Demystifying Topological Message-Passing with Relational Structures: A Case Study on Oversquashing in Simplicial Message-Passing

Abstract

Topological deep learning (TDL) has emerged as a powerful tool for modeling higher-order interactions in relational data. However, phenomena such as oversquashing in topological message-passing remain understudied and lack theoretical analysis. We propose a unifying axiomatic framework that bridges graph and topological message-passing by viewing simplicial and cellular complexes and their message-passing schemes through the lens of relational structures. This approach extends graph-theoretic results and algorithms to higher-order structures, facilitating the analysis and mitigation of oversquashing in topological message-passing networks. Through theoretical analysis and empirical studies on simplicial networks, we demonstrate the potential of this framework to advance TDL.

1 Introduction

Recent years have witnessed a growing recognition that traditional machine learning, rooted in Euclidean spaces, often fails to capture the complex structure and relationships present in real-world data. This shortcoming has driven the development of geometric deep learning (GDL) (Bronstein et al., 2021) and, more recently, topological deep learning (TDL) (Hajij et al., 2023), for handling non-Euclidean and relational data. TDL, in particular, has emerged as a promising frontier for relational learning, that extends beyond graph neural networks (GNNs). TDL offers tools to capture and analyze higher-order interactions and topological features in complex data and higher-order structures, such as simplicial complexes, cell complexes, and sheaves (Hajij et al., 2023). However, the TDL field is young, and the TDL community has yet many open theoretical and practical questions relating to, e.g., oversquashing and rewiring (research directions 2 and 9 of Papamarkou et al., 2024).

Oversquashing is a challenging failure mode in GNNs, where information struggles to propagate across long paths due to the compression of an exponentially growing number of messages into fixed-size vectors (Alon & Yahav, 2021). This phenomenon has been examined through various perspectives, including curvature (Topping et al., 2022), graph expansion (Banerjee et al., 2022), effective resistance (Black et al., 2023), and spectral properties (Karhadkar et al., 2023). Despite the potential of higher-order message passing architectures—such as simplicial neural networks (Ebli et al., 2020), message passing simplicial networks (Bodnar et al., 2021b), and CW networks (Bodnar et al., 2021a)—there remains a lack of unified frameworks for analyzing and mitigating oversquashing in these settings.

In this paper, we take a first step toward studying oversquashing in TDL by showing that simplicial complexes and their message passing schemes can be interpreted as relational structures, making it possible to extend key GNN insights and tools to higher-order message passing architectures. The conceptual framework and theoretical results developed in this paper address pressing questions in the TDL community (e.g., research directions 2 and 9 of Papamarkou et al., 2024).

Contributions. Our contributions are threefold:

- **Axiomatic**: We provide a unifying view of simplicial complexes and their message passing schemes through the lens of relational structures.
- **Theoretical**: We introduce *influence graphs* which enable novel extensions of prior graph analyses to higher order structures, where existing methods for analysis do not apply. We

extend graph-theoretic concepts and results on oversquashing to relational structures, analyzing network sensitivity (Lemma 3.2), local geometry (Proposition 3.4), and the impact of network depth (Theorem 3.5) and hidden dimensions (Section 3.4).

- **Practical**: We propose a heuristic to extend oversquashing-mitigation techniques from graph-based models to relational structures.

Related Work. Our work sits at the intersection of graph neural networks, topological deep learning, relational learning, and the study of oversquashing and graph rewiring in graph neural networks. We review related work in Appendix A.

The rest of this paper is organized as follows: Section 2 provides the axiomatic groundwork for relating simplicial and relational message passing. Section 3 presents our theoretical analysis of oversquashing in this context. Section 4 introduces a heuristic rewiring strategy to mitigate oversquashing in relational message passing. Section 5 presents our experimental results, followed by a discussion and conclusions in Section 6.

2 SIMPLICIAL COMPLEXES ARE RELATIONAL STRUCTURES

In this section, we reinterpret simplicial complexes and message passing through the lens of *relational structures*. We begin by recalling simplicial complexes and a representative simplicial message passing scheme, then reframe these notions within a relational framework. We illustrate the definitions in this section with a small worked example in Appendix H.

2.1 SIMPLICIAL COMPLEXES AND MESSAGE PASSING

Simplicial complexes are mathematical structures that generalize graphs to higher dimensions, capturing relationships among vertices, edges, triangles, and higher-dimensional objects.

Definition 2.1 (Simplicial Complex, Nanda, 2021). *Let V be a non-empty set. A simplicial complex $\mathcal{K}$ is a collection of non-empty subsets of V that contains all the singleton subsets of V and is closed under the operation of taking non-empty subsets.*

A member $\sigma = \{v_0, v_1, \ldots, v_d\} \in \mathcal{K}$ with cardinality $|\sigma| = d + 1$ is called a *d-simplex*. Geometrically, 0-simplices are vertices, 1-simplices are edges, 2-simplices are triangles, and so on.

Definition 2.2 (Boundary Incidence Relation). *We say that τ covers σ, written $\sigma \prec \tau$, iff $\sigma \subsetneq \tau$ and there is no $\delta \in \mathcal{K}$ such that $\sigma \subsetneq \delta \subsetneq \tau$.*

The incidence relations from Definition 2.2 can be used to construct four types of (local) adjacencies.

Definition 2.3. *Consider a simplex $\sigma \in \mathcal{K}$. Four types of adjacent simplices can be defined:*

1. Boundary adjacency*: $\mathcal{B}(\sigma) = \{\tau \colon \tau \prec \sigma\}$;*
2. Co-boundary adjacency*: $\mathcal{C}(\sigma) = \{\tau \colon \sigma \prec \tau\}$;*
3. Lower adjacency*: $\mathcal{N}_\downarrow(\sigma) = \{\tau \colon \exists \delta \text{ such that } \delta \prec \tau \text{ and } \delta \prec \sigma\}$;*
4. Upper adjacency*: $\mathcal{N}_\uparrow(\sigma) = \{\tau \colon \exists \delta \text{ such that } \tau \prec \delta \text{ and } \sigma \prec \delta\}$.*

In Figure 1, we illustrate an example of a simplicial complex and its adjacency relations.

We now, following Bodnar et al. (2021b, Section 4), review a general scheme for message passing on simplicial complexes. While the field of topological deep learning is relatively young, it has seen rapid growth in recent years. In Appendix A, we provide references for topological message passing architectures that fit the general scheme that we review below. For comprehensive overviews of topological message passing schemes and architectures, we refer readers to Papillon et al. (2023b), Hajij et al. (2023) and Giusti (2024). We also refer readers to Appendix F.5 where we list the update rules for all the graph and topological message passing models in this work, which sheds light on the specifics of some architectures that fit these schemes.

Simplicial message passing generalizes graph message passing by exchanging features between adjacent simplices, extending pairwise connections between nodes in a graph to higher-dimensional adjacency relationships between simplices in a simplicial complexes Message passing schemes on simplicial complexes iteratively update feature vectors assigned to simplices by exchanging messages between adjacent simplices. We now review a typical simplicial message passing scheme following Bodnar et al. (2021b): For a simplicial complex $\mathcal{K}$, we denote the feature vector of a

(a) Simplicial Complex $\mathcal{K}$ (b) Adjacency Relations on $\mathcal{K}$

$\rightarrow$ Boundary $\rightarrow$ Co-boundary $\rightarrow$ Lower $\rightarrow$ Upper

Figure 1: The left panel shows a simplicial complex $\mathcal{K}$ consisting of five vertices, five edges, and one triangle. The right panel shows the corresponding adjacency relations depicted as arrows to each simplex $\sigma \in \mathcal{K}$ emanating from each of its adjacent simplices in $\mathcal{B}(\sigma), \mathcal{C}(\sigma), \mathcal{N}_\downarrow(\sigma), \mathcal{N}_\uparrow(\sigma)$.

simplex $\sigma \in \mathcal{K}$ as $\mathbf{h}_\sigma \in \mathbb{R}^p$. At each iteration (layer) t, the feature vectors $\mathbf{h}_\sigma^{(t)}$ of simplices $\sigma \in \mathcal{K}$ are updated by aggregating messages from adjacent simplices. For a simplex $\sigma \in \mathcal{K}$, the messages passed from adjacent simplices are defined as follows:

$$
\begin{aligned}
\mathbf{m}_\mathcal{B}^{(t+1)}(\sigma) &= \mathrm{AGG}_{\tau \in \mathcal{B}(\sigma)}\left(M_\mathcal{B}(\mathbf{h}_\sigma{}^{(t)}, \mathbf{h}_\tau{}^{(t)})\right), \\
\mathbf{m}_\mathcal{C}^{(t+1)}(\sigma) &= \mathrm{AGG}_{\tau \in \mathcal{C}(\sigma)}\left(M_\mathcal{C}(\mathbf{h}_\sigma{}^{(t)}, \mathbf{h}_\tau{}^{(t)})\right), \\
\mathbf{m}_\downarrow^{(t+1)}(\sigma) &= \mathrm{AGG}_{\tau \in \mathcal{N}_\downarrow(\sigma)}\left(M_\downarrow(\mathbf{h}_\sigma{}^{(t)}, \mathbf{h}_\tau{}^{(t)}, \mathbf{h}_{\sigma \cap \tau}{}^{(t)})\right), \\
\mathbf{m}_\uparrow^{(t+1)}(\sigma) &= \mathrm{AGG}_{\tau \in \mathcal{N}_\uparrow(\sigma)}\left(M_\uparrow(\mathbf{h}_\sigma{}^{(t)}, \mathbf{h}_\tau{}^{(t)}, \mathbf{h}_{\sigma \cup \tau}{}^{(t)})\right).
\end{aligned}
\tag{1}
$$

Here AGG is an aggregation function (e.g., sum or mean), and $M_\mathcal{B}, M_\mathcal{C}, M_\downarrow, M_\uparrow$ are message functions (e.g., linear or MLP). Then, an update operation UPDATE (e.g., MLP) incorporates these four different types of incoming messages:

$$
\mathbf{h}_\sigma^{(t+1)} = \mathrm{UPDATE}\left(\mathbf{h}_\sigma^{(t)}, \mathbf{m}_\mathcal{B}^{(t+1)}(\sigma), \mathbf{m}_\mathcal{C}^{(t+1)}(\sigma), \mathbf{m}_\downarrow^{(t+1)}(\sigma), \mathbf{m}_\uparrow^{(t+1)}(\sigma)\right).
\tag{2}
$$

Finally, after the last iteration, a read-out function is applied to process the features to perform a desired task, such as classification or regression.

2.2 RELATIONAL STRUCTURES AND MESSAGE PASSING

We model simplicial complexes and the above message passing scheme using *relational structures*.

Definition 2.4 (Relational Structure, Hodges, 1993). *A relational structure* $\mathcal{R} = (\mathcal{S}, R_1, \ldots, R_k)$ *consists of a finite set* $\mathcal{S}$ *of* entities, *and relations* $R_i \subseteq \mathcal{S}^{n_i}$, *where* n_i *is the arity of* R_i.

We note that modeling simplicial complexes as relational structures generalizes a powerful perspective in which simplicial complexes and similar constructs are treated as augmented Hasse diagrams—a viewpoint that has shown both practical and theoretical advantages, as demonstrated, for example, by Hajij et al. (2023), Eitan et al. (2024), and Papillon et al. (2024).

We now introduce a general scheme for message passing on relational structures which encompasses the simplicial message passing scheme from Section 2.1. This scheme is an extension of the relational graph convolution model from Schlichtkrull et al. (2018) which allows for relations of different arities, not just binary relations.

Definition 2.5 (Relational Message Passing Model). *A relational message passing model* *on a relational structure* $\mathcal{R} = (\mathcal{S}, R_1, \ldots, R_k)$ *consists of:*

- Feature vectors*: $\mathbf{h}_\sigma^{(t)} \in \mathbb{R}^{p_t}$ for each $\sigma \in \mathcal{S}$ at layer $t \geq 0$, initialized as $\mathbf{h}_\sigma^{(0)} = \mathbf{x}_\sigma$ (input features). Here, p_t denotes the dimensionality of the feature vectors at layer t.*
- Message functions*: $\boldsymbol{\psi}_i^{(t)} \colon \mathbb{R}^{p_t} \times \cdots \times \mathbb{R}^{p_t} \to \mathbb{R}^{p_{i,t}}$ (with n_i arguments) for each relation R_i, where $i = 1, \ldots, k$. Each message function takes n_i input feature vectors (corresponding to the*

target simplex and its $n_i - 1$ related simplices) and outputs a message vector of dimension $p_{i,t}$.
The parameter n_i represents the arity of the relation R_i.

- Update function $\phi^{(t)} \colon \mathbb{R}^{p_{1,t}} \times \cdots \times \mathbb{R}^{p_{k,t}} \to \mathbb{R}^{p_{t+1}}$. *The output dimension p_{t+1} specifies the dimensionality of the feature vectors at layer $t + 1$.*

- Shift operators $\mathbf{A}^{R_i} \in \mathbb{R}_{\geq 0}^{|\mathcal{S}|^{n_i}}$ *associated with each relation R_i, for $i = 1, \ldots, k$. For each $\sigma \in \mathcal{S}$ and $\boldsymbol{\xi} = (\xi_1, \ldots, \xi_{n_i-1}) \in \mathcal{S}^{n_i-1}$ with $(\sigma, \boldsymbol{\xi}) \in R_i$, the element $\mathbf{A}^{R_i}_{\sigma, \boldsymbol{\xi}}$ represents the strength of the signal passed from $\boldsymbol{\xi}$ to σ. More specifically, for any combination of entities $(\zeta_1, \zeta_2, \ldots, \zeta_{n_i}) \in \mathcal{S}^{n_i}$ where the relation R_i does not hold among the entities $\zeta_1, \zeta_2, \ldots, \zeta_{n_i}$ (i.e., $(\zeta_1, \zeta_2, \ldots, \zeta_{n_i}) \notin R_i$), the tensor $\mathbf{A}^{R_i}$ satisfies $\mathbf{A}^{R_i}_{\zeta_1, \zeta_2, \ldots, \zeta_{n_i}} = 0$.*

The update rule is given by:

$$\mathbf{h}^{(t+1)}_\sigma = \phi^{(t)} \left(\mathbf{m}^{(t)}_{\sigma,1}, \ldots, \mathbf{m}^{(t)}_{\sigma,k} \right), \tag{3}$$

where for each $i = 1, \ldots, k$, the message $\mathbf{m}^{(t)}_{\sigma,i}$ received by σ over R_i is computed as:

$$\mathbf{m}^{(t)}_{\sigma,i} = \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-1}} \mathbf{A}^{R_i}_{\sigma,\boldsymbol{\xi}} \, \boldsymbol{\psi}^{(t)}_i \left(\mathbf{h}^{(t)}_\sigma, \mathbf{h}^{(t)}_{\xi_1}, \ldots, \mathbf{h}^{(t)}_{\xi_{n_i-1}} \right). \tag{4}$$

Remark 2.6. *The shift operators in Definition 2.5 extend the definition of graph shift operators (Mateos et al., 2019; Gama et al., 2020; Dasoulas et al., 2021) from graphs to relational structures. Whereas relations are combinatorial data structures that specify whether entities are connected, the shift operators are numerical representations of the relations that assign weights or strengths to the connections specified by the relations.*

In the context of message passing, a simplicial complex $\mathcal{K}$ can be viewed as a relational structure $\mathcal{R}(\mathcal{K}) = (\mathcal{S}, R_1, \ldots, R_5)$, where $\mathcal{S} = \mathcal{K}$ are the entities, and R_i are relations defined as follows: $R_1 = \{(\sigma) : \sigma \in \mathcal{K}\}$ (identity), $R_2 = \{(\sigma, \tau) : \sigma \in \mathcal{K}, \tau \in \mathcal{B}(\sigma)\}$ (boundary), $R_3 = \{(\sigma, \tau) : \sigma \in \mathcal{K}\tau \in \mathcal{C}(\sigma)\}$ (co-boundary), $R_4 = \{(\sigma, \tau, \delta) : \sigma \in \mathcal{K}, \tau \in \mathcal{N}_\downarrow(\sigma), \delta = \sigma \cap \tau\}$ (lower adjacency), $R_5 = \{(\sigma, \tau, \delta) : \sigma \in \mathcal{K}, \tau \in \mathcal{N}_\uparrow(\sigma), \delta = \sigma \cup \tau\}$ (upper adjacency). The message functions $\boldsymbol{\psi}^{(t)}_i$ correspond to $M_\mathcal{B}, M_\mathcal{C}, M_\downarrow, M_\uparrow$, the update function $\phi^{(t)}$ to UPDATE, and aggregation uses shift operators $\mathbf{A}^{R_i}$. This establishes an equivalence between message passing on the simplicial complex $\mathcal{K}$ and the relational structure $\mathcal{R}(\mathcal{K})$.

Remark 2.7. *The relational message passing scheme in Definition 2.5 encompasses relational graph neural networks (Schlichtkrull et al., 2018), simplicial neural networks (Bodnar et al., 2021b), higher-order graph neural networks (Morris et al., 2019), and cellular complex neural networks (Bodnar et al., 2021a). We demonstrate how higher-order graphs fit the relational framework in Appendix G. The relational structure point of view aligns with, and extends, a point of view where simplicial structures are similar objects are studied via their augmented Hasse diagrams. Refer to, for example, Hajij et al. (2023), Eitan et al. (2024), and Papillon et al. (2024).*

> **Takeaway Message 1 (Axiomatic)**
>
> Simplicial complexes can be represented as *relational structures*, where the entities are simplices, and the relations capture the adjacency among simplices of different dimensions. Simplicial message passing is an instance of relational message passing on these structures.

3 Oversquashing in Relational Message-Passing

The existing literature on oversquashing in GNNs does not directly address relational message passing. In this section, we address that gap by deriving new extensions of key results on oversquashing in GNNs to relational message passing. We illustrate the definitions in this section with a small worked example in Appendx H.

In our analysis of relational structures and message passing schemes, we naturally encounter matrices and graphs that capture the aggregated influence of the underlying shift operators. For convenience, we introduce notation for these matrices and graphs. For each relation R_i of arity n_i with

shift operator $\mathbf{A}^{R_i}$, we define the matrix $\tilde{\mathbf{A}}^{R_i} \in \mathbb{R}_{\geq 0}^{|\mathcal{S}| \times |\mathcal{S}|}$ as:

$$\tilde{\mathbf{A}}_{\sigma,\tau}^{R_i} = \sum_{j=1}^{n_i-1} \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-2}} \mathbf{A}_{\sigma,\xi_1,\ldots,\xi_{j-1},\tau,\xi_j,\ldots,\xi_{n_i-2}}^{R_i}, \quad \sigma, \tau \in \mathcal{S}. \tag{5}$$

This matrix captures all possible ways an entity τ can influence entity σ via the relation R_i. Specifically, it sums over all positions j where τ can appear among the arguments of the shift operator $\mathbf{A}^{R_i}$, and over all possible combinations of the other entities $\boldsymbol{\xi}$.

