

000 001 DASHCO: DATA-AWARE SAT HEURISTICS COM- 002 BINATIONS OPTIMIZATION VIA LARGE LANGUAGE 003 MODELS 004 005

006 **Anonymous authors**
007 Paper under double-blind review
008
009
010
011
012

ABSTRACT

013 The performance of Conflict-Driven Clause Learning solvers hinges on internal
014 heuristics, yet the heterogeneity of SAT problems makes a single, universally optimal
015 configuration unattainable. While prior automated methods can find specialized
016 configurations for specific problem families, this dataset-specific approach
017 lacks generalizability and requires costly re-optimization for new problem types.
018 We introduce DASHCO, a framework that addresses this challenge by learning
019 a generalizable mapping from instance features to tailored heuristic ensembles,
020 enabling a train-once, adapt-broadly model. Our framework uses a Large Lan-
021 guage Model, guided by systematically defined Problem Archetypes, to generate
022 a diverse portfolio of specialized heuristic ensembles and subsequently learns an
023 adaptive selection mechanism to form the final mapping. Experiments show that
024 DASHCO achieves superior performance and, most notably, demonstrates robust
025 out-of-domain generalization where non-adaptive methods show limitations. Our
026 work establishes a more scalable and practical path toward automated algorithm
027 design for complex, configurable systems.
028
029

1 INTRODUCTION

030 The Boolean Satisfiability (SAT) problem is a cornerstone of computational complexity theory Cook
031 (2023) and a problem of immense practical importance Crawford & Baker (1994). Its applications
032 are vast, with modern solvers enabling crucial advances in diverse fields such as formal verifica-
033 tion Prasad et al. (2005), planning Rintanen (2012), and program analysis Harris et al. (2010).
034 While SAT is NP-complete, powerful solvers based on the Conflict-Driven Clause Learning (CDCL)
035 paradigm Audemard & Simon (2009) can often solve massive industrial instances with remarkable
036 efficiency. However, this success is not uniform. The performance of these solvers is critically sensi-
037 tive to the internal algorithmic heuristics that guide their search. Developing effective heuristics has
038 consequently been a central focus of SAT research for decades Audemard & Simon (2012); Liang
039 et al. (2018), traditionally relying on a long and costly process of manual design.
040

041 The challenge of designing effective heuristics has led to the development of hyper-heuristics Burke
042 et al. (2013), which automate the process of selecting or generating algorithms. With the advent of
043 LLMs Achiam et al. (2023); Comanici et al. (2025), this field has seen a surge of innovation. Frame-
044 works like FunSearch Romera-Paredes et al. (2024), Evolution of Heuristics (EoH) Liu et al. (2024),
045 and ReEvo Ye et al. (2024) have demonstrated that LLMs can generate novel and effective heuris-
046 tics for various combinatorial optimization (CO) problems by treating algorithm design as a program
047 search task within an evolutionary framework van Stein & Bäck (2024). While the recent AutoSAT
048 framework Sun et al. (2024) has made important progress by applying this paradigm to SAT solvers,
049 a significant limitation underlies these pioneering works: they are inherently dataset-specific. Their
050 methodology is tailored to a particular training distribution, yielding a single, static heuristic con-
051 figuration optimized for that specific class of problems. Consequently, the resulting solver lacks
052 generalizability, and the expensive search process must be repeated for each new problem family.
053

This limitation is particularly severe because the SAT problem space is enormously heterogeneous,
serving as a universal language for problems from diverse domains like Minesweeper Kaye (2000),
cryptographic analysis Soos et al. (2009), and logistics planning Rintanen (2012). Moreover, the

054 performance of a complex system like a CDCL solver is determined by the intricate and often
 055 non-intuitive **interplay** of its multiple, interacting heuristics for tasks such as restarts Audemard
 056 & Simon (2012); Liang et al. (2018) and phase selection. The synergistic or conflicting effects of
 057 these combinations are notoriously difficult to predict, a challenge that even human experts struggle
 058 to navigate. This implies that simply selecting the best-performing heuristic for each component in
 059 isolation is unlikely to yield a globally optimal solver. It is therefore crucial to holistically explore
 060 the vast combinatorial space of heuristic configurations, rather than merely optimizing individual
 061 components. The goal of such exploration should not be to find a single “best” configuration, as
 062 the optimal choice is deeply instance-dependent given the problem space’s heterogeneity Ansótegui
 063 et al. (2009). Rather, the challenge is to develop a system that can adapt its strategy based on the
 064 problem at hand.

065 To address these fundamental challenges of generalizability and adaptation, we propose **DASHCO**
 066 (**Data-Aware SAT Heuristics Combinations Optimization**). Instead of seeking a single, univer-
 067 sally optimal solver, DASHCO introduces a new paradigm: its objective is to learn a generalizable
 068 mapping from instance characteristics to tailored heuristic configurations. The core of our frame-
 069 work is its data-aware nature, which enables a train-once, adapt-broadly model. DASHCO first
 070 leverages an LLM, guided by high-level Problem Archetypes, to automatically generate a diverse
 071 portfolio of specialized heuristic ensembles, each tailored to different problem structures. Subse-
 072 quently, it learns an adaptive selection mechanism that partitions the instance space based on per-
 073 formance, creating a map to dynamically choose the best-suited ensemble for any new SAT instance.
 074 This approach transforms the problem from finding one optimal point to learning a function over the
 075 entire problem space.

076 We summarize our main contributions as follows:

- 077 • We introduce DASHCO, a novel framework that shifts the paradigm from dataset-specific
 078 optimization to generalizable, data-aware algorithm design, directly addressing the critical
 079 limitation of prior work.
- 080 • We propose a methodology where Problem Archetypes guide an LLM to generate a di-
 081 verse portfolio of specialized heuristic ensembles, enabling an automated exploration of
 082 the complex interactions between different heuristic components.
- 083 • Through extensive experiments, we demonstrate that DASHCO significantly outperforms
 084 baselines. Crucially, it exhibits superior **out-of-domain generalization** compared to meth-
 085 ods that learn a single, non-adaptive configuration, validating the effectiveness and robust-
 086 ness of our approach.

