STACKED FILTERS STATIONARY FLOW FOR HARDWARE-ORIENTED ACCELERATION OF DEEP CONVOLUTIONAL NEURAL NETWORKS

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ABSTRACT
To address memory and computation resource limitations for hardware-oriented acceleration of deep convolutional neural networks (CNNs), we present a computation flow, stacked filters stationary flow (SFS), and a corresponding data encoding format, relative indexed compressed sparse filter format (CSF), to make the best of data sparsity, and simplify data handling at execution time. And we also propose a three dimensional Single Instruction Multiple Data (3D-SIMD) processor architecture which takes full advantage of these two features. Comparing with the state-of-the-art result [Han et al., 2016b], our method achieved 1.11× improvement in reducing the storage required by AlexNet, and 1.09× improvement in reducing the storage required by SqueezeNet, without loss of accuracy on the ImageNet dataset. Moreover, using this approach, chip area for logics handling irregular sparse data access can be saved.

1 INTRODUCTION
CNNs have achieved substantial progress during the past years. But hardware resource limitations have hindered their wide usage in embedded devices. Various efforts have been made to address this issue, such as ShiftCNN [Gudovskiy & Rigazio, 2017], Ristretto [Gysel, 2016], Eyeriss [Chen et al., 2017], Deep Compression [Han et al., 2016b] and EIE [Han et al., 2016a]. Through compressing deep neural networks with pruning, trained quantization and Huffman coding, Deep Compression [Han et al., 2016b], the best paper of ICLR 2016, achieved state-of-the-art result in reducing storage requirement of neural networks without affecting their accuracy.

In spite of the great progress achieved till now, there are still many problems to be solved. The first problem is manipulating compressed sparse data consumes considerable extra clock cycles and need extra logics (about 19.1% chip area in [Han et al., 2016a]). Several sparse matrix encoding formats have been proposed, such as CSC, CSR and CISR [Fowers et al., 2014]. But existing encoding formats complex the computation at runtime due to its irregular memory access characteristics. This results in inefficiency in parallelizing computation and bigger chip area. Therefore, it would be desirable if the encoding can be directly used during run time. The second one is, for deeply compressed sparse networks, the computational efficiency of recently proposed hardware acceleration designs, such as Eyeriss [Chen et al., 2017], DVAS [Moons & Verhelst, 2015], ENVISION [Moons et al., 2017], DNPU [Shin et al., 2017], etc., is fairly low. In this paper we present a novel computation flow SFS, and a corresponding memory layout and encoding format CSF which achieves the desirable goal that it can be directly used at run time. We also propose a 3D-SIMD processor architecture to illustrate how to accelerate deep CNNs by taking advantage of SFS flow and CSF format.

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2 Stack Filters Stationary Flow (SFS) and Relative Indexed Compressed Sparse Filter (CSF) Format

Computations of convolutional and fully connected layers in CNNs can be unified into one formula Eq[1] (ignoring biases). Eq. 2-6 illustrate our approach. \( V_o, V_i \) and \( W_f \) are the matrices of output feature maps, input feature maps and filters, respectively. \( S, C, K, M, M', m \) is a given stride size, channel number, filter kernel size, total filter number, number of batches and batch size (Eq[6]), Filters are firstly grouped into \( M' \) batches (Eq[2]), and each \( W_{f'}^{(n)} \) is then reshaped to \( W_f^{(M')} \) (Eq[5]). One channel of feature data will convolute with \( m \) filters from the same channel at a time (Eq[4]). At the end of computation, \( V_o'^{(n)} \) should be concatenated back to \( V_o \) (Eq[5]).

\[
V_o[cho][x][y] = \sum_{chi=0}^{C-1} \sum_{r=0}^{K-1} \sum_{c=0}^{K-1} W_f[cho][chi][r][c] \times V_i[chi][Sx+r][Sy+c] \tag{1}
\]
\[
W_f = [W_f^{(0)}, W_f^{(1)}, ..., W_f^{(M'-1)}], W_{f'} = [W_f^{(0)}, W_f^{(1)}, ..., W_f^{(M'-1)}] \tag{2}
\]
\[
W_{f'}^{(n)}[chi][r][c][j] = W_f^{(n)}[j][chi][r][c] \tag{3}
\]
\[
V_o' = [V_o'^{(0)}, V_o'^{(1)}, ..., V_o'^{(M'-1)}] \tag{4}
\]
\[
V_o = [V_o^{(0)}, V_o^{(1)}, ..., V_o^{(M'-1)}] \tag{5}
\]
\[
M' = M/m, 0 \leq n < M', 0 \leq j < m, 0 \leq cho < M. \tag{6}
\]

