Communication Efficient Distributed Learning for Kernelized Contextual Bandits

Abstract

We tackle the communication efficiency challenge of learning kernelized contextual
bandits in a distributed setting. Despite the recent advances in communication-
efficient distributed bandit learning, existing solutions are restricted to simple
models like multi-armed bandits and linear bandits, which hamper their practical
utility. In this paper, instead of assuming the existence of a linear reward mapping
from the features to the expected rewards, we consider non-linear reward map-
pings, by letting agents collaboratively search in a reproducing kernel Hilbert space
(RKHS). This introduces significant challenges in communication efficiency as
distributed kernel learning requires the transfer of raw data, leading to a communi-
cation cost that grows linearly w.r.t. time horizon T. We addresses this issue by
equipping all agents to communicate via a common Nyström embedding that gets
updated adaptively as more data points are collected. We rigorously proved that
our algorithm can attain sub-linear rate in both regret and communication cost.

1 Introduction

Contextual bandit algorithms have been widely used for a variety of real-world applications, including
recommender systems [13], display advertisement [14] and clinical trials [8]. While most existing
bandit solutions assume a centralized setting (i.e., all the data reside in and all the actions are
taken by a central server), there is increasing research effort on distributed bandit learning lately
[21, 7, 16, 10, 12], where N clients, under the coordination of a central server, collaborate to minimize
the overall cumulative regret incurred over a finite time horizon T. In many distributed application
scenarios, communication is the main bottleneck, e.g., communication in a network of mobile devices
can be slower than local computation by several orders of magnitude [9]. Therefore, it is vital
for distributed bandit learning algorithms to attain sub-linear rate (w.r.t. time horizon T) in both
cumulative regret and communication cost.

However, prior works in this line of research are restricted to linear models [21], which could oversim-
plify the problem and thus leads to inferior performance in practice. In centralized setting, kernelized
bandit algorithms, e.g., KernelUCB [20] and IGP-UCB [5], are proposed to address this issue by
modeling the unknown reward mapping as a non-parametric function lying in a reproducing kernel
Hilbert space (RKHS), i.e., the expected reward is linear w.r.t. an action feature map of possibly
infinite dimensions. Despite the strong modeling capability of kernel method, collaborative explo-
ration in the RKHS gives rise to additional challenges in designing a communication efficient bandit
algorithm. Specifically, unlike distributed linear bandit where the clients can simply communicate
the $d \times d$ sufficient statistics [21], where d is the dimension of the input feature vector, the *joint
kernelized estimation* of the unknown reward function requires communicating either 1) the $p \times p$
sufficient statistics in the RKHS, where p is the dimension of the RKHS that is possibly infinite, or 2)
the set of input feature vectors that grows linearly w.r.t. T. Neither of them is practical due to the
huge communication cost.

In this paper, we propose the first communication efficient algorithm for distributed kernel bandits, which tackles the aforementioned challenge via a low-rank approximation of the empirical kernel matrix. In particular, we extended the Nyström method [15] to distributed learning for kernelized contextual bandits. In this solution, all clients first project their local data to a finite RKHS spanned by a common dictionary, i.e., a small subset of the original dataset, and then they only need to communicate the embedded statistics for collaborative exploration. To ensure effective regret reduction after each communication round, as well as ensuring the dictionary remains representative for the entire distributed dataset throughout the learning process, the frequency of dictionary update and synchronization of embedded statistics is controlled by measuring the amount of new information each client has gained since last communication. We rigorously prove that the proposed algorithm incurs an $O(N^2\gamma_{NT}^3)$ communication cost, where γ_{NT} is the maximum information gain that is known to be $O\big(\log(NT)\big)$ for kernels with exponentially decaying eigenvalues, which includes the most commonly used Gaussian kernel, while attaining the optimal $O(\sqrt{NT}\gamma_{NT})$ cumulative regret.

2 Related Works

To balance exploration and exploitation in stochastic linear contextual bandits, LinUCB algorithm [13, 1] is commonly used, which selects arm optimistically w.r.t. a constructed confidence set on the unknown linear reward function. By using kernels and Gaussian processes, studies in [18, 20, 5] further extend UCB algorithms to non-parametric reward functions in RKHS, i.e., the feature map associated with each arm is possibly infinite.

Recent years have witnessed increasing research efforts in distributed bandit learning, where multiple agents collaborate in minimizing the cumulative regret. They mainly differ in the relations of learning problems solved by the agents (i.e., homogeneous vs., heterogeneous) and the type of communication network (i.e., peer-to-peer (P2P) vs., star-shaped). Most of these works assume linear reward functions, and the clients communicate by transferring the $O(d^2)$ sufficient statistics. Korda et al. [11] considered a peer-to-peer (P2P) communication network and assumed that the clients form clusters, i.e., each cluster is associated with a unique bandit problem. Huang et al. [10] considered a star-shaped communication network as in our paper, but their proposed phase-based elimination algorithm only works in the fixed arm set setting. The closest works to ours are [21, 7, 12], which proposed event-triggered communication protocols to obtain sub-linear communication cost over time for distributed linear bandits with a time-varying arm set. In comparison, distributed kernelized contextual bandits still remain under-explored. The only existing work in this direction [6] considered heterogeneous agents, where each agent is associated with an additional feature describing the task similarity between agents. However, they assumed a local communication setting, where the agent immediately shares the new raw data point to its neighbors after each interaction, and thus the communication cost is still linear over time.

Another closely related line of works is kernelized bandits with approximation, where Nyström method is adopted to improve computation efficiency in a centralized setting. Calandriello et al. [2] proposed an algorithm named BKB, which uses Ridge Leverage Score sampling (RLS) to re-sample a new dictionary from the updated dataset after each interaction with the environment. A recent work by Zenati et al. [22] further improved the computation efficiency of BKB by adopting an online sampling method to update the dictionary. However, both of them updated the dictionary at each time step to ensure the dictionary remains representative w.r.t. the growing dataset, and therefore are not applicable to our problem. This is because the dataset is stored cross clients in a distributed manner, and projecting the dataset to the space spanned by the new dictionary requires communication with all clients, which is prohibitively expensive in terms of communication. Calandriello et al. [3] also proposed a variant of BKB, named BBKB, for batched Gaussian process optimization. BBKB only needs to update the dictionary occasionally according to an adaptive schedule, and thus partially addresses the issue mentioned above. However, as BBKB works in a centralized setting, their adaptive schedule can be computed based on the whole batch of data, while in our decentralized setting, each client can only make the update decision according to the data that is locally available. Moreover, in BBKB, all the interactions are based on a fixed model estimation over the whole batch, which is mentioned in their Appendix A.4 as a result of an inherent technical difficulty. In comparison, our proposed method effectively addresses this difficulty with improved analysis, and thus allows each client to utilize newly collected data to update its model estimation on the fly.