088 2 PRELIMINARIES & RELATED WORK

089 2.1 SAT AND CDCL-BASED SOLVERS

090 Let $V = \{x_1, \dots, x_n\}$ be a finite set of Boolean variables. The corresponding set of literals is
 091 defined as $L = V \cup \{\neg v \mid v \in V\}$. A *clause* C is a finite subset of L , representing the disjunction of
 092 its literals, with the constraint that for any variable $v \in V$, both v and $\neg v$ cannot be simultaneously
 093 present in C . A formula F in Conjunctive Normal Form (CNF) is a set of clauses, $\{C_1, \dots, C_m\}$,
 094 representing their conjunction. A *truth assignment* (or interpretation) is a function $\tau : V \rightarrow \{\top, \perp\}$
 095 that maps each variable to a truth value. The satisfaction of a formula under τ , denoted by $\tau \models F$,
 096 is defined hierarchically:

- 097 1. A literal $l \in L$ is satisfied by τ ($\tau \models l$) if $l = v$ and $\tau(v) = \top$, or if $l = \neg v$ and $\tau(v) = \perp$.
- 098 2. A clause C is satisfied by τ ($\tau \models C$) if there exists at least one literal $l \in C$ such that
 $\tau \models l$.
- 099 3. A formula F is satisfied by τ ($\tau \models F$) if for all clauses $C_i \in F$, $\tau \models C_i$.

100 The Boolean Satisfiability (SAT) problem is the computational task of determining whether a given
 101 CNF formula F is satisfiable, i.e., whether there exists any truth assignment τ such that $\tau \models F$.

102 Modern SAT solvers are predominantly based on the Conflict-Driven Clause Learning (CDCL)
 103 framework. A CDCL-based solver iteratively builds a partial assignment by making decisions and

108 applying Boolean Constraint Propagation (BCP). When a conflict arises, the solver analyzes the
 109 cause, learns a new clause to prevent the same conflict from recurring, and backtracks. This process
 110 is profoundly influenced by a complex interplay of internal heuristics. These strategies guide
 111 key aspects of the search, including which variable to decide on next (branching), when to aban-
 112 don an unpromising search path and restart, and how to diversify the search by adjusting variable
 113 phases. The specific combination and implementation of these heuristics are what differentiate mod-
 114 ern solvers and are the primary determinants of their performance.

115 2.2 LLM-BASED HEURISTIC GENERATION

116 The emergence of LLMs has opened a new frontier for automated algorithm design. By leveraging
 117 their powerful code generation and reasoning capabilities, researchers have started to automate the
 118 discovery of heuristics. A prominent approach is to frame heuristic generation as a program search
 119 problem within an evolutionary framework. FunSearch established this paradigm by evolving pro-
 120 grams to find new mathematical discoveries Romera-Paredes et al. (2024). This was extended to CO
 121 problems by frameworks like EoH Liu et al. (2024) and ReEvo Ye et al. (2024), which use evolu-
 122 tionary algorithms to prompt an LLM to iteratively refine and improve heuristic code for problems
 123 like TSP and online bin packing. ReEvo notably introduced a reflective step, where the LLM pro-
 124 vides textual feedback to guide the evolutionary search, emulating a verbal gradient Ye et al. (2024).
 125 MEoH Yao et al. (2025) models automated heuristic design as a multi-objective optimization prob-
 126 lem, using a dominance-dissimilarity mechanism with an LLM to generate a set of heuristics that
 127 balance performance and efficiency.

128 AutoSAT Sun et al. (2024) was the first to apply this paradigm to the intricate environment of
 129 SAT solvers. Recognizing that generating a competitive solver from scratch is infeasible due to
 130 code complexity, AutoSAT proposed a modular framework where an LLM optimizes specific, pre-
 131 defined heuristic functions within an existing solver. It successfully demonstrated that an LLM could
 132 enhance a baseline CDCL solver to achieve competitive performance.

133 While these methods are powerful, they generally produce a single, universally applied heuristic
 134 or configuration, overlooking the instance-specific nature of algorithm performance. An emerging
 135 paradigm in automated algorithm design seeks to overcome this limitation by partitioning a problem
 136 class into subclasses based on instance features. This allows for the creation of specialized heuristics
 137 tailored to the unique characteristics of each subclass. Our work, DASHCO, applies this principle
 138 of data-awareness to the multi-heuristic, complex environment of CDCL solvers. In doing so, we
 139 bridge the gap between the universal optimization of frameworks like AutoSAT and a more granular,
 140 data-centric approach to algorithm design.

141 2.3 ALGORITHM SELECTION AND PORTFOLIO SOLVERS

142 A major paradigm for tackling instance heterogeneity is portfolio-based algorithm selection, pi-
 143oneered by the influential SatZilla framework Xu et al. (2008). SatZilla leverages a portfolio of
 144 diverse, human-designed solvers and uses machine learning models to predict the best-performing
 145 one for a given instance based on its features. DASHCO inherits this data-driven philosophy but
 146 introduces a fundamental novelty: rather than selecting from a portfolio of pre-existing solvers, it
 147 first uses an LLM to *automatically generate* a new portfolio of fine-grained heuristic ensembles.
 148 DASHCO is thus not only an algorithm selector but also an automated portfolio generator, a key
 149 distinction that significantly expands the space of possible solver configurations.

150 3 METHODOLOGY

151 To address the challenge of dataset-specific optimization and the lack of generalizability in prior
 152 work, we propose DASHCO. The core of our methodology is a paradigm shift: instead of repeatedly
 153 executing an expensive search for a single, specialized solver for each new problem family, our goal
 154 is to construct a single, robust, and adaptive framework that generalizes across them.

155 Specifically, DASHCO’s objective is not to find one *best* heuristic ensemble, but to learn a rich *map-
 156 ping* from the instance feature space to the space of high-performance heuristic configurations. This
 157 is achieved by first creating a diverse portfolio of specialized heuristic ensembles and then learn-

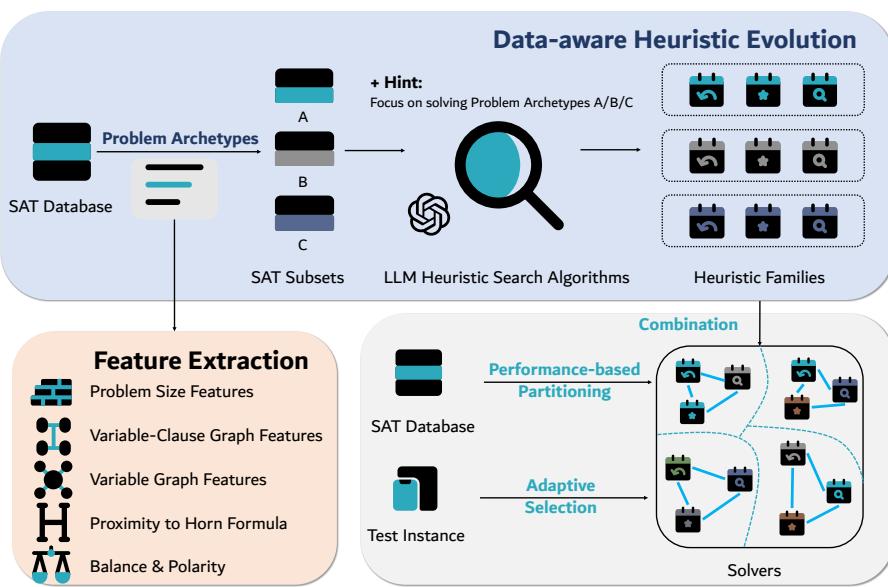


Figure 1: An overview of the DASHCO framework, illustrating the three primary stages: (1) Data-Aware Heuristic Evolution, (2) Instance Space Partitioning, and (3) Adaptive Heuristic Selection.

