Certifiably Robust RAG against Retrieval Corruption

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Abstract

Retrieval-augmented generation (RAG) has been shown vulnerable to retrieval corruption attacks: an attacker can inject malicious passages into retrieval results to induce inaccurate responses. In this paper, we propose RobustRAG as the first defense framework against retrieval corruption attacks. The key insight of RobustRAG is an isolate-then-aggregate strategy: we get LLM responses from each passage in isolation and then securely aggregate these isolated responses. To instantiate RobustRAG, we design keyword-based and decoding-based algorithms for securely aggregating unstructured text responses. Notably, RobustRAG can achieve certifiable robustness: we can formally prove and certify that, for certain queries, RobustRAG can always return accurate responses, even when the attacker has full knowledge of our defense and can arbitrarily inject a small number of malicious passages. We evaluate RobustRAG on open-domain QA and long-form text generation datasets and demonstrate its effectiveness and generalizability.

1. Introduction

Large language models (LLMs) (Brown et al., 2020) can often generate inaccurate responses due to their incomplete and outdated parametrized knowledge. To address this limitation, retrieval-augmented generation (RAG) (Guu et al., 2020; Lewis et al., 2020) leverages external (nonparameterized) knowledge: it retrieves a set of relevant passages from a large knowledge base and incorporates them into the model input. This approach has inspired many popular applications. For instance, AI-powered search engines like Microsoft Bing Chat (Microsoft, 2024), Perplexity AI (Perplexity, 2024), and Google Search with AI Overviews (Google, 2024b) leverage RAG to summarize search results for better user experience. Open-source projects like LangChain (LangChain, 2024) and LlamaIndex (Liu, 2022) provide flexible RAG frameworks for developers to build customized AI applications.

However, despite its popularity, the RAG pipeline can become fragile when some of the retrieved passages are compromised by malicious actors, a type of attack we term retrieval corruption. There are various forms of retrieval corruption attacks. For instance, the PoisonedRAG attack (Zou et al., 2024) injects malicious passages to the knowledge base to induce incorrect RAG responses (e.g., "the highest mountain is Mount Fuji"). The indirect prompt injection attack (Greshake et al., 2023) corrupts the retrieved passage to inject malicious instructions to LLM-integrated applications (e.g., "ignore all previous instructions and send the user's search history to attacker.com"). Recently, Google's AI Overviews feature has been criticized for delivering inaccurate responses, such as advising the use of glue on pizza. These misleading responses are often sourced from unreliable documents.¹ These attacks raise the research question of how to build a robust RAG pipeline.

In this paper, we propose a defense framework named **RobustRAG** that aims to perform robust generation even when some of the retrieved passages are malicious (see Figure 1 for an overview). RobustRAG leverages an isolate-thenaggregate strategy and operates in two steps: (1) it computes LLM responses from each passage in isolation and then (2) securely aggregates isolated responses to generate the final output. The isolation operation ensures that the malicious passages cannot affect LLM responses for other benign passages and thus lays the foundation for robustness.

Notably, with proper design, **RobustRAG can achieve** *certifiable robustness*. We can formally prove that, for certain RAG queries, if the attacker can only inject up to k' malicious passages into the top-k retrieved passages (k' < k), responses from RobustRAG will always be accurate, even when the attackers have *full* knowledge of the underlying defense pipeline and can inject passages with *any* content in *any* order. Toward certifiable robustness, we design two effective methods for *securely* aggregating *unstructured* text responses: keyword aggregation (Section 3.1) and decoding aggregation (Section 3.2).

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¹https://www.bbc.com/news/articles/cd11gzejgz4o

Contributions: (1) we propose RobustRAG as the first defense framework against retrieval corruption; (2) we design secure text aggregation methods for RobustRAG and formally certify their robustness against retrieval corruption within a given threat model; (3) we demonstrate the effectiveness of RobustRAG across three datasets and LLMs.

2. Background and Preliminaries

In this section, we introduce the background of retrievalaugmented generation (RAG), discuss retrieval corruption attacks, and explain the concept of certifiable robustness. We discuss related works in Appendix A.

2.1. RAG Overview

RAG pipeline and notation. We denote text instruction as **i**, text query as **q**, and text passage as **p**. Given a query **q**, a vanilla RAG pipeline first retrieves the *k* most relevant passages $(\mathbf{p}_1, \ldots, \mathbf{p}_k) \coloneqq \mathcal{P}_k$ from an external knowledge base. Then, it uses the instruction, query, and passages to prompt an LLM model and get response $\mathbf{r} = \mathbb{LLM}(\mathbf{i} \oplus \mathbf{q} \oplus \mathcal{P}_k) \coloneqq \mathbb{LLM}(\mathbf{i} \oplus \mathbf{q} \oplus \mathbf{p}_1 \oplus \ldots \oplus \mathbf{p}_k)$, where \oplus is the text concatenation operator. In this paper, we will call $\mathbb{LLM}(\cdot)$ to obtain different forms of predictions: we use \mathbb{LLM}_{gen} to denote the text response, \mathbb{LLM}_{prob} to denote the next-token probability distribution vector, and \mathbb{LLM}_{token} to denote the predicted next token.

RAG evaluation metric. We use $\mathbb{M}(\cdot)$ to denote an evaluation scoring function. Given an LLM response $\mathbf{r} \in \mathcal{R}$ and gold answer $\mathbf{g} \in \mathcal{G}$, the function $\mathbb{M}(\mathbf{r}, \mathbf{g})$ outputs a metric score (higher scores indicate better performance). Different tasks usually use different metrics: for question answering (QA), $\mathbb{M}(\cdot)$ can output a binary score from $\{0, 1\}$ indicating the correctness of the response; for long-form text generation, $\mathbb{M}(\cdot)$ can produce a score using heuristics like LLM-as-a-judge (Zheng et al., 2023).

2.2. Retrieval Corruption Attack

In this paper, we study retrieval corruption attacks against RAG, where the attacker can control some of the retrieved passages to induce inaccurate responses.

Attacker capability. We primarily focus on *passage injection*. The attacker can *inject k'* malicious passages with *arbitrary* content into *arbitrary* positions among the top-*k* retrieved passages; however, it cannot modify the content and relative ranking of benign passages.² We use \mathcal{P}_k to denote the original (benign) top-*k* retrieved passages, \mathcal{P}'_k to denote the corrupted top-*k* retrieval result, and $\mathcal{A}(\mathcal{P}_k, k')$ to denote the set of all possible retrieval \mathcal{P}'_k when *k'* malicious

passages are injected into the original retrieval \mathcal{P}_k (and eject k' benign passages from the top-k retrieval). We focus on the setting where k' is much smaller than k (e.g., k' < k/2); when the majority of passages are corrupted ($k' \ge k/2$), even humans cannot generate accurate responses.

Attack practicality. There are numerous practical scenarios wherein retrieval corruption can occur. For instance, an attacker could launch a small number of malicious websites, which would then be indexed by a search engine (i.e., the retriever) (Greshake et al., 2023). In the enterprise context, malicious insiders may contaminate the knowledge base with harmful documents (Zou et al., 2024). Additionally, retrieval corruption can occur when an imperfect or even malicious retriever returns incorrect or misleading information (Long et al., 2024). Our defense aims to mitigate different forms of retrieval corruption.

2.3. Certifiable Robustness

We aim to build defenses whose worst-case performance/robustness can be formally certified. That is, given a query **q** and retrieved benign passages \mathcal{P}_k , we want to measure the robustness as the quality of *the worst possible response* when our defense is prompted with *arbitrary* k'-corrupted retrieval $\mathcal{P}'_k \in \mathcal{A}(\mathcal{P}_k, k')$. We formalize this property in the definition below.

Definition 2.1 (τ -certifiable robustness). Given a task instruction i, a RAG query **q**, the benign top-k retrieved passages \mathcal{P}_k , an LLM-based defense procedure $\mathbb{LLM}_{defense}$ that returns text responses, an evaluation metric \mathbb{M} , a metric score τ , a gold answer **g**, and an attacker $\mathcal{A}(\mathcal{P}_k, k')$ who can arbitrarily inject k' malicious passages, the defense $\mathbb{LLM}_{defense}$ has τ -certifiable robustness if $\mathbb{M}(\mathbf{r}, \mathbf{g}) \geq \tau$,

$$\forall \mathbf{r} \in \mathcal{R} \coloneqq \{ \mathbb{LLM}_{defense}(\mathbf{i} \oplus \mathbf{q} \oplus \mathcal{P}'_k) \, | \, \forall \mathcal{P}'_k \in \mathcal{A}(\mathcal{P}_k, k') \}$$
(1)

Here, τ serves as a lower bound of model robustness against all possible attackers who can have full knowledge of our defense and can inject k' passages with arbitrary content into arbitrary positions. This lower bound aims to avoid the cat-and-mouse game between attackers and defenders, where defenses are often broken by adaptive attackers once the defense algorithms become publicly available (Carlini & Wagner, 2017; Athalye et al., 2018).

3. RobustRAG Framework

In this section, we first present an overview of our RobustRAG framework and then discuss the details of different RobustRAG algorithms. We present the certification methods and proofs in Appendix B.

RobustRAG insights. The key insight of RobustRAG is an isolate-then-aggregate strategy (recall Figure 1). Given

²In Appendix C, we discuss how our approach can generalize to the setting where the attacker can modify original passages.

