Exploiting Data Sparsity in Secure Cross-Platform Social Recommendation

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Abstract

Social recommendation has shown promising improvements over traditional sys-1 tems since it leverages social correlation data as an additional input. Most existing 2 works assume that all data are available to the recommendation platform. However, З 4 in practice, user-item interaction data (e.g., rating) and user-user social data are usually generated by different platforms, both of which contain sensitive infor-5 mation. Therefore, How to perform secure and efficient social recommendation 6 across different platforms, where the data are highly-sparse in nature remains an 7 important challenge. In this work, we bring secure computation techniques into 8 social recommendation, and propose S³Rec, a sparsity-aware secure cross-platform 9 social recommendation framework. As a result, S³Rec can not only improve the 10 recommendation performance of the rating platform by incorporating the sparse 11 social data on the social platform, but also protect data privacy of both platforms. 12 Moreover, to further improve model training efficiency, we propose two secure 13 sparse matrix multiplication protocols based on homomorphic encryption and 14 private information retrieval. Our experiments on two benchmark datasets demon-15 strate that S³Rec improves the computation time and communication size of the 16 state-of-the-art model by about $40 \times$ and $423 \times$ in average, respectively. 17

18 **1** Introduction

The recent advances of social recommendation have achieved remarkable performances in recommendation tasks [10, 26]. Unlike traditional methods, social recommendation leverages user-item rating data (e.g. from Netflix) with user-user social data (e.g. from Facebook) to facilitate model training. The cross-platform nature, the high sparsity and sensitivity of recommendation/social data make social recommendation hard-to-deploy in the real world [5]. Therefore, the main problem is,

24 How to perform secure and efficient social recommendation across different platforms, where the 25 data are highly-sparse in nature?

26 Specifically, we focus on the problem of collaborative social recommendation in the two-party model, 27 where one party (denoted as P_0) holds user-item rating data, and the other party (denoted as P_1) 28 holds user-user social data. We also assume that the adversaries are semi-honest and non-colluding, 29 which is commonly used in literature [8]. That is, the adversary will not deviate from the pre-defined 30 protocol, but will try to learn as much information as possible from its received messages.

Choices of privacy enhancing techniques. Currently, many anonymization techniques have been used in publishing recommendation data, such as *k-anonymity* and *differential privacy* [9]. On the other hand, cryptographic methods like *secure multiparty computation* (MPC) [9] and *homomorphic encryption* (HE) have been proposed to enable calculation on the protected data. Since k-anonymity has been demonstrated risky in practice (e.g., the re-identification attack on Netflix Prize dataset [20]), and differential privacy introduces random noises to the dataset which eventually affects model accuracy, we consider they are not the ideal choice for our framework. Instead, we choose a combination of cryptographic tools (i.e., MPC and HE, but mainly MPC) which allows multiple parties to jointly compute a function depending on their private inputs while providing security guarantees.

Choice of social recommendation model. In literature, many social recommendation models have been proposed [7, 16, 25] using matrix factorization or neural networks. Existing MPC-based neural network protocols [19, 27] usually suffer from accuracy loss and inefficiency due to their approximation of non-linear operations. Especially for the case of social recommendation, training data could exceed to millions, and this makes NN-based model a less ideal choice. Therefore, we choose the classic social recommendation model, Soreg [16], as a typical example, and present how to build a secure and efficient version of Soreg under cross-domain social recommendation scenario.

Dealing with sparse data in secure machine learning. One important property of social recommendation data is its high sparsity. Take LibraryThing dataset [29] for example, its social matrix density is less than 0.02%. Recently, Schoppmann et al. introduced the ROOM framework [24] for secure computation over sparse data. However, their solution only works on column-sparse or row-sparse data, and in addition, it requires secure matrix multiplication protocol (for instance, based on Beaver's multiplication triple). Different from their work, in this paper, we propose a PIR-based matrix multiplication which does not reply on pre-generated correlated randomness.

Our framework. In this paper, we propose S^3 Rec, a sparsity-aware secure cross-platform social recommendation framework. Starting with the classic Soreg model, we observe that the training process of Soreg involves two types of calculation terms: (1) the *rating term* which could be calculated by P_0 locally, and (2) the *social term* which needs to be calculated by P_0 and P_1 collaboratively. Therefore, the key to S^3 Rec is designing secure and efficient protocols for calculating the social term.

