Safety Guarantees for Neural Network Dynamic Systems via Stochastic Barrier Functions

Abstract

Neural Networks (NNs) have been successfully employed to represent the state evolution of complex dynamical systems. Such models, referred to as NN dynamic models (NNDMs), use iterative noisy predictions of NN to estimate a distribution of system trajectories over time. Despite their accuracy, safety analysis of NNDMs is known to be a challenging problem and remains largely unexplored. To address this issue, in this paper, we introduce a method of providing safety guarantees for NNDMs. Our approach is based on stochastic barrier functions, whose relation with safety are analogous to that of Lyapunov functions with stability. We first show a method of synthesizing stochastic barrier functions for NNDMs via a convex optimization problem, which in turn provides a lower bound on the system's safety probability. A key step in our method is the employment of the recent convex approximation results for NNs to find piece-wise linear bounds, which allow the formulation of the barrier function synthesis problem as a sum-of-squares optimization program. If the obtained safety probability is above the desired threshold, the system is certified. Otherwise, we introduce a method of generating controls for the system that robustly minimize the unsafety probability in a minimally-invasive manner. We exploit the convexity property of the barrier function to formulate the optimal control synthesis problem as a linear program. Experimental results illustrate the efficacy of the method. Namely, they show that the method can scale to multi-dimensional NNDMs with multiple layers and hundreds of neurons per layer, and that the controller can significantly improve the safety probability.

1 Introduction

Safety is a major concern for autonomous dynamical systems, especially in *safety-critical* applications. Examples include UAVs, autonomous cars, and surgical robots. To this end, various methods are developed to provide safety guarantees for such systems [1–3]. Most of these methods, however, assume (simple) analytical models, which are often unavailable or too expensive to obtain due to complexity (nonlinearities) in real-world systems [4]. This has given rise to recent efforts to harvest the representational power of *neural networks* (NNs) to predict how the state of a system evolves over time, referred to as *NN dynamic models* (NNDMs) [5, 6]. Despite their efficient and accurate representation of realistic dynamics, it has been difficult to make formal claims about the behavior of NNDMs over time. This poses a major challenge in safety-critical domains: *how to provide safety guarantees for NNDMs?*

In this paper, we propose a framework to provide safety certificates for NNDMs. The approach is based on stochastic barrier functions [7, 8], whose role for safety is analogous to that of Lyapunov functions for stability, and which can guarantee forward invariance of a safe set. Specifically, we show a method of synthesizing barrier functions for NNDMs via an efficient convex optimization method, which in turn provides a lower bound on the safety probability of the system trajectories.

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The benefit of our approach is that it does not require unfolding of the NNDM to compute this probability; instead, we base our reasoning on non-negative super-martingale inequalities to bound the probability of leaving the safe set. The key to this is to utilize the recent convex approximation results for NNs to obtain piece-wise linear bounds, and a formulation of the optimization problem that enables the use of efficient solvers. If the obtained safety probability is above a desired threshold, the system can be certified. For the cases that such safety certificates cannot be provided, we introduce a method of generating controls for the system, which robustly maximizes the safety probability in a minimally-invasive manner. We exploit the convexity property of the barrier function to formulate the control synthesis problem as a linear program. Experimental results on various NNDMs trained using state-of-the-art RL techniques [9] illustrate the efficacy of the method. Namely, they show that the method can scale to multi-dimensional models with various layers and hundreds of neurons per layer, and that the controller can significantly improve the safety probability.

In summary, this paper makes the following main contributions: (i) a method of computing a lower bound for the safety probability of NNDMs via stochastic barrier functions and efficient nonlinear optimization, thereby providing safety guarantees, (ii) a correct-by-construction control synthesis framework for NNDMs to ensure safety while being *minimally-invasive*, i.e., intervening only when needed, (iii) demonstration of the efficacy and scalability of the framework in benchmarks on various NNDMs with multiple architectures, i.e., multiple hidden layers and hundreds of neurons per layer.

Related Work In recent years, NNs have been successfully employed to model complex dynamical systems [5, 6]. Such models use iterative noisy predictions of an NN to estimate the distribution of the system over time. They provide means of predicting state evolution of the systems that have no analytical form, e.g., soft robotics [10]. In addition, because of their ability to quickly unfold trajectories, NNDMs have been used to enhance controller training in Reinforcement Learning (RL) frameworks [9, 11]. In these works, specifically, a NNDM of the system is trained in closed-loop with an NN controllers, which can be concatenated in a single NN representing the dynamics of the closed-loop system. Despite their popularity and increased usage, little work has gone to safety analysis of NNDMs. This paper aims to fill this gap.

Barrier functions provide a formal methodology to prove safety of dynamical systems [12, 13]. The state-of-the-art to find stochastic barrier functions is arguably to rely on Sum-of-Squares (SOS) optimization [8]. Barrier certificates obtained via SOS optimization have been studied for deterministic, uncontrolled nonlinear systems [14], and control-affine stochastic systems [15, 8], as well as for hybrid models [12, 13]. Unlike the existing work that focus on closed-form models, this work is the first to consider barrier function for models represented by NNs. In addition, the control synthesis formulation in this paper relaxes the limitations of the existing SOS methods, which generally leads to a non-convex optimization problem [8], by reducing it to a simple linear program.