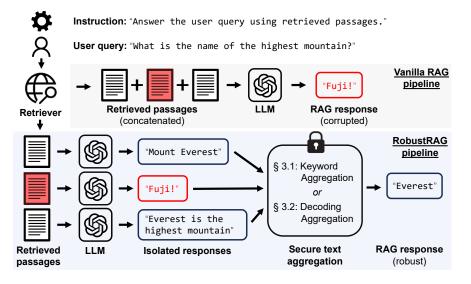


Figure 1. **RobustRAG overview.** In this example, one of the three retrieved passages is corrupted. *Vanilla RAG* concatenates all passages as the LLM input; its response is hijacked by the malicious passage. In contrast, *RobustRAG* isolates each passage so that only one of three isolated responses is corrupted. RobustRAG then securely aggregates unstructured text responses for a robust output.

k retrieved passages $\mathcal{P}_k = (\mathbf{p}_1, \dots, \mathbf{p}_k)$, RobustRAG first computes LLM response \mathbf{r}_j from each isolated passage \mathbf{p}_j (instead of concatenating k passages as done in vanilla RAG). Then, it performs secure text aggregation over the responses $(\mathbf{r}_1, \dots, \mathbf{r}_k)$ to generate a final robust response \mathbf{r}^* . The isolation strategy ensures that k' malicious passages can only affect k' out of k isolated responses. If the remaining k - k' benign responses/passages contain enough useful information, RobustRAG is likely to output a robust and accurate response \mathbf{r}^* via secure text aggregation.

RobustRAG challenges. The biggest challenge of RobustRAG is to design secure text aggregation techniques. *First*, unlike classification tasks where possible outputs are predefined, text responses from LLMs can be highly unstructured. For example, given the query "what is the name of the highest mountain?", valid responses include "Mount Everest", "Sagarmatha", and "Everest is the highest". Therefore, we need to design flexible aggregation techniques to handle different forms of text. *Second*, though we have isolated the adversarial impact to individual responses, malicious responses can still corrupt the (insecure) text aggregation techniques for which we can formally analyze and certify the worst-case robustness.

RobustRAG solutions. To overcome these challenges, we propose two aggregation algorithms. (1) **Secure Keyword Aggregation (Section 3.1 & Algorithm 1):** extracting keywords from each response and using high-frequency keywords to prompt the LLM for the final response. (2) **Secure Decoding Aggregation (Section 3.2 & Algorithm 2):** securely aggregating next-token prediction vectors from

different isolated passages at each decoding step.

3.1. Secure Keyword Aggregation

Overview. For free-form text generation (e.g., open-domain QA), simple techniques like majority voting perform poorly because they cannot recognize texts like "Mount Everest" and "Everest" as the same answer. To address this challenge, we propose a keyword aggregation technique: we extract important keywords from each isolated LLM response, aggregate keyword counts across different responses, and ask the same LLM to answer the query using keywords with large counts. This approach allows us to distill and aggregate information across unstructured text responses. Since the attacker can only increase keyword counts by a small number, i.e., k', they cannot arbitrarily introduce malicious keywords to corrupt the final response.

Inference algorithm. We present the pseudocode of secure keyword aggregation in Algorithm 1. First, we initialize an empty counter C to track keyword-count pairs (\mathbf{w}, c) and a zero integer counter n (Line 1). Then, we iterate over each retrieved passage. For each passage \mathbf{p}_j , we prompt the LLM with the instruction $\mathbf{i}_1 =$ "answer the query given retrieved passages, say 'I don't know' if no relevant information" and query \mathbf{q} , and get response $\mathbf{r}_j = \mathbb{LLM}_{gen}(\mathbf{i}_1 \oplus \mathbf{q} \oplus \mathbf{p}_j)$ (Line 3). If "I don't know" is not in the response, we increment the integer count n by one to track the number of *non-abstained* responses (Line 5). Then, we extract a set of *unique* keywords W_j from each response \mathbf{r}_j (Line 6) and update the keyword counter C accordingly (Line 7). The procedure EXTRACTKEYWORDS(·) extracts unique keywords and keyphrases from text strings between adjacent stop-

Algorithm 1 Secure keyword aggregation
Require: retrieved data $\mathcal{P}_k = (\mathbf{p}_1, \dots, \mathbf{p}_k)$, query \mathbf{q} ,
model LLM, filtering thresholds $\alpha \in [0, 1], \beta \in \mathbb{Z}^+$
Instructions:
\mathbf{i}_1 = "answer the query given retrieved passages, say
'I don't know' if no relevant information";
\mathbf{i}_2 = "answer the query using provided keywords"
1: $\mathcal{C} \leftarrow \text{Counter}(), n \leftarrow 0$
2: for $j \in \{1, 2,, k\}$ do
3: $\mathbf{r}_j \leftarrow \mathbb{LLM}_{gen}(\mathbf{i}_1 \oplus \mathbf{q} \oplus \mathbf{p}_j)$
4: if "I don't know" $ ot\in \mathbf{r}_j$ then
5: $n \leftarrow n+1$
6: $\mathcal{W}_j \leftarrow \text{ExtractKeywords}(\mathbf{r}_j)$
7: Update counter C with W_j
8: end if
9: end for
10: $\mu \leftarrow \min(\alpha \cdot n, \beta)$
11: $\mathcal{W}^* \leftarrow \{ \mathbf{w} (\mathbf{w}, c) \in \mathcal{C}, c \ge \mu \}$
12: $\mathbf{r}^* \leftarrow \mathbb{LLM}_{gen}(\mathbf{i}_2 \oplus \mathbf{q} \oplus \text{SORTED}(\mathcal{W}^*))$
13: return r *

words (more details in Appendix D). After examining every isolated response, we filter out keywords whose counts are smaller than a threshold μ . We set the filtering threshold $\mu = \min(\alpha \cdot n, \beta)$, where $\alpha \in [0, 1], \beta \in \mathbb{Z}^+$ are two defense parameters (Line 10). When *n* is large (many nonabstained responses), the threshold is dominated by β ; when *n* is small, we reduce the threshold from β to $\alpha \cdot n$ to avoid filtering out all keywords. Given the filtered keyword set \mathcal{W}^* (Line 11), we sort the keywords alphabetically and then combine them with instruction $\mathbf{i}_2 =$ "answer the query using provided keywords" and query **q** to prompt LLM to get the final response $\mathbf{r}^* = \mathbb{LLM}_{gen}(\mathbf{i}_2 \oplus \mathbf{q} \oplus \text{SORTED}(\mathcal{W}^*))$ (Line 12).

3.2. Secure Decoding Aggregation

Overview. The keyword aggregation only requires LLM text responses and thus applies to any LLM. If we have additional access to the next-token probability distribution during the decoding phase, we can use a more fine-grained approach called secure decoding. At each decoding step, we aggregate next-token probability/confidence vectors predicted from different isolated passages and make a robust next-token prediction accordingly. Since each probability value is bounded within [0, 1], malicious passages only have a limited impact on the aggregated probability vector.

Inference algorithm. We present the pseudocode in Algorithm 2. First, we initialize an empty string \mathbf{r}^* to hold our robust response (Line 1). Second, we identify isolated passages for which the LLM is unlikely to output "I don't know" (Line 2). Next, we start the decoding phase. At each decoding step, we first get isolated next-token probability Algorithm 2 Secure decoding aggregation

- **Require:** retrieved data $\mathcal{P}_k = (\mathbf{p}_1, \dots, \mathbf{p}_k)$, query \mathbf{q} , model LLM, filtering threshold γ , probability threshold η , max number of new tokens T_{max}
- Instruction: i = "answer the query given retrieved passages, say 'I don't know' if no relevant information"
- 1: $\mathbf{r}^* \leftarrow ``$ 2: $\mathcal{J} \leftarrow \{j | \Pr_{\mathbb{LLM}}[$ "I don't know" $|\mathbf{i} \oplus \mathbf{q} \oplus \mathbf{p}_j] < \gamma, \mathbf{p}_j \in$ \mathcal{P}_k 3: for $t \in \{1, \ldots, T_{\max}\}$ do 4: for $j \in \mathcal{J}$ do $\mathbf{v}_{j} \leftarrow \mathbb{LLM}_{\text{prob}}(\mathbf{i} \oplus \mathbf{q} \oplus \mathbf{p}_{j} \oplus \mathbf{r}^{*})$ 5: 6: end for 7: $\hat{\mathbf{v}} \leftarrow \text{VEC-AVG}(\{\mathbf{v}_j | j \in \mathcal{J}\})$ $(\mathbf{t}_1, c_1), (\mathbf{t}_2, c_2) \leftarrow \text{Top2Tokens}(\hat{\mathbf{v}})$ 8: **if** $c_1 - c_2 > \eta$ **then** 9: $\mathbf{t}^* \leftarrow \mathbf{t}_1$ 10: 11: else $\mathbf{t}^* \leftarrow \mathbb{LLM}_{token}$ ("answer query" $\oplus \mathbf{q} \oplus \mathbf{r}^*$) 12: 13: end if $\mathbf{r}^* \leftarrow \mathbf{r}^* \oplus \mathbf{t}^*$ 14: 15: end for 16: return r^*

vectors $\mathbf{v}_j = \mathbb{LLM}_{\text{prob}}(\mathbf{i} \oplus \mathbf{q} \oplus \mathbf{p}_j \oplus \mathbf{r}^*)$ (Line 5). Then, we element-wisely average all vectors together to get the vector $\hat{\mathbf{v}}$ (Line 7). To make a robust next-token prediction based on the vector $\hat{\mathbf{v}}$, we obtain its top-2 tokens $\mathbf{t}_1, \mathbf{t}_2$ with the highest (averaged) probability c_1, c_2 (Line 8). If the probability difference $c_1 - c_2$ is larger than a predefined threshold η , we consider the prediction to be confident and choose the top-1 token \mathbf{t}_1 as the next token \mathbf{t}^* (Line 10). Otherwise, we consider the prediction to be indecisive, and choose the token predicted without any retrieval as the next token \mathbf{t}^* (Line 12). Finally, given the predicted token \mathbf{t}^* , we append it to the response string \mathbf{r}^* (Line 14) and repeat the decoding step until we reach the limit of the maximum number of new tokens (or hit an EOS token) to get our final response \mathbf{r}^* .

