Safety Guarantees for Neural Network Dynamic Systems via Stochastic Barrier Functions

Abstract

Neural Networks (NNs) have been successfully employed to represent the state evolution of complex dynamical systems. Such models, referred to as NN dynamic models (NNDMs), use iterative noisy predictions of NN to estimate a distribution of system trajectories over time. Despite their accuracy, safety analysis of NNDMs is known to be a challenging problem and remains largely unexplored. To address this issue, in this paper, we introduce a method of providing safety guarantees for NNDMs. Our approach is based on stochastic barrier functions, whose relation with safety are analogous to that of Lyapunov functions with stability. We first show a method of synthesizing stochastic barrier functions for NNDMs via a convex optimization problem, which in turn provides a lower bound on the system's safety probability. A key step in our method is the employment of the recent convex approximation results for NNs to find piece-wise linear bounds, which allow the formulation of the barrier function synthesis problem as a sum-of-squares optimization program. If the obtained safety probability is above the desired threshold, the system is certified. Otherwise, we introduce a method of generating controls for the system that robustly minimize the unsafety probability in a minimally-invasive manner. We exploit the convexity property of the barrier function to formulate the optimal control synthesis problem as a linear program. Experimental results illustrate the efficacy of the method. Namely, they show that the method can scale to multi-dimensional NNDMs with multiple layers and hundreds of neurons per layer, and that the controller can significantly improve the safety probability.

1 Introduction

Safety is a major concern for autonomous dynamical systems, especially in *safety-critical* applications. Examples include UAVs, autonomous cars, and surgical robots. To this end, various methods are developed to provide safety guarantees for such systems [1-3]. Most of these methods, however, assume (simple) analytical models, which are often unavailable or too expensive to obtain due to complexity (nonlinearities) in real-world systems [4]. This has given rise to recent efforts to harvest the representational power of *neural networks* (NNs) to predict how the state of a system evolves over time, referred to as *NN dynamic models* (NNDMs) [5, 6]. Despite their efficient and accurate representation of realistic dynamics, it has been difficult to make formal claims about the behavior of NNDMs over time. This poses a major challenge in safety-critical domains: *how to provide safety guarantees for NNDMs?*

In this paper, we propose a framework to provide safety certificates for NNDMs. The approach is based on stochastic barrier functions [7, 8], whose role for safety is analogous to that of Lyapunov functions for stability, and which can guarantee forward invariance of a safe set. Specifically, we show a method of synthesizing barrier functions for NNDMs via an efficient convex optimization method, which in turn provides a lower bound on the safety probability of the system trajectories.

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The benefit of our approach is that it does not require unfolding of the NNDM to compute this probability; instead, we base our reasoning on non-negative super-martingale inequalities to bound the probability of leaving the safe set. The key to this is to utilize the recent convex approximation results for NNs to obtain piece-wise linear bounds, and a formulation of the optimization problem that enables the use of efficient solvers. If the obtained safety probability is above a desired threshold, the system can be certified. For the cases that such safety certificates cannot be provided, we introduce a method of generating controls for the system, which robustly maximizes the safety probability in a minimally-invasive manner. We exploit the convexity property of the barrier function to formulate the control synthesis problem as a linear program. Experimental results on various NNDMs trained using state-of-the-art RL techniques [9] illustrate the efficacy of the method. Namely, they show that the method can scale to multi-dimensional models with various layers and hundreds of neurons per layer, and that the controller can significantly improve the safety probability.

In summary, this paper makes the following main contributions: (i) a method of computing a lower bound for the safety probability of NNDMs via stochastic barrier functions and efficient nonlinear optimization, thereby providing safety guarantees, (ii) a correct-by-construction control synthesis framework for NNDMs to ensure safety while being *minimally-invasive*, i.e., intervening only when needed, (iii) demonstration of the efficacy and scalability of the framework in benchmarks on various NNDMs with multiple architectures, i.e., multiple hidden layers and hundreds of neurons per layer.

Related Work   In recent years, NNs have been successfully employed to model complex dynamical systems [5, 6]. Such models use iterative noisy predictions of an NN to estimate the distribution of the system over time. They provide means of predicting state evolution of the systems that have no analytical form, e.g., soft robotics [10]. In addition, because of their ability to quickly unfold trajectories, NNDMs have been used to enhance controller training in Reinforcement Learning (RL) frameworks [9, 11]. In these works, specifically, a NNDM of the system is trained in closed-loop with an NN controllers, which can be concatenated in a single NN representing the dynamics of the closed-loop system. Despite their popularity and increased usage, little work has gone to safety analysis of NNDMs. This paper aims to fill this gap.

Barrier functions provide a formal methodology to prove safety of dynamical systems [12, 13]. The state-of-the-art to find stochastic barrier functions is arguably to rely on Sum-of-Squares (SOS) optimization [8]. Barrier certificates obtained via SOS optimization have been studied for deterministic, uncontrolled nonlinear systems [14], and control-affine stochastic systems [15, 8], as well as for hybrid models [12, 13]. Unlike the existing work that focus on closed-form models, this work is the first to consider barrier function for models represented by NNs. In addition, the control synthesis formulation in this paper relaxes the limitations of the existing SOS methods, which generally leads to a non-convex optimization problem [8], by reducing it to a simple linear program.

