DISCRETE NEURAL ALGORITHMIC REASONING

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ABSTRACT

Neural algorithmic reasoning aims to capture computations with neural networks via learning the models to imitate the execution of classic algorithms. While common architectures are expressive enough to contain the correct model in the weights space, current neural reasoners are struggling to generalize well on out-of-distribution data. On the other hand, classic computations are not affected by distributional shifts as they can be described as transitions between discrete computational states. In this work, we propose to force neural reasoners to maintain the execution trajectory as a combination of finite predefined states. To achieve that, we separate discrete and continuous data flows and describe the interaction between them. Trained with supervision on the algorithm's state transitions, such models are able to perfectly align with the original algorithm. To show this, we evaluate our approach on multiple algorithmic problems and get perfect test scores both in single-task and multitask setups. Moreover, the proposed architectural choice allows us to prove the correctness of the learned algorithms for any test data.

1 INTRODUCTION

Learning to capture algorithmic dependencies in data and to perform algorithmic-like computations
with neural networks are core problems in machine learning, studied for a long time using various approaches (Roni Khardon, 1994; Graves et al., 2014; Zaremba & Sutskever, 2014; Reed & De Freitas, 2015; Kaiser & Sutskever, 2015; Veličković et al., 2020b).

Neural algorithmic reasoning (Veličković & Blundell, 2021) is a research area focusing on building
models capable of executing classic algorithms. Relying on strong theoretical guarantees of algorithms to work correctly on any input of any size and distribution, this setting provides unlimited
challenges for out-of-distribution generalization of neural networks. Prior work explored this setup
using the CLRS-30 benchmark (Veličković et al., 2022), which covers classic algorithms from the
Introduction to Algorithms textbook (Cormen et al., 2009) and uses graphs as a universal tool to
encode data of various types. Importantly, CLRS-30 also provides the decomposition of classic algorithms into subroutines and simple transitions between consecutive execution steps, called hints,
which can be used during training in various forms.

The core idea of the CLRS-30 benchmark is to understand how neural reasoners generalize well beyond the training distribution, namely on larger graphs. Classic algorithms possess strong gen-eralization due to the guarantee that correct execution steps never encounter 'out-of-distribution' states, as all state transitions are predefined by the algorithm. In contrast, when encountering in-puts from distributions that significantly differ from the train data, neural networks are usually not capable of robustly maintaining internal calculations in the desired domain. Consequently, due to the complex and diverse nature of all possible data that neural reasoners can be tested on, the gen-eralization performance of such models can vary depending on particular test distribution (Mahdavi et al., 2023).

Given that, it is becoming important to interpret internal computations of neural reasoners to find errors or to prove the correctness of the learned algorithms (Georgiev et al., 2021).

Interpretation methods have been actively developing recently due to various real-world applications
 of neural networks and the need to debug and maintain systems based on them. Especially, the
 Transformer architecture (Vaswani et al., 2017) demonstrates state-of-the-art performance in natural
 language processing and other modalities, representing a field for the development of interpretability

methods (Elhage et al., 2021; Weiss et al., 2021; Zhou et al., 2024; Lindner et al., 2024). Based on active research on a computational model behind the transformer architecture, recent works propose a way to learn models that are fully interpretable by design (Friedman et al., 2023).

We found the ability to design models that are interpretable in a simple and formalized way to be crucial for neural algorithmic reasoning as it is naturally related to the goal of learning to perform computations with neural networks.

In this paper, we propose to force neural reasoners to follow the execution trajectory as a combination of finite predefined states, which is important for both generalization ability and interpretability of neural reasoners. To achieve that, we start with an attention-based neural network and describe three building blocks to enhance its generalization abilities: feature discretization, hard attention and separating discrete and continuous data flows. In short, all mentioned blocks are connected:

- State discretization does not allow the model to use complex and redundant dependencies in data;
- Hard attention is needed to ensure that attention weights will not be annealed for larger graphs. Also, hard attention limits the set of possible messages that each node can receive;
- Separating discrete and continuous flows is needed to ensure that state discretization does not lose information about continuous data.

Then, we build fully discrete neural reasoners for different algorithmic tasks and demonstrate their
ability to perfectly mimic ground-truth algorithm execution. As a result, we achieve perfect test
scores on the multiple algorithmic tasks with guarantees of correctness on any test data. Moreover,
we demonstrate that a single network is capable of executing all covered algorithms in a multitask
manner with perfect generalization too.

In summary, we consider the proposed blocks as a crucial component for robust and interpretable
 neural reasoners and demonstrate that trained with hint supervision, discretized models perfectly
 capture the dynamic of the underlying algorithms and do not suffer from distributional shifts.

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2 BACKGROUND

084 2.1 ALGORITHMIC REASONING

Performing algorithmic-like computations usually requires the execution of sequential steps and the number of such steps depends on the input size. To imitate such computations, neural networks are expected to be based on some form of recurrent unit, which can be applied to a particular problem instance several times (Kaiser & Sutskever, 2015; Zaremba & Sutskever, 2014; Vinyals et al., 2015; Veličković et al., 2020b).

The CLRS Algorithmic Reasoning Benchmark (CLRS-30) (Veličković et al., 2022) defines a general paradigm of algorithmic modeling based on Graph Neural Networks (GNNs), as graphs can naturally represent different input types and manipulations over such inputs. Also, GNNs are proven to be well-suited for neural execution (Xu et al., 2020; Dudzik & Veličković, 2022).

