Faster Stochastic Algorithms for Minimax Optimization under Polyak-Łojasiewicz Conditions

Abstract

This paper considers stochastic first-order algorithms for minimax optimization under Polyak-Łojasiewicz (PL) conditions. We propose SPIDER-GDA for solving the finite-sum problem of the form $\min_x \max_y f(x,y) \triangleq \sum_{i=1}^{n} f_i(x,y)$, where the objective function $f(x,y)$ is μ_x-PL in x and μ_y-PL in y; and each $f_i(x,y)$ is L-smooth. We prove SPIDER-GDA could find an ϵ-approximate solution within $\mathcal{O}\left((n + \sqrt{n}\,\kappa_x\kappa_y^2)\log(1/\epsilon)\right)$ stochastic first-order oracle (SFO) complexity, which is better than the state-of-the-art method whose SFO upper bound is $\mathcal{O}\left((n + n^{2/3}\kappa_x\kappa_y^2)\log(1/\epsilon)\right)$, where $\kappa_x \triangleq L/\mu_x$ and $\kappa_y \triangleq L/\mu_y$. For the ill-conditioned case, we provide an accelerated algorithm to reduce the computational cost further. It achieves $\tilde{\mathcal{O}}\left((n + \sqrt{n}\,\kappa_x\kappa_y)\log(1/\epsilon)\right)$ SFO upper bound when $\kappa_x \geq \sqrt{n}$. Our ideas also can be applied to the more general setting that the objective function only satisfies PL condition for one variable. Numerical experiments validate the superiority of proposed methods.

1 Introduction

This paper focuses on smooth minimax optimization problem of the form

$$\min_{x \in \mathbb{R}^{d_x}} \max_{y \in \mathbb{R}^{d_y}} f(x,y) \triangleq \frac{1}{n}\sum_{i=1}^{n} f_i(x,y), \tag{1}$$

which covers a lot of important applications in machine learning such as reinforcement learning [10, 39], AUC maximization [13, 22, 45], imitation learning [5, 29], robust optimization [11], causal inference [26], game theory [6, 27] and so on.

We are interested in the minimax problems under PL conditions [9, 29, 42]. The PL condition [32] was originally proposed to relax the strong convexity in minimization problem that is sufficient for achieving the global linear convergence rate for first-order methods. In machine learning community, it has been successfully used to analyze the convergence behavior for overparameterized neural networks [21], robust phase retrieval [37] and a plenty of fundamental models [17]. There are many popular minimax formulations only satisfy PL condition, but lack strong convexity (or strong concavity). The examples include PL-game [29], robust least square [42], deep AUC maximization [22] and generative adversarial imitation learning of LQR [5, 29].

Yang et al. [42] showed that the alternating gradient descent ascent (AGDA) algorithm linearly converges to the saddle point when the objective function satisfies two-sided PL condition. They also

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Table 1: We present the comparison of SFO complexities under two-sided PL condition.

Algorithm	Complexity	Reference
GDA/AGDA	$\mathcal{O}\left(n\kappa_x\kappa_y^2 \log\left(1/\epsilon\right)\right)$	Theorem B.1, [42]
SVRG-AGDA	$\mathcal{O}\left((n + n^{2/3}\kappa_x\kappa_y^2)\log\left(1/\epsilon\right)\right)$	[42]
SVRG-GDA	$\mathcal{O}\left((n + n^{2/3}\kappa_x\kappa_y^2)\log\left(1/\epsilon\right)\right)$	Theorem C.1
SPIDER-GDA	$\mathcal{O}\left((n + \sqrt{n}\kappa_x\kappa_y^2)\log\left(1/\epsilon\right)\right)$	Theorem 4.1
AccSPIDER-GDA	$\begin{cases} \tilde{\mathcal{O}}\left(\sqrt{n}\kappa_x\kappa_y \log^2\left(1/\epsilon\right)\right), & 0 < \sqrt{n} \leq \kappa_x; \\ \tilde{\mathcal{O}}\left(n\kappa_y \log^2\left(1/\epsilon\right)\right), & \kappa_x < \sqrt{n} \leq \kappa_x\kappa_y; \\ \mathcal{O}\left((n + \sqrt{n}\kappa_x\kappa_y^2)\log\left(1/\epsilon\right)\right), & \kappa_x\kappa_y < \sqrt{n}. \end{cases}$	Corollary 5.1

proposed the SVRG-AGDA method[1] for the finite-sum problem (1), which could find ϵ-approximate solution within $\mathcal{O}\left((n + n^{2/3}\kappa_x\kappa_y^2)\log(1/\epsilon)\right)$ stochastic first-order oracle (SFO) calls,[2] where κ_x and κ_y are the condition numbers with respect to PL condition for x and y respectively. The variance reduced technique in the SVRG-AGDA leads to better a convergence rate than full batch AGDA whose SFO complexity is $\mathcal{O}\left(n\kappa_x\kappa_y^2 \log(1/\epsilon)\right)$. However, there are still some open questions left. Firstly, Yang et al. [42]'s theoretical analysis heavily relies on the alternating update rules. It remains interesting whether a simultaneous version of GDA (or its stochastic variants) also has similar convergence results. Secondly, it is unclear whether the SFO upper bound obtain by SVRG-AGDA can be improved by designing more efficient algorithms.