As to encoding format, our approach is to further rearrange the memory layout of the \( m \) filters illustrated in figure [1], storing the elements column by column. So in computation flow SFS, when each element in the feature map multiplies with a column of data from \( m \) filters, the filter weights could be loaded sequentially. The first line in figure [1] illustrates the changing. If there is any weight value equals to 0, just remove the column contains that value, add one to the relative index in the next column, and subtract one to the pointer next to it. The non-zero value number (includes padding zeros) of a column is given by the next pointer. Column pointer is 0 means all the values in the column before the pointer equal to 0. Relative column pointer is not needed when parameters are stored in files.

![Figure 1: Memory layout for m filters with kernel size 3 from a single channel.](image1)

![Figure 2: 3D-SIMD processor architecture.](image2)
SFS flow and CSF encoding format are the key features of the proposed 3D-SIMD processor architecture, see figure 2. In this architecture, after feature data are loaded into the line buffer and window registers from a single channel of input feature map, and m filter data from the same channel are loaded into the local buffer, each element in the window will multiply with a column of data from m filters at the same position. So data in CSF format can be directly used at run time, and all zeros are skipped as designed. No complex sparse data handling logic is required.

4 Result

The distribution of continuous zero numbers after applying the changing is first evaluated. As figure 3 shows, the distribution narrows to the left. It means that fewer bits is needed to store the relative index values, and there will be fewer padding zeros. This will further reduce storage space. The distribution of continuous nonzero numbers after applying the changes is also evaluated. The distribution also narrows to the left. It means the computation load balance during execution will be better comparing to the reference work. The effect of batch size m is also analyzed. It shows that there do exist an optimum batch size for each layer. Smaller batch size requires smaller local buffer, but the input feature map need to be loaded more times. For simplicity, all the experiments in this section use filter number as the batch size. That means, all the filters and input feature maps are loaded only one time and the output feature maps are saved one time for one reference.

Table 1: Extra space (in bits) improvement and total storage requirement improvement

<table>
<thead>
<tr>
<th>Network</th>
<th>Nonzeros</th>
<th>Extra space</th>
<th>Improvement</th>
<th>Total</th>
</tr>
</thead>
<tbody>
<tr>
<td>AlexNet by Han et al. 2016b</td>
<td>30592056</td>
<td>41325956</td>
<td>1.21×</td>
<td>1.11×</td>
</tr>
<tr>
<td>AlexNet by SFS+CSF</td>
<td>30592056</td>
<td>34138970</td>
<td>1.11×</td>
<td>1.09×</td>
</tr>
<tr>
<td>SqueezeNet by Han et al. 2016b</td>
<td>3327368</td>
<td>1737628</td>
<td>1.11×</td>
<td>1.09×</td>
</tr>
<tr>
<td>SqueezeNet by SFS+CSF</td>
<td>3327368</td>
<td>1307160</td>
<td>1.33×</td>
<td>1.09×</td>
</tr>
</tbody>
</table>

SqueezeNet by SFS+CSF 3327368 1307160 1.33× SqueezeNet by (Han et al., 2016b) 3327368 1737628

AlexNet by SFS+CSF 30592056 34138970 1.11× AlexNet by (Han et al., 2016b) 30592056 41325956

Figure 3: Distributions of continuous zero numbers in different layers of Alexnet, comparing with ref (Han et al., 2016b).

Table 1 illustrate the improvement of extra space needed to store index and padding zeros, and the total improvement of storage requirement after applying our method comparing to former works.

5 Conclusion

In this paper, we propose a computation flow SFS, and its corresponding data encoding format CSF. And we also proposed a 3D-SIMD processor architecture to make the best of these two features. Experimental results show that our approach narrowed the distribution of the numbers of continuous zeros and non-zeros to the smaller number direction. This helps to further compress the network parameters by about 8% to 10% and balance computation load at run time. By adopting the proposed 3D-SIMD processor architecture, chip area for logics handling irregular memory access of sparse data can be saved, for example, about 19.1% chip area in EIE (Han et al. 2016a) for pointer read can be saved. Moreover, directly using the encoded data as it is at runtime can also save zero bypassing clock cycles, and data looking up time (Moons & Verhelst 2015), (Moons et al., 2017), (Han et al. 2016a).
REFERENCES