The CLRS-30 benchmark covers different algorithms over various domains (arrays, strings, graphs) 096 and formulates them as algorithms over graphs. Also, CLRS-30 proposes to utilize the decom-097 position of the algorithmic trajectory execution into simple logical steps, called hints. Using this 098 decomposition is expected to better align the model to desired computations and prevent it from utilizing hidden non-generalizable dependencies of a particular train set. Prior work demonstrates a 099 wide variety of additional inductive biases for models towards generalizing computations, including 100 different forms of hint usage (Veličković et al., 2022; Bevilacqua et al., 2023), biases from standard 101 data structures (Jürß et al., 2024; Jain et al., 2023), knowledge transfer and multitasking (Xhonneux 102 et al., 2021; Ibarz et al., 2022; Numeroso et al., 2023), etc. Also, recent studies demonstrate sev-103 eral benefits of learning neural reasoners end-to-end without any hints at all (Mahdavi et al., 2023; 104 Rodionov & Prokhorenkova, 2023). 105

The recently proposed SALSA-CLRS benchmark (Minder et al., 2023) enables a more thorough
 OOD evaluation compared to CLRS-30 with increased test sizes (up to 100-fold train-to-test scaling, compared to 4-fold for CLRS-30) and diverse test distributions. Despite significant gains in the

108 performance of neural reasoners in recent work, current models still struggle to generalize to out-109 of-distribution (OOD) test data (Mahdavi et al., 2023; Georgiev et al., 2023; Minder et al., 2023). 110 While de Luca & Fountoulakis (2024) prove by construction the ability of the transformer-based 111 neural reasoners to perfectly simulate graph algorithms (with minor limitations occurring from the 112 finite precision), it is still unclear if generalizable and interpretable models can be obtained via learning. Importantly, the issues of OOD generalization are induced not only by the challenges 113 of capturing the algorithmic dependencies in the data but also by the need to carefully operate 114 with continuous inputs. For example, investigating the simplest scenario of learning to emulate 115 the addition of real numbers, Klindt (2023) demonstrates the failure of some models to exactly 116 imitate the desired computations due to the nature of gradient-based optimization. This limitation 117 can significantly affect the performance of neural reasoners on adversarial examples and larger input 118 instances when small errors can be accumulated.

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2.2 TRANSFORMER INTERPRETABILITY AND COMPUTATION MODEL

Transformer (Vaswani et al., 2017) is a neural network architecture for processing sequential data.
The input to the transformer is a sequence of tokens from a discrete vocabulary. The input layer maps each token to a high-dimensional embedding and aggregates it with the positional encoding. The key components of each layer are attention blocks and MLP with residual connections. Providing a detailed description of mechanisms learned by transformer models (Elhage et al., 2021) is of great interest due to their widespread applications.

RASP (Weiss et al., 2021) is a programming language proposed as a high-level formalization of the computational model behind transformers. The main primitives of RASP are elementwise sequence functions, *select* and *aggregate* operations, which conceptually relate to computations performed by different blocks of the model. Later, Lindner et al. (2024) presented Tracr, a compiler for converting RASP programs to the weights of the transformer model, which can be useful for evaluating interpretability methods.

While RASP might have limited expressibility, it supports arbitrary complex continuous functions which in theory can be represented by transformer architecture, but are difficult to learn. Also, RASP is designed to formalize computations over sequences of fixed length. Motivated by that, Zhou et al. (2024) proposed RASP-L, a restricted version of RASP, which aims to formalize the computations that are easy to learn with transformers in a size-generalized way. The authors also conjecture that the length-generalization of transformers on algorithmic problems is related to the 'simplicity' of solving these problems in RASP-L language.

Another recent work (Friedman et al., 2023) describes Transformer Programs: constrained transformers that can be trained using gradient-based optimization and then automatically converted into a discrete, human-readable program. Built on RASP, transformer programs are not designed to be size-invariant.

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3 DISCRETE NEURAL ALGORITHMIC REASONING

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3.1 ENCODE-PROCESS-DECODE PARADIGM

¹⁵⁰Our work follows the encode-process-decode paradigm (Hamrick et al., 2018), which is usually ¹⁵²employed for step-by-step neural execution.

All input data is represented as a graph G with an adjacency matrix A and node and edge features that are first mapped with a simple linear encoder to high-dimensional vectors of size h. Let us denote node features at a time step t $(1 \le t \le T)$ as $X^t = (x_1^t, \ldots, x_n^t)$ and edge features as $E_t = (e_1^t, \ldots, e_m^t)$. Then, the processor, usually a single-layer GNN, recurrently updates these features, producing node and edge features for the next step:

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$$X^{t+1}, E^{t+1} = \operatorname{Processor}(X^t, E^t, A).$$

The processor network can operate on the original graph defined by the task (for graph problems) or
 on the fully connected graph. For the latter option, the information about the original graph can be
 encoded into the edge features.

The number of processor steps T can be defined automatically by the processor or externally (e.g., as the number of steps of the original algorithm). After the last step, the node and edge features are mapped with another linear layer, called the decoder, to the output predictions of the model.

If the model is trained with hint supervision, the changes of node and edge features at each step are
 expected to be related to the original algorithm execution. In this sense, the processor network is
 aimed to mimic the algorithm's execution in the latent space.

169 170 3.2 DISCRETE NEURAL ALGORITHMIC REASONERS

In this section, we describe the constraints for the processor that allow us to achieve a fully inter pretable neural reasoner. We start with Transformer Convolution (Shi et al., 2020) with a single attention head.

As mentioned above, at each computation step t $(1 \le t \le T)$, the processor takes the highdimensional embedding vectors for node and edge features as inputs and then outputs the representations for the next execution step.

Each node feature vector x_i is projected into $query(Q_i)$, $key(K_i)$, and $value(V_i)$ vectors via learnable parameter matrices W_Q , W_K , and W_V , respectively. Edge features e_{ij} are projected into $key(K_{ij})$ vector with a matrix W_K^E . Then, for each directed edge from node j to node i in the graph G, we compute the attention coefficient

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186 187 188 where $\langle a, b \rangle$ denotes the dot product. Then, each node *i* normalizes all attention coefficients across its neighbors with the softmax function and temperature τ and receives the aggregated message:

 $\alpha_{ij} = \frac{\langle Q_j, K_i + K_{ij} \rangle}{\sqrt{h}},$

$$\hat{\alpha}_{ij} = \frac{\exp(\alpha_{ij}/\tau)}{\sum_{k \in \mathcal{N}(i)} \exp(\alpha_{ik}/\tau)}, \quad M_i = \sum_{k \in \mathcal{N}(i)} \hat{\alpha}_{ik} V_k, \tag{1}$$

where $\mathcal{N}(i)$ denotes the set of all incoming neighbors of node *i* and M_i is the message sent to the *i*-th node.