To begin with, we first let both parties perform local calculation. Then both parties invoke a secure 60 social term calculation protocol and let P_0 finally receive the plaintext social term, and update the 61 model accordingly. In this way, the security of our protocol relies significantly on the secure social 62 term calculation protocol (for simplicity, we refer this protocol as the 'ST-MPC' protocol), and we 63 propose a secure instantiation and prove its security. Similarly, the efficiency of S³Rec relies heavily 64 on the performance of ST-MPC, and at the core, it relies on the efficiency of a matrix multiplication 65 protocol. The naïve secure matrix multiplication protocol is traditionally evaluated through Beaver's 66 triples [3], and has $O(km^2)$ asymptotic communication complexity, where k is the dimension of 67 latent factors and m is the number of users. To improve the communication efficiency, we propose 68 two secure sparse matrix multiplication protocols for ST-MPC, based on two sparsity settings: (1) 69 insensitive sparsity, which is a weaker variant of matrix multiplication where we assume both parties 70 know the locations of non-zero values in the sparse matrix, and (2) sensitive sparsity, which is 71 also a weaker variant of matrix multiplication, but stronger than (1), and we assume 'only' the 72 number of zeros is public. Nevertheless, we present secure constructions for MatrixMul in both 73 cases by leveraging two cryptography primitives called Private Information Retrieval (PIR) [1] and 74 Homomorphic Encryption (HE) [22]. PIR can hide the locations of the non-zero values in the sparse 75 matrix while HE enables additions and multiplications on ciphertexts. To this end, we drop the 76 77 communication complexity of secure MatrixMul to O(km) for the insensitive sparsity case and to 78 $O(\alpha km)$ for the sensitive sparsity case, where α denotes the density of user social matrix.

⁷⁹ **Summary of our experimental results.** We conduct experiments on two popularly used dataset, ⁸⁰ i.e., Epinions [17] and LibraryThing [29]. The results demonstrate that (1) S³Rec achieves the same ⁸¹ performance as existing social recommendation models, and (2) S³Rec improves the computation ⁸² time and communication size of the state-of-the-art (SeSoRec) by about $40 \times$ and $423 \times$ in average.

Contributions. We summarize our main contributions below: (1) We propose S³Rec, a privacypreserving cross-platform social recommendation framework, which relies on a general protocol for calculating the social term securely; (2) We propose two secure sparse matrix multiplication protocols based on different sparsity visibility, i.e., insensitive sparsity and sensitive sparsity. We prove that both protocols are secure under semi-honest adversaries; and (3) We empirically evaluate the performance of S³Rec on benchmark datasets.

2 Preliminaries and Recommendation Model

90 2.1 Preliminaries

91 We begin with the syntax., we use [n] to denote the set $\{1, ..., n\}$, and |x| to denote the length of x.

- In terms of MPC, we denote a secret shared value of x in \mathbb{Z}_N as [x], where N is a positive integer.
- Also, we let $[\![x]\!]_0$ denote P_0 's share, and $[\![x]\!]_1$ denote P_1 's share, where $[\![x]\!] = [\![x]\!]_0 + [\![x]\!]_1 \in \mathbb{Z}_N$.

Multi-Party Computation (MPC). MPC is a 94 cryptographic tool which enables multiple par-95 ties (say, n parties) to jointly compute a function 96 $f(x_1, ..., x_n)$, where x_i is *i*-th party's private in-97 put. MPC protocols ensure that, at the end of 98 the protocol, parties eventually learn nothing but 99 their own input and the function output. MPC 100 has been widely-used in secure machine learn-101 ing systems such as PrivColl [28] and CrypT-102 Flow [13], most of which support a wide range 103 of linear (e.g. addition, multiplication) and non-104 linear functions (e.g. equality test, comparison). 105 Here, we present three popular MPC protocols, 106 which we will use later in our protocol, 107

$$\begin{array}{ll} {}_{108} & - \ \llbracket z \rrbracket \leftarrow \mathsf{Add}(\llbracket x \rrbracket, \llbracket y \rrbracket) \colon P_0 \ \mathsf{lets} \ \llbracket z \rrbracket_0 = \llbracket x \rrbracket_0 + \\ {}_{109} & \ \llbracket y \rrbracket_0, \ \mathsf{and} \ P_1 \ \mathsf{lets} \ \llbracket z \rrbracket_1 = \llbracket x \rrbracket_1 + \llbracket y \rrbracket_1; \end{array}$$

110 $- [[z]] \leftarrow Mul([[x]], [[y]])$: Multiplication is evaluated via Beaver's Triples [3]. $\llbracket \mathbf{Z} \rrbracket \leftarrow \mathsf{Matrix}\mathsf{Mul}(\mathbf{X},\mathbf{Y})$ 1: P_0, P_1 jointly generate km^2 Beaver's triples Offline Completed 2: $\forall x_{i,j} \in \mathbf{X}, P_0 \text{ invokes } [x_{i,j}] \leftarrow \mathsf{Shr}(x_{i,j})$ 3: $\forall y_{i,j} \in \mathbf{Y}, P_1 \text{ invokes } [\![y_{i,j}]\!] \leftarrow \mathsf{Shr}(y_{i,j})$ foreach $i \in [k], j \in [m]$, let $[z_{i,j}] = 0$, 4: foreach $a \in [m], b \in [m],$ 5: 6: $\llbracket \mathsf{tmp} \rrbracket \leftarrow \mathsf{Mul}(\llbracket x_{i,a} \rrbracket, \llbracket y_{b,j} \rrbracket)$ $\llbracket z_{i,j} \rrbracket \leftarrow \mathsf{Add}(\llbracket \mathsf{tmp} \rrbracket, \llbracket z_{i,j} \rrbracket)$ 7: endfor 8: endfor 9: 10 : return **Z**

Figure 1: Matrix multiplication protocol.

