000 001 002 UNIVERSAL LENGTH GENERALIZATION WITH TURING PROGRAMS

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ABSTRACT

Length generalization refers to the ability to extrapolate from short training sequences to long test sequences and is a challenge for current large language models. While prior work has proposed some architecture or data format changes to achieve length generalization, these proposals typically apply to a limited set of tasks. Building on prior scratchpad and Chain-of-Thought (CoT) techniques, we propose *Turing Programs*, a novel CoT strategy that decomposes an algorithmic task into steps mimicking the computation of a Turing Machine. This framework is both universal, as it can accommodate any algorithmic task, and simple, requiring only copying text from the context with small modifications. We show that by using Turing Programs, we obtain robust length generalization on a range of algorithmic tasks: addition, multiplication and in-context SGD. We then demonstrate that transformers achieve length generalization on random Turing Programs, suggesting that length generalization is possible for any algorithmic task. Finally, we theoretically prove that transformers can implement Turing Programs, constructing a simple RASP (Weiss et al. [Weiss et al.](#page-12-0) [\(2021\)](#page-12-0)) program that simulates an arbitrary Turing machine.

1 INTRODUCTION

031 032 033 034 035 036 037 038 039 040 041 042 Transformer-based language models have shown impressive abilities in natural language generation, reading comprehension, code-synthesis, instruction-following, commonsense reasoning, and many other tasks [Brown et al.](#page-10-0) [\(2020\)](#page-10-0); [Chen et al.](#page-10-1) [\(2021\)](#page-10-1); [Chowdhery et al.](#page-10-2) [\(2023\)](#page-10-2); [Lewkowycz et al.](#page-11-0) [\(2022\)](#page-11-0); [Gunasekar](#page-11-1) [et al.](#page-11-1) [\(2023\)](#page-11-1); [Touvron et al.](#page-12-1) [\(2023\)](#page-12-1). Despite these impressive abilities, transformers struggle with *length generalization*, which refers to the ability to generalize to longer sequences than seen during training [Abbe et al.](#page-9-0) [\(2023\)](#page-9-0); [Anil et al.](#page-10-3) [\(2022\)](#page-10-3); [Jelassi et al.](#page-11-2) [\(2023\)](#page-11-2); [Zhou et al.](#page-12-2) [\(2023\)](#page-12-2). This limitation raises a central question about transformers: are they capable of actually learning an algorithm or do they solve algorithmic tasks by resorting to memorization or shortcuts [Liu et al.](#page-11-3) [\(2022\)](#page-11-3)?

043 044 045 046 047 048 049 Recently, several works have reported poor length generalization of transformers on a wide range of algorithmic tasks [Anil](#page-10-3) [et al.](#page-10-3) [\(2022\)](#page-10-3); [Delétang et al.](#page-10-4) [\(2022\)](#page-10-4); [Dziri et al.](#page-10-5) [\(2024\)](#page-10-5); [Zhang](#page-12-3) [et al.](#page-12-3) [\(2022\)](#page-12-3). In parallel, a myriad of papers [Jelassi et al.](#page-11-2) [\(2023\)](#page-11-2); [Kazemnejad et al.](#page-11-4) [\(2024\)](#page-11-4); [Shen et al.](#page-12-4) [\(2023\)](#page-12-4); [Zhou et al.](#page-12-2) [\(2023;](#page-12-2) [2024\)](#page-12-5) have optimized the data formats choice (see [Section 3](#page-3-0) for details) to improve the length generalization of transform-

050 051 052 053 ers when trained to perform multi-digit addition of two numbers. While the recent progress is impressive[—Zhou et al.](#page-12-5) [\(2024\)](#page-12-5) achieve almost perfect accuracy on addition with 100-digit operands while trained on 40-digit, all these "tricks" are specific to the case of addition and may not generalize to other tasks. In contrast, our goal is to develop a technique that is *general* enough to enable length generalization on *any* algorithmic task.

054 055 056 057 058 059 060 061 062 To achieve this, we introduce *Turing Programs*, a novel scratchpad technique that may be applied to general algorithmic tasks. This technique is motivated by the operations of a Turing Machine, a mathematical model of computation that is capable of implementing any computable algorithm. A Turing machine consists of a "tape" with symbols and a "head" that, at each step, moves left or right on the tape, and can read and write symbols in a single tape cell. Therefore,

Table 1: Length generalization results on various problems with Turing Programs. We use $x \to y$ to denote training on $n = x$ and generalizing to $n = y$.

063 064 065 066 067 068 when a Turing Machine processes an input, the tape at each step is a copy of the previous one up to a few changes. Our Turing Programs follow this philosophy by decomposing an algorithmic task into a series of steps. At each step we update a "tape" by copying the previous tape with a few elementary changes. We refer the reader to [Figure 1](#page-0-0) for the correspondence between Turing Machines and Turing Programs and to Figures [2](#page-3-1) and [4,](#page-6-0) for examples of Turing Programs.

069 070 071 072 073 074 075 076 077 Using the Turing Programs technique, we show that transformers enhanced with the Hard-ALiBi positional encoding [Jelassi et al.](#page-11-5) [\(2024\)](#page-11-5)—a recent encoding that achieves state-of-the-art length generalization on copying—are capable of length generalization on a wide range of algorithmic tasks. Our method achieves non-trivial length generalization on addition, multiplication and simulation of SGD steps (see [Table 1\)](#page-1-0). Additionally, we show that transformers can be trained to execute random Turing machines, extrapolating from 50 to over 100 input tokens, suggesting that our method can work for general algorithmic tasks. To our knowledge, these are the first results showing nontrivial length generalization on multiplication, and the first attempt to study length generalization on complex algorithms like SGD. We hope that this recipe will be further used to unlock novel length generalization on other algorithmic tasks.

- **078 079** Our key contributions are summarized as follows:
- **080 081** – In [Section 3,](#page-3-0) we present length generalization results on multi-digit addition using a Turing Program and Hard-ALiBi positional encoding.
- **082 083 084** – In [Section 4,](#page-5-0) we present the Turing Program framework in full generality and its connections to Turing machines. Additionally, we theoretically prove that transformers can implement Turing Programs, constructing a RASP program [Weiss et al.](#page-12-0) [\(2021\)](#page-12-0) simulating Turing machines.
- **085 086 087** – In [Section 5,](#page-7-0) we demonstrate that Turing Programs are general and lead to novel length generalization results in unexplored algorithmic tasks: multiplication by 1 or 3-digit operand, SGD for linear regression and Turing Machine simulation.
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RELATED WORK

092 093 094 096 Length generalization remains an important challenge for large language models as underlined in several works [Delétang et al.](#page-10-4) [\(2022\)](#page-10-4); [Dziri et al.](#page-10-5) [\(2024\)](#page-10-5); [Hupkes et al.](#page-11-6) [\(2020\)](#page-11-6); [Schwarzschild et al.](#page-12-6) [\(2021\)](#page-12-6); [Zhang et al.](#page-12-3) [\(2022\)](#page-12-3). Despite their advanced reasoning capabilities, Transformer-based large language models struggle to process longer sequences than they were trained on [Anil et al.](#page-10-3) [\(2022\)](#page-10-3). The main approaches for improving length generalization focus on changing the positional encoding and optimizing the data format.