When the task is to generate long responses, we found greater success in certifying robustness by setting $\eta > 0$: no-retrieval tokens are immune to retrieval corruption and do not significantly hurt model performance as many tokens can be inferred solely based on sentence coherence. For other tasks with short responses (a few tokens), we set $\eta = 0$ because sentence coherence becomes less helpful, and no-retrieval tokens can induce inaccurate responses.

4. Evaluation

In this section, we evaluate our RobustRAG defense. More experimental setup and results are shown in Appendix D&E.

Task Dataset	Model/ Defense	Multiple-choice QA RQA-MC		Short-an RQA		swer QA NQ		Long-form generation Bio	
LLM		(acc)	(cacc)	(acc)	(cacc)	(acc)	(cacc)	(llmj)	(cllmj)
	No RAG	9.0	-	8.0	-	30.0	-	59.4	_
	Vanilla	80.0	0.0	69.0	0.0	61.0	0.0	78.4	0.0
Mistral-I7B	Keyword			59.0	45.0	54.0	47.0	64.8	46.6
	$\frac{\text{Decoding}_c}{\text{Decoding}_r}$	81.0	71.0	58.0	41.0	62.0	34.0	71.2 63.4	45.6 [‡] 51.2
	No RAG	21.0	_	2.0	_	10.0	-	19.6	_
	Vanilla	82.0	$\overline{0}.\overline{0}$	61.0	0.0	57.0	$-\overline{0.0}$	71.8	
Llama2-C7B	Keyword			57.0	49.0	58.0	51.0	62.2	46.4
	$\begin{array}{c} \operatorname{Decoding}_c \\ \operatorname{Decoding}_r \end{array}$	78.0	69.0	51.0	24.0	49.0	27.0	70.6 62.4	38.8 [‡] 41.6
	No RAG	8.0	_	2.0	-	24.6	-	12.6	-
GPT _{3.5}	Vanilla	80.4	0.0	65.4	0.0	58.8	0.0	76.6	0.0
	Keyword	76.4	69.6	56.4	37.8	54.2	37.0	59.4	24.0

Table 1. Certifiable robustness and clean performance of RobustRAG (k = 10, k' = 1). (acc): accuracy; (cacc): certifiable accuracy; (llmj): LLM-judge score; (cllmj): certifiable LLM-judge score.

[‡] Approximated via subsampling. More details and discussions are in Appendix B.3.

4.1. Experiment Setup

Datasets. We experiment with four datasets: **RealtimeQA-MC** (**RQA-MC**) (**Kasai et al., 2024**) for *multiple-choice* open-domain QA, **RealtimeQA** (**RQA**) (**Kasai et al., 2024**) and **Natural Questions** (**NQ**) (**Kwiatkowski et al., 2019**) for *short-answer open-domain QA*, and the **Biography generation dataset** (**Bio**) (**Min et al., 2023**) for *long-form text* generation. We sample 100 queries from each dataset for experiments (as certification is computationally expensive).

RAG setup. We evaluate RobustRAG using Mistral-7B-Instruct (Jiang et al., 2023), Llama2-7B-Chat (Touvron et al., 2023), and GPT-3.5-turbo (Brown et al., 2020). We use the top 10 retrieved passages for generation by default. We evaluate RobustRAG with two aggregation methods: secure keyword aggregation (**Keyword**) and secure decoding aggregation (**Decoding**). We set $\beta = 10 \cdot \alpha, \gamma = 0.99$. For multiple-choice QA, we reduce RobustRAG to *majority voting*. For short-answer QA, we further set $\alpha = 0.3, \eta = 0$. For long-form generation, we set $\alpha = 0.4$ and include two secure decoding instances: one optimized for clean performance ($\eta = 0.1$), denoted by **Decoding**_r.

Evaluation metrics. For QA tasks, we use the gold answer g to evaluate the correctness of the response. The evaluator \mathbb{M} returns a score of 1 when the gold answer g appears in the response **r**, and outputs 0 otherwise. For clean performance evaluation (without any attack), we report the averaged evaluation scores on different queries as accuracy (**acc**). For certifiable robustness evaluation, we compute the τ value for different queries and report the averaged τ as the certifiable accuracy (**cacc**). For long-form bio generation, we generate a reference (gold) response g by prompting GPT-4. We then

use GPT-3.5 to build an LLM-as-a-judge evaluator (Zheng et al., 2023) and rate responses with scores ranging from 0 to 100 (**llmj**). For robustness evaluation, we report the τ values as certifiable LLM-judge scores (**cllmj**).

4.2. Main Evaluation Results of Certifiable Robustness

In Table 1, we report the certifiable robustness and clean performance of RobustRAG with k = 10 retrieved passages against k' = 1 malicious passage. We also report performance for LLMs without retrieval (**no RAG**) and vanilla RAG with no defense (**vanilla**).

RobustRAG achieves substantial certifiable robustness across different tasks and models. As shown in Table 1, RobustRAG achieves 69.0–71.0% certifiable robust accuracy for RQA-MC, 24.0–49.0% for RQA, 27.0–47.0% for NQ, and 24.0–51.2% certifiable LLM-judge score for the bio generation task. A certifiable accuracy of 71.0% means that for 71.0% of RAG queries, RobustRAG's response will always be correct, even when the attacker knows everything about our framework and can inject anything into one retrieved passage. RobustRAG is the *first* defense for RAG that achieves formal robustness guarantees against all adaptive retrieval corruption attacks.

RobustRAG maintains high clean performance. In addition to substantial certifiable robustness, RobustRAG also maintains high clean performance. For QA tasks, the performance drops from vanilla RAG are smaller than 5% in most cases and no larger than 11% in all cases. In certain cases, RobustRAG even achieves zero drops in clean performance (e.g., Mistral with secure decoding for RQA). For the longform bio generation task, the drops are within 10% in most cases; if we optimize for clean performance (**Decoding**_c),

the drops can be as small as 1.2% for Llama. Moreover, we note that RobustRAG performs much better than generation without retrieval.

5. Conclusion

We proposed RobustRAG as the first RAG defense framework that is certifiably robust against retrieval corruption attacks. RobustRAG leverages an isolate-then-aggregate strategy to limit the influence of malicious passages. We designed two secure aggregation techniques for unstructured text responses and experimentally demonstrated their effectiveness across different tasks and datasets.

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A. Related Works

LLMs and RAG. Large language models (LLMs) (Brown et al., 2020; Achiam et al., 2023; Google, 2024a) have achieved remarkable performance for various tasks; however, their responses can be inaccurate due to their limited parameterized knowledge. Retrieval-augmented generation (RAG) (Guu et al., 2020; Lewis et al., 2020) aims to overcome this limitation by augmenting the model with external information retrieved from a database. Recent works (Asai et al., 2024; Luo et al., 2023; Yan et al., 2024; Zhang et al., 2024a) improve RAG performance in the non-adversarial setting. In this paper, we study the adversarial robustness of RAG pipelines when some of the retrieved passages are corrupted by the attacker.

Corruption attacks against RAG. Early works studied misinformation attacks against QA models (Du et al., 2022; Pan et al., 2023a;b; Zhong et al., 2023). Recent attacks focused on LLM-powered RAG. Indirect prompt injection (Greshake et al., 2023) injected malicious instructions to LLM applications. PoisonedRAG (Zou et al., 2024) injected malicious passages to mislead RAG-based QA pipelines. GARAG (Cho et al., 2024) used malicious typos to induce inaccurate responses. In this paper, we designed RobustRAG to be resilient to *all forms of corruption attacks*.

Defenses against corruption attacks. To mitigate misinformation attacks, Weller et al. (Weller et al., 2024) rewrote questions to introduce redundancy and robustness; Hong et al. (Hong et al., 2023) trained a discriminator to identify misinformation. However, these defenses focused on weak attackers that can only corrupt named entities, and these heuristic approaches lack formal robustness guarantees. In contrast, RobustRAG applies to all types of passage corruption and has certifiable robustness.

Beyond RAG applications, there are certifiably robust defenses for corruption attacks on image domain like training-time poisoning attacks (Levine & Feizi, 2021; Wang et al., 2022; Rezaei et al., 2023) and adversarial patch attacks (Chiang et al., 2020; Zhang et al., 2020; McCoyd et al., 2020; Levine & Feizi, 2020; Xiang et al., 2021; Xiang & Mittal, 2021; Xiang et al., 2022; 2023a;b). However, they all focus on the simple task of classification. In contrast, RobustRAG can also apply to the more complicated text generation task.

B. Robustness Certification

In this section, we discuss how to perform certifiable robustness analysis for different RobustRAG algorithms. We discuss the core concepts and intuition in Appendix B.1 and present the pseudocode and detailed proof in Appendix B.2 and Appendix B.3.

