On Producing Shortest Cost-Optimal Plans

Anonymized for blind review

Abstract

Cost-optimal planning is at the heart of planning research, with many existing planners that produce provably optimal solutions. While some applications pose additional restrictions, such as producing shortest (in the number of actions) among the cost-optimal plans, standard cost-optimal planning does not provide such a guarantee. We discuss two possible approaches to produce provably the shortest among the cost-optimal plans, one corresponding to an instantiation of cost-algebraic A*, the other based on a cost transformation. We formally prove that the new cost-transformation method indeed produces the shortest among the cost-optimal plans and empirically compare the performance of the approaches in different configurations.

1 Introduction

The aim of classical planning is to determine a plan, which is a sequence of deterministic actions that leads from a given initial state to a goal state. Each action is associated with a non-negative cost and the cost of a plan is defined as the sum of these action costs. In the special case of *optimal* planning, only the plans with *minimal cost among all plans* are considered valid solutions.

Alternative optimization criteria have been considered in the literature: in multi-objective search (Stewart and White 1991; Mandow and la Cruz 2005) there are multiple – typically incomparable and conflicting – cost functions and the aim is to identify all plans with non-dominated cost vectors. Other recent planning variants where the solution consists of entire sets of plans are diverse planning (Nguyen et al. 2012) with the aim to identify a set of qualitatively different plans, top-k planning (Katz et al. 2018) with the aim to identify a set of plans such that no better plan outside the set exists, and top-quality planning (Katz, Sohrabi, and Udrea 2020) with the aim to identify all plans of some bounded quality.

In this paper, we consider the problem of finding a cost-optimal plan that has a minimal number of actions among all cost-optimal plans, or shortest cost-optimal plan. This is for example required as a subroutine for top-quality planning (Katz, Sohrabi, and Udrea 2020) but also an interesting property if the plan gets presented to humans, as it is for example the case in human-in-the-loop planning or explainable AI planning where the optimality is measured with respect to the observer model (Chakraborti et al. 2021).

Note that the problem of finding a shortest cost-optimal plan is different from the multi-objective case. We indeed have two cost measures here, namely plan cost and plan length, but these are not pari passu. If there is a plan π of length 1 (i.e. a single action application) and cost 5 but there is another plan of cost 4 then π is not considered a valid solution, even if the other plan consists of hundreds of actions. So we do not give up the requirement of cost-optimal plans, we only have the additional requirement that among these we want a shortest one in the number of actions.

This concept of shortest cost-optimal plans is loosely related to the concept of strongly optimal plans in stubborn set pruning (Wehrle and Helmert 2014). These are optimal plans with a minimal number of *zero-cost* actions. The difference is that we consider the number of *all* actions, independent of their cost.

Edelkamp, Jabbar, and Lluch Lafuente (2005) already considered the problem of finding optimal paths relative to cost notions that can be very different from the standard sum of action costs. They introduce so-called *cost algebras*, which are a very general concept covering a wide range of cost measures, including the number of actions as well as the sum of action costs. Moreover, both these cost measures are *strictly isotone*, a property that Edelkamp, Jabbar, and Lluch Lafuente identified to be sufficient for their prioritized Cartesian product being again a cost algebra. This prioritized Cartesian product corresponds exactly to the case we consider in this paper: Prefer cheaper plans, and among equally cheap plans, prefer the shorter one.

In this paper, we will consider their cost-algebraic A* algorithm in the specific instantiation for our cost functions and compare this native support empirically to an alternative proposed approach that uses a cost-transformation with standard scalar A*.

2 Shortest Cost-Optimal Planning Problem

For this paper, it is sufficient to consider tasks directly via their transition system $\Pi = \langle S, A, T, cost, s_0, S_* \rangle$, where S is a finite set of *states*, $s_0 \in S$ is the *initial* state and $S_* \subseteq S$ is the set of *goal states*. Set A is the finite set of *actions*, where the *action cost function* $cost : A \to \mathbb{R}^{0+}$ assigns every action a non-negative cost. The *transition relation* $T \subseteq S \times A \times S$ is deterministic, i.e. for every state s and action a, there is at most one s' with $(s, a, s') \in T$. If

there is such an s', we say that a is *applicable* in s and that s' is the successor state of applying a in s. An *s-plan* is a sequence of actions that is consecutively applicable in state s and where the final state is a goal state. A *plan* is an s_0-plan. We will sometimes focus on *simple* applicable action sequences, i.e. applicable action sequences that traverse no state more than once.

For defining the cost of an action sequence, we extend the cost function from single actions to action sequences as $cost(\langle a_1, \ldots, a_n \rangle) := \sum_{i=1,\ldots,n} cost(a_i)$. With $|\pi|$ we denote the *length* of action sequence π, i.e. $|\langle a_1, \ldots, a_n \rangle| = n$.

We denote by $\leq_\Pi$ the partial order of action sequences defined by their cost, i.e., $\pi \leq_\Pi \pi'$ iff $cost(\pi) \leq cost(\pi')$. By $\preceq_\Pi$ we denote the partial order defined by their costs and length, i.e., $\pi \preceq_\Pi \pi'$ if $cost(\pi) < cost(\pi')$ or if $cost(\pi) = cost(\pi')$ and $|\pi| \leq |\pi'|$.

In standard cost-optimal planning we must for a given task Π identify a plan that is minimal among all plans with respect to $\leq_\Pi$ or detect that the task is unsolvable, i.e. it has no plan. In contrast, in the *shortest cost-optimal planning problem* we must find a plan for Π that is minimal among all plans with respect to $\preceq_\Pi$ or again detect Π as unsolvable.

3 Cost-Algebraic Approach

As already mentioned in the introduction, our shortest cost-optimal planning problem is a special case of the problems covered by cost-algebraic heuristic search (Edelkamp, Jabbar, and Lluch Lafuente 2005).

We will now briefly describe the specific instantiation of the cost-algebraic A* algorithm for our scenario.

