#### **000 001 002** STEPPROOF: STEP-BY-STEP VERIFICATION OF NATURAL LANGUAGE MATHEMATICAL PROOFS

Anonymous authors

Paper under double-blind review

#### ABSTRACT

Interactive theorem provers (ITPs) are powerful tools for the formal verification of mathematical proofs down to the axiom level. However, their lack of a natural language interface remains a significant limitation. Recent advancements in large language models (LLMs) have enhanced the understanding of natural language inputs, paving the way for autoformalization—the process of translating natural language proofs into formal proofs that can be verified. Despite these advancements, existing autoformalization approaches are limited to verifying complete proofs and lack the capability for finer, sentence-level verification. To address this gap, we propose StepProof, a novel autoformalization method designed for granular, step-by-step verification. StepProof breaks down complete proofs into multiple verifiable subproofs, enabling sentence-level verification. Experimental results demonstrate that StepProof significantly improves proof success rates and efficiency compared to traditional methods. Additionally, we found that minor manual adjustments to the natural language proofs, tailoring them for step-level verification, further enhanced StepProof's performance in autoformalization.

**024 025 026**

#### 1 INTRODUCTION

**027 028**

**029 030 031 032 033 034 035** Mathematics is the basic tool for the development of science, and the reliability of its conclusions affects the stable growth of various disciplines. With the development of the mathematical edifice, mathematical proofs have become more complex and lengthy. The verification of mathematical proof often requires years of careful verification to ensure the accuracy of the work. However, reading a mathematical work requires a large amount of knowledge, and in the face of the many branches of mathematics today, traditional manual verification has become increasingly disastrous. Thus, an idea arose to validate mathematical work written in natural language automatically.

**036 037 038 039 040 041 042 043** At present, there are two kinds of automatic verification of mathematical proof. One is to write the mathematical certificate into a machine code that can be verified by a specific expert system, which is called the interactive theorem prover [\(Harrison et al., 2014\)](#page-8-0). After more than 40 years of development, the interactive theorem prover has begun to take shape and has been used in the verification of many mathematical works [\(Maric, 2015\)](#page-9-0). However, because such machine-verifiable proofs need to be written in a specific programming language, whose learning cost is relatively high, interactive theorem provers are only used by a small number of mathematicians at present [\(Nawaz](#page-9-1) [et al., 2019\)](#page-9-1).

**044 045 046 047 048 049 050 051** On the other hand, after the advent of the large language model[\(Zhao et al., 2023\)](#page-10-0), through prompt engineering [\(Liu et al., 2023\)](#page-9-2) and few shot learning[\(Wang et al., 2020b\)](#page-9-3), large language model can be applied to many natural language processing tasks, and achieved considerable performance[\(Achiam](#page-8-1) [et al., 2023\)](#page-8-1). However, due to the hallucination problem of large language model[\(Ji et al., 2023\)](#page-9-4), its performance in mathematics and strong logic-related work is not good enough[\(Huang et al., 2023\)](#page-9-5), so the lack of reliability of large language model to verify mathematical work is easy to cause a lot of errors. In this context, a method that combines large language models and theorem provers gradually comes into people's view, which is called autoformalization verification[\(Li et al., 2024\)](#page-9-6).

**052 053** At present, the existing automatic formal verification work generally adopts a FULL-PROOF strategy. Although such a strategy has achieved some impressive performance in some studies, there are still many problems in the stability of its proof and the fine-grained verification. Aiming at the **054 055 056** problems existing in FULL-PROOF, we innovatively propose a new automatic formal proof strategy, STEP-PROOF. And the performance of StepProof is better than traditional methods in several aspects.

**057 058 059 060 061 062** In summary, our contributions mainly include the following points: 1. We pioneered a novel natural language mathematical verification method StepProof, which realizes informal mathematical proof verification at the sentence level. 2. We were the first to realize the test of automatic formalization capabilities on small open-source LLMs. 3. Compared with existing methods, the StepProof method has been significantly improved in all aspects of performance.

**063 064 065 066 067 068 069** In this paper, we will first make a brief summary of the existing relevant works and the technical background in Chapter 2. In Chapter 3, we will give a detailed introduction to the two types of strategies, FULL-PROOF and STEP-PROOF, and point out many problems faced by traditional methods. In Chapter 4, we set up a series of experiments to verify the performance of STEP-PROOF in verification tasks and its superiority over FULL-PROOF. At the same time, we also carried out a detailed analysis of some phenomena observed in the experiment to further explain the reason for the advantages. In the end, we analyze and look forward to the current limitations and future development direction.

