

Bridging the Cyber and Physical with a Verifiable, Executable Language for Robotics

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Abstract—Building reliable Cyber-Physical Systems (CPS) often requires formal verification, which is capable of producing the rigorous safety guarantees necessary to provide confidence in its design. However, verification may be challenging to perform on real-world systems, and as a result, a simplified model may be verified instead. This results in a disconnect between verification and reality, leading to diminished confidence in the verification result. We have developed a method that enables CPS designers to both formally verify and implement CPS software in a single language. By combining the industry-tested modeling and compilation of synchronous programming with the verification rigor of refinement types, we strive to give CPS designers additional confidence in designing reliable, real-world CPS.

I. INTRODUCTION

Designing high-reliability CPS often entails formal verification, as it can provide stronger safety guarantees than simply testing experimentally. However, formal verification requires that the system being verified accurately represents the real system. Unfortunately, this is not always straightforward. Programming languages for CPS design require well-understood formal semantics for verification, while those designed for implementation require expressiveness for producing executable software. However, there are challenges in reconciling the two. Real implementations of CPS may be complex and difficult to characterize, making verification difficult unless performed on an abstract model. This may introduce uncertainty in the verification result. We present a method for verification and implementation in a single language, by combining the industry-tested modeling and compilation capabilities of synchronous programming with formal verification through an extended type system. With this method, we have formally verified and implemented a vehicle collision-avoidance braking controller on a real robot [1].

A. Synchronous Modeling and Execution

Synchronous programming languages treat CPS as collections of functions that operate on streams of incoming data, such as sensor values, to produce output streams, such as actuator commands [2], [3]. By design, synchronous programs have limited memory and must react on the same time base as their inputs, which makes them useful for modeling embedded systems. Real-world implementations of synchronous languages have already been employed in industry [4], [5], modeling critical systems such as avionics and integrated circuits [6]. We extend an existing synchronous programming

language already capable of producing executable simulations [7], [8], with verification and the ability to operate real robots.

B. Verification through Refinement Types

Formal verification is introduced via refinement types, similar to those found in non-synchronous languages [9], [10]. Refinement types extend a language’s type system by allowing type specifications to be conditioned on program values. For instance, a value that must only be a positive integer may be specified as $\{v : \text{int} \mid v > 0\}$, where values of the type can only be an integer v where $v > 0$. The incorporation of logical predicates into type specifications and the ability to refer to program terms make refinement types a powerful specification tool. Furthermore, refinement types still inherit the compile-time verification and modularity of normal types [9]. Although type-checking has already been applied to safety-critical verification tasks, such as those in robotics [11] and networking [12], extending refinement type checking to streams presents some unique challenges, which we address in our project. Our contribution includes the translation of refinement types in synchronous programs into constraints, which can then be checked for compatibility using a Satisfiability Modulo Theory (SMT) solver such as Z3 [13].

C. Related Work

Hybrid automata [14] and differential dynamic logic [15] are commonly used for modeling and verifying CPS. Although they can model real-world systems [16], their nondeterministic nature precludes directly generating executable code, though some translation may be possible [17]. Nevertheless, there exists a gap between the CPS that is implemented on a real system and the model CPS that is verified.

Other approaches to synchronous program verification have been explored in existing projects. For instance, Kind [18] and Kind 2 [19], use an approach similar to bounded model checking to synthesize invariants. Similarly, the Simulink-based verification tool CoCoSim [20] translates Simulink models and assume-guarantee statements into synchronous programs verifiable using the previous method. Our approach differs in that the constraints generated for type checking allow for a more thorough search of the state space, but may require invariants to avoid spurious counterexamples.

Verifiable and executable CPS have been explored in other works. VeriPhy [17] proposes a method for compiling differ-

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ential dynamic logic specifications into runtime monitors that allow a system to recover from anomalies [17]. Koord [21], [22] verifies coordination-level multi-agent robotics systems and executes on real robots. Though similar in use case, our project fulfills a distinct role relative to these works, and they may well complement one another. For instance, one may use our work to design and verify the fallback controller used in VeriPhy, or the low-level controllers used in Koord.

