Learning on the Edge: Online Learning with Stochastic Feedback Graphs

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Abstract

The framework of feedback graphs is a generalization of sequential decision-making with bandit or full information feedback. In this work, we study an extension where the directed feedback graph is stochastic, following a distribution similar to the classical Erdős-Rényi model. Specifically, in each round every edge in the graph is either realized or not with a distinct probability for each edge. We prove regret bounds of order $\min\left\{\min_\varepsilon \sqrt{(\alpha_\varepsilon/\varepsilon)T},\ \min_\varepsilon (\delta_\varepsilon/\varepsilon)^{1/3}T^{2/3}\right\}$ (ignoring logarithmic factors), where α_ε and δ_ε are graph-theoretic quantities measured on the support of the stochastic feedback graph $\mathcal{G}$ with edge probabilities thresholded at ε. Our result, which holds without any preliminary knowledge about $\mathcal{G}$, requires the learner to observe only the realized out-neighborhood of the chosen action. When the learner is allowed to observe the realization of the entire graph (but only the losses in the out-neighborhood of the chosen action), we derive a more efficient algorithm with a tighter bound featuring a dependence on weighted versions of the independence and weak domination numbers.

1 Introduction

In this work we study an online learning framework for decision-making with partial feedback. In each decision round, the learner chooses an action in a fixed set and is charged a loss. In our setting, the loss of any action in all decision rounds is preliminarily chosen by an adversary, but the feedback received by the learner at the end of each round t is stochastic. More specifically, the loss of each action i (including I_t, the one selected by the learner at round t) is independently observed with a certain probability $p(I_t, i)$, where the probabilities $p(i, j)$ for all pairs i, j are fixed but unknown.

This feedback model can be viewed as a stochastic version of the feedback graph model for online learning [Mannor and Shamir, 2011], where the feedback received by the learner at the end of each round is determined by a directed graph defined over the set of actions. In this model, the learner deterministically observes the losses of all the actions in the out-neighborhood of the action selected in that round. In certain applications, however, deterministic feedback is not realistic. Consider for instance a sensor network for monitoring the environment, where the learner can decide which sensor to probe in order to maximize some performance measure. Each probed sensor may also receive readings of other sensors, but whether a sensor successfully transmits to another sensor depends on a number of environmental factors, which include the position of the two sensors, but also their internal state (e.g., battery levels) and the weather conditions. Due to the variability of some of these factors, the possibility of reading from another sensor can be naturally modeled as a stochastic event.

Submitted to 36th Conference on Neural Information Processing Systems (NeurIPS 2022). Do not distribute.

Online learning with adversarial losses and stochastic feedback graphs has been studied before, but under fairly restrictive assumptions on the probabilities $p(i,j)$. Let $\mathcal{G}$ be a stochastic feedback graph, represented by its probability matrix $p(i,j)$ for $i,j \in V$ where V is the action set. When $p(i,j) = \varepsilon$ for all distinct $i,j \in V$ and for some $\varepsilon > 0$, then $\mathcal{G}$ becomes an Erdős-Rényi random graph model. Under the assumption that ε is known and $p(i,i) = 1$ for all $i \in V$ (all self-loops occur w.p. 1), Alon et al. [2017] show that the optimal regret after T rounds is of order $\sqrt{T/\varepsilon}$, up to logarithmic factors. This result has been extended by Kocák et al. [2016a], who prove a regret bound of order $\sqrt{\sum_t (1/\varepsilon_t)}$ when the parameter ε_t of the random graph is unknown and allowed to change over time. However, their result holds only under rather strong assumptions on the sequence ε_t for $t \geq 1$. In a recent work, Ghari and Shen [2022] show a regret bound of order $(\alpha/\varepsilon)\sqrt{KT}$, ignoring logarithmic factors, when each (unknown) probability $p(i,j)$ in $\mathcal{G}$ is either zero or at least ε for some known $\varepsilon > 0$, and all self-loops (i,i) have probability $p(i,i) \geq \varepsilon$. Here α is the independence number (computed ignoring edge orientations) of the support graph $\mathrm{supp}(\mathcal{G})$; i.e., the directed graph with adjacency matrix $A(i,j) = \mathbb{I}_{\{p(i,j)>0\}}$. Their bound holds under the assumption that $\mathrm{supp}(\mathcal{G})$ is preliminarily known to the learner.

Our analysis does away with a crucial assumption that was key to prove all previous results. Namely, we do not assume any special property of the matrix $\mathcal{G}$, and we do not require the learner to have any preliminary knowledge of this matrix. The fact that positive edge probabilities are not bounded away from zero implies that the learner must choose a threshold $\varepsilon \in (0,1]$ below which the edges are deemed to be too rare to be exploitable for learning. If ε is too small, then waiting for rare edges slows down learning. On the other hand, if ε is too big, then the feedback becomes sparse and the regret increases. To formalize the intuition of rare edges, we introduce the notion of thresholded graph $\mathrm{supp}(\mathcal{G}_\varepsilon)$ for any $\varepsilon > 0$. This is the directed graph with adjacency matrix $A(i,j) = \mathbb{I}_{\{p(i,j)\geq\varepsilon\}}$. As the thresholded graph is a deterministic feedback graph G, we can refer to Alon et al. [2015] for a characterization of minimax regret R_T based on whether G is not observable (R_T of order T), weakly observable (R_T of order $\delta^{1/3}T^{2/3}$), or strongly observable (R_T of order $\sqrt{\alpha T}$).[1] Here α and δ are, respectively, the independence and the weak domination number of G; see Section 2 for definitions. Let α_ε and δ_ε respectively denote the independence number and the weak domination number of $\mathrm{supp}(\mathcal{G}_\varepsilon)$. As α_ε and δ_ε both grow when ε gets larger, the ratios $\alpha_\varepsilon/\varepsilon$ and $\delta_\varepsilon/\varepsilon$ capture the trade-off involved in choosing ε. We define the optimal values for ε as follows:

$$\varepsilon_s^* = \operatorname*{arg\,min}_{\varepsilon\in(0,1]} \left\{ \frac{\alpha_\varepsilon}{\varepsilon} \,:\, \mathrm{supp}(\mathcal{G}_\varepsilon) \text{ is strongly observable} \right\}, \tag{1}$$

$$\varepsilon_w^* = \operatorname*{arg\,min}_{\varepsilon\in(0,1]} \left\{ \frac{\delta_\varepsilon}{\varepsilon} \,:\, \mathrm{supp}(\mathcal{G}_\varepsilon) \text{ is observable} \right\}. \tag{2}$$