3 Preliminaries

In this section, we first formulate the problem of distributed kernelized contextual bandits. Then, as a starting point, we propose and analyze a naive UCB-type algorithm for distributed kernelized contextual bandit problem, named DisKernelUCB. This demonstrates the challenges in designing a communication efficient algorithm for this problem, and also lays down the foundation for further improvement on communication efficiency in Section 4.

3.1 Distributed Kernelized Contextual Bandit Problem

Consider a learning system with 1) N clients that are responsible for taking actions and receiving feedback from the environment, and 2) a central server that coordinates the communication among the clients. The clients cannot directly communicate with each other, but only with the central server, i.e., a star-shaped communication network. Following prior works [21, 7], we assume the N clients interact with the environment in a round-robin manner for a total number of T rounds.

Specifically, at round $l \in [T]$, each client $i \in [N]$ chooses an arm $\mathbf{x}_t$ from a candidate set $\mathcal{A}_t$, and then receives the corresponding reward feedback $y_t = f(\mathbf{x}_t) + \eta_t \in \mathbb{R}$, where the subscript $t := N(l-1)+i$ indicates this is the t-th interaction between the learning system and the environment, and we refer to it as time step t [1]. Note that $\mathcal{A}_t$ is a time-varying subset of $\mathcal{A} \subseteq \mathbb{R}^d$ that is possibly infinite, f denotes the unknown reward function shared by all the clients, and η_t denotes the noise.

Denote the sequence of indices corresponding to the interactions between client i and the environment up to time t as $\mathcal{N}_t(i) = \{1 \leq s \leq t : i_s = i\}$ (if $s \bmod N = 0$, then $i_s = N$; otherwise $i_s = s \bmod N$) for $t = 1, 2, \ldots, NT$. By definition, $|\mathcal{N}_{Nl}(i)| = l, \forall l \in [T]$, i.e., the clients have equal number of interactions at the end of each round l.

Kernelized Reward Function   We consider an unknown reward function f that lies in a RKHS, denoted as $\mathcal{H}$, such that the reward can be equivalently written as

$$y_t = \theta_\star^\top \phi(\mathbf{x}_t) + \eta_t,$$

where $\theta_\star \in \mathcal{H}$ is an unknown parameter, and $\phi : \mathbb{R}^d \to \mathcal{H}$ is a known feature map associated with $\mathcal{H}$. We assume η_t is zero-mean R-sub-Gaussian conditioned on $\sigma\big((\mathbf{x}_s, \eta_s)_{s \in \mathcal{N}_{t-1}(i_t)}, \mathbf{x}_t\big), \forall t$, which denotes the σ-algebra generated by client i_t's previously pulled arms and the corresponding noise. In addition, there exists a positive definite kernel $k(\cdot, \cdot)$ associated with $\mathcal{H}$, such that $k(\mathbf{x}, \mathbf{x}') = \phi(\mathbf{x})^\top \phi(\mathbf{x}')$, and we assume $k(\cdot, \cdot)$ is bounded by κ^2.

Regret and Communication Cost   The goal of the learning system is to minimize the cumulative (pseudo) regret for all N clients, i.e., $R_{NT} = \sum_{t=1}^{NT} r_t$, where $r_t = \max_{\mathbf{x} \in \mathcal{A}_t} \phi(\mathbf{x})^\top \theta_\star - \phi(\mathbf{x}_t)^\top \theta_\star$. Meanwhile, the learning system also wants to keep the communication cost C_{NT} low, which is measured by the total number of scalars being transferred across the system up to time step NT.

3.2 Distributed Kernel UCB

As a starting point to studying the communication efficient algorithm in Section 4 and demonstrate the challenges in designing a communication efficient distributed kernelized contextual bandit algorithm, here we first introduce and analyze a naive algorithm where the N clients collaborate on learning the exact parameters of kernel bandit, i.e., the mean and variance of estimated reward. We name this algorithm Distributed Kernel UCB, or DisKernelUCB for short, and its description is given in Algorithm 1.

Arm Selection   For each round $l \in [T]$, when client $i \in [N]$ interacts with the environment, i.e., the t-th interaction between the learning system and the environment where $t = N(l-1) + i$, it chooses arm $\mathbf{x}_t \in \mathcal{A}_t$ based on the UCB of the mean estimator (line 5):

$$\mathbf{x}_t = \arg\max_{\mathbf{x} \in \mathcal{A}_t} \hat{\mu}_{t-1,i}(\mathbf{x}) + \alpha_{t-1,i} \hat{\sigma}_{t-1,i}(\mathbf{x}) \tag{1}$$

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[1]The meaning of index t is slightly different from prior works, e.g. DisLinUCB in [21], but this is only to simplify the use of notation and does not affect the theoretical results

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Algorithm 1 Distributed Kernel UCB (DisKernelUCB)

1: **Input** threshold $D > 0$
2: **Initialize** $t_{\text{last}} = 0$, $\mathcal{D}_0(i) = \Delta\mathcal{D}_0(i) = \emptyset, \forall i \in [N]$
3: **for** round $l = 1, 2, ..., T$ **do**
4:    **for** client $i = 1, 2, ..., N$ **do**
5:       Client i chooses arm $\mathbf{x}_t \in \mathcal{A}_t$ according to Eq (1) and observes reward y_t, where $t = N(l-1) + i$
6:       Client i updates $\mathbf{K}_{\mathcal{D}_t(i),\mathcal{D}_t(i)}, \mathbf{y}_{\mathcal{D}_t(i)}$, where $\mathcal{D}_t(i) = \mathcal{D}_{t-1}(i) \cup \{t\}$; and its upload buffer $\Delta\mathcal{D}_t(i) = \Delta\mathcal{D}_{t-1}(i) \cup \{t\}$
      // Global Synchronization
7:       **if** the event $\mathcal{U}_t(D)$ defined in Eq (2) is true **then**
8:          **Clients** $\forall j \in [N]$: send $\{(\mathbf{x}_s, y_s)\}_{s\in\Delta\mathcal{D}_t(j)}$ to server, and reset $\Delta\mathcal{D}_t(j) = \emptyset$
9:          **Server**: aggregates and sends back $\{(\mathbf{x}_s, y_s)\}_{s\in[t]}$; sets $t_{\text{last}} = t$
10:        **Clients** $\forall j \in [N]$: update $\mathbf{K}_{\mathcal{D}_t(j),\mathcal{D}_t(j)}, \mathbf{y}_{\mathcal{D}_t(i)}$, where $\mathcal{D}_t(j) = [t]$

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where $\hat{\mu}_{t,i}(\mathbf{x})$ and $\hat{\sigma}_{t,i}^2(\mathbf{x})$ denote client i's local estimated mean reward for arm $\mathbf{x} \in \mathcal{A}$ and its variance, and $\alpha_{t-1,i}$ is a carefully chosen scaling factor to balance exploration and exploitation (see Lemma 3.1 for proper choice).