Many recent studies have focused on the verification of NNs. However, most of them only aim to provide certification guarantees against adversarial examples [16], with approaches ranging from SMT [17], and convex relaxations [18] to linear programming [19]. Such methods, because of their local nature, cannot be directly employed to perform safety analysis on the entire state space and over the trajectories as required for NNDMs. Other methods have been employed for trajectory-level properties of NNs based on solution of the Bellman equations [20], formal (finite) abstractions [21], and mixed integer programming [22], which does not support noisy dynamics. In contrast, our approach supports additive noise and only requires to check static properties on the NN by showing that there exists a function whose composition with the NN forms a stochastic barrier function.

2   Problem Formulation

In this work, we are interested in providing safety guarantees for dynamical systems that are described by NNs. Specifically, we consider the following discrete-time stochastic process whose time evolution is given by iterative predictions of an NN

$$\mathbf{x}_{k+1} = f^w(\mathbf{x}_k) + \mathbf{v}_k, \qquad k \in \mathbb{N}, \quad \mathbf{x}_k, \mathbf{v}_k \in \mathbb{R}^n, \tag{1}$$

where f^w is a trained, fully-connected, feed-forward NN with general continuous activation functions, and w represents the maximum likelihood weights. Term $\mathbf{v}_k$ is a random variable modelling an additive noise of stationary Gaussian distribution with zero mean and covariance matrix $R \in \mathbb{R}^{n \times n}$, i.e., the probability density function of $\mathbf{v}_k$ is $\mathcal{N}(\bar{x} \mid 0, R)$.

Remark 1. *System* (1) *can be viewed as a closed-loop system, where f^w is the composition of a NNDM (that represents the open-loop system) with an NN controller. NNDMs such as the one in System* (1) *are increasingly employed in robotics for both model representation and NN controller training. These models are often obtained by state-of-the-art model-based RL techniques (e.g., [9, 23, 11]) or perturbations (step response) of the actual system (e.g., [10]).*

The evolution of System (1) can be characterized by the predictive posterior distribution $p_{\mathbf{x}}(\bar{x} \mid x)$, which describes the probability density of $\mathbf{x}_{k+1}$ when $\mathbf{x}_k = x \in \mathbb{R}^n$. This distribution is induced by noise $\mathbf{v}_k$ and hence is Gaussian with mean $f^w(x)$ and covariance R, i.e., $p_{\mathbf{x}}(\bar{x} \mid x) = \mathcal{N}(\bar{x} \mid f^w(x), R)$. For a subset of states $X \subseteq \mathbb{R}^n$ and (initial) state $x \in \mathbb{R}^n$, we call $T(X \mid x) = \int_X p_{\mathbf{x}}(\bar{x} \mid x)d\bar{x}$ the *stochastic kernel* of System (1). From the definition of T it follows that, for a given initial condition $\mathbf{x}_0 = x_0 \in \mathbb{R}^n$, $\mathbf{x}_k$ is a Markov process with a well-defined probability measure Pr uniquely generated by the stochastic kernel T [24, Proposition 7.45]. For $X_0, X_k \subseteq \mathbb{R}^n$, this probability measure is inductively defined as

$$\mathrm{Pr}[\mathbf{x}_0 \in X_0] = \mathbf{1}_{X_0}(x_0), \qquad \mathrm{Pr}[\mathbf{x}_k \in X_k \mid \mathbf{x}_{k-1} = x] = T(X_k \mid x),$$

where indicator function $\mathbf{1}_{X_0}(x_0) = 1$ if $x_0 \in X_0$, otherwise 0. The definition of Pr allows one to make probabilistic statements over the trajectories of System (1). In this work, we are particularly interested in two main problems: safety certification and safe controller synthesis for System (1).

Problem 1 (Safety Certificate). *Let $X_{\mathrm{s}} \subset \mathbb{R}^n$ and $X_0 \subseteq X_{\mathrm{s}}$ be basic closed semi-algebraic sets representing respectively the safe set and the set of initial states. Denote the probability that the system initialized in X_0 remains in the safe set in the next $N \in \mathbb{N}_{\geq 0}$ steps by*

$$P_{\mathrm{s}}(X_{\mathrm{s}}, X_0, N) = Pr[\forall x_0 \in X_0, \forall k \leq N, \mathbf{x}_k \in X_{\mathrm{s}}]. \tag{2}$$

Then, given threshold $\delta_{\mathrm{s}} \in [0, 1]$, provide a safety certificate for System (1) *by proving that $P_{\mathrm{s}}(X_{\mathrm{s}}, X_0, N) \geq \delta_{\mathrm{s}}$.*

Problem 1 seeks to compute the probability that $\mathbf{x}_k$ remains within a safe set, e.g., it avoids obstacles or any undesirable states. Computation of this probability is particularly challenging because System (1) is stochastic and NN (f^w) is generally a non-convex function. The assumption that X_{s} and X_0 are basic closed semi-algebraic sets is not limiting, as these sets are defined as an intersection of a finite number of polynomial inequalities and, for instance, include convex polytopes and spectrahedra [25].

In the case that a safety certificate cannot be generated, we are interested in synthesizing a controller that *by-design* guarantees safety. To this end, we extend System (1) with an affine input

$$\mathbf{x}_{k+1} = f^w(\mathbf{x}_k) + g\mathbf{u}_k + \mathbf{v}_k, \tag{3}$$

where $g : \mathbb{R}^c \to \mathbb{R}^n$ is given, $\mathbf{u}_k \in U \subset \mathbb{R}^c$ is a control action, and U is bounded. Our goal is to synthesize a feedback controller $\pi : \mathbb{R}^n \to U$ such that $\mathbf{u}_k = \pi(\mathbf{x}_k)$ guarantees safety of System (3).

