Error Correction Code Transformer

Abstract

Error correction code is a major part of the physical communication layer, ensuring the reliable transfer of data over noisy channels. Recently, neural decoders were shown to outperform classical decoding techniques. However, the existing neural approaches present strong overfitting, due to the exponential training complexity, or a restrictive inductive bias, due to reliance on Belief Propagation. Recently, Transformers have become methods of choice in many applications, thanks to their ability to represent complex interactions between elements. In this work, we propose to extend for the first time the Transformer architecture to the soft decoding of linear codes at arbitrary block lengths. We encode each channel's output dimension to a high dimension for better representation of the bits' information to be processed separately. The element-wise processing allows the analysis of channel output reliability, while the algebraic code and the interaction between the bits are inserted into the model via an adapted masked self-attention module. The proposed approach demonstrates the power and flexibility of Transformers and outperforms existing state-of-the-art neural decoders by large margins, at a fraction of their time complexity.

1 Introduction

Reliable digital communication is of major importance in the modern information age and involves the design of codes to be robustly decoded under noisy transmission channels. While optimal decoding is defined by the NP-hard maximum likelihood rule, the efficient decoding of algebraic block codes is an open problem. Recently, powerful learning-based techniques have been introduced to this field.

At present, deep learning models that implement parameterized versions of the legacy Belief Propagation (BP) decoders are dominant [22, 23, 21, 24, 4]. This approach is appealing, since the exponential size of the input space is reduced by ensuring that certain symmetry conditions are met. In this case, it is sufficient to train the network on a noisy version of the zero codeword. However, this approach remains extremely restrictive, due to the strong inductive bias induced by the Tanner graph.

Model-free decoders [26, 11, 16] employ variants of generic neural networks and may potentially benefit from the application of powerful architectures that have emerged in recent years in various fields. However, such decoders suffer from the curse of dimensionality, since the training grows exponentially with the number of information bits [34]. A major challenge of model-free decoders, therefore, is the difficulty of the network in learning the code and identifying its most reliable components, especially for the basic multi-layer perceptron architectures suggested in the literature.

In this work, we propose a *model-free* decoder built upon the Transformer architecture [32]. As far as we can ascertain, this is the first adaptation of Transformers to error correction codes. Our Transformer employs a high-dimensional *scaled* element-wise embedding of the input which can be seen as positional and reliability encoding. Our decoder employs an *adapted-mask* self-attention mechanism, in which the interaction between bits follows the code's parity check matrix, enabling the incorporation of domain knowledge into the model.

Submitted to 36th Conference on Neural Information Processing Systems (NeurIPS 2022). Do not distribute.

Applied to a wide variety of codes, our method outperforms the state-of-the-art learning-based solutions by very large margins, even with very shallow architectures, enabling potential deployment. This is the first time a decoder designed de-novo outperforms neural BP-based decoders.

2 Related Works

Over the past few years, the emergence of deep learning has demonstrated the advantages of Neural Networks in many communication applications, such as channel equalization, modulation, detection, quantization, compression, and *decoding* [15]. In [26] the entire coding-decoding channel transmission pipeline was abstracted as a fully parameterized autoencoder. In many scenarios, neural network-based methods outperform existing decoders, as well as popular codes [23, 11].

Two main classes of Neural decoders can be established in the current literature [28]. Model-free decoders employ general types of neural network architectures, as in [6, 11, 16, 2]. The results obtained in [11] for polar codes ($n = 16$) are similar to the maximum a posteriori (MAP) decoding results. However, the exponential number of possible codewords make the decoding of larger codes unfeasible. In [2], a preprocessing of the channel output allows the decoder to remain provably invariant to the codeword and enables the deployment of any kind of model without any overfitting cost. The model-free approaches generally make use of stacked fully connected networks or recurrent neural networks in order to simulate the iterative process existing in many legacy decoders. However, these architectures have difficulties in learning the code and analyzing the reliability of the output, and generally require prohibitive parameterization or expensive graph permutation preprocessing [2].

The second class is that of model-based decoders [22, 23] implementing augmented parameterized versions of classical BP decoders, where the Tanner graph is unfolded into a neural network in which weights are assigned to each variable edge, resulting in improvement in comparison to the baseline BP method. This approach remains limited since it is built on top of a restrictive model. Recently, [28] presented a data-driven framework for codeword permutation selection as a pre-decoding stage. A self-attention mechanism is applied over several permutation-specific embeddings to estimate the best permutation (among a given set) to be decoded. A combination of the BP model with a hyper-graph network was shown to enhance performance [24]. Further improvements over the latter were obtained by introducing an autoregressive signal based on previous state conditioning and SNR estimation [25].

Transformer neural networks were originally introduced for machine translation [32] and now dominate the performance tables for most applications in the field of Natural Language Processing (NLP), e.g., [32, 8]. Transformer encoders primarily rely on the self-attention operation in conjunction with feed-forward layers, allowing manipulation of variable-size sequences and learning of long-range dependencies. Many works extended the Transformer architecture from NLP to various applications, such as computer vision [9], speech analysis [20] and molecular properties prediction [7].

3 Background

We provide the necessary background on error correction coding and the Transformer architecture.