We aggregate these matrices over all relations to form the *aggregated influence matrix* $\tilde{\mathbf{A}} \in \mathbb{R}_{\geq 0}^{|\mathcal{S}| \times |\mathcal{S}|}$:

$$\tilde{\mathbf{A}} = \sum_{i=1}^{k} \tilde{\mathbf{A}}^{R_i}. \tag{6}$$

Next, we define the *augmented influence matrix* $\mathbf{B}$, which plays the role of an augmented adjacency matrix in our analysis:

$$\mathbf{B} = \gamma \mathbf{I} + \tilde{\mathbf{A}}, \tag{7}$$

where $\gamma = \max_\sigma \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-1}} \tilde{\mathbf{A}}_{\sigma,\boldsymbol{\xi}}$ is the maximum row sum of $\tilde{\mathbf{A}}$.

Lastly, we introduce graphs that capture the aggregated message passing dynamical structure implied by the relational structure and the message passing scheme.

Definition 3.1 (Influence Graph). *Given a relational structure $\mathcal{R} = (\mathcal{S}, R_1, \ldots, R_k)$ and a relational message passing scheme with update rule given by Equation 3, and given $\mathbf{Q} \in \{\tilde{\mathbf{A}}, \mathbf{B}\}$, where $\tilde{\mathbf{A}}$ and $\mathbf{B}$ are defined by Equations 6 and 7 respectively, we define the* influence graph *$\mathcal{G}(\mathcal{S}, \mathbf{Q}) = (\mathcal{S}, \mathcal{E}, w)$ as follows: The set of entities (i.e., nodes) is $\mathcal{S}$. The set of edges $\mathcal{E}$ consists of directed edges from entity τ to entity σ for each pair $(\sigma, \tau) \in \mathcal{S} \times \mathcal{S}$ with $\mathbf{Q}_{\sigma,\tau} > 0$. Each edge from τ to σ is assigned a weight $w_{\tau \to \sigma} = \mathbf{Q}_{\sigma,\tau}$.*

As we will see next, these graphs make it possible to leverage and extend graph-theoretic concepts, results, and intuition to understand and analyze the behavior of relational message passing schemes.

3.1 SENSITIVITY ANALYSIS

We now analyze the sensitivity of relational message passing to changes in the input features. This analysis is crucial for understanding how information propagates through the network and for identifying potential bottlenecks or oversquashing effects. We begin with a standard assumption about the boundedness of the Jacobians of the message and update functions.

Assumption 1 (Bounded Jacobians). *All message functions $\psi_i^{(\ell)}$ and update functions $\phi^{(\ell)}$ are differentiable with bounded Jacobians: There exist constants $\beta_i^{(\ell)}$ and $\alpha^{(\ell)}$ such that $\|\partial\psi_i^{(\ell)}/\partial\mathbf{h}_\sigma\|_1 \leq \beta_i^{(\ell)}$ for any input feature vector $\mathbf{h}_\sigma$, and $\|\partial\phi^{(\ell)}/\partial\mathbf{m}_j\|_1 \leq \alpha^{(\ell)}$ for any message input $\mathbf{m}_j$. We write $\beta^{(\ell)} = \max_i \beta_i^{(\ell)}$.*

Our main result on sensitivity is the following, which is a novel extension of GNN sensitivity analysis results (e.g., Topping et al., 2022, Lemma 1 and Di Giovanni et al., 2023, Theorem 3.2) to relational (and topological) message passing. We provide the proof in Appendix C.1.

Lemma 3.2 (Sensitivity Bound for Relational Message Passing). *Consider a relational structure $\mathcal{R} = (\mathcal{S}, R_1, \ldots, R_k)$ with update rule given by Equation 3 and satisfying Assumption 1. Then, for any $\sigma, \tau \in \mathcal{S}$ and $t > 0$, the Jacobian at layer t with respect to the input features ($t = 0$) satisfies*

$$\left\|\frac{\partial\mathbf{h}_\sigma^{(t)}}{\partial\mathbf{h}_\tau^{(0)}}\right\|_1 \leq \left(\prod_{\ell=0}^{t-1} \alpha^{(\ell)}\beta^{(\ell)}\right)\left(\mathbf{B}^t\right)_{\sigma,\tau}. \tag{8}$$

Thus, the bound on the Jacobian of the σ-feature with respect to the input τ-feature depends on the (σ, τ)-entry of the t-th matrix power $\mathbf{B}^t$, which reflects the number and strength of t-length paths from τ to σ in the graph $\mathcal{G}(\mathcal{S}, \mathbf{B})$. Structural properties of $\mathcal{G}(\mathcal{S}, \mathbf{B})$ that lead to small values of $\left(\mathbf{B}^t\right)_{\sigma,\tau}$, such as bottlenecks or long distances between nodes, therefore contribute to the phenomenon of oversquashing, where the influence of distant entities is diminished.

As demonstrated throughout this work, our result offers a systematic framework for extending other theoretical findings on oversquashing in graphs, which do not directly apply to simplicial complexes and similar relational structures. This includes the influential works by Topping et al. (2022), Di Giovanni et al. (2023), and Fesser & Weber (2023). By leveraging our axiomatic framework, we derive principled extensions on the impact of local geometry (Section 3.2), depth (Section 3.3), and hidden dimensions (Section 3.4) in higher-order message passing, addressing settings where prior results are not applicable. Additionally, this result offers a clear approach for deriving analogs of key quantities such as curvature (Definition 3.3), and can serve as a guide for future work.

3.2 THE IMPACT OF LOCAL GEOMETRY

Lemma 3.2 shows that the entries of the matrix $\mathbf{B}^t$, which encode the number and strength of connections in a relational message passing scheme, control feature sensitivity. Prior works relate similar bounds to notions of discrete curvature for unweighted undirected graphs, such as balanced Forman curvature (Topping et al., 2022), Ollivier-Ricci curvature (Nguyen et al., 2023), and augmented Forman curvature (Fesser & Weber, 2023), via counting local motifs, such as triangles and squares, in the underlying graphs. Following this approach, we derive a result analogous to Fesser & Weber (2023, Proposition 3.4), introducing a motif-counting quantity inspired by the augmented Forman curvature, but adapted for the particular weighted directed graphs arising in our setting.

Definition 3.3. *Let $\mathcal{G} = (\mathcal{S}, \mathcal{E}, w)$ be a weighted directed graph with entities (nodes) $\mathcal{S}$, edges $\mathcal{E}$, and edge weights $w \colon \mathcal{E} \to \mathbb{R}_{\geq 0}$. For each entity $\tau \in \mathcal{S}$, define the* weighted out-degree $w_\tau^{\mathrm{out}} = \sum_{(\tau \to \sigma) \in \mathcal{E}} w_{\tau \to \sigma}$ *and the* weighted in-degree $w_\tau^{\mathrm{in}} = \sum_{(\sigma \to \tau) \in \mathcal{E}} w_{\sigma \to \tau}$. *For an edge $(\tau \to \sigma) \in \mathcal{E}$, define the* weighted triangle count $w_T = \sum_{\xi \in \mathcal{S}} w_{\tau \to \xi} \cdot w_{\xi \to \sigma}$ *and the* weighted quadrangle count $w_F = \sum_{\xi_1, \xi_2 \in \mathcal{S}} w_{\tau \to \xi_1} \cdot w_{\xi_1 \to \xi_2} \cdot w_{\xi_2 \to \sigma}$. *Then, the* extended Forman curvature *of the edge $(\tau \to \sigma)$ is defined as:*

$$\mathrm{EFC}_{\mathcal{G}}(\tau, \sigma) = 4 - w_\tau^{\mathrm{out}} - w_\sigma^{\mathrm{in}} + 3w_T + 2w_F. \tag{9}$$

We immediately get the following result, inspired by Nguyen et al. (2023, Proposition 4.4) and Fesser & Weber (2023, Proposition 3.4), and which is exemplary of results connecting sensitivity analysis to notions of discrete curvature. We provide the proof in Appendix C.2.

Proposition 3.4. *Consider a relational structure $\mathcal{R} = (\mathcal{S}, R_1, \ldots, R_k)$ with update rule given by Equation 3 and satisfying Assumption 1. Denote $\mathcal{G} = \mathcal{G}(\mathcal{S}, \mathbf{B})$. Then, for any $\sigma, \tau \in \mathcal{S}$ with an edge $(\tau \to \sigma) \in \mathcal{G}$, the following holds:*

$$\left\| \frac{\partial \mathbf{h}_\sigma^{(2)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \leq \frac{1}{3} \left(\prod_{\ell=0}^{1} \alpha^{(\ell)} \beta^{(\ell)} \right) \left[\mathrm{EFC}_{\mathcal{G}}(\tau, \sigma) + w_\tau^{\mathrm{out}} + w_\sigma^{\mathrm{in}} - 4 \right]. \tag{10}$$

In principle, a similar result using balanced Forman curvature (as in Topping et al., 2022, Theorem 4) is possible using our framework, and we leave that extension for future work. Connections to Ollivier-Ricci curvature are discussed in Appendix B.1.

We present experimental analyses related to curvature on relational structures in Section D.3 (edge curvature distribution) and Appendices D.1 (edge curvature visualization) and D.2 (weighted curvature). We further propose a relational extension of curvature-based rewiring techniques in Section 4 and empirically analyze the impact of relational rewiring using real-world and synthetic benchmarks in Sections 5.1 and 5.2, respectively.

3.3 THE IMPACT OF DEPTH

To facilitate our analysis of depth, we make the following non-restrictive assumptions:

Assumption 2 (Row-Normalized Shift Operators)**.** *Each shift operator $\mathbf{A}^{R_i}$ associated with relation R_i is row-normalized, such that for all $\sigma \in \mathcal{S}$,*

$$\sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i - 1}} A_{\sigma, \boldsymbol{\xi}}^{R_i} = \begin{cases} 1, & \text{if } \sum_{\boldsymbol{\xi}} A_{\sigma, \boldsymbol{\xi}}^{R_i} \neq 0, \\ 0, & \text{if } \sum_{\boldsymbol{\xi}} A_{\sigma, \boldsymbol{\xi}}^{R_i} = 0. \end{cases} \tag{11}$$

Assumption 3 (Bounded $\alpha^{(\ell)}$ and $\beta^{(\ell)}$)**.** *There exist constants $\alpha_{\max} > 0$ and $\beta_{\max} > 0$ such that for all layers ℓ, $\alpha^{(\ell)} \leq \alpha_{\max}$ and $\beta^{(\ell)} \leq \beta_{\max}$.*

We now present our main result on the impact of depth in relational message passing, extending a previous result by Di Giovanni et al. (2023, Theorem 4.1) to our setting. We provide the proof in Appendix C.3. By *the combinatorial distance* from τ to σ in the graph $\mathcal{G}(\mathcal{S}, \tilde{\mathbf{A}})$, we mean the smallest number of edges in a directed path from τ to σ in the graph. Similarly, by *combinatorial length* of a directed path, we mean the number of edges in the path.

Theorem 3.5 (Impact of Depth on Relational Message Passing). *Consider a relational structure $\mathcal{R} = (\mathcal{S}, R_1, \ldots, R_k)$ with update rule given by Equation 3 and satisfying Assumptions 1, 2, and 3. Let $\sigma, \tau \in \mathcal{S}$ be entities such that the combinatorial distance from τ to σ in the graph $\mathcal{G}(\mathcal{S}, \tilde{\mathbf{A}})$ is r. Denote by $\omega_\ell(\sigma, \tau)$ the number of directed paths from τ to σ of combinatorial length at most ℓ in $\mathcal{G}(\mathcal{S}, \tilde{\mathbf{A}})$. Then, for any $0 \leq m < r$, there exists a constant $C > 0$, depending only on $\alpha_{\max}$, $\beta_{\max}$, k, and m, but not on r nor the specific relations in $\mathcal{R}$, such that*

$$\left\| \frac{\partial \mathbf{h}_\sigma^{(r+m)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \leq C \omega_{r+m}(\sigma, \tau)(2\alpha_{\max}\beta_{\max}M)^r,$$

where $M = \max_{\sigma, \tau} \tilde{\mathbf{A}}_{\sigma, \tau}$.

This result indicates that the sensitivity can decay exponentially with depth when $M < 1/(2\alpha_{\max}\beta_{\max})$, particularly when the number of walks $\omega_t(\sigma, \tau)$ is limited by the structure of $\mathcal{G}(\mathcal{S}, \tilde{\mathbf{A}})$. Such exponential decay is a characteristic of the oversquashing phenomenon, where information from distant nodes becomes increasingly compressed, reducing its influence on the output.

We present experimental validation of this result in Section 5.2.

3.4 THE IMPACT OF HIDDEN DIMENSIONS

In situations where the Lipschitz constants of the message and update functions from Assumption 1 are affected by hyperparameters, such as the widths of neural networks implementing said functions, one can have $\beta_i^{(\ell)} = O(p_{i,\ell})$ and $\alpha^{(\ell)} = O(p_{\ell+1})$. This is the case, for instance, when the message and update functions are shallow neural networks (see Appendix B.2 and Appendix C.4).

Writing $p_\ell' = \max_i p_{i,\ell}$ and substituting $\beta^{(\ell)} = O(p_\ell')$ and $\alpha^{(\ell)} = O(p_{\ell+1})$ into the bound from Lemma 3.2, one gets:

$$\left\| \frac{\partial \mathbf{h}_\sigma^{(t)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \leq C \cdot \left(\prod_{\ell=0}^{t-1} p_{\ell+1} \cdot p_\ell' \right) \left(\mathbf{B}^t \right)_{\sigma, \tau}, \tag{12}$$

where C is a constant independent of the layer widths, p_ℓ' is the maximum dimension of the message vectors at layer ℓ, and $p_{\ell+1}$ is the output dimension of the update function at layer ℓ.

This implies that low hidden dimensions in the message and update functions contribute to a low sensitivity upper bound, which can exacerbate the oversquashing problem. Increasing the hidden dimensions will raise the upper bound, which can help improve the model's ability to propagate information effectively and enhance performance on tasks. However, increasing the hidden dimensions risks overfitting due to the increased model complexity (Bartlett et al., 2017).

We present experimental validation of this result in Section 5.2.

> **Takeaway Message 2 (Theoretical)**
>
> By reformulating higher order structures as relational structures, key results on oversquashing in graph neural networks extend to relational message passing schemes through the *aggregated influence matrix* and the *influence graph*. This conceptual framework enables analysis of the impact of local geometry, depth, and hidden dimensions in relational message passing schemes, just as in graph neural networks.

4 REWIRING HEURISTICS FOR RELATIONAL STRUCTURES

Inspired by First-Order Spectral Rewiring (FoSR) (Karhadkar et al., 2023), we propose a rewiring heuristic that integrates additional connections into a relational structure without altering its original

connections. To capture the overall connectivity of a relational structure, we define the collapsed adjacency matrix, which counts the number of direct connections between entities.

Definition 4.1 (Collapsed Adjacency Matrix). *Given a relational structure $\mathcal{R} = (\mathcal{S}, R_1, \ldots, R_k)$, the* collapsed adjacency matrix $\mathbf{A}^{\text{col}}$ *for the structure $\mathcal{R}$ is defined by:*

$$\mathbf{A}^{\text{col}}_{\sigma, \tau} = \sum_{i=1}^{k} \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i - 1}} \mathbf{1}_{\{(\sigma, \boldsymbol{\xi}) \in R_i, \tau = \xi_j \text{ for some } j \in \{1, \ldots, n_i\}\}}, \quad \sigma, \tau \in \mathcal{S}.$$

This matrix captures direct connections between entities through any relation, effectively collapsing the relational structure into a graph. Our proposed relational rewiring algorithm is as follows.

Algorithm 1 Relational Rewiring Algorithm

Require: Relational structure $\mathcal{R} = (\mathcal{S}, R_1, \ldots, R_k)$; graph rewiring algorithm REWIREALGO
 1: Construct the collapsed adjacency matrix $\mathbf{A}^{\text{col}}$ (Definition 4.1)
 2: Build the graph $\mathcal{G}^{\text{col}} = (\mathcal{S}, \mathbf{A}^{\text{col}})$
 3: Apply REWIREALGO to $\mathcal{G}^{\text{col}}$ to obtain additional edges E_{new}
 4: Define a new relation $R_{k+1} = E_{\text{new}}$
 5: Update the relational structure: $\mathcal{R}' = (\mathcal{S}, R_1, \ldots, R_k, R_{k+1})$

Adding new connections (E_{new}) without removing existing ones improves the model capacity to capture long-range dependencies while preserving the original relational structure. We experimentally analyze the impact of relational rewiring using real-world and synthetic benchmarks in Sections 5.1 and 5.2, respectively. Future work could explore rewiring algorithms that remove or reclassify edges, such as spectral pruning (Jamadandi et al., 2024), either by re-labeling the edges as new relations or by deletion.

> **Takeaway Message 3 (Practical)**
>
> Graph rewiring techniques for improving information flow and mitigating oversquashing can be adapted to relational structures. This approach improves long-range connectivity and enhances message propagation while maintaining the integrity of the original connections.

5 EXPERIMENTS AND RESULTS

5.1 REAL-WORLD BENCHMARK: GRAPH CLASSIFICATION

We run an empirical analysis with real-world datasets to compare the performance of different graph, relational graph, and simplicial message passing models, and the impact of relational rewiring on said models. We provide more details in Appendix E.

Task and Datasets. We use graph classification tasks ENZYMES, IMDB-B, MUTAG, NCI1, and PROTEINS from the TUDataset (Morris et al., 2020) for evaluation.

Graph Lifting. For the simplicial message passing models, graphs are separately treated as simplicial complexes with only graph nodes (0-dimensional simplices) and upper adjacencies, and also lifted into clique complexes (see appendix).

Models. We evaluate three types of models: a) *Graph message passing models*: SGC (Wu et al., 2019), GCN (Kipf & Welling, 2017), and GIN (Xu et al., 2019); b) *Relational graph message passing models*: RGCN (Schlichtkrull et al., 2018) and RGIN; c) *Topological message passing models*: SIN (Bodnar et al., 2021b), CIN (Bodnar et al., 2021a), and CIN++ (Giusti et al., 2024).

Relational Rewiring. We apply relational rewiring for 40 iterations (Section 4) using three choices for REWIREALGO: SDRF (Topping et al., 2022), FoSR (Karhadkar et al., 2023), and AFRC (Fesser & Weber, 2023). Due to computational constraints, we run the three choices on all datasets except IMDB with Clique graph lifting, where we only run FoSR. We use fixed, dataset- and model-agnostic hyperparameters, diverging from prior work where hyperparameter sweeps are carried out. It is

important to note that hyperparameter tuning can significantly impact performance on downstream tasks, as highlighted by, e.g., Tori et al. (2024).

Training. Models are trained for up to 500 epochs, with early stopping and learning rate decay based on a validation set. Additional details can be found in Appendix E. Results are reported as *mean ± standard error* over 10 trials.