Problem 2 (Safe Control Synthesis). *Let $X_{\mathrm{s}} \subset \mathbb{R}^n$ and $X_0 \subseteq X_{\mathrm{s}}$ be basic closed semi-algebraic sets representing respectively the safe set and the set of initial states, and $N \in \mathbb{N}_{\geq 0}$ a time horizon. Then, given threshold δ_{s}, synthesize a feedback controller π such that $P_{\mathrm{s}}(X_{\mathrm{s}}, X_0, N, \pi) \geq \delta_{\mathrm{s}}$, where $P_{\mathrm{s}}(X_{\mathrm{s}}, X_0, N, \pi)$ is the safety probability obtained for control-affine System* (3).

Remark 2. *There may exist many possible controllers that are solutions to Problem 2. Since System* (1) *represents a closed-loop system with a given controller, one may want to specifically synthesize a safety controller that is "minimally-invasive," i.e., it minimally intervenes to guarantee safety. In Section 5, we elaborate on this notion and introduce methods of generating such controllers.*

Overall Approach Our approaches to Problems 1 and 2 are based on (stochastic) barrier functions. Specifically, we aim to synthesize barrier functions that can serve as a safety certificate by providing a lower bound for P_{s}. To achieve this, we rely on recent results for local convex relaxation of an NN to find piece-wise linear over- and under-approximations of f^w. This allows us to formulate the problem of finding a stochastic barrier function for System (1) as an SOS optimization problem.

Our (minimally-invasive) control synthesis method (Problem 2) focuses on injecting controls to System (3) only in the regions of the state space that potentially contribute to unsafety of the system. To this end, we exploit two main properties of the barrier function: the required (super-) martingale property and the convexity property of SOS polynomials. The former allows us to design a method of identifying the "unsafe" regions, and the latter enables us to formulate a Linear Programming (LP) optimization problem to find controllers that robustly minimize the probability of unsafety.

3 Preliminaries

In this section, we provide a brief background on stochastic barrier functions and non-negative polynomials, which are central components of our framework.

3.1 Stochastic Barrier Functions

Consider stochastic process $\mathbf{x}_{k+1} = F(\mathbf{x}_k, \mathbf{v}_k)$, where $\mathbf{x}_k \in X \subseteq \mathbb{R}^n$, $\mathbf{v}_k \in V \subseteq \mathbb{R}^m$ is a well-defined stochastic process, and $F : X \times V \to X$ is a locally Lipschitz continuous function. Then, given initial set $X_0 \subset X$, safe set $X_s \subseteq X$, and unsafe set $X_u \subset X$, a twice differentiable function $B : X \to \mathbb{R}$ is called a *stochastic barrier function* or simply barrier function if the following conditions hold for $\alpha \geq 1$, $\beta, \eta \in [0, 1]$

$$B(x) \geq 0 \quad \forall x \in X \tag{4a}$$
$$B(x) \leq \eta \quad \forall x \in X_0 \tag{4b}$$
$$B(x) \geq 1 \quad \forall x \in X_u \tag{4c}$$
$$E[B(F(x,v)) \mid x] \leq B(x)/\alpha + \beta \quad \forall x \in X_s \tag{4d}$$

If such a B exists, then for any $N \in \mathbb{N}_{\geq 0}$, it follows, assuming $\alpha = 1$ without loss of generality, that

$$\Pr[\forall k \leq N, \ \mathbf{x}_k \in X_s \mid \mathbf{x}_0 \in X_0] \geq 1 - (\eta + \beta N). \tag{5}$$

Intuitively, Conditions (4a)-(4d) allow one to use non-negative super-martingale inequalities to bound the probability that the system reaches the unsafe region. This results in the probability bound in (5) (see [26, Chapter 3, Theorem 3]). A major benefit of using barrier functions for safety analysis is that while Conditions (4b)-(4d) are static properties, they allow one to make probabilistic statements on the time evolution of the system without the need to evolve the system over time.

3.2 Non-negative Polynomials

Synthesis of barrier functions can be formulated as a nonlinear optimization problem. To solve it efficiently via convex programming, we formulate the problem using SOS polynomials, which is a sufficient condition for polynomial non-negativity [27]. This formulation allows the use of efficient semidefinite programming solvers, which have polynomial worst-case complexity [28].

Definition 1 (SOS Polynomial). *A multivariate polynomial $\lambda(x)$ is a Sum-Of-Squares (SOS) for $x \in \mathbb{R}^n$ if there exists some polynomials λ_i, $i = 1, \ldots, r$, for some $r \in \mathbb{N}$, such that $\lambda(x) = \sum_{i=1}^{r} \lambda_i^2(x)$. If $\lambda(x)$ is an SOS, then $\lambda(x) \geq 0$ for all $x \in \mathbb{R}^n$. The set of all SOS polynomials is denoted by Λ.*

Now, consider a basic closed semi-algebraic set $X = \{x \in \mathbb{R}^n \mid h_i(x) \geq 0 \ \forall i \in \{1, \ldots, l\}\}$, which is defined as the intersection of l polynomial ($h_i(x)$) inequalities. Then, Proposition 1 allows one to enforce non-negativity on these semi-algebraic sets.

