# On Sparse Canonical Correlation Analysis

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## Abstract

The classical Canonical Correlation Analysis (CCA) identifies the correlations between two sets of multivariate variables based on their covariance, which has been widely applied in diverse fields such as computer vision, natural language processing, and speech analysis. Despite its popularity, CCA can encounter challenges in explaining correlations between two variable sets within high-dimensional data contexts. Thus, this paper studies Sparse Canonical Correlation Analysis (SCCA) that enhances the interpretability of CCA. We first show that SCCA generalizes three well-known sparse optimization problems, sparse PCA, sparse SVD, and sparse regression, which are all classified as NP-hard problems. This result motivates us to develop strong formulations and efficient algorithms. Our main contributions include (i) the introduction of a combinatorial formulation that captures the essence of SCCA and allows the development of approximation algorithms; (ii) the establishment of the complexity results for two low-rank special cases of SCCA; and (iii) the derivation of an equivalent mixed-integer semidefinite programming model that facilitates a specialized branch-and-cut algorithm with analytical cuts. The effectiveness of our proposed formulations and algorithms is validated through numerical experiments.

## 1 Introduction

The Canonical Correlation Analysis (CCA), proposed by H. Hotelling [\[23\]](#page-10-0), aims to identify the correlations between two sets of multivariate variables based on their covariance. Since then, CCA has become a powerful statistical technique used for multivariate data analysis, with its applications across diverse fields such as computer vision [\[24\]](#page-11-0), natural language processing [\[38\]](#page-11-1), and speech analysis [\[21\]](#page-10-1). Despite its popularity, CCA can encounter challenges in explaining correlations between two variable sets within high-dimensional data contexts, such as genomic datasets [\[36\]](#page-11-2). In contrast, Sparse Canonical Correlation Analysis (SCCA), which seeks sparse linear combinations of these variable sets, offers substantially enhanced interpretability [\[41,](#page-11-3) [42,](#page-11-4) [44\]](#page-11-5).

Formally, this paper studies the SCCA problem:

<span id="page-0-0"></span>
$$
v^* := \max_{\boldsymbol{x} \in \mathbb{R}^n, \boldsymbol{y} \in \mathbb{R}^m} \left\{ \boldsymbol{x}^\top \boldsymbol{A} \boldsymbol{y} : \boldsymbol{x}^\top \boldsymbol{B} \boldsymbol{x} \leq 1, \boldsymbol{y}^\top \boldsymbol{C} \boldsymbol{y} \leq 1, \|\boldsymbol{x}\|_0 \leq s_1, \|\boldsymbol{y}\|_0 \leq s_2 \right\},\qquad \text{(SCCA)}
$$

where  $s_1 \le n$ ,  $s_2 \le m$  are positive integers and  $\begin{pmatrix} B & A \\ A^{\top} & C \end{pmatrix}$  $A^\top$   $C$ denotes a covariance matrix of  $(n + m)$ random variables. Specifically,  $B$  and  $C$  are the covariance matrices of the n and m random variables, respectively, and  $\vec{A} \in \mathbb{R}^{n \times m}$  is the cross-covariance matrix between n and m random variables. Hence,  $\begin{pmatrix} B & A \\ A^{\top} & C \end{pmatrix}$  $A^\top$   $C$ , **B**, **C** are positive semidefinite matrices of size  $(n + m)$ , *n*, and *m*, respectively. Here, matrices  $B, C$  can be singular, i.e., some random variables may be dependent on others. In fact, the covariance matrices  $B, C$  are often low-rank, especially within the high-dimension low-sample size data context (see, e.g., the gene expression data in [\[41\]](#page-11-3)).

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The [SCCA](#page-0-0) problem generalizes three widely-studied sparsity-constrained optimization problems as special cases, which are sparse PCA [\[2,](#page-10-2) [13,](#page-10-3) [27\]](#page-11-6), sparse SVD [\[28,](#page-11-7) [41\]](#page-11-3), and sparse regression [\[22,](#page-10-4) [3\]](#page-10-5). To be specific, when  $n = m$ ,  $s_1 = s_2$ ,  $B$ ,  $C$  are identity matrices, and  $A$  is a positive semidefinite matrix, [SCCA](#page-0-0) reduces to the classic sparse PCA problem; when  $B, C$  are identity matrices, SCCA becomes the sparse SVD problem; and when  $\vec{A}$  is rank-one, Section [3](#page-4-0) shows that [SCCA](#page-0-0) is equivalent to two sparse linear regression subproblems.

#### 1.1 Main contributions

[SCCA](#page-0-0) is generally NP-hard, given that its special cases, sparse PCA, sparse SVD, and sparse regression are all classified as NP-hard problems. We are motivated to develop efficient formulations and algorithms for [SCCA](#page-0-0) through a mixed-integer optimization lens. The main contributions, along with the structure of the remainder of this paper, are the following:

- (i) In Section [2,](#page-2-0) we present an exact semidefinite programming (SDP) reformulation and derive a closed-form optimal value of classic CCA problem. We also develop an equivalent combinatorial formulation of [SCCA,](#page-0-0) which allows the development of approximation algorithms;
- (ii) When the covariance matrix  $\begin{pmatrix} B & A \\ A^{\top} & C \end{pmatrix}$  $A^\top$   $C$  is low-rank, Section [3](#page-4-0) studies the complexity of two special cases of [SCCA.](#page-0-0) This motivates us to develop a polynomial-time exact algorithm of complexity  $\mathcal{O}(n^3 + m^3)$  for solving [SCCA](#page-0-0) to global optimiality when the sparsity levels (i.e.,  $s_1$  and  $s_2$ ) meet or exceed the ranks of **B** and **C**;
- (iii) Section [4](#page-6-0) derives an equivalent mixed-integer SDP (MISDP) reformulation for [SCCA.](#page-0-0) When applying the Benders decomposition approach, instead of solving the large-scale SDPs, we design a customized branch-and-cut algorithm with closed-form cuts, which can successfully solve [SCCA](#page-0-0) to optimality; and
- (iv) Section [5](#page-6-1) numerically test the proposed formulations and algorithms. It is noted that our polynomial-time exact algorithm can solve real-world instances with  $n = 19,672$  and  $m =$ 2, 149 variables in seconds, provided that both  $s_1$  and  $s_2$  are at least the ranks of **B** and **C**.

Our analyses and results can be extended to [SCCA](#page-0-0) with multiple pairs of basis vectors  $(x, y)$ , allowing for a more flexible and comprehensive exploration of correlations among data sets. The detailed formulations of multiple [SCCA,](#page-0-0) along with the computational results, are provided in Appendix [F.](#page-20-0)

#### 1.2 Relevant literature

*SCCA.* To the best of our knowledge, the work [\[36\]](#page-11-2) was the first paper that introduced the concept of SCCA to select only small subsets of variables to better explain the relationship between many genetic loci and gene expression phenotypes. A handful subset of features enhances interpretability, a desirable property, especially in complex data analysis, which has been successfully demonstrated in Sparse PCA [\[25\]](#page-11-8). To obtain sparse canonical loadings  $(x, y)$ , [\[39\]](#page-11-9) first applied elastic net penalty to the classical CCA via an iterative regression procedure. In a seminal work on SCCA [\[41\]](#page-11-3), the authors proposed a rigorous formulation by enforcing the  $\ell_1$  constraints on variables  $(x, y)$  and developed a penalized matrix decomposition method to solve the penalized CCA problem. Then, extensive research has focused on various penalty norm functions to obtain sparse canonical loadings (see, e.g., [\[20,](#page-10-6) [26,](#page-11-10) [39,](#page-11-9) [42,](#page-11-4) [10\]](#page-10-7)). In particular, [\[10\]](#page-10-7) penalized multiple canonical loadings by  $\ell_1$  norm and computed the sparse solution by the linearized Bregman method. It should be noted that under the assumption that the leading canonical loadings are sparse, [\[7,](#page-10-8) [17,](#page-10-9) [18\]](#page-10-10) established theoretical guarantees of iterative approaches for estimating sparse solutions. Another research direction in SCCA introduced penalty functions based on group structural information of input data and developed group SCCA methods [\[29,](#page-11-11) [30\]](#page-11-12). For a comprehensive overview of CCA and SCCA methods, we refer readers to the survey by [\[44\]](#page-11-5) and the references therein. These approaches, however, do not strictly enforce the exact sparsity requirement but only approximate the sparsity requirement (i.e., the  $\ell_0$  norm) by a convex function. Another relevant work [\[40\]](#page-11-13) introduced binary variables to recast SCCA as a mixed-integer nonconvex program under the assumption of positive definite matrices  $B, C$ , based on which they designed a branch-and-bound algorithm. Different from the literature, our work does not require positive definiteness assumption of matrices  $B, C$ , and we are able to

obtain mixed-integer conic and semidefinite programming reformulations, allowing for better exact and approximation algorithms.

*Connections to and differences with sparse PCA and sparse SVD.* Analogous to [SCCA,](#page-0-0) both sparse PCA [\[13,](#page-10-3) [25\]](#page-11-8) and sparse SVD [\[28\]](#page-11-7) select small subsets of variables to improve the interpretability of dimensionality reduction methods: PCA and SVD. Considerable investigation has been conducted on solving sparse PCA and sparse SVD from three angles: convex relaxations [\[12](#page-10-11)[–14\]](#page-10-12), approximation algorithms [\[6,](#page-10-13) [9,](#page-10-14) [28\]](#page-11-7), and exact algorithms [\[2,](#page-10-2) [27,](#page-11-6) [28\]](#page-11-7). As mentioned before, in sparse PCA and sparse SVD, the covariance matrices  $B, C$  are identity. Such a setting dramatically simplifies the subset selection problems of sparse PCA and sparse SVD compared to that of SCCA, as in these problems, it suffices to focus on the selection of a submatrix of the matrix  $\vec{A}$ . Specifically, it is shown in  $[11, 27, 35]$  $[11, 27, 35]$  $[11, 27, 35]$  $[11, 27, 35]$  $[11, 27, 35]$  that sparse PCA reduces to selecting a principal submatrix of  $\vec{A}$  to maximize the largest eigenvalue(s) and sparse SVD reduces to selecting a possibly non-symmetric submatrix of  $\bm{A}$ to maximize the largest singular value(s) [\[28\]](#page-11-7). Quite differently, the combinatorial reformulation [\(1\)](#page-3-0) of [SCCA](#page-0-0) aims to simultaneously select a sized-( $s_1 \times s_1$ ) principal submatrix of B, a sized-( $s_2 \times s_2$ ) principal submatrix of C, and a sized-( $s_1 \times s_2$ ) submatrix of A. These fundamental differences in the underlying formulations of sparse PCA and sparse SVD preclude the direct application of their existing algorithms to the [SCCA.](#page-0-0)

Notations: The following notation is used throughout the paper. We use bold lower-case letters (e.g.,  $x$ ) and bold upper-case letters (e.g.,  $X$ ) to denote vectors and matrices, respectively, and we use corresponding non-bold letters (e.g.,  $x_i$ ) to denote their components. We let  $S^n, S^n_+, S^n_{++}$  denote the set of all the  $n \times n$  symmetric real matrices, the set of all the  $n \times n$  symmetric positive semidefinite matrices, and the set of all the  $n \times n$  symmetric positive definite matrices, respectively. We let I denote the identity matrix and let 0 denote the vector or matrix with all-zero entries. We let  $\mathbb{R}^n_+$  denote the set of all *n*-dimensional nonnegative vectors. We let  $[n] = \{1, 2, \dots, n\}$ ,  $[s, n] = \{s, s+1, \dots, n\}$ . Given a matrix  $A \in \mathbb{R}^{n \times m}$  and two subsets  $S \subseteq [n]$ ,  $T \subseteq [m]$ , we let  $A^{\dagger}$  denote the pseudo inverse of matrix A, let  $A_{S,T}$  denote a submatrix of A with rows and columns indexed by sets  $S, T$ , respectively, and let  $(\bm{A}_{S,T})^\dagger$  denote the pseudo inverse of submatrix  $\bm{A}_{S,T}.$  For a set  $S$  and an integer k, we define the set  $S + k = \{i + k | i \in S\}$ . Given a vector  $\boldsymbol{a} \in \mathbb{R}^n$  and a subset  $S \subseteq [n]$ , we let  $\boldsymbol{a}_S$ denote a subvector of  $a$  in the subset S. We define  $[\lambda]_+ = \max\{\lambda, 0\}$ . We let  $\sigma_{\max}(\cdot)$  denote the largest singular value function and let  $\lambda_{\text{max}}(\cdot)$  denote the largest eigenvalue value function.

## <span id="page-2-0"></span>2 A combinatorial reformulation of SCCA

This section introduces an equivalent combinatorial optimization reformulation of [SCCA.](#page-0-0) This reformulation serves as the foundation for developing two effective approximation algorithms.

#### 2.1 An exact semidefinite programming representation of CCA

To begin with, let us focus on the classic CCA problem, which refers to [SCCA](#page-0-0) without zero-norm constraints, as defined below:

<span id="page-2-1"></span>
$$
\max_{\boldsymbol{x}\in\mathbb{R}^n,\boldsymbol{y}\in\mathbb{R}^m}\left\{\boldsymbol{x}^{\top}\boldsymbol{A}\boldsymbol{y}:\boldsymbol{x}^{\top}\boldsymbol{B}\boldsymbol{x}\leq1,\boldsymbol{y}^{\top}\boldsymbol{C}\boldsymbol{y}\leq1\right\}.
$$
 (CCA)

This formulation of [CCA](#page-2-1) can be regarded as a quadratically constrained quadratic program concerning the variables  $\begin{pmatrix} x \\ y \end{pmatrix}$  $\boldsymbol{y}$  $\mathcal{L} \in \mathbb{R}^{n \times m}$ . We next define three-block matrices of size  $(n+m)$  below that aid in the presentation of our results.

<span id="page-2-2"></span>
$$
\tilde{A} = \begin{pmatrix} 0 & A/2 \\ A^\top/2 & 0 \end{pmatrix}, \ \tilde{B} = \begin{pmatrix} B & 0 \\ 0 & 0 \end{pmatrix}, \ \tilde{C} = \begin{pmatrix} 0 & 0 \\ 0 & C \end{pmatrix}.
$$

By introducing a size- $(n + m)$  matrix variable  $X = \begin{pmatrix} x \\ y \end{pmatrix}$  $\boldsymbol{y}$  $\setminus x$  $\boldsymbol{y}$  $\int_0^{\top}$  and removing the rank-one constraint on  $X$ , we can obtain an SDP relaxation of  $(CCA)$ , as described below

$$
\max_{\substack{\boldsymbol{X}\in\mathcal{S}_{+}^{m+n}}}\left\{\text{tr}\left(\tilde{\boldsymbol{A}}\boldsymbol{X}\right):\text{tr}\left(\tilde{\boldsymbol{B}}\boldsymbol{X}\right)\leq 1,\text{tr}\left(\tilde{\boldsymbol{C}}\boldsymbol{X}\right)\leq 1\right\}.
$$
\n(SDP Relaxation)

Next, let us present a key lemma regarding properties of block matrices being positive semidefinite, fundamental for reformulating the [SCCA.](#page-0-0)

<span id="page-3-1"></span>**Lemma 1** ([\[16\]](#page-10-16)) For any symmetric matrix  $\begin{pmatrix} B & A \\ A^T & C \end{pmatrix}$  $A^\top$   $C$  $\Big) \in \mathcal{S}^{n+m}$ , the followings are equivalent:

- *(i) The block matrix is positive semidefinite;*
- *(ii)*  $B \in \mathcal{S}_{+}^{n}$ ,  $(I BB^{\dagger})A = 0$ ,  $C A^{\top}B^{\dagger}A \in \mathcal{S}_{+}^{m}$ ; and
- *(iii)*  $C \in \mathcal{S}_{+}^{m}$ ,  $(I CC^{\dagger})A^{\top} = 0$ ,  $B AC^{\dagger}A^{\top} \in \mathcal{S}_{+}^{n}$ .