We adopt the convention that the minimum of an empty set is infinity and the relative $\arg\min$ is set to 0. The $\arg\min$ are well defined: there are at most K^2 values of ε for which the support of $\mathcal{G}_\varepsilon$ varies, and the minimum is attained in one of these values. For simplicity, we let $\alpha^* = \alpha_{\varepsilon_s^*}$ and $\delta^* = \delta_{\varepsilon_w^*}$. Our first result can be informally stated as follows.

Theorem 1 (Informal)**.** *Consider the problem of online learning with an unknown stochastic feedback graph $\mathcal{G}$ on T time steps. If $\mathrm{supp}(\mathcal{G}_\varepsilon)$ is not observable for $\varepsilon = \tilde{\Theta}(K/T)$, then any learning algorithm suffers regret linear in T. Otherwise, there exists an algorithm whose regret satisfies (ignoring polylog factors in K and T)*

$$R_T \leq \min\left\{ \sqrt{\frac{\alpha^*}{\varepsilon_s^*}T}, \left(\frac{\delta^*}{\varepsilon_w^*}\right)^{1/3} T^{2/3} \right\}.$$

This bound is tight (up to polylog factors).

This result shows that, without any preliminary knowledge of $\mathcal{G}$, we can obtain a bound that optimally trades off between the strongly observable rate $\sqrt{(\alpha^*/\varepsilon_s^*)T}$, for the best threshold ε for which $\mathrm{supp}(\mathcal{G}_\varepsilon)$ is strongly observable, and the (weakly) observable rate $(\delta^*/\varepsilon_w^*)^{1/3}T^{2/3}$,

[1] All these rates ignore logarithmic factors.

for the best threshold ε for which $\mathrm{supp}(\mathcal{G}_\varepsilon)$ is (weakly) observable. Note that this result improves on Ghari and Shen [2022] bound $(\alpha^*/\varepsilon_s^*)\sqrt{KT}$, who additionally assume ε_s and $\mathrm{supp}(\mathcal{G}_{\varepsilon_s})$ are both preliminarily available to the learner. On the other hand, the algorithm achieving the bound of Theorem 1 need not receive any information (neither prior nor during the learning process) besides the stochastic feedback.

We obtain positive results in Theorem 1 via an elaborate reduction to online learning with deterministic feedback graphs. Our algorithm works in two phases: first, it learns the edge probabilities in a round-robin procedure, then it commits to a carefully chosen estimate of the feedback graph and feeds it to an algorithm for online learning with deterministic feedback graphs. There are two main technical challenges the algorithm faces: on the one hand, it needs to switch from the first to the second phase at the right time in order to achieve the optimal regret. On the other hand, in order for the reduction to work, it needs to simulate the behaviour of a deterministic feedback graph using only feedback from a stochastic feedback graph (with unknown edge probabilities). We complement the positive results in Theorem 1 with matching lower bounds that are obtained by a suitable modification of the hard instances in Alon et al. [2015, 2017] so as to consider stochastic feedback graphs.

Our last result is an algorithm that, at the cost of an additional assumption on the feedback (i.e., the learner additionally observes the realization of the entire feedback graph at the end of each round), has regret which is never worse and may be considerably better than the regret of the algorithm in Theorem 1. While the bounds in Theorem 1 are tight up to log factors, we show that the factors α^*/ε_s^* and δ^*/ε_w^* can be improved for specific feedback graphs. Specifically, we design weighted versions of the independence and weak domination numbers, where the weights of a given node depend on the probabilities of seeing the loss of that node. On the technical side, we design a new importance-weighted estimator which uses a particular version of upper confidence bound estimates of the edge probabilities $p(i,j)$, rather than the true edge probabilities, which are unknown. We show that the cost of using this estimator is of the same order as the regret bound achievable had we known $p(i,j)$. Additionally, the algorithm that obtains these improved bounds is more efficient than the algorithm of Theorem 1. The improvement in efficiency comes from the following idea: we start with an optimistic algorithm that assumes that the support of $\mathcal{G}$ is strongly observable and only switches to the assumption that the support of $\mathcal{G}$ is (weakly) observable when it estimates that the regret under this second assumption is smaller. The algorithm learns which regime is better by keeping track of a bound on the regret of the optimistic algorithm while simultaneously estimating the regret in the (weakly) observable case, which it can do efficiently.