Proposition 1 (Putinar's Certificate, [29]). *Let $X \subset \mathbb{R}^n$ be a compact basic semi-algebraic set defined as $X = \{x \in \mathbb{R}^n \mid h_i(x) \geq 0 \ \forall i \in \{1, \ldots, l\}\}$, where $h_i(x)$ is a polynomial. Then, polynomial $\gamma(x)$ is non-negative on X if, for some $\lambda_0(x), \lambda_i(x) \in \Lambda$, $\gamma(x) = \lambda_0(x) + \sum_{i=1}^{l} \lambda_i(x)h_i(x)$.*

Corollary 1. *Let $\gamma(x)$ be a non-negative polynomial on $X = \{x \in \mathbb{R}^n \mid h_i(x) \geq 0 \ \forall i \in \{1, \ldots, l\}\}$, and $h(x)$ be an l-dimensional vector of polynomials with $h_i(x)$ as its i-th dimension for all $i \in \{1, \ldots, l\}$. Further, let $\mathcal{L}$ be an l-dimensional vector of SOS polynomials, i.e., the i-th element of $\mathcal{L}$ is $\lambda_i(x) \in \Lambda$. Then, it holds that $\gamma(x) - \mathcal{L}^T h(x) \in \Lambda$.*

4 Barrier Functions for Neural Network Dynamical Models

In this section, we introduce a formulation for synthesizing a barrier function that guarantees safety of System (1). Our approach relies on the local convex relaxations of f^w, which we use to formulate an SOS optimization problem that generates a valid barrier function for System (1).

4.1 Local Convex Relaxation of NNDMs

Our approach uses a convex relaxation of f^w. Specifically, we utilize the recent relaxation results for NNs [18] to compute piece-wise linear functions that under- and over-approximate f^w. The advantage of this method is that it is computationally efficient and produces fairly tight bounds for small neighborhoods. To this end, we first partition X_s to a finite set of regions $Q = \{q_1, ..., q_{|Q|}\}$, where $q_i \subseteq X_s$ for all $i \in \{1, \ldots, |Q|\}$ is a basic (closed) semi-algebraic set. Such a decomposition can be obtained using, e.g., a grid or half-spaces. Then, for every $q \in Q$, we compute linear bounds

$$\underline{f}_q(x) = \underline{A}_q x + \underline{b}_q, \quad \overline{f}_q(x) = \overline{A}_q x + \overline{b}_q \quad \text{such that} \quad \forall x \in q, \quad \underline{f}_q(x) \leq f^w(x) \leq \overline{f}_q(x). \quad (6)$$

The above linear functions can be produced in several ways. One method is to forward propagate the bounds in the input domain of $f^w(x)$ using Interval Bound Propagation (IBP) based approaches [30, 31]. These approaches are efficient but generally lead to loose approximations. Another method is Linear Relaxation based Perturbation Analysis (LiRPA) algorithms, which perform a symbolic back propagation followed by a forward evaluation of the bounds [32, 33, 18]. Such LiRPA-based algorithms produce tighter bounds but can be slow. In our implementation, we use the off-the-shelf toolbox [33] to compute $\overline{f}_q(x)$ and $\underline{f}_q(x)$ in (6) for its trade-off between speed and accuracy.

4.2 Barrier Function (Certificate) Synthesis

Recall that a valid barrier function $B(x)$ must satisfy Conditions (4b)-(4d). With the above local linear bounds for f^w, Condition (4d) can be redefined over the partitioned regions in Q. That is, Condition (4d) can be replaced with the following constraints: for each $q \in Q$,

$$E[B(f_q^w(x) + \mathbf{v}) \mid x] \leq B(x)/\alpha + \beta \quad \text{and} \quad \underline{f}_q(x) \leq f_q^w(x) \leq \overline{f}_q(x) \quad \forall x \in q. \quad (7)$$

In this formulation, $f_q^w(x)$ is a (free) variable that is bounded by two linear functions, avoiding the non-convexity of $f_q^w(x)$. Consequently, an efficient evaluation of the constraint becomes feasible.

Remark 3. *Note that there are in fact three constraints in (7) (one super-martingale condition on B and two bounds on $f_q^w(x)$). If they are treated separately, it leads to a conservative barrier function because each constraint needs to hold for all $x \in q$ separately. A less conservative approach is to pose all constraints for the same x at the same time. In Theorem 1, we show how this can be achieved.*

As commonly practiced in the literature, e.g., [14, 8, 34], we restrict our search for a $B(x)$ to the class of polynomial functions. Then, the degree of the polynomial becomes a design parameter. In the following lemma, we show that the degree of the polynomial must be at least 2 for the polynomial to be a valid candidate barrier function for System (1). That is, there does not exist a linear function satisfying Conditions (4b)-(4d) for $\eta < 1$, which implies that the lower bound in (5) is trivially 0.

Lemma 1. *For $k \in \mathbb{R}^n$ and $c \in \mathbb{R}$, consider linear function $B(x) = k^T x + c$. Then, $B(x)$ satisfies Conditions (4b)-(4d) iff $k = 0$ and $c = 1$.*

The proof of Lemma 1 is in Appendix A.1. It relies on the fact that a linear function is either a constant or always eventually violates Condition (4a). Having established that the polynomial degree must be at least 2, we further restrict $B(x)$ to SOS polynomials. Theorem 1 shows a formulation of the barrier function synthesis problem as an SOS program, which can be solved in polynomial time. In particular, the theorem shows how to formulate Conditions (4b)-(4c) and (7) as SOS constraints.