Many recent studies have focused on the verification of NNs. However, most of them only aim to provide certification guarantees against adversarial examples [16], with approaches ranging from SMT [17], and convex relaxations [18] to linear programming [19]. Such methods, because of their local nature, cannot be directly employed to perform safety analysis on the entire state space and over the trajectories as required for NNDMs. Other methods have been employed for trajectory-level properties of NNs based on solution of the Bellman equations [20], formal (finite) abstractions [21], and mixed integer programming [22], which does not support noisy dynamics. In contrast, our approach supports additive noise and only requires to check static properties on the NN by showing that there exists a function whose composition with the NN forms a stochastic barrier function.

2   Problem Formulation

In this work, we are interested in providing safety guarantees for dynamical systems that are described by NNs. Specifically, we consider the following discrete-time stochastic process whose time evolution is given by iterative predictions of an NN

$$\mathbf{x}_{k+1} = f^w(\mathbf{x}_k) + \mathbf{v}_k, \qquad k \in \mathbb{N}, \quad \mathbf{x}_k, \mathbf{v}_k \in \mathbb{R}^n, \qquad (1)$$

where f^w is a trained, fully-connected, feed-forward NN with general continuous activation functions, and w represents the maximum likelihood weights. Term $\mathbf{v}_k$ is a random variable modelling an additive noise of stationary Gaussian distribution with zero mean and covariance matrix $R \in \mathbb{R}^{n \times n}$, i.e., the probability density function of $\mathbf{v}_k$ is $\mathcal{N}(\bar{x} \mid 0, R)$.

Remark 1. *System (1) can be viewed as a closed-loop system, where f^w is the composition of a NNDM (that represents the open-loop system) with an NN controller. NNDMs such as the one in System (1) are increasingly employed in robotics for both model representation and NN controller training. These models are often obtained by state-of-the-art model-based RL techniques (e.g., [9, 23, 11]) or perturbations (step response) of the actual system (e.g., [10]).*

The evolution of System (1) can be characterized by the predictive posterior distribution $p_{\mathbf{x}}(\bar{x} \mid x)$, which describes the probability density of $\mathbf{x}_{k+1}$ when $\mathbf{x}_k = x \in \mathbb{R}^n$. This distribution is induced by noise $\mathbf{v}_k$ and hence is Gaussian with mean $f^w(x)$ and covariance R, i.e., $p_{\mathbf{x}}(\bar{x} \mid x) = \mathcal{N}(\bar{x} \mid f^w(x), R)$. For a subset of states $X \subseteq \mathbb{R}^n$ and (initial) state $x \in \mathbb{R}^n$, we call $T(X \mid x) = \int_X p_{\mathbf{x}}(\bar{x} \mid x) d\bar{x}$ the *stochastic kernel* of System (1). From the definition of T it follows that, for a given initial condition $\mathbf{x}_0 = x_0 \in \mathbb{R}^n$, $\mathbf{x}_k$ is a Markov process with a well-defined probability measure Pr uniquely generated by the stochastic kernel T [24, Proposition 7.45]. For $X_0, X_k \subseteq \mathbb{R}^n$, this probability measure is inductively defined as

$$\Pr[\mathbf{x}_0 \in X_0] = \mathbf{1}_{X_0}(x_0), \qquad \Pr[\mathbf{x}_k \in X_k \mid \mathbf{x}_{k-1} = x] = T(X_k \mid x),$$

where indicator function $\mathbf{1}_{X_0}(x_0) = 1$ if $x_0 \in X_0$, otherwise 0. The definition of Pr allows one to make probabilistic statements over the trajectories of System (1). In this work, we are particularly interested in two main problems: safety certification and safe controller synthesis for System (1).

Problem 1 (Safety Certificate). *Let $X_{\mathrm{s}} \subset \mathbb{R}^n$ and $X_0 \subseteq X_{\mathrm{s}}$ be basic closed semi-algebraic sets representing respectively the safe set and the set of initial states. Denote the probability that the system initialized in X_0 remains in the safe set in the next $N \in \mathbb{N}_{\geq 0}$ steps by*

$$P_{\mathrm{s}}(X_{\mathrm{s}}, X_0, N) = Pr[\forall x_0 \in X_0, \forall k \leq N, \mathbf{x}_k \in X_{\mathrm{s}}]. \tag{2}$$

Then, given threshold $\delta_{\mathrm{s}} \in [0, 1]$, provide a safety certificate for System (1) by proving that $P_{\mathrm{s}}(X_{\mathrm{s}}, X_0, N) \geq \delta_{\mathrm{s}}$.

Problem 1 seeks to compute the probability that $\mathbf{x}_k$ remains within a safe set, e.g., it avoids obstacles or any undesirable states. Computation of this probability is particularly challenging because System (1) is stochastic and NN (f^w) is generally a non-convex function. The assumption that X_{s} and X_0 are basic closed semi-algebraic sets is not limiting, as these sets are defined as an intersection of a finite number of polynomial inequalities and, for instance, include convex polytopes and spectrahedra [25].