**070 071**

### 2 RELATED WORK

**072 073**

**074 075 076 077 078 079 080 081 082 Theorem Prover:** Theorem provers can be roughly divided into two types, interactive theorem provers (ITPs) in which the user can enter and verify the existing proof[\(Asperti, 2009\)](#page-8-2), and automated theorem provers (ATPs) that try to prove the statements fully automatically[\(Harrison, 2013\)](#page-8-3). The two types of theorem provers are not mutually exclusive [\(Nawaz et al., 2019\)](#page-9-1). Most ITPs such as Isabelle[\(Paulson, 1994\)](#page-9-7), Coq [\(Huet et al., 1997\)](#page-9-8) and Lean [\(De Moura et al., 2015\)](#page-8-4) are supported by ATPs that try to automatically prove "obviously" intermediate steps in the proof entered by the user. The entered proofs are rigorously verified back to the axioms of mathematics. Different ITPs use different axiomatic foundations, e.g. set theory, first-order logic, higher-order logic, etc. Each ITP use its own language and syntax, which makes the learning cost of theorem prover high, and precludes ITPs from being widely used [\(Yushkovskiy, 2018\)](#page-10-1).

**083 084 085 086 087 088 089 090 091 092 093** Large Language Model: In recent years, large language models (LLMs) have achieved outstanding performance in many downstream tasks of natural language processing. LLMs such as Llama[\(Dubey](#page-8-5) [et al., 2024\)](#page-8-5), GPT series[\(Ye et al., 2023\)](#page-9-9) and GLM 4 [\(GLM et al., 2024\)](#page-8-6) are trained on large databases to understand user input in natural language and produce the related output. While large language models perform well in general tasks such as translation, their performance in dealing with logic problems has been limited. A lot of work has also shown that large language models are prone to a variety of problems when dealing with logic problems. Although in the subsequent iteration of the model, the developers provided a large amount of high-quality logic-related data to improve the logic capability of the model, the effect that could be achieved was still very limited[\(Lappin, 2024\)](#page-9-10). Therefore, it has become a trend to add an expert system to the model to improve its accuracy, such as RAG[\(Fan et al., 2024\)](#page-8-7), which is currently commonly used in question-answering systems.

**094 095 096 097 098 099 100 101 102 103 104 105 106 107** Autoformalization: The definition of automatic formalization is very broad, but can be roughly seen as understanding and extracting translation from natural language to obtain the required structured data or formal language, such as entity relation extraction[\(Nguyen & Grishman, 2015\)](#page-9-11). Early automatic formalization work involved extracting logical propositions from natural language in addition to entity relation extraction[\(Singh et al., 2020;](#page-9-12) [Lu et al., 2022\)](#page-9-13). However, the main problem in this kind of work is that the extracted logical propositions lack corresponding symbols for derivation and application, so the output of the automatic formalization output cannot be directly applied. With the launch of pre-trained language models such as transformer and bert, language models have a stronger understanding ability. [Wang et al.](#page-9-14) [\(2020a\)](#page-9-14) conducted an early automatic formalization attempt for the theorem prover Mizar. However, due to the limited size and training expectations of the models at that time, the effects they could achieve were very limited. With the rapid expansion of the scale of models and predictions, many new and better performance automatic formalization work has emerged, such as the Majority Voting method by [Lewkowycz et al.](#page-9-15) [\(2022\)](#page-9-15), the DSP method by [Jiang et al.](#page-9-16) [\(2022\)](#page-9-16)., and the method proposed b[yZhou et al.](#page-10-2) [\(2024\)](#page-10-2), in DTV. They further improved the automatic formalization of the system by combining large models with some syntaxmodifying filters. However, on the one hand, all these works are tested in the closed-source large

 model Minerva, which lacks the testing work of open source model and small model. Meanwhile, all these works adopt the strategy of FULL-PROOF, which has poor controllability for the formal output of the model, and it is difficult to locate the error point of non-formal proof. Qinghua et al. proposed a problem location method based on the FULL-PROOF strategy in SlideRule. However, this method relies heavily on the format and quality of the generated formal content, so it cannot achieve 100% problem locations detected. To solve these problems, we propose StepProof, an automatic formalization strategy that can realize sentence-level verification, to realize the verification of natural language mathematical proofs.