For one-sided PL condition, we desire to find the stationary point of $g(x) = \max_{y \in \mathbb{R}^{d_y}} f(x, y)$, since the saddle point may not exist. Nouiehed et al. [29] proposed the multi-step GDA method that achieves the ϵ-stationary point within $\tilde{\mathcal{O}}(\kappa_y^2 L\epsilon^{-2})$ numbers of full gradient iterations. The similar complexity also can be obtained by AGDA [42]. Recently, Yang et al. [44] proposed the smoothed-AGDA that improves the upper bound into $\mathcal{O}(\kappa_y L\epsilon^{-2})$. Both multi-step GDA Nouiehed et al. [29] and smoothed-AGDA [43] can be extended to online setting, but the formulation (1) with finite-sum structure has not been explored.

In this paper, we introduce a variance reduced first-order method, called SPIDER-GDA, which constructs the gradient estimator by stochastic recursive gradient and the iterations are based on simultaneous gradient descent ascent. We prove that SPIDER-GDA could achieve ϵ-approximate solution of the two-sided PL problem of the form (1) within $\mathcal{O}\left((n + \sqrt{n}\,\kappa_x\kappa_y^2)\log(1/\epsilon)\right)$ SFO calls, which has better dependency on n than SVRG-AGDA [42]. We also provide an acceleration framework to improve first-order methods for solving ill-conditioned minimax problems under PL conditions. The accelerated SPIDER-GDA (AccSPIDER-GDA) could achieve ϵ-approximate solution within $\tilde{\mathcal{O}}\left((n + \sqrt{n}\,\kappa_x\kappa_y)\log(1/\epsilon)\right)$ SFO calls when $\kappa_x \geq \sqrt{n}$, which is the best known SFO upper bound for this problem. We summarize our main results and compare them with related work in Table 1. Without loss of generality, we always suppose $\kappa_x \geq \kappa_y$. Furthermore, the proposed algorithms also work for minimax problem with one-sided PL condition. We present the results for this case in Table 2.

2 Related Work

The minimax optimization problem (1) can be viewed as the following minimization problem

$$\min_{x \in \mathbb{R}^{d_x}} \left\{ g(x) \triangleq \max_{y \in \mathbb{R}^{d_y}} f(x, y) \right\}.$$

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[1] Yang et al. [42] name this method as variance-reduced-AGDA (VR-AGDA). Here we call it SVRG-AGDA to distinguish with other variance reduced algorithms.

[2] The original analysis[42] provided an SFO upper bound $\mathcal{O}\left((n + n^{2/3}\max\{\kappa_x^3, \kappa_y^3\})\log(1/\epsilon)\right)$, which can be refined to $\mathcal{O}\left((n + n^{2/3}\kappa_x\kappa_y^2)\log(1/\epsilon)\right)$ by some little modification in the proof.

Table 2: We present the comparison of SFO complexities under one-sided PL condition.

Algorithm	Complexity	Reference
Multi-Step GDA	$\tilde{\mathcal{O}}(n\kappa_y^2 L\epsilon^{-2})$	[29]
GDA/AGDA	$\mathcal{O}\left(n\kappa_y^2 L\epsilon^{-2}\right)$	Theorem B.2, [42]
Smooothed-AGDA	$\mathcal{O}\left(n\kappa_y L\epsilon^{-2}\right)$	[44]
SVRG-GDA	$\mathcal{O}\left(n + n^{2/3}\kappa_y^2 L\epsilon^{-2}\right)$	Theorem F.1
SPIDER-GDA	$\mathcal{O}\left(n + \sqrt{n}\kappa_y^2 L\epsilon^{-2}\right)$	Theorem 6.1
AccSPIDER-GDA	$\begin{cases} \tilde{\mathcal{O}}\left(\sqrt{n}\kappa_y L\epsilon^{-2}\right), & 0 < \sqrt{n} \leq \kappa_y; \\ \tilde{\mathcal{O}}\left(nL\epsilon^{-2}\right), & \kappa_y < \sqrt{n} \leq \kappa_y^2; \\ \mathcal{O}\left(n + \sqrt{n}\kappa_y^2 L\epsilon^{-2}\right), & \kappa_y^2 < \sqrt{n}. \end{cases}$	Theorem 6.2

A natural way to solve such problem is the multi-step GDA algorithm [19, 23, 29, 33] that contains double-loop iterations in which the outer loop can be regarded as running inexact gradient descent on $g(x)$ and the inner loop finds the approximate solution to $\max_{y\in\mathbb{R}^{d_y}} f(x, y)$ for a given x. Another class of methods is the two-timescale (alternating) GDA algorithm [9, 19, 41, 42] that only has single-loop iterations which update two variables with different stepsizes. The two-timescale GDA method can be implemented more easily and typically performs better than multi-step GDA empirically. Its convergence rate also can be established by analyzing function $g(x)$ but the analysis is more challenging than the multi-step GDA.

