ON THE CONVERGENCE OF TSETLIN MACHINES FOR THE AND AND THE OR OPERATORS

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ABSTRACT

The Tsetlin Machine (TM) is an innovative machine learning algorithm rooted in propositional logic, achieving state-of-the-art performance in various pattern recognition tasks. While previous studies analyzed its convergence properties for the 1-bit and XOR operators, this work extends the analysis to the AND and OR operators, completing the study of fundamental digital operations. Our findings demonstrate that the TM almost surely converges to reproduce the AND and OR operators when trained on noise-free data over an infinite time horizon. Notably, the analysis of the OR operator uncovers a distinct property: the ability of the TM to represent two sub-patterns jointly within a single clause, contrasting with its behavior in the XOR case. Furthermore, we investigate the TM's behavior for AND/OR/XOR operators with noisy training samples, including mislabeled samples and irrelevant inputs. With wrong labels, the TM does not converge to the intended operators but can still learn efficiently. With irrelevant variables, the TM converges to the intended operators almost surely. Together, these analyses provide a comprehensive theoretical foundation for the TM's convergence properties across basic Boolean operators.

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1 INTRODUCTION

A Tsetlin Machine (TM) (Granmo, 2018) organizes clauses, each of which is associated with a 031 team of Tsetlin Automata (TAs) (Tsetlin, 1961), to collaboratively capture distinct sub-patterns for 032 a certain class. A TA, which is the core learning entity of TM, is a kind of learning automata (Zhang 033 et al., 2020; Yazidi et al., 2019; Omslandseter et al., 2022) that selects the current action based on 034 past experiences learned from the environment in order to obtain the maximum reward. In a TM, 035 a clause is a conjunction of literals, where a literal is a Boolean input or its negation. A clause is used to represent a sub-pattern. Once distinct sub-patterns are learned by a number of clauses, the 037 overall pattern recognition task is completed by a voting scheme from the clauses. The TM has 038 several advantages, such as transparent inference and learning (Bhattarai et al., 2024; Abeyrathna et al., 2023; Rafiev et al., 2022), and hardware friendliness (Maheshwari et al., 2023; Rahman et al., 2022; Morris et al., 2022). 040

041 The TM, together with its variations (Granmo et al., 2019; Abeyrathna et al., 2021; Dar-042 shana Abeyrathna et al., 2020; Sharma et al., 2023), has been employed in many applications, such 043 as word sense disambiguation (Yadav et al., 2021c), aspect-based sentiment analysis (Yadav et al., 044 2021b), novelty detection (Bhattarai et al., 2021), text classification (Yadav et al., 2021a) with en-045 hanced interpretability (Yadav et al., 2022), and solving contextual bandit problems (Seraj et al., 2022). These studies indicate that TMs obtain better or competitive performance compared with 046 most of the state-of-the-art techniques. At the same time, the transparency of learning is maintained 047 with smaller memory footprint and higher computational efficiency. 048

The TM convergence properties of the 1-bit operator and XOR operator were analyzed in (Zhang et al., 2022) and (Jiao et al., 2022), respectively. In (Zhang et al., 2022), TM's almost surely convergence to the identity/NOT operator with 1-bit input was confirmed, revealing the role of the hyperparameter s. In (Jiao et al., 2022), TM's convergence to the XOR operator with 2-bit input was proven, highlighting the functionality of the hyperparameter T. In this paper, we first focus on analyzing the AND and OR operators in the noise-free training samples, followed by an examina-

tion of the convergence properties of AND, OR, and XOR with noisy training samples, including the presence of wrong labels and irrelevant inputs.

This paper differs from prior work in several key aspects. While (Zhang et al., 2022) used stationary distribution analysis of discrete-time Markov chains (DTMC), the current study focuses on absorbing states. For XOR (Jiao et al., 2022), where sub-patterns are bit-wise exclusive, TM learns and converges to sub-patterns individually. In contrast, the OR operator's sub-patterns share features (e.g., $[x_1 = 1, x_2 = 1]$ and $[x_1 = 1, x_2 = 0]$ share $x_1 = 1$), allowing joint representation. We show that TM can effectively learn and represent these shared features, making the convergence process distinct. Additionally, this paper examines the role of Type II feedback, omitted in the prior work, and analyzes convergence property under noise.

064 It is worth noting that learning 2-bit operators, both with and without noise, is a well-solved prob-065 lem with transparent solutions. Since the 1980s, numerous studies in concept learning and probably 066 approximately correct (PAC) learning have extensively explored this topic. For instance, it has 067 been shown in Valiant (1984); Haussler et al. (1994) that 2-DNF formulas are both properly and 068 efficiently PAC learnable, with sample complexity scaling logarithmically in the input dimension. 069 More generally, transparent algorithms for learning k-DNF formulas have been proposed in Marc-070 hand & Shawe-Taylor (2002), and the problem of learning conjunctions under noise has been studied 071 in Mansour & Parnas (1998). While many elegant methods exist for learning conjunctions or disjunctions, their existence does not necessarily imply that the TM converges to such operators in 072 the same manner. TM employs a unique approach, learning from samples to construct conjunctive 073 expressions and coordinating these expressions across various sub-patterns, which merits its own 074 dedicated investigation. 075

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2 NOTATIONS OF THE TM

To make the article self-contained, we present the notations of TM. For more details of the inference and training concept, please refer to Appendix A.

The input of a TM is denoted as $\mathbf{X} = [x_1, x_2, \dots, x_o]$, where $x_k \in \{0, 1\}, k = 1, 2, \dots, o$, and o is the number of features. A literal is either the x_k in the original form or its negation $\neg x_k$. A clause is a conjunction of literals, and each literal is associated with a TA. The TA is a 2-action learning automaton whose job is to decide whether to Include/Exclude its literal in/from the clause, and the decision is determined by the current state of the TA. A clause is associated with 2o TAs, forming a TA team. A TA team is denoted in general as $\mathcal{G}_j^i = \{TA_{k'}^{i,j}|1 \le k' \le 2o\}$, where k' is the index of the TA, j is the index of the TA team/clause (multiple TA teams form a TM), and i is the index of the TM/class to be identified (A TM identifies a class, multiple TMs identify multiple classes).

Suppose we are investigating the i^{th} TM whose job is to identify class *i*, and that the TM is composed of *m* TA teams. Then $C_j^i(\mathbf{X})$ can be used to denote the output of the j^{th} TA team, which is a conjunctive clause:

For training :
$$C_j^i(\mathbf{X}) = \begin{cases} \left(\bigwedge_{k \in \xi_j^i} x_k\right) \land \left(\bigwedge_{k \in \bar{\xi}_j^i} \neg x_k\right), & \text{for } \xi_j^i, \ \bar{\xi}_j^i \neq \emptyset, \\ 1, & \text{for } \xi_j^i, \ \bar{\xi}_j^i = \emptyset. \end{cases}$$
 (1)

For testing :
$$C_j^i(\mathbf{X}) = \begin{cases} \left(\bigwedge_{k \in \xi_j^i} x_k\right) \land \left(\bigwedge_{k \in \bar{\xi}_j^i} \neg x_k\right), & \text{for } \xi_j^i, \ \bar{\xi}_j^i \neq \emptyset, \\ 0, & \text{for } \xi_j^i, \ \bar{\xi}_j^i = \emptyset. \end{cases}$$
 (2)

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In Eqs. (1) and (2), ξ_j^i and $\bar{\xi}_j^i$ are defined as the sets of indexes for the literals that have been included in the clause. ξ_j^i contains the indexes of included original (non-negated) inputs, x_k , whereas $\bar{\xi}_j^i$ contains the indexes of included negated inputs, $\neg x_k$.

Each clause represents a sub-pattern associated with class i by including a literal (a feature or its negation) if it contributes to the sub-pattern, or excluding it when deemed irrelevant. Multiple

clauses, i.e., the TA teams, are assembled into a complete TM to sum up the outputs of the clauses $f_{\sum}(\mathcal{C}^{i}(\mathbf{X})) = \sum_{j=1}^{m} C_{j}^{i}(\mathbf{X})$, where $\mathcal{C}^{i}(\mathbf{X})$ is the set of clauses for class *i*. The output of the TM

is further determined by the unit step function: $\hat{y}^i = \begin{cases} 0, & \text{for } f_{\sum}(\mathcal{C}^i(\mathbf{X})) < Th \\ 1, & \text{for } f_{\sum}(\mathcal{C}^i(\mathbf{X})) \ge Th \end{cases}$, where Th is a predefined threshold for classification. This is indeed a voting scheme. For example, the classifier $(x_1 \land \neg x_2) + (\neg x_1 \land x_2)$ captures the XOR-relation when Th = 1, meaning if any sub-pattern is satisfied, the input will be identified as following the XOR logic.

Note that the TM can assign a polarity to each TA team (Granmo, 2018), and one can refer to
 Appendix A for more information. In this study, for ease of analysis, we consider only positive
 polarity clauses. Nevertheless, this does not change the nature of TM learning.

For training, the labeled data ($\mathbf{X} = [x_1, x_2, ..., x_o]$, y^i) is given to TM, and each TA is guided by Type I and Type II Feedback defined in Tables 1 and 2, respectively. Type I Feedback is triggered when the training sample has a positive label: $y^i = 1$, while Type II feedback is utilized when $y^i = 0$. The parameter, s, controls the granularity of the clauses. NA in these tables means not applicable. Examples demonstrating TA state transitions per feedback tables can be found in Section 3.1 in (Zhang et al., 2022). In brief, Type I feedback is to reinforce true positive and Type II feedback is to fight against false negative.

Value of t	Value of the clause $C_i^i(\mathbf{X})$		1	0	
Value of th	e Literal $x_k/\neg x_k$	1	0	1	0
	P(Reward)	$\frac{s-1}{s}$	NA	0	0
Include Litera	P(Inaction)	$\frac{1}{s}$	NA	$\frac{s-1}{s}$	$\frac{s-1}{s}$
	P(Penalty)	Ŏ	NA	$\frac{1}{s}$	$\frac{1}{s}$
	P(Reward)	0	$\frac{1}{s}$	$\frac{1}{s}$	$\frac{1}{s}$
Exclude Litera	P(Inaction)	$\frac{1}{s}$	$\frac{s-1}{s}$	$\frac{s-1}{s}$	$\frac{s-1}{s}$
	P(Penalty)	$\frac{s-1}{s}$	Ŏ	Ŏ	Ŏ

Table 1: Type I Feedback — Feedback upon receiving a sample with label $y^i = 1$ (Granmo, 2018).

Value of the clause $C_i^i(\mathbf{X})$			1	0		
Value of the L	Literal $x_k / \neg x_k$	1	0	1	0	
	P(Reward)	0	NA	0	0	
Include Literal	P(Inaction)	1.0	NA	1.0	1.0	
	P(Penalty)	0	NA	0	0	
	P(Reward)	0	0	0	0	
Exclude Literal	P(Inaction)	1.0	0	1.0	1.0	
	P(Penalty)	0	1.0	0	0	

Table 2: Type II Feedback — Feedback upon receiving a sample with label $y^i = 0$ (Granmo, 2018).

To avoid the situation that a majority of the TA teams learn only a subset of sub-patterns, forming an incomplete representation¹, the hyperparameter T is used to regulate the resource allocation. The strategy works as follows (Granmo, 2018):

Generating Type I Feedback. If the label of the training sample X is $y^i = 1$, we generate, in probability, *Type I Feedback* for each clause $C_j^i \in C^i$ according to:

$$u_1 = \frac{T - \max(-T, \min(T, f_{\sum}(\mathcal{C}^i(\mathbf{X}))))}{2T}.$$
(3)

¹For example, for the XOR operator, one should avoid the situation that a majority of TA teams converge to $\neg x_1 \land x_2$ to represent the sub-pattern of [0, 1], and ignore the other sub-pattern [1, 0].

162 Generating Type II Feedback. If the label of the training sample X is $y^i = 0$, we generate, again, in probability, *Type II Feedback* to each clause $C_j^i \in C^i$ according to:

$$u_2 = \frac{T + \max(-T, \min(T, f_{\sum}(\mathcal{C}^i(\mathbf{X}))))}{2T}.$$
(4)

Briefly speaking, when the number of clauses representing one sub-pattern reaches T, learning from that sub-pattern will stop as the probability of triggering update is 0 according to Eq. (3) for positive polarity. The same concept applies according to Eq. (4) for negative polarity.

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3 CONVERGENCE ANALYSIS OF THE AND OPERATOR

A TM has converged when the states of its TAs do not change any longer. We assume that the training samples are noise free, i.e., $P(y = 1|x_1 = 1, x_2 = 1) = 1$, $P(y = 0|x_1 = 0, x_2 = 1) = 1$, $P(y = 0|x_1 = 1, x_2 = 0) = 1$, $P(y = 0|x_1 = 1, x_2 = 0) = 1$. We also assume the training samples are independently drawn at random, and the above four cases will appear with non-zero probability, which means that all of the four types of samples will appear for infinite times.

178 Because the considered AND operator has only one pattern of input, i.e., $x_1 = 1, x_2 = 1$, that will 179 trigger a true output, we employ one clause in this TM, and we thus can ignore the indices of the classes and the clauses in our notation in the proof. After simplification, $TA_k^{i,j}$ becomes TA_k , and 181 C_1^1 becomes C. Since there are two input parameters, namely x_1 and x_2 , we implement four TAs 182 in the clause, i.e., TA_1 , TA_2 , TA_3 , and TA_4 . TA_1 has two actions, i.e., including or excluding x_1 . 183 Similarly, TA₂ corresponds to including or excluding $\neg x_1$. TA₃ and TA₄ determine the behavior of 184 x_2 and $\neg x_2$, respectively. Once the TM can converge correctly to the intended operation, the actions 185 of TA₁, TA₂, TA₃, and TA₄ should be I, E, I, and E. Here we use "I" and "E" as abbreviations for 186 include and exclude respectively.

Theorem 1. Any clause will converge almost surely to $x_1 \wedge x_2$ given noise free AND training samples in infinite time when $u_1 > 0$ and $u_2 > 0$.

Due to page limit, the complete proof of Theorem 1 can be found in Appendix B. We here outline the concept and the main steps of the proof.

The condition $u_1 > 0$ and $u_2 > 0$ guarantees that all types of samples are always given and no specific type is blocked by Eqs. (3) and (4) during training. The goal of the proof is to show that the system transitions will guarantee that there is a unique absorbing state of the TM and the absorbing state has the actions of TA₁, TA₂, TA₃, and TA₄ to be I, E, I, E, respectively, corresponding to the propositional expression $x_1 \wedge x_2$.

To simplify the complex analysis of joint TA transitions, we use quasi-stationary analysis by freezing the transitions of the TAs for the first input bit and focusing on the transitions of the second input bit. Clearly, there are four possibilities for the first bit x_1 . We name them as <u>cases</u>, as: **Case 1**: TA₁ = E, TA₂ = I, i.e., include $\neg x_1$. **Case 2**: TA₁ = I, TA₂ = E, i.e., include x_1 . **Case 3**: TA₁ = E, TA₂ = E, i.e., exclude both x_1 and $\neg x_1$. **Case 4**: TA₁ = I, TA₂ = I, i.e., include both x_1 and $\neg x_1$.

203 In each of the above four cases, we analyze individually the transition of TA_3 with a given current 204 action, for different actions of TA₄, and vice versa. We index the possibilities as situations: Situ-205 ation 1. We study the transition of TA_3 when it has "Include" as its current action, given different 206 actions of TA_4 (i.e., when the action of TA_4 is frozen as "Include" or "Exclude"). Situation 2. We study the transition of TA_3 when it has "Exclude" as its current action, given different actions of 207 TA_4 (i.e., when the action of TA_4 is frozen as "Include" or "Exclude"). Situation 3. We study the 208 transition of TA_4 when it has "Include" as its current action, given different actions of TA_3 (i.e., 209 when the action of TA_3 is frozen as "Include" or "Exclude"). Situation 4. We study the transition 210 of TA_4 when it has "Exclude" as its current action, given different actions of TA_3 (i.e., when the 211 action of TA_3 is frozen as "Include" or "Exclude"). 212

Within each of the situation, there are 8 possible <u>instances</u>, determined by 4 possible combinations of the input samples of x_1 and x_2 , and the two possible TA actions, Include and Exclude. As an example, we randomly select an instance in Case 1, Situation 1. The selected instance is when the training sample is ($[x_1 = 1, x_2 = 1], y = 1$), and TA₄ is E. For this instance, the training sample will trigger Type I feedback because y = 1. Based on the current status of the TAs, the clause is in the form $C = \neg x_1 \land x_2$, with value 0. In Situation 1, the studied TA is TA₃, its corresponding literal is thus x_2 , with value 1. Given y = 1, clause value 0, literal value 1, we go to Table 1, the third column of transition probabilities for "Include Literal", and find the transition of TA_3 to be: the penalty probability $\frac{1}{s}$ and the inaction probability $\frac{s-1}{s}$. To indicate the transitions of TA_3 , we have plotted the transition diagram in Fig. 1, where P and R represent Reward and Penalty respectively. Note that the overall transition probability is $u_1 \frac{1}{s}$, where u_1 is defined in Eq. (3). Here, we have assumed $u_1 > 0$.

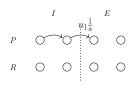


Figure 1: Transition of TA_3 when its current action is Include, TA_1 , TA_2 , and TA_4 's actions are Exclude, Include, and Exclude, respectively, upon a training sample $(x_1 = 1, x_2 = 1, y = 1)$.

To complete the quasi-stationary analysis of TA_3 and TA_4 , we must in total analyze $4 \times 4 \times 8 = 128$ transition instances, similar to the diagram in Fig. 1.