First, we need a heuristic for the cost-algebraic problem. Such a cost-algebraic heuristic maps every state to a pair $\langle h_c, h_d \rangle \in \mathbb{R}^{0+} \times \mathbb{R}^{0+}$. The heuristic is *admissible* if for every state s it holds that $h_c \leq c^*$ and $h_d \leq d^*$, where c^* is the cost of a cheapest plan for state s and d^* is the length of a shortest among these cheapest plans.

We assume that the reader is familiar with the standard A* algorithm (Hart, Nilsson, and Raphael 1968). It maintains an open list of search nodes, where each search node is associated with a state and with an implicit path from the initial state to this state. In each iteration, A* pops a node from the open list, which in the standard scalar case is ordered by the f-values of the nodes. Here, $f = g_c + h_c$, where g_c is the cost of the associated path from the initial node to the associated state and h_c is the heuristic estimate for this state. If the popped node is a goal node, the search terminates. Otherwise it expands the node by generating all successor nodes and adding the expanded state to the closed list. Each generated node that is not associated with a closed state is added to the open list. Also nodes with closed states are added to open, if the associated path to the state is cheaper than the previous one. This is called reopening and only necessary if the heuristic does not have the additional property of being consistent (a case that we do not consider in this paper).

The cost-algebraic A* variant extends standard A* in two ways: the open list prioritizes by $\langle f_c, f_d \rangle$, where $f_c = g_c + h_c$ and $f_d = g_d + h_d$ with cost-algebraic heuristic estimate $\langle h_c, h_d \rangle$, g_c being the cost of the associated path and g_d being the length of the associated path. So f_c corresponds to the f-value in the standard case and f_d is the analogous value in terms of plan length. Pair $\langle f_c, f_d \rangle$ has a higher priority than $\langle f'_c, f'_d \rangle$ if $f_c < f'_c$ or if $f_c = f'_c$ and $f_d < f'_d$.

The second extension is the reopening criterion. A state is not only reopened if the search encounters a cheaper path to the state but also if it encounters an equally cheap but shorter path to the state.

The results by Edelkamp, Jabbar, and Lluch Lafuente imply that with these two changes, the algorithm is guaranteed to produce shortest cost-optimal plans.

To determine an admissible cost-algebraic heuristic in practice, we can use any (standard) admissible heuristic to determine the cost estimate h_c. For the distance estimate h_d, it is not necessary to only consider minimum-cost plans but also a lower-bound on the length of *any* plan will be an admissible estimate (but possibly lower than necessary). For this reason, we can simply use a standard admissible heuristic, but evaluate it with respect to a different cost function that assigns every action a cost of 1.

Of course the computation of an additional distance estimate causes computational overhead and it is not immediately obvious that it will pay off: since the the first component f_c dominates the ordering of the open list, the A* search will have to consider all action sequences where f_c is strictly lower than the cost of an cost-optimal plan. Hence, the distance information can only be beneficial for the ordering on the last f_c layer. For this reason, we will also consider the extreme case where we set the second heuristic component h_d to constant 0. This is still an admissible estimate, so also this much cheaper configuration guarantees shortest cost-optimal plans.

4 Cost Transformation Approach

The second approach that we consider will use standard cost-optimal solvers but ensures shortest cost-optimal plans by means of a transformation of the action cost function. A different cost function does not alter the set of applicable action sequences but will only influence what plans are considered optimal.

Theorem 1 *For every planning task Π there is a task Π' that only differs from Π in the action cost function, such that for all applicable simple action sequences π and π' it holds that $\pi \preceq_\Pi \pi'$ iff $\pi \leq_{\Pi'} \pi'$.*

Proof: Consider planning task $\Pi = \langle S, A, T, cost, s_0, S_* \rangle$. The length of any applicable simple action sequence is bounded by the number of states $|S|$. Let $\delta > 0$ be smaller or equal to the smallest non-zero cost difference between any two applicable simple action sequences in Π. For integer action costs we can set $\delta = 1$. For real-valued costs, there also is such a δ because there is only a finite number of simple applicable action sequences. We define $\epsilon := \frac{\delta}{|S|} > 0$ and a new cost function $cost_{+\epsilon}$ based on the action cost function $cost$ of Π and ϵ as $cost_{+\epsilon}(a) = cost(a) + \epsilon$. Let Π' be defined as Π only with cost function $cost_{+\epsilon}$, i.e. $\Pi' = \langle S, A, T, cost_{+\epsilon}, s_0, S_* \rangle$.

Let π and π' be two applicable simple action sequences. We distinguish two cases.

If $cost(\pi) = cost(\pi')$, then $\pi \preceq_\Pi \pi'$ iff $|\pi| \leq |\pi'|$ iff $cost_{+\epsilon}(\pi) = cost(\pi) + \epsilon|\pi| \leq cost(\pi') + \epsilon|\pi'| = cost_{+\epsilon}(\pi')$ iff $\pi \leq_{\Pi'} \pi'$.

For the case that $cost(\pi) \neq cost(\pi')$, we show both directions separately, using the following observation: Since both sequences are simple, their length and hence also the difference in their length is bound by $|S| - 1$. Hence, we have $\epsilon(|\pi| - |\pi'|) < \epsilon|S| = \delta$ and therefore $\epsilon|\pi| - \delta < \epsilon|\pi'|$.

If $\pi \preceq_\Pi \pi'$ and $cost(\pi) \neq cost(\pi')$ then $cost(\pi) < cost(\pi')$ and hence $cost(\pi) \leq cost(\pi') - \delta$, by the definition of δ. Together with the observation we get $cost_{+\epsilon}(\pi) = cost(\pi) + \epsilon|\pi| \leq cost(\pi') - \delta + \epsilon|\pi| \leq cost(\pi') + \epsilon|\pi'| = cost_{+\epsilon}(\pi')$ and therefore $\pi \leq_{\Pi'} \pi'$.