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098 099 100 101 102 103 104 105 106 107 Positional encodings for length generalization. [Shaw et al.](#page-12-7) Shaw et al. [\(2018\)](#page-12-7) were early to notice that the weak length generalization of Transformers was due to the choice of absolute positional encoding. Following this, many alternatives were proposed to replace the absolute positional encoding: relative positional encodings, which focus on the relative distances between tokens [Dai et al.](#page-10-6) [\(2019\)](#page-10-6); and weighted attention mechanisms in place of position embeddings [Chi et al.](#page-10-7) [\(2022\)](#page-10-7); [Jelassi et al.](#page-11-2) [\(2023\)](#page-11-2); [Li et al.](#page-11-7) [\(2023\)](#page-11-7); [Press et al.](#page-12-8) [\(2021\)](#page-12-8); [Raffel et al.](#page-12-9) [\(2020\)](#page-12-9). These alternatives showed substantial improvements in length generalization on natural language processing tasks. On the other hand, [Kazemnejad et al.](#page-11-4) [\(2024\)](#page-11-4) found that a causal language model with no positional encoding can length generalize better than some of these specialized positional encodings on algorithmic tasks. In this work, we apply the Hard-ALiBi positional encoding [Jelassi et al.](#page-11-5) [\(2024\)](#page-11-5), that achieved state-of-the-art length generalization on the specific task of copying, to more general algorithmic tasks.

108 109 110 111 112 113 114 115 116 117 118 119 Data formatting for length generalization. A wide range of data formatting methods have been introduced to achieve length extrapolation in algorithmic tasks. Scratchpad and Chain-of-Thought formats were proposed to learn arithmetic either through finetuning or in-context learning [Anil et al.](#page-10-3) [\(2022\)](#page-10-3); [Zhou et al.](#page-12-2) [\(2023\)](#page-12-2). When training from scratch, some other proposed techniques to improve length generalization on addition include: reversed formatting and random space augmentation [Shen](#page-12-4) [et al.](#page-12-4) [\(2023\)](#page-12-4), adding padding to the sequence [Jelassi et al.](#page-11-2) [\(2023\)](#page-11-2), and setting index hints in front of each digit [Zhou et al.](#page-12-2) [\(2023\)](#page-12-2). Closer to our work, several works [Anil et al.](#page-10-3) [\(2022\)](#page-10-3); [Dziri et al.](#page-10-5) [\(2024\)](#page-10-5); [Hu et al.](#page-11-8) [\(2024\)](#page-11-8); [Kazemnejad et al.](#page-11-4) [\(2024\)](#page-11-4); [Lanchantin et al.](#page-11-9) [\(2024\)](#page-11-9) report that training or finetuning a model on scratchpad data does not yield any significant length generalization improvement. In our work, we demonstrate that length generalization is possible using a combination of a particular scratchpad variant and a favorable positional encoding. Additionally, we develop Turing Programs, a novel scratchpad strategy that is general and may be applied to achieve length generalization on any algorithmic task.

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121 122 123 124 125 126 127 128 129 130 131 132 133 Neural networks and Turing Machines. Many prior works designed architectures inspired by Turing Machines [Dehghani et al.](#page-10-8) [\(2018\)](#page-10-8); [Graves et al.](#page-10-9) [\(2014\)](#page-10-9); [Kaiser & Sutskever](#page-11-10) [\(2015\)](#page-11-10). From a theoretical perspective, some works proved the Turing completeness of RNNs [Chen et al.](#page-10-10) [\(2017\)](#page-10-10); [Siegelmann & Sontag](#page-12-10) [\(1992\)](#page-12-10), transformers [Bhattamishra et al.](#page-10-11) [\(2020\)](#page-10-11); [Chung & Siegelmann](#page-10-12) [\(2021\)](#page-10-12); [Pérez et al.](#page-12-11) [\(2019\)](#page-12-11); [Wei et al.](#page-12-12) [\(2022a\)](#page-12-12); [Merrill & Sabharwal](#page-11-11) [\(2023\)](#page-11-11) and even linear next-token predictors [Malach](#page-11-12) [\(2023\)](#page-11-12) under a wide range of assumptions. Lastly, another line of work characterizes the computational model that Transformers express: [Weiss et al.](#page-12-0) [\(2021\)](#page-12-0) introduce RASP, a human-readable programming language that can be implemented by transformers, [Lindner et al.](#page-11-13) [\(2024\)](#page-11-13) show how human-readable programs are compiled into transformer models and other works [Giannou et al.](#page-10-13) [\(2023\)](#page-10-13); [Jojic et al.](#page-11-14) [\(2023\)](#page-11-14) study how transformers can emulate computer programs. Closer to our work, [Zhou et al.](#page-12-5) [\(2024\)](#page-12-5) hypothesize that Transformers can length generalization on any algorithmic task that may written as a "simple" RASP program. In this work, we construct a simple RASP program that generates Turing Programs to simulate arbitrary Turing machines.

2 SETTING

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137 138 139 140 In this section, we present the length generalization problem and some instances where it appears. Then, we discuss scratchpad prompting [Nye et al.](#page-11-15) [\(2021\)](#page-11-15), a technique that lets the model generate solution steps before producing the final answer. Finally, we introduce various positional encoding methods and discuss their implications on length generalization.

2.1 LENGTH GENERALIZATION

143 144 145 146 147 Many sequence modeling tasks have problem instances of different lengths. Shorter instances are often easier to state, process and handle, and require less compute to find the answer. By contrast, longer instances are more challenging to parse and require more compute to solve. Reasoning tasks such as multi-hop reasoning, program execution, deductive reasoning, and theorem proving fit in this category.

148 149 150 151 152 Algorithmic reasoning tasks consist of inputs that are sequences of tokens describing the task (e.g. addition, multiplication) and outputs that are the corresponding solutions. We assume that the language model is allowed to generate (many) intermediate tokens before outputting the answer. Then formally, the *length generalization* problem consists of training a language model on inputs of length $\leq L$ and solving problems of length $>L$ at test time.

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- **154** 2.2 SCRATCHPAD

155 156 157 158 159 160 It has been shown in prior work that the performance of LLMs on algorithmic tasks can be greatly improved by generating step-by-step solutions instead of immediately outputting the final answer [Wei](#page-12-13) [et al.](#page-12-13) [\(2022b\)](#page-12-13). Among the multiple methods described in the literature, we focus on the scratchpad method [Nye et al.](#page-11-15) [\(2021\)](#page-11-15). Given an algorithmic task, this method encodes the intermediate steps of the algorithm as text and trains the model to emit them to a buffer that is referred to as the "scratchpad".

161 Nye et al. [Nye et al.](#page-11-15) [\(2021\)](#page-11-15) showed that scratchpad finetuning can be used to achieve strong indistribution performance on execution based tasks such as code execution and computing polynomials. **162 163 164** They also report modest length generalization results on integer arithmetic. The limitation of scratchpad training for length generalization is further highlighted in [Anil et al.](#page-10-3) [\(2022\)](#page-10-3); [Dziri et al.](#page-10-5) [\(2024\)](#page-10-5); [Hu et al.](#page-11-8) [\(2024\)](#page-11-8); [Kazemnejad et al.](#page-11-4) [\(2024\)](#page-11-4); [Lanchantin et al.](#page-11-9) [\(2024\)](#page-11-9).

165 166 167 168 169 170 171 172 In this paper, we revisit the use of scratchpad training to achieve length generalization on algorithmic tasks. We begin with the observation that the scratchpad technique can be realized as an iterative sequence of copying operations, where at each iteration the input is slightly modified. Building on previous works showing that with the right positional encoding, transformers can achieve length generalization on the copying operation [Jelassi et al.](#page-11-5) [\(2024\)](#page-11-5), we hypothesize that combining the scratchpad technique with a favorable positional encoding can unlock length generalization capa-bilities. We verify this hypothesis in [Section 3](#page-3-0) and [Section 5,](#page-7-0) but first we review various choices of positional encoding.