# 3 STEPPROOF

 In this chapter, we will provide a detailed introduction to the workflow design of StepProof. We will also compare the STEP-PROOF strategy used by StepProof with the FULL-PROOF strategy adopted by existing autoformalization systems, highlighting the problems with traditional strategies and the advantages of STEP-PROOF over FULL-PROOF.

 

 

3.1 FULL-PROOF

 Current research on natural language proofs formal verification predominantly employs the FULL-PROOF strategy, as seen in methods like DSP and DTV. The workflow of FULL-PROOF method can be roughly illustrated as the left of Figure [1.](#page-2-0) Users submit a provable problem along with its proof process. The problem is first formalized, then the informal problem, the formalized problem, and the entire informal proof are packaged as inputs to a large language model for formalization. After obtaining the formal proof, it is combined with the formalized problem and input into a theorem prover for rule-based formal verification. The verification results are then returned to the user. Despite the clarity and simplicity of the FULL-PROOF workflow, it has significant drawbacks.



<span id="page-2-0"></span> Figure 1: Full-Proof Strategy generate from the whole proof and only provide proof result instead of detailed feedback to help user improve the proof. While Step-Proof separate the whole proof into sub-proof to proof from the bottom to the top and enable the user to get more detailed feedback and fine-grained operation.

 First, in the FULL-PROOF automated formalization process, the highly structured and formalized nature of the input and output, coupled with the numerous similarities in solving mathematical equations, often leads to excessive noise in the output. To obtain the desired formalized content, numerous filters must be set up, which results in considerable waste and contamination of generated content. This can also cause generation loops, where the same content is repeatedly generated, a common issue in FULL-PROOF strategies.

 Additionally, the length of proofs varies, and LLMs in FULL-PROOF struggle to adjust the output's max new tokens effectively based on input length. This leads to shorter proofs not being  truncated in time, thus generating repetitive or noisy content, and longer proofs lacking sufficient max\_new\_tokens.

 Lastly, the stability of FULL-PROOF generation is poor. For instance, a full proof might be almost entirely correct except for a minor error in a small step, leading to the failure of the entire formal proof. Users attempting regeneration may find previously correct parts presenting erroneous. Thus, FULL-PROOF requires highly accurate one-time generation of the entire content.

 Moreover, the correlation between formal and informal content generated by FULL-PROOF is weak, making it difficult for users to map formal feedback to corresponding informal content. Although LLMs can generate formal proofs with annotations to map back to informal proofs as Qinghua et al proposed the failure step detection method in SlideRule, this approach is unstable in practice and increases the required token count.

3.2 STEP-PROOF

 



<span id="page-3-0"></span>Figure 2: User Interface of StepProof

**216 217 218 219** To address the numerous issues faced by the FULL-PROOF strategy, we innovatively propose STEP-PROOF. STEP-PROOF employs a step-by-step generation and verification strategy, offering better performance and stability compared to FULL-PROOF. The workflow of STEP-PROOF is illustrated in the left of Figure [1.](#page-2-0)

**220 221 222 223 224 225 226** Unlike the one-time generation and verification of FULL-PROOF, STEP-PROOF assumes each sentence in the proof is a verifiable sub-proposition. Each step is formalized and pushed onto a formal proof stack, where it is verified along with other sub-propositions in the stack. Upon successful verification, the formalized proof and informal proof are packaged as inputs for the next step. For failed steps, the Step Proof allows users to backtrack, retaining previously verified steps and only clearing the erroneous step from the stack. The StepProof can then either re-formalize or optimize the existing informal step as needed.