Additional related work The problem of adversarial online learning with feedback graphs was introduced by Mannor and Shamir [2011], in the special case where all nodes in the feedback graph have self-loops. The results of Alon et al. [2015] (also based on prior work by Alon et al. [2013], Kocák et al. [2014]) have been recently slightly improved by Chen et al. [2021], with tighter constants in the regret bound. Variants of the adversarial setting have been studied by Feng and Loh [2018], Arora et al. [2019], Rangi and Franceschetti [2019] and Van der Hoeven et al. [2021], who study online learning with feedback graphs and switching costs and online multiclass classification with feedback graphs, respectively. There is also a considerable amount of work in the stochastic setting [Liu et al., 2018, Cortes et al., 2019, Li et al., 2020]. Finally, Erez and Koren [2021] designed a best-of-both worlds learning algorithm achieving nearly optimal (up to polylogarithmic factors in T) regret bounds in the adversarial and stochastic settings. Following Mannor and Shamir [2011], we can consider a more general scenario where the feedback graph is not fixed but changes over time, resulting in a sequence $G_1, \ldots, G_T$ of feedback graphs [Alon et al., 2013, 2017, Kocák et al., 2014, 2016b, Cohen et al., 2016]. Further work takes also into account a time-varying probability for the revelation of side-observations [Kocák et al., 2016a]. While the idea of a general probabilistic feedback graph has been already considered in the stochastic setting (Li et al. [2020], Cortes et al. [2020]), the recent work by Ghari and Shen [2022] seems to be the first one in the adversarial setting that generalizes from the Erdős-Rényi model to a more flexible distribution where they allow "edge-specific" probabilities. We remark, however, that the assumptions of Ghari and Shen [2022] exclude some important instances of feedback graphs. For example, we cannot hope to employ their algorithm for efficiently solving the revealing

action problem (see for example [Alon et al., 2015]). In a spirit similar to ours, Resler and Mansour [2019] studied a version of the problem where the topology of the graph is fixed and known a priori, but the feedback received by the learner is perturbed when traversing edges.

2 Problem Setting

A feedback graph over a set $V = [K]$ of actions is any directed graph $G = (V, E)$, possibly with self-loops. For any vertex $i \in V$, we use $N_G^{\text{in}}(i) = \{j \in V : (j, i) \in E\}$ to denote the in-neighborhood of i and $N_G^{\text{out}}(i) = \{j \in V : (i, j) \in E\}$ to denote its out-neighborhood (we may omit the subscript when the graph is clear from the context). The independence number $\alpha(G)$ of a feedback graph G is the cardinality of the largest subset S of V such that, for all distinct $i, j \in S$, it holds that $(i, j) \notin E$ and $(j, i) \notin E$. We also use the following terminology for directed graphs $G = (V, E)$ [Alon et al., 2015]. Any $i \in V$ is: observable if $N_G^{\text{in}}(i) \neq \emptyset$, strongly observable if $i \in N_G^{\text{in}}(i)$ or $V \setminus \{i\} \subseteq N_G^{\text{in}}(i)$, and weakly observable if it is observable but not strongly. The graph G is: observable if all $i \in V$ are observable, strongly observable if all $i \in V$ are strongly observable, and weakly observable if it is observable but not strongly. The weak domination number $\delta(G)$ of G is the cardinality of the smallest subset S of V such that for all weakly observable $i \in V \setminus S$ there exists $j \in S$ such that $(j, i) \in E$.

In the online learning problem with a stochastic feedback graph, an oblivious adversary privately chooses a stochastic feedback graph $\mathcal{G}$ and a sequence $\ell_1, \ell_2, \ldots$ of loss functions $\ell_t \colon V \to [0, 1]$. At each round $t = 1, 2, \ldots$, the learner selects an action $I_t \in V$ to play and, independently, the adversary draws a feedback graph G_t from $\mathcal{G}$ (denoted by $G_t \sim \mathcal{G}$). The learner then incurs loss $\ell_t(I_t)$ and observes the feedback $\big\{(i, \ell_t(i)) : i \in N_{G_t}^{\text{out}}(I_t)\big\}$. In some cases we consider a richer feedback, where at the end of each round t the learner also observes the realized graph G_t. The learner's performance is measured using the standard notion of regret,

$$R_T = \max_{k \in V} \mathbb{E}\left[\sum_{t=1}^{T} \big(\ell_t(I_t) - \ell_t(k)\big)\right]$$

where $I_1, \ldots, I_T$ are the actions played by the learner, and the expectation is computed over both the sequence $G_1, \ldots, G_T$ of feedback graphs drawn i.i.d. from $\mathcal{G}$ and the learner's internal randomization.

Fix any stochastic feedback graph $\mathcal{G} = \{p(i, j) : i, j \in V\}$. We sometimes use e to denote a pair (i, j), in which case we write p_e to denote the probability $p(i, j)$. When $G_t = (V, E_t)$ is drawn from $\mathcal{G}$, each pair $(i, j) \in V \times V$ independently becomes an edge (i.e., $(i, j) \in E_t$) with probability $p(i, j)$. For any $\varepsilon > 0$, we define the thresholded version $\mathcal{G}_\varepsilon$ of $\mathcal{G}$ represented by $\{p'(i, j) : i, j \in V\}$, where $p'(i, j) = p(i, j)\mathbb{I}_{\{p(i,j) \geq \varepsilon\}}$. We also define the support feedback graph of $\mathcal{G}$ as the graph $\text{supp}(\mathcal{G}) = (V, E)$ having $E = \{(i, j) \in V \times V : p(i, j) > 0\}$. To keep the notation tidy, we write $\alpha(\mathcal{G})$ instead of $\alpha(\text{supp}(\mathcal{G}))$ and similarly for δ. We denote by $e_k \in \mathbb{R}^K$ the basis vector in direction k.

3 Block Decomposition Approach

In this section, we present an algorithm for online learning with stochastic feedback graphs via a reduction to online learning with deterministic feedback graphs. Our algorithm EdgeCatcher (Algorithm 3) has an initial exploration phase followed by a commit phase. In the exploration phase, the edge probabilities are learned online in a round-robin fashion. A carefully designed stopping criterion then triggers the commit phase, where we feed the support of the estimated stochastic feedback graph to an algorithm for online learning algorithm with (deterministic) feedback graphs.