Results. Table 1 shows test accuracies for the TUDataset experiments. Rewiring generally boosts performance for base graphs across models and datasets, particularly for RGCN, which improves significantly on all datasets. GCN and GIN also see notable gains on MUTAG and ENZYMES, respectively. For clique complex lifted graphs, results are similar. CIN++ performs exceptionally well after rewiring on NCI1 and ENZYMES. However, models like GIN and RGIN show modest improvements or slight decreases on certain datasets. Relational models on lifted graphs (RGCN, RGIN) perform similarly to topological models (SIN, CIN) and sometimes outperform them, such as RGCN on PROTEINS. Additionally, for ENZYMES, MUTAG, NCI1, and PROTEINS, in 75% of cases, the best-performing rewiring algorithm for Lift=None and Lift=Clique are the same. (Cf. Appendix E.3.) In conclusion, the impact of rewiring with our dataset- and model-agnostic choice of hyperparameters varies across datasets, with relational and topological models performance responding to rewiring similarly to graph models.

Lift	Model	ENZYMES No Rew.	ENZYMES Best Rew.	IMDB-B No Rew.	IMDB-B FoSR	MUTAG No Rew.	MUTAG Best Rew.	NCI1 No Rew.	NCI1 Best Rew.	PROTEINS No Rew.	PROTEINS Best Rew.
None	SGC	18.3 ± 1.2	21.5 ± 1.6	49.5 ± 1.5	50.0 ± 1.8	64.5 ± 5.8	70.0 ± 2.6	55.2 ± 1.0	54.4 ± 0.7	62.2 ± 1.4	65.0 ± 1.5
	GCN	32.2 ± 2.0	30.7 ± 1.5	49.1 ± 1.4	47.9 ± 1.0	71.0 ± 3.8	83.0 ± 1.5	48.3 ± 0.6	49.1 ± 0.8	73.1 ± 1.2	75.8 ± 1.7
	GIN	47.2 ± 1.9	50.0 ± 2.7	71.7 ± 1.5	67.1 ± 1.5	83.0 ± 3.1	88.0 ± 2.4	77.2 ± 0.5	77.9 ± 0.5	70.6 ± 1.4	72.2 ± 0.8
	RGCN	33.8 ± 1.6	42.5 ± 1.3	47.6 ± 1.4	68.0 ± 1.3	72.5 ± 2.5	83.5 ± 1.8	53.2 ± 0.7	63.4 ± 1.1	71.9 ± 1.6	75.6 ± 1.2
	RGIN	46.8 ± 1.8	49.8 ± 2.0	69.6 ± 1.6	48.9 ± 2.9	81.5 ± 1.7	85.5 ± 2.0	76.8 ± 1.1	77.0 ± 0.7	70.8 ± 1.2	72.4 ± 1.4
	SIN	47.5 ± 2.3	46.8 ± 2.1	70.0 ± 1.4	63.0 ± 2.7	88.5 ± 3.0	85.5 ± 1.7	77.0 ± 0.6	76.4 ± 0.4	70.2 ± 1.3	73.2 ± 1.5
	CIN	50.0 ± 1.9	49.0 ± 2.0	58.1 ± 4.0	58.4 ± 2.7	86.5 ± 1.8	87.0 ± 2.4	51.4 ± 2.5	66.2 ± 2.0	70.7 ± 1.0	71.0 ± 1.4
	CIN++	48.5 ± 1.9	51.0 ± 1.5	66.6 ± 3.7	56.0 ± 3.9	85.0 ± 3.4	91.0 ± 2.3	60.8 ± 3.8	64.8 ± 3.1	67.9 ± 1.9	71.4 ± 1.4
Clique	SGC	14.5 ± 1.4	16.8 ± 0.9	48.7 ± 2.2	47.8 ± 1.6	70.0 ± 3.3	69.5 ± 2.6	50.0 ± 1.3	56.8 ± 0.8	59.9 ± 1.8	59.1 ± 1.4
	GCN	30.7 ± 1.2	30.2 ± 2.4	64.0 ± 3.1	65.5 ± 3.1	67.0 ± 3.5	81.5 ± 2.9	48.4 ± 0.4	49.6 ± 0.6	69.9 ± 0.6	75.0 ± 1.4
	GIN	44.0 ± 1.7	48.5 ± 2.2	69.1 ± 1.2	70.8 ± 1.1	83.0 ± 2.8	82.5 ± 2.6	78.8 ± 0.7	78.2 ± 0.6	68.7 ± 1.4	72.8 ± 1.2
	RGCN	48.8 ± 1.2	45.2 ± 1.5	71.0 ± 1.0	69.7 ± 1.5	79.5 ± 1.7	81.5 ± 3.8	72.9 ± 0.8	75.0 ± 0.9	72.4 ± 1.6	74.2 ± 1.2
	RGIN	50.8 ± 1.5	55.8 ± 2.5	71.6 ± 0.9	69.0 ± 1.4	86.0 ± 2.3	85.0 ± 2.4	79.2 ± 0.6	79.5 ± 0.4	71.5 ± 1.5	71.8 ± 1.7
	SIN	51.0 ± 2.4	46.5 ± 1.2	53.0 ± 1.9	64.0 ± 2.3	87.0 ± 3.2	83.5 ± 1.7	76.6 ± 1.3	75.4 ± 0.7	66.9 ± 1.3	70.4 ± 1.2
	CIN	49.8 ± 1.9	46.7 ± 1.3	52.6 ± 2.4	68.1 ± 1.6	85.5 ± 2.8	86.5 ± 2.6	51.8 ± 2.3	72.5 ± 0.8	70.7 ± 1.2	70.3 ± 0.8
	CIN++	50.5 ± 2.1	52.7 ± 1.6	62.8 ± 3.8	64.7 ± 1.5	90.5 ± 2.2	84.5 ± 3.3	61.5 ± 4.6	76.8 ± 0.4	68.3 ± 1.3	71.9 ± 1.0

Table 1: Test accuracy for TUDataset experiments. Each value is presented as the mean ± standard error across ten trials. The best-performing result for each dataset is highlighted in gold, while the second-best is in silver. The results after rewiring are shown with green text if the mean increased and red text if the mean decreased.

5.2 SYNTHETIC BENCHMARK: RINGTRANSFER

We confirm the theoretical results from Section 3 using the RINGTRANSFER benchmark, a graph feature transfer task designed to tease out the effect of long-range dependencies in message-passing models using rings of growing size. We follow the experimental setup of Karhadkar et al. (2023) and Di Giovanni et al. (2023), and provide more details in Appendix E.2. We test the impact of neural network hidden dimensions (Section 3.4), relational structure depth (Section 3.3), and relational structure local geometry (Sections 3.4 and 4) on task performance by varying the hidden dimensions, ring sizes, and rewiring iterations. The results, consistent with the theory, demonstrate that increasing network hidden dimensions improves performance up to a point, after which it declines, potentially due to overfitting. Larger ring sizes lead to performance deterioration, as the effects of long range dependencies and bottlenecks start to take over. At the same time, rewiring improves performance by facilitating communication between distant nodes and mitigating oversquashing. As illustrated in Figure 2, message passing on graphs and simplicial complexes demonstrate similar trends, consistent with our theoretical predictions.

5.3 ADDITIONAL EXPERIMENTS AND ANALYSES

We report additional analyses in Appendix D.1 and Appendix D.2. There, we visualize the curvature of relational structures for dumbbell graphs and their corresponding clique complexes. We also

(a) Hidden Dimensions   (b) Size   (c) Rewiring

Figure 2: Performance on RINGTRANSFER obtained by varying model hidden dimensions (left), ring size (middle), and number of rewiring iterations (right).

reports a statistically significant linear relationship between the weighted curvature of graphs and their lifted clique complexes. These interesting patterns merit further investigation.

6 DISCUSSION AND CONCLUSIONS

This work addresses pressing questions about oversquashing in topological networks and higher-order generalizations of rewiring algorithms raised by the TDL community (Questions 2 and 9 of Papamarkou et al., 2024). We introduce a theoretical framework for unifying graph and topological message passing via relational structures, extending key graph-theoretic results on oversquashing and rewiring strategies to higher-order networks such as simplicial complexes via *influence graphs* that capture the aggregated message passing dynamical structure on relational structures. Our approach applies broadly to other message-passing schemes, including relational GNNs, high-order GNNs, and CW networks, providing a foundation for future theoretical and empirical research. Empirical results on real-world datasets show that simplicial networks respond to rewiring similarly to graph networks, and synthetic benchmarks further confirm our theoretical findings.

Certain aspects are worthy of further investigation. In particular, we compare message passing on graphs and their clique complexes through proxies (e.g., performance on tasks), as the significant differences in size and structure make direct empirical comparisons, e.g., of curvature, less theoretically rigorous. While we observe statistically significant patterns when comparing weighted curvatures, further theoretical and empirical investigation is needed. Furthermore, the rewiring algorithms we applied our relational rewiring heuristic to were not originally designed with weighted directed influence graphs in mind. Potentially, further improvements could be obtained by implementing algorithms specifically tailored for rewiring weighted directed graphs.

For future work, exploring global geometric properties of relational structures, studying oversmoothing, and empirically analyzing more relational message-passing schemes are promising directions. Developing theoretical tools and tailored rewiring heuristics for weighted directed graphs will be crucial, as will be tools for direct comparisons of message-passing across different relational structures. Furthermore, systematically and empirically assessing our framework's higher-order extensions of more state-of-the-art (SoTA) graph rewiring solutions is essential; Appendix A lists some recent graph rewiring works, which could serve as potential candidates for future work in that direction. By unifying topological message passing into message passing on relational structures and generalizing graph-based analysis to this setting, we hope that the present work can aid in both rigorous analysis and direct comparison between different higher order message passing schemes.

Lastly, for practitioners, we recommend topological message passing as yet another relational learning tool with relational rewiring as a preprocessing step.

Reproducibility Statement The implementation and scripts to replicate our experiments are available at https://anonymous.4open.science/r/ Simplicial-Oversquashing-C4CD/. Experimental settings and implementation details are described in Section 5, and Appendices E and F.

REFERENCES

Uri Alon and Eran Yahav. On the bottleneck of graph neural networks and its practical implications. In *International Conference on Learning Representations*, 2021. URL https://openreview.net/forum?id=i80OPhOCVH2.

Alessia Antelmi, Gennaro Cordasco, Mirko Polato, Vittorio Scarano, Carmine Spagnuolo, and Dingqi Yang. A survey on hypergraph representation learning. *ACM Computing Surveys*, 56 (1):1–38, 2023.

Hugo Attali, Davide Buscaldi, and Nathalie Pernelle. Delaunay graph: Addressing over-squashing and over-smoothing using delaunay triangulation. In *Forty-first International Conference on Machine Learning*.

Pradeep Kr. Banerjee, Kedar Karhadkar, Yu Guang Wang, Uri Alon, and Guido Montúfar. Over-squashing in GNNs through the lens of information contraction and graph expansion. In *2022 58th Annual Allerton Conference on Communication, Control, and Computing (Allerton)*, pp. 1–8, 2022. doi: 10.1109/Allerton49937.2022.9929363.

Federico Barbero, Cristian Bodnar, Haitz Sáez de Ocáriz Borde, and Pietro Lio. Sheaf attention networks. In *NeurIPS 2022 Workshop on Symmetry and Geometry in Neural Representations*, 2022. URL https://openreview.net/forum?id=LIDvgVjpkZr.

Peter L Bartlett, Dylan J Foster, and Matus J Telgarsky. Spectrally-normalized margin bounds for neural networks. In *Advances in Neural Information Processing Systems*, volume 30. Curran Associates, Inc., 2017. URL https://proceedings.neurips.cc/paper_files/paper/2017/file/b22b257ad0519d4500539da3c8bcf4dd-Paper.pdf.

Mitchell Black, Zhengchao Wan, Amir Nayyeri, and Yusu Wang. Understanding oversquashing in GNNs through the lens of effective resistance. In *Proceedings of the 40th International Conference on Machine Learning*, volume 202 of *Proceedings of Machine Learning Research*, pp. 2528–2547. PMLR, 2023. URL https://proceedings.mlr.press/v202/black23a.html.

Cristian Bodnar, Fabrizio Frasca, Nina Otter, Yu Guang Wang, Pietro Liò, Guido Montúfar, and Michael M. Bronstein. Weisfeiler and Lehman go cellular: CW networks. In *Advances in Neural Information Processing Systems*, 2021a. URL https://openreview.net/forum?id=uVPZCMVtsSG.

Cristian Bodnar, Fabrizio Frasca, Yuguang Wang, Nina Otter, Guido Montúfar, Pietro Lió, and Michael Bronstein. Weisfeiler and Lehman go topological: Message passing simplicial networks. In *Proceedings of the 38th International Conference on Machine Learning*, volume 139 of *Proceedings of Machine Learning Research*, pp. 1026–1037. PMLR, 2021b. URL https://proceedings.mlr.press/v139/bodnar21a.html.

Cristian Bodnar, Francesco Di Giovanni, Benjamin Paul Chamberlain, Pietro Lio, and Michael M. Bronstein. Neural sheaf diffusion: A topological perspective on heterophily and oversmoothing in GNNs. In Alice H. Oh, Alekh Agarwal, Danielle Belgrave, and Kyunghyun Cho (eds.), *Advances in Neural Information Processing Systems*, 2022. URL https://openreview.net/forum?id=vbPsD-BhOZ.

Michael M. Bronstein, Joan Bruna, Taco Cohen, and Petar Veličković. Geometric deep learning: Grids, groups, graphs, geodesics, and gauges, 2021. URL https://arxiv.org/abs/2104.13478.

Benjamin Chamberlain, James Rowbottom, Davide Eynard, Francesco Di Giovanni, Xiaowen Dong, and Michael Bronstein. Beltrami flow and neural diffusion on graphs. In *Advances in Neural Information Processing Systems*, volume 34, pp. 1594–1609. Curran Associates, Inc., 2021. URL https://proceedings.neurips.cc/paper_files/paper/2021/file/0cbed40c0d920b94126eaf5e707be1f5-Paper.pdf.

Zhengdao Chen, Soledad Villar, Lei Chen, and Joan Bruna. On the equivalence between graph isomorphism testing and function approximation with GNNs. In *Advances in Neural Information Processing Systems*, volume 32. Curran Associates, Inc., 2019. URL https://proceedings.neurips.cc/paper_files/paper/2019/file/71ee911dd06428a96c143a0b135041a4-Paper.pdf.

George Dasoulas, Johannes F. Lutzeyer, and Michalis Vazirgiannis. Learning parametrised graph shift operators. In *International Conference on Learning Representations*, 2021. URL https://openreview.net/forum?id=0OlrLvrsHwQ.

Francesco Di Giovanni, Lorenzo Giusti, Federico Barbero, Giulia Luise, Pietro Liò, and Michael M. Bronstein. On over-squashing in message passing neural networks: The impact of width, depth, and topology. In *Proceedings of the 40th International Conference on Machine Learning*, volume 202 of *Proceedings of Machine Learning Research*, pp. 7865–7885. PMLR, 2023. URL https://proceedings.mlr.press/v202/di-giovanni23a.html.

Vijay Prakash Dwivedi and Xavier Bresson. A generalization of transformer networks to graphs. *AAAI 2021 Workshop on Deep Learning on Graphs: Methods and Applications (DLG-AAAI 2021)*, 2020. URL https://arxiv.org/abs/2012.09699.

Stefania Ebli, Michaël Defferrard, and Gard Spreemann. Simplicial neural networks. In *TDA & Beyond*, 2020. URL https://openreview.net/forum?id=nPCt39DVIfk.

Marzieh Eidi and Jürgen Jost. Ollivier Ricci curvature of directed hypergraphs. *Scientific Reports*, 10(1):12466, 2020. URL https://doi.org/10.1038/s41598-020-68619-6.

Yam Eitan, Yoav Gelberg, Guy Bar-Shalom, Fabrizio Frasca, Michael Bronstein, and Haggai Maron. Topological blind spots: Understanding and extending topological deep learning through the lens of expressivity. *arXiv preprint arXiv:2408.05486*, 2024.

Bahare Fatemi, Perouz Taslakian, David Vazquez, and David Poole. Knowledge hypergraph embedding meets relational algebra. *Journal of Machine Learning Research*, 24(105):1–34, 2023.

Lukas Fesser and Melanie Weber. Mitigating over-smoothing and over-squashing using augmentations of Forman-Ricci curvature. In *The Second Learning on Graphs Conference*, 2023. URL https://openreview.net/forum?id=bKTkZMRtfC.

Lukas Fesser and Melanie Weber. Mitigating over-smoothing and over-squashing using augmentations of forman-ricci curvature, 2024. URL https://arxiv.org/abs/2309.09384.

Matthias Fey and Jan Eric Lenssen. Fast graph representation learning with pytorch geometric. *arXiv preprint arXiv:1903.02428*, 2019.

Rémi Flamary, Nicolas Courty, Alexandre Gramfort, Mokhtar Z. Alaya, Aurélie Boisbunon, Stanislas Chambon, Laetitia Chapel, Adrien Corenflos, Kilian Fatras, Nemo Fournier, Léo Gautheron, Nathalie T.H. Gayraud, Hicham Janati, Alain Rakotomamonjy, Ievgen Redko, Antoine Rolet, Antony Schutz, Vivien Seguy, Danica J. Sutherland, Romain Tavenard, Alexander Tong, and Titouan Vayer. Pot: Python optimal transport. *Journal of Machine Learning Research*, 22(78): 1–8, 2021. URL http://jmlr.org/papers/v22/20-451.html.

Fernando Gama, Alejandro Ribeiro, and Joan Bruna. Stability of graph neural networks to relative perturbations. In *ICASSP 2020-2020 IEEE International Conference on Acoustics, Speech and Signal Processing (ICASSP)*, pp. 9070–9074. IEEE, 2020.

Lorenzo Giusti. Awesome topological deep learning, 2024. URL https://github.com/lrnzgiusti/awesome-topological-deep-learning. Accessed: Sep. 28, 2024.

Lorenzo Giusti, Claudio Battiloro, Lucia Testa, Paolo Di Lorenzo, Stefania Sardellitti, and Sergio Barbarossa. Cell attention networks. In *2023 International Joint Conference on Neural Networks (IJCNN)*, pp. 1–8. IEEE, 2023.

Lorenzo Giusti, Teodora Reu, Francesco Ceccarelli, Cristian Bodnar, and Pietro Liò. Topological message passing for higher - order and long - range interactions. In *2024 International Joint Conference on Neural Networks (IJCNN)*, pp. 1–8, 2024. doi: 10.1109/IJCNN60899.2024.10650343.

Christopher Wei Jin Goh, Cristian Bodnar, and Pietro Lio. Simplicial attention networks. In *ICLR 2022 Workshop on Geometrical and Topological Representation Learning*, 2022. URL `https://openreview.net/forum?id=ScfRNWkpec`.

Aric Hagberg, Pieter J Swart, and Daniel A Schult. Exploring network structure, dynamics, and function using networkx. 2008.

Mustafa Hajij, Kyle Istvan, and Ghada Zamzmi. Cell complex neural networks. In *TDA & Beyond*, 2020. URL `https://openreview.net/forum?id=6Tq18ySFpGU`.

