# RDAS: A Low Latency and High Throughput Raw Data Engine for Machine Learning Systems

Anonymous Author(s) Affiliation Address email

## Abstract

 In the era of large pretrained models, a key challenge in deep learning is the underutilization of fine-grained raw data, often replaced by information-lossy normalized data. To bridge this gap, we introduce the Raw Data Aggregation System for Machine Learning (RDAS). RDAS offers a seamless data interface, enabling machine learning systems to directly access unstructured, high-resolution raw event data with minimal latency. At the heart of RDAS lies the Message Book Model, an innovative data representation framework that underpins the system's ability to handle event data at nanosecond precision. RDAS is structured around three conceptual layers: (i) the Message Layer, featuring dual message aggregators for sequential and random access, which compile raw messages into timestamp-specific message book snapshots; (ii) the Feature Layer, which derives user-specified data features from the message book for any given moment; and (iii) the Verification Layer, tasked with real-time error monitoring and integrity assurance of the message book. A C++ implementation of these layers ensures RDAS's exceptional performance. To validate its effectiveness, we applied RDAS in an Internet of Things (IoT) scenario, demonstrating significant performance enhancements over existing methods in terms of data throughput and latency. Our results underscore RDAS's potential to revolutionize data processing in machine learning, offering a pathway to leverage the full spectrum of raw data's granularity and richness.

# 21 1 Introduction

 Recent advancements in machine learning (ML) have seen large pretrained models, such as those used in natural language processing (NLP), computer vision (CV), and multi-modality fields, achieve unparalleled performance. A key characteristic of these models is their substantial data requirements. For example, ChatGPT, a model renowned for its capabilities, was trained on an extensive dataset comprising 570GB of text data from the Internet. The effectiveness of these models is further enhanced by Transformer-based architectures, which are known for scaling efficiently with increased data size [\[6\]](#page-9-0). Given the ongoing trend towards larger models, it is reasonable to anticipate that current data scales will continue to grow to meet these evolving requirements.

 While the success of large pretrained models in NLP and CV is noteworthy, their expansion into other domains like time series prediction [\[9\]](#page-9-1) and DNA understanding [\[5,](#page-9-2) [28\]](#page-10-0) presents new challenges. The

key issue lies in the limited quantity and granularity of training data available in these fields. For

example, popular datasets in time series prediction, such as ETTh1, ETTh2, ETTm1, ETTm2, ILI,

Traffic, and ECL, are relatively small, typically under 500MB, with only a few reaching between 1 to

10GB. These datasets, predominantly normalized and simplified from their original, more complex

forms, result in significant information loss. ETTh1, for instance, is an hourly electricity dataset

 derived from higher-frequency sensor data. This loss is a major drawback, as raw datasets, often unstructured like event messages, contain richer details than their normalized counterparts.

 To address this, we propose an integrated approach encompassing end-to-end data acquisition and processing. This method differs from traditional practices by processing raw data into structured form dynamically during model training. Unlike the static nature of preprocessed, normalized data, this dynamic approach allows for flexible and adaptive data transformation. This could include batch- specific normalization adjustments based on prior training results or on-the-fly data augmentations. Ultimately, this end-to-end process aims to harness the full potential of raw data, preserving its fine-grained nature for more effective model training.

 However, the advantage of raw data's granularity comes with its own set of challenges, notably its unstructured nature and inherent heterogeneity. This heterogeneity significantly complicates the data acquisition and processing. Take autonomous driving systems as an example: they rely on a diverse array of sensors, including laser, image, inertial measurement units (IMUs), and odometry sensors [\[14\]](#page-9-3). Each sensor type generates data streams that vary in format, type, and resolution, adding layers of complexity [\[7\]](#page-9-4) [\[8\]](#page-9-5) [\[17\]](#page-9-6) [\[10\]](#page-9-7) [\[26\]](#page-10-1).

 To effectively manage the complexities of raw data acquisition and processing, thereby enabling ML systems to access the most detailed information in raw data, we propose abstracting diverse systems into a unified framework. This model conceptualizes the process as a competition among heterogeneous entities for a variety of resources, governed by multiple constraints like time of arrival, importance, ranking, cost, and gain. An entity could be, for example, a robot awaiting a task or a request for computational resources in a cluster. We envision this as N connected queues, where N represents the number of constraints, and refer to it as a 'message book' for clarity.