3.1 Estimating the Edge Probabilities

As a first step we design a routine, RoundRobin (Algorithm 1), that sequentially estimates the stochastic feedback graph until a certain stopping criterion is met. The stopping criterion depends on a nonnegative function Φ that takes in input a stochastic feedback graph $\mathcal{G}$ together with a time horizon. Let $\hat{\tau} \leq T/K$ be the index of the last iteration of the outer

Algorithm 1: RoundRobin

Environment: stochastic feedback graph $\mathcal{G}$, sequence of losses $\ell_1, \ell_2, \dots, \ell_T$;
Input: time horizon T, stopping function Φ, actions $V = \{1, 2, \dots, K\}$;
$n_e \leftarrow 0$, for all $e \in V^2$;
for each $\tau = 1, 2, \dots, \lfloor T/K \rfloor$ **do**
  for each $i = 1, 2, \dots K$ **do**
    Play action i and observe $N^{\text{out}}_{G_t}(i)$ from $G_t \sim \mathcal{G}$; `// t is the time step`
    $n_e \leftarrow n_e + 1$ for all $e \in N^{\text{out}}_{G_t}(i)$;
  $\hat{p}^\tau_e \leftarrow n_e/\tau$ for all edges $e \in V^2$;
  $\varepsilon_\tau \leftarrow 60 \ln(KT)/\tau$;
  $\hat{\mathcal{G}}_\tau \leftarrow \big(V, \{e \in V^2 : \hat{p}^\tau_e \geq \varepsilon_\tau\}\big)$ with weights $\hat{p}^\tau_e$; `// estimated feedback graph`
  if $\Phi(\hat{\mathcal{G}}_\tau, T) \leq \tau K$ **then**
    output $\hat{\mathcal{G}}_\tau$, ε_τ;
output $\hat{\mathcal{G}}_\tau$, ε_τ;

for loop in Algorithm 1. We want to make sure that, for all $\tau \leq \hat{\tau}$, the stochastic feedback graphs $\hat{\mathcal{G}}_\tau$ are valid estimates of the underlying $\mathcal{G}$ up to a $\Theta(\varepsilon_\tau)$ precision. To formalize this notion of approximation, we introduce the following definition.

Definition 1 (ε-good approximation)**.** *A stochastic feedback graph $\hat{\mathcal{G}} = \{\hat{p}_e : e \in V^2\}$ is an ε-good approximation of $\mathcal{G} = \{p_e : e \in V^2\}$ for some $\varepsilon \in (0, 1]$, if the following holds:*

1. *All the edges $e \in \text{supp}(\mathcal{G})$ with probability $p_e \geq 2\varepsilon$ belong to $\text{supp}(\hat{\mathcal{G}})$;*
2. *For all edges $e \in \text{supp}(\hat{\mathcal{G}})$ with $p_e \geq \varepsilon/2$ it holds that $|\hat{p}_e - p_e| \leq p_e/2$;*
3. *No edge $e \in V^2$ with probability $p_e < \varepsilon/2$ belongs to $\text{supp}(\hat{\mathcal{G}})$.*

We can now state the following theorem; we defer the proof in Appendix A. The proof follows from an application of the multiplicative Chernoff bound on edge weights.

Theorem 2. *If* RoundRobin *(Algorithm 1) is run on the stochastic feedback graph $\mathcal{G}$, then, with probability at least $1 - 1/T$, the estimate $\hat{\mathcal{G}}_\tau$ is an ε_τ-good approximation of $\mathcal{G}$ for all $\tau \leq \hat{\tau}$, where $\hat{\tau} \leq T/K$ is the index of the last iteration of the outer for loop in Algorithm 1.*

3.2 Block Decomposition: Reduction to Deterministic Feedback Graph

As a second step, we present BlockReduction (Algorithm 2) which reduces the problem of online learning with stochastic feedback graph to the corresponding problem with deterministic feedback graph. Surprisingly enough, in order for this reduction to work, we do not need the exact edge probabilities: an ε-good approximation is sufficient for this purpose.

The intuition behind BlockReduction is simple: given that each edge e in $\text{supp}(\mathcal{G}_\varepsilon)$ appears in G_t with probability $p_e \geq \varepsilon$ at each time step t, if we wait for $O\big((1/\varepsilon)\ln(T)\big)$ time steps it will appear at least once with high probability. Applying a union bound over all edges, we can argue that considering $\Delta = O\big((1/\varepsilon)\ln(KT)\big)$ realizations of the stochastic feedback graph, then all the edges in $\text{supp}(\mathcal{G}_\varepsilon)$ are realized at least once with high probability.

Imagine now to play a certain action a consistently during a block B_τ of Δ time steps. We want to reconstruct the average loss suffered by a' in B_τ:

$$c_\tau(a') = \sum_{t \in B_\tau} \frac{\ell_t(a')}{\Delta} \quad , \tag{3}$$

and we want to do it for all a' in the out-neighborhood of a. Let $\Delta^\tau_{(a,a')}$ be the number of times that the loss of a' is observed by the learner; i.e., the number of times that (a, a') is realized in the Δ time steps. With this notation, we can define the natural estimator $\hat{c}_\tau(a')$:

$$\hat{c}_\tau(a') = \sum_{t \in B_\tau} \ell_t(a') \frac{\mathbb{I}_{\{(a,a') \in E_t\}}}{\Delta^\tau_{(a,a')}} \quad . \tag{4}$$

Algorithm 2: BlockReduction

Environment: stochastic feedback graph $\mathcal{G}$, sequence of losses $\ell_1, \ell_2, \dots, \ell_T$;

Input: time horizon T, threshold ε, estimate $\hat{\mathcal{G}}$ of $\mathcal{G}$, learning algorithm $\mathcal{A}$;

