & = \alpha_2 \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\|, \end{aligned}$$

646 which implies that

$$\hat{\eta}_k = \beta \tilde{\eta}_k > \frac{\alpha_2 \beta \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\|}{\|\nabla f(\tilde{\mathbf{x}}_{k+1}) - \nabla f(\mathbf{y}_k) - \mathbf{B}_k(\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k)\|}.$$

647 This proves the first inequality in (14).

648 To show the second inequality in (14), first note that $\tilde{\mathbf{x}}_{k+1}$ and $\hat{\mathbf{x}}_{k+1}$ are the inexact solutions of the
 649 linear system of equations

$$(\mathbf{I} + \tilde{\eta}_k \mathbf{B}_k)(\mathbf{x} - \mathbf{y}_k) = -\tilde{\eta}_k \mathbf{g}_k \quad \text{and} \quad (\mathbf{I} + \hat{\eta}_k \mathbf{B}_k)(\mathbf{x} - \mathbf{y}_k) = -\hat{\eta}_k \mathbf{g}_k,$$

650 respectively. Let $\tilde{\mathbf{x}}_{k+1}^*$ and $\hat{\mathbf{x}}_{k+1}^*$ be the exact solutions of the above linear systems, that is, $\tilde{\mathbf{x}}_{k+1}^* =$
 651 $\mathbf{y}_k - \tilde{\eta}_k (\mathbf{I} + \tilde{\eta}_k \mathbf{B}_k)^{-1} \mathbf{g}_k$ and $\hat{\mathbf{x}}_{k+1}^* = \mathbf{y}_k - \hat{\eta}_k (\mathbf{I} + \hat{\eta}_k \mathbf{B}_k)^{-1} \mathbf{g}_k$. We first establish the following
 652 inequality between $\|\tilde{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\|$ and $\|\hat{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\|$:

$$\|\tilde{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\| \leq \frac{1}{\beta} \|\hat{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\|. \quad (68)$$

653 This follows from

$$\|\tilde{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\| = \|\tilde{\eta}_k (\mathbf{I} + \tilde{\eta}_k \mathbf{B}_k)^{-1} \mathbf{g}_k\| \leq \tilde{\eta}_k \|(\mathbf{I} + \hat{\eta}_k \mathbf{B}_k)^{-1} \mathbf{g}_k\| = \frac{\tilde{\eta}_k}{\hat{\eta}_k} \|\hat{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\| = \frac{1}{\beta} \|\hat{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\|,$$

654 where we used the fact that $(\mathbf{I} + \tilde{\eta}_k \mathbf{B}_k)^{-1} \preceq (\mathbf{I} + \hat{\eta}_k \mathbf{B}_k)^{-1}$ in the first inequality. Furthermore, we
 655 can show that

$$(1 - \alpha_1) \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| \leq \|\hat{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\| \leq (1 + \alpha_1) \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|, \quad (69)$$

$$(1 - \alpha_1) \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\| \leq \|\tilde{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\| \leq (1 + \alpha_1) \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\|. \quad (70)$$

656 We will only prove (69) in the following, as (70) can be proved similarly. Note that since $(\hat{\eta}_k, \hat{\mathbf{x}}_{k+1})$
 657 satisfies the condition in (11), we can write

$$\|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k + \hat{\eta}_k (\mathbf{g}_k + \mathbf{B}_k(\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k))\| = \|(\mathbf{I} + \hat{\eta}_k \mathbf{B}_k)(\hat{\mathbf{x}}_{k+1} - \hat{\mathbf{x}}_{k+1}^*)\| \leq \alpha_1 \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|.$$

658 Moreover, since $\mathbf{B}_k \succeq 0$, we have $\|\hat{\mathbf{x}}_{k+1} - \hat{\mathbf{x}}_{k+1}^*\| \leq \|(\mathbf{I} + \hat{\eta}_k \mathbf{B}_k)(\hat{\mathbf{x}}_{k+1} - \hat{\mathbf{x}}_{k+1}^*)\| \leq \alpha_1 \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|$.
 659 Thus, by the triangle inequality, we obtain

$$\begin{aligned} \|\hat{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\| & \leq \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| + \|\hat{\mathbf{x}}_{k+1}^* - \hat{\mathbf{x}}_{k+1}\| \leq (1 + \alpha_1) \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|. \\ \|\hat{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\| & \geq \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\| - \|\hat{\mathbf{x}}_{k+1}^* - \hat{\mathbf{x}}_{k+1}\| \geq (1 - \alpha_1) \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|. \end{aligned}$$

660 which proves (69). Finally, by combining (68), (69) and (70), we conclude that

$$\|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\| \leq \frac{1}{1 - \alpha_1} \|\tilde{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\| \leq \frac{1}{(1 - \alpha_1)\beta} \|\hat{\mathbf{x}}_{k+1}^* - \mathbf{y}_k\| \leq \frac{1 + \alpha_1}{(1 - \alpha_1)\beta} \|\hat{\mathbf{x}}_{k+1} - \mathbf{y}_k\|.$$

661 This completes the proof.

662 C Hessian Approximation Update

663 In this section, we fully describe our Hessian approximation update in Section 3.2. We first prove
 664 Lemma 2 in Section C.1.

665 C.1 Proof of Lemma 2

666 We decompose the sum $\sum_{k=0}^{N-1} \frac{1}{\hat{\eta}_k^2}$ as

$$\sum_{k=0}^{N-1} \frac{1}{\hat{\eta}_k^2} = \frac{1}{\hat{\eta}_0^2} + \sum_{1 \leq k \leq N-1, k \in \mathcal{B}} \frac{1}{\hat{\eta}_k^2} + \sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \frac{1}{\hat{\eta}_k^2} \quad (71)$$

667 Recall that we have $\hat{\eta}_k = \eta_k$ for $k \notin \mathcal{B}$. Hence, we can further bound the last term by

$$\begin{aligned} \sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \frac{1}{\hat{\eta}_k^2} &= \sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \frac{1}{\eta_k^2} \leq \sum_{k=1}^{N-1} \frac{1}{\eta_k^2} \\ &= \frac{1}{\eta_1^2} + \sum_{1 \leq k \leq N-2, k \in \mathcal{B}} \frac{1}{\eta_{k+1}^2} + \sum_{1 \leq k \leq N-2, k \notin \mathcal{B}} \frac{1}{\eta_{k+1}^2}. \end{aligned}$$

668 Recall that we have $\eta_{k+1} = \hat{\eta}_k$ if $k \in \mathcal{B}$ and $\eta_{k+1} = \hat{\eta}_k/\beta$ otherwise. Hence, we further have

$$\begin{aligned} \sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \frac{1}{\hat{\eta}_k^2} &\leq \frac{1}{\eta_1^2} + \sum_{1 \leq k \leq N-2, k \in \mathcal{B}} \frac{1}{\eta_{k+1}^2} + \sum_{1 \leq k \leq N-2, k \notin \mathcal{B}} \frac{1}{\eta_{k+1}^2} \\ &= \frac{1}{\eta_1^2} + \sum_{1 \leq k \leq N-2, k \in \mathcal{B}} \frac{1}{\hat{\eta}_k^2} + \sum_{1 \leq k \leq N-2, k \notin \mathcal{B}} \frac{\beta^2}{\hat{\eta}_k^2} \\ &\leq \frac{1}{\eta_1^2} + \sum_{1 \leq k \leq N-1, k \in \mathcal{B}} \frac{1}{\hat{\eta}_k^2} + \sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \frac{\beta^2}{\hat{\eta}_k^2}. \end{aligned}$$

669 By moving the last term to the left-hand side and dividing both sides by $1 - \beta^2$, we obtain

$$\sum_{1 \leq k \leq N-1, k \notin \mathcal{B}} \frac{1}{\hat{\eta}_k^2} \leq \frac{1}{1 - \beta^2} \left(\frac{1}{\eta_1^2} + \sum_{1 \leq k \leq N-1, k \in \mathcal{B}} \frac{1}{\hat{\eta}_k^2} \right). \quad (72)$$

670 Furthermore, since $\eta_1 \geq \hat{\eta}_0$, we have $\frac{1}{\eta_1^2} \leq \frac{1}{\hat{\eta}_0^2}$. Hence, by combining (71) and (72), we get

$$\sum_{k=0}^{N-1} \frac{1}{\hat{\eta}_k^2} \leq \frac{2 - \beta^2}{1 - \beta^2} \left(\frac{1}{\hat{\eta}_0^2} + \sum_{1 \leq k \leq N-1, k \in \mathcal{B}} \frac{1}{\hat{\eta}_k^2} \right) \leq \frac{2 - \beta^2}{(1 - \beta^2)\sigma_0^2} + \frac{2 - \beta^2}{1 - \beta^2} \sum_{0 \leq k \leq N-1, k \in \mathcal{B}} \frac{1}{\hat{\eta}_k^2}, \quad (73)$$

671 where in the last inequality we used the fact that $\hat{\eta}_k = \sigma_0$ if $0 \notin \mathcal{B}$. Finally, (16) follows from
 672 Lemma 1 and (73).

673 C.2 The computational cost of Euclidean projection

674 Recall that $\mathcal{Z} \triangleq \{\mathbf{B} \in \mathbb{S}_+^d : 0 \preceq \mathbf{B} \preceq L_1 \mathbf{I}\}$. As described in [34] Section D.1], the Euclidean
 675 projection on \mathcal{Z} has a closed form solution. Specifically, Given the input $\mathbf{A} \in \mathbb{S}^d$, we first need
 676 to perform the eigendecomposition $\mathbf{A} = \mathbf{V} \mathbf{\Lambda} \mathbf{V}^\top$, where \mathbf{V} is an orthogonal matrix and $\mathbf{\Lambda} =$
 677 $\text{diag}(\lambda_1, \dots, \lambda_d)$ is a diagonal matrix. Then the Euclidean projection of \mathbf{A} onto \mathcal{Z} is given by
 678 $\mathbf{V} \hat{\mathbf{\Lambda}} \mathbf{V}^\top$, where $\hat{\mathbf{\Lambda}}$ is a diagonal matrix with the diagonals being $\hat{\lambda}_k = \min\{L_1, \max\{0, \lambda_k\}\}$ for
 679 $1 \leq k \leq d$. Since the eigendecomposition requires $\mathcal{O}(d^3)$ arithmetic operations in general, the cost
 680 of computing the Euclidean projection can be prohibitive.