3.1 Coding

We assume a standard transmission that uses a linear code C. The code is defined by the binary generator matrix G of size $k \times n$ and the binary parity check matrix H of size $(n - k) \times n$ defined such that $GH^T = 0$ over the two elements' Galois field.

The input message $m \in \{0, 1\}^k$ is encoded by G to a codeword $x \in C \subset \{0, 1\}^n$ satisfying $Hx = 0$ and transmitted via a Binary-Input Symmetric-Output channel (e.g. an AWGN channel). Let y denote the channel output represented as $y = x_s + z$, where x_s denotes the Binary Phase Shift Keying (BPSK) modulation of x (i.e., over $\{\pm 1\}$), and z is random noise independent of the transmitted x.

The main goal of the decoder $f : \mathbb{R}^n \to \mathbb{R}^n$ is to provide a soft approximation $\hat{x} = f(y)$ of the codeword. An illustration of the coding framework is presented in Figure 1.

We wish to employ the Transformer architecture as part of a model-free approach. We, therefore, follow the preprocessing and post-processing of [2] to remain provably invariant to the transmitted

Figure 1: Illustration of the communication system. Our work focuses on the design and training of the parameterized decoder f_θ.

codeword and to avoid overfitting. This step involves a transformation of the channel output y, which entails a loss of information without implying any intrinsic performance penalty in decoding, i.e., the preprocessing coupled with an appropriately designed decoder can achieve MMSE decoding.

The preprocessing replaces y with a vector of dimensionality $2n - k$ defined as

$$\tilde{y} = h(y) = [|y|, s(y)], \tag{1}$$

where, $[\cdot, \cdot]$ denotes vector concatenation, $|y|$ denotes the absolute value (magnitude) of y and $s(y) = H y_b \in \{0, 1\}^{n-k}$ denotes the binary code *syndrome* obtained via the multiplication of the parity check matrix by the binary mapping y_b of y such that $y_b = 0.5(1 - \text{sign}(y))$.

The post-processing predicts the transmitted codeword by multiplying the elements of y with the predicted multiplicative noise $\tilde{z}$ defined such that $y = x_s \tilde{z}$. Namely, the prediction takes the form

$$\hat{x} = y \cdot f(h(y)). \tag{2}$$

In our setting, the decoder is parameterized by θ such that $\hat{x} = y \cdot f_\theta(|y|, H y_b)$.

3.2 Transformers

The Transformer was introduced as a novel, attention-based building block for machine translation [32]. The input sequence is first embedded into a high-dimensional space, coupled with positional embedding for each element. The embeddings are then propagated through multiple normalized self-attention and feed-forward blocks.

The self-attention mechanism introduced by Transformers is based on a trainable associative memory with (key, value) vector pairs, where a query vector $q \in \mathbb{R}^d$ is matched against a set of k key vectors using scaled inner products, as follows

$$A(Q, K, V) = \text{Softmax}\left(\frac{QK^T}{\sqrt{d}}\right) V, \tag{3}$$

where $Q \in \mathbb{R}^{N \times d}$, $K \in \mathbb{R}^{k \times d}$ and $V \in \mathbb{R}^{k \times d}$ represent the packed N queries, k keys and values tensors respectively. Keys, queries and values are obtained using linear transformations of the sequence's elements. A multi-head self-attention layer is defined by extending the self-attention mechanism using h attention *heads*, i.e. h self-attention functions applied to the input, reprojected to values via a $dh \times D$ linear layer.

4 Error Correction Code Transformer

We present the elements of the proposed Transformer decoder, the complete architecture, and the training procedure.

4.1 Positional Reliability Encoding

We first consider each dimension of $\{\tilde{y}_i\}_{i=1}^{2n-k}$ separately and project each one to a high d dimensional embedding $\{\phi_i\}_{i=1}^{2n-k}$ such that

$$\phi_i = \begin{cases} |y_i| W_i, & \text{if } i \leq n \\ \left(1 - 2(s(y))_{i-n+1}\right) W_i, & \text{otherwise} \end{cases} \tag{4}$$

where $\{W_j \in \mathbb{R}^d\}_{j=1}^{2n-k}$ denotes the one-hot encoding defined according to the bit position. The embedding is modulated by the magnitude and syndrome values, such that less reliable elements (i.e. low magnitude) collapse to the origin. This property becomes especially appealing when applied to the standard dot product of the self-attention module: for the first layer and two distinct information bits embedding ϕ_i, ϕ_j we have

$$\langle \phi_i, \phi_j \rangle = \begin{cases} |y_i||y_j|\langle W_i, W_j \rangle & i, j \leq n \\ |y_i|(1 - 2(s(y))_{j-n+1})\langle W_i, W_j \rangle, & i \leq n < j, \end{cases} \tag{5}$$

where an unreliable information bit vanishes and a non-zero syndrome entails a negative scaling, potentially reducing the impact on the softmax aggregation, as illustrated in Sec. 6.

In contrast to [2], which requires permutation of the code in order to artificially provide to the fully connected network indications about the most reliable channel outputs, in our construction, the channel output reliability is directly obtained and maintained in the network via the scaled bit-wise embedding. The proposed scaled encoding can be thought of as a *positional* encoding according to the input *reliability* since the bit positions are fixed. We note that representing the codeword by a set of one-hot vectors is very different from any of the existing methods.