For undirected graphs, we consider two separate edges in each direction. Also, for each node, we consider a self-loop connecting the node to itself. For multi-head attention, each head l separately computes the messages M_i^l which are then concatenated.

Similar to Transformer Programs, we enforce attention to be hard attention. We found this property important not only for interpretability but also for size generalization, as hard attention allows us to overcome the annealing of the attention weights for arbitrarily large graphs and strictly limits the set of messages that each vertex can receive.

After message computation, node and edge features are updated depending on the current values and sent messages using feed-forward MLP blocks:

$$\hat{x}_i^{t+1} = \text{FFN}_{nodes}([x_i^t, M_i^t]),$$
$$\hat{e}_{ij}^{t+1} = \text{FFN}_{edges}([e_{ij}^t, \hat{\alpha}_{ji}V_i, \hat{\alpha}_{ij}V_j]).$$

We also enforce all node and edge features to be from a fixed finite set, which we call *states*. We ensure such property by adding discrete bottlenecks at the end of the processor block:

We implement discretization by projecting the features to the vectors of size k which we force to be one-hot using annealing Gumbel-Softmax (Jang et al., 2017) during training and the argmax at the inference.

212 3.3 CONTINUOUS INPUTS

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Clearly, most of the algorithmic problems operate with continuous or unbounded inputs (e.g., weights on edges). Usually, all input data is encoded into node and edge features and the processor operates over the resulting vectors. The proposed discretization of such vectors would lead to

the loss of information necessary for performing correct execution steps. One possible option to operate with such inputs (we will call them *scalars*, meaning both continuous or size-dependent integer inputs, such as node indexes) is Neural Execution Engines (Yan et al., 2020), which allows one to operate with bit-wise representations of integer and (in theory) real numbers. Such representations are bounded by design, but fully discrete and interpretable.

221 We propose another option: to maintain scalar inputs (denote them by S) separately from the node 222 and edge features and use them only as edge priorities s_{ij} in the attention block. If scalars are related 223 to the nodes, we assign them to edges depending on the scalar of the sender or receiver node. Now, 224 we can consider the hard attention block as a selector which for each node selects the best edge 225 based on an ordered pair of 'states priority' (attention weights described above, which depend only 226 on states of the corresponding nodes and edges) and s_{ij} . We note that this selector is related to the theoretical primitive *select_best* from RASP. We implement this simply by augmenting key vectors 227 K_{ij} of each edge with the indicator if the given edge has the "best" (min or max) scalar among the 228 other edges to node j. Thus, scalars affect only the attention weights, not the messages and the node 229 states. 230

For multiple different scalar inputs (e.g., weighted edges and node indexes), we use multi-head attention, where each head operates with separate scalars.

Given that, the interface of the proposed processor can be described as

$$X^{t+1}, E^{t+1} = \operatorname{Processor}(X^t, E^t, A, S),$$

where X^t, X^{t+1} and E^t, E^{t+1} are from the fixed sets. State sets are independent of the execution step t and the input graph (including scalar inputs S).

239 3.4 MANIPULATIONS OVER CONTINUOUS INPUTS

The proposed selector offers a read-only inter-241 face to scalar inputs, which is not expressive 242 enough for most of the algorithms. However, 243 we note that the algorithms can be described as 244 discrete manipulations over input data. For ex-245 ample, the Dijkstra algorithm (Dijkstra, 1959) 246 takes edge weights as inputs and uses them 247 to find the shortest path distances. Computed 248 distances can affect the consequent execution 249 steps. We note that such distances can be de-250 scribed as the sum of the weights of the edges that form the shortest path to the given node. In 251 other words, the produced scalars depend only 252 on input scalars and discrete execution states. 253

To avoid the challenges of learning continuous updates with high precision (Klindt, 2023),
we propose to learn discrete manipulations with
scalars. The updated scalars can then be used
with the described selector in the next steps.

In our experiments, we use a scalar updater
capable of incrementing, moving, and adding
scalars depending on discrete node/edge states:



Figure 1: An illustration of the proposed separation between discrete and continuous data flows. Scalars can only affect the attention weights (Green) and can be modified with actions via ScalarUpdate (Blue).

$$s_i^{t+1} = \operatorname{inc}(x_i^t) + \operatorname{keep}(x_i^t) \cdot s_i^t + \sum_{j \in \mathcal{N}(i)} \operatorname{push}(e_{ji}^t) \cdot s_{ji}^t,$$
$$s_{ij}^{t+1} = \operatorname{inc}(e_{ij}^t) + \operatorname{keep}(e_{ij}^t) \cdot s_{ij}^t + \operatorname{push}(x_i^t) \cdot s_i^t,$$

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where s_i are node-related scalars, s_{ij} are edge-related scalars, and inc, keep, push are 0-1 functions representing if scalar in each node/edge should be incremented, kept, or pushed to any of its neighbors. We implement these functions as simple linear projections of node/edge features with consecutive discretization.

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Finally, our proposed method (see Figure 1) can be described as:

$$\begin{aligned} X^{t+1}, E^{t+1} &= \operatorname{Processor}(X^t, E^t, A, S^t), \\ S^{t+1} &= \operatorname{ScalarUpdate}(X^{t+1}, E^{t+1}, A, S^t). \end{aligned}$$

The proposed neural reasoners are fully discrete and can be interpreted by design. Moreover, the proposed selector block guarantees the predicted behavior of the message passing for any graph size, as it compares discrete state importance and uses continuous scalars only to break ties between equally important states.

4 EXPERIMENTS

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In this section, we perform experiments to evaluate how the proposed discretization affects the performance of neural reasoners on diverse algorithmic reasoning tasks. Our main questions are:

- 1. Can the proposed discrete neural reasoners capture the desired algorithmic dynamics with hint supervision?
- 2. How does discretization affect OOD and size-generalization performance of neural reasoners?
- 3. Is the proposed model capable of multi-task learning?

Also, we are interested if discrete neural reasoners can be learned without hints and how they tend to utilize the given amount of node and edge states to solve the problem. We discuss no-hint experiments separately in Section 6.