The variance reduction is a popular technique to improve the efficiency of stochastic optimization algorithms [2–4, 7, 8, 12, 15, 16, 25, 30, 31, 34–36, 40, 46, 47]. It is shown that solving nonconvex minimization problems with stochastic recursive gradient estimator [12, 15, 31, 40, 48] has the optimal SFO complexity. In the context of minimax optimization, the variance reduced algorithms also obtain the best-known SFO complexities in several settings [1, 14, 23, 24, 38, 42]. Specifically, the (near) optimal SFO algorithm for several convex-concave minimax problem has been proposed [14, 24], but the optimality for the more general case is still unclear [23, 42].

The Catalyst acceleration [18] is a useful approach to reduce the computational cost of ill-conditioned optimization problems, which is based on a sequence of inexact proximal point iterations. Lin et al. [20] first introduced Catalyst into minimax optimization. Later, Luo et al. [24], Tominin et al. [38], Yang et al. [43] designed the accelerated stochastic algorithms for convex-concave and nonconvex-concave problem. Recently, Yang et al. [44] also applied this technique to one-sided PL setting.

3   Notation and Preliminaries

First of all, we present the definition of saddle point.

Definition 3.1. *We say* $(x^*, y^*) \in \mathbb{R}^{d_x} \times \mathbb{R}^{d_y}$ *is a saddle point of function* $f : \mathbb{R}^{d_x} \times \mathbb{R}^{d_y} \to \mathbb{R}$ *if it holds that* $f(x^*, y) \leq f(x^*, y^*) \leq f(x, y^*)$ *for any* $x \in \mathbb{R}^{d_x}$ *and* $x \in \mathbb{R}^{d_y}$.

Then we formally define the Polyak-Łojasiewicz (PL) condition [32] as follows.

Definition 3.2. *We say a differentiable function* $h : \mathbb{R}^d \to \mathbb{R}$ *satisfies* μ*-PL for some* $\mu > 0$ *if* $\|\nabla h(z)\|^2 \geq 2\mu\big(h(z) - \min_{z'\in\mathbb{R}^d} h(z')\big)$ *holds for any* $z \in \mathbb{R}^d$.

Note that the PL condition does not require the strongly convexity and it can be satisfied even if the function is nonconvex [17].

We are interested in the finite-sum minimax optimization problem (1) under following assumptions.

Assumption 3.1. *We suppose each component $f_i : \mathbb{R}^{d_x} \times \mathbb{R}^{d_y} \to \mathbb{R}$ is L-smooth, i.e., there exists a constant $L > 0$ such that $\|\nabla f_i(x, y) - \nabla f_i(x', y')\|^2 \le L^2 \big(\|x - x'\|^2 + L^2\|y - y'\|^2\big)$ holds for any $x, x' \in \mathbb{R}^{d_x}$ and $y, y' \in \mathbb{R}^{d_y}$.*

Assumption 3.2. *We suppose the differentiable function $f : \mathbb{R}^{d_x} \times \mathbb{R}^{d_y} \to \mathbb{R}$ satisfies two-sided PL condition, i.e., there exist constants $\mu_x > 0$ and $\mu_y > 0$ such that $f(\cdot, y)$ is μ_x-PL for any $y \in \mathbb{R}^{d_y}$ and $-f(x, \cdot)$ is μ_y-PL for any $x \in \mathbb{R}^{d_x}$.*

Under Assumption 3.1 and 3.2, we define the condition numbers of problem (1) with respect to PL conditions for x and y as $\kappa_x \triangleq L/\mu_x$ and $\kappa_y \triangleq L/\mu_y$ respectively.

We also introduce the following assumption for the existence of saddle point.

Assumption 3.3 (Yang et al. [42]). *We suppose the function $f : \mathbb{R}^{d_x} \times \mathbb{R}^{d_y} \to \mathbb{R}$ has at least one saddle point (x^*, y^*). We also suppose that for any fixed $y \in \mathbb{R}^{d_y}$, the problem $\min_{x \in \mathbb{R}^{d_x}} f(x, y)$ has a nonempty solution set and an optimal value; and for any fixed $x \in \mathbb{R}^{d_x}$, the problem $\max_{y \in \mathbb{R}^{d_y}} f(x, y)$ has a nonempty solution set and a finite optimal value.*

The goal of solving minimax problem under two-sided PL condition is finding an ϵ-approximate solution or ϵ-saddle point that is defined as follows.

Definition 3.3. *We say x is an ϵ-approximate solution of problem (1) if it holds that $g(x) - g(x^*) \le \epsilon$, where $g(x) = \max_{y \in \mathbb{R}^{d_y}} f(x, y)$.*

Definition 3.4. *Under Assumption 3.3, we say (x, y) is an ϵ-saddle point of problem (1) if it holds that $\|x - x^*\|^2 + \|y - y^*\|^2 \le \epsilon$.*

We allow the saddle point does not exist for the problem with one-sided PL condition. In such case, it is guaranteed that $g(x) \triangleq \max_{y \in \mathbb{R}^{d_y}} f(x, y)$ is differentiable [29, Lemma A.5] and we target to find an ϵ-stationary point of $g(x)$.