To facilitate further discussion, for time step $t \in [NT]$, we denote the sequence of time indices for the data points that have been used to update client i's local estimate as $\mathcal{D}_t(i)$, which include both data points collected locally and those shared by the other clients. If the clients never communicate, $\mathcal{D}_t(i) = \mathcal{N}_t(i), \forall t, i$; otherwise, $\mathcal{N}_t(i) \subset \mathcal{D}_t(i) \subseteq [t]$, with $\mathcal{D}_t(i) = [t]$ recovering the centralized setting, i.e., each new data point collected from the environment immediately becomes available to all the clients in the learning system. The design matrix and reward vector for client i at time step t are denoted by $\mathbf{X}_{\mathcal{D}_t(i)} = [\mathbf{x}_s]_{s\in\mathcal{D}_t(i)}^\top \in \mathbb{R}^{|\mathcal{D}_t(i)|\times d}, \mathbf{y}_{t,i} = [y_s]_{s\in\mathcal{D}_t(i)}^\top \in \mathbb{R}^{|\mathcal{D}_t(i)|}$, respectively. By applying the feature map $\phi(\cdot)$ to each row of $\mathbf{X}_{\mathcal{D}_t(i)}$, we obtain $\mathbf{\Phi}_{\mathcal{D}_t(i)} \in \mathbb{R}^{|\mathcal{D}_t(i)|\times p}$, where p is the dimension of $\mathcal{H}$ and is possibly infinite. Since the reward function is linear in $\mathcal{H}$, client i can construct the Ridge regression estimator $\hat{\theta}_{t,i} = (\mathbf{\Phi}_{\mathcal{D}_t(i)}^\top \mathbf{\Phi}_{\mathcal{D}_t(i)} + \lambda I)^{-1} \mathbf{\Phi}_{\mathcal{D}_t(i)}^\top \mathbf{y}_{t,i}$, where $\lambda > 0$ is the regularization coefficient. This gives us the following estimated mean reward and variance in primal form for any arm $\mathbf{x} \in \mathcal{A}$:

$$\hat{\mu}_{t,i}(\mathbf{x}) = \phi(\mathbf{x})^\top \mathbf{A}_{t,i}^{-1} \mathbf{b}_{t,i}$$

$$\hat{\sigma}_{t,i}(\mathbf{x}) = \sqrt{\phi(\mathbf{x})^\top \mathbf{A}_{t,i}^{-1} \phi(\mathbf{x})}$$

where $\mathbf{A}_{t,i} = \mathbf{\Phi}_{\mathcal{D}_t(i)}^\top \mathbf{\Phi}_{\mathcal{D}_t(i)} + \lambda \mathbf{I}$ and $\mathbf{b}_{t,i} = \mathbf{\Phi}_{\mathcal{D}_t(i)}^\top \mathbf{y}_{t,i}$. Then using the kernel trick, we can obtain their equivalence in the dual form that only involves entries of the kernel matrix, and avoids directly working on $\mathcal{H}$ which is possibly infinite:

$$\hat{\mu}_{t,i}(\mathbf{x}) = \mathbf{K}_{\mathcal{D}_t(i)}(\mathbf{x})^\top (\mathbf{K}_{\mathcal{D}_t(i),\mathcal{D}_t(i)} + \lambda I)^{-1} \mathbf{y}_{\mathcal{D}_t(i)}$$

$$\hat{\sigma}_{t,i}(\mathbf{x}) = \lambda^{-1/2}\sqrt{k(\mathbf{x}, \mathbf{x}) - \mathbf{K}_{\mathcal{D}_t(i)}(\mathbf{x})^\top (\mathbf{K}_{\mathcal{D}_t(i),\mathcal{D}_t(i)} + \lambda I)^{-1} \mathbf{K}_{\mathcal{D}_t(i)}(\mathbf{x})}$$

where $\mathbf{K}_{\mathcal{D}_t(i)}(\mathbf{x}) = \mathbf{\Phi}_{\mathcal{D}_t(i)}\phi(\mathbf{x}) = [k(\mathbf{x}_s, \mathbf{x})]_{s\in\mathcal{D}_t(i)}^\top \in \mathbb{R}^{|\mathcal{D}_t(i)|}$, and $\mathbf{K}_{\mathcal{D}_t(i),\mathcal{D}_t(i)} = \mathbf{\Phi}_{\mathcal{D}_t(i)}^\top \mathbf{\Phi}_{\mathcal{D}_t(i)} = [k(\mathbf{x}_s, \mathbf{x}_{s'})]_{s,s'\in\mathcal{D}_t(i)} \in \mathbb{R}^{|\mathcal{D}_t(i)|\times|\mathcal{D}_t(i)|}$.

Communication Protocol To reduce the regret in future interactions with the environment, the N clients need to collaborate via communication, and a carefully designed communication protocol is essential in ensuring the communication efficiency. In prior works like DisLinUCB [21], after each round of interaction with the environment, client i checks whether the event $\{(|\mathcal{D}_t(i)| - |\mathcal{D}_{t_{\text{last}}}(i)|)\log(\frac{\det(\mathbf{A}_{t,i})}{\det(\mathbf{A}_{t_{\text{last}},i})}) > D\}$ is true, where t_{last} denotes the time step of last global synchronization. If true, a new global synchronization is triggered, such that the server will require all clients to upload their sufficient statistics since t_{last}, aggregate them to compute $\{\mathbf{A}_t, \mathbf{b}_t\}$, and then synchronize the aggregated sufficient statistics with all clients, i.e., set $\{\mathbf{A}_{t,i}, \mathbf{b}_{t,i}\} = \{\mathbf{A}_t, \mathbf{b}_t\}, \forall i \in [N]$.

Using kernel trick, we can obtain an equivalent event-trigger in terms of the kernel matrix,

$$\mathcal{U}_t(D) = \left\{ (|\mathcal{D}_t(i_t)| - |\mathcal{D}_{t_{\text{last}}}(i_t)|) \log \left(\frac{\det(\mathbf{I} + \lambda^{-1}\mathbf{K}_{\mathcal{D}_t(i_t),\mathcal{D}_t(i_t)})}{\det(\mathbf{I} + \lambda^{-1}\mathbf{K}_{\mathcal{D}_t(i_t)\backslash\Delta\mathcal{D}_t(i_t),\mathcal{D}_t(i_t)\backslash\Delta\mathcal{D}_t(i_t)})} \right) > D \right\}. \tag{2}$$

where $D > 0$ denotes the predefined threshold value. If event $\mathcal{U}_t(D)$ is true (line 7), a global synchronization is triggered (line 7-10), where the local datasets of all N clients are synchronized to $\{(\mathbf{x}_s, y_s)\}_{s \in [t]}$. We should note that the transfer of raw data $(\mathbf{x}_s, y_s)$ is necessary for the update of the kernel matrix and reward vector in line 6 and line 10, which will be used for arm selection at line 5. This is an inherent disadvantage of kernelized estimation in distributed settings, which, as we mentioned in Section 2, is also true for the existing distributed kernelized bandit algorithm [6]. Lemma 3.1 below shows that in order to obtain the optimal order of regret, DisKernelUCB incurs a communication cost linear in T (proof given in the appendix), which is expensive for an online learning problem.