- **227 228 229 230 231** STEP-PROOF offers several advantages over FULL-PROOF. First, it only needs to formalize single sentences in context, resulting in shorter, less noisy, and more stable output. The step-by-step generation strategy also means that each step's length is relatively fixed, eliminating the need for adjusting max new tokens and allowing the use of smaller max new tokens for formalizing longer theorems.
- **232 233 234 235 236 237** Moreover, STEP-PROOF's incremental generation and verification tolerate step errors well. Only the specific erroneous step needs to be retracted, rather than regenerating the entire content, enhancing robustness and efficiency. Finally, Step Proof ensures high correspondence between each informal and formal proof step, providing users with finer operational granularity. For instance, users can suspend a correct but incomplete step and assume it is correct to proceed, a functionality almost impossible under the FULL-PROOF strategy.
- **238 239 240 241 242 243 244 245 246 247** We also designed a simple and user-friendly interactive interface for the user, as shown in Figure [2.](#page-3-0) Users can complete the natural language proof of the problem through interface interaction, and each step of the proof will be formally verified after submission, and the verified proof step will be marked in green in the background to indicate that the current step has passed the verification and is reliable. If the current step does not pass the automatic validation but the user thinks it is true and wants to continue, you can select HOLD, and the current step will be highlighted in yellow in the background to indicate that the current step is in a suspended state. After completing all the proof steps, the user can input QED to indicate that the system has completed the proof, and the system will combine all the steps to perform the final verification of the proof target. Through the interactive interface, the user can also realize the PDF export of the proof process.
- **248 249**

**250**

**252**

#### 4 EXPERIMENT

**251** 4.1 EXPERIMENT SETUP

**253 254 255** To validate the performance advantages of the STEP-PROOF strategy over FULL-PROOF, and to compare StepProof's overall performance against existing automated formal proof methods, we conducted both strategy performance tests and baseline tests using the same dataset and model settings.

**256 257 258 259 260** For the test dataset, we used GSM8K, which includes a large number of informal mathematical problems and their correct informal proofs. These informal proofs can be easily segmented into a series of sub-steps. We chose Llama3 8B-Instruct as the large language model for automated formalization. Existing autoformalization tests use closed-source models, and we aim to fill this gap by using an open-source small LLMs.

**261 262 263 264 265 266 267** We set the temperature to 0.3 to balance stability and flexibility. Given the small parameter count of 8B, we used a single example for the few-shot in strategy performance tests. With proof steps in GSM8K [\(Cobbe et al., 2021\)](#page-8-8) averaging 4-5 steps, we set the max new tokens for FULL-PROOF to 1024, and 256 for each step in Step Proof. The test environment consisted of a single NVIDIA A4000 16GB, 8 cores of AMD5800X, and 32GB DDR4 3200 RAM. Isabelle2024 was used as the formal theorem prover with only *Main* prove library as the theorem base<sup>[1](#page-4-0)</sup>, with Isabelle-client [\(Shminke, 2022\)](#page-9-17) as the testing service proxy.

<span id="page-4-0"></span>**268 269** <sup>1</sup>Here, we only use Main as the proof library, considering that the use of different libraries will greatly affect the formal verification ability of the theorem prover, in order to provide a relatively standard index. In StepProof, we do not introduce other libraries to further improve the proof ability of the theorem prover. **270 271 272 273** In the strategy performance tests, we evaluated the performance of FULL-PROOF and STEP-PROOF on the GSM8K test set using the following metrics: 1. One-attempt generation proof pass rate  $r_p$ . 2. Average formalization time for passed proofs  $\mu_f$ . 3. Variance in formalization time  $\sigma_f^2$ . 4. Average proof time for passed proofs  $\mu_p$ . 5. Variance in proof time  $\sigma_p^2$ .

**274 275 276 277 278** In the baseline tests, we compared the multi-attempt test results of GSM8K by [Lewkowycz et al.](#page-9-15) [\(2022\)](#page-9-15) in Majority Voting and [Zhou et al.](#page-10-2) [\(2024\)](#page-10-2) in DTV with our results. We evaluated the performance based on the number of attempts and multi-round proof pass rate, allowing up to 10 retries for each failed step in each formalization attempt.

**279 280 281 282 283** At the same time, considering that StepProof has better granularity than traditional proof tasks, we not only evaluated the overall proof passing rate but also counted the proportion of step-proof passing. For example, if a 6-step proofs can be verified to be true in 3 steps, then we will mark that the step pass rate of the proof is 0.5. In this way, we quantify the formal proof capability of the method more comprehensively rather than the Proof Passing Rate.

**284 285 286 287** In addition, to verify the influence of the writing method of non-formal proof on the passing rate of StepProof formal proof, we extracted 100 questions from the Number theory of MATH[\(Hendrycks](#page-9-18) [et al., 2021\)](#page-9-18) and made simple manual modifications to make the proof step more consistent with the proof requirement of StepProof.