Based on the analysis of the 128 transition instances, we can summarize the transitions of TA_3 and TA_4 . By observing the transition directions, we can conclude that there is a unique absorbing state for TA_3 and TA_4 , given TA_1 and TA_2 being frozen as I, and E respectively. The absorbing state is when TA₃ and TA₄ are in I and E respectively. Once this step is completed, we must freeze TA₃ and TA_4 , and study the transitions of TA_1 and TA_2 in the same way. Thereafter, we can conclude that the system has a unique absorbing state, which is TA₁, TA₂, TA₃, and TA₄ being in I, E, I, E respectively, in the full dynamics of the system.

CONVERGENCE ANALYSIS OF THE OR OPERATOR

We assume the training samples for the OR operator are noise free (i.e., Eq. (5)), and are independently drawn at random. All those four cases will appear with non-zero probability.

$$P(y = 1|x_1 = 1, x_2 = 1) = 1, P(y = 1|x_1 = 0, x_2 = 1) = 1,$$

$$P(y = 1|x_1 = 1, x_2 = 0) = 1, P(y = 0|x_1 = 0, x_2 = 0) = 1.$$
(5)

Theorem 2. The clauses in a TM can almost surely learn the 2-bit OR logic given noise free training samples (shown in Eq. (5)) in infinite time, when $T \leq \lfloor \frac{m}{2} \rfloor$.

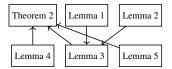


Figure 2: The dependence for the proof of the Theorem 2.

The proof of the theorem requires Lemma 1-Lemma 5 and their dependence is shown in Fig. 2. Clearly, there are three sub-patterns for the OR operator. In Lemma 1, we will show that any clause is able to converge to an intended sub-pattern when the training sample of only one sub-pattern is given, and when $u_1 > 0$ and $u_2 > 0$. In Lemma 2, we will show that the TM will become recurrent (not absorbing) when more sub-patterns jointly appear in the training samples and when $u_1 > 0$ and $u_2 > 0$. These two lemmas will be utilized in the proof of Lemma 3. Lemma 2 also reveals the recurrent nature of TM for the OR operator when the functionality of T is not enabled, i.e., when $u_1 > 0$ and $u_2 > 0$. This confirms the necessity of enabling the functionality of T in order to converge to an absorbing state that fulfills the OR operator, to be indicated by Lemma 3-Lemma 5. Specifically, Lemma 3-Lemma 5 analyze the system behavior when T is enabled and how T should be configured for the TM to converge to the OR operator. They guarantee that when the system

270 arrives an absorbing state, any sample from the intended sub-patterns will offer a vote sum no less 271 than T while the sample from the unintended sub-pattern has a vote sum 0. Then the OR operator 272 can be inferred by setting Th = T. In what follows, we will present and prove the lemmas.

273 **Lemma 1.** For any one of the three sub-patterns of x_1 and x_2 resulting in y = 1, shown in Eqs. 274 (6)-(8), the TM can converge to the intended sub-pattern when noise free training samples following 275 this sub-pattern are given, and when $u_1 > 0$, $u_2 > 0$. 276

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 $P(y = 1 | x_1 = 1, x_2 = 1) = 1, P(y = 0 | x_1 = 0, x_2 = 0) = 1,$ $P(y = 1 | x_1 = 0, x_2 = 1) = 1, P(y = 0 | x_1 = 0, x_2 = 0) = 1.$ (6)

$$P(y = 1|x_1 = 0, x_2 = 1) = 1, P(y = 0|x_1 = 0, x_2 = 0) = 1,$$
 (7)

$$P(y=1|x_1=1, x_2=0) = 1, P(y=0|x_1=0, x_2=0) = 1.$$
 (8)

282 The proof of Lemma 1 involves demonstrating convergence for three sub-patterns: those governed 283 by Eqs. (6), (7), and (8). These analyses build upon the convergence proofs for the XOR and AND 284 operators. For the sub-pattern in Eq. (6), transition diagrams in Appendix B confirm that the TAs 285 converge to $TA_1 = I$, $TA_2 = E$, $TA_3 = I$, and $TA_4 = E$, when input samples $[x_1 = 0, x_2 = 1]$ and 286 $[x_1 = 1, x_2 = 0]$ are excluded. The other two sub-patterns are proven using similar principles. Full 287 details are provided in Appendix C.

288 From Lemma 1, we show that the clauses converge to the intended sub-pattern if the training sam-289 ples following this particular sub-pattern are given. From Lemma 2, we will show that the system 290 becomes recurrent if any two or more sub-patterns of training samples are given. Specifically, we 291 show the TM is recurrent given samples following Eq. (5) and Eqs. (9)-(11), when $u_1 > 0$, $u_2 > 0$. 292

$$P(y=1|x_1=1, x_2=1) = P(y=1|x_1=1, x_2=0) = P(y=0|x_1=0, x_2=0) = 1, \quad (9)$$

$$P(y=1|x_1=1, x_2=1) = P(y=1|x_1=0, x_2=1) = P(y=0|x_1=0, x_2=0) = 1, (10)$$

$$P(y=1|x_1=1, x_2=0) = P(y=1|x_1=0, x_2=1) = P(y=0|x_1=0, x_2=0) = 1.$$
(11)

Lemma 2. The TM becomes recurrent if any two or more of the three sub-patterns jointly appear in the training samples, as shown in Eqs. (5), (9)-(11), when $u_1 > 0$, $u_2 > 0$.

299 **Proof of Lemma 2:** To show the recurrent property when samples following Eq. (9) are given, we 300 need to show that the absorbing states for Eq. (6) disappear when $([x_1 = 1, x_2 = 0], y = 1)$ is given 301 in addition, and the same applies for Eq. (8) when $([x_1 = 1, x_2 = 1], y = 1)$ is given.

302 We first show that the absorbing state of $TA_1 = I$, $TA_2 = E$, $TA_3 = I$, $TA_4 = E$, for sub-pattern 303 $[x_1 = 1, x_2 = 1], y = 1$) as shown in Eq. (6), disappears when sub-pattern $[x_1 = 1, x_2 = 0], y = 1$ 304 1) is given in addition. Indeed, TA₃ will move toward E when $([x_1 = 1, x_2 = 0], y = 1)$ is given, 305 because a penalty is given to TA_3 as shown in Fig. 3. 306

		Ι	1	E
			$u_{1s}^{\frac{1}{s}}$	
P	0	-10-	70	0
R	0	0	0	0

Figure 3: Transition of TA_3 when its current action is Include, TA_1 , TA_2 , and TA_4 's actions are 313 Include, Exclude, and Exclude, respectively, upon a training sample ($x_1 = 1, x_2 = 0, y = 1$). 314

Clearly, when $([x_1 = 1, x_2 = 0], y = 1)$ is given in addition, TA₃ has a non-zero probability to 316 move towards "Exclude". Therefore, "Include" is not the only direction that TA_3 moves to upon 317 the new input. In other words, $([x_1 = 1, x_2 = 0], y = 1)$ will make the state $TA_1 = I$, $TA_2 = E$, 318 $TA_3 = I$, $TA_4 = E$, not absorbing any longer. For other states, the newly added training sample 319 will not remove any transition from the previous case. For this reason, the system will not have any 320 new absorbing state. Therefore, when $([x_1 = 1, x_2 = 0], y = 1)$ is given in addition, the absorbing 321 state disappears and the system will not have any new absorbing state. 322

Following the same concept, we show that the absorbing state for $([x_1 = 1, x_2 = 0], y = 1)$ 323 shown in Eq. (8), i.e., $TA_1 = I$, $TA_2 = E$, $TA_3 = E$, $TA_4 = I$, disappears when sub-pattern

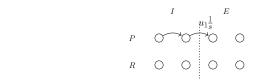


Figure 4: Transition of TA₄ when its current action is Include, TA₁, TA₂, and TA₃'s actions are Include, Exclude, and Exclude, respectively, upon a training sample ($x_1 = 1, x_2 = 1, y = 1$).

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 $([x_1 = 1, x_2 = 1], y = 1)$ is given in addition. Indeed, TA₄ will also move towards E when $([x_1 = 1, x_2 = 1], y = 1)$ is given, as shown in Fig. 4.

Understandably, because of the newly added sub-patterns, the absorbing states in Eqs. (6) and (8)
 disappear and no new absorbing states are generated. In other words, the TM trained based on
 samples from Eq. (9) becomes recurrent.

Following the same concept, we can show that the system becomes recurrent for Eqs. (5), (10), and (11) as well. For the sake of conciseness, we will not provide the details here. In general, any newly added sub-pattern will involve a probability for the learnt sub-pattern to move outside the learnt state, making the system recurrent.

Lemma 2 tells us that if we always give TM the training samples from all sub-patterns without blocking the learnt patterns by using T via Eqs. (3) and (4), the system is recurrent. In other words, if we want to have the TM converge to the OR operator in an absorbing state, it is critical to utilize the feature of T to block any incoming training samples from updating the learnt sub-patterns. Specifically, we need to configure T (1) so that the absorbing states exist and (2) confirm that the absorbing states follows the OR operator. In what follows, we will, through Lemmas 3-5, show how T via Eqs. (3) and (4) can guarantee the convergence and how the value of T should be configured.

351 Let's revisit the functionality of T. T can block the training samples from updating a learnt sub-352 pattern (clauses that have converged to one of the absorbing states) so that the clauses that have not 353 converged can be guided to learn the other unlearned sub-patterns. More specifically, if the vote sum of the clauses reaches T for a certain sub-pattern, the new training samples of this sub-pattern will 354 be blocked by the TM. There are three sub-patterns in OR operator. When the sum of clauses for 355 each of the three sub-patterns reaches T, all training samples for Type I feedback are blocked. At 356 the same time, if all samples for Type II feedback will not trigger any update to the states of TAs, 357 the TM is absorbed. In Lemma 3, we detail the necessity and sufficiency of the absorbing state. 358

Lemma 3. The system is absorbed if and only if (1) the vote sum of any sample from intended sub-patterns reaches T, i.e., $f_{\Sigma}(C_i(\mathbf{X})) = T$, $\forall \mathbf{X} = [x_1 = 0, x_2 = 1]$ or $[x_1 = 0, x_2 = 1]$ or $[x_1 = 0, x_2 = 1]$, and (2) no clause is formed only by a negated literal or negated literals.

Proof of Lemma 3: In Lemma 2, the TM is recurrent if the functionality of T is disabled (i.e., $u_1 > 0, u_2 > 0$). Therefore, for the OR operator to converge, the functionality of T is critical to block any feedback in order to form an absorbing state.

By design, TM will either be updated via Type I feedback or Type II feedback. We show via (1) the condition when Type I feedback is blocked and then show via (2) when any update from Type II feedback is not given. When both types of feedback are blocked, the system will not be updated anymore and thus absorbed.

To prove (1) in Lemma 3, we show that the system is not absorbed when 0 or 1 intended sub-pattern is blocked by T. When 2 intended sub-patterns are blocked, the system will guide the clauses to learning the remaining intended sub-pattern. Only when all 3 intended sub-patterns are blocked by T, the system will stop updating based on Type I feedback.

Clearly, when no intended sub-pattern is blocked by T, the training samples given to the system follow Eq. (5). Following this type of training samples, it has already been shown in Lemma 2 that the TM is recurrent. When only 1 intended sub-pattern is blocked by T, the system is updated based on Eqs. (9), (10), or (11), which is also recurrent. We look at the cases when two intended sub-patterns are blocked by T but the third is not blocked. In other words, the vote sum for any two intended sub-patterns reaches at least T, and the sum for the remaining sub-pattern is less than T. In this case, only one type of the samples from Eqs. (6) or (7) or (8) will be given to the TM. Based on Lemma 1, we understand that all clauses, including the ones that follow the two blocked sub-patterns, will be reinforced to learn the unblocked sub-pattern. This is due to the fact that only the samples following the unblocked sub-pattern are given to the TMs. In this circumstance, as soon as the unblocked sub-pattern also has T clauses, i.e., when all three sub-patterns are blocked by T at the same time, Type I feedback will be blocked completely.

- 386 Note that the samples from the unblocked sub-pattern will encourage the learnt clauses (the clauses 387 that follow sub-patterns with vote sum T) move out from the learnt sub-patterns, making the vote sum of learnt sub-patterns being lower than T again. If this happens before the vote sum of the 388 to-be-learnt sub-pattern reaches T, two sub-patterns will be unblocked and the system becomes one 389 of three cases described by Eqs. (9), (10) or (11). In other words, even if an absorbing state exists 390 when two intended sub-patterns are already blocked by T, the system may not monotonically move 391 towards the absorbing state. Nevertheless, as soon as all three intended sub-patterns are blocked by 392 reaching T, the Type I feedback will be blocked. 393
- 394 Here we prove (2) in Lemma 3. Type II feedback is only triggered by training sample ($[x_1 = 0, x_2]$ 395 $x_2 = 0$], y = 0 in the OR operator. For Type II feedback, based on Table 2, any transition is only triggered as a penalty when excluded literal has 0 value and the clause is evaluated as 1. Specifically 396 for the OR operation, this only happens when $C = \neg x_1 \land \neg x_2$ or $C = \neg x_1$ or $C = \neg x_2$. For 397 $C = \neg x_1 \land \neg x_2$, based on the Type II feedback, the TA with the action "excluding x_1 " and the TA 398 with the action "excluding x_2 " will be penalized. In other words, the actions of the TAs for x_1 and 399 x_2 will be encouraged to move from exclude to include side. As soon as any one of TAs for x_1 or 400 x_2 (or occasionally both of them) becomes included, the clause will become $C = x_1 \wedge \neg x_1 \wedge \neg x_2$ 401 or $C = \neg x_1 \land x_2 \land \neg x_2$ (or occasionally $C = x_1 \land \neg x_1 \land x_2 \land \neg x_2$). In this case, input $[x_1 = 0, x_1 \land x_2 \land \neg x_2]$ 402 $x_2 = 0$ will always result in 0 as the output of the clause and then the Type II feedback will not 403 update the system any longer. Following the same concept, for $C = \neg x_2$, the Type II feedback will 404 encourage the excluded x_1 to be included so that the clause becomes $C = x_1 \wedge \neg x_2$. The same 405 applies to $C = \neg x_1$, which will eventually become $C = \neg x_1 \land x_2$ upon Type II feedback. When all clauses in $C = \neg x_2$ or $C = \neg x_1$ are also updated to $C = x_1 \land \neg x_2$ or $C = \neg x_1 \land x_2$, no Type II 406 feedback is triggered up on any input sample. 407
- 408 We summarize the requirements for an absorbing state:

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- For any sample following $\mathbf{X} = [x_1 = 1, x_2 = 1]$, or $\mathbf{X} = [x_1 = 1, x_2 = 0]$, or $\mathbf{X} = [x_1 = 0, x_2 = 1]$, the vote sum of clauses, i.e., $f_{\Sigma}(C^i(\mathbf{X}))$ must be at least T, no matter in which form the clauses are constructed. This will block any Type I feedback.
- There are no clauses with literal(s) in only negated form, such as $C = \neg x_1$ or $C = \neg x_2$ or or $C = \neg x_1 \land \neg x_2$. This guarantees no transition happens upon any Type II feedback.
- In Lemma 3, we find the conditions of the absorbing state. In the next Lemma, we will show how to set up the value of T so that the vote sum for each intended sub-pattern can indeed reach T.
- Lemma 4. $T \leq \lfloor m/2 \rfloor$ is required so that the vote sum of any sample from intended sub-patterns can reach T.
- 422 **Proof of Lemma 4:** There are three intended sub-patterns in the OR operator. Given m clauses in total, to make sure each one has at least T votes, we have $3T \leq m$. This requires $T \leq \lfloor m/3 \rfloor$ 423 for any integer. However, the nature of the OR operator offers the possibility to represent 2 sub-424 patterns jointly. For example, T clauses in the form of x_1 will result in the vote sum as T for both 425 $[x_1 = 1, x_2 = 0]$ and $[x_1 = 1, x_2 = 1]$. If there are other T clauses to represent the remaining 426 sub-pattern, in total 2T clauses can offer the vote sum as T for all intended sub-patterns. We thus 427 have $2T \le m$, giving $T \le |m/2|$ for any integer. Note that the fact that two sub-patterns can be 428 jointly represented has been observed and confirmed in experiments shown in Appendix F. 429
- 430 When we have a smaller T, different sub-patterns may be represented by distinct clauses, offering 431 more flexibility. However, when $T > \lfloor m/2 \rfloor$, there will always be one or two sub-patterns that cannot obtain a sum of T clauses. For this reason, the maximum integer value is $T = \lfloor m/2 \rfloor$.

In Lemma 5, we show that the input sample $[x_1 = 0, x_2 = 0]$ will not give a vote sum greater than or equal to T. This is to avoid any possible false positive upon input $[x_1 = 0, x_2 = 0]$ in testing.

Lemma 5. When absorbing, the sample from unintended sub-pattern, i.e., $[x_1 = 0, x_2 = 0]$, will not give any vote sum greater than or equal to T.

437 Proof of Lemma 5: Obviously, to have a positive output form $[x_1 = 0, x_2 = 0]$, the clause should **438** be in the form of $C = \neg x_1$ or $C = \neg x_2$ or $C = \neg x_1 \land \neg x_2$. It has already shown in the proof of **439** Lemma 3 that Type II feedback will eliminate such clauses. In fact, when the system is absorbed, **440** no clause will be in the form of $C = \neg x_1$ or $C = \neg x_2$ or $C = \neg x_1 \land \neg x_2$. For this reason, **441** $[x_1 = 0, x_2 = 0]$ will never result in a sum of clause outputs greater than or equal to T.