Inspired by [Lemma 1,](#page-3-1) we hereby establish the equivalence between [CCA](#page-2-1) and its [SDP Relaxation.](#page-2-2) Remarkably, both of these problems achieve the same optimal value, namely  $\sigma_{\max}(\sqrt{B^\dagger A} \sqrt{C^\dagger})$ .

<span id="page-3-2"></span>Proposition 1 *For the [CCA](#page-2-1) problem, we have the following results.*

- *(i) Both [CCA](#page-2-1) and its [SDP Relaxation](#page-2-2) have an optimal value*  $\sigma_{\text{max}}($ √  $B^{\dagger}A$ √ C†)*;*
- *(ii) A pair of optimal solutions*  $(x^*, y^*)$  *to [CCA](#page-2-1) satisfies*

$$
\boldsymbol{x}^* = \sqrt{\boldsymbol{B}^\dagger} \boldsymbol{q}, \enskip \boldsymbol{y}^* = \sqrt{\boldsymbol{C}^\dagger} \boldsymbol{p},
$$

where  $\bm{q} \in \mathbb{R}^n, \bm{p} \in \mathbb{R}^m$  denote a pair of leading singular vectors of matrix  $\sqrt{\bm{B}^{\dagger}}\bm{A}$ √ C† *; and*

*(iii) An optimal solution* X<sup>∗</sup> *to the [SDP Relaxation](#page-2-2) is*

$$
\boldsymbol{X}^* = \begin{pmatrix} \boldsymbol{x}^* \\ \boldsymbol{y}^* \end{pmatrix} \begin{pmatrix} \boldsymbol{x}^* \\ \boldsymbol{y}^* \end{pmatrix}^\top.
$$

*Proof.* See Appendix [A.1.](#page-12-0) □

[Proposition 1](#page-3-2) motivates the following observation on the optimal values of [CCA](#page-2-1) and [SCCA.](#page-0-0)

Observation 1 *The optimal value of [CCA](#page-2-1) is upper bounded by* 1*, so is the optimal value of [SCCA.](#page-0-0)*

It is noteworthy that the results presented in [Proposition 1](#page-3-2) are established through a distinct methodology. This methodology leverages the positive semidefinite condition of block matrices, as shown in Lemma [1,](#page-3-1) and incorporates duality theory. This approach differs from most prior research [\[31,](#page-11-15) [37,](#page-11-16) [44\]](#page-11-5), which proved Part (i) of [Proposition 1](#page-3-2) by relying on the singular value decomposition and assuming that matrices  $\bm{B}$  and  $\bm{C}$  are positive definite (i.e., full rank). To the best of our knowledge, [\[10\]](#page-10-7) showed parts (i) and (ii) of [Proposition 1](#page-3-2) for a special low-rank [CCA](#page-2-1) problem, where the authors assumed that the covariance matrices are defined as  $A = UV^{\top}$ ,  $B = UU^{\top}$ , and  $C = VV^{\top}$ . Remarkably, [Proposition 1](#page-3-2) extends this result to a more general scenario where  $B$  and  $C$  are not constrained to be strictly positive definite and  $\vec{A}$  is not constrained to directly depend on  $B, C$ , allowing for rank deficiencies and flexible data structure.

#### 2.2 An equivalent formulation of [SCCA](#page-0-0)

In this subsection, we transform [SCCA](#page-0-0) into a combinatorial optimization problem, according to the insights provided by [Proposition 1.](#page-3-2)

<span id="page-3-3"></span>Theorem 1 *[SCCA](#page-0-0) is equivalent to the following combinatorial optimization:*

$$
v^* = \max_{S_1 \subseteq [m], |S_1| \leq s_1, S_2 \subseteq [n], |S_2| \leq s_2} \left\{ \sigma_{\max} \left( \sqrt{\left( \mathbf{B}_{S_1, S_1} \right)^{\dagger}} \mathbf{A}_{S_1, S_2} \sqrt{\left( \mathbf{C}_{S_2, S_2} \right)^{\dagger}} \right) \right\}.
$$
 (1)

*Proof.* See Appendix [A.2.](#page-13-0) □

The combinatorial formulation [\(1\)](#page-3-0) presents significant computational difficulties when attempting to solve [SCCA.](#page-0-0) The primary obstacles are two-fold: first, simultaneously selecting submatrices from the matrices  $A, B, C$  requires a sophisticated optimization across multiple dimensions. Second,

<span id="page-3-0"></span>

the selection criterion is particularly complex, as it involves optimizing the largest singular value of the product of the selected submatrix of  $\boldsymbol{A}$  and the square root of pseudo-inverse submatrices of  $B$  and  $C$ . These complexities necessitate effective optimization solution procedures to address the high-dimensional and non-convex nature of the problem.

Motivated by [Theorem 1,](#page-3-3) we customize the greedy and local search algorithms for SCCA [\(1\)](#page-3-0) that has been widely used in the literature to solve special cases of [SCCA,](#page-0-0) such as sparse PCA and sparse SVD in literature (see, e.g., [\[27,](#page-11-6) [28\]](#page-11-7)). The detailed implementations can be found in Appendix [B.](#page-15-0)

## <span id="page-4-0"></span>3 Low-rank SCCA

In practice, it is common that the sample covariance matrix  $\begin{pmatrix} B & A \\ A^{\top} & C \end{pmatrix}$  $A^\top$   $C$  exhibits low-rank characteristics. This phenomenon is especially prominent when dealing with high-dimensional, low-sample size data, e.g., the real gene expression data in [\[41\]](#page-11-3). In this section, we study two special cases of low-rank [SCCA](#page-0-0) and their computational complexities. Specifically, we develop a polynomial-time exact algorithm of complexity  $\mathcal{O}(n^3 + m^3)$  for solving [SCCA](#page-0-0) to global optimiality when sparsity levels (i.e.,  $s_1$  and  $s_2$ ) exceed or equal the ranks of B and C. Besides, we recast [SCCA](#page-0-0) into mixed-integer convex quadratic programming when matrix  $\vec{A}$  is rank-one.

#### 3.1 Special Case I: [SCCA](#page-0-0) with low-rank covariance matrices

In this section, we show that the computational complexity of [SCCA](#page-0-0) is contingent upon the ranks of the covariance matrices  $B$  and  $C$ . To be more precise, when the sparsity level  $s_1$  (or  $s_2$ ) is equal to or greater than the rank r (or  $\hat{r}$ ) of the covariance matrix B (or C), the imposition of a zero-norm constraint over  $x$  (or  $y$ ) in [SCCA](#page-0-0) becomes redundant. Consequently, lower ranks in the covariance matrices correspond to better computational complexity in solving [SCCA.](#page-0-0)

<span id="page-4-1"></span>**Theorem 2** *Suppose*  $r = \text{rank}(B)$  *and*  $\hat{r} = \text{rank}(C)$ *, then [SCCA](#page-0-0) takes a complexity of*  $\mathcal{O}(n^{r-1}m^{\hat{r}-1}+n^{r-1}+m^{\hat{r}-1})$ . The following results hold:

*(i) When*  $s_1 \geq r$  *and*  $s_2 \geq \hat{r}$ *, the [SCCA](#page-0-0) problem is equivalent to [CCA,](#page-2-1) i.e.,* 

<span id="page-4-2"></span>
$$
v^* = \max_{\boldsymbol{x} \in \mathbb{R}^n, \boldsymbol{y} \in \mathbb{R}^m} \left\{ \boldsymbol{x}^\top \boldsymbol{A} \boldsymbol{y} : \boldsymbol{x}^\top \boldsymbol{B} \boldsymbol{x} \leq 1, \boldsymbol{y}^\top \boldsymbol{C} \boldsymbol{y} \leq 1 \right\};
$$
 (2)

*(ii)* When  $s_1 \geq r$  *and*  $s_2 < \hat{r}$ *, the [SCCA](#page-0-0) problem can be reduced to* 

<span id="page-4-3"></span>
$$
v^* = \max_{\boldsymbol{x} \in \mathbb{R}^n, \boldsymbol{y} \in \mathbb{R}^m} \left\{ \boldsymbol{x}^\top \boldsymbol{A} \boldsymbol{y} : \boldsymbol{x}^\top \boldsymbol{B} \boldsymbol{x} \leq 1, \boldsymbol{y}^\top \boldsymbol{C} \boldsymbol{y} \leq 1, \|\boldsymbol{y}\|_0 \leq s_2 \right\};
$$
(3)

*(iii)* When  $s_1 < r$  *and*  $s_2 \geq \hat{r}$ *, the [SCCA](#page-0-0) problem can be reduced to* 

<span id="page-4-4"></span>
$$
v^* = \max_{\boldsymbol{x} \in \mathbb{R}^n, \boldsymbol{y} \in \mathbb{R}^m} \left\{ \boldsymbol{x}^\top \boldsymbol{A} \boldsymbol{y} : \boldsymbol{x}^\top \boldsymbol{B} \boldsymbol{x} \leq 1, \boldsymbol{y}^\top \boldsymbol{C} \boldsymbol{y} \leq 1, \|\boldsymbol{x}\|_0 \leq s_1 \right\}.
$$
 (4)

*Proof.* See Appendix [A.3.](#page-13-1) □

The results in [Theorem 2](#page-4-1) build on the covariance structure of the data matrix. Specifically, if B and C are of rank r and  $\hat{r}$ , respectively, there are only r and  $\hat{r}$  linearly independent vectors in the subspaces corresponding to  $B$  and  $C$ . Thus, the cosine of the principal angle can always be represented by these r and  $\hat{r}$  vectors. As a result, the canonical directions of CCA consist of only r and  $\hat{r}$  nonzero elements. We further make the following remarks about [Theorem 2:](#page-4-1)

- (i) [Theorem 2](#page-4-1) implies the complexity of solving [SCCA,](#page-0-0) as summarized in the corollary below.
- (ii) Inspired by Part (i) of [Theorem 2,](#page-4-1) we also develop a polynomial-time exact algorithm yielding an optimal solution to [SCCA](#page-0-0) [\(1\)](#page-3-0) when  $s_1 \geq r$  and  $s_2 \geq \hat{r}$ . The detailed implementation can be found in Algorithm [1,](#page-5-0) which successfully solves some large instances with up to  $n = 19,672$ and  $m = 2, 149$  variables in seconds in our numerical experiments; and
- (iii) The proof of [Theorem 2](#page-4-1) implies that [CCA](#page-2-1) always admits an optimal sparse solution  $(x^*, y^*)$ satisfying  $\|\boldsymbol{x}^*\|_0 \leq r$  and  $\|\boldsymbol{y}^*\|_0 \leq \hat{r}$ . We show in [Proposition 1](#page-3-2) that the [SDP Relaxation](#page-2-2) of [CCA](#page-2-1) is exact. Therefore, as a side product, we provide the first-known sufficient condition about (i.e.,  $s_1 \ge r$  and  $s_2 \ge \hat{r}$ ) when the convex [SDP Relaxation](#page-2-2) matches [SCCA.](#page-0-0)

**Corollary 1** *Suppose*  $r = \text{rank}(B)$  *and*  $\hat{r} = \text{rank}(C)$ *. There exists an algorithm that can find an optimal solution to [SCCA](#page-0-0) in*  $\mathcal{O}(n^{r-1}m^{\hat{r}-1})$  *time complexity.* 

**Proposition 2** *Suppose*  $r = \text{rank}(B)$  *and*  $\hat{r} = \text{rank}(C)$ *. Then Algorithm [1](#page-3-0) returns an optimal solution to [SCCA](#page-0-0)* [\(1\)](#page-3-0) *in*  $\mathcal{O}(n^3 + m^3)$  *time complexity when*  $s_1 \geq r$  *and*  $s_2 \geq \hat{r}$ *.* 

*Proof.* Following the proof of [Theorem 2,](#page-4-1) we can show that the output solution of Algorithm [1](#page-3-0) is optimal to [SCCA](#page-0-0) [\(1\)](#page-3-0). In addition, the Step 2 of Algorithm [1](#page-3-0) needs computing the eigendecomposition of matrix  $B\in\mathcal S^n_+$ , which takes a time of  $\mathcal O(n^3)$ . Given a matrix  $\bm Q\in\mathbb R^{n\times (n-r)},$  it also takes a time of  $\mathcal{O}(n^3)$  to find its  $(n - r)$  linearly independent rows at Step 3 through the QR decomposition [\[19\]](#page-10-17). Hence, Algorithm [1](#page-3-0) takes a complexity of  $\mathcal{O}(n^3 + m^3)$ ). □

**Algorithm 1** An exact algorithm for SCCA [\(1\)](#page-3-0) when  $s_1 \geq r$  and  $s_2 \geq \hat{r}$ 

<span id="page-5-0"></span>1: **Input:** Matrices  $A \in \mathbb{R}^{m \times n}$ ,  $B \in \mathcal{S}_{+}^{m}$ ,  $C \in \mathcal{S}_{+}^{m}$  and integers  $s_1 \in [r, n]$ ,  $s_2 \in [\hat{r}, m]$ 

2: Compute the eigenvectors  $Q \in \mathbb{R}^{n \times (n-r)}$  of  $B$  that correspond to its  $(n-r)$  zero eigenvalues,

3: Find  $(n - r)$  linearly independent rows in Q, and collect their indices into a subset  $T_1^* \subseteq [n]$ 

4: Perform the same procedure on matrix C to obtain the subset  $T_2^* \subseteq [m]$ 

5: Define the subsets  $S_1^* = [n] \setminus T_1^*$  and  $S_2^* = [m] \setminus T_2^*$ , and compute

$$
v^{*} = \sigma_{\max} \left( \sqrt{\left( \boldsymbol{B}_{S_1^{*}, S_1^{*}} \right)^{\dagger}} \boldsymbol{A}_{S_1^{*}, S_2^{*}} \sqrt{\left( \boldsymbol{C}_{S_2^{*}, S_2^{*}} \right)^{\dagger}} \right)
$$

6: **Output:** An optimal solution  $(S_1^*, S_2^*)$  and optimal value  $v^*$ 

#### <span id="page-5-4"></span>3.2 Special Case II: [SCCA](#page-0-0) with a rank-one cross-covariance matrix

In this subsection, we study the other interesting low-rank special case of [SCCA](#page-0-0) where the crosscovariance matrix  $\vec{A}$  is rank-one. For this special case, we prove its NP-hardness with reduction to the sparse regression problem. We further demonstrate that rank-one [SCCA](#page-0-0) can be simplified to solving two Mixed-Integer Convex Quadratic Programs (MICQPs), which can be more scalable than directly solving [SCCA.](#page-0-0) Our numerical findings confirm this improved scalability.