 The message book abstraction offers several advantages for complex systems management. First, its intuitive nature facilitates easy understanding and implementation. Second, it is a versatile abstraction, applicable across various data types and systems in different domains. For example, in cloud computing, the message book can represent tasks awaiting execution, with computational power and time as the contested resources and submission time and task priority as constraints. The primary goal in such a system is maximizing the allocation of computational resources over time. This concept extends beyond cloud computing to encompass IoT systems, single-robot systems, and large-scale decentralized robotics systems. Third, the message book's construction is driven purely by data, relying on the most fundamental and unprocessed message data, thereby eliminating the need for complex data preprocessing. These attributes render the message book abstraction particularly well-suited for creating efficient, low-latency data visualization tools and data access interfaces in ML systems.

 The implementation of a message book data structure is crucial for efficiently managing complex systems. Firstly, without such a model, it becomes challenging to determine the order in which participants access resources, as this decision relies heavily on various constraint functions. Secondly, given the heterogeneity, volume, and dispersed nature of the data, directly analyzing every minute activity of each participant is impractical. The message book addresses this by offering a unified representation that consolidates this diverse, granular data in real-time. This occurs concurrently with the data loading and model training processes in a ML system. As a result, the message book not only simplifies the understanding of the system's real-time dynamics but also serves as an effective intermediary for subsequent tasks. These tasks can range from visualizing system statistics to aiding in model training, thereby providing a foundational tool for various downstream applications.

81 Constructing the message book requires the system to adeptly route, filter, process, and aggregate all the activity information generated during its operation. We conceptualize each activity within the system as an event, with each occurring event represented as a message initiation and transmission process. Given that the message book's state is continuously evolving with the system's operation, it is crucial to model both the message book and the event message manipulations in a streaming fashion. The core research challenge we address involves developing an event-based data engine. This engine is designed to intelligently aggregate the most detailed event messages into a dynamic, general-purpose message book data structure, operating in a streaming manner. The proposed data engine is tailored to facilitate both real-time and historical data retrieval, offering efficient space and time complexity.

 We organize the rest of the paper as follows. Section 2 mentioned prior work related to data processing systems, event processing systems, etc. Section 3 demonstrates the system architecture and implementation of RDAS. Section 4 shows experiments and results. Section 5 is the conclusion.

# 94 2 Related Work

 Data loading frameworks Common data loading frameworks such as PyTorch Datasets & DataLoad-96 ers<sup>[1](#page-2-0)</sup> and Tensorflow tf.data<sup>[2](#page-2-1)</sup> are in a different position in the ML pipeline. They are for loading structured data from permanent storage such as local disk and remote storage such as S3. RDAS is to address the first-mile problem of how to acquire and aggregate raw data from data sources into a structured representation. These data loading frameworks can be naturally the next component that connects to RDAS in the ML pipeline.

 Data processing system. General-purpose data processing systems include Pandas [\[11\]](#page-9-8), Dask [\[20\]](#page-10-2), Numpy [\[15\]](#page-9-9). However, they are for structured data that is generated by processing the raw unstructured data. There are also some big data frameworks that can process raw data such as Hadoop [\[24\]](#page-10-3), Flink [\[1\]](#page-9-10), Spark [\[27\]](#page-10-4), Storm [\[2\]](#page-9-11) and Hive [\[23\]](#page-10-5). But, they all serve as basic infrastructures and building blocks to construct data processing pipelines. They lack the higher level design to meaningfully process the raw data. This is where RDAS steps in.

 Event processing system. In an IoT system, information sharing and communication are often modeled as the distribution of real-time event messages. [\[19\]](#page-10-6) proposes an event processing solution to detect vehicle speed violations. [\[12\]](#page-9-12) proposes an event processing engine to process events from data streams for supply chain management purposes. [\[22\]](#page-10-7) proposes an event processing architecture to monitor the automotive manufacturing process. However, all these solutions are domain-specific. They are not a generic framework that is suitable for various kinds of domains in ML.

