Differential Privacy of Dirichlet Posterior Sampling

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Abstract

1 Introduction

 The Bayesian framework provides a way to perform statistical analysis by combining prior beliefs with real-life evidence. At a high level, the belief and the evidence are assumed to be described by probabilistic models. As we receive new data, our belief is updated accordingly via the Bayes' theorem, resulting in the so-called posterior belief. The posterior tells us how much we are uncertain about the model's parameters.

 The Dirichlet distribution is usually chosen as the prior when performing Bayesian analysis on discrete variables, as it is a conjugate prior to the categorical and multinomial distributions. Specifically, Dirichlet distributions are often used in discrete mixture models, where a Dirichlet prior is put on the mixture weights [\[LW92;](#page-9-0) [MMR05\]](#page-9-1). Such models have applications in NLP [\[PB98\]](#page-10-0), biophysical systems [\[Hin15\]](#page-9-2), accident analysis [\[de 06\]](#page-9-3), and genetics [\[BHW00;](#page-9-4) [PM01;](#page-10-1) [CWS03\]](#page-9-5). In all of these studies, samplings from Dirichlet posteriors arise when performing Markov chain Monte Carlo methods for approximate Bayesian inference.

 Dirichlet posterior sampling also appears in other learning tasks. For example, in Bayesian active learning, it arises in Gibbs sampling, which is used to approximate the posterior of the classifier over the labeled sample [\[NLYCC13\]](#page-10-2). In Thompson sampling for multi-armed bandits, one repeatedly draws a sample from the Dirichlet posterior of each arm, and picks the arm whose sample maximizes the reward [\[ZHGSY20;](#page-10-3) [AAFK20;](#page-9-6) [NIK20\]](#page-10-4). And in Bayesian reinforcement learning, state-transition probabilities are sampled from the Dirichlet posterior over past observed states [\[Str00;](#page-10-5) [ORR13\]](#page-10-6). Dirichlet posterior sampling can also be used for data synthesis. Suppose that we have a histogram

 (x_1, \ldots, x_d) of actual data. An approximate discrete distribution of this histogram can be obtained by

31 drawing a sample Y from Dirichlet $(x_1 + \alpha_1, \ldots, x_d + \alpha_d)$, where $\alpha_1, \ldots, \alpha_d$ are prior parameters.

Then synthetic data is produced by repeatedly drawing from Multinomial (Y) . There are many

studies on data synthesis that followed this approach [\[AV08;](#page-9-7) [MKAGV08;](#page-9-8) [RWZ14;](#page-10-7) [PG14;](#page-10-8) [SJGLY17\]](#page-10-9).

In the above examples, the data that we integrate into these tasks might contain sensitive information.

Thus it is important to ask: how much of the information is protected from the Dirichlet samplings?

The goal of this study is to find an answer to this question.

 The mathematical framework of differential privacy (DP) [\[DMNS06\]](#page-9-9) allows us to quantify how much 38 the privacy of the Dirichlet posterior sampling is affected by the prior parameters $\alpha_1, \ldots, \alpha_d$. In the definition of DP, the privacy of a randomized algorithm is measured by how much its distribution changes upon perturbing a single data point of the input. Nonetheless, this notion might be too strict for the Dirichlet distribution, as a small perturbation of a near-zero parameter can cause a large distribution shift. Thus, it might be more appropriate to rely on one of several relaxed notions of DP, such as approximate differential privacy, Rényi differential privacy, or concentrated differential privacy. It is natural to wonder if the Dirichlet posterior sampling satisfies any of these definitions.

1.1 Overview of Our results

 This study focuses on the privacy and utility of Dirichlet posterior sampling. In summary, we provide a closed-form privacy guarantee of the Dirichlet posterior sampling, which in turn allows us to effectively analyze its utility in various settings.

 [§3](#page-4-0) Privacy. We study the role of the prior parameters in the privacy of the Dirichlet posterior sampling. Theorem [1](#page-4-1) is our main result, where we provide a guaranteed upper bound for truncated concentrated differential privacy (tCDP) of the Dirichlet posterior sampling. In addition, we convert the tCDP guarantee into an approximate differential privacy guarantee in Corollary [2.](#page-6-0)

 [§4](#page-6-1) Utility. Using the tCDP guarantee, we investigate the utility of Dirichlet posterior sampling applied in two specific applications:

- In Section [4.1,](#page-6-2) we consider one-time sampling from a Multinomial-Dirichlet distribution. But instead of directly sampling from this distribution, we sample from another distribution with larger prior parameters. The accuracy is then measured by the KL-divergence between the original and the private distributions.
- In Section [4.2,](#page-7-0) we use the Dirichlet posterior sampling for a private release of a normalized histogram. In this case, the accuracy is measured by the mean-squared error between the sample and the original normalized histogram.

 In both tasks, we compute the sample size that guarantees the desired level of accuracy. In the case of private histogram publishing, we also compare the Dirichlet posterior sampling to the Gaussian mechanism.

1.2 Related work

 There are several studies on the differential privacy of posterior sampling. Wang, Fienberg, and Smola [\[WFS15\]](#page-10-10) showed that any posterior sampling with the log-likelihood bounded by B is $4B$ - differentially private. However, the likelihoods that we study are not bounded away from zero; they 69 have the form $\prod_i p_i^{x_i}$ which becomes small when one of the p_i 's is close to zero. Dimitrakakis, Nelson, Zhang, Mitrokotsa, and Rubinstein [\[DNZMR17\]](#page-9-10) showed that if the condition on the log-likelihood is relaxed to the Lipschitz continuity with high probability, then one can obtain the approximate DP. Nonetheless, with the Dirichlet density, it is difficult to compute the probability of events in which the Lipschitz condition is satisfied.