Proof of Theorem 2: Based on Lemma 3–Lemma 5, we understand that if $T \le \lfloor m/2 \rfloor$ holds, Type I feedback will eventually be blocked and Type II feedback will eventually only give "inaction" feedback. In this situation, no actual transition will be triggered and thus the system reaches the absorbing state. Before absorbed, the system moves back and forth in the intermediate states. Once absorbed, samples from any one of the intended sub-patterns will result in a vote sum to no less than T and the unintended sub-pattern will have a vote sum to 0. We thus have the OR logic almost surely by setting a threshold Th = T and conclude the proof.

Now let's study a simple example with m = 2, T = 1. Here, $C_1 = x_1$ and $C_2 = x_2$ can be an instance for an absorbing case. $C_1 = x_1$ and $C_2 = \neg x_1 \land x_2$ also works. Clearly, the clauses can be in various forms, as long as the conditions in Lemma 3 fulfill. These converged clauses are not necessarily in the exact form of the three sub-patterns, which is distinct to that of the XOR operator.

453 **Remark 1.** Although both AND and OR operators converge, the approaches are different. For AND 454 operator, the system is converged because the clauses become eventually absorbed to the intended 455 pattern upon Type I and Type II feedback, even if the functionality of T is disabled $(u_1 > 0$ and $u_2 > 0$). As the TM enables the functionality of T by default, the system will be absorbed when 456 T clauses converge to $x_1 \wedge x_2$, before all clauses converge to this pattern. However, for the OR 457 operator, the functionality of T is critical because the TM is recurrent if $u_1 > 0$ and $u_2 > 0$. 458 The absorbing state of the OR operator is achieved because the functionality of T blocks all Type I 459 feedback and Type II feedback gives only "Inaction" feedback. The concept of convergence for the 460 OR operator is similar to that of XOR, but the form of clauses after absorbing varies due to the 461 possible joint representation of sub-patterns in OR. 462

Remark 2. When T is greater than half of the number of the clauses, i.e., $T > \lfloor m/2 \rfloor$, the system will not have an absorbing state. We conjecture that the system can still learn the sub-patterns in an unbalanced manner, as long as T is not configured too close to the total number of clauses m.

Given $T > \lfloor m/2 \rfloor$, Type I feedback cannot be completely blocked and the TM is recurrent. Nevertheless, if T is not close to m, there will be clauses that possibly learn distinct sub-patterns. In addition, Type II feedback can avoid the form of $C = \neg x_1$ or $C = \neg x_2$ or $C = \neg x_1 \land \neg x_2$ from happening. Therefore, with Th > 0, the TM may still learn the OR operator with high probability.

To validate the theoretical analyses, we present in Appendix F the experiment results² for both the AND and the OR operators, confirming the correctness of the above theorems.

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5 REVISIT THE XOR OPERATOR

Let us revisit the proof of XOR operator. As stated in (Jiao et al., 2022), when the system is absorbed, 476 the clauses follow the format $C = x_1 \wedge \neg x_2$ or $C = \neg x_1 \wedge x_2$ precisely. In other words, a clause 477 with just one literal, such as $C = x_1$, cannot absorb the system. The main reason is that the 478 sub-patterns in XOR operator are mutual exclusive, i.e., the sub-patterns cannot be merged in any 479 way. Although Type I feedback can be blocked when T clauses follow one sub-pattern using one 480 literal, the Type II feedback can reinforce the other missing literal to be included. For example, 481 when T clauses happens to converge to $C = x_1$, the Type I feedback from any input samples of 482 $([x_1 = 1, x_2 = 0], y = 1)$ will be blocked. In this situation, the unblocked Type II feedback from 483 $([x_1 = 1, x_2 = 1], y = 0)$ will encourage the clause to include $\neg x_2$. This is because upon a sample 484

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²The code for the experiments of this paper can be found at https://github.com/JaneGlim/ Convergence-of-Tsetline-Machine-for-the-AND-OR-operators.

486 $([x_1 = 1, x_2 = 1], y = 0)$, we have Type II feedback, $C = x_1 = 1$, and the studied literal is 487 $\neg x_2 = 0$. When the TA for excluding $\neg x_2$ is considered, a big penalty, i.e., 1, is given to the TA, 488 making it moving towards action *Included*, and thus $C = x_1$ eventually becomes $C = x_1 \land \neg x_2$. 489 Following the same concept, we can analyze the development for $C = \neg x_1, C = x_2$, and $C = \neg x_2$, 490 which will eventually converge to $C = \neg x_1 \land x_2$ or $C = x_1 \land \neg x_2$, upon Type II feedback.

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6 CONVERGENCE ANALYSIS UNDER RANDOM NOISE

494 We studied the convergence properties of AND, OR, and XOR operators under training samples 495 with noise following the noise type named noisy completely at random Frénay & Verleysen (2013), 496 categorized as wrong labels (in Appendix D) and irrelevant input variables (in Appendix E). A 497 wrong label refers to an input that should be labeled as 1 but is instead labeled as 0, or vice versa. 498 An irrelevant input variable, on the other hand, is one that does not contribute to the classification. 499 We demonstrate that, with wrong labels, the TM does not converge to the intended operators but 500 can still learn efficiently. With irrelevant variables, the TM converges to the intended operators almost surely. Experimental results confirmed these findings (Appendix G). We summarize the main 501 findings below. The proof and the experiment results can be found in the corresponding appendices. 502

Theorem 3. The TM is recurrent given training samples with wrong labels for the AND, OR, and
 XOR operators.

Remark 3. The recurrent property of TM indicates that there is a non-zero probability that it cannot learn the intended operator. The primary reason for the recurrent behavior when wrong labels are present is the statistically conflicting labels for the same input samples. These inconsistency causes the TAs within a clause to learn conflicting outcomes for the same input. When a clause learns to evaluate an input as 1 based on Type I feedback, samples with a label of 0 for the same input prompt it to learn the input as 0 during Type II feedback. This conflict in labels confuses the TM, leading to back-and-forth learning.

Remark 4. Note that although wrong labels will make the TM not converge (not absorbing with 100% accuracy for the intended logic), via experiments, we can still find that the TM are able to learn the operators efficiently, as shown in Appendix G. This property aligns with the concept of PAC learnable Mansour & Parnas (1998) or ϵ -optimality Zhang et al. (2020), although a formal proof remains an open question.

Theorem 4. The clauses in a TM can almost surely learn the 2-bit AND logic given training samples with k irrelevant input variables in infinite time, $0 < k < \infty$, when $T \le m$.

Theorem 5. The clauses in a TM can almost surely learn the 2-bit XOR and OR logic given training samples with k irrelevant input variables in infinite time, $0 < k < \infty$, when $T \leq \lfloor m/2 \rfloor$.

522 When the number of irrelevant variables is large, the training set may not cover all possible examples 523 due to the required exponential space. Although not yet theoretically proven, polynomial space for 524 training samples seems feasible for TM, which has been confirm by experiments (Appendix G.3). 525 This is because the TM can independently update the actions of a TA within a clause, as long as the 526 clause value and the literal value are determined by the training sample. In other words, once the clause value and the literal value are known, the transitions triggered by Type I and Type II feedback 527 are fully determined. As a result, the TM does not need to observe all possible combinations of 528 irrelevant inputs to learn effectively. Instead, as long as the statistical irrelevance of certain inputs is 529 demonstrated in the training samples, the corresponding TA transitions will be triggered accordingly. 530 This enables the TM to learn without requiring exhaustive coverage of the input space. 531

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7 CONCLUSIONS

In this article, we prove the convergence of the TM for the AND and OR operators with noise free
training samples. Our proof for the OR operator highlights the TM's ability to learn joint subpatterns, showcasing the efficiency of its learning process. Additionally, we analyze the behavior of
the TM for the AND, OR, XOR operators in the presence of random noise within the training data.
Combined with the convergence proofs in (Zhang et al., 2022) and (Jiao et al., 2022), this work
concludes the analysis of TM convergence for fundamental digital operators.

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648 A APPENDIX: BRIEF OVERVIEW OF THE TM

We present the basics of TM here. Those who already are familiar with the concept and notations of TM can ignore this appendix.

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A.1 BASIC CONCEPT OF THE TM

The input of a TM is denoted as $\mathbf{X} = [x_1, x_2, \dots, x_o]$, where $x_k \in \{0, 1\}, k = 1, 2, \dots, o$, and o is the number of features. A literal is either the x_k in the original form or its negation $\neg x_k$. A clause is a conjunction of literals, and each literal is associated with a TA. The TA is a 2-action learning automaton whose job is to decide whether to Include/Exclude its literal in/from the clause, and the decision is determined by the current state of the TA.

Figure 5 illustrates the structure of a TA with two actions and 2N states, where N is the number of states for each action. This study considers N as a finite number. When the TA is in any state between 0 to N - 1, the action "Include" is selected. The action becomes "Exclude" when the TA is in any state between N to 2N - 1. The transitions among the states are triggered by a reward or a penalty that the TA receives from the environment, which, in this case, is determined by different types of feedback defined in the TM (to be explained later).

A clause is associated with 2*o* TAs, forming a TA team. A TA team is denoted in general as $\mathcal{G}_j^i = \{\mathrm{TA}_{k'}^{i,j} | 1 \le k' \le 2o\}$, where k' is the index of the TA, j is the index of the TA team/clause (multiple TA teams form a TM), and i is the index of the TM/class to be identified (A TM identifies a class, multiple TMs identify multiple classes).

671 Suppose we are investigating the i^{th} TM whose job is to identify class *i*, and that the TM is composed 672 of *m* TA teams. Then $C_j^i(\mathbf{X})$ can be used to denote the output of the j^{th} TA team, which is a 673 conjunctive clause:

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For training :
$$C_j^i(\mathbf{X}) = \begin{cases} \left(\bigwedge_{k \in \xi_j^i} x_k\right) \land \left(\bigwedge_{k \in \bar{\xi}_j^i} \neg x_k\right), & \text{for } \xi_j^i, \ \bar{\xi}_j^i \neq \emptyset, \\ 1, & \text{for } \xi_j^i, \ \bar{\xi}_j^i = \emptyset. \end{cases}$$
 (12)

For testing :
$$C_j^i(\mathbf{X}) = \begin{cases} \begin{pmatrix} \bigwedge_{k \in \xi_j^i} x_k \\ 0, & \end{cases} \land \begin{pmatrix} \bigwedge_{k \in \bar{\xi}_j^i} \neg x_k \\ 0, & \text{for } \xi_j^i, \ \bar{\xi}_j^i \neq \emptyset, \end{cases}$$
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In Eqs. (12) and (13), ξ_j^i and $\bar{\xi}_j^i$ are defined as the sets of indexes for the literals that have been included in the clause. ξ_j^i contains the indexes of included original (non-negated) inputs, x_k , whereas $\bar{\xi}_j^i$ contains the indexes of included negated inputs, $\neg x_k$. Note that in propositional logic, an empty clause is typically defined as having a value of 1. However, empirical results indicate that TMs generally achieve higher test accuracy on new data when empty clauses are 0-valued. Therefore, during TM training, an "empty" clause outputs 1 to encourage the TAs to include literals, following the feedback mechanisms of the TM. In contrast, during TM testing, an "empty" clause outputs 0, indicating that it does not influence the final classification decision since it does not represent any specific sub-pattern.

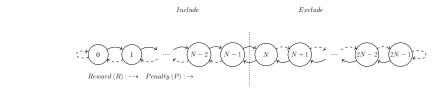




Figure 5: A two-action Tsetlin automaton with 2N states Jiao et al. (2022).

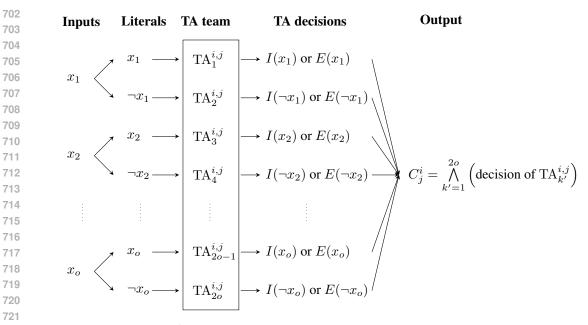


Figure 6: A TA team G_j^i consisting of 2*o* TAs Zhang et al. (2022). Here $I(x_1)$ means "include x_1 " and $E(x_1)$ means "exclude x_1 ".

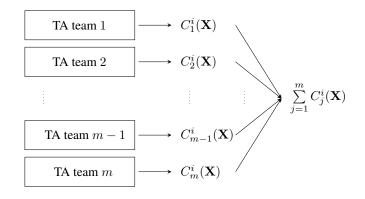


Figure 7: TM voting architecture Jiao et al. (2022).

Figure 6 illustrates the structure of a clause and its relationship to its literals. Here, for ease of notation, we define I(x) = x, $I(\neg x) = \neg x$, and $E(x) = E(\neg x) = 1$ in the analysis of the training procedure, with the latter meaning that an excluded literal does not contribute to the output.

Multiple clauses, i.e., the TA teams in conjunctive form, are assembled into a complete TM. There are two architectures for clause assembling: Disjunctive Normal Form Architecture and Voting Architecture. In this study, we focus on the latter one, as shown in Figure 7. The voting consists of summing the outputs of the clauses:

$$f_{\sum}(\mathcal{C}^{i}(\mathbf{X})) = \sum_{j=1}^{m} C_{j}^{i}(\mathbf{X}), \tag{14}$$

⁷⁵¹ where $C^i(\mathbf{X})$ is the set of trained clauses for class *i*.

The output of the TM, in turn, is decided by the unit step function:

$$\hat{y}^{i} = \begin{cases} 0 & \text{for } f_{\sum}(\mathcal{C}^{i}(\mathbf{X})) < Th \\ 1 & \text{for } f_{\sum}(\mathcal{C}^{i}(\mathbf{X})) \ge Th \end{cases}$$
(15)

where Th is a predefined threshold for classification. For example, the classifier $(x_1 \land \neg x_2) + (\neg x_1 \land x_2)$ captures the XOR-relation given Th = 1, meaning if any sub-pattern is satisfied, the input will be identified as following the XOR logic.

759 Note that for the voting architecture, the TM can assign a polarity to each TA team (Granmo, 2018). 760 Specifically, TA teams with odd indices have positive polarity, learning from training samples with 761 label 1, while those with even indices have negative polarity, learning from training samples with 762 label 0. The only difference between these polarities is that the output of a clause associated with an 763 even-indexed TA team will be flipped to its negative. The voting consists of summing the polarized 764 clause outputs, and the threshold Th is set to zero. For example, for the XOR operator with four 765 clauses, the learned clauses with positive polarity can be $C_1 = x_1 \wedge \neg x_2$ and $C_3 = \neg x_1 \wedge x_2$, while the ones with negative polarity can be $C_2 = x_1 \wedge x_2$ and $C_4 = \neg x_1 \wedge \neg x_2$. In this case, when the 766 testing sample $[x_1 = 1, x_2 = 0]$ arrives, the sum of the clause values is 1. On the contrary, when 767 the testing sample $[x_1 = 0, x_2 = 0]$ arrives, the sum of the clause values is -1. In this way, with 768 Th = 0, the system's decision range and tolerance is expected to be larger. 769

In this study, we consider only positive polarity clauses. The reason is two-folds: firstly, in the AND/OR case, once the TM has learned out the pattern that outputs 1, it also has learned the pattern that outputs 0, as they are complementary. Therefore, the learning/reasoning process of TM can be explained from the perspective of learning the pattern that outputs 1. Secondly, for the sake of easy analysis and better understanding.

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A.2 TRAINING PROCESS OF THE TM

777 The training process is built on letting all the TAs take part in a decentralized game. Training data 778 $(\mathbf{X} = [x_1, x_2, ..., x_o], y^i)$ is obtained from a data set \mathcal{S} , distributed according to the probability 779 distribution $P(\mathbf{X}, y^i)$. In the game, each TA is guided by Type I Feedback and Type II Feedback 780 defined in Table 3 and Table 4, respectively. Type I Feedback is triggered when the training sample 781 has a positive label, i.e., $y^i = 1$, meaning that the sample belongs to class i. When the training 782 sample is labeled as not belonging to class i, i.e., $y^i = 0$, Type II Feedback is utilized for generating 783 feedback. Examples demonstrating TA state transitions per feedback tables can be found in Section 784 3.1 in (Zhang et al., 2022). In brief, Type I feedback is to reinforce true positive and Type II feedback 785 is to fight against false negative.

The parameter, s, controls the granularity of the clauses and a larger s encourages more literals to be included in each clause. A more detailed analysis on parameter s can be found in (Zhang et al., 2022).

Value of the clause $C_i^i(\mathbf{X})$			l	0		
Value of the Literal $x_k/\neg x_k$			0	1	0	
	P(Reward)	$\frac{s-1}{s}$	NA	0	0	
Include Literal	P(Inaction)	$\frac{1}{s}$	NA	$\frac{s-1}{s}$	$\frac{s-1}{s}$	
	P(Penalty)	Ŏ	NA	$\frac{1}{s}$	$\frac{1}{s}$	
	P(Reward)	0	$\frac{1}{s}$	$\frac{1}{s}$	1	
Exclude Literal	P(Inaction)	1	$\frac{s-1}{s}$	$\frac{s-1}{s}$	$\frac{s-1}{s}$	
	P(Penalty)	$\frac{s-1}{s}$	Ő	Ő	Ő	

Table 3: Type I Feedback — Feedback upon receiving a sample with label y = 1, for a single TA to decide whether to Include or Exclude a given literal $x_k/\neg x_k$ into C_j^i . NA means not applicable (Granmo, 2018).