Theorem 1 (NNDM Barrier Certificate)**.** *Consider SOS polynomial function $B(x)$, and safe set $X_s = \{x \in \mathbb{R}^n : h_s(x) \geq 0\}$, initial set $X_0 = \{x \in \mathbb{R}^n : h_0(x) \geq 0\}$, unsafe set $X_u = \mathbb{R}^n \setminus X_s = \{x \in \mathbb{R}^n : h_u(x) \geq 0\}$, and partition region $q = \{x \in \mathbb{R}^n : h_q(x) \geq 0\}$ for all $q \in Q$. Let $\mathcal{L}_s(x)$, $\mathcal{L}_0(x)$ and $\mathcal{L}_u(x)$ be vectors of SOS polynomials with the same dimensions as h_s, h_0, and h_u, respectively. Likewise, let $\mathcal{L}_{q,x}(x)$ and $\mathcal{L}_{q,y}(x)$ be vectors of SOS polynomials with the same dimension as h_q. Then, a stochastic barrier certificate $B(x)$ for System (1) with time horizon*

$N \in \mathbb{N}_{\geq 0}$ *can be obtained by solving the following SOS optimization problem for* $\eta, \beta \in [0, 1]$:

$$\min_{\beta, \eta} \quad \eta + \beta N \qquad \textit{subject to:}$$

$$B(x) \in \Lambda, \tag{8a}$$

$$-B(x) - \mathcal{L}_0^T(x)h_0(x) + \eta \in \Lambda, \tag{8b}$$

$$B(x) - \mathcal{L}_{\mathrm{u}}^T(x)h_{\mathrm{u}}(x) - 1 \in \Lambda, \tag{8c}$$

$$E[B(y + v) \mid x] \leq B(x)/\alpha + \beta - \mathcal{L}_{q,x}^T(x)h_q(x) -$$
$$\mathcal{L}_{q,y}^T(x)(\overline{f}_q(x) - y)(y - \underline{f}_q(x)) \in \Lambda \qquad \forall q \in Q \tag{8d}$$

which guarantees safety probability $P_{\mathrm{s}}(X_{\mathrm{s}}, X_0, N) \geq 1 - (\eta + \beta N)$.

The proof of Theorem 1 is in Appendix A.2. The key idea is that an SOS polynomial function is a linear combination of monomials for a given degree $m \geq 2$. Hence, Condition (4a) can be written as Condition (8a). Using Corollary 1, we can enforce non-negativity of a polynomial function on semi-algebraic sets using Lagrangian multipliers $\mathcal{L}$ in terms of vectors of SOS polynomials and the corresponding polynomials in $h(x)$. Hence, we can write Conditions (4b)-(4c) for basic closed semi-algebraic sets as Conditions (8b)-(8c), which are guaranteed to only hold over their corresponding basic closed semi-algebraic sets. Finally, Condition (8d) is derived from Condition (4d) for System (1) by using a similar technique to restrict the super-martingale inequality to the corresponding local linear bounds of $f^w(x)$ in (7) in each region $q \in Q$. Therefore, if a feasible solution exists, the SOS optimization problem in Theorem 1 yields a safety certificate for System (1).

5 Minimally-invasive Barrier Controller Synthesis

Here, we focus on Problem 2. The assumption is that a barrier function $B(x)$ is synthesized for System (1) via Theorem 1, which has a safety probability that is lower than threshold δ_{s}. Hence, our goal is to design a feedback controller for System (3) that guarantees safety. Ideally, this controller is minimal in the interruptions it causes to the system evolution. Below, we first provide a definition for a minimally-invasive controller, and then show how it can be synthesized.

Recall that a feedback controller $\pi : \mathbb{R}^n \to U$ is a function that assigns a control value to each state. Let $X_\pi \subseteq X_{\mathrm{s}}$ be the set of states, to which π assigns non-zero control values, i.e., $\pi(x) \neq 0$ for all $x \in X_\pi$. Further, we denote by $V(X)$ the volume of set $X \subset \mathbb{R}^n$. Then, we say a feedback controller π is *minimally-invasive* if it minimizes $V(X_\pi)$. Moreover, we say π is *ϵ-minimally-invasive* if $V(X_\pi) - \min_{\pi'} V(X_{\pi'}) \leq \epsilon$. In this work, we seek ϵ-*minimally-invasive* controller with an arbitrary small ϵ. To this end, we make the observation that if the discretization of X_{s} discussed in Section 4.1 is uniform, $V(X_\pi)$ for a given π is directly proportional to the number of discrete regions in Q which π applies a non-zero control. We denote by Q_π the set of such regions. Then, a feedback controller π that minimizes $|Q_\pi|$ is ϵ-*minimally-invasive*, and as the discretization becomes finer, ϵ monotonically approaches zero. Therefore, given barrier function $B(x)$, our first goal is to identify the minimal number of regions in Q that require a non-zero controller to guarantee safety. Recall that, given $B(x)$, we require the safety probability to be $1 - (\eta + \beta N) \geq \delta_{\mathrm{s}}$. While η is related to the value of $B(x)$ in the initial set, β is a compensation needed to turn a sub-martingale inequality to a super-martingale inequality (Condition (8d) or (4d)). We can compute this compensation for each region $q \in Q$ by evaluating Condition (8d) or (4d). Let β_q denote the computed compensation term for region q. Then, region q requires a non-zero controller if $\beta_q > (1 - \delta_{\mathrm{s}} - \eta)/N$.

We now turn our attention to designing a controller that reduces β_q. In our approach, we tap directly into the convex nature of the SOS polynomial $B(x)$. Lemma 2 shows that a controller that drives System (3) to a point closest to the minimum point of $B(x)$, also robustly minimizes β_q.