We observe that [SCCA](#page-0-0) can be separable over variables  $x$  and  $y$  for the rank-one  $A$ . In fact, suppose that  $A = ab^{\top}$ , then [SCCA](#page-0-0) is equivalent to

<span id="page-5-3"></span>
$$
v^* = \max_{\boldsymbol{x} \in \mathbb{R}^n, \boldsymbol{y} \in \mathbb{R}^m} \left\{ \boldsymbol{x}^\top \boldsymbol{a} \boldsymbol{b}^\top \boldsymbol{y} : \boldsymbol{x}^\top \boldsymbol{B} \boldsymbol{x} \leq 1, \boldsymbol{y}^\top \boldsymbol{C} \boldsymbol{y} \leq 1, \|\boldsymbol{x}\|_0 \leq s_1, \|\boldsymbol{y}\|_0 \leq s_2 \right\}
$$
(5)

which can be equivalently the product of the optimal values of the following two subproblems:

<span id="page-5-1"></span>
$$
v_x = \max_{\boldsymbol{x} \in \mathbb{R}^n} \{ \boldsymbol{a}^\top \boldsymbol{x} : \boldsymbol{x}^\top \boldsymbol{B} \boldsymbol{x} \le 1, \|\boldsymbol{x}\|_0 \le s_1 \},
$$
  
\n
$$
v_y = \max_{\boldsymbol{y} \in \mathbb{R}^m} \{ \boldsymbol{b}^\top \boldsymbol{y} : \boldsymbol{y}^\top \boldsymbol{C} \boldsymbol{y} \le 1, \|\boldsymbol{y}\|_0 \le s_2 \}.
$$
 (6)

That is, the identity  $v^* = v_x v_y$  holds. Next, we show that each subproblem in [\(6\)](#page-5-1) can be reduced to the classic sparse regression problem [\[1,](#page-10-18) [33\]](#page-11-17) and is thus NP-hard as shown below.

## <span id="page-5-2"></span>**Theorem 3** When matrix  $A = a^{\top}b$  is rank-one, each maximization problem in [\(6\)](#page-5-1) is NP-hard.

#### *Proof.* See Appendix [A.4.](#page-14-0) □

Theorem [3](#page-5-2) links the maximization problem [\(6\)](#page-5-1) and the well-known sparse regression problem, implying that even solving the rank-one SCCA problem [\(5\)](#page-5-3) is NP-hard. However, it also motivates us to adapt existing mixed-integer optimization techniques from sparse regression (see, e.g.,  $[1, 4, 43]$  $[1, 4, 43]$  $[1, 4, 43]$  $[1, 4, 43]$  $[1, 4, 43]$ ) to tackle each subproblem in [\(6\)](#page-5-1). By introducing binary variables to model the zero-norm constraint, we derive equivalent MICQP formulations for subproblems [\(6\)](#page-5-1) in Appendix [C.](#page-15-1) There are two types of formulations depending on whether matrices  $\bm{B}$  and  $\bm{C}$  are positive definite, which build on the Big-M and perspective techniques, respectively.

#### <span id="page-6-0"></span>4 Reformulating [SCCA](#page-0-0) as a mixed-integer semidefinite program (MISDP)

While the combinatorial formulation [\(1\)](#page-3-0) is elegant in its structure, it poses significant challenges when attempting to solve it to optimality using branch-and-bound based methods. To fill this gap, in this section, we derive an equivalent MISDP formulation for [SCCA,](#page-0-0) amenable for developing exact methods.

First, it is noted that an optimal solution  $(x^*, y^*)$  to [SCCA](#page-0-0) is always bounded that satisfies  $\|x^*\|_2^2 \le$  $M_1$  and  $||\mathbf{y}^*||_2^2 \leq M_2$ , where we specify the construction of the coefficients  $M_1$  and  $M_2$  in Appendix [C.1.](#page-16-0) Such bounds are essential to the derivation of the MISDP. It is convenient to introduce the following notation about  $\{M_{ii}\}_{i\in[n+m]}\$ :

$$
M_{ii} = M_1, \forall i \in [n], \ M_{ii} = M_2, \forall i \in [n+1, n+m].
$$

<span id="page-6-3"></span>Theorem 4 *The [SCCA](#page-0-0) is equivalent to the following MISDP:*

<span id="page-6-2"></span>
$$
v^= \max_{\mathbf{X} \in \mathcal{S}_+^{n+m}, \mathbf{z} \in \mathcal{Z}} \{ \text{tr}(\tilde{\mathbf{A}}\mathbf{X}) : \text{tr}(\tilde{\mathbf{B}}\mathbf{X}) \le 1, \text{tr}(\tilde{\mathbf{C}}\mathbf{X}) \le 1, X_{ii} \le M_{ii} z_i, \forall i \in [n+m] \}.
$$
 (7)

where the feasible set is defined as  $\mathcal{Z} = \{ \bm{z} \in \{0,1\}^{n+m} : \sum_{i \in [n]} z_i \leq s_1, \sum_{i \in [n+1,n+m]} z_i \leq s_2 \}.$ 

*Proof.* See Appendix [A.5.](#page-15-2) □

<span id="page-6-4"></span>
$$
f_{\rm{max}}
$$

Note that the proposed MISDP formulation [\(7\)](#page-6-2) is of size  $(n+m)\times(n+m)$  since our matrix variable

X replaces 
$$
\begin{pmatrix} x \\ y \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix}^{\top}
$$
 in SCCA.

We have formulated [SCCA](#page-0-0) as a mixed-integer convex optimization problem in [Theorem 4.](#page-6-3) Unfortunately, no commercial solvers can efficiently solve MISDP problems. We derive an equivalent mixed-integer linear program of [SCCA](#page-0-0) with exponentially many linear constraints and an efficient separation oracle based on the approach introduced by [\[15\]](#page-10-20), which allows us to develop a tailored branch-and-cut algorithm. First, by separating the binary variables  $z$ , we rewrite the MISDP [\(7\)](#page-6-2) as

$$
v^* = \max_{\mathbf{z} \in \mathcal{Z}, v} \{ v : v \le f(\mathbf{z}) \},\tag{8}
$$

where the function  $f(z)$  is defined as

<span id="page-6-5"></span>
$$
f(\boldsymbol{z}) = \max_{\boldsymbol{X} \in \mathcal{S}_+^{n+m}} \left\{ \text{tr}(\tilde{\boldsymbol{A}}\boldsymbol{X}) : \text{tr}(\tilde{\boldsymbol{B}}\boldsymbol{X}) \leq 1, \text{tr}(\tilde{\boldsymbol{C}}\boldsymbol{X}) \leq 1, X_{ii} \leq M_{ii} z_i, \forall i \in [n+m] \right\}.
$$
 (9)

For any feasible solution  $\hat{z} \in \mathcal{Z}$  of the problem [\(8\)](#page-6-4), by leveraging the concavity of function  $f(\cdot)$ , the linear inequality

$$
v \leq f(\hat{z}) + \partial f(\hat{z})^\top (z - \hat{z})
$$

cuts off the solution  $\hat{z}$  unless it happened to be optimal in [\(8\)](#page-6-4), which paves the way for a delayed cut-generation procedure within a branch-and-bound framework. As the linear inequality of the above type needs to be added dynamically given different solutions  $\hat{z}$  at each iteration, it calls for an efficient evaluation of function  $f(\hat{z})$  and its subgradient. To speed up the computation, we derive the closed-form expression for both of them. The detailed derivations can be found in Appendix [D.](#page-18-0)

Strategies to improve computational speed in practice: First, we provide a variable-fixing method that can identify some binary variables of the MISDP [\(7\)](#page-6-2) being one at optimality. Removing these pre-selected variables from the feasible set reduces the problem size of [SCCA.](#page-0-0) Second, we enhance the branch-and-cut algorithm with a high-quality warm start solution obtained from the local search algorithm. Third, by relaxing the binary variables in the MISDP [\(7\)](#page-6-2) to be continuous or computing [CCA,](#page-2-1) we can obtain an upper bound of [SCCA,](#page-0-0) and the gap between this bound and the local search output gives an initial gap at the root node. Finally, at each iteration, the branching node is selected based on its potential to decrease the current upper bound instead of random branching.

## <span id="page-6-1"></span>5 Numerical results

This section tests the numerical performance of our formulations and algorithms on synthetic and real data. All the experiments are conducted in Python 3.6 with calls to Gurobi 9.5.2 and MOSEK 10.0.29 on a PC with 10-core CPU, 16-core GPU, and 16GB of memory. The codes and data used in our experiments are available at <https://github.com/yongchunli-13/SCCA.git>.

#### <span id="page-7-2"></span>5.1 Experimental setup

Synthetic data generation: Before we present the empirical results, we first describe the properties of the synthetic data which shall be used throughout this section. By following [\[32\]](#page-11-19), given parameters  $(n, m, s_1, s_2)$ , we first synthetically generate positive definite matrices  $B^* \in S_{++}^n$  and  $C^* \in S_{++}^n$  by  $B^* = \hat{B}\hat{B}^\top + I$  and  $C^* = \hat{C}\hat{C}^\top + I$ , respectively, where  $\hat{B}$  and  $\hat{C}$  consist of elements generated from a normal distribution  $\mathcal{N}(0, 1)$ . Then, we let  $\mathbf{A}^* \in \mathbb{R}^{n \times m} = \lambda \mathbf{B}^* \mathbf{u} \mathbf{v}^\top \mathbf{C}^*$ , where we generate  $\lambda$ uniformly from  $(0, 1)$ , and vectors  $u, v$  are generated from a normal distribution  $\mathcal{N}(0, 1)$  that satisfy  $\|u\|_0 = s_1$ ,  $\|v\|_0 = s_2$ ,  $\pmb{u}^\top \pmb{B}^* \pmb{u} = 1$  and  $\pmb{v}^\top \pmb{C}^* \pmb{v} = 1$ . Next, we sample  $N = 5,000$  data samples from a normal distribution  $\mathcal{N}\left(0, \begin{pmatrix} B^* & A^* \\ (A^*)^\top & C^* \end{pmatrix}\right)$  $\begin{pmatrix} B^* & A^* \ (A^*)^\top & C^* \end{pmatrix}$  and compute their sample covariance matrix to obtain the testing data  $\begin{pmatrix} B & A \\ A^{\top} & C \end{pmatrix}$  $A^\top$   $C$ .

Real data: To obtain a comprehensive understanding of the overall performance of our algorithms, we further conduct experiments on six UCI datasets [\[5\]](#page-10-21) with sizes ranging from 34 to 385 variables. The dataset is split into the first n variables and the remaining  $m$  variables to construct the sample covariance matrices  $A, B, C$ . Besides, we examine the performance of the proposed algorithms on the real breast cancer dataset [\[8\]](#page-10-22) that contain  $n = 19,672$  and  $m = 2,149$  variables. The information on each dataset is summarized in Appendix [E.](#page-20-1)

Throughout, the computational time is in seconds, the time limit is one hour, and the dashed line "-" denotes the unsolved case within the time limit. Note that we let LB denote the lower bound obtained from the approximation algorithm, and we let UB denote the upper bound obtained from convex relaxations of [SCCA.](#page-0-0) Besides, we define  $\text{gap}(\%) = 100 \times (\text{UB} - v^*)/v^*$  to be the optimality gap, and we replace  $v^*$  with the best lower bound when  $v^*$  is not available. We define **MIPGap(%)** to be the gap of exact algorithms at termination. Notably, the complexity analysis of the [SCCA](#page-0-0) problems in Section [3](#page-4-0) indicates that its solution process depends on the ranks of the data matrices. Therefore, we present the numerical results under both full-rank and low-rank cases for a comprehensive evaluation.

#### 5.2 Illustration of the impact sparsity levels on SCCA

In this subsection, we apply the local search algorithm to evaluating the performance of [SCCA](#page-0-0) against different sparsity levels  $s_1$  and  $s_2$ . Specifically, for a given dataset, we compute the ratio of correlations between [SCCA](#page-0-0) and [CCA](#page-2-1) for various  $s_1, s_2$  parameters. We test the real UCI data and synthetic data, and the results are displayed in [Figure 1](#page-7-0) and [Figure 2.](#page-7-1) This visualization provides insights for the maximum sparsity [SCCA](#page-0-0) can achieve while maintaining the correlation of full data. For the real UCI data, [SCCA](#page-0-0) almost recovers the correlation of [CCA](#page-2-1) when  $s_1 \approx n/2$  and  $s_2 \approx m/2$ .

<span id="page-7-1"></span>

<span id="page-7-0"></span>Figure 1: On UCI data with  $n, m = 28, 29$  Figure 2: On synthetic data with  $n, m = 80, 80$ 

#### 5.3 Solving SCCA with full-rank matrices

The numerical results on synthetic and real data are presented in [Table 1](#page-8-0) and [Table 2,](#page-8-1) respectively, which include multiple instances with various parameters  $(n, m, s_1, s_2)$ . First, we observe that the greedy and local search algorithms are scalable, and their outputs match the optimal values for most solved testing cases. That is, they achieve zero optimality gaps on these cases. In the "Convex relaxation" column, we compute an upper bound by solving either the continuous relaxation of MISDP [\(7\)](#page-6-2) or [CCA,](#page-2-1) and we use [CCA](#page-2-1) for  $n \ge 40$  and  $m \ge 40$  cases for efficiency. It is seen that the upper bound maintains an optimality gap at most 2.78%. Then, we apply the branch-and-cut algorithm to solve [SCCA](#page-0-0) to optimality, which can handle the case up to a size of  $n = m = 120$  in

[Table 1.](#page-8-0) The unsolved case in [Table 1](#page-8-0) may be because the initial gap is weak at the root node, implying that the branch-and-cut algorithm explores a considerable amount of nodes before termination (see, e.g.,  $n = m = 80$  and  $s_1 = s_2 = 10$ . Hence, the branch-and-cut algorithm may struggle with large-scale instances with weak initial gaps.