To avoid the situation that a majority of the TA teams learn only one sub-pattern (or a subset of sub-patterns) while ignore other sub-patterns, forming an incomplete representation³, the hyperparameter T is used to regulate the resource allocation. If the votes, i.e., the summation $f_{\sum}(\mathcal{C}^i(\mathbf{X}))$, for a certain sub-pattern **X** already reach a total of T or more, neither rewards nor penalties are provided to the TAs when more training samples of this particular sub-pattern are given. In this way, we

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³For example, for the XOR operator, we should avoid the situation that a majority of TA teams learn subpattern $x_1 = 0$ and $x_2 = 1$ and ignore sub-pattern $x_1 = 1$ and $x_2 = 0$, making the learning outcome biased/unbalanced. A proper configuration of T can avoid this situation.

810	Value of the	clause $C_i^i(\mathbf{X})$		1	0	
811	Value of the L	Value of the Literal $x_k/\neg x_k$			1	0
812		P(Reward)	0	NA	0	0
813	Include Literal	P(Inaction)	1.0	NA	1.0	1.0
814		P(Penalty)	0	NA	0	0
815		P(Reward)	0	0	0	0
816	Exclude Literal	P(Inaction)	1.0	0	1.0	1.0
817		P(Penalty)	0	1.0	0	0
818						

Table 4: Type II Feedback — Feedback upon receiving a sample with label y = 0, for a single TA to decide whether to Include or Exclude a given literal $x_k/\neg x_k$ into C_i^i . NA means not applica-ble (Granmo, 2018).

can ensure that each specific sub-pattern can be captured by a limited number, i.e., T, of available clauses, allowing sparse sub-pattern representations among competing sub-patterns. In more details, the strategy works as follows:

Generating Type I Feedback. If the label of the training sample X is $y^i = 1$, we generate, in probability, Type I Feedback for each clause $C_j^i \in C^i$. The probability of generating Type I Feedback is (Granmo, 2018):

$$u_1 = \frac{T - \max(-T, \min(T, f_{\sum}(\mathcal{C}^i(\mathbf{X}))))}{2T}.$$
(16)

Generating Type II Feedback. If the lable of the training sample X is $y^i = 0$, we generate, again, in probability, Type II Feedback to each clause $C_j^i \in C^i$. The probability is (Granmo, 2018):

$$u_{2} = \frac{T + \max(-T, \min(T, f_{\sum}(\mathcal{C}^{i}(\mathbf{X}))))}{2T}.$$
(17)

After Type I Feedback or Type II Feedback is generated for a clause, each individual TA within each clause is given reward/penalty/inaction according to the probability defined, and then the states of the corresponding TAs are updated.

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B APPENDIX: DETAILED PROOF OF THE CONVERGENCE OF THE AND OPERATOR

Proof: In this Appendix, we will prove Theorem 1. The condition $u_1 > 0$ and $u_2 > 0$ guarantees that all types of samples for AND operator, following Eq. (18), are always given and no specific type is blocked during training. The goal of the proof is to show that the system transitions will guarantee the actions of TA₁, TA₂, TA₃, and TA₄ to be I, E, I, E, and these actions correspond to the unique absorbing state of the system.

872	D(1 1 1)	(10)
873	$P(y=1 x_1=1, x_2=1) = 1,$	(18)
874	$P(y=0 x_1=0, x_2=1) = 1,$	
875	$P(y=0 x_1=1, x_2=0) = 1,$	
876	$P(y=0 x_1=0, x_2=0) = 1.$	
877		

In Subsections B.1, we will describe the transitions of the system in an exhaustive manner. Thereafter, in the Subsection B.2, we summarize the transitions in Subsection B.1 and reveal the absorbing state of the system, which is the intended AND operator.

B.1 THE TRANSITIONS OF THE TAS

In order to analyze the transitions of the system, we freeze the transition of the two TAs for the first bit of the input and study the transition of the second bit of input. Clearly, there are four <u>cases</u> for the first bit, x_1 , as:

- Case 1: $TA_1 = E$, $TA_2 = I$, i.e., include $\neg x_1$.
 - Case 2: $TA_1 = I$, $TA_2 = E$, i.e., include x_1 .
 - Case 3: $TA_1 = E$, $TA_2 = E$, i.e., exclude both x_1 and $\neg x_1$.
 - Case 4: $TA_1 = I$, $TA_2 = I$, i.e., include both x_1 and $\neg x_1$.

In what follows, we will analyze the transition of the TAs for x_2 , given the TAs of x_1 frozen in the above four distinct cases, one by one.

B.1.1 CASE 1: INCLUDE $\neg x_1$

In this subsection, we assume that the TAs for first bit is frozen as $TA_1 = E$ and $TA_2 = I$, and thus the overall joint actions of TAs for the first bit give " $\neg x_1$ ". In this case, we have 4 <u>situations</u> to study, detailed below:

- Situation1: We study the transition of TA₃ when it has "Include" as its current action, given different actions of TA₄ (i.e., when the action of TA₄ is frozen as "Include" or "Exclude".).
- Situation 2: We study the transition of TA_3 when it has "Exclude" as its current action, given different actions of TA_4 (i.e., when the action of TA_4 is frozen as "Include" or "Exclude".).
- Situation 3: We study the transition of TA_4 when it has "Include" as its current action, given different actions of TA_3 (i.e., when the action of TA_3 is frozen as "Include" or "Exclude".).
- Situation 4: We study the transition of TA_4 when it has "Exclude" as its current action, given different actions of TA_3 (i.e., when the action of TA_3 is frozen as "Include" or "Exclude".).
- In what follows, we will go through, exhaustively, the four situations.
- 915 B.1.1.1 Study TA₃ with Action Include
- 916 917 Here we study the transitions of TA_3 when its current action is *Include*, given different actions of TA_4 and input samples. For ease of expressions, the self-loops of the transitions are not depicted

in the transition diagram. Clearly, this situation has 8 instances, depending on the variations of the training samples and the status of TA_4 , where the first four correspond to the instances with $TA_4 = E$ while the remaining four represent the instances with $TA_4 = I$.

Now we study the first instance, with $x_1 = 1$, $x_2 = 1$, y = 1, and $TA_4 = E$. Clearly, this training sample will trigger Type I feedback because y = 1. Together with the current status of the other TAs, the clause is determined to be $C = \neg x_1 \land x_2 = 0$ and the literal is $x_2 = 1$. From Table 3, we know that the penalty probability is $\frac{1}{s}$ and the inaction probability is $\frac{s-1}{s}$. To indicate the transitions, we have plotted the diagram, with the transitions for penalty (P) below. Note that the overall transition probability is $u_1 \frac{1}{s}$, where u_1 is defined in Eq. (3). Here, we have assumed $u_1 > 0$.

007	probability is $u_1 \frac{1}{s}$, where u_1 is defined in E	q. (3). He	re, w	e hav	e assume	$d u_1 > 0.$
927				Ι		E	
928	Condition: $x_1 = 1, x_2 = 1, y = 1,$				$u_1 \frac{1}{s}$		
929	$TA_4 = E.$	P	O	~~O	0	0	
930	Thus, Type I, $x_2 = 1$,				0		
931	$C = \neg x_1 \land x_2 = 0.$	R	0	0	0	0	
932					:		
933	We here continue with analyzing another						
934	training samples: $x_1 = 1, x_2 = 0, y = 0$,						
935	Type II feedback because $y = 0$. The clause						
936	study TA ₃ , the corresponding literal is $x_2 =$						
937	Table 4 and find the probability of "Inaction						ransition diagram does not
	have any arrow, indicating that there is "No	trans	sition	" for	TA_3 .		
938				Ι		E	
939	Condition: $x_1 = 1, x_2 = 0, y = 0,$:		
940	$TA_4 = E.$	P	0	0	0	0	No transition
941	Thus, Type II, $x_2 = 0$,						i to transition
942	$C = \neg x_1 \land x_2 = 0.$	R	0	0	0	0	
943					:		
944	The same analytical principle applies for al					s, and we	e therefore will not explain
945	them in detail. Instead, we just list the trans	ition	diagi	ams.	•		
946				Ι	1	3	
947	Condition: $x_1 = 0, x_2 = 1, y = 0,$:		
948	$TA_4 = E.$	P	0	0	0	0	No transition
949	Thus, Type II, $x_2 = 1$, $C = \neg x_1 \land x_2 = 1$.	_	~	~	_	0	
950	$C = \neg x_1 \land x_2 = 1.$	R	0	0	0	0	
951				I	. 1	7	
952	Condition: $x_1 = 0, x_2 = 0, y = 0,$			1	. 1	-	
	$TA_4 = E.$	P	0	0	0	0	No transition
953	Thus, Type II, $x_2 = 0$,		-	-		-	No transition
954	$C = \neg x_1 \land x_2 = 0.$	R	0	0	0	0	
955							
956	Condition: $x_1 = 1, x_2 = 1, y = 1,$			Ι	1	E	
957	TA ₄ = I.		\sim	\rightarrow	u_{1s}	0	
958	Thus, Type I, $x_2 = 1$,	P	0	0	0	0	
959	$C = \neg x_1 \land x_2 \land \neg x_2 = 0.$	R	\cap	\cap	0	0	
960	$C = \omega_1 / \omega_2 / \omega_2 = 0.$	11	U	0	1	0	
961				I	1	3	
962	Condition: $x_1 = 1, x_2 = 0, y = 0,$:		
963	$TA_4 = I.$	P	0	0	0	0	No transition
964	Thus, Type II, $x_2 = 0$,						
965	$C = \neg x_1 \land x_2 \land \neg x_2 = 0.$	R	0	\bigcirc	i O	0	
966					:	_	
967	Condition: $x_1 = 0, x_2 = 1, y = 0,$			Ι		E	
	$TA_4 = I.$	P	\cap	0	\cap	0	
968	Thus, Type II, $x_2 = 1$,	1	0	\cup		\cup	No transition
969	$C = \neg x_1 \wedge x_2 \wedge \neg x_2 = 0.$	R	0	0	0	0	
970			0	~		-	

972 Ι E Condition: $x_1 = 0, x_2 = 0, y = 0,$ 973 0 0 0 0 $TA_4 = I.$ P974 No transition Thus, Type II, $x_2 = 0$, 975 $R \bigcirc \bigcirc \bigcirc$ $C = \neg x_1 \land x_2 \land \neg x_2 = 0.$ Ο 976 977 B.1.1.2 Study TA₃ with Action Exclude 978 Here we study the transitions of TA_3 when its current action is *Exclude*, given different actions of 979 TA_4 and input samples. This situation has 8 instances, depending on the variations of the training 980 samples and the status of TA_4 . In this subsection and the following subsections, we will not plot the 981 transition diagrams for "No transition". 982 Ι E983 Condition: $x_1 = 1, x_2 = 1, y = 1$, 984 $TA_4 = E.$ P0 0 0 0 985 Thus, Type I, $x_2 = 1$, $C = \neg x_1 = 0.$ 0 0 986 987 IE988 Condition: $x_1 = 0, x_2 = 0, y = 0,$ $u_2 \times 1$ 989 $TA_4 = E.$ P \bigcirc \bigcirc \mathcal{O} 990 Thus, Type II, $x_2 = 0$, 991 $C = \neg x_1 = 1.$ R0 0 Ο 0 992 I E993 Condition: $x_1 = 1, x_2 = 1, y = 1$, 994 $TA_4 = I.$ P \bigcirc O i O0 995 Thus, Type I, $x_2 = 1$, 996 $C = \neg x_1 \land \neg x_2 = 0.$ R0 0 O 997 998 ECondition: $x_1 = 0, x_2 = 0, y = 0,$ 999 $u_2 \times 1$ $TA_4 = I.$ Ρ 1000 Thus, Type II, $x_2 = 0$, 1001 $C = \neg x_1 \land \neg x_2 = 1.$ RΟ 0 0 0 1002 1003 B.1.1.3 Study TA₄ with Action Include 1004 Here we list the transitions for TA₄ when its current action is Include. 1005 1006 Condition: $x_1 = 1, x_2 = 1, y = 1$, 1007 $TA_3 = E.$ 0 1008 Thus, Type I, $\neg x_2 = 0$, 1009 $R \bigcirc \bigcirc \bigcirc$ $C = \neg x_1 \land \neg x_2 = 0.$ 0 1010 1011 $u_1 \frac{1}{s}$ Condition: $x_1 = 1, x_2 = 1, y = 1$, 1012 $TA_3 = I.$ 1013 Thus, Type I, $\neg x_2 = 0$, 1014 $C = \neg x_1 \land x_2 \land \neg x_2 = 0.$ \bigcirc 0 \bigcirc 1015 1016 *B.1.1.4* Study TA₄ with Action Exclude 1017 1018 Here we list the transitions for TA_4 when its current action is *Exclude*. 1019 ECondition: $x_1 = 1, x_2 = 1, y = 1$, 1020 $\begin{array}{cccc} P & \bigcirc & \bigcirc & \bigcirc & \bigcirc & \bigcirc \\ R & \bigcirc & \bigcirc & \bigcirc & \bigcirc & \bigcirc \\ u_1 \frac{1}{8} \end{array}$ $TA_3 = E.$ 1021 Thus, Type I, $\neg x_2 = 0$, 1022 $C = \neg x_1 = 0.$ 1023 1024 1025

1026 Ι ECondition: $x_1 = 0, x_2 = 1, y = 0,$ $u_2 \times 1$ 1027 $TA_3 = E.$ PΟ O \mathcal{O} 1028 Thus, Type II, $\neg x_2 = 0$, 1029 $C = \neg x_1 = 1.$ R0 Ο Ο \bigcirc 1030 1031 Ι E Condition: $x_1 = 1, x_2 = 1, y = 1$, 1032 $TA_3 = I.$ PΟ Ο 0 0 1033 Thus, Type I, $\neg x_2 = 0$, 1034 RΟ $C = \neg x_1 \land x_2 = 0.$ Ο \cap 'Or 1035 $u_{1\frac{1}{s}}$ 1036 ECondition: $x_1 = 0, x_2 = 1, y = 0,$ 1037 $u_2 \times 1$ $TA_3 = I.$ 1038 PThus, Type II, $\neg x_2 = 0$, 1039 $C = \neg x_1 \land x_2 = 1.$ 0 RΟ Ο Ο 1040 1041 1042 **B.1.2** CASE 2: INCLUDE x_1 1043 1044 For Case 2, we assume that the actions of the TAs for the first bit are frozen as $TA_1 = I$ and 1045 $TA_2 = E$, and thus the overall joint action for the first bit is " x_1 ". Similar to Case 1, we also have 4 situations. 1046 1047 B.1.2.1 Study TA₃ with Action Include 1048 E1049 Condition: $x_1 = 1, x_2 = 1, y = 1$, 1050 $TA_4 = E.$ PΟ 0 Ο 0 1051 Thus, Type I, $x_2 = 1$, 1052 $C = x_1 \wedge x_2 = 1.$ $R \simeq O$ 0 Ο 1053 1054 ECondition: $x_1 = 1, x_2 = 1, y = 1$, 1055 $u_{1\frac{1}{2}}$ $TA_4 = I.$ Ο 1056 Thus, Type I, $x_2 = 1$, 1057 $C = x_1 \wedge x_2 \wedge \neg x_2 = 0.$ R Ο Ο Ο 0 1058 1059 B.1.2.2 Study TA₃ with Action Exclude 1060 1061 Condition: $x_1 = 1, x_2 = 1, y = 1$, $u_1 \frac{s-1}{s}$ 1062 $TA_4 = E.$ P \bigcirc C 1063 Thus, Type I, $x_2 = 1$, $C = x_1 = 1.$ 1064 R 0 0 \bigcirc 0 1065 E Ι 1066 Condition: $x_1 = 1, x_2 = 0, y = 0$, $u_2 \times 1$ 1067 $TA_4 = E.$ P \bigcirc С 1068 Thus, Type II, $x_2 = 0$, 1069 $C = x_1 = 1.$ RΟ 0 0 \bigcirc 1070 Ι E1071 Condition: $x_1 = 1, x_2 = 1, y = 1$, 1072 $TA_4 = I.$ Ρ Ο Ο Ο Ο 1073 Thus, Type I, $x_2 = 1$, 1074 R $C = x_1 \land \neg x_2 = 0.$ Ο Ο C $u_1 \frac{1}{s}$ 1075 1076 ECondition: $x_1 = 1, x_2 = 0, y = 0$, 1077 $u_2 \times 1$ $TA_4 = I.$ P \bigcirc 1078 Thus, Type II, $x_2 = 0$, 1079 $C = x_1 \land \neg x_2 = 1.$ R0 Ο \bigcirc Ο

1080 B.1.2.3 Study TA₄ with Action Include 1081 1082 Ι ECondition: $x_1 = 1, x_2 = 1, y = 1$, $u_1 \frac{1}{s}$ 1083 $TA_3 = E.$ 0 1084 Thus, Type I, $\neg x_2 = 0$, 1085 $C = x_1 \wedge \neg x_2 = 0.$ R0 0 0 Ο 1086 E1087 Condition: $x_1 = 1, x_2 = 1, y = 1$, $u_1 \frac{1}{e}$ 1088 $TA_3 = I.$ Ο 1089 Thus, Type I, $\neg x_2 = 0$, 1090 $C = x_1 \wedge x_2 \wedge \neg x_2 = 0.$ \bigcirc \bigcirc RΟ Ο 1091 1092 *B.1.2.4* Study TA₄ with Action Exclude 1093 I E1094 Condition: $x_1 = 1, x_2 = 1, y = 1$, 1095 0 0 0 $TA_3 = E.$ 0 Thus, Type I, $\neg x_2 = 0$, 1096 $C = x_1 = 1.$ 0 1097 1098 Ι 1099 Condition: $x_1 = 1, x_2 = 1, y = 1$, 1100 $TA_3 = I.$ P0 0 0 0 1101 Thus, Type I, $\neg x_2 = 0$, 1102 0 0 0 R $C = x_1 \wedge x_2 = 1.$ 1103 1104 1105 1106 1107 **B.1.3** Case 3: Exclude Both $\neg x_1$ and x_1 1108 1109 For Case 3, we assume that the actions of TAs for the first bit are frozen as $TA_1 = E$ and $TA_2 = E$, 1110 with 4 situations. Note that in the training process, when all literals are excluded, C is assigned to 1. 1111 B.1.3.1 Study TA₃ with Action Include 1112 1113 ECondition: $x_1 = 1, x_2 = 1, y = 1$, 1114 0 0 $TA_4 = E.$ 1115 Thus, Type I, $x_2 = 1$, 1116 $R \stackrel{O}{u_1 \frac{s-1}{s}}$ $C = x_2 = 1.$ Ο 0 1117 1118 ECondition: $x_1 = 1, x_2 = 1, y = 1$, 1119 1120 $TA_4 = I.$ \bigcirc Thus, Type I, $x_2 = 1$, 1121 C = 0.0 0 Ο 0 1122 1123 B.1.3.2 Study TA₃ with Action Exclude 1124 1125 Condition: $x_1 = 1, x_2 = 1, y = 1$, 1126 $TA_4 = E.$ 1127 Thus, Type I, $x_2 = 1$, 1128 C = 1.Ο R0 Ο \bigcirc 1129 1130 Condition: $x_1 = 1, x_2 = 0, y = 0$, 1131 $u_2 \times 1$ $TA_4 = E.$ Ρ 1132 Thus, Type II, $x_2 = 0$, 1133 C = 1.R Ο 0 \cap Ο