In the case that a safety certificate cannot be generated, we are interested in synthesizing a controller that *by-design* guarantees safety. To this end, we extend System (1) with an affine input

$$\mathbf{x}_{k+1} = f^w(\mathbf{x}_k) + g\mathbf{u}_k + \mathbf{v}_k, \tag{3}$$

where $g : \mathbb{R}^n \to \mathbb{R}^c$ is given, $\mathbf{u}_k \in U \subset \mathbb{R}^c$ is a control action, and U is bounded. Our goal is to synthesize a feedback controller $\pi : \mathbb{R}^n \to U$ such that $\mathbf{u}_k = \pi(\mathbf{x}_k)$ guarantees safety of System (3).

Problem 2 (Safe Control Synthesis). *Let $X_{\mathrm{s}} \subset \mathbb{R}^n$ and $X_0 \subseteq X_{\mathrm{s}}$ be basic closed semi-algebraic sets representing respectively the safe set and the set of initial states, and $N \in \mathbb{N}_{\geq 0}$ a time horizon. Then, given threshold δ_{s}, synthesize a feedback controller π such that $P_{\mathrm{s}}(X_{\mathrm{s}}, X_0, N, \pi) \geq \delta_{\mathrm{s}}$, where $P_{\mathrm{s}}(X_{\mathrm{s}}, X_0, N, \pi)$ is the safety probability obtained for System (3).*

Remark 2. *Note that there may exist many possible controllers that are solutions to Problem 2. Since System (1) represents a closed-loop system with a given controller, one may want to specifically synthesize a safety controller that is "minimally-invasive," i.e., it minimally intervenes to guarantee safety. In Section 5, we elaborate on this notion and introduce methods of generating such controllers.*

Overall Approach   Our approaches to Problems 1 and 2 are based on (stochastic) barrier functions. Specifically, we aim to synthesize barrier functions that can serve as a safety certificate by providing a lower bound for P_{s}. To achieve this, we rely on recent results for local convex relaxation of an NN to find piece-wise linear over- and under-approximations of f^w. This allows us to formulate the problem of finding a stochastic barrier function for System (1) as an SOS optimization problem.

Our (minimally-invasive) control synthesis method (Problem 2) focuses on injecting controls to System (3) only in the regions of the state space that potentially contribute to unsafety of the system. To this end, we exploit two main properties of the barrier function: the required (super-) martingale property and the convexity property of SOS polynomials. The former allows us to design a method of identifying the "unsafe" regions, and the latter enables us to formulate a Linear Programming (LP) optimization problem to find controllers that robustly minimize the probability of unsafety.

3 Preliminaries

In this section, we provide a brief background on stochastic barrier functions and non-negative polynomials, which are central components of our framework.

3.1 Stochastic Barrier Functions

Consider stochastic process $\mathbf{x}_{k+1} = F(\mathbf{x}_k, \mathbf{v}_k)$, where $\mathbf{x}_k \in X \subseteq \mathbb{R}^n$, $\mathbf{v}_k \in V \subseteq \mathbb{R}^m$ is a well-defined stochastic process, and $F : X \times V \to X$ is a locally Lipschitz continuous function. Then, given initial set $X_0 \subset X$, safe set $X_{\mathrm{s}} \subseteq X$, and unsafe set $X_{\mathrm{u}} \subset X$, a twice differentiable function $B : X \to \mathbb{R}$ is called a *stochastic barrier function* or simply barrier function if the following conditions hold for $\alpha \geq 1$ and $\beta \in [0, 1]$

$$B(x) \geq 0 \quad \forall x \in X \quad \text{(4a)} \qquad\qquad B(x) \geq 1 \quad \forall x \in X_{\mathrm{u}} \quad \text{(4c)}$$

$$B(x) \leq \eta \quad \forall x \in X_0 \quad \text{(4b)} \qquad E[B(F(x,v)) \mid x] \leq B(x)/\alpha + \beta \quad \forall x \in X_{\mathrm{s}} \quad \text{(4d)}$$

If such a B exists, then for any $N \in \mathbb{N}_{\geq 0}$, it follows, assuming $\alpha = 1$ without loss of generality, that

$$\Pr[\forall k \leq N, \ \mathbf{x}_k \in X_{\mathrm{s}} \mid \mathbf{x}_0 \in X_0] \geq 1 - (\eta + \beta N). \tag{5}$$

Intuitively, Conditions (4a)-(4d) allow one to use non-negative super-martingale inequalities to bound the probability that the system reaches the unsafe region. This results in the probability bound in (5) (see [26, Chapter 3, Theorem 3]). A major benefit of using barrier functions for safety analysis is that while Conditions (4b)-(4d) are static properties, they allow one to make probabilistic statements on the time evolution of the system without the need to evolve the system over time.

3.2 Non-negative Polynomials

Synthesis of barrier functions can be formulated as a nonlinear optimization problem. To solve it efficiently via convex programming, we formulate the problem using SOS polynomials, which is a sufficient condition for polynomial non-negativity [27]. This formulation allows the use of efficient semidefinite programming solvers, which have polynomial worst-case complexity [28].