Definition 3.5. *If the function $g : \mathbb{R}^{d_x} \to \mathbb{R}$ is differentiable, we say x is an ϵ-stationary point of g if it holds that $\|\nabla g(x)\| \le \epsilon$.*

4 A Faster Algorithm for the Two-Sided PL Condition

We first consider the two-sided PL conditioned minimax problem of the finite-sum form (1) under Assumption 3.1, 3.2 and 3.3. We propose a novel stochastic algorithm, which we refer to as SPIDER-GDA. The detailed procedure of our method is presented in Algorithm 1. SPIDER-GDA constructs the stochastic recursive gradient estimators as follows:

$$G_x(x_{t,k}, y_{t,k}) = \frac{1}{B} \sum_{i \in S_x} \big(\nabla_x f_i(x_{t,k}, y_{t,k}) - \nabla_x f_i(x_{t,k-1}, y_{t,k-1}) + G_x(x_{t,k-1}, y_{t,k-1})\big),$$

$$G_y(x_{t,k}, y_{t,k}) = \frac{1}{B} \sum_{i \in S_y} \big(\nabla_y f_i(x_{t,k}, y_{t,k}) - \nabla_y f_i(x_{t,k-1}, y_{t,k-1}) + G_y(x_{t,k-1}, y_{t,k-1})\big).$$

It simultaneously updates two variables **x** and **y** by estimators G_x and G_y with different stepsizes $\tau_x = \Theta(1/(\kappa_y^2 L))$ and $\tau_y = \Theta(1/L)$ respectively. Huang et al. [15], Luo et al. [23], Xian et al. [41] have studied the SPIDER-type algorithm for nonconvex-strongly-concave problem and showed it converges to the stationary point of $g(x) \triangleq \max_{y \in \mathbb{R}^{d_y}} f(x, y)$ sublinearly. However, solving the problem minimax problems with two-sided PL condition desires stronger linear convergence rate, which leads to our theoretical analysis be different from previous work.

We measure the convergence of SPIDER-GDA by the following Lyapunov function

$$\mathcal{V}_{t,k} \triangleq g(x_{t,k}) - g(x^*) + \frac{\lambda \tau_x}{\tau_y} \big(g(x_{t,k}) - f(x_{t,k}, y_{t,k})\big),$$

where $x^* \in \arg\min_{x \in \mathbb{R}^{d_x}} g(x)$ and $\lambda = \Theta(\kappa_y^2)$. We can establish recursion for $\mathcal{V}_{t,k}$ as follows

$$\mathbb{E}[\mathcal{V}_{t,K}] \le \mathbb{E}\left[\mathcal{V}_{t,0} - \frac{\tau_x}{16}\left(2 - \frac{M}{B}\right) \sum_{k=0}^{K-1} \|G_x(x_{t,k}, y_{t,k})\|^2 - \frac{\lambda \tau_x}{16}\left(2 - \frac{M}{B}\right) \sum_{k=0}^{K-1} \|G_y(x_{t,k}, y_{t,k})\|^2\right].$$

Algorithm 1 SPIDER-GDA $(f, (x_0, y_0), T, K, M, B, \tau_x, \tau_y)$

1: $\tilde{x}_0 = x_0, \tilde{y}_t = y_0$

2: **for** $t = 0, 1, \ldots, T-1$ **do**

3:     $x_{t,0} = \tilde{x}_t, y_{t,0} = \tilde{y}_t$

4:     **for** $k = 0, 1, \ldots, K-1$ **do**

5:        **if** $\mod(k, M) = 0$ **then**

6:          $G_x(x_{t,k}, y_{t,k}) = \nabla_x f(x_{t,k}, y_{t,k})$

7:          $G_y(x_{t,k}, y_{t,k}) = \nabla_y f(x_{t,k}, y_{t,k})$

8:        **else**

9:          draw mini-batches S_x and S_y independently with both sizes of B.

10:          $G_x(x_{t,k}, y_{t,k}) = \frac{1}{B} \sum_{i \in S_x} [\nabla_x f_i(x_{t,k}, y_{t,k}) - \nabla_x f_i(x_{t,k-1}, y_{t,k-1}) + G_x(x_{t,k-1}, y_{t,k-1})]$

11:          $G_y(x_{t,k}, y_{t,k}) = \frac{1}{B} \sum_{i \in S_y} [\nabla_y f_i(x_{t,k}, y_{t,k}) - \nabla_y f_i(x_{t,k-1}, y_{t,k-1}) + G_y(x_{t,k-1}, y_{t,k-1})]$

12:        **end if**

13:        $x_{t,k+1} = x_{t,k} - \tau_x G_x(x_{t,k}, y_{t,k})$

14:        $y_{t,k+1} = x_{y,k} + \tau_y G_y(x_{t,k}, y_{t,k})$

15:     **end for**

16:     choose $(\tilde{x}_{t+1}, \tilde{y}_{t+1})$ from $\{(x_{t,k}, y_{t,k})\}_{k=0}^{K-1}$ uniformly at random.

