Meta-sketch: A Neural Data Structure for Estimating Item Frequencies of Data Streams

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Abstract

1	To estimate item frequencies of data streams with limited space, sketches are widely
2	used in real applications, including real-time web analytics, network monitoring,
3	and self-driving. Sketches can be viewed as a model which maps the identifier of a
4	stream item to the corresponding frequency domain. Starting from the premise, we
5	envision a neural data structure, which we term the <i>meta-sketch</i> , to go beyond the
6	basic structure of conventional sketches. The meta-sketch learns basic sketching
7	abilities from meta-tasks constituted with synthetic datasets following Zipf distribu-
8	tions in the pre-training phase, and can be fast adapted to real (skewed) distributions
9	in the adaption phase. Extensive experiments demonstrate the performance gains
10	of the meta-sketch and offer insights into our proposals.

11 **1 Introduction**

Estimating item frequency is a basic topic in data stream processing, which finds applications in
the fields of networking, databases, and machine learning, such as real-time data analyzing [1–4],
network traffic monitoring [5–7], natural language processing [8] and search ranking [9]. Towards
infinite data streams, a common class of solutions [10–15] use a compact structure taking sublinear
space for counting the number of occurrences of each stream item, called the *sketch*.

Under the prevalent evidence of skewed distributions in data streams, basic sketches achieve the space 17 compactness by hashing and approximately aggregating stream items. Basic sketches, including 18 CM-sketch [10], C-sketch [11] and CU-sketch [12], use a 2D array of counters as the core structure. 19 To optimize the sketching performance, there arise *augmented sketches* [13, 14], which attach filters to 20 basic sketches, to capture the preliminary patterns of skewed distributions (e.g., high/low-frequency 21 items). By separately maintaining the filtered high/low-frequency items, augmented sketches strive 22 to eliminate the estimation error incurred by hash collisions between the high- and low-frequency 23 items. Further, learned augmented sketches [15] improve the filters of the augmented sketches by 24 memorizing short-term high/low-frequency items via a pre-trained neural network (NN in short) 25 classifier. But it is not clear how the pre-trained NN can be adapted to dynamic streaming scenarios, 26 where the correspondence between items and frequencies varies. In a nutshell, sketches are structures 27 compactly summarizing stream distributions to count item frequencies with limited space budgets. 28 From the retrospective analysis of sketches, an observation can be drawn that the evolution of 29 sketches conforms with the exploitation of data distributions. It is thus a natural evolution to consider 30

a sketch that generally and automatically captures more distribution patterns with limited space

³² budgets. In this paper, we envision a novel neural sketch, called the *meta-sketch*, with techniques

33 of meta-learning and memory-augmented neural networks. The meta-sketch learns the sketching

³⁴ abilities from automatically generated meta-tasks. Depending on the types of meta-tasks, we study

two versions of the meta-sketch, called *basic* and *advanced meta-sketches*.

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Figure 1: The Framework of the Meta-sketch

36 The basic meta-sketch implements the simulation of basic sketches, through the training process with basic meta-tasks following Zipf distributions, which are prevalent in the scenes of real data streams [16-37 20]. The advanced meta-sketch extends the basic version to fast adapt to the specific runtime of stream 38 processing, through the training with adaptive meta-tasks, which are generated by online sampling 39 of real data streams. Our work follows a typical setting where the distribution of item frequencies 40 follows a skewed distribution, but the correspondence between items and frequencies varies. For 41 example, in software-defined networks (SDN), sketches are deployed to programmable switches to 42 collect per-flow statistics, where IP packets follow *heavy-tailed* distributions [15,21]. In distributed 43 databases, it gives advances to collect statistics of data shards to optimize data placement and 44 45 query caching, where query phrases follow approximate *Zipf* distributions [15]. Given that the item population follows a specific distribution, the local distributions, i.e., item-frequency correspondences 46 on shards or flows, are different. Instead of retraining learned augmented sketches on each local 47 distribution, the advanced-sketch can be quickly adapted to different local distributions once trained. 48

As a member of the neural data structure family [15, 22–24], the meta-sketch significantly differs from conventional sketches, in terms of the structure and working mechanism. The meta-sketch utilizes NN's powerful encoding/decoding capabilities to perceive data distributions and express and compress explicit or implicit information to retrieve item frequencies with better accuracies. Meanwhile, the meta-sketch is differentiable to fully perceive frequency patterns for self-optimization.

Our contributions are as follows. 1) We propose the meta-sketch, the first neural data structure for the problem of item frequency estimation, based on meta-learning. 2) The basic meta-sketch acquires sketching abilities by learning from synthetic datasets, and outperforms basic sketches in real datasets. The advanced meta-sketch automatically encompasses the ability analogous to the auxiliary structures deliberately devised in (learned) augmented sketches, yet yielding better accuracies and robustness when adapted to dynamic scenes. 3) Through extensive empirical studies on real and synthetic datasets, we evaluate our proposed meta-sketches and analyze the mechanism of major modules.