<span id="page-5-0"></span>4.2 EXPERIMENT RESULTS



<span id="page-5-1"></span>Table 1: Performance Test of Full-Proof and Step-Proof

In strategy performance tests, the STEP-PROOF strategy outperformed the FULL-PROOF strategy across the board on the GSM8K test set. As shown in Table [4.2,](#page-5-0) the STEP-PROOF strategy improved the one-attempt proof pass rate by 15.1% compared to FULL-PROOF. In terms of average formalization time, STEP-PROOF required 38.9% less time than FULL-PROOF. For average proof time, STEP-PROOF achieved a 39.5% performance improvement over FULL-PROOF. Additionally, STEP-PROOF showed more stability in both formalization and proof time compared to FULL-PROOF. These results confirm that our strategy offers better performance, efficiency and stability.



Table 2: Baseline Test in GSM8K

**316 317 318 319** In the baseline test as shown in Table [4.2,](#page-5-1) Step Proof surpassed DTV in multi-round verification tests on GSM8K, achieving a 10.3% performance improvement. Moreover, StepProof required fewer attempts compared to DTV<sup>[2](#page-5-2)</sup>, demonstrating its superior proof capability and further validating the advantages of the StepProof methods.

**320 321** Introducing more libraries in practice will greatly improve the proof ability of the theorem prover, and also improve the proof ability of the whole system to some extent.

<span id="page-5-2"></span>**<sup>322</sup> 323** <sup>2</sup>Don't Trust: Verify (DTV) originally used two close source models–GPT3.5 as the problem generation model and Minerva 8B as the proof generation model, while due to the Minerva being inaccessible and GPT3.5 being costing, we use the same method in DTV, but replace the LLM into Llama3.



 

<span id="page-6-0"></span>Table 3: Step Passing Rate Distribution in GSM8K

In the step pass rate test as shown in Table [3,](#page-6-0) we found that StepProof was able to perform some degree of validation for nearly half of the proofs after 10 rounds of trying. In LLAMA3 8B, 38.1% of the proofs completed more than half of the verification, and 27.9% of the proofs completed all of the verification. Compared with a single attempt, this is a significant improvement. At the same time, we propose a new index for a more comprehensive evaluation of the proof of automatic formal verification methods.



<span id="page-6-1"></span>

In our test to verify the influence of informal proof writing on the proof pass rate (as shown in the Table [4\)](#page-6-1), we found that the proof pass rate was significantly improved after simple fitting of the informal proof. This shows that compared with FULL-PROOF, STEP-PROOF adopts proof mathematics that is more suitable for step verification, which will be more conducive to improving the pass rate of automatic formal verification.

#### 4.3 EXPERIMENT ANALYSIS



<span id="page-6-2"></span>Figure 3: Time Cost of Two Strategies in Formalization and Proof

 From Figure [3,](#page-6-2) we find that compared with the FULL-PROOF Strategy, STEP-PROOF takes less time to formalize and prove and is more stable. On the one hand, the generation strategy of Step-Proof makes the content generated in a single attempt less and more stable. Stable step content reduces the number of false proofs and the time to repeatedly prove successful content. Therefore, in terms of both generation efficiency and proof efficiency, STEP-PROOF is superior to FULL-PROOF.

 To investigate the relation between step proof pass rate and number of attempts, we plotted Figure [4.](#page-7-0) It shows that most steps can be proven with relatively few attempts, with only a small fraction



<span id="page-7-0"></span>Figure 4: Passed Proofs and Steps in Different Attempts

requiring multiple tries. We believe the main limitation to step pass rate lies not in the model's formalization ability, but in whether the informal proof steps are suitable for conversion into provable formal steps. We found that many steps in the test set cannot be formalized into provable steps. These unformalizable informal steps significantly limit the further performance improvement of Step Proof. Therefore, to better achieve verification in practical applications of StepProof, each proof sub-step should be ensured to be a formally verifiable sub-proposition as much as possible.

 In addition, we also found that the processing capacity of LLMs is poor when faced with the processing of continuous equations, and it is easy to fall into the generation loop during the generation process, resulting in the poor quality of the generated formal proof. Therefore, in practical applications, it is necessary to split the steps as much as possible, rather than pursue once completion. At the same time, since the automatic proof ability of automated theorem proving is also limited, resolving the proof process is not only conducive to the formalization process, but also very helpful to the verification process.