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2.3 POSITIONAL ENCODINGS

176 177 178 179 180 181 The inability of transformers to extrapolate to longer sequences has been primarily attributed to the positional encoding [Shaw et al.](#page-12-7) [\(2018\)](#page-12-7); [Shen et al.](#page-12-4) [\(2023\)](#page-12-4). In this section, we review the different positional encoding schemes and in [Section 3,](#page-3-0) we report their length generalization performance. We review here specific choices for positional encodings that are known to perform well for length generalization, and discuss additional encoding schemes (such as absolute and relative positional encodings) in Appendix [A.](#page-13-0)

No Positional Encoding (NoPE). Decoder-only models with causal attention, as shown by [Haviv](#page-11-16) [et al.](#page-11-16) [\(2022\)](#page-11-16), acquire positional understanding, without explicit positional encoding. [Kazemnejad et al.](#page-11-4) [\(2024\)](#page-11-4) shows that a model without positional encoding extrapolate better than those with specialized positional encodings on some algorithmic tasks.

187 188 189 190 191 ALiBi. [Press et al.](#page-12-8) [\(2021\)](#page-12-8) introduces this additive positional encoding where the bias function follows $b(i, j) = -r|i - j|$, where $r > 0$ is some hyperparameter. This scheme has led to state-ofthe-art length generalization on natural language tasks. However, [Jelassi et al.](#page-11-5) [\(2024\)](#page-11-5) notices that it struggles at length generalization on the copy task and hypothesize that it is due to the slow decay of r.

193 194 195 196 197 198 199 200 Hard-ALiBi. [Jelassi et al.](#page-11-5) [\(2024\)](#page-11-5) introduce Hard-ALiBi, an additive positional encoding where the bias satisfies $b(i, j) = -\infty$ for $j \leq i - m$ and $b(i, j) = 0$ for $j > i - m$, for some hyperparameter $m > 0$. Intuitively, with this hard thresholding, tokens can only attend to the m closest tokens. Different heads may have different values of m and some heads use $m = \infty$ which corresponds to softmax attention with no positional embedding at all (allowing for propagation of global information). The authors demonstrate empirically that models equipped with the Hard-ALiBi positional encoding achieve remarkable length generalization on the copy task. In this work, we use the Hard-ALiBi position encoding to enable length generalization on algorithmic tasks as we show below.

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3 LENGTH GENERALIZATION ON ADDITION

203 204 205 206 207 208 209 210 211 212 In this section, we address the length generalization problem for addition. We first review prior results on this problem and describe the techniques used in these works. We then demonstrate that Transformers trained with Turing Program scratchpads and Hard-ALiBi positional encoding achieve good length generalization performance, extrapolating from length-50 to length-100 addition. This is a remarkable improvement over previous length generalization results using the "vanilla" scratchpad technique (e.g. [Nye et al.](#page-11-15)

213 214 215 Figure 2: Turing Program for addition, text in performance. As mentioned, there is a long list of works understition for the semperalization on [\(2021\)](#page-11-15)), which showed weak length generalization addition (see [Appendix B](#page-13-1) for a complete review). Notably, [Zhou et al.](#page-12-5) [\(2024\)](#page-12-5) report somewhat better length generalization results compared to our results. However, we note that these results rely on

226 227 228 229 230 Figure 3: (a): Comparison of different positional encodings and data formats for length generalization on addition. We see significant extrapolation to longer sequence lengths with Hard-ALiBi and scratchpad. The shade shows the 95% confidence intervals. (b): Hard-ALiBi with Turing Program, trained with 5 different initialization seeds. To clarify, the randomness used to plot the 95% confidence intervals in Figure [3a](#page-4-0) comes from the samples we draw to calculate the accuracy once a seed is fixed, not from different training seeds.

particular choices for the formatting of the input and the output, which are "tailored" for the task of multi-digit addition.

3.1 LENGTH GENERALIZATION ON ADDITION WITH TURING PROGRAMS AND HARD-ALIBI

In this section, we present our Turing Program scratchpad strategy for addition and report length generalization results.

3.1.1 EXPERIMENTAL SETUP

242 243 244 245 246 247 248 249 250 251 252 253 254 Data. We adopt the scratchpad format and write all the steps into one sequence, where steps are separated by a separator token. [Figure 2](#page-3-1) shows our scratchpad strategy for getting length generalization on addition^{[1](#page-4-1)}. If not specified otherwise, our token space is of size 24 and made of $V = \{0, \ldots, 9, +, a, \ldots, j, \wedge, \langle |\text{BOS}| \rangle, \langle |\text{EOS}| \rangle, \langle |\text{SEP}| \rangle\}$. All the digits are sampled uniformly as follows: we first sample the length of each operand (between 2 and $L = 50$) and then independently sample each digit. Finally, we "pack the context" with i.i.d. sequences during training, i.e. we fill the context with multiple independent samples of the task (similarly to [Zhou et al.](#page-12-2) [\(2023\)](#page-12-2)). We set the training context length to 500. At test time, we evaluate our models using a sliding window: we generate tokens until the training context length (500) is filled, and then each additional token is generated by feeding the context of the most recent 500 tokens, effectively dropping all past tokens^{[2](#page-4-2)}. This way, we are able to generate very long sequences of tokens without training or evaluating on long context windows. To evaluate the accuracy at a given length, we test the model's output on 288 examples. We report the accuracy of exactly matching the desired output.

Model and Training. Our base model is a 150M parameter Transformer with $L = 12$ layers, a $D = 1024$ hidden size, feedforward layer with a hidden dimension of 4096 and $H = 16$ attention heads. The backbone of our model is based on the GPT-NeoX architecture [Black et al.](#page-10-14) [\(2022\)](#page-10-14). We pick a context length of 500 tokens. We use the AdamW optimizer [Loshchilov & Hutter](#page-11-17) [\(2017\)](#page-11-17) to train the model with a weight decay value of 0.1 and no dropout, for 200,000 steps. The learning rate schedule incorporates an initial 100-step linear warm-up, followed by a linear decay, starting at $7e-5.$

Hard-ALiBi positional encoding. Similarly to [Jelassi et al.](#page-11-5) [\(2024\)](#page-11-5), we use M masked heads and $(H - M)$ NoPE heads. In the masked heads, we respectively set the hyperparameter m to 1, 2,... and M. We swept over $\{3, 4, 5, 6, 7, 8\}$ and found that $M = 6$ is the best choice.

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¹In the experiments, we use a slightly more compact version of the scratchpad, where each examples is represented as $$4324+139|432e+13j(1,3)|43c+1d(0,63)|4d+b(0,463)|e+^{\hat{(0,4463)}|4463}.$

² For efficiency reasons, once we reach the context length we advance the "window" by 20 tokens.

270 271 3.1.2 RESULTS

272 273 274 275 276 277 278 279 280 In Figure [3a](#page-4-0) we show the length generalization performance of transformers trained to perform multidigit addition using the scratchpad described above. We compare the performance of different choices of positional encodings, as well as comparing to the performance on addition without scratchpad (directly outputting the answer). We observe that by using Hard-ALiBi together with scratchpad, transformers are able to generalize well beyond the length of the training examples. In particular, the Hard-ALiBi model achieves a 98% accuracy at length 100. As shown by Figure [9](#page-15-0) in the appendix, the model also length generalizes well when the operands are of different lengths. Finally, in Figure [3b](#page-4-0) we analyze the robustness of length generalization performance to different choices of initialization seed. We observe that, while there is significant variance in performance when testing on longer sequences, our method is more robust compared to prior results (as reported in [Zhou et al.](#page-12-5) [\(2024\)](#page-12-5)).