1134			Ι		E	
1135	Condition: $x_1 = 0, x_2 = 0, y = 0,$ TA ₄ = E.	P	\sim		$u_1 \times$	< 1 ~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~~
1136	Thus, Type II, $x_2 = 0$,	Р	0	04	0.	0
1137	C = 1.	R	0	0	0	0
1138 1139				:	_	
1140	Condition: $x_1 = 1, x_2 = 1, y = 1,$		Ι		E	
1141	$TA_4 = I.$	P	0	0	0	0
1142	Thus, Type I, $x_2 = 1$,	P	\sim	~	\sim	*
1143	C = 0.	R	0	0	0	$u_1 \frac{1}{s}$
1144			Ι		E	
1145	Condition: $x_1 = 1, x_2 = 0, y = 0,$. :	$u_2 \times$	< 1
1146 1147	$TA_4 = I.$ Thus, Type II, $x_2 = 0$,	P	0	Or	Or	\mathcal{O}
1147	C = 1.	R	0	0	0	0
1149						
1150	Condition: $x_1 = 0, x_2 = 0, y = 0,$		Ι		E $u_2 \times$	
1151	$TA_4 = I.$	P	0	0		\mathcal{O}
1152	Thus, Type II, $x_2 = 0$,		0	_	0	0
1153	C = 1.	R	0	0	0	0
1154 1155	B.1.3.3 Study TA ₄ with Action Include					
1156			Ι		Ε	
1157	Condition: $x_1 = 1, x_2 = 1, y = 1,$		-		$u_1 \frac{1}{s}$	
1158	$TA_3 = E.$ Thus, Type I, $\neg x_2 = 0$,	P	0	*0	70	0
1159	$C = \neg x_2 = 0.$	R	0	0	\bigcirc	0
1160	2		0		Ŭ	Ŭ
1161 1162	Condition: $x_1 = 1, x_2 = 1, y = 1,$		Ι		$u_1 \frac{1}{s}$ E	
1163	$TA_3 = I.$	P	0	*0	\sim	0
1164	Thus, Type I, $\neg x_2 = 0$,				_	
1165	$C = \neg x_2 \land x_2 = 0.$	R	0	0	0	0
1166	<i>B.1.3.4</i> Study TA ₄ with Action Exclude					
1167			Ι		E	
1168 1169	Condition: $x_1 = 1, x_2 = 1, y = 1,$		-	:	1	
1170	$TA_3 = E.$	P	0	0	0	0
1171	Thus, Type I, $\neg x_2 = 0$, C = 1.	R	0	0	\circ	4OP
1172	0 – 1.		0		0	$u_1 \frac{1}{s}$
1173	Condition: $x_1 = 0, x_2 = 1, y = 0,$		Ι		E	
1174	TA ₃ = E.	Ρ	\bigcirc	~~: 	$u_2 \times$	\sim
1175	Thus, Type II, $\neg x_2 = 0$,	1	0		0	0
1176 1177	C = 1.	R	0	0	0	0
1178			Ι		E	
1179	Condition: $x_1 = 1, x_2 = 1, y = 1,$:		
1180	$TA_3 = I.$	P	0	0	0	0
1181	Thus, Type I, $\neg x_2 = 0$, C = 1.	R	0	0	0	YOP.
1182						$u_1 \frac{1}{s}$
1183 1184	Condition: $x_1 = 0, x_2 = 1, y = 0,$		Ι		E	
1185	$\begin{array}{l} \text{Condition:} \ x_1 = 0, \ x_2 = 1, \ y = 0, \\ \text{TA}_3 = \text{I.} \end{array}$	Р	0	0	$\sim u_2 \times$	\sim
1186	Thus, Type II, $\neg x_2 = 0$,		0	Ŭ	0	~
1187	C = 1.	R	0	0	0	0

1188 B.1.4 CASE 4: INCLUDE BOTH $\neg x_1$ A	AND x_1
1190	
	f TAs for the first bit are frozen as $TA_1 = I$ and $TA_2 = I$,
1192 and thus $C = 0$ always. Similarly, we al	so have 4 situations, detailed below.
1193 $B.1.4.1$ Study TA ₃ with Action Include	
1194	I E
1195 Condition: $x_1 = 1, x_2 = 1, y = 1,$	$u_1 \frac{1}{s}$
1196 $TA_4 = E.$ 1197 Thus, Type I, $x_2 = 1$,	P O O O
1198 $C = 0.$	$R \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc$
1199	
1200 Condition: $x_1 = 1, x_2 = 1, y = 1$,	
1201 $TA_4 = I.$	P
1202 Thus, Type I, $x_2 = 1$,	
1203 $C = 0.$	$R \bigcirc \bigcirc \bigcirc \bigcirc$
1204 1205 B.1.4.2 Study TA ₃ with Action Exclude	:
1200	
1206 1207 Condition: $x_1 = 1, x_2 = 1, y = 1,$	I E
1207 Condition: $x_1 = 1, x_2 = 1, y = 1, TA_4 = E.$	$P \bigcirc \bigcirc \bigcirc \bigcirc$
1209 Thus, Type I, $x_2 = 1$,	
1210 $C = 0.$	$R \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc \square_{n_1 \underline{1}}$
1211	. <i>w</i> 1 _S
1212 Condition: $x_1 = 1, x_2 = 1, y = 1$,	$I \qquad E$.
1213 $TA_4 = I.$	$P \bigcirc \bigcirc \bigcirc \bigcirc$
1214 Thus, Type I, $x_2 = 1$,	
1215 $C = 0.$	$R \bigcirc \bigcirc \bigcirc \qquad \bigcirc \qquad $
1216 1217 B.1.4.3 Study TA ₄ with Action Include	- 3
1217 D.1.4.5 Shudy 1114 with Action Include	
1219 Condition: $x_1 = 1, x_2 = 1, y = 1$,	$I \qquad E$
1220 $TA_3 = E.$	P
1221 Thus, Type I, $\neg x_2 = 0$,	
1222 $C = 0.$	$R \bigcirc \bigcirc \bigcirc \bigcirc$
1223	I E
1224 Condition: $x_1 = 1, x_2 = 1, y = 1$,	$u_1 \frac{1}{s}$
1225 $TA_3 = I$.	
1226 Thus, Type I, $\neg x_2 = 0$, 1227 $C = 0$.	$R \bigcirc \bigcirc \bigcirc \bigcirc$
$\begin{array}{llllllllllllllllllllllllllllllllllll$	
1229	I E
1231 Condition: $x_1 = 1, x_2 = 1, y = 1$,	:
1232 $TA_3 = E.$	$P \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc$
1233 Thus, Type I, $\neg x_2 = 0$,	
1234 $C = 0.$	$R \bigcirc \bigcirc \bigcirc \qquad \bigcirc \qquad \bigcirc \qquad \bigcirc \qquad \bigcirc \qquad \bigcirc \qquad \qquad \bigcirc \qquad \qquad \bigcirc \qquad \qquad$
1235	I E
1236 Condition: $x_1 = 1, x_2 = 1, y = 1,$:
1237 $TA_3 = I$.	$P \bigcirc \bigcirc \bigcirc \bigcirc \bigcirc$
1238 Thus, Type I, $\neg x_2 = 0$, 1239 $C = 0$.	$R \bigcirc \bigcirc \bigcirc \qquad \bigcirc \qquad \bigcirc \qquad \bigcirc \qquad \qquad \bigcirc \qquad \qquad \qquad \qquad \qquad \qquad \qquad $
1239 $C = 0.$ 1240	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$
	y the transitions of $T\Lambda_{a}$ and $T\Lambda_{b}$ for all the cases (all nos

So far, we have gone through, exhaustively, the transitions of TA_3 and TA_4 for all the cases (all possible training samples and system states). Hereafter, we can summarize the direction of transitions

1242 and study the convergence properties of the system for the given training samples, to be detailed in 1243 the next subsection. 1244 1245 1246 1247 SUMMARIZE OF THE DIRECTIONS OF TRANSITIONS IN DIFFERENT CASES B.2 1248 1249 1250 Based on the analysis above, we summarize here what happens to TA_3 and TA_4 , given different 1251 status (Cases) of TA_1 and TA_2 . More specifically, we will summarize here the directions of the transitions for the TAs. For example, "TA₃ \Rightarrow E" means that TA₃ will move towards the action 1252 "Exclude", while "TA₄ \Rightarrow E or I" means TA₄ transits towards either "Exclude" or "Include". 1253 1254 **Scenario 1:** Study $TA_3 = I$ and $TA_4 = I$. 1255 Case 1, we have: Case 3, we have: 1256 $TA_3 \Rightarrow E.$ $TA_3 \Rightarrow E.$ 1257 $TA_4 \Rightarrow E.$ $TA_4 \Rightarrow E.$ 1258 Case 2, we have: Case 4, we have: 1259 $TA_3 \Rightarrow E.$ $TA_3 \Rightarrow E.$ 1260 $TA_4 \Rightarrow E.$ $TA_4 \Rightarrow E.$ 1261 1262 From the facts presented above, we can confirm that regardless the state of TA_1 and TA_2 , if $TA_3 =$ 1263 I and $TA_4 = I$, they (TA₃ and TA₄) will eventually move out of their states. 1264 1265 **Scenario 2:** Study $TA_3 = I$ and $TA_4 = E$. 1266 Case 3, we have: Case 1, we have: 1267 $TA_3 \Rightarrow E.$ $TA_3 \Rightarrow I.$ 1268 $TA_4 \Rightarrow E \text{ or } I.$ $TA_4 \Rightarrow E \text{ or } I.$ 1269 Case 2, we have: Case 4, we have: 1270 $TA_3 \Rightarrow I.$ $TA_3 \Rightarrow E.$ 1271 $TA_4 \Rightarrow E.$ $TA_4 \Rightarrow E.$ 1272 1273 For Scenario 2 Case 2, we can observe that if $TA_3 = I$, $TA_4 = E$, $TA_1 = I$, and $TA_2 = E$, TA_3 will 1274 move deeper to "include" and TA_4 will go deeper to "exclude". It is not difficult to derive also that 1275 TA_1 will move deeper to "include" and TA_2 will transfer deeper to "exclude" in this circumstance. This tells us that the TAs in states $TA_3 = I$, $TA_4 = E$, $TA_1 = I$, and $TA_2 = E$, reinforce each other 1276 1277 to move deeper to their corresponding directions and they therefore construct an absorbing state of the system. If it is the only absorbing state, we can conclude that the TM converge to the intended 1278 "AND" operation. 1279 1280 In Scenario 2, we can observe for Cases 1, 3, and 4, the actions for TA_3 and TA_4 are not ab-1281 sorbing because the TAs will not be reinforced to move monotonically deeper to the states of the 1282 corresponding actions for difference cases. 1283 For Scenario 2, Case 3, TA_4 has two possible directions to transit, I or E, depending on the input 1284 of the training sample. For action exclude, it will be reinforced when training sample $x_1 = 1$ and 1285 $x_2 = 1$ is given, based on Type I feedback. However, TA₄ will transit towards "include" side when 1286 training sample $x_1 = 0$ and $x_2 = 1$ is given, due to Type II feedback. Therefore, the direction of the 1287 transition for TA_4 is I or E, depending on the training samples. In the following paragraphs, when 1288 "or" appears in the transition direction, the same concept applies. 1289 1290 **Scenario 3:** Study $TA_3 = E$ and $TA_4 = I$. 1291 Case 1, we have: Case 3, we have: 1292 $TA_3 \Rightarrow E \text{ or } I.$ $TA_3 \Rightarrow E \text{ or } I.$ 1293 $TA_4 \Rightarrow E.$ $TA_4 \Rightarrow E.$ 1294 Case 2, we have: Case 4, we have: 1295

 $\begin{array}{ll} \mathrm{TA}_3 \Rightarrow \mathrm{E} \text{ or } \mathrm{I}. & \mathrm{TA}_3 \Rightarrow \mathrm{E}. \\ \mathrm{TA}_4 \Rightarrow \mathrm{E}. & \mathrm{TA}_4 \Rightarrow \mathrm{E}. \end{array}$

¹²⁹⁶ In Scenario 3, we can see that the actions for $TA_3 = E$ and $TA_4 = I$ are not absorbing because the TAs will not be reinforced to move deeper to the states of the corresponding actions.

Scenario 4: Study $TA_3 = E$ and $TA_4 = E$.

1300	a i i	a a i
1301	Case 1, we have:	Case 3, we have:
	$TA_3 \Rightarrow I \text{ or } E.$	$TA_3 \Rightarrow I.$
1302	$TA_4 \Rightarrow I \text{ or } E.$	$TA_4 \Rightarrow I \text{ or } E.$
1303	Case 2, we have:	Case 4, we have:
1304	$TA_3 \Rightarrow I.$	$TA_3 \Rightarrow E.$
1305	$TA_4 \Rightarrow E.$	$TA_3 \Rightarrow E.$ $TA_4 \Rightarrow E.$
1306	$IA_4 \rightarrow L.$	$1\Lambda_4 \rightarrow L.$

1307 In Scenario 4, we see that, the actions for $TA_3 = E$ and $TA_4 = E$ seem to be an absorbing state, 1308 because the states of TAs will move deeper in Case 4. After a revisit of the condition for Case 4, i.e., 1309 include both $\neg x_1$ and x_1 , we understand that this condition is not absorbing. In fact, when TA_1 and 1310 TA_2 both have "Include" as their actions, they monotonically move towards "Exclude". Therefore, 1311 from the overall system's perspective, the system state $TA_1 = I$, $TA_2 = I$, $TA_3 = E$, and $TA_4 = E$ 1312 is not absorbing. For the other cases in this scenario, there is no absorbing state.

Based on the above analysis, we understand that there is only one absorbing condition in the system, namely, $TA_1 = I$, $TA_2 = E$, $TA_3 = I$, and $TA_4 = E$, for the given training samples with AND logic. The same conclusion applies when we freeze the transition of the two TAs for the second bit of the input and study behavior of the first bit of input. Therefore, we can conclude that the TM with a clause can learn to be the intended AND operator, almost surely, in infinite time horizon. We thus complete the proof of Theorem 1.