ing an intelligent selection mechanism. Once constructed, this framework can be deployed on new, unseen SAT instances from diverse problem families, dynamically selecting a suitable configuration without the need for re-optimization. This “train-once, adapt-broadly” approach is designed to be more practical and efficient for real-world applications. The framework operates in three main stages: (1) Data-Aware Heuristic Evolution, to build the rich portfolio of heuristic components; (2) Instance Space Partitioning, to learn the mapping between instance types and optimal ensembles; and (3) Adaptive Heuristic Selection, to apply this mapping to new instances.

3.1 HEURISTIC MODULES IN CDCL-BASED SOLVERS

Following the modular design of AutoSAT, we focus on optimizing a set of critical, independent heuristic functions within a CDCL-based solver. For this work, we target three key heuristics:

1. **Restart Policy (`restart`):** This module implements a crucial strategy to prevent the solver from becoming trapped in unproductive regions of the search space. A restart policy determines when to abandon the current search path and backtrack to the top decision level. While the current partial assignment is discarded, all learned clauses are retained, allowing the solver to begin a new search attempt with more information. Modern policies are often dynamic, adapting their restart frequency to the search progress to effectively combat heavy-tailed runtime distributions Luby et al. (1993); Audemard & Simon (2012).
2. **Phase Selection (`rephase`):** This heuristic acts as a diversification mechanism by managing the default polarity (true or false) assigned to variables. During branching, the solver often uses a variable’s last assigned value as a default choice. Over time, these saved phases can lead to search stagnation. The `rephase` function is called periodically to alter these saved phases, for instance by resetting them or flipping them. This forces the solver to explore different branches first, effectively diversifying the search and pushing it into new areas of the solution space Jeroslow & Wang (1990).
3. **Variable Bumping (`bump_var`):** This module is a core component of adaptive branching strategies. When the solver encounters a conflict, it analyzes the implication graph to identify the variables responsible for the contradiction. The `bump_var` function is then invoked to increase a numerical score, often called an ‘activity score’, associated with each of these variables. By elevating the scores of conflict-prone variables, this mechanism dynamically influences the main branching heuristic to prioritize them in future decisions.

216 This effectively focuses the search on the most constrained and active parts of the problem
 217 space Biere et al. (2009); Liang et al. (2016).
 218

219 Our goal is to find not just one optimal implementation for each of these, but a diverse set of effective
 220 implementations that can be combined into powerful ensembles.
 221

222 3.2 FEATURE EXTRACTION FOR SAT INSTANCES

224 A critical precursor to any data-aware method is the ability to characterize instances with a quantitative
 225 feature vector. We define a 37-dimensional feature vector $v(j)$ that maps a given SAT instance j
 226 to a vector in \mathbb{R}^{37} , designed to capture a comprehensive set of its structural and statistical properties.
 227

228 Our feature set is constructed by augmenting an established feature template with new global de-
 229 scriptors. The foundation consists of the complete 33-feature set adapted from the well-established
 230 algorithm selection framework SATzilla, excluding only the most computationally expensive fea-
 231 tures that require pre-solving. To further enhance the descriptive power, we then introduced 4 ad-
 232 ditional lightweight, high-level features that capture global properties such as overall polarity bias
 233 and constraint density. This approach combines a proven, powerful feature set with novel global
 234 metrics, balancing expressive capability with the efficiency required for a dynamic, per-instance
 235 selection model. A detailed breakdown of all 37 features is provided in the Appendix.
 236

237 3.3 DATA-AWARE HEURISTIC EVOLUTION

238 To build our library of heuristics, we introduce a guided search process centered around a set of
 239 predefined **Problem Archetypes**. The purpose of these archetypes is to guide the LLM to search
 240 for heuristics along high-level, human-understandable directions. This ensures that the resulting
 241 heuristics are differentiated, each specialized for a particular type of problem structure, which is
 242 crucial for building a powerful and diverse portfolio for later combination.
 243

244 **Defining Problem Archetypes and Search Environments.** A key design choice is to define these
 245 archetypes manually (e.g., ‘highly-constrained problems’ or ‘instances with heterogeneous clause
 246 structures’), rather than using automated clustering at this stage. This is because our goal is to
 247 create effective, semantic prompts to guide the LLM. These interpretable concepts serve as more
 248 potent guidance for the LLM’s creative process than the abstract centroids produced by a black-box
 249 clustering algorithm.
 250

251 These archetypes serve a crucial dual role in creating specialized search environments. For each
 252 textual archetype d_i , we first leverage the LLM to identify a relevant subset of features from
 253 our 38-dimensional space that best characterize it. The corresponding data subset $I_i \subseteq I_{train}$ is
 254 then curated by splitting the training set at the 50% threshold based on the values of these LLM-
 255 selected features. This process creates a targeted training and evaluation environment for that spe-
 256 cific archetype. Simultaneously, the textual description of the archetype is injected as a direct hint
 257 into the LLM’s prompt to guide the creative code generation, for example: Please note that
 258 we are focusing on highly-constrained problems.
 259

260 **Guided Evolutionary Search.** Let $\mathcal{D} = \{(d_i, I_i)\}_{i=1}^p$ be the set of established pairs, where each
 261 d_i is an optimization direction and I_i is its corresponding data subset. For each pair $(d_i, I_i) \in \mathcal{D}$,
 262 we conduct an independent evolutionary search for each heuristic module. During this search, the
 263 textual hint from d_i directs the LLM to generate heuristics specialized for that direction, with their
 264 performance evaluated exclusively on the dedicated data subset I_i . After this guided search, the total
 265 set of generated heuristic ensembles, formed by the Cartesian product $\mathcal{H} = \mathcal{L}_{restart} \times \mathcal{L}_{rephase} \times$
 266 \mathcal{L}_{bump_var} , undergoes a pruning step where low-performing combinations are filtered out, resulting
 267 in a smaller, high-quality portfolio \mathcal{H}' .
 268

269 3.4 INSTANCE SPACE PARTITIONING VIA PERFORMANCE-BASED CLUSTERING

270 The Cartesian product of generated heuristics can yield a large number of candidate ensembles,
 271 posing a computational challenge for the partitioning stage. To manage this, we constrain the number
 272 of components per heuristic type (e.g., $k \leq 3$) and perform an aggressive pruning step. All generated
 273

270 ensembles are evaluated on a benchmark subset, and only a fixed number of the top-performing
 271 candidates are retained in the final portfolio, \mathcal{H}' .
 272

273 Given this pruned portfolio of high-quality ensembles \mathcal{H}' , the next stage is to understand which
 274 ensemble is best for which kind of instance. To achieve this, we partition our training set of instances,
 275 I_{train} , based on performance.