74 In the case that the sufficient statistics x has finite ℓ^1 -sensitivity, Foulds, Geumlek, Welling and Chaudhuri [\[FGWC16\]](#page-9-11) suggested adding Laplace noises to x. Suppose that y is the output; they ⁷⁶ showed that sampling from $p(\theta|\mathbf{y})$ is differentially private and as asymptotically efficient as sampling 77 from $p(\theta|\mathbf{x})$. However, for a small sample size, the posterior over the noisy statistics might be too far away from the actual posterior. Bernstein and Sheldon [\[BS18\]](#page-9-12) thus proposed to approximate the 79 joint distribution $p(\theta, \mathbf{x}, \mathbf{y})$ using Gibbs sampling, which is then integrated over x to obtain a more accurate posterior over y.

⁸¹ Geumlek, Song, and Chaudhuri [\[GSC17\]](#page-9-13) were the first to study the posterior sampling with the

82 RDP. Even though they provided a general framework to find (λ, ϵ) -RDP guarantees for exponential

83 families, explicit forms of ϵ and the upper bound of λ were not given. In contrast, our tCDP guarantees 84 of the Dirichlet posterior sampling imply an explicit expression for ϵ , and also an upper bound for λ .

85 The privacy of data synthesis via sampling from Multinomial(Y), where Y is a discrete distri-⁸⁶ bution drawn from the Dirichlet posterior, was first studied by Machanavajjhala, Kifer, Abowd, ⁸⁷ Gehrke, and Vilhuber [\[MKAGV08\]](#page-9-8). They showed that the data synthesis is (ε,δ)-*probabilistic* DP, 88 which implies (ε, δ) -approximate DP. However, as their privacy analysis includes the sampling from 89 Multinomial(Y), their privacy guarantee depends on the number of synthetic samples. In contrast, ⁹⁰ we show that the one-time sampling from the Dirichlet posterior is approximate DP, which by the 91 post-processing property allows us to sample from Multinomial(Y) as many times as we want while ⁹² retaining the same privacy guarantee. ⁹³ The Dirichlet mechanism was first introduced by Gohari, Wu, Hawkins, Hale, and Topcu [\[GWHHT21\]](#page-9-14).

94 Originally, the Dirichlet mechanism takes a discrete distribution $\mathbf{p} := (p_1, \ldots, p_d)$ and draws one 95 sample Y ∼ Dirichlet(rp_1, \ldots, rp_d). Note the absence of the prior parameters, which makes Y an unbiased estimator of p. But this comes with a cost, as the worst case of privacy violation occurs when almost all of the parameters are close to zero. The authors avoided this issue by restricting 98 the input space to a subset of the unit simplex, with some of the p_i 's bounded below by a fixed positive constant. This results in complicated expressions for the privacy guarantees as they involve a minimization problem over the restricted domain. In this study, we take a different approach by adding prior parameters to the Dirichlet mechanism. As a result, we obtain a biased algorithm that requires no assumption on the input space and has simpler forms of privacy guarantees.

¹⁰³ 1.3 Notations

104 We let $\mathbb{R}^d_{\geq 0}$ be the set of d-tuples of non-negative real numbers and $\mathbb{R}^d_{>0}$ be the set of d-tuples of 105 positive real numbers. We assume that all vectors are d-dimensional where $d \geq 2$. The notations for 106 all vectors are always in bold. Specifically, $\mathbf{x} := (x_1, \dots, x_d) \in \mathbb{R}_{\geq 0}^d$ consists of sample statistics of 107 the data and $\alpha := (\alpha_1, \dots, \alpha_d) \in \mathbb{R}_{>0}^d$ consists of the prior parameters. The vector $\mathbf{p} := (p_1, \dots, p_d)$ 108 always satisfies $\sum_i p_i = 1$. The number of observations is always N. We also denote $x_0 \coloneqq \sum_i x_i$ 109 and $\alpha_0 := \sum_i \alpha_i$. For any vectors \mathbf{x}, \mathbf{x}' and scalar $r > 0$, we write $\mathbf{x} + \mathbf{x}' := (x_1 + x'_1, \dots, x_d + x'_d)$ 110 and $r\mathbf{x} \coloneqq (rx_1, \dots, rx_d)$. For any positive reals x and x', the notation $x \propto x'$ means $x = Cx'$ for some constant $C > 0$, $x \approx x'$ means $cx' \le x \le Cx'$ for some $c, C > 0$, and $x \le x'$ means $x \le Cx'$ 111 for some $C > 0$. Lastly, $\|\mathbf{x}\|_{\infty} := \max_i |x_i|$ is the ℓ^{∞} norm of **x**.

¹¹³ 2 Background

¹¹⁴ 2.1 Privacy models

115 **Definition 2.1** (Pure and Approximate DP [\[DMNS06\]](#page-9-9)). A randomized mechanism $M : \mathcal{X}^n \to \mathcal{Y}$ 116 is (ε, δ) -differentially private $((\varepsilon, \delta)$ -DP) if for any datasets x, x' differing on a single entry, and all 117 events $E \subset \mathcal{Y}$,

$$
\mathbb{P}[M(x) \in E] \le e^{\varepsilon} \mathbb{P}[M(x') \in E] + \delta.
$$

118 If M is $(\varepsilon, 0)$ -DP, then we say that it is ε -differential privacy $(\varepsilon$ -DP).

¹¹⁹ The term *pure differential privacy* (pure DP) refers to -differential privacy, while *approximate* 120 *differential privacy* (approximate DP) refers to (ε, δ) -DP when $\delta > 0$.