Mustafa Hajij, Ghada Zamzmi, Theodore Papamarkou, Nina Miolane, Aldo Guzmán-Sáenz, and Karthikeyan Natesan Ramamurthy. Higher-order attention networks. *arXiv preprint arXiv:2206.00606*, 2(3):4, 2022.

Mustafa Hajij, Ghada Zamzmi, Theodore Papamarkou, Nina Miolane, Aldo Guzmán-Sáenz, Karthikeyan Natesan Ramamurthy, Tolga Birdal, Tamal K. Dey, Soham Mukherjee, Shreyas N. Samaga, Neal Livesay, Robin Walters, Paul Rosen, and Michael T. Schaub. Topological deep learning: Going beyond graph data, 2023. URL `https://arxiv.org/abs/2206.00606`.

Charles R. Harris, K. Jarrod Millman, Stéfan J. van der Walt, Ralf Gommers, Pauli Virtanen, David Cournapeau, Eric Wieser, Julian Taylor, Sebastian Berg, Nathaniel J. Smith, Robert Kern, Matti Picus, Stephan Hoyer, Marten H. van Kerkwijk, Matthew Brett, Allan Haldane, Jaime Fernández del Río, Mark Wiebe, Pearu Peterson, Pierre Gérard-Marchant, Kevin Sheppard, Tyler Reddy, Warren Weckesser, Hameer Abbasi, Christoph Gohlke, and Travis E. Oliphant. Array programming with NumPy. *Nature*, 585(7825):357–362, September 2020. doi: 10.1038/s41586-020-2649-2. URL `https://doi.org/10.1038/s41586-020-2649-2`.

Wilfrid Hodges. *Model theory*. Cambridge University Press, 1993.

Keke Huang, Yu Guang Wang, Ming Li, and Pietro Lio. How universal polynomial bases enhance spectral graph neural networks: Heterophily, over-smoothing, and over-squashing. In *Forty-first International Conference on Machine Learning*, 2024a. URL `https://openreview.net/forum?id=Z2LH6Va7L2`.

Xingyue Huang, Miguel Romero Orth, Pablo Barceló, Michael M Bronstein, and İsmail İlkan Ceylan. Link prediction with relational hypergraphs. *arXiv preprint arXiv:2402.04062*, 2024b.

J. D. Hunter. Matplotlib: A 2d graphics environment. *Computing in Science & Engineering*, 9(3): 90–95, 2007. doi: 10.1109/MCSE.2007.55.

Adarsh Jamadandi, Celia Rubio-Madrigal, and Rebekka Burkholz. Spectral graph pruning against over-squashing and over-smoothing, 2024. URL `https://arxiv.org/abs/2404.04612`.

Kedar Karhadkar, Pradeep Kr. Banerjee, and Guido Montúfar. FoSR: First-order spectral rewiring for addressing oversquashing in GNNs. In *The Eleventh International Conference on Learning Representations*, 2023. URL `https://openreview.net/forum?id=3YjQfCLdrzz`.

Thomas N. Kipf and Max Welling. Semi-supervised classification with graph convolutional networks. In *International Conference on Learning Representations*, 2017. URL `https://openreview.net/forum?id=SJU4ayYgl`.

Siu Kwan Lam, Antoine Pitrou, and Stanley Seibert. Numba: a llvm-based python jit compiler. In *Proceedings of the Second Workshop on the LLVM Compiler Infrastructure in HPC*, LLVM '15, New York, NY, USA, 2015. Association for Computing Machinery. ISBN 9781450340052. doi: 10.1145/2833157.2833162. URL `https://doi.org/10.1145/2833157.2833162`.

Hao Li, Chen Li, Jianfei Zhang, Yuanxin Ouyang, and Wenge Rong. Addressing over-squashing in gnns with graph rewiring and ordered neurons. In *2024 International Joint Conference on Neural Networks (IJCNN)*, pp. 1–8. IEEE, 2024.

Qimai Li, Zhichao Han, and Xiao-Ming Wu. Deeper insights into graph convolutional networks for semi-supervised learning. In *Proceedings of the AAAI conference on artificial intelligence*, volume 32, 2018. URL `https://doi.org/10.1609/aaai.v32i1.11604`.

Richard Liaw, Eric Liang, Robert Nishihara, Philipp Moritz, Joseph E Gonzalez, and Ion Stoica. Tune: A research platform for distributed model selection and training. *arXiv preprint arXiv:1807.05118*, 2018.

Yang Liu, Chuan Zhou, Shirui Pan, Jia Wu, Zhao Li, Hongyang Chen, and Peng Zhang. Curvdrop: A ricci curvature based approach to prevent graph neural networks from over-smoothing and over-squashing. In *Proceedings of the ACM Web Conference 2023*, pp. 221–230, 2023.

Yu Liu, Quanming Yao, and Yong Li. Generalizing tensor decomposition for n-ary relational knowledge bases. In *Proceedings of the web conference 2020*, pp. 1104–1114, 2020.

Haggai Maron, Heli Ben-Hamu, Hadar Serviansky, and Yaron Lipman. Provably powerful graph networks. In *Advances in Neural Information Processing Systems*, volume 32. Curran Associates, Inc., 2019. URL https://proceedings.neurips.cc/paper_files/paper/2019/file/bb04af0f7ecaee4aae62035497da1387-Paper.pdf.

Gonzalo Mateos, Santiago Segarra, Antonio G Marques, and Alejandro Ribeiro. Connecting the dots: Identifying network structure via graph signal processing. *IEEE Signal Processing Magazine*, 36(3):16–43, 2019.

Philipp Moritz, Robert Nishihara, Stephanie Wang, Alexey Tumanov, Richard Liaw, Eric Liang, Melih Elibol, Zongheng Yang, William Paul, Michael I Jordan, et al. Ray: A distributed framework for emerging {AI} applications. In *13th USENIX symposium on operating systems design and implementation (OSDI 18)*, pp. 561–577, 2018.

Christopher Morris, Martin Ritzert, Matthias Fey, William L Hamilton, Jan Eric Lenssen, Gaurav Rattan, and Martin Grohe. Weisfeiler and Leman go neural: Higher-order graph neural networks. In *Proceedings of the AAAI conference on artificial intelligence*, volume 33, pp. 4602–4609, 2019. URL https://doi.org/10.1609/aaai.v33i01.33014602.

Christopher Morris, Nils M. Kriege, Franka Bause, Kristian Kersting, Petra Mutzel, and Marion Neumann. TUDataset: A collection of benchmark datasets for learning with graphs. In *ICML 2020 Workshop on Graph Representation Learning and Beyond (GRL+ 2020)*, 2020. URL http://eprints.cs.univie.ac.at/6469/.

Florentin Münch. Ollivier curvature, betweenness centrality and average distance, 2022. URL https://arxiv.org/abs/2209.15564.

V. Nanda. Computational algebraic topology lecture notes, 2021. Available online.

Khang Nguyen, Nong Minh Hieu, Vinh Duc Nguyen, Nhat Ho, Stanley Osher, and Tan Minh Nguyen. Revisiting over-smoothing and over-squashing using Ollivier-Ricci curvature. In *Proceedings of the 40th International Conference on Machine Learning*, volume 202 of *Proceedings of Machine Learning Research*, pp. 25956–25979. PMLR, 2023. URL https://proceedings.mlr.press/v202/nguyen23c.html.

Maximilian Nickel, Kevin Murphy, Volker Tresp, and Evgeniy Gabrilovich. A review of relational machine learning for knowledge graphs. *Proceedings of the IEEE*, 104(1):11–33, 2015.

Hoang Nt and Takanori Maehara. Revisiting graph neural networks: All we have is low-pass filters. *arXiv preprint arXiv:1905.09550*, 2019. URL https://arxiv.org/abs/1905.09550.

Kenta Oono and Taiji Suzuki. Graph neural networks exponentially lose expressive power for node classification. In *International Conference on Learning Representations*, 2020. URL https://openreview.net/forum?id=S1ldO2EFPr.

The pandas development team. pandas-dev/pandas: Pandas, February 2020. URL https://doi.org/10.5281/zenodo.3509134.

Theodore Papamarkou, Tolga Birdal, Michael M. Bronstein, Gunnar E. Carlsson, Justin Curry, Yue Gao, Mustafa Hajij, Roland Kwitt, Pietro Lio, Paolo Di Lorenzo, Vasileios Maroulas, Nina Miolane, Farzana Nasrin, Karthikeyan Natesan Ramamurthy, Bastian Rieck, Simone Scardapane, Michael T Schaub, Petar Veličković, Bei Wang, Yusu Wang, Guowei Wei, and Ghada Zamzmi.

Position: Topological deep learning is the new frontier for relational learning. In *Proceedings of the 41st International Conference on Machine Learning*, volume 235 of *Proceedings of Machine Learning Research*, pp. 39529–39555. PMLR, 2024. URL `https://proceedings.mlr.press/v235/papamarkou24a.html`.

Mathilde Papillon, Mustafa Hajij, Audun Myers, Helen Jenne, Johan Mathe, Theodore Papamarkou, Aldo Guzmán-Sáenz, Neal Livesay, Tamal Dey, Abraham Rabinowitz, Aiden Brent, Alessandro Salatiello, Alexander Nikitin, Ali Zia, Claudio Battiloro, Dmitrii Gavrilev, German Magai, Gleb Bazhenov, Guillermo Bernardez, Indro Spinelli, Jens Agerberg, Kalyan Nadimpalli, Lev Telyatninkov, Luca Scofano, Lucia Testa, Manuel Lecha, Maosheng Yang, Mohammed Hassanin, Odin Hoff Gardaa, Olga Zaghen, Paul Hausner, Paul Snopoff, Rubén Ballester, Sadrodin Barikbin, Sergio Escalera, Simone Fiorellino, Henry Kvinge, Karthikeyan Natesan Ramamurthy, Paul Rosen, Robin Walters, Shreyas N. Samaga, Soham Mukherjee, Sophia Sanborn, Tegan Emerson, Timothy Doster, Tolga Birdal, Abdelwahed Khamis, Simone Scardapane, Suraj Singh, Tatiana Malygina, Yixiao Yue, and Nina Miolane. ICML 2023 Topological Deep Learning Challenge: Design and results. In *Proceedings of 2nd Annual Workshop on Topology, Algebra, and Geometry in Machine Learning (TAG-ML)*, volume 221 of *Proceedings of Machine Learning Research*, pp. 3–8. PMLR, 2023a. URL `https://proceedings.mlr.press/v221/papillon23a.html`.

Mathilde Papillon, Sophia Sanborn, Mustafa Hajij, and Nina Miolane. Architectures of topological deep learning: A survey of message-passing topological neural networks. *arXiv preprint arXiv:2304.10031*, 2023b.

Mathilde Papillon, Guillermo Bernárdez, Claudio Battiloro, and Nina Miolane. Topotune: A framework for generalized combinatorial complex neural networks. *arXiv preprint arXiv:2410.06530*, 2024.

Adam Paszke, Sam Gross, Francisco Massa, Adam Lerer, James Bradbury, Gregory Chanan, Trevor Killeen, Zeming Lin, Natalia Gimelshein, Luca Antiga, et al. Pytorch: An imperative style, high-performance deep learning library. *Advances in neural information processing systems*, 32, 2019.

F. Pedregosa, G. Varoquaux, A. Gramfort, V. Michel, B. Thirion, O. Grisel, M. Blondel, P. Prettenhofer, R. Weiss, V. Dubourg, J. Vanderplas, A. Passos, D. Cournapeau, M. Brucher, M. Perrot, and E. Duchesnay. Scikit-learn: Machine learning in Python. *Journal of Machine Learning Research*, 12:2825–2830, 2011.

Hongbin Pei, Yu Li, Huiqi Deng, Jingxin Hai, Pinghui Wang, Jie Ma, Jing Tao, Yuheng Xiong, and Xiaohong Guan. Multi-track message passing: Tackling oversmoothing and oversquashing in graph learning via preventing heterophily mixing. In *Forty-first International Conference on Machine Learning*.

Chendi Qian, Andrei Manolache, Kareem Ahmed, Zhe Zeng, Guy Van den Broeck, Mathias Niepert, and Christopher Morris. Probabilistic task-adaptive graph rewiring. In *ICML 2023 Workshop on Differentiable Almost Everything: Differentiable Relaxations, Algorithms, Operators, and Simulators*, 2023.

Joshua Robinson, Rishabh Ranjan, Weihua Hu, Kexin Huang, Jiaqi Han, Alejandro Dobles, Matthias Fey, Jan Eric Lenssen, Yiwen Yuan, Zecheng Zhang, et al. Relational deep learning: Graph representation learning on relational databases. In *NeurIPS 2024 Third Table Representation Learning Workshop*, 2024.

T. Konstantin Rusch, Michael M. Bronstein, and Siddhartha Mishra. A survey on oversmoothing in graph neural networks. *arXiv preprint arXiv.2303.10993*, 2023. URL `https://doi.org/10.48550/arXiv.2303.10993`.

Michael Schlichtkrull, Thomas N Kipf, Peter Bloem, Rianne Van Den Berg, Ivan Titov, and Max Welling. Modeling relational data with graph convolutional networks. In *The semantic web: 15th international conference, ESWC 2018, Heraklion, Crete, Greece, June 3–7, 2018, proceedings 15*, pp. 593–607. Springer, 2018.

Dazhong Shen, Chuan Qin, Qi Zhang, Hengshu Zhu, and Hui Xiong. Handling over-smoothing and over-squashing in graph convolution with maximization operation. *IEEE Transactions on Neural Networks and Learning Systems*, 2024a.

Xu Shen, Pietro Lio, Lintao Yang, Ru Yuan, Yuyang Zhang, and Chengbin Peng. Graph rewiring and preprocessing for graph neural networks based on effective resistance. *IEEE Transactions on Knowledge and Data Engineering*, 2024b.

Dai Shi, Andi Han, Lequan Lin, Yi Guo, and Junbin Gao. Exposition on over-squashing problem on GNNs: Current methods, benchmarks and challenges. *arXiv preprint arXiv.2311.07073*, 2023. URL https://doi.org/10.48550/arXiv.2311.07073.

Dai Shi, Andi Han, Lequan Lin, Yi Guo, Zhiyong Wang, and Junbin Gao. Design your own universe: A physics-informed agnostic method for enhancing graph neural networks. *International Journal of Machine Learning and Cybernetics*, pp. 1–16, 2024.

Stevan Stanovic, Benoit Gaüzère, and Luc Brun. Graph neural networks with maximal independent set-based pooling: Mitigating over-smoothing and over-squashing. *Pattern Recognition Letters*, 2024.

Julian Suk, Lorenzo Giusti, Tamir Hemo, Miguel Lopez, Konstantinos Barmpas, and Cristian Bodnar. Surfing on the neural sheaf. In *NeurIPS 2022 Workshop on Symmetry and Geometry in Neural Representations*, 2022. URL https://openreview.net/forum?id=xOXFkyRzTlu.

Li Sun, Zhenhao Huang, Hua Wu, Junda Ye, Hao Peng, Zhengtao Yu, and S Yu Philip. Deepricci: Self-supervised graph structure-feature co-refinement for alleviating over-squashing. In *2023 IEEE International Conference on Data Mining (ICDM)*, pp. 558–567. IEEE, 2023.

The GUDHI Project. *GUDHI User and Reference Manual*. GUDHI Editorial Board, 3.1.1 edition, 2020. URL https://gudhi.inria.fr/doc/3.1.1/.

Jake Topping, Francesco Di Giovanni, Benjamin Paul Chamberlain, Xiaowen Dong, and Michael M. Bronstein. Understanding over-squashing and bottlenecks on graphs via curvature. In *International Conference on Learning Representations*, 2022. URL https://openreview.net/forum?id=7UmjRGzp-A.

Floriano Tori, Vincent Holst, and Vincent Ginis. The effectiveness of curvature-based rewiring and the role of hyperparameters in gnns revisited. *arXiv preprint arXiv:2407.09381*, 2024.

Shikhar Vashishth, Soumya Sanyal, Dasgupta S Nitin, and Partha Talukdar. Composition-based multi-relational graph convolutional networks. In *International Conference on Learning Representations*, 2019.

Petar Veličković, Guillem Cucurull, Arantxa Casanova, Adriana Romero, Pietro Liò, and Yoshua Bengio. Graph attention networks. In *International Conference on Learning Representations*, 2018.

Michael L. Waskom. seaborn: statistical data visualization. *Journal of Open Source Software*, 6 (60):3021, 2021. doi: 10.21105/joss.03021. URL https://doi.org/10.21105/joss.03021.

Felix Wu, Amauri Souza, Tianyi Zhang, Christopher Fifty, Tao Yu, and Kilian Weinberger. Simplifying graph convolutional networks. In *Proceedings of the 36th International Conference on Machine Learning*, volume 97 of *Proceedings of Machine Learning Research*, pp. 6861–6871. PMLR, 2019. URL https://proceedings.mlr.press/v97/wu19e.html.

Keyulu Xu, Weihua Hu, Jure Leskovec, and Stefanie Jegelka. How powerful are graph neural networks? In *International Conference on Learning Representations*, 2019. URL https://openreview.net/forum?id=ryGs6iA5Km.

Naganand Yadati. Neural message passing for multi-relational ordered and recursive hypergraphs. *Advances in Neural Information Processing Systems*, 33:3275–3289, 2020.

A   RELATED WORK

Topological Networks   Topological deep learning (TDL) integrates algebraic topology with neural networks to create message-passing schemes that are more expressive than graph neural networks. An early contribution was the simplicial Weisfeiler-Lehman (SWL) test introduced by Bodnar et al. (2021a), which extends the Weisfeiler-Lehman (WL) test from graphs to simplicial complexes. Simplicial neural networks (SNNs) built on SWL generalize graph isomorphism networks (GIN) (Xu et al., 2019), offering provably stronger expressiveness. CW networks (CWNs) (Bodnar et al., 2021b) extend message passing to cell complexes, achieving greater power than the traditional WL test and exceeding the 3-WL test. These hierarchical, geometrically-grounded representations enable effective handling of higher-order interactions. Hajij et al. (2020) proposed a general message-passing scheme for cell complexes, though it lacks formal analysis of expressiveness and complexity. In contrast, Bodnar et al. (2022) and Suk et al. (2022) introduced neural sheaf diffusion models, which learn sheaf structures over graphs, particularly excelling in heterophilic graph tasks. Attention mechanisms have been integrated into topological message passing in works such as Goh et al. (2022) for simplicial complexes, Barbero et al. (2022) for sheaves, and Giusti et al. (2023) for cellular complexes. Additionally, Hajij et al. (2022) extends message-passing to combinatorial complexes. Our work represents a first step toward extending the theory of oversquashing and oversmoothing, widely studied in graph neural networks, to these topological networks.

We note that our relational structures approach generalizes a powerful perspective in which simplicial complexes and similar constructs are treated as augmented Hasse diagrams—a viewpoint that has shown both practical and theoretical advantages. Recent works have leveraged this perspective: Hajij et al. (2023) provided a general description of topological message passing schemes, Eitan et al. (2024) explored the expressivity limits of topological message passing, and Papillon et al. (2024) introduced a framework for systematically transforming any graph neural network into a topological analog. Our work aligns with and contributes to this growing body of research.