Lemma 2. *Given a convex barrier function* $B(x)$, *let* $x^* = \arg\min_x B(x)$, $x' = f^w(x) + gu + v$ *as in System (3), and* $\beta_q \geq E[B(x') \mid x] - B(x)$ *for all* x *in region* $q \subseteq X_{\mathrm{s}}$. *Then, the controller that minimizes* $\|\tilde{x}' - x^*\|_1$, *where* $\tilde{x}' = f^w(x) + gu$, *also robustly minimizes* β_q, *i.e., minimizes an upper bound of* β_q.

Proof. Given that B is a polynomial and x' is affine in both u and v, we can write $E[B(x') \mid x] = B(E[x' \mid x]) + \xi$, where ξ, called Jensen's gap [35], is a non-negative polynomial of x, u, and

moments $E[v^d]$, where $2 \le d \le m$ is an even integer. This polynomial can always be upper bounded by a constant $\overline{\xi} \ge 0$ for bounded x and u. Then, with $\tilde{x}' = E[x']$, the min β_q is upper bounded by

$$\min_u \beta_q \le \min_u \max_{x \in q} \left(B(\tilde{x}' \mid x) - B(x) + \overline{\xi} \right) \le \min_u \max_{x \in q} B(\tilde{x}' \mid x) - \min_{x \in q} B(x) + \overline{\xi}.$$

The second inequality holds because B is a non-negative convex function. Since $B(\tilde{x}')$ monotonically decreases as $\tilde{x}' \to x^*$, $\max_{x \in q} B(\tilde{x}' \mid x)$ is minimized by $\arg \min_u \|f^w(x) + gu - x^*\|_1 \ \forall x \in q$. $\quad\square$

We note that, in the proof of Lemma 2, the moments $E[v^d]$ appear in every monomial of polynomial ξ (Jensen's gap). Hence, for v with variance less than 1, Jensen's gap is very small and vanishes as the variance becomes smaller, making the controller in the lemma optimal. Furthermore, even though Lemma 2 uses L_1 norm for $\|\tilde{x}' - x^*\|_1$, the statement also holds for L_2 and L_∞ norms. We specifically use L_1 to be able to compute the controller via an LP as stated in Theorem 2.

Theorem 2 (Controller Synthesis). *Consider convex barrier function $B(x)$ and its minimum argument x^* as specified in Lemma 2, and linear bounds $\underline{f}_q(x)$ and $\overline{f}_q(x)$ for $f^w(x)$ in region $q \subseteq \mathbb{R}^n$ as in (6). A controller that robustly minimizes β_q for System (3) is the solution to the following LP optimization problem, where $\theta = (\theta_1, \theta_2, \dots, \theta_n) \in \mathbb{R}_{\ge 0}^n$, and the vector inequality relations are element-wise*

$$\min_{\theta_i, z', z'', u} \sum_{i=1}^n \theta_i \quad \text{subject to:}$$

$$
\begin{aligned}
y - x^* &\le \theta & \text{(10a)} & \qquad y \ge \underline{f}_q(z) + gu & \text{(10d)}\\
x^* - y &\le \theta & \text{(10b)} & \qquad y \le \overline{f}_q(z) + gu & \text{(10e)}\\
z &= z' - z'' & \text{(10c)} & \qquad z', z'' \ge 0, \ z \in q, \ u \in U & \text{(10f)}
\end{aligned}
$$

Proof. By Lemma 2 it holds that to robustly minimizes β_q, it suffices to minimize $\|\tilde{x}' - x^*\|_1$. This norm is the sum of n absolute values, each of which can be set to θ_i with linear Constraints (10a) and (10b). Hence, the minimization of $\|\tilde{x}' - x^*\|_1$ is equivalent to minimization of $\sum_{i=1}^n \theta_i$ subject to (10a) and (10b). The bounds on the dynamics of System (3) can be expressed as linear Conditions (10d) and (10e), each of which has to hold for all $z \in q$, i.e., z is a free variable. To turn the optimization into an LP, z is expressed as the difference of two non-negative decision variables z' and z'' in Constraint (10c). $\quad\square$

Remark 4. *Note that the controller obtained via the LP is a vector of scalars for each region q. This leads to a feedback controller π over the regions $q \in Q$. It may be possible to express u as SOS polynomials, which would be a feedback controller over x. That could potentially provide a tighter bound for β_q at the cost of extra computation since SOS optimization is more expensive than LP.*

5.1 Control Synthesis Algorithm

Our framework for computing a barrier certificate as well as minimally-invasive controllers for NNDMs is outlined in Algorithm 1. First, the safe set X_s is uniformly partitioned into a finite set of regions $|Q|$, and bounds of f^w are computed as piece-wise linear functions per Section 4.1. Then, for a given polynomial degree $m \ge 2$, a barrier certificate along with its corresponding safety probability bound are computed in accordance with Theorem 1. If this probability is greater than or equal to δ_s, the NNDM is certified, and the algorithm terminates. Otherwise, it iteratively searches for ϵ-minially-invasive controllers. In each iteration, upper bound $\overline{\eta}$ is set on η (initially with its largest possible value $1 - \delta_s$ and then reduced by $\Delta\eta > 0$), and a new barrier function is computed with this bound, where the objective function is minimization of β to ensure ϵ-minimally-invasive controllers. Then, for each $q \in Q$, β_q is checked against the threshold. If exceeding, a controller is computed for q in accordance Theorem 2. This process repeats until a satisfactory P_s is obtained or $\overline{\eta} < 0$.

We highlight two important properties of this algorithm: it clearly terminates in finite time, and if $P_s \ge \delta_s$, the synthesized controller is correct-by-construction.