$\Delta \leftarrow \lceil \frac{2}{\varepsilon} \ln(KT) \rceil$, $N \leftarrow \lfloor T/\Delta \rfloor$, $\hat{G} \leftarrow \mathrm{supp}(\hat{\mathcal{G}})$;

Initialize $\mathcal{A}$ with time horizon N and graph $\hat{G}$;

$B_\tau \leftarrow \{(\tau-1)\Delta + 1, \dots, \tau\Delta\}$, for all $\tau = 1, \dots, N$;

for each round $\tau = 1, 2, \dots, N$ **do**

  Draw action a_τ from the probability distribution over actions output by $\mathcal{A}$;

  for each round $t \in B_\tau$ **do**

    Play action a_τ and observe the revealed feedback; // $G_t \sim \mathcal{G}$

  For all $a' \in N^{\mathrm{out}}_{\hat{G}}(a_\tau)$, compute $\hat{c}_\tau(a')$ as in (4), and feed them to $\mathcal{A}$;

Play arbitrarily the remaining $T - \Delta N$ rounds;

Conditioning on the event $\mathcal{E}^\tau_{(a,a')}$ that the edge (a, a') in $\hat{G}$ is observed at least once in block B_τ, we show in Lemma 2 in Appendix A that $\hat{c}_\tau(a')$ is an unbiased estimator of $c_\tau(a')$.

Therefore, we can construct unbiased estimators of the average of the losses on the blocks as if the stochastic feedback graph were deterministic. This allows us to reduce the original problem to that of online learning with deterministic feedback graph on the meta-instance given by the blocks. The details of BlockReduction are reported in Algorithm 2, while the theoretical properties are summarized in the next result, whose proof can be found in Appendix A.

Theorem 3. *Consider the problem of online learning with stochastic feedback graph $\mathcal{G}$, and let $\hat{\mathcal{G}}$ be an ε-good approximation of $\mathcal{G}$. Let $\mathcal{A}$ be an algorithm for online learning with arbitrary deterministic feedback graph G with regret bound $R^{\mathcal{A}}_N(G)$ over any sequence of losses in $[0, 1]$. Then the regret of* BlockReduction *(Algorithm 2) run with input $(T, \varepsilon/2, \hat{\mathcal{G}}, \mathcal{A})$ is at most $\Delta R^{\mathcal{A}}_N\big(\mathrm{supp}(\hat{\mathcal{G}})\big) + \Delta$, where $N = \lfloor T/\Delta \rfloor$ and $\Delta = \lceil \frac{4}{\varepsilon} \ln(KT) \rceil$.*

For online learning with deterministic feedback graphs we use the variants of Exp3.G contained in Alon et al. [2015]. Together with Theorem 3, this gives the following corollary; the details of the proof are in Appendix A.

Corollary 1. *Consider the problem of online learning with stochastic feedback graph $\mathcal{G}$, and let $\hat{\mathcal{G}}$ be an ε-good approximation of $\mathcal{G}$ for $\varepsilon \geq 1/T$ and with support $\hat{G}$.*

If $\hat{G}$ is strongly observable with independence number α, then the regret of BlockReduction *run with parameter $\varepsilon/2$ using* Exp3.G *for strongly observable graphs as base algorithm $\mathcal{A}$ satisfies: $R_T \leq 4C_s \sqrt{(\alpha/\varepsilon)T}\big(\ln(KT)\big)^{3/2}$, where $C_s > 0$ is a constant in the regret bound of $\mathcal{A}$.*

If $\hat{G}$ is weakly observable with weak domination number δ, then the regret of BlockReduction *run with parameter $\varepsilon/2$ using* Exp3.G *for weakly observable graphs as base algorithm $\mathcal{A}$ satisfies: $R_T \leq 4C_w \left((\delta/\varepsilon)\big(\ln(KT)\big)^2\right)^{1/3} T^{2/3}$, where $C_w > 0$ is a constant in the regret bound of $\mathcal{A}$.*

3.3 Explore then Commit to a Graph

We are now ready to combine the two routines we presented, RoundRobin and BlockReduction, in our final learning algorithm, EdgeCatcher. EdgeCatcher has two phases: in the first phase, RoundRobin is used to quickly obtain an ε-good approximation $\hat{\mathcal{G}}$ of the underlying feedback graph $\mathcal{G}$, for a suitable ε. In the second phase, the algorithm commits to $\hat{\mathcal{G}}$ and feeds it to BlockReduction. The crucial point is when to commit to a certain (estimated) stochastic feedback graph. If we commit too early, we might not observe a denser support graph, which implies missing out on a richer feedback. If we wait for too long, then the exploration phase ends up dominating the regret. To balance this trade-off, we use the

Algorithm 3: EdgeCatcher

Environment: stochastic feedback graph $\mathcal{G}$, sequence of losses $\ell_1, \ell_2, \dots, \ell_T$;
Input: time horizon T and actions $V = \{1, 2, \dots, K\}$;
Let Φ defined as in Equation (5);
Run RoundRobin(T, Φ, V) and obtain $\hat{\mathcal{G}}$ and $\hat{\varepsilon}$;
Compute $\hat{\varepsilon}_s^*$ and $\hat{\varepsilon}_w^*$ for graph $\hat{\mathcal{G}}$ as in Equations (1) and (2);
Let $\hat{\varepsilon}^*$ be the best threshold as in Equation (5) and $G = \mathrm{supp}(\hat{\mathcal{G}}_{\hat{\varepsilon}^*})$;
if $\hat{\varepsilon}^* = \hat{\varepsilon}_s^*$ **then** Let $\mathcal{A}$ be Exp3.G for strongly observable feedback graph;
else Let $\mathcal{A}$ be Exp3.G for weakly observable feedback graph;
Let $T' = T - \hat{\tau}K$ be the remaining time steps; `// $\hat{\tau}$ as in` RoundRobin
Run BlockReduction$(T', \hat{\varepsilon}^*/2, G, \mathcal{A})$;

stopping function Φ. This function takes as input a probabilistic feedback graph together with a time horizon and outputs the regret bound that BlockReduction would guarantee on this pair. It is defined as

$$\Phi(\mathcal{G}, T) = \min\left\{4C_s\sqrt{\frac{\alpha^*}{\varepsilon_s^*}T}\big(\ln(KT)\big)^{3/2},\ 4C_w\left(\frac{\delta^*}{\varepsilon_w^*}\big(\ln(KT)\big)^2\right)^{1/3}T^{2/3}\right\}\ . \tag{5}$$

