

# A Survey of Networking Cipher Algorithms and How Natural Language Can Be Used to Enhance Them

Anonymous ACL submission

## Abstract

This paper provides a survey of several of the networking cipher algorithms and proposes a method for integrating natural language processing (NLP) as a protective agent for them. Two main proposals are covered for the use of NLP in networking. First, NLP is considered as the weakest link in a networking encryption model; and, second, as a hefty deterrent when combined as an extra layer over what could be considered a strong type of encryption – the stream cipher. This paper summarizes how languages can be integrated into symmetric encryption as a way to assist in the encryption of vulnerable streams that may be found under attack due to the natural frequency distribution of letters or words in a local language stream.

## 1 Introduction

A stream cipher can be illustrated in many ways. In its purest algorithmic form, a stream cipher is a type of symmetric encryption algorithm (Burnett and Paine, 2001). A symmetric algorithm achieves encryption by using the same cryptographic keys in order to encrypt or decrypt a message where a shared secret is shared by the sender and the receiver. The sharing of a secret, as most of us know from typical childhood “keep a secret” games, is not secure. And, as a result of their lack of security, stream ciphers must be considered attackable and in need of a stronger defense against attacks and greater security.

Algorithms are the key to privacy but by their nature are public and easy to read. The public availability of algorithms along with the simple frequency of a local language can prove to be devastating for a stream cipher algorithmic modeler. At times, safe stream ciphering can be considered almost as an n-complete problem due to the numerous attacks that have occurred in the past towards them.

The idea that the algorithms should all be public and the secrecy should reside exclusively in

the keys is called Kerckhoffs’ principle, all serious cryptographers subscribe to this idea. (Stallings, 2017) Knowing the cryptography relies on the keys as its secrecy, an attacker will often times focus on breaking the key that is generated by a key generation algorithm. Key generation algorithms are directly used in the majority of stream ciphers and can be considered the weakest link for transferring data due to the aforementioned details where secrecy lies within a key.

One way to prevent attackers from using publicized symmetric algorithm knowledge and key decryption techniques that break stream ciphers is to provide an extra layer of security on top of the currently available layers. The layer of security should be simple to understand while at the same time robust enough to be applied to any cipher stream available. Several methods (Malladi et al., 2002) have been proposed and are used for strengthening security such as randomness, bit shifting, and the use of digits. Contrastingly, a common framework, while seemingly easy-to-decrypt and insecure, could be the use of a language for encryption. Most networking stream attacks, such as the commonly implemented replay attack (El Abbadi and Jamouli, 2021), use the knowledge of the local language at hand to stage attacks; they normally do not consider the idea of another language being used as an extra layer of encryption.

Stream ciphers, as opposed to block ciphers, hide the pre-known fact that a message will be sent using the local language and; thus, are less prone to the simpler attacks. This paper explains several stream ciphers and how the addition of a foreign language as an extra layer on top of the current stream cipher’s capabilities can serve as an extra deterrent for attacks.

## 2 Stream ciphers

A stream cipher (or pseudo-random generator) is an algorithm that takes a short random string, and

042  
043  
044  
045  
046  
047  
048  
049  
050  
051  
052  
053  
054  
055  
056  
057  
058  
059  
060  
061  
062  
063  
064  
065  
066  
067  
068  
069  
070  
071  
072  
073  
074  
075  
076  
077  
078  
079  
080  
081

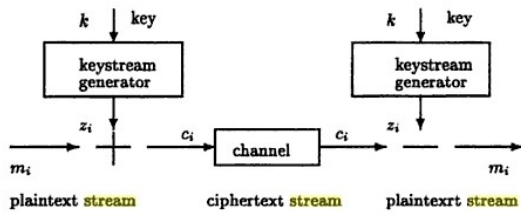


Figure 1: Key generation in a stream cipher. (Cusick et al., 2004)

expands it into a much longer string, that still looks random to adversaries with limited resources.<sup>1</sup>. Stream cipher algorithms are typically used as a mechanism for encryption on devices such as wireless routers where encryption is required in order to not expose the data packets that are being passed as messages. Since the data that is being passed back and forth is passed randomly and in real time, data transfer can be considered a stream of packets from one endpoint to another. A stream cipher specifies a device with internal memory that enciphers the  $j$  digit  $M_j$  of the message stream into the  $j$  digit  $C_j$  of the cipher text stream by means of a function which depends on the secret key and the internal state of the stream cipher at time  $j$ . The sequence  $Z_0, Z_1, Z_2, \dots, Z_n$  which controls the enciphering is called the *key stream* or *running key*. The deterministic automation which produces the key stream from the actual key  $k$  and the internal state is called the *running-key generator*, or *key-stream generator*. (Cusick et al., 2004)

The key-stream generator generates the running key sequence described above as the key stream. The key-stream generator combines digit by digit the key sequence, or the running key, on top of the plain text sequence in order to obtain the ciphered text that can be considered somewhat easy to attack due to the fact that the text, although in a ciphered format, is normally produced using letters and/or words from the local language where the data stream occurs.