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4 TURING PROGRAMS

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286 287 288 289 290 In [Section 3,](#page-3-0) we showed that Transformers with Hard-ALiBi trained on a specific choice of scratchpad format can length generalize to sequences that are $2\times$ longer. On closer inspection, each line in the scratchpad in [Figure 2](#page-3-1) is a slightly modified copy of the previous one where a few elementary changes are applied, e.g. removing one digit for each operand and updating the intermediate result/carry. Since Hard-ALiBi yields robust length generalization on copying, this may explain why we achieve better extrapolation than previous works that trained their models with scratchpad.

291 292 293 294 295 296 297 298 In this section, we generalize this approach and claim that every algorithmic task can be written as a sequence of *modified copy* operations: i.e. copy operations with small and localized modifications. Such decomposition follows immediately from the standard construction of a Turing Machine, a universal model of computation. We therefore refer to this scratchpad strategy as a *Turing Program*. We start this section introducing the standard definition of a Turing Machine, and then present Turing Programs, our scratchpad strategy for achieving length generalization on any algorithmic task. Lastly, we present our main theoretical result: Transformers can implement Turing Programs over long sequences of inputs.

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4.1 BACKGROUND: TURING MACHINES

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306 307 308 309 Definition 4.1 *A Turing Machine is specified as a quadruple* $T = (Q, \Sigma, s, \delta)$ *where: 1) Q is a finite set of states, 2)* Σ *is a finite set of symbols, 3)* $s \in Q$ *is the initial state and* $f \in Q$ *is the final state, 4)* δ is a transition function determining the next move: δ : $(Q \times \Sigma) \rightarrow (\Sigma \times \{L, R\} \times Q)$.

310 311 312 313 *At the first iteration, the machine is set to state* $s \in Q$ *, the head is on the first (leftmost) cell of the tape, and the input is written on the tape from left to right. At each iteration, the head is on the i-th cell in the tape, is in state* $q \in Q$ *and reads the i-th symbol on the tape* α . *Then, if* $\delta(q, \alpha) = (\alpha', D, q')$ *, the head writes the symbol* α' *, moves in the direction* $D \in \{L, R\}$ *, and the machine changes its state to* q ′ *. If the machine reaches the state* f*, it stops, and its "output" is written on the tape.*

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316 317 318 319 320 Turing Machines are a powerful model for solving algorithmic tasks since (a) the framework is *universal* i.e. it is possible to write any algorithmic task in the Turing Machine formalism, (b) Turing Machines can solve a wide range of algorithmic problems—ranging from simple arithmetic to determining whether a number is a prime [Agrawal et al.](#page-10-15) [\(2004\)](#page-10-15)—in a polynomial number of steps. In the next section, we show how to use the Turing Machine formalism to obtain a novel scratchpad strategy that unlocks length generalization on any algorithmic task.

³²² 323 ³We assume that the tape is unbounded from the right side, but bounded from the left. Namely, there are infinitely many cells to the right of the input, but no empty cells to the left. This is computationally equivalent to a tape that is infinite from both sides.

324 325 326 4.2 TURING PROGRAMS: A UNIVERSAL SCRATCHPAD STRATEGY FOR LENGTH GENERALIZATION

327 328 329 330 The left panel of [Figure 1](#page-0-0) represents the simulation of a Turing Machine and shows how the state, the head and the tape evolves with time. Note that at each time step, the state of the tape is a copy of the previous tape with a few elementary changes such as a move of the head, an edit of a single symbol and a change of state.

331 332 333 334 335 The steps in a Turing Machine simulation are similar to a scratchpad strategy where each string is a copy of the previous one with a few modifications. Therefore, we claim that for any algorithmic task that can be solved by a Turing-computable algorithm, there is a corresponding scratchpad strategy for solving this problem (as demonstrated in the right panel of [Figure 1\)](#page-0-0). We refer to this novel scratchpad strategy as *Turing Programs*.

336 337 338 339 340 341 Turing Programs decompose an algorithmic task into a series of intermediate reasoning steps. Each step is a "tape" that maintains the state of the machine, and the next step is a copy of the previous tape with a few elementary changes, such as trimming of digits and update of carry/intermediate result as in the case of addition and multiplication (see Figures [2](#page-3-1) and [4\)](#page-6-0) or update of the parameters in the case of SGD on linear regression (see [Subsection 5.2\)](#page-8-0). In [Section 5,](#page-7-0) we show how to use Turing Programs to unlock novel length generalization results on challenging algorithmic tasks.

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4.3 THEORY: TURING PROGRAMS IN RASP

345 346 347 348 349 350 351 352 353 To further motivate the use of Turing Programs to achieve length generalization on arbitrary algorithms, we prove that transformers can implement Turing Programs over long sequences of inputs. In particular, we show that Turing Programs can be implemented in RASP [Weiss et al.](#page-12-0) [\(2021\)](#page-12-0), a programming language that was suggested as an abstract description of the operations of a transformer. Following [Zhou et al.](#page-12-2) [\(2023\)](#page-12-2), we use a restricted version of RASP that does not allow direct index operations, as [Zhou et al.](#page-12-2) [\(2023\)](#page-12-2) hypothesized that RASP programs with index arithmetics may not length generalize^{[4](#page-6-1)}. Therefore, our result should be viewed as a length-generalization-friendly construction of a transformer that can execute (most) Turing Programs (and hence, can simulate most Turing machines).

354 355 356 357 358 359 360 361 362 363 364 365 366 367 368 369 370 To avoid index operations, we leverage the n -gram hashing mechanism suggested by [Jelassi et al.](#page-11-2) [\(2023\)](#page-11-2) as a basis for the copying ability of transformers. In their construction, copying a string from the input was achieved by storing a sequence of n preceding tokens (n-gram) at each position, and iteratively retrieving the next token after the uniquely matched n -gram. Our Turing Program construction is very similar, except that instead of copying a string from the input, we copy the next state of the Turing machine as computed from the previous string. As in the construction of [Jelassi et al.](#page-11-2) [\(2023\)](#page-11-2), our RASP program is limited to operating on inputs that have no repeated n -grams (i.e., no sequence of n tokens appears twice in the input), which can be guaranteed with high probability for uniformly random sequences of tokens of length \leq exp(n). Additionally, we require that the Turing machine does not generate repeated n -grams when

Figure 4: Turing Program for 3-digit multiplication. At each step, we update three information: the head position, the result of the "local" multiplication, the carry and the intermediate result of the "global" multiplication.

371 372 processing the input, and that all the operations of the Turing machine are applied in-memory^{[5](#page-6-2)}. Under these assumptions, we get the following result:

⁴Our RASP program does not follow all the restrictions of the RASP-L language suggested in [Zhou et al.](#page-12-2) [\(2023\)](#page-12-2), as we do not restrict the tokens to have int8 values.

³⁷⁵ 376 377 ⁵Namely, we assume that the head of the Turing machine does not go beyond the input sequence. We believe that this restriction may be removed at the cost of constructing a more complex RASP program. While this may seem like a limiting restriction, we note that this limitation can be easily mitigated by padding the input with random tokens.