 

## 5 LIMITATIONS & FUTURE RESEARCH

- At present, the performance of StepProof in small models has been verified in various aspects. However, due to the lack of equipment and funds, the performance of StepProof in larger models has not been verified, although the previous relevant work verified the conclusion that FULL-PROOF strategy has better performance in larger models. However, the STEP-PROOF strategy has not yet been tested on a larger model.
- In addition, StepProof is strict for users to enter proof steps, while FULL-PROOF is more flexible and can incorporate some non-proof explanatory language or prompt words into the proof. However, StepProof will make mistakes in the step proof because the prompt word is not provable.
- Finally, StepProof pays more attention to the sequential proof with steps, but when faced with some structured proof methods, its performance is still limited, while FULL-PROOF can better capture the overall proof structure.
- In the future, our work will mainly start from the following two points: 1. At present, we are writing a corpus for StepProof for automatic formal verification tasks oriented to step verification, and we hope that a more targeted corpus will help improve the step formalization ability of the model. 2. We will implement specific structured proof based on the structured design of the system to further improve the integrity and flexibility of StepProof proof.
- 
- 6 CONCLUSION
- 
- In this paper, we innovatively propose a new automatic formal proof method called StepProof. Step-Proof implements sentence-level formal verification of natural language mathematical proofs, al-

**432 433 434** lowing users to conduct more flexible formal verification. At the same time, compared with the traditional FullProof strategy, StepProof has been significantly improved in formalization, proof efficiency and proof accuracy.

**435 436 437 438 439** In addition, StepProof can preserve the contents of the proof that has already been verified, providing more information than the traditional Full-Proof strategy that can only indicate the passage and failure of the proof. We used a small model on the GSM8K data set to test the proof pass rate, and its performance reached the level of state-of-the-art.

**440 441 442 443 444** We also conducted a test on the Number Theory dataset of MATH to test the effect of formal content writing on the proof pass rate. The experimental results show that by optimizing the non-formal proof for step verification, the passing rate of the non-formal proof can be significantly improved. We will further optimize the architecture of StepProof in future work to make it more flexible to handle formal verification of various non-formal mathematical proofs.