4.1 DATASETS

We perform our experiments on the problems from the recently proposed SALSA-CLRS benchmark (Minder et al., 2023), namely BFS, DFS, Prim, Dijkstra, Maximum Independent Set (MIS), and Eccentricity. We believe that the proposed method is not limited by the covered problems, but we leave the implementation of the required data flows (e.g., edge-based reasoning (Ibarz et al., 2022), graph-level hints, interactions between different scalars) for future work.

The train dataset of SALSA-CLRS consists of random graphs with at most 16 nodes sampled from the Erdös-Rényi (ER) distribution with parameter p chosen to be as low as possible while graphs remain connected with high probability. The test set consists of sparse graphs of sizes from 16 to 1600 nodes.

We slightly modify the hints from the benchmark without conceptual changes (e.g., we have modified the hints for the DFS problem to remove graph-level hints). Discrete states are fully described by the non-scalar hints and scalars are exactly the hints of the *scalar* type (we refer to Veličković et al. (2022) for the details on hint design).

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4.2 **BASELINES AND EVALUATION**

310 We compare the performance of our proposed discrete model with two baseline sparse models, GIN 311 (Xu et al., 2019) and Pointer Graph Network (Veličković et al., 2020a). We report both node-level 312 and graph-level metrics for the baselines and our model. Also, we compare our model with Triplet-313 GMPNN (Ibarz et al., 2022) and two recent approaches, namely Hint-ReLIC (Bevilacqua et al., 314 2023) and G-ForgetNet (Bohde et al., 2024), which demonstrate state-of-the-art performance of 315 hint-based neural algorithmic reasoning. However, as these methods are evaluated on the CLRS-30 316 benchmark and their code is not yet publicly available, we can only compare them on the corre-317 sponding tasks (BFS, DFS, Dijkstra, Prim) and CLRS-30 test data, namely ER graphs with p = 0.5318 of size 64. Note that this test data is more dense than that of SALSA-CLRS, meaning shorter roll-319 outs for given tasks. Also, only node-level metrics have been reported for these methods.

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- 321 4.3 MODEL DETAILS
- For our experiments, we use the model described in Section 3. We use one attention head for each scalar value. The number of processor steps is defined externally as the length of the ground truth

algorithm trajectory, which is consistent with the prior work. We use one architecture (except the task-depended encoders/decoders), including the ScalarUpdate module for all the problems.

We recall that neural reasoners can operate either on the base graph (which is defined by the problem) or on more dense graphs with the original graph encoded into edge features. SALSA-CLRS proposes to enhance the size-generalization abilities of neural reasoners with increased test sizes (up to 100fold train-to-test scaling), so we use the base graph for message-passing, similar to the SALSA-CLRS baselines. We also add a virtual node that communicates with all nodes of the graph.

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4.4 TRAINING DETAILS AND HYPERPARAMETERS

We train each model using the Adam optimizer, learning rate $\eta = 0.001$, using teacher forcing, batch size of 32 graphs with 1000 optimization steps, and evaluate the resulting model. We anneal softmax temperatures in the discrete bottlenecks (attention weights, ScalarUpdate operations) geometrically from 3.0 to 0.01, decreasing the temperature at each training step. We report all hyperparameters in the source code.¹

During training, we minimize the standard hints and output losses: scalar hints are optimized with MSE loss, and other types of hints are optimized with cross-entropy and categorized cross-entropy losses (Veličković et al., 2022; Ibarz et al., 2022). Note that we do not supervise any additional details in model behavior, e.g., selecting the most important neighbor in the attention block, the exact operations with scalars, etc.

344 For multitask experiments, we follow the setup 345 proposed by Ibarz et al. (2022) and train 346 a single processor with task-dependent en-347 coders/decoders to imitate all covered algo-348 rithms simultaneously. We make 10000 op-349 timization steps on the accumulated gradients across each task and keep all hyperparameters 350 the same as in the single task. 351

Our models are trained on a single A100 GPU,
requiring less than 1 hour for single-task and
5-6 hours for multitask training.

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4.5 RESULTS

358 We found learning with teacher forcing suitable for discrete neural reasoners, as discretiza-359 tion blocks allow us to perform the exact tran-360 sitions between the states. Trained with step-361 wise hint supervision, discrete neural reason-362 ers are able to perfectly align with the origi-363 nal algorithm and generalize on larger test data 364 without any performance loss. We report the 365 evaluation results in Tables 1 and 2. Also, our 366 multitask experiments show that the proposed 367 discrete models are capable of multitask learn-368 ing and demonstrate the perfect generalization 369 scores in a multitask manner too.

Recall that we have three key components of our contribution: feature discretization, hard attention, and separating discrete and continuous data flows. To evaluate the importance of each component for generalization capabilities of the proposed models, we conduct an ablation study, the details can be found in Appendix A. In

Table 1: Node \setminus graph level test scores for the
proposed discrete reasoner and the baselines on
SALSA-CLRS test data. Scores are averaged
across 5 different seeds, standard deviation is
omitted. For the eccentricity problem, only the
graph-level metric is applicable.