Algorithm 2 Projection-Free Online Learning

```

1: Input: Initial point  $\mathbf{w}_0 \in \mathcal{B}_R(0)$ , step size  $\rho > 0$ ,  $\delta > 0$ 
2: for  $t = 0, 1, \dots, T-1$  do
3:   Query the oracle  $(\gamma_t, \mathbf{s}_t) \leftarrow \text{SEP}(\mathbf{w}_t; \delta_t)$ 
4:   if  $\gamma_t \leq 1$  then   # Case I: we have  $\mathbf{w}_t \in \mathcal{C}$ 
5:     Set  $\mathbf{x}_t \leftarrow \mathbf{w}_t$  and play the action  $\mathbf{x}_t$ 
6:     Receive the loss  $\ell_t(\mathbf{x}_t)$  and the gradient  $\mathbf{g}_t = \nabla \ell_t(\mathbf{x}_t)$ 
7:     Set  $\tilde{\mathbf{g}}_t \leftarrow \mathbf{g}_t$ 
8:   else   # Case II: we have  $\mathbf{w}_t/\gamma_t \in \mathcal{C}$ 
9:     Set  $\mathbf{x}_t \leftarrow \mathbf{w}_t/\gamma_t$  and play the action  $\mathbf{x}_t$ 
10:    Receive the loss  $\ell_t(\mathbf{x}_t)$  and the gradient  $\mathbf{g}_t = \nabla \ell_t(\mathbf{x}_t)$ 
11:    Set  $\tilde{\mathbf{g}}_t \leftarrow \mathbf{g}_t + \max\{0, -\langle \mathbf{g}_t, \mathbf{x}_t \rangle\} \mathbf{s}_t$ 
12:   end if
13:   Update  $\mathbf{w}_{t+1} \leftarrow \frac{R}{\max\{\|\mathbf{w}_t - \rho \tilde{\mathbf{g}}_t\|_2, R\}} (\mathbf{w}_t - \rho \tilde{\mathbf{g}}_t)$    # Euclidean projection onto  $\mathcal{B}_R(0)$ 
14: end for

```

C.3 Online Learning with an Approximate Separation Oracle

To set the stage for our Hessian approximation matrix update, we first describe a projection-free online learning algorithm in a general setup. Specifically, the online learning protocol is as follows: For rounds $t = 0, 1, \dots, T-1$, a learner chooses an action $\mathbf{x}_t \in \mathcal{C}$ from a convex set \mathcal{C} and then observes a loss function $\ell_t : \mathbb{R}^n \rightarrow \mathbb{R}$. We measure the performance of an online learning algorithm by the dynamic regret [39, 42] defined by

$$\text{D-Reg}_T(\mathbf{u}_1, \dots, \mathbf{u}_{T-1}) \triangleq \sum_{t=0}^{T-1} \ell_t(\mathbf{x}_t) - \sum_{t=0}^{T-1} \ell_t(\mathbf{u}_t),$$

where $\{\mathbf{u}_t\}_{t=1}^T$ is a sequence of comparators. Moreover, we assume that the convex set \mathcal{C} is contained in the Euclidean ball $\mathcal{B}_R(0)$ for some $R > 0$, and we assume $0 \in \mathcal{C}$ without loss of generality.

Most existing online learning algorithms are projection-based, that is, they require computing the Euclidean projection on the action set \mathcal{C} . However, as we have seen in Section C.2, computing the projection is computationally costly in our setting. Inspired by the work in [40], we will describe an online learning algorithm that relies on an approximate separation oracle defined in Definition 3.

Definition 3. The oracle $\text{SEP}(\mathbf{w}; \delta)$ takes $\mathbf{w} \in \mathcal{B}_R(0)$ and $\delta > 0$ as input and returns a scalar $\gamma > 0$ and a vector $\mathbf{s} \in \mathbb{R}^n$ with one of the following possible outcomes:

- Case I: $\gamma \leq 1$ which implies that $\mathbf{w} \in \mathcal{C}$;
- Case II: $\gamma > 1$ which implies that $\mathbf{w}/\gamma \in \mathcal{C}$ and $\langle \mathbf{s}, \mathbf{w} - \mathbf{x} \rangle \geq \gamma - 1 - \delta \quad \forall \mathbf{x} \in \mathcal{C}$.

In summary, the oracle $\text{SEP}(\mathbf{w}; \delta)$ has two possible outcomes: it either certifies that \mathbf{w} is feasible, i.e., $\mathbf{w} \in \mathcal{C}$, or it produces a scaled version of \mathbf{w} that is in \mathcal{C} and gives an approximate separating hyperplane between \mathbf{w} and the set \mathcal{C} .

The full algorithm is shown in Algorithm 2. The key idea here is to introduce surrogate loss functions $\tilde{\ell}_t(\mathbf{w}) = \langle \tilde{\mathbf{g}}_t, \mathbf{w} \rangle$ on the larger set $\mathcal{B}_R(0)$ for $0 \leq t \leq T-1$, where $\tilde{\mathbf{g}}_t$ is the surrogate gradient to be defined later. On a high level, we will run online projected gradient descent with $\tilde{\ell}_t(\mathbf{w})$ to update the auxiliary iterates $\{\mathbf{w}_t\}_{t \geq 0}$ (note that the projection on $\mathcal{B}_R(0)$ is easy to compute), and then produce the actions $\{\mathbf{x}_t\}_{t \geq 0}$ for the original problem by calling the $\text{SEP}(\mathbf{w}_t; \delta)$ oracle in Definition 3. The follow lemma shows that the immediate regret $\tilde{\ell}_t(\mathbf{w}_t) - \tilde{\ell}_t(\mathbf{x})$ can serve as an upper bound on $\ell_t(\mathbf{x}_t) - \ell_t(\mathbf{x})$ for any $\mathbf{x} \in \mathcal{C}$.

Lemma 10. Let $\{\mathbf{x}_t\}_{t=0}^{T-1}$ be the iterates generated by Algorithm 2. Then we have $\mathbf{x}_t \in \mathcal{C}$ for $t = 0, 1, \dots, T-1$. Also, for any $\mathbf{x} \in \mathcal{C}$, we have

$$\langle \mathbf{g}_t, \mathbf{x}_t - \mathbf{x} \rangle \leq \langle \tilde{\mathbf{g}}_t, \mathbf{w}_t - \mathbf{x} \rangle + \max\{0, -\langle \mathbf{g}_t, \mathbf{x}_t \rangle\} \delta_t \quad (74)$$

$$\leq \frac{1}{2\rho} \|\mathbf{w}_t - \mathbf{x}\|_2^2 - \frac{1}{2\rho} \|\mathbf{w}_{t+1} - \mathbf{x}\|_2^2 + \frac{\rho}{2} \|\tilde{\mathbf{g}}_t\|_2^2 + \max\{0, -\langle \mathbf{g}_t, \mathbf{x}_t \rangle\} \delta_t, \quad (75)$$

and

$$\|\tilde{\mathbf{g}}_t\| \leq \|\mathbf{g}_t\| + |\langle \mathbf{g}_t, \mathbf{x}_t \rangle| \|\mathbf{s}_t\|. \quad (76)$$

Subroutine 2 Online Learning Guided Hessian Approximation Update

```

1: Input: Initial matrix  $\mathbf{B}_0 \in \mathbb{S}^d$  s.t.  $0 \preceq \mathbf{B}_0 \preceq L_1 \mathbf{I}$ , step size  $\rho > 0$ ,  $\delta > 0$ ,  $\{q_t\}_{t=1}^{T-1}$ 
2: Initialize: set  $\mathbf{W}_0 \leftarrow \frac{2}{L_1}(\mathbf{B}_0 - \frac{L_1}{2}\mathbf{I})$ ,  $\mathbf{G}_0 \leftarrow \frac{2}{L_1}\nabla\ell_0(\mathbf{B}_0)$  and  $\tilde{\mathbf{G}}_0 \leftarrow \mathbf{G}_0$ 
3: for  $t = 1, \dots, T-1$  do
4:   Query the oracle  $(\gamma_t, \mathbf{S}_t) \leftarrow \text{SEP}(\mathbf{W}_t; \delta_t, q_t)$ 
5:   if  $\gamma_t \leq 1$  then # Case I
6:     Set  $\hat{\mathbf{B}}_t \leftarrow \mathbf{W}_t$  and  $\mathbf{B}_t \leftarrow \frac{L_1}{2}\hat{\mathbf{B}}_t + \frac{L_1}{2}\mathbf{I}$ 
7:     Set  $\mathbf{G}_t \leftarrow \frac{2}{L_1}\nabla\ell_t(\mathbf{B}_t)$  and  $\tilde{\mathbf{G}}_t \leftarrow \mathbf{G}_t$ 
8:   else # Case II
9:     Set  $\hat{\mathbf{B}}_t \leftarrow \mathbf{W}_t/\gamma_t$  and  $\mathbf{B}_t \leftarrow \frac{L_1}{2}\hat{\mathbf{B}}_t + \frac{L_1}{2}\mathbf{I}$ 
10:    Set  $\mathbf{G}_t \leftarrow \frac{2}{L_1}\nabla\ell_t(\mathbf{B}_t)$  and  $\tilde{\mathbf{G}}_t \leftarrow \mathbf{G}_t + \max\{0, -\langle \mathbf{G}_t, \mathbf{B}_t \rangle\} \mathbf{S}_t$ 
11:   end if
12:   Update  $\mathbf{W}_{t+1} \leftarrow \frac{\sqrt{d}}{\max\{\sqrt{d}, \|\mathbf{W}_t - \rho\tilde{\mathbf{G}}_t\|_F\}} (\mathbf{W}_t - \rho\tilde{\mathbf{G}}_t)$  # Euclidean projection onto  $\mathcal{B}_{\sqrt{d}}(0)$ 
13: end for

```

710 *Proof.* By the definition of SEP in Definition 3, we can see that $\mathbf{x}_t \in \mathcal{C}$ for all $t = 1, \dots, T$. We
 711 now show that both (74) and (76) hold. We distinguish two cases depending on the outcomes of
 712 $\text{SEP}(\mathbf{w}_t; \delta_t)$.

- 713 • If $\gamma_t \leq 1$, then we have $\mathbf{x}_t = \mathbf{w}_t$ and $\tilde{\mathbf{g}}_t = \mathbf{g}_t$. In this case, (74) and (76) trivially hold.
- 714 • If $\gamma_t > 1$, then $\mathbf{x}_t = \mathbf{w}_t/\gamma_t$ and $\tilde{\mathbf{g}}_t = \mathbf{g}_t + \max\{0, -\langle \mathbf{g}_t, \mathbf{x}_t \rangle\} \mathbf{s}_t$. We can then write

$$\begin{aligned}
 \langle \tilde{\mathbf{g}}_t, \mathbf{w}_t - \mathbf{x} \rangle &= \langle \mathbf{g}_t + \max\{0, -\langle \mathbf{g}_t, \mathbf{x}_t \rangle\} \mathbf{s}_t, \mathbf{w}_t - \mathbf{x} \rangle \\
 &= \langle \mathbf{g}_t, \gamma_t \mathbf{x}_t - \mathbf{x} \rangle + \max\{0, -\langle \mathbf{g}_t, \mathbf{x}_t \rangle\} \langle \mathbf{s}_t, \mathbf{w}_t - \mathbf{x} \rangle \\
 &\geq \langle \mathbf{g}_t, \mathbf{x}_t - \mathbf{x} \rangle + (\gamma_t - 1) \langle \mathbf{g}_t, \mathbf{x}_t \rangle + \max\{0, -\langle \mathbf{g}_t, \mathbf{x}_t \rangle\} (\gamma_t - 1 - \delta_t) \\
 &= \langle \mathbf{g}_t, \mathbf{x}_t - \mathbf{x} \rangle - \max\{0, -\langle \mathbf{g}_t, \mathbf{x}_t \rangle\} \delta_t + (\gamma_t - 1) \max\{0, \langle \mathbf{g}_t, \mathbf{x}_t \rangle\} \\
 &\geq \langle \mathbf{g}_t, \mathbf{x}_t - \mathbf{x} \rangle - \max\{0, -\langle \mathbf{g}_t, \mathbf{x}_t \rangle\} \delta_t,
 \end{aligned}$$

715 which leads to (74) after rearranging. Also, by the triangle inequality we obtain

$$\|\tilde{\mathbf{g}}_t\| \leq \|\mathbf{g}_t\| + \max\{0, -\langle \mathbf{g}_t, \mathbf{x}_t \rangle\} \|\mathbf{s}_t\| \leq \|\mathbf{g}_t\| + |\langle \mathbf{g}_t, \mathbf{x}_t \rangle| \|\mathbf{s}_t\|,$$

716 which proves (76).