Theorem 4.1 *Let* T *be a Turing Machine s.t. 1)* T *does not generate repeated* n*-grams and 2)* T *operates in-memory. Then, there exists a RASP program* P *of size (number of lines)* O(n) *s.t. for every input* x *without repeated n-grams, P correctly simulates* T *for* $\exp(n)$ *steps.*

We give the full code for the construction of such RASP programs in Appendix [D.](#page-14-0)

5 LENGTH GENERALIZATION ON OTHER ALGORITHMIC TASKS

Building upon the encouraging length generalization results on addition from [Section 3](#page-3-0) and the Turing Programs framework from [Section 4,](#page-5-0) we show that Transformers enhanced with Hard-ALiBi may achieve robust length generalization on complex algorithmic tasks. We show that our framework achieves length generalization on multiplication by 1-digit and 3-digit operands, on SGD applied to linear regression, and finally, on next-state prediction of a random Turing Machine.

5.1 MULTIPLICATION BY A FIXED-LENGTH OPERAND

Prior work. Multiplication is known to be a challenging task for length generalization and very few works report positive length generalization results on this task. On pretrained models, [Zhou](#page-12-2) [et al.](#page-12-2) [\(2023\)](#page-12-2) shows that elaborate prompting techniques slightly improve the length generalization of Codex on $(n \le 3)$ -multiplication. [Dziri et al.](#page-10-5) [\(2024\)](#page-10-5) show that even fine-tuned GPT-3 struggles with performing 3-digit multiplication. On randomly intialized networks, [Lee et al.](#page-11-18) [\(2023\)](#page-11-18) show that models can learn in-distribution the 2-digit multiplication in a sample efficient way using scratchpad. [Shen et al.](#page-12-4) [\(2023\)](#page-12-4) shows that with padding and reversed products it is possible to perfectly learn in-distribution 12-digit multiplication. [Jelassi et al.](#page-11-2) [\(2023\)](#page-11-2) focuses on 3-digit multiplication and shows that when training on (5×3) -digit-multiplication and adding a few examples of (35×3) -digitmultiplication, the model length generalizes to (35×3) -digit-multiplication. In summary, prior work mainly focused on in-distribution learning of multiplication and did not manage to obtain length generalization results.

Figure 5: (a): Comparison of different positional encodings and data formats for length generalization on $(n \times 1)$ -digit-multiplication using the same hyperparameters. The shade shows the 95% confidence intervals. (b): Comparison of different positional encodings and data formats for length generalization on $(n \times 3)$ -digitmultiplication. We see that directly outputting the answer has zero accuracy at length 40 already.

423 424 425 426 427 428 Data setup. Our experimental setup is similar to the one in [Section 3.](#page-3-0) We focus on multiplication by a fixed-length operand, i.e. $(n \times k)$ -digit-multiplication where the first operand has variable length n and the second operand always has a fixed length $k \in \{1,3\}$ across all examples. We adopt the scratchpad format and write all the steps into one sequence, where steps are separated by a separator token. The Turing Program for multiplication is described in [Figure 4.](#page-6-0)^{[6](#page-7-1)} Our token space is similar to the token space used in Section [3,](#page-3-0) using a $*$ symbol instead of $+$ and using an additional separator

⁴²⁹ 430 431 ⁶In the experiments, we use a slightly more compact version of the scratchpad, where each examples is represented as \$4324*135|432e*135(0540∼054,0)|43c*135(0270∼032,40)|4d*135(0405 ∼043,740)|e*135(0540∼058,3740)|[∧] *135(0000∼005,83740)|[∧] *135(0000∼000,5837 40)|583740.

432 433 434 435 token ∼. All the digits are sampled uniformly as follows: we first sample the length of the first operand (between 2 and 50) and then independently sample each digit. The remaining details of the training/test protocols are similar to those in [Section 3.](#page-3-0)

436 437 438 439 440 441 442 443 Results. Figure [5](#page-7-2) reports our length generalization results on $(n \times 1)$ and $(n \times 3)$ multiplications. We obtain robust length generalization by a factor \times 2 (from 50 to 100-digit numbers) on $(n \times 1)$ and $(n \times 3)$ multiplication. We note that, up to length 100, $(n \times 1)$ and $(n \times 3)$ multiplication perform roughly the same $((n \times 1)$ has accuracy 97.1% and $(n \times 3)$ has accuracy 96.8%), which demonstrates the generality of our Turing Programs framework. Both results are achieved with $M = 7$ masked heads and peak learning rate 0.0003. The head numbers were again chosen by sweeping over candidate numbers as before while the learning rates were chosen from the candidate set $\{7e-5, e-4, 3e-4\}.$

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5.2 SGD ON LINEAR REGRESSION

446 447 448 449 450 In this section, we train a model to perform SGD and demonstrate its ability to length generalize. While in previous examples we varied the number of digits in the operands, here we instead vary the number of examples.

452 453 454 455 456 457 458 459 Problem Description. Let $D =$ $\{(\vec{x}_i, y_i)\}_{i=0,\dots,n-1}$ with $\vec{x}_i \in \mathbb{R}$ 2 and $y_i \in \mathbb{R}$ be a dataset of size *n*. Given initial weights $\vec{w}_0 \in \mathbb{R}^n$, we can obtain the final weight \vec{w}_{n-1} by performing gradient descent: $\vec{w}_{i+1} = \vec{w}_i - \lambda \nabla_{w_i} (y_i - \vec{w}_i \cdot \vec{x}_i)^2$, where λ is the learning rate. For our experiment, we pick $\lambda = 0.5$ and $\vec{w}_0 = 0$.

Figure 6: Length generalization on running the SGD algorithm, varying the number of examples.

460 461 462 463 464 Tokenization and Data. We divide the interval $[-1, 1]$ into 200 discrete tokens $\{a_0, a_1, ..., a_{199}\}$. As an input, the model receives a sequence of n examples, each encoded as two input coordinate and one output (label) value. The model then needs to compute the iterates of the SGD algorithm when processing the data examples, starting from the last data point, and output the resulting weight vector \vec{w}_{n-1} . A detailed description of the Turing Program for solving SGD is detailed in [Appendix E.](#page-19-0)

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466 467 468 469 470 471 472 473 474 Results. Unlike previous experiments, where we report accuracy w.r.t. exact string matching, here we allow the network to err by two quantization unit, counting any output that is within $2/100$ from the ground-truth output (in ℓ_{∞} norm) as correct. In other words, we disregard errors that may occur to differences in quantization of the real-valued iterates of SGD. As shown by the blue curve in Figure [6,](#page-8-1) training the transformer to perform SGD on dataset of sizes $n \leq 50$ generalizes with accuracy $> 95\%$ to datasets of size $n = 80$. Our Hard-ALiBi model has $M = 7$ masked heads, a context length of 600, and was trained with peak learning rate $7e-5$ for 400, 000 steps with a batchsize of 16. For comparison, we also trained a model to directly compute the final answer as shown by the red curve in Figure 6 . We observe that training the model to immediately output the answer significantly degrades its performance.

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5.3 TURING SIMULATIONS

478 479 480 In this section, we test the length generalization of transformers trained to predict the next state of an arbitrary, randomly generated, Turing machine. Our experimental setup is similar to the one in [Section 3](#page-3-0) except for the data as detailed below.