# 3 Methodology

## 3.1 Message Book

 To maintain a simple and efficient data structure, we use queues to implement the message book data 116 structure. If there are  $N$  constraints, there are  $N$  different queues. Each constraint corresponds to one queue and we store the specific entity objects in the queue whose constraint has the lowest priority. The entire message book data structure is a hierarchical structure. Each element of a queue with a specific constraint is a queue whose constraint is of a lower priority. For example, Figure [1](#page-3-0) shows a minimal example of the message book. In this example, there are two constraints and constraint 1 has a higher priority than constraint 2, meaning that constraint 1 has to be met first. To be more intuitive, taking game matching as an example, constraint 1 could be the players' skill levels and constraint 2 could be the players' geographical distance. Each entity object stored in the message book represents a player in this case. Game matching for a player needs to first find other players with the same skill levels then among which find the player with the smallest distance. In this case, each element of the constraint 1 queue is a Struct that includes some basic information for that level such as the total number of players that are in that level, and a pointer to the queue of players in that level. In the queue of players, each element is a Struct representing a player and they are sorted according to constraint 2, which is the geographical distance that is of a lower priority. We use C++ standard library to implement the message book. The details are in ??.

 Traverse The traversing of all entities (players) follows the priority of the constraints that are level first and entity second. This process starts from the index of the first level. For each level, it starts from the first entity and traverses all entities following the pointer to the next entity in each entity's entry. Then, it follows the pointer to the next level in the current level's entry and traverses all entities 135 of the next level. This process goes on until it traverses all levels' entities. The traversing has  $O(n)$ time complexity.

 Random access to existing entities Each time a new level or a new entity is added to the message book, their id-to-index mapping is saved in the hash map. Thus, for the random access to any specific 139 entity or level, it is  $O(1)$  time complexity using the hash map.

<span id="page-2-1"></span><span id="page-2-0"></span>[https://pytorch.org/tutorials/beginner/basics/data\\_tutorial.html](https://pytorch.org/tutorials/beginner/basics/data_tutorial.html)  $^{2}$ <https://www.tensorflow.org/guide/data>



<span id="page-3-0"></span>Figure 1: A Minimal example of the message book implementation. Left: a visualization of the message book with two constraints. Constraint 1 has a higher priority than constraint 2. Right: the implementation of the message book using two queues. The top corresponds to constraint 1 and the bottom corresponds to constraint 2. Only the queue of the least prioritized constraint(bottom queue) stores the entity objects. The top queue only stores level information.

 Add entities or levels To add a new entity to the queue of Constraint 1 or add a new level to the queue of Constraint 2, the algorithm first gets the index of free entry in the queue by popping an element from the corresponding FreeVector. Then, it creates a Struct representing the new entity 143 or new level in the location indicated by the index in the corresponding queue. This is  $O(1)$  time complexity.

 Delete entities or levels To delete an entity or a level, RDAS gets its index in the corresponding queue using the corresponding hash map. Then, it pushes that index into the corresponding FreeVector. 147 This is  $O(1)$  time complexity.

#### 3.2 Data Transmission Format and Channels

 The domains of machine learning systems that our system is for are those whose raw data are event data. For instance, in transportation systems, all raw data are event data that depict vehicles', road sensors', and traffic lights' status at each time point. In cloud systems, all the raw data are event data about different machines, different components, or different software of the system. In IoT or robotics, all the raw data are the events generated by different sensors, devices, or robots. Because of the blooming of edge device computation power and high-speed, low-latency communication tech- nologies, highly decentralized IoT systems and robotics systems are becoming a reality. We choose a data format that is suitable to the current and future trends of heterogeneous event data transmission scenarios in various domains. In this case, we need a suitable data storage and transmission format that is capable of providing the properties of large volume, high frequency, low latency and high compatibility with network transmission of the event messages data because all the sensors and other system participants transmit data over the network [\[4\]](#page-9-13) [\[16\]](#page-9-14) [\[25\]](#page-10-8) [\[13\]](#page-9-15). In RDAS, we store all the raw event messages in *pcap* (Packet Capture) format. Pcap is a direct capture of the data packet (event messages) that the senders send over the network [\[21\]](#page-10-9). It provides a single truth of source and an unbeatable granularity of up to nanosecond regarding the timestamp [\[3\]](#page-9-16). These features ensure that it has high credibility and enough capacity for any highly heterogeneous raw data source requiring various degrees of transmission latency, transmission frequency, and communication credibility.