717 Finally, from the update rule of \mathbf{w}_{t+1} , for any $\mathbf{x} \in \mathcal{C} \subset \mathcal{B}_R(0)$ we have $\langle \mathbf{w}_t - \rho\tilde{\mathbf{g}}_t - \mathbf{w}_{t+1}, \mathbf{w}_{t+1} - \mathbf{x} \rangle \geq 0$, which further implies that

$$\langle \tilde{\mathbf{g}}_t, \mathbf{w}_t - \mathbf{x} \rangle \leq \langle \tilde{\mathbf{g}}_t, \mathbf{w}_t - \mathbf{w}_{t+1} \rangle + \frac{1}{\rho} \langle \mathbf{w}_t - \mathbf{w}_{t+1}, \mathbf{w}_{t+1} - \mathbf{x} \rangle \quad (77)$$

$$= \langle \tilde{\mathbf{g}}_t, \mathbf{w}_t - \mathbf{w}_{t+1} \rangle + \frac{1}{2\rho} \|\mathbf{w}_t - \mathbf{x}\|_2^2 - \frac{1}{2\rho} \|\mathbf{w}_{t+1} - \mathbf{x}\|_2^2 - \frac{1}{2\rho} \|\mathbf{w}_t - \mathbf{w}_{t+1}\|_2^2 \quad (78)$$

$$\leq \frac{1}{2\rho} \|\mathbf{w}_t - \mathbf{x}\|_2^2 - \frac{1}{2\rho} \|\mathbf{w}_{t+1} - \mathbf{x}\|_2^2 + \frac{\rho}{2} \|\tilde{\mathbf{g}}_t\|_2^2. \quad (79)$$

719 Combining (74) and (79) leads to (75). \square

720 C.4 Projection-free Hessian Approximation Update

721 Now we are ready to describe our Hessian approximation matrix update, which is an specific
 722 instantiation of the general projection-free online learning algorithm described in Algorithm 2. The
 723 full algorithm is described in Subroutine 2.

724 Recall that $\mathcal{Z} = \{\mathbf{B} \in \mathbb{S}_+^d : 0 \preceq \mathbf{B} \preceq L_1 \mathbf{I}\}$ in our online learning problem in Section 3.2. Since
 725 the projection-free scheme in Subroutine 2 requires the set \mathcal{C} to contain the origin, we consider the
 726 transform $\hat{\mathbf{B}} \triangleq \frac{2}{L_1}(\mathbf{B} - \frac{L_1}{2}\mathbf{I})$ and define $\hat{\mathcal{Z}} \triangleq \{\hat{\mathbf{B}} \in \mathbb{S}^d : -\mathbf{I} \preceq \hat{\mathbf{B}} \preceq \mathbf{I}\} = \{\mathbf{B} \in \mathbb{S}^d : \|\hat{\mathbf{B}}\|_{\text{op}} \leq 1\}$.

727 We note that $0 \in \hat{\mathcal{Z}}$ and $\hat{\mathcal{Z}} \subset \mathcal{B}_{\sqrt{d}}(0) = \{\mathbf{W} \in \mathbb{S}^d : \|\mathbf{W}\|_F \leq \sqrt{d}\}$. Moreover, we can see that
 728 the approximate separation oracle $\text{SEP}(\mathbf{W}; \delta, q)$ defined in Definition 2 corresponds to the oracle in
 729 Definition 3. We defer the specific implementation details to Section E.2.

D Proof of Theorem 1

Regarding the choices of the hyper-parameters, we consider Algorithm 1 with the line search scheme in Subroutine 1, where $\alpha_1, \alpha_2 \in (0, 1)$ with $\alpha_1 + \alpha_2 < 1$ and $\beta \in (0, 1)$, and with the Hessian approximation update in Subroutine 2 where $\rho = \frac{1}{128}$, $q_t = p/2.5(t+1) \log^2(t+1)$ for $t \geq 1$, and $\delta_t = 1/(\sqrt{t} + 2 \ln(t+2))$ for $t \geq 0$. In the following, we first provide a proof sketch of Theorem 1. The complete proofs of the lemmas shown below will be provided in the subsequent sections.

Proof Sketch. To begin with, throughout the proof, we assume that every call of the SEP oracle in Definition 2 is successful during the execution of Algorithm 1. Indeed, by using the union bound, we can bound the failure probability by $\sum_{t=1}^{T-1} q_t \leq \frac{p}{2.5} \sum_{t=2}^{\infty} \frac{1}{t \log^2 t} \leq p$. In particular, we note that Subroutine 2 ensures that $0 \preceq \mathbf{B}_k \preceq L_1 \mathbf{I}$ for any $k \geq 0$.

We first prove Part (a) of Theorem 1 which relies on the following lemma.

Lemma 11. For $k \in \mathcal{B}$, we have $\ell_k(\mathbf{B}_k) \triangleq \frac{\|\mathbf{w}_k - \mathbf{B}_k \mathbf{s}_k\|^2}{\|\mathbf{s}_k\|^2} \leq L_1^2$.

We combine Lemma 2 and Lemma 11 to derive

$$\sum_{k=0}^{N-1} \frac{1}{\hat{\eta}_k^2} \leq \frac{2 - \beta^2}{(1 - \beta^2)\sigma_0^2} + \frac{2 - \beta^2}{(1 - \beta^2)\alpha_2^2\beta^2} \sum_{k \in \mathcal{B}} \frac{\|\mathbf{w}_k - \mathbf{B}_k \mathbf{s}_k\|^2}{\|\mathbf{s}_k\|^2} \leq \frac{2 - \beta^2}{(1 - \beta^2)\sigma_0^2} + \frac{(2 - \beta^2)L_1^2}{(1 - \beta^2)\alpha_2^2\beta^2} N.$$

By further using (15) and the elementary inequality that $\sqrt{a+b} \leq \sqrt{a} + \sqrt{b}$, we obtain

$$f(\mathbf{x}_N) - f(\mathbf{x}^*) \leq \frac{C_4 L_1 \|\mathbf{z}_0 - \mathbf{x}^*\|^2}{N^2} + \frac{C_5 \|\mathbf{z}_0 - \mathbf{x}^*\|^2}{\sigma_0 N^{2.5}}, \quad (80)$$

where $C_4 = C_1 \sqrt{\frac{2-\beta^2}{(1-\beta^2)\sigma_0^2} + \frac{(2-\beta^2)}{(1-\beta^2)\alpha_2^2\beta^2}}$ and $C_5 = C_1 \sqrt{\frac{2-\beta^2}{(1-\beta^2)\sigma_0^2}}$

Next, we divide the proof of Part (b) of Theorem 1 into the following steps.

Step 1: We first use regret analysis to control the cumulative loss $\sum_{t=0}^{T-1} \ell_t(\mathbf{B}_t)$ incurred by our online learning algorithm in Subroutine 2. In particular, we prove a dynamic regret bound, where we compare the cumulative loss of our algorithm against the one achieved by the sequence $\{\mathbf{H}_t\}_{t=0}^{T-1}$.

Lemma 12. We have

$$\sum_{t=0}^{T-1} \ell_t(\mathbf{B}_t) \leq 256 \|\mathbf{B}_0 - \mathbf{H}_0\|_F^2 + 4 \sum_{t=0}^{T-1} \ell_t(\mathbf{H}_t) + 2L_1^2 \sum_{t=0}^{T-1} \delta_t^2 + 512L_1 \sqrt{d} \sum_{t=0}^{T-1} \|\mathbf{H}_{t+1} - \mathbf{H}_t\|_F,$$

where $\mathbf{H}_t \triangleq \nabla^2 f(\mathbf{y}_t)$.

Step 2: In light of Lemma 12, it suffices to upper bound the cumulative loss $\sum_{t=0}^{T-1} \ell_t(\mathbf{H}_t)$ and the path-length $\sum_{t=0}^{T-1} \|\mathbf{H}_{t+1} - \mathbf{H}_t\|_F$ in the following lemma. To achieve this, we use the stability properties of our algorithm in (22) and Lemma 8, which is most technical part of the proof.

Lemma 13. We have

$$\sum_{t=0}^{T-1} \ell_t(\mathbf{H}_t) \leq \frac{C_3}{4} L_2^2 \|\mathbf{z}_0 - \mathbf{x}^*\|^2 \quad \text{and} \quad \sum_{t=0}^{T-1} \|\mathbf{H}_{t+1} - \mathbf{H}_t\|_F \leq C_2 \sqrt{d} L_2 \left(1 + \log \frac{A_N}{A_1}\right) \|\mathbf{z}_0 - \mathbf{x}^*\|, \quad (81)$$

where C_2 is defined in (55) and $C_3 = \frac{(1+\alpha_1)^2}{\beta^2(1-\alpha_1)^2(1-\sigma^2)}$.

Step 3: Thus, we obtain an upper bound on $\sum_{t=0}^{T-1} \ell_t(\mathbf{B}_t)$ by combining Lemma 12 and Lemma 13. Finally, in the following lemma, we prove an upper bound on $\frac{1}{A_N}$ by further using Lemma 2 and Proposition 1.

759 **Lemma 14.** *We have*

$$\frac{1}{A_N} \leq \frac{1}{N^{2.5}} \left(M + C_{10} L_1 L_2 d \|\mathbf{z}_0 - \mathbf{x}^*\| \log^+ \left(\frac{\max\{\frac{L_1}{\alpha_2 \beta}, \frac{1}{\sigma_0}\} N^{2.5}}{\sqrt{M}} \right) \right)^{\frac{1}{2}},$$

760 *where we define $\log^+(x) \triangleq \max\{\log(x), 0\}$,*

$$M = \frac{C_6}{\sigma_0^2} + C_7 L_1^2 + C_8 \|\mathbf{B}_0 - \mathbf{H}_0\|_F^2 + C_9 L_2^2 \|\mathbf{z}_0 - \mathbf{x}^*\|^2 + C_{10} L_1 L_2 d \|\mathbf{z}_0 - \mathbf{x}^*\|,$$

761 *and C_i ($i = 6, \dots, 10$) are absolute constants given by*

$$C_6 = \frac{4C_1^2(2 - \beta^2)}{1 - \beta^2}, \quad C_7 = \frac{5C_6}{\alpha_2^2 \beta^2}, \quad C_8 = \frac{256C_6}{\alpha_2^2 \beta^2}, \quad C_9 = \frac{C_3 C_6}{\alpha_2^2 \beta^2}, \quad C_{10} = \frac{512C_2 C_6}{\alpha_2^2 \beta^2}.$$

762 Therefore, Part (b) of Theorem 1 immediately follows from Proposition 1. □

763 In the remaining of this section, we present the proofs for the above lemmas that we used to prove the
764 results in Theorem 1.