**445 446**

**453**

#### **REFERENCES**

- <span id="page-8-1"></span>**447 448 449 450** Josh Achiam, Steven Adler, Sandhini Agarwal, Lama Ahmad, Ilge Akkaya, Florencia Leoni Aleman, Diogo Almeida, Janko Altenschmidt, Sam Altman, Shyamal Anadkat, et al. Gpt-4 technical report. *arXiv preprint arXiv:2303.08774*, 2023.
- <span id="page-8-2"></span>**451 452** Andrea Asperti. A survey on interactive theorem proving. *URL: http://www. cs. unibo. it/˜ asperti/SLIDES/itp. pdf*, 2009.
- <span id="page-8-8"></span><span id="page-8-4"></span>**454 455 456 457** Karl Cobbe, Vineet Kosaraju, Mohammad Bavarian, Mark Chen, Heewoo Jun, Lukasz Kaiser, Matthias Plappert, Jerry Tworek, Jacob Hilton, Reiichiro Nakano, Christopher Hesse, and John Schulman. Training verifiers to solve math word problems. *arXiv preprint arXiv:2110.14168*, 2021.
	- Leonardo De Moura, Soonho Kong, Jeremy Avigad, Floris Van Doorn, and Jakob von Raumer. The lean theorem prover (system description). In *Automated Deduction-CADE-25: 25th International Conference on Automated Deduction, Berlin, Germany, August 1-7, 2015, Proceedings 25*, pp. 378–388. Springer, 2015.
- <span id="page-8-5"></span>**463 464 465** Abhimanyu Dubey, Abhinav Jauhri, Abhinav Pandey, Abhishek Kadian, Ahmad Al-Dahle, Aiesha Letman, Akhil Mathur, Alan Schelten, Amy Yang, Angela Fan, et al. The llama 3 herd of models. *arXiv preprint arXiv:2407.21783*, 2024.
- <span id="page-8-7"></span>**466 467 468 469 470** Wenqi Fan, Yujuan Ding, Liangbo Ning, Shijie Wang, Hengyun Li, Dawei Yin, Tat-Seng Chua, and Qing Li. A survey on rag meeting llms: Towards retrieval-augmented large language models. In *Proceedings of the 30th ACM SIGKDD Conference on Knowledge Discovery and Data Mining*, pp. 6491–6501, 2024.
- <span id="page-8-6"></span>**471 472 473 474 475 476 477 478 479** Team GLM, Aohan Zeng, Bin Xu, Bowen Wang, Chenhui Zhang, Da Yin, Diego Rojas, Guanyu Feng, Hanlin Zhao, Hanyu Lai, Hao Yu, Hongning Wang, Jiadai Sun, Jiajie Zhang, Jiale Cheng, Jiayi Gui, Jie Tang, Jing Zhang, Juanzi Li, Lei Zhao, Lindong Wu, Lucen Zhong, Mingdao Liu, Minlie Huang, Peng Zhang, Qinkai Zheng, Rui Lu, Shuaiqi Duan, Shudan Zhang, Shulin Cao, Shuxun Yang, Weng Lam Tam, Wenyi Zhao, Xiao Liu, Xiao Xia, Xiaohan Zhang, Xiaotao Gu, Xin Lv, Xinghan Liu, Xinyi Liu, Xinyue Yang, Xixuan Song, Xunkai Zhang, Yifan An, Yifan Xu, Yilin Niu, Yuantao Yang, Yueyan Li, Yushi Bai, Yuxiao Dong, Zehan Qi, Zhaoyu Wang, Zhen Yang, Zhengxiao Du, Zhenyu Hou, and Zihan Wang. Chatglm: A family of large language models from glm-130b to glm-4 all tools, 2024.
- <span id="page-8-3"></span>**480 481** John Harrison. A survey of automated theorem proving. *Sat*, 17:20–18, 2013.
- <span id="page-8-0"></span>**482 483 484 485** John Harrison, Josef Urban, and Freek Wiedijk. History of interactive theorem proving. In Jörg H. Siekmann (ed.), *Computational Logic*, volume 9 of *Handbook of the History of Logic*, pp. 135–214. North-Holland, 2014. doi: https://doi.org/10.1016/ B978-0-444-51624-4.50004-6. URL [https://www.sciencedirect.com/science/](https://www.sciencedirect.com/science/article/pii/B9780444516244500046) [article/pii/B9780444516244500046](https://www.sciencedirect.com/science/article/pii/B9780444516244500046).