61 2 Meta-sketch Structure

62 2.1 Preliminaries

We consider a standard data stream scenario [19]. Suppose a data stream $S_N : \{e_1, ..., e_N\}$ with Nitems and n distinct items. Each item $e_i \in S_N$ takes a value from the item domain $\mathbb{X} = \{x_1, ..., x_n\}$ where $x_i \neq x_j$. The frequency f_i is equal to the number of times that item x_i appears in S_N .

To leverage learning techniques for item frequency estimation, a naïve way is to train a NN model 66 (e.g., MLP/LSTM) that learns/memorizes the mapping relationship between items and frequencies 67 with multiple training iterations, similar to [15, 22, 24]. However, it violates the typical setting of 68 69 stream processing where item observations are transient and are therefore handled in one pass [18]. 70 More, the costly procedure has to be repeated from the scratch for a new data stream. Inspired by the 71 meta-bloom filter [23], we consider a case of one-shot learning (fitting for one-pass stream processing) 72 by using meta-learning [25, 26] and memory-augmented networks [27, 28]. Meta-learning employs sampled meta-tasks to learn the ability to solve a class of domain tasks rather than memorizing patterns 73 for a specific task. The memory-augmented networks incorporate external memories into NN models, 74 significantly enhancing the potentials of NN models with more learnable parameters. Meanwhile, it 75 performs efficient and explicit operations (i.e., reading and storing) for external memories, allowing 76 NN models to process information similarly to conventional data structures. 77

The framework of the meta-sketch consists of 4 functional modules, *Embedding* (\mathcal{F}_E), *Sparse* addressing (\mathcal{F}_{Sa}), *Compressed storage* matrix (M), and *Decoding* (\mathcal{F}_{dec}), as shown in Figure 1. Like traditional sketches, the meta-sketch encodes and memorizes online stream items in one pass, and answers queries by decoding corresponding item-frequency information from the structure.

Thus, we define 2 operations, *Store* and *Query*. Specifically, the *Store* operation first passes each incoming stream item to \mathcal{F}_E for the embedding representation, and then writes the embedding vector into M, according to the address derived by \mathcal{F}_{Sa} . When estimating the frequency of an item, the ⁸⁵ *Query* operation calculates the item's address in M via \mathcal{F}_{Sa} , reads the corresponding information

vector from M, and decodes the item frequency by \mathcal{F}_{dec} from the retrieved information vector.

87 2.2 Modules

Embedding. The module \mathcal{F}_E has two purposes: **1**) performing representational transformation for an incoming item e_i and mapping it into a dense embedding vector z_i that holds implicit features about item-frequency distributions and serves as the basis for identifying stream items; **2**) decoupling the embedding vector z_i to obtain a refined vector r_i , which is used to derive the address for reading/writing on the compressed storage matrix M.

Accordingly, \mathcal{F}_E consists of the embedding network g_{emb} and the address network g_{add} . We assume 93 that an item $e_i \in S_N$ is numerically encoded for the unique identification, following the conventions 94 of stream processing [18, 19]. Thus, we have $z_i, r_i \leftarrow \mathcal{F}_E(e_i)$, where $z_i \leftarrow g_{emb}(e_i)$ and $r_i \leftarrow \mathcal{F}_E(e_i)$ 95 $g_{add}(z_i)$. Here, $z_i \in \mathbb{R}^{l_z}$ is an embedding vector of length l_z , and $r_i \in \mathbb{R}^{l_r}$ is a refined vector of 96 length l_r . The vector z_i serves multiple intents: 1) it makes a basis for deriving the address of an item 97 in \mathcal{F}_{Sa} ; 2) it serves as the compressed vector of an item written into M; 3) it works as a partial input 98 of \mathcal{F}_{dec} for decoding the item frequency; 4) it also plays the role of perceiving/compressing patterns 99 of a specific frequency distribution, as discussed in Section 5. In addition, to enhance the addressing 100 functionality and eliminate other interference factors, we decouple z_i to generate a refined vector r_i , 101 instead of using z_i directly for the addressing. 102

103 **Sparse addressing.** The module \mathcal{F}_{Sa} aims to derive the address a_i for storing the embedding vector 104 z_i into the storage matrix: $a_i \leftarrow \mathcal{F}_{Sa}(r_i)$. In terms of functionality, \mathcal{F}_{Sa} is analogous to the hash 105 functions of traditional sketches, except that \mathcal{F}_{Sa} is parameterized and differentiable. Specifically, 106 the addressing of the meta-sketch is done via a 3D addressing matrix A of parameters to be learned 107 and a sparse SoftMax function: $a_i \leftarrow SparseMax(r_i^T A)$, where $A \in \mathbb{R}^{d_1 \times l_r \times d_2}$. Then, the batch 108 matrix multiplication of A and the transpose of r_i results in the addressing vector $a_i \in \mathbb{R}^{d_1 \times 1 \times d_2}$.