4.2 Code-Aware Self-Attention

In order to detect and correct errors, a decoder must *analyze and compare* the received word bits via the parity check equations, such that a non-zero syndrome indicates that channel errors have occurred. In our Transformer-based architecture, bit comparisons become natural via the self-attention *interaction* mechanism.

However, comparing every pair of elements, as is generally performed in Transformer architectures, is sub-optimal, since not every bit is necessarily related to all the others. We propose to make use of the self-attention mechanism in order to incorporate fundamental domain knowledge about the relevant code. Specifically, we use it to indicate that the syndrome values should only be dependent on the corresponding parity check bits.

Thus, for a given algebraic code defined by the matrix H we propose to define a function $g(H) : \{0,1\}^{(n-k)\times k} \to \{-\infty, 0\}^{2n-k \times 2n-k}$ defining the mask to be applied to the self-attention mechanism, such that

$$A_H(Q, K, V) = \text{Softmax}\left(\frac{QK^T + g(H)}{\sqrt{d}}\right)V. \tag{6}$$

We propose to build the *symmetric* mask such that it contains information about *every* pairwise bit relation, as follows. The mask is first initialized as the identity matrix. For each row i of the binary parity check matrix H, we unmask at the locations of *every pair* of ones in the row, since those bits are connected and may impact each other as well as the syndrome in the decoding procedure. We also unmask the location of pairs of ones with the corresponding syndrome bit at $n + i$, since they define the parity check equations. Illustrations of typical self-attention maps are given in Appendix A.

Algorithm 1: Mask construction Pseudo Code

```
1  function g(H)
2      l,n = H.shape
3      k = n-l
4      mask = eye(2n-k)
5      for i in range(0,n-k) do
6          idx = where(H[i]==1)
7          for j in idx do
8              mask[n+i,j] = mask[j,n+i] = 1
9              for k in idx do
10                 mask[j,k] = mask[k,j] = 1

11     return -∞(¬ mask)
```

Figure 2: We present the corresponding parity-check matrix (a), the induced Tanner graph (c) and the proposed masking (b) on the Hamming(7,4) code.

We summarize the construction of the mask in Algorithm 1. This construction enables greater freedom in decoding than the relations enabled by the Tanner graph, since related bits may impact each other *beyond* the parity check equations as depicted in Figure 2. While regular Transformers can be assimilated to a neural network applied on a complete graph, the proposed mask can be seen as the adjacency matrix of the Tanner graph extended to a two rings connectivity. In contrast to BP, which collapses the information via interleaved variable and check layers, the masked self-attention allows the simultaneous cross-analysis of related elements.

Most importantly, since the mask is fixed and computed once, the self-attention quadratic complexity bottleneck $\mathcal{O}(n^2 d)$ is now reduced to the density of the code $\mathcal{O}(\sum_{ij}(H)_{ij}d)$. This property is especially appealing for low-density codes (e.g. [10]), while in all our experiments the original complexity of the self-attention is reduced by 84% on average, as presented in Sec. 6.

4.3 Architecture and Training

The initial encoding is defined as a d dimensional one hot encoding of the $2n - k$ input elements. The decoder is defined as a concatenation of N decoding layers composed of self-attention and feed-forward layers interleaved by normalization layers. The output module is defined by two fully connected layers. The first layer reduces the element-wise embedding to a one-dimensional $2n - k$ vector and the second to a n dimensional vector representing the soft decoded noise. An illustration of the model is given in Figure 3.

The dimension of the feed-forward network is four times that of the embedding [32] and is composed of GEGLU layers [30], with layer normalization set to the pre-layer norm setting, as in [18, 35]. We use an eight-head self-attention module in all experiments. colorblue The impact of the number of heads on the accuracy is further explored in Appendix C.

The training objective is the cross-entropy function where the goal is to learn to predict the *multiplicative* noise $\tilde{z}$ [2]. Denoting the *soft* multiplicative noise $\tilde{z}_s$ such that $y = x_s \tilde{z}_s$, we obtain $\tilde{z}_s = \tilde{z}_s x_s^2 = y x_s$. Thus, the binary multiplicative noise to be predicted is defined by $\tilde{z} = \text{bin}(y \cdot x_s)$, such that the loss computed for a single received word y is

$$\mathcal{L} = -\sum_{i=1}^{n} \tilde{z}_i \log(f_\theta(y)) + (1 - \tilde{z}_i)\log(1 - f_\theta(y)). \tag{7}$$

The estimated hard-decoded codeword is straightforwardly obtained as $\hat{x}_b = \text{bin}(\text{sign}(f_\theta(y) \cdot y))$.