765 D.1 Proof of Lemma 11

766 Recall that $\mathbf{w}_k \triangleq \nabla f(\tilde{\mathbf{x}}_{k+1}) - \nabla f(\mathbf{y}_k)$ and $\mathbf{s}_k \triangleq \tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k$ for $k \in \mathcal{B}$. We can write
767 $\nabla f(\tilde{\mathbf{x}}_{k+1}) - \nabla f(\mathbf{y}_k) = \bar{\mathbf{H}}_k(\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k)$ by using the fundamental theorem of calculus, where
768 $\bar{\mathbf{H}}_k = \int_0^1 \nabla^2 f(t\tilde{\mathbf{x}}_{k+1} + (1-t)\mathbf{y}_k) dt$. Since we have $0 \preceq \nabla^2 f(\mathbf{x}) \preceq L_1 \mathbf{I}$ for all $\mathbf{x} \in \mathbb{R}^d$ by
769 Assumption 1, it implies that $0 \preceq \bar{\mathbf{H}}_k \preceq L_1 \mathbf{I}$. Moreover, since $0 \preceq \mathbf{B}_k \preceq L_1 \mathbf{I}$, we further have
770 $-L_1 \mathbf{I} \preceq \bar{\mathbf{H}}_k - \mathbf{B}_k \preceq L_1 \mathbf{I}$, which yields $\|\bar{\mathbf{H}}_k - \mathbf{B}_k\|_{\text{op}} \leq L_1$. Thus, we have

$$\|\mathbf{w}_k - \mathbf{B}_k \mathbf{s}_k\| = \|(\bar{\mathbf{H}}_k - \mathbf{B}_k)(\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k)\| \leq L_1 \|\tilde{\mathbf{x}}_k - \mathbf{x}_k\|,$$

771 which proves that $\ell_k(\mathbf{B}_k) \leq L_1^2$.

772 D.2 Proof of Lemma 12

773 To prove Lemma 12, we first present the following lemma showing a smooth property of the loss
774 function ℓ_k . The proof is similar to [34, Lemma 15].

775 **Lemma 15.** *For $k \in \mathcal{B}$, we have*

$$\nabla \ell_k(\mathbf{B}) = \frac{1}{\|\mathbf{s}_k\|^2} (-\mathbf{s}_k(\mathbf{w}_k - \mathbf{B}\mathbf{s}_k)^\top - (\mathbf{w}_k - \mathbf{B}\mathbf{s}_k)\mathbf{s}_k^\top). \quad (82)$$

776 Moreover, for any $\mathbf{B} \in \mathbb{S}^d$, it holds that

$$\|\nabla \ell_k(\mathbf{B})\|_F \leq \|\nabla \ell_k(\mathbf{B})\|_* \leq 2\sqrt{\ell_k(\mathbf{B})}, \quad (83)$$

777 where $\|\cdot\|_F$ and $\|\cdot\|_*$ denote the Frobenius norm and the nuclear norm, respectively.

778 *Proof.* It is straightforward to verify the expression in (82). The first inequality in (83) follows from
779 the fact that $\|\mathbf{A}\|_F \leq \|\mathbf{A}\|_*$ for any matrix $\mathbf{A} \in \mathbb{S}^d$. For the second inequality, note that

$$\begin{aligned} \|\nabla \ell_k(\mathbf{B})\|_* &\leq \frac{1}{\|\mathbf{s}_k\|^2} (\|\mathbf{s}_k(\mathbf{w}_k - \mathbf{B}\mathbf{s}_k)^\top\|_* + \|(\mathbf{w}_k - \mathbf{B}\mathbf{s}_k)\mathbf{s}_k^\top\|_*) \\ &\leq \frac{2}{\|\mathbf{s}_k\|^2} \|\mathbf{w}_k - \mathbf{B}\mathbf{s}_k\| \|\mathbf{s}_k\| = \frac{2\|\mathbf{w}_k - \mathbf{B}\mathbf{s}_k\|}{\|\mathbf{s}_k\|} = 2\sqrt{\ell_k(\mathbf{B})}, \end{aligned}$$

780 where in the first inequality we used the triangle inequality, and in the second inequality we used the
781 fact that the rank-one matrix $\mathbf{u}\mathbf{v}^\top$ has only one nonzero singular value $\|\mathbf{u}\| \|\mathbf{v}\|$. □

782 We will also need the following helper lemma.

783 **Lemma 16.** *If the real number x satisfies $x \leq A + B\sqrt{x}$, then we have $x \leq 2A + B^2$.*

784 *Proof.* From the assumption, we have

$$\left(\sqrt{x} - \frac{B}{2}\right)^2 \leq A + \frac{B^2}{4}.$$

785 Hence, we obtain

$$x \leq \left(\sqrt{A + \frac{B^2}{4}} + \frac{B}{2}\right)^2 \leq 2A + B^2.$$

786 □

787 Before proving Lemma 12 we also present the following lemma that bounds the loss in each round.

788 **Lemma 17.** For any $\mathbf{H} \in \mathcal{Z}$, we have

$$\ell_t(\mathbf{B}_t) \leq 4\ell_t(\mathbf{H}) + 64L_1^2\|\mathbf{W}_t - \hat{\mathbf{H}}\|_F^2 - 64L_1^2\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}\|_F^2 + 2L_1^2\delta_t^2.$$

789 *Proof.* By letting $\mathbf{x}_t = \hat{\mathbf{B}}_t$, $\mathbf{x} = \hat{\mathbf{H}} \triangleq \frac{2}{L_1}(\mathbf{H} - \frac{L_1}{2}\mathbf{I})$, $\mathbf{g}_t = \mathbf{G}_t \triangleq \frac{2}{L_1}\nabla\ell_t(\mathbf{B}_t)$, $\tilde{\mathbf{g}}_t = \tilde{\mathbf{G}}_t$, $\mathbf{w}_t = \mathbf{W}_t$
790 in Lemma 10 we obtain:

791 (i) $\hat{\mathbf{B}}_t \in \hat{\mathcal{Z}}$, which means that $\|\hat{\mathbf{B}}_t\|_{\text{op}} \leq 1$.

792 (ii) It holds that

$$\langle \mathbf{G}_t, \hat{\mathbf{B}}_t - \hat{\mathbf{H}} \rangle \leq \frac{1}{2\rho}\|\mathbf{W}_t - \hat{\mathbf{H}}\|_F^2 - \frac{1}{2\rho}\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}\|_F^2 + \frac{\rho}{2}\|\tilde{\mathbf{G}}_t\|_F^2 + \max\{0, -\langle \mathbf{G}_t, \hat{\mathbf{B}}_t \rangle\}\delta_t, \quad (84)$$

$$\|\tilde{\mathbf{G}}_t\|_F \leq \|\mathbf{G}_t\|_F + |\langle \mathbf{G}_t, \hat{\mathbf{B}}_t \rangle|\|\mathbf{S}_t\|_F. \quad (85)$$

793 First, note that $\|\mathbf{S}_t\|_F \leq 3$ by Definition 2 and $|\langle \mathbf{G}_t, \hat{\mathbf{B}}_t \rangle| \leq \|\mathbf{G}_t\|_*\|\hat{\mathbf{B}}_t\|_{\text{op}} \leq \|\mathbf{G}_t\|_*$. Together with
794 (85), we get

$$\|\tilde{\mathbf{G}}_t\|_F \leq \|\mathbf{G}_t\|_F + 3\|\mathbf{G}_t\|_* \leq 4\|\mathbf{G}_t\|_* \leq \frac{16}{L_1}\sqrt{\ell_t(\mathbf{B}_t)}, \quad (86)$$

795 where we used the fact that $\mathbf{G}_t = \frac{2}{L_1}\nabla\ell_t(\mathbf{B}_t)$ and Lemma 15 in the last inequality. Furthermore,
796 since ℓ_t is convex, we have

$$\ell_t(\mathbf{B}_t) - \ell_t(\mathbf{H}) \leq \langle \nabla\ell_t(\mathbf{B}_t), \mathbf{B}_t - \mathbf{H} \rangle = \left(\frac{L_1}{2}\right)^2 \langle \mathbf{G}_t, \hat{\mathbf{B}}_t - \hat{\mathbf{H}} \rangle,$$

797 where we used $\mathbf{G}_t = \frac{2}{L_1}\nabla\ell_t(\mathbf{B}_t)$, $\hat{\mathbf{B}}_t \triangleq \frac{2}{L_1}(\mathbf{B}_t - \frac{L_1}{2}\mathbf{I})$, and $\hat{\mathbf{H}} \triangleq \frac{2}{L_1}(\mathbf{H} - \frac{L_1}{2}\mathbf{I})$. Therefore, by
798 combining (84) and (86) we get

$$\ell_t(\mathbf{B}_t) - \ell_t(\mathbf{H}) \leq \frac{L_1^2}{8\rho}\|\mathbf{W}_t - \hat{\mathbf{H}}\|_F^2 - \frac{L_1^2}{8\rho}\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}\|_F^2 + \frac{\rho}{8}L_1^2\|\tilde{\mathbf{G}}_t\|_F^2 + \frac{L_1^2}{4}\|\mathbf{G}_t\|_*\delta_t \quad (87)$$

$$\leq \frac{L_1^2}{8\rho}\|\mathbf{W}_t - \hat{\mathbf{H}}\|_F^2 - \frac{L_1^2}{8\rho}\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}\|_F^2 + 32\rho\ell_t(\mathbf{B}_t) + L_1\sqrt{\ell_t(\mathbf{B}_t)}\delta_t. \quad (88)$$

799 Note that $\ell_t(\mathbf{B}_t)$ appears on both sides of (88). By further applying Lemma 16, we obtain

$$\ell_t(\mathbf{B}_t) \leq 2\ell_t(\mathbf{H}) + \frac{L_1^2}{4\rho}\|\mathbf{W}_t - \hat{\mathbf{H}}\|_F^2 - \frac{L_1^2}{4\rho}\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}\|_F^2 + 64\rho\ell_t(\mathbf{B}_t) + L_1^2\delta_t^2.$$

800 Since $\rho = 1/128$, by rearranging and simplifying terms in the above inequality, we obtain

$$\ell_t(\mathbf{B}_t) \leq 4\ell_t(\mathbf{H}) + 64L_1^2\|\mathbf{W}_t - \hat{\mathbf{H}}\|_F^2 - 64L_1^2\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}\|_F^2 + 2L_1^2\delta_t^2.$$

801 □

802 *Proof of Lemma 12* We let $\mathbf{H}_t = \nabla^2 f(\mathbf{y}_t)$ for $t = 0, 1, \dots, T-1$. Thus, we get

$$\begin{aligned}\ell_t(\mathbf{B}_t) &\leq 4\ell_t(\mathbf{H}_t) + 64L_1^2\|\mathbf{W}_t - \hat{\mathbf{H}}_t\|_F^2 - 64L_1^2\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}_t\|_F^2 + 2L_1^2\delta_t^2 \\ &= 4\ell_t(\mathbf{H}_t) + 64L_1^2\|\mathbf{W}_t - \hat{\mathbf{H}}_t\|_F^2 - 64L_1^2\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}_{t+1}\|_F^2 + 2L_1^2\delta_t^2 \\ &\quad + 64L_1^2(\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}_{t+1}\|_F^2 - \|\mathbf{W}_{t+1} - \hat{\mathbf{H}}_t\|_F^2).\end{aligned}$$

803 Furthermore, note that

$$\begin{aligned}&\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}_{t+1}\|_F^2 - \|\mathbf{W}_{t+1} - \hat{\mathbf{H}}_t\|_F^2 \\ &= (\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}_{t+1}\|_F + \|\mathbf{W}_{t+1} - \hat{\mathbf{H}}_t\|_F)(\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}_{t+1}\|_F - \|\mathbf{W}_{t+1} - \hat{\mathbf{H}}_t\|_F) \\ &\leq 4\sqrt{d}\|\hat{\mathbf{H}}_{t+1} - \hat{\mathbf{H}}_t\|_F = \frac{8\sqrt{d}}{L_1}\|\mathbf{H}_{t+1} - \mathbf{H}_t\|_F,\end{aligned}$$