Homomorphic Encryption (HE) scheme. HE is essentially a specific type of encryption scheme which allows manipulation on encrypted data. More specifically, HE involves a key pair (pk, sk), where the public key pk is used for encryption and the secret key sk is used for decryption. In this work, we use an additive HE scheme (i.e., Paillier [22]) which allows the following operations:

116 $-z \leftarrow \mathsf{Enc}_{\mathsf{pk}}(x) \oplus \mathsf{Enc}_{\mathsf{pk}}(y)$: addition between two ciphertexts, where $z = \mathsf{Enc}_{\mathsf{pk}}(x+y)$;

117 $-z \leftarrow \text{Enc}_{pk}(x) \otimes y$: multiplication between a ciphertext and a plaintext, where $z = \text{Enc}_{pk}(x \cdot y)$.

Private Information Retrieval (PIR). Now, we introduce single-server PIR [1]. In this setting, we 118 assume there is a server and a client, where the server holds a database $DB = \{d_1, ..., d_n\}$ with n 119 elements, and the client wants to retrieve DB_i while hiding the query index *i* from the server. Roughly, 120 a PIR protocol consists of a tuple of algorithm (PIR.Query, PIR.Response, PIR.Extract). First, the 121 client generates a query $q \leftarrow \mathsf{PIR}.\mathsf{Query}(i)$ from an index i, and then sends query q to the server. 122 The server then is able to generate a response $r \leftarrow \mathsf{PIR.Response}(\mathsf{DB},q)$ based on the query and 123 database DB, and returns r to the client. Finally, the client extracts the result from server's response 124 $\mathsf{DB}_i \leftarrow \mathsf{PIR}.\mathsf{Extract}(r).$ 125



Figure 2: An overview of Private Information Retrieval

126 2.2 Recommendation model

Recall that we assume there are two platforms, a rating platform P_0 , and a social platform P_1 . We assume P_0 holds a private rating matrix $\mathbf{R} \in \mathbb{R}^{m \times n}$, and P_1 holds a private user social matrix

Input: Private rating matrix R from platform P_0 , private user social matrix S from platform P_1 .					
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Figure 3: Our proposed $S^{3}Rec$ framework, where MatrixMul stands for secure matrix multiplication protocol, Add stands for secure add protocol, Rec stands for reconstruction protocol for secret sharing.

 $\mathbf{S} \in \mathbb{R}^{m \times m}$, where *n* and *m* denote the number of items and their common users, respectively. Also, we denote the user latent factor matrix as $\mathbf{U} \in \mathbb{R}^{k \times m}$ and item latent factor matrix as $\mathbf{V} \in \mathbb{R}^{k \times n}$, 129 130 where k is the dimension of latent factors. We further define an indication matrix $\mathbf{I} \in \mathbb{R}^{m \times n}$, where 131 $I_{i,j}$ denotes whether user *i* has rated item *j*. 132

Existing work [25] summarizes factorization based social recommendation models as the combination 133 of a "basic factorization model" and a "social information model". To date, different kinds of social 134 information models have been proposed [16, 12], and their common intuition is that users with social 135 relations tend to have similar preferences. In this work, we focus on the classic social recommendation 136 model, i.e., Soreg [16], which aims to learn U and V by minimizing the following objective function, 137

$$\sum_{i=1}^{m} \sum_{j=1}^{n} \frac{1}{2} I_{i,j} \left(r_{i,j} - \mathbf{u}_{*,i}^{T} \mathbf{v}_{*,j} \right)^{2} + \frac{\lambda}{2} \sum_{i=1}^{m} \|\mathbf{u}_{*,i}\|_{F}^{2} + \frac{\lambda}{2} \sum_{j=1}^{n} \|\mathbf{v}_{*,j}\|_{F}^{2} + \frac{\gamma}{2} \sum_{i=1}^{m} \sum_{f=1}^{m} s_{i,f} \|\mathbf{u}_{*,i} - \mathbf{u}_{*,f}\|_{F}^{2},$$
(1)

where the first term is the basic factorization model, the last term is the social information model, and 138 the middle two terms are regularizers, $\|\cdot\|_F^2$ is the Frobenius norm, λ and γ are hyper-parameters. If we denote $\mathbf{D} \in \mathbb{R}^{m \times m}$ as a diagonal matrix with diagonal element $d_b = \sum_{c=1}^m s_{b,c}$ and $\mathbf{E} \in \mathbb{R}^{m \times m}$ as a diagonal matrix with diagonal element $e_i = \sum_{b=1}^m s_{b,i}$. The gradients of \mathcal{L} in Eq. (1) with 139 140 141 respect to U and V are, 142

Rating term: computed by P_0 locally

$$\frac{\partial \mathcal{L}}{\partial \mathbf{U}} = \underbrace{-\mathbf{V}\left(\left(\mathbf{R} - \mathbf{U}^T \mathbf{V}\right)^T \circ \mathbf{I}\right) + \lambda \mathbf{U}}_{\text{Beting terms computed by } \mathcal{P}_0 \text{ and } \mathcal{P}_1 \text{ collaboratively}} + \underbrace{\frac{\gamma}{2} \mathbf{U}(\mathbf{D}^T + \mathbf{E}^T) - \gamma \mathbf{U}\mathbf{S}^T}_{\text{Social terms computed by } \mathcal{P}_0 \text{ and } \mathcal{P}_1 \text{ collaboratively}}, \quad (2)$$