The Adam optimizer [17] is used with 128 samples per minibatch, for 1000 epochs, with 1000 minibatches per epoch. For $N = 10$ architectures we trained the models for 1500 epochs. We note that using more epochs can improve performance. However, the current setting is already more than enough to reach new SOTA performance. Due to the construction of the model and its input preprocessing, the zero codeword is enough for training. The additive Gaussian noise is sampled randomly per batch in the $\{3, \ldots, 7\}$ normalized SNR (i.e. E_b/N_0) range. Additional experiments with the non-Gaussian Rayleigh fading channel are provided in Appendix E. We initialized the learning rate to 10^{-4} coupled with a cosine decay scheduler down to $5 \cdot 10^{-7}$ at the end of the training. No warmup was employed [35]. Training time range from 12 to 24 hours depending on the code length, and no optimization of the self-attention mechanism have been employed. Training and

Figure 3: Illustration of the proposed Transformer architecture. The main differences from other Transformers are the initial scaled bit-embedding, the construction of the code-aware masked self-attention, and the output module.

experiments have been performed on a 12GB Titan V GPU. The training time was in the range of 19-40, 40-102, and 56-171 seconds per epoch for the $N = 2, 6, 10$ architectures, respectively.

5 Experiments

To evaluate our method, we train the proposed architecture with three classes of linear block codes: Low-Density Parity Check (LDPC) codes [10], Polar codes [1] and Bose–Chaudhuri–Hocquenghem (BCH) codes [3]. All parity check matrices are taken from [12]. The proposed architecture is defined solely by the number of encoder layers N and the dimension of the embedding d. We compare our method to the BP algorithm, to the augmented hypernetwork BP algorithm of [24] (hyp BP), to the RNN architecture of [2] and to the very recent SOTA performance of [25] (AR BP). All the results were obtained from the corresponding papers. A comparison to the SCL Polar decoder [31] is given in Appendix D. The results are reported as bit error rates (BER) for different normalized SNR values (dB). We follow the testing benchmark of [24, 25]. During testing, our decoder decodes at least 10^5 random codewords, to obtain at least 500 frames with errors at each SNR value.

The results are reported in Tab. 1, where we present the negative natural logarithm of the BER. For each code we present the results of the BP-based competing methods for 5 and 50 iterations (first and second rows), corresponding to a neural network with 10 and 100 layers, respectively. We present our framework performance for six different architectures with $N = \{2, 6\}$ $d = \{32, 64, 128\}$ respectively (first to third rows). Since LDPC codes are optimized for BP decoding [29], we also present for these codes the performance of a $N = 10, d = 128$ architecture. As can be seen, our approach outperforms current SOTA results by very large margins on several codes, at a fraction of the number of iterations, and is able to outperform legacy methods (e.g. BP) *at convergence*, even with extremely shallow architectures (e.g. $N = 2, d = 32$). We remark that the proposed method seems to perform better for high rate (i.e. $\frac{k}{n}$) codes.

We provide plots of more BER values for some of the codes in Figure 4(a,b) for Polar(64,32) and BCH(63,51) codes, respectively. Our method is able to outperform the converged AR-BP by up to two orders of magnitude for high SNR values. Performance on larger codes is studied in Appendix H.

In Fig. 4(c), We compare our method with the best model of [2] on the only code provided BCH(127,64). Their network is built as a five-layer stacked GRU model, with five iterations with an embedding of $5(2n - k) = 950$ trained for 5000 epochs, with a batch size of 1000. As can be seen, our method outperforms this model by up to 211%, with only half the number of layers, one-tenth as many parameters, and a fraction of the training requirements.

Table 1: A comparison of the negative natural logarithm of Bit Error Rate (BER) for three normalized SNR values of our method with literature baselines. Higher is better.

Concurrent results are obtained after $L = 5$ BP iterations in first row (i.e. 10 layers neural network) and *at convergence* results in second row obtained after $L = 50$ BP iterations (i.e. 100 layers neural network). Best results in **bold**, second best is underlined, and the minimal Transformer architecture to outperform every other competing method is in *italic*.

Our performance is presented for seven different architectures: for $N = \{2, 6\}$, we present results for $d = \{32, 64, 128\}$ (first to third rows), and for $N = 10$ we run only the $d = 128$ configuration.