¹²¹ In contrast to pure and approximate DP, the next definitions of differential privacy are defined in terms of the Rényi divergence between $M(x)$ and $M(x')$:

123 **Definition 2.2** (Rényi Divergence [\[Rén61\]](#page-10-11)). Let P and Q be probability distributions. For $\lambda \in (1, \infty)$ 124 the Rényi divergence of order λ between P and Q is defined as

$$
D_{\lambda}(P\|Q) \coloneqq \frac{1}{\lambda - 1} \log \int P(y)^{\lambda} Q(y)^{1 - \lambda} dy = \frac{1}{\lambda - 1} \log \left(\mathop{\mathbb{E}}_{y \sim P} \left[\frac{P(y)^{\lambda - 1}}{Q(y)^{\lambda - 1}} \right] \right)
$$

125 **Definition 2.3** (tCDP and zCDP [\[BDRS18;](#page-9-15) [BS16\]](#page-9-16)). A randomized mechanism $M : \mathcal{X}^n \to \mathcal{Y}$ is 126 ω -truncated ρ -concentrated differentially private ((ρ , ω)-tCDP) if for any datasets x, x' differing on a

127 single entry and for all $\lambda \in (1, \omega)$,

$$
D_{\lambda}(M(x)||M(x')) \leq \lambda \rho.
$$

128 If M is (ρ, ∞) -tCDP, then we say that it is ρ -zero-concentrated differential privacy (ρ -zCDP).

129 Note that both tCDP and zCDP have the composition and post-processing properties. Intuitively, ρ controls the expectation and standard deviation of the privacy loss random variable: $Z = \log \frac{P[M(x)=Y]}{P[M(x')=Y]},$ 131 where Y has density $M(x)$, and ω controls the number of standard deviations for which Z concen-132 trates like a Gaussian. A smaller ρ and larger ω correspond to a stronger privacy guarantee. It turns ¹³³ out that tCDP implies approximate DP:

134 Lemma 1 (From tCDP to Approximate DP [\[BDRS18\]](#page-9-15)). Let $\delta > 0$. If M is a (ρ, ω) -tCDP mechanism, 135 *then it also satisfies* (ε, δ) -DP with

$$
\varepsilon = \begin{cases} \rho + 2\sqrt{\rho \log(1/\delta)} & \text{if } \log(1/\delta) \le (\omega - 1)^2 \rho \\ \rho \omega + \frac{\log(1/\delta)}{\omega - 1} & \text{if } \log(1/\delta) > (\omega - 1)^2 \rho \end{cases}.
$$

¹³⁶ 2.2 Dirichlet distribution

- 137 For $\alpha \in \mathbb{R}^d_{>0}$, the Dirichlet distribution Dirichlet (α) is a continuous distribution of d-dimensional
- 138 probability vectors i.e. vectors whose coordinate sum is equal to 1. The density function of Y \sim 139 Dirichlet(α) is given by:

$$
p(\mathbf{y}) = \frac{1}{B(\boldsymbol{\alpha})} \prod_{i=1}^{d} y_i^{\alpha_i - 1},
$$

140 where $B(\alpha)$ is the *beta function*, which can be written in terms of the gamma function:

$$
B(\alpha) = \frac{\prod_{i} \Gamma(\alpha_i)}{\Gamma(\sum_{i} \alpha_i)}.
$$
 (1)

¹⁴¹ 2.3 Dirichlet posterior sampling

142 We consider the prior Dirichlet(α) and the likelihood of the form $p(\mathbf{x}|\mathbf{y}) \propto \prod_{i=1}^{d} y_i^{x_i}$ where $x \in \mathbb{R}^d_{\geq 0}$ consists of sample statistics of the dataset. The *Dirichlet posterior sampling* is a one-time ¹⁴⁴ sampling:

$$
\mathbf{Y} \sim \text{Dirichlet}(\mathbf{x} + \boldsymbol{\alpha}).
$$

145 There is a modification of the sampling which introduces a concentration parameter $r > 0$, and 146 instead we sample from Dirichlet($r\mathbf{x} + \alpha$) [\[GSC17;](#page-9-13) [GWHHT21\]](#page-9-14). Smaller values of r make the 147 sampling more private, and larger values of r make Y a closer approximation of x. Even though the 148 case $r = 1$ is the main focus of this study, our main privacy results can be easily extended to other 149 values of r as we will see at the end of Section [3.1.](#page-4-2)

150 Consider a special case where $x = p$ is an empirical distribution derived from the dataset, and we 151 want Y to be a private approximation of p; the sampling Y \sim Dirichlet(rp + α) is called the ¹⁵² *Dirichlet mechanism* [\[GWHHT21\]](#page-9-14). It is interesting to note that the Dirichlet mechanism is a form of 153 the exponential mechanism [\[MT07\]](#page-10-12): let $r > 0$ be the privacy parameter, Dirichlet(α) be the prior, ¹⁵⁴ and the negative KL-divergence be the score function of the exponential mechanism. Then the output ¹⁵⁵ Y of this mechanism is distributed according to the following density function:

$$
\frac{\exp(-r \operatorname{D}_{\text{KL}}(\mathbf{p}, \mathbf{y})) \prod_i y_i^{\alpha_i - 1}}{\int \exp(-r \operatorname{D}_{\text{KL}}(\mathbf{p}, \mathbf{y})) \prod_i y_i^{\alpha_i - 1} d\mathbf{y}} \propto \exp\left(r \sum_{i, p_i \neq 0} p_i \log(y_i/p_i)\right) \prod_i y_i^{\alpha_i - 1}
$$

$$
\propto \prod_{i, p_i \neq 0} y_i^{rp_i} \prod_i y_i^{\alpha_i - 1} = \prod_i y_i^{rp_i + \alpha_i - 1},
$$

156 which is exactly the density function of Dirichlet($r\mathbf{p} + \alpha$).