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$$\frac{\partial \mathcal{L}}{\partial \mathbf{V}} = \underbrace{-\mathbf{U}\left(\left(\mathbf{R} - \mathbf{U}^T \mathbf{V}\right)^T \circ \mathbf{I}\right) + \lambda \mathbf{V}}_{\text{Rating term: computed by} P_0 \text{ locally}}.$$
(3)

3 Framework 144

We summarize our proposed $S^{3}Rec$ framework in Figure 3. To begin with, we assume that party 145 P_0 holds the rating matrix **R** and P_1 holds the social matrix **S**. At first, P_0 randomly initializes $\mathbf{U} \leftarrow \mathbb{R}^{k \times m}$ and $\mathbf{V} \leftarrow \mathbb{R}^{k \times n}$. Then, for each iteration (while the model dose not coverage), we let 146 147 P_0 and P_1 jointly evaluate the social term defined in Eq 2. P_0 then locally calculates the rating term 148 in Eq 2 and Eq 3, as well as $\partial \mathcal{L}/\partial \mathbf{U}$ and $\partial \mathcal{L}/\partial \mathbf{V}$. Party P_0 then locally updates U and V accordingly 149 and ends the iteration. 150

Communication efficiency. In our framework, the only communication between two parties occurs 151 in the ST-MPC protocol. Since we choose additive secret sharing, the Add protocol contains only 152 local computation, we claim that the communication efficiency of S³Rec significantly relies on 153

$\textbf{Protocol:} \ \llbracket \textbf{Z} \rrbracket \leftarrow MatrixMul(\textbf{X},\textbf{Y})$
1: P_0, P_1 jointly generate km^2 Beaver's triples
Offline Completed
2: $\forall (i,j) \in l_y, P_1$ pushes $y_{i,j}$ into $T_y(j)$
3: for $j \in [k]$ do
4: P_1 lets $T_y = \emptyset$
5: $\forall a \in [m], b \in [m], \text{ if } (i, a) \in l_y, P_0 \text{ pushes } x_{a,b} \text{ into } T_x(a)$
6: for $j \in [m]$ do
7: Both parties let $[\![z_{i,j}]\!] = 0$, then, for all values $v \in T_x(i), u \in T_y(j)$
8: $P_0 \text{ invokes } \llbracket v \rrbracket \leftarrow Shr(v), P_1 \text{ invokes } \llbracket v \rrbracket \leftarrow Shr(u)$
9: $\llbracket z_{i,j} \rrbracket = Add(Mul(\llbracket v \rrbracket, \llbracket u \rrbracket), \llbracket z_{i,j} \rrbracket)$
10: endfor
11: endfor

Figure 5: Dense-sparse MatrixMul(\mathbf{X}, \mathbf{Y}) with insensitive sparsity, that is, $\mathbf{X} \in \mathbb{R}^{k \times m}, \mathbf{Y} \in \mathbb{R}^{m \times m}$, and \mathbf{Y} 's location vector l_y is public.

the efficiency of matrix multiplication protocol. We give a popular MatrixMul protocol in Figure 154 1 and analyze its efficiency in our framework. The protocol in Figure 1 requires $km^2 \log_2 N$ bit 155 online communication, where m is the number of users and k is the dimension of latent factors. As 156 for the usual case where the number of users is $\approx 10^4$, k = 10, and $\log_N = 64$, one invocation of 157 MatrixMul protocol would have a total communication of around 7.4GB. Considering 100 iterations 158 of our framework, this leads to ≈ 1491 GB communication, which is impractical. Fortunately, the 159 social matrices (**D**, **E**, and **S**) are highly sparse in social recommendation. In the following section, 160 we propose a PIR-based sparse matrix multiplication protocol with better communication efficiency. 161

162 3.1 Secure sparse matrix multiplication

Essentially, any matrix could be represented by a value vector and a location vector, where the value vector contains all non-zero values and the location vector contains locations of those values. That is, a sparse matrix $\mathbf{Y} \in \mathbb{R}^{m \times m}$ can be represented by a pair of vectors $(l_y \in \mathbb{N}_{m^2}^t, v_y \in \mathbb{R}^t)$, where t is the number of non-zero values in \mathbf{Y} .

167 **Dense-sparse matrix multiplication.** Considering the case 168 where $\mathbf{X} \in \mathbb{R}^{k \times m}$ is the dense matrix from P_0 and $\mathbf{Y} \in \mathbb{R}^{m \times m}$ 169 is the sparse matrix from P_1 . Now we consider the following 170 two cases.

171 *Case 1: insensitive sparsity, i.e., insensitive* l_y *and sensitive* 172 v_y . This refers to the case where the locations of zero values 173 are public or contain no sensitive information. Take the social 174 matrices (**D** and **E**) for example, both of them are diagonal, and 175 thus the location vector is insensitive while the value vector is 176 still sensitive.