276 First, we evaluate every heuristic ensemble $h_i \in \mathcal{H}$ on every instance $j \in I_{train}$. Let $p(h_i, j)$ be the
 277 performance metric (e.g., PAR-2 score) of ensemble h_i on instance j .
 278

279 Next, for each instance j , we identify its optimal ensemble, $h^*(j)$, from our library:
 280

$$h^*(j) = \arg \min_{h_i \in \mathcal{H}} p(h_i, j) \quad (1)$$

282 This allows us to partition the instance set I_{train} into disjoint clusters, where each cluster C_i is
 283 associated with a single best-performing ensemble h_i :
 284

$$C_i = \{j \in I_{train} \mid h^*(j) = h_i\} \quad (2)$$

286 This process creates a direct mapping from a region in the instance feature space (represented by the
 287 instances in C_i) to an optimal solver configuration h_i . For each resulting cluster C_i , we compute its
 288 feature space centroid, \bar{v}_i , by averaging the feature vectors of all its member instances.
 289

290 3.5 ADAPTIVE HEURISTIC SELECTION FOR INSTANCES 291

292 With the partitioned instance space and associated optimal ensembles, we can now perform adaptive
 293 selection for any new, unseen test instance j_{new} . We first extract its feature vector $v(j_{new})$. We
 294 then calculate the distance (using normalized Euclidean distance) between $v(j_{new})$ and each cluster
 295 centroid \bar{v}_i . The heuristic ensemble h_k associated with the closest centroid \bar{v}_k is selected as the
 296 most suitable configuration for solving j_{new} . This allows the solver to dynamically adapt its strategy
 297 based on the data characteristics of the problem at hand.
 298

299 4 EXPERIMENTS

300 To evaluate the effectiveness of DASHCO, we conduct a series of experiments designed to assess its
 301 performance against baseline and state-of-the-art SAT solvers.
 302

304 4.1 EXPERIMENTAL SETUP

306 **Environment and Parameters.** All solvers are implemented in C++ and compiled with g++ 12.3.0.
 307 The LLM interaction and the evolutionary framework are managed in Python. Experiments were
 308 conducted on servers equipped with AMD Ryzen 9 5950X 16-core processors and 128GB of RAM.
 309 For all heuristic generation tasks, we utilized the GPT-4o model with a temperature of 0.8 to encour-
 310 age diverse outputs. Considering the computational expense of solving SAT problems and the scale
 311 of our benchmarks, the timeout for solving any single SAT instance was set to 1000 seconds.
 312

313 **Backbone Solver.** To ensure a methodologically consistent and comparable basis, we follow the
 314 precedent set by AutoSAT Sun et al. (2024) in selecting our backbone solver. Our framework is
 315 built upon *EasySAT*, a lightweight and modular CDCL solver. As established in prior work, this
 316 choice provides a clean and capable baseline that is well-suited for modification by LLMs, striking
 317 a practical balance between solver functionality and the token-context limitations inherent in current
 318 LLM-based code generation.
 319

320 **Datasets.** Our experimental design is structured to rigorously evaluate the generalization capabili-
 321 ties of DaSATHco. We construct a single, heterogeneous training set by combining 24 unclassified
 322 instances from the SAT Competitions of 2022 and 2023 with 10 instances from each of the *Coins-*
 323 *Grid*, *LangFord*, and *PRP* benchmarks. Our evaluation then proceeds in two distinct settings:
 324

- 325 • **In-Domain Generalization:** We test performance on held-out instances from the *Coins-*
 326 *Grid*, *LangFord*, and *PRP* families, which were partially represented in the training set.
 327

Table 1: Main performance comparison across In-Domain and Out-of-Domain datasets. For each solver, we report both the PAR-2 score (lower is better) and the number of solved instances (Solved, higher is better). The timeout is 1000s. Best results in each category are in bold.

		EasySAT		MiniSat		AutoSAT		DASHCO		Kissat	
Dataset	# Inst.	PAR-2	#S	PAR-2	#S	PAR-2	#S	PAR-2	#S	PAR-2	#S
<i>In-Domain Benchmarks</i>											
CoinsGrid	52	1757.5	7	1985.1	4	1711.1	9	1399.5	16	1404.2	16
LangFord	64	1915.9	4	2000.0	0	1850.2	8	1750.9	12	1610.1	14
PRP	144	1970.4	3	1935.8	9	1843.8	15	1739.6	20	1570.0	38
<i>Out-of-Domain Benchmarks</i>											
CNP	50	594.9	38	1150.2	22	614.4	39	550.7	41	270.7	44
Zamkeller	48	830.5	30	1950.6	5	764.8	32	611.4	35	24.8	48
KnightTour	22	1733.4	3	1900.7	2	1684.6	4	1369.0	7	1638.5	4

- **Out-of-Domain Generalization:** To evaluate performance on entirely novel problem structures, we use test sets from the *CNP* (*Chromatic Number of the Plane*), *Zamkeller*, and *KnightTour* families, which were completely excluded from the training process.

Baselines. We compare the performance of DASHCO against a set of representative baselines:

- **EasySAT**: The lightweight, modular CDCL solver that serves as the direct backbone for our modifications. Its performance represents the starting point before any LLM-based optimization.
- **AutoSAT** Sun et al. (2024): A state-of-the-art framework representing the prior paradigm of dataset-specific optimization. To ensure a fair comparison of generalization capabilities, we adapt its methodology to our experimental setting. Instead of running its search process on each individual problem family, we run AutoSAT on the same single, heterogeneous training set used by DASHCO. This evaluates its ability to find a single "average-best" configuration for a diverse set of problems. Consequently, the performance reported here is not directly comparable to its original publication, where it was optimized on specialized datasets, and may be lower.
- **MiniSat** Sorensson & Een (2005): A classic and highly-optimized CDCL solver, which serves as a robust and widely-recognized traditional baseline.
- **Kissat** Biere et al. (2024): A state-of-the-art, highly-engineered CDCL solver that represents the pinnacle of modern manual heuristic design and frequently wins SAT competitions. It serves as a top-tier benchmark for performance. We deploy the 4.0.0 version of kissat which submitted to the SAT Competition 2024 and won 3 gold medals.