Lemma 3.1 (Regret and Communication Cost of DisKernelUCB). *With threshold $D = \frac{T}{N\gamma_{NT}}$,*
$\alpha_{t,i} = \sqrt{\lambda}\|\theta_\star\| + R\sqrt{4\ln N/\delta + 2\ln\det(\mathbf{I} + \mathbf{K}_{\mathcal{D}_t(i),\mathcal{D}_t(i)}/\lambda)}$*, we have*

$$R_{NT} = O\big(\sqrt{NT}(\|\theta_\star\|\sqrt{\gamma_{NT}} + \gamma_{NT})\big),$$

with probability at least $1 - \delta$, and

$$C_{NT} = O(TN^2d).$$

where $\gamma_{NT} := \max_{\mathcal{D}\subset\mathcal{A}:|\mathcal{D}|=NT} \frac{1}{2}\log\det(\mathbf{K}_{\mathcal{D},\mathcal{D}}/\lambda + \mathbf{I})$ is the maximum information gain after NT interactions [5]. It is problem-dependent and can be bounded for specific arm set $\mathcal{A}$ and kernel function $k(\cdot,\cdot)$. For example, $\gamma_{NT} = O(d\log(NT))$ for linear kernel and $\gamma_{NT} = O(\log(NT)^{d+1})$ for Gaussian kernel.

Remark 1. *In the distributed linear bandit problem, to attain $O(d\sqrt{NT}\ln(NT))$ regret, DisLinUCB [21] requires a total number of $O(N^{0.5}d\log(NT))$ synchronizations, and DisKernelUCB matches this result under linear kernel, as it requires $O(N^{0.5}\gamma_{NT})$ synchronizations. We should note that the communication cost for each synchronization in DisLinUCB is fixed, i.e., $O(Nd^2)$ to synchronize the sufficient statistics with all the clients, so in total $C_{NT} = O(N^{1.5}d^3\ln(NT))$. However, this is not the case for DisKernelUCB that needs to send raw data, because the communication cost for each synchronization in DisKernelUCB is not fixed, but depends on the number of unshared data points on each client. Even if the total number of synchronizations is small, DisKernelUCB could still incur $C_{NT} = O(TN^2d)$ in the worse case. Consider the extreme case where synchronization only happens once, but it happens near NT, then we still have $C_{NT} = O(TN^2d)$. The time when synchronization gets triggered depends on $\{\mathcal{A}_t\}_{t \in [NT]}$, which is out of the control of the algorithm. Therefore, in the following section, to improve the communication efficiency of DisKernelUCB, we propose to let each client communicate embedded statistics in some small subspace during each global synchronization.*

4 Approximated Distributed Kernel UCB

In this section, we propose and analyze a new algorithm that improves the communication efficiency of DisKernelUCB using the Nyström approximation, such that the clients only communicate the embedded statistics during event-triggered synchronizations. We name this algorithm Approximated Distributed Kernel UCB, or Approx-DisKernelUCB for short. Its description is given in Algorithm 2.

4.1 Algorithm

Arm selection For each round $l \in [T]$, when client $i \in [N]$ interacts with the environment, i.e., the t-th interaction between the learning system and the environment where $t := N(l-1) + i$, instead of using the UCB for the exact estimator in Eq (1), client i chooses arm $\mathbf{x}_t \in \mathcal{A}_t$ that maximizes the UCB for the approximated estimator (line 5):

$$\mathbf{x}_t = \arg\max_{\mathbf{x}\in\mathcal{A}_{t,i}} \tilde{\mu}_{t-1,i}(\mathbf{x}) + \alpha_{t-1,i}\tilde{\sigma}_{t-1,i}(\mathbf{x}) \tag{3}$$

where $\tilde{\mu}_{t-1,i}(\mathbf{x})$ and $\tilde{\sigma}_{t-1,i}(\mathbf{x})$ are approximated using Nyeström method, and the statistics used to compute these approximations are much more efficient to communicate as they scale with the maximum information gain γ_{NT} instead of T.

Algorithm 2 Approximated Distributed Kernel UCB (Approx-DisKernelUCB)

1: **Input:** threshold $D > 0$, regularization parameter $\lambda > 0$, $\delta \in (0,1)$ and kernel function $k(\cdot, \cdot)$.
2: **Initialize** $\tilde{\mu}_{0,i}(\mathbf{x}) = 0, \tilde{\sigma}_{0,i}(\mathbf{x}) = \sqrt{k(\mathbf{x},\mathbf{x})}, \mathcal{N}_0(i) = \mathcal{D}_0(i) = \emptyset, \forall i \in [N]; \mathcal{S}_0 = \emptyset, t_{\text{last}} = 0$
3: **for** round $l = 1, 2, ..., T$ **do**
4:     **for** client $i = 1, 2, ..., N$ **do**
5:         [Client i] selects arm $\mathbf{x}_t \in \mathcal{A}_t$ according to Eq (3) and observes reward y_t, where $t :=$ $N(l-1) + i$
6:         [Client i] updates $\mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}_{t_{\text{last}}}}^\top \mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}_{t_{\text{last}}}}$ and $\mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}_{t_{\text{last}}}}^\top \mathbf{y}_{\mathcal{D}_t(i)}$ using $\left(\mathbf{z}(\mathbf{x}_t; \mathcal{S}_{t_{\text{last}}}), y_t\right)$; sets $\mathcal{N}_t(i) = \mathcal{N}_{t-1}(i) \cup \{t\}$, and $\mathcal{D}_t(i) = \mathcal{D}_{t-1}(i) \cup \{t\}$
        // Global Synchronization
7:     **if** the event $\mathcal{U}_t(D)$ defined in Eq (4) is true **then**
8:         [Clients $\forall i$] sample $\mathcal{S}_{t,i} = \text{RLS}(\mathcal{N}_t(i), \bar{q}, \tilde{\sigma}_{t_{\text{last}},i}^2)$, and send $\{(\mathbf{x}_s, y_s)\}_{s \in \mathcal{S}_{t,i}}$ to server
9:         [Server] aggregates and sends $\{(\mathbf{x}_s, y_s)\}_{s \in \mathcal{S}_t}$ back to all clients, where $\mathcal{S}_t = \cup_{i \in [N]}\mathcal{S}_{t,i}$
10:        [Clients $\forall i$] compute and send $\{\mathbf{Z}_{\mathcal{N}_t(i);\mathcal{S}_t}^\top \mathbf{Z}_{\mathcal{N}_t(i);\mathcal{S}_t}, \mathbf{Z}_{\mathcal{N}_t(i);\mathcal{S}_t}^\top \mathbf{y}_{\mathcal{N}_t(i)}\}$ to server
11:        [Server] aggregates $\sum_{i=1}^N \mathbf{Z}_{\mathcal{N}_t(i);\mathcal{S}_t}^\top \mathbf{Z}_{\mathcal{N}_t(i);\mathcal{S}_t}, \sum_{i=1}^N \mathbf{Z}_{\mathcal{N}_t(i);\mathcal{S}_t}^\top \mathbf{y}_{\mathcal{N}_t(i)}$ and sends it back
12:        [Clients $\forall i$] updates $\mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}_t}^\top \mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}_t}$ and $\mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}_t}^\top \mathbf{y}_{\mathcal{D}_t(i)}$; sets $\mathcal{D}_t(i) = \cup_{i=1}^N \mathcal{N}_t(i) = [t]$ and $t_{\text{last}} = t$