Method	BP[27]			Hyp BP[24]			AR BP[25]			Ours N=2			Ours N=6			Ours N=10		
	4	5	6	4	5	6	4	5	6	4	5	6	4	5	6	4	5	6
Polar(64,32)	3.52	4.04	4.48	4.25	5.49	7.02	4.77	6.30	8.19	4.27	5.44	6.95	*5.71*	*7.63*	*9.94*			
	4.26	5.38	6.50	4.59	6.10	7.69	5.57	7.43	9.82	4.57	5.86	7.50	6.48	8.60	11.43			
										4.87	6.2	7.93	**6.99**	**9.44**	**12.32**			
Polar(64,48)	4.15	4.68	5.31	4.91	6.48	8.41	5.25	6.96	9.00	4.92	6.46	8.41	*5.82*	*7.81*	10.24			
	4.74	5.94	7.42	4.92	6.44	8.39	5.41	7.19	9.30	5.14	6.78	8.9	6.15	8.20	10.86			
										5.36	7.12	*9.39*	**6.36**	**8.46**	**11.09**			
Polar(128,64)	3.38	3.80	4.15	3.89	5.18	6.94	4.02	5.48	7.55	3.51	4.52	5.93	4.47	6.34	8.89			
	4.10	5.11	6.15	4.52	6.12	8.25	4.84	6.78	9.30	3.83	5.16	7.04	*5.12*	*7.36*	*10.48*			
										4.04	5.52	7.62	**5.92**	**8.64**	**12.18**			
Polar(128,86)	3.80	4.19	4.62	4.57	6.18	8.27	4.81	6.57	9.04	4.30	5.58	7.34	*5.36*	*7.45*	*10.22*			
	4.49	5.65	6.97	4.95	6.84	9.28	5.39	7.37	10.13	4.49	5.90	7.75	5.75	8.16	11.29			
										4.75	6.25	8.29	**6.31**	**9.01**	**12.45**			
Polar(128,96)	3.99	4.41	4.78	4.73	6.39	8.57	4.92	6.73	9.30	4.56	5.98	7.93	*5.39*	*7.62*	*10.45*			
	4.61	5.79	7.08	4.94	6.76	9.09	5.27	7.44	10.2	4.69	6.20	8.30	5.88	8.33	11.49			
										4.88	6.58	8.93	**6.31**	**9.12**	**12.47**			
LDPC(49,24)	5.30	7.28	9.88	5.76	7.90	11.17	6.05	8.13	11.68	4.51	6.07	8.11	5.74	8.13	11.30	*6.68*	*9.53*	*13.30*
	6.23	8.19	11.72	6.23	8.54	11.95	6.58	9.39	12.39	4.58	6.18	8.46	5.91	8.42	11.90			
										4.71	6.38	8.73	6.13	8.71	12.10			
LDPC(121,60)	4.82	7.21	10.87	5.22	8.29	13.00	5.22	8.31	13.07	3.88	5.51	8.06	4.98	7.91	12.70	*5.73*	*9.35*	*15.01*
	-	-	-	-	-	-	-	-	-	3.89	5.55	8.16	5.02	7.94	12.72			
										3.93	5.66	8.51	5.17	8.31	13.30			
LDPC(121,70)	5.88	8.76	13.04	6.39	9.81	14.04	6.45	10.01	14.77	4.63	6.68	9.73	6.11	9.62	15.10	*7.13*	**11.50**	**17.92**
	-	-	-	-	-	-	-	-	-	4.64	6.71	9.77	6.28	10.12	15.57			
										4.67	6.79	9.98	6.40	*10.21*	*16.11*			
LDPC(121,80)	6.66	9.82	13.98	6.95	10.68	15.80	7.22	11.03	15.90	5.27	7.59	10.08	6.92	10.74	15.10	**8.12**	**12.90**	17.82
	-	-	-	-	-	-	-	-	-	5.29	7.63	10.90	7.17	*11.21*	*16.31*			
										5.30	7.65	11.03	*7.41*	11.51	16.44			
MacKay(96,48)	6.84	9.40	12.57	7.19	10.02	13.16	7.43	10.65	14.65	4.95	6.67	8.94	6.88	9.86	13.40	*8.39*	**12.24**	**16.41**
	-	-	-	-	-	-	-	-	-	5.04	6.80	9.23	7.10	10.12	14.21			
										5.17	7.07	9.64	7.38	*10.72*	*14.83*			
CCSDS(128,64)	6.55	9.65	13.78	6.99	10.57	15.27	7.25	10.99	16.36	4.35	6.01	8.30	6.34	9.80	14.40	*8.02*	*12.60*	*17.75*
	-	-	-	-	-	-	-	-	-	4.41	6.09	8.49	6.65	10.40	15.46			
										4.59	6.42	9.02	6.88	10.90	15.90			
BCH(31,16)	4.63	5.88	7.60	5.05	6.64	8.80	5.48	7.37	9.61	4.51	5.74	7.35	*5.74*	*7.42*	*9.59*			
	-	-	-	-	-	-	-	-	-	4.78	6.15	7.98	5.85	7.52	10.08			
										5.18	6.82	8.91	**6.39**	**8.29**	**10.66**			
BCH(63,36)	3.72	4.65	5.66	3.96	5.35	7.20	4.33	5.94	8.21	3.79	4.87	6.35	4.42	5.91	8.01			
	4.03	5.42	7.26	4.29	5.91	8.01	4.57	6.39	8.92	4.05	5.28	7.01	4.62	6.24	8.44			
										4.21	5.50	7.25	**4.86**	***6.65***	***9.10***			
BCH(63,45)	4.08	4.96	6.07	4.48	6.07	8.45	4.80	6.43	8.69	4.47	5.88	7.81	*5.16*	*7.02*	*9.75*			
	4.36	5.55	7.26	4.64	6.27	8.51	4.97	6.90	9.41	4.66	6.16	8.17	5.41	7.49	10.25			
										4.79	6.39	8.49	**5.60**	**7.79**	**10.93**			
BCH(63,51)	4.34	5.29	6.35	4.64	6.08	8.16	4.95	6.69	9.18	4.60	6.05	8.05	*5.20*	7.08	*9.65*			
	4.5	5.82	7.42	4.80	6.44	8.58	5.17	7.16	9.53	4.78	6.34	8.49	5.46	7.57	10.51			
										5.01	6.72	9.03	**5.66**	**7.89**	**11.01**			

217 6 Analysis

218 We study the impact of the proposed embedding and the masking procedure; furthermore, we analyze
219 and compare the complexity of our method with other methods.