¹⁵⁷ 2.4 Polygamma functions

¹⁵⁸ In most of this study, we take advantage of several nice properties of the log-gamma function and its 159 derivatives. Specifically, $\psi(x) \coloneqq \frac{d}{dx} \log \Gamma(x)$ is concave and increasing, while its derivative $\psi'(x)$ is 160 positive, convex, and decreasing. In addition, ψ' can be approximated by the reciprocals:

$$
\frac{1}{x} + \frac{1}{2x^2} < \psi'(x) < \frac{1}{x} + \frac{1}{x^2},
$$

161 which implies that $\psi'(x) \approx \frac{1}{x^2}$ as $x \to 0$ and $\psi'(x) \approx \frac{1}{x}$ as $x \to \infty$.

¹⁶² 3 Main privacy results

¹⁶³ 3.1 Truncated concentrated differential privacy

Theorem 1. Let $\alpha \in \mathbb{R}_{>0}^d$ and $\alpha_m \coloneqq \min_i \alpha_i$. Let $\gamma \in (0, \alpha_m)$. Let $\Delta_2, \Delta_\infty > 0$ be constants that *satisfy* $\sum_i (x_i - x'_i)^2 \leq \tilde{\Delta}_2^2$ and $\max_i |x_i - x'_i| \leq \tilde{\Delta}_{\infty}$ whenever $\mathbf{x}, \mathbf{x}' \in \mathbb{R}_{\geq 0}^2$ are sample statistics *of any two datasets differing on a single entry. The one-time sampling from* Dirichlet $(x + \alpha)$ *is* (ρ, ω) -tCDP, where $\omega = \frac{\gamma}{\Delta_{\infty}} + 1$ and

$$
\rho = \frac{1}{2} \Delta_2^2 \psi'(\alpha_m - \gamma). \tag{2}
$$

168 Note that (ρ, ∞) -tCDP is not obtainable, as the ratio between two Dirichlet densities blows up as 169 $\omega \to \infty$. We present here a short proof that skips some calculations (see Appendix 1 for a full proof).

proof. Consider any $\lambda \in \left(1, \frac{\gamma}{\Delta_{\infty}} + 1\right)$. Let $\mathbf{u} := \mathbf{x} + \boldsymbol{\alpha}$ and $\mathbf{u}' := \mathbf{x}' + \boldsymbol{\alpha}'$. Let $P(\mathbf{y})$ be the density 171 of Dirichlet (u) and $P'(\mathbf{y})$ be the density of Dirichlet (u'). A quick calculation shows that:

$$
\mathbb{E}_{\mathbf{y} \sim P(\mathbf{y})} \left[\frac{P(\mathbf{y})^{\lambda - 1}}{P'(\mathbf{y})^{\lambda - 1}} \right] = \frac{B(\mathbf{u}')^{\lambda - 1}}{B(\mathbf{u})^{\lambda - 1}} \cdot \frac{B(\mathbf{u} + (\lambda - 1)(\mathbf{u} - \mathbf{u}'))}{B(\mathbf{u})}.
$$
 (3)

¹⁷² We take the logarithm on both sides and apply the second-order Taylor expansion to the following 173 $G(u_i, u'_i)$ and $H(u_i, u'_i)$ terms that appear on the right-hand side. As a result, there exist ξ between 174 $u_i + (\lambda - 1)(u_i - u'_i)$ and u_i , and ξ' between u_i and u'_i such that

$$
G(u_i, u'_i) := (\lambda - 1)(\log \Gamma(u'_i) - \log \Gamma(u_i))
$$

= -(\lambda - 1)(x_i - x'_i)\psi(u_i) + \frac{1}{2}(\lambda - 1)(x_i - x'_i)^2\psi'(\xi')

H(u_i, u'_i) := \log \Gamma(u_i + (\lambda - 1)(u_i - u'_i)) - \log \Gamma(u_i)

= (\lambda - 1)(x_i - x'_i)\psi(u_i) + \frac{1}{2}(\lambda - 1)^2(x_i - x'_i)^2\psi'(\xi), (5)

175 Note that ψ' is increasing. If $x_i > x'_i$, then ξ and ξ' are bounded below by $u'_i \ge \alpha_m$. On the 176 other hand, if $x_i \leq x'_i$, then ξ and ξ' are bounded below by $u_i - (\lambda - 1)|u_i - u'_i|$. The condition 177 $\lambda < \frac{\gamma}{\Delta_{\infty}} + 1$ guarantees that $u_i - (\lambda - 1)|u_i - u'_i| > \alpha_m - \gamma$. All cases considered, we have

$$
G(u_i, u'_i) + H(u_i, u'_i) \le \frac{1}{2} ((\lambda - 1) + (\lambda - 1)^2)(x_i - x'_i)^2 \psi'(\alpha_m - \gamma)
$$