Note that the dependence of Φ on the feedback graph $\mathcal{G}$ is contained in the topological parameters α^* and δ^* and the corresponding thresholds ε_s^* and ε_w^*, defined in Equations (1) and (2). If we apply Φ to a stochastic feedback graph that is observable w.p. zero, its value is conventionally set to infinity. Note that, otherwise, the min is achieved for a specific ε^* and a specific $\mathcal{G}^* = \mathcal{G}_{\varepsilon^*}$. In Appendix C, we provide a sequence of lemmas (Lemmas 3 and 4 in particular) showing that, if RoundRobin outputs an ε-good approximation of the graph, then the regret is bounded by a multiple of the stopping criterion evaluated at $\mathcal{G}$. Combined with Theorem 2, which tells us that RoundRobin does in fact output an ε-good approximation of the graph with high probability, this proves our main result for this section:

Theorem 4. *Consider the problem of online learning with stochastic feedback graph $\mathcal{G}$ on T time steps. If $\mathrm{supp}(\mathcal{G}_{\varepsilon(K,T)})$ is observable for $\varepsilon(K,T) = CK(\ln(KT))^2/T$ for a given constant $C > 0$, then there exists an algorithm whose regret R_T satisfies (ignoring polylog factors in K and T): $R_T \le \min\left\{\sqrt{(\alpha^*/\varepsilon_s^*)T}, \left(\delta^*/\varepsilon_w^*\right)^{1/3}T^{2/3}\right\}$.*

4 Lower Bounds

In this section, we provide lower bounds that match the regret bound guaranteed by EdgeCatcher, shown in Theorem 4, up to polylogarithmic factors in K and T. These lower bounds are valid even if the learner is allowed to observe the realization of the entire feedback graph at every time step, and knows a priori the "correct" threshold ε to work with. Theorem 5 summarizes the lower bounds whose proofs can be found in Appendix B.

Theorem 5 (Informal)**.** *Let $\mathcal{A}$ be a possibly randomized algorithm for the online learning problem with stochastic feedback graphs. Consider any directed graph $G = (V, E)$ with $|V| \ge 2$ and any $\varepsilon \in (0, 1]$. There exists a stochastic feedback graph $\mathcal{G}$ with $\mathrm{supp}(\mathcal{G}) = G$ and, for a sufficiently large time horizon T, there is a sequence $\ell_1, \dots, \ell_T$ of loss functions on which the expected regret of $\mathcal{A}$ with respect to the stochastic generation of $G_1, \dots, G_T \sim \mathcal{G}$ is*

- $\Omega(\sqrt{(\alpha(\mathcal{G}_\varepsilon)/\varepsilon)T})$ *if G is strongly observable;*
- $\tilde{\Omega}((\delta(\mathcal{G}_\varepsilon)/\varepsilon)^{1/3}T^{2/3})$ *if G is weakly observable;*
- $\Omega(T)$ *if G is not observable.*

The lower bound in the non-observable case is the same as Alon et al. [2015, Theorem 6] with a deterministic feedback graph. The remaining lower bounds are nontrivial adaptations of the corresponding bounds for the deterministic case by Alon et al. [2015, 2017]. The main technical hurdle is due to the stochastic nature of the feedback graph, which needs

to be taken into account in the proofs. The rationale behind the constructions used for proving the lower bounds is as follows: since each edge is realized only with probability ε, any algorithm requires $1/\varepsilon$ rounds in expectation in order to observe the loss of an action in the out-neighborhood of the played action, whereas one round would suffice with a deterministic feedback graph. Note that Theorem 5 implies that, if $\mathrm{supp}(\mathcal{G}_{\varepsilon(K,T)})$ is not observable for $\varepsilon(K,T)$ as in Theorem 4, then there is no hope to achieve sublinear regret, as the lower bounds for both strongly and weakly observable supports are linear in T for all $\varepsilon \leq \varepsilon(K,T)$.

5 Refined Graph-Theoretic Parameters

Although the results from Section 3 are worst-case optimal up to log factors, we may find that the factors $\sqrt{\alpha(G_\varepsilon)/\varepsilon}$ and $(\delta(G_\varepsilon)/\varepsilon)^{1/3}$ for strongly and weakly observable $\mathrm{supp}(\mathcal{G}_\varepsilon) = G_\varepsilon$, respectively, may be improved upon in certain cases. In particular, we show that, under additional assumptions on the feedback that we receive, we can obtain better regret bounds. To understand our results, we need some initial definitions. The weighted independence number for a graph $H = (V, E)$ and positive vertex weights $w(i)$ for $i \in V$ is defined as

$$\alpha_{\mathsf{w}}(H, w) = \max_{S \in \mathcal{I}(H)} \sum\nolimits_{i \in S} w(i) \ ,$$