				Greedy			Local search		Convex relaxation		Branch-and-cut			
$\, n$	$\,m$	S <sub>1</sub>	$s_2$	LB	time	LB	time	UВ	$gap(\%)$	time	$v^*$	$\overline{\text{MIPGap}(\%)}$	time	
20	20			$0.244$ $0.01$		0.244		$0.02\,10.256$	.23		0.244	0.00	9	
20	20	10	10	$0.275$ $0.02$   $0.275$				$0.04 \mid 0.278$	1.23		1 0.275	0.00	4	
40	40	5	5	$0.695$ $0.03$   0.695				0.0510.701	0.83		110.695	0.00		
40	40	10		10 0.705 0.06 0.705				$0.12\,10.708$	0.45		110.705	0.00	7	
60	60	5	5	$0.885$ $0.0410.885$				0.09 0.887	0.28		1 0.885	0.00		
60	60	10		10 0.884 0.09 0.884				0.1910.887	0.28		1 0.884	0.00	8	
80	80	5.	5.	$0.633$ $0.06$   0.633				0.1210.650	2.78		1 0.633		0.00 1705	
80	80	10		10 0.631 0.13 0.631				$0.26 \mid 0.644$	2.02		110.643	1.85		
100	100	5	5	$0.942$ $0.0910.942$				$0.16 \mid 0.944$	0.23		1 0.942	0.00	$\overline{4}$	
100	100	10		$10 0.940 \t0.17 0.940$				$0.34 \mid 0.942$	0.23		1 0.940	0.00	15	
120	120	5	5	$0.845$ 0.11		0.845		$0.27 \mid 0.853$	0.97		110.845	0.00	924	
	120 120	-10		10 0.848 0.21		0.848		$0.43 \mid 0.856$	0.85		10.855	0.80		

<span id="page-8-0"></span>Table 1: Evaluation of algorithms on synthetic data

<span id="page-8-1"></span>Table 2: Evaluation of algorithms on six UCI datasets

				Greedy			Local search		Convex relaxation		Branch-and-cut			
$\, n$	$m\,$	$s_1$	$s_{2}$	LB	time	LB	time	$\hat{v}$	$gap(\%)$	time	$v^*$	$\overline{\text{MIPGap}(\%)}$	time	
17	17		$\overline{\mathcal{L}}$	$0.970$ 0.01		0.971	0.04	0.984	1.35		0.980	0.00		
17	17	10	10 I	$0.981$ $0.02$   0.983			0.14	0.984	0.09		0.983	0.00		
28	29	5.	5	$0.761$ $0.02$   0.761			0.05	0.769	1.11		0.761	0.00	6	
28	29	10	10 I	$0.766$ $0.0410.766$			0.071	0.769	0.45		0.766	0.00	22	
32	32	5	5	$0.991$ $0.02$   0.991			0.04	0.993	0.17		0.991	0.00		
32	32	10	10	$0.992$ $0.04$		0.992	0.12	0.993	0.05		0.992	0.00		
38	39	5	5		1 0.02	0.02	0.06	0.16	$0.00\ 0.01$			0.00	2	
38	39	10	10		0.04	0.05	0.07	0.47	$0.00\ 0.01$			0.00	$\mathfrak{D}$	
64	64	5	5	$0.998$ $0.05$		0.998	0.39	0.999	0.07		0.998	0.00		
64	64	10	10 I	0.99900.09		0.999	0.59	0.999	0.02		0.999	0.00	$\mathfrak{D}_{\mathfrak{p}}$	
192	193	5	5		0.21		0.46		0.00			0.00	36	
	192 193	10	10		0.36		0.68		0.00			0.00	37	

#### 5.4 Solving SCCA with low-rank matrices

Despite the high dimensions of the breast cancer dataset [\[8\]](#page-10-22), the resulting sample covariance matrices B and C have a rank of 89, i.e.,  $r = \hat{r} = 89$ . When  $s_1 \ge r$  $s_1 \ge r$  $s_1 \ge r$  and  $s_2 \ge \hat{r}$ , we apply Algorithm 1 to solving [SCCA](#page-0-0) that returns optimal solutions and values in one second, as shown in [Table 3.](#page-8-2) If  $s_1 < r$ and  $s_2 < \hat{r}$ , Algorithm [1](#page-5-0) cannot applied and the proposed branch-and-cut algorithm is hard to scale to this dataset. Hence, we only consider using approximation algorithms to solve [SCCA](#page-0-0) for small  $s_1$  and  $s_2$  cases in [Table 4.](#page-8-3) To be specific, we randomly sample n and m variables from the breast cancer data to construct testing cases. As displayed in [Table 4,](#page-8-3) the greedy and local search algorithms run fast, and the local search algorithm slightly outperforms the greedy output.

<span id="page-8-2"></span>Table 3: Solving SCCA by Algorithm [1](#page-5-0) on breast cancer data when  $s_1 \geq r$  and  $s_2 \geq \hat{r}$ 

<span id="page-8-3"></span>Table 4: Approximation algorithms on breast cancer data when  $s_1 \leq r$  and  $s_2 \leq \hat{r}$ 

			Algorithm 1					Greedy	Local search		
$\it n$ m	S <sub>1</sub>		$s_2$  v* MIPGap(%) time		$n_{\rm}$	m			$s_1$ $s_2$ LB time LB		time
		100 100	0.00		100.					$100$ 10 10 0.983 0.17 0.985	0.63
$19,672$ 2,149 $\frac{150}{200}$ 200			0.00						500 500 10 10 0.991	1 0.993	h
			0.00						800 800 10 10 0.993	5 0.993	37
		250 250	0.00		1000 1000 20 20						146

The SCCA  $(5)$  problem with a rank-one matrix  $\vec{A}$  can be more tractable, as we can equivalently decompose it into two MICQPs (see Appendix [C\)](#page-15-1). By approximating  $A$  with a rank-one matrix

including leading singular value and vectors, [Table 5](#page-9-0) presents the numerical results for solving rank-one SCCA [\(5\)](#page-5-3). The continuous relaxation of the MICQPs also provides an upper bound and is denoted by the **Perspective relaxation** in [Table 5.](#page-9-0) We see that the perspective relaxation is computationally efficient and yields small optimality gaps. Besides, we can directly solve two MICQPs below via Gurobi to find the optimal value of rank-one SCCA [\(5\)](#page-5-3), i.e.,  $v^* = v_x v_y$ , where the performance can be found in the last column of [Table 5.](#page-9-0) We can address the rank-one SCCA [\(5\)](#page-5-3) problem up to size  $200 \times 200$  within one hour, improving the problem-solving capacity compared to the size-120  $\times$  120 full-rank SCCA [\(5\)](#page-5-3) in [Table 1.](#page-8-0) However, it should be pointed out that the [SCCA](#page-0-0) may not be mixed-integer convex quadratic representable in general.

Greedy Local Search Perspective relaxation SCCA [\(5\)](#page-5-3)<br>LB time LB time UB gap(%) time  $v^*$  MIPGap(%)  $\begin{array}{|l|c|c|c|c|c|c|c|c|} \hline n & m & s_1 & s_2 & \text{LB} & \text{time} & \text{LB} & \text{time} & \text{UB} & \text{gap}(\%) & \text{time} & v^* \ \hline 50 & 50 & 5 & 5 & 0.753 & 0.04 & 0.753 & 0.07 & 0.757 & 0.57 & 1 & 0.753 \hline \end{array}$  $MIPGap(\%)$  time 50 50 5 5 0.753 0.04 0.753 0.07 0.757 0.57 1 0.753 0.00 1 50 50 10 10 0.753 0.07 0.753 0.15 0.757 0.46 1 0.753 0.00 9<br>100 100 5 5 0.975 0.09 0.975 0.16 0.977 0.21 1 0.975 0.00 2  $\begin{array}{cccc} |0.975 \ 0.09 \ 0.975 \ 0.17 \ 0.966 \ 0.33 \ 0.969 \ 0.35 \end{array}$   $\begin{array}{cccc} 0.21 & 1 & 0.975 & 0.00 & 2 \\ 1 & 0.966 & 0.00 & 25 \end{array}$  $100 \t100 \t10 \t10 \t0.966 \t0.17 \t0.966 \t0.33 \t0.969 \t0.35 \t1 \t0.966$ <br>150 150 5 5 0.850 0.15 0.850 0.26 0.859 1.08 1 0.850  $15 | 0.850 | 0.15 | 0.850 | 0.26 | 0.859 | 1.08 1 | 0.850 | 0.00 9$ 150 150 10 10 0.857 0.27 0.857 0.53 0.867 1.13 1 0.857 0.00 167 200 200 5 5  $|0.810 \t0.23|0.810 \t0.37|0.828 \t2.19$  2  $|0.810 \t0.00 \t55$ <br>200 200 10 10 0.816 0.39 0.816 0.74 0.833 2.11 2 0.816 0.00 1692 200 200 10 10 0.816 0.39 0.816 0.74 0.833 2.11 2 0.816 0.00 1692

<span id="page-9-0"></span>Table 5: Solving SCCA on synthetic data with a rank-one matrix  $\boldsymbol{A}$ 

#### 5.5 Experimental comparison of SCCA algorithms

This section compares the proposed local search algorithm with the SCCA methods of [\[10,](#page-10-7) [37,](#page-11-16) [41\]](#page-11-3) in correlation (Corr) value, the zero norm of x (denoted by S.x), the zero norm of y (denoted by  $S(y)$ , and running time. The computational results on synthetic, UCI, and breast cancer data are presented in [Table 6,](#page-9-1) where we highlight the best correlation and sparsity results in bold. Unlike the local search algorithm, these existing methods do not strictly enforce the exact sparsity requirement, i.e., the zero-norm constraints on variables  $x, y$ . Consequently, the local search algorithm achieves the best sparsity for nearly all testing cases. More importantly, the local search algorithm yields a larger correlation value than these existing methods in 16 out of 22 testing cases. Finally, the running time of the local search algorithm dominates that of [\[10,](#page-10-7) [37\]](#page-11-16).

<span id="page-9-1"></span>Table 6: Comparison of [SCCA](#page-0-0) algorithms in Correlation, sparsity, and time

					Local search				[41]			$\overline{37}$			$\overline{110}$			
$\it{n}$	$\boldsymbol{m}$		$S_1$ $S_2$					Corr S.x S.y time Corr S.x S.y time Corr S.x S.y time Corr S.x S.y time										
20	20	$\overline{5}$		5 0.244	5			50.04 0.200	9		11 0.01 0.239	13	14		2 0.256	18	16	15
20	20	10		10 0.275	10			100.0610.212	7		8 0.01 0.259	19	14	$\mathcal{D}_{\mathcal{L}}$	0.278	18	18	12
40	40	5		5 0.695	5			5 0.08 0.594	23		30 0.03 0.659	26	19		2 0.696	17	18	13
40	40	10		10 0.705	10			10 0.17 0.597	17		15 0.01 0.660	19	22	$\mathcal{D}_{\mathcal{L}}$	0.704	16	20	15
60	60	5		5 0.885	5			5 0.13 0.794	37		45 0.01 0.865	27	25		3 0.880	10	8	23
60	60	10		10 0.884	10			10 $0.2710.777$	33		36 0.01 0.847	31	31		3 0.879	12	15	23
80	80	5.		5 0.633	5			5 0.17 0.534	55		43 0.03 0.606	36	32		3 0.634	33	23	34
80	80	10		10 0.631	10			10 0.37 0.528	53		38 0.02 0.603	29	41		3 0.632	36	41	32
100	100	5		5 0.942	5			5 0.22 0.813	58		63 0.02 0.902	46	33		4 0.941	5	.5	38
100	100	10		10 0.940	10			10 0.49 0.768	54		55 0.02 0.884	47	37		4 0.938	10	10	37
120	120	5		5 0.845	5			5 0.42 0.681	83		74 0.02 0.804	37	43		4 0.833	10	8	47
120	120			10 10 0.848	10			10 $0.59 0.720$	71		76 0.02 0.821	39	36		4 0.840	13	12	47
$\overline{17}$	17	5	$\overline{\mathcal{F}}$	0.971	5			5 0.04 0.794	6		9 0.01 0.742	10	$\overline{8}$		2 0.970	14	16	$\overline{15}$
28	29	5	5	0.761	5			5 0.05 0.704	23		29 0.01 0.667	3	$\overline{7}$	$\mathcal{D}_{\mathcal{L}}$	0.744	24	19	15
32	32	5	5	0.991	5			50.04(0.730)	31		23 0.01 0.904	15	10	$\overline{c}$	0.906	18	15	43
38	39	5	5	1	5			5 0.15 0.994	8		9 0.43 0.997	38	23	3		24	32	29
64	64	5	5	0.998	5			5 0.38 0.897	64		64 0.01 0.990	7	6	$\overline{c}$	0.997	33	31	19
192	193	5	5	1	5			5 0.45 0.103	82		50 0.21 0.856	13	12	10	1	11	6	45
100	100			10 10 0.985	10			10 0.63 0.936	18		41 0.01 0.686	42	43		4 0.952	$\overline{75}$	76	$\overline{36}$
500	500			10 10 0.993	10	10		60.977			78 200 0.78 0.871 173 371				31 0.935	74	44	50
800	800			10100.993	10	10		37 0.974 133 316 1.94 0.879 280 526							50 0.990	84	63	57
1000				1000 10 10 0.995	10	10		11 0.980 166 401 7.30 0.848 401 558							58 0.985	80	59	60

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## Appendices

## A Proofs of technical results

#### <span id="page-12-0"></span>A.1 Proof of [Proposition 1](#page-3-2)

*Proof.* The proof includes three parts.

Part (i). To prove the equivalence between [CCA](#page-2-1) and its [SDP Relaxation,](#page-2-2) let us introduce the Lagrangian multiplies  $\theta_1 \geq 0, \theta_2 \geq 0$  corresponding to two constraints in [SDP Relaxation,](#page-2-2) which leads to the following Lagrangian dual problem

$$
\min_{\theta_1 \ge 0, \theta_2 \ge 0} \left\{ \theta_1 + \theta_2 : \theta_1 \tilde{B} + \theta_2 \tilde{C} \succeq \tilde{A} \right\} = \min_{\theta_1 \ge 0, \theta_2 \ge 0} \left\{ \theta_1 + \theta_2 : \begin{pmatrix} \theta_1 B & \frac{A}{-2} \\ \frac{A^{\top}}{-2} & \theta_2 C \end{pmatrix} \succeq 0 \right\} \tag{10}
$$

where the equation results from the definition of block matrices  $\vec{A}, \vec{B}$ , and  $\vec{C}$ . Given the nonzero matrices  $A \neq 0, B \neq 0, C \neq 0$  and positive semidefinite matrices  $B \succeq 0, C \succeq 0$ , following [Lemma 1,](#page-3-1) we must have  $\theta_2 C - A^{\top} (\theta_1 B)^{\dagger} A/4 \succeq 0$  and  $\theta_1 B - A(\theta_2 C)^{\dagger} A^{\top}/4 \succeq 0$ , implying that either  $\theta_1 = 0$  or  $\theta_2 = 0$  is infeasible to the minimization problem above. That is,  $\theta_1 > 0$  and  $\theta_2 > 0$ must hold.