The setting of d_1 and d_2 determines the size of address space for storing the embedding vectors. 109 Typical addressing methods [23, 28] use a 2D matrix $(l_r \times d_2)$ for recording the mapping of an 110 embedding vector to a slot (d_2 is the number of slots). In contrast, we add one more dimension d_1 111 to simulate the multi-hash setting of traditional sketches, in view of that a 2D addressing matrix 112 can reach a differentiable simulation of a hash function [23, 24]. Matrix A simulates multiple hash 113 functions, yielding robust frequency decoding and the rationality of the learning optimization. Note 114 that each 2D slice A^* of A is stacked from d_2 -unit vectors $b_i \in \mathbb{R}^{l_r}$ by normalizing the parameters 115 of A at each gradient update of the training process. Normalized A can avoid overflowing when 116 compressing its size by reducing data precisions and enhance the interpretability (see Section 5). 117

In addition, we utilize sparse SoftMax [29, 30] instead of SoftMax to normalize the address a_i . It brings the following benefits by constraining some bits of a_i to zero, which 1) promotes quick derivation during the back-propagation; 2) reduces the overhead of storage matrix accessing by skipping the slots of M corresponding to the "0" bits of a_i ; 3) leads to de-noising with the vector compression.

Compressed storage matrix. We use a matrix $M \in \mathbb{R}^{d_1 \times l_z \times d_2}$ to store an embedding vector $z_i \in \mathbb{R}^{l_z}$ in accordance to its address $a_i \in \mathbb{R}^{d_1 \times 1 \times d_2}$. The functionality of M is similar to the 2D array of counters in traditional sketches, yet yielding better capabilities in the storage compression. Traditional sketches store item counts. Differently, M stores embedding vectors, which have richer information compression capabilities, due to the diversity of value changes on different bits.

Decoding. Given a query item x_i , the module \mathcal{F}_{dec} , consisting of one NN component g_{dec} , decodes 128 the information corresponding to x_i , in order to obtain the estimated frequency f_i . The vector fed 129 into g_{dec} is the concatenation of vector $\{M \ominus a_i\}$, vector z_i , and the current number of items (i.e., N) 130 recorded in a counter, $\hat{f}_i \leftarrow g_{dec}(\{M \ominus a_i\}, z_i, N)$. The operator \ominus refers to the reading operation 131 for the storage matrix. The basic form of \ominus gives the operation as $M \ominus a_i = M a_i^{T2}$ [27,28]. For 132 optimization, we consider two optimized forms of \ominus , inspired by the "count-min" mechanism of the 133 CM-sketch. The first one gives the minimum value of each row in Ma_i^T , aiming to remove the noise 134 of other items. The second one gives the minimum value of each row in $Ma_i^T \circ \frac{1}{z_i}$, a normalized 135

¹In this paper, we control $l_r : l_z \approx 1 : 5$ to compress A.

 $^{^{2}}a_{i}^{T}$ means transpose operation for dim 1 and d_{2}

form of Ma_i^T . Here, \circ denotes the Hadamard product, and z_i requires broadcast operations to comply with its requirements. So, $\{M \ominus a_i\}$ refers to the concatenation of vectors generated by the basic form and the two optimized forms. Please refer to supplement materials for more details.

139 2.3 Operations

Operation Store is performed by feeding an incoming item e_i to \mathcal{F}_E and \mathcal{F}_{Sa} to obtain embedding vector z_i and address a_i , and then additively writing z_i to M, weighted by $a_i: M \leftarrow M + z_i a_i$. Here, other writing types [23, 26–28] can also be employed, but simple additive writing is more efficient and allows to compute gradients in parallel [23]. In addition, additive writing also allows to define an optional *Delete* operation for the meta-sketch (see the supplement materials).

Operation Query estimates the frequency of a given query item x_i . First, z_i and a_i are obtained, similar to that of operation Store. Then, the vectors $\{M \ominus a_i\}$ are retrieved from M and N can be easily obtained by a small counter. Finally, $\{M \ominus a_i\}$, z_i and N are jointly fed into g_{dec} to get the estimated frequency \hat{f}_i of x_i as the returned value. The two operations are shown in Algorithm 1.

	Algorithm 1: Operations	Algorithm 2: Training Framework
1 2 3 149 5 6 7 8 9		Data: Meta-sketch with all learnable parameters θ , Meta-task sampler R ; 1 while i not reach max training steps do 2 Sample a meta-task $t_i : \{s_i, q_i\} \sim R$ and count N ; 3 for $e_j^{(i)} \in s_i$ do Store $(e_j^{(i)}, M)$; end 4 for $x_j^{(i)}, f_j^{(i)} \in q_i$ do $\hat{f}_j^{(i)} \leftarrow Query(x_j^{(i)}, M, N)$; $\mathcal{L} + = LossFun(f_j^{(i)}, \hat{f}_j^{(i)})$; 5 Backprop through: $d\mathcal{L}/d\theta$ and update parameters: $\theta \leftarrow Optimizer(\theta, d\mathcal{L}/d\theta)$; 6 Normalize A; 7 Clear M ; 8 end

150 3 Meta-sketch training

151 3.1 Training Framework

The meta-sketch employs an efficient one-shot meta-training method [31]. The training process thus contains two phases, *pre-training* and *adaption* phases. In the pre-training phase, the meta-sketch learns an initial set of module parameters, including g_{emb} , g_{add} , A, and g_{dec} . The pre-training goes offline across training units, i.e., basic meta-tasks, to acquire the ability of stream frequency estimation. Then, in the adaption phase, the pre-trained meta-sketch goes fast across a set of lightweighted training units, i.e., adaptive meta-tasks, to quickly acquire the task-specific knowledge, i.e., parameters for sketching real data streams at runtime.