The "channel" from Figure 1 above shows the typical stream's flow as it passes through the stream ciphering process. The ciphered text, meant to be secure, has been found to be vulnerable to attacks due to the frequency of letters in a common alphabet such as English. (Singh and Johari, 2015) If the ciphered text's secret key is used more than once, while appearing random to its adversaries, it can be

<sup>1</sup>[https://en.wikipedia.org/wiki/Stream\\_cipher](https://en.wikipedia.org/wiki/Stream_cipher)

easily decrypted by a skilled cryptographer even though stream ciphers operate with a time-varying transformation on the individual plain text digits. Stream ciphers depend on a pre-agreed secret for their key encryption. That idea in itself could be considered a security breach since both parties have to maintain the same secret. In this paper, the focus is on the encrypted stream and how to avoid attacks that use cryptic algorithms to decrypt the streams with prior knowledge of a particular language, especially when the same key is used more than once.

There are two major components of a stream cipher algorithm: 1) a short input string (referred to as the **key** in Figure 1 and 2) a long output string called the *key stream*. Stream ciphers can be used for shared-key encryption, by using the output stream as a one-time-pad. (Burnett and Paine, 2001) The stream cipher can deploy random digits or letters for its encryption and decryption process, this is known as a synchronous stream ciphering process (Rueppel, 1986). Additionally, there is another model called *self-synchronous stream ciphering* (Lamba, 2010) that calculates ciphered digits using the previous cipher text's digits which automatically synchronize the key generator when receiving the digits. Both stream ciphering approaches can be considered part of the stream cipher paradigm. In this paper, an additional ciphering mechanism is described to further encrypt the cipher for heightened security that uses natural language as an extra layer to the key stream.

## 2.1 Stream cipher word frequency

The random digits (numerical or alphabetical) that are formed as part of the encrypted stream in a stream cipher are usually in a local language known to the cryptographer. For that reason, a cryptographer's attempt to decrypt a stream cipher that has been created using an alphabet known by both parties and used multiple times can be considered vulnerable. An attack could be formed that uses the easy-to-discern frequency of digits that focuses on the higher occurring digits, or letters, in the local language.

A good example of the weakness of a stream cipher that would typically use a local alphabet for its digits is the RS4 encryption algorithm described in Section ?? . Here, we can assume that the algorithm is easier to attack due to the knowledge of the actual language at hand. A more concrete

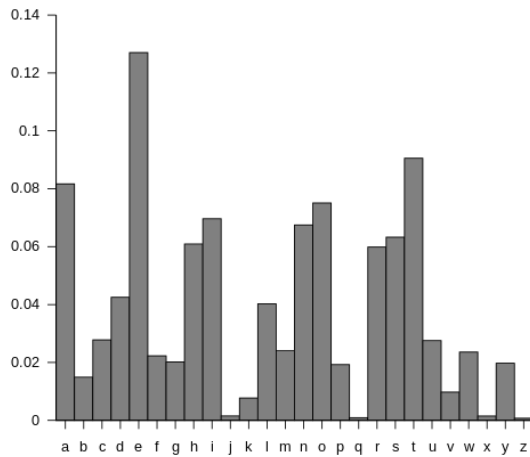


Figure 2: English language letter frequency.<sup>2</sup>

example of local language vulnerabilities could be found in a city such as Frankfurt, Germany where an attacker would probably attack a wireless network using the German language due to the fact that more than ninety percent of Frankfurt’s inhabitants use German (or Dutch) as their language of choice. On the contrary, the same principle may not be applicable for cities of higher immigration such as Miami, Florida, USA where the spoken language (Spanish) is not the official language of the country (English).

Figure 2 displays the typical frequency distribution of letters in a word of the English language and gives conclusive notions that, by using the knowledge of the digit, or letter, distribution in a language, an attacker may be able to establish an attack model paradigm with ease.

It is clear that a stream cipher whose ciphered output is generated using the English language would, judging from Figure 2 above, probably contain the letter “e” within its context. Therefore, by using the fact that certain letters are more likely to be included in a stream, attacks are normally crafted using higher occurring letters from the local language’s alphabet.