According to [Lemma 1,](#page-3-1) the block matrix  $\begin{pmatrix} B & A \\ A^{\top} & C \end{pmatrix}$  $A^\top$   $C$ ) is positive semidefinite, implying that  $(I CC^{\dagger}$ ) $A^{\top} = 0$ ,  $(I - BB^{\dagger})A = 0$ . Then, it is easy to show

<span id="page-12-1"></span>
$$
\left(\boldsymbol{I}-\theta_2\boldsymbol{C}(\theta_2\boldsymbol{C})^{\dagger}\right)\frac{\boldsymbol{A}^{\top}}{2}=\boldsymbol{0},\forall\theta_2>0.
$$

Given  $\theta_1, \theta_2 > 0$  and using [Lemma 1,](#page-3-1) the result above allows us to further simplify the right-hand side minimization problem in [\(10\)](#page-12-1) to

$$
\min_{\theta_1 \ge 0, \theta_2 \ge 0} \left\{ \theta_1 + \theta_2 : 4\theta_1 \theta_2 \mathbf{B} \succeq \mathbf{A} \mathbf{C}^\dagger \mathbf{A}^\top \right\} \n= \min_{\theta_1 \ge 0, \theta_2 \ge 0} \left\{ \theta_1 + \theta_2 : 4\theta_1 \theta_2 \ge \sigma_{\text{max}}^2 \left( \sqrt{\mathbf{B}^\dagger} \mathbf{A} \sqrt{\mathbf{C}^\dagger} \right) \right\} = \sigma_{\text{max}} \left( \sqrt{\mathbf{B}^\dagger} \mathbf{A} \sqrt{\mathbf{C}^\dagger} \right),
$$

where the first equation is because

$$
\begin{aligned} &4\theta_1\theta_2\bm{B}\succeq \bm{A}\bm{C}^\dagger\bm{A}^\top \Longleftrightarrow 4\theta_1\theta_2\bm{I}\succeq \sqrt{\bm{\Lambda}^{-1}}\bm{Q}^\top \bm{A}\bm{C}^\dagger\bm{A}^\top\bm{Q}\sqrt{\bm{\Lambda}^{-1}}\\ &\Longleftrightarrow 4\theta_1\theta_2\geq \lambda_{\max}\left(\sqrt{\bm{\Lambda}^{-1}}\bm{Q}^\top \bm{A}\bm{C}^\dagger\bm{A}^\top\bm{Q}\sqrt{\bm{\Lambda}^{-1}}\right)\\ &\Longleftrightarrow 4\theta_1\theta_2\geq \lambda_{\max}\left(\sqrt{\bm{C}^\dagger}\bm{A}^\top\bm{B}^\dagger\bm{A}\sqrt{\bm{C}^\dagger}\right)\Longleftrightarrow 4\theta_1\theta_2\geq \sigma_{\max}^2\left(\sqrt{\bm{B}^\dagger}\bm{A}\sqrt{\bm{C}^\dagger}\right), \end{aligned}
$$

where we let  $B = Q\Lambda Q^{\top}$  denote the eigendecomposition of matrix B with  $\Lambda$  containing all the positive eigenvalues.

As a result, the dual problem of [SDP Relaxation](#page-2-2) admits an optimal value of  $\sigma_{\rm max}\left(\sqrt{\pmb B^\dagger}\pmb A\right)$ √  $\overline{C^\dagger}\big),$ which gives an upper bound of the [CCA](#page-2-1) and its [SDP Relaxation.](#page-2-2) Next, we construct their optimal solutions, which exactly attain this upper bound. Thus, this upper bound is achievable and equals their optimal values.

Part (ii). For the [CCA,](#page-2-1) let us consider a part of optimal solutions  $(x^*, y^*)$  below

$$
\boldsymbol{x}^* = \sqrt{\boldsymbol{B}^{\dagger}} \boldsymbol{q}, \ \ \boldsymbol{y}^* = \sqrt{\boldsymbol{C}^{\dagger}} \boldsymbol{p},
$$

with  $q \in \mathbb{R}^n$ ,  $p \in \mathbb{R}^m$  denoting a pair of leading singular vectors of matrix  $\sqrt{B^{\dagger}}A$ √  $C^{\dagger}.$ First,  $(x^*, y^*)$  is feasible to the [CCA](#page-2-1) as

$$
(\boldsymbol{x}^*)^{\top} \boldsymbol{B} \boldsymbol{x}^* = \boldsymbol{q}^{\top} \sqrt{\boldsymbol{B}^{\dagger}} \boldsymbol{B} \sqrt{\boldsymbol{B}^{\dagger}} \boldsymbol{q} \leq \boldsymbol{q}^{\top} \boldsymbol{q} = 1, \quad (\boldsymbol{y}^*)^{\top} \boldsymbol{C} \boldsymbol{y}^* = \boldsymbol{p}^{\top} \sqrt{\boldsymbol{C}^{\dagger}} \boldsymbol{C} \sqrt{\boldsymbol{C}^{\dagger}} \boldsymbol{p} \leq \boldsymbol{p}^{\top} \boldsymbol{p} = 1,
$$

where the inequalities stem from the facts that  $I \succeq$  $B^\dagger B$  $\bm{B}^\dagger$  and  $\bm{I} \succeq$  $C^\dagger C$  $C^{\dagger}.$ 

On the other hand, according to the definitions of q, p, we can show that  $(x^*, y^*)$  is optimal to the [CCA,](#page-2-1) i.e., √ √ √

$$
(\boldsymbol{x}^*)^\top \boldsymbol{A} \boldsymbol{y}^* = \boldsymbol{q}^\top \sqrt{\boldsymbol{B}^{\dagger}} \boldsymbol{A} \sqrt{\boldsymbol{C}^{\dagger}} \boldsymbol{p} = \sigma_{\max} \left( \sqrt{\boldsymbol{B}^{\dagger}} \boldsymbol{A} \sqrt{\boldsymbol{C}^{\dagger}} \right).
$$

**Part (iii).** In a similar vein, we can show that  $X^* = \begin{pmatrix} x^* & 0 \\ 0 & x^* \end{pmatrix}$  $\left(\begin{matrix} x^* \ y^* \end{matrix}\right) \left(\begin{matrix} x^* \ y^* \end{matrix}\right)$  $\begin{bmatrix} x^* \ y^* \end{bmatrix}^\top$  is optimal to [SDP Relaxation](#page-2-2) with the optimal value  $\sigma_{\text{max}} \left( \sqrt{B^{\dagger}} A \right)$ √  $\overline{C^\dagger}\big)$ . □

#### <span id="page-13-0"></span>A.2 Proof of [Theorem 1](#page-3-3)

*Proof.* By introducing the subsets  $(S_1, S_2)$  to denote the supports of variables  $(x, y)$  in [SCCA,](#page-0-0) then we can remove the zero-norm constraints on  $(x, y)$  and reformulate [SCCA](#page-0-0) as

$$
v^* = \max_{\substack{S_1 \subseteq [m], |S_1| \leq s_1, \boldsymbol{x} \in \mathbb{R}^{|S_1|}, \ S_2 \subseteq [n], |S_2| \leq s_2}} \left\{ \boldsymbol{x}^\top \boldsymbol{A}_{S_1, S_2} \boldsymbol{y} : \boldsymbol{x}^\top \boldsymbol{B}_{S_1, S_1} \boldsymbol{x} \leq 1, \boldsymbol{y}^\top \boldsymbol{C}_{S_2, S_2} \boldsymbol{y} \leq 1 \right\}. \tag{11}
$$

Following from the Part (i) in [Proposition 1,](#page-3-2) we can show that for any subsets  $S_1 \subseteq [n], S_2 \subseteq [m]$ , the following identity holds.

<span id="page-13-2"></span>
$$
\max_{\boldsymbol{x}\in\mathbb{R}^{|S_1|},\boldsymbol{y}\in\mathbb{R}^{|S_2|}}\left\{\boldsymbol{x}^\top\boldsymbol{A}_{S_1,S_2}\boldsymbol{y}:\boldsymbol{x}^\top\boldsymbol{B}_{S_1,S_1}\boldsymbol{x}\leq 1,\boldsymbol{y}^\top\boldsymbol{C}_{S_2,S_2}\boldsymbol{y}\leq 1\right\}\\=\sigma_{\max}\left(\sqrt{(\boldsymbol{B}_{S_1,S_1})^\dagger}\boldsymbol{A}_{S_1,S_2}\sqrt{(\boldsymbol{C}_{S_2,S_2})^\dagger}\right).
$$

Plugging the result above into the inner maximization problem in [\(11\)](#page-13-2), we complete the proof.  $\square$ 

#### <span id="page-13-1"></span>A.3 Proof of [Theorem 2](#page-4-1)

*Proof.* The proof is split into three parts.

**Part (i).** It suffices to prove that [CCA](#page-2-1) admits an optimal solution  $(x^*, y^*)$  satisfying  $\|x^*\|_0 \leq r$  and  $||y^*||_0 \leq \hat{r}$ . Then,  $(x^*, y^*)$  is also feasible and optimal to [SCCA,](#page-0-0) which implies the equivalence between [SCCA](#page-0-0) and [CCA.](#page-2-1)

First, according to Part (ii) in [Proposition 1,](#page-3-2) we can obtain a closed-form optimal solution  $(\hat{x}, \hat{y})$  for the [CCA.](#page-2-1) By adjusting  $(\hat{x}, \hat{y})$ , we will construct a new optimal sparse solution  $(x^*, y^*)$  satisfying  $||\boldsymbol{x}^*||_0 \leq r$  and  $||\boldsymbol{y}^*||_0 \leq \hat{r}$ .

For matrix  $B \in \mathcal{S}_{+}^n$ , we let  $\{q_i\}_{i \in [n-r]} \in \mathbb{R}^n$  denote the eigenvectors corresponding to  $(n-r)$  zero eigenvalues of B. Thus,  $\{q_i\}_{i\in[n-r]}$  are orthonormal. There exists a size- $(n-r)$  subset  $S \subseteq [n]$ such that the subvectors  $\{(q_i)_S\}_{i\in[n-r]}$  are linearly independent, where  $(q_i)_S$  denotes the subvector of  $q_i$  indexed by S for each  $i \in [n-r]$ . As a result, there exist a vector  $(\gamma_1, \dots, \gamma_{n-r})^\top$  such that

<span id="page-13-3"></span>
$$
\hat{\boldsymbol{x}}_S = \sum_{i \in [n]} \gamma_i(\boldsymbol{q}_i)_S. \tag{12}
$$

Let us now construct solution  $x^*$ 

$$
\boldsymbol{x}^* = \hat{\boldsymbol{x}} - \sum_{i \in [n-r]} \gamma_i \boldsymbol{q}_i,
$$

where  $x_i^* = 0$  for all  $i \in S$  based on the equation [\(12\)](#page-13-3) and  $|S| = n - r$ , implying  $||x^*||_0 \le r$ . In addition, we show that the new solution  $x^*$  is still optimal to [CCA.](#page-2-1) First,  $x^*$  is feasible since

$$
(\boldsymbol{x}^*)^\top \boldsymbol{B} (\boldsymbol{x}^*) = \hat{\boldsymbol{x}}^\top \boldsymbol{B} \hat{\boldsymbol{x}} \leq 1,
$$

where the equation is due to  $Bq_i = 0$  for all  $i \in [n - r]$ . Given the positive semidefinite block matrix  $\begin{pmatrix} B & A \\ A^T & C \end{pmatrix}$  $A^\top$   $C$ ), using Part (ii) of [Lemma 1,](#page-3-1) the identity  $(I - BB^{\dagger})A = 0$  is equivalent to  $\sum_{i \in [n-r]} q_i q_i^{\top} A = 0$ . Then, for each  $i \in [n-r]$ , multiplying  $q_i^{\top}$  on both sides of this equation leads to

$$
\boldsymbol{q}_i^\top \bigg(\sum_{j\in [n-r]} \boldsymbol{q}_j \boldsymbol{q}_j^\top \boldsymbol{A}\bigg) \boldsymbol{A} = \boldsymbol{q}_i^\top \boldsymbol{0} \Longrightarrow \boldsymbol{q}_i^\top \boldsymbol{A} = \boldsymbol{0},
$$

where the result follows from  $q_i^{\top}q_j = 0$  for any  $i \neq j$ . Then, we can show the optimality of the new solution  $x^*$ :

$$
(\boldsymbol{x}^*)^{\top}\boldsymbol{A}\hat{\boldsymbol{y}} = \hat{\boldsymbol{x}}^{\top}\boldsymbol{A}\hat{\boldsymbol{y}} + \sum_{i \in [n-r]} \beta_i \boldsymbol{q}_i^{\top}\boldsymbol{A}\hat{\boldsymbol{y}} = \hat{\boldsymbol{x}}^{\top}\boldsymbol{A}\hat{\boldsymbol{y}}.
$$

Similarly, we can also construct an optimal sparse solution  $y^*$  by leveraging  $\hat{y}$  and eigenvectors of zero eigenvalues of C such that  $||y^*||_0 \leq s_2$ .

Therefore, there exists an optimal solution  $(x^*, y^*)$  to the [CCA](#page-2-1) whose zero norms are bounded from above by r,  $\hat{r}$ , respectively. Adding the constraints  $||x||_0 \le r$ ,  $||y||_0 \le \hat{r}$  to the [CCA](#page-2-1) does not affect the optimality, which gives an equivalent formulation [\(2\)](#page-4-2) of [CCA.](#page-2-1)

**Part (ii).** Suppose that  $(\tilde{x}, \tilde{y})$  denotes an optimal solution to problem [\(3\)](#page-4-3). When  $s_1 \geq r$ , following the proof of Part (I),  $\tilde{x}$ , we can construct another optimal solution  $x^*$  whose zero norm is bounded by r and  $(x^*, \tilde{y})$  is feasible and optimal to [SCCA.](#page-0-0)

**Part (iii).** Similarly, we can reduce [SCCA](#page-0-0) to problem [\(4\)](#page-4-4). We thus complete the proof.  $\Box$ 

#### <span id="page-14-0"></span>A.4 Proof of [Theorem 3](#page-5-2)

*Proof.* Let us first consider the maximization problem over  $x$  in [\(6\)](#page-5-1), i.e.,

<span id="page-14-1"></span>
$$
v_x = \max_{\boldsymbol{x} \in \mathbb{R}^n} \{ \boldsymbol{a}^\top \boldsymbol{x} : \boldsymbol{x}^\top \boldsymbol{B} \boldsymbol{x} \leq 1, \|\boldsymbol{x}\|_0 \leq s_1 \}.
$$
 (13)

Then, we derive a combinatorial optimization reformulation of problem [\(13\)](#page-14-1) based on the result below.