= $\frac{1}{2} \lambda (\lambda - 1)(x_i - x'_i)^2 \psi'(\alpha_m - \gamma).$

178 Denoting $u_0 := \sum_i u_i$ and $u'_0 := \sum_i u'_i$, the same argument shows that $G(u_0, u'_0) + H(u_0, u'_0) > 0$. ¹⁷⁹ Therefore,

$$
D_{\lambda}(P(\mathbf{y})||P'(\mathbf{y})) = \frac{1}{\lambda - 1} \left(\sum_{i} (G(u_i, u'_i) + H(u_i, u'_i)) - G(u_0, u'_0) - H(u_0, u'_0) \right)
$$

$$
< \frac{1}{\lambda - 1} \sum_{i} (G(u_i, u'_i) + H(u_i, u'_i))
$$

$$
\leq \frac{1}{2} \lambda \sum_{i} (x_i - x'_i)^2 \psi'(\alpha_m - \gamma) \leq \frac{1}{2} \lambda \Delta_2^2 \psi'(\alpha_m - \gamma).
$$

Figure 1: Left: the actual values of $\rho = \frac{1}{2} D_2(P || P')$ and the worst case $(\rho, 2)$ -tCDP guarantees [\(2\)](#page-4-3) at $\Delta_2^2 = \Delta_{\infty} = 1$. Here, P and P' are Dirichlet posterior densities over $\mathbf{x} = (11, 8, 65, 25, 38, 0)$, ${\bf x}' = (11, 8, 65, 25, 38, 1)$, and ${\bf \alpha} = (\alpha, \dots, \alpha)$. Right: comparison between (ε, δ) -DP guarantees of the Dirichlet posterior samplings [\(8\)](#page-6-3) with different uniform priors: $\alpha = (\alpha, \dots, \alpha)$.

¹⁸⁰ The guaranteed upper bound [\(2\)](#page-4-3) is independent of the sample statistics. As a result, the bound applies 181 even in worst settings i.e., when $x_i = 0$ and $x'_i = \Delta_{\infty}$, or vice versa, for some *i*. As we can see in 182 Figure [1,](#page-5-0) the upper bound is a close approximation to the actual value of ρ when $x_6 = 0$ and $x'_6 = 1$. 183 However, being a sample independent bound, the difference becomes substantial when all x_i 's are ¹⁸⁴ large. There is one way to get around this issue: if there is no privacy violation in assuming that 185 the sample statistics are always bounded below by some threshold τ , then we can incorporate the 186 threshold into the prior (thus $\psi'(\alpha_m - \gamma)$ in [\(2\)](#page-4-3) is replaced by $\psi'(\alpha_m + \tau - \gamma)$).

187 The parameter γ allows us to adjust the moment bound ω as desired. Even though a higher ω usually 188 leads to a better privacy guarantee, there are two downsides to picking γ close to α_m in this case. 189 First, note that ρ contains $\psi'(\alpha_m - \gamma)$; as $\gamma \to \alpha_m$, the value of ρ diverges to ∞ , leading to a weaker privacy guarantee instead. Second, as the Taylor approximation [\(5\)](#page-4-4) is accurate when u_i is close to 191 $u_i + (\lambda - 1)(u_i - u'_i)$, having a large value of λ would push the guaranteed upper bound away from 192 the actual privacy loss. Thus it is recommended to pick γ so that $\gamma/\Delta_{\infty} \geq 1$ and $\alpha_m - \gamma \gg 0$. 193 Alternatively, we can choose the value of γ that minimizes ε when converting from tCDP to (ε, δ) -DP ¹⁹⁴ using Lemma [1—](#page-3-0)this method will be explored in the next subsection.

195 Theorem [1](#page-4-1) can be easily applied to sampling from Dirichlet($r\mathbf{x} + \alpha$). Replacing x with rx, we have 196 Δ_2 replaced by $r\Delta_2$ and Δ_∞ replaced by $r\Delta_\infty$. Consequently, the sampling is $(\rho, \frac{\gamma}{r\Delta_\infty} + 1)$ -tCDP, 197 where $\rho = \frac{1}{2}r^2\Delta_2^2\psi'(\alpha_m - \gamma)$. In Appendix 4, we analyze the scaling of r in conjunction with α_m 198 at a fixed privacy budget ρ .

¹⁹⁹ 3.2 Approximate differential privacy

200 We now convert the tCDP guarantee to an approximate DP guarantee. Let $\delta \in (0, 1)$. Using Lemma [1,](#page-3-0) 201 the Dirichlet posterior sampling with Dirichlet(α) as the prior is (ε, δ) -DP with

$$
\varepsilon = \begin{cases}\n\rho(\gamma) + 2\sqrt{\rho(\gamma)\log(1/\delta)} & \text{if } \log(1/\delta) \le \gamma^2 \rho(\gamma)/\Delta_{\infty}^2 \\
\rho(\gamma)\left(\frac{\gamma}{\Delta_{\infty}} + 1\right) + \frac{\log(1/\delta)\Delta_{\infty}}{\gamma} & \text{if } \log(1/\delta) > \gamma^2 \rho(\gamma)/\Delta_{\infty}^2\n\end{cases},\n\tag{6}
$$

202 where $\rho(\gamma) = \frac{1}{2} \Delta_2^2 \psi'(\alpha_m - \gamma)$.

203 We try to minimize ϵ by adjusting the value of γ . First, we consider the case $\log(1/\delta) \le \gamma^2 \rho(\gamma)/\Delta_{\infty}^2$. 204 Since $\rho(\gamma)$ is a strictly increasing function of γ , both $\rho(\gamma) + 2\sqrt{\rho(\gamma) \log(1/\delta)}$ and $\gamma^2 \rho(\gamma)/\Delta_{\infty}^2$ 205 are both strictly increasing function of γ . Therefore, *ε* is minimized at the minimum possible 206 value of γ in this case, that is, at the unique γ_M that satisfies $\log(1/\delta) = \gamma_M^2 \rho(\gamma_M)/\Delta_{\infty}^2 =$ 207 $\frac{1}{2} \gamma_M^2 \Delta_2^2 \psi'(\alpha_m - \gamma_M)/\Delta_\infty^2$.