where $\mathcal{I}(H)$ denotes the family of independent sets of H. We consider two different weight assignments computed in terms of any stochastic feedback graph $\mathcal{G}$ with edge probabilities $p(i,j)$ and $\mathrm{supp}(\mathcal{G}) = G$. They are defined as $w^-_{\mathcal{G}}(i) = \big(\min_{j \in N^{\mathrm{in}}_G(i)} p(j,i)\big)^{-1}$ and $w^+_{\mathcal{G}}(i) = \big(\min_{j \in N^{\mathrm{out}}_G(i)} p(i,j)\big)^{-1}$. Then, the two corresponding weighted independence numbers are $\alpha^-_{\mathsf{w}}(\mathcal{G}) = \alpha_{\mathsf{w}}(G, w^-_{\mathcal{G}})$ and $\alpha^+_{\mathsf{w}}(\mathcal{G}) = \alpha_{\mathsf{w}}(G, w^+_{\mathcal{G}})$. The parameter of interest for the results in this section is $\alpha_{\mathsf{w}}(\mathcal{G}) = \alpha^-_{\mathsf{w}}(\mathcal{G}) + \alpha^+_{\mathsf{w}}(\mathcal{G})$. For more details on the weighted independence number see Appendix D. We also use the following definitions of weighted weak domination number δ_{w} for a graph H and positive vertex weights w, and self-observability parameter σ:

$$\delta_{\mathsf{w}}(H, w) = \min_{D \in \mathcal{D}(H)} \sum\nolimits_{i \in D} w(i) \ , \qquad \sigma(\mathcal{G}) = \sum\nolimits_{i : i \in N^{\mathrm{in}}_G(i)} (p(i,i))^{-1} \ ,$$

where $\mathcal{D}(H)$ denotes the family of weakly dominating sets in H. In this section, we focus on the weighted weak domination number $\delta_{\mathsf{w}}(\mathcal{G}) = \delta_{\mathsf{w}}(G, w^+_{\mathcal{G}})$. To gain some intuition on the graph-theoretic parameters introduced above, consider the graph with only self-loops, also used in Example 1 below. If all $p(i,i) = \varepsilon$, the learner needs to pull a single arm $1/\varepsilon$ times for one observation in expectation, and K/ε times to see the losses of all arms once. However, when the edge probabilities are different we need to pull arms for $\sum_{i=1}^K 1/p(i,i)$ times. The weighted independence number, weighted weak domination and self-observability generalize this intuition and tell us how many observations the learner needs to see all losses at least once in expectation. We now state the main result of this section.

Theorem 6 (Informal)**.** *There exists an algorithm with per-round running time of $O(K^4)$ and whose regret is bounded (ignoring logarithmic factors) by*

$$\min\Bigg\{T, \min_{\varepsilon}\Big\{\sqrt{T\alpha_{\mathsf{w}}(\mathcal{G}_\varepsilon)} \ : \ \mathrm{supp}(\mathcal{G}_\varepsilon) \text{ is strongly observable}\Big\}, \\ \min_{\varepsilon}\Big\{T^{2/3}\left(\delta_{\mathsf{w}}(\mathcal{G}_\varepsilon)\right)^{1/3} + \sqrt{\sigma(\mathcal{G}_\varepsilon)T} \ : \ \mathrm{supp}(\mathcal{G}_\varepsilon) \text{ is observable}\Big\}\Bigg\} \ ,$$

The regret bound in Theorem 6 follows from Theorem 11 in Appendix C. The running time bound is determined by approximating δ_{w} for all K^2 possible thresholds. In each of the thresholded graphs, we can compute a $(\ln(K)+1)$-approximation for the weighted weak domination number in $O(K^2)$ time by reduction to set cover [Vazirani, 2001]. Doing so only introduces an extra factor of order $(\ln(K))^{1/3}$ in the regret bound.

An important property of the bound in Theorem 6 is that it is never worse than the bounds obtained before. The following example shows that the regret bound in Theorem 6 can also be better than previously obtained regret bounds.

Example 1. Consider a stochastic feedback graph $\mathcal{G}$ for the K-armed bandit setting: $p(i,i) = \varepsilon_i \in (0,1]$ for all $i \in V$ and $p(i,j) = 0$ for all $i \neq j$. In this case, the regret of EdgeCatcher is $\tilde{O}\big(\sqrt{KT/(\min_i \varepsilon_i)}\big)$. On the other hand, Theorem 6 provides the bound $\tilde{O}\big(\sqrt{T\sum_i(1/\varepsilon_i)}\big)$, as $\alpha_{\mathsf{w}}(\mathcal{G}) = 2\sum_i 1/\varepsilon_i$. In the special case when $\varepsilon_i = \varepsilon \in (0,1]$ for some $i \in V$ while $\varepsilon_j = 1$ for all $j \neq i$, then the regret of EdgeCatcher is $\tilde{O}(\sqrt{KT/\varepsilon})$, while Theorem 6 guarantees a $\tilde{O}(\sqrt{(K+1/\varepsilon)T})$ regret bound. □