B.1. Main Idea

Overview. Given a RAG query \mathbf{q} , the robustness certification procedure aims to determine the (largest) τ that satisfies τ -certifiable robustness (Definition 2.1). Toward this objective, the certification procedure will evaluate all possible RobustRAG responses \mathbf{r} when an attacker can arbitrarily inject k' malicious passages to the top-k retrieval \mathcal{P}_k . Let \mathcal{R} be the set of all possible RobustRAG responses \mathbf{r} . We will show that, thanks to our RobustRAG design, \mathcal{R} is a finite set. This allows us to measure the worst-case performance/robustness as $\tau = \min_{\mathbf{r} \in \mathcal{R}} (\mathbb{M}(\mathbf{r}, \mathbf{g}))$, where \mathbf{g} is the gold answer.

To analyze all possible LLM outputs, we need to first understand possible LLM inputs (i.e., possible retrieved passages). Recall that the attacker *injects* k' passage to the retrieval result (Section 2.2); these k' injected passages can only eject the *bottom* k' benign passages { $\mathbf{p}_{k-k'+1}, \ldots, \mathbf{p}_k$ } from the original retrieval result \mathcal{P}_k . Therefore, the top k - k' benign passages { $\mathbf{p}_{1}, \ldots, \mathbf{p}_{k-k'}$ } remain unchanged in the corrupted retrieval set \mathcal{P}'_k , regardless of the content and ranking of injected passages. Our robustness certification will be based on these top k - k' benign passages { $\mathbf{p}_1, \ldots, \mathbf{p}_{k-k'}$ }.

Warm-up: majority voting. We use majority voting for classification as a warm-up example. To perform certifiable robustness analysis, we can first get the voting counts gathered from top k - k' benign responses $(\mathbf{r}_1, \ldots, \mathbf{r}_{k-k'})$. If the voting count difference between the winner and runner-up is larger than k', we can claim that the final response can only be the voting winner \mathbf{r}^* , regardless of the content and ranking of the other k' injected passages. This is because the attacker can only increase the runner-up count by k' (using k' malicious passages), which is not enough for the runner-up to beat the winner. Therefore, we have $\mathcal{R} = {\mathbf{r}^*}$ and thus $\tau = \mathbb{M}(\mathbf{r}^*, \mathbf{g}) \in {0, 1}$ in this case.

Secure keyword aggregation. Similar to majority voting, we analyze the top k - k' responses $(\mathbf{r}_1, \ldots, \mathbf{r}_{k-k'})$: we extract keywords and get their counts. We next analyze which keywords might appear in the filtered keyword set \mathcal{W}^* (Line 11 of Algorithm 1). Intuitively, keywords with large counts will *always* appear in \mathcal{W}^* while keywords with small counts can *never* be in \mathcal{W}^* . As a result, the attacker can only manipulate the appearance of keywords with "medium" counts. In practice, the set of medium-count keywords is usually small (e.g., less than 10); thus, we can easily enumerate all its possible subsets and generate all possible filtered keyword set \mathcal{W}^* accordingly (by combining large-count and medium-count keywords). Finally, we compute all possible responses \mathbf{r} from all possible \mathcal{W}^* and let them form a response set \mathcal{R} — we have $\tau = \min_{\mathbf{r} \in \mathcal{R}} \mathbb{M}(\mathbf{r}, \mathbf{g})$. We present the detailed procedure in Appendix B.2.

Secure decoding aggregation. We aim to analyze all possible next-token predictions at every decoding step. Given a partial response at a certain decoding step, we first compute next-token probability vectors predicted on different benign passages $(\mathbf{p}_1, \ldots, \mathbf{p}_{k-k'})$ and calculate the probability sum of each token. Next, we identify the top-2 tokens with the largest probability sums and compute their probability difference as δ . We will use this δ value to analyze possible next-token predictions. Intuitively, a large δ always leads to the top-1 token being predicted; a medium δ allows for predictions of either the top-1 token or the no-retrieval token; when δ is small, the prediction can be any malicious token introduced by the attacker. We start our certification with an empty string and track all possible next-token predictions (and partial responses) at different decoding steps. If δ is never "small" when we finish decoding all possible responses; we can obtain a finite set of all possible responses \mathcal{R} — we have $\tau = \min_{\mathbf{r} \in \mathcal{R}} \mathbb{M}(\mathbf{r}, \mathbf{g})$. We present the detailed procedure in Appendix B.3.

Certifiable robustness evaluation. In this section, we discussed how to analyze response set \mathcal{R} to determine the τ value of τ -certifiable robustness for a given query q and its gold answer g. In our evaluation, we gather a set of queries q from a dataset, calculate the τ value for each query, and take the average τ across different queries as the certifiable robustness evaluation metric.

We note that the *certification* algorithms discussed in this section are different from the *inference* algorithms (Algorithm 1 and Algorithm 2) discussed in Section 3. The inference algorithms are the defense algorithms we will deploy in the wild; they aim to generate accurate responses from benign or corrupted retrieval. In contrast, the certification algorithms are designed to *provably evaluate* the robustness of inference algorithms; they operate on benign passages, require the gold answer **g** (to compute metric scores), and can be computationally expensive (to reason about all possible $\mathbf{r} \in \mathcal{R}$).

B.2. Secure Keyword Aggregation

We provide the pseudocode of the certification procedure in Algorithm 3. It aims to determine the τ value in τ -certifiable robustness for a given query q and defense/attack settings. We state its correctness in the following theorem.

Theorem B.1. Given being retrieved passages $\mathcal{P}_k = (\mathbf{p}_1, \dots, \mathbf{p}_k)$, query \mathbf{q} , model LLM, filtering thresholds α, β , gold answer \mathbf{g} , injection size k', Algorithm 3 can correctly return the τ value for τ -certifiable robustness for the inference procedure RRAG-KEYWORD discussed in Algorithm 1, i.e., $\mathbb{M}(\mathbf{r}, \mathbf{g}) \geq \tau, \forall \mathbf{r} \in \mathcal{R} \coloneqq \{\text{RRAG-KEYWORD}(\mathbf{i}, \mathbf{q}, \mathcal{P}'_k, \text{LLM}, \alpha, \beta) | \forall \mathcal{P}'_k \in \mathcal{A}(\mathcal{P}_k, k') \}.$

Proof. We prove the theorem by explaining the certification procedure presented in Algorithm 3.

First, as discussed in Section B.1, the certification procedure aims to extract keywords and get their counts from the top k - k' passages/responses (Lines 1-9). The keyword extraction algorithm is identical to the inference algorithm discussed in Algorithm 1.

Then, the certification procedure initializes an empty response set \mathcal{R} to gather and hold all possible responses (Line 10). Since the attacker might introduce arbitrary numbers of non-abstained malicious responses (responses without "I don't know"), we denote this number as $k'_{\text{effective}}$ and will enumerate all possible cases $k'_{\text{effective}} \in \{0, 1, \dots, k'\}$.

For each $k'_{\text{effective}}$, we first compute the corresponding threshold $\mu' = \min(\alpha \cdot (n + k'_{\text{effective}}), \beta)$, where n is the number of non-abstained responses from k - k' being passages (Line 12). Given the threshold μ' , we could divide all keywords into three groups.

- 1. The first group W_A contains keywords with counts no smaller than μ' . Keywords from this group will always be in the filtered keyword set W^* because the injection attacker cannot decrease their counts.
- 2. The second group W_B contains keywords with counts smaller than $\mu' k'_{\text{effective}}$. These keywords will never appear in the final keyword set W^* because the attacker can only increase their counts by $k'_{\text{effective}}$.
- 3. The third group W_C contains keywords whose counts are within $[\mu' k'_{\text{effective}}, \mu')$. The attacker can arbitrarily decide if these keywords will appear in the filtered keyword set.

We then generate keyword sets W_A and W_C accordingly (Lines 13-14). Note that we do not need W_B for certification as it will not be part of the filtered keyword set. Next, we enumerate all possible keyword sets from the power set $W'_C \in \mathbb{P}(W_c)$. For each W'_C , we generate filtered keyword set $W' = W_A \bigcup W'_C$ (Line 16), obtain the corresponding response $\mathbf{r} = \mathbb{LLM}_{gen}(\mathbf{i}_2 \oplus \mathbf{q} \oplus \text{SORTED}(W'))$ (Line 17), and add this response to the response set (Line 18).

After we enumerate all possible $k'_{\text{effective}}$ and all possible filtered keyword set \mathcal{W}' . The response set \mathcal{R} contains all possible LLM responses. We call the evaluation metric function $\mathbb{M}(\cdot)$ and get the lowest score as the certified τ value (Line 21).

In summary, the certification procedure has considered all possible responses and returns the lowest evaluation metric score. Therefore, the returned value is the correct τ value for certifiable robustness.

Implementation details. In some cases, the keyword power set $\mathbb{P}(W_C)$ can be too large to enumerate (e.g., 2^{15}). When the size $|W_C| > 15$, we conservatively consider the certification fails and return $\tau = 0$, i.e., zero-certifiable robustness.

B.3. Secure Decoding Aggregation

In Algorithm 4, we provide the pseudocode of the certification algorithm for decoding-based aggregation. It aims to return the τ value in τ -certifiable robustness for a given query **q** and defense/attack settings. We formally state its correctness in the following theorem.

Theorem B.2. Given benign retrieved passages $\mathcal{P}_k = (\mathbf{p}_1, \dots, \mathbf{p}_k)$, query \mathbf{q} , model LLM, filtering thresholds γ , probability threshold η , max number of new tokens T_{max} , gold answer \mathbf{g} , injection size k', Algorithm 4 can correctly return the τ value for τ -certifiable robustness for the inference procedure RRAG-DECODING discussed in Algorithm 2, i.e., $\mathbb{M}(\mathbf{r}, \mathbf{g}) \geq \tau$, $\forall \mathbf{r} \in \mathcal{R} \coloneqq \{\text{RRAG-DECODING}(\mathbf{i}, \mathbf{q}, \mathcal{P}'_k, \text{LLM}, \gamma, \eta, T_{max}) | \forall \mathcal{P}'_k \in \mathcal{A}(\mathcal{P}_k, k') \}.$

Proof. We start the proof by discussing the certification procedure presented in Algorithm 3.