Definition 1 (SOS Polynomial). *A multivariate polynomial $\lambda(x)$ is a Sum-Of-Squares (SOS) for $x \in \mathbb{R}^n$ if there exists some polynomials λ_i, $i = 1, \ldots, r$, for some $r \in \mathbb{N}$, such that $\lambda(x) = \sum_{i=1}^r \lambda_i^2(x)$. If $\lambda(x)$ is an SOS, then $\lambda(x) \geq 0$ for all $x \in \mathbb{R}^n$. The set of all SOS polynomials is denoted by Λ.*

Now, consider a basic closed semi-algebraic set $X = \{x \in \mathbb{R}^n \mid h_i(x) \geq 0 \ \forall i \in \{1, \ldots, l\}\}$, which is defined as the intersection of l polynomial ($h_i(x)$) inequalities. Then, Proposition 1 allows one to enforce non-negativity on these semi-algebraic sets.

Proposition 1 (Putinar's Certificate, [29]). *Let $X \subset \mathbb{R}^n$ be a compact basic semi-algebraic set defined as $X = \{x \in \mathbb{R}^n \mid h_i(x) \geq 0 \ \forall i \in \{1, \ldots, l\}\}$, where $h_i(x)$ is a polynomial. Then, polynomial $\gamma(x)$ is non-negative on X if, for some $\lambda_i(x) \in \Lambda$, $\gamma(x) = \lambda_0(x) + \sum_{i=1}^l \lambda_i(x)h_i(x)$.*

Corollary 1. *Let $\gamma(x)$ be a non-negative polynomial on $X = \{x \in \mathbb{R}^n \mid h_i(x) \geq 0 \ \forall i \in \{1, \ldots, l\}\}$, and $h(x)$ be an l-dimensional vector with $h_i(x)$ as its i-th dimension for all $i \in \{1, \ldots, l\}$. Further, let $\mathcal{L}$ be an l-dimensional vector of SOS polynomials, i.e., the i-th element of $\mathcal{L}$ is $\lambda_i(x) \in \Lambda$. Then, it holds that $\gamma(x) - \mathcal{L}^T h(x) \in \Lambda$.*

4 Barrier Functions for Neural Network Dynamical Models

In this section, we introduce a formulation for synthesizing a barrier function that guarantees safety of System (1). Our approach relies on the local convex relaxations of f^w, which we use to formulate an SOS optimization problem that generates a valid barrier function for System (1).

4.1 Local Convex Relaxation of NNDMs

Our approach uses a convex relaxation of f^w. Specifically, we utilize the recent relaxation results for NNs [18] to compute piece-wise linear functions that under- and over-approximate f^w. The advantage of this method is that it is computationally efficient and produces fairly tight bounds for

small neighborhoods. To this end, we first partition X_{s} to a finite set of regions $Q = \{q_1, ..., q_{|Q|}\}$, where $q_i \subseteq X_{\mathrm{s}}$ for all $i \in \{1, \ldots, |Q|\}$ is a basic (closed) semi-algebraic set. Such a decomposition can be obtained using, e.g., a grid or half-spaces. Then, for every $q \in Q$, we compute linear bounds

$$\underline{f}_q(x) = \underline{A}_q x + \underline{b}_q, \quad \overline{f}_q(x) = \overline{A}_q x + \overline{b}_q \quad \text{such that} \quad \forall x \in q, \quad \underline{f}_q(x) \le f^w(x) \le \overline{f}_q(x). \quad (6)$$

The above linear functions can be produced in several ways. One method is to forward propagate the bounds in the input domain of $f^w(x)$ using Interval Bound Propagation (IBP) based approaches [30, 31]. These approaches are efficient but generally lead to loose approximations. Another method is Linear Relaxation based Perturbation Analysis (LiRPA) algorithms, which perform a symbolic back propagation followed by a forward evaluation of the bounds [32, 33, 18]. Such LiRPA-based algorithms produce tighter bounds but can be slow. In our implementation, we use the off-the-shelf toolbox [33] to compute $\overline{f}_q(x)$ and $\underline{f}_q(x)$ in (6) for its trade-off between speed and accuracy.

4.2 Barrier Function (Certificate) Synthesis

Recall that a valid barrier function $B(x)$ must satisfy Conditions (4b)-(4d). With the above local linear bounds for f^w, Condition (4d) can be redefined over the partitioned regions in Q. That is, Condition (4d) can be replaced with the following constraints: for each $q \in Q$,

$$E[B(f_q^w(x) + \mathbf{v}) \mid x] \le B(x)/\alpha + \beta \quad \text{and} \quad \underline{f}_q(x) \le f_q^w(x) \le \overline{f}_q(x) \quad \forall x \in q. \quad (7)$$

In this formulation, $f_q^w(x)$ is a (free) variable that is bounded by two linear functions, avoiding the non-convexity of $f_q^w(x)$. Consequently, an efficient evaluation of the constraint becomes feasible.

Remark 3. *Note that there are in fact three constraints in (7) (one super-martingale condition on B and two bounds on $f_q^w(x)$). If they are treated separately, it leads to a conservative barrier function because each constraint needs to hold for all $x \in q$ separately. A less conservative approach is to pose all constraints for the same x at the same time. In Theorem 1, we show how this can be achieved.*

As commonly practiced in the literature, e.g., [14, 8, 34], we restrict our search for a $B(x)$ to the class of polynomial functions. Then, the degree of the polynomial becomes a design parameter. In the following lemma, we show that the degree of the polynomial must be at least 2 for the polynomial to be a valid candidate barrier function for System (1). That is, there does not exist a linear function satisfying Conditions (4b)-(4d) for $\eta < 1$, which implies that the lower bound in (5) is trivially 0.