Our protocol mainly works as follows. First, P_0 and P_1 parse **X** and **Y** into two tables T_x and T_y separately, where the value set of each bin in T_x is a subset of one row in **X**, that is, $T_x(i) \subseteq$ $x_{i,*}$. Similarly, bin set in T_y is a subset of one column in **Y**, $T_y(i) \subseteq y_{*,i}$. The intuition behind is to use bins to contain only



Figure 4: Matrix multiplication with insensitive sparsity.

the necessary values needed to calculate the output value (which means filter out the zero multiplies in each bin). Take the first bin for example (that is, $T_x(0)$ and $T_y(0)$), for $j \in [m]$, $T_x(0)$ contains all $x_{0,j}$ where $y_{j,0}$ is a non-zero value, and $T_y(0)$ contains all non-zero $y_{j,0}$. In order to get the final result, we perform the secure inner product protocol on $T_x(0)$ and $T_y(0)$, and denote the result as $[z_{0,0}]$. We show the high level idea in Figure 4. By doing this, our protocol concretely consumes $k|l_y|$ Beaver's triples and therefore has $O(k|l_y|)$ online communication complexity. Figure 5 shows **Protocol:** $[\![\mathbf{Z}]\!] \leftarrow \mathsf{MatrixMul}(\mathbf{X}, \mathbf{Y})$

1: P_0 generates an additive HE key pair (pk, sk), then sends pk to P_1 Offline Completed 2: $\forall i \in [k], j \in [m], P_0$ lets $e_{i,j} = \mathsf{Enc}(\mathsf{pk}, x_{i,j})$, and lets **E** be the encrypted matrix 3: $\forall (i,j) \in l_y, P_1 \text{ pushes } y_{i,j} \text{ into } T_y(j), \text{ also, } P_1 \text{ invokes } q_{i,j} \leftarrow \mathsf{PIR.Query}(i+jk)$ 4 : P_1 sends the query set (denoted as **q**) to P_0 5: $\forall q_{i,j} \in \mathbf{q}, P_0 \text{ invokes } r_{i,j} \leftarrow \mathsf{PIR}.\mathsf{Response}(\mathbf{E}, q_{i,j}).$ 6: P_0 sends the response set (denoted as **r**) to P_1 7: $\forall r_{i,j} \in \mathbf{r}, P_1 \text{ invokes } e_{i,j} \leftarrow \mathsf{PIR}.\mathsf{Extract}(r_{i,j}) \text{ and pushes } e_{i,j} \text{ to } T'_e(i)$ 8: **for** $i \in [k], j \in [m]$ **do** P_0 lets $\beta_{i,j} = \mathsf{Enc}_{\mathsf{pk}}(0)$ 9: 10: $\forall v \in T'_e(i), u \in T_y(j), P_0 \text{ invokes } \beta_{i,j} = v \otimes u \oplus \beta_{i,j}$ P_0 samples random numbers $g_{i,j} \leftarrow \mathbb{S}\mathbb{Z}_{\delta}$, then lets $\beta_{i,j} = g_{i,j} \oplus \beta_{i,j}$ 11: P_0 sends $\beta_{i,j}$ to P_1 , then lets $[z_{i,j}]_0 = -g_{i,j}$ 12 : 13: endfor 14: P_0 decrypts all receivng messages and lets $[z_{i,j}]_1 = \mathsf{Dec}_{\mathsf{sk}}(\beta_{i,j})$ 15 : return $\llbracket Z \rrbracket$

Figure 6: Dense-sparse MatrixMul with sensitive sparsity, that is, $\mathbf{X} \in \mathbb{R}^{k \times m}$, $\mathbf{Y} \in \mathbb{R}^{m \times m}$, and \mathbf{Y} 's location vector l_u is private.

- the technical details of our proposed protocol for case 1. For Line 1 in ST-MPC (Figure 3), clearly 188
- both parties know that **D** and **E** are diagonal matrices, that is, $|l_u| = m$. Therefore, our proposed 189 protocol in Figure 4 can drop the complexity from $O(km^2)$ to O(km).
- 190
- **Lemma 1.** The protocol in Figure 5 is secure in the MPC-hybrid model. 191
- *Proof.* Please find the proof in the Technical Appendix. 192

Case 2: sensitive sparsity, i.e., sensitive l_y and sensitive v_y . 193 For a more general case, where both the location vector and 194 the value vector contain sensitive information. Take the social 195 matrix S for instance, its location vector indicates the existence 196 of a social relation between two users, its value vector further 197 shows the strength of their relation, and both of which are 198 sensitive. 199

In this case, both the dense matrix **X** and the entire sparse ma-200 trix **Y** are sensitive. Following the idea in case 1, the matrix 201 multiplication protocol should first generate T_x, T_y according 202 to v_x, v_y and l_y , and then perform the inner product multipli-203 cation for each aligned bins in T_x, T_y . Still, P_1 can generate T_y according to its own inputs v_y, l_y . However, P_0 cannot 204 205 generate T_x directly, since v_x is kept by itself while l_y is held 206 by P_1 . We make a communication and computation trade-off 207 by leveraging PIR techniques, and as a result, our PIR-based 208



Figure 7: Matrix multiplication with sensitive sparsity.