Specifically, Nyström method works by projecting some original dataset $\mathcal{D}$ to the subspace defined by a small representative subset $\mathcal{S} \subseteq \mathcal{D}$, which is called the dictionary. The orthogonal projection matrix is defined as

$$\mathbf{P}_{\mathcal{S}} = \mathbf{\Phi}_{\mathcal{S}}^\top \left(\mathbf{\Phi}_{\mathcal{S}}\mathbf{\Phi}_{\mathcal{S}}^\top\right)^{-1}\mathbf{\Phi}_{\mathcal{S}} = \mathbf{\Phi}_{\mathcal{S}}^\top \mathbf{K}_{\mathcal{S},\mathcal{S}}^{-1}\mathbf{\Phi}_{\mathcal{S}} \in \mathbb{R}^{p \times p}$$

We then take eigen-decomposition of $\mathbf{K}_{\mathcal{S},\mathcal{S}} = \mathbf{U}\mathbf{\Lambda}\mathbf{U}^\top$ to rewrite the orthogonal projection as $\mathbf{P}_{\mathcal{S}} = \mathbf{\Phi}_{\mathcal{S}}^\top \mathbf{U}\mathbf{\Lambda}^{-1/2}\mathbf{\Lambda}^{-1/2}\mathbf{U}^\top \mathbf{\Phi}_{\mathcal{S}}$, and define the Nyström embedding function

$$z(\mathbf{x}; \mathcal{S}) = \mathbf{P}_{\mathcal{S}}^{1/2}\phi(\mathbf{x}) = \mathbf{\Lambda}^{-1/2}\mathbf{U}^\top\mathbf{\Phi}_{\mathcal{S}}\phi(\mathbf{x}) = \mathbf{K}_{\mathcal{S},\mathcal{S}}^{-1/2}\mathbf{K}_{\mathcal{S}}(\mathbf{x})$$

which maps the data point $\mathbf{x}$ from $\mathbb{R}^d$ to $\mathbb{R}^{|\mathcal{S}|}$.

Therefore, we can approximate the Ridge regression estimator in Section 3.2 as $\tilde{\theta}_{t,i} = \tilde{\mathbf{A}}_{t,i}^{-1}\tilde{\mathbf{b}}_{t,i}$, where $\tilde{\mathbf{A}}_{t,i} = \mathbf{P}_{\mathcal{S}}\mathbf{\Phi}_{\mathcal{D}_t(i)}^\top \mathbf{\Phi}_{\mathcal{D}_t(i)}\mathbf{P}_{\mathcal{S}} + \lambda\mathbf{I}$, and $\tilde{\mathbf{b}}_{t,i} = \mathbf{P}_{\mathcal{S}}\mathbf{\Phi}_{\mathcal{D}_t(i)}^\top \mathbf{y}_{\mathcal{D}_t(i)}$, and thus the approximated mean reward and variance in Eq (3) can be expressed as

$$\tilde{\mu}_{t,i}(\mathbf{x}) = \phi(\mathbf{x})^\top \tilde{\mathbf{A}}_{t,i}^{-1}\tilde{\mathbf{b}}_{t,i}$$
$$\tilde{\sigma}_{t,i}(\mathbf{x}) = \sqrt{\phi(\mathbf{x})^\top \tilde{\mathbf{A}}_{t,i}^{-1}\phi(\mathbf{x})}$$

and their kernelized representation are (see appendix for detailed derivation)

$$\tilde{\mu}_{t,i}(\mathbf{x}) = z(\mathbf{x}; \mathcal{S})^\top \left(\mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}}^\top \mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}} + \lambda\mathbf{I}\right)^{-1}\mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}}^\top \mathbf{y}_{\mathcal{D}_t(i)}$$
$$\tilde{\sigma}_{t,i}(\mathbf{x}) = \lambda^{-1/2}\sqrt{k(\mathbf{x},\mathbf{x}) - z(\mathbf{x};\mathcal{S})^\top \mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}}^\top \mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}}[\mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}}^\top \mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}} + \lambda\mathbf{I}]^{-1}z(\mathbf{x}|\mathcal{S})}$$

where $\mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}} \in \mathbb{R}^{|\mathcal{D}_t(i)| \times |\mathcal{S}|}$ is obtained by applying $z(\cdot; \mathcal{S})$ to each row of $\mathbf{X}_{\mathcal{D}_t(i)}$, i.e., $\mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}} = \mathbf{\Phi}_{\mathcal{D}_t(i)}\mathbf{P}_{\mathcal{S}}^{1/2}$. We can see that the computation of $\tilde{\mu}_{t,i}(\mathbf{x})$ and $\tilde{\sigma}_{t,i}(\mathbf{x})$ only requires the embedded statistics: matrix $\mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}}^\top \mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}} \in \mathbb{R}^{|\mathcal{S}| \times |\mathcal{S}|}$ and vector $\mathbf{Z}_{\mathcal{D}_t(i);\mathcal{S}}^\top \mathbf{y}_{\mathcal{D}_t(i)} \in \mathbb{R}^{|\mathcal{S}|}$, which, as we will show later, makes joint kernelized estimation among N clients much more efficient in communication.

After obtaining the new data point $(\mathbf{x}_t, y_t)$, client i immediately updates both $\tilde{\mu}_{t-1,i}(\mathbf{x})$ and $\tilde{\sigma}_{t-1,i}(\mathbf{x})$ using the newly collected data point $(\mathbf{x}_t, y_t)$, i.e., by projecting $\mathbf{x}_t$ to the finite dimensional RKHS spanned by $\mathbf{\Phi}_{\mathcal{S}_{t_{\text{last}}}}$ (line 6). Recall that, we use $\mathcal{N}_t(i)$ to denote the sequence of indices for data collected by client i, and denote by $\mathcal{D}_t(i)$ the sequence of indices for data that has been used to update client i's model estimation $\tilde{\mu}_{t,i}$. Therefore, both of them need to be updated to include time step t.