1350 **PROOF OF LEMMA 1** С 1351

1352 The probability of the training samples for the noise-free OR operator can be presented by the 1353 following equations. 135/

1554		(10)
1355	$P(y=1 x_1=1, x_2=1) = 1,$	(19)
1356	$P(y=1 x_1=0, x_2=1) = 1,$	
1357	$P(y=1 x_1=1, x_2=0) = 1,$	
1358	$P(y=0 x_1=0, x_2=0) = 1.$	
1359		

1360 Clearly, there are three sub-patterns of x_1 and x_2 that will give y = 1, i.e., $[x_1 = 1, x_2 = 1]$, 1361 $[x_1 = 1, x_2 = 0]$, and $[x_1 = 0, x_2 = 1]$. More specifically, Eq. (19) can be split into three cases, corresponding to the three sub-patterns: 1362

$$\begin{array}{l} P\left(y=1|x_{1}=1,x_{2}=1\right)=1, \\ P\left(y=0|x_{1}=0,x_{2}=0\right)=1, \\ P\left(y=1|x_{1}=0,x_{2}=1\right)=1, \\ P\left(y=0|x_{1}=0,x_{2}=0\right)=1, \\ \end{array} \tag{21}$$
 and
$$\begin{array}{l} P\left(y=1|x_{1}=1,x_{2}=0\right)=1, \\ P\left(y=0|x_{1}=0,x_{2}=0\right)=1. \\ P\left(y=0|x_{1}=0,x_{2}=0\right)=1. \end{array}$$
 In what follows, we will show the convergence of the three sub-patterns, i.e., Lemma 1. \\ The convergence analyses of the above three sub-patterns can be derived by reusing the analyses

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1374 1375 of the sub-patterns of the XOR operator plus the AND operator. For the sub-pattern described by 1376 Eq. (20), we can confirm that the TAs will indeed converge to $TA_1 = I$, $TA_2 = E$, $TA_3 = I$, 1377 and $TA_4 = E$, by studying the transition diagrams in Subsection B when input samples of $[x_1 = 0, x_2 = 0]$ 1378 $x_2 = 1$ and $[x_1 = 1, x_2 = 0]$ are removed. In this case, the directions of the transitions for different scenarios are summarized below. 1379

1380 **Scenario 1:** Study $TA_3 = I$ and $TA_4 = I$.

```
1381
                                               Case 3, we have:
            Case 1, we have:
1382
            TA_3 \Rightarrow E.
                                               TA_3 \Rightarrow E.
1383
                                              TA_4 \Rightarrow E.
            TA_4 \Rightarrow E.
1384
            Case 2, we have:
                                               Case 4, we have:
1385
                                               TA_3 \Rightarrow E.
            TA_3 \Rightarrow E.
1386
                                              TA_4 \Rightarrow E.
           TA_4 \Rightarrow E.
1387
1388
           Scenario 2: Study TA_3 = I and TA_4 = E.
1389
1390
            Case 1, we have:
                                               Case 3, we have:
           TA_3 \Rightarrow E.
1391
                                               TA_3 \Rightarrow I.
           TA_4 \Rightarrow E.
1392
                                               TA_4 \Rightarrow E.
            Case 2, we have:
1393
                                               Case 4, we have:
           \mathrm{TA}_3 \Rightarrow \mathrm{I}.
1394
                                               TA_3 \Rightarrow E.
           TA_4 \Rightarrow E.
1395
                                              TA_4 \Rightarrow E.
1396
1397
           Scenario 3: Study TA_3 = E and TA_4 = I.
1398
                                               Case 3, we have:
1399
           Case 1, we have:
                                               TA_3 \Rightarrow E \text{ or } I.
           TA_3 \Rightarrow E \text{ or } I.
1400
                                              TA_4 \Rightarrow E.
           TA_4 \Rightarrow E.
1401
                                               Case 4, we have:
            Case 2, we have:
1402
                                              TA_3 \Rightarrow E.
           TA_3 \Rightarrow E.
1403
                                              TA_4 \Rightarrow E.
           TA_4 \Rightarrow E.
```

1404 **Scenario 4:** Study $TA_3 = E$ and $TA_4 = E$. 1405 Case 1, we have: Case 3, we have: 1406 $TA_3 \Rightarrow I \text{ or } E.$ $TA_3 \Rightarrow I.$ 1407 $TA_4 \Rightarrow E.$ $TA_4 \Rightarrow E.$ 1408 Case 4, we have: Case 2, we have: 1409 $TA_3 \Rightarrow E.$ $TA_3 \Rightarrow I.$ 1410 $TA_4 \Rightarrow E.$ $TA_4 \Rightarrow E.$ 1411 1412 Comparing the analysis with the one in Subsection B.2, there is apparently another possible absorbing case, which can be observed in Scenario 2, Case 3, where $TA_3 = I$ and $TA_4 = E$, given 1413 $TA_1 = E$ and $TA_2 = E$. However, given $TA_3 = I$ and $TA_4 = E$, the TAs for the first bit, i.e., 1414 $TA_1 = E$ and $TA_2 = E$, will not move only towards Exclude. Therefore, they do not reinforce 1415 each other to move to deeper states for their current actions. For this reason, the system in $TA_3 = I$, 1416 $TA_4 = E$, $TA_1 = E$, and $TA_2 = E$, is not in an absorbing state. In addition, given $TA_3 = I$ and 1417 $TA_4 = E$, TA_1 and TA_2 with actions E and E will transit towards I and E, encouraging the overall 1418 system to move towards I, E, I, and E. Consequently, the system state with $TA_1 = I$, $TA_2 = E$, 1419 $TA_3 = I$, and $TA_4 = E$ is still the only absorbing case for the given training samples following 1420 Eq. (20). 1421 For Eq. (21), similar to the proof of in Lemma 1 in (Jiao et al., 2022), we can derive that the TAs will 1422 converge in $TA_1 = E$, $TA_2 = I$, $TA_3 = I$, and $TA_4 = E$. The transition diagrams for the samples 1423 of Eq. (21) are in fact a subset of the ones presented in Subsection 3.2.1 and Appendix 2 of (Jiao 1424 et al., 2022), when the input samples of $[x_1 = 1 \text{ and } x_2 = 1]$ are removed. We summarize below 1425 only the directions of transitions. 1426 The directions of the transitions of the TAs for the second input bit, i.e., $x_2/\neg x_2$, when the TAs 1427 for the first input bit are frozen, are summarized as follows (based on the subset of the transition 1428 diagrams in Subsection 3.2.1 of (Jiao et al., 2022)). 1429 1430 **Scenario 1:** Study $TA_3 = I$ and $TA_4 = I$. 1431 Case 3: we have Case 1: we have 1432 $TA_3 \rightarrow E$ $TA_3 \rightarrow E$ 1433 $TA_4 \rightarrow E$ $TA_4 \rightarrow E$ 1434 Case 2: we have Case 4: we have 1435 $TA_3 \rightarrow E$ $TA_3 \rightarrow E$ $TA_4 \rightarrow E$ 1436 $TA_4 \rightarrow E$ 1437 1438 **Scenario 2:** Study $TA_3 = I$ and $TA_4 = E$. 1439 Case 1: we have Case 3: we have 1440 $TA_3 \rightarrow I$ $TA_3 \rightarrow I$ 1441 $TA_4 \rightarrow E$ $TA_4 \rightarrow E$ 1442 Case 2: we have Case 4: we have 1443 $\mathrm{TA}_3 \to E$ $TA_3 \rightarrow E$ 1444 $\mathrm{TA}_4 \to E$ $TA_4 \rightarrow E$ 1445 1446 **Scenario 3:** Study $TA_3 = E$ and $TA_4 = I$. 1447 1448 Case 1: we have Case 3: we have 1449 $TA_3 \rightarrow I$, or E $TA_3 \rightarrow I$, or E 1450 $TA_4 \rightarrow E$ $TA_4 \rightarrow E$ 1451 Case 2: we have Case 4: we have $TA_3 \rightarrow E$ $TA_3 \rightarrow E$ 1452 $TA_4 \rightarrow E$ $TA_4 \rightarrow E$ 1453 1454 1455 **Scenario 4:** Study $TA_3 = E$ and $TA_4 = E$. 1456

1458 Case 3: we have Case 1: we have 1459 $TA_3 \rightarrow I$ $TA_3 \rightarrow I$ 1460 $TA_4 \rightarrow E$ $TA_4 \rightarrow E$ Case 2: we have Case 4: we have 1461 $TA_3 \rightarrow E$ $TA_3 \rightarrow E$ 1462 $TA_4 \rightarrow E$ $TA_4 \rightarrow E$ 1463 1464 1465 The directions of the transitions of the TAs for the first input bit, i.e., $x_1/\neg x_1$, when the TAs for 1466 the second input bit are frozen, are summarized as follows (based on the subset of the transition 1467 diagrams in Appendix 2 of (Jiao et al., 2022)). 1468 **Scenario 1:** Study $TA_1 = I$ and $TA_2 = I$. 1469 Case 1: we have Case 3: we have 1470 $TA_1 \rightarrow E$ $TA_1 \rightarrow E$ 1471 $TA_2 \rightarrow E$ $TA_2 \rightarrow E$ 1472 Case 4: we have Case 2: we have 1473 $TA_1 \rightarrow E$ $TA_1 \rightarrow E$ 1474 $\mathrm{TA}_2 \to E$ $TA_2 \rightarrow E$ 1475 1476 **Scenario 2:** Study $TA_1 = I$ and $TA_2 = E$. 1477 1478 Case 3: we have Case 1: we have 1479 $\mathrm{TA}_1 \to E$ $TA_1 \rightarrow E$ $\mathrm{TA}_2 \to E$ $\mathrm{TA}_2 \to \mathrm{E}$ 1480 Case 4: we have Case 2: we have 1481 $TA_1 \rightarrow E$ $TA_1 \rightarrow E$ 1482 $TA_2 \rightarrow E$ $TA_2 \rightarrow E$ 1483 1484 1485 **Scenario 3:** Study $TA_1 = E$ and $TA_2 = I$. 1486 Case 1: we have Case 3: we have 1487 $TA_1 \rightarrow I$, or E $\mathrm{TA}_1 \to \mathrm{I}$ 1488 $TA_2 \rightarrow E$ $TA_2 \rightarrow I$ 1489 Case 4: we have Case 2: we have 1490 $TA_1 \rightarrow E$ $TA_1 \rightarrow E$ 1491 $TA_2 \rightarrow I$ $TA_2 \rightarrow E$ 1492 1493 **Scenario 4:** Study $TA_1 = E$ and $TA_2 = E$. 1494 Case 3: we have Case 1: we have 1495 $TA_1 \rightarrow I$, or E $\mathrm{TA}_1 \to E$ 1496 $\mathrm{TA}_2 \to E$ $TA_2 \rightarrow E$ 1497 Case 2: we have Case 4: we have 1498 $TA_1 \rightarrow E$ $TA_1 \rightarrow E$ 1499 $TA_2 \rightarrow E$ $TA_2 \rightarrow I$ 1500 1501 By analyzing the transitions of TAs for the two input bits with samples following Eq. (21), we can 1502 conclude that $TA_1 = E$, $TA_2 = I$, $TA_3 = I$, and $TA_4 = E$ is an absorbing state, as the actions of 1503 TA_1-TA_4 reinforce each other to transit to deeper states for the current actions upon various input 1504 samples. There are a few other cases in different scenarios that seem to be absorbing, but in fact 1505 not. For example, the status $TA_3 = I$ and $TA_4 = E$ seems also absorbing in Scenario 2, Case 3, 1506 i.e., when $TA_1 = E$ and $TA_2 = E$ hold. However, to make $TA_1 = E$ and $TA_2 = E$ absorbing, 1507 the condition is $TA_3 = I$ and $TA_4 = I$, or $TA_3 = E$ and $TA_4 = E$. Clearly, the status $TA_3 = I$ 1508 and $TA_4 = I$ is not absorbing. For $TA_3 = E$ and $TA_4 = E$ to be absorbing, it is required to have 1509 $TA_1 = I$ and $TA_2 = I$ to be absorbing, or $TA_1 = I$ and $TA_2 = E$ to be absorbing, which are not

true. Therefore, all those absorbing-like states are not absorbing. In fact, when $TA_3 = I$, $TA_4 = E$, TA₁ = E, and TA₂ = E hold, the condition TA₃ = I, TA₄ = E will reinforce TA₁ and TA₂ to move towards E, I, which is the absorbing state of the system. Based on the above analysis on the transition directions, we can thus confirm the convergence of TM when training samples from Eq. (21) are given.

1515 1516 1517	Following the same principle, we can also confirm that the TAs will converge to $TA_1 = I$, $TA_2 = E$, $TA_3 = E$, and $TA_4 = I$ when training samples from Eq. (22) are given, according to the proof of
1518	Lemma 2 in (Jiao et al., 2022).
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¹⁵⁶⁶ D APPENDIX: ANALYSIS OF THE TM WITH WRONG TRAINING LABELS

In this appendix, we analyze the transition properties of the TM when training samples contain wrong labels.

1571 There are two types of wrong labels:

- Inputs labeled as 0, which should be 1.
- Inputs labeled as 1, which should be 0.

We begin by examining the first type of wrong label, followed by the second type, and then address the general case.

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D.1 THE AND OPERATOR WITH THE FIRST TYPE OF WRONG LABELS

To formally define training samples with the first type of wrong label, we use the following formulas:

1581 $P(y = 1 | x_1 = 1, x_2 = 1) = a, a \in (0, 1)$ (23)1582 $P(y = 0 | x_1 = 1, x_2 = 1) = 1 - a,$ 1583 $P(y = 0 | x_1 = 0, x_2 = 1) = 1,$ 1584 $P(y = 0 | x_1 = 0, x_2 = 1) = 1,$ 1585 $P(y = 0 | x_1 = 1, x_2 = 0) = 1,$ 1586 $P(y = 0 | x_1 = 0, x_2 = 0) = 1.$

In this case, the label for training samples representing the intended logic $[x_1 = 1, x_2 = 1]$ is y = 1 with probability a and y = 0 with probability 1 - a. In other words, in addition to the training samples detailed in Subsection B, a new training sample will appear to the system, namely $([x_1 = 1, x_2 = 1], y = 0)$.

Lemma 6. *The TM exhibits recurrence for the training samples defined in Eq. (23).* 1592

Proof: To prove this lemma, we analyze the TM's transitions as follows. First, we examine the transitions assuming $u_1 > 0$ and $u_2 > 0$, similar to the analysis in Subsection B, as detailed in Subsection D.1.1. Next, we study the impact of T to determine whether it leads to convergence (absorption), as discussed in Subsection D.1.2.

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D.1.1 TRANSITION OF TM WITH AND OPERATOR GIVEN u_1 > 0 and u_2 > 0
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Following the approach in Subsection B, we examine the transitions of TA₃ and TA₄ when the additional training sample ($[x_1 = 1, x_2 = 1], y = 0$) is introduced, considering Cases 1 to 4 as defined in Subsection B. Since y = 0 for this sample, only Type II feedback can be triggered to cause transitions. As TA₃ is responsible for the literal x_2 , which is always 1 for this sample, Type II feedback does not trigger any transitions for TA₃. Therefore, we focus on studying the potential transitions of TA₄ in the four cases defined in Subsection B.1.

1605 In Case 1, where $TA_1 = E$ and $TA_2 = I$, the clause value will always be 0 for the training sample 1606 because $\neg x_1$ is included in the clause, regardless of the action TA_4 takes. According to the Type II 1607 feedback transition table, no transition occurs when C = 0, so no transitions are triggered for TA_4 . 1608 Similarly, in Case 4, where $TA_1 = I$ and $TA_2 = I$, the clause value will always be 0 due to the 1609 presence of $x_1 \land \neg x_1$ in the clause. As a result, there are no transitions for TA_4 .

1610 In Case 2, where $TA_1 = I$ and $TA_2 = E$, the literal x_1 will always appear in the clause. When 1611 $TA_4 = I$, the clause includes the literal $\neg x_2$, which results in a clause value of 0. Therefore, no 1612 transition is triggered. However, when $TA_4 = E$, the literal x_1 will always appear in the clause, 1613 and the value of x_2 is 1, making the clause value 1 regardless of TA_3 's action (whether it includes 1614 or excludes x_2). According to the Type II feedback table, with the literal value of $\neg x_2$ being 0 and 1615 the clause value being 1, the transition for $TA_4 = E$ is:

1616IE1617Condition:
$$x_1 = 1, x_2 = 1, y = 0.$$
IE1618Thus, Type II, $\neg x_2 = 0,$ PO1619C = 1.ROO

1620 In Case 3, where $TA_1 = E$ and $TA_2 = E$, the clause value is fully determined by the TA_3 and TA_4 . 1621 In this case, whenever TA_4 's action is include, the clause value is 0 for this given sample because it 1622 includes a literal which is $\neg x_2$. For this reason, there is no transition for TA_4 . When TA_4 's action 1623 is exclude, the clause value is always 1, regardless the action of TA_3 . When it is include, the clause 1624 value is 1 has a literal value 1 because its corresponding literal is x_2 . When it is exclude, all literals 1625 are excluded and then the clause value becomes 1 by definition. Checking the transitions of TA_4 , 1626 we can conclude the following graph:

1627 In Case 3, where $TA_1 = E$ and $TA_2 = E$, the clause value is fully determined by TA_3 and TA_4 . 1628 When TA_4 's action is to include, the clause value is 0 for this sample because it includes the literal 1629 $\neg x_2$, resulting in no transition for TA_4 . However, when TA_4 's action is to exclude, the clause value 1630 is always 1, regardless of TA_3 's action. Specifically, when TA_4 includes x_2 , the clause value is 1, 1631 as the literal value of x_2 is 1. When it is exclude, all literals are excluded and then the clause value 1632 becomes 1 by definition. By examining the transitions of TA_4 , we can summarize the following 1633 graph:

1634				I	E
1635	Condition: $x_1 = 1, x_2 = 1, y = 0.$			\vdots u_2	
1636	Thus, Type II, $\neg x_2 = 0$,	P	0	OFOF	0
1637	C = 1.	R	\circ	0 0	\circ
1638			$\mathbf{\circ}$		\cup

We summarize the directions of the transitions when the new wrongly labeled sample is added, withthe newly added actions highlighted in red.

Scenario 1: Study $TA_3 = I$ and $TA_4 = I$.

1642 Case 1, we have: Case 3, we have: 1643 $TA_3 \Rightarrow E.$ $TA_3 \Rightarrow E.$ 1644 $TA_4 \Rightarrow E.$ $TA_4 \Rightarrow E.$ 1645 Case 2, we have: Case 4, we have: 1646 $TA_3 \Rightarrow E.$ $TA_3 \Rightarrow E.$ 1647 $TA_4 \Rightarrow E.$ $TA_4 \Rightarrow E.$ 1648 1649 1650 **Scenario 2:** Study $TA_3 = I$ and $TA_4 = E$. 1651 Case 1, we have: Case 3, we have: 1652 $TA_3 \Rightarrow E.$ $TA_3 \Rightarrow I.$ 1653 $TA_4 \Rightarrow E \text{ or } I.$ $TA_4 \Rightarrow E \text{ or } I.$ 1654 Case 2, we have: Case 4, we have: 1655 $TA_3 \Rightarrow I.$ $TA_3 \Rightarrow E.$ 1656 $TA_4 \Rightarrow E \text{ or } I.$ $TA_4 \Rightarrow E.$ 1657 1658 **Scenario 3:** Study $TA_3 = E$ and $TA_4 = I$. 1659 Case 1, we have: Case 3, we have: 1661 $TA_3 \Rightarrow E \text{ or } I.$ $TA_3 \Rightarrow E \text{ or } I.$ 1662 $TA_4 \Rightarrow E.$ $TA_4 \Rightarrow E.$ Case 2, we have: Case 4, we have: 1663 $TA_3 \Rightarrow E \text{ or } I.$ $TA_3 \Rightarrow E.$ 1664 $TA_4 \Rightarrow E.$ $TA_4 \Rightarrow E.$ 1665 1666 **Scenario 4:** Study $TA_3 = E$ and $TA_4 = E$. 1667 1668 Case 1, we have: Case 3, we have: $TA_3 \Rightarrow I.$ 1669 $TA_3 \Rightarrow I \text{ or } E.$ $TA_4 \Rightarrow I \text{ or } E.$ $TA_4 \Rightarrow I \text{ or } E.$ 1670 Case 2, we have: Case 4, we have: 1671 $TA_3 \Rightarrow I.$ $TA_3 \Rightarrow E.$ 1672 $TA_4 \Rightarrow E \text{ or } I.$ $TA_4 \Rightarrow E.$ 1673

1674 Clearly, the only absorbing state $(TA_3 = I \text{ and } TA_4 = E)$ becomes recurrent due to the newly added transition (the red I for TA₄). As a result, the system is recurrent when $u_1 > 0$ and $u_2 > 0$.