A cipher that is streamed, specifically a streaming cipher that uses the same key and input data will produce identical key streams if used with the same key and input data over successive operations. Since the key stream is frequently combined with plain text using an invertible operation, this means that successive cipher texts can be combined to produce a combination of the plain text. (Burnett and Paine, 2001) That makes an attack, such as a replay attack (El Abbadi and Jamouli, 2021), a good

candidate for attack because one could identify the commonality of a repeating stream using easy-to-obtain tools such as Wireshark<sup>3</sup> and Aircrack-Ng<sup>4</sup>. A cryptographer could use those tools and a wifi stream to apply a special type of algorithm implementing the replay functionality that combines the knowledge of local language digit occurrences and possible input phrases to stage an attack.

In addition to simple repetition detection, local language frequency gives way to a high amount of redundancy. In the case of messages with a high amount of redundancy (like in natural language or other data formats), error propagation may be sufficient to detect modifications to a streamed message, but in general an additional cryptographic operation is needed to guarantee the integrity of a message. (Hudde, 2009) Stream ciphers are normally processed in real time and the size and quantity of data this is passed via the two endpoints of a stream are normally unknown. Ideally the algorithm that produces ciphered text in a stream would be random enough such that simple word frequency tactics and reasoning would not be enough for a cryptographer to decrypt. However, due to the easily attainable algorithms that are highly publicized and other general factors that apply to most stream cipher algorithms, stream ciphers are still vulnerable to attack and require an extra layer of security.

## 2.2 Stream cipher vulnerability

As described, stream ciphers and their counterparts, block ciphers, are vulnerable due to word frequency probability, local language use, and repetitiveness. The stream cipher key is dependent on the key generator which may produce output of a particular stream cipher that could be considered less weak due to its key. If a key has been generated using a weak algorithm, then the attacks described earlier can be executed with ease. Since many of the key generation algorithms are already published, certain algorithms have been proven to be more vulnerable.

### 2.2.1 RC4 algorithms in stream ciphers

In a strong key stream generator, each bit of the output will depend on the entire key for its value, and the relationship between the key and a given bit (or set of bits) should be extremely complicated. (Wash, 2001) The most widely used stream cipher is the RC4 stream cipher. RC4 is currently found in

<sup>3</sup><https://www.wireshark.org/>

<sup>4</sup><https://www.aircrack-ng.org/>

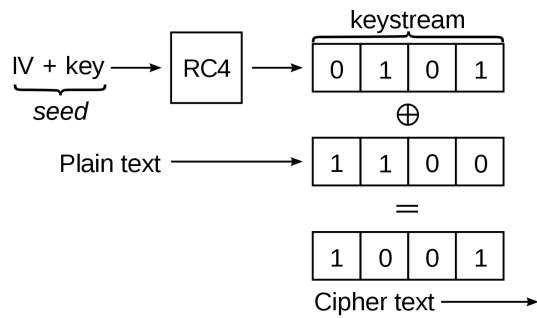


Figure 3: RC4 Stream Ciphers for WEP.<sup>6</sup>

various applications. In the stream cipher context, RC4 can be commonly found in a wireless protocol called wired equivalent privacy (WEP). WEP has already been considered a vulnerable protocol due to its stream cipher key vulnerability; newer protocols such as wifi protected access (WPA) have already been introduced to replace WEP. WEP is especially vulnerable when the beginning of the output key stream is not discarded, or nonrandom or related keys are used; some ways of using RC4 can lead to very insecure cryptosystems.<sup>5</sup>

RC4 generates a key stream using an internal state algorithm that has a permutation of 256 possible bytes with two 8-bit pointers. The pointers randomly swap bytes pointed to in order to XOR message bytes. RC4 can be considered a simple and quite elegant algorithm. Nonetheless, its simplicity combined with the three factors mentioned at the beginning of this section make it vulnerable to attacks such as the bit-flipping attack that use the knowledge of the algorithm to decipher the text in the stream. It can be understood from Figure 3 that an RC4 application can be deciphered by knowledge of the algorithm easily found on the internet or other publications. A denial of service attack (DOS) (Schuba et al., 1997) could be used to insert plain text that would produce a predictable output exposing the stream cipher's algorithm and, thus, makes it easier for an attacker to attack stream ciphered text. For example, previous work (Klein, 2008) presented an analysis of an RC4 stream cipher showing more correlations between the RC4 key stream and the key and was able to crack an RC4 encrypted algorithm for WEP in under a minute.

RC4 attacks are now commonplace and almost any primitive hacker can use the knowledge of algo-

rithms such as the RC4 algorithm to attack a stream such as the streams found in wireless WEP technologies for many commonly sold routers on the market. Many variations of the RC4 algorithm have been implemented and, unfortunately, successfully decrypted without the knowledge of the key. The RC4 creates a one-time key of about 24 bits for its security. 24 bit length really cannot be considered safe. The fact that the RC4 algorithm is readily available combined with its key shortness and use of local language digits make it highly vulnerable. The RC4 algorithm is just one of the many algorithms for stream ciphers that can be considered weak due to the vulnerabilities described.