804 where in the last inequality we used the fact that $\hat{\mathbf{H}}_t, \hat{\mathbf{H}}_{t+1}, \mathbf{W}_{t+1} \in \mathcal{B}_{\sqrt{d}}(0)$ and the triangle
805 inequality. Therefore, we get

$$\ell_t(\mathbf{B}_t) \leq 4\ell_t(\mathbf{H}_t) + 64L_1^2\|\mathbf{W}_t - \hat{\mathbf{H}}_t\|_F^2 - 64L_1^2\|\mathbf{W}_{t+1} - \hat{\mathbf{H}}_{t+1}\|_F^2 + 2L_1^2\delta_t^2 + 512L_1\sqrt{d}\|\mathbf{H}_{t+1} - \mathbf{H}_t\|_F.$$

806 By summing the above inequality from $t = 0$ to $T-1$, we get

$$\sum_{t=0}^{T-1} \ell_t(\mathbf{B}_t) \leq 64L_1^2\|\mathbf{W}_0 - \hat{\mathbf{H}}_0\|_F^2 + 4 \sum_{t=0}^{T-1} \ell_t(\mathbf{H}_t) + 2L_1^2 \sum_{t=0}^{T-1} \delta_t^2 + 512L_1\sqrt{d} \sum_{t=0}^{T-1} \|\mathbf{H}_{t+1} - \mathbf{H}_t\|_F.$$

807 Finally, we use the fact that $\mathbf{W}_0 \triangleq \frac{2}{L_1}(\mathbf{B}_0 - \frac{L_1}{2}\mathbf{I})$, and $\hat{\mathbf{H}}_0 \triangleq \frac{2}{L_1}(\mathbf{H}_0 - \frac{L_1}{2}\mathbf{I})$ to obtain Lemma 12. \square

808 D.3 Proof of Lemma 13

809 By Assumption 2, we have $\|\mathbf{w}_t - \mathbf{H}_t \mathbf{s}_t\| = \|\nabla f(\tilde{\mathbf{x}}_{t+1}) - \nabla f(\mathbf{y}_t) - \nabla f(\mathbf{y}_t)(\tilde{\mathbf{x}}_{t+1} - \mathbf{y}_t)\| \leq$
810 $\frac{L_2}{2}\|\tilde{\mathbf{x}}_{t+1} - \mathbf{y}_t\|^2$. Thus,

$$\ell_t(\mathbf{H}_t) = \frac{\|\mathbf{w}_t - \mathbf{H}_t \mathbf{s}_t\|^2}{\|\mathbf{s}_t\|^2} \leq \frac{L_2^2}{4}\|\tilde{\mathbf{x}}_{t+1} - \mathbf{y}_t\|^2 \leq \frac{(1 + \alpha_1)^2 L_2^2}{4\beta^2(1 - \alpha_1)^2}\|\tilde{\mathbf{x}}_{t+1} - \mathbf{y}_t\|^2,$$

811 where we used Lemma 1 in the last inequality. Also, Since $a_k \geq \eta_k$ for all $k \geq 0$, by (22) we get

$$\sum_{k=0}^{N-1} \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2 \leq \sum_{k=0}^{N-1} \frac{a_k^2}{\eta_k^2} \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2 \leq \frac{1}{1 - \sigma^2} \|\mathbf{z}_0 - \mathbf{x}^*\|^2.$$

812 Hence, we have

$$\begin{aligned}\sum_{t=0}^{T-1} \ell_t(\mathbf{H}_t) &\leq \frac{(1 + \alpha_1)^2 L_2^2}{4\beta^2(1 - \alpha_1)^2} \sum_{k \in \mathcal{B}} \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2 \leq \frac{(1 + \alpha_1)^2 L_2^2}{4\beta^2(1 - \alpha_1)^2} \sum_{k=0}^{N-1} \|\tilde{\mathbf{x}}_{k+1} - \mathbf{y}_k\|^2 \\ &\leq \frac{(1 + \alpha_1)^2 L_2^2 \|\mathbf{z}_0 - \mathbf{x}^*\|^2}{4\beta^2(1 - \alpha_1)^2(1 - \sigma^2)},\end{aligned}$$

813 which proves the first inequality in (81).

814 Furthermore, by Assumption 2, we have

$$\|\mathbf{H}_{t+1} - \mathbf{H}_t\|_F = \|\nabla^2 f(\mathbf{y}_{t+1}) - \nabla^2 f(\mathbf{y}_t)\|_F \leq \sqrt{d}\|\nabla^2 f(\mathbf{y}_{t+1}) - \nabla^2 f(\mathbf{y}_t)\|_{\text{op}} \leq \sqrt{d}L_2\|\mathbf{y}_{t+1} - \mathbf{y}_t\|.$$

815 Hence, by using the triangle inequality, we can bound

$$\sum_{t=0}^{T-1} \|\mathbf{H}_{t+1} - \mathbf{H}_t\|_F \leq \sqrt{d}L_2 \sum_{k=0}^{N-1} \|\mathbf{y}_{k+1} - \mathbf{y}_k\| \leq \sqrt{d}L_2 C_2 \left(1 + \log \frac{A_N}{A_1}\right) \|\mathbf{z}_0 - \mathbf{x}^*\|,$$

816 where we used Lemma 8 in the last inequality.

817 D.4 Proof of Lemma 14

818 We combine Lemma 12 and Lemma 13 to get

$$\begin{aligned} \sum_{k \in \mathcal{B}} \frac{\|\mathbf{w}_k - \mathbf{B}_k \mathbf{s}_k\|^2}{\|\mathbf{s}_k\|^2} &= \sum_{t=0}^{T-1} \ell_t(\mathbf{B}_t) \leq 256 \|\mathbf{B}_0 - \mathbf{H}_0\|_F^2 + C_3 L_2^2 \|\mathbf{z}_0 - \mathbf{x}^*\|^2 + 2L_1^2 \sum_{t=0}^{T-1} \delta_t^2 \\ &\quad + 512C_2 L_1 L_2 d \left(1 + \log \frac{A_N}{A_1}\right) \|\mathbf{z}_0 - \mathbf{x}^*\|. \end{aligned}$$

819 Since $\delta_t = 1/(\sqrt{t+2} \ln(t+2))$, we have

$$\sum_{t=0}^{T-1} \delta_t^2 = \sum_{t=2}^{T+1} \frac{1}{t \ln^2 t} \leq \frac{1}{2 \ln^2 2} + \int_2^{T+1} \frac{1}{t \ln^2 t} dt = \frac{1}{2 \ln^2 2} + \frac{1}{\ln 2} - \frac{1}{\ln(T+1)} \leq 2.5.$$

820 Hence, it further follows from (15) and Lemma 2 that

$$\begin{aligned} \frac{N^5}{A_N^2} &\leq 4C_1^2 \sum_{k=0}^{N-1} \frac{1}{\hat{\eta}_k^2} \\ &\leq \frac{4C_1^2(2-\beta^2)}{(1-\beta^2)\sigma_0^2} + \frac{4C_1^2(2-\beta^2)}{(1-\beta^2)\alpha_2^2\beta^2} \sum_{k \in \mathcal{B}} \frac{\|\mathbf{w}_k - \mathbf{B}_k \mathbf{s}_k\|^2}{\|\mathbf{s}_k\|^2} \\ &\leq \frac{C_6}{\sigma_0^2} + C_7 L_1^2 + C_8 \|\mathbf{B}_0 - \mathbf{H}_0\|_F^2 + C_9 L_2^2 \|\mathbf{z}_0 - \mathbf{x}^*\|^2 \\ &\quad + C_{10} L_1 L_2 d \left(1 + \log \frac{A_N}{A_1}\right) \|\mathbf{z}_0 - \mathbf{x}^*\|. \end{aligned}$$

821 To simplify the notation, define

$$M = \frac{C_6}{\sigma_0^2} + C_7 L_1^2 + C_8 \|\mathbf{B}_0 - \mathbf{H}_0\|_F^2 + C_9 L_2^2 \|\mathbf{z}_0 - \mathbf{x}^*\|^2 + C_{10} L_1 L_2 d \|\mathbf{z}_0 - \mathbf{x}^*\|.$$

822 Let A_N^* be the number that achieves the equality

$$\frac{N^5}{(A_N^*)^2} = M + C_{10} L_1 L_2 d \|\mathbf{z}_0 - \mathbf{x}^*\| \log \frac{A_N^*}{A_1}.$$

823 We can see that $A_N \geq A_N^*$. Thus, we instead try to construct a lower bound on A_N^* . If $A_N^* \leq A_1$,
824 then $\log(A_N^*/A_1) \leq 0$ and furthermore

$$\frac{N^5}{(A_N^*)^2} \leq M \quad \Rightarrow \quad A_N^* \geq \frac{1}{\sqrt{M}} N^{2.5}.$$

825 Otherwise, assume that $A_N^* > A_1$. Then $\log(A_N^*/A_1) > 0$ and we first show an upper bound on
826 A_N^* :

$$\frac{N^5}{(A_N^*)^2} = M + C_8 L_1 L_2 d \|\mathbf{z}_0 - \mathbf{x}^*\| \log \frac{A_N^*}{A_1} \geq M \quad \Rightarrow \quad A_N^* \leq \frac{1}{\sqrt{M}} N^{2.5}.$$

827 This in turn leads to a lower bound on A_N^* :

$$\frac{N^5}{(A_N^*)^2} = M + C_8 L_1 L_2 d \|\mathbf{z}_0 - \mathbf{x}^*\| \log \frac{A_N^*}{A_1} \leq M + C_8 L_1 L_2 d \|\mathbf{z}_0 - \mathbf{x}^*\| \log \left(\frac{\max\{\frac{L_1}{\alpha_2 \beta}, \frac{1}{\sigma_0}\} N^{2.5}}{\sqrt{M}} \right),$$

828 where we also used the fact that $A_1 = \hat{\eta}_1 \geq \min\{\sigma_0, \frac{\alpha_2 \beta}{L_1}\}$. Thus, we get

$$\frac{1}{A_N} \leq \frac{1}{A_N^*} \leq \frac{1}{N^{2.5}} \left(M + C_{10} L_1 L_2 d \|\mathbf{z}_0 - \mathbf{x}^*\| \log \left(\frac{\max\{\frac{L_1}{\alpha_2 \beta}, \frac{1}{\sigma_0}\} N^{2.5}}{\sqrt{M}} \right) \right)^{\frac{1}{2}}.$$

Subroutine 3 LinearSolver($\mathbf{A}, \mathbf{b}; \alpha$)

```

1: Input:  $\mathbf{A} \in \mathbb{S}_+^d, \mathbf{b} \in \mathbb{R}^d, 0 < \alpha < 1$ 
2: Initialize:  $\mathbf{s}_0 \leftarrow \mathbf{0}, \mathbf{r}_0 \leftarrow \mathbf{b} - \mathbf{A}\mathbf{s}_0, \mathbf{p}_0 \leftarrow \mathbf{r}_0$ 
3: for  $k = 0, 1, \dots$  do
4:   if  $\|\mathbf{r}_k\|_2 \leq \alpha \|\mathbf{s}_k\|_2$  then
5:     Return  $\mathbf{s}_k$ 
6:   end if
7:    $\alpha_k \leftarrow \langle \mathbf{r}_k, \mathbf{A}\mathbf{r}_k \rangle / \langle \mathbf{A}\mathbf{p}_k, \mathbf{A}\mathbf{p}_k \rangle$ 
8:    $\mathbf{s}_{k+1} \leftarrow \mathbf{s}_k + \alpha_k \mathbf{p}_k$ 
9:    $\mathbf{r}_{k+1} \leftarrow \mathbf{r}_k - \alpha_k \mathbf{A}\mathbf{p}_k$ 
10:  Compute and store  $\mathbf{A}\mathbf{r}_{k+1}$ 
11:   $\beta_k \leftarrow \langle \mathbf{r}_{k+1}, \mathbf{A}\mathbf{r}_{k+1} \rangle / \langle \mathbf{r}_k, \mathbf{A}\mathbf{r}_k \rangle$ 
12:   $\mathbf{p}_{k+1} \leftarrow \mathbf{r}_{k+1} + \beta_k \mathbf{p}_k$ 
13:  Compute and store  $\mathbf{A}\mathbf{p}_{k+1} \leftarrow \mathbf{A}\mathbf{r}_{k+1} + \beta_k \mathbf{A}\mathbf{p}_k$ 
14: end for

```

E Characterizing the Computational Cost

In this section, we first specify the implementation details of the LinearSolver oracle in Definition [1] and the SEP oracle in Definition [2]. Then in Section [E.3], we present the proof of Theorem [2].