If $\pi \leq_{\Pi'} \pi'$ then $cost_{+\epsilon}(\pi) \leq cost_{+\epsilon}(\pi')$. By the definition of $cost_{+\epsilon}$, we get

$$cost(\pi) + \epsilon|\pi| \leq cost(\pi') + \epsilon|\pi'|. \qquad (*)$$

If $|\pi| \geq |\pi'|$ this directly implies $cost(\pi) \leq cost(\pi')$ and with $cost(\pi) \neq cost(\pi')$ that $\pi \preceq_\Pi \pi'$.

If $|\pi| < |\pi'|$, then our observation above changes into $\epsilon(|\pi'| - |\pi|) < \epsilon|S| = \delta$. Assume to the contrary that $cost(\pi) > cost(\pi')$. As before, we then have $cost(\pi) \geq cost(\pi') + \delta$ and $cost(\pi) - cost(\pi') \geq \delta$. From (**?**) we have $cost(\pi) - cost(\pi') \leq \epsilon(|\pi'| - |\pi|)$ and therefore $\delta \leq cost(\pi) - cost(\pi') \leq \epsilon(|\pi'| - |\pi|)$, contradicting $\epsilon(|\pi'| - |\pi|) < \delta$. Therefore, we have $cost(\pi) \leq cost(\pi')$ and with $cost(\pi) \neq cost(\pi')$ we get $\pi \preceq_\Pi \pi'$. □

Based on Theorem 1, one can find shortest cost-optimal plans by finding cost-optimal plans for the transformed cost task. The restriction to simple action sequences is not a limitation for this purpose: shortest cost-optimal plans for the original task are always simple because otherwise there were shorter plans of equal or lower cost (cutting out the cycle); cost-optimal plans in the transformed task are simple because there are no 0-cost actions in this task (which could be used to form a cycle without increasing the cost).

The proof of the theorem is in some respect unsatisfactory, because we only establish the *existence* of a suitable value ε. For the standard case of integer action costs, we suggested in the proof the value $\varepsilon = 1/|S|$, but since $|S|$ is typically very high, ϵ would be very small. This is problematic because planners often only support integer costs, so we need to scale up the cost function $cost_{+\epsilon}$ to integer values, which can lead to extremely high action costs and overflow problems in the planner. For this reason, we are interested in a cost function that serves the same purpose but with lower action costs.

For finding shortest optimal problem, we actually do not need the full generality of Theorem 1 because we do not need to cover *all* applicable simple action sequences but only those relevant to the search.

For this reason, we first revisit the proof of Theorem 1 for a task with integer costs to identify a sufficient requirements for ϵ and each two applicable simple action sequences that we compare during the search:

- If $cost(\pi) = cost(\pi')$ then any positive ϵ is sufficient.
- For $\pi \preceq_\Pi \pi'$ implying $\pi \leq_{\Pi'} \pi'$ if $cost(\pi) \neq cost(\pi')$, we only need that $cost(\pi') - \delta + \epsilon|\pi|$ is smaller than $cost(\pi') + \epsilon|\pi'|$. Hence it is sufficient, if $\epsilon(|\pi| - |\pi'|) < \delta$. Since with integer costs $\delta = 1$ is sufficient for the proof if the length difference of the compared action sequences is strictly smaller than $1/\epsilon$.
- For $\pi \leq_{\Pi'} \pi'$ implying $\pi \preceq_\Pi \pi'$ if $cost(\pi) \neq cost(\pi')$, we only exploit from ϵ that $\epsilon(|\pi'| - |\pi|) < \delta$. Again it is sufficient if the length difference of the compared action sequences is strictly smaller than $1/\epsilon$.

So overall, it is sufficient, if $\epsilon > 0$ and the length difference of any compared action sequences is smaller than $1/\epsilon$. This also makes intuitively sense because it means that the additional action cost we incur overall for accounting for the length of the sequence will never outweigh the original cost of any action application, at least on the action sequences we consider.

What state sequences must cost-algebraic A* consider? Obviously, those that would actually be popped from the open list in cost-algebraic A*are sufficient. Can we characterize these sequences? For standard scalar A*, it is well known that it expands all sequences that have a lower f-value than the optimal solution cost C^*, a subset of the sequences, where the f-value equals C^* (last f-layer; depending on tie-breaking) but no sequences with higher f-value. Since in cost-algebraic A* the f_c-value for the cost dominates the ordering of the open list, it also must expand all sequences with $f_c < C^*$ but from the nodes on the last f_c-layer (with $f_c = C^*$), it will only consider a subset of those, where f_c is less than or equal to the length of a shortest optimal plan.

Of course, we do not have much information about the heuristics for cost and distance, but we know that their estimates are non-negative. So we can conclude that cost-algebraic A* will never pop an action sequence π from the open list for which $cost(\pi) > C^*$.

Let us now combine these two lines of thought and let M be some integer that is strictly larger than the length of any applicable simple action sequence π with $cost(\pi) \leq C^*$. Number M is larger than the length difference of any two action sequences we might need to compare (because it covers to the extreme case of comparing the empty plan with π). From a theoretical perspective, we can always use $|S|$ for M (as we do in the proof of Theorem 1) but this more precise definition allows us to exploit additional insights we might have into the problem at hand. With our previous consideration, any $\epsilon > 0$ with $M \leq 1/\epsilon$ will work in the transformation (still considering tasks with integer costs). In particular, we can set ϵ to $1/M$. With this choice, we get $cost_{+\epsilon}(a) = cost(a) + 1/M$ and $Mcost_{+\epsilon}(a) = Mcost(a) + 1$. Since all action costs with $cost_{+\epsilon}$ are positive, scaling them with a constant factor does not affect the ordering of action sequences by their cost. Overall, we observe that with integer action costs this function $Mcost_{+\epsilon}$ is a cost function suitable for the cost-transformation-based approach to finding shortest optimal plans. Moreover, it can also be used with planners that only support integer action costs, because it maps only into the integers.

In what follows, we will frequently refer to this cost function. Therefore, we include it in the following definition.