For detailed surveys of topological deep learning architectures, we refer to Papillon et al. (2023b) and Giusti (2024). We also refer the reader to the recent position paper by Papamarkou et al. (2024) on open problems in TDL.

Oversquashing, Oversmoothing, and Graph Rewiring   Oversquashing refers to the phenomenon where information from distant nodes is compressed into fixed-size vectors during message passing, limiting the ability of a model to capture long-range dependencies. In GNNs, this has been extensively studied, with works such as those of Alon & Yahav (2021), Topping et al. (2022), and Di Giovanni et al. (2023) naming and relating the phenomenon to the geometric properties of graphs. To address this issue, numerous techniques have been proposed, including spatial and curvature-based rewiring methods such as SDRF (Topping et al., 2022), BORF (Nguyen et al., 2023), and AFRC (Fesser & Weber, 2023), which modify the graph structure to improve connectivity and alleviate oversquashing. Spectral rewiring approaches, such as FOSR (Karhadkar et al., 2023), optimize the graph spectral gap to enhance long-range message propagation, while implicit methods such as Graph Beltrami Diffusion (Chamberlain et al., 2021) and Graph Transformers (Dwivedi & Bresson, 2020) allow information flow across a fully connected graph without explicitly modifying the topology.

Related to oversquashing, another significant challenge in GNNs is oversmoothing, where node features become indistinguishable as they are excessively aggregated through layers of message passing. Li et al. (2018), Oono & Suzuki (2020), and Nt & Maehara (2019) have identified and theoretically analyzed oversmoothing, showing that it limits the effectiveness of deep GNNs. Generally, there is a trade-off between mitigating oversquashing and avoiding oversmoothing, and various rewiring techniques aim to balance these competing objectives.

Following foundational work on understanding and mitigating oversquashing and oversmoothing in graph neural networks, the field has rapidly evolved with various innovative and sophisticated approaches in recent years. For example, Liu et al. (2023) introduced CurvDrop, a Ricci curvature-based, topology-aware method to address both failure modes of message passing, while Sun et al. (2023) proposed DeepRicci, a self-supervised Riemannian model designed to alleviate oversquashing. Qian et al. (2023) developed a probabilistic framework leveraging differentiable k-subset sampling, and Shen et al. (2024a) introduced a technique based on maximization operations in graph

convolution. Shen et al. (2024b) used effective resistance for graph rewiring and preprocessing. Li et al. (2024) combined graph rewiring with ordered neurons, while Stanovic et al. (2024) proposed maximal independent set-based pooling to mitigate both oversquashing and oversmoothing. Shi et al. (2024) presented a physics-informed, agnostic method targeting both failure modes, and Pei et al. addressed them through multi-track message passing. Attali et al. employed Delaunay triangulation of features, while Huang et al. (2024a) proposed particular spectral filters to improve spectral GNNs. Collectively, these works demonstrate the breadth and depth of recent advancements, highlighting a vibrant and rapidly evolving research landscape.

However, the study of oversquashing and oversmoothing in topological message passing remains largely unexplored. While works such as those of Bodnar et al. (2021b), Bodnar et al. (2021a), and Giusti et al. (2024) have alluded to the potential of topological message passing to capture group interactions and long-range dependencies, a theoretical analysis of oversquashing and oversmoothing in higher-order structures like simplicial or cellular complexes is, to the best of our knowledge, absent from the literature. Our work aims to fill this gap, providing a rigorous first step in studying oversquashing in topological message passing. Moreover, our framework for extending graph rewiring to relational rewiring (Section 4), along with future innovations from the topological deep learning community, promises a wealth of topological and relational rewiring techniques for further assessment and refinement.

For more exhaustive expositions on oversquashing and oversmoothing, we refer the reader to the excellent recent surveys of Shi et al. (2023) and Rusch et al. (2023).

Relational Learning   Relational graph neural networks (R-GNNs) extend traditional GNNs to handle multi-relational data, particularly in the context of knowledge graphs, where nodes represent entities and edges capture diverse types of relationships between them. Knowledge graphs, such as those used in knowledge representation and reasoning tasks (Nickel et al., 2015), are inherently multi-relational and benefit significantly from RGNNs, which model the varying nature of relations explicitly. The relational graph convolutional network (R-GCN) (Schlichtkrull et al., 2018) is a foundational approach that assigns distinct transformations to each relation type, allowing for efficient representation learning on multi-relational graphs. Extensions such as Relational Graph Attention Networks (RGAT) (Veličković et al., 2018) incorporate attention mechanisms, enabling the model to focus on the most important relations during message passing. Additionally, Composition-based Graph Convolutional Networks (CompGCN) (Vashishth et al., 2019) apply compositional operators to better capture interactions in knowledge graphs.

Recent works in relational learning have introduced various extensions for modeling multi-relational and higher-order interactions. Hypergraph-based approaches, such as Fatemi et al. (2023) and Huang et al. (2024b), unify relational reasoning with hypergraph structures, enabling the modeling of higher-order relationships. Message-passing frameworks like the one proposed by Yadati (2020) extend traditional GNN paradigms to ordered and recursive hypergraphs. Additionally, tensor decomposition methods such as GETD (Liu et al., 2020) represent hypergraphs as high-dimensional tensors, allowing efficient encoding of hyper-relational data. We refer the reader to the excellent survey by Antelmi et al. (2023) for a comprehensive overview of relational hypergraphs and their representation learning techniques. Robinson et al. (2024) recently extended relational learning to relational databases consisting of data laid out across multiple tables. Lastly, topological deep learning, which is one focus of this work, is a new frontier in relational learning (Papamarkou et al., 2024).

Our work unifies relational graph neural networks and topological neural networks by viewing complexes as relational structures, bridging the gap between the two fields.

Graph Lifting   Graph lifting transforms a graph into a higher-dimensional structure to enable more expressive message-passing schemes. For example, higher-order graph neural networks (k-GNNs) (Morris et al., 2019) lift graphs by representing k-node subgraphs as entities, capturing more complex relationships between nodes. Similarly, Chen et al. (2019) and Maron et al. (2019) proposed lifting graphs to higher-order structures to improve graph isomorphism testing. In topological message passing, graphs are lifted into structures like clique complexes to capture interactions beyond pairwise relationships. Recently, this topic has gained increased attention, as highlighted by the ICML Topological Deep Learning Challenge 2024 (Papillon et al., 2023a), which emphasized the

development of topological lifting techniques across various data structures, including graphs, hypergraphs, and simplicial complexes. Among these, high-order graphs, simplicial complexes, and cellular complexes fit naturally within the relational structure framework and can be analyzed uniformly through this lens.

B ADDITIONAL REMARKS

B.1 REMARKS FOR SECTION 3.2

Remark B.1 (Ollivier-Ricci Curvature for Weighted Directed Graphs). *We recall the definition of Ollivier-Ricci curvature (ORC) for weighted directed graphs from Eidi & Jost (2020). Consider a weighted directed graph $\mathcal{G} = (\mathcal{S}, \mathcal{E}, w)$, where $\mathcal{S}$ is the set of entities (nodes), $\mathcal{E}$ is the set of directed edges, and $w : \mathcal{E} \to \mathbb{R}_{\geq 0}$ assigns non-negative weights to edges. We abuse notation, and for any two entities $\xi, \eta \in \mathcal{S}$ without an edge $(\xi \to \eta) \in \mathcal{E}$, we write $w_{\xi \to \eta} = 0$. For each entity $\sigma \in \mathcal{S}$, denote the* weighted out-degree *and* weighted in-degree*, as in Definition 3.3, by:*

$$w_{\sigma}^{\text{out}} = \sum_{(\sigma \to \eta) \in \mathcal{E}} w_{\sigma \to \eta}, \quad w_{\sigma}^{\text{in}} = \sum_{(\xi \to \sigma) \in \mathcal{E}} w_{\xi \to \sigma}.$$

For a directed edge $(\sigma \to \tau) \in \mathcal{E}$ with non-zero weight $w_{\sigma \to \tau} > 0$, we define the probability measures:

$$\mu_{\tau}^{\text{out}}(\xi) = \frac{w_{\tau \to \xi}}{w_{\tau}^{\text{out}}}, \quad \mu_{\sigma}^{\text{in}}(\xi) = \frac{w_{\xi \to \sigma}}{w_{\sigma}^{\text{in}}}, \quad \xi \in \mathcal{S}.$$

The Ollivier-Ricci curvature *of an edge $(\sigma \to \tau)$ with non-zero weight $w_{\sigma \to \tau}$ is then defined as:*

$$k(\sigma, \tau) = 1 - \frac{W\left(\mu_{\sigma}^{\text{in}}, \mu_{\tau}^{\text{out}}\right)}{w_{\sigma \to \tau}},$$

where $W\left(\mu_{\sigma}^{\text{in}}, \mu_{\tau}^{\text{out}}\right)$ is the (directed) Wasserstein distance between the measures μ_{σ}^{in} and μ_{τ}^{out}, defined by:

$$W\left(\mu_{\sigma}^{\text{in}}, \mu_{\tau}^{\text{out}}\right) = \inf_{\pi \in \Pi(\mu_{\sigma}^{\text{in}}, \mu_{\tau}^{\text{out}})} \sum_{\xi, \eta \in \mathcal{S}} \pi(\xi, \eta) d(\xi, \eta),$$

where $\Pi(\mu_{\sigma}^{\text{in}}, \mu_{\tau}^{\text{out}})$ is the set of all joint probability measures on $\mathcal{S} \times \mathcal{S}$ with marginals μ_{σ}^{in} and μ_{τ}^{out}, and $d(\xi, \eta)$ is the distance from ξ (incoming neighbor of σ) to η (outgoing neighbor of τ) in the graph $\mathcal{G}$ (the shortest directed path distance based on edge weights).

Remark B.2. *Assume that the edge $(\sigma \to \tau) \in \mathcal{E}$ exists and has non-zero weight $w_{\sigma \to \tau} > 0$ in the influence graph $\mathcal{G}(\mathcal{S}, \mathbf{B})$. Then, the sensitivity bound for information flow from τ to σ from Lemma 3.2 can be related to the Ollivier-Ricci curvature $k(\sigma, \tau)$ of the edge $(\sigma \to \tau)$ as follows:*

$$\left\| \frac{\partial \mathbf{h}_{\sigma}^{(2)}}{\partial \mathbf{h}_{\tau}^{(0)}} \right\|_1 \leq \left(\prod_{\ell=0}^{1} \alpha^{(\ell)} \beta^{(\ell)} \right) w_{\tau}^{\text{out}} w_{\sigma}^{\text{in}} \left(1 - \frac{w_{\sigma \to \tau}}{w_{\max}^3} (1 - k(\sigma, \tau)) \right),$$

where $w_{\max}^3$ is the maximum weighted 3-step path from incoming neighbors of σ to outgoing neighbors of τ. This result indicates that lower Ollivier-Ricci curvature (smaller $k(\sigma, \tau)$) leads to reduced sensitivity, thereby contributing to oversquashing. The connection leverages the existence of the reversed edge $(\sigma \to \tau)$ in addition to the edge $(\tau \to \sigma)$ over which information flows, which is not guaranteed to be the case in a generic relational message passing scheme, but is the case for undirected graphs where $k(\sigma, \tau) = k(\tau, \sigma)$. This result aligns with Theorem 4.5 from Nguyen et al. (2023) but without requiring their assumption of linearity in the message and update functions. Analyzing oversquashing for connections beyond 2-steps requires stronger assumptions, such as the influence graph $\mathcal{G}(\mathcal{S}, \mathbf{B})$ being strongly connected. These assumptions present obstructions to analyzing oversquashing in general relational structures with Ollivier-Ricci curvature.

We provide the proof in Appendix C.2.

B.2 REMARKS FOR SECTION 3.4

Remark B.3 (Lipschitz Constants for MLP Message and Update Functions). *Consider the following message and update functions at layer t:*

Message function:

$$\boldsymbol{\psi}_i^{(t)}\left(\mathbf{h}_\sigma^{(t)}, \mathbf{h}_{\xi_1}^{(t)}, \ldots, \mathbf{h}_{\xi_{n_i-1}}^{(t)}\right) = \mathbf{W}_i^{(t)} \begin{bmatrix} \mathbf{h}_\sigma^{(t)} \\ \mathbf{h}_{\xi_1}^{(t)} \\ \vdots \\ \mathbf{h}_{\xi_{n_i-1}}^{(t)} \end{bmatrix},$$

where $\mathbf{W}_i^{(t)}$ is a weight matrix of appropriate dimensions, and $[\cdot]$ denotes column-wise concatenation.

Update function:

$$\boldsymbol{\phi}^{(t)}\left(\mathbf{m}_{\sigma,1}^{(t)}, \ldots, \mathbf{m}_{\sigma,k}^{(t)}\right) = \mathbf{f}\left(\mathbf{W}^{(t)} \begin{bmatrix} \mathbf{g}\left(\mathbf{m}_{\sigma,1}^{(t)}\right) \\ \vdots \\ \mathbf{g}\left(\mathbf{m}_{\sigma,k}^{(t)}\right) \end{bmatrix}\right),$$

where $\mathbf{f}$ and $\mathbf{g}$ are the component-wise applications of non-linear functions f and g with bounded derivatives C_f and C_g, respectively. I.e., $f, g : \mathbb{R} \to \mathbb{R}$, and $|f'(x)| \leq C_f$ and $|g'(x)| \leq C_g$ for all x.

Assume that the entries of all weight matrices $\mathbf{W}_i^{(t)}$ and $\mathbf{W}^{(t)}$ are bounded in absolute value by a constant $C_w > 0$. Then, the Lipschitz constants $\beta_i^{(t)}$ and $\alpha^{(t)}$ satisfy:

$$\beta_i^{(t)} \leq C_w p_{i,t},$$

and

$$\alpha^{(t)} \leq C_w C_f C_g p_{t+1},$$

where $p_{i,t}$ is the dimension of the message vector $\mathbf{m}_{\sigma,i}^{(t)}$, and p_{t+1} is the output dimension of the update function ϕ_t.

We provide the proof in Appendix C.4.

C PROOFS

C.1 PROOFS FOR SECTION 3.1

Proof of Lemma 3.2. First, we compute the Jacobian of Equation 3:

$$\frac{\partial \mathbf{h}_\sigma^{(s+1)}}{\partial \mathbf{h}_\tau^{(0)}} = \sum_{i=1}^{k} \frac{\partial \boldsymbol{\phi}^{(s)}}{\partial \mathbf{m}_i^{(s)}} \frac{\partial \mathbf{m}_{\sigma,i}^{(s)}}{\partial \mathbf{h}_\tau^{(0)}}$$

$$= \sum_{i=1}^{k} \left(\frac{\partial \boldsymbol{\phi}^{(s)}}{\partial \mathbf{m}_i^{(s)}}\right) \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-1}} \mathbf{A}_{\sigma,\boldsymbol{\xi}}^{R_i} \left[\frac{\partial \boldsymbol{\psi}_i^{(s)}}{\partial \mathbf{h}_\sigma^{(s)}} \frac{\partial \mathbf{h}_\sigma^{(s)}}{\partial \mathbf{h}_\tau^{(0)}} + \sum_{j=1}^{n_i-1} \frac{\partial \boldsymbol{\psi}_i^{(s)}}{\partial \mathbf{h}_{\xi_j}^{(s)}} \frac{\partial \mathbf{h}_{\xi_j}^{(s)}}{\partial \mathbf{h}_\tau^{(0)}}\right].$$

By the submultiplicative and additive properties of the induced 1-norm (maximum absolute column sum) and the boundedness of the Jacobians of the functions $\phi^{(s)}$ and $\psi_i^{(s)}$ (Assumption 1), we get:

$$\left\| \frac{\partial \mathbf{h}_\sigma^{(s+1)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \leq \sum_{i=1}^k \left\| \frac{\partial \phi^{(s)}}{\partial \mathbf{m}_i^{(s)}} \right\|_1 \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-1}} \mathbf{A}_{\sigma,\boldsymbol{\xi}}^{R_i} \left[\left\| \frac{\partial \psi_i^{(s)}}{\partial \mathbf{h}_\sigma^{(s)}} \right\|_1 \left\| \frac{\partial \mathbf{h}_\sigma^{(s)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 + \sum_{j=1}^{n_i-1} \left\| \frac{\partial \psi_i^{(s)}}{\partial \mathbf{h}_{\xi_j}^{(s)}} \right\|_1 \left\| \frac{\partial \mathbf{h}_{\xi_j}^{(s)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \right]$$

$$\leq \alpha^{(s)} \sum_{i=1}^k \beta_i^{(s)} \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-1}} \mathbf{A}_{\sigma,\boldsymbol{\xi}}^{R_i} \left[\left\| \frac{\partial \mathbf{h}_\sigma^{(s)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 + \sum_{j=1}^{n_i-1} \left\| \frac{\partial \mathbf{h}_{\xi_j}^{(s)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \right]$$

$$\leq \alpha^{(s)} \beta^{(s)} \sum_{i=1}^k \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-1}} \mathbf{A}_{\sigma,\boldsymbol{\xi}}^{R_i} \left[\left\| \frac{\partial \mathbf{h}_\sigma^{(s)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 + \sum_{j=1}^{n_i-1} \left\| \frac{\partial \mathbf{h}_{\xi_j}^{(s)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \right]$$

$$\leq \alpha^{(s)} \beta^{(s)} \left[\gamma \left\| \frac{\partial \mathbf{h}_\sigma^{(s)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 + \sum_{i=1}^k \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-1}} \mathbf{A}_{\sigma,\boldsymbol{\xi}}^{R_i} \sum_{j=1}^{n_i-1} \left\| \frac{\partial \mathbf{h}_{\xi_j}^{(s)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \right].$$

Here, we used that the entries of $\mathbf{A}^{R_i}$ are nonnegative, and $\sum_{i=1}^k \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-1}} \mathbf{A}_{\sigma,\boldsymbol{\xi}}^{R_i} \leq \gamma$.

We now prove the lemma using induction on the layer t. For the base case $t = 1$, we get

$$\left\| \frac{\partial \mathbf{h}_\sigma^{(1)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \leq \alpha^{(0)} \beta^{(0)} \left[\gamma \left\| \frac{\partial \mathbf{h}_\sigma^{(0)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 + \sum_{i=1}^k \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-1}} \mathbf{A}_{\sigma,\boldsymbol{\xi}}^{R_i} \sum_{j=1}^{n_i-1} \left\| \frac{\partial \mathbf{h}_{\xi_j}^{(0)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \right]$$

$$= \alpha^{(0)} \beta^{(0)} \left[\gamma \mathbf{I}_{\sigma,\tau} + \sum_{i=1}^k \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-1}} \mathbf{A}_{\sigma,\boldsymbol{\xi}}^{R_i} \sum_{j=1}^{n_i-1} \mathbf{I}_{\xi_j,\tau} \right]$$

$$= \alpha^{(0)} \beta^{(0)} \left[\gamma \mathbf{I}_{\sigma,\tau} + \sum_{i=1}^k \tilde{\mathbf{A}}_{\sigma,\tau}^{R_i} \right]$$

$$= \alpha^{(0)} \beta^{(0)} \left(\gamma \mathbf{I} + \sum_{i=1}^k \tilde{\mathbf{A}}^{R_i} \right)_{\sigma,\tau}$$

$$= \alpha^{(0)} \beta^{(0)} \left(\mathbf{B}^1 \right)_{\sigma,\tau},$$

where we used that $\tilde{\mathbf{A}}_{\sigma,\tau}^{R_i} = \sum_{j=1}^{n_i-1} \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-2}} \mathbf{A}_{\sigma,\xi_1,\ldots,\xi_{j-1},\tau,\xi_j,\ldots,\xi_{n_i-2}}^{R_i}$. This proves the base case.