 Besides the foundational data storage format for the raw event message data, we also define a concept called data transmission channel in RDAS. There are four different channels as shown by Figure [2,](#page-4-0) which are actually four parallel data feeds when the system is operating. They are Alpha channel 1, Alpha channel 2, Beta channel 1 and Beta channel 2. Alpha channel only includes incremental information so that the transmission can have extreme compactness and high frequency. The Alpha channel is necessary when the frequency of new updates is very high because it is impossible to store the snapshot of the entire message book's status for every new update regarding the scale of the storage that is needed and the large percentage of redundant data in each snapshot. Beta channel only includes snapshot information the sender sends at a slightly lower frequency so that the transmission data volume is not too large. In the meantime, any device listening to this feed at any time can get a relatively new snapshot of the sender's state and start evolving the state from this snapshot using information from the Alpha channel. To ensure a high degree of robustness of the data transmission,  both Alpha channel and Beta channel is a set of two parallel channels. Senders send the same packets simultaneously on channel 1 and channel 2. Receivers can use channel 1 and channel 2 to do cross verification or if some packets are corrupted on one channel, the system can recover them from the other channel. It is also noteworthy that the use of Alpha channel and Beta channel is flexible. For senders that only need to send stateless information, the use of Beta channel is optional. For senders that have a very limited amount of state data and have no requirement on high-frequency transmission, they can choose to only use the Beta channel to broadcast their full state periodically or whenever there is an update. Such a high degree of flexibility is an essential property that makes RDAS suitable to highly heterogeneous systems.

#### 3.3 Core Data Processing Pipeline

 The data processing pipeline of RDAS con- sists of four phases: channel merging, chan- nel routing, channel decoding, knowledge distilling, and verification. The first three phases support the message layer. The last phase supports the feature layer and the ver- ification layer. All these phases are piped to each other in a streaming manner, which means that it is not the case where a phase first processes all the data and then hands them over all at once to the next phase. It is the alternative case where when each phase processes a very small piece of the data, it immediately hands it over to the next phase



<span id="page-4-0"></span>Figure 2: Two types of data channels: Alpha for high-frequency incremental information; Beta for lowfrequency snapshot information. Each channel has a backup channel for robustness.

 and then starts processing the next piece. It is easy to see that in this streaming manner, all the events can be processed in the same chronological order as in that they took place and the large data size induced along the temporal dimension does not impair the processing speed.

 In the first phase, channel merging, the system merges all channels from all senders together into one monolithic stream. This process makes sure all the network packets in this stream are in chronological order using their timestamp of nanosecond granularity. In the second phase, channel routing, the system routes each packet to a specialized channel decoder that is responsible for decoding the data of a specific channel. The routing is based on the unique IP address of the sender. The third phase, channel decoding, as shown by Figure [3](#page-5-0) is a major part of the pipeline and handled by a channel decoder. There is a recovery mode and an incremental mode of the channel decoder and these two modes can switch back and forth into each other. It is usually in the recovery mode when the decoder is at the start of the decoding process or the decoded stream has corruptions for some reason in the middle of the decoding process and the decoder needs to be re-calibrate it to a correct checkpoint of the states of the sender. The incremental mode is to apply incremental changes to a base snapshot of states so that the system can maintain an evolving and always on-sync message book. There are two sub-phases in the decoder. We call the one before the decoder starts processing the packet on-packet-start, which is an API allowing the user to define and conduct any task before the decoder processes the packet. RDAS generates message book snapshots through this API. We call another one after the decoder processes the packet on-packet-end. It is also an API to allow users to define and conduct any task before the channel decoding phase ends for this specific packet.

 When the channel decoder is in recovery mode and a packet comes in, after the on-packet-start phase, the decoder checks the packet if it is from an Alpha channel or a Beta channel. If it is from an Alpha channel, the decoder would push it into the buffer queue without processing it. If it is from a Beta channel, the decoder hands it to a message parser that parses the content of the packet and produces a set of atomic operations. An atomic operation is a standardized and generic operation whose definition is a function. For example, some of the atomic operations are Add Event, Modify Event, and Cancel Event. The data engine uses a single set of available atomic operations of limited size for all kinds of data. Then, the decoder applies the set of atomic operations to the message book data structure so that the state of the message book incorporates the new information from the newcomer packet. Usually, in the recovery mode, because the message book is either in an empty state or corrupted state, the atomic operations are just Add Events to fill the empty message book or overwrite the existing message book.



<span id="page-5-0"></span>Figure 3: The feed decoding process: (1) two modes, recovery mode, and incremental mode. In recovery mode, the system uses the messages from the Beta channel to recover a snapshot of the message book from scratch. In the incremental model, the system applies incremental messages to the snapshot to evolve the message book. (2) The system preserves two user APIs on-packet-begin and on-packet-end, to generate useful information on the fly such as snapshots and features and do verification at any timestamp.