Communication Protocol With the approximated estimator, the size of message being communicated across the learning system is reduced. However, a carefully designed event-trigger is still required to minimize the total number of global synchronizations up to time NT. Since the clients can no longer evaluate the exact kernel matrices in Eq (2), we instead use the event-trigger in Eq (4), which can be computed using the approximated variance from last global synchronization as,

$$\mathcal{U}_t(D) = \left\{ \sum_{s \in \mathcal{D}_t(i) \setminus \mathcal{D}_{t_{\text{last}}}(i)} \tilde{\sigma}^2_{t_{\text{last}},i}(\mathbf{x}_s) > D \right\} \tag{4}$$

Similar to Algorithm 1, if Eq (4) is true, global synchronization is triggered, where both the dictionary and the embedded statistics get updated. During synchronization, each client first samples a subset $\mathcal{S}_t(i)$ from $\mathcal{N}_t(i)$ (line 8) using Ridge Leverage Score sampling (RLS) [2, 3], which is given in Algorithm 3, and then sends $\{(\mathbf{x}_s, y_s)\}_{s \in \mathcal{S}_t(i)}$ to the server. The server aggregates the received local subsets to construct a new dictionary $\{(\mathbf{x}_s, y_s)\}_{s \in \mathcal{S}_t}$, where $\mathcal{S}_t = \cup_{i=1}^N \mathcal{S}_t(i)$, and then sends it back to all N clients (line 9). Finally, the N clients use this updated dictionary to re-compute the embedded statistics of their local data, and then synchronize it with all other clients via the server (line 10-12).

Algorithm 3 Ridge Leverage Score Sampling (RLS)

1: **Input:** dataset $\mathcal{D}$, scaling factor $\bar{q}$, (possibly delayed and approximated) variance function $\tilde{\sigma}^2(\cdot)$
2: **Initialize** new dictionary $\mathcal{S} = \emptyset$
3: **for** $s \in \mathcal{D}$ **do**
4:      Set $\tilde{p}_s = \bar{q}\tilde{\sigma}^2(\mathbf{x}_s)$
5:      Draw $q_s \sim \text{Bernoulli}(\tilde{p}_s)$
6:      If $q_s = 1$, add s into $\mathcal{S}$
7: **Output:** $\mathcal{S}$

Intuitively, in Algorithm 2, the clients first agree upon a common dictionary $\mathcal{S}_t$ that serves as a good representation of the whole dataset at the current time t, and then project their local data to the subspace spanned by this dictionary before communication, in order to avoid directly sending the raw data as in Algorithm 1. Then using the event-trigger, each client monitors the amount of new knowledge it has gained through interactions with the environment from last synchronization. When there is a sufficient amount of new knowledge, it will inform all the other clients to perform a synchronization. As we will show in the following section, the size of $\mathcal{S}_t$ scales linearly w.r.t. the maximum information gain γ_{NT}, and therefore it improves both the local computation efficiency on each client, and the communication efficiency during the global synchronization.

4.2 Theoretical Analysis

Denote the sequence of time steps when global synchronization is performed, i.e., the event $\mathcal{U}_t(D)$ in Eq (4) is true, as $\{t_p\}_{p=1}^B$, where $B \in [NT]$ denotes the total number of global synchronizations. Note that in Algorithm 2, the dictionary is only updated during global synchronization, e.g., at time t_p, the dictionary $\{(\mathbf{x}_s, y_s)\}_{s \in \mathcal{S}_{t_p}}$ is sampled from the whole dataset $\{(\mathbf{x}_s, y_s)\}_{s \in [t_p]}$ in a distributed manner, and remains fixed for all the interactions happened at $t \in [t_p + 1, t_{p+1}]$. Moreover, at time t_p, all the clients synchronize their embedded statistics, so that $\mathcal{D}_{t_p}(i) = [t_p], \forall i \in [N]$.

Since Algorithm 2 enables local update on each client, for time step $t \in [t_p + 1, t_p]$, new data points are collected and added into $\mathcal{D}_t(i)$, such that $\mathcal{D}_t(i) \supseteq [t_p]$. This decreases the approximation accuracy of $\mathcal{S}_{t_p}$, as new data points may not be well approximated by $\mathcal{S}_{t_p}$. For example, in extreme cases, the new data could be orthogonal to the dictionary. To formally analyze the accuracy of the dictionary, we adopt the definition of ϵ-accuracy from [4]. Denote by $\bar{\mathbf{S}}_{t,i} \in \mathbb{R}^{|\mathcal{D}_t(i)| \times |\mathcal{D}_t(i)|}$ a diagonal matrix, with its s-th diagonal entry equal to $\frac{1}{\sqrt{\tilde{p}_s}}$ if $s \in \mathcal{S}_{t_p}$ and 0 otherwise. Then if

$$(1 - \epsilon_{t,i})(\mathbf{\Phi}_{\mathcal{D}_t(i)}^\top \mathbf{\Phi}_{\mathcal{D}_t(i)} + \lambda \mathbf{I}) \preceq \mathbf{\Phi}_{\mathcal{D}_t(i)}^\top \bar{\mathbf{S}}_{t,i}^\top \bar{\mathbf{S}}_{t,i} \mathbf{\Phi}_{\mathcal{D}_t(i)} + \lambda \mathbf{I} \preceq (1 + \epsilon_{t,i})(\mathbf{\Phi}_{\mathcal{D}_t(i)}^\top \mathbf{\Phi}_{\mathcal{D}_t(i)} + \lambda \mathbf{I}),$$

we say the dictionary $\{(\mathbf{x}_s, y_s)\}_{s \in \mathcal{S}_{t_p}}$ is $\epsilon_{t,i}$-accurate w.r.t. dataset $\{(\mathbf{x}_s, y_s)\}_{s \in \mathcal{D}_t(i)}$. In the following, we analyze the $\epsilon_{t,i}$-accuracy of the dictionary for all t, i.

At the time steps when global synchronization happens, i.e., t_p for $p \in [B]$, $\mathcal{S}_{t_p}$ is sampled from $[t_p] = \mathcal{D}_{t_p}(i)$ using approximated variance $\tilde{\sigma}^2_{t_{p-1},i}$. In this case, the accuracy of the dictionary only

depends on the RLS procedure, and Calandriello et al. [3] have already showed that the following guarantee on the accuracy and size of dictionary holds $\forall t \in \{t_p\}_{p \in [B]}$.