### 2.2.2 LFSR algorithms in stream ciphers

The LFSR (Linear Feedback Shift Register) algorithm (Wang and McCluskey, 1988) is yet another, considerably insecure, algorithm that can be used in stream ciphers to generate a key. LFSR depends on a previous state by applying a linear function to it. The most common linear function is to take the previous state's bit pattern and XOR it with some bits to modify the overall state. LFSR eventually repeats because its registers have a finite number of states and, due to the states finiteness, could be considered less secure when states are cycled repeatedly. Nonetheless, if a LFSR algorithm is chosen with a strategic security plan in mind, it could appear randomly acyclic when under attack. LFSRs have long been used as pseudo-random number generators for use in stream ciphers (especially in military cryptography), due to the ease of construction from simple electromechanical or electronic circuits, long periods, and very long periods, and very uniformly distributed output streams.<sup>7</sup> A skilled attacker could decrypt an LFSR quite easily using output text combined with the simulation of a receiver to gain access to encrypted information. One such attack is known as the correlation attack (Meier and Staffelbach, 1988).

A correlation attack can be devised to understand the boolean, cyclic nature of the LFSR algorithm. Predictive possibility tables can be drawn that take the possible input and output in order for the hijacker to be able to decrypt the stream cipher using Boolean logic like Figure 4. So, the decryption would intercept the stream cipher, apply the key stream generation algorithm table using statistical probability, and gain access to the stream. In order

<sup>5</sup>[https://en.wikipedia.org/wiki/Wired\\_Equivalent\\_Privacy](https://en.wikipedia.org/wiki/Wired_Equivalent_Privacy)

<sup>7</sup>[https://en.wikipedia.org/wiki/Linear-feedback\\_shift\\_register](https://en.wikipedia.org/wiki/Linear-feedback_shift_register)

Boolean function output table

$x_1$	$x_2$	$x_3$	$F(x_1, x_2, x_3)$
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	0
1	0	1	0
1	1	0	1
1	1	1	1

Figure 4: A boolean table for a correlation attack.

to statistically decrypt a stream cipher algorithm, the cryptanalyst would only have to apply the correlative technique in a key generation algorithm with an algorithm such as the Geffe generator.<sup>8</sup>

If the stream cipher’s algorithm is implemented using LFSR, the key stream may be too vulnerable and easy to attack, even with another layer of non-local language applied. While a natural language layer could be applied to a stream cipher with LFSR, the simple fact of the repetitiveness in LFSRs cycle make it easier to attain the correct keys. Nonetheless, if one could find a correlation between the output of one of the shift registers and the key stream, then one can try to find the initial state of this LFSR independently of the other LFSRs.<sup>9</sup> Correlation attacks are the most common way to attack LFSR key generations and serve as an example of the weakness of stream ciphers. Correlation attacks can be considered extremely dangerous and stream ciphers extremely susceptible; extreme care must be taken when designing stream ciphers in order to protect against correlation attacks.

### 3 Natural language encryption

Natural languages are based on the day-to-day conversations that we experience and can be considered as pieces of information that help us as humans

<sup>8</sup>[http://en.wikipedia.org/wiki/Correlation\\_attack](http://en.wikipedia.org/wiki/Correlation_attack).

<sup>9</sup><http://www.eit.lth.se/fileadmin/eit/courses/edi051/projects/corattack/CorrAt.pdf>

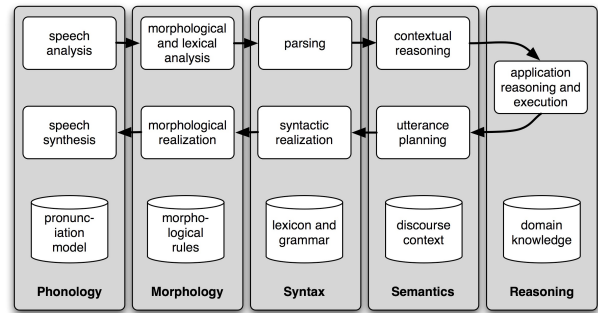


Figure 5: NLP in software (Bird et al., 2009).

to communicate more effectively within our domain. Natural languages are governed by implied rules for which natural selection inherently defines. (Pinker and Bloom, 1990). While we can attempt to define those rules using techniques such as finite state transducers (FST) (Mohri, 1997), it can be assumed that natural language rules are nearly impossible to approximate via mathematics or grammatical structure, at least with one-hundred percent accuracy. This motivates the study of their use in cryptography as a stronger cipher because they are complex and difficult to solve even by those highly trained in statistical digit, or letter, probability. Overall, it makes practical sense that a key generated with an extra layer of natural language may be more secure due to its grammatical and mathematical incorrectness that make prediction of the key more complex.