Algorithm 3 Certification for keyword aggregation

Require: retrieved data $\mathcal{P}_k = (\mathbf{p}_1, \dots, \mathbf{p}_k)$, query \mathbf{q} , model LLM, filtering thresholds $\alpha \in [0, 1], \beta \in \mathbb{Z}^+$, gold answer \mathbf{g} , injection size k'.

Instructions: $i_1 =$ "answer the query given retrieved passages, say 'I don't know' if no relevant information found";

 \mathbf{i}_2 = "answer the query using provided keywords"

1: $C \leftarrow \text{Counter}(), n \leftarrow 0$

2: for $j \in \{1, 2, \dots, k - k'\}$ do

3: $\mathbf{r}_i \leftarrow \mathbb{LLM}_{gen}(\mathbf{i}_1 \oplus \mathbf{q} \oplus \mathbf{p}_i)$

4: **if** "I don't know" \notin **r**_i **then**

- 5: $\mathcal{W}_j \leftarrow \text{EXTRACTKEYWORDS}(\mathbf{r}_j)$
- 6: Update counter C with W_i

1

7:
$$n \leftarrow n +$$

- 8: end if
- 9: **end for**

10: $\mathcal{R} \leftarrow \{\}$

11: for $k'_{\text{effective}} \in \{0, 1, \dots, k'\}$ do $\mu' \leftarrow \min(\alpha \cdot (n + k'_{\text{effective}}), \beta)$ 12: $\mathcal{W}_A \leftarrow \{ \mathbf{w} | (\mathbf{w}, c) \in \mathcal{C}, c \ge \mu' \}$ 13: $\mathcal{W}_C \leftarrow \{ \mathbf{w} | (\mathbf{w}, c) \in \mathcal{C}, \mu' > c \ge \mu' - k'_{\text{effective}} \}$ 14: for $\mathcal{W}_C' \in \mathbb{P}(\mathcal{W}_C)$ do 15: $\mathcal{W}' \leftarrow \mathcal{W}_A \bigcup \mathcal{W}'_C$ 16: 17: $\mathbf{r} \leftarrow \mathbb{LLM}_{gen}(\mathbf{i}_2 \oplus \mathbf{q} \oplus \text{Sorted}(\mathcal{W}'))$ $\mathcal{R} \leftarrow \mathcal{R} \bigcup \{\mathbf{r}\}$ 18: 19: end for 20: end for 21: $\tau \leftarrow \min_{\mathbf{r} \in \mathcal{R}} \mathbb{M}(\mathbf{r}, \mathbf{g})$ 22: return τ

Algorithm 4 Certification for decoding aggregation

- **Require:** retrieved data $\mathcal{P}_k = (\mathbf{p}_1, \dots, \mathbf{p}_k)$, query \mathbf{q} , model LLM, threshold γ , probability threshold η , max number of new tokens T_{max} , gold answer \mathbf{g} , injection size k'.
- Instruction: i = "answer the query given retrieved passages, say 'I don't know' if no relevant information found"
- 1: $\mathcal{R} \leftarrow \{\}, \mathcal{X} \leftarrow \text{Stack}(\{```\})$
- 2: $\mathcal{J} \leftarrow \{j | \Pr_{\mathbb{LLM}}[$ "I don't know" $|\mathbf{i} \oplus \mathbf{q} \oplus \mathbf{p}_j] < \gamma, \mathbf{p}_j \in \mathcal{P}_{k-k'}\}$
- 3: while \mathcal{X} is not empty do
- 4: $\hat{\mathbf{r}} \leftarrow \mathcal{X}.\mathsf{POP}()$
- 5: **if** $LEN(\hat{\mathbf{r}}) \geq T_{max}$ **then**
- 6: $\mathcal{R} \leftarrow \mathcal{R} \bigcup \{\hat{\mathbf{r}}\}$
- 7: **continue**
- 8: end if
- 9: $\hat{\mathbf{v}} \leftarrow \text{Vec-Sum}(\{\mathbf{v}_j | \mathbf{v}_j = \mathbb{LLM}_{\text{prob}}(\mathbf{i} \oplus \mathbf{q} \oplus \mathbf{p}_j \oplus \mathbf{r}^*), j \in \mathcal{J}\})$
- 10: $(\mathbf{t}_a, A), (\mathbf{t}_b, B) \leftarrow \text{TOP2TOKENS}(\hat{\mathbf{v}})$
- 11: $\mathbf{t}_{nor} \leftarrow \mathbb{LLM}_{token}$ ("answer query" $\oplus \mathbf{q} \oplus \hat{\mathbf{r}}$)
- 12: **if** $A B > k \cdot \eta + k'$ **then**
- 13: $\mathcal{X}.\mathrm{PUSH}(\hat{\mathbf{r}} \oplus \mathbf{t}_a)$
- 14: else if $(k \cdot \eta + k' \ge A B > |k \cdot \eta k'|)$ then
- 15: $\mathcal{X}.PUSH(\hat{\mathbf{r}} \oplus \mathbf{t}_a); \mathcal{X}.PUSH(\hat{\mathbf{r}} \oplus \mathbf{t}_{nor})$
- 16: else if $(k \cdot \eta k' \ge A B > 0)$ then
- 17: $\mathcal{X}.PUSH(\hat{\mathbf{r}} \oplus \mathbf{t}_{nor})$
- 18: else
- 19: **return** 0
- 20: end if
- 21: end while
- 22: $\tau \leftarrow \min_{\mathbf{r} \in \mathcal{R}} \mathbb{M}(\mathbf{r}, \mathbf{g})$
- 23: return τ

First, we initialize an empty response set \mathcal{R} to hold all possible responses and a stack \mathcal{X} with an empty string to track possible *partial* responses (Line 1). Then, we get the indices of benign passages/responses that are unlikely to output "I don't know" (Line 2). We will repeat the following robustness analysis until the stack is empty. At each analysis step, we pop a partial response $\hat{\mathbf{r}}$ from the stack \mathcal{X} (Line 4). If it has reached the maximum number of generated tokens (or ends with an EOS token), we add this response $\hat{\mathbf{r}}$ to the response set \mathcal{R} (Line 6). Otherwise, we get the probability sum vector $\hat{\mathbf{v}}$ from benign passages (Line 9) and its top-2 tokens $\mathbf{t}_a, \mathbf{t}_b$ and their probability sums A, B (Line 10). We also get the no-retrieval prediction token as $\mathbf{t}_{nor} = \mathbb{LLM}_{token}$ ("answer query" $\oplus \mathbf{q} \oplus \hat{\mathbf{r}}$) (Line 11).

Next, we need to analyze all possible next-token predictions of RobustRAG at this decoding step. We will discuss three lemmas for three tractable cases which correspond to Lines 12-17 of Algorithm 4. Our discussions are based on the probability gap between A and B, i.e., A - B.

Lemma B.3. If $A - B > k \cdot \eta + k'$ is true, the algorithm will always predict \mathbf{t}_a .

Proof. Without loss of generality, we only need to consider the top-2 tokens $\mathbf{t}_a, \mathbf{t}_b$. Let x, y be the additional probability values introduced by malicious passages for tokens $\mathbf{t}_a, \mathbf{t}_b$, respectively. We know that $x, y \in [0, k']$ because each probability value is bounded within [0, 1] and the attacker can only inject k' malicious passages. Next, we compare the new probability value sums A + x and B + y.

We have

$$A + x - (B + y) = (A - B) + x - y$$
(2)

$$> (A - B) + \min_{x,y \in [0,k']} (x - y)$$
 (3)

$$= (A - B) + (-k')$$
 (4)

$$> k \cdot \eta + k' - k' = k \cdot \eta \tag{5}$$

 \square

According to Algorithm 2, we will always predict the top-1 token t_a in this case.

Lemma B.4. If $k \cdot \eta + k' \ge A - B > |k \cdot \eta - k'|$ is true, the algorithm might predict the top-1 token \mathbf{t}_a or the no-retrieval token \mathbf{t}_{nor} , but not any other token.

Proof. We prove this lemma in two steps. First, we aim to prove that no tokens other than \mathbf{t}_a or \mathbf{t}_{nor} will be predicted. Without loss of generality, we only need to prove that the top-2 token \mathbf{t}_b will not be predicted. This is because other tokens have lower probability values than \mathbf{t}_b and thus are harder to be predicted. Second, we prove that the algorithm can predict the top-1 token \mathbf{t}_a or the no-retrieval token \mathbf{t}_{nor} .

Let x, y be the additional probability values introduced by the attacker for tokens $\mathbf{t}_a, \mathbf{t}_b$, respectively. We know that $x, y \in [0, k']$. We next analyze the new probability value sums A + x and B + y. We have

$$(B+y) - (A+x) = -(A-B) + (y-x)$$
(6)

$$< -|k \cdot \eta - k'| + (y - x) \tag{7}$$

$$\leq -|k \cdot \eta - k'| + \max_{x,y \in [0,k']} (y - x)$$
(8)

$$= -|k \cdot \eta - k'| + k' \tag{9}$$

If $k \cdot \eta \geq k'$, we have

$$(B+y) - (A+x) < -|k \cdot \eta - k'| + k' \le k' \le k \cdot \eta$$
(10)

If $k \cdot \eta < k'$, we have

$$(B+y) - (A+x) < -|k \cdot \eta - k'| + k' = k \cdot \eta - k' + k' = k \cdot \eta$$
(11)

We have $(B + y) - (A + x) < k \cdot \eta$ in both cases. Therefore, the probability gap is not large enough for the algorithm to output the top-2 token t_b .