Lemma 4.1 (Lemma 2 of [3]). *Under the condition that $\bar{q} = 6\frac{1+\epsilon}{1-\epsilon}\log(4NT/\delta)/\epsilon^2$, for some $\epsilon \in [0, 1)$, with probability at least $1 - \delta$, we have $\forall t \in \{t_p\}_{p \in [B]}$ that the dictionary $\{(\mathbf{x}_s, y_s)\}_{s \in \mathcal{S}_t}$ is ϵ-accurate w.r.t. $\{(\mathbf{x}_s, y_s)\}_{s \in \mathcal{D}_t(i)}$, and $\frac{1-\epsilon}{1+\epsilon}\sigma_t^2(\mathbf{x}) \leq \tilde{\sigma}_t^2(\mathbf{x}) \leq \frac{1+\epsilon}{1-\epsilon}\sigma_t^2(\mathbf{x}), \forall \mathbf{x} \in \mathcal{A}$. Moreover, the size of dictionary $|\mathcal{S}_t| \leq 3(1 + \kappa^2/\lambda)\frac{1+\epsilon}{1-\epsilon}\bar{q}\tilde{d}$, where $\tilde{d} := Tr(\mathbf{K}_{[NT],[NT]}(\mathbf{K}_{[NT],[NT]} + \lambda\mathbf{I})^{-1})$ denotes the effective dimension of the problem, and it is known that $\tilde{d} = O(\gamma_{NT})$ [5].*

Lemma 4.1 guarantees that for all $t \in \{t_p\}_{p \in [B]}$, the dictionary has a constant accuracy, i.e., $\epsilon_{t,i} = \epsilon, \forall i$. In addition, since the dictionary is fixed for $t \notin \{t_p\}_{p \in [B]}$, its size $\mathcal{S}_t = O(\gamma_{NT}), \forall t \in [NT]$.

Then for time steps $t \notin \{t_p\}_{p \in [B]}$, due to the local update, the accuracy of the dictionary will degrade. However, thanks to our event-trigger in Eq (4), the extent of such degradation can be controlled, i.e., a new dictionary update will be triggered before the previous dictionary becomes completely irrelevant. This is shown in Lemma 4.2 below.

Lemma 4.2. *Under the condition that $\{(\mathbf{x}_s, y_s)\}_{s \in \mathcal{S}_{t_p}}$ is ϵ-accurate w.r.t. $\{(\mathbf{x}_s, y_s)\}_{s \in \mathcal{D}_{t_p}(i)}, \forall t \in [t_p + 1, t_{p+1}], i \in [N], \mathcal{S}_{t_p}$ is $\left(\epsilon + 1 - \frac{1}{1+\frac{1+\epsilon}{1-\epsilon}D}\right)$-accurate w.r.t. $\mathcal{D}_t(i)$.*

Combining Lemma 4.1 and Lemma 4.2, we have that with $\bar{q} = 6\frac{1+\epsilon}{1-\epsilon}\log(4NT/\delta)/\epsilon^2$ and threshold $D > 0$, Algorithm 2 guarantees that the dictionary is accurate with constant $\epsilon_{t,i} := \left(\epsilon + 1 - \frac{1}{1+\frac{1+\epsilon}{1-\epsilon}D}\right)$, and its size $|\mathcal{S}_t| = O(\gamma_{NT})$ for all $t \in [NT]$. Then we can construct the following confidence ellipsoid for the unknown parameter $\theta_\star$. Its proof is given in appendix.

Lemma 4.3 (Confidence Ellipsoid of Approximated Estimator). *Under the same conditions as Lemma 4.1 and Lemma 4.2, with probability at least $1 - \delta$, we have $\forall t, i$ that*

$$\|\tilde{\theta}_{t,i} - \theta_\star\|_{\tilde{\mathbf{A}}_{t,i}} \leq \left(\frac{1}{\sqrt{-\epsilon + \frac{1}{1+\frac{1+\epsilon}{1-\epsilon}D}}} + 1\right)\sqrt{\lambda}\|\theta_\star\| + 2R\sqrt{\ln N/\delta + \gamma_{NT}} := \alpha_{t,i}.$$

Using Lemma 4.3, we obtain the regret and communication cost upper bound of Approx-DisKernelUCB, which is given in Theorem 4.4 below.

Theorem 4.4 (Regret and Communication Cost of Approx-DisKernelUCB). *Under the same condition as Lemma 4.1 and Lemma 4.2, and by setting $D = \frac{1}{N}, \epsilon < \frac{1}{3}$, we have*

$$R_{NT} = O\left(\sqrt{NT}(\|\theta_\star\|\sqrt{\gamma_{NT}} + \gamma_{NT})\right)$$

with probability at least $1 - \delta$, and

$$C_{NT} = O\left(N^2\gamma_{NT}^3\right)$$

Here we provide a proof sketch for Theorem 4.4, and the complete proof can be found in appendix.

Proof Sketch. Similar to the analysis of DisKernelUCB in Section 3.2 and DisLinUCB from [21], the cumulative regret incurred by Approx-DisKernelUCB can be decomposed in terms of 'good' and 'bad' epochs, and bounded separately. Here an epoch refers to the time period in-between two consecutive global synchronizations, e.g., the p-th epoch refers to $[t_{p-1} + 1, t_p]$. Now consider an imaginary centralized agent that has immediate access to each data point in the learning system, and denote by $A_t = \sum_{s=1}^{t}\phi_s\phi_s^\top$ for $t \in [NT]$ the matrix constructed by this centralized agent. We call the p-th epoch a good epoch if $\ln(\frac{\det(A_{t_p})}{\det(A_{t_{p-1}})}) \leq 1$, otherwise it is a bad epoch. Note that $\ln(\frac{\det(A_{t_1})}{\det(\lambda\mathbf{I})}) + \ln(\frac{\det(A_{t_2})}{\det(A_{t_1})}) + \cdots + \ln(\frac{\det(A_{NT})}{\det(A_{t_B})}) = \ln(\frac{\det(A_{NT})}{\det(\lambda\mathbf{I})}) = \ln(\det(I + \lambda^{-1}\mathbf{K}_{[NT],[NT]})) \leq 2\gamma_{NT}$, where the last equality is due to the matrix determinant lemma, and the last inequality is by the definition of the maximum information gain γ_{NT} in Lemma 3.1. Then based on the pigeonhole principle, there can be at most $2\gamma_{NT}$ bad epochs.

By combining Lemma 4.1 and Lemma 4.3, we can bound the cumulative regret incurred during all good epochs, i.e., $R_{good} = O(\sqrt{NT}\gamma_{NT})$, which matches the optimal regret attained by the KernelUCB algorithm in centralized setting.