Lemma 1. *For $k \in \mathbb{R}^n$ and $c \in \mathbb{R}$, consider linear function $B(x) = k^T x + c$. Then, $B(x)$ satisfies Conditions (4b)-(4d) iff $k = 0$ and $c = 1$.*

The proof of Lemma 1 is in Appendix A.1. It relies on the fact that a linear function is either a constant or always eventually violates Condition (4a). Having established that the polynomial degree must be at least 2, we further restrict $B(x)$ to SOS polynomials. Theorem 1 shows a formulation of the barrier function synthesis problem as an SOS program, which can be solved in polynomial time. In particular, the theorem shows how to formulate Conditions (4b)-(4c) and (7) as SOS constraints.

Theorem 1 (NNDM Barrier Certificate). *Consider SOS polynomial function $B(x)$, and safe set $X_{\mathrm{s}} = \{x \in \mathbb{R}^n : h_{\mathrm{s}}(x) \ge 0\}$, initial set $X_0 = \{x \in \mathbb{R}^n : h_0(x) \ge 0\}$, unsafe set $X_{\mathrm{u}} = \mathbb{R}^n \setminus X_{\mathrm{s}} = \{x \in \mathbb{R}^n : h_{\mathrm{u}}(x) \ge 0\}$, and partition region $q = \{x \in \mathbb{R}^n : h_q(x) \ge 0\}$ for all $q \in Q$. Let $\mathcal{L}_{\mathrm{s}}(x)$, $\mathcal{L}_0(x)$ and $\mathcal{L}_{\mathrm{u}}(x)$ be vectors of SOS polynomials with the same dimensions as h_{s}, h_0, and h_{u}, respectively. Likewise, let $\mathcal{L}_q(x)$ and $\mathcal{L}_{q,y}(x)$ be vectors of SOS polynomials with the same dimension as h_q. Then, a stochastic barrier certificate $B(x)$ for System (1) with time horizon $N \in \mathbb{N}_{\ge 0}$ can be obtained by solving the following SOS optimization problem*

$$\min_{\beta, \eta} \quad \eta + \beta N \quad \text{subject to:}$$

$$B(x) \in \Lambda, \quad (8a)$$

$$-B(x) - \mathcal{L}_0^T(x) h_0(x) + \eta \in \Lambda, \quad (8b)$$

$$B(x) - \mathcal{L}_{\mathrm{u}}^T(x) h_{\mathrm{u}}(x) - 1 \in \Lambda, \quad (8c)$$

$$E[B(y + v) \mid x] \le B(x)/\alpha + \beta - \mathcal{L}_q^T(x) h_q(x) -$$
$$\mathcal{L}_{q,y}^T(x)(\overline{f}_q(x) - y)(y - \underline{f}_q(x)) \in \Lambda \quad \forall q \in Q \quad (8d)$$

which guarantees safety probability $P_{\mathrm{s}}(X_{\mathrm{s}}, X_0, N) \ge 1 - (\eta + \beta N)$.

221 The proof of Theorem 1 is in Appendix A.2. The key idea is that an SOS polynomial function
222 is a linear combination of monomials for a given degree $m \geq 2$. Hence, Condition (4a) can be
223 written as Condition (8a). Using Corollary 1, we can enforce non-negativity of a polynomial function
224 on semi-algebraic sets using Lagrangian multipliers $\mathcal{L}$ in terms of vectors of SOS polynomials
225 and the corresponding polynomials in $h(x)$. Hence, we can write Conditions (4b)-(4c) for basic
226 closed semi-algebraic sets as Conditions (8b)-(8c), which are guaranteed to only hold over their
227 corresponding basic closed semi-algebraic sets. Finally, Condition (8d) is derived from Condition
228 (4d) for System (1) by using a similar technique to restrict the super-martingale inequality to the
229 corresponding local linear bounds of $f^w(x)$ in (7) in each region $q \in Q$. Therefore, if a feasible
230 solution exists, the SOS optimization problem in Theorem 1 yields a safety certificate for System (1).

5 Minimally-invasive Barrier Controller Synthesis

232 Here, we focus on Problem 2. The assumption is that a barrier function $B(x)$ is synthesized for
233 System (1) via Theorem 1, which has a safety probability that is lower than threshold δ_s. Hence, our
234 goal is to design a feedback controller for System (3) that guarantees safety. Ideally, this controller is
235 minimal in the interruptions it causes to the system evolution. Below, we first provide a definition for
236 a minimally-invasive controller, and then show how it can be synthesized.