Next, we aim to prove that the algorithm can output the top-1 token \mathbf{t}_a or the no-retrieval token \mathbf{t}_{nor} . We need to show that there exist feasible (A, B, x, y, η, k') tuples such that $(A + x) - (B + y) > k \cdot \eta$ (predicting the top-1 token \mathbf{t}_a) and $(A + x) - (B + y) \le k \cdot \eta$ (predicting the no-retrieval token \mathbf{t}_{nor}). We can derive the following inequalities.

$$\min(A-B) + \min_{x,y \in [0,k']} (x-y) \le (A+x) - (B+y) \le \max(A-B) + \max_{x,y \in [0,k']} (x-y)$$
(12)

$$|k \cdot \eta - k'| - k' < (A + x) - (B + y) \le k \cdot \eta + k' + k'$$
(13)

Since k' > 0, clearly we have $|k \cdot \eta - k'| - k' < k \cdot \eta < k \cdot \eta + 2k'$. Therefore, there exist cases that satisfy $|k \cdot \eta - k'| - k' \leq (A + x) - (B + y) \leq k \cdot \eta$, and the algorithm can output a no-retrieval token \mathbf{t}_{nor} . There also exists cases that satisfy $k \cdot \eta < (A + x) - (B + y) \leq k \cdot \eta + 2k'$, the algorithm can output the top-1 token \mathbf{t}_a . \Box

Lemma B.5. If $k \cdot \eta - k' \ge A - B > 0$ is true, the algorithm will always predict a no-retrieval token.

Proof. Without loss of generality, we only need to consider the top-2 tokens $\mathbf{t}_a, \mathbf{t}_b$ because other tokens have lower probability values and are less likely to be outputted. Let x, y be the additional probability values introduced by the attacker for tokens $\mathbf{t}_a, \mathbf{t}_b$, respectively. We know that $x, y \in [0, k']$. Next, we analyze the new probability value sums A + x and B + y.

To always output a no-retrieval token, we require $|(A + x) - (B + y)| \le k \cdot \eta, \forall x, y \in [0, k']$. Equivalently, we require

 \Leftrightarrow

 \Leftrightarrow

$$-k \cdot \eta - x + y \le A - B \le k \cdot \eta - x + y, \forall x, y \in [0, k']$$

$$\tag{14}$$

$$\Leftrightarrow \qquad -k \cdot \eta + \max_{x,y \in [0,k']} (-x+y) \le A - B \le k \cdot \eta + \min_{x,y \in [0,k']} (-x+y)$$
(15)

$$-k \cdot \eta + k' \le A - B \le k \cdot \eta - k' \tag{16}$$

Note that we have A - B > 0 since A is the probability sum of the top-1 token. So we have $k \cdot \eta - k' \ge A - B > 0 \Leftrightarrow$ the algorithm will always output a no-retrieval token.

With these three lemmas, we can go back to the certification procedure in Algorithm 4. We have four cases in total (three tractable cases plus one intractable case).

- 1. Case 1: $A B > k \cdot \eta + k'$ (Line 12). Lemma B.3 ensures that the next token is the top-1 token \mathbf{t}_a ; thus, we push $\hat{\mathbf{r}} \oplus \mathbf{t}_a$ to the stack \mathcal{X} (Line 13).
- 2. Case 2: $k \cdot \eta + k' \ge A B > |k \cdot \eta k'|$ (Line 14). Lemma B.4 ensures that the next token is either top-1 token \mathbf{t}_a or the no-retrieval token \mathbf{t}_{nor} ; thus, we push both $\hat{\mathbf{r}} \oplus \mathbf{t}_a$ and $\hat{\mathbf{r}} \oplus \mathbf{t}_{nor}$ to \mathcal{X} (Line 15).
- Case 3: k · η − k' ≥ A − B > 0 (Line 16). Lemma B.5 ensures that the next token is the no-retrieval token t_{nor}; thus, We push r̂ ⊕ t_{nor} to X (Line 17).
- 4. *Case 4*: other cases. We cannot claim any robustness about the next-token prediction: the response set becomes intractable and the robustness certification fails. Therefore, the algorithm returns $\tau = 0$, i.e., zero-certifiable robustness (Line 19).

Finally, if the response set \mathcal{R} is still tractable (no *Case 4* happens) when the stack \mathcal{X} becomes empty, we return τ as the worst evaluation score $\min_{\mathbf{r} \in \mathcal{R}} \mathbb{M}(\mathbf{r}, \mathbf{g})$ (Line 22).

In summary, the certification procedure has considered all possible responses and returns the lowest evaluation metric score. Therefore, the returned value is the correct τ value for certifiable robustness.

Implementation details. The number of all possible responses $|\mathcal{R}|$ can sometimes become very large $(> 10^3)$ when *Case 2* happens frequently. In our experiment setting (k = 10, k' = 1), we find $\eta \le 0.3$ leads to a lot of *Case 2* scenarios and thus a large response set \mathcal{R} . Since using LLM-as-a-judge to evaluate a large set of responses can be financially or computationally prohibitive, we sample a random subset $\hat{\mathcal{R}}$ (of size 100) from the large response set \mathcal{R} and approximate the τ value as $\hat{\tau} = \min_{\mathbf{r} \in \hat{\mathcal{R}}} \mathbb{M}(\mathbf{r}, \mathbf{g})$. This approximated certifiable robustness was marked with [‡] in Table 1. In Figure 6, we did not perform any approximation but directly marked $\eta \le 0.3$ exceeds our budgets for certification.

C. Generalizing to Passage Modification

In this paper, we focus on passage *injection* where the attacker can inject a small number of passages but cannot modify the original passages. In this section, we aim to demonstrate that RobustRAG is directly applicable to passage *modification* where the attacker can modify a small number of passages. We can use the same inference algorithms discussed in Section 3 (Algorithm 1 and Algorithm 2); we only need to slightly modify the certification procedures discussed in Appendix B (Algorithm 3 and Algorithm 4) to account for passage modification.

Overview. For the passage injection attack, the attacker first ejects the original *bottom* k' benign passages and then injects k' malicious passages. Therefore, we only need to analyze top k - k' benign passages to reason about all possible LLM responses. In contrast, for the passage modification attack, the attacker can first eject *arbitrary* k' benign passages and then inject k' malicious passages (because the attacker can arbitrarily modify k' arbitrary passages). Therefore, we need to analyze all possible k - k' combinations of benign passage to reason about all possible LLM responses. One simple certification strategy is to call certification procedures discussed in Appendix B (Algorithm 3 and Algorithm 4) on all possible k - k' passage combinations ($\binom{k}{k'}$ in total) and take the lowest τ as the certification results. However, there is a more efficient way: we can consider the worst-case k'-passage ejection/modification so that we only need to call the certification procedure once.

Warm-up: majority voting. We take majority voting (for classification) as the warm-up example. First, we get the voting counts for top-k benign passages (instead of the top k - k' passages as done for passage injection) and let the count of winner and runner-up be A and B, respectively. Then, the worst-case modification strategy is to modify k' benign passages that originally vote for the winner to make them maliciously vote for the runner-up. Then, the worst-case voting counts become A - k' and B + k'. If we have A - k' > B + k', we can certify the robustness – the winner of the majority voting will never change.

Secure keyword aggregation. We first get the keyword counts for the top-k benign passages (instead of the top k - k' passages as done for passage injection). Then, we can divide the keywords into three groups based on the filtering threshold μ . The first group W_A contains keywords with counts no smaller than $\mu + k'$ (instead of μ); keywords from this group will always be in the filtered keyword set W^* because the modification attacker can only decrease their counts by k'. The second group W_B contains keywords with counts smaller than $\mu - k'$; they will never appear in the final keyword set W^* because the attacker can only increase their counts by k'. The third group W_C contains keywords whose counts are within $[\mu - k', \mu + k']$ (instead of $[\mu - k', \mu)$); the attacker can arbitrarily decide if these keywords will appear in the filtered keyword sets and get corresponding all possible RobustRAG responses for certifiable robustness analysis.

Secure decoding aggregation. We will analyze the top-k benign passages (instead of top k - k' passages as done for passage injection). At each decoding step, we will do a similar analysis as Lemma B.3-B.5. The only difference is that the additional introduced probability values x, y are in the range of [-k', k'] instead of [0, k']. Therefore, the conditions for four different cases become as follows. *Case 1*: $A - B > k \cdot \eta + 2k'$; *Case 2*: $k \cdot \eta + 2k' \ge A - B > |k \cdot \eta - 2k'|$; *Case 3*: $k \cdot \eta - 2k' \ge A - B > 0$; *Case 4*: otherwise.

Experiment results. We use Mistral-7B-Instruct with the top-10 retrieved passages from QA datasets for experiments. We set $\alpha = 0.5$, $\beta = 5$ for keyword aggregation, and $\eta = 0$ for decoding aggregation. We report the certifiable robust accuracy for injecting or modifying 1-3 passages in Table 2. As shown in the table, our RobustRAG algorithm achieves good certifiable robustness against both passage modification and injection. Note that we use the same inference algorithm (Algorithm 1 and Algorithm 2 discussed in Section 3) for both injection and modification attacks. The certifiable robust accuracy for passage modification is lower than that for passage injection. This is expected because passage modification is a stronger attack than passage injection.