<span id="page-14-2"></span>Claim 1 *For any subset*  $S \subseteq [n]$ ,  $\max_{\mathbf{x} \in \mathbb{R}^{|S|}} {\{\mathbf{a}_S^\top \mathbf{x} : \mathbf{x}^\top \mathbf{B}_{S,S} \mathbf{x} \leq 1\}} = \sqrt{\mathbf{a}_S^\top (\mathbf{B}_{S,S})^\dagger \mathbf{a}_S}.$ 

*Proof.* Given  $A = ab^{\top}$ , since the matrix  $\begin{pmatrix} B & ab^{\top} \\ b^{\top}c & C \end{pmatrix}$  $\begin{pmatrix} B & ab^{\top} \\ b^{\top}a & C \end{pmatrix}$  is positive semidefinite, using [Lemma 1,](#page-3-1) the identity  $(I - B_{S,S} B_{S,S}^{\dagger}) a_S b^{\top} = 0$  must hold for any subset S. As a result, we have  $a_S$  –  $B_{S,S}B_{S,S}^{\dagger}a_S = 0$  as vector *b* is nonzero.

Next, the Lagrangian dual of the problem  $\max_{\bm x\in\mathbb R^{|S|}}\{\bm a_{S}^\top\bm x:\bm x^\top\bm B_{S,S}\bm x\leq 1\}$  can be written as  $\max_{\boldsymbol{x}\in\mathbb{R}^{|S|}}\{\boldsymbol{a}_{S}^{\top}\boldsymbol{x}:\boldsymbol{x}^{\top}\boldsymbol{B}_{S,S}\boldsymbol{x}\leq 1\}=\min_{\mu\geq 0}\max_{\boldsymbol{x}\in\mathbb{R}^{|S|}}\boldsymbol{a}_{S}^{\top}\boldsymbol{x}+\mu-\mu\boldsymbol{x}^{\top}\boldsymbol{B}_{S,S}\boldsymbol{x}$ 

$$
= \min_{\mu \geq 0} \mu + \frac{\boldsymbol{a}_{S}^{\top} \boldsymbol{B}_{S,S}^{\dagger} \boldsymbol{a}_{S}}{4 \mu} = \sqrt{\boldsymbol{a}_{S}^{\top} (\boldsymbol{B}_{S,S})^{\dagger} \boldsymbol{a}_{S}},
$$

where the second equation builds on the identity  $a_S - B_{S,S} B_{S,S}^{\dagger} a_S = 0$  and optimal solution  $\boldsymbol{x}^* = \frac{\boldsymbol{B}_{S,S}^\dagger \boldsymbol{a}_S}{\sqrt{\boldsymbol{a}_S^\top (\boldsymbol{B}_{S,S})^\dagger \boldsymbol{a}_S}}$ .  $\Diamond$ 

Suppose that an optimal solution to problem  $(13)$  admits the support  $S^*$ . According to Claim [1,](#page-14-2) we have

$$
v_x = \max_{S \subseteq [n], |S| \leq s} \sqrt{\boldsymbol{a}_S^{\top}(\boldsymbol{B}_{S,S})^{\dagger} \boldsymbol{a}_S} = \sqrt{\boldsymbol{a}_{S^*}^{\top}(\boldsymbol{B}_{S^*,S^*})^{\dagger} \boldsymbol{a}_{S^*}}.
$$

On the other hand, the Lagrangian dual of problem [\(13\)](#page-14-1) can be written as

$$
\begin{aligned} v_x &\leq \min_{\lambda \in \mathbb{R}_+} \max_{\boldsymbol{x} \in \mathbb{R}^n} \{ \boldsymbol{a}^\top \boldsymbol{x} + \lambda - \lambda \boldsymbol{x}^\top \boldsymbol{B} \boldsymbol{x} : \|\boldsymbol{x}\|_0 \leq s_1 \} \\ &= \min_{\lambda \in \mathbb{R}_+} \max_{S \subseteq [n], |S| \leq s} \lambda + \frac{\boldsymbol{a}_S^\top (\boldsymbol{B}_{S,S})^\dagger \boldsymbol{a}_S}{4 \lambda} \\ &\leq \max_{S \subseteq [n], |S| \leq s} \lambda^* + \frac{\boldsymbol{a}_S^\top (\boldsymbol{B}_{S,S})^\dagger \boldsymbol{a}_S}{4 \lambda^*} = \sqrt{\boldsymbol{a}_{S^*}^\top (\boldsymbol{B}_{S^*,S^*})^\dagger \boldsymbol{a}_{S^*}} \leq v_x, \end{aligned}
$$

where the first equation is due to Claim [1,](#page-14-2) the second inequality is by plugging the feasible solution  $\lambda^* = \frac{\sqrt{\boldsymbol{a}_{S^*}^\top(\boldsymbol{B}_{S^*,S^*})^\dagger \boldsymbol{a}_{S^*}}}{2}$  $\frac{1}{2}$  into minimization problem, and the last equation is from the optimality of subset  $S^*$ . Since both left-hand and right-hand sides above equal  $v_x$ , the strong duality of problem [\(13\)](#page-14-1) holds, and all the inequalities above must attain the equalities. That is, problem [\(13\)](#page-14-1) is equivalent to

$$
v_x = \min_{\lambda \in \mathbb{R}_+} \max_{\boldsymbol{x} \in \mathbb{R}^n} \{ \boldsymbol{a}^\top \boldsymbol{x} + \lambda - \lambda \boldsymbol{x}^\top \boldsymbol{B} \boldsymbol{x} : \|\boldsymbol{x}\|_0 \leq s_1 \}.
$$

Since the outer minimization is a one-dimensional convex program that can be solved efficiently, as a result, for any given  $\lambda > 0$ , the inner maximization is equivalent to solving

<span id="page-15-4"></span><span id="page-15-3"></span>
$$
\max_{\boldsymbol{x}\in\mathbb{R}^n} \{\boldsymbol{a}^\top \boldsymbol{x} - \lambda \boldsymbol{x}^\top \boldsymbol{B} \boldsymbol{x} : \|\boldsymbol{x}\|_0 \le s_1\}.
$$
\n(14)

Next, let us consider the NP-hard sparse regression problem (see, e.g., [\[34\]](#page-11-20)), which admits

$$
\min_{\boldsymbol{\beta} \in \mathbb{R}^n} \left\{ \|\boldsymbol{v} - \boldsymbol{U}\boldsymbol{x}\|_2^2 : \|\boldsymbol{x}\|_0 \leq s \right\} \Longleftrightarrow \max_{\boldsymbol{x} \in \mathbb{R}^n} \left\{ 2\boldsymbol{v}^\top \boldsymbol{U}\boldsymbol{x} - \boldsymbol{x}^\top \boldsymbol{U}^\top \boldsymbol{U}\boldsymbol{\beta} : \|\boldsymbol{x}\|_0 \leq s \right\}, \qquad (15)
$$

where data matrix  $U$  consists of observations of n variables and vector  $v$  denotes the corresponding response variables.

Suppose that in the problem [\(14\)](#page-15-3), let us define  $\lambda \mathbf{B} = \mathbf{U}^\top \mathbf{U}$  and  $\mathbf{a} = 2\mathbf{U}^\top \mathbf{v}$ . Then using the singular value decomposition of matrix  $U$ , we see that the following equation still holds.

$$
\boldsymbol{a}_{S}-\boldsymbol{B}_{S,S}\boldsymbol{B}_{S,S}^{\dagger}\boldsymbol{a}_{S}=\boldsymbol{0},\forall S\subseteq[n].
$$

Thus, for any given  $\lambda > 0$ , the maximization problem [\(14\)](#page-15-3) is equivalent to the sparse regression problem [\(15\)](#page-15-4). This shows that problem [\(13\)](#page-14-1) is NP-hard.

Similarly, the maximization problem over  $y$  in [\(6\)](#page-5-1) can also be reduced to the sparse regression problem.  $\Box$ 

#### <span id="page-15-2"></span>A.5 Proof of [Theorem 4](#page-6-3)

*Proof.* For the SCCA [\(11\)](#page-13-2), according to [Proposition 1,](#page-3-2) the inner maximization problem admits an exact semidefinite programming formulation. Using the variables  $z \in \mathcal{Z}$  to describe the set constraints in SCCA [\(11\)](#page-13-2), we can reformulate it as

<span id="page-15-5"></span>
$$
v^* = \max_{\mathbf{z} \in \mathcal{Z}} \max_{\mathbf{X} \in \mathcal{S}_+^{n+m}} \left\{ \text{tr}(\tilde{\mathbf{A}}\mathbf{X}) : \text{tr}(\tilde{\mathbf{B}}\mathbf{X}) \le 1, \text{tr}(\tilde{\mathbf{C}}\mathbf{X}) \le 1, \right\}
$$
\n
$$
X_{ii}(1 - z_i) = 0, \forall i \in [m+n] \right\}.
$$
\n(16)

[Proposition 3](#page-16-1) shows that there is an optimal solution  $(x^*, y^*)$  to [SCCA](#page-0-0) that satisfies  $\|x^*\|_2^2 \le M_1$ and  $||y^*||_2^2 \le M_2$ . Based on this, we can construct an optimal solution  $(z^*, \mathbf{X}^*)$  for SCCA [\(16\)](#page-15-5) by letting

$$
\boldsymbol{X}^* = \begin{pmatrix} \boldsymbol{x}^* \\ \boldsymbol{y}^* \end{pmatrix} \begin{pmatrix} \boldsymbol{x}^* \\ \boldsymbol{y}^* \end{pmatrix}^\top, z_i = \begin{cases} 1 & \text{if } x_i^* \neq 0 \\ 0 & \text{if } x_i^* = 0 \end{cases}, \forall i \in [n], z_{i+n} = \begin{cases} 1 & \text{if } y_i^* \neq 0 \\ 0 & \text{if } y_i^* = 0 \end{cases}, \forall i \in [m],
$$

where the optimal solution  $X^*$  satisfies the following inequalities

$$
X_{ii}^* = (x_i^*)^2 \le M_1 z_i, \forall i \in [n], \ \ X_{(i+n)(i+n)}^* = (y_i^*)^2 \le M_2 z_{i+n}, \forall i \in [m].
$$

This allows us to recast the SCCA [\(16\)](#page-15-5) into an MISDP formulation [\(7\)](#page-6-2).  $\Box$ 

## <span id="page-15-0"></span>B Implementations of greedy and local search algorithms

This section presents the detailed implementations of greedy and local search algorithms based on the combinatorial formulation [\(1\)](#page-3-0) of [SCCA.](#page-0-0)

## <span id="page-15-1"></span>C Mixed-integer convex quadratic programming reformulations

This section shows that each subproblem in [\(6\)](#page-5-1) can be equivalently formulated by a MICQP. Therefore, SCCA  $(5)$  is mixed-integer convex quadratic representable when matrix  $\vec{A}$  is rank-one.

Algorithm 2 Greedy algorithm for SCCA [\(1\)](#page-3-0)

<span id="page-16-2"></span>1: **Input:** Matrices  $A \in \mathbb{R}^{m \times n}$ ,  $B \in \mathcal{S}_{+}^{m}$ ,  $C \in \mathcal{S}_{+}^{m}$  and integers  $s_1 \in [n]$ ,  $s_2 \in [m]$ 2: Compute  $(i^*, j^*) \in \text{argmax}_{i \in [m], j \in [n]} \sqrt{(B_{ii})^{\dagger}} A_{ij} \sqrt{(C_{jj})^{\dagger}}$ 3: Define subsets  $\hat{S}_1 = \{i^*\}$  and  $\hat{S}_2 = \{j^*\}$ 4: for  $\ell = 2, \cdots, \max\{s_1, s_2\}$  do<br>5: if  $\ell \le \min\{s_1, s_2\}$  then if  $\ell \leq \min\{s_1, s_2\}$  then 6: i  $\mathcal{A}^* \in \mathop{\rm argmax}_{i \in [n] \setminus \hat{S}_1} \sigma_{\max} \left( \sqrt{(\boldsymbol{B}_{\hat{S}_1 \cup \{i\}, \hat{S}_1 \cup \{i\}})^\dagger} \boldsymbol{A}_{\hat{S}_1 \cup \{i\}, \hat{S}_2} \sqrt{(\boldsymbol{C}_{\hat{S}_2, \hat{S}_2})^\dagger} \right)$ 7: Update  $\hat{S}_1 = \hat{S}_1 \cup \{i^*\}$  $s\colon \qquad \quad j^*\in\mathop{\rm argmax}_{j\in[m]\backslash{\hat S_2}}\sigma_{\max}\left(\sqrt{(B_{{\hat S_1},{\hat S_1}})^{\dag}}\bm A_{{\hat S_1},{\hat S_2\cup\{j\}}}\sqrt{(C_{{\hat S_2\cup\{j\}},\hat{S_2\cup\{j\}}})^{\dag}}\right)$ 9: else if  $s_1 \leq s_2$  then  $10:$  $\{*\in\mathop{\rm argmax}_{j\in[m]\setminus{\hat{S}_2}}\sigma_{\max}\left(\sqrt{(\boldsymbol{B}_{\hat{S}_1,\hat{S}_1})^{\dagger}}\boldsymbol{A}_{\hat{S}_1,\hat{S}_2\cup\{j\}}\sqrt{(\boldsymbol{C}_{\hat{S}_2\cup\{j\},\hat{S}_2\cup\{j\}})^{\dagger}}\right)$ 11: Update  $\hat{S}_2 = \hat{S}_2 \cup \{j^*\}$ 12: else  $13:$  $\mathbf{A}^* \in \mathop{\rm argmax}_{i \in [n] \setminus \hat{S}_1} \sigma_{\max} \left( \sqrt{(\boldsymbol{B}_{\hat{S}_1 \cup \{i\}, \hat{S}_1 \cup \{i\}})^\dagger} \boldsymbol{A}_{\hat{S}_1 \cup \{i\}, \hat{S}_2} \sqrt{(\boldsymbol{C}_{\hat{S}_2, \hat{S}_2})^\dagger} \right)$ 14: Update  $\hat{S}_1 = \hat{S}_1 \cup \{i^*\}$ 15: end if 16: end for 17: **Output:**  $\hat{S}_1, \hat{S}_2$ 

Algorithm 3 Local search algorithm for SSVD [\(1\)](#page-3-0)