approach has lower concrete communication, and overall is faster than the baseline protocol. 209

We show the high-level idea of our PIR-based protocol in Figure 7. The intuition behind is to let P_1 210 obliviously filter each bin in T_x since both value vector and location vector are sensitive. In summary, 211 first P_0 encrypts all the values in T_x , the encrypted table is denoted as T_e . Then P_1 and P_0 invoke PIR 212 protocol, where P_0 acts as server and sets T_e as PIR database, P_1 acts as client and parses l_y to many 213 PIR queries. At the end of PIR protocol, P_1 receives the encrypted and filtered table T'_e . Afterwards 214 P_1 performs secure inner product evaluation. By doing this, the communication complexity drops 215

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from $O(km^2)$ to $O(\alpha km)$, compared with the simple solution. The details of our protocol are shown in Figure 6. For Line 2 in ST-MPC (Figure 3), the social matrix (**S**) is sparse in nature, and thus our proposed protocol in Figure 7 can significantly improve its efficiency. In summary, with our proposed two secure MatrixMul protocols, one can securely calculate the social term efficiently. For instance, again considering the social recommendation with $\approx 10^4$ users, our proposal only requires a total of ≈ 3.6 GB communication for each iteration.

Lemma 2. The protocol in Figure 6 is secure in the PIR-hybrid model with the leakage of $|l_u|$.

223 Proof. Please find the proof in the Technical Appendix.

224 **3.2** Security discussions of the social term

In S³Rec, P_0 and P_1 securely calculate the social term $ST = \gamma U(D^T + E^T)/2 - \gamma US^T$ and reveal 225 the result to P_0 (Eq.(2)). We denote the ideal functionality of secure calculating the social term as \mathcal{F}_{st} . 226 During each epoch in Algorithm 3, P_0 sends U to \mathcal{F}_{st} , and P_1 sends D, E, S to \mathcal{F}_{st} accordingly. From 227 the view of P_0 , after each iteration, it additionally learns the output of \mathcal{F}_{st} , that is, the social term ST. 228 The security of S³Rec relies on whether P_0 can resolve the social matrix S^T given its own inputs U 229 and the social term ST. We claim that this is difficult because, the number of equations (#epoch, 100 230 in our experiments) is much smaller than that of the variables (#user, much more than 100 in practice), 231 which indicates that there are infinite solutions for this. Moreover, the reveal of the social term to P_0 232 could be avoided by taking the whole model training procedure as an MPC functionality and designing 233 a complicated protocol for it. Inevitably, such protocol introduces impractical communication costs, 234 and we leave how to solve this efficiently as a future work. 235

236 4 Experiments

Our experiments intend to answer the following questions. Q1: How do the social recommendation models using both rating data on P_0 and social data on P_1 outperform the model that only uses rating data on P_0 (Section 4)? Q2: How does our model perform compared with SeSoRec (Section 4)? Q3: How does the social data sparsity affect the performance of SeSoRec and our model (Section 4)?

Implementation and setup. We run our experiments on a machine with 4-Core 2.4GHz Intel Core 241 i5 with 16G memory, we compile our program using a modern C++ compiler (with support for C++ 242 standard 17). In addition, our tests were run in a local network, with \approx 3ms network latency. For 243 additive HE scheme, we choose the implementation of libpaillier¹. Also, we use Seal-PIR² with same 244 parameter setting as the original paper [1]. For security, we choose 128-bit computational security 245 and 40-bit statistical security as recommended by NIST [2]. Similarly we leverage the generic ABY 246 library³ to implement SeSoRec [5] and MPC building blocks such as addition, multiplication, and 247 truncation. In particular, we choose 64-bit secret sharing in all our experiments. 248

Dataset. We choose two popular benchmark datasets to evaluate the performance of our proposed model, i.e., Epinions [17] and LibraryThing (Lthing) [29], both of which are popularly used for evaluating social recommendation tasks. Following existing work [5], we remove the users and items that have less than 15 interactions for both datasets. We summarize the statistics of both datasets after process in Table 1. Notice that we assume users' rating data are located at P_0 , users' social data are located at P_1 , and P_0 and P_1 share the same user set.

Table 1: Dataset statistics

Dataset	#user	#item	#rating	rating density	#social relation	social density
Epinions	11,500	7,596	283,319	0.32%	275,117	0.21%
Lthing	15,039	14.957	529,992	0.24%	44,710	0.02%

¹libpaillier: http://acsc.cs.utexas.edu/libpaillier/, GPL license

²Seal-PIR: https://github.com/microsoft/SealPIR, MIT license

³ABY: https://github.com/encryptogroup/ABY, LGPL license

Table 2: Comparison results of different models in terms of model accuracy (in RMSE), running time (in seconds), and communication size (in GB), on Epinions and Lthing datasets.

Models		Epinio	ons dataset			Lthir	ig dataset	
Widdels	MF	Soreg	SeSoRec	S ³ Rec	MF	Soreg	SeSoRec	S ³ Rec
RMSE	1.193	1.062	1.062	1.062	0.927	0.908	0.908	0.908
Offline Time	-	-	7,271	10.86	-	-	14,450	8.912
Total Time	3.846	40.50	7,799	419.9	9.596	57.76	16,084	262.1
Offline Comm.	-	-	788.3	0	-	-	1,348	0
Total Comm.	-	-	798.6	3.552	-	-	1,365	2.201

Table 3: Comparison results by varying social data sparsity on Epinions and Lthing datasets.