II. METHODS

We chose to extend the Zélus language [7], a language primarily designed for simulating synchronous hybrid automata, and its compiler [8]. Although its support for hybrid systems allows one to model systems with continuous-time dynamics, we note that our current work verifies only discrete-time components. We base the extended type system on the refinement types found in Liquid Haskell [10] and its translation of type specifications into automatically checkable constraints. We then bridge the gap between the synchronous program and the real world using a shared-variable protocol similar to the CyPhyHouse multi-agent robotics platform [22]. Combining these components, we ensure a seamless verification to execution workflow in a single language. Technical details regarding the language and implementation can be found in our workshop paper [1].

A. Type-Checking Synchronous Programs

We developed a set of typing rules modeled after existing refinement type systems [9], [10], [23]. A unique challenge that arises from combining synchronous programs with refinement types is ensuring that types on streams indeed describe all possible values of the stream. Unlike ordinary refinement types, our type system must also account for temporal behavior such as an if-statement being conditioned on a stream of boolean values. In this case, the program repeatedly switches between the two branches, which themselves exhibit temporal behavior. We note that, in the discrete subset of Zélus, streams can be modeled inductively, with one or more initial conditions and an evolution function that determines the stream’s subsequent value using a finite number of previous values. Thus, universal temporal properties, specifying that something always or never happens, can be conveniently encoded as SMT constraints, though inductive invariants may be necessary to rule out false counterexamples. Although only a subset of temporal logic, many safety properties of note can be encoded as a universal temporal property in practice. The typing rules we developed form the foundation of our type checker, which serves as an interface between the Zélus program augmented with refinement types and the Z3 SMT solver used to verify that all constraints are compatible.

B. Verifying a Real-World Example

We demonstrated the design process of a simple robot program modeled after a vehicle’s adaptive cruise control system [16]. In the various scenarios we sought to verify, the robot with the formally-verified controller (the “ego” vehicle) follows another vehicle along a single lane, which may

choose to accelerate or decelerate without the knowledge of the ego vehicle. The ego vehicle is specified to avoid a collision at all times. For simplicity, the ego can only either accelerate or brake at its maximum rate (a “bang-bang” controller). However, we note that more complex controllers can still be verified, given the appropriate invariants are chosen. The key insight behind choosing invariants often entails explicitly defining an inductive property of the system. For instance, the inductive invariant we use in this scenario assumes that the ego is never in a situation where a collision is unavoidable. Verification then becomes proving that the controller in the current time step sustains this property. Combined with verifying that the system began in a safe configuration, property is thus proven inductively.

C. Executing Synchronous Programs on Real Robots

Although the Zélus compiler in its unmodified state is able to generate executable simulations, these simulations have no interactions with the physical world, save for synchronizing simulations to real time. We extend Zélus with an inter-process communications scheme based on Lightweight Communications and Marshalling (LCM) [24] to share variable values between the synchronous program and low-level drivers for various robot peripherals. The low-level drivers implement simple motor speed controllers and sensor sampling for convenience, though it is possible to implement parts of these in the verified synchronous code to shrink the trusted base. Our current experiments employ a small differential-drive wheeled robot equipped with LiDAR and wheel speed sensors, making it well-suited for simulating ground vehicles. We have published experimental results from a simplified version of the aforementioned adaptive cruise control scenario [1], and we have results from the full system in submission.

III. DISCUSSION AND FUTURE WORK

The project’s primary goal is to develop a trusted chain from system specification and modeling to implementation on real-world CPS. Our methods extend the modeling capabilities of synchronous languages with specifications through type annotations, allowing for compile-time verification of safety properties. The present work trusts that the compiler correctly translates the verified synchronous program into safe executables. There have been developments in formally-verified compilation of synchronous languages [25], which if implemented would further enhance trust in our methods. We also assume that the user has correctly specified the environment in which the CPS operates. As a result, we carry assumptions made about the environment’s behavior to runtime. A possible solution could be to incorporate runtime monitoring, at the software [17] or hardware [26] level, which can leverage our generated verification constraints to check that verification assumptions are met in the real system. Finally, we would like to extend verification to hybrid components of the Zélus language, giving the user even more expressiveness in verifying and implementing safe CPS.

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