1677 1678

D.1.2 TRANSITION OF TM WITH AND OPERATOR WHEN T CAN BLOCK TYPE I FEEDBACK

1679 Based on the above analysis, we understand that the system is recurrent when $u_1 > 0$ and $u_2 > 0$. 1680 Next, we will examine whether there is any possibility of the system becoming absorbing when T 1681 can block Type I feedback.

1682 Clearly, when T clauses learn the intended pattern $\mathbf{X} = [x_1 = 1, x_2 = 1]$, i.e., when $f_{\sum}(\mathcal{C}^i(\mathbf{X})) =$ 1683 $T, u_1 = 0$ holds, and Type I feedback is blocked. In this situation, only Type II feedback can occur. 1684 Due to the presence of the wrong label, i.e., $([x_1 = 1, x_2 = 1], y = 0)$, Type II feedback triggers transitions in the TAs that have already learned the intended logic ($([x_1 = 1, x_2 = 1], y = 1)$). 1685 For example, Type II feedback will cause a transition in TAs of a learned clause $C = x_1 \wedge x_2$, 1686 making the clause deviate from its learned state (e.g., changing from $x_1 \wedge x_2$ to $x_1 \wedge x_2 \wedge \neg x_2$). 1687 Once this happens, $u_1 > 0$ holds, and Type I feedback is triggered by samples of ($[x_1 = 1, x_2 =$ 1688 1, y = 1, encouraging TAs in this clause to move back toward the action Exclude. Thus, even 1689 when T blocks all Type I feedback samples (setting $u_1 = 0$), the system remains recurrent due to 1690 the wrong label and Type II feedback. Notably, no value of $f_{\Sigma}(\mathcal{C}^i(\mathbf{X}))$ can make both $u_1 = 0$ and 1691 $u_2 = 0$ simultaneously⁴. Therefore, Type I and Type II feedback cannot be blocked simultaneously, ensuring the system is recurrent. 1693

1694 1695

D.2 THE AND OPERATOR WITH THE SECOND TYPE OF WRONG LABELS

To properly define the training samples with the second type of wrong label, we employ the following formulas:

1699
$$P(y=1|x_1=1, x_2=1) = 1,$$
(24)1700 $P(y=0|x_1=1, x_2=0) = a, a \in (0,1)$ 1701 $P(y=1|x_1=1, x_2=0) = 1 - a,$ 1702 $P(y=0|x_1=0, x_2=1) = 1,$ 1703 $P(y=0|x_1=0, x_2=0) = 1.$

1705 In this case, clearly, label of the training samples $[x_1 = 1, x_2 = 0]$ are wrongly labeled as 1 with 1706 probability 1 - a. In other words, in addition to the training samples detailed in Subsection B, a new 1707 (wrongly labeled) training sample will appear to the system, namely $([x_1 = 1, x_2 = 0], y = 1)$.

Lemma 7. *The TM is recurrent for the training samples given by Eq. (24).*

Proof: Similar to the proof of Lemma 6, we first consider the transitions of TM with $u_1 > 0$ and $u_2 > 0$, and then examine the impact of T for the system transition.

1712 Clearly, when $u_1 > 0$, $u_2 > 0$ holds, there is a non-zero probability that training sample ($[x_1 = 1, x_2 = 0], y = 1$) will appear to the system. The appearance of this sample will involve transition of TA₃ moving from action Include toward Exclude, as shown in Fig. 3, making the system recurrent.

1715 1716 1717 1717 1718 1719 Now we study if the functionality of T can offer system absorption. When T clauses learn the intended pattern $\mathbf{X} = [x_1 = 1, x_2 = 1]$, i.e., $f_{\sum}(\mathcal{C}^i(\mathbf{X})) = T$, $u_1 = 0$ holds, and thus Type I feedback is blocked for this training sample. In this situation, the TM can only see the training samples following:

$P(y = 0 x_1 = 1, x_2 = 0) = a, a \in (0, 1)$	(25)
$P(y = 1 x_1 = 1, x_2 = 0) = 1 - a,$	
$P(y=0 x_1=0, x_2=1) = 1,$	

⁴In this study, we focus only on positive polarity thus $u_2 > 0$ always holds. When negative polarity is enabled (i.e., when a set of clauses learns sub-patterns with label y = 0), u_2 becomes 0 when T clauses learn a sample with y = 0. However, it remains true that no value of $f_{\sum}(\mathcal{C}^i(\mathbf{X}))$ can make both u_1 and u_2 equal to 0 simultaneously.

Following the same concept as the proof of Lemma 6, we can conclude that the TM is recurrent for the samples in Eq. (25). Clearly, the system is recurrent, regardless of the value of u_1 . Therefore, we can conclude that the TM is recurrent for the training samples described in Eq.(24).

Following the same principle, we can also prove that the TM is recurrent when other training samples, i.e., $[x_1 = 0, x_2 = 1]$, and $[x_1 = 0, x_2 = 0]$, or their combinations, have wrong labels. We thus can conclude that the TM is recurrent for the second type of wrong labels.

So far, we have proven that the TM is recurrent when only one type of wrong label exists for the AND operator. It is straightforward to conclude that the TM remains recurrent when both types of wrong labels are present. The key reason is that adding both types of wrong labels does not eliminate any transitions between system states in recurrent systems. Therefore, the TM is recurrent for training samples with general wrong labels for the AND operator. Using the same reasoning, we can extend this conclusion to the XOR and OR operators. Thus, the following theorem holds.

Theorem 6. *The TM is recurrent given training samples with wrong labels for the AND, OR, and XOR operators.*

Remark 5. The primary reason for the recurrent behavior of the TM when wrong labels are present is the introduction of statistically conflicting labels for the same input samples. These inconsistency causes the TAs within a clause to learn conflicting outcomes for the same input due to the corresponding Type I and Type II feedback for label 1 and 0 respectively. When a clause learns to evaluate an input as 1 based on Type I feedback, samples with a label of 0 for the same input prompt it to learn the input as 0 during Type II feedback. This conflict in labels confuses the TM, leading to back-and-forth learning.

Remark 6. Note that although wrong labels will make the TM not converge (not absorbing with 100% accuracy for the intended logic), via simulations, we can still find that the TM can learn the operators efficiently, which is to be demonstrated in Appendix G, especially when the probability of wrong label is small. Interestingly, when the probability of the second type of wrong label is large, TM will consider it as a sub-pattern, and learn it, which aligns with the nature of learning.

1782 APPENDIX: ANALYSIS OF THE TM WITH AN IRRELEVANT INPUT Ε 1783 VARIABLE 1784

1785 In this appendix, we examine the impact of irrelevant input noise on the TM. Irrelevant noise refers 1786 to an input bit with a random value that does not affect the classification result. For instance, in the 1787 AND operator, a third input bit, x_3 , may appear in the training sample with random 1 and 0 values, 1788 but its value does not influence the output of the AND operator. In other words, the output is entirely 1789 determined by the values of x_1 and x_2 . Formally, we have: 1790

1790		(00)
1791	$P(y = 1 x_1 = 1, x_2 = 1, x_3 = 0 \text{ or } 1) = 1,$	(26)
1792	$P(y=0 x_1=1,x_2=0,x_3=0 \ or \ 1)=1,$	
1793	$P(y = 0 x_1 = 0, x_2 = 1, x_3 = 0 \text{ or } 1) = 1,$	
1794	$P(y=0 x_1=0, x_2=0, x_3=0 \text{ or } 1)=1.$	
1795		
1796	Here $x_3 = 0$ or 1 means $P(x_3 = 0) = a$, $P(x_3 = 1) = 1 - a$, $a \in (0, 1)$.	
1797		
1798	E.1 CONVERGENCE ANALYSIS OF THE AND OPERATOR WITH IRRELEVANT VARIABLE	
1799	Theorem 7. The clauses in a TM can almost surely learn the AND logic given training samp	las in
1800		ies in
1801	Eq. (26) in infinite time, when $T \leq m$.	
1001		

Proof: The proof of Theorem 7 consists of two steps: (1) Identifying a set of absorbing conditions 1803 and confirming that the TM, when in these conditions, satisfies the requirements of the AND operator. (2) Demonstrating that any state of the TM that deviates from the conditions defined in step (1) 1805 is not absorbing.

1806 The TM will be absorbed when the following conditions fulfill: 1807

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1808 1. Condition to block Type I feedback: For any input sample $\mathbf{X} = [x_1 = 1, x_2 = 1, x_3]$ 1809 regardless of whether $x_3 = 1$ or 0, the TM has at least T clauses that output 1. 1810 2. Conditions to guarantee no action upon Type II feedback: 1811 (a) $\mathbf{W}\mathbf{h}$ use in the TM: The literals that a a ta a la la ta a a1a.

1812	(a) When x_3 or $\neg x_3$ appears in a clause in the TM: The literals that are included in the
	clause for the first two input variables must result in a clause value of 0 for the input
1813	samples $\mathbf{X} = [x_1 = 0, x_2 = 1, x_3], \mathbf{X} = [x_1 = 1, x_2 = 0, x_3]$ and $\mathbf{X} = [x_1 = 0, x_2 = 1, x_3], \mathbf{X} = [x_1 = 1, x_2 = 0, x_3]$
1814	$0, x_2 = 0, x_3$]. This ensures that $C = 0$ for these input samples, regardless of the value
1815	of x_3 , thereby preventing transitions caused by any Type II feedback. The portion of
1816	the clause involving the first two input variables can be, e.g., $x_1 \wedge x_2$ or $x_1 \wedge \neg x_1 \wedge x_2$,
1817	while the overall clauses can be, e.g., $C = x_1 \wedge x_2 \wedge x_3$, or $C = x_1 \wedge \neg x_1 \wedge x_2 \wedge \neg x_3$,
1818	as long as the resulted clause value is 0 for those input samples.
1819	(b) When x_2 or $\neg x_2$ does NOT appear in a clause in the TM. There is no clause that is in

(b) When x_3 or $\neg x_3$ does NOT appear in a clause in the TM: There is no clause that is in the form of $C = x_1, C = x_2, C = x_1 \land \neg x_2, C = \neg x_1 \land x_2, C = \neg x_1, C = \neg x_2$, or $C = \neg x_1 \land \neg x_2.$

Clearly, when the above conditions fulfill, the system is in absorption because no feedback appears to the system. Additionally, this absorbing state follows AND operator. Based on the statement of 1824 the condition to block Type I feedback, there are at least T clauses that output 1 for input sample 1825 $\mathbf{X} = [x_1 = 1, x_2 = 1, x_3]$, regardless $x_3 = 1$ or 0. Studying the conditions for Type II feedback, we 1826 can conclude that the clause outputs 0 for all input samples $\mathbf{X} = [x_1 = 1, x_2 = 0, x_3], \mathbf{X} = [x_1 = 1, x_2 = 0, x_3]$ 1827 $0, x_2 = 1, x_3$, or $\mathbf{X} = [x_1 = 0, x_2 = 0, x_3]$. We can then setup the Th = T to confirm the AND 1828 logic. 1829

The next step is to show that any state of the TM deviating from the above conditions is not absorb-1830 ing. To demonstrate this, we can simply confirm that transitions, which might change the current 1831 actions of the TAs, will occur due to updates from Type I or Type II feedback. 1832

1833 When literal x_3 or literal $\neg x_3$ is included as a part of the clause, there is non-zero probability for C = 0 due to the randomness of input variable x_3 . As a result, Type I Feedback will encourage 1834 the TA for the included literal x_3 or $\neg x_3$ to move away from its current action, thus preventing the 1835 system from becoming absorbing.

For the case where literal x_3 or literal $\neg x_3$ is not included in the clause, the system operates purely based on the first two input variables, namely x_1 and x_2 . According our previous analysis for the noise free AND case (Theorem 1), there is only one absorbing status, which is $C = x_1 \land x_2$. However, this absorbing state disappears because Type I feedback will encourage the excluded literal x_3 to be included when $x_3 = 1$, and similarly encourage the excluded literal $\neg x_3$ to be included when $x_3 = 0$. Once either x_3 or $\neg x_3$ is included, the analysis in the previous paragraph applies, and thus the system is not absorbing.

1843 From the above discussion, it is clear that Type I feedback is the key driver of action changes in **1844** non-absorbing cases. If Type I feedback is not blocked, the system cannot reach an absorbing state. **1845** Therefore, blocking Type I feedback is critical for achieving convergence. The condition T < m **1846** is to guarantee that T should not be greater than the total number of clauses, making it feasible to **1847** block Type I feedback.

Remark 7. Due to the existence of the noisy input x_3 , the system requires the functionality of T to block Type I feedback in order to converge. This contrasts with the noise-free case, where the TM will almost surely converge to the AND operator even when Type I feedback is consistently present $(u_1 > 0)$.

⁸⁵³ E.2 CONVERGENCE ANALYSIS OF THE OR OPERATOR WITH IRRELEVANT VARIABLE

1855 For the OR case, we have

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1050	D(1) 1 1 1 0 1 1	(07)
1857	$P(y = 1 x_1 = 1, x_2 = 1, x_3 = 0 \text{ or } 1) = 1,$	(27)
1858	$P(y = 1 x_1 = 1, x_2 = 0, x_3 = 0 \text{ or } 1) = 1,$	
1859	$P(y = 1 x_1 = 0, x_2 = 1, x_3 = 0 \text{ or } 1) = 1,$	
1860	$P(y = 0 x_1 = 0, x_2 = 0, x_3 = 0 \text{ or } 1) = 1.$	
1961		

Theorem 8. The clauses in a TM can almost surely learn the OR logic given training samples in Eq. (27) in infinite time, when $T \leq \lfloor m/2 \rfloor$.

Proof: The proof of Theorem 8 follows a similar structure to that of the AND case and involves two steps: (1) Identifying a set of absorbing conditions and verifying that, under these conditions, the TM satisfies the requirements of the OR operator. (2) demonstrating that any state of the TM deviating from these conditions is not absorbing.

- 1. Condition to block Type I feedback: For any input sample $\mathbf{X} = [x_1 = 1, x_2 = 1, x_3]$, $\mathbf{X} = [x_1 = 1, x_2 = 0, x_3]$, and $\mathbf{X} = [x_1 = 0, x_2 = 1, x_3]$ regardless of whether $x_3 = 1$ or 0, the TM has at least T clauses that output 1.
- 2. Conditions to guarantee no action upon Type II feedback:
- (a) When x₃ or ¬x₃ appears in a clause in the TM: The literals included in the clause for the first two input variables must ensure a clause value of 0 for the input samples X = [x₁ = 0, x₂ = 0, x₃]. This is to guarantee that C = 0 for those input samples, irrespective of the value of x₃, thereby preventing any transitions caused by Type II feedback. The portion of the clause involving the first two input variables can take the form such as x₁, x₁∧¬x₂, x₁∧x₂, x₁∧¬x₁∧x₂. Correspondingly, the overall clauses can take the form such as C = x₁ ∧ ¬x₃, C = x₁ ∧ ¬x₂ ∧ x₃, or C = x₁ ∧ ¬x₁ ∧ x₂ ∧ ¬x₃, as long as the resulted clause value is 0 for those input samples.
- 1883
- (b) When x_3 or $\neg x_3$ does not appear in a clause in the TM: There are no clauses with literal(s) in only negated form, such as $C = \neg x_1, C = \neg x_2$, or $C = \neg x_1 \land \neg x_2$.

1885 Clearly, when the above conditions fulfill, the system is absorbing because no feedback appears to 1886 the system. Additionally, this absorbing state adheres to the OR operator. Based on the condition 1887 required to block Type I feedback, there are at least T clauses that output 1 for input sample $\mathbf{X} =$ 1888 $[x_1 = 1, x_2 = 1, x_3], \mathbf{X} = [x_1 = 1, x_2 = 0, x_3], \text{ or } \mathbf{X} = [x_1 = 0, x_2 = 1, x_3]$ regardless of whether 1889 $x_3 = 1$ or 0. Analyzing the conditions for Type II feedback, we find that the clause outputs 0 for all 1890 input samples $\mathbf{X} = [x_1 = 0, x_2 = 0, x_3]$. We can then setup the Th = T to confirm the OR logic. The next step is to demonstrate that any state of the TM that deviates from the above conditions outlined above is not absorbing. To do this, we can confirm that transitions which may alter the current actions of the TAs will occur due to updates from Type I and Type II feedback.

When literal x_3 or literal $\neg x_3$ is included in the clause, there is non-zero probability for C = 0 due to the randomness of input variable x_3 . In this case, Type I Feedback will move the included literal x_3 or $\neg x_3$ towards action Exclude, preventing the system from being absorbing.