Metrics. We evaluate solver performance using two standard metrics in the SAT competition: the number of solved instances within the timeout, and the Penalized Average Runtime with a factor of 2 (PAR-2) score. The PAR-2 score is the average runtime across a set of instances, but with a heavy penalty for any instance that is not solved within the 1000s timeout. Specifically, unsolved instances are assigned a runtime of twice the timeout (2000s).

4.2 PERFORMANCE COMPARISON

We evaluate the performance of DASHCO against the baselines on both in-domain and out-of-domain datasets, with detailed results presented in Table 1. The findings clearly demonstrate the effectiveness of our data-aware, portfolio-based approach and offer critical insights when compared against different design paradigms.

Overall, DASHCO consistently and significantly outperforms the EasySAT, MiniSat, and AutoSAT baselines across both evaluation settings. The comparison with AutoSAT is particularly illuminating, as it confirms that for a diverse problem set, our adaptive selection from a specialized portfolio is more effective than relying on a single, average-best configuration produced by prior paradigms.

378 The comparison with Kissat highlights the core strengths of our framework. As expected, Kissat,
 379 a top-tier solver, shows exceptional performance on many benchmarks, particularly on well-known
 380 families like Zamkeller. However, the goal of our work is not to surpass a manually-honed solver
 381 on every problem, but to demonstrate a more generalizable and automated design paradigm. The
 382 results strongly support this goal. Most notably, on the out-of-domain KnightTour dataset, DASHCO
 383 decisively outperforms all other solvers, including Kissat. This result is a powerful validation of our
 384 central thesis: for novel problem structures, a static set of highly-tuned heuristics can be suboptimal,
 385 whereas DASHCO’s adaptive mechanism can dynamically select a more suitable configuration from
 386 its generated portfolio. Furthermore, on the in-domain CoinsGrid benchmark, DASHCO achieves
 387 performance on par with Kissat, demonstrating that our automated framework can generate and
 388 select configurations that are competitive with the state-of-the-art.
 389

390 These results lead to a clear conclusion: while highly-engineered solvers like Kissat represent
 391 the peak of performance for specific problem distributions, our automated, data-aware framework
 392 presents a robust and promising path towards building solvers that can generalize more effectively
 393 across a wide and unpredictable landscape of SAT instances.

394 4.3 ABLATION STUDY

395 To understand the contribution of the key components of our framework, we conduct an ablation
 396 study on two representative datasets: CoinsGrid for in-domain generalization, and Zamkeller for
 397 out-of-domain (OOD) generalization. We compare our full DASHCO model against several ablated
 398 variants:
 399

- 400 • **w/o Data-Aware Generation:** The portfolio is generated without the guidance of Problem
 401 Archetypes. The adaptive selection is then applied to this non-specialized portfolio.
- 402 • **Random Selection:** For each instance, we randomly select an ensemble from the fully
 403 generated, specialized portfolio.
- 404 • **Single Best Selection:** We select the single ensemble that performs best on average across
 405 the entire training set and apply it to all instances, simulating a non-adaptive paradigm.
- 406 • **Oracle:** This represents the theoretical performance upper bound of our portfolio, where
 407 for each test instance, we assume an oracle perfectly selects the best-performing ensemble
 408 from our generated portfolio.

410 The results are presented in Table 2. The full DASHCO model achieves the best performance among
 411 all practical variants. The importance of our data-aware generation is evident when comparing the
 412 full model to the ‘w/o Data-Aware Generation’ variant, which shows a clear performance drop.

413 The analysis of the selection mechanism reveals a crucial insight. The *Single Best Selection* ap-
 414 proach performs significantly worse than the adaptive DASHCO, particularly on the OOD dataset
 415 Zamkeller, highlighting the failure of a non-adaptive strategy to generalize. Interestingly, the ‘Or-
 416 acle’ results demonstrate that there is still considerable room for improvement in the selection mech-
 417 anism. The gap between DASHCO and the Oracle suggests that while our nearest-centroid selector is
 418 effective, more sophisticated selection models could unlock even greater performance from the gen-
 419 erated portfolio, representing a promising avenue for future work. Nonetheless, both the data-aware
 420 portfolio generation and the dynamic selection mechanism are shown to be critical to DASHCO’s
 421 robust performance.

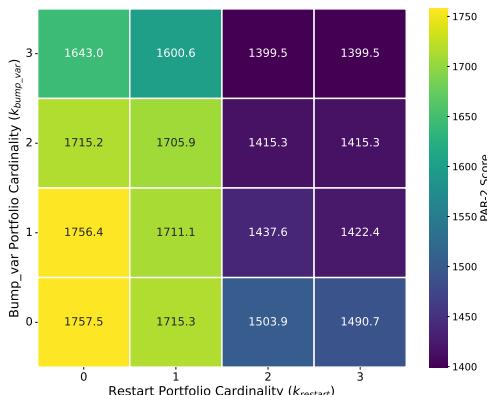
423 4.4 ANALYSIS OF PORTFOLIO SCALE AND DIVERSITY

424 The size of the heuristic portfolio is a critical hyperparameter in our framework. We focused our sen-
 425 sitivity analysis on the two most influential modules: restart and bump_var, fixing the rephase com-
 426 ponent to its default implementation. We evaluated the performance of DASHCO on the CoinsGrid
 427 dataset while varying the portfolio cardinalities for these two heuristics ($k_{restart}$ and k_{bump_var}),
 428 with k ranging from 0 to 3.

429 Figure 2 presents a heatmap of the results, which reveals several key insights. First, a significant
 430 performance improvement is observed when moving from a single heuristic ($k = 1$ for either dimen-
 431 sion) to combinations of multiple diverse heuristics ($k \geq 2$). For instance, the performance of the

432
 433
 434
 435
 436 Table 2: Ablation study on a representative
 437 In-Domain (CoinsGrid) and Out-of-Domain
 438 (Zamkeller) dataset. We report the PAR-2 score
 439 (lower is better) and the number of solved in-
 440 stances (#S, higher is better).