To derive these tighter bounds we exploit the additional assumption that the realized feedback graph G_t is observed at the end of each round. This allows us to *simultaneously* estimate the feedback graph and control the regret rather than perform these two tasks sequentially as was done in Section 3. In particular, we use this extra information to construct a novel importance weighted estimator for the loss, which for $t > 1$ is defined to be:

$$\tilde{\ell}_t(i) = \frac{\ell_t(i)}{\hat{P}_t(i)} \mathbb{I}_{\{i \in N^{\text{out}}_{G_t}(I_t)\} \cap \{(I_t,i) \in \hat{H}_t\}} \ , \tag{6}$$

where $\hat{P}_t(i) = \sum_{j \in N^{\text{in}}_{\hat{G}_t}(i)} \pi_t(j)\hat{p}_t(j,i)$ is the estimated probability of observing the loss of arm i in round t, $\pi_t(i)$ is the distribution we sample I_t from, and $\hat{G}_t$ is the support of the estimated graph $\hat{\mathcal{G}}_t$. Note that we ignore losses that we receive due to missing edges in $\hat{G}_t$. We show that we pay an additive term in the regret for wrongly estimating an edge, which is why it is important to control which edges are in $\hat{G}_t$. Ideally, we would use $P_t(i) = \sum_{j \in N^{\text{in}}_{\hat{G}_t}(i)} \pi_t(j)p(j,i)$ rather than $\hat{P}_t(i)$, as this is the true probability of observing the loss of arm i in round t. However, since we do not have access to $p(j,i)$ we use instead an upper confidence estimate of $p(j,i)$ for rounds $t \geq 2$:

$$\hat{p}_t(j,i) = \tilde{p}_t(j,i) + \sqrt{\frac{2\tilde{p}_t(j,i)}{t-1}\ln(3K^2T^2)} + \frac{3}{t-1}\ln(3K^2T^2) \ ,$$

where $\tilde{p}_t(j,i) = \frac{1}{t-1}\sum_{s=1}^{t-1} \mathbb{I}_{\{(j,i) \in E_s\}}$. We denote by $\hat{\mathcal{G}}_t^{\text{UCB}}$ the stochastic graph with edge probabilities $\hat{p}_t(j,i)$. Note that the support of $\hat{\mathcal{G}}_t^{\text{UCB}}$ is a complete graph because $\hat{p}_t(j,i) > 0$ for all $(j,i) \in V \times V$. These estimators for the edge probabilities are sufficiently good for our purposes whenever event $\mathcal{K}$ occurs, which we define as the event that, for all $t \geq 2$,

$$|\tilde{p}_t(j,i) - p(j,i)| \leq \sqrt{\frac{2\tilde{p}_t(j,i)}{t-1}\ln(3K^2T^2)} + \frac{3}{t-1}\ln(3K^2T^2), \quad \forall (j,i) \in V \times V \ .$$

An important property of $\tilde{\ell}_t$ can be found in Lemma 1 below, which tells us that we may treat $\tilde{\ell}_t$ as if event $\mathcal{K}$ is always realized, i.e., $\hat{p}_t(j,i)$ is always an upper bound estimator on $p(j,i)$. The proof of Lemma 1 is implied by Lemma 6 in Appendix C.

Lemma 1 (Informal)**.** *Let e_k denote the basis vector with $e_k(i) = \mathbb{I}_{\{i=k\}}$ as i-th entry for each $i \in [K]$. The loss estimate $\tilde{\ell}_t$ defined in* (6) *satisfies*

$$R_T = \tilde{O}\left(\mathbb{E}\left[\sum_{t=2}^{T}\sqrt{\sum_{i=1}^{K}\frac{\pi_t(i)}{(t-1)\hat{P}_t(i)}} \;\middle|\; \mathcal{K}\right] + \max_{k \in V}\mathbb{E}\left[\sum_{t=2}^{T}\sum_{i=1}^{K}\big(\pi_t(i) - e_k(i)\big)\tilde{\ell}_t(i) \;\middle|\; \mathcal{K}\right]\right). \tag{7}$$

Lemma 1 shows that we only suffer $\tilde{O}\Big(\sqrt{\sum_{t=2}^{T}\sum_{i=1}^{K}\frac{\pi_t(i)}{P_t(i)}}\Big)$ additional regret compared to when we know $p(j,i)$. Lemma 1 also shows that $\tilde{\ell}_t$ behaves nicely in the sense that, conditioned on $\mathcal{K}$, we have $\tilde{\ell}_t(i) \leq \frac{\ell_t(i)}{P_t(i)}\mathbb{I}_{\{i \in N^{\text{out}}_{G_t}(I_t)\} \cap \{(I_t,i) \in \hat{H}_t\}}$. This is a crucial property to bound the regret of our algorithm. We show that, with a modified version of Exp3.G, the second sum on the right-hand side of (7) is bounded by a term of order $\sqrt{\sum_{t=2}^{T}\sum_{i=1}^{K}\frac{\pi_t(i)}{P_t(i)}}$, meaning that the regret is also bounded similarly. Our final step is to prove that the above term is bounded in terms of the minimum of the weighted independence number and the weighted weak domination number plus self-observability.

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Checklist

1. For all authors...
   (a) Do the main claims made in the abstract and introduction accurately reflect the paper's contributions and scope? [Yes] All the results claimed in the abstract have proofs either in the main body or appendix
   (b) Did you describe the limitations of your work? [Yes] For each result we state the parameters under which the result holds.
   (c) Did you discuss any potential negative societal impacts of your work? [No] Our work is primarily theoretical and we do not foresee any potential negative societal impacts of our work.
   (d) Have you read the ethics review guidelines and ensured that your paper conforms to them? [Yes]
2. If you are including theoretical results...
   (a) Did you state the full set of assumptions of all theoretical results? [Yes]
   (b) Did you include complete proofs of all theoretical results? [Yes] Although some proofs may only appear in the appendix.
3. If you ran experiments...
   (a) Did you include the code, data, and instructions needed to reproduce the main experimental results (either in the supplemental material or as a URL)? [N/A]
   (b) Did you specify all the training details (e.g., data splits, hyperparameters, how they were chosen)? [N/A]
   (c) Did you report error bars (e.g., with respect to the random seed after running experiments multiple times)? [N/A]
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4. If you are using existing assets (e.g., code, data, models) or curating/releasing new assets...
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(a) Did you include the full text of instructions given to participants and screenshots, if applicable? [N/A]

(b) Did you describe any potential participant risks, with links to Institutional Review Board (IRB) approvals, if applicable? [N/A]

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