Our analysis deviates from that of DisKernelUCB in the bad epochs, because of the difference in the event-trigger. Previously, the event-trigger of DisKernelUCB directly bounds the cumulative regret each client incurs during a bad epoch, i.e., $\sum_{t \in \mathcal{D}_{t_p}(i) \setminus \mathcal{D}_{t_{p-1}}(i)} \hat{\sigma}_{t-1,i}(\mathbf{x}_t) \leq$

$\sqrt{(|\mathcal{D}_{t_p}(i)| - |\mathcal{D}_{t_{p-1}}(i)|) \log \left(\det(\mathbf{I} + \lambda^{-1}\mathbf{K}_{\mathcal{D}_t(i_t),\mathcal{D}_t(i_t)}) / \det(\mathbf{I} + \lambda^{-1}\mathbf{K}_{\mathcal{D}_t(i_t) \setminus \Delta\mathcal{D}_t(i_t),\mathcal{D}_t(i_t) \setminus \Delta\mathcal{D}_t(i_t)}) \right)} <$
$\sqrt{D}$. However, the event trigger of Approx-DisKernelUCB only bounds part of it, i.e., $\sum_{t \in \mathcal{D}_{t_p}(i) \setminus \mathcal{D}_{t_{p-1}}(i)} \tilde{\sigma}_{t-1,i}(\mathbf{x}_t) \leq \sqrt{(|\mathcal{D}_{t_p}(i)| - |\mathcal{D}_{t_{p-1}}(i)|)D}$, which leads to $R_{bad} = O(\sqrt{T}\gamma_{NT}N\sqrt{D})$ that is slightly worse than that of DisKernelUCB, i.e., a $\sqrt{T}$ factor in place of the $\sqrt{\gamma_{NT}}$ factor. By setting $D = 1/N$, we have $R_{NT} = O(\sqrt{NT}\gamma_{NT})$. Note that, to make sure $\epsilon_{t,i} = \left(\epsilon + 1 - \frac{1}{1+\frac{1+\epsilon}{1-\epsilon}\frac{1}{N}}\right) \in [0,1)$ is still well-defined, we can set $\epsilon < 1/3$.

For communication cost analysis, we bound the total number of epochs B by upper bounding the total number of summations like $\sum_{s=t_{p-1}+1}^{t_p} \hat{\sigma}_{t_{p-1}}^2(\mathbf{x}_s)$, over the time horizon NT. Using Lemma 4.1, our event-trigger in Eq (4) provides a lower bound $\sum_{s=t_{p-1}+1}^{t_p} \hat{\sigma}_{t_{p-1}}^2(\mathbf{x}_s) \geq \frac{1-\epsilon}{1+\epsilon}D$. Then in order to apply the pigeonhole principle, we continue to upper bound the summation over all epochs,
$\sum_{p=1}^{B} \sum_{s=t_{p-1}+1}^{t_p} \hat{\sigma}_{t_{p-1}}^2(\mathbf{x}_s) = \sum_{p=1}^{B} \sum_{s=t_{p-1}+1}^{t_p} \hat{\sigma}_{s-1}^2(\mathbf{x}_s) \frac{\hat{\sigma}_{t_{p-1}}^2(\mathbf{x}_s)}{\hat{\sigma}_{s-1}^2(\mathbf{x}_s)}$ by deriving a uniform bound for the ratio $\frac{\hat{\sigma}_{t_{p-1}}^2(\mathbf{x}_s)}{\hat{\sigma}_{s-1}^2(\mathbf{x}_s)} \leq \frac{\hat{\sigma}_{t_{p-1}}^2(\mathbf{x}_s)}{\hat{\sigma}_{t_p}^2(\mathbf{x}_s)} \leq 1 + \sum_{s=t_{p-1}+1}^{t_p} \hat{\sigma}_{t_{p-1}}^2(\mathbf{x}_s) \leq 1 + \frac{1+\epsilon}{1-\epsilon} \sum_{s=t_{p-1}+1}^{t_p} \tilde{\sigma}_{t_{p-1}}^2(\mathbf{x}_s)$
in terms of the communication threshold D on each client. This leads to the following upper bound about the total number of epochs $B \leq \frac{1+\epsilon}{1-\epsilon}[\frac{1}{D} + \frac{1+\epsilon}{1-\epsilon}(N + \kappa^2/(\lambda D))]2\gamma_{NT}$, and with $D = 1/N$, we have $C_{NT} \leq B \cdot N\gamma_{NT}^2 = O(N^2\gamma_{NT}^3)$, which completes the proof. $\qquad\square$

Remark 2. *Compared with DisKernelUCB's $O(TN^2d)$ communication cost, Approx-DisKernelUCB removes the linear dependence on T, but introduces an additional γ_{NT}^3 dependence due to the communication of the embedded statistics. In situations where $\gamma_{NT} \ll T^{1/3}d^{1/3}$, DisKernelUCB is preferable. As mentioned in Lemma 3.1, the value of γ_{NT}, which affects how much the data can be compressed, depends on the specific arm set of the problem and the kernel function of the choice. By Mercer's Theorem, one can represent the kernel using its eigenvalues, and γ_{NT} characterizes how fast its eigenvalues decay. Vakili et al. [19] showed that for kernels whose eigenvalues decay exponentially, i.e., $\lambda_m = O(\exp(-m^{\beta_e}))$, for some $\beta_e > 0$, $\gamma_{NT} = O(\log^{1+\frac{1}{\beta_e}}(NT))$. In this case, Approx-DisKernelUCB is far more efficient than DisKernelUCB. This includes Gaussian kernel, which is widely used for GPs and SVMs. For kernels that have polynomially decaying eigenvalues, i.e., $\lambda_m = O(m^{-\beta_p})$, for some $\beta_p > 1$, $\gamma_{NT} = O(T^{\frac{1}{\beta_p}} \log^{1-\frac{1}{\beta_p}}(NT))$. Then as long as $\beta_p > 3$, Approx-DisKernelUCB still enjoys reduced communication cost.*

Due to space limit, we leave our empirical evaluations for the proposed algorithm in the appendix.

Conclusion

In this paper, we proposed the first communication efficient algorithm for distributed kernel bandits using Nyström approximation. Clients in the learning system project their local data to a finite RKHS spanned by a shared dictionary, and then communicate the embedded statistics for collaborative exploration. To ensure communication efficiency, the frequency of dictionary update and synchronization of embedded statistics are controlled by an event-trigger. The algorithm is proved to incur $O(N^2\gamma_{NT}^3)$ communication cost, while attaining the optimal $O(\sqrt{NT}\gamma_{NT})$ cumulative regret.

We should note that the total number of synchronizations required by Approx-DisKernelUCB is $N\gamma_{NT}$, which is $\sqrt{N}$ worse than DisKernelUCB. An important future direction of this work is to investigate whether this part can be further improved. It is also interesting to extend the proposed algorithm to P2P setting, at the absence of a central server to coordinate the update of the shared dictionary and the exchange of embedded statistics. Due to the delay in propagating messages, it may be beneficial to utilize possible local structures in the network of clients, and approximate each block of the kernel matrix separately [17], i.e., each block corresponds to a group of clients, instead of directly approximating the complete matrix.

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