The training units, i.e., meta-tasks, are crucial for both phases. The training process of the meta-sketch on a single meta-task is equivalent to simulating storing and querying an instance of data streams while computing the estimation error to optimize the learnable parameters. Thus, a meta-task t_i consists of a store set s_i (also called a support set) and a query set q_i . The store set s_i can be viewed as an instance of data streams, $s_i : \{e_1^{(i)}, \dots, e_{N_i}^{(i)}\}$, where N_i is the number of stream items in s_i . The query set q_i can be represented by a set of items from the stream instance with paired frequencies in the store set s_i , formally, $q_i : \{(x_1^{(i)} : f_1^{(i)}), \dots, (x_{n_i}^{(i)} : f_{n_i}^{(i)})\}$, where n_i is the number of distinct items in s_i . In this work, we define two types of meta-tasks, *basic* (Section 3.2) and *adaptive* (Section 3.3) meta-tasks, corresponding to the pre-training and adaption phases, respectively.

The two training phases, that are based on different types of meta-tasks, follow the same training framework, as shown in Algorithm 2, except for the sampler and initial parameters. To optimize on reducing both absolute and relative frequency estimation errors³, we devise an adaptive hybrid loss function [32] for the meta-sketch: $\frac{1}{2\sigma_1^2}(f_i - \hat{f}_i)^2 + \frac{1}{2\sigma_2^2}|f_i - \hat{f}_i|/f_i + \log\sigma_1\sigma_2$, where σ_1 and σ_2 are

learned parameters, and f_i and \hat{f}_i are the true and estimated frequencies of item x_i , respectively.

³Average Absolute Error: $AAE = \frac{1}{n} \sum_{i=1}^{n} |f_i - \hat{f}_i|$; Average Relative Error: $ARE = \frac{1}{n} \sum_{i=1}^{n} \frac{|f_i - \hat{f}_i|}{f_i}$.

173 3.2 Basic Meta-task Generation

In the pre-training phase, basic meta-tasks should make the meta-sketch to simulate traditional sketches and preserve certain generality without relying too much on the patterns of specific distributions (Section 5). Therefore, we generate meta-tasks based on the Zipf distribution, which is found to

be prevalent in real scenes of data streams [16–20].

A meta-task is essentially a data stream instance with item size n, which can be determined by the total number of items N and the relative frequency distribution p. Alternatively, we can generate meta-tasks by presupposing different n, \bar{f} and p, where \bar{f} is the frequency mean, since $N = \bar{f} \times n$. Thus, basic meta-task generation is based on a sampler $R : \{I, L, P\}$, as follows.

An **item pool** I is a subset of the item domain X. The cardinality of I is in relevance to the identification capability of the meta-sketch. If the item domain is known a-priori, it can be directly taken as the item pool. Otherwise, in applications where the item domain is only partially known or even unknown, the item pool can be constructed by sampling from the historical records. Even in the case that the item pool does not completely cover the item domain, the "missing" item can still be identified, due to the homogeneity of the domain-specific embedding space, given that the number of distinct items does not meet the item pool capacity |I|.

A **frequency mean range** L is the range for the frequency mean \overline{f} . One can get the value of \overline{f} by statistics of each sampled stream instance and extract the minimum and maximum \overline{f} s to build L.

A **distribution pool** P consists of many instances generated according to different parameters of relative frequency distributions. In this paper, we consider a family of *Zipf* distributions [33] with varied parameter α , as the base for constructing P. α can be selected from a wide range to have a good coverage of different distributions.

Notice that the meta-tasks are for the meta-sketch to learn the sketching ability, instead of spoonfeeding the meta-sketch to mechanically memorize the parameters of R. It means that the trained meta-sketch has the generalization ability to handle the case not covered in R (see Section 4.2).

The generation of a meta-task t_i can be done based on sampler R, as follows. We first randomly sample a subset of n_i items from I, and a frequency mean $f_i \in L$. Then, we sample a distribution instance $p_i \in P$ and make the n_i items' frequencies conform to p_i and \bar{f}_i . For example, the frequencies of n_i items can be set as $n_i \times \bar{f}_i \times p_i$, where $p_i \sim Zipf(\alpha)$ is a random variable. The above steps are repeated until the store set s_i and query set q_i are built.