Figure 4: BER for various values of SNR for (a) Polar(64,32) and (b) BCH(63,51) codes. Comparison with the best model of Bennatan et al. [2] on BCH(127,64) code (c).

Figure 5: Analysis of the scaled embedding values of the first column of the masked self-attention map with corrupted and uncorrupted codeword (CW) for the standardized Hamming code, using the (a) $(N, d) = (4, 32)$ and (b) $(N, d) = (2, 32)$ architecture. The impact of the noisy embedding on the self-attention's first parity-check element (index 7) is visible.

6.1 Impact of the Reliability Embedding

We depict in Figure 5 the impact of the *scaled* embedding on the self-attention mechanism. We chose the popular Hamming$(7, 4)$ code, corrupting the zero codeword with additive noise at the *first* bit (zero bit index), involving a non-zero syndrome at the first parity check equation. We present the *masked* softmax values of the self-attention map at index zero, i.e. we visualize the influence of the first (unreliable) bit on the self-attention aggregation (first column of the map). As can be seen, the first bit embedding is impactless when it is corrupted and then detectable, while its impact on the syndrome embedding is considerably increased (index 7). Once the NN corrects the bit (last layer(s)), the values return to normal. A similar analysis for larger codes can be found in in Appendix B.

6.2 Ablation Study: Masking

We present in Figure 6 the impact of the proposed mask on the convergence. We show an *unmasked* architecture, where we leave the network to learn the code by itself (i.e. $g(H) = 0$), and compare it with our proposed masking approach. The benefit obtained with the mask is clear and understandable in light of the difficulties the Transformer may have in learning the code by itself. Therefore, the connectivity provided by the proposed masking framework reduces the loss by 76%, 69% and 66% for the BCH$(63, 36)$, POLAR$(64, 32)$ and LDPC$(49, 24)$ codes, respectively.

Figure 6: The effect of the proposed masking procedure on the convergence for BCH$(63, 36)$, POLAR$(64, 32)$ and LDPC$(49, 24)$ using $(N, d) = (6, 64)$.

Figure 7: Sparsity and complexity ratio of the proposed self-attention map for several codes.

6.3 Complexity Analysis

We first present, in Figure 7, the sparsity ratio of the *masked* self-attention map, as well as the induced complexity ratio of the *symmetric* map for several codes presented in the Results section. The ratio is computed with respect to the $\mathcal{O}(n^2)$ legacy self-attention map size. We can see that sparsity can reach up to more than 80% while the computation ratios range from 5 to 35% only. This improvement can lead to a sizable reduction in power and computation time on a *dedicated* hardware.

The complexity of the network is defined by $\mathcal{O}(N(d^2(2n - k) + hd))$, where $h \ll n^2$ denotes the *fixed* number of computations of the self-attention module. The real-world complexity for different codes is explored in Appendix G. As with most neural decoding approaches, while being extremely effective, the proposed method may remain inefficient compared to non-learning solutions, such as [27], in terms of memory requirement, power consumption, and computational resources, limiting their deployment potential. Acceleration of the proposed Transformer via architectural modification or NN acceleration (e.g. pruning, quantization) [33, 19] is left for future work.

In comparison, the existing AR-BP SOTA method of [25] requires $\mathcal{O}(2N(nd_v d_f d_g + n^2 d_v + (n - k)^2 + d_f))$ more operations than Hyper BP[24], where d_v, d_f, d_g are the number of variable nodes, the capacity of the hypernetwork and primary network, respectively. Typical networks f and g are a 128 dimensional FC network with 4 layers and a 16 dimensional FC network with 2 layers respectively. With these hyperparameters, the $nd_v d_f d_g$ part of the complexity is approximately $nd_v 128^2$. Moreover, the method of [25] cannot be parallelized, due to the hypernetwork structure, i.e. each sample defines a different set of parameters. Thus, our network far surpasses the existing state-of-the-art error rates, with much less time complexity. Hyper-BP[24] is slightly less computationally intensive, since it adds a complexity term of $\mathcal{O}(2Nnd_v d_f d_g)$ on top of BP, but this approach as well as the original model-based network of [22] are far less competitive in terms of accuracy.

7 Conclusions

We present a novel Transformer architecture for decoding algebraic block codes. The proposed model allows effective representation of interactions, based on the high-dimensional embedding of the channel output and a code-dependent masking scheme of the self-attention module. The proposed framework is more efficient than standard Transformers, due to the sparse design of existing block codes, making its training and deployment extremely affordable. Even with a limited number of computational blocks, it outperforms popular (neural) decoders on a broad range of code families. We believe this new architecture will allow the development and adoption of new families of performing Transformer-based codes capable of setting new standards in this field.

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A   Self-Attention Maps Visualization

In this section, we present the self-attention maps for various codes obtained at the $E_b/N_0 = 5$ noise level. Figure 8, presents the self-attention maps and the corresponding input for the Hamming BCH(7,4) code with $N = 6, d = 32$, similarly to the results presented in Section 6.1.

For better visual analysis we present in Figure 10 the self-attention maps for two longer codes. Interestingly, in the early stage of the decoding, ECCT seems to focus its processing of the syndrome. At deeper layers towards the final prediction, the focus shifts to the information bits.