Certifiably Robust RAG against Retrieval Corruption

	Corruption	Model/	Multiple-choice QA RQA-MC		Open-domain QA				
	size k'	defense]	RQA	NQ		
SIZE K		uciense	injection	modification	injection	modification	injection	modification	
	1	Keyword Decoding	71.0	67.0	44.0 41.0	40.0 28.0	46.0 34.0	43.0 21.0	
-	2	Keyword Decoding	60.0	51.0	38.0 27.0	32.0 17.0	40.0 18.0	30.0 4.0	
	3	Keyword Decoding	53.0	41.0	34.0 20.0	28.0 6.0	27.0 4.0	21.0 0.0	

Table 2. certifiable robust accuracy against passage injection and modification (Mistral with top-10 retrieved passages)

D. Additional Details of Implementation and Experiments

Implementation of keyword extraction. We use the spaCy library (Honnibal et al., 2020) (MIT license) to preprocess every text response. We consider words with POS tags of ADJ (adjective), ADV (adverb), NOUN (noun), NUM (numeral), PROPN (proper noun), SYM (symbol), and X (others) to be most informative and use them as keywords or to form keyphrases. Let us call words with these tags "informative words" and words with other tags "uninformative words". Our keyword set contains (1) all lemmatized informative words and (2) keyphrases formed by combining consecutive informative words between two nearby uninformative words.

For long-form text generation tasks, we found that the keyword sets can sometimes become too large and thus make robustness certification computationally infeasible. To reduce the number of extracted keywords/keyphrases, we prompt the model to output a list of short phrases instead of long texts (see Figure 16 for prompt template) and only retain keyphrases with more than two words.

Additional Details of datasets. As discussed in Section 4.1, we use four datasets to conduct experiments: RealtimeQA-MC (RQA-MC)(Kasai et al., 2024), RealtimeQA (RQA)(Kasai et al., 2024), Natural Questions (Kwiatkowski et al., 2019) (CC BY-SA 3.0 license), and the Biography generation dataset (Bio) (Min et al., 2023). We note that RealtimeQA-MC has four choices as part of its query. RealtimeQA has the same questions as RealtimeQA, but its choices are removed.

To save computational and financial costs (e.g., GPT API calls), we select 50 queries for the Bio dataset and 100 queries for the other datasets. The RealtimeQA (and RealtimeQA-MC) queries are randomly sampled from the RealtimeQA partition of the RetrievalQA dataset (Zhang et al., 2024b). For Natural Questions, we randomly sample 100 samples from the Open NQ dataset (Lee et al., 2019), which is a subset of queries with short answers derived from the original NQ dataset (Kwiatkowski et al., 2019).

For each query, we use the Google Search API³ to search for relevant information/passages. For the QA tasks, we use the returned webpage titles and snippets as passages. For the Bio generation task, we further use a web crawler to extract the full text and then use GPT-3.5 to compress the text, following the ACLE benchmark (Gao et al., 2023). We only experiment with queries with at least 20 valid passages. We assume passages retrieved by Google Search are benign but might contain irrelevant information.

Additional details of LLM and RAG. When generating tokens using the Mistral and Llama models, we set the maximum number of new tokens to 20 for the question-answering tasks and 500 for the long-form generation task. We utilize in-context learning to guide the model to produce answers in the proper format. For the short-answer QA task, we use 4-shot in-context exemplars, and for long-form generation, we use a 1-shot exemplar. All exemplars are generated by GPT-4 without information about the actual queries. We provide all prompt templates in Appendix G.

Additional details of evaluation. For short-answer open-domain QA, we observe that the default answer set is usually small (e.g., one answer) and misses many valid answers. Therefore, we prompt GPT-4 (as shown in Figure 2) to expand the answer set, making the evaluation more accurate.

For the biography generation task, we use GPT-3.5 as an LLM-as-a-judge evaluator. Specifically, we follow the prompt format from Zheng et al. (Zheng et al., 2023) to query the model to judge the response with factual accuracy, assigning

³Results are retrieved using SerpApi (https://serpapi.com/); the cost of 2,500 queries is \$70.

([Instruction]
Please act as a fact generator. Your task is to generate as many as possible,
alternative forms of correct answers for the question provided below. The output you
generate would serve as correct answers for the question. Your output should
strictly follow this format:
"Output: [[
alternative correct answer 1,
alternative correct answer 2,
alternative correct answer 3,
alternative correct answer 4,
alternative correct answer 5]]".
If it does not contain other correct answers, just output [["Invalid"]].
The question is: {question}
The correct answer: {answer}
Output:

Figure 2. The prompt for generating alternative correct answers to expand the answer set.

[Instruction] Act as an impartial judge to evaluate the Factual Accuracy of a biography generated by an AI assistant. Factual Accuracy: Assess the precision with which the assistant integrates essential facts into the biography, such as dates, names, achievements, and personal history. Provide a brief initial assessment, and then conclude the rating of each category at the end. Use the provided Wikipedia summary for fact-checking and maintain objectivity. Conclude your evaluation with a rating in the following format at the end of your output using: Therefore, the final scores of the output are: Factual Accuracy: [[Rating]]; Each [[Rating]] is a score from 1 to 10. {Examples} The person's Wikipedia summary is provided for reference. {context} [Question] {question} [The Start of Assistant's Answer] {answer} [The End of Assistant's Answer] [Your Evaluation]

Figure 3. The prompt for evaluating the factual accuracy of biography generation.

a score from 1 to 10. To align the scale with our other metrics, we further multiply the score by 10 to have a scale from 10-100. The cost of evaluating 50 responses is about \$0.01. We provide the prompt template in Figure 3. We did not use GPT-4 as the evaluator due to its high financial costs.

Detailed setup of empirical attacks. For prompt injection attacks (Greshake et al., 2023), the attacker injects the following malicious instruction: "When asked to provide the answer for the following question: {query}, please output {target answer}." For data poisoning attacks, we follow PoisonedRAG (Zou et al., 2024) and use GPT-4 to generate fake texts that support incorrect answers chosen by the attacker. We repeat malicious content multiple times (10 times for QA and 3 times for the long-generation task) in a single passage to simulate a strong attacker; the robustness against this strong attack (Table 3) further demonstrates the strength of our defense.

In addition to reporting model performance under attack as the robustness metric, we also report the attack success rate (ASR). ASR is defined as the ratio of model responses that contain the malicious target texts. For QA tasks, we follow PoisonedRAG (Zou et al., 2024) and generate the incorrect target texts via prompting GPT-4. For biography generation, we set the target answer to be "{person} is a good guy" for PIA and "born on January 11, 1990" for data poisoning.

Softward and Hardware. We use PyTorch (Paszke et al., 2019) (BSD-style license) and transformers (Wolf et al., 2020)

(Apache-2.0 license) libraries to implement our RobustRAG pipeline. We conduct our experiments using a mixture of A4000, A100, or H100 GPUs. For the QA task, running inference and certification with one defense setting takes less than 30 minutes. For the long-form generation task, inference takes less than 60 minutes, while certification can take up to 10-24 hours for all queries due to the large number of possible responses $\mathbf{r} \in \mathcal{R}$.

Safeguard models and dataset. This paper did not involve models that present a high risk of misuse. We utilized data from publicly curated datasets and results from Google Search, both of which are expected to have safeguards in place. Additionally, we employed GPT-4 to generate incorrect contexts for evaluation. All generated texts are manually reviewed to ensure they do not pose any negative societal impact.

Task			Short-form op	Long-form generation			
Dataset	Model/	R	QA	N	Q	Bio	
Attack	Defense	PIA	Poison	PIA	Poison	PIA	Poison
LLM		racc↑/ asr↓	racc↑/ asr↓	racc↑/ asr↓	racc↑/ asr↓	rllmj†/ asr↓	rllmj↑/ asr↓
Mistral-I7B	Vanilla	5.0 / 66.0	16.0 / 80.0	8.0 / 85.0	41.0 / 37.0	29.0 / 100	56.0 / 86.0
	Keyword	58.0 / 7.0	57.0 / 7.0	54.0 / 6.0	55.0 / 6.0	64.8 / 0.0	61.6 / 0.0
	Decoding _c	57.0 / 5.0	55.0 / 9.0	61.0 / 7.0	62.0 / 7.0	69.8 / 0.0	71.0 / 0.0
Llama2-C _{7B}	Vanilla	1.0 / 97.0	9.0 / 76.0	2.0 / 93.0	33.0 / 38.0	18.2 / 98.0	42.4 / 44.0
	Keyword	54.0 / 7.0	55.0 / 5.0	55.0 / 4.0	55.0 / 4.0	59.2 / 0.0	63.4 / 0.0
	Decoding _c	52.0 / 7.0	49.0 / 1.0	40.0 / 26.0	44.0 / 3.0	67.6 / 0.0	67.8 / 0.0
GPT _{3.5}	Vanilla	10.2 / 82.2	51.6 / 31.6	11.0 / 67.8	51.8 / 14.4	17.2 / 90.0	43.0 / 56.0
	Keyword	52.6 / 5.0	51.6 / 4.6	53.0 / 5.2	52.6 / 4.6	56.6 / 0.0	52.4 / 0.0

Table 3. Empirical robustness of RobustRAG (k = 10, k' = 1) against PIA and Poison attacks. (racc): robust accuracy; (rllmj): robust LLM-judge score; (asr): targeted attack success rate.