17: **end for**

18: **return** $(\tilde{x}_T, \tilde{y}_T)$

Using the above inequality by setting $M = B = \sqrt{n}$ leads to the estimators $G_x(\tilde{x}_t, \tilde{y}_t)$ and $G_y(\tilde{x}_t, \tilde{y}_t)$ be sufficiently close to the exact gradient and converge to zero linearly, which indicates $g(\tilde{x}_t)$ also converges to $g(x^*)$ linearly. We formally provide the convergence result for SPIDER-GDA in the following theorem and its detailed proof is shown in appendix.

Theorem 4.1. *Under Assumption 3.1, 3.2 and 3.3, we run Algorithm 1 with $M = B = \sqrt{n}$,*
$\tau_y = 1/(5L), \lambda = 32L^2/\mu_y^2, \tau_x = \tau_y/(24\lambda), K = \lceil 4224/(\mu_x \tau_x) \rceil$ and $T = \lceil \log(1/\epsilon) \rceil$. Then the
output $(\tilde{x}_T, \tilde{y}_T)$ satisfies $g(\tilde{x}_T) - g(x^*) \leq \epsilon$ and $g(\tilde{x}_T) - f(\tilde{x}_T, \tilde{y}_T) \leq \epsilon$ in expectation; and it takes
no more than $\mathcal{O}\big((n + \sqrt{n}\kappa_x \kappa_y^2) \log(1/\epsilon)\big)$ SFO calls.

Corollary 4.1. *Under the setting of Theorem 4.1, the output $(\tilde{x}_T, \tilde{y}_T)$ is also a $(16L^2\epsilon/(\mu_x^2 \mu_y))$-*
saddle point in expectation.

Our results provide an SFO upper bound of $\mathcal{O}((n + \sqrt{n}\kappa_x \kappa_y^2) \log(1/\epsilon))$ for finding an ε-approximate solution that is better than the complexity $\mathcal{O}((n + n^{2/3}\kappa_x \kappa_y^2) \log(1/\epsilon))$ derived from SVRG-AGDA [42]. It is possible to use SVRG-type [16, 46] estimators to replace the stochastic recursive estimators in Algorithm 1, which results the algorithm SVRG-GDA. We can prove that SVRG-GDA also has $\mathcal{O}((n + n^{2/3}\kappa_x \kappa_y^2) \log(1/\epsilon))$ SFO upper bound that matches the theoretical result of SVRG-AGDA. We provide the details in Appendix C.

5   Further Acceleration with Catalyst

Both the proposed SPIDER-GDA (Algorithm 1) and existing SVRG-AGDA [42] have the complexities more heavily depend on the condition number of y than the condition number of x. It is natural to ask can we make the dependency of two condition numbers balanced like the results in the strongly-convex-strongly-concave case [20, 23, 38]. In this section, we show it is possible by introducing the Catalyst acceleration.

Algorithm 2 AccSPIDER-GDA

1: $u_0 = x_0$

2: **for** $k = 0, 1, \ldots, K-1$ **do**

3:     $(x_{k+1}, y_{k+1}) = \text{SPIDER-GDA}\big(f(x,y) + \frac{\beta}{2}\|x - u_k\|^2, (x_k, y_k), T, K, M, B, \tau_x, \tau_y\big)$

4:     $u_{k+1} = x_{k+1} + \gamma(x_{k+1} - x_k)$

5: **end for**

6: **option I** (two-sided PL): **return** (x_K, y_K)

7: **option II** (one-sided PL): **return** $(\hat{x}, \hat{y})$ chosen uniformly at random from $\{(x_k, y_k)\}_{k=0}^{K-1}$

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149   We proposed the accelerated SPIDER-GDA (AccSPIDER-GDA) in Algorithm 2 for reducing the
150 computational cost further. Each iteration of the algorithm solve the following sub-problem

$$\min_{x \in \mathbb{R}^{d_x}} \max_{y \in \mathbb{R}^{d_y}} F_k(x, y) \triangleq \min_{x \in \mathbb{R}^{d_x}} \left\{ g(x) + \frac{\beta}{2}\|x - u_k\|_2^2 \right\}. \tag{2}$$

151 by SPIDER-GDA (Algorithm 1). AccSPIDER-GDA has the following convergence result if the
152 sub-problem attain the required accuracy.