 When the channel decoder is in incremental mode, the process is almost the same as the recovery mode. The decoder first checks the buffer queue. If it is not empty, the decoder processes packets in

 the queue one by one until the queue is empty. After emptying the buffer queue, the decoder starts to handle the input packet. If the input packet is from a Beta channel, the decoder discards it. If it is

from an Alpha channel, the message parser parses it to a set of atomic operations then the decoder

applies them to the message book to incorporate new information from the packet.

 The last phase in the pipeline is knowledge distillation and verification. It supports both the feature layer and the verification layer. This phase usually does not happen at the end of decoding the current packet. It happens at the start of decoding the next packet to guarantee the captured snapshot is the closest timestamp to the boundary timestamp. The feature generation and message book verification happen through the on-packet-begin API so that users have the flexibility to tailor these two processes. This phase is an essential bridge users can use to connect the message book to various sorts of downstream tasks such as data visualization terminals, analytical models, logging devices, etc.

#### 3.4 Time Traveling

 The time-traveling feature is a representative feature of RDAS based on its core data processing pipeline. It allows the query of a snapshot of the data stream at any timestamp with a constant maximum loading time no matter how far the snapshot's timestamp is from the start of the stream. The mechanism behind the time traveler is that the system first generates and caches a bunch of snapshots in the permanent storage at a frequency much lower than the original raw data granularity, for instance, every 10 minutes. Then, when the user jumps to any specific timestamp, RDAS first loads the snapshot from the cache that is the closest to that timestamp and then applies all incremental messages between the snapshot timestamp and the target timestamp to the message book to generate the target snapshot the query requests for.

# 4 Experiment

### 4.1 Case Study

 We demonstrate the efficacy of RDAS by building a simulated high-frequency robotics logistics system and using it to acquire and parse the unstructured raw event data into the message book representation that is ready to be connected with the later data loading and model training pipeline. The reason that we choose to simulate a scenario in such a domain is that its raw data has high heterogeneity and high granularity, which provides enough complexities to test our system. Besides, the logistics industry is one of the kinds of complex systems that foresees to eventually transition from a labor-intensive industry into an automation industry. In 2021, there are about 1.3 million

 delivery drivers and about 40 million packages are in delivery every day in the US. In China, those numbers are about 4 million and 3 billion. However, this industry receives many critiques from both workers about its intense working schedule and clients about its far-from-satisfaction efficiency. Many delivery companies are considering or attempting to automate the delivery process by adopting delivery robots.

 In an efficient logistics system with all delivery drivers as delivery robots, millions of robots should be accurately and optimally dispatched to deliver tens of millions of packages based on their capacity and package pick-up distance and deliver them to the client. It is essentially a message book model with very high precision in timestamps and large enough capacity to handle millions of events sent from different robots in both real-time and offline analysis.

 In this scenario, there are two types of participants in the system. The first one is the local package dispatch centers that receive packages from higher-level dispatch centers from time to time. Whenever a batch of packages arrives at a local dispatch center, it sends out a Ready-to-Give (RtG) request to the system. The second one is the delivery robots that send a Ready-to-Take (RtT) request associated with the local dispatch center that is the nearest to them anytime they have capacity and are ready to fetch and deliver some new packages from the local dispatch center. During the operation, the system maintains a message book for each local dispatch center in parallel. In the message book, the constraints that decide the order of events are the distance and the time of arrival (ToA) of the request. The distance is discrete and represents the straight line distance to that specific dispatch center. Thus, there will be two types of distances. The first type is the distance between the destination of the delivery and the dispatch center. The second is the distance between the delivery robot and the dispatch center. There are different zones for different distance ranges. For instance, Zone 1 is within [0km, 3km), Zone 2 is within [3km, 6km), Zone 3 is within [6km, 9km), etc.

 During the operation of the system, all the delivery robots keep sending location/distance information periodically (for example, every 10 min) to the local dispatch center that they belong to. They are all in the message book no matter whether they are on their way to deliver packages or to the dispatch center to fetch new packages. Besides, all the packages are in the message book until the delivery is complete. They also have two states in the message book, either being-delivered or to-be-delivered. Once the dispatch center assigns a package to a delivery robot, the package's state changes from to-be-delivered to being-delivered and it leaves the message book once the delivery is complete.