Algorithm 1: NNDM Controller Synthesis

Input : NN f^w, initial set X_0, safe set X_s, noise covariance R, polynomial degree m, threshold δ_s, and step size $\Delta\eta$

Output : feedback controller π, and safety probability bound P_s

1   $Q, \underline{f}_q, \overline{f}_q \leftarrow$ PARTITIONANDCOMPUTEBOUNDS(X_s, f^w), $k \leftarrow 0$

2   $\eta, \beta, B(x) \leftarrow$ COMPUTEBARRIERCERTIFICATE$(Q, \underline{f}_q, \overline{f}_q, R, m)$ // Theorem 1

3   **while** $P_s < \delta_s$ *and* $\overline{\eta} > 0$ **do**

4      $\overline{\eta} \leftarrow 1 - (\delta_s - k\Delta\eta)$, $k \leftarrow k+1$ // set upper bound on η

5      $\eta, \beta, B(x) \leftarrow$ COMPUTEBARRIERCERTIFICATE$(Q, \underline{f}_q, \overline{f}_q, R, m, \overline{\eta})$ // Theorem 1

6   **for** $q \in Q$ **do**

7         $\pi \leftarrow (q, 0)$

8   **if** $\beta_q > (1 - \delta_s - \eta)/N$ **then**

9            $\pi \leftarrow (q, u_q) \leftarrow$ COMPUTECONTROL$(\arg\min B(x), q, \underline{f}_q, \overline{f}_q, U)$ // Theorem 2

10            $\beta_q \leftarrow$ UPDATEBETA$(B(x), u_q)$ // evaluate Condition (8d)

11   $P_s \leftarrow 1 - (\eta + N \max_{q \in Q} \beta_q)$

12   **return** π, P_s

6   Case Studies

We demonstrate the efficacy of our framework on several case studies. We first show that our verification method is able to produce non-trivial safety guarantees for NNDMs trained using imitation learning techniques [36] from data gathered by rolling out an RL agent trained with state-of-the-art RL techniques. We then show that our control approach significantly increases the safety probability.

Models We trained the Pendulum and Cartpole agents from the OpenAI gym environment. For training, we collected input-output states of the system under an expert controller, either a look-up table or an NN controller. We picked the best controller available for the agents from the OpenAI Leaderboard [37]. We also trained four-dimensional husky models. We initially trained a NNDM along with NN controller on the trajectory data generated from a physics simulator as outlined in [9]. We then used this NNDM to generate input-output data and trained closed-loop NNDMs as specified above. In total, we trained 8 fully-connected multi-layer perceptron NNDM architectures with upto 5 hidden layers and 256 neurons, each with ReLU activation function. The dimensionality and architecture information for each model is shown in Table 1.

Table 1: Benchmark results for barrier certificate and control synthesis on various NNDMs. Threshold $\delta_s = 0.95$ was used for the minimally-invasive control synthesis. Architecture $h \times [r]$ denotes a NNDM with h hidden layers, each with r neurons. The dimensionality of each system is n.

Model	n	$h \times [r]$	$\|Q\|$	Verification						Control					
				Interval Bounds			Linear Bounds			Interval Bounds			Linear Bounds		
				β	P_s	Time (min)	β	P_s	Time (min)	β	P_s	Time (min)	β	P_s	Time (min)
Pendulum	2	1 x [64]	120	0.54	0.459	0.09	0.002	0.997	0.09	10^{-6}	0.999	0.03	-	-	-
			240	0.444	0.555	0.12	0.002	0.997	0.14	10^{-6}	0.999	0.11	-	-	-
			480	0.155	0.844	0.37	0.003	0.997	0.42	10^{-6}	0.999	0.35	-	-	-
		2 x [64]	120	0.575	0.424	0.05	0.191	0.785	0.06	10^{-6}	0.999	0.04	10^{-6}	0.999	0.05
			240	0.471	0.526	0.13	0.153	0.844	0.14	10^{-6}	0.998	0.1	10^{-6}	0.998	0.12
			480	0.196	0.801	0.39	0.075	0.922	0.46	10^{-6}	0.997	0.32	10^{-6}	0.997	0.34
		3 x [64]	120	0.689	0.304	0.05	0.384	0.603	0.06	10^{-6}	0.994	0.03	10^{-6}	0.986	0.04
			240	0.555	0.439	0.13	0.288	0.707	0.14	10^{-6}	0.995	0.11	10^{-6}	0.995	0.13
			480	0.357	0.636	0.39	0.208	0.791	0.44	10^{-6}	0.994	0.35	10^{-6}	0.998	0.46
		5 x [64]	480	0.878	0.027	0.41	0.735	0.172	0.45	-	-	-	10^{-6}	0.951	0.36
			960	0.635	0.27	1.44	0.603	0.302	1.56	10^{-6}	0.905	0.36	10^{-6}	0.951	1.59
			1920	0.449	0.447	5.95	0.421	0.504	6.61	10^{-6}	0.905	1.35	10^{-6}	0.951	5.33
Cartpole	4	1 x [128]	960	-	-	-	0.623	0.375	17.99	-	-	-	10^{-6}	0.998	13.35
			1920	-	-	-	0.414	0.577	65.14	-	-	-	10^{-6}	0.992	35.89
			3840	0.794	0.195	215.22	0.382	0.602	222.84	10^{-6}	0.98	156.86	10^{-6}	0.984	160.99
		2 x [128]	960	-	-	-	0.811	0.181	19.84	-	-	-	10^{-6}	0.993	13.3
			1920	-	-	-	0.779	0.216	49.79	-	-	-	10^{-6}	0.997	36.89
			3840	0.877	0.102	230.49	0.664	0.333	239.21	10^{-6}	0.979	150.31	10^{-6}	0.997	162.57
Husky	4	1 x [256]	900	0.672	0.199	15.56	0.431	0.366	16.65	10^{-6}	0.872	11.45	10^{-6}	0.951	11.81
			1800	0.667	0.211	53.33	0.425	0.382	53.55	10^{-6}	0.878	33.15	10^{-6}	0.951	33.73
			2250	0.645	0.288	68.4	0.367	0.494	81.32	10^{-6}	0.933	42.73	10^{-6}	0.951	46.49
			4800	0.622	0.331	307.61	0.347	0.538	345.28	10^{-6}	0.953	299.97	10^{-6}	0.951	304.67
		2 x [256]	1800	-	-	-	0.544	0.011	74.69	-	-	-	10^{-6}	0.8	35.22
			2250	-	-	-	0.502	0.222	74.83	-	-	-	10^{-6}	0.951	51.76
			4800	0.845	0.062	369.88	0.415	0.384	406.69	10^{-6}	0.907	297.57	10^{-6}	0.951	306.41