153 **Lemma 5.1.** *Under Assumption 3.1, 3.2 and 3.3, we run Algorithm 2 by $\beta = 2L$, $\gamma = 0$ and the*
154 *appropriate setting for the sub-problem solver such that $\|x_k - \tilde{x}_k\|^2 + \|y_k - \tilde{y}_k\|^2 \leq \delta$, where*
155 *$(\tilde{x}_k, \tilde{y}_k)$ is a saddle point of F_{k-1} $(k \geq 1)$ and $\delta = \mu_x \epsilon / (11(\mu_x + 4L)L)$. Then it holds that*

$$\mathbb{E}[g(x_k) - g(x^*)] \leq \left(1 - \frac{\mu_x}{2\beta + \mu_x}\right)^k \big(g(x_0) - g(x^*)\big) + \frac{\epsilon}{2}.$$

156   The setting $\beta = \Theta(L)$ in Lemma 5.1 guarantees the sub-problem (2) has condition number $\tilde{\kappa}_x = \mathcal{O}(1)$
157 for x. It is more well-conditioned on x, we prefer to address the following equivalent problem

$$\max_{y \in \mathbb{R}^{d_y}} \min_{x \in \mathbb{R}^{d_x}} F_k(x, y) = - \min_{y \in \mathbb{R}^{d_y}} \max_{x \in \mathbb{R}^{d_x}} \{-F_k(x, y)\}.$$

158 Using Corollary 4.1 (in the view of changing the roles of x and y), we can find a δ-approximate
159 saddle point of $-F_k(x, y)$ within $\mathcal{O}((n + \sqrt{n}\kappa_y \tilde{\kappa}_x^2) \log(1/\delta)) = \mathcal{O}((n + \sqrt{n}\kappa_y) \log(1/\delta))$ SFO
160 calls in expectation. Additionally, Lemma 5.1 means Algorithm 2 requires $\mathcal{O}(\kappa_x \log(1/\epsilon))$ number
161 of iterations to find an ϵ-approximate solution of the problem. Thus, the total complexity for
162 AccSPIDER-GDA becomes $\mathcal{O}((n\kappa_x + \sqrt{n}\kappa_x \kappa_y) \log(1/\epsilon) \log(1/\delta))$.

163 **Theorem 5.1.** *Under Assumption 3.1, 3.2 and 3.3, if we let $\gamma = 0$, $\beta = 2L$ and use Algorithm*
164 *1 to solve each sub-problem $\max_{y \in \mathbb{R}^{d_y}} \min_{x \in \mathbb{R}^{d_x}} F_k(x, y)$ with $M = B = \sqrt{n}$, $\tau_x = 1/(10L)$,*
165 *$\lambda = 128$, $\tau_y = \tau_x/(24\lambda)$, $K = \lceil 4224/(\mu_y \tau_y) \rceil$, $T = \lceil \log(1/\hat{\delta}) \rceil$ and precision $\hat{\delta}$ is defined by*

$$\hat{\delta} = \min \left\{ \frac{1}{6}, \frac{(\beta - L)\mu_x \mu_y \delta}{192\beta^2(g(x_0) - g(x^*))}, \frac{(\beta - L)\mu_y \delta}{96\beta^2 \epsilon} \right\},$$

166 *where δ is followed by the definition in Lemma 5.1. Then Algorithm 2 can return x_K such that*
167 *$g(x_K) - g(x^*) \leq \epsilon$ in expectation with no more than $\mathcal{O}((n\kappa_x + \sqrt{n}\kappa_x \kappa_y) \log(1/\epsilon) \log(\kappa_x \kappa_y/\epsilon))$*
168 *SFO calls.*

169   Furthermore, Theorem 5.1 implies the more general result as follows.

170 **Corollary 5.1.** *Under Assumption 3.2, 3.1 and 3.3, running Algorithm 2 by appropriate parameters*
171 *could achieve the output x_K such that $g(x_K) - g(x^*) \leq \epsilon$ in expectation within the following SFO*
172 *complexity*

$$\begin{cases} \tilde{\mathcal{O}}\left(\sqrt{n}\kappa_x \kappa_y \log^2(1/\epsilon)\right), & 0 < \sqrt{n} \leq \kappa_x; \\ \tilde{\mathcal{O}}\left(n\kappa_y \log^2(1/\epsilon)\right), & \kappa_x < \sqrt{n} \leq \kappa_x \kappa_y; \\ \mathcal{O}\left((n + \sqrt{n}\kappa_x \kappa_y^2) \log(1/\epsilon)\right), & \kappa_x \kappa_y < \sqrt{n}. \end{cases}$$

173   Lemma 5.1 does not rely on the choice of sub-problem solver, we can apply the acceleration
174 framework in Algorithm 2 by replacing SPIDER-GDA with other algorithms. We summarize the
175 SFO complexities for the acceleration of different algorithms in Table 3.

Table 3: Accelerated results for different methods under two-sided PL condition.