Definition 1 *For task* $\Pi = \langle S, A, T, cost, s_0, S_* \rangle$ *and* $M \in \mathbb{N}_{>0}$*, we define cost function* $cost_M$ *as* $cost_M(a) = Mcost(a) + 1$ *and task* Π_M *as* $\Pi_M = \langle S, A, T, cost_M, s_0, S_* \rangle$.

So far our discussion was driven by the behaviour of cost-algebraic A* and indeed this will be relevant for Theorem 3 that establishes that in a certain configuration both approaches will behave equally. However, we can also take a step back because after the cost-transformation we have a purely scalar problem. What would happen if we decrease M below the previous criterion?

Consider an applicable action sequence π of length $|\pi| \geq M$ and another such sequence π' with $cost(\pi') > cost(\pi)$ and length $|\pi'| < |\pi| - M(cost(\pi') - cost(\pi))$. If both sequences are actual plans, π would be better wrt. the optimization criterion because of its lower cost. But under cost $cost_M$, we have $\pi' \leq_{\Pi_M} \pi$, so A* potentially would not consider π but terminate with π'. Is this always a problem or put differently, are there applicable action sequences a search could skip while still finding an optimal solution? Obviously, all action sequences that cannot be extended to an optimal plan can be pruned from the search space without harming completeness or the optimality guarantee. Due to the monotonicity of the plan length, this is the case for all action sequences that are longer than the length L of a shortest optimal plan.

Hence, if we are only interested into finding a shortest optimal plan for task π and we know that the solution will be shorter than some $M \in \mathbb{N}$ then we can simply solve the induced problem Π_M with standard A*.

Heuristics

If we use a heuristic search planner to solve the cost-transformed task, we have some freedom what cost function we use for the heuristic computation.

The first and most obvious option we will consider is to compute the heuristic estimates wrt. the transformed cost function.

However, with this cost function there can be a huge number of action sequences that all exhibit different costs. This can slow down the heuristic computation. For this reason, we will also consider the alternative approach of computing heuristic estimates with the original cost function and scaling them by M. These estimates are admissible wrt. to the transformed cost:

Theorem 2 *Consider task* Π *with cost function* $cost$ *and a* $M \in \mathbb{N}_{>0}$*. If* h *is an admissible estimate for state* s *in* Π *then* Mh *is an admissible estimate for* s *in* Π_M.

Proof: Let π be a plan for state s that is optimal with respect to cost function $cost_M$ and let h be an admissible estimate for s wrt. to cost function $cost$. Since π is an s-plan, it holds that $cost(\pi) \geq h$. As $cost_M(\pi) = Mcost(\pi) + |\pi| \geq Mcost(\pi) \geq Mh$, value Mh is an admissible estimate for s wrt. cost function $cost_M$. □

Interestingly, in the typical setup with integer action costs and integer heuristic estimates, A* with this second heuristic approach leads to the same search behaviour as cost-algebraic A* using the same heuristic for cost and the blind estimator that always sets h_d to 0 for distance.

Theorem 3 *Consider task* $\Pi = \langle S, A, T, cost, s_0, S_* \rangle$ *with* $cost : A \to \mathbb{N}_0$*, heuristic* $h_c : S \to \mathbb{N}_0$ *that is admissible wrt. cost and heuristic* h_d *with* $h_d(s) = 0$ *for all* $s \in S$*. Let* l *be the length of the longest action sequence associated with at search node expanded by cost-algebraic* A* *on* Π *with the given heuristics.*

For $M > l$*, standard* A* *on* Π_M *with heuristic* Mh_c *will expand nodes in the same order as cost-algebraic* A* *on* Π *with heuristics* h_c *and* h_d *if the algorithms use the same-tie breaking.*

Proof: To see why this is the case, we consider the ordering of the open list and the reopening condition. Note that in both variants of A* the paths associated with search nodes are always simple paths.

Let π and π' be applicable simple action sequences with final state s and s', respectively and let (the nodes associated with) π and π' be in the open list.

If the cost-algebraic open list selects π rather than π' from the open list, then either (a) $cost(\pi) + h_c(s) < cost(\pi') + h_c(s')$ or (b) $cost(\pi) + h_c(s) = cost(\pi') + h_c(s')$ and $|\pi| + 0 \leq |\pi'| + 0$. We will show that in both cases $cost_M(\pi) + Mh_c(s) < cost_M(\pi') + Mh_c(s')$ and thus scalar A* will also prioritize π over π' in the open list.

For (a), since π gets expanded, we know that $|\pi| \leq l < M$, and since $cost(\pi) + h_c(s) < cost(\pi') + h_c(s')$ and action costs and heuristic estimates are integers, we know that $cost(\pi) + h_c(s) + 1 \leq cost(\pi') + h_c(s')$. Hence, $cost_M(\pi) + Mh_c(s) = Mcost(\pi) + |\pi| + Mh_c(s) < Mcost(\pi) + M + Mh_c(s) = M(cost(\pi) + h_c(s) + 1) \leq M(cost(\pi') + h_c(s')) \leq Mcost(\pi') + |\pi'| + Mh_c(s') = cost_M(\pi') + Mh_c(s')$ and scalar A* equally must prefer π.

For (b), we get $cost_M(\pi) + Mh_c(s) = Mcost(\pi) + |\pi| + Mh_c(s) = M(cost(\pi) + h_c(s)) + |\pi| = M(cost(\pi') + h_c(s')) + |\pi| = (*)$. If cost-algebraic A* must prefer π because $|\pi| < |\pi'|$, then $(*) < M(cost(\pi') + h_c(s')) + |\pi'| = cost_M(\pi') + Mh_c(s')$ and also scalar A* must prefer π'. If cost-algebraic A* prefers π only due to tie-breaking ($|\pi| = |\pi'|$) then $(*) = cost_M(\pi') + Mh_c(s')$ and scalar A* prefers π because it uses the same tie-breaking.