482 483 484 485 Data setup. We first sample a random Turing Machine T with 5 states, 15 input symbols and a random transition function (i.e., for every pair of state and symbol we randomly draw a triplet of state, symbol and move-direction). During training, each input example is generated as follows: we randomly choose an input sequence length L between 2 and 50, then randomly choose L tokens, a random position for the head and a random state for the machine. At each step of training, we

486 487 488 489 generate in an online manner a batch of size 16 of Turing simulations from T and focus on learning 1-step prediction: given the input tape, the model has to generate the output of the transition function followed by the next state of the tape. At test time, we evaluate the model on tapes of length $L \geq 50$. Further details are in [Appendix F.](#page-20-0)

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491 492 493 494 495 496 497 498 499 500 501 502 503 504 505 506 Results. [Figure 7](#page-9-1) shows that transformers enhanced with Hard-ALiBi predict almost perfectly the 1-step Turing Machine transition of tapes that are $2 \times$ to $3 \times$ longer than those seen during training. Trained with a peak learning rate of 7e-5, the models have $M = 8$ masked heads and a context length of 450. This experiment suggests that transformers may length generalize on *arbitrary* Turing Programs^{[7](#page-9-2)}. However, this admittedly does not imply that transformers can successfully execute Turing Programs for multiple steps, as accumulating errors might cause the programs to fail. That said, we note that in many cases we get length generalization with virtually zero error, suggesting that multiple steps of the machine can be execute while still maintaining accurate performance. The per-

Figure 7: Length generalization performance on 10 different randomly generated Turing machines.

507 508 509 formance of different positional encodings and data formats for Turing simulation can be found in Appendix [C.](#page-14-1) We observed that both directly outputting the answer and using alternative positional encodings significantly degraded the performance of length generalization.

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6 DISCUSSION AND LIMITATIONS

513 514 515 516 517 518 519 520 521 522 523 524 Studying and improving the length generalization abilities of transformers on algorithmic tasks has been the focus of various recent works. In parallel, it has been established experimentally that the ability of language models to solve algorithmic tasks is greatly enhanced when allowing them to use scratchpad/CoT data. Additionally, recent theoretical works demonstrate that transformers can use CoT to simulate arbitrary algorithms Merrill $\&$ Sabharwal [\(2023\)](#page-11-11), establishing that they are computationally "universal". These results motivate us to study whether transformers are universal *learners*, able to learn from examples to execute arbitrary algorithms. Since algorithms are typically defined over arbitrary sequence lengths, we use length generalization as a measure of whether the model has learned the *true* algorithm. To establish this, we use the key observation that transformers can length generalize on the copying operation. Since executing an algorithm can be implemented as a sequence of "smart" copy operations, the copying ability of transformers can be leveraged to achieve non-trivial length generalization performance on a wide range of algorithmic tasks.

525 526 527 528 529 530 That said, we acknowledge that our work still falls short of demonstrating that transformers can robustly length generalize on *any* algorithmic task. In some of our results, the extrapolation to longer sequence length is not robust, and degradation in performance may appear shortly after moving out-of-distribution. Additionally, our results rely on potentially very long and cumbersome CoT data, in a way that is not necessarily useful for real-world applications of language models. Thus, we view our results as theoretical evidence that length generalization is possible, and leave the development of more practical and robust methods for real-world length generalization to future work.

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⁵³⁸ 539 ⁷We note that, formally, the experiment demonstrates the ability of transformers to learn in the "average case", but does not rule out the possibility that some "worst case" Turing Programs have much more restricted length generlization.

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702 703 A ADDITIONAL POSITIONAL ENCODINGS REVIEW

704 705 706 707 708 709 Absolute Positional Encoding (APE). APE consists in maintaining a positional vector p_i for each position i . This vector is either predefined via a sinusoidal function [Vaswani et al.](#page-12-15) [\(2017\)](#page-12-15) or learned [Devlin et al.](#page-10-16) [\(2018\)](#page-10-16). Then, p_i is added to the token embedding e_i before being processed by the Transformer. Prior work observed that this positional encoding does not generalize well to longer sequences in both natural language [Press et al.](#page-12-8) [\(2021\)](#page-12-8) and algorithmic tasks [Jelassi et al.](#page-11-2) [\(2023\)](#page-11-2); [Kazemnejad et al.](#page-11-4) [\(2024\)](#page-11-4).

Additive Relative Positional Encoding (RPE). [Shaw et al.](#page-12-7) [\(2018\)](#page-12-7) were the first to integrate positional encodings at the level of each attention layer (instead of doing it at the input level). [Raffel](#page-12-9) [et al.](#page-12-9) [\(2020\)](#page-12-9) built upon this approach and added scalar biases to pre-softmax logits as follows:

$$
A = XW_Q(XW_K)^\top + B,\tag{1}
$$

715 716 717 718 719 720 721 where X, W_Q, W_K are the input and query and key weight matrices. The bias matrix $B \in \mathbb{R}^{n \times n}$ is induced by some positional encoding function $b \colon \mathbb{N}^{*2} \to \mathbb{R}$. For instance, the T5 relative positional encoding [Raffel et al.](#page-12-9) [\(2020\)](#page-12-9) set $b(i, j) = f(i - j)$, where f is some function. Most of the subsequent positional encodings such as ALiBi [Press et al.](#page-12-8) [\(2021\)](#page-12-8), Kerple [Chi et al.](#page-10-7) [\(2022\)](#page-10-7), Randomized Positional Encoding [Ruoss et al.](#page-12-16) [\(2023\)](#page-12-16) and Fire [Li et al.](#page-11-7) [\(2023\)](#page-11-7) rely on changing the pre-softmax logits and differ in their definition of b.

Rotary Positional Encoding (RoPE). RoPE [Su et al.](#page-12-17) [\(2024\)](#page-12-17) encodes position information in attention logits by applying a rotation transformation to the query and key vectors based on their relative positions. Despite being widely used, RoPE exhibits limited length generalization [Press et al.](#page-12-8) [\(2021\)](#page-12-8); [Kazemnejad et al.](#page-11-4) [\(2024\)](#page-11-4).

B PRIOR RESULTS ON MULTI-DIGIT ADDITION

In this section, we summarize the methods proposed by prior work to get length generalization on addition along with their corresponding performance. In what follows, we indicate in red the positional encoding and in green the data format used in these works. We also take as a running example the addition 576+361=937.

733 734 735 736 737 – [Lee et al.](#page-11-18) [\(2023\)](#page-11-18) from 7 to 7-digit $(1.0 \times)$. APE + Reversed format. They train their models by reversing each operand as $675+163=739$. Therefore, the causal model that processes information from left to right can start with the least significant digit and proceed to the most significant digit, which corresponds to the algorithm for addition. They do not achieve any length generalization.

738 739 740 741 – [Kazemnejad et al.](#page-11-4) [\(2024\)](#page-11-4) from 8 to 9-digit $(1.125 \times)$: NoPE + Reversed format. They show that a model without positional encoding trained on reversed additions like 675+163=739 outperforms those with specialized positional encodings like T5's relative positional [Raffel et al.](#page-12-9) [\(2020\)](#page-12-9) or RoPE [Su et al.](#page-12-17) [\(2024\)](#page-12-17).

742 743 744 745 746 – [Shen et al.](#page-12-4) [\(2023\)](#page-12-4) from 10 to 11-digit $(1.1 \times)$: NoPE + Reversed format + random space augmentation. They introduced random spacing between digits, aiming to alleviate the model's reliance on absolute positional information. Combining this with the reversed format, the running example becomes 6 75+16 3=739. They show that NoPE Transformers length generalize from 10 to 11 digit-addition.

747 748 749 750 751 – [Zhou et al.](#page-12-2) [\(2023\)](#page-12-2) from 30 to 45 digits $(1.5 \times)$: NoPE + Index Hints. They define "index hints", a formatting that consists in adding a letter in front of each digit in the addition to indicate their position. For instance, the running example becomes $a5b7c6+a3b6c1=a9b3c7$. This change is applied during training and inference and enables transformers to execute indexing via induction heads [Olsson et al.](#page-12-18) [\(2022\)](#page-12-18).