TASK	Size	GIN	PGN	DNAR (OURS)
BFS	16	98.8 \ 92.5	100. \ 100.	100. \ 100.
	80	95.3 \ 59.4	99.8 \ 88.1	100. \ 100.
	160	95.1 \ 37.8	99.6 \ 66.3	100. \ 100.
	800	86.9 \ 0.9	98.7 \ 0.2	100. \ 100.
	1600	$86.5 \setminus 0.0$	$98.5 \setminus 0.0$	100. \ 100.
DFS	16	$41.5 \setminus 0.0$	82.0 \ 19.9	100. \ 100.
	80	$30.4 \setminus 0.0$	$38.4 \setminus 0.0$	$100. \setminus 100.$
	160	$20.0 \setminus 0.0$	$26.9 \setminus 0.0$	$100. \setminus 100.$
	800	$19.5 \setminus 0.0$	$24.9 \setminus 0.0$	$100. \setminus 100.$
	1600	$17.8 \setminus 0.0$	$23.1 \setminus 0.0$	$100. \setminus 100.$
SP	16	$95.2 \setminus 49.8$	99.3 \ 89.5	100. \ 100.
	80	$62.4 \ 0.0$	94.2 \ 3.3	100. \ 100.
	160	$53.3 \setminus 0.0$	$92.0 \setminus 0.0$	$100. \setminus 100.$
	800	$40.4 \setminus 0.0$	$87.1 \setminus 0.0$	$100. \setminus 100.$
	1600	$36.9 \setminus 0.0$	$84.5 \setminus 0.0$	$100. \setminus 100.$
Prim	16	89.6 \ 29.7	96.4 \ 69.9	100. \ 100.
	80	$51.6 \ 0.0$	79.7 \ 0.0	100. \ 100.
	160	$49.5 \ 0.0$	75.6 \ 0.0	100. \ 100.
	800	$45.0 \setminus 0.0$	$69.5 \setminus 0.0$	$100. \setminus 100.$
	1600	$43.2 \setminus 0.0$	$66.8 \setminus 0.0$	100. \setminus 100.
MIS	16	79.9 \ 3.3	99.8 \ 98.6	100. \ 100.
	80	$79.9 \setminus 20.0$	$99.4 \setminus 88.9$	$100. \setminus 100.$
	160	$78.2 \setminus 0.0$	$99.4 \setminus 76.2$	$100. \setminus 100.$
	800	$83.4 \setminus 0.0$	$98.8 \setminus 18.0$	$100. \setminus 100.$
	1600	$79.2 \setminus 0.0$	$98.9 \setminus 5.2$	$100. \setminus 100.$
Ecc.	16	25.3	100.	100.
	80	23.8	100.	100.
	160	26.1	100.	100.
	800	17.1	100.	100.
	1600	16.0	83.0	100.

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¹https://anonymous.4open.science/r/F4CA/

Table 2: Node-level test scores for the proposed discrete reasoner and the baselines on CLRS-30 test
 data. Test graphs are of size 64. Scores are averaged across 5 different seeds.

TASK	TRIPLET-GMPNN	HINT-RELIC	G-ForgetNet	DNAR (OURS)
BFS	99.73 ± 0.0	99.00 ± 0.2	99.96 ± 0.0	$100. \pm 0.0$
DFS	47.79 ± 4.2	$100. \pm 0.0$	74.31 ± 5.0	$100.\pm0.0$
Dijkstra	96.05 ± 0.6	97.74 ± 0.5	99.14 ± 0.1	$100.\pm0.0$
MST-PRIM	86.39 ± 1.3	87.97 ± 2.9	95.19 ± 0.3	$100.\pm0.0$

short, our additional experiments demonstrate that the proposed processor without the discretization performs non-above the baselines level and that removing hard-attention or discrete *ScalarUpdate* module strictly limits the generalization capabilities of the proposed model.

5 INTERPE

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INTERPETABILITY AND TESTING

In addition to the empirical evaluation of the trained models on diverse test data, the discrete and size-independent design of the proposed models allows us to interpret and test them manually. The main idea is to show that the model will perform the exact discrete state transitions (including discrete operations with scalars) as the original algorithm.

First, we note that due to the hard attention, each node receives exactly one message. Also, the message depends only on the discrete states of the corresponding nodes and edges. Thereby, as each node and edge states after a single processor step depend only on the current states and received message, all the possible options can be directly enumerated and tested if all states change to the correct ones.

The only remaining part to fully interpret the whole model is the attention block. We note that 404 our implementation of the *select_best* selector (Section 3.3) does not necessarily produce the top-1 405 choice over the ordered pairs of 'states priority' and scalars s_{ii} as it simply augments key vectors 406 with indicators if the given edge has the best scalar among others. For example, it may happen 407 that for some state, the maximum attention weight is achieved for an edge without the indicator. 408 However, given the finite number of discrete states, we can manually check if the mentioned "best" 409 indicator increases the attention weight between every pair of states. Combined with hard attention, 410 this would imply that the attention block attends depending on the predefined states and uses scalar 411 priorities only to break ties. Please see the particular example with more detailed explanation for the 412 BFS problem in the Appendix D.

Given that, we can unit-test all possible state transitions and attention blocks. With full coverage of such tests, we can guarantee the correct execution of the desired algorithm for *any* test data. We tested our trained models from Section 4 manually verifying state transitions. As a result, we confirm that the attention block indeed operates as *select_best* selector, as the model actually uses these indicators to increase the attention weights. Thus, we can guarantee that for any graph size, the model will mirror the desired algorithm, which is correct for any test size.

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6 TOWARDS NO-HINT DISCRETE REASONERS

In this section, we discuss the challenges of training discrete reasoners without hints, which can be useful when tackling new algorithmic problems.

Training deep discretized models is known to be challenging: without hyperparameter search, discrete models are only slightly improved over the untrained models. Therefore, we focus only on the BFS algorithm, as it is well-aligned with the message-passing framework, has short roll-outs, and can be solved with small states count (note that for no-hint models node/edge states count is a hyperparameter due to the absence of the ground truth states trajectory).

We recall that the output of the BFS problem is the exploration tree pointing from each node to its parent. Each node chooses as a parent the neighbor from the previous distance layer with the smallest index.

Table 3: BFS node/graph level scores of the *best_no_hint_model* for different graph sizes.

	5	16	64
best_no_hint_model	99 / 86	94 / 34	79/

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We perform hyperparameter search over the training sizes (using ER graphs with p = 0.5 and $n \in [4, 16]$), discrete node states count (from 2 to 6 states), softmax temperature annealing schedules ([3, 0.01], [3, 0.1], [3, 1]). For each hyperparameter choice, we train 5 models with different seeds. We validate the resulting models on the graphs of size 16. The best resulting model is obtained with the training size 5 and 4 node states. The trained models never achieved the perfect validation scores, see Table 3 for the results.