For the other direction, we again distinguish the situations (1) where the A* open list must prioritize π over π' and (2) where it only does so because of tie-breaking. We show that in both cases, cost-algebraic A* will behave equally.

In situation (1), we have $cost_M(\pi) + Mh_c(s) < cost_M(\pi') + Mh_c(s')$. Expressing $cost_M$ in terms of $cost$, we get $Mcost(\pi) + |\pi| + Mh_c(s) < Mcost(\pi') + |\pi'| + Mh_c(s')$. Let us consider three possible cases: (a) $cost(\pi) + h_c(s) < cost(\pi') + h_c(s')$, (b) $cost(\pi) + h_c(s) = cost(\pi') + h_c(s')$, and (c) $cost(\pi) + h_c(s) > cost(\pi') + h_c(s')$.

For (a), the cost-algebraic approach trivially also must prioritize π over π'. For (b), we have $Mcost(\pi) + Mh_c(s) = Mcost(\pi') + Mh_c(s')$ and therefore $|\pi| < |\pi'|$ and the cost-algebraic approach must prioritizes π over π'. For (c), we have $cost(\pi) + h_c(s) \geq 1 + cost(\pi') + h_c(s')$ because action costs and heuristic estimates are integers. Thus,

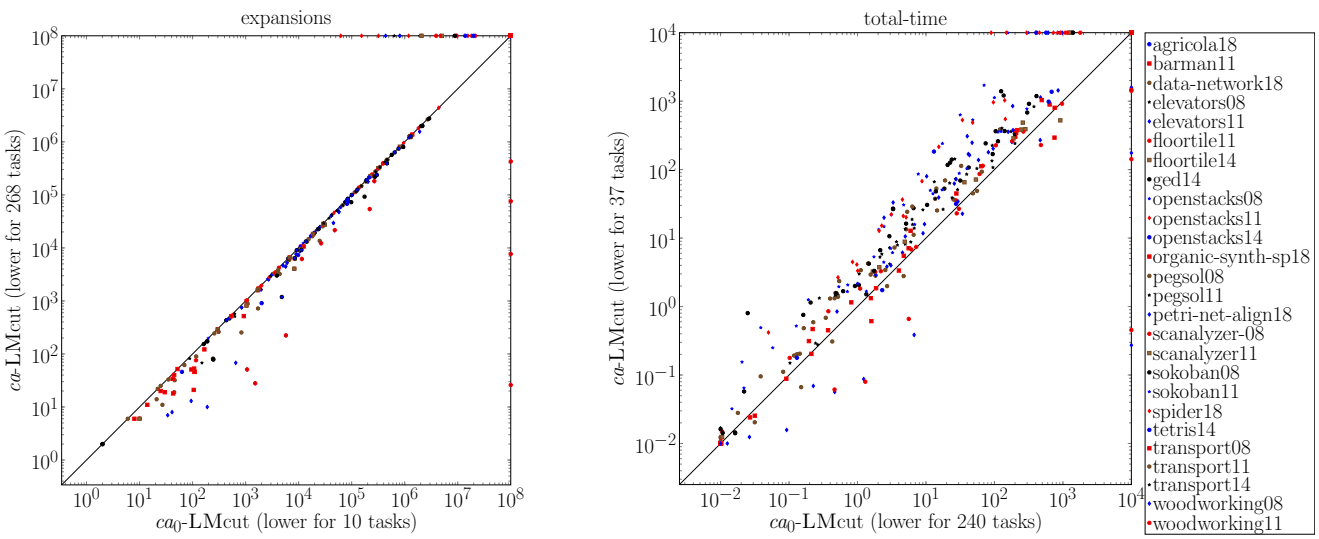

Figure 1: Number of expanded nodes (left) and total time (right), comparing the cost-algebraic approach (*ca*) with informed distance estimates to the (ca_0) approach using constant 0.

$Mcost(\pi) + |\pi| + Mh_c(s) > Mcost(\pi) + Mh_c(s) \geq M + Mcost(\pi') + Mh_c(s') > Mcost(\pi') + |\pi'| + Mh_c(s')$, contradicting $cost_M(\pi) + Mh_c(s) < cost_M(\pi') + Mh_c(s')$.

Overall, only cases (a) and (b) are consistent with situation (1) and in both cases, π must be preferred over π' also by the cost-algebraic search.

In situation (2), we have $cost_M(\pi) + Mh_c(s) = cost_M(\pi') + Mh_c(s')$ or equally $Mcost(\pi) + |\pi| + Mh_c(s) = Mcost(\pi') + |\pi'| + Mh_c(s')$. We distinguish the same cases (a)–(c) as for situation (1). For (a), $cost(\pi) + h_c(s) + 1 \leq cost(\pi') + h_c(s')$ because action costs and heuristic estimates are integer. This is the case iff

$$1 \leq cost(\pi') + h_c(s') - cost(\pi) - h_c(s). \qquad (**)$$

From $Mcost(\pi) + |\pi| + Mh_c(s) = Mcost(\pi') + |\pi'| + Mh_c(s')$ we get $|\pi| - |\pi'| = Mcost(\pi') + Mh_c(s') - Mcost(\pi) - Mh_c(s) \geq M$ (last inequality from **). Since π is longer than any plan expanded by cost-algebraic A*, it won't be expanded. Since π would be preferred over π', this means that π and π' cannot both be in the open list of cost-algebraic A* and (a) is inconsistent with situation (2). For (b) we have again $Mcost(\pi) + Mh_c(s) = Mcost(\pi') + Mh_c(s')$, which in situation (2) implies $|\pi| = |\pi'|$. Hence cost-algebraic A* will pick π due to the same tie-breaking. For (c) we get the analogous contradiction between $cost_M(\pi) + Mh_c(s) > cost_M(\pi') + Mh_c(s')$ and $cost_M(\pi) + Mh_c(s) = cost_M(\pi') + Mh_c(s')$ as in situation (1).

Overall, only case (b) is relevant for situation (2) and also there cost-algebraic A* prefers π over π' due to tie-breaking.