For the induction step, assume the bound holds for t. We now compute:

$$\left\| \frac{\partial \mathbf{h}_\sigma^{(t+1)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \leq \alpha^{(t)} \beta^{(t)} \left[\gamma \left\| \frac{\partial \mathbf{h}_\sigma^{(t)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 + \sum_{i=1}^k \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-1}} \mathbf{A}_{\sigma,\boldsymbol{\xi}}^{R_i} \sum_{j=1}^{n_i-1} \left\| \frac{\partial \mathbf{h}_{\xi_j}^{(t)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \right]$$

$$\leq \left(\prod_{\ell=0}^t \alpha^{(\ell)} \beta^{(\ell)} \right) \left[\gamma \left(\mathbf{B}^t \right)_{\sigma,\tau} + \sum_{i=1}^k \sum_{\boldsymbol{\xi} \in \mathcal{S}^{n_i-1}} \mathbf{A}_{\sigma,\boldsymbol{\xi}}^{R_i} \sum_{j=1}^{n_i-1} \left(\mathbf{B}^t \right)_{\xi_j,\tau} \right]$$

$$= \left(\prod_{\ell=0}^t \alpha^{(\ell)} \beta^{(\ell)} \right) \left[\left(\gamma \mathbf{I} \mathbf{B}^t \right)_{\sigma,\tau} + \left(\left(\sum_{i=1}^k \tilde{\mathbf{A}}^{R_i} \right) \mathbf{B}^t \right)_{\sigma,\tau} \right]$$

$$= \left(\prod_{\ell=0}^t \alpha^{(\ell)} \beta^{(\ell)} \right) \left(\left(\gamma \mathbf{I} + \sum_{i=1}^k \tilde{\mathbf{A}}^{R_i} \right) \mathbf{B}^t \right)_{\sigma,\tau}$$

$$= \left(\prod_{\ell=0}^t \alpha^{(\ell)} \beta^{(\ell)} \right) \left(\mathbf{B}^{t+1} \right)_{\sigma,\tau}.$$

To see this, note that:

$$\left(\left(\sum_{i=1}^{k}\tilde{\mathbf{A}}^{R_i}\right)\mathbf{B}^t\right)_{\sigma,\tau} = \sum_{i=1}^{k}\sum_{\nu\in\mathcal{S}}\tilde{\mathbf{A}}_{\sigma,\nu}^{R_i}\left(\mathbf{B}^t\right)_{\nu,\tau}$$

$$= \sum_{i=1}^{k}\sum_{j=1}^{n_i-1}\sum_{\boldsymbol{\xi}\in\mathcal{S}^{n_i-2}}\sum_{\nu\in\mathcal{S}}\mathbf{A}_{\sigma,\xi_1,\dots,\xi_{j-1},\nu,\xi_j,\dots,\xi_{n_i-2}}^{R_i}\left(\mathbf{B}^t\right)_{\nu,\tau}$$

$$= \sum_{i=1}^{k}\sum_{\boldsymbol{\xi}\in\mathcal{S}^{n_i-1}}\mathbf{A}_{\sigma,\boldsymbol{\xi}}^{R_i}\sum_{j=1}^{n_i-1}\left(\mathbf{B}^t\right)_{\xi_j,\tau}.$$

This completes the proof. $\square$

C.2 Proofs for Section 3.2 and Appendix B.1

Proof of Proposition 3.4. In the influence graph $\mathcal{G}(\mathcal{S},\mathbf{B})$, the edge weights correspond to the entries of $\mathbf{B}$, that is, $w_{\tau\to\sigma} = \mathbf{B}_{\sigma,\tau}$. Therefore, the (σ,τ) entry of $\mathbf{B}^2$ is

$$(\mathbf{B}^2)_{\sigma,\tau} = \sum_{\xi\in\mathcal{S}}\mathbf{B}_{\sigma,\xi}\mathbf{B}_{\xi,\tau} = \sum_{\xi\in\mathcal{S}}w_{\xi\to\sigma}\cdot w_{\tau\to\xi} = w_T.$$

From Lemma 3.2 with $t=2$, we have

$$\left\|\frac{\partial\mathbf{h}_\sigma^{(2)}}{\partial\mathbf{h}_\tau^{(0)}}\right\|_1 \le \left(\prod_{\ell=0}^{1}\alpha^{(\ell)}\beta^{(\ell)}\right)(\mathbf{B}^2)_{\sigma,\tau} = \left(\prod_{\ell=0}^{1}\alpha^{(\ell)}\beta^{(\ell)}\right)w_T.$$

Rewriting the curvature formula to solve for w_T, and since $w_F \ge 0$, we get

$$w_T = \frac{1}{3}\left(\mathrm{EFC}_\mathcal{G}(\tau,\sigma) + w_\tau^{\mathrm{out}} + w_\sigma^{\mathrm{in}} - 4 - 2w_F\right)$$

$$\le \frac{1}{3}\left(\mathrm{EFC}_\mathcal{G}(\tau,\sigma) + w_\tau^{\mathrm{out}} + w_\sigma^{\mathrm{in}} - 4\right).$$

Substituting back, we obtain the desired inequality. This completes the proof. $\square$

Proof of Remark B.2. Consider the influence graph $\mathcal{G}(\mathcal{S},\mathbf{B}) = (\mathcal{S},\mathcal{E},w)$ derived from the matrix $\mathbf{B}$, where the edge weights correspond to the entries of $\mathbf{B}$: $w_{\xi\to\eta} = \mathbf{B}_{\eta,\xi}$. Assume that the reversed edge $(\sigma\to\tau)\in\mathcal{E}$ exists with weight $w_{\sigma\to\tau} > 0$. From Lemma 3.2 with $t=2$, we have:

$$\left\|\frac{\partial\mathbf{h}_\sigma^{(2)}}{\partial\mathbf{h}_\tau^{(0)}}\right\|_1 \le \left(\prod_{\ell=0}^{1}\alpha^{(\ell)}\beta^{(\ell)}\right)(\mathbf{B}^2)_{\sigma,\tau} = \left(\prod_{\ell=0}^{1}\alpha^{(\ell)}\beta^{(\ell)}\right)\sum_{\xi\in\mathcal{S}}w_{\xi\to\sigma}\cdot w_{\tau\to\xi} = \left(\prod_{\ell=0}^{1}\alpha^{(\ell)}\beta^{(\ell)}\right)w_T,$$

where $w_T = \sum_{\xi}w_{\tau\to\xi}\cdot w_{\xi\to\sigma}$.

We construct a *transference plan* to transport mass from μ_σ^{in} to μ_τ^{out}: We do not move $\frac{w_T}{w_\tau^{\mathrm{out}}w_\sigma^{\mathrm{in}}}$, but move the rest of the mass with a cost at most w_{max}^3 from the incoming neighbors of σ to the outgoing neighbors of τ, and we get:

$$W\left(\mu_\sigma^{\mathrm{in}},\mu_\tau^{\mathrm{out}}\right) \le \left(1 - \frac{w_T}{w_\sigma^{\mathrm{in}}w_\tau^{\mathrm{out}}}\right)w_{\mathrm{max}}^3.$$

The Ollivier-Ricci curvature of the edge $(\sigma\to\tau)$ is:

$$k(\sigma,\tau) = 1 - \frac{W\left(\mu_\sigma^{\mathrm{in}},\mu_\tau^{\mathrm{out}}\right)}{w_{\sigma\to\tau}}.$$

Substituting and rearranging, we obtain:

$$w_T \le w_\sigma^{\mathrm{in}}w_\tau^{\mathrm{out}}\left(1 - \frac{w_{\sigma\to\tau}}{w_{\mathrm{max}}^3}(1 - k(\sigma,\tau))\right).$$

We thus get:

$$\left\|\frac{\partial\mathbf{h}_\sigma^{(2)}}{\partial\mathbf{h}_\tau^{(0)}}\right\|_1 \le \left(\prod_{\ell=0}^{1}\alpha^{(\ell)}\beta^{(\ell)}\right)w_T \le \left(\prod_{\ell=0}^{1}\alpha^{(\ell)}\beta^{(\ell)}\right)w_\tau^{\mathrm{out}}w_\sigma^{\mathrm{in}}\left(1 - \frac{w_{\sigma\to\tau}}{w_{\mathrm{max}}^3}(1 - k(\sigma,\tau))\right).$$

This completes the proof. $\square$

C.3 PROOFS FOR SECTION 3.3

Proof of Theorem 3.5. We start with the bound from Lemma 3.2 and expand the right hand side using the binomial theorem:

$$
\left\| \frac{\partial \mathbf{h}_\sigma^{(r+m)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \leq \left(\prod_{\ell=0}^{r+m-1} \alpha^{(\ell)}\beta^{(\ell)} \right) (\gamma \mathbf{I} + \tilde{\mathbf{A}})_{\sigma,\tau}^{r+m}
$$

$$
= \left(\prod_{\ell=0}^{r+m-1} \alpha^{(\ell)}\beta^{(\ell)} \right) \sum_{i=0}^{r+m} \binom{r+m}{i} \gamma^{r+m-i} \left(\tilde{\mathbf{A}}^i \right)_{\sigma,\tau}.
$$

Since the combinatorial distance from τ to σ in the graph $\mathcal{G}(\mathcal{S}, \tilde{\mathbf{A}})$ is r, the first $r-1$ terms of the sum vanish:

$$
\left\| \frac{\partial \mathbf{h}_\sigma^{(r+m)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \leq \left(\prod_{\ell=0}^{r+m-1} \alpha^{(\ell)}\beta^{(\ell)} \right) \sum_{i=r}^{r+m} \binom{r+m}{i} \gamma^{r+m-i} \left(\tilde{\mathbf{A}}^i \right)_{\sigma,\tau}.
$$

Using $\tilde{\mathbf{A}}_{\sigma,\tau}^i \leq \omega_i(\sigma,\tau) M^i \leq \omega_{r+m}(\sigma,\tau) M^i$ and letting $q = i - r$:

$$
\left\| \frac{\partial \mathbf{h}_\sigma^{(r+m)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \leq \left(\prod_{\ell=0}^{r+m-1} \alpha^{(\ell)}\beta^{(\ell)} \right) M^r \omega_{r+m}(\sigma,\tau) \sum_{q=0}^{m} \binom{r+m}{r+q} \gamma^{m-q} M^q.
$$

We can bound $\binom{r+m}{r+q}$ as follows:

$$
\binom{r+m}{r+q} = \frac{(r+m)(r-1+m)\cdots(1+m)}{(r+q)(r-1+q)\cdots(1+q)} \binom{m}{q} \leq \frac{(r+m)(r-1+m)\cdots(1+m)}{r!} \binom{m}{q}
$$

$$
\leq \left(1 + \frac{m}{r}\right)\cdots\left(1 + \frac{m}{1}\right) \binom{m}{q} \leq \left(1 + \frac{m}{m+1}\right)^{r-m} (1+m)^m \binom{m}{q}.
$$

Substituting this bound, we get:

$$
\left\| \frac{\partial \mathbf{h}_\sigma^{(r+m)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \leq \left(\prod_{\ell=0}^{r+m-1} \alpha^{(\ell)}\beta^{(\ell)} \right) M^r \omega_{r+m}(\sigma,\tau) \left(1 + \frac{m}{m+1}\right)^{r-m} (1+m)^m \sum_{q=0}^{m} \binom{m}{q} \gamma^{m-q} M^q
$$

$$
= \left(\prod_{\ell=0}^{r+m-1} \alpha^{(\ell)}\beta^{(\ell)} \right) M^r \omega_{r+m}(\sigma,\tau) \left(1 + \frac{m}{m+1}\right)^{r-m} (1+m)^m (\gamma + M)^m
$$

$$
= \left(\prod_{\ell=r}^{r+m-1} \alpha^{(\ell)}\beta^{(\ell)} \right) \left(\frac{(1+m)^2}{2m+1} (\gamma + M) \right)^m \omega_{r+m}(\sigma,\tau) \left(\prod_{\ell=0}^{r-1} \alpha^{(\ell)}\beta^{(\ell)} \right) \left(\left(1 + \frac{m}{m+1}\right) M \right)^r.
$$

Using $1 + \frac{m}{m+1} \leq 2$, $M \leq k$, and $\gamma \leq k$:

$$
\left\| \frac{\partial \mathbf{h}_\sigma^{(r+m)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \leq \left(\prod_{\ell=r}^{r+m-1} \alpha^{(\ell)}\beta^{(\ell)} \right) (2k(1+m))^m \omega_{r+m}(\sigma,\tau) \left(\prod_{\ell=0}^{r-1} \alpha^{(\ell)}\beta^{(\ell)} \right) (2M)^r.
$$

Define $C = (\alpha_{\max}\beta_{\max})^m (2k(1+m))^m$. This depends only on $\alpha_{\max}$, $\beta_{\max}$, k, and m.

Finally, we can write:

$$
\left\| \frac{\partial \mathbf{h}_\sigma^{(r+m)}}{\partial \mathbf{h}_\tau^{(0)}} \right\|_1 \leq C \omega_{r+m}(\sigma,\tau)(\alpha_{\max}\beta_{\max})^r (2M)^r. \tag{13}
$$

This completes the proof. □

C.4 PROOFS FOR SECTION 3.4 AND APPENDIX B.2

Proof of Remark B.3. We derive the bounds for $\beta_i^{(t)}$ and $\alpha^{(t)}$ separately.

Derivation of $\beta_i^{(t)}$:

The message function is linear:

$$\boldsymbol{\psi}_i^{(t)}\left(\mathbf{h}_\sigma^{(t)}, \mathbf{h}_{\xi_1}^{(t)}, \ldots, \mathbf{h}_{\xi_{n_i-1}}^{(t)}\right) = \mathbf{W}_i^{(t)}\mathbf{H}_{\sigma,\xi}^{(t)},$$

where $\mathbf{H}_{\sigma,\xi}^{(t)} = \begin{bmatrix} \mathbf{h}_\sigma^{(t)} \\ \mathbf{h}_{\xi_1}^{(t)} \\ \vdots \\ \mathbf{h}_{\xi_{n_i-1}}^{(t)} \end{bmatrix}$ is the concatenated feature vector.

We need to compute the Lipschitz constant $\beta_i^{(t)}$ of the message function $\boldsymbol{\psi}_i^{(t)}$ with respect to each neighbor's feature vector $\mathbf{h}_{\xi_j}^{(t)}$.

The Jacobian of $\boldsymbol{\psi}_i^{(t)}$ with respect to $\mathbf{h}_{\xi_j}^{(t)}$ is:

$$\frac{\partial \boldsymbol{\psi}_i^{(t)}}{\partial \mathbf{h}_{\xi_j}^{(t)}} = \mathbf{W}_i^{(t)(:,\mathcal{I}_j)},$$

where $\mathbf{W}_i^{(t)(:,\mathcal{I}_j)}$ denotes the columns of $\mathbf{W}_i^{(t)}$ corresponding to $\mathbf{h}_{\xi_j}^{(t)}$.

Since the entries of $\mathbf{W}_i^{(t)}$ are bounded by C_w, and $\mathbf{h}_{\xi_j}^{(t)} \in \mathbb{R}^{p_t}$, the matrix $\frac{\partial \boldsymbol{\psi}_i^{(t)}}{\partial \mathbf{h}_{\xi_j}^{(t)}}$ is of size $p_{i,t} \times p_t$ with entries bounded by C_w.

Using the induced matrix 1-norm:

$$\left\|\frac{\partial \boldsymbol{\psi}_i^{(t)}}{\partial \mathbf{h}_{\xi_j}^{(t)}}\right\|_1 = \max_{1 \le l \le p_t} \sum_{k=1}^{p_{i,t}} \left|\left(\frac{\partial \boldsymbol{\psi}_i^{(t)}}{\partial \mathbf{h}_{\xi_j}^{(t)}}\right)_{k,l}\right|.$$

Since each entry $\left|\left(\frac{\partial \boldsymbol{\psi}_i^{(t)}}{\partial \mathbf{h}_{\xi_j}^{(t)}}\right)_{k,l}\right| \le C_w$, the sum over k is bounded by $C_w p_{i,t}$. Therefore,

$$\left\|\frac{\partial \boldsymbol{\psi}_i^{(t)}}{\partial \mathbf{h}_{\xi_j}^{(t)}}\right\|_1 \le C_w p_{i,t}.$$

Thus, the Lipschitz constant $\beta_i^{(t)}$ satisfies:

$$\beta_i^{(t)} \le C_w p_{i,t}.$$

Derivation of $\alpha^{(t)}$:

The update function is given by:

$$\phi^{(t)}\left(\mathbf{m}_{\sigma,1}^{(t)}, \ldots, \mathbf{m}_{\sigma,k}^{(t)}\right) = \mathbf{f}\left(\mathbf{W}^{(t)}\mathbf{M}_\sigma^{(t)}\right),$$

where $\mathbf{M}_\sigma^{(t)} = \begin{bmatrix} \mathbf{g}\left(\mathbf{m}_{\sigma,1}^{(t)}\right) \\ \vdots \\ \mathbf{g}\left(\mathbf{m}_{\sigma,k}^{(t)}\right) \end{bmatrix}.$

We need to compute the Lipschitz constant $\alpha^{(t)}$ of the update function $\phi^{(t)}$ with respect to each input message $\mathbf{m}_{\sigma,i}^{(t)}$.

First, compute the Jacobian of $\phi^{(t)}$ with respect to $\mathbf{m}_{\sigma,i}^{(t)}$:

$$\frac{\partial \phi^{(t)}}{\partial \mathbf{m}_{\sigma,i}^{(t)}} = \frac{\partial \phi^{(t)}}{\partial \mathbf{M}_{\sigma}^{(t)}} \frac{\partial \mathbf{M}_{\sigma}^{(t)}}{\partial \mathbf{m}_{\sigma,i}^{(t)}}.$$

Compute $\frac{\partial \mathbf{M}_{\sigma}^{(t)}}{\partial \mathbf{m}_{\sigma,i}^{(t)}}$:

$$\frac{\partial \mathbf{M}_{\sigma}^{(t)}}{\partial \mathbf{m}_{\sigma,i}^{(t)}} = \begin{bmatrix} \mathbf{0} \\ \vdots \\ \mathrm{diag}\left(g'\left(\mathbf{m}_{\sigma,i}^{(t)}\right)\right) \\ \vdots \\ \mathbf{0} \end{bmatrix},$$

where the non-zero block $\mathrm{diag}\left(g'\left(\mathbf{m}_{\sigma,i}^{(t)}\right)\right)$ is at position i.