Then, we select the best-performing models and try to analyze the mistakes of the resulting models 445 and reverse-engineer how they utilize the given states. First, we look at the node states after the last 446 step of the processor and note that the states correspond to the distances from the starting node. More 447 formally, we note that the model with four states uses the first state for the starting node, the second 448 state for its neighbors, the third state for nodes at distance two from the starting node, and the last 449 state for all other nodes and such states-based classification of distance has accuracy > 98% when 450 tested on 1000 random graphs with 16 nodes. Then, we note that for the nodes that are from the first 451 four distance layers from the starting node, the pointers are predicted with 100% accuracy and these 452 pointers are computed layer-by-layer as in the ground truth algorithm (we refer to Appendix B for 453 illustrations). The mistakes of the model are on the distance ≥ 4 from the starting node (we did not 454 reverse-engineer the specific logic of computations on larger distances).

We found this behavior well-aligned to the BFS algorithm and indicating the possibility of achieving the perfect validation score with enough states count. However, this algorithm does not generalize since it fails at distances larger than those encountered during training.

On the other hand, one can demonstrate that for a small enough state count (for BFS, it is two node states and two edge states) and diverse enough validation data, the perfect validation performance implies that the learned solution will generalize to any graph size.

Therefore, we highlight the need to achieve perfect validation performance with models that use as few states as possible, which corresponds to the minimum description length (MDL) theory (Myung, 2000; Rissanen, 2006) and is related to the notion of Kolmogorov complexity (Kolmogorov, 1963).

Finally, we note that for sequential problems, such as DFS, obtaining a good no-hint model can
be even more challenging and can require additional effort. One possible way to overcome this
limitation is to implement a curriculum learning setup.

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7 LIMITATIONS AND FUTURE WORK

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Limitations In this work, we propose a method to learn robust neural reasoners that demonstrate perfect generalization performance and are interpretable by design. In this section, we describe some limitations of our work and important directions for future research.

475 First, several proposed design choices strictly reduce the expressive power of the model. For exam-476 ple, due to the hard attention, the proposed model is unable to compute the mean value from all the 477 neighbors in a single message-passing step, which is trivial for attention-based models (note that this can be computed in several message-passing steps). Thus, the model in its current form is unable to 478 express transitions between hints for some algorithms from the CLRS-30 in a single processor step. 479 However, we believe that the expressivity of the proposed model can be enhanced with additional 480 architectural modifications (e.g., edge-based reasoning (Ibarz et al., 2022), global states, interactions 481 between different scalars) that can be combined with the proposed discretization ideas. 482

483 Second, while we report the perfect scores for the covered tasks, we cannot guarantee that the train 484 ing will converge to the correct model for any initialization/training data distributions. However, we
 485 empirically found the proposed method to be quite robust to various architecture/training hyperparameter choices.

Future work Our method is based on the particular architectural choice and actively utilizes the attention mechanism. However, the graph deep learning field is rich in various architectures exploiting different inductive biases and computation flows. The proposed separation between discrete states and continuous inputs may apply to other models, however, any particular construction can require additional efforts.

Also, we provide only one example of the ScalarUpdate block. We believe that utilizing a general architecture (e.g., some form of discrete Neural Turing Machine (Graves et al., 2014; Gulcehre et al., 2016)) capable of executing a wider range of manipulation is of interest for future work.

With the development of neural reasoners and their ability to execute classic algorithms on abstract data, it is becoming more important to investigate how such models can be applicable in real-world scenarios according to the Neural Algorithmic Reasoning blueprint (Veličković & Blundell, 2021) and transfer their knowledge to high-dimensional noisy data with intrinsic algorithmic dependencies. While there are several examples of NAR-based models tackling real-world problems (Beurer-Kellner et al., 2022; Numeroso et al., 2023), there are no established benchmarks for extensive evaluation and comparison of different approaches.

Lastly, we leave for future work a deeper investigation of learning interpretable neural reasoners without hints, which we consider essential from both theoretical perspective and practical applications, e.g., combinatorial optimization.

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8 CONCLUSION

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In this paper, we force neural reasoners to maintain the execution trajectory as a combination of finite predefined states. To achieve that, we separate discrete and continuous data flows and describe the interaction between them. The obtained discrete reasoners are interpretable by design. Moreover, trained with hint supervision, such models perfectly capture the dynamic of the underlying algorithms and do not suffer from distributional shifts. We consider discretization of hidden representations as a crucial component for robust neural reasoners.

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648 A ABLATION STUDY

 Recall that we have three key components of our contribution: feature discretization, hard attention, and separating discrete and continuous data flows. In this section, we study the importance of these components for generalization capabilities of the proposed models.

Discrete bottlenecks First, we evaluate the model without all discrete bottlenecks: the result is a simple Transformer Convolution processor, which performs comparable to other baseline models, see Table 4.

Table 4: Node \setminus graph level test scores for our base model without all discrete bottlenecks. Scores are averaged across 5 different seeds, standard deviation is omitted.

Size	16	80	160	800	1600
BFS	99.9 \ 99.3	99.7 \ 88.2	99.5 \ 57.9	$98.4 \setminus 0.0$	$97.2 \setminus 0.0$
DFS	$79.2 \ 6.8$	$41.1 \setminus 0.0$	$28.1 \setminus 0.0$	$24.7 \setminus 0.0$	$21.9 \setminus 0.0$
SP	99.3 \ 88.7	$94.1 \setminus 12.4$	$90.3 \setminus 0.0$	$86.9 \setminus 0.0$	$82.4 \setminus 0.0$
Prim	95.1 \ 72.7	$82.6 \ 0.0$	$79.7 \setminus 0.0$	$68.1 \setminus 0.0$	$66.0 \setminus 0.0$
MIS	$99.8 \setminus 98.6$	$99.6 \setminus 86.1$	$99.2 \setminus 69.0$	$97.1 \setminus 11.9$	$96.3 \setminus 0.0$
Ecc	79.2	41.1	28.1	24.7	21.9

Hard attention To highlight the importance of the hard attention for strong size generalization, we train the proposed model but with the regular attention mechanism on the BFS task. The resulting model also demonstrates the perfect scores for the given test data. However, the standard test data (the Erdös-Renyi graphs with low edge probability) does not contain nodes with a large number of neighbors, while such nodes can be problematic due to the annealing of the attention weights. Thus, we additionally test the resulting model on the complete bipartite graphs $K_{2,n-2}$ for different n. For each n, we assign the starting node from the smaller component and test if the second node in this component correctly selects its parent in the BFS tree, see Figure 2.