Method	Before Acceleration	After Acceleration	
GDA	$\mathcal{O}(n\kappa_x\kappa_y^2\log(1/\epsilon))$	$\tilde{\mathcal{O}}(n\kappa_x\kappa_y\log^2(1/\epsilon))$	
SVRG-GDA	$\mathcal{O}((n+n^{2/3}\kappa_x\kappa_y^2)\log(1/\epsilon))$	$\begin{cases} \tilde{\mathcal{O}}(n^{2/3}\kappa_x\kappa_y\log^2(1/\epsilon)), & 0 < n^{1/3} \leq \kappa_x; \\ \tilde{\mathcal{O}}\left(n\kappa_y\log^2(1/\epsilon)\right), & \kappa_x < n^{1/3} \leq \kappa_x\kappa_y; \\ \text{no acceleration}, & \kappa_x\kappa_y < n^{1/3}. \end{cases}$	
SPIDER-GDA	$\mathcal{O}((n+\sqrt{n}\kappa_x\kappa_y^2)\log(1/\epsilon))$	$\begin{cases} \tilde{\mathcal{O}}(\sqrt{n}\kappa_x\kappa_y\log^2(1/\epsilon)), & 0 < \sqrt{n} \leq \kappa_x; \\ \tilde{\mathcal{O}}\left(n\kappa_y\log^2(1/\epsilon)\right), & \kappa_x < \sqrt{n} \leq \kappa_x\kappa_y; \\ \text{no acceleration}, & \kappa_x\kappa_y < \sqrt{n}. \end{cases}$	

6 Extension to One-Sided PL Condition

In this section, we show the idea if SPIDER-GDA and its Catalyst acceleration also work for one-sided PL condition. We relax Assumption 3.2 and 3.3 to the following one.

Assumption 6.1. *We suppose that* $-f(x,\cdot)$ *is* μ_y*-PL for any* $x \in \mathbb{R}^{d_x}$*; the problem* $\max_{y\in\mathbb{R}^{d_y}} f(x,y)$ *has a nonempty solution set and an optimal value ;* $g(x) \triangleq \max_{y\in\mathbb{R}^{d_y}} f(x,y)$ *is lower bounded, i.e., we have* $g^* = \inf_{x\in\mathbb{R}^{d_x}} g(x) > -\infty$.

We first show that the SFO complexity of SPIDER-GDA outperforms SVRG-GDA by a factor of $\mathcal{O}(n^{1/6})$ in Theorem 6.1.

Theorem 6.1. *Under Assumption 3.1 and 6.1 , Let* $T = 1$ *and* $M, B, \tau_x, \tau_y, \lambda$ *as defined in Theorem 4.1 and* $K = \lceil 64/(\tau_x\epsilon^2)\rceil$*, then Algorithm 1 can guarantee the output* $\hat{x}$ *to satisfy* $\|\nabla g(\hat{x})\| \leq \epsilon$ *in expectation with no more than* $\mathcal{O}(n + \sqrt{n}\kappa_y^2 L\epsilon^{-2})$ *SFO calls.*

The AccSPIDER-GDA also performs better than SPIDER-GDA in one-sided PL condition for ill conditioned case. In the following lemma, we show that AccSPIDER-GDA could find an approximate stationary point if we solve the sub-problem sufficiently accurate.

Lemma 6.1. *Under Assumption 3.1 and 6.1, if it holds true that* $\|x_k - \tilde{x}_k\|^2 + \|y_k - \tilde{y}_k\|^2 \leq \delta$ *for some saddle point* $(\tilde{x}_k, \tilde{y}_k)$ *of* F_{k-1} $(k \geq 1)$*, where* $\delta = \epsilon^2/(16L^2(1/2\mu_y + 11))$*. Let* $\beta = 2L$*, then for the output* $(\hat{x}, \hat{y})$ *of Algorithm 2, it holds true that*

$$\mathbb{E}\|\nabla g(\hat{x})\|^2 \leq \frac{8\beta(g(x_0) - g^*)}{K} + \frac{\epsilon^2}{2}.$$

Compared with two-sided PL condition, the analysis of AccSPIDER-GDA is more complicated since the precision δ_k at each round are different. By choosing the parameters of the algorithm carefully, we obtain the following results.

Theorem 6.2. *Under Assumption 3.1 and 6.1, if we run Algorithm 2 by* $\gamma = 0, \beta = 2L$ *and use the sub-problem Algorithm 1 to solve with* $M, B, \tau_x, \tau_y, \lambda, K$ *as Theorem 5.1 and* $T = \lceil\log(1/\delta_k)\rceil$*, where*

$$\delta_k = \begin{cases} \min\left\{\frac{1}{4}, \frac{(\beta-L)\mu_y\delta}{16\beta^2\|x_k-x_{k-1}\|^2}\right\}, & k \geq 1; \\ \frac{\delta\mu_y}{2(g(x_0)-g(x^*))}, & k = 0, \end{cases}$$

and δ *is followed by the definition in Lemma 6.1, then Algorithm 2 can find* $\hat{x}$ *such that* $\|\nabla g(\hat{x})\| \leq \epsilon$ *in expectation within* $\mathcal{O}((n + \sqrt{n}\kappa_y L\epsilon^{-2})\log(\kappa_y/\epsilon))$ *SFO calls.*

We can directly set $\beta = 0$ for Algorithm 2 in the case of very large n. The comparison the complexities of algorithm for one-sided PL condition in Table 2. Besides, the algorithms SPIDER-GDA, GDA and SVRG-GDA can also be accelerated with Catalyst framework and we present the corresponding results in Table 4.