752 753 754 755 – [Zhou et al.](#page-12-5) [\(2024\)](#page-12-5) from 40 to 100 digits $(2.5 \times)$: Fire [Li et al.](#page-11-7) [\(2023\)](#page-11-7) + Randomized positional encoding [Ruoss et al.](#page-12-16) [\(2023\)](#page-12-16) + Reversed format + Index Hints . They use a combination of two positional encodings: Fire [Li et al.](#page-11-7) [\(2023\)](#page-11-7), a additive relative positional encoding that has obtained strong length generalization on natural language benchmarks and Randomized positional encoding [Ruoss et al.](#page-12-16) [\(2023\)](#page-12-16): a technique that samples encodings from a range exceeding test-time lengths.

Figure 8: Comparison of different positional encodings for length generalization on a randomly generated Turing machine using the same hyperparameters (peak learning rate of $7e - 5$, batch size of 16, trained for 200, 000 steps).

 60 80 100 120 140 Length of Tape

HAlibi Alibi NoPE RoPE

The goal is to ensure that Transformers can adapt to larger positional encodings during training and not encounter any OOD encoding at test-time. With reversed format and index hints, the data format looks like a6b7c5+a1b6c3=a7b3c9. By using all these modifications, they reach state-of-the-art performance on length generalization for addition. However, these choices seem to be specific to the addition case and hard to transfer to other algorithmic tasks.

C ADDITIONAL EXPERIMENTAL RESULTS

 0^{+}_{40}

Accuracy (%)

Additional experimental results for Turing simulation are shown in Figure [8.](#page-14-2)

 To the best of our knowledge, [Zhou et al.](#page-12-5) [\(2024\)](#page-12-5) achieved length generalization mainly for addition when the two summands had the same length. Our method generalizes even when the two summands have different lengths. For $L_1, L_2 \in \{17, 32, 47, 62, 77, 92\}$, we sampled 96 addition examples where the first summand has length L_1 and the second summand has length L_2 . The accuracy for each combination is shown in Figure [9.](#page-15-0) We see that it generalizes well beyond the trained distribution $(L_1, L_2 \leq 50).$