1: **Input:** Matrices  $A \in \mathbb{R}^{m \times n}$ ,  $B \in \mathcal{S}_{+}^{m}$ ,  $C \in \mathcal{S}_{+}^{m}$  and integers  $s_1 \in [n]$ ,  $s_2 \in [m]$ [2](#page-16-2): Initialize  $(\hat{S}_1, \hat{S}_2)$  as the output of greedy Algorithm 2 3: do 4: for each pair  $(i_1, j_1) \in \hat{S}_1 \times ([n] \setminus \hat{S}_1)$  do  $\quad \text{if} \quad \quad \sigma_{\max} \left( \sqrt{ ( B_{ \hat{S}_1 \cup \{ j_1 \} \setminus \{ i_1 \} }, \hat{S}_1 \cup \{ j_1 \} \setminus \{ i_1 \} }, \hat{S}_1 \cup \{ j_1 \} \setminus \{ i_1 \} , \hat{S}_2} \sqrt{ ( C_{ \hat{S}_2 , \hat{S}_2 } )^\dagger } \right)$  $>$  $\sigma_{\max}\left(\sqrt{(\bm{B}_{\hat{S}_1,\hat{S}_1})^{\dagger}}\bm{A}_{\hat{S}_1,\hat{S}_2}\sqrt{(\bm{C}_{\hat{S}_2,\hat{S}_2})^{\dagger}}\right)$  then 6: Update  $\hat{S}_1 = \hat{S}_1 \cup \{j_1\} \setminus \{i_1\}$ 7: end if 8: end for 9: for each pair  $(i_2, j_2) \in \hat{S}_2 \times ([m] \setminus \hat{S}_2)$  do  $\text{10:} \qquad \text{if} \qquad \sigma_{\max}\left( \sqrt{(B_{\hat{S}_1 \cup \{j_1 \} \setminus \{i_1 \} }, \hat{S}_1 \cup \{j_1 \} \setminus \{i_1 \} } )^\dagger A_{\hat{S}_1 \cup \{j_1 \} \setminus \{i_1 \} }, \hat{S}_2 \sqrt{(C_{\hat{S}_2, \hat{S}_2} )^\dagger} \right)$  $>$  $\sigma_{\max}\left(\sqrt{(B_{\hat{S}_{1},\hat{S}_{1}})^{\dagger}}\bm{A}_{\hat{S}_{1},\hat{S}_{2}}\sqrt{(C_{\hat{S}_{2},\hat{S}_{2}})^{\dagger}}\right)$  then 11: Update  $\hat{S}_2 = \hat{S}_2 \cup \{j_2\} \setminus \{i_2\}$ 12: end if 13: end for 14: while there is still an improvement 15: **Output:**  $\hat{S}_1$ ,  $\hat{S}_2$ 

### <span id="page-16-0"></span>C.1 Valid inequalities for SCCA

Before deriving the formulations, we first prove that there exists a bounded optimal solution  $(x^*, y^*)$ of the [SCCA.](#page-0-0) To be specific, we show that there exists an optimal solution  $(x^*, y^*)$  of the [SCCA](#page-0-0) satisfying the constraints  $||x^*||_2^2 \leq M_1$  and  $||y^*||_2^2 \leq M_2$ , where  $M_1$  and  $M_2$  are finite-valued parameters.

<span id="page-16-1"></span>**Proposition 3** The [SCCA](#page-0-0) admits an optimal solution  $(x^*, y^*)$  satisfying  $||x^*||_2^2 \le M_1$  and  $||y^*||_2^2 \le$  $M_2$ , where  $M_1 = 1/\lambda_r(B) + 1/(\lambda_r(B)s_{\min}(B))$  and  $M_2 = 1/\lambda_{\hat{r}}(C) + 1/(\lambda_{\hat{r}}(C)s_{\min}(C))$  with  $\lambda_r(B), \lambda_{\hat{r}}(C)$  *being the smallest nonzero eigenvalues of matrices* **B**, C and  $s_{\min}(R)$  *being the smallest nonzero singular value of all the submatrices of the zero eigenvectors of matrix* R*.*

*Proof.* Let  $(x^*, y^*)$  denote an optimal solution to [SCCA.](#page-0-0) We bound  $||x^*||_2$  first and the same technique can be also straightforwardly applied to bound  $||y^*||_2$ .

For matrix  $B \in \mathcal{S}_{+}^n$  of rank r, we let  $\{q_i\}_{i \in [n]} \in \mathbb{R}^n$  denote the eigenvectors corresponding to *n* eigenvalues  $\lambda$  of B such that  $\lambda_1 \geq \ldots \geq \lambda_r > \lambda_{r+1} = \ldots = \lambda_n = 0$ . Thus,  $\{q_i\}_{i \in [n]}$  are orthonormal and span the space of  $\mathbb{R}^n$ . Hence, there exists  $\alpha \in \mathbb{R}^n$  such that  $x^* = \sum_{i \in [n]} \alpha_i \mathbf{q}_i$ . Given that  $(x^*)^\top B x^* \leq 1$ , we have

$$
\sum_{i \in [r]} \alpha_i^2 \lambda_i \le 1.
$$

Hence, the values of  $\{\alpha_i\}_{i\in[r]}$  are bounded. On the other hand, let us define a subset  $S \subseteq [n]$  of size at most  $s_1$  such that  $x_i^* \neq 0$  for each  $i \in S$  and  $x_j^* = 0$  for each  $j \in [n] \setminus S$ . Then for each  $j \in [n] \setminus S$ , we arrive at the following linear system:

$$
\sum_{j \in [r+1,n]} \alpha_j \hat{q}_j = -\sum_{i \in [r]} \alpha_i \hat{q}_i,\tag{17}
$$

<span id="page-17-0"></span>,

where  $\hat{q}_i$  denote a subvector of  $q_i$  with indices  $[n] \setminus S$  for each  $i \in [n]$ . For a fixed  $\{\alpha_i\}_{i \in [r]}$ , since the linear system [\(17\)](#page-17-0) is nonempty, we let  $\overline{Q}\overline{\alpha} = \overline{q}$  denote its minimal linear subsystem such that a submatrix  $\bar{Q}$  is non-singular and the index set  $\hat{S}$  of  $\bar{\alpha}$  is a subset of  $[n] \setminus S$ . Thus, we can construct an alternative solution  $\hat{\alpha}$  such that

$$
\hat{\alpha}_i = \begin{cases} \alpha_i, & \text{if } i \in [r], \\ (\bar{Q}^{-1}\bar{q})_i, & \text{if } i \in \hat{S}, \\ 0, & \text{otherwise,} \end{cases}
$$

and  $\hat{\boldsymbol{x}} = \sum_{i \in [n]} \hat{\alpha}_i \boldsymbol{q}_i$ . According to [Lemma 1,](#page-3-1) we have

$$
\hat{\boldsymbol{x}}^\top \boldsymbol{B} \hat{\boldsymbol{x}} \leq 1, \hat{\boldsymbol{x}}^\top \boldsymbol{A} \boldsymbol{y}^* = (\boldsymbol{x}^*)^\top \boldsymbol{A} \boldsymbol{y}^*
$$

i.e.,  $(\hat{x}, y^*)$  is also optimal to [SCCA.](#page-0-0) Hence,

$$
\|\hat{\bm{x}}\|_2 \leq \sqrt{\|\bar{\bm{Q}}^{-1}\bar{\bm{q}}\|_2^2 + \sum_{i \in [r]} \alpha_i^2}
$$

Note that  $\sum_{i \in [r]} \alpha_i^2 \leq 1/\lambda_r$  and

$$
\|\bar{\bm{Q}}^{-1}\bar{\bm{q}}\|_2^2 \leq \|\bar{\bm{Q}}^{-1}\|_2^2 \|\bar{\bm{q}}\|_2^2 \leq \frac{1}{s_{\min}(\bm{B})}\frac{1}{\lambda_r}
$$

where  $s_{\min}(B)$  denotes the smallest nonzero singular values of all the submatrices of  $[q_{r+1}, \ldots, q_n]$ . In summary, we have

$$
\|\hat{\boldsymbol{x}}\|_2 \leq \sqrt{1/\lambda_r + 1/(\lambda_r s_{\min}(\boldsymbol{B}))}.
$$

This completes the proof. □

The proof of Proposition [3](#page-16-1) is straightforward in the case when  $B$  and  $C$  are of full rank as in this case the feasible region is a bounded set. In order to prove the result in the case when  $\bf{B}$  is not full-rank, one has to show that it is possible to construct sparse solutions that are not "too far" away.

In fact, the bounds  $M_1, M_2$  in [Proposition 3](#page-16-1) also hold for any given feasible subsets  $(S_1, S_2)$  of SCCA [\(1\)](#page-3-0).

<span id="page-17-1"></span>**Corollary 2** *For any given feasible subsets*  $(S_1, S_2)$  $(S_1, S_2)$  $(S_1, S_2)$  *of SCCA* 1, *there exists* a SCCA feasible solution  $(x, y)$  such that the supports of  $x, y$  are  $S_1, S_2$ , respectively and we have that  $||x||_2^2 \le M_1$  and  $||y||_2^2 \leq M_2$ , where  $M_1, M_2$  are defined in [Proposition 3.](#page-16-1)

#### C.2 Equivalent mixed-integer convex quadratic program of rank-one SCCA

When matrix A is rank-one, let us focus on analyzing the subproblem over x in [\(6\)](#page-5-1), i.e.,

<span id="page-17-2"></span>
$$
v_x = \max_{\boldsymbol{x} \in \mathbb{R}^n} \{ \boldsymbol{a}^\top \boldsymbol{x} : \boldsymbol{x}^\top \boldsymbol{B} \boldsymbol{x} \le 1, \|\boldsymbol{x}\|_0 \le s_1 \}.
$$
 (18)

Then the second subproblem over  $\boldsymbol{\psi}$  in [\(6\)](#page-5-1) simply follows.

According to [Corollary 2,](#page-17-1) introducing the binary variables  $z^1 \in \{0,1\}^n$  can reformulate the problem [\(18\)](#page-17-2) as

$$
v_x = \max_{\boldsymbol{x} \in \mathbb{R}^n, \boldsymbol{z}^1 \in \{0,1\}^n} \bigg\{ \boldsymbol{a}^\top \boldsymbol{x} : \boldsymbol{x}^\top \boldsymbol{B} \boldsymbol{x} \leq 1, x_i \leq \sqrt{M_1} z_i^1, \forall i \in [n], \sum_{i \in [n]} z_i^1 \leq s_1 \bigg\}.
$$

When matrix B is positive definite, there is a positive vector  $b \in \mathbb{R}_{++}^n$  and a positive semidefinite matrix  $\hat{B}$  such that  $B = \hat{B} + \text{Diag}(b)$ . Given this equation, by leveraging the perspective techniques (see, e.g., [\[1,](#page-10-18) [43\]](#page-11-18)), we can derive another equivalent MICQP formulation of the problem [\(18\)](#page-17-2):

$$
v_x = \max_{\boldsymbol{x} \in \mathbb{R}^n, \boldsymbol{z}^1 \in \{0,1\}^n, \boldsymbol{\mu} \in \mathbb{R}_+^n} \left\{ \boldsymbol{a}^\top \boldsymbol{x} : \boldsymbol{x}^\top \hat{\boldsymbol{B}} \boldsymbol{x} + \sum_{i \in [n]} \mu_i \leq 1, x_i^2 \leq \mu_i z_i^1, \forall i \in [n], \sum_{i \in [n]} z_i^1 \leq s_1 \right\}.
$$

which is often stronger than the above formulation.

## <span id="page-18-0"></span>D A branch-and-cut algorithm with closed-form cuts

By dualizing the inner maximization problem over  $X$  in the MISDP [\(7\)](#page-6-2), in this subsection, we derive an equivalent mixed-integer linear program for [SCCA,](#page-0-0) which motivates us to develop a branch-and-cut algorithm.

By introducing the Lagrangian multipliers  $(\theta_1, \theta_2, \lambda)$ , the Lagrangian dual of the maximization problem [\(9\)](#page-6-5) can be written as

<span id="page-18-1"></span>
$$
f(z) = \min_{\substack{\theta_1 \ge 0, \theta_2 \ge 0, \ X \in \mathcal{S}_+^{n+m}}} \max_{\mathbf{X} \in \mathcal{S}_+^{n+m}} \text{tr}(\tilde{\mathbf{A}}\mathbf{X}) - \theta_1 \text{tr}(\tilde{\mathbf{B}}\mathbf{X}) - \theta_2 \text{tr}(\tilde{\mathbf{C}}\mathbf{X}) + \theta_1 + \theta_2,
$$
  
\n
$$
- \sum_{i \in [n+m]} \lambda_i X_{ii} + \sum_{i \in [n+m]} \lambda_i M_{ii} z_i
$$
  
\n
$$
= \min_{\substack{\theta_1 \ge 0, \theta_2 \ge 0, \ \lambda \in \mathbb{R}_+^{n+m}}} \left\{ \theta_1 + \theta_2 + \sum_{i \in [n+m]} \lambda_i M_{ii} z_i : \begin{pmatrix} \theta_1 \mathbf{B} & -\mathbf{A}/2 \\ -\mathbf{A}^\top / 2 & \theta_2 \mathbf{C} \end{pmatrix} \succeq -\text{Diag}(\boldsymbol{\lambda}) \right\},
$$
\n(19)

where the strong duality holds due to the function  $f(z)$  being concave, bounded, and thus continuous in the set  $\hat{z}$  and Slater condition holds for any interior point z in the set  $\hat{z}$ .

Below, we derive the closed-form expression of the function  $f(z)$  with the given binary variable  $z \in \mathcal{Z}$ . This allows us to reformulate SCCA [\(8\)](#page-6-4) as a mixed-integer linear program with exponentially many linear constraints and an efficient separation oracle.