E.1 Implementation of the LinearSolver Oracle

We implement the LinearSolver oracle by running the conjugate residual (CR) method [38] to solve the linear system $\mathbf{A}\mathbf{s} = \mathbf{b}$. In particular, we initialize the CT method with $\mathbf{s}_0 = \mathbf{0}$ and returns the iterate \mathbf{s}_k once we have $\|\mathbf{A}\mathbf{s}_k - \mathbf{b}\| \leq \alpha \|\mathbf{s}_k\|$.

The following lemma provides the convergence guarantee of the CR method, which will be later used in the proof of Theorem [2].

Lemma 18. *Let \mathbf{s}^* be any optimal solution of $\mathbf{A}\mathbf{s}^* = \mathbf{b}$ and let $\{\mathbf{s}_k\}$ be the iterates generated by Subroutine [3]. Then we have*

$$\|\mathbf{r}_k\|_2 = \|\mathbf{A}\mathbf{s}_k - \mathbf{b}\|_2 \leq \frac{\lambda_{\max}(\mathbf{A}) \|\mathbf{s}^*\|_2}{(k+1)^2}.$$

E.2 Implementation of SEP Oracle

We implement the SEP oracle in Definition [2] by running the classical Lanczos method, with a random start, where the initial vector is chosen randomly and uniformly from the unit sphere (see, e.g., [45, 46]). For completeness, the full algorithm is shown in Subroutine [4].

To prove the correctness of our algorithm, we first recall a classical result in [41] on the convergence behavior of the Lanczos method.

Proposition 3 ([41] Theorem 4.2). *Consider a symmetric matrix \mathbf{W} and let $\lambda_1(\mathbf{W})$ and $\lambda_d(\mathbf{W})$ denote its largest and smallest eigenvalues, respectively. Then after k iterations of the Lanczos method with a random start, we find unit vectors $\mathbf{u}^{(1)}$ and $\mathbf{u}^{(d)}$ such that*

$$\begin{aligned} \mathbb{P}(\langle \mathbf{W}\mathbf{u}^{(1)}, \mathbf{u}^{(1)} \rangle &\leq \lambda_1(\mathbf{W}) - \epsilon(\lambda_1(\mathbf{W}) - \lambda_d(\mathbf{W}))) \leq 1.648\sqrt{de}^{-\sqrt{\epsilon}(2k-1)}, \\ \mathbb{P}(\langle \mathbf{W}\mathbf{u}^{(d)}, \mathbf{u}^{(d)} \rangle &\geq \lambda_d(\mathbf{W}) + \epsilon(\lambda_1(\mathbf{W}) - \lambda_d(\mathbf{W}))) \leq 1.648\sqrt{de}^{-\sqrt{\epsilon}(2k-1)}, \end{aligned}$$

As a corollary, to ensure that, with probability at least $1 - q$,

$$\langle \mathbf{W}\mathbf{u}^{(1)}, \mathbf{u}^{(1)} \rangle > \lambda_1(\mathbf{W}) - \epsilon(\lambda_1(\mathbf{W}) - \lambda_d(\mathbf{W})) \text{ and } \langle \mathbf{W}\mathbf{u}^{(d)}, \mathbf{u}^{(d)} \rangle < \lambda_d(\mathbf{W}) + \epsilon(\lambda_1(\mathbf{W}) - \lambda_d(\mathbf{W})),$$

the number of iterations can be bounded by $\lceil \frac{1}{4}\epsilon^{-1/2} \log(11d/q^2) + \frac{1}{2} \rceil$.

Lemma 19. *Let γ and \mathbf{S} be the output of SEP($\mathbf{W}; \delta, q$) in Subroutine [4]. Then with probability at least $1 - q$, they satisfy one of the following properties:*

- Case I: $\gamma \leq 1$, then we have $\|\mathbf{W}\|_{\text{op}} \leq 1$;

Subroutine 4 SEP($\mathbf{W}; \delta, q$)

```

1: Input:  $\mathbf{W} \in \mathbb{S}^d$ ,  $\delta > 0$ ,  $q \in (0, 1)$ 
2: Initialize: sample  $\mathbf{v}_1 \in \mathbb{R}^d$  uniformly from the unit sphere,  $\beta_1 \leftarrow 0$ ,  $\mathbf{v}_0 \leftarrow 0$ 
3: Set the number of iterations  $N_1 \leftarrow \min\left\{\left\lceil \log \frac{11d}{q^2} + \frac{1}{2} \right\rceil, d\right\}$ 
4: for  $k = 1, \dots, N_1$  do
5:   Set  $\mathbf{w}_k \leftarrow \mathbf{W}\mathbf{v}_k - \beta_k \mathbf{v}_{k-1}$ 
6:   Set  $\alpha_k \leftarrow \langle \mathbf{w}_k, \mathbf{v}_k \rangle$  and  $\mathbf{w}_k \leftarrow \mathbf{w}_k - \alpha_k \mathbf{v}_k$ 
7:   Set  $\beta_{k+1} \leftarrow \|\mathbf{w}_k\|$  and  $\mathbf{v}_{k+1} \leftarrow \mathbf{w}_k / \beta_{k+1}$ 
8: end for
9: Form a tridiagonal matrix  $\mathbf{T} \leftarrow \text{tridiag}(\beta_{2:N_1}, \alpha_{1:N_1}, \beta_{2:N_1})$ 
10: # Use the tridiagonal structure to compute eigenvectors of  $\mathbf{T}$ 
11: Compute  $(\hat{\lambda}_1, \mathbf{z}^{(1)}) \leftarrow \text{MaxEvec}(\mathbf{T})$  and  $(\hat{\lambda}_d, \mathbf{z}^{(d)}) \leftarrow \text{MinEvec}(\mathbf{T})$ 
12: Set  $\mathbf{u}^{(1)} \leftarrow \sum_{k=1}^{N_1} z_k^{(1)} \mathbf{v}_k$  and  $\mathbf{u}^{(d)} \leftarrow \sum_{k=1}^{N_1} z_k^{(d)} \mathbf{v}_k$ 
13: Set  $\hat{\lambda}_{\max} \leftarrow \max\{\hat{\lambda}_1, -\hat{\lambda}_d\}$ 
14: if  $\hat{\lambda}_{\max} \leq 1/2$  then # Case I:  $\gamma \leq 1$ , which implies  $\|\mathbf{W}\|_{\text{op}} \leq 1$ 
15:   Return  $\gamma = 2\hat{\lambda}_{\max}$  and  $\mathbf{S} = 0$ 
16: else if  $\hat{\lambda}_{\max} \geq 2$  then # Case II:  $\gamma > 1$  and  $\mathbf{S}$  defines a separating hyperplane
17:   if  $\hat{\lambda}_1 > -\hat{\lambda}_d$  then
18:     Return  $\gamma = 2\hat{\lambda}_{\max}$  and  $\mathbf{S} = 3\mathbf{u}^{(1)}(\mathbf{u}^{(1)})^\top$ 
19:   else
20:     Return  $\gamma = 2\hat{\lambda}_{\max}$  and  $\mathbf{S} = -3\mathbf{u}^{(d)}(\mathbf{u}^{(d)})^\top$ 
21:   end if
22: else #  $\frac{1}{2} < \hat{\lambda}_{\max} < 2$ 
23:   Set the number of iterations  $N_2 \leftarrow \min\left\{\left\lceil \frac{1}{4\sqrt{2}\delta} \log \frac{44d}{q^2} + \frac{1}{2} \right\rceil, d\right\}$ 
24:   for  $k = N_1 + 1, \dots, N_2$  do
25:     Set  $\mathbf{w}_k \leftarrow \mathbf{W}\mathbf{v}_k - \beta_k \mathbf{v}_{k-1}$ 
26:     Set  $\alpha_k \leftarrow \langle \mathbf{w}_k, \mathbf{v}_k \rangle$  and  $\mathbf{w}_k \leftarrow \mathbf{w}_k - \alpha_k \mathbf{v}_k$ 
27:     Set  $\beta_{k+1} \leftarrow \|\mathbf{w}_k\|$  and  $\mathbf{v}_{k+1} \leftarrow \mathbf{w}_k / \beta_{k+1}$ 
28:   end for
29:   Form a tridiagonal matrix  $\mathbf{T} \leftarrow \text{tridiag}(\beta_{2:N_2}, \alpha_{1:N_2}, \beta_{2:N_2})$ 
30:   Compute  $(\tilde{\lambda}_1, \tilde{\mathbf{z}}^{(1)}) \leftarrow \text{MaxEvec}(\mathbf{T})$  and  $(\tilde{\lambda}_d, \tilde{\mathbf{z}}^{(d)}) \leftarrow \text{MinEvec}(\mathbf{T})$ 
31:   Set  $\tilde{\mathbf{u}}^{(1)} \leftarrow \sum_{k=1}^{N_2} \tilde{z}_k^{(1)} \mathbf{v}_k$  and  $\tilde{\mathbf{u}}^{(d)} \leftarrow \sum_{k=1}^{N_2} \tilde{z}_k^{(d)} \mathbf{v}_k$ 
32:   Set  $\tilde{\lambda}_{\max} = \max\{\tilde{\lambda}_1, -\tilde{\lambda}_d\}$ 
33:   if  $\tilde{\lambda}_{\max} \leq 1 - \delta$  then
34:     Return  $\gamma = \tilde{\lambda}_{\max} + \delta$  and  $\mathbf{S} = 0$ 
35:   else if  $\tilde{\lambda}_1 \geq -\tilde{\lambda}_d$  then
36:     Return  $\gamma = \tilde{\lambda}_{\max} + \delta$  and  $\mathbf{S} = \tilde{\mathbf{u}}^{(1)}(\tilde{\mathbf{u}}^{(1)})^\top$ 
37:   else
38:     Return  $\gamma = \tilde{\lambda}_{\max} + \delta$  and  $\mathbf{S} = -\tilde{\mathbf{u}}^{(d)}(\tilde{\mathbf{u}}^{(d)})^\top$ 
39:   end if
40: end if