237 Recall that a feedback controller $\pi : \mathbb{R}^n \to U$ is a function that assigns a control value to each state.
238 Let $X_\pi \subseteq X_s$ be the set of states, to which π assigns non-zero control values, i.e., $\pi(x) \neq 0$ for all
239 $x \in X_\pi$. Further, we denote by $V(X)$ the volume of set $X \subset \mathbb{R}^n$. Then, we say a feedback controller
240 π is *minimally-invasive* if it minimizes $V(X_\pi)$. Moreover, we say π is *ϵ-minimally-invasive* if
241 $V(X_\pi) - \min_{\pi'} V(X_{\pi'}) \leq \epsilon$. In this work, we seek *ϵ-minimally-invasive* controller with an arbitrary
242 small ϵ. To this end, we make the observation that if the discretization of X_s discussed in Section 4.1
243 is uniform, $V(X_\pi)$ for a given π is directly proportional to the number of discrete regions in Q to
244 which π applies a non-zero control. We denote by Q_π the set of such regions. Then, a feedback
245 controller π that minimizes $|Q_\pi|$ is *ϵ-minimally-invasive*, and as the discretization becomes finer, ϵ
246 monotonically approaches zero. Therefore, given barrier function $B(x)$, our first goal is to identify
247 the minimal number of regions in Q that require a non-zero controller to guarantee safety. Recall
248 that, given $B(x)$, we require the safety probability to be $1 - (\eta + \beta N) \geq \delta_s$. While η is related to
249 the value of $B(x)$ in the initial set, β is a compensation needed to turn a sub-martingale inequality to
250 a super-martingale inequality (Condition (8d) or (4d)). We can compute this compensation for each
251 region $q \in Q$ by evaluating Condition (8d) or (4d). Let β_q denote the computed compensation term
252 for region q. Then, region q requires a non-zero controller if $\beta_q > (1 - \delta_s - \eta)/N$.

253 We now turn our attention to designing a controller that reduces β_q. In our approach, we tap directly
254 into the convex nature of the SOS polynomial $B(x)$. Lemma 2 shows that a controller that drives
255 System (3) to a point closest to the minimum point of $B(x)$, also robustly minimizes β_q.

Lemma 2. *Given a convex barrier function $B(x)$, let $x^* = \arg\min_x B(x)$, $x' = f^w(x) + gu + v$*
257 *as in System (3), and $\beta_q \geq E[B(x') \mid x] - B(x)$ for all x in region $q \subseteq X_s$. Then, the controller*
258 *that minimizes $\|\tilde{x}' - x^*\|_1$, where $\tilde{x}' = f^w(x) + gu$, also robustly minimizes β_q, i.e., minimizes an*
259 *upper bound of β_q.*

Proof. Given that B is a polynomial and x' is affine in both u and v, we can write $E[B(x') \mid x] =$
261 $B(E[x'] \mid x) + \xi$, where ξ, called Jensen's gap, is a non-negative polynomial of x, u, and moments
262 $E[v^d]$, where $2 \leq d \leq m$ is an even integer. This polynomial can always be upper bounded by a
263 constant $\bar{\xi} \geq 0$ for bounded x and u. Then, with $\tilde{x}' = E[x']$, the min β_q is upper bounded by

$$\min_u \beta_q \leq \min_u \max_{x \in q} \left(B(\tilde{x}' \mid x) - B(x) + \bar{\xi} \right) \leq \min_u \max_{x \in q} B(\tilde{x}' \mid x) - \min_{x \in q} B(x) + \bar{\xi}.$$

264 The second inequality holds because B is a non-negative convex function. Since $B(\tilde{x}')$ monotonically
265 decreases as $\tilde{x}' \to x^*$, $\max_{x \in q} B(\tilde{x}' \mid x)$ is minimized by $\arg\min_u \|f^w(x) + gu - x^*\|_1$ $\forall x \in q$. $\square$

266 We note that, in the proof of Lemma 2, the moments $E[v^d]$ appear in every monomial of polynomial
267 ξ (Jensen's gap). Hence, for v with variance less than 1, Jensen's gap is very small and vanishes
268 as the variance becomes smaller, making the controller in the lemma optimal. Furthermore, even
269 though Lemma 2 uses L_1 norm for $\|\tilde{x}' - x^*\|_1$, the statement also holds for L_2 and L_∞ norms. We
270 specifically use L_1 to be able to compute the controller via an LP as stated in Theorem 2.

Theorem 2 (Controller Synthesis). *Consider convex barrier function $B(x)$ and its minimum argument x^* as specified in Lemma 2, and linear bounds $\underline{f}_q(x)$ and $\overline{f}_q(x)$ for $f^w(x)$ in region $q \subseteq \mathbb{R}^n$ as in (6). A controller that robustly minimizes β_q for System (3) is the solution to the following LP optimization problem, where $\theta = (\theta_1, \theta_2, \ldots, \theta_n)$, and the vector inequality relations are element-wise:*

$$\min_{\theta_i, z', z'', u} \sum_{i=1}^{n} \theta_i \quad subject\ to:$$

$$y - x^* \leq \theta \quad (10a) \qquad\qquad y \geq \underline{f}_q(z) + gu \quad (10d)$$

$$x^* - y \leq \theta \quad (10b) \qquad\qquad y \leq \overline{f}_q(z) + gu \quad (10e)$$

$$z = z' - z'' \quad (10c) \qquad\qquad z', z'' \geq 0,\ z \in q,\ u \in U \quad (10f)$$

The proof of Theorem 2 is in Appendix A.3. It directly follows from Lemma 2, requiring the minimization of $\|\tilde{x}' - x^*\|_1$, which can be reformulated as the minimization of the sum of the components of θ with linear constraints (10a) and (10b). Condition (10c) allows free variable z to be expressed with two non-negative decision variables z' and z''. Finally, Conditions (10d) and (10e) enforce the linear bounds on the dynamics of System (3) in region q.