#### Proposition 4 *The SCCA* [\(8\)](#page-6-4) *is equivalent to*

<span id="page-18-2"></span>
$$
v^* = \max_{\mathbf{z} \in \mathcal{Z}, v} \left\{ v : v \le \sigma_{\max} \left( \sqrt{(\mathbf{B}_{S_1, S_1})^{\dagger} \mathbf{A}_{S_1, S_2} \sqrt{(\mathbf{C}_{S_2, S_2})^{\dagger}}} \right) + \sum_{i \in S_1 \cup (S_2 + n)} \lambda^* M_{ii} z_i : \forall S_1 \subseteq [n], |S_1| \le s_1, S_2 \subseteq [m], |S_2| \le s_2 \right\},
$$
\n(20)

where for a pair of subsets  $(S_1, S_2)$ , the scalar  $\lambda^*$  is defined as the largest positive eigenvalue of matrix  $\boldsymbol{D}_2^\top \boldsymbol{D}_1^{-1} \boldsymbol{D}_2 - \boldsymbol{D}_3$  with

$$
\boldsymbol{D}_1 = \begin{pmatrix} \theta_1^* \boldsymbol{B}_{S_1,S_1} & -\boldsymbol{A}_{S_1,S_2}/2 \\ -\boldsymbol{A}_{S_1,S_2}^{\top}/2 & \theta_2^* \boldsymbol{C}_{S_2,S_2} \end{pmatrix}, \ \ \boldsymbol{D}_2 = \begin{pmatrix} \theta_1^* \boldsymbol{B}_{S_1,[n]\backslash S_1} & -\boldsymbol{A}_{S_1,[m]\backslash S_2}/2 \\ -\boldsymbol{A}_{S_2,[n]\backslash S_1}^{\top}/2 & \theta_2^* \boldsymbol{C}_{S_2,[m]\backslash S_2} \end{pmatrix},
$$

*and*

$$
\bm{D}_3=\left(\begin{array}{cc}\theta_1^* \bm{B}_{[n]\backslash S_1,[n]\backslash S_1} & -\bm{A}_{[n]\backslash S_1,[m]\backslash S_2}/2 \\ -\bm{A}_{[n]\backslash S_1,[m]\backslash S_2}^\top/2 & \theta_2^* \bm{C}_{[m]\backslash S_2,[m]\backslash S_2} \end{array}\right),
$$
 where  $\theta_1^*=\theta_2^*=\sigma_{\max}\left(\sqrt{(\bm{B}_{S_1,S_1})^{\dagger}}\bm{A}_{S_1,S_2}\sqrt{(\bm{C}_{S_2,S_2})^{\dagger}}\right)/2.$ 

*Proof.* First, for any binary variable  $z \in \mathcal{Z}$ , suppose  $S_1 = \{i : z_i = 1, \forall i \in [n]\}, S_2 = \{i - n : z_i = 1, \forall i \in [n]\}$  $z_i = 1, \forall i \in [n+1, n+m]$ , and  $T \subseteq [n+m]$  denotes the support of z. Then following the proof of [Proposition 1,](#page-3-2) we can construct a rank-one optimal solution  $X^* = \begin{pmatrix} x^* & 0 \\ 0 & x^* \end{pmatrix}$  $\left(\begin{matrix} x^* \ y^* \end{matrix}\right) \left(\begin{matrix} x^* \ y^* \end{matrix}\right)$  $\begin{bmatrix} x^* \\ y^* \end{bmatrix}^\top$  to the maximization problem below that admits the optimal value  $\sigma_{\max}\left(\sqrt{(B_{S_1,S_1})^\dagger}A_{S_1,S_2}\sqrt{(C_{S_2,S_2})^\dagger}\right)$ , i.e.,

$$
\max_{\mathbf{X}\in\mathcal{S}_{+}^{n+m}}\{\text{tr}(\tilde{\mathbf{A}}\mathbf{X}) : \text{tr}(\tilde{\mathbf{B}}\mathbf{X}) \leq 1, \text{tr}(\tilde{\mathbf{C}}\mathbf{X}) \leq 1, X_{ii} = 0, \forall i \in [n+m] \setminus T\}
$$
\n
$$
= \sigma_{\max}\left(\sqrt{(\mathbf{B}_{S_1, S_1})^{\dagger}} \mathbf{A}_{S_1, S_2} \sqrt{(\mathbf{C}_{S_2, S_2})^{\dagger}}\right) \geq f(\mathbf{z}),
$$

where the inequality is because the maximization problem above relaxes the valid constraints  $X_{ii} \leq$  $M_{ii}$  for all  $i \in T$  in maximization problem [\(9\)](#page-6-5). The result in [Corollary 2](#page-17-1) suggests that  $x^*$ ,  $y^*$  can be bounded and their two norms must not exceed  $M_1, M_2$ , which means that the optimal solution  $X^*$  satisfies the  $X_{ii} \leq M_{ii}$  for all  $i \in T$ . Therefore,  $X^*$  is feasible and optimal to maximization problem [\(9\)](#page-6-5) and we have that

$$
f(\boldsymbol{z}) = \sigma_{\max}\left(\sqrt{(\boldsymbol{B}_{S_1,S_1})^{\dagger}} \boldsymbol{A}_{S_1,S_2}\sqrt{(\boldsymbol{C}_{S_2,S_2})^{\dagger}}\right).
$$

According to strong duality, the minimization problem [\(19\)](#page-18-1) admits an optimal value  $\sigma_{\max} \left( \sqrt{(\boldsymbol{B}_{S_1,S_1})^{\dagger}} \boldsymbol{A}_{S_1,S_2} \sqrt{(\boldsymbol{C}_{S_2,S_2})^{\dagger}} \right)$ . Next, we construct its optimal solution  $(\theta_1^*,\theta_2^*,\boldsymbol{\lambda}^*)$ .

For any given  $\epsilon > 0$ , we let  $\theta_1^* = f(z)/2$ ,  $\theta_2^* = f(z)/2$ ,  $\hat{\lambda}_i(\epsilon) = \frac{\epsilon}{M_{ii}|T|}$  for all  $i \in T$ , and  $\hat{\lambda}_i(\epsilon) = \lambda^*(\epsilon)$  for all  $i \in [n] \setminus T$ , where

$$
\lambda^*(\epsilon) = \left[\lambda_{\max}\left(\boldsymbol{D}_2^\top \left(\boldsymbol{D}_1 + \mathrm{Diag}\left(\hat{\boldsymbol{\lambda}}_T(\epsilon)\right)\right)^{-1}\boldsymbol{D}_2 - \boldsymbol{D}_3\right)\right]_+.
$$

It is easy to compute that  $\theta_1^* + \theta_2^* + \sum_{i \in [n+m]} \hat{\lambda}_i(\epsilon) M_{ii} z_i = f(z) + \epsilon$ . Thus, for any  $\epsilon > 0$ , if  $(\theta_1^*, \theta_2^*, \hat{\lambda}(\epsilon))$  were feasible, then it is an  $\epsilon$ -optimal solution to the minimization problem [\(19\)](#page-18-1). It remains to verify the feasibility of the solution  $(\theta_1^*, \theta_2^*, \hat{\bm{\lambda}}(\epsilon))$ , i.e., checking the constraint below

$$
\begin{pmatrix} \theta_1^* \mathbf{B} & -\mathbf{A}/2 \\ -\mathbf{A}^\top/2 & \theta_2^* \mathbf{C} \end{pmatrix} + \text{Diag}\left(\hat{\boldsymbol{\lambda}}(\epsilon)\right) \succeq 0.
$$

By performing the permutation of the rows and columns of the above matrix, it is sufficient to show that the new block matrix

<span id="page-19-0"></span>
$$
\begin{pmatrix} \mathbf{D}_1 + \text{Diag}\left(\hat{\boldsymbol{\lambda}}_T(\epsilon)\right) & \mathbf{D}_2 \\ \mathbf{D}_2^\top & \mathbf{D}_3 + \lambda^*(\epsilon)\mathbf{I} \end{pmatrix} \succeq 0, \tag{21}
$$

is positive semidefinite.

Since  $\begin{pmatrix} B_{S_1,S_1} & -A_{S_1,S_2}/2 \\ -A^{\top} & 0 & C_{S,S_1} \end{pmatrix}$  $-{\bm A}_{S_1,S_2}^\top/2$   ${\bm C}_{S_2,S_2}$  is a principal submatrix of a positive semidefinite matrix  $\begin{pmatrix} B & -A/2 \end{pmatrix}$  $-{\bm A}^\top/2$   $C$ ), it is also positive semidefinite. According to [Lemma 1](#page-3-1) and the fact that  $\theta_1^* = \theta_2^* = \sigma_{\max} \left( \sqrt{(\boldsymbol{B}_{S_1,S_1})^{\dagger}} \boldsymbol{A}_{S_1,S_2} \sqrt{(\boldsymbol{C}_{S_2,S_2})^{\dagger}} \right) / 2$ , the matrix  $\boldsymbol{D}_1$  is also positive semidefinite. As  $\epsilon > 0$ , the matrix  $\mathcal{D}_1 + \text{Diag}(\hat{\lambda}_T(\epsilon))$  must be positive definite, which means that

$$
\left(\boldsymbol{I}-\left(\boldsymbol{D}_1+\text{Diag}\left(\hat{\boldsymbol{\lambda}}_T(\epsilon)\right)\right)\left(\boldsymbol{D}_1+\text{Diag}\left(\hat{\boldsymbol{\lambda}}_T(\epsilon)\right)\right)^{-1}\right)\boldsymbol{D}_2=\mathbf{0}.
$$

Besides, according to the definition of  $\lambda^*(\epsilon)$ , we obtain

$$
\boldsymbol{D}_3+\lambda^*(\epsilon)\boldsymbol{I}-\boldsymbol{D}_2^\top\left(\boldsymbol{D}_1+\text{Diag}\left(\hat{\boldsymbol{\lambda}}_T(\epsilon)\right)\right)^{-1}\boldsymbol{D}_2\succeq \boldsymbol{0}.
$$

Taking these results together, according to [Lemma 1,](#page-3-1) the constraint in [\(21\)](#page-19-0) must hold for a given solution  $(\theta_1^*, \theta_2^*, \hat{\lambda}(\epsilon))$ . Since the objective value corresponding to  $(\theta_1^*, \theta_2^*, \hat{\lambda}(\epsilon))$  is at most  $\epsilon$  larger than the optimal value of problem [\(19\)](#page-18-1), letting  $\epsilon \to 0$  and using the closedness of the feasible set in problem [\(19\)](#page-18-1), we can confirm the optimality of  $(\theta_1^*, \theta_2^*, \lambda^*)$  with  $\lambda_i^* = 0$  for all  $i \in T$  and  $\lambda_i^* = \lambda^*$ for all  $i \in [n] \setminus T$ .

Given the closed-form optimal solution to problem [\(19\)](#page-18-1), the rest of the proof follows from [\[28,](#page-11-7) theorem 7].

We note that SCCA [\(20\)](#page-18-2) can be implemented via a delayed cut-generation procedure. That is, at each feasible branch-and-bound node with a binary solution  $\hat{z}$ , let  $S_1 = \{i : \hat{z}_i = 1, \forall i \in [n]\}\$ and  $S_2 = \{i - n : \hat{z}_i = 1, \forall i \in [n+1, n+m]\}$ . Then we can compute the corresponding scalar  $\lambda^*$  and generate the following valid inequality based on [\(20\)](#page-18-2):

$$
v \leq \sigma_{\max} \left( \sqrt{(\boldsymbol{B}_{S_1,S_1})^{\dagger}} \boldsymbol{A}_{S_1,S_2} \sqrt{(\boldsymbol{C}_{S_2,S_2})^{\dagger}} \right) + \sum_{i \in S_1 \cup (S_2 + n)} \lambda^* M_{ii} z_i.
$$

## <span id="page-20-1"></span>E Data description

Dataset	$\#$ of variables	$#$ of samples	$\boldsymbol{n}$	m	rank $r$	rank $\hat{r}$
dermatology	34	366	17	17	17	17
spambase	57	4601	28	29	28	29
digits	64	1797	32	32	32	32
$b$ <i>uzz</i> .	77	583250	38	39	38	39
gas	128	2565	64	64	64	64
slice	385	53500	192	193	192	193
breast cancer	21821	89	19,672	2,149	89	89

Table 7: Description of UCI and breast cancer datasets used

## <span id="page-20-0"></span>F Multiple Sparse Canonical Correlation Analysis

The multiple CCA problem can be formulated as follows:

$$
\max_{\bm{x}\in\mathbb{R}^{n\times k},\bm{y}\in\mathbb{R}^{m\times k}}\left\{\mathrm{tr}(\bm{x}^{\top}\bm{A}\bm{y}):\bm{x}^{\top}\bm{B}\bm{x}=\bm{I}_{k},\bm{y}^{\top}\bm{C}\bm{y}=\bm{I}_{k}\right\},
$$

where k denotes the number of pairs of basis vectors and  $I_k$  denotes the identity matrix of size k.

As  $x, y$  can be matrices, we propose adding row sparse constraints to extend [SCCA](#page-0-0) for multiple vectors, which is defined as:

$$
\max_{{\boldsymbol x}\in\R^{n\times k},{\boldsymbol y}\in\R^{m\times k}}\big\{\operatorname{tr}({\boldsymbol x}^\top{\boldsymbol A}{\boldsymbol y}): {\boldsymbol x}^\top{\boldsymbol B}{\boldsymbol x} = {\boldsymbol I}_k, {\boldsymbol y}^\top{\boldsymbol C}{\boldsymbol y} = {\boldsymbol I}_k, \|{\boldsymbol x}\|_0\le s_1, \|{\boldsymbol y}\|_0\le s_2\big\},
$$

where we let  $||x||_0$  and  $||y||_0$  denote the number of nonzero rows of x and y, respectively.

This multiple SCCA model can (i) compute the multiple weight vectors  $(x, y)$  simultaneously and (ii) enforce the sparsity and orthogonality strictly. To be specific, the constraints  $x^\top B x = I_k, y^\top C y =$  $I_k$  ensure the orthogonal left- and right-canonical loading vectors in multiple SCCA. By the definition of row sparsity, the resultant multiple left- and right-basis vectors, i.e., the columns of x and y, share the same nonzero rows, respectively.

More importantly, the row-sparsity enables us to readily extend the proposed algorithms to solve multiple SCCA. We have tested them on UCI data, and the computational results are presented in [Table 8.](#page-21-0) As k increases, it takes branch-and-cut a longer time to return an optimal solution.

					Greedy		Local search			Convex relaxation		Branch-and-cut			
$\boldsymbol{n}$	$m\,$	$s_1$	$S_2$	k <sub>1</sub>	LB	time	LB	time	UB	$\text{gap}(\%)$ time		$v^*$	$MIPGap(\%)$	time	
17	17	5	5	2 <sup>1</sup>	$1.907$ 0.01		1.935	0.06	1.957		$1.14 \ 0.01$	1.935	0.00	2	
17	17	10	10	$\left 3\right $			$2.879$ 0.02 2.884	0.09	2.898			$0.45$ $0.01$   2.884	0.00	6	
28	29	.5	5	2 <sup>1</sup>	$1.182$ 0.02 1.233				$0.09 \mid 1.358$	$10.19$ $0.01$		1.233	0.00	234	
28	29	10	10		$3 \mid 1.579 \quad 0.04 \mid 1.586$			0.14	1.685			$6.23$ $0.01$   1.587	5.33		
32	32	5	.5	2 <sub>1</sub>			$1.906$ 0.02 1.906	0.04	1.935			$1.55$ 0.01   1.916	0.00	14	
32	32	10	10	3	2.736 0.04 2.741			0.19	2.770			$1.05 \quad 0.01 \mid 2.742$	$0.00^{\circ}$	3093	
38	39	5.	5 <sup>5</sup>	$\mathfrak{D}$		2 0.03	2	0.25	2		$0.00 \quad 0.01$	2	0.00	8	
38	39	10	10 <sup>3</sup>			3 0.05	$\mathcal{E}$	0.59	$\mathbf{3}$		$0.00 \ 0.01$	3	0.00	10	
64	64	5	5	$2^{1}$	$1.947$ $0.05$   1.991			0.34	1.997		$0.29$ $0.01$	1.993	0.15		
64	64	10	10		3 2.983 0.09 2.989			0.71	2.993			$0.14 \quad 0.02 \mid 2.989$	0.14		
192	193	5	5	2 <sup>1</sup>	$1.911 \quad 0.21 \quad 1.991$			2.36	1.995			$0.22 \quad 0.04 \mid 1.991$	0.21		
192	193	10	10		$3 \mid 2.907 \quad 0.38 \mid 2.951$			5.78	2.977			$0.90 \quad 0.04 \mid 2.954$	0.78		

<span id="page-21-0"></span>Table 8: Evaluation of our algorithms for solving multiple SCCA on UCI datasets

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