Model Variant	CoinsGrid		Zamkeller	
	PAR-2	#S	PAR-2	#S
DASHCO	1399.5	16	611.4	35
w/o Data-Aware	1671.0	10	722.4	33
Random	1740.2	8	780.5	29
Single Best	1503.9	12	764.8	32
Oracle	1157.8	19	467.6	41



441
 442
 443
 444
 445
 446
 447
 448
 449
 450
 451
 452
 453
 454 portfolio with $(k_{restart} = 2, k_{bump_var} = 2)$ is substantially better than portfolios with only one varied component, such as $(k_{restart} = 2, k_{bump_var} = 1)$. This demonstrates the powerful synergistic effects that arise from combining a diverse set of specialized heuristics, validating the core premise of our portfolio-based approach. Second, the heatmap shows that simply increasing the number of heuristics does not guarantee monotonic improvement, highlighting the complex interplay between them. The optimal performance is achieved with the portfolio of size $(k_{restart} = 2, k_{bump_var} = 3)$ or $(k_{restart} = 3, k_{bump_var} = 3)$. This illustrates that while a richer portfolio is generally beneficial, the quality and compatibility of the added heuristics are crucial. The results confirm that a diverse, multi-component portfolio is essential for achieving top performance.

463 4.5 OVERHEAD ANALYSIS

464 We analyze the computational overhead of DASHCO in terms of its one-time offline costs and per-
 465 instance online costs. The offline cost involves generating and compiling the heuristic portfolio, a
 466 process made manageable by the lightweight EasySAT backbone. The per-instance online overhead
 467 for solving a new instance is negligible, consisting of two fast operations: feature extraction and
 468 adaptive selection. The calculation of the feature vector is computationally inexpensive and can be
 469 pre-computed for known benchmarks. Subsequently, our adaptive selection via a nearest-centroid
 470 search is extremely efficient, consistently taking less than a second. This minimal online cost makes
 471 DASHCO highly practical. Future work could explore the trade-off between this efficient selector
 472 and more sophisticated models, such as LLMs.

475 5 CONCLUSION

476 In this work, we introduced DASHCO, a novel framework that addresses the critical generalizability
 477 limitations in automated SAT solver design by shifting from dataset-specific optimization to a scal-
 478 able “train-once, adapt-broadly” paradigm. Our methodology leverages a Large Language Model,
 479 guided by Problem Archetypes, to generate a diverse portfolio of specialized heuristic ensembles
 480 and then learns an adaptive mechanism to select the best configuration for new instances. Exper-
 481 iments confirm that this approach not only improves performance but, more importantly, exhibits
 482 robust out-of-domain generalization. This validates that learning a mapping from instance features
 483 to a generated portfolio of solvers is a more effective and practical paradigm for the automated de-
 484 sign of complex, configurable systems like SAT solvers. Future work could focus on automating the
 485 discovery of these archetypes and extending this paradigm to other domains.

486 REPRODUCIBILITY STATEMENT
487488 To ensure the reproducibility of our work, we commit to making our source code publicly available
489 upon acceptance of this paper. All SAT instances used for training and evaluation are from publicly
490 available SAT Competition benchmarks and previous work. We will provide detailed lists and in-
491 structions for obtaining them in our repository. More details of our implementation can be found in
492 the Appendix.494 REFERENCES
495496 Josh Achiam, Steven Adler, Sandhini Agarwal, Lama Ahmad, Ilge Akkaya, Florencia Leoni Ale-
497 man, Diogo Almeida, Janko Altenschmidt, Sam Altman, Shyamal Anadkat, et al. Gpt-4 technical
498 report. *arXiv preprint arXiv:2303.08774*, 2023.499 Carlos Ansótegui, Maria Luisa Bonet, and Jordi Levy. On the structure of industrial sat instances. In
500 *International Conference on Principles and Practice of Constraint Programming*, pp. 127–141.
501 Springer, 2009.502 503 Gilles Audemard and Laurent Simon. Predicting learnt clauses quality in modern sat solvers. In
504 *IJCAI*, volume 9, pp. 399–404, 2009.505 506 Gilles Audemard and Laurent Simon. Glucose 2.1: Aggressive, but reactive, clause database man-
507 agement, dynamic restarts (system description). In *Pragmatics of SAT 2012 (POS'12)*, 2012.508 509 Armin Biere, Marijn Heule, Hans van Maaren, and Toby Walsh. Conflict-driven clause learning
510 sat solvers. *Handbook of Satisfiability, Frontiers in Artificial Intelligence and Applications*, pp.
131–153, 2009.511 512 Armin Biere, Tobias Faller, Katalin Fazekas, Mathias Fleury, Nils Froleyks, and Florian Pollitt.
513 CaDiCaL, Gimsatul, IsaSAT and Kissat entering the SAT Competition 2024. In Marijn Heule,
514 Markus Iser, Matti Järvisalo, and Martin Suda (eds.), *Proc. of SAT Competition 2024 – Solver,
515 Benchmark and Proof Checker Descriptions*, volume B-2024-1 of *Department of Computer Sci-
516 ence Report Series B*, pp. 8–10. University of Helsinki, 2024.517 Edmund K Burke, Michel Gendreau, Matthew Hyde, Graham Kendall, Gabriela Ochoa, Ender
518 Özcan, and Rong Qu. Hyper-heuristics: A survey of the state of the art. *Journal of the Op-
519 erational Research Society*, 64(12):1695–1724, 2013.520 521 Gheorghe Comanici, Eric Bieber, Mike Schaeckermann, Ice Pasupat, Noveen Sachdeva, Inderjit
522 Dhillon, Marcel Blistein, Ori Ram, Dan Zhang, Evan Rosen, et al. Gemini 2.5: Pushing the
523 frontier with advanced reasoning, multimodality, long context, and next generation agentic capa-
524 bilities. *arXiv preprint arXiv:2507.06261*, 2025.525 526 Stephen A Cook. The complexity of theorem-proving procedures. In *Logic, automata, and compu-
527 tational complexity: The works of Stephen A. Cook*, pp. 143–152. 2023.528 529 James M Crawford and Andrew B Baker. Experimental results on the application of satisfiability
algorithms to scheduling problems. In *AAAI*, volume 2, pp. 1092–1097, 1994.530 531 William R Harris, Sriram Sankaranarayanan, Franjo Ivančić, and Aarti Gupta. Program analysis via
532 satisfiability modulo path programs. In *Proceedings of the 37th annual ACM SIGPLAN-SIGACT
symposium on Principles of programming languages*, pp. 71–82, 2010.533 534 Robert G Jeroslow and Jinchang Wang. Solving propositional satisfiability problems. *Annals of
535 mathematics and Artificial Intelligence*, 1(1):167–187, 1990.536 537 Richard Kaye. Minesweeper is np-complete. *The Mathematical Intelligencer*, 22(2):9–15, 2000.538 539 Jia Hui Liang, Vijay Ganesh, Pascal Poupart, and Krzysztof Czarnecki. Learning rate based branch-
ing heuristic for sat solvers. In *International Conference on Theory and Applications of Satisfia-
bility Testing*, pp. 123–140. Springer, 2016.