NLP is the process of a computer extracting meaningful information from natural language input and/or producing natural language output. NLP, as Figure 5’s complexity shows, is often considered the problem and not the solution due to the difficulty of the task of accepting natural input and producing natural output that are governed by implicit language grammatical models that may not be traceable to any group of persons. Notwithstanding, if a stream cipher is created by the implementation of a natural language that is typically spoken where the stream is being transmitted, the likelihood that the stream will be decrypted using the local natural language is higher than if it were to use a non-local language. In this section, an introduction to the idea of encryption by using an atypical natural language is proposed.

#### 3.1 Plain text language encoding

A stream cipher that is used for encoding performs its encryption at the level of individual letters or

bits. Typically, a cipher, whether a stream or block cipher, uses *plain text* letters to encrypt a message. A cryptanalyst is considered an expert at decoding *plain text* letters. It is not a surprise, then, that plain texts are often used as targets for decryption algorithms that a cryptanalyst may routinely use. Plain text taken from everyday sources such as newspapers, recorded telephone conversations, and wireless traffic can be considered a prime target for an attack. The knowledge that a specific target may be written in plain text combined with the fact that a target's implemented language is probably the most common language used within the target's geographic location allows cryptographers to devise plain-text algorithms using bits or letters from the local language. For example, the following scheme could be used as a way of encrypting letters in the English language:

A B C D E F GH I J KLMNOPQR S T  
UVWXYZ QWE R T YU I O P A S D F G H  
J K L Z X C V B N M

This general system, according to Stallings (Stallings, 2017), is called a **mono-alphabetic substitution**, with the key being the 26-letter string corresponding to the full alphabet. The encryption key in this example is "QWERTYUIOPASDFGHJKLZXCVBNM". For the key above, the plain text word: "ATTACK" would be transformed into the ciphered text word: "QZZQEA". Plain text can be described as the typical writings that we see written in the media that surrounds us and is normally near grammatically correct. In order to understand plain text, plain text's reader would have to have a basic knowledge of grammatical rules that govern the language that the plain text is written in. Stream cipher encoding which uses plain text is insecure when using a typical language. If the plain text is encoded using a highly redundant language – such as English or any other natural language – it can be extracted without knowledge of the key. (Burnett and Paine, 2001) Ideally, if a sender and a receiver would like to communicate using ciphers and plain text, an encrypted layer must be applied to the plain text in order to make the cipher less vulnerable to attack. One such case where plain text was practically undecipherable was found in a study<sup>10</sup> done of the Al-Qaeda group in the UK that used a combination of known natural languages from countries where the group exists such as Pakistan, Yemen, and Sudan. It was almost

<sup>10</sup><http://www.wnd.com/2009/11/116381>

impossible for the local cryptographers who were accustomed to decrypting messages sent in the native local language, English, to decrypt messages encrypted with natural languages from other countries. The encoded messages were finally decrypted by employing cryptographers from the aforementioned countries. Between them, the code-breakers spoke all the dialects that form the basis for the code. Several of them have high-value skills in computer technology. The local language, native to the Al-Qaeda group, was used as a way of encrypting plain-text messages that could not be understood by the local, mostly English native speakers, inhabitants. Plain text can seem somewhat simple to decrypt. But, if the plain text is written in a language that is not known to the reader and if that language is written in a natural (unstructured) form, it would be much more difficult to decrypt.

### 3.2 Natural language layer for ciphers

The presence of a natural language can be seen as the weak link of a stream or block cipher. While it may be difficult to determine the text of an encrypted message, given the natural language of a base encryption, a cryptographer can use word frequency algorithms, such as the Berlekamp-Massey Algorithm (Atti et al., 2006), to exploit one weakness in the decryption process. In that respect, NLP can be seen as a weakness in stream ciphers; however, NLP can also be used for heightened security.

Alternative language constructs can be used as a way of obfuscating encrypted keys. In order to hide the keys, a layer of encryption for increased security can be applied in a language spoken by non-natives to enhance the quality of the algorithm. In NLP, the term "noise" can be defined as the extra phonetics or disturbance inherent to a language that makes the language hard to understand. Languages such as German could be considered "noisy" forms of the English language. (Bakir, 2016) With sufficient distortion, or noise, a language can be undecipherable and nearly impossible to dissemble. One may consider this technique as a form of "scrambling" (Phillips et al., 1971). For example, in the United States, it is known that the Federal Bureau of Investigation has scrambled mobile phone signals when conducting investigations. (Zou et al., 2019) The noise that one hears when a mobile phone signal is scrambled makes conversations nearly impossible to understand.