208 Now we consider the second case, when $\gamma < \gamma_M$. As $\rho(\gamma)$ is an increasing positive convex function 209 of γ , the function

$$
f(\gamma) := \frac{1}{2} \Delta_2^2 \psi'(\alpha_m - \gamma) \left(\frac{\gamma}{\Delta_{\infty}} + 1\right) + \frac{\log(1/\delta) \Delta_{\infty}}{\gamma}; \qquad \gamma \in (0, \gamma_M],
$$
 (7)

210 is also convex in γ , and thus has a unique minimizer $\gamma_m \in (0, \gamma_M]$. Comparing to the first case, we 211 have $f(\gamma_m) \le f(\gamma_M) = \rho(\gamma_M) + 2\sqrt{\rho(\gamma_M) \log(1/\delta)}$. We then conclude that $\varepsilon = f(\gamma_m)$.

Theorem 2. *Let* $\alpha \in \mathbb{R}^2_{>0}$ *and denote* $\alpha_m = \min_i \alpha_i$. *Let* $\Delta_2, \Delta_\infty > 0$ *be constants that satisfy* $\sum_i (x_i - x'_i)^2 \leq \Delta_2^2$ *and* $\max_i |x_i - x'_i| \leq \Delta_\infty$ *whenever* $\mathbf{x}, \mathbf{x}' \in \mathbb{R}^d_{>0}$ *are sample stat* 213 $\sum_i (x_i - x'_i)^2 \leq \Delta_2^2$ and $\max_i |x_i - x'_i| \leq \Delta_\infty$ whenever $\mathbf{x}, \mathbf{x}' \in \mathbb{R}_{\geq 0}^d$ are sample statistics of any 214 *two datasets differing on a single entry. For any* $\delta \in (0,1)$, let γ_M *be the solution to the equation* 215 $log(1/\delta) = \frac{1}{2}\gamma^2 \Delta_2^2 \psi'(\alpha_m - \gamma)/\Delta_{\infty}^2$. The one-time sampling from Dirichlet($\mathbf{x} + \boldsymbol{\alpha}$) is (ε, δ) -DP, ²¹⁶ *where*

$$
\varepsilon = \min_{\gamma \in (0, \gamma_M]} f(\gamma). \tag{8}
$$

217 Figure [1](#page-5-0) shows how δ decays as a function of $ε$ at three different values of $α_m$.

²¹⁸ 4 Utility

²¹⁹ Using the results from the previous section, we analyze the Dirichlet posterior sampling's utility in ²²⁰ two specific tasks.

²²¹ 4.1 Multinomial-Dirichlet sampling

222 Suppose that we are observing N trials, each of which has d possible outcomes. For each $i \in$ 223 $\{1, \ldots, d\}$, let x_i be the number of times the *i*-th outcome was observed. Then we have the 224 multinomial likelihood $p(\mathbf{x}|\mathbf{y}) \propto \prod_i y_i^{x_i}$. From this, we sample from the Dirichlet posterior:

$$
\mathbf{Y} \sim \text{Dirichlet}(\mathbf{x} + \boldsymbol{\alpha}).\tag{9}
$$

225 Suppose that we want to sample from a true distribution $P_X \sim$ Dirichlet($x + \alpha$), but for privacy example from $Q_x \sim \text{Dirichlet}(\mathbf{x} + \alpha')$ where $\alpha'_i > \alpha_i$ for all i. The utility of the 227 privacy scheme is then measured by the KL-divergence between P_x and Q_x . Assuming that x is an 228 observation of Multinomial (p) , the following Theorem tells us that, on average, the KL-divergence 229 is small when the sample size is large, and the p_i 's are evenly distributed.

230 Theorem 3. Let $\mathbf{p} := (p_1, \ldots, p_d)$ where $p_i > 0$ for all i and $\sum_i p_i = 1$. Define a random *variable* **X** ∼ Multinomial(**p**). Let P **x** ∼ Dirichlet(**X** + α) and Q **x** ∼ Dirichlet(**X** + α') where $\alpha'_i \geq \alpha_i \geq 1$ *for all i. The following estimate holds:*

$$
\mathbb{E}_{\mathbf{X}}[\mathcal{D}_{\mathrm{KL}}(P_{\mathbf{X}} \| Q_{\mathbf{X}})] \le \frac{1}{N+1} \sum_{i} (\alpha_i' - \alpha_i)^2 \cdot \frac{1}{p_i}.
$$
 (10)

233 The proof is given in Appendix 2. Let us consider a simple privacy scheme where we fix $s > 0$ and 234 let $\alpha'_i = \alpha_i + s$ for all *i*. Thus [\(10\)](#page-6-4) becomes:

$$
\mathbb{E}_{\mathbf{X}}[\mathcal{D}_{\mathrm{KL}}(P_{\mathbf{X}} \| Q_{\mathbf{X}})] \le \frac{G(\mathbf{p})s^2}{N+1},\tag{11}
$$