203 3.3 Adaptive Meta-task Generation

While processing real data streams, we can get the item set I_r and its distribution p_r by online 204 sampling. I_r and p_r are then used for generating the set of adaptive meta-tasks. For each adaptive 205 meta-task, an item subset is sampled from I_r , and the relative frequency corresponding to each item 206 is sampled from p_r . The process is similar to the generation of basic meta-tasks. The only difference 207 from basic meta-task generation is that, there is no distribution pool anymore, because the real data 208 stream is unique. Also, we intentionally randomize the correspondence between an item and its real 209 relative frequency on the original data records. It is equivalent to constructing meta-tasks where 210 the item frequencies dynamically change. For example, the frequency of an item may first increase, 211 then suddenly drop [21]. With adaptive meta-tasks, the meta-sketch learns to quickly adapt to the 212 distribution p_r , while being flexible against the item frequency change. The detailed algorithms of 213 generating basic/adaptive meta-tasks are shown in supplement materials. 214

215 4 Experiments

216 4.1 Basic Setup

Dataset. We use two real datasets. *Word-query* is a streaming record of search queries, where each query contains multiple words (e.g., "News today") [15]. *IP-trace* consists of IP packets, where each packet is identified by a unique source/destination address pair (e.g., 192.168.1.1/12.13.41.4) [21]. We assume that query phrases and IP addresses are numerically encoded, similar to [15].

Table 1: Results of Basic Meta-sketch (T_r)

			Word-quer	ý	IP-trace					
Method	Metrics	n=5K,	n=10K,	n=20K,	n=40K,	n=5K,	n=10K,	n=20K,	n=40K,	
		B=9KB	B=11KB	B=13KB	B=15KB	B=9KB	B=11KB	B=13KB	B=15KB	
Dasia MS	ARE	12.3	14.74	10.98	13.79	3.00	1.51	2.97	1.13	
Dasic Mis	AAE	31.54	38.54	40.63	53.67	5.57	5.01	6.94	5.56	
CS	ARE	32.94	57.97	98.01	162.43	6.08	9.94	15.57	24.49	
C3	AAE	57.54	101.44	172.44	282.59	10.42	16.82	26.46	41.91	
CMS	ARE	21.34	48.33	111.82	239.11	8.12	16.07	32.77	65.19	
CIVIS	AAE	38.04	84.62	195.61	416.01	13.67	27.39	55.29	110.65	

Table 2: Results of Basic Meta-sketch (T_s)

										(0)				
	Method	Metrics	n=5K,B=9KB			n=10K,B=11KB			n=20K,B=13KB			n=40K,B=15KB		
Ì			0.5	1.1	1.5	0.5	1.1	1.5	0.5	1.1	1.5	0.5	1.1	1.5
Î	Basic MS	ARE	0.43	1.05	2.63	0.73	3.25	3.14	0.47	1.67	1.35	0.43	2.58	9.65
	(Word-query)	AAE	24.7	17.72	8.93	31.24	27.02	9.41	27.29	22.19	9.2	25.04	26.95	19.87
	Basic MS	ARE	0.59	2.27	9.38	0.73	0.86	1.02	0.72	1.73	7.52	0.73	0.79	2.33
	(IP-trace)	AAE	26.45	21.49	14.73	38.33	19.32	7.95	35.48	22.28	15.74	39.57	21.75	14.06
CS	CS	ARE	1.98	6.72	10.99	2.7	12.12	16.9	3.73	20.8	27.46	5.17	37.96	43.76
	0.5	AAE	74.96	47.98	15.89	102.05	75.83	23.8	140.65	118.29	38.7	194.32	198.4	59.96
CMS	CME	ARE	4.96	7.52	5.47	9.27	15.85	9.44	17.29	32.7	16.38	32.24	66.35	27.89
	AAE	187.52	53.81	8.17	350.08	99.82	13.58	651.63	185.54	22.88	1213.38	347.32	38.18	

Baseline. We hereby evaluate the basic and advanced meta-sketches. From now on, we use MS to represent the term meta-sketch for brevity. We compare basic MS (after the pre-training phase) with CM-sketch (CMS) and C-sketch (CS). We compare the advanced MS (after the adaptation phase) with learned augmented sketch (LS) and cold filter (CF), which are two variants of CM/C sketches with auxiliary structures. According to the default setting [10, 11], the number of hash functions for all sketches is 3. We adopt two commonly accepted metrics for evaluating the accuracies of stream frequency estimation, AAE and ARE³.

Parameters. We implement g_{emb} or g_{add} in MLP with 2-layers of sizes 128 and 48, followed by batch normalization, and g_{dec} in an MLP with 3-layers of 256 with residual connections. We use the *relu* function for layer connections. The space budget *B* is spent on storing *M*, the same as the setting in neural data structures [23]. Other modules, like hashing libraries, are commonly accepted as reusable and amortizable resources for multi-deployment of sketches [21, 23]. Note that due to space limitations, the details and methods of parameter settings of M(A), the ablation experiments and some parameter discussions are shown in the supporting material.