This paper proposes the addition of natural lan-

505 guage to block and stream ciphers by using a non-  
506 native “noisy” layer to scramble text message to  
507 the point where a local cryptographer would have  
508 a harder time decrypting the message, similar to  
509 the scrambled mobile phone message described  
510 in the previous paragraph. With the addition use  
511 of a foreign language for use of scrambling the  
512 text, the attacking cryptographer would first have  
513 to decrypt a message and then translate it, the trans-  
514 lation would be in two or more languages make it  
515 very difficult for the most state-of-the-art machine  
516 translation systems like Google Translate<sup>11</sup>.

517 Since translation of two or more encrypted lan-  
518 guages added as layers to stream ciphers would re-  
519 quire that a parallel key known by both the sender  
520 and receiver is established and agreed upon. The  
521 parallel key combined with a non-native language  
522 is considered as the "noise" of an already encrypted  
523 stream. If the noise caused by the encrypted natural  
524 language is sufficient enough to scramble encrypted  
525 messages, the type of security can be considered  
526 an addition to current standardized layers. One ex-  
527 ample of how this has been done in the past is the  
528 use of language mixing by terrorist. (Guidère et al.,  
529 2009) Consider the following example:

- 530 • **native language:** bob is a joker
- 531 • **simple encryption algorithm:** b=a, o=c, i=r,  
532 s=z, a=q, j=g, k=e, e=x, r=t
- 533 • **result:** aca rz q gcext

534 For a cryptographer, the example above would  
535 take seconds to decrypt. But, if an additional lan-  
536 guage was added as an extra layer of encryption,  
537 and if the language was a mixture of two or more  
538 languages, the message would be tougher to de-  
539 crypt. Below is the same example using *Spanglish*,  
540 the mixture of Spanish with English, a language  
541 without official rules spoken in several parts of the  
542 world. (Ardila, 2005)

- 543 • **mixed language:** bob es un joker
- 544 • **simple encryption algorithm:** b=a, o=c, i=r,  
545 s=z, a=q, j=g, k=e, e=x, r=t, u=h, n=l
- 546 • **result:** aca xz hl gcext

547 While the example above is simple in nature, the  
548 decryption technique is more difficult to decipher  
549 due to the language not only lacking grammatical

<sup>11</sup><https://translate.google.com>

550 sense but also having no meaning after decryption.  
551 An attack on the encrypted text above, for example,  
552 would be difficult for a person who deals only in  
553 English. Additionally, if we assume that the cryp-  
554 tographer were from a country where English and  
555 Spanish were not spoken (Russia for example), the  
556 decryption above would be even more difficult.

557 The idea of using natural language in stream ci-  
558 phers will deviate cryptographers to break it ,and if  
559 even they break, it will be hard for them to under-  
560 stand because of the ignorance of the natural lan-  
561 guage that we are using in the cipher (i.e. ARABIC,  
562 Chinese or Japanese, Italian or Greek languages)  
563 (Mahmud, 2008). In order to apply a natural lan-  
564 guage to a stream cipher, a dependency must be  
565 established and the encoding language set as a part  
566 of the encryption. The application of the language  
567 on top of the stream layer requires that a Unicode  
568 representation deemed as input for the second lan-  
569 guage is created. After the representation has been  
570 combined to the stream, an XOR operation is per-  
571 formed on the binary Unicode representation of the  
572 input in the second natural language and a binary  
573 key is then used to generate an encrypted output.  
574 Decryption is finalized by the receiving end using  
575 the reverse order.

576 While NLP can be used as a key deterrent against  
577 attacks, it is still not full proof. It is important,  
578 for a better level of security, that the generated  
579 keys not be repeated twice. Repetition avoidance  
580 applies to stream ciphers specifically because of the  
581 encryption cycle that occurs. It would also be wise  
582 that the stream cipher’s encryption algorithm and  
583 its language counterpart use languages that are not  
584 so typical to a specific region. For example, if a key  
585 generator algorithm created for a wireless router is  
586 made in Spain, it would not be wise to create the  
587 ciphered text using an algorithm that translates to a  
588 nearly similar alphabet such as Spain’s neighboring  
589 France.