Finally, let n be some node generated by a search. If the cost-algebraic approach reopens n, then n was previously reached by a path π' and now is reached by a path π such that either (a) $cost(\pi) < cost(\pi')$ or (b) $cost(\pi) = cost(\pi')$ and $|\pi| < |\pi'|$. In both cases, we have $cost_M(\pi) < cost_M(\pi')$ and therefore A* will reopen n.

For the other direction, if A* reopens n, then n was previously reached by a path π' and now is reached by a path π such that $cost_M(\pi) < cost_M(\pi')$. Expressing $cost_M$ in terms of $cost$, we get $Mcost(\pi) + |\pi| < Mcost(\pi') + |\pi'|$. As before, let us consider three possible cases (a) $cost(\pi) < cost(\pi')$, (b) $cost(\pi) = cost(\pi')$, and (c) $cost(\pi) > cost(\pi')$. As before, for (c) we get a contradiction. For (a), the cost-algebraic approach will also reopen n. For (b), we have $Mcost(\pi) = Mcost(\pi')$ and therefore $|\pi| < |\pi'|$ and the cost-algebraic approach will reopen n.

□

5 Experimental Evaluation

To empirically compare the various suggested methods, we have implemented these methods on top of the Fast Downward planning system (Helmert 2006). The code will be available upon acceptance. The experiments were performed on Intel(R) Xeon(R) Gold 6248 CPU @2.50GHz machines, with the time and memory limit of 30min and 3.5GB, respectively. The benchmark set consists of all STRIPS benchmarks with non-unit costs (with the exception of parcprinter[1]) from optimal tracks of International Planning Competitions 1998-2018, a total of 587 tasks in 28 domains. We exclude unit-cost tasks because for these all optimal plans are shortest optimal plans, so there is not requirement for a special algorithm. As all tasks have integer costs and Fast Downward supports integer costs only, in order to obtain a cost transformation into integer values, we assume a

[1]The parcprinter domain has actions with costs in hundreds of thousands. After cost transformation the domain would have actions with costs in billions.

	LMcut		M&S		CEGAR		h_{max}		iPDB	
Coverage	ca_0	ca	ca_0	ca	ca_0	ca	ca_0	ca	ca_0	ca
data-network18 (20)	12	12	9	9	14	14	11	11	**12**	11
elevators08 (30)	**22**	19	15	15	20	20	17	17	23	23
elevators11 (20)	**18**	17	12	12	17	17	**14**	13	18	18
floortile11 (20)	**7**	6	4	4	2	2	6	6	2	2
floortile14 (20)	**6**	5	2	2	0	0	5	5	0	0
openstacks08 (30)	**22**	16	22	22	**22**	20	**22**	21	22	22
openstacks11 (20)	**17**	10	17	17	**17**	15	**17**	16	17	17
openstacks14 (20)	**3**	1	3	3	3	3	3	3	3	3
organic-s-sp18 (20)	14	14	**7**	6	9	9	**19**	18	**7**	6
pegsol08 (30)	27	27	**29**	28	28	28	27	27	29	29
pegsol11 (20)	17	17	**19**	18	18	18	17	17	19	19
petri-net-align18 (20)	**9**	8	4	4	1	1	**11**	9	0	0
scanalyzer-08 (30)	**11**	9	12	12	12	12	9	9	13	13
scanalyzer11 (20)	**7**	6	9	9	9	9	6	6	10	10
sokoban08 (30)	**29**	28	27	27	23	23	28	28	**30**	29
sokoban11 (20)	20	20	20	20	20	20	20	20	**20**	19
spider18 (20)	**11**	8	13	13	**11**	10	**9**	8	**13**	3
tetris14 (17)	5	5	7	7	**9**	8	**9**	7	1	1
transport11 (20)	6	6	6	6	6	6	6	6	**12**	11
transport14 (20)	6	6	7	7	7	7	**8**	6	9	9
woodworking08 (30)	13	**17**	14	14	10	**11**	10	10	12	12
woodworking11 (20)	8	**12**	9	9	5	**6**	5	5	7	7
Sum other (90)	30	30	34	34	30	30	34	34	37	37
Sum (587)	**320**	299	**301**	298	**293**	289	**313**	302	**316**	301

Table 1: Per-domain coverage for the cost algebra methods.

bound of $M = 10000$ on the plan length. Transformed action costs $cost_M(a)$ are therefore $Mcost(a)+1$. As there are no existing planners that produce shortest optimal plans, the suggested methods are compared to each other, with a variety of admissible heuristics. The heuristics we experiment with are the blind heuristic, LMcut (Helmert and Domshlak 2009), merge-and-shrink abstraction (denoted by M&S) (Helmert, Haslum, and Hoffmann 2007), counterexample-guided Cartesian abstraction refinement (denoted by CEGAR) (Seipp and Helmert 2018), h_{max} (Bonet and Geffner 2001), and pattern database heuristic iPDB (Haslum et al. 2007). We measure the search progress in terms of the number of expanded nodes as well as the number of heuristic evaluations; for the required time, we distinguish search time (used by the A*algorithms) and total time total time (including heuristic pre-computations); coverage refers to the number of problems solved.

5.1 Cost-Algebraic Approach

As already discussed in Section 3, using a distance estimator in cost-algebraic A* can only be beneficial for the ordering of the last f_c layer, so it is worth evaluating whether the additional effort for a distance estimator still pays off.

For this reason, we compare the cost-algebraic approach (*ca*) with informative heuristic functions for both, cost and distance estimates, to the one that only uses an informed heuristic for cost estimates but constant 0 for the distance estimate (ca_0). For simplicity, we use the same heuristic for both cost and distance estimates, simply replacing all action costs with 1 for the distance estimate. Figure 1 (left) compares *ca*-LMcut to ca_0-LMcut in terms of expanded nodes: when using an additional heuristic for distance estimation the number of expanded nodes indeed decreases. However, as per-node evaluation time also increases for most tasks this does not pay off in terms of total time (Figure 1, right). This translates directly to the coverage of the two approaches: *ca*-LMcut solves 299 tasks overall, compared to 320 of ca_0-LMcut. The only domains where *ca*-LMcut has a better coverage are the woodworking domains (17 vs. 13 and 12 vs. 8). Per-domain coverage results for multiple heuristics are depicted in Table 1. The best results per domain and heuristic are bolded. Note that the results are consistent across the various heuristics. Notably, the largest difference in coverage is on the spider domain, 13 for ca_0-iPDB vs. 3 for *ca*-iPDB.