540 Jia Hui Liang, Chanseok Oh, Minu Mathew, Ciza Thomas, Chunxiao Li, and Vijay Ganesh. Machine
 541 learning-based restart policy for cdcl sat solvers. In *International Conference on Theory and*
 542 *Applications of Satisfiability Testing*, pp. 94–110. Springer, 2018.

543 Fei Liu, Xiali Tong, Mingxuan Yuan, Xi Lin, Fu Luo, Zhenkun Wang, Zhichao Lu, and Qingfu
 544 Zhang. Evolution of heuristics: Towards efficient automatic algorithm design using large language
 545 model. *arXiv preprint arXiv:2401.02051*, 2024.

546 Michael Luby, Alistair Sinclair, and David Zuckerman. Optimal speedup of las vegas algorithms.
 547 *Information Processing Letters*, 47(4):173–180, 1993.

548 Mukul R Prasad, Armin Biere, and Aarti Gupta. A survey of recent advances in sat-based formal
 549 verification. *International Journal on Software Tools for Technology Transfer*, 7(2):156–173,
 550 2005.

551 Jussi Rintanen. Planning as satisfiability: Heuristics. *Artificial intelligence*, 193:45–86, 2012.

552 Bernardino Romera-Paredes, Mohammadamin Barekatain, Alexander Novikov, Matej Balog,
 553 M Pawan Kumar, Emilien Dupont, Francisco JR Ruiz, Jordan S Ellenberg, Pengming Wang,
 554 Omar Fawzi, et al. Mathematical discoveries from program search with large language models.
 555 *Nature*, 625(7995):468–475, 2024.

556 Mate Soos, Karsten Nohl, and Claude Castelluccia. Extending sat solvers to cryptographic problems.
 557 In *International Conference on Theory and Applications of Satisfiability Testing*, pp. 244–257.
 558 Springer, 2009.

559 Niklas Sorensson and Niklas Een. Minisat v1. 13-a sat solver with conflict-clause minimization.
 560 *SAT*, 2005(53):1–2, 2005.

561 Yiwen Sun, Furong Ye, Xianyin Zhang, Shiyu Huang, Bingzhen Zhang, Ke Wei, and Shaowei
 562 Cai. Autosat: Automatically optimize sat solvers via large language models. *arXiv preprint*
 563 *arXiv:2402.10705*, 2024.

564 Niki van Stein and Thomas Bäck. Llamea: A large language model evolutionary algorithm for
 565 automatically generating metaheuristics. *IEEE Transactions on Evolutionary Computation*, 2024.

566 Lin Xu, Frank Hutter, Holger H Hoos, and Kevin Leyton-Brown. Satzilla: Portfolio-based algorithm
 567 selection for sat. *Journal of Artificial Intelligence Research*, 32:565–606, 2008.

568 Shunyu Yao, Fei Liu, Xi Lin, Zhichao Lu, Zhenkun Wang, and Qingfu Zhang. Multi-objective
 569 evolution of heuristic using large language model. In *Proceedings of the AAAI Conference on*
 570 *Artificial Intelligence*, volume 39, pp. 27144–27152, 2025.

571 Haoran Ye, Jiarui Wang, Zhiguang Cao, Federico Berto, Chuanbo Hua, Haeyeon Kim, Jinkyoo Park,
 572 and Guojie Song. Reevo: Large language models as hyper-heuristics with reflective evolution.
 573 *Advances in neural information processing systems*, 37:43571–43608, 2024.

582 A DETAILED SAT INSTANCE FEATURES

583 Our 37-dimensional feature vector is constructed from a combination of 33 features adapted from
 584 the SATzilla template and four new lightweight additions. The complete list is detailed below.

585 • Problem Size Features

- 586 – num_variables: The number of variables, denoted as v .
- 587 – num_clauses: The number of clauses, denoted as c .
- 588 – var_clause_ratio: The ratio of clauses to variables (c/v).

589 • Variable-Clause Graph Features

- 590 – Variable Nodes Degree (Frequency): avg_var_frequency, var_degree_cv,
 591 var_degree_min, max_var_frequency, var_frequency_entropy.

594 – Clause Nodes Degree (Length): `avg_clause_length`, `clause_degree_cv`,
 595 `min_clause_length`, `max_clause_length`, `clause_length_entropy`.
 596

- 597 • **Variable Graph Features**
 - 598 – Nodes Degree Statistics: `var_graph_degree_mean`, `var_graph_degree_cv`,
 599 `var_graph_degree_min`, `var_graph_degree_max`.

600 • **Balance and Polarity Features (Distributional)**

- 601 – Per-Clause Literal Ratio: `clause_pos_neg_ratio_mean`,
 602 `clause_pos_neg_ratio_cv`, `clause_pos_neg_ratio_entropy`.
- 603 – Per-Variable Occurrence Ratio: `var_pos_neg_ratio_mean`,
 604 `var_pos_neg_ratio_min`,
 605 `var_pos_neg_ratio_max`, `var_pos_neg_ratio_entropy`.
- 606 – Clause Type Fraction: `binary_clause_fraction`,
 607 `ternary_clause_fraction`.

608 • **Proximity to Horn Formula Features**

- 609 – `horn_clause_fraction`: The fraction of clauses that are Horn clauses.
- 610 – Horn Variable Occurrences: `horn_var_occurrence_mean`,
 611 `horn_var_occurrence_cv`,
 612 `horn_var_occurrence_min`,
 613 `horn_var_occurrence_max`, `horn_var_occurrence_entropy`.

614 • **Additional High-Level Features (Global)**

- 615 – `positive_literal_ratio`: The global ratio of positive literals to total literals.
- 616 – `balanced_var_ratio`: The proportion of variables with an equal number of positive and negative occurrences.
- 617 – `polarity_bias`: A measure of the overall tendency towards positive or negative literals.
- 618 – `constraint_density`: A measure of problem constraint level.