For the case where literal x_3 or literal $\neg x_3$ is not included as a part of the clause, the system operates purely based on the first two input variables, namely x_1 and x_2 . Based on our previous analysis for the noise free OR case shown in Lemma 2, the system is recurrent. This recurrent behavior will eventually lead the system to a state where the excluded literal, either x_3 or $\neg x_3$, is encouraged to be included. For example, if the TM has a clause $C = x_1 \land x_2$, upon a training sample $\mathbf{X} = [x_1 = 1, x_2 = 1, x_3 = 0]$, the Type I feedback will encourage the excluded literal $\neg x_3$ to be included. Once one of the excluded literal, x_3 or $\neg x_3$, is included, the analysis in the previous paragraph applies, meaning the system is not absorbing.

Clearly, if Type I feedback is not blocked, the system will not be absorbing. As blocking Type I feedback is critical, condition $T \le \lfloor m/2 \rfloor$ is necessary, refer to Lemma 4.

1907 When T clauses have learned the intended sub-patterns of OR operation, the Type I feedback will 1908 be blocked. At the same time, Type II feedback will eliminate all clauses that output 1 for input 1909 sample following $\mathbf{X} = [x_1 = 0, x_2 = 0, x_3]$, removing false positives. At this point, the system has 1910 converged. The presence of x_3 does not change the convergence feature, but it adds more dynamics 1911 to the TM.

1913 E.3 CONVERGENCE ANALYSIS OF THE XOR OPERATOR WITH IRRELEVANT VARIABLE

Theorem 9. The clauses in a TM can almost surely learn the XOR logic given training samples in Eq. (28) in infinite time, when $T \leq \lfloor m/2 \rfloor$.

Eq. (28) in infinite time, when $T \leq \lfloor m/2 \rfloor$.

 $P(y=0|x_1=1, x_2=1, x_3=0 \text{ or } 1) = 1,$

 $P(y = 1 | x_1 = 1, x_2 = 0, x_3 = 0 \text{ or } 1) = 1,$

 $P(y = 1 | x_1 = 0, x_2 = 1, x_3 = 0 \text{ or } 1) = 1,$

 $P(y=0|x_1=0, x_2=0, x_3=0 \text{ or } 1) = 1.$

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The proof for XOR follows the same principles as the AND and OR cases, and therefore, we do not present it explicitly here.

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1927 E.4 CONVERGENCE ANALYSIS OF THE OPERATORS WITH MULTIPLE IRRELEVANT
 1928 VARIABLES

In the previous subsections, we demonstrated that if a single irrelevant bit is present in the training samples, the system will almost surely converge to the intended operators. This conclusion can be readily extended to scenarios involving multiple irrelevant variables. Here, "multiple irrelevant variables" refers to the presence of additional variables, beyond x_3 , in the training samples that do not contribute to the classification.

Theorem 10. The clauses in a TM can almost surely learn the 2-bit AND logic given training samples with k irrelevant input variables in infinite time, $0 < k < \infty$, when $T \le m$.

Theorem 11. The clauses in a TM can almost surely learn the 2-bit XOR and OR logic given training samples with k irrelevant input variables in infinite time, $0 < k < \infty$, when $T \le |m/2|$.

Proof: The proofs of Theorems 10 and 11 are straightforward. It suffices to verify whether the conditions for blocking Type I and Type II feedback remain valid when multiple irrelevant variables are present.

1943 The condition for blocking Type I feedback remains valid because Type I feedback is only determined by the first two input bits and is not a function of the irrelevant variables. For Type II feedback, its effect depends on whether the literals for the irrelevant inputs are present in the clause. In cases where the literals of the irrelevant bits are not included in the clause, the analysis holds, as those literals are absent. When the literals of the irrelevant bits are included, their number does not impact the analysis. This is because the clause value is entirely determined by the first two bits, and the clause value remains C = 0, regardless of the number of irrelevant variables.

F APPENDIX: EXPERIMENT RESULTS WITH NOISE-FREE TRAINING SAMPLES

To validate the theoretical analyses, we here present the experiment results⁵ for both the AND and the OR operators.

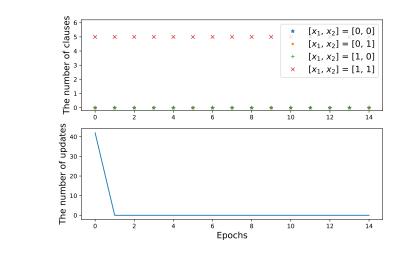


Figure 8: The convergence of a TM with 7 clauses when T = 5 for the AND operator.

Figure 8 shows the convergence of TM for the AND operator when m = 7, T = 5, s = 4, and N = 50. More specifically, we plot the number of clauses that learn the AND operator, namely, $x_1 = x_2 = 1$, and the number of system updates as a function of epochs. From these figures, we can clearly observe that after a few epochs, the TM has 5 clauses that learn the AND operator and then the system stops updating because no update is triggered anymore. Note that if we control T so that $u_1 > 0$ always holds, all clauses will converge to the AND operator, which has been validated via experiments. These observations confirm Theorem 1. Although the theorem says it may require infinite time in principle, the actual convergence can be much faster.

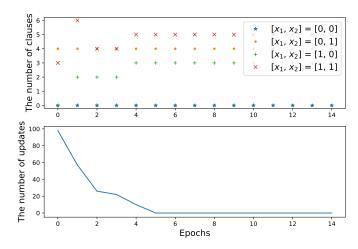


Figure 9: The convergence of a TM with 7 clauses when T = 3 for the OR operator.

In Fig. 9, we illustrate the number of clauses in distinct sub-patterns when we employ m = 7, T = 3, s = 4, and N = 50 for the OR operator. Based on the analytical result, i.e., Theorem 2, the system will be absorbed, where each sub-pattern will have at least 3 clauses and no update will happen afterwards. From the figure, we can clearly observe that after a few epochs, the system

⁵The code for validating the convergence of AND and OR operators can be found at https://github.com/JaneGlim/Convergence-of-Tsetline-Machine-for-the-AND-OR-operators.

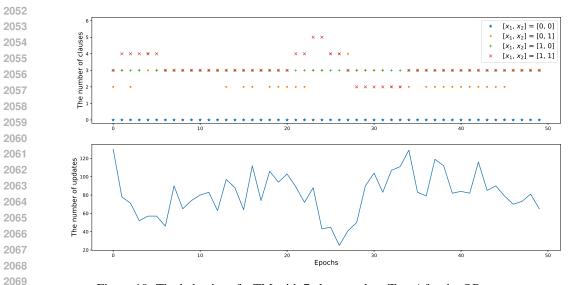


Figure 10: The behavior of a TM with 7 clauses when T = 4 for the OR operator.

2072 becomes indeed absorbed as no updates are observed. When absorbed, two intended sub-patterns 2073 have 3 clauses individually, one intended sub-pattern has 4 clauses, while the unintended sub-pattern 2074 has 0 clause, which coincides with the theorem. After checking the converged actions of the TAs, we 2075 find the list of the converged clauses: $C_1 = x_1, C_2 = x_1, C_3 = x_2, C_4 = x_1 \land \neg x_2, C_5 = x_1 \land x_2, C_5 = x_2 \land x_2, C_5 = x_2 \land x_2, C_5 = x_2 \land x_2 \land x_3 \land x_4 \land$ $C_6 = x_2$, and $C_7 = \neg x_1 \wedge x_2$, which explained the number of the converged clauses in different 2076 sub-patterns shown in the figure. Clearly, in this example, some clauses, such as C_1 and C_3 , can 2077 cover multiple sub-patterns. This indicates that in real world applications, if distinct sub-patterns 2078 have certain bits in common, which can be used to differentiate it from other classes, it is possible 2079 for TM to learn those bits as jointly features, confirming the efficiency of the TM. 2080

Note that there are many other possible absorbing states that are different from the shown example, which have been observed when we run multiple instances of the experiments. As long as at least Tclauses can cover each intended sub-pattern in the OR operator, the system converges.

In Fig. 10, the configuration is identical to that in Fig. 9 other than T = 4. In this case, as stated in Remark 2, the system will not become absorbing, but will still cover the intended sub-patterns with high probability. From this figure, we can observe that at least two clauses are able to cover each intended sub-pattern and the unintended sub-pattern has zero clause. At the same time, the TAs will update their states along epochs, which can be seen in the bottom figure. It is worth mentioning that we have occasionally observed in other rounds of experiments, that one sub-pattern is covered by only 1 clause. In this case, it is still possible to set up $Th \ge 1$ to have successful classification. Nevertheless, there is no guarantee that at least one clause will follow each intended sub-pattern in this configuration.

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²¹⁰⁶ G APPENDIX: EXPERIMENT RESULTS WITH NOISY TRAINING SAMPLES

This Appendix presents the experimental results for the operators under noisy conditions. First, we show the results when incorrect labels are present, followed by the results involving irrelevant variables. The final subsection addresses a case where both incorrect labels and irrelevant variables are present.

2113 G.1 EXPERIMENT RESULTS FOR WRONG LABELS

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To evaluate the performance of the TM when exposed to mislabeled samples, we introduced incorrectly labeled data into the system. The key observation is that the TM does not converge to the intended logic, meaning it does not absorb into a state where the correct logic is consistently represented. However, with carefully chosen hyperparameters, the TM can still learn the intended logic with high probability.

To demonstrate the TM's behavior, we first conduct experiments on the OR operator, which satisfies the following equation:

2123	$P(y = 1 x_1 = 1, x_2 = 1) = 90\%,$	(29)
2124	$P(y=1 x_1=1, x_2=0) = 90\%,$	
2125	$P(y=1 x_1=1, x_2=0) = 90\%,$	
2126	$P(y=0 x_1=0, x_2=0) = 1.$	
2127	$1 (g - 0 x_1 - 0, x_2 - 0) - 1.$	

2128 In this scenario, 10% of the input samples that should be labeled as 1 were incorrectly labeled as 2129 0. To train the TM and evaluate its performance, we used the following hyperparameters: T = 4, 2130 Th = 2, s = 3, c = 7, and N = 100. Fig. 11 shows the number of updates and the number of 2131 clauses that learn distinct sub-patterns, as a function of epochs. As shown in Fig. 11, the number 2132 of updates is big, and thus the system did not converge. Nevertheless, when examining the number 2133 of clauses associated with each sub-pattern, we observed that each sub-pattern was covered by at 2134 least two clauses, ensuring that the OR operator remained valid. Similar results were observed in 2135 experiments conducted on the AND and XOR operators.

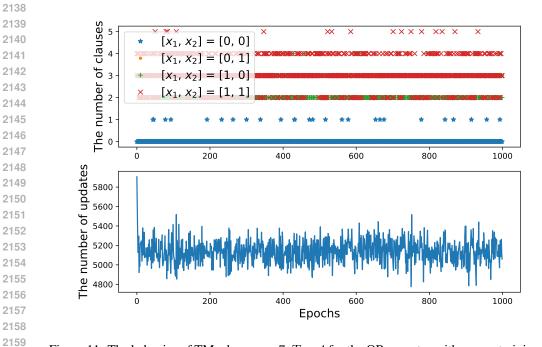


Figure 11: The behavior of TM when m = 7, T = 4 for the OR operator with wrong training labels.

Interestingly, when the proportion of mislabeled samples increases to an extreme level, where inputs that should be labeled as 0 are instead labeled as 1, the TM begins to treat the noise as a sub-pattern. For instance, consider the AND operator with input $\mathbf{X} = [x_1 = 0, x_2 = 1]$, which is mislabeled as 1 in 90% of the cases, as shown in Eq. (30). Using the hyperparameters T = 3, s = 3.0, c = 7, and N = 100, we observed from experiments that the TM generates three clauses with an output of 1 for $\mathbf{X} = [x_1 = 0, x_2 = 1]$ and another three clauses with an output of 1 for $\mathbf{X} = [x_1 = 1, x_2 = 1]$. This behavior indicates that the TM has incorporated the noise as a learned sub-pattern. Such outcomes align with the TM's underlying principle of learning, where it identifies and models sub-patterns associated with the label 1.

(30)

$P(y=1 x_1=1, x_2=1) = 1,$
$P(y=1 x_1=1, x_2=0) = 1,$
$P(y=0 x_1=0, x_2=1) = 90\%,$
$P(y=0 x_1=0, x_2=0) = 1.$



G.2 EXPERIMENT RESULTS FOR IRRELEVANT VARIABLE

2177 To confirm the convergence property of TM with irrelevant variable, we setup the experiments for 2178 the AND, OR, and XOR operators when one irrelevant variable, namely, x_3 , exists. The probability 2179 of x_3 being one in the training and testing samples is 50%.

For the AND operator, we use the hyperparameters c = 5, T = 2, s = 3, Th = 2, and N = 100. Fig. 12 illustrates the convergence of TM for the AND operator in the presence of an irrelevant bit. The results confirm that the TM can correctly learn the AND operator without uncertainty, validating the correctness of Theorem 7.

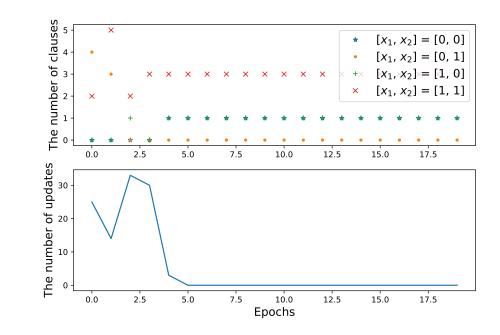


Figure 12: Convergence of TM when m = 5, T = 2 for the AND operator with an irrelevant label.

2209 Interestingly, upon convergence, the form of the included literals varies. For instance, with the 2210 aforementioned hyperparameters, we observe that the converged TM includes two clauses of the 2211 form $x_1 \wedge x_2 \wedge x_3$ and another two clauses of the form $x_1 \wedge x_2 \wedge \neg x_3$. This suggests that, instead 2212 of excluding the irrelevant bit x_3 , the TM includes at least T clauses containing x_3 and at least T 2213 clauses containing $\neg x_3$, which ensures correct classification regardless of the value of x_3 . However, when the hyperparameters are set to c = 1, T = 1, s = 3, Th = 1, and N = 100, where only a

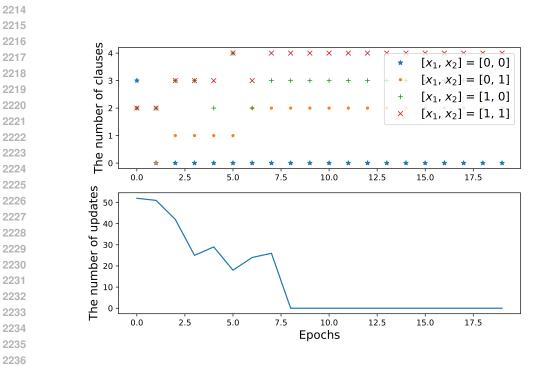


Figure 13: Convergence of TM when m = 5, T = 2 for the OR operator with an irrelevant label.

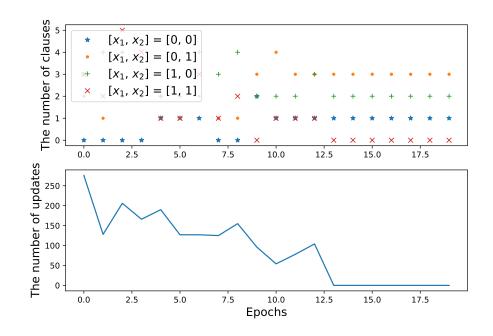


Figure 14: Convergence of TM when m = 7, T = 2 for the XOR operator with an irrelevant label.

single clause exists in the TM, the converged clause takes the form $x_1 \wedge x_2$, excluding the literals x_3 and $\neg x_3$.

As T increases (T > m/2), we observe that convergence becomes challenging. This difficulty arises because the TM cannot simultaneously learn T clauses containing x_3 and another T clauses

containing $\neg x_3$. In such cases, the TM must rely on T clauses in the form $x_1 \land x_2$ to achieve convergence, which can be particularly demanding.

For the OR operator, we use the hyperparameters c = 5, T = 2, s = 3, Th = 2, and N = 100. Figure 13 illustrates the convergence of the TM for the OR operator in the presence of an irrelevant bit. The results confirm that the TM successfully learns the OR operator without ambiguity, validating the correctness of Theorem 8.

From the experimental results, we also observe that there are multiple possible absorbing states, as long as the absorbing conditions are satisfied. Additionally, the TM is capable of presenting two sub-patterns simultaneously. Depending on the hyperparameter configuration, x_3 and $\neg x_3$ may be included in the clauses, provided that T clauses can align with each intended sub-pattern, which ensures correct classification regardless of the value of x_3 .

2280 We have also studied the XOR operator. The convergence instance is shown in Fig. 14, confirming 2281 Theorem 9. Here we use m = 7, T = 2, s = 3, Th = 2.

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2283 G.3 EXPERIMENT RESULTS FOR BOTH WRONG LABELS AND IRRELEVANT VARIABLES

In this experiment, we assess the performance of the TM in the presence of both mislabeled data and irrelevant variables. Specifically, we evaluate the TM's ability to learn the XOR operator when 40% of the samples are incorrectly labeled, and 10 irrelevant variables are added. The input comprises 12 bits, with only the first two bits determining the output based on the XOR logic.

The hyperparameters are configured as follows: T = 15, s = 3.9, c = 20, and N = 100 with polarity enabled. Experimental results reveal that the TM successfully learns the XOR operator in 99% of 200 independent runs. These findings demonstrate the robustness of the TM training in noisy environments.

In another experiment, we configured the TM to learn a noisy XOR function with 2 useful input bits and 18 irrelevant input bits (hyper parameters: N = 128, m = 20, T = 10, s = 3, label noise 0.1). Remarkably, the TM was still able to learn the XOR operator with 100% accuracy using just 5000 training samples. If all possible input combinations were required in the training samples, we would require $2^{20} = 1048576$ samples. Clearly, the TM does not rely on the entire combinatorial input space to learn effectively.

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