Compute $\frac{\partial \phi^{(t)}}{\partial \mathbf{M}_{\sigma}^{(t)}}$:

$$\frac{\partial \phi^{(t)}}{\partial \mathbf{M}_{\sigma}^{(t)}} = \mathrm{diag}\left(f'\left(\mathbf{W}^{(t)}\mathbf{M}_{\sigma}^{(t)}\right)\right)\mathbf{W}^{(t)}.$$

Therefore,

$$\frac{\partial \phi^{(t)}}{\partial \mathbf{m}_{\sigma,i}^{(t)}} = \mathrm{diag}\left(f'\left(\mathbf{W}^{(t)}\mathbf{M}_{\sigma}^{(t)}\right)\right)\mathbf{W}^{(t)(:,\,\mathcal{I}_i)}\,\mathrm{diag}\left(g'\left(\mathbf{m}_{\sigma,i}^{(t)}\right)\right),$$

where $\mathbf{W}^{(t)(:,\,\mathcal{I}_i)}$ denotes the columns of $\mathbf{W}^{(t)}$ corresponding to $\mathbf{m}_{\sigma,i}^{(t)}$.

The matrix $\frac{\partial \phi^{(t)}}{\partial \mathbf{m}_{\sigma,i}^{(t)}}$ is of size $p_{t+1} \times p_{i,t}$.

Since $|f'(x)| \leq C_f$, $|g'(x)| \leq C_g$, and $|(\mathbf{W}^{(t)})_{k,l}| \leq C_w$, the entries of $\frac{\partial \phi^{(t)}}{\partial \mathbf{m}_{\sigma,i}^{(t)}}$ are bounded by $C_w C_f C_g$.

Using the induced matrix 1-norm:

$$\left\| \frac{\partial \phi^{(t)}}{\partial \mathbf{m}_{\sigma,i}^{(t)}} \right\|_1 = \max_{1 \leq l \leq p_{i,t}} \sum_{k=1}^{p_{t+1}} \left| \left(\frac{\partial \phi^{(t)}}{\partial \mathbf{m}_{\sigma,i}^{(t)}} \right)_{k,l} \right|.$$

Each column l sums over k up to p_{t+1} entries, each bounded by $C_w C_f C_g$. Therefore,

$$\left\| \frac{\partial \phi^{(t)}}{\partial \mathbf{m}_{\sigma,i}^{(t)}} \right\|_1 \leq C_w C_f C_g p_{t+1}.$$

Thus, the Lipschitz constant $\alpha^{(t)}$ satisfies:

$$\alpha^{(t)} \leq C_w C_f C_g p_{t+1}.$$

$\square$

D SUPPLEMENTARY ANALYSES

D.1 GRAPH LIFTING EXAMPLE AND CURVATURE

As shown in Section D.3, the edge curvature distribution is non-trivially impacted by the graph lifting procedure. We explore this more in this section.

One possible explanation for the general positive distribution shift is the widenning of bottleneck regions as well as the addition of many nodes and edges in densely connected regions. We observe this qualitatively in Figure 3. We can see that the narrow path is now widened into two nodes instead of one. We also see that the clique regions gain a lot more nodes and edges than the path. The Ollivier-Ricci curvature becomes much more red, while the balanced Forman curvature and augmented Forman curvature maintain a higher number of small, negatively curved edges. This example is consistent with the observations from Figure 5.

In the augmented Forman curvature plot, we can see that the curvature for edge `DI` becomes more negative. While this edge may not propagate information very well due to its negative curvature, there are now many edges for information to flow around this edge. This provides some qualitative evidence that incorporating global structure into analysis of graph lifting could present an important future direction.

D.2 Graph Lifting and Weighted Curvature

The addition of many nodes and edges, as well as the shift in the curvature distribution makes direct comparisons between graphs and their corresponding lifts more challenging. While this widened bottleneck may alleviate oversquashing, measures like algebraic connectivity may not measure this effect since widened bottlenecks are counteracted by the addition of many nodes. Similarly, average curvature could be biased by the addition of many positively curved edges located in densely connected regions. We instead propose the *betweenness weighted curvature*

$$\text{wc} = \sum_{e \in E(G)} \text{bc}(e) \, \text{curv}(e), \tag{14}$$

where bc denotes betwenness centrality and curv denotes the Ollivier-Ricci curvature. This measure, originally in introduced by Münch (2022), places more weight on bottleneck edges which gives a weighted average of the curvature in the graph. This measure has the added benefit that it captures both local information propagation through curvature as well as global information propagation through betweenness centrality. To account for the influence of more positively weighted edges, we also consider the *negative betweenness weighted curvature*

$$\text{nwc} = \sum_{e \in E(G)} \text{bc}(e) \, \text{curv}(e) \chi_{\text{curv}(e)<0}, \tag{15}$$

where χ is the indicator function.

In Figure 4, we create a scatter plot of the weighted curvature of a graph and its corresponding clique complex: $(\text{wc}(G), \text{wc}(C))$. We can see that the weighted curvature generally becomes less negative after complex construction for Ollivier-Ricci and balanced Forman curvature, which may suggest that oversquashing is alleviated. However, the augmented Forman curvature makes the weighted curvature more negative for all graphs in the MUTAG dataset. This is consistent with Figure 5 which produced a small number of negatively curved edges for augmented Forman curvature. Interestingly, there is a strong linear correlation in each of the scatter plots. Further study is required to understand this trend.

D.3 Graph Lifting and Edge Curvature Distribution

Understanding the impact on information propagation of lifting a graph to a clique complex is crucial. Graph lifting often adds numerous edges and nodes, complicating theoretical analysis, as the structural differences between the relational structures corresponding to a graph and its clique complex can be significant. As a first step, we analyze the edge curvature distribution to gain empirical insight into the clique graph lift. Figure 5 shows the kernel density estimate (KDE) of edge curvature on the MUTAG dataset. Ollivier-Ricci curvature (Nguyen et al., 2023) increases uniformly, while balanced Forman curvature (Topping et al., 2022) generally increases with some edges becoming slightly more negative. A more pronounced shift is observed with augmented Forman curvature (Fesser & Weber, 2023), but the overall effect on information propagation remains unclear, as global graph structure is not captured in this analysis.

(a) ORC

(b) BFC

(c) AFC

Figure 3: Long dumbbell graph before and after lifting to its clique complex. Edges are colored based on their curvature: Ollivier Ricci curvature (top), balanced Forman curvature (middle), and augmented Forman curvature (bottom). Note that this uses Boundary, Co-boundary, Lower, and Upper relations. These are presented as edges for visual clarity.

(a) ORC WC   (b) ORC NWC

(c) BFC WC   (d) BFC NWC

(e) AFC WC   (f) AFC NWC

Figure 4: Weighted curvature scatter plots for MUTAG: Ollivier Ricci curvature (top), balanced Forman curvature (middle), and augmented Forman curvature (bottom); betweeness weighted curvature (left) and negative betweenness weighted curvature (right).

(a) ORC   (b) BFC   (c) AFC

Figure 5: Edge curvature distribution across all graphs in the MUTAG dataset: Ollivier-Ricci curvature (left), Balanced Forman curvature (middle), and Augmented Forman curvature (right).

Figure 6: NeighborsMatch Graph with 3 cliques of 5 nodes.

D.4 SYNTHETIC BENCHMARK: NEIGHBORSMATCH

The neighbors match experiment is a graph transfer task introduced by Alon & Yahav (2021) to test oversquashing in GNNs. We build on the implementation of Karhadkar et al. (2023) which adapts this benchmark to test rewiring. Each graph in the dataset is a path of cliques as shown in Figure 6. Each of the green nodes is assigned a distinct random label which is represented by a one-hot encoding. The root node, colored red, is also assigned a random label which is encoded using a one-hot encoding. The goal is for the neural network to learn which green node has the same one-hot encoding as the root node based solely on the output of the neural network at the root node.

In this experiment, we choose the number of nodes in each clique to be 5 to avoid excessively large clique complexes when performing graph lifting. Note that we see similar trends as in the RingTransfer experiment. However, this creates a much larger clique complex which may present a greater challenge for rewiring methods. We can see in Figure 7 that the Clique complex requires more rewiring iterations to achieve the same performance as its graph counterpart. This is likely related to the clique complex having many more nodes than the original graph. The original graph has 15 nodes and 32 edges, whereas the clique complex has 77 nodes and 764 edges.

Figure 7: Neighbors Match Rewiring experiment.

We also benchmark trees and trees with cycles attached to the leaves on the NeighborsMatch benchmark (see Figure 8). We note that a tree as an acyclic connected graph has Betti numbers $b_0 = 1$ (number of connected components) and $b_1 = 0$ (number of independent loops or cycles). On the other hand, a tree with n attached cycles has Betti numbers $b_0 = 1$ and $b_1 = n$. As such, the two structures are topologically very distinct when considered as 1-simplicial complexes: the tree is contractible and has a trivial fundamental group, while the tree with attached cycles is not contractible and has a non-trivial fundamental group corresponding to a free group with n generators. However,

their relational message passing structures are quite similar, and they demonstrate almost identical prototypical oversquashing trends (see Figure 9).

(a) Tree with no attached cycles   (b) Tree with attached cycles

Figure 8: Tree without and with attached cycles have very distinct topologies.

Figure 9: Neighbors Match Rewiring experiment on trees and trees with cycles, with increasing tree depths. "None" and "Clique" indicate whether the underlying graph has been lifted to its clique complex. "Line" and "Cycle" indicate whether lines or cycles are attached to the leaves.

D.5 REAL-WORLD BENCHMARK: GRAPH CLASSIFICATION (ZINC)

Complex	Rewiring	GIN	R-GIN	SIN	CIN++
None	FoSR	1.056 ± 0.010	1.485 ± 1.045	0.377 ± 0.005	0.374 ± 0.008
	None	0.405 ± 0.001	0.368 ± 0.003	0.370 ± 0.003	0.339 ± 0.005
Clique	FoSR	0.844 ± 0.011	0.420 ± 0.002	0.369 ± 0.003	0.379 ± 0.008
	None	0.409 ± 0.006	0.326 ± 0.003	0.331 ± 0.004	0.320 ± 0.006
Ring	FoSR	0.641 ± 0.011	0.232 ± 0.002	0.212 ± 0.003	0.211 ± 0.002
	None	0.332 ± 0.006	0.215 ± 0.003	0.223 ± 0.004	0.214 ± 0.003

Table 2: ZINC (MAE). Lower values are better.

In this section we test our methods on the ZINC dataset. Due to increased computational demands, we do not perform a hyperparameter sweep for this experiment and instead use the hyperparameters from the TUDataset experiments. We also test this on the ring lifting procedure. The ring lift consists of adding cycles in the graph as 2-cells in the complex as in Bodnar et al. (2021a). In this work,

we restrict the rings to have size at most 7. Structures like carbon rings are important in molecule datasets and prediction tasks can benefit from encoding larger rings in the lifting procedure. The results are presented in Table 2, where we see a clear improvement from the Ring lift across all models and rewirings tested for the ZINC dataset.

D.6 REAL-WORLD BENCHMARK: NODE CLASSIFICATION

In this section, we test the lifting and rewiring methods on node classification tasks. The results of this experiment are shown in Table 3. The rewirings considered included FoSR and the pruning algorithm introduced in Appendix D.8. The best results for each dataset and model are highlighted in gold. Notably, CIN++ benefited from lifting on all datasets tested.

(a) GIN

Complex	Rewiring	CORNELL	TEXAS	WISCONSIN	CITESEER	CORA
None	FoSR	41.6 ± 3.5	57.9 ± 2.7	49.6 ± 2.5	69.3 ± 0.8	82.4 ± 0.7
	Prune	**45.8 ± 4.0**	55.3 ± 4.9	48.1 ± 3.9		
	None	33.7 ± 3.3	60.0 ± 5.2	49.2 ± 2.9	**69.9 ± 0.8**	82.7 ± 0.9
Clique	FoSR	40.5 ± 3.5	**64.7 ± 5.3**	**56.5 ± 2.6**	68.6 ± 0.8	**84.4 ± 0.6**
	Prune	36.8 ± 2.6	61.6 ± 3.3	47.3 ± 3.7		
	None	44.7 ± 3.6	58.9 ± 2.3	55.8 ± 2.2	69.0 ± 0.8	83.2 ± 0.6

(b) RGIN

Complex	Rewiring	CORNELL	TEXAS	WISCONSIN	CITESEER	CORA
None	FoSR	54.2 ± 3.3	51.1 ± 3.3	**71.2 ± 2.4**	73.2 ± 0.6	81.9 ± 0.8
	Prune	52.6 ± 3.5	65.3 ± 3.2	68.5 ± 2.1		
	None	**59.5 ± 5.0**	**67.9 ± 3.2**	69.2 ± 2.8	72.2 ± 1.0	86.4 ± 0.7
Clique	FoSR	51.1 ± 3.8	62.1 ± 3.7	68.5 ± 3.3	**73.4 ± 1.0**	85.7 ± 0.6
	Prune	47.9 ± 2.3	64.2 ± 3.6	60.4 ± 2.0		
	None	50.5 ± 3.5	67.4 ± 3.4	68.5 ± 3.8	72.1 ± 0.9	**86.6 ± 0.6**

(c) SIN

Complex	Rewiring	CORNELL	TEXAS	WISCONSIN	CITESEER	CORA
None	FoSR	30.5 ± 3.5	34.7 ± 4.1	30.8 ± 3.2	54.3 ± 1.1	50.6 ± 2.1
	Prune	**44.7 ± 4.2**	55.8 ± 4.9	**53.1 ± 2.9**		
	None	42.6 ± 2.9	52.6 ± 4.3	50.4 ± 2.8	**66.9 ± 0.8**	**81.6 ± 0.8**
Clique	FoSR	36.3 ± 2.7	49.5 ± 3.0	47.3 ± 3.0	60.2 ± 1.2	68.2 ± 0.9
	Prune	40.5 ± 2.9	**56.3 ± 3.0**	51.5 ± 3.1		
	None	36.8 ± 2.9	55.8 ± 2.7	45.8 ± 3.8	63.2 ± 0.8	77.7 ± 0.7

(d) CIN++

Complex	Rewiring	CORNELL	TEXAS	WISCONSIN	CITESEER	CORA
None	FoSR	41.6 ± 2.0	55.3 ± 2.7	51.9 ± 4.1	44.9 ± 1.1	64.0 ± 1.8
	Prune	42.6 ± 3.5	49.5 ± 4.7	61.2 ± 2.3		
	None	43.2 ± 4.4	59.5 ± 2.6	54.2 ± 3.8	44.7 ± 1.7	63.6 ± 2.0
Clique	FoSR	47.9 ± 4.5	**66.3 ± 3.7**	64.6 ± 1.6	**64.4 ± 1.0**	78.6 ± 0.8
	Prune	**52.6 ± 3.5**	57.9 ± 3.2	63.8 ± 3.5		
	None	50.5 ± 2.7	65.3 ± 2.6	**66.2 ± 2.8**	63.2 ± 0.9	**79.4 ± 0.7**

Table 3: Node classification experiments.

D.7 ABLATION TESTS

We perform ablation studies to evaluate the effect of increasing the number of layers and rewiring iterations, as well as the effect of increasing the hidden dimensions and rewiring iterations, on the performance of CIN++ on the graph classification MUTAG dataset with FoSR rewiring. The results are presented in Tables 4 and 5. This analysis complements the experiments in Section 5.2 by testing on real world data. We note that the results of the ablation studies for the MUTAG dataset do not show a clear pattern across all rewiring iterations. However, we broadly see that rewiring

too much tends to hurt performance as they begin to oversmooth information. We also generally see that increasing the hidden dimension tends to improve performance. Due to time constraints and limited compute, we could only test this one one dataset and perform 10 trials for each set of hyperparamters. With more trials, the standard error will decrease and make the patterns more clear. Nguyen et al. (2023) and Fesser & Weber (2024) performed similar ablations and considered 100 trials for each set of hyperparameters. As in the experiments from Section 5 and the results reported in Rusch et al. (2023) and Nguyen et al. (2023), we expect that increasing the hidden dimension will improve classification accuracy. Similarly, increasing layers will initially improve performance until the model becomes too deep and experiences oversmoothing.

Lift	Rewire Iterations Layers	0	10	20	40
none	2	86.0 ± 2.9	85.0 ± 2.4	75.5 ± 1.7	85.5 ± 3.1
	4	87.5 ± 2.3	81.5 ± 2.6	81.0 ± 3.2	77.5 ± 2.9
	6	89.0 ± 2.7	82.5 ± 2.8	78.5 ± 3.7	80.0 ± 2.8
clique	2	85.5 ± 2.2	84.0 ± 1.8	81.0 ± 2.2	83.5 ± 3.0
	4	86.0 ± 1.9	78.5 ± 3.6	83.5 ± 2.9	79.0 ± 2.1
	6	84.5 ± 2.7	80.5 ± 2.0	82.5 ± 3.1	80.5 ± 2.3

Table 4: Impact of increasing the number of layers and rewiring iterations for FoSR on the MUTAG dataset with CIN++.

Lift	Rewire Iterations Hidden Dimension	0	10	20	40
none	16	81.5 ± 2.5	77.5 ± 3.0	78.5 ± 3.7	74.0 ± 1.8
	32	83.0 ± 2.9	77.5 ± 2.8	84.5 ± 2.6	82.0 ± 4.2
	64	85.0 ± 1.8	82.0 ± 2.1	83.0 ± 1.5	82.5 ± 2.4
clique	16	87.5 ± 3.0	79.5 ± 3.6	77.0 ± 2.6	79.0 ± 1.6
	32	83.5 ± 2.2	81.0 ± 2.2	80.5 ± 4.1	78.5 ± 3.4
	64	89.0 ± 2.6	81.0 ± 2.8	75.5 ± 3.8	83.5 ± 2.8

Table 5: Impact of increasing the hidden dimensions and rewiring iterations for FoSR on the MUTAG dataset with CIN++.

D.8 SIMPLEX PRUNING

In this section we analyze the impact of pruning edges in the influence graph. For the pruning algorithm, we consider removing the edge with the highest balanced Forman curvature. More precisely, we first convert the simplex into the influence graph and collapse multi-edges. Then we compute the balanced Forman curvature of each edge in the graph. Finally, we remove the edge with the highest curvature. This process is repeated 40 times. When removing an edge with higher arity $e = (\sigma, \tau, \delta)$, the entire higher-order edge is removed from the simplex. When an edge (σ, τ) with higher multiplicity is selected, all instances of that edge are removed from the graph.

The results of this experiment are shown in Table 6. In five out of eight experiments, the performance improves with some amount of pruning. Notably, GIN and CIN++ attain their highest performance with 40 pruning iterations. This suggests that rewiring algorithms that implement edge pruning and edge addition could be the most promising candidates for simplicial rewiring algorithms. This is in part due to the addition of many nodes and edges in the graph lifting procedure.