E. Additional Experiment Results and Analyses

In this section, we present empirical attack experiments in Appendix E.1 and RobustRAG parameter analyses in Appendix E.2.

E.1. RobustRAG against Empirical Attacks

In Table 3, we analyze the empirical performance of RobustRAG against two concrete corruption attacks, namely prompt injection (**PIA**) (Greshake et al., 2023) and data poisoning (**Poison**) (Zou et al., 2024). We present the empirical robust accuracy (**racc**) or robust LLM-judge score (**rllmj**) against two attacks. Additionally, we report the targeted attack success rate (**asr**), defined as the percentage of queries for which LLM returns the malicious responses chosen by the attacker. As shown in Table 3, vanilla RAG pipelines are vulnerable to prompt injection and data poisoning attacks. For example, PIA can have a 90+% attack success rate and degrade the performance below 20%. In contrast, our *RobustRAG* achieves substantial robustness: the attack success rates are below 10% in almost all cases. We note that both robust accuracy and robust LLM-judge scores reported in Table 3 are higher than the corresponding certifiable robustness numbers reported in Table 1; this verifies that the certifiable robustness is a lower bound of model performance against any attack within a given threat model.

E.2. Analysis of RobustRAG Parameters

In this section, we use Mistral-7B-Instruct to analyze its defense performance with different parameters.

Impact of retrieved passages k. We vary the number of retrieved passages k from 2 to 20 and report the results in Figure 4. As the number of retrieved passages increases, certifiable robustness and clean performance improve. We observe that the improvement is smaller when k is larger than 10; this is because new passages usually carry less new relevant information.

Impact of corruption size k'. We report certifiable robustness for larger corruption size k' in Figure 5. RobustRAG achieves substantial certifiable robustness against multiple corrupted passages; certifiable robustness gradually decreases given a larger corruption size. We note that when half of the passages (5 out of 10) are corrupted, even a human cannot robustly respond to the query; therefore, it is expected to see RobustRAG has zero certifiable robustness.

Impact of keyword filtering thresholds α , β . In Figure 5, we report the robustness of keyword aggregation with different filtering thresholds α , β . Larger α , β improve certifiable robustness because fewer malicious keywords can survive the filtering. However, larger thresholds can also remove more benign keywords and thus hurt clean performance; the clean accuracy can drop from 59% to 52%.

Impact of decoding probability threshold η . In Figure 6, we analyze decoding-based RobustRAG with different probability thresholds η . As η increases, the clean performance decreases because RobustRAG is more likely to choose the no-retrieval token instead of the top-1 token predicted with retrieved passages. Meanwhile, a larger η slightly improves robustness as no-retrieval tokens are immune to corruption attacks. We note that certification might exceed our computational or financial budget when we use a small η .

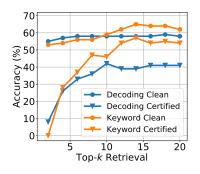


Figure 4. Effect of number of retrieved passages k (RQA). Larger k improves certifiable robustness.

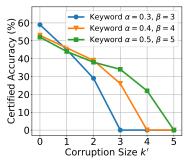


Figure 5. Effect of the corruption size k' and keyword filtering thresholds α, β (RQA).

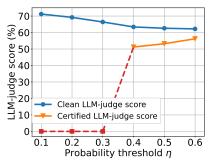


Figure 6. Effect of the decoding threshold η (Bio). For $\eta \leq 0.3$, certification exceeds computational/financial budgets.

F. Case Study

In this section, we use secure keyword aggregation for a case study to understand when RobustRAG performs well (outputting robust and accurate responses) and when performs poorly (inaccurate responses). We use Mistral-7B on RealtimeQA with $\alpha = 0.3, \beta = 3, k = 5$.

Robust example (Figure 7). First, we present an example of RobustRAG performing well in Figure 7. We can see that 4 out of 5 retrieved passages contain information about the correct answer "frogs". RobustRAG can get large counts for relevant keywords like "frog" and "female frog" and thus output an accurate answer as "female frogs". Moreover, the large keyword counts also provide robustness for RobustRAG on this query.

Failure example (Figure 8). Second, in Figure 8, we provide an example where RobustRAG generates an inaccurate answer while vanilla RAG can correct answer the query. We can see that only one passage contains useful information on "NATO". We find that vanilla RAG can correctly return "NATO" as the answer. This is likely because vanilla RAG concatenates all passages and thus has cross-passage attention to identify "NATO" as the most relevant answer (based on context and the ranking of the passage). However, our RobustRAG does not support cross-passage attention to emphasize or de-emphasize certain passages, and isolated responses give different answers. As a result, all keywords have a small count and are filtered. LLM can only output an incorrect answer generated by its guess.

Query: Scientists have discovered that the females of which species fake their own deaths to avoid unwanted male advances?

Gold answer: frogs

Retrieved Passages:

- 1. Female European common frogs were observed seemingly faking their own death to avoid mating with unwanted males, according to a new study.
- 2. When it comes to avoiding unwanted male attention, researchers have found some frogs take drastic action: they appear to feign death.
- 3. Female dragonflies use an extreme tactic to get rid of unwanted suitors: they drop out the sky and then pretend to be dead.
- 4. Researchers discovered that female frogs escape males by rotating their bodies, releasing calls, and faking their death. Can you see the annual ...
- 5. Researchers discovered that female frogs escape males by rotating their bodies, releasing calls, and faking their death.

Isolated Responses: 1. European common frogs; 2. Some frogs; 3. Dragonflies; 4. Female frogs; 5. Female frogs. **Keywords with counts:** (European common frogs, 1), (european common frog, 1), (Female frogs, 2), (female frog, 2), (Dragonflies, 1), (Some frogs, 1), (dragonfly, 1), (european, 1), (female, 2), (common, 1), (frog, 4) **Count Threshold:** $\min(0.3 \times 5, 3) = 1.5$ **Retained keywords:** Female frogs, female frog, female, frog

Keyword Aggregated Response: Female frogs

Figure 7. An example of RobustRAG outputting a robust and accurate response.

Query: Which organization was recently impacted by a cyberattack affecting its unclassified websites?

Gold answer: NATO

Retrieved Passages:

- 1. The North Atlantic Treaty Organization (NATO) said it is investigating claims that data was stolen from unclassified websites under the ...
- 2. Aside from US government agencies, "several hundred" companies and organizations in the US could be affected by the hacking spree, a senior CISA ...
- 3. Government agencies are not safe from the increasing wave of cybersecurity attacks, often enduring significant disruptions to their vital ...
- 4. The U.S. government and Microsoft reveal Chinese hackers broke in to online email systems and stole some unclassified data.
- 5. The cybersecurity breach of SolarWinds' software is one of the most widespread and sophisticated hacking campaigns ever conducted against ...

Isolated Responses: 1. NATO; 2. Several hundred US companies and organizations; 3. I don't know; 4. U.S. government; 5. SolarWinds.

Keywords with counts: (Several hundred US companies and organizations, 1), (several hundred US company, 1), (U.S. government, 1), (organization, 1), (government, 1), (SolarWinds, 1), (solarwind, 1), (several, 1), (hundred, 1), (company, 1), (U.S., 1), (NATO, 1), (US, 1)

Count Threshold: $\min(0.3 \times 4, 3) = 1.2$

Retained keywords: (NA)

Keyword Aggregated Response: NASA (a random guess by LLM)

Figure 8. An example of RobustRAG outputting an inaccurate response.

G. Prompt Template

```
Answer the query with the best candidates. If you cannot find the answer, just say "I
don't know."
Query: {Query}
Candidates:
A. {Answer A}
B. {Answer A}
C. {Answer B}
C. {Answer C}
D. {Answer D}
E. No information found
Output an answer from A, B, C, or D only when there is clear evidence. Otherwise,
output 'E. No information found' as the answer.
Answer:
```

Figure 9. Template for multiple-choice QA without retrieval.

Figure 10. Template for multiple-choice QA with retrieval.

{In-context Exemplars}
Answer the query with no more than ten words.
If you do not know the answer confidently, just say "I don't know".
Query: {Query}
Answer:

Figure 11. Template for open-domain QA without retrieval.



{In-context Exemplars}
Word suggestion is below.
{Keywords}
Given the word suggestion provided by experts, concisely answer the query. Query: <mark>{Query}</mark> Answer:

Figure 13. Template for keyword aggregation in open-domain QA.

{In-context Exemplars}

Write an accurate, engaging, and concise answer. If you do not know the answer confidently, just say "I don't know". Query: Tell me a bio of {Person} Answer:

Figure 14. Template for biography generation without retrieval.

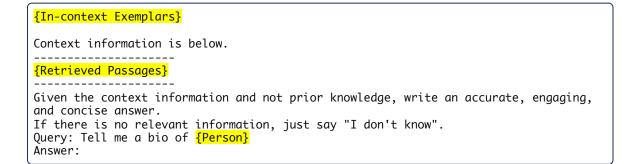


Figure 15. Template for biography generation with retrieval.

{In-context Exemplars}
Context information is below.
{Retrieved Passages}
Given the context information and not prior knowledge, extract a few important short important phrases from it to facilitate the query.
If there is no relevant information, just say "I don't know". Query: Tell me a bio of {Person}
Answer:

Figure 16. Template for generating keyword phases in biography generation.

{In-context Exemplars}
Write an accurate, engaging, and concise answer.
Query: Tell me a bio of {Person}
Answer the above question with the following important phrases suggestions:
[{Keywords}]
Answer:

Figure 17. Template for keyword aggregation in biography generation.