**489**

<span id="page-9-4"></span>**495 496 497**

**504**

<span id="page-9-1"></span>**520**

<span id="page-9-18"></span>

- <span id="page-9-5"></span>**490 491 492** Jie Huang, Xinyun Chen, Swaroop Mishra, Huaixiu Steven Zheng, Adams Wei Yu, Xinying Song, and Denny Zhou. Large language models cannot self-correct reasoning yet. *arXiv preprint arXiv:2310.01798*, 2023.
- <span id="page-9-8"></span>**493 494** Gérard Huet, Gilles Kahn, and Christine Paulin-Mohring. The coq proof assistant a tutorial. *Rapport Technique*, 178, 1997.
- **498** Ziwei Ji, Nayeon Lee, Rita Frieske, Tiezheng Yu, Dan Su, Yan Xu, Etsuko Ishii, Ye Jin Bang, Andrea Madotto, and Pascale Fung. Survey of hallucination in natural language generation. *ACM Computing Surveys*, 55(12):1–38, 2023.
- <span id="page-9-16"></span>**499 500 501** Albert Q Jiang, Sean Welleck, Jin Peng Zhou, Wenda Li, Jiacheng Liu, Mateja Jamnik, Timothee´ Lacroix, Yuhuai Wu, and Guillaume Lample. Draft, sketch, and prove: Guiding formal theorem provers with informal proofs. *arXiv preprint arXiv:2210.12283*, 2022.
- <span id="page-9-10"></span>**502 503** Shalom Lappin. Assessing the strengths and weaknesses of large language models. *Journal of Logic, Language and Information*, 33(1):9–20, 2024.
- <span id="page-9-15"></span>**505 506 507 508** Aitor Lewkowycz, Anders Andreassen, David Dohan, Ethan Dyer, Henryk Michalewski, Vinay Ramasesh, Ambrose Slone, Cem Anil, Imanol Schlag, Theo Gutman-Solo, Yuhuai Wu, Behnam Neyshabur, Guy Gur-Ari, and Vedant Misra. Solving quantitative reasoning problems with language models, 2022. URL <https://arxiv.org/abs/2206.14858>.
- <span id="page-9-6"></span>**509 510** Zhaoyu Li, Jialiang Sun, Logan Murphy, Qidong Su, Zenan Li, Xian Zhang, Kaiyu Yang, and Xujie Si. A survey on deep learning for theorem proving. *arXiv preprint arXiv:2404.09939*, 2024.
- <span id="page-9-2"></span>**511 512 513 514** Pengfei Liu, Weizhe Yuan, Jinlan Fu, Zhengbao Jiang, Hiroaki Hayashi, and Graham Neubig. Pretrain, prompt, and predict: A systematic survey of prompting methods in natural language processing. *ACM Computing Surveys*, 55(9):1–35, 2023.
- <span id="page-9-13"></span>**515 516 517 518** Xuantao Lu, Jingping Liu, Zhouhong Gu, Hanwen Tong, Chenhao Xie, Junyang Huang, Yanghua Xiao, and Wenguang Wang. Parsing natural language into propositional and first-order logic with dual reinforcement learning. In *Proceedings of the 29th International Conference on Computational Linguistics*, pp. 5419–5431, 2022.
- <span id="page-9-0"></span>**519** Filip Maric. A survey of interactive theorem proving. *Zbornik radova*, 18(26):173–223, 2015.
- **521 522** M Saqib Nawaz, Moin Malik, Yi Li, Meng Sun, and M Lali. A survey on theorem provers in formal methods. *arXiv preprint arXiv:1912.03028*, 2019.
- <span id="page-9-11"></span>**523 524 525** Thien Huu Nguyen and Ralph Grishman. Relation extraction: Perspective from convolutional neural networks. In *Proceedings of the 1st workshop on vector space modeling for natural language processing*, pp. 39–48, 2015.
- <span id="page-9-7"></span>**526 527** Lawrence C Paulson. *Isabelle: A generic theorem prover*. Springer, 1994.
- <span id="page-9-17"></span>**528** Boris Shminke. Python client for isabelle server. *arXiv preprint arXiv:2212.11173*, 2022.
- <span id="page-9-12"></span>**529 530 531** Hrituraj Singh, Milan Aggrawal, and Balaji Krishnamurthy. Exploring neural models for parsing natural language into first-order logic. *arXiv preprint arXiv:2002.06544*, 2020.
- <span id="page-9-14"></span>**532 533 534** Qingxiang Wang, Chad Brown, Cezary Kaliszyk, and Josef Urban. Exploration of neural machine translation in autoformalization of mathematics in mizar. In *Proceedings of the 9th ACM SIG-PLAN International Conference on Certified Programs and Proofs*, pp. 85–98, 2020a.
- <span id="page-9-3"></span>**535 536 537** Yaqing Wang, Quanming Yao, James T Kwok, and Lionel M Ni. Generalizing from a few examples: A survey on few-shot learning. *ACM computing surveys (csur)*, 53(3):1–34, 2020b.
- <span id="page-9-9"></span>**538 539** Junjie Ye, Xuanting Chen, Nuo Xu, Can Zu, Zekai Shao, Shichun Liu, Yuhan Cui, Zeyang Zhou, Chao Gong, Yang Shen, et al. A comprehensive capability analysis of gpt-3 and gpt-3.5 series models. *arXiv preprint arXiv:2303.10420*, 2023.
- <span id="page-10-1"></span> Artem Yushkovskiy. Comparison of two theorem provers: Isabelle/hol and coq. *arXiv preprint arXiv:1808.09701*, 2018.
- <span id="page-10-0"></span> Wayne Xin Zhao, Kun Zhou, Junyi Li, Tianyi Tang, Xiaolei Wang, Yupeng Hou, Yingqian Min, Beichen Zhang, Junjie Zhang, Zican Dong, et al. A survey of large language models. *arXiv preprint arXiv:2303.18223*, 2023.

#### <span id="page-10-2"></span> Jin Peng Zhou, Charles Staats, Wenda Li, Christian Szegedy, Kilian Q Weinberger, and Yuhuai Wu. Don't trust: Verify–grounding llm quantitative reasoning with autoformalization. *arXiv preprint arXiv:2403.18120*, 2024.