Remark 4. *Note that the controller obtained via the LP is a vector of scalars for each region q. This leads to a feedback controller π over the regions $q \in Q$. It may be possible to express u as SOS polynomials, which would be a feedback controller over x. That could potentially provide a tighter bound for β_q at the cost of extra computation since SOS optimization is more expensive than LP.*

5.1 Control Synthesis Algorithm

Our framework for computing a barrier certificate as well as minimally-invasive controllers for NNDMs is outlined in Algorithm 1. First, the safe set X_s is uniformly partitioned into a finite set of regions $|Q|$, and bounds of f^w are computed as piece-wise linear functions per Section 4.1. Then, for a given polynomial degree $m \geq 2$, a barrier certificate along with its corresponding safety probability bound are computed in accordance with Theorem 1. If this probability is greater than or equal to δ_s, the NNDM is certified, and the algorithm terminates. Otherwise, it iteratively searches for ϵ-minially-invasive controllers. In each iteration, upper bound $\overline{\eta}$ is set on η (initially with its largest possible value $1 - \delta_s$ and then reduced by $\Delta\eta > 0$), and a new barrier function is computed with this bound, where the objective function is minimization of β to ensure ϵ-minimally-invasive controllers. Then, for each $q \in Q$, β_q is checked against the threshold. If exceeding, a controller is computed for q in accordance with Theorem 2. This process repeats until a satisfactory P_s is obtained or $\overline{\eta} < 0$.

We highlight two important properties of this algorithm: it clearly terminates in finite time, and if $P_s \geq \delta_s$, the synthesized controller is correct-by-construction.

Algorithm 1: NNDM Controller Synthesis

    Input   : NN f^w, safe set X_s, noise covariance R, polynomial degree m, threshold δ_s, and step size $\Delta\eta$

    Output: feedback controller π, and safety probability bound P_s

1 $Q, \underline{f}_q, \overline{f}_q \leftarrow$ PARTITIONANDCOMPUTEBOUNDS(X_s, f^w), $k \leftarrow 0$

2 $\eta, \beta, B(x) \leftarrow$ COMPUTEBARRIERCERTIFICATE$(Q, \underline{f}_q, \overline{f}_q, R, m)$ // Theorem 1

3   **while** $P_s < \delta_s$ *and* $\overline{\eta} > 0$ **do**

4 $\overline{\eta} \leftarrow 1 - (\delta_s - k\Delta\eta)$, $k \leftarrow k + 1$ // set upper bound on η

5 $\eta, \beta, B(x) \leftarrow$ COMPUTEBARRIERCERTIFICATE$(Q, \underline{f}_q, \overline{f}_q, R, m, \overline{\eta})$ // Theorem 1

6   **for** $q \in Q$ **do**

7 $\pi \leftarrow (q, 0)$

8   **if** $\beta_q > (1 - \delta_s - \eta)/N$ **then**

9 $\pi \leftarrow (q, u_q) \leftarrow$ COMPUTECONTROL$(\arg\min B(x), q, \underline{f}_q, \overline{f}_q, U)$ // Theorem 2

10 $\beta_q \leftarrow$ UPDATEBETA$(B(x), u_q)$ // evaluate Condition (8d)

11 $P_s \leftarrow 1 - (\eta + N \max_{q \in Q} \beta_q)$

12 **return** π, P_s

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6 Case Studies

We demonstrate the efficacy of our framework on several case studies. We first show that our verification method is able to produce non-trivial safety guarantees for NNDMs trained using imitation learning techniques [35] from data gathered by rolling out an RL agent trained with state-of-the-art RL techniques. We then show that our control approach significantly increases the safety probability.

Models We trained the Pendulum and Cartpole agents from the OpenAI gym environment. For training, we collected input-output states of the system under an expert controller, either a look-up table or an NN controller. We picked the best controller available for the agents from the OpenAI Leaderboard [36]. We also trained four-dimensional husky models. We initially trained a NNDM along with NN controller on the trajectory data generated from a physics simulator as outlined in [9]. We then used this NNDM to generate input-output data and trained closed-loop NNDMs as specified above. In total, we trained 8 fully-connected multi-layer perceptron NNDM architectures with upto 5 hidden layers and 256 neurons, each with ReLU activation function. The dimensionality and architecture information for each model is shown in Table 1.

Implementation and Experimental Setup We implemented our algorithms in Python and Julia. For collecting data and training, we used the TensorFlow framework [37]. To compute the linear approximations of the NNDMs, we utilized α-CROWN [33]. For the optimization, we used Julia's *SumOfSquares.jl* package [38, 39]. The optimizations are all performed single-threaded on a computer with 3.9 GHz 8-Core CPU and 64 GB of memory. The code is available in the supplementary material.

All the results are shown in Table 1. The computation times indicate the time to synthesize the barrier and the controller using Theorems 1 and 2 respectively. The barrier degree is 4 for all the case studies. Below we provide a brief discussion on the results. For more in-depth discussions and details on the discretization and hyperparameters, see Appendix B in the supplementary material.