The state space of the Pendulum is pole angle θ and angular velocity $\dot{\theta}$ with safe set defined as $[-\pi/15, \pi/15] \times [-1, 1]$ and initial set with $\theta_0 \in [-\pi/36, \pi/36]$. For Cartpole, the state space is cart position x and velocity $\dot{x}$, and pole angle θ and angular velocity $\dot{\theta}$. The safe sets are defined for cart position as $x \in [-1, 1]$ and $\theta \in [-\pi/15, \pi/15]$, with the initial set $\theta_0 \in [-\pi/36, \pi/36]$. For Husky, the states we observe are position x and y, orientation θ and linear velocity v. The task is to stay within a lane and hence y should be between $[-1, 1]$, and the initial set is defined as any position $[x, y]$ that is within a circle of radius 0.1 centered around the origin.

Implementation and Experimental Setup: We implemented our algorithms in Python and Julia. For collecting data and training, we used the TensorFlow framework [38]. To compute the linear approximations of the NNDMs, we utilized α-CROWN [33]. For the optimization, we used Julia's *SumOfSquares.jl* package [39, 40]. The optimizations are all performed single-threaded on a computer with 3.9 GHz 8-Core CPU and 64 GB of memory. The code is available in the supplementary material.

Benchmarks: To the best of our knowledge, there is no other work on safety verification of NNDMs, which makes it difficult to compare our framework with other approaches. As a baseline, however, we compare our approach based on linear under and over approximation of the NNDM against interval (global) bounding approaches cf. Remark 5 in Appendix A.2. This results in a relaxed optimization problem, but at the cost of more conservative results. The results are shown in Table 1. The computation times indicate the time to synthesize the barrier and the controller using Theorems 1 and 2, respectively. The barrier degree is 4 for all the case studies. Below we provide a brief discussion on the results. For more in-depth discussions and details on the discretization and hyperparameters, see Appendix B in the supplementary material.

Verification: As expected, the safety probability always increases with $|Q|$ (finer partitions), since the linear and interval bounds for f^w become more accurate. This behavior is encountered for all the systems and architectures, i.e., regardless of the number of hidden layers, or neurons per layer. However, it is also observed that this trend of finer discretizations and increased safety probability consistently comes at the cost of an increased computation time. This is due to the fact that the number of partitions directly dictate the number of constraints in the optimization problems. Note that the probability of safety using interval bounds is strictly worse than linear approximations for all the experiments. For some instances we could not certify the system using interval bounds. Using linear bounds, we are able to compute non-trivial probability of safety for all the case-studies and observe the discretization vs accuracy trend. Take for instance the 1-layer Cartpole model, where an increase in $|Q|$ from 960 to 3840 causes an increase in P_s from 0.375 to 0.602, at the expense of an increased computation time of more than 12 times. This behavior is also true for the Husky model, where increasing $|Q|$ from 1800 to 4800 for the 2-layer model, increases P_s from 0.011 to 0.384 at the price of a more than 5-fold increase in the computation time.

Control Synthesis: When the certification probability is below $\delta_s = 0.95$, the control strategy is utilized to increase the probability of safety. Using interval bounds, controllers had to be synthesized for all case-studies, which illustrates the looseness of the bounds. In fact, in several case studies, the controller was not able to meet the desired threshold. For linear bounds, we observe that for all NNDMs, whose $P_s < \delta_s$, the controller synthesis algorithm is able to produce $P_s \geq \delta_s$, showing the effectiveness of our control method. The exception is for the Husky model with 2-hidden layers and $|Q| = 1800$, where the controller increases P_s from 0.011 to 0.800 (73-fold increase) yet cannot meet the 0.95 threshold. Nonetheless, for all the linear bound cases, the controller reduces β to 10^{-6}.

7 Conclusion

In this work, we introduced a methodology for providing safety guarantees for NNDMs using stochastic barrier functions. We showed that the problem of finding a stochastic barrier function for a NNDM can be relaxed to an SOS optimization problem. Furthermore, we derived a novel framework for synthesis of an affine controller with the goal of maximizing the safety probability in a minimally-invasive manner. Experiments showcase that, along with a minimally-invasive controller, we can guarantee safety for various standard reinforcement learning problems.

A bottleneck to our approach is the discretization step, which is required to obtain piece-wise linear under- and over-approximations of the neural network, and can be expensive for higher dimensional spaces. A future direction is to relax the reliance on this step. Another future direction is to consider a more expressive class of barrier functions than SOS polynomials.

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