Figure 2: Complete bipartite graphs $K_{2,n-2}$ used to evaluate the effect of the attention weights annealing. The black node is the starting node. The highlighted edge (Green) is the ground truth pointer from the bottom node to its parent from the BFS tree.

Our experiments demonstrate that the model without hard attention fails to predict the correct pointer for larger graphs due to the attention weight annealing, see Table 5. We note that the models from Section 4 are provably correct on any test data.

Table 5: Attention weights for the ground truth pointer (green pointer from Figure 2) for different graph sizes; (+)-) denotes if the correct pointer was predicted.

	-	-			
SIZE	16	80	160	800	1600
ATTENTION WEIGHT	0.97 (+)	0.86 (+)	0.76 (+)	0.38 (-)	0.24 (-)

Scalar updater As demonstrated in Klindt (2023), simple neural networks trained to sum two real numbers fail to learn the structure of the task and struggle to extrapolate beyond the training data distributions. In this section, we study how these limitations affect the overall performance of neural reasoners to highlight the importance of the proposed discrete manipulations with scalars.

First, we recall that usually all input data is encoded into node and edge features and the processor operates over the resulting vectors. Then, hints of type scalar are directly predicted from the

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702 node/edge features. To evaluate the effect of non-discrete ScalarUpdate modules, we simply re-703 place the proposed discrete ScalarUpdate module with a single-layer transformer convolution net-704 work, which inputs scalars and node/edge states and outputs the scalars for the next step, keeping 705 the remaining architecture the same as in the main experiments. We train the resulting model on 706 Dijkstra and MST problems.

707 Additionally, we evaluate the non-discrete *ScalarUpdate* module in a more straightforward setup. 708 Similarly to Klindt (2023), we train a 2-layer MLP to add two real numbers and use the resulting 709 model as a *ScalarUpdate* module for the Dijkstra algorithm. We manually use the learned addition 710 module when the node distances are updated (e.g., the distance of the node u is updated with the 711 sum of the distance of v and the edge (v, u) cost), and use the ground truth scalars for other scalar 712 updates. Our experiments demonstrate that the resulting model outperforms the baselines on the test size of 16 nodes, but does not generalize well on larger graphs, see the evaluation results in Table 6. 713

Table 6: Node \setminus graph level test scores for the proposed model with ScalarUpdater replaced by 715 a regular attention-based network trained to predict hints of type scalar. 'Addition only' means that 716 ScalarUpdater is replaced by a 2-layer MLP trained to predict the sum of two numbers (other 717 values are taken from the ground truth). 718

719	Size	16	80	160	800	1600
720	Dijkstra	99.3 \ 94.6	$60.7 \setminus 0.0$	$42.8 \setminus 0.0$	$19.0 \setminus 0.0$	$11.8 \setminus 0.0$
721	MST	99.8 \ 98.1	$98.2 \setminus 54.1$	$97.2 \setminus 28.1$	$95.5 \setminus 0.0$	91.73 \ 0.0
722	DIJKSTRA (ADDITION ONLY)	99.8 \ 96.6	95.3 \ 71.0	86.5 \ 46.3	41.6 \ 3.1	$22.2 \setminus 0.0$

724 We note that all the resulting models demonstrate perfect scores when evaluated with teacher-forced ground truth scalars. Thus, all state transitions are learned correctly and imperfect test scores are 725 fully described by the errors in manipulations with continuous values. 726

727 To summarize, small errors in manipulations with scalars (even restricted on the simplest addition 728 sub-task) strictly affect the overall performance of the model, highlighting the importance of the 729 proposed discrete manipulations with scalars.

В **STATE USAGE FOR NO-HINT MODELS**

733 In this section, we provide several illustrations of the node states and the dynamics of the pointer 734 prediction updates (Figures 3-5). Our analysis suggests that no-hint models with K states tend to 735 use states as distances from the starting node, with the distances $\geq K$ merged to the same state. 736 Also, pointer predictions for the first K BFS layers are correct and computed layer-by-layer as in 737 the ground truth algorithm. The mistakes of the model are at the later layers and some pointers at the 738 later layers are computed before that in the ground truth algorithm (Figure 5c). For the simplicity of 739 illustrations, we use the model with 3 discrete states.



753 Figure 3: Node states and the predicted pointers after the last processor step of the DNAR model 754 (with 3 states), trained without hints. Different colors represent different states. The green node is 755 the starting node.



Figure 4: Node states and the predicted pointers after the last processor step of the DNAR model (with 3 states), trained without hints. Different colors represent different states. The green node is the starting node.



(b) Predicted pointers after the first processor step (c) Predicted pointers after the second processor step (self-loops are omitted) (self-loops are omitted)

Figure 5: Node states and the dynamics of the pointer prediction updates of the DNAR model (with 3 states), trained without hints. Different colors represent different states. The green node is the starting node.

C EXTENDED ScalarUpdate MODULE

In this section, we investigate if the proposed *ScalarUpdate* module can be successfully extended to support more complex manipulations with scalars.

First, we note that simple manipulations with scalars cover a significant part of the classical algorithms. In this work, we use the minimum set of the required functions, but it can be directly extended by other functions. Importantly, as *ScalarUpdate* can be viewed as a separate module, we can separately check if it is possible to train it with any given set of predefined manipulations for any problem only with supervision on the results.

Let us formalize the problem: each input of the ScalarUpdate module can be described as an object with a discrete state s_i (from a fixed predefined set) and several scalar values (we consider two scalars x_i and y_i). Note that we omit the separation between nodes and edges and consider objects with several scalar values. For each discrete state s_i , there exists a ground truth update of scalars, e.g., $f(s, x, y) = x + \cos(y)$. The output of the scalar updater can be viewed as a sum:

$$ScalarUpdate(s, x, y) = \sum_{g \in \mathcal{OPS}} g(x, y) \cdot \operatorname{activate}_{g}(s)$$

where OPS is a predefined set of operations and $activate_g$ is a 0-1 function representing if a specific operation should be applied. Note that $activate_g$ depends only on a discrete state of the input.