Figure 8: Illustration of the input $h(y)$ (magnitude and syndrome concatenation) and of the self-attention maps at different layers of the network, for the Hamming BCH(7,4) code with $N = 6, d = 32$. The self-attention heads are averaged and the white lines denote the separation between the magnitude and syndrome elements.

B   Impact of the Reliability Embedding

Similarly to Section 6.1, we present in Figure 9 the impact of the reliability embedding on the self-attention map on a larger code. We chose the BCH$(31, 16)$ code which is still short enough for providing a clear visualization, and similarly to Section 6.1, the zero codeword is corrupted with additive noise at the first bit (zero bit index), involving a non-zero syndrome at the first parity check equation (i.e. the 32-th element of the embedding). As can be seen, the first bit embedding is impact-less when it is corrupted and then detectable, while its impact on the syndrome embedding is considerably increased. Once the network corrects the bit (last layers(s)), the values return to normal.

Figure 9: Analysis of the scaled embedding values of the first column of the masked self-attention map with corrupted and uncorrupted codeword (CW) for the standardized BCH$(31, 16)$ code, using the $(N, d) = (4, 32)$ architecture.

Figure 10: Illustration of the input $h(y)$ (magnitude and syndrome concatenation) and of the self-attention maps at different layers of the network, for POLAR(62,32) code with $N = 6, d = 32$ (top) and BCH(63,36) code with $N = 4, d = 32$ (bottom). The self-attention heads are averaged and the white lines denote the separation between the magnitude and syndrome elements. One can observe in the first layers that high attention values are assigned to the syndrome part.

C   Impact of the Multi-Head Self-Attention

We present in table 2 the impact of the number of heads in the self-attention mechanism on performance. Using multiple heads is clearly beneficial for the model's performance. This is consistent with the notion that multi-heads enable richer interpretations of the sequence by separating sections of the embedding and learning different aspects of the interplay between related tokens [32].

Table 2: A comparison of the negative natural logarithm of Bit Error Rate (BER) for three normalized SNR values (4,5,6) of our method for different number of self-attention heads h. Higher is better. All the models are $N = 6, d = 32$ ECCT.

Method	$h = 1$			$h = 4$			$h = 8$		
	4	5	6	4	5	6	4	5	6
Polar(64,32)	5.33	7.01	9.06	5.65	7.49	9.89	5.71	7.63	9.94
LDPC(49,24)	5.44	7.62	10.70	5.69	7.96	11.21	5.74	8.13	11.30
BCH(63,36)	4.33	5.76	7.79	4.44	5.94	8.11	4.42	5.91	8.01

D   Comparison with Successive Cancellation List (SCL) Polar Decoder

Table 3 compares the performance of our model to the SOTA SCL Polar decoder [31] for several Polar Codes. The SCL decoder has a time and space complexity of $\mathcal{O}(LN \log N)$ and $\mathcal{O}(LN)$, respectively. We tested the SCL algorithm for $L = \{1, 4\}$ and we sampled *only* 10^5 noisy codewords because of the high complexity and non-parallel application of the SCL algorithm. The proposed *shallow* ECCTs are able to compete and even surpass the SCL for some of the codes and SNRs. Increasing the capacity of the network, especially with more layers, is expected to lead to better results as demonstrated for LDPC codes. Similarly, SCL with bigger lists would obtain improved accuracy.

Table 3: A comparison of the negative natural logarithm of Bit Error Rate (BER) for three normalized SNR values (4,5,6) between the proposed method with $N = 6, d = 128$ and the SOTA SC-L algorithm. Higher is better.

Method	SC-$L = 1$			SC-$L = 4$			ECCT		
	4	5	6	4	5	6	4	5	6
Polar(64,32)	7.28	9.84	12.50	7.55	10.21	13.60	6.99	9.44	12.32
Polar(64,48)	6.16	8.22	11.30	6.17	8.22	11.30	6.36	8.46	11.09
Polar(128,64)	8.42	11.10	-	8.46	11.15	-	5.92	8.64	12.18
Polar(128,86)	7.36	9.76	12.40	7.37	9.76	12.45	6.31	9.01	12.45
Polar(128,96)	6.78	9.13	11.92	6.80	9.14	12.01	6.31	9.12	12.47

E   Non Gaussian Channel

In this Section, we test our framework on a non-gaussian Rayleigh fading channels, whcih are often used for simulating the propagation environment of a signal, e.g., for wireless devices.

In this fading model, the transmission of the codeword $x \in \{0, 1\}^n$ is defined as $y = hx_s + z$, where h is an n-dimensional i.i.d. Rayleigh distributed vector with a scale parameter α, and $z \sim \mathcal{N}(0, \sigma^2 I_n)$.

In our simulations, we assume a *high* scale $\alpha = 1$ in order to easily compare and reproduce the results, while the level of the Gaussian noise and of the testing procedure remains exactly the same as

described in the paper. The overall variance of the transmitted codeword y in the Rauleigh channel is roughly *twice* the AWGN's on the tested SNR range.

The results are presented in Figure 11. As can be observed, our method is still able to learn to decode even under these very noisy fading channels.

Figure 11: BER comparison between BP and the proposed ECCT $N = 6, d = 32$ for the Rayleigh fading channel for various values of SNR for (a) Polar(64,32), (b) BCH(63,36), and (c) LDPC(49,24) codes.