235 where $G(\mathbf{p}) := \sum_i 1/p_i$. Now we take into account the privacy parameters. Let $\rho = \Delta_2^2 \psi'(\alpha_m - \gamma)$ 236 and $ρ' = Δ₂²ψ'(α'_{m} - γ)$, where $α_m = min_i α_i$, $α'_{m} = min_i α'_{i}$, and $γ < α_m$. Here, we approximate 237 the values of $\psi'(\alpha_m - \gamma)$ and $\psi'(\alpha'_m - \gamma)$ under two regimes:

238 **High-privacy regime:** $\alpha'_m - \gamma > 1$. We have $\psi'(\alpha'_m - \gamma) \approx 1/(\alpha'_m - \gamma)$, which implies 239 $\alpha'_m - \gamma \approx \Delta_2^2/\rho'$. We also have $\alpha_m - \gamma \approx \Delta_2^2/\rho$ for $\alpha_m - \gamma \ge 1$ and $\alpha_m - \gamma > (\alpha_m - \gamma)^2 \approx \Delta_2^2/\rho$ 240 for $\alpha - \gamma < 1$. Thus we have the following bound for the right-hand side of [\(11\)](#page-6-5):

$$
\frac{G(\mathbf{p})s^2}{N+1} = \frac{G(\mathbf{p})(\alpha_m' - \alpha_m)^2}{N+1} \lesssim \frac{\Delta_2^4 G(\mathbf{p})}{N+1} \left(\frac{1}{\rho'} - \frac{1}{\rho}\right)^2 < \frac{\Delta_2^4 G(\mathbf{p})}{\rho'^2 (N+1)}.
$$
\n(12)

241 Consequently, we have $D_{KL}(P||Q) < \epsilon$ for $N = \Omega\left(\frac{\Delta_2^4 G(\mathbf{p})}{\rho'^2 \epsilon}\right)$.

242 Low-privacy regime: $1 > \alpha'_m - \gamma > 0$. This is similar as above, except we have $\alpha'_m - \gamma \approx$ 243 $\Delta_2/\rho'^{1/2}$ and $\alpha_m - \gamma \approx \Delta_2/\rho^{1/2}$. Similar computation as [\(12\)](#page-6-6) shows that $D_{KL}(P||Q) < \epsilon$ when 244 $N = \Omega\left(\frac{\Delta_2^2 G(\mathbf{p})}{\rho' \epsilon}\right)$.

245 We observe that, in both regimes, the sample size scales faster with respect to ϵ with a higher value of 246 $G(\mathbf{p})$, which is associated with a higher number of outcomes d, and more concentrated multinomial 247 parameter p; this agrees with the result of our simulation in Appendix 3. Moreover, for small ρ' the 248 sample size scales as $1/\rho^2$, while for large ρ' the sample size scales as $1/\rho'$.

²⁴⁹ 4.2 Private normalized histograms

250 Let $\mathbf{x} = (x_1, \dots, x_d)$ be a histogram of N observations and $\mathbf{p} := \mathbf{x}/N$. We can privatize p by 251 sampling a probability vector: Y \sim Dirichlet(x + α). Note that Y is a biased estimator of p.

252 Denoting $\alpha_0 := \sum_i \alpha_i$, the bias of each component of Y is given by $\mathbb{E}[Y] - p_i$. Hence,

$$
|\text{Bias}(Y_i)| = \left|\frac{x_i + \alpha_i}{N + \alpha_0} - p_i\right| = \frac{|x_i\alpha_0 - N\alpha_i|}{N(N + \alpha_0)} \le \frac{N\alpha_0}{N(N + \alpha_0)} = \frac{\alpha_0}{N + \alpha_0}
$$

253 Since $Y_i \sim \text{Beta}(x_i + \alpha_i, N + \alpha_0 - x_i - \alpha_i)$ is $\frac{1}{4(N + \alpha_0 + 1)}$ -sub-Gaussian [\[MA17\]](#page-9-17), we have,

$$
\mathbb{P}[|Y_i - p_i| > t + |\text{Bias}(Y_i)|] \le \mathbb{P}[|Y_i - \mathbb{E}[Y_i]| + |\text{Bias}(Y_i)| > t + |\text{Bias}(Y_i)|]
$$

= $\mathbb{P}[|Y_i - \mathbb{E}[Y_i]| > t]$
 $\le 2e^{-2t^2(N + \alpha_0 + 1)}.$

254 With the union bound, we plug in $t = \sqrt{\frac{\log(2d/\beta)}{2(N+\alpha_0+1)}}$, for any $\beta \in (0, 1)$, to obtain the following ²⁵⁵ accuracy guarantee of the private normalized histogram:

256 Theorem 4. Let Y ∼ Dirichlet($x + \alpha$), where $x \in \mathbb{R}_{\geq 0}^d$ and $\alpha \in \mathbb{R}_{>0}^d$, and $p \coloneqq x/N$. For any 257 $β ∈ (0, 1)$ *, with probability at least* $1 − β$ *, the following inequality holds:*

$$
\|\mathbf{Y} - \mathbf{p}\|_{\infty} \le \sqrt{\frac{\log(2d/\beta)}{2(N + \alpha_0 + 1)}} + \frac{\alpha_0}{N + \alpha_0}.
$$
\n(13)

.