Matria	Madala	Epinions			Lthing		
Metric	widdels	0.4	0.6	0.8	0.4	0.6	0.8
Total time (Seconds)	SesoRec S ³ Rec	7,799 366.3	7,799 381.2	7,799 401.8	16,084 194	16,084 217	16,084 238
	(Improvement)	(21.29x)	(20.46x)	(19.41x)	(82.91x)	(74.12x)	(67.58x)
Total communication	SesoRec S ³ Rec	798 3.12	798 3.29	798 3.46	1,366 1.62	1,366 1.82	1,366 2.01
	(Improvement)	(255x)	(243x)	(231x)	(843x)	(751x)	(680x)

Comparison Methods. We compare S^{3} Rec with the following classic and state-of-the-art models:

- *MF* [18] is a classic matrix factorization model that only uses rating data on P_0 , i.e., when $\gamma = 0$ for S³Rec.

- Soreg [16] is a classic social recommendation model, which does not consider data privacy and assumes both rating data and social data are available on P_0 .

SeSoRec [5] tries to solve the privacy-preserving cross-platform social recommendation problem,
 but suffers from security and efficiency problem.

Hyper-parameters. For all the model, during comparison, we set k = 10. We tune learning rate θ and regularizer parameter λ in $\{10^{-3}, 10^{-2}, ..., 10^1\}$ to achieve their best values. We also report the effect of K on model performance.

Metrics. We will evaluate both accuracy and efficiency of our proposed model. For accuracy, we choose Root Mean Square Error (RMSE) as the evaluation metric, since ratings range in [0, 5]. For efficiency, we report the computation time (in seconds) and the communication size between P_0 and P_1 (in gigabytes), if has, for all the models. We use five-fold cross-validation during experiments.

Performance Comparison. We first compare the model performances in terms of accuracy (RMSE)
 and efficiency (total time and communication). Table 2 shows the time and communication for each
 epoch, where time is shown in seconds, and communication is shown in GB.

From those Tables, we find that: (1) the use of social information can indeed improve the recom-272 mendation performance of the rating platform, e.g., 1.193 vs. 1.062 and 0.927 vs. 0.098 in terms of 273 RMSE on Epinions and Lthing, respectively. This result is consistent with existing work from [16, 5]; 274 (2) despite the same RMSE as SeSoRec and Soreg, S^{3} Rec significantly improves the efficiency of 275 276 SeSoRec, especially on the more sparse Lthing dataset, reducing the total time for one epoch from around 4.5 hours to around 4.5 minutes, and reducing the total communication from nearly 1.3TB to 277 around 2.2GB. This yields an improvement of $18.57 \times$ faster, and $224.8 \times$ less communication on 278 Epinions and $61.37 \times$ faster and $620.2 \times$ less communication on Lthing, respectively. 279

Effect of Social Data Sparsity. Next, we try to study the effect of social data sparsity on training efficiency. In order to do this, we sample the social relation of both datasets with a rate of 0.8, 0.6, and 0.4. As the result, the RMSEs of both SeSoRec and S³Rec decrease to 1.0932, 1.1373,

Models	k = 10	$\begin{array}{c} \text{SeSoRec} \\ k = 15 \end{array}$	k = 20	k = 10	$S^{3}Rec$ k = 15	k = 20
Offline Time	7,271	12,651	17,676	10.86	9.667	9.815
Total Time	7,799	13,565	19,585	419.9	449.6	527.4
Offline Comm.	788.3	1,182	1,577	0	0	0.
Total Comm.	798.6	1,198	1,597	3.552	3.552	3.552

Table 4: Effect of k on running time and communication size on Epinions dataset

1.1751 on Epinions dataset, and 0.9112, 0.9187, 0.9210 on Lthing dataset. The rational behind is that recommendation performance decreases with the number of social relations. We also report the efficiency of both models on Epinions and Lthing datasets in Table 3. From it, we can find that the computation time and communication size of SeSoRec are constant no mater what the sample rate is. In contrast, the computation time and communication size of S³Rec decrease linearly with sample rate. This result benefits from that S³Rec can deal with sparse social data with our proposed sparse matrix multiplication protocols.

Effect of k. For efficiency, we report the running time and communication size of SeSoRec and 290 PriorRec w.r.t k in Table 4, where we use the Epinions dataset. From it, we can get that in average, 291 S^{3} Rec improves SeSoRec **18.6x** in terms of total running time and **225x** in terms of communication. 292 More specifically, we observe that (1) the total running time of both SeSoRec and PriorRec increase 293 with k, but the increase rate of $S^{3}Rec$ is slower than that of SeSoRec; (2) the communication size 294 of SeSoRec increases with k, in contrast, the communication size of S³Rec is constant. This result 295 demonstrates that our proposed S³Rec has better scalability than SeSoRec in terms of both running 296 time and communication size. 297

298 5 Related Work

Traditional recommender systems that only consider user-item rating information suffer from severe data sparsity problem [18]. On the one hand, researchers extensively incorporate other kinds of information, e.g., social [25], review [23], location [14], and time [6], to further improve recommendation performance. On the other hand, existing studies begin to explore information on multiple platforms or domains to address the data sparsity problem, i.e., cross-platform and cross-domain recommendation [15, 30]. However, most of them cannot solve the data isolation problem in practice.