590 It can be noted that by frequency alone stream  
591 ciphers are considered vulnerable. In some ciphers,  
592 such properties of the natural language plain text  
593 are preserved in the cipher text, and these patterns  
594 have the potential to be exploited in a cipher text-  
595 only attack. Language models typically written in  
596 published algorithms can be trained to learn ciphers.  
597 While research is still ongoing, some language al-  
598 gorithms can learn by repetition. Therefore, the  
599 pure repetitiveness of certain words such as the ar-  
600 ticle “the” in English can serve as a weak point in

601 a stream cipher text encryption. Cryptographers  
602 dedicate themselves to finding patterns in common  
603 texts that render symbolic patterns. By applying  
604 the NLP technique described here, decryption be-  
605 comes more difficult due to the language barrier  
606 that a cryptographer would probably display. Con-  
607 trastingly, multi-lingual cryptographers are more  
608 likely to find patterns in ciphered texts that have  
609 been encrypted with non-native languages due to  
610 the fact that they are probably more likely to have  
611 seen specific data points within language patterns  
612 that serve as key indicators that a stream may have  
613 been encrypted using another language.

614 The insertion of a distinct language in a stream  
615 is not difficult to perform. The most important  
616 role that language plays in the stream cipher is the  
617 protective role of defense. As is typical in stream ci-  
618 phers, both the sender and receiver must be aware  
619 of the language applied and its rules should be  
620 made clear before a key is generated. When ap-  
621 plying the XOR described above as a binary set, if  
622 one of the words does not match a set pattern, the  
623 decryption algorithm may be thrown off and more  
624 difficult to read. While this may sound simple to do,  
625 local languages, by their sheer use, are less likely  
626 to be bound by rules which make them less useful  
627 in general. Regardless, if a common language can  
628 be understood in a local area, rules can be applied  
629 to inject the proper encryption. The parallel key  
630 (along with key stream bits) for this type of ciphers  
631 can be the languages name itself or other world of  
632 common interest between two parties may be used.  
633 (Mahmud, 2008)

634 The XOR operation can be considered the single  
635 most important part of applying a NLP technique  
636 to a stream. An XOR operation is also a key fo-  
637 cus of attackers. When adding the language in as  
638 an extra layer of protection, the key generation al-  
639 gorithm must be careful that a replay attack can't  
640 reproduce through redundancy techniques a way of  
641 combining series of messages. For that reason, it  
642 is more secure to add a non-local language into the  
643 XOR operation. Randomness plays an important  
644 key in any key generation technique for stream ci-  
645 phers. Hence, a naturally spoken language should  
646 be clearly known by both the sending and receiving  
647 algorithms in order to avoid simplistic yet mean-  
648 ingful collisions that can be translated using a key  
649 deciphering algorithm.

650 The principle vulnerability in a stream cipher,  
651 and the reason why the XOR operation is the most

652 important, is the frequency at which letters or sym-  
653 bols occur within the encryption language. The  
654 final binary added on as a layer discussed in this  
655 paper should help to disqualify stream ciphering en-  
656 cryption detection algorithms. The likelihood of at-  
657 tack would highly decrease if a key is created with  
658 high security by using a key that is not repeated and  
659 random along with the extra layer of security that  
660 languages provide. An attacker would have to have  
661 great knowledge of languages and decryption in or-  
662 der to recognize patterns that may occur; especially,  
663 if the XOR operation implies a mixture of language  
664 similar to those used by the military gotten from  
665 native American tribes(Marshall, 2012).

#### 666 4 Stream reliability and Conclusion

667 Application and protocol designers, even those with  
668 experience and training in cryptography, cannot  
669 be expected to always identify accurately the re-  
670 quirements that must be met for a mode to be used  
671 securely or the conditions that apply to the applica-  
672 tion at hand. (?) Private enterprises such as Google  
673 and Microsoft receive millions of attacks a year.  
674 Whether an enterprise level user or a simple home  
675 user, network security, no matter at what level, can  
676 be attributed to a price with information containing  
677 a value. The protection of that information really  
678 depends on its value. Credit card numbers may be  
679 considered more important than a user id for an  
680 adventure gaming website. Higher valued items  
681 and messages are retrieved via network streams  
682 of data and are captured and decrypted by skilled  
683 cryptographers. The heightened sense of security  
684 towards streams must be considered important.

685 Attacks are direct and easy to accomplish with  
686 the current attacker tools available. Wireless WEP  
687 attacks have proven to be as simple as inserting  
688 a disc or usb into a laptop and pressing enters.  
689 Although the latest wireless networks seem to be  
690 more secure and robust, keys are retrieved through  
691 cryptology and it is inevitable that algorithms will  
692 be created to decrypt the most difficult encryption.  
693 But, if tactics such as the NLP layer described in  
694 this paper are employed, a cryptographer's job can  
695 be made considerably more difficult.