235 4.2 Basic Meta-sketch

Settings. For each dataset, we train the basic MSs under 4 item pools with $\{5K, 10K, 20K, 40K\}$ 236 different items, respectively. The meta-task sampler are with Zipf distributions. We build the 237 distribution pools set with $\alpha \in [0.8, 1.3]$ and set frequency mean range L = [50, 500]. For basic 238 meta-sketch training, the default maximum number of training steps ϕ is 5 million, the learning rate 239 is 0.0001, and the Adam optimizer is used. For evaluation, we consider two types of tasks, T_r and 240 241 T_s . T_r are directly obtained by random sampling on two real data streams with different values of n, i.e., the number of distinct items. Note that the frequency distributions of T_r are not necessarily obey 242 Zipf distributions. T_s are the synthetic tasks, where the item frequency follows the Zipf distribution 243 with $\alpha \in \{0.5, 1.1, 1.5\}$. To evaluate the generability and stability of basic MS, both $T_s(0.5)$ and 244 $T_s(1.5)$'s distributions are not covered by the distribution pool of the meta-task samplers. 245

Performance. Table 1 shows the performance of all competitors based on real dataset T_r . It shows that the basic MS outperforms traditional basic sketches, i.e., CMS and CS, on all testing cases. For example, the results on IP-trace show that, when n=40K and B=15KB, the ARE of basic MS is 1.13, while AREs of CMS and CS are 65.19 and 24.49, respectively. The advantage of meta-sketch is significant when testing on T_s with different α s, as shown in Table 2. Note that we use random choices to simulate the ideal hash functions for traditional sketches like [15], so that CS and CMS have the same result on test tasks with the same α in both datasets.

We show the trend of ARE w.r.t. the space budget, in Figure 2 (T_r , n=5K, Word-query). Compared to the dramatic performance degrading of traditional sketches, basic MS holds stable performance. We show that the trend of ARE w.r.t. the number of distinct items in Figure 3 (T_r , B=9KB, Word-query). Compared to traditional sketches, the ARE of basic MS increases sub-linearly w.r.t. the value of n. Note that AAE has similar results for the above experiments, see the supplement materials.

Generalization. We test the generality of basic MS to new items that are not in the item pool of the meta-task sampler in Figure 4(a). We make the experiments (n=5K, B=9KB, Word-query) by replacing some items in T_r with new items, and vary the fraction of new items to observe the trend of the performance. It shows that the ARE/AAE moderately increases w.r.t. the ratio



Figure 4: Generality of Meta-sketch

IP-trace Method Metrics Word-query n=40K n=5Kn=5Kn=40K n=10K. n=20K n=10K n=20K B=9KB B=9KB B=11KB B=13KB B=15KB B=11KB B=13KB B=15KB 1.38 2.29 Advanced ARE 3.05 2.83 4.06 5.20 0.87 0.89 MS AAE $\overline{21.42}$ 26.11 35.00 43.81 3 77 4.46 5.13 6.55 ARE 3.58 14.53 141.70 1127.11 0.85 2.74 4.20 16.71 CF90 AAE 21.13 59 18 381.63 2217.28 1.32 3.01 7.71 31.20 7.95 29.02 139.87 541.37 1.51 3.10 8.95 46.79 ARE CF70 AAE 29.02 76.58 295.63 970.94 2.57 5.51 16.83 82.84 ARE 91.16 138.64 244.24 407.83 12.62 33.50 103.76 155.61 CF40 AAE 174.86 252.22 421.85 693.47 24.16 60.79 175.14 279.72 20.52 48.69 266.50 8.34 17.09 35.22 77.79 ARE 111.85 LCMS(1%) AAE 37.80 81.93 194.15 451.28 13.72 28.39 59.10 129.86 ARE 25.53 40.84 67.21 104.54 5.20 7.80 11.33 17.12 LCS(1%) 44.53 122.57 180.56 8.78 18.97 28.38 AAE 78.17 13.10

Table 3: Results of Advanced Meta-sketch

of new items. The performance is acceptable considering the fact that the item domain is often 262 stable in practical applications. We then test the generality of meta-sketches to varied frequency 263 means that are not in range L of the meta-task sampler, as shown in Figure 4(b). The experiment 264 (n=5K, B=9KB, Word-query) is done by sampling a series of T_s tasks with frequency means in 265 $\{500, 5K, 50K, 500K, 5000K\}$. It shows that as the mean of the true frequencies increases, the 266 estimated frequencies of the meta-sketch increase linearly, so that the ARE keeps stable. 267

4.3 Advanced Meta-sketch 268

Settings. The generation of adaptive meta-tasks is similar to that of basic meta-tasks (Section 3.2), ex-269 cept that each item pool reads real frequency distributions for the adaption as described in Section 3.3. 270 In the adaption phase, the maximum number of training steps is $0.002 * \phi$. 271

Performance. Table 3 compares the performance of advanced MS with traditional sketches and their 272 variants, LS and CF, on real dataset T_r . We implement two LSs according to [15], learned CM-sketch 273 (LCMS) and learned C-sketch (LCS), following the default setting that (top 1%) high-frequency 274 items are separately stored. For CF, we follow the parameter setting in [14], and use CF40, CF70, and 275 CF90 for setting the filter percentages to 40%, 70%, and 90% of the total size, respectively. It shows 276 that the advanced MS achieves a better performance than LSs and CFs. Also, AAE/ARE of advanced 277 MS increases more moderately w.r.t. the number of distinct items n, compared to its competitors. 278