Overall, the reduction in the search effort from using an informed distance estimator does not appear to outweigh the additional computation cost, both time- and coverage-wise, for all tested heuristics. For this reason, in the following experiments we will only compare to the better-performing variant ca_0.

5.2 Cost-Transformation Approach

For the cost-transformation-based approach, there is only a single heuristic for estimating the cost to go, but no separate distance estimator. One obvious option for the heuristic is to compute it directly from the cost-transformed task (= the standard behaviour of Fast Downward). We denote this configuration by *ct*. The other option is to use the cost-transformed task for the computation of g-values in the search but to compute the heuristic estimates on the original task and scaling them with constant factor M. The resulting estimates are admissible for the cost-transformed task as established by Theorem 2. We denote this second approach of searching on the cost-transformed task while computing heuristic values on the original task and scaling them for the search by *wct*.

Table 2 shows the per-domain coverage comparison. Observe that, with the exception of h_{max} heuristic, these two methods are quite complementary across the tested heuristics, and the difference in coverage on a particular domain can be quite large.

While the heuristic computed directly on the cost-transformed task (*ct*) is more informed, the transformed cost function can have a negative impact on the computation-time of some heuristics. We observe that the overhead incurred by the cost transformation is very different for the different heuristics but also with the same heuristic, the balance can be on the positive or the negative side, depending on the domain. Both effects are very pronounced in the LMcut heuristic, hence we present more details for this heuristic. Figure 2 shows the details of the comparison in terms of the number of heuristic evaluations per time unit, as well as the number of expanded nodes. We see the dominance of the

Coverage	blind		LMcut		M&S		CEGAR		h_{max}		iPDB	
	ct	wct	ct	wct	ct	wct	ct	wct	ct	wct	ct	wct
data-network18 (20)	7	7	12	12	9	9	13	**14**	11	11	12	12
elevators08 (30)	14	14	21	**22**	15	**16**	**21**	20	17	17	**25**	23
elevators11 (20)	12	12	16	**18**	13	13	**18**	17	14	14	**19**	18
floortile11 (20)	2	2	7	7	4	4	2	2	6	6	**3**	2
openstacks08 (30)	22	22	15	**22**	22	22	**25**	22	22	22	16	**22**
openstacks11 (20)	17	17	10	**17**	17	17	**19**	17	17	17	11	**17**
openstacks14 (20)	3	3	1	**3**	3	3	**4**	3	3	3	2	**3**
organic-s-sp18 (20)	10	10	13	**14**	6	6	9	9	19	19	6	**7**
pegsol08 (30)	27	27	27	27	29	29	28	28	28	28	**30**	29
pegsol11 (20)	17	17	17	17	19	19	18	18	18	18	**20**	19
petri-net-align18 (20)	4	4	8	**9**	0	**4**	1	**2**	10	**11**	0	0
scanalyzer-08 (30)	12	12	**15**	10	**13**	12	12	12	9	9	13	13
scanalyzer11 (20)	9	9	**12**	7	**10**	9	9	9	6	6	10	10
sokoban08 (30)	22	22	29	29	27	27	**24**	23	28	28	30	30
spider18 (20)	11	11	10	**11**	13	13	11	11	9	9	3	**13**
woodworking08 (30)	8	8	**17**	13	**15**	14	**11**	10	11	11	**13**	12
woodworking11 (20)	3	3	**12**	8	**10**	9	**6**	5	6	6	**8**	7
Sum other (187)	71	71	73	73	76	76	72	72	82	82	79	79
Sum (587)	271	271	315	**319**	301	**302**	**303**	294	316	**317**	300	**316**

Table 2: Per-domain coverage for the cost transformation methods.

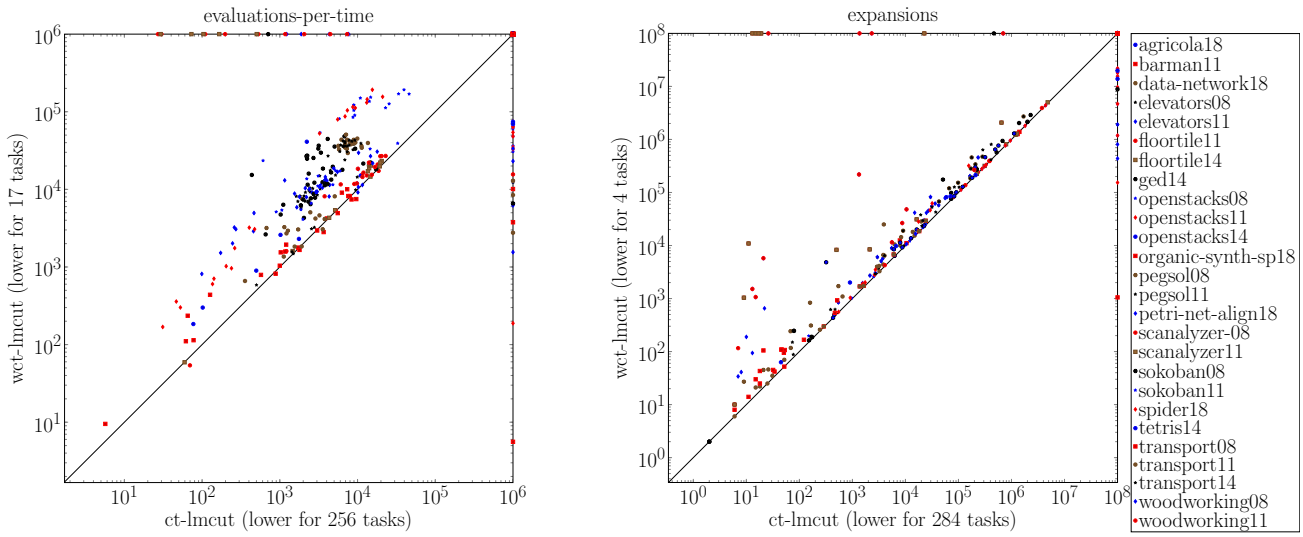

Figure 2: Number of heuristic evaluations per time (left) and expansions (right), comparing the both cost-transformation configurations *ct* and *wct* with the LMcut heuristic.