```

} *Lanczos method*

854 • *Case II: $\gamma > 1$, then we have $\|\mathbf{W}/\gamma\|_{\text{op}} \leq 1$, $\|\mathbf{S}\|_F = 3$ and $\langle \mathbf{S}, \mathbf{W} - \hat{\mathbf{B}} \rangle \geq \gamma - 1$ for any*
855 *$\hat{\mathbf{B}}$ such that $\|\hat{\mathbf{B}}\|_{\text{op}} \leq 1$.*

856 *Proof.* Note that in Subroutine 4, we first run the Lanczos method for $\left\lceil \epsilon^{-1/2} \log \frac{11d}{q^2} + \frac{1}{2} \right\rceil$ iterations,
857 where $\epsilon = \frac{1}{4}$. Thus, by Proposition 3, with probability at least $1 - q/2$ we have

$$\hat{\lambda}_1 \triangleq \langle \mathbf{W}\mathbf{u}^{(1)}, \mathbf{u}^{(1)} \rangle \geq \lambda_1(\mathbf{W}) - \frac{1}{4}(\lambda_1(\mathbf{W}) - \lambda_d(\mathbf{W})), \quad (89)$$

$$\hat{\lambda}_d \triangleq \langle \mathbf{W}\mathbf{u}^{(d)}, \mathbf{u}^{(d)} \rangle \leq \lambda_d(\mathbf{W}) + \frac{1}{4}(\lambda_1(\mathbf{W}) - \lambda_d(\mathbf{W})). \quad (90)$$

858 Combining (89) and (90), we get

$$\frac{1}{2}(\lambda_1(\mathbf{W}) - \lambda_d(\mathbf{W})) \leq \hat{\lambda}_1 - \hat{\lambda}_d \quad \Rightarrow \quad \lambda_1(\mathbf{W}) - \lambda_d(\mathbf{W}) \leq 2(\hat{\lambda}_1 - \hat{\lambda}_d).$$

By plugging the above inequality back into (89) and (90), we further have

$$\lambda_1(\mathbf{W}) \leq \hat{\lambda}_1 + \frac{1}{4}(\lambda_1(\mathbf{W}) - \lambda_d(\mathbf{W})) \leq \hat{\lambda}_1 + \frac{1}{2}(\hat{\lambda}_1 - \hat{\lambda}_d), \quad (91)$$

$$\lambda_d(\mathbf{W}) \geq \hat{\lambda}_d - \frac{1}{4}(\lambda_1(\mathbf{W}) - \lambda_d(\mathbf{W})) \geq \hat{\lambda}_d - \frac{1}{2}(\hat{\lambda}_1 - \hat{\lambda}_d). \quad (92)$$

Let $\hat{\lambda}_{\max} = \max\{\hat{\lambda}_1, -\hat{\lambda}_d\}$. By (91) and (92), we can further bound the eigenvalues of \mathbf{W} by

$$\lambda_1(\mathbf{W}) \leq \hat{\lambda}_{\max} + \frac{1}{2} \cdot 2\hat{\lambda}_{\max} = 2\hat{\lambda}_{\max} \quad \text{and} \quad \lambda_d(\mathbf{W}) \geq -\hat{\lambda}_{\max} - \frac{1}{2} \cdot 2\hat{\lambda}_{\max} = -2\hat{\lambda}_{\max}. \quad (93)$$

Hence, we can see that $\|\mathbf{W}\|_{\text{op}} = \max\{\lambda_1(\mathbf{W}), -\lambda_d(\mathbf{W})\} \leq 2\hat{\lambda}_{\max}$. Now we distinguish three cases.

(a) If $\hat{\lambda}_{\max} \leq \frac{1}{2}$, then we are in **Case I** and the ExtEvec oracle outputs $\gamma = 2\hat{\lambda}_{\max} \leq 1$ and $\mathbf{S} = \mathbf{0}$. In this case, we indeed have $\|\mathbf{W}\|_{\text{op}} \leq \gamma \leq 1$.

(b) If $\hat{\lambda}_{\max} \geq 2$, then we are in **Case II**. In addition, if $\hat{\lambda}_1 \geq -\hat{\lambda}_d$, then the ExtEvec oracle returns $\gamma = 2\hat{\lambda}_{\max}$ and $\mathbf{S} = 3\mathbf{u}^{(1)}(\mathbf{u}^{(1)})^\top$. Similarly, if $-\hat{\lambda}_d > \hat{\lambda}_1$, then the ExtEvec oracle returns $\gamma = 2\hat{\lambda}_{\max}$ and $\mathbf{S} = -3\mathbf{u}^{(d)}(\mathbf{u}^{(d)})^\top$. Without loss of generality, consider the case where $\hat{\lambda}_1 \geq -\hat{\lambda}_d$. Since $\|\mathbf{W}\|_{\text{op}} \leq 2\hat{\lambda}_{\max} = \gamma$, we have $\|\mathbf{W}/\gamma\|_{\text{op}} \leq 1$. Also, since \mathbf{u}_1 is a unit vector, we have $\|\mathbf{S}\|_F = 3\|\mathbf{u}^{(1)}\|^2 = 3$. Finally, for any $\hat{\mathbf{B}}$ such that $\|\hat{\mathbf{B}}\|_{\text{op}} \leq 1$, we have

$$\langle \mathbf{S}, \mathbf{W} - \hat{\mathbf{B}} \rangle = 3(\mathbf{u}^{(1)})^\top \mathbf{W} \mathbf{u}^{(1)} - 3(\mathbf{u}^{(1)})^\top \hat{\mathbf{B}} \mathbf{u}^{(1)} \geq 3\hat{\lambda}_{\max} - 3 \geq 2\hat{\lambda}_{\max} - 1 = \gamma - 1,$$

where we used the fact that $\hat{\lambda}_{\max} \geq 2$ in the last inequality.

(c) If $\frac{1}{2} < \hat{\lambda}_{\max} < 2$, we continue to run the Lanczos method for a total number of $\left\lceil \frac{1}{4}\epsilon^{-1/2} \log \frac{11d}{q^2} + \frac{1}{2} \right\rceil$ iterations, where $\epsilon = \frac{1}{8}\delta$. Thus, by Proposition 3, with probability at least $1 - q/2$ we have

$$\tilde{\lambda}_1 \triangleq \langle \mathbf{W} \tilde{\mathbf{u}}^{(1)}, \tilde{\mathbf{u}}^{(1)} \rangle \geq \lambda_1(\mathbf{W}) - \frac{1}{8}\delta(\lambda_1(\mathbf{W}) - \lambda_d(\mathbf{W})), \quad (94)$$

$$\tilde{\lambda}_d \triangleq \langle \mathbf{W} \tilde{\mathbf{u}}^{(d)}, \tilde{\mathbf{u}}^{(d)} \rangle \leq \lambda_d(\mathbf{W}) + \frac{1}{8}\delta(\lambda_1(\mathbf{W}) - \lambda_d(\mathbf{W})). \quad (95)$$

Let $\tilde{\lambda}_{\max} = \max\{\tilde{\lambda}_1, -\tilde{\lambda}_d\}$. Since we have $\lambda_1(\mathbf{W}) \leq 2\hat{\lambda}_{\max} \leq 4$ and $\lambda_d(\mathbf{W}) \geq -2\hat{\lambda}_{\max} \geq -4$, the above implies that $\tilde{\lambda}_1 \geq \lambda_1(\mathbf{W}) - \delta$ and $\tilde{\lambda}_d \leq \lambda_d(\mathbf{W}) + \delta$. Hence, we can see that $\|\mathbf{W}\|_{\text{op}} = \max\{\lambda_1(\mathbf{W}), -\lambda_d(\mathbf{W})\} \leq \hat{\lambda}_{\max} + \delta$. We further consider two subcases.

(c1) If $\tilde{\lambda}_{\max} \leq 1 - \delta$, then we are in **Case I** and the ExtEvec oracle outputs $\gamma = \tilde{\lambda}_{\max} + \delta$ and $\mathbf{S} = \mathbf{0}$. In this case, we indeed have $\|\mathbf{W}\|_{\text{op}} \leq \gamma \leq 1$.

(c2) If $\tilde{\lambda}_{\max} > 1 - \delta$, then we are in **Case II**. In addition, if $\tilde{\lambda}_1 \geq -\tilde{\lambda}_d$, then the ExtEvec oracle returns $\gamma = \tilde{\lambda}_{\max} + \delta$ and $\mathbf{S} = \tilde{\mathbf{u}}^{(1)}(\tilde{\mathbf{u}}^{(1)})^\top$. Similarly, if $-\tilde{\lambda}_d > \tilde{\lambda}_1$, then the ExtEvec oracle returns $\gamma = \tilde{\lambda}_{\max} + \delta$ and $\mathbf{S} = -\tilde{\mathbf{u}}^{(d)}(\tilde{\mathbf{u}}^{(d)})^\top$. Without loss of generality, consider the case where $\tilde{\lambda}_1 \geq -\tilde{\lambda}_d$. Since $\|\mathbf{W}\|_{\text{op}} \leq \tilde{\lambda}_{\max} + \delta = \gamma$, we have $\|\mathbf{W}/\gamma\|_{\text{op}} \leq 1$. Also, since $\tilde{\mathbf{u}}^{(1)}$ is a unit vector, we have $\|\mathbf{S}\|_F = \|\tilde{\mathbf{u}}^{(1)}\|^2 = 1$. Finally, for any $\hat{\mathbf{B}}$ such that $\|\hat{\mathbf{B}}\|_{\text{op}} \leq 1$, we have

$$\langle \mathbf{S}, \mathbf{W} - \hat{\mathbf{B}} \rangle = (\tilde{\mathbf{u}}^{(1)})^\top \mathbf{W} \tilde{\mathbf{u}}^{(1)} - (\tilde{\mathbf{u}}^{(1)})^\top \hat{\mathbf{B}} \tilde{\mathbf{u}}^{(1)} \geq \tilde{\lambda}_{\max} - 1 = \gamma - 1 - \delta.$$

This completes the proof. \square

888 E.3 Proof of Theorem 2

889 We divide the proof of Theorem 2 into the following three lemmas.

890 **Lemma 20.** *If we run Algorithm 1 as specified in Theorem 1 for N iterations, then the total number*
 891 *of line search steps can be bounded by $2N + \log_{1/\beta}(\sigma_0 L_1 / \alpha_2)$. As a corollary, the total number of*
 892 *gradient queries is bounded by $3N_\epsilon + \log_{1/\beta}(\frac{\sigma_0 L_1}{\alpha_2})$.*

893 *Proof.* In our backtracking scheme, the number of steps in each iteration is given by $\log_{1/\beta}(\eta_k / \hat{\eta}_k) +$
 894 1. Also note that $\eta_{k+1} \leq \hat{\eta}_k / \beta$ for all $k \geq 0$. Thus, we have

$$\begin{aligned} \sum_{k=0}^{N-1} \left(\log_{1/\beta} \frac{\eta_k}{\hat{\eta}_k} + 1 \right) &= N + \log_{1/\beta} \frac{\sigma_0}{\hat{\eta}_0} + \sum_{k=0}^{N-2} \log_{1/\beta} \frac{\eta_{k+1}}{\hat{\eta}_{k+1}} \\ &\leq N + \log_{1/\beta} \frac{\sigma_0}{\hat{\eta}_0} + \sum_{k=0}^{N-2} \left(\log_{1/\beta} \frac{\hat{\eta}_k}{\hat{\eta}_{k+1}} + 1 \right) \\ &\leq 2N - 1 + \log_{1/\beta} \frac{\sigma_0}{\hat{\eta}_{N-1}} \end{aligned}$$

895 Furthermore, since $\hat{\eta}_k \geq \alpha_2 \beta / L_1$ for all $k \geq 0$, we arrive at the conclusion. \square

896 **Lemma 21.** *The total number of matrix-vector product evaluations in the LinearSolver oracle is*
 897 *bounded by $N_\epsilon + C_{11} \sqrt{\sigma_0 L_1} + C_{12} \sqrt{\frac{L_1 \|\mathbf{z}_0 - \mathbf{x}^*\|^2}{2\epsilon}}$, where C_{11} and C_{12} are absolute constants.*

898 *Proof.* The following proof loosely follows the strategy in [31]. We first bound the number of steps
 899 required by Subroutine 3 before it terminates.