## References

- 696
- 697 Alfredo Ardila. 2005. Spanglish: an anglicized spanish  
698 dialect. *Hispanic Journal of Behavioral Sciences*,  
699 27(1):60–81.
- 700 Nadia Ben Atti, Gema M Diaz-Toca, and Henri Lom-  
701 bardi. 2006. The berlekamp-massey algorithm revis-  
702 ited. *Applicable Algebra in Engineering, Communi-  
703 cation and Computing*, 17(1):75–82.
- 704 Cigdem Bakir. 2016. Automatic voice and speech recog-  
705 nition system for the german language with deep  
706 learning methods. *International Journal of Applied  
707 Mathematics Electronics and Computers*, (Special  
708 Issue-1):399–403.
- 709 Steven Bird, Ewan Klein, and Edward Loper. 2009. *Nat-  
710 ural language processing with Python: analyzing text  
711 with the natural language toolkit*. " O'Reilly Media,  
712 Inc."
- 713 Steve Burnett and Stephen Paine. 2001. *RSA Security's  
714 official guide to cryptography*. McGraw-Hill, Inc.
- 715 Thomas W Cusick, Cunsheng Ding, and Ari R Renvall.  
716 2004. *Stream ciphers and number theory*. Elsevier.
- 717 Reda El Abbadi and Hicham Jamouli. 2021. Takagi-  
718 sugeno fuzzy control for a nonlinear networked sys-  
719 tem exposed to a replay attack. *Mathematical Prob-  
720 lems in Engineering*, 2021.
- 721 Mathieu Guidère, Newton Howard, and Shlomo Arg-  
722 mon. 2009. Rich language analysis for counterterror-  
723 ism. In *Computational Methods for Counterterror-  
724 ism*, pages 109–120. Springer.
- 725 Hans Christoph Hudde. 2009. Building stream ciphers  
726 from block ciphers and their security. *Seminararbeit  
727 Ruhr-Universität Bochum*.
- 728 Andreas Klein. 2008. Attacks on the rc4 stream cipher.  
729 *Designs, codes and cryptography*, 48(3):269–286.
- 730 Chhattar Singh Lamba. 2010. Design and analysis of  
731 stream cipher for network security. In *2010 Second  
732 International Conference on Communication Soft-  
733 ware and Networks*, pages 562–567. IEEE.
- 734 Maqsood Mahmud. 2008. Natural language (arabic)  
735 as a strengthening layer for stream ciphers in wire-  
736 less networks. In *Proceedings of the 17th IASTED  
737 International Conference*, volume 609, page 130.
- 738 Sreekanth Malladi, Jim Alves-Foss, and Robert B Heck-  
739 endorn. 2002. On preventing replay attacks on secu-  
740 rity protocols. Technical report, IDAHO UNIV  
741 MOSCOW DEPT OF COMPUTER SCIENCE.
- 742 Suzanne Marshall. 2012. A hidden story: American  
743 indian code talkers. *DttP*, 40:27.
- 744 Willi Meier and Othmar Staffelbach. 1988. Fast cor-  
745 relation attacks on stream ciphers. In *Workshop on  
746 the Theory and Application of Cryptographic Tech-  
747 niques*, pages 301–314. Springer.
- Mehryar Mohri. 1997. Finite-state transducers in lan-  
guage and speech processing. *Computational linguis-  
tics*, 23(2):269–311.
- VJ Phillips, MH Lee, and JE Thomas. 1971. Speech  
scrambling by the re-ordering of amplitude samples.  
*Radio and Electronic Engineer*, 41(3):99–112.
- Steven Pinker and Paul Bloom. 1990. Natural language  
and natural selection. *Behavioral and brain sciences*,  
13(4):707–727.
- Rainer A Rueppel. 1986. Stream ciphers. In *Analysis  
and Design of Stream Ciphers*, pages 5–16. Springer.
- Christoph L Schuba, Ivan V Krsul, Markus G Kuhn,  
Eugene H Spafford, Aurobindo Sundaram, and Diego  
Zamboni. 1997. Analysis of a denial of service attack  
on tcp. In *Proceedings. 1997 IEEE Symposium on  
Security and Privacy (Cat. No. 97CB36097)*, pages  
208–223. IEEE.
- Laukendra Singh and Rahul Johari. 2015. Clct: cross  
language cipher technique. In *International Sympo-  
sium on Security in Computing and Communication*,  
pages 217–227. Springer.
- William Stallings. 2017. *Network Security Essentials*.  
Pearson Education Limited.
- L-T Wang and Edward J. McCluskey. 1988. Linear feed-  
back shift register design using cyclic codes. *IEEE  
Transactions on Computers*, 37(10):1302–1306.
- Rick Wash. 2001. Lecture notes on stream ciphers and  
rc4. *Reserve University*, pages 1–19.
- Ying Zou, Ge Zhang, and Leian Liu. 2019. Research on  
image steganography analysis based on deep learn-  
ing. *Journal of Visual Communication and Image  
Representation*, 60:266–275.