#### 4.2 Experiment Setup

 The experiment is twofold. One is to demonstrate RDAS's large throughput, the other is to show the low latency and the low space cost of RDAS's time-traveling feature. We use ROS to simulate the dynamics of the aforementioned logistics system and generate the event message data. We assume there are always 2000 delivery robots and 20000 packages on the message book and the robots and the dispatch center send messages at some preset frequency. We built a separate ROS program to uniformly sample and generate these messages from all the possible message types to simulate reality. There are in total six types of event messages that affect the message book status. Three are for delivery robots and the other three are for packages. For each data stream, the system generates a snapshot every 10 minutes and caches it to the disk.

 For delivery robots, the first type of message is AddRobot. When a new robot becomes online, the robot sends this message. Its fields include distance, capacity, state, last\_update\_time. The second type is ModifyRobot. When a robot's status changes between Occupied and Empty, it sends out this message. The third type of message is RemoveRobot, when a robot becomes offline, it sends this message. For packages, when a new package arrives at the dispatch center and waits for delivery, the dispatch center sends out an AddPackage message. When a package's status changes from to-be-delivered to being-delivered, the dispatch center sends an ModifyPackage message. Lastly, when a package delivery is complete, the dispatch center sends a RemovePackage message to remove the delivered package from the queue.

 For the throughput experiment, the baseline system is the same as RDAS except that the baseline system does not have the incremental message aggregating mechanism and message book representa- tion. Instead, it assumes each message includes the full picture. We assume there are two message data streams for RDAS and the baseline system respectively. In each time unit, two streams have the same number of messages and the same amount of information such as the number of robots and

 packages and their status. The only difference is that one stream uses incremental messages and the other uses full-picture messages. We measure the total size of the messages each system processes within one second. Under the same network bandwidth, a system processes more information (higher throughput) at each time unit if it requires a smaller total message size to recover the same amount of information. Thus, we represent the throughput by taking the inverse of the total message size per second with a scale.

 The time-traveling feature is a representative feature that can demonstrate the superb performance of RDAS. We conduct experiments to measure the latency and the space cost of the time-traveling mechanism. To have a low latency for this mechanism, there are two types of overhead to consider. The first is the time cost to generate the snapshot cache. The second is the latency to apply the incremental messages between the loaded snapshot time and the target timestamp to generate a new snapshot on the fly. We measure both of them. Furthermore, we also measure the size of disk space to store the snapshot cache. The experiment demonstrates that the random access time traveler enables an O(1) latency with respect to the number of minutes away from the initialization point of the data stream and the size of the disk space to save the snapshot cache is within a reasonable range.

#### 4.3 Results

 We can see from Table [1](#page-7-0) and Figure [4](#page-7-1) that with the incremental message aggregation and the message book representation, RDAS's throughput is several magnitudes higher than the baseline system. When 339 the data stream's frequency is at  $10^3$ hz, RDAS's throughput is about 6 times larger than the baseline. 340 The discrepancy keeps increasing when the stream frequency increases. When the frequency is  $10^7$ hz, RDAS's throughput is about 400 times higher than the baseline.

<span id="page-7-0"></span>Table 1: Throughput of RDAS and the baseline system for data stream of different frequencies. It is a scaled value of the inverse of the total message size per second.

Data stream frequency					
				System $10^3$ hz $10^4$ hz $10^5$ hz $10^6$ hz $10^7$ hz	
RDAS. Baseline 0.15	$\mathbf{I}$	1.47	14.87	23.66	163.68 1850.35 5989.47 10366.29 26.18

 To demonstrate RDAS's low la- tency, we first run the system several times to repeatedly mea- sure the time it takes to process 10000 packets and take an aver- age. The final result is that it takes about 0.01 seconds to pro- cess 10000 network packets. If the data simulation program gen- erates the data stream at a 100hz frequency, a day(24 hours) of the streamed Pcap data has 8.64 mil- lion messages and the whole snap- shot cache generation for them only takes 8.64 seconds. Even the stream is of 100000hz frequency, meaning on average, there are

 100000 messages every second, which is already very unlikely in

 real-world robotics and IoT sys-tems or systems in other domains,



<span id="page-7-1"></span>Figure 4: Throughput (the larger the better): starting from the stream frequency of  $10<sup>3</sup>$ hz, the throughput of RDAS is about 6 times larger than that of the baseline system. When the frequency increases to  $10^{\circ}$ hz, the RDAS is about 400 times larger than the baseline. RDAS's throughput increases much faster than the baseline's when the frequency increases.