F Robustness to the Modulation

In this Section, we test our framework on a modulation other than the Binary Phase-shift keying. The 16 Quadrature amplitude modulation (16QAM) is common for wireless fading. We simulate the Bit-Interleaved Coded Modulation (BICM) demapping [36, 5] as follows

$$y = Demap\big(Map(x) + z\big), \tag{8}$$

where Map and $Demap$ represent the 4-bits per symbol constellation (16QAM) mapping and demapping, respectively, and z denotes the $n/4$-dimensional *complex* AWGN noise, with noise levels that are defined using the same protocol as in our main experiments. This protocol has been implemented using NVIDIA's Sionna library [13, 14].

The results are presented in Figure 12. Evidently, our model is able to learn accurate decoding under different modulation even on low SNRs.

Figure 12: BER comparison between BP and the proposed ECCT $N = 6, d = 32$ under 16QAM modulation for various values of SNR for (a) Polar(64,32), (b) Polar(64,48), and (c) Polar(128,96) codes.

G Comparison with Augmented Neural BP

In this Section, we compare the complexity and the performance of ECCT with the complexity of the Neural BP (NBP) [22], which does not suffer from the high computational cost of the later hyper-network based models [24].

We recall that the time complexity of the (neural) BP is defined as $\mathcal{O}(L(\sum_i^c d_i^2 + \sum_i^v \tilde{d}_i))$, where d_i denotes the degree of the i-th parity check node, $\tilde{d}_i$ denotes the degree of the i-th variables nodes, and L is the number of iterations.

We present in Table 4 the numerical time complexity values of different ECCTs and the neural BP decoder. The (neural) BP is obtained at convergence with $L = 50$ as reported in [24]. These results are the best results that the neural BP architecture can reach. The inability of the (neural) BP to further improve, was the motivation for the development of more computationally expensive augmented BP methods [24, 25]

The shallower ECCTs (i.e. $N = 2, d = 32$) are comparable to the complexity of BP for most of the codes and is even much more efficient for several codes such as BCH codes (the level of sparsity of each code is the major cause of the difference in complexity). For LDPC codes, for which BP is specially fitted and reliable [29], BP's complexity is much lower than ours. However, as described in Section 6.3, many acceleration methods can greatly improve the complexity. For example, low-rank approximations [33] can reduce the quadratic complexity in d to *linear*.

In the current setting, one can observe that, at similar complexity, $N = 2$ ECCTs are too shallow and thus cannot always reach NBP's accuracies. The most impactful architectural parameter on the reliability of the ECCT is the number of layers N, while the embedding size d has the most influence on the complexity (quadratic). Thus, one can compare, for example, the results for the BCH(63,45) code, for which BP has a similar complexity to the $N = 6, \boldsymbol{d = 16}$ ECCT. The negative logarithm of the BER for $E_b/N_0 \in \{4, 5, 6\}$ is $4.98, 6.77, 9.26$, respectively, for ECCT, while the neural-BP of [22] reaches at convergence $4.49, 6.01, 8.20$ [24], a 12% improvement on average. Similarly, on the Polar(128,96) code, the accuracy of the $N = 6, \boldsymbol{d = 12}, \boldsymbol{h = 4}$ ECCT is $4.98, 6.85, 9.25$ while NBP reaches $4.63, 6.31, 8.54$, an 8% improvement on average.

Table 4: Numerical comparison of the complexity of the (Neural) BP of [22] and of the proposed ECCT for different N and $d = \{32, 64, 128\}$.

Method	NBP[22]	ECCT N=2			ECCT N=6		
	50	32	64	128	32	64	128
POLAR(64,32)	231600	638336	1669888	4912640	1915008	5009664	14737920
POLAR(64,48)	244600	388864	1105408	3521536	1166592	3316224	10564608
POLAR(128,64)	1074400	2236672	5259776	13665280	6710016	15779328	40995840
POLAR(128,86)	1305200	1503104	3702528	10190336	4509312	11107584	30571008
POLAR(128,96)	1244800	1245184	3145728	8912896	3735552	9437184	26738688
LDPC(49,24)	78400	407680	1118464	3449344	1223040	3355392	10348032
LDPC(121,60)	435600	1965696	4676864	12335616	5897088	14030592	37006848
LDPC(121,70)	363000	1803776	4312064	11442176	5411328	12936192	34326528
LDPC(121,80)	290400	1654656	3972864	10599936	4963968	11918592	31799808
MACKAY(96,48)	100800	1483776	3557376	9474048	4451328	10672128	28422144
CCSDS(128,64)	230400	2445312	5677056	14499840	7335936	17031168	43499520
BCH(63,36)	533900	486144	1340928	4156416	1458432	4022784	12469248
BCH(63,45)	413600	362880	1057536	3442176	1088640	3172608	10326528
BCH(63,51)	501000	262144	831488	2891776	786432	2494464	8675328

H Performance on Larger Codes

We present in Figure13 the performance of the proposed ECCT on larger codes.

(a)   (b)

Figure 13: ECCT $N = 6, d = 32$ performance for various values of normalized SNR for (a) Polar(512,384), (b) LDPC (529,440) codes.