900 **Lemma 22.** *Suppose $\mathbf{A} \succeq \mathbf{I}$. Then Subroutine 3 terminates after at most $\left\lceil \sqrt{\frac{\alpha+1}{\alpha} \lambda_{\max}(\mathbf{A})} - 1 \right\rceil$*
 901 *iterations.*

902 *Proof.* Note that $\|\mathbf{s}_k\|_2 \geq \|\mathbf{s}^*\|_2 - \|\mathbf{s}_k - \mathbf{s}^*\|_2$. Also, since $\mathbf{A} \succeq \mathbf{I}$, we have $\|\mathbf{s}_k - \mathbf{s}^*\|_2 \leq$
 903 $\|\mathbf{A}(\mathbf{s}_k - \mathbf{s}^*)\|_2 = \|\mathbf{r}_k\|_2$. Therefore, we have

$$\|\mathbf{r}_k\|_2 \leq \alpha \|\mathbf{s}_k\|_2 \iff \|\mathbf{r}_k\|_2 \leq \alpha \|\mathbf{s}^*\|_2 - \alpha \|\mathbf{r}_k\|_2 \iff \|\mathbf{r}_k\|_2 \leq \frac{\alpha}{\alpha+1} \|\mathbf{s}^*\|_2.$$

904 By using Lemma 18, we only need $k \geq \sqrt{\frac{\alpha+1}{\alpha} \lambda_{\max}(\mathbf{A})} - 1$ to achieve $\|\mathbf{A}\mathbf{s}_k - \mathbf{b}\| \leq \alpha \|\mathbf{s}_k\|$. \square

905 Moreover, when the step size is smaller enough, we can show that Subroutine 3 will terminate in one
 906 iteration.

907 **Lemma 23.** *Let $\mathbf{A} = \mathbf{I} + \eta \mathbf{B}$. When $\eta \leq \frac{\alpha}{2L_1}$, Algorithm 3 terminates in one iteration.*

908 *Proof.* From the update rule of Subroutine 3, we can compute that $\mathbf{s}_1 = \frac{\mathbf{b}^\top \mathbf{A} \mathbf{b}}{\|\mathbf{A} \mathbf{b}\|_2^2} \mathbf{b}$, which implies

$$\|\mathbf{s}_1\| = \|\mathbf{b}\| \cdot \frac{\|\mathbf{A}^{1/2} \mathbf{b}\|^2}{(\mathbf{A}^{1/2} \mathbf{b})^\top \mathbf{A} (\mathbf{A}^{1/2} \mathbf{b})} \geq \frac{\|\mathbf{b}\|}{\lambda_{\max}(\mathbf{A})} \geq \frac{\|\mathbf{b}\|}{1 + \eta L_1}.$$

909 On the other hand, we also have

$$\|\mathbf{r}_1\| \leq \|\mathbf{A} \mathbf{b} - \mathbf{b}\| = \eta \|\mathbf{B} \mathbf{b}\| \leq \eta L_1 \|\mathbf{b}\|. \quad (96)$$

910 Moreover, when $\eta \leq \frac{\alpha}{2L_1}$, we have $\eta L_1 \leq \frac{\alpha}{1 + \eta L_1}$, which implies that $\|\mathbf{r}_1\| \leq \alpha \|\mathbf{s}_1\|$. \square

911 Now we upper bound the total number of matrix-vector products in Algorithm 1. When $\mathbf{A} = \mathbf{I} + \eta_+ \mathbf{B}_k$
 912 where $\eta_+ = \eta_k \beta^i$. We can store the vector $\mathbf{B}_k \mathbf{b}$ at the beginning and reuse it to compute \mathbf{s}_1 when the
 913 step size $\eta_+ < \frac{\alpha_1}{2L_1}$. And when $\beta^i \eta_k L_1 \geq \frac{\alpha_1}{2}$, it holds that

$$1 + \beta^i \eta_k L_1 \leq \frac{\alpha_1 + 2}{\alpha_1} \beta^i \eta_k L_1.$$

Thus, at the k -th iteration, the number of matrix-vector products can be bounded by

$$\begin{aligned} \text{MV}_k &\leq 1 + \sum_{i \geq 0, \eta_k \beta^i \geq \frac{\alpha_1}{2L_1}} \sqrt{\frac{\alpha_1 + 1}{\alpha_1}} (1 + \eta_k \beta^i L_1) \\ &\leq 1 + \sum_{i \geq 0, \eta_k \beta^i \geq \frac{\alpha_1}{2L_1}} \frac{\alpha_1 + 2}{\alpha_1} \sqrt{\beta^i \eta_k L_1} \\ &\leq 1 + \frac{\alpha_1 + 2}{\alpha_1} \frac{1}{1 - \sqrt{\beta}} \sqrt{\eta_k L_1}. \end{aligned}$$

Furthermore, we can bound that

$$\sum_{k=0}^{N-1} \sqrt{\eta_k} \leq \sqrt{\sigma_0} + \sum_{k=1}^{N-1} \sqrt{\eta_k} \leq \sqrt{\sigma_0} + \frac{1}{\sqrt{\beta}} \sum_{k=0}^{N-2} \sqrt{\hat{\eta}_k} \leq \sqrt{\sigma_0} + \frac{2(2 - \sqrt{\beta})}{\sqrt{\beta}(1 - \sqrt{\beta})} \sqrt{A_{N-1}} \quad (97)$$

Note that $\epsilon < f(x_{N-1}) - f(\mathbf{x}^*) \leq \frac{\|\mathbf{z}_0 - \mathbf{x}^*\|^2}{2A_{N-1}}$. Hence, we have $A_{N-1} \leq \frac{\|\mathbf{z}_0 - \mathbf{x}^*\|^2}{2\epsilon}$. Thus, we can bound the total number of matrix-vector product evaluations by

$$\begin{aligned} \text{MV} &= \sum_{k=0}^{N_\epsilon-1} \text{MV}_k \leq N_\epsilon + \frac{\alpha_1 + 2}{\alpha_1} \frac{1}{1 - \sqrt{\beta}} \left(\sqrt{\sigma_0 L_1} + \frac{2(2 - \sqrt{\beta})}{\sqrt{\beta}(1 - \sqrt{\beta})} \sqrt{\frac{L_1 \|\mathbf{z}_0 - \mathbf{x}^*\|^2}{2\epsilon}} \right), \\ &= N_\epsilon + C_{11} \sqrt{\sigma_0 L_1} + C_{12} \sqrt{\frac{L_1 \|\mathbf{z}_0 - \mathbf{x}^*\|^2}{2\epsilon}}, \end{aligned}$$

where we define $C_{11} = \frac{\alpha_1 + 2}{\alpha_1} \frac{1}{1 - \sqrt{\beta}}$ and $C_{12} = \frac{\alpha_1 + 2}{\alpha_1} \frac{1}{1 - \sqrt{\beta}} \frac{2(2 - \sqrt{\beta})}{\sqrt{\beta}(1 - \sqrt{\beta})}$. \square

Lemma 24. The total number of matrix-vector product evaluations in the SEP oracle is bounded by $\mathcal{O}(N_\epsilon^{1.25} (\log N_\epsilon)^{0.5} \log(\frac{\sqrt{d} N_\epsilon}{p}))$.

Proof. Note that we have $N_t \leq \left\lceil \frac{1}{4\sqrt{2}\delta_t} \log \frac{44d}{q_t^2} + \frac{1}{2} \right\rceil$ in Subroutine 4, where $\delta_t = 1/(\sqrt{t+2} \log(t+2))$ and $q_t = p/(2.5(t+1) \log^2(t+1))$. Thus, we have

$$N = \sum_{t=0}^{T-1} N_t \leq \sum_{t=0}^{T-1} \frac{(t+2)^{0.25} \log^{0.5}(t+2)}{2\sqrt{2}} \log \frac{2.5\sqrt{44d}(t+1) \log^2(t+1)}{p} \quad (98)$$

$$= \mathcal{O} \left(N_\epsilon^{1.25} \sqrt{\log N_\epsilon} \log \frac{\sqrt{d} N_\epsilon}{p} \right). \quad (99)$$

\square

F Experiments

In our experiments, we consider the logistic regression problem. Below we provide more details about the data generation scheme as well as the implementation of Nesterov's accelerated gradient method, BFGS, and our proposed A-QPNE algorithm.

Dataset generation. The dataset consists of n data points $\{(\mathbf{a}_i, y_i)\}_{i=1}^n$, where $\mathbf{a}_i \in \mathbb{R}^d$ is the i -th feature vector and $y_i \in \{-1, 1\}$ is its corresponding label. The labels $\{y_i\}_{i=1}^n$ are generated by

$$y_i = \text{sign}(\langle \mathbf{a}_i^*, \mathbf{x}^* \rangle), \quad i = 1, 2, \dots, n,$$

where $\mathbf{a}_i^* \in \mathbb{R}^{d-1}$ and $\mathbf{x}^* \in \mathbb{R}^{d-1}$ are the underlying true feature vector and the underlying true parameter, respectively. Moreover, each entry of \mathbf{a}_i^* and \mathbf{x}^* is drawn independently according to the standard normal distribution $\mathcal{N}(0, 1)$. Note that the true feature vectors $\{\mathbf{a}_i^*\}_{i=1}^n$ are not given in our dataset; instead, we generate $\{\mathbf{a}_i\}_{i=1}^n$ by adding noises and appending an extra dimension to $\{\mathbf{a}_i^*\}_{i=1}^n$. Specifically, we let $\mathbf{a}_i = [\mathbf{a}_i^* + \mathbf{n}_i + \mathbf{1}; 1]^\top \in \mathbb{R}^d$, where $\mathbf{n}_i \sim \mathcal{N}(0, \sigma^2 \mathbf{I})$ is the i.i.d. Gaussian noise

935 vector and $\mathbf{1} \in \mathbb{R}^{d-1}$ denotes the all-one vector. In our experiment, we set $n = 2,000$, $d = 150$ and
936 $\sigma = 0.8$.

937 **NAG.** We implemented a monotone variant of the Nesterov accelerated gradient method as described
938 in [43] Section 10.7.4]. Moreover, we determine the step size using a backtracking line search scheme.

939 **BFGS.** We implemented the classical BFGS algorithm, where the step size is determined by the
940 Moré–Thuente line search scheme.

941 **A-QPNE (our method).** We implemented our proposed A-QPNE method following the pseudocode
942 in Algorithm 1 where the line search scheme is given in Subroutine 1 and the Hessian approximation
943 update is given in Subroutine 2. Moreover, the implementations of the LinearSolver oracle and the
944 SEP oracle are given by Subroutines 3 and 4, respectively.