 RDAS only takes 2.4 hours to cache a day of snapshot data. Considering the snapshot cache genera- tion is an offline task, such a level of time complexity is already sufficient. It is also noteworthy that in these settings, the latency for the time traveler to generate a requested snapshot on the fly has a reasonably short upper limit which is 1 second.

 We also measure the random access waiting time in querying snapshots of the data stream at any timestamp to demonstrate RDAS's low latency. We construct a baseline system using the same codebase as RDAS except that the baseline does not have the random access message aggregating feature and it can only start building the target snapshot from the initialization point of the data stream. By comparing their random access waiting time, we can clearly see from Table [2](#page-8-0) and Figure [5](#page-8-1) that RDAS has a constant bound but the baseline keeps increasing, causing long latency when the target

timestamp is far away from the initialization point.

<span id="page-8-0"></span>Table 2: Random access waiting time (milliseconds) for RDAS is a piece-wise linear function with the peaks always at a similar constant number. However, it is an increasing linear function for baseline with no bounds, which indicates a much higher and eventually unacceptably long latency for the user to obtain a snapshot.



 Lastly, we examine the disk space consumption of the snap- shot cache. Although the snap- shot size varies in different sys- tems in different domains, our logistics robotics system exam- ple, with 2000 delivery robots and 20000 packages on the mes- sage book for each dispatch cen- ter at any time, is already am- bitious and thus, representative enough. We assume that with- out compromising the user experi- ence, a 1-second latency (waiting time) is a reasonable upper limit when the user jumps around to query snapshots at random times- tamps. It means we need to at least cache 1 snapshot for every 1 million messages because the



<span id="page-8-1"></span>Figure 5: Random access waiting time: On the data stream of 10000hz frequency, RDAS's random access waiting time always has a nearly constant bound at around 6000ms. However, the baseline model has a linearly increasing waiting time with respect to the further and further timestamps.

 processing speed is 1 million messages per second. Hence, disk consumption is a meaningful and critical metric. We generate a day of snapshot cache by generating one snapshot for every 1 million messages. The average size of each snapshot file is 100KB. Following the same deduction as above, if the simulated data stream is of 100hz frequency, the snapshot cache of one day is only about 860KB and it is about 860MB if the frequency is 100000hz. This shows that the space cost of RDAS is very low.

## 5 Conclusion

 We propose RDAS, a general-purpose raw data acquisition and processing system for machine learning systems in various domains. It provides a unified data interface to bridge the unstructured, high-resolution and heterogeneous raw data of up to nanosecond granularity to common data loading and model training pipeline. It features a novel data representation mechanism, the message book, with the incremental message aggregation mechanism and a low-latency random access time traveler. RDAS allows users to quickly query a snapshot of the message book at an arbitrary timestamp. The experiments demonstrate RDAS has a high throughput, low latency, and consumes reasonably small disk space with full support to high-resolution raw event data. Future works could provide a programming language agnostic interface that enables quick integration into any existing machine learning frameworks.

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# A Additional Background

 Robotics system. ROS (Robotics Operating System) [\[18\]](#page-9-17) is a system providing communication layers on top of machine operating systems to a computation cluster. Heterogeneous robots can use it to communicate with each other or a centralized cloud system. However, ROS only provides an underlying communication infrastructure and it does not deal with a higher level regarding data processing, which is where RDAS can step in.

**Data format.** Some other file formats are popular in data-logging scenarios. For instance, Ros Bag 489 file format <sup>[3](#page-11-0)</sup> is the format that Ros uses to store Ros messages in files. However, it is bound to Ros's 90 ecosystem. Another general-purpose message recording data format is the MCAP format <sup>4</sup>. However, both Bag and MCAP's disadvantage compared the PCAP format our data engine uses is that they capture data from the application layer while PCAP captures data from the lower network layer according to the 7-layer OSI model [\[29\]](#page-10-10). Capturing from the network layer is much faster. Besides, network layer has the additional network information that helps network monitor and analysis. It is crucial for the scenarios that put the requirement for low latency, high throughput and high robustness to extreme.

<span id="page-11-0"></span><https://wiki.ros.org/Bags/Format/2.0>

<span id="page-11-1"></span><https://foxglove.dev/blog/introducing-the-mcap-file-format>

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