Table 1: Benchmark results for barrier certificate and control synthesis on various NNDMs. Threshold $\delta_{\mathrm{s}} = 0.95$ was used for the minimally-invasive control synthesis. Architecture $h \times [r]$ denotes a NNDM with h hidden layers, each with r neurons.

Model	Dim.	Architecture	$|Q|$	Verification			Control		
				β	P_{s}	Time (min)	β	P_{s}	Time (min)
Pendulum	2	$1 \times [64]$	120	0.002	0.997	0.09	-	-	-
			240	0.002	0.997	0.14	-	-	-
			480	0.003	0.997	0.42	-	-	-
		$2 \times [64]$	120	0.191	0.785	0.06	10^{-6}	0.999	0.05
			240	0.153	0.844	0.14	10^{-6}	0.998	0.12
			480	0.075	0.922	0.46	10^{-6}	0.997	0.34
		$3 \times [64]$	120	0.384	0.603	0.06	10^{-6}	0.986	0.04
			240	0.288	0.707	0.14	10^{-6}	0.995	0.13
			480	0.208	0.791	0.44	10^{-6}	0.998	0.46
		$5 \times [64]$	480	0.735	0.172	0.45	10^{-6}	0.951	0.36
			960	0.603	0.302	1.56	10^{-6}	0.951	1.59
			1920	0.421	0.504	6.61	10^{-6}	0.951	5.33
Cartpole	4	$1 \times [128]$	960	0.623	0.375	17.99	10^{-6}	0.998	13.35
			1920	0.414	0.577	65.14	10^{-6}	0.992	35.09
			3840	0.382	0.602	222.84	10^{-6}	0.984	160.99
		$2 \times [128]$	960	0.811	0.181	19.84	10^{-6}	0.993	13.30
			1920	0.576	0.402	51.36	10^{-6}	0.997	36.95
			3840	0.664	0.333	239.21	10^{-6}	0.997	162.57
Husky	4	$1 \times [256]$	900	0.431	0.366	16.65	10^{-6}	0.951	11.81
			1800	0.425	0.382	53.55	10^{-6}	0.951	33.73
			2250	0.367	0.494	81.32	10^{-6}	0.951	46.49
			4800	0.347	0.538	345.28	10^{-6}	0.951	304.67
		$2 \times [256]$	1800	0.544	0.011	74.69	10^{-6}	0.800	35.22
			2250	0.502	0.222	74.83	10^{-6}	0.951	51.76
			4800	0.415	0.384	406.69	10^{-6}	0.951	306.41

Figure 1: Comparison of β_q values for the Pendulum NNDM with 2 layers and 64 neurons per layer before and after applying the controller for various $|Q|$. Shade of gray correspond to the values of β_q.

Verification As expected, the safety probability almost always increases with $|Q|$ (finer partitions) since the linear bounds for f^w become more accurate. The only exception is Cartpole 2-layer NNDM, which decreases in P_s from $|Q|$ of 1920 to 3840. After a closer inspection, we observed that the linear bounds computed for 3840 discretization are looser than 1920, resulting in lower P_s. This is possible because, to compute the linear bounds, α-CROWN uses a projected gradient descent approach, which does not guarantee to converge to global optima. Another trend in the results is that finer discretizations come at the cost of an increased computation time. For instance, on the 2-layer Husky model, increasing $|Q|$ from 1800 to 4800 increases P_s from 0.011 to 0.384, at the price of a more than 5-fold increase in the computation time.

Control Synthesis We observe that for all NNDMs, whose P_s is lower than $\delta_s = 0.95$, except one case, the controller synthesis algorithm is able to produce $P_s \geq \delta_s$, showing the effectiveness of our control method. The exception is for Husky with 2-hidden layers and $|Q| = 1800$, where the controller increases P_s from 0.011 to 0.800 (73-fold increase) yet cannot meet the 0.95 threshold. In fact, for all the models, the controller reduces β to 10^{-6}. Figure 1 illustrates the minimally-invasive aspect of the controllers. It shows the β_q values for the Pendulum NNDM with 2 layers before and after the application of the controllers with three different sizes of $|Q|$. For $|Q| = 120$, all the regions require controllers, but for $|Q| = 240$ and 480, respectively, only 69.2% and 23.75% of the regions require controllers; hence, the minimally-invasive controller focuses specifically on those regions and does not affect those with already-small β_q.

7 Conclusion

In this work, we introduced a methodology for providing safety guarantees for NNDMs using stochastic barrier functions. We showed that the problem of finding a stochastic barrier function for a NNDM can be relaxed to an SOS optimization problem. Furthermore, we derived a novel framework for synthesis of an affine controller with the goal of maximizing the safety probability in a minimally-invasive manner. Experiments showcase that, along with a minimally-invasive controller, we can guarantee safety for various standard reinforcement learning problems.

A bottleneck to our approach is the discretization step, which is required to obtain piece-wise linear under- and over-approximations of the neural network, and can be expensive for higher dimensional spaces. A future direction is to relax the reliance on this step. Another future direction is to consider a more expressive class of barrier functions than SOS polynomials.

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