For our additional experiment, we train several different *ScalarUpdate* modules with the extended set of operations:

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• $g_1(x,y) = x;$

• $g_2(x,y) = \cos(x);$

• $q_0(x,y) = 1;$

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- $g_3(x,y) = x \cdot y;$
- $g_4(x,y) = \operatorname{atan2}(x,y)$

to learn the following set of the ground truth updates simultaneously:

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833• $f_0(x, y) = x;$ 834• $f_1(x, y) = \cos(x);$ 835• $f_2(x, y) = \cos(x) + x \cdot y;$ 836• $f_3(x, y) = \operatorname{atan2}(x, y);$ 837• $f_4(x, y) = 1 + x + \operatorname{atan2}(x, y).$

In particular, we consider a set of 16 discrete states (numbered from 0 to 15 and sampled uniformly)
and the ground truth scalar update is derived from these states by taking the remainder of the division
by 5 (updates count).

The learnable parameters of the *ScalarUpdate* are state's embeddings and linear projections for
each indicator. We train *ScalarUpdate* to minimize the MSE loss between the ground truth and
predicted outputs with 5000 optimization steps. Additionally, we train a non-discrete scalar updater
(2-layer MLP), similar to our ablation experiments. We refer to the source code for the experiment
details.

Inspired by Klindt (2023), we generate training scalars X and Y from Uniform[0.5, 1.0] and generate test set by sampling scalars from the Uniform[0., 0.5] distribution.

We report the evaluation results in Table 7. The proposed discrete ScalarUpdate module successfully learned the correct operations for updates $f_1, ..., f_4$ for all seeds and 3 times out of 5 for f_0 (note that the model was trained to predict different manipulations for different states simultaneously). For unsuccessful runs, when f_0 was not learned correctly, the learned operation for f_0 was g_2 for some states (i.e. $x \cdot y$ instead of x), which can be explained by optimization challenges as the distribution of y is close to 1.

Our experiment demonstrates that the proposed *ScalarUpdate* module can be extended to support a wider range of manipulations with scalars. We note that this complicates the optimization problem of selecting the correct operations/operands from the operations results (e.g., such decomposition might not be unique).

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D INTERPETABILITY AND TESTING DETAILS

In this section, we provide additional details on interpretability and testing of the proposed discrete reasoners.

Table 7: MSE for train/test distributions for the discrete and non-discrete *ScalarUpdate* modules
 and different operations.

		f_0	f_1	f_2	f_3	f_4
	discrete	0.01 / 0.1	0. / 0.	0 / 0	0 / 0	0 / 0
As an example, consider the BFS algorithm. First, recall the pseudocode of the algorithm: Starting_node \leftarrow visited All_other_nodes \leftarrow not_visited for steep in range(T) do for node U in a graph do if U is visited on previous steps then continue end if if U has a neighbor P that visited on previous steps then $U \leftarrow visited$ on this step U select the smallest-indexed such neighbor P as parent: Edge $(U, P) \leftarrow pointer$ Self-loop $(U, U) \leftarrow not_pointer$ end for end for $end for down the step t+1 (denoted by U_{t+1}) is the funof U_t and V_t, where V is the node that sends the message to U on step t:U_{t+1} = StateUpdate(U_t, message_from_V_t)How does the node U select a node that will send a message to it? For any node V connected theindicator if each node has the smallest (or largest) scalar among all neighbors of U with thediscrete state as V. Then, the node U selects the node V with the largest attention scores.In our (slightly simplified) case, the attention scores only depend on the tuples (U_{state}, Vindicator i_1_u_has_he_smalles_index) and there are only 8 such tuples. We can directly upup these attention scores and verify the required invariants, e.g.,Attention(not_visited, visited, smallest) > Attention(not_visited, *any_other*),$ which would imply that the not_visited node will receive the message from the smallest-indiv visited = StateUpdate(not_visited, message_from_visited) $visited = StateUpdate(not_visited, message_from_visited)$ $visited = StateUpdate(not_visited, message_from_visited)$ The main idea is that due to the finite states count and discrete manipulations with scalars, ther only finite amounts of such checks that can cover all possible state transitions and all of them si be evaluated only once.	non-discrete	5.7 <i>e</i> -6 / 0.03	5.1 <i>e</i> -6 / 0.001	1.1e-5 / 0.007	1.0e-5 / 0.08	2.0e-5 / 0.03
As an example, consider the BFS algorithm. First, recall the pseudocode of the algorithm: Starting_node \leftarrow visited All_other_nodes \leftarrow not_visited for step in range(T) do for node U in a graph do if U is visited on previous steps then continue end if if U has a neighbor P that visited on previous steps then $U \leftarrow visited$ on this step U select the smallest-indexed such neighbor P as parent: Edge (U, P) \leftarrow pointer Self-loop (U, U) \leftarrow not_pointer end if end for return a BFS tree described by pointers Now let us describe how we can verify that the trained DNAR model will perfectly imitate algorithm for any test data. First, we note that for each node U, the node state on the step $t+1$ (denoted by U_{t+1}) is the fun of U_t and V_t , where V is the node that sends the message to U on step t : $U_{t+1} = StateUpdate(U_t, message_from_V_t)$ How does the node uselect a node that will send a message to i? For any node V connected the node U computes attention scores depending on discrete states of V of each node and a disindicator if each node has the smallest (or largest) scalar among all neighbors of U with the discrete state as V. Then, the node U selects the node V with the largest attention score. In our (slightly simplified) case, the attention scores only depend on the tuples (U_{state} , V indicator_if_uhas_the_smallest_index) and there are only 8 such tuples. We can directly put these attention scores and verify the required invariants, e.g., Attention(not_visited, visited, smallest) > Attention(not_visited, *any_other*), which would imply that the not_visited node will receive the message from met smallest-ind visited = StateUpdate(not_visited, message_from_not_visited) $visited = StateUpdate(not_visited, message_from_not_visited)$ $visited = StateUpdate(not_visited, message_from_not_visited)$ The main idea is that due to the finite states count and discrete manipulations with scalars, ther only finite amounts of such checks that can cover all possible state tran						
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