B ALGORITHM OF DASHCO

Algorithm 1 DASHCO Framework Overview

625

626

627 1: **Input:** Training instance set I_{train} , test instance j_{new} .

628 2: **Output:** Solved result for j_{new} . ▷ Stage 1: Data-Aware Heuristic Evolution

629 3: Define a set of Problem Archetypes $D = \{d_1, \dots, d_p\}$.

630 4: **for** each archetype $d_i \in D$ **do**

631 5: Create data subset $I_i \subseteq I_{train}$ by filtering.

632 6: **end for**

633 7: $\mathcal{H} \leftarrow \mathcal{L}_{restart} \times \mathcal{L}_{rephase} \times \mathcal{L}_{bump_var}$

634 8: $\mathcal{H}' \leftarrow \text{Prune}(\mathcal{H}, I_{train})$ by removing ensembles with performance below a predefined threshold. ▷ Stage 2: Offline Instance Space Partitioning

635

636 9: **for** each instance $j \in I_{train}$ **do**

637 10: $h^*(j) \leftarrow \arg \min_{h_k \in \mathcal{H}'} \text{performance}(h_k, j)$.

638 11: **end for**

639 12: **for** each unique h_k that is optimal for some instance **do**

640 13: $C_k \leftarrow \{j \in I_{train} | h^*(j) = h_k\}$.

641 14: $\bar{v}_k \leftarrow \frac{1}{|C_k|} \sum_{j \in C_k} v(j)$.

642 15: **end for** ▷ Stage 3: Online Adaptive Selection

643 16: $v_{new} \leftarrow v(j_{new})$.

644 17: $k \leftarrow \arg \min_i \text{distance}(v_{new}, \bar{v}_i)$.

645 18: $h_{selected} \leftarrow h_k$.

646 19: Solve j_{new} using the solver configured with $h_{selected}$.

648 **C DATASET CHARACTERISTICS**
649650 Our experimental design aims to evaluate the generalization capability of DASHCO. We construct a
651 single, heterogeneous **training set** composed of instances from multiple sources, including unclas-
652 sified problems from the SAT Competitions of 2022 and 2023, as well as the SCPC family. Our
653 evaluation is then divided into two settings:
654655

- 656 • **In-Domain Generalization:** We test on unseen instances from families that were partially
657 represented in the training set. These benchmarks include *CoinsGrid*, which originates
658 from a puzzle about arranging coins on a grid under specific constraints; *LangFord*, a com-
659 binatorial challenge of arranging paired numbers such that each pair of number k is sepa-
660 rated by exactly k other items; and *PRP* (*Profitable Robust Production*), which models an
661 industrial task of finding a robust production plan under uncertainty. We select 10 instances
662 from those datasets to compose our training set.

663 - 664 • **Out-of-Domain (OOD) Generalization:** To evaluate true generalization capabilities, we
665 use entire problem families that were completely held out during the heuristic evolution
666 process. These OOD test sets include *CNP* (*Chromatic Number of the Plane*), a classic
667 graph coloring problem; *Zamkeller*, a complex permutation problem concerning subse-
668 quences; and *KnightTour*, which seeks a path for a knight to visit every square on a chess-
669 board exactly once.

670 Table 3: Statistical characteristics of the evaluation datasets.
671672

Dataset	# Inst.	Variables (Mean \pm Std)	Clauses (Mean \pm Std)
<i>In-Domain Benchmarks</i>			
CoinsGrid	52	530807 ± 513663	3825868 ± 3701594
LangFord	64	312492 ± 213284	2734786 ± 1972624
PRP	144	499206 ± 324889	3337426 ± 2175367
<i>Out-of-Domain Benchmarks</i>			
CNP	50	9890 ± 11139	86724 ± 77379
Zamkeller	48	21435 ± 19119	265218 ± 283330
KnightTour	22	135288 ± 191062	5742107 ± 9872215

683 **D DEFINITION OF PAR-2 METRIC**
684685 The PAR-2 score is calculated as:
686

687
$$\text{PAR-2} = \frac{1}{N} \sum_{i=1}^N t'_i$$

688

689 where N is the number of instances, and t'_i for instance i is its actual runtime if solved, or 2000s if
690 unsolved. A lower PAR-2 score is better, as it indicates a solver is both fast and robust. Crucially,
691 the PAR-2 score also serves as the primary fitness metric that guides the LLM-driven evolutionary
692 search for better heuristics.
693694 **E EXAMPLE PROMPT TEMPLATE**
695696 Our data-aware guidance mechanism is designed to be agnostic to the specific underlying LLM-
697 based heuristic search algorithm. It can be integrated as a modular hint into various existing frame-
698 works, such as AutoSAT or ReEvo. The following provides an example of how our data-aware hint
699 can be incorporated into an *Advisor* prompt, using a structure similar to that of AutoSAT. The key
700 addition is the highlighted text, which provides the LLM with the specific Problem Archetype it
701 should optimize for.

702
703**Example Advisor Prompt Template**704
705
706
707
708

You are a SAT solver researcher trying to write the `{} task {}` to help SAT solver escape from local optimum. Your goal is to write a `{} task {}` for the SAT solver that will help it restart the search and escape from local optimum, after reading and understanding the `<key_code>` of SAT solver below. **Please note, the SAT problem instances we are targeting have the following characteristics: `{} problem_archetype_description {}`.**

Your answer must follow the following JSON format:

```

710  {
711      "description": "Provide a ...",
712      "modification_direction": ["some possible directions..."]
713  }
714
715 Tips:

```

1. You must traverse all possible positions of `<key_code>` if you want to modify the `{} task {}`.
2. You need to give us some advice to modify the `{} task {}`. e.g. some potential directions to change the heuristics.
3. Notice that, you can only change `{} task {}`.
4. `{} other_tips {}`

key_code of SAT solver is:

```

724     """
725     {{ origin_key_code }}
726     """

```

Take a deep breath and think it step by step. Then respond strictly in JSON format!

F HYPERPARAMETER SETTINGS

Table 4 lists the main hyperparameters used across our experiments to ensure reproducibility.

Table 4: Main hyperparameters used in our experiments.

Hyperparameter	Value
<i>General Experimental Settings</i>	
Solver Timeout	1000s
Random Seed	42
<i>LLM and Evolutionary Search</i>	
LLM Model	GPT-4o
Temperature	0.8
Generations	3
Population Size	2
<i>DASHCO Framework Settings</i>	
Number of Problem Archetypes	3
Heuristics per Type (k)	3
Pruned Portfolio Size	9

G USAGE OF LLMs

We take advantage of LLMs to improve the writing of this paper.

750
751
752
753
754
755