Furthermore, we compare the performance of the advanced MS and the LS under dynamic streaming 279 scenarios, as shown in Figure 5. We select a set of T_r (n=5K,B=9KB,Word-query), and gradually 280 shuffle the correspondence between items and frequencies. Here, the shuffle ratio is increased from 0 281 to 100%. It shows that the average ARE of advanced MS only slightly fluctuates between 3.26 and 282 4.0, and the average AAE is in the range of 21.28 and 21.68. In contrast, AAE of LCS or LCMS starts 283 above 37, and increase significantly w.r.t. the increase of the shuffle ratio. Actually, the classifier of 284 LS tends to incur more errors due to the gradual shift of high- and low-frequency items, resulting in 285 an increased number of hash collisions, thus deteriorating the estimation accuracy. 286

5 Analysis 287

The meta-sketch is trained based on meta-tasks, consisting of various stream distributions. We 288 expected that the meta-sketch can learn the ability to sketch item frequencies. Somehow, it is 289 unavoidable that the meta-sketch's ability is limited by patterns of given meta-tasks. Thus, setting 290 up the two training phases benefits the balance of the trade-offs. In the pre-training phase, we select 291 the most representative Zipf distribution to form basic meta-tasks, making the basic meta-sketch 292 adaptable to a wide range of data streams. In the adaptation phase, we sample adaptive meta-tasks 293 from raw data streams to make the advanced meta-sketch more specialized. Next, we analyze the 294



working mechanism of the three modules of the meta-sketch as well as their roles in acquiring thetwo abilities.

Sparse Addressing Module. We take a 2D slice A^* (size is $l_r \times d_2$) of the A matrix to analyze the process of a refined vector r getting addressing a through this module. First, we have $a \leftarrow$ $SparseMax(r^TA^*) \Rightarrow a \leftarrow SparseMax(\langle r \cdot b_1, r \cdot b_2, ..., r \cdot b_{d_2} \rangle)$. Since b_i are unit vectors, we can get $a \leftarrow SparseMax(|r|c), c = \langle cos\theta_1, cos\theta_2, ..., cos\theta_{d_2} \rangle$, where θ_i is the angle between r and b_i . We continue to transform the form to get addressing $a \leftarrow Sparsegen(c; u; \frac{|r|-1}{|r|})$ [30], where uis a component-wise transformation function applied on c. in this paper, we set u(c)=c.

Based on the principle of Sparsegen [30], |r| mainly affects the sparsity (i.e., the proportion of 303 non-zero bits in the vector) of a during training process, while c determines the positions and values 304 of non-sparse bits. The Figure 6 shows a strong correlation between the average |r| and the sparsity 305 of a during training from scratch (n=5K, B=9KB, Word-query, Basic MS). Since the embedding 306 vector z does not directly participate in the addressing process, the average |z| remains stable. Further, 307 we observe that the sparsity of a will eventually converge to around 1, which means that each item 308 is generally stored in a slot corresponding to the refined vector r and the unit vector in A^* with the 309 maximum cosine similarity. 310

Therefore, the role of A^* is to map refined vectors to the addressing vectors. The d_2 unit vectors in A^{*} are the reference standard for mapping, which is equivalent to the mutually exclusive d_2 -divisions of the refined vector space. Follow this point, we construct two matrices K^* and R^* of the same size as A^* . Among them, the d_2 unit vectors in K^* come from the cluster centers of the sampled refined vectors. To achieve mutually exclusive division, we perform Kmeans clustering with $K = d_2$ and *Cosine similarity* criterion. Then, we normalize the resulting d_2 cluster centers and stack them as K^* . In contrast, the unit vectors in R^* are entirely randomly generated.

Figure 7 (a) shows the results of replacing A^* on the trained meta-sketch with K^* and R^* . The 318 319 meta-sketch with R^* shows the worst performance, but the performance of the meta-sketch with K^* is close to the original A^* . Furthermore, We count the number of items mapped in every slot of A^* . 320 K^* , R^* and show their standard deviation in Figure 7 (b). The standard deviation of R^* is much 321 322 higher than A^* and K^* , and a better meta-sketch tends to store items more evenly in each slot. Thus, The addressing module simulates the traditional sketch mechanism. Its principal function is to store 323 the embedding vectors of items as evenly as possible in multiple memory slots, and an item is written 324 to only one slot. 325

Embedding Module. The major source of conflicts in the meta-sketch is the stacking of different embedding vectors in a single slot. Thus, the sparsity of the embedding vector becomes an important indicator to determine the degree of conflicts. Figure 8 shows the relation between the sparsity of embedding vectors and the stream distributions (n=5K, B=9KB, Word-query, advanced MS). We select the meta-tasks under *Zipf*, *Triangular*, and *Uniform* distributions with different skewness levels (the definition of skewness and corresponding distribution parameters are shown in the supplement materials). The results show that the sparsity of the embedding vector is positively proportional to



the skewness of a distribution. Therefore, we speculate that the meta-sketch memorizes the pattern information of the distribution being adapted by self-tuning the sparsity of embedding vectors.