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D RASP TURING PROGRAMS
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       HOU ET AL(2023))
       import numpy as np
       def full(x, const):
            return np. full_like (x, \text{ const}, \text{ dtype}=int)def indices(x):
            return np. arange (len(x), dtype=int)
       def tok_map(x, func):
            return np. array (\lceil \text{func}(xi) \rceil for xi in x ]). astype (int)
       def seq_map(x, y, func):
            return np. array (\lceil \text{func}(xi, yi) \rceil for xi, yi in zip(x, y)). astype (int)
       def select (k, q, pred, causal=True):
            s = len(k)A = np \tvert z \text{eros}((s, s), dy \text{pe} = \text{bool})for qi in range(s):
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           return aggr(s elect (k, q, pred), v, default=default, reduction=reduction)HOU ET AL(2023))
           import operator as op
       import numpy as np
       # D e f i n e comparison o p e r a t o r s
       equals, leq, lt, geq, gt = op.eq, op.le, op.lt, op.ge, op.gt
       def shift\_right(x, n, default=0):
           # shifts sequence x to the right by n positions
           return kqv(indices(x) + n, indices(x), x, equals, default=default)
       def cumsum (bool array):
           # r e t u r n s number o f p r e v i o u s True e l e m e n t s i n b o o l _ a r r a y
           return sel_width (self (bool_array, boolean_array, lambda k, q: k))def where (condition, x_i if, y_e else):
           # e q u i v a l e n t t o np . where ( c o n d i t i o n , x _ i f , y _ e l s e )
           x_m asked = seq_map(x_i if, condition, lambda x, m: x if m else 0)
           y_masked = seq_map(y_else, condition, lambda y, m: y if not m else 0)
           return seq_map (x_mmasked, y_masked, lambda x, y: x if y == 0 else y)
       def mask(x, bool\_mask, mask\_val=0):
           # e q u i v a l e n t t o x * bool_mask + d e f a u l t *(~ bool_mask )
           return where (bool\_mask, x, full(x, mask\_val))def maximum(x):
           return kqv(x, x, x, lambda k, q: True, reduction='max')
       def minimum(x):
           return -maximum(-x)def argmax(x):
           mm = maximum(x)return \kappa qv (mm, x, indices (x), reduction = 'max')
       def argmin(x):
           return argmax(-x)def num prev(x, queries):
           # o u t p u t [ i ] = number o f p r e v i o u s e l e m e n t s o f x e q u a l t o q u e r i e s [ i ] , i n c l u s i v e
           return sel_width(s elect(x, queries, equals))def has\_seen(x, queries):
           return kqv(x, queries, full(x, 1), equals, default=0)def firsts (x, queries, default = -1):
           # f i n d t h e i n d e x o f t h e f i r s t o c c u r r e n c e o f each query [ i ] i n x
           \# out[i] := np.flatnonzero(x[:i+1] == queries[i]).min()
           return kqv(x, queries, indices(x), equals, default=default, reduction='min')def lasts (x, queries, default = -1:
           # f i n d t h e i n d e x o f t h e l a s t o c c u r r e n c e o f each query [ i ] i n x
           # out[i] := np. flatnonzero(x[:i+1]) == queue[i]). max()return kqv(x, queries, indices(x), equals, default=default, reduction='max')
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       def index_select(x, idx, default = 0):
           # i n d e x e s i n t o s e q u e n c e x , v i a i n d e x s e q u e n c e i d x
           # i . e . , r e t u r n x [ i d x ] i f i d x [ i ] <= i e l s e d e f a u l t
           return kqv(indices(x), idx, x, equals, default=default)def first_true (x, default = -1):
           # r e t u r n s t h e i n d e x o f t h e f i r s t t r u e v a l u e i n x
           seen\_true = kqv(x, full(x, 1), full(x, 1), equals, default=0)first\_occ = seq\_map (seen_true, shift_right (seen_true, 1), lambda curr, prev: cur
           return kqv(first\_occ, full(x, 1), indices(x), equals, default=default)def induct_kqv(k, q, v, offset, default=0, null_val=-999):
           # get value of v at index of: first occurrence of q[i] found in k (if found) + o
           # ( excludes the last OFFSET tokens of k from matching)
           # null_val is a special token that cannot appear in k or q; used to prevent acci
           indices_to_opy = first s (shift\_right(k, offset, default = null_val), q, default = null\text{copied}_\text{values} = \text{index}_\text{select}(v, \text{ indices}_\text{top} \text{copy}, \text{default} = \text{default})return copied_values
       def induct(k, q, offset, default=0, null_val = -999):
           return induct_kqv(k, q, k, offset=offset, default=default, null_val=null_val)
       def induct_prev(k, q, offset, default=0, null_val = -999):
           # A version of induct for negative offsets.
           indices_to_opy = firsts (k, q, default = null_val) + offsetcopied_values = index_select(k, indices_to_copy, default=default)
           return copied_values
       D.3 UTILITY FUNCTIONS
       def prefix_fill(x, n, value):
           ones = full(x, 1)no\_fill = shift\_right (ones, n)return where (no\_fill, x, full(x, value))def where 3 (cond, x, y, z):
           out = where (cond == 0, x, y)
           return where \text{(cond == 2, z, out)}D.4 TURING MACHINE TRANSITION FUNCTION
       sep = 0bos = 1eos = 2
       empt = 3alphabet = list(range(4, 16))stategrace = list (range (16, 32))
       state_transition = \{a: \{s: np.random choice(state_space) \text{ for } s \text{ in state_space} \} for a
       symbol_transition = {a: {s: np.random.choice(alphabet) for s in state_space} for a in
       move_direction = {a: {s: np.random.choice([0, 1]) for s in state_space} for a in alp
       def next_state (state, token):
           if token in state_transition.keys() and state in state_space:
                return stat = transition \lceil \text{token} \rceil state \rceile l s e :
                return 0
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      def next_symbol(state, token):
           if token in alphabet and state in state_space:
                return symbol_transition [token] [state]
           elif token == bos:
                return bos
           elif token == eos:
                return eos
           e l s e :
                return 0
      def move (state, token):
           if token in alphabet and state in state_space:
                return move_direction [token] [state]
           elif token == bos:
                return 1
           e l s e :
                return 0
      D.5 COMPUTATION OF NEXT TAPE STATE
      def get\_next(x, x\_left, x\_right):
           # compute t h e n e x t s t a t e o f head and new symbol , w i t h o u t moving t h e head
           x_{\text{1}} \text{state} = \text{seq\_map}(x, x_{\text{1}} \text{left}, \text{next\_state})x_symbol = seq_map(x_right, x, next_symbol)
           x_move_R = seq_map(x, x_{\text{left}}, \text{move})is head = tok_map(x, lambda z: z in state_space)
           is head right = tok_map(x-right, lambda z: z in state space)
           x_n = x + x + y is head, x_n = x + yx_n = x + x + y is head right, x_n is x_n = x + yx_move_R = x_move_R & is headreturn is head, x next, x move Rdef select_move_token(no_head_around, head_left_move_left, head_left_move_right, hea
           LEFT TOKEN = full (no head around, 0)
           CUR_TOKEN = full (no\_head\_around, 1)RIGHT_TOKEN = full (no\_head\_around, 2)out = CUR TOKEN
           out = where (head_left_move_right | is_head_move_left, LEFT_TOKEN, out)
           out = where (head\_right\_move\_left \mid i_s\_head\_move\_right, RIGHT_TOKEN, out)
           return out
       def move_head ( cur_state, right_state):
           is head, cur next, move R = cur state
           right_is_{head}, right_{next}, right_{next}, right_{model}left\_is\_head, left\_next, left\_memory\_R = shift\_right ( is\_head , 1), shift\_right ( cur\_ino_head_around = (-let\_is\_head & -right\_is\_head & -is\_head)head[left_move_left = left_is_head \& ~left_move_R
           head [left_move_right = left_is_head & left_move_R
           head\_right\_move\_left = right\_is\_head \& \sim right\_move\_Rhead\_right\_move\_right = right\_is\_head \& right\_move\_Ris head move left = is head \& ~move R
           is head move right = is head & move R
           x_sel_move = select_move_token(no_head_around, head_left_move_left, head_left_mo
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            return where3(x_{\text{self} \text{move}}, \text{left} \text{next}, \text{current}, \text{right} \text{next})def next\_tape(x, shift):
            # compute the state of the head, after shifting by some n >= 2x_{-} = shift_right(x, shift)
            x left = shift right (x, \text{shift+1})x _right = shift _right (x, \text{shift} -1)x right right = shift right (x, \text{ shift } -2)# compute the next state (before moving the head) for current tape and right tape
            cur\_state = get\_next(x_, x\_left , x\_left , x\_right )right\_state = get\_next(x\_right, x_, x\_right\_right\_rightx_n = x = move\_head(cur\_state, right\_state)return x_next
       D.6 HASHING FUNCTIONS
       MAX INT = 32def hash_n_gram(x, n):
            out = xb e f o r e _ l a s t _ s e p = tok_map(x, lambda z : z = 0)
            shifted = shift\_right(x, 1)for i in range(n):
                 shifted_is\_sep = tok\_map(\{shifted\}, \{lambda\} \, z := 0)before\_last\_sep = shifted\_is\_sep \mid before\_last\_septo\_add = seq\_map(\text{shifted}, \text{before}\_\text{last}\_\text{sep}, \text{lambda } a, b: a*(1-b))# add t o hash
                 out = seq\_map (out, to_add, lambda a, b: b + MAX_INT * a)
                 shifted = shift\_right( shifted, 1)return out
       def hash_n_gram_iter(x, n):
            is\_sep = tok\_map(x, lambda z: z == 0)sep\_cs = cumsum(i s\_sep)x_{\text{hash}} = \text{hash}_n_{\text{gram}}(x, n)return seq_map (sep_ccs, x_h hash, lambda a, b: a + (MAX_mTx**n)*b)
       D.7 NEXT-TOKEN PREDICTION FOR TURING PROGRAMS
       def next\_token\_turng(x):
            x_n = x_t - t ape 2 = \text{next\_t} ape (x, 2)x_n = x_t - \tan x = n \cdot x_t = \tan x (x, 3)
            x_n = x_t + \text{ap}e_3 = \text{prefix\_fill}(x_n = x_t + \text{ap}e_3, 2, \text{em}e_1)k = \text{hash\_n\_gram\_iter}(\text{x\_next\_tape\_3, 1})q = hash_n_gram\_iter(x, 1)v = x_n e^{i\omega t} = tape 2out = kqv(k, q, v, equals, reduction='max')return out [-1]E SGD TURING PROGRAM DESCRIPTION
```
1079 We briefly describe here the Turing Program we used in [Subsection 5.2.](#page-8-0) Beyond the numerical tokens "a0, a1, a2,... a199", we include tokens "\$, d, yp, g , cur , |" to aid the calculation. A typical CoT for a gradient descent then looks like the following:

 In the above example, the first line provides a dataset of size three where "d a179 a166 , a76" denotes the first example ("a179"and "a166" are the coordinates of \vec{x} , "a76" is the value of y, and "d" is a token that denotes the start of an example). From the second line onward, we perform gradient descent starting from the last data point, working backward: On the second line, the original dataset is copied, while the "a100" following "yp" is the predicted value of y given the initial weight and the last feature vector "a77 a139", the "g a101 a99" says that $\lambda \nabla_{w_i} ||y_i - \vec{w_i} \cdot \vec{x_i}||$ has value "a101 a99", and "cur a99 a101" means that the current weight after update is "a99 a101". After a example's gradient is calculated, we delete that example.

F TURING PROGRAMS FOR SIMULATING TURING MACHINES

 We use the tokens space $a_1, a_2, \ldots, b_1, b_2, \ldots, s_1, s_2, L, R|, (0, c_1)$ ζ , <|BOS|>, <|EOS|>, <|SEP|>}, where the a_j 's are input symbols, the b_j 's are symbols substituting the a_j 's when the head is pointing to them and $(,), |, \sim, L, R$ are symbols used to encode the transitions. For instance, the transition (s_1, a_6, L) means that the Turing machines moves to state s_1 , edits the tape by writing a_6 and moves the head to the left.

G COMPARISON WITH PAST METHODS

 In this section, we show the performance of some of the methods mentioned in Appendix [B](#page-13-1) under our experimental condition. We consider three data formats:

- Reversed format in [Shen et al.](#page-12-4) [\(2023\)](#page-12-4).
- Index hints in [Zhou et al.](#page-12-2) [\(2023\)](#page-12-2)
	- Index hints + Reversed format in [Zhou et al.](#page-12-2) [\(2023\)](#page-12-2)

 Moreover, we consider three positional encodings: ALiBi, NoPE, and RoPE. We performed the addition experiments under the exact hyperparameter setting of Figure [3.](#page-4-0) The results are shown in Figure [10.](#page-21-0)

 Figure 10: Comparison of different positional encodings and data formats for addition. All hyperparameters were held fixed: learning rate of 7e − 5, batch size of 16, and trained for 200k steps.

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