258 Given $\epsilon > 0$, we use [\(13\)](#page-7-1) to find a lower bound for N that gives $||\mathbf{Y} - \mathbf{p}||_{\infty} < \epsilon$ w.p. $1 - \beta$ when 259 Y is sampled with ρ -tCDP. For simplicity, we consider a uniform prior: $\alpha_i = \alpha > 0$ for all i. 260 Thus, $\rho = \frac{1}{2} \Delta_2^2 \psi'(\alpha - \gamma)$, where γ might be chosen according to Corollary [2.](#page-6-0) We consider the two

²⁶¹ following regimes:

262 **High-privacy regime:** $\alpha - \gamma > 1$. In this case, $\psi'(\alpha - \gamma) \approx 1/(\alpha - \gamma)$. From $\rho = \frac{1}{2} \Delta_2^2 \psi'(\alpha - \gamma)$, 263 we have $\alpha \approx \frac{\Delta_2^2}{2\rho + \gamma}$. Replacing α_0 by $d\alpha$ in [\(13\)](#page-7-1) yields the sample size:

$$
N = \Omega \left(\frac{\log(2d/\beta)}{\epsilon^2} + \frac{d}{\epsilon} \left(\frac{\Delta_2^2}{2\rho} + \gamma \right) \right),\tag{14}
$$

²⁶⁴ for the desired accuracy.

265 Low-privacy regime: $\alpha - \gamma < 1$. This is the same as above, except now we have $\psi'(\alpha - \gamma) \approx$ 266 $1/(\alpha - \gamma)^2$, which implies $\alpha \approx \Delta_2/(2\rho)^{1/2} + \gamma$. The sample size that guarantees the desired ²⁶⁷ accuracy is:

$$
N = \Omega \left(\frac{\log(2d/\beta)}{\epsilon^2} + \frac{d}{\epsilon} \left(\frac{\Delta_2}{\sqrt{2\rho}} + \gamma \right) \right).
$$
 (15)

268 Let us compare this result to the Gaussian mechanism, which adds a noise $\mathbf{Z} \sim N(0, \sigma^2 I_d)$ to the 269 normalized histogram p directly. Thus the ℓ_2 -sensitivity in this case is Δ_2/N . We have that the 270 Gaussian mechanism is ρ-zCDP where $ρ = \frac{Δ^2}{2N^2σ^2}$ [\[BS16\]](#page-9-16). Using the same argument as above, with 271 probability at least $1 - \beta$, the following inequality holds for all *i*:

$$
\|\mathbf{Z}\|_{\infty} \le \sqrt{\frac{\log(2d/\beta)\Delta_2^2}{N^2 \rho}}.
$$
\n(16)

Figure 2: The ℓ^{∞} -accuracy, as a function of N, of Dirichlet posterior sampling ($\gamma = 1$) and Gaussian mechanisms for private normalized histograms ($\Delta_2^2 = 2$ and $\Delta_{\infty} = 1$). For each N, d and ρ , we generated the inputs x_1, \ldots, x_{200} , where $x_k \sim \text{Multinomial}(q_k)$ and $q_k \sim \text{Dirichlet}(5, \ldots, 5)$.

272 Hence, the sample size of $N = \Omega\left(\sqrt{\log(2d/\beta)\Delta_2^2/\rho\epsilon^2}\right)$ guarantees the desired accuracy. Compar-273 ing this to [\(14\)](#page-7-2), if we assume $\epsilon < 1$, the AM-GM inequality tells us that

$$
\frac{\log(2d/\beta)}{\epsilon^2} + \frac{d\Delta_2^2}{\rho \epsilon} > \frac{\log(2d/\beta)}{\epsilon^2} + \frac{\Delta_2^2}{\rho} \ge 2\sqrt{\frac{\log(2d/\beta)\Delta_2^2}{\rho \epsilon^2}}.
$$
 (17)

²⁷⁴ The inequality [\(17\)](#page-8-0) implies that the Gaussian mechanism requires less sample than the Dirichlet ²⁷⁵ mechanism in order to guarantee the same level of accuracy. The Gaussian mechanism is also better in the low-privacy regime as the ρ in [\(15\)](#page-7-3) satisfies $\sqrt{\rho} < \rho$ and $\Delta_2 \approx \Delta_2^2$, leading to the same 277 inequality [\(17\)](#page-8-0). Nonetheless, the decay in [\(16\)](#page-7-4) is linear in d, while that in [\(13\)](#page-7-1) has $\alpha_0 = d\alpha$ in 278 the denominators. This observation suggests that, when x is a sparse histogram i.e. when $N \leq d$, 279 the ℓ^{∞} -accuracy of the Dirichlet mechanism is smaller than that of the Gaussian mechanism. This 280 conclusion is supported by our simulation in Figure [2.](#page-8-1) We see that the ℓ^{∞} -accuracy of the Dirichlet 281 mechanism is smaller than that of the Gaussian mechanism for small N when $d = 1000$. The code ²⁸² for all experiments in this study can be found in the supplemental material.

²⁸³ Potential negative societal impacts

284 It is important to note that, when ρ becomes unacceptably large (e.g., $\rho = 10^4$), the sampling is far away from being private. Thus any organization that deploys the posterior sampling on sensitive data must not vacuously refer to this study and claim that its algorithm is private. It is the organization's responsibility to fully publish the prior parameters, and educate its users/customers on differential privacy and how the privacy guarantees are calculated.

 It is desirable that differentially private algorithms are accurate for the task at hand, especially when the data is used for important decision-making. Thus, one needs to make sure that there is enough sample to achieve the desired level of accuracy. For a large differentially private system, privacy budgets need to be allocated to the parts that require accurate outputs.

 Lastly, one must be careful with the choice of prior parameters; if a uniform prior is used, smaller groups will suffer a relatively larger statistical bias. As a result, private statistics of small populations (such as ethnic or racial minorities) will be relatively less accurate. One way to get around this issue is to (privately) impose larger prior parameters on larger populations.

References

Checklist