Decoding Module. The decoding module, as the deepest NNs in the meta-sketch, integrates various information to predict the item frequency and achieves generalization ability. To verify this, we adapt the advanced MS (n=5K, B=9KB, Word-query) to a special adaptive meta-task. The meta-task was sampled from the real data stream but with a fixed item size (5000) and frequency mean (250). Meanwhile, we do not change the correspondence between items and frequencies. Such meta-task forces the meta-sketch to pay more attention to the fixed patterns and thus limit its generalization.

Thus, we train the advanced MS with (or without) freezing the decoding module parameters based 341 on the above meta-task. Figure 9 (a) shows the performance changes of the three models (advanced 342 MS as baseline) on the evaluation tasks (T_r) of different item sizes. Without the frozen decoding 343 module, the meta-sketch loses generalization ability at extended item sizes other than 5000. On the 344 contrary, the meta-sketch with the frozen decoding module still retains the generalization ability and 345 further utilizes the data stream pattern compared to the advanced MS, achieving the best performance. 346 Similarly, as shown in Figure 9 (b), the meta-sketch without the frozen decoding module also loses a 347 certain generalization ability in terms of frequency mean. 348

Actually, the above meta-task (termed as the *stable case*) can be viewed as a special case of an ordinary 349 adaptive meta-task (termed as the *unstable case*). As a matter of fact, augmented sketches utilize 350 frequency patterns similar to the stable case. For example, the learned augmented sketch memorizes 351 (relatively) stable correspondence between items and frequencies, for filtering high-frequency items. 352 353 To understand the meta-sketch's self-optimizing mechanism from the unstable case to the stable case, we analyze the storage of high/low-frequency items between multiple slots and a single slot in the 354 memory. In Figure 10 (a), we show density heat-maps of low-frequency (below the top 20% high 355 frequencies) items, stored by meta-sketches of stable and unstable cases on a 2D slice $(\bar{d}_1=2)$ of the 356 storage matrix M, where the x-axis is the index of slots. The two heat-maps show that the meta-sketch 357 under the stable case can store the low-frequency items concentratedly in some slots to avoid the 358 conflicts with high-frequency items. Interestingly, the meta-sketch does not intentionally do this like 359 augmented sketches. Instead, it is achieved by self-optimization during the training. Furthermore, 360 Figure 10 (b) shows the relation between the sparsity of the embedding vector of items stored in a 361 single slot and the frequency order, where the x-axis represents the frequencies in the ascending order. 362 We speculate that the meta-sketch autonomously adjusts the sparsity of the embedding vector within 363 a single slot in the stable case, so that the high/low-frequency items are automatically separated. 364

365 6 Conclusion

In this paper, we propose a neural data structure, called the meta-sketch, for estimating item fre-366 quencies in data streams. Unlike traditional sketches, the meta-sketch utilizes meta-learning and 367 memory-augmented neural networks. The meta-sketch is pre-trained with Zipf distributions and can 368 be fast adapted to specific runtime streams. We study a series of techniques for constructing the 369 meta-sketch. We also devise the generation of basic and adaptive meta-tasks corresponding to the 370 pre-training and adaption phases, respectively. Extensive empirical studies on real datasets are done 371 to evaluate our proposals. In the future, it is interesting to extend our proposal to other sketching 372 tasks that are supported by traditional sketches. 373

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478 479 480 481 482	 (d) Did you include the total amount of compute and the type of resources used (e.g., type of GPUs, internal cluster, or cloud provider)? [Yes] All of our experiments are implemented in python and run at a NVIDIA DGX workstation with CPU E5-2698 (2.20GHz, 20 cores), and 4 NVIDIA V100 GPUs (5120 CUDA cores and 16GB GPU memory on each GPU).
483	4. If you are using existing assets (e.g., code, data, models) or curating/releasing new assets
484	(a) If your work uses existing assets, did you cite the creators? [Yes]
485	(b) Did you mention the license of the assets? [N/A]
486	(c) Did you include any new assets either in the supplemental material or as a URL? [N/A]
487	
488 489	(d) Did you discuss whether and how consent was obtained from people whose data you're using/curating? [N/A]
490 491	(e) Did you discuss whether the data you are using/curating contains personally identifiable information or offensive content? [N/A]
492	5. If you used crowdsourcing or conducted research with human subjects
493	(a) Did you include the full text of instructions given to participants and screenshots, if
494	applicable? [N/A]
495	(b) Did you describe any potential participant risks, with links to Institutional Review
496	Board (IRB) approvals, if applicable? [N/A]
497 498	(c) Did you include the estimated hourly wage paid to participants and the total amount spent on participant compensation? [N/A]