heuristic computation used in *ct* clearly reflected in the number of expansions (right). But if we have a look at the time required for the heuristic computation (left), we observe that the computation on the original cost function with subsequent scaling is almost consistently faster. The speedup is most prominent on the openstacks domains, which also benefit most from the *wct* configuration in terms of coverage.

Overall, the direct computation is better-informed but

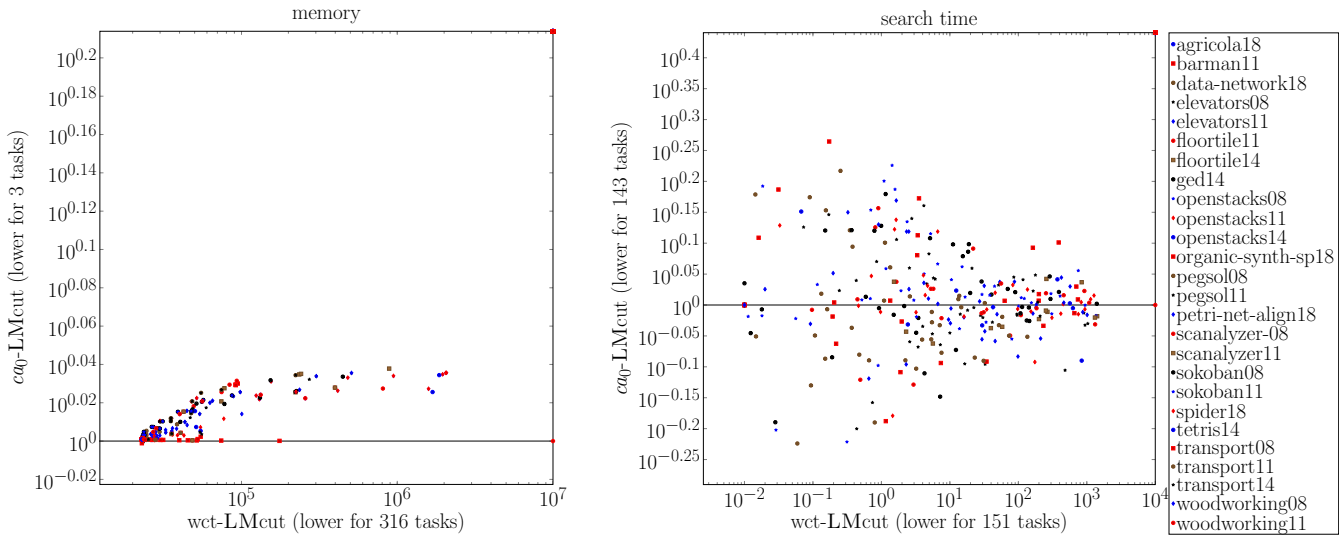

Figure 3: Memory (left) and search time (right), comparing cost algebraic configuration ca_0 to cost-transformation configuration *wct*.

	LMcut		M&S		CEGAR		h_{max}		iPDB	
Coverage	ca_0	wct	ca_0	wct	ca_0	wct	ca_0	wct	ca_0	wct
Sum (587)	**320**	319	301	**302**	293	**294**	313	**317**	**316**	**316**

Table 3: Comparison ca_0 and wct.

can have a negative impact on the computation time with some heuristics. Whether the potentially lower computation time of the scaling approach outweighs the somewhat lower heuristic guidance depends on the balance of the two effects (speedup and impact on guidance), which varies from domain to domain.

5.3 Cost-Algebraic vs. Cost-Transformation

Theorem 3 guarantees that the search behavior of ca_0 and *wct* to be exactly the same, that is, they examine the same nodes in the same order. However, the computational overhead can still be different. For this reason, we also experimentally compare these two configurations to each other.

We indeed can empirically verify that both methods always result in the exact same number of expansions. Figure 3 shows the memory consumption and the search time comparison between the two approaches for the LMcut heuristic. We see that the direct cost-algebraic implementation (ca_0) consistently suffers from the memory overhead for representing the pairs for ordering the open list. Still, Table 3, summarizing the overall coverage of these two approaches, shows that the impact on the overall coverage is low. Both approaches are very much on par, with the only larger coverage difference occurring with the h_{max} heuristic, on pergsol and woodworking domains.

6 Conclusions and Future Work

We consider the problem of finding shortest cost-optimal plans, a useful problem for which to date no planners existed. We have explored the existing algorithm based on cost algebra, with different distance estimators and proposed an alternative approach based on cost transformation. We have formally proved the search-equivalence of a cost algebra based approach to a cost transformation based one. Finally, we have implemented and empirically compared these approaches, identifying their respective strong sides and shortcomings.

As this work is the first that is explicitly covering the shortest optimal planning problem, we believe that this is the beginning of a long journey. One interesting avenue for future work is understanding whether one of these methods will excel on a given task and heuristic, without performing any search. Another possible direction is an efficient heuristic computation, either simultaneously computing heuristics for both cost and distance or computing a heuristic for a transformed cost task. It would also be interesting to explore whether existing search pruning techniques (Pochter, Zohar, and Rosenschein 2011; Alkhazraji et al. 2012) require any adaptation in order to be applied to the cost algebra based A* in a shortest optimal planning setting.

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