Table 4: Acceleration for different methods under one-sided PL condition.

Method	Before Acceleration	After Acceleration
GDA	$\mathcal{O}\left(n\kappa_y^2 L\epsilon^{-2}\right)$	$\tilde{\mathcal{O}}\left(n\kappa_y L\epsilon^{-2}\log(1/\epsilon)\right)$
SVRG-GDA	$\mathcal{O}\left(n + n^{2/3}\kappa_y^2 L\epsilon^{-2}\right)$	$\begin{cases} \tilde{\mathcal{O}}\left(n^{2/3}\kappa_y L\epsilon^{-2}\log(1/\epsilon)\right), & 0 < n^{1/3} \leq \kappa_y; \\ \tilde{\mathcal{O}}\left(nL\epsilon^{-2}\log(1/\epsilon)\right), & \kappa_y < n^{1/3} \leq \kappa_y^2; \\ \text{no acceleration}, & \kappa_y^2 < n^{1/3}. \end{cases}$
SPIDER-GDA	$\mathcal{O}\left(n + \sqrt{n}\kappa_y^2 L\epsilon^{-2}\right)$	$\begin{cases} \tilde{\mathcal{O}}\left(\sqrt{n}\kappa_y L\epsilon^{-2}\log(1/\epsilon)\right), & 0 < \sqrt{n} \leq \kappa_y; \\ \tilde{\mathcal{O}}\left(nL\epsilon^{-2}\log(1/\epsilon)\right), & \kappa_y < \sqrt{n} \leq \kappa_y^2; \\ \text{no acceleration}, & \kappa_y^2 < \sqrt{n}. \end{cases}$

7 Experiments

In this section, we conduct the numerical experiments to show the advantage of proposed algorithms. We consider the following two player Polyak-Łojasiewicz game

$$\min_{x\in\mathbb{R}^d}\max_{y\in\mathbb{R}^d} f(x,y) \triangleq \frac{1}{2}x^\top P x - \frac{1}{2}y^\top Q y + x^\top R y.$$

where

$$P = \frac{1}{n}\sum_{i=1}^n p_i p_i^\top, \quad Q = \frac{1}{n}\sum_{i=1}^n q_i q_i^\top \quad \text{and} \quad R = \frac{1}{n}\sum_{i=1}^n r_i r_i^\top.$$

We independently sample p_i, q_i and r_i from $\mathcal{N}(0,\Sigma_P)$, $\mathcal{N}(0,\Sigma_Q)$ and $\mathcal{N}(0,\Sigma_R)$ respectively. We set the covariance matrix Σ_P as the form of $UDU^\top$ such that $U \in \mathbb{R}^{d\times r}$ is column orthogonal matrix and $D \in \mathbb{R}^{r\times r}$ is diagonal with $r < d$. The diagonal elements of D are distributed uniformly in the interval $[\mu, L]$ with $0 < \mu < L$. The matrix Σ_Q is set by the similar way to Σ_P. We also let $\Sigma_R = 0.1VV^\top$, where each element of V is sampled from $\mathcal{N}(0,1)$ independently. Since the covariance matrices Σ_P and Σ_Q are rank-deficient, it is guarantee that both P and Q are singular. Hence, the objective function is not strongly-convex and not strongly-concave, but it satisfies the two-sided PL-condition [17]. We set $n = 6000, d = 10, r = 5$ and conduct two cases for different condition numbers.

We compare the proposed SPIDER-GDA (Algorithm 1) and AccSPIDER-GDA (Algorithm 2) with the baseline algorithm SVRG-AGDA [42]. We let $B = 1$ and $M = n$ for all of these algorithms and both of the stepsizes for x and y are tuned from $\{10^{-1}, 10^{-2}, 10^{-3}, 10^{-4}, 10^{-5}\}$. For AccSPIDER, we set $\beta = L/(20n)$ and $\gamma = 0.999$. We present the results of the number of SFO calls against the norm of gradient and the distance to the saddle point in Figure 1 and Figure 2. It is clear that our algorithms outperform than baselines.

8 Conclusion and Future Works

In this paper, we have investigated stochastic GDA algorithms for PL conditioned minimax problems. We have proposed the SPIDER-GDA algorithm that reduce the dependency of the sample numbers in SFO complexity. Moreover, we have introduced Catalyst scheme for further acceleration for ill-conditioned problems. Our ideas both work for the two-sided PL and one-sided cases. However, it is unclear whether our algorithms obtain the optimal SFO complexity. The tightness for the minimax problem under PL conditions is still an open problem.

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(a) Distance to saddle point

(b) Norm of gradient

Figure 1: SFO complexity comparison of SVRG-AGDA, SPIDER-GDA and AccSPIDER-GDA for the case of $\boldsymbol{\mu} = 10^{-5}$ and $L = 1$.

(a) Distance to saddle point

(b) Norm of gradient

Figure 2: SFO complexity comparison of SVRG-AGDA, SPIDER-GDA and AccSPIDER-GDA for the case of $\mu = 10^{-9}$ and $L = 1$.

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