So far, there has been several works that may be applied for privacy-preserving cross-domain 305 recommendations. For example, [21] applied garbled circuits for secure matrix factorization, and it 306 has high security but low efficiency. Chai et al. [4] adopted homomorphic encryption for federated 307 matrix factorization, but it assumes the existence of a semi-honest server and is not provable secure. 308 [11] uses differential privacy to protect user location privacy using transfer learning technique, which 309 is not provable secure and does not suitable to our problem. The most similar work to ours is SeSoRec 310 [5], however, it suffers from two main shortcomings: (1) as admitted by SeSoRec, it improves 311 efficiency by sacrificing security. That is, it reveals the sum of two rows or two columns of the input 312 matrix. We emphasis that this raises serious security concern in the social recommendation since 313 one may infer detailed social relations from the element-wise sum of two rows/columns of the user 314 315 social matrix, especially when social relations are binary values; (2) SeSoRec treats the social data as a dense matrix and thus still has serious efficiency issue under the practical sparse social data setting. 316

317 6 Conclusion

This paper aims to solve the data isolation problem in cross-platform social recommendation. To do this, we proposed S^3 Rec, a sparsity-aware secure cross-platform social recommendation framework. S^3 Rec conducts social recommendation task and preserves data privacy at the same time. We also proposed two secure sparse matrix multiplication protocols to improve the model training efficiency. Experiments conducted on two datasets demonstrate that S^3 Rec improves the computation time and communication size by around $40 \times$ and $423 \times$ on average, compared with the state-of-the-art work.

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388 Checklist

389	1.	For a	Il authors
390 391		(a)	Do the main claims made in the abstract and introduction accurately reflect the paper's contribu- tions and scope? [Yes]
392		(b)	Did you describe the limitations of your work? [Yes]
393		(-)	Our proposed sparse-aware secure cross-domain social recommendation model is customized to
394			a classic factorization based social recommendation models i.e. Soreg [16] We did not choose
205			the state-of-the-art social recommendation model, e.g. deen neural network based model [7]
395			since it contains complicated non-linear computations and will suffer from serious efficiency
390			since it complete the implete the implete to implete the set of th
397		(-)	Did was discuss any network is a section as sisted investor of some marks [New]
398		(c)	Did you discuss any potential negative societal impacts of your work? [Yes]
399			Our proposed model can be used to build cross-domain social recommendation model while
400			protecting the data privacy of both domains. The proposed sparse-aware matrix multiplication
401			protocols can be widely used to other machine learning scenarios besides social recommendation
402			model. This paper could shed light on the future research direction of recommendation area. To
403			our knowledge, our work does not have any potential negative societal impacts.
404		(d)	Have you read the ethics review guidelines and ensured that your paper conforms to them? [Yes]
405	2.	If yo	u are including theoretical results
406		(a)	Did you state the full set of assumptions of all theoretical results? [Yes] We assume that all
407			participants are semi-honest, we also prove the security of our protocols using simulation-based
408			technique.
409		(b)	Did you include complete proofs of all theoretical results? [Yes] Please find the proof details in
410			the technical appendix.
411	3.	If yo	u ran experiments
412 413		(a)	Did you include the code, data, and instructions needed to reproduce the main experimental results (either in the supplemental material or as a URL)? [Yes]
414		(h)	Did you specify all the training details (e.g. data splits hyperparameters how they were chosen)?
415		(0)	[Yes]
416 417		(c)	Did you report error bars (e.g., with respect to the random seed after running experiments multiple times)? [N/A]
418 419		(d)	Did you include the total amount of compute and the type of resources used (e.g., type of GPUs, internal cluster, or cloud provider)? [Yes]
420	4.	If yo	u are using existing assets (e.g., code, data, models) or curating/releasing new assets
421		(a)	If your work uses existing assets, did you cite the creators? [Yes]
422		(b)	Did you mention the license of the assets? [Yes]
423		(c)	Did you include any new assets either in the supplemental material or as a URL? [N/A]
424		(d)	Did you discuss whether and how consent was obtained from people whose data you're us-
425		(4)	ing/curating? [N/A]
426		(e)	Did you discuss whether the data you are using/curating contains personally identifiable informa-
420		(U)	tion or offensive content? $[N/\Delta]$
-27	F	T£	
428	э.	пуо	u used crowdsourcing or conducted research with numan subjects
429 430		(a)	Did you include the full text of instructions given to participants and screenshots, if applicable? [N/A]
431		(b)	Did you describe any potential participant risks, with links to Institutional Review Board (IRB)
432			approvals, if applicable? [N/A]
433		(c)	Did you include the estimated hourly wage paid to participants and the total amount spent on
434			participant compensation? [N/A]