E ADDITIONAL EXPERIMENTAL DETAILS

For all experiments, we split the dataset randomly using 80% for training, 10% for validation, and 10% for testing. For classificiation, we use the cross entropy loss. Each experiment is run with early stopping using the validation set. Also, if there is no improvement for 10 iterations, the learning rate

| Complex | Prune Iterations
Model | 0 | 10 | 20 | 40 |
|---------|---------------------------|---|----|----|----|
| none | gin | 83.0 ± 3.1 | 87.5 ± 3.1 | 89.0 ± 3.1 | 86.0 ± 2.1 |
| | rgin | 81.5 ± 1.7 | 84.0 ± 2.8 | 91.0 ± 2.2 | 86.0 ± 1.2 |
| | sin | 88.5 ± 3.0 | 83.0 ± 2.5 | 86.0 ± 2.1 | 86.5 ± 2.5 |
| | cin++ | 85.0 ± 3.4 | 91.5 ± 1.8 | 88.0 ± 2.0 | 84.5 ± 2.5 |
| clique | gin | 83.0 ± 2.8 | 83.5 ± 1.7 | 82.5 ± 3.1 | 90.5 ± 1.4 |
| | rgin | 86.0 ± 2.3 | 81.0 ± 2.4 | 82.5 ± 2.5 | 85.0 ± 1.7 |
| | sin | 87.0 ± 3.2 | 79.5 ± 2.2 | 83.5 ± 2.2 | 84.5 ± 3.4 |
| | cin++ | 90.5 ± 2.2 | 84.5 ± 3.7 | 78.5 ± 3.2 | 91.0 ± 1.8 |

Table 6: MUTAG pruning experiment. Removing edges with highest balanced Forman Curvature.

is decreased. For all rewiring methods, we only consider adding edges and no edge removal. Each rewiring method is run for 40 iterations.

E.1 REAL-WORLD BENCHMARK: GRAPH CLASSIFICATION ON TUDATASET

To train and evaluate these method on real world data, we use the TUDataset (Morris et al., 2020). These graph classification tasks are commonly used to benchmark message passing neural networks. We have included summary statistics for MUTAG, NCI1, ENZYMES, PROTEINS, and IMDB-BINARY in Table 7.

Dataset	# Graphs	Classes	Avg. Nodes	Avg. Edges	Features
MUTAG	188	2	17.93	19.79	7
NCI1	4110	2	29.87	32.30	37
ENZYMES	600	6	32.63	62.14	3
PROTEINS	1113	2	39.06	72.82	3
IMDB-BINARY	1000	2	19.77	96.53	0

Table 7: Realworld Dataset Statistics.

E.2 SYNTHETIC BENCHMARK: RINGTRANSFER

The RingTransfer benchmark is a graph transfer task introduced by Bodnar et al. (2021a). Each graph in the dataset is a cycle graph of size $2k$ for some $k \geq 1$. A node i is randomly selected as the *root node* and the node $j = i + k(\bmod)2k$ on the opposite side of the cycle is designated as the feature node. All features are initialized to zero except for the feature node which contains a one-hot label among the 5 classes. The task is to then predict the label by reading the feature stored in the root node. This is shown for a cycle graph with 10 nodes in Figure 10, where the red node is the root node and the green node is the opposite node. Note that the size of the ring determines the number of layers necessary for a model to complete the task successfully. By increasing the size of the ring, we can test how depth and the task distance impact training performance. When increasing the ring size, the model depth is set to $2k + 1$ so that task information remains within the model's receptive field.

Parameter	Value
Number of Graphs	1000
Trials	10
Default Neighbors Match Cliques	3
Default Neighbors Match Clique Size	5
Default Ring Transfer Nodes	10

Table 8: Synthetic Parameters

Figure 10: Graph for the RingTransfer experiment.

In Tables 9 and 10, we provide all the results for Table 1 from Section 5. We note that in Table 1, we only report the Lift=None results for FoSR, and not AFR4 or SDRF, for fair comparison with Lift=Clique, where we only report the results for FoSR. We note that for $24/32 = 75\%$ of the results on ENZYME, MUTAG, NCI1, and PROTEINS, the best performing rewiring algorithms for Lift=None and Lift=Clique agree. In Figure 11, we compare model performance on graphs and the corresponding clique complexes, and note how graph, relational graph, and topological models all respond similarly to relational rewiring.

Figure 11: Comparison of model performance between graph representations (Lift=None) and corresponding clique complex representations (Lift=Clique). Each point corresponds to a specific model on a given dataset, comparing the performance with and without rewiring.

(a) **SGC**

Lift	Rewiring	ENZYMES	IMDB-B	MUTAG	NCI1	PROTEINS
None	None	18.3 ± 1.2	49.5 ± 1.5	64.5 ± 5.8	55.2 ± 1.0	62.2 ± 1.4
	AFR4	17.0 ± 1.7	48.6 ± 2.3	65.5 ± 3.5	54.4 ± 0.7	65.0 ± 1.5
	FoSR	18.2 ± 1.7	50.0 ± 1.8	52.5 ± 6.6	49.3 ± 0.8	64.8 ± 1.1
	SDRF	21.5 ± 1.6	49.7 ± 1.7	70.0 ± 2.6	52.4 ± 0.7	59.9 ± 1.4
Clique	None	14.5 ± 1.4	48.7 ± 2.2	70.0 ± 3.3	50.0 ± 1.3	59.9 ± 1.8
	AFR4	16.3 ± 1.0	> 1 day	62.5 ± 3.0	56.8 ± 0.8	59.1 ± 1.4
	FoSR	16.0 ± 1.4	47.8 ± 1.6	60.5 ± 5.8	50.0 ± 0.6	50.0 ± 4.3
	SDRF	16.8 ± 0.9	> 12 hours	69.5 ± 2.6	50.8 ± 0.7	54.9 ± 3.2

(b) **GCN**

Lift	Rewiring	ENZYMES	IMDB-B	MUTAG	NCI1	PROTEINS
None	None	32.2 ± 2.0	49.1 ± 1.4	71.0 ± 3.8	48.3 ± 0.6	73.1 ± 1.2
	AFR4	30.5 ± 2.5	48.6 ± 1.0	69.0 ± 2.7	48.4 ± 0.6	72.6 ± 1.5
	FoSR	27.2 ± 2.5	47.9 ± 1.0	83.0 ± 1.5	49.1 ± 0.8	75.8 ± 1.7
	SDRF	30.7 ± 1.5	46.7 ± 1.2	72.5 ± 3.5	48.9 ± 0.5	74.5 ± 1.5
Clique	None	30.7 ± 1.2	64.0 ± 3.1	67.0 ± 3.5	48.4 ± 0.4	69.9 ± 0.6
	AFR4	28.2 ± 2.2	> 1 day	66.5 ± 3.0	48.3 ± 0.6	72.5 ± 0.9
	FoSR	26.0 ± 2.1	65.5 ± 3.1	81.5 ± 2.9	48.8 ± 1.1	75.0 ± 1.4
	SDRF	30.2 ± 2.4	> 12 hours	72.0 ± 2.6	49.6 ± 0.6	72.1 ± 1.2

(c) **GIN**

Lift	Rewiring	ENZYMES	IMDB-B	MUTAG	NCI1	PROTEINS
None	None	47.2 ± 1.9	71.7 ± 1.5	83.0 ± 3.1	77.2 ± 0.5	70.6 ± 1.4
	AFR4	50.0 ± 2.7	69.7 ± 1.8	88.0 ± 2.4	77.9 ± 0.5	70.1 ± 0.9
	FoSR	33.3 ± 1.7	67.1 ± 1.5	75.5 ± 1.7	64.8 ± 0.9	72.2 ± 0.8
	SDRF	45.5 ± 2.1	65.8 ± 2.0	84.5 ± 2.8	76.3 ± 1.0	72.0 ± 1.5
Clique	None	44.0 ± 1.7	69.1 ± 1.2	83.0 ± 2.8	78.8 ± 0.7	68.7 ± 1.4
	AFR4	48.5 ± 2.2	> 1 day	82.5 ± 2.6	78.2 ± 0.6	67.6 ± 0.8
	FoSR	39.3 ± 1.4	70.8 ± 1.1	79.0 ± 2.6	63.5 ± 0.7	72.8 ± 1.2
	SDRF	41.7 ± 2.2	> 12 hours	73.5 ± 3.4	76.9 ± 0.7	69.2 ± 0.8

(d) **RGCN**

Lift	Rewiring	ENZYMES	IMDB-B	MUTAG	NCI1	PROTEINS
None	None	33.8 ± 1.6	47.6 ± 1.4	72.5 ± 2.5	53.2 ± 0.7	71.9 ± 1.6
	AFR4	36.2 ± 2.3	48.1 ± 1.0	69.5 ± 3.4	54.3 ± 1.1	72.8 ± 1.2
	FoSR	39.3 ± 1.6	68.0 ± 1.3	83.5 ± 1.8	62.4 ± 0.5	71.2 ± 1.8
	SDRF	42.5 ± 1.3	59.1 ± 2.2	76.0 ± 1.9	63.4 ± 1.1	75.6 ± 1.2
Clique	None	48.8 ± 1.2	71.0 ± 1.0	79.5 ± 1.7	72.9 ± 0.8	72.4 ± 1.6
	AFR4	42.3 ± 2.4	> 1 day	78.0 ± 2.1	75.0 ± 0.9	74.2 ± 1.2
	FoSR	39.3 ± 2.5	69.7 ± 1.5	79.0 ± 3.2	64.5 ± 0.6	70.9 ± 1.4
	SDRF	45.2 ± 1.5	> 12 hours	81.5 ± 3.8	68.5 ± 0.6	70.6 ± 1.6

(e) **RGIN**

Lift	Rewiring	ENZYMES	IMDB-B	MUTAG	NCI1	PROTEINS
None	None	46.8 ± 1.8	69.6 ± 1.6	81.5 ± 1.7	76.8 ± 1.1	70.8 ± 1.2
	AFR4	49.8 ± 2.0	73.2 ± 1.4	85.5 ± 2.0	77.0 ± 0.7	71.1 ± 1.3
	FoSR	48.2 ± 1.4	48.9 ± 2.9	85.5 ± 2.8	55.1 ± 2.6	72.4 ± 1.4
	SDRF	49.7 ± 2.1	50.4 ± 3.2	85.0 ± 1.7	52.8 ± 2.7	70.9 ± 1.2
Clique	None	50.8 ± 1.5	71.6 ± 0.9	86.0 ± 2.3	79.2 ± 0.6	71.5 ± 1.5
	AFR4	55.8 ± 2.5	> 1 day	85.0 ± 2.4	79.5 ± 0.4	71.0 ± 1.1
	FoSR	46.2 ± 1.4	69.0 ± 1.4	79.0 ± 2.4	72.7 ± 0.6	71.8 ± 1.7
	SDRF	45.5 ± 1.6	> 12 hours	83.0 ± 3.2	76.4 ± 0.9	69.1 ± 2.4

Table 9: Baseline and rewiring results for SGC, GCN, GIN, RGCN, and RGIN. Numbers highlighted in blue correspond to no rewiring and numbers highlighted in red are the best among the rewiring methods.

(a) **SIN**

Lift	Rewiring	ENZYMES	IMDB-B	MUTAG	NCI1	PROTEINS
None	None	47.5 ± 2.3	70.0 ± 1.4	88.5 ± 3.0	77.0 ± 0.6	70.2 ± 1.3
	AFR4	44.0 ± 2.5	71.0 ± 1.0	85.5 ± 1.7	76.4 ± 0.4	69.6 ± 1.3
	FoSR	45.7 ± 2.6	63.0 ± 2.7	85.5 ± 2.8	61.3 ± 2.4	73.2 ± 1.5
	SDRF	46.8 ± 2.1	62.8 ± 3.3	80.0 ± 2.1	54.7 ± 3.5	70.2 ± 1.8
Clique	None	51.0 ± 2.4	53.0 ± 1.9	87.0 ± 3.2	76.6 ± 1.3	66.9 ± 1.3
	AFR4	46.5 ± 1.2	> 1 day	83.5 ± 1.7	75.4 ± 0.7	69.6 ± 1.0
	FoSR	38.2 ± 2.2	64.0 ± 2.3	83.5 ± 2.7	65.2 ± 0.7	70.4 ± 1.2
	SDRF	44.8 ± 1.9	> 12 hours	80.0 ± 3.5	73.9 ± 0.4	68.7 ± 1.2

(b) **CIN**

Lift	Rewiring	ENZYMES	IMDB-B	MUTAG	NCI1	PROTEINS
None	None	50.0 ± 1.9	58.1 ± 4.0	86.5 ± 1.8	51.4 ± 2.5	70.7 ± 1.0
	AFR4	48.2 ± 1.9	58.8 ± 4.3	87.0 ± 2.4	53.3 ± 2.6	71.0 ± 1.4
	FoSR	49.0 ± 2.0	58.4 ± 2.7	81.0 ± 1.8	66.2 ± 2.0	70.9 ± 1.3
	SDRF	44.0 ± 1.8	55.3 ± 2.5	82.5 ± 3.0	59.0 ± 4.3	69.5 ± 1.5
Clique	None	49.8 ± 1.9	52.6 ± 2.4	85.5 ± 2.8	51.8 ± 2.3	70.7 ± 1.2
	AFR4	46.7 ± 1.3	> 1 day	86.5 ± 2.6	49.0 ± 0.5	69.2 ± 1.3
	FoSR	37.2 ± 1.8	68.1 ± 1.6	82.5 ± 3.0	66.9 ± 0.9	70.3 ± 0.8
	SDRF	42.0 ± 2.3	> 12 hours	78.5 ± 4.1	72.5 ± 0.8	67.9 ± 1.8

(c) **CIN++**

Lift	Rewiring	ENZYMES	IMDB-B	MUTAG	NCI1	PROTEINS
None	None	48.5 ± 1.9	66.6 ± 3.7	85.0 ± 3.4	60.8 ± 3.8	67.9 ± 1.9
	AFR4	51.0 ± 1.5	61.0 ± 4.1	91.0 ± 2.3	54.8 ± 3.7	70.7 ± 1.3
	FoSR	50.5 ± 1.8	56.0 ± 3.9	82.0 ± 2.5	64.8 ± 3.1	71.4 ± 1.4
	SDRF	45.8 ± 2.2	53.3 ± 2.3	87.5 ± 2.3	58.0 ± 4.4	69.3 ± 1.6
Clique	None	50.5 ± 2.1	62.8 ± 3.8	90.5 ± 2.2	61.5 ± 4.6	68.3 ± 1.3
	AFR4	52.7 ± 1.6	> 1 day	84.5 ± 3.3	60.6 ± 4.8	66.7 ± 1.3
	FoSR	43.0 ± 2.1	64.7 ± 1.5	78.5 ± 2.4	72.9 ± 0.5	71.9 ± 1.0
	SDRF	47.7 ± 1.5	> 12 hours	81.0 ± 3.0	76.8 ± 0.4	70.0 ± 1.7

Table 10: Baseline and rewiring results for SIN, CIN, and CIN++. Numbers highlighted in blue correspond to no rewiring and numbers highlighted in red are the best among the rewiring methods.

F   Implementation Details

F.1   Model Details

Each model consists of multiple graph-convolution layers. For more details about the convolutional layers, refer to Appendix F.5. The hyperparameters are presented in Table 11. For the real world experiments, each model has a neural network head consisting of a pooling operation followed by a two layer neural network with ReLU and dropout. The only exception to this is the SGC model which uses pooling and a single linear layer for the head. For the synthetic graph transfer experiments, there is no neural network head. Instead, the neural network just returns the feature of the root node at the final layer.

Parameter	Value
Layers	4
Hidden	64
Dropout	0.5
Pooling	mean
Multidimensional	False
Max Dimension	2

Table 11: Model Parameters

Parameter	Value
Optimizer	Adam
Batch Size	64
Learning Rate	0.001
Max Epochs	500
Stop Criteria	Validation
Stop Factor	1.01
Stop Patience	100

Table 12: Optimization Parameters

F.2   Hardware

Experiments were run on a system with Intel(R) Xeon(R) Gold 6152 CPUs @ 2.10GHz and NVIDIA GeForce RTX 2080 Ti GPUs.

F.3   Libraries

Code Use: Rewiring algorithms were adapted from the works of Karhadkar et al. (2023) and Fesser & Weber (2023). We used PyTorch Geometric Benchmarks for neural network hyperparameters. We build on the work of Giusti et al. (2024) for clique graph lifting and ring transfer. The real world graph datasets were sourced from the work of Morris et al. (2020). The SGC model was adapted from the work of Wu et al. (2019). For curvature implementations, we adapt code from Nguyen et al. (2023) for Ollivier-Ricci and Fesser & Weber (2023) for BFC and AFC.

Python libraries: Torch Geometric (Fey & Lenssen, 2019). Torch (Paszke et al., 2019). Scikit-Learn (Pedregosa et al., 2011). Ray (Moritz et al., 2018). Ray Tune (Liaw et al., 2018). Seaborn (Waskom, 2021). Matplotlib (Hunter, 2007). Networkx (Hagberg et al., 2008). Pandas pandas development team (2020). Numpy Harris et al. (2020). Numba (Lam et al., 2015). Gudhi (The GUDHI Project, 2020). POT (Flamary et al., 2021).

F.4   Graph Lifting

We lift graphs to clique complexes for relational message passing using Algorithm 3. We refer to this graph lifting as the *Clique* lift. Note that when representing a graph in a simplicial complex (eg

max_dim = 0), we simply add the $\uparrow$ relation to the edges. We refer to this graph lifting as the *None* lift. The resulting complex has trivial cell intersections and does not contain cell unions. This is outlined in Algorithm 2.

Algorithm 2 Graph Lifting (None)

Input: Unweighted undirected graph $G = (V, E)$, node features $(\mathbf{x}_v)_{v \in V}$
Output: Relational structure $(\mathcal{S}, \{R_{\text{upper}}\})$ and features $(\mathbf{x}_\sigma)_{\sigma \in \mathcal{S}}$
  1: Initialize $\mathcal{S} = \emptyset$ ▷ Entities
  2: **for** each node $v \in V$ **do**
  3:      $\sigma \leftarrow \{v\}$
  4:      $\mathcal{S} \leftarrow \mathcal{S} \cup \{\sigma\}$
  5:      $\mathbf{x}_\sigma \leftarrow \mathbf{x}_v$ ▷ Features
  6: **end for**
  7: Initialize $R_{\text{upper}} = \emptyset$ ▷ Relations
  8: **for** each edge $\{u, v\} \in E$ **do**
  9:      $\sigma, \tau \leftarrow u, v$
10:      $R_{\text{upper}} \leftarrow R_{\text{upper}} \cup \{(\sigma, \tau), (\tau, \sigma)\}$
11: **end for**
12: **return** $(\mathcal{S}, \{R_{\text{upper}}\})$ and $(\mathbf{x}_\sigma)_{\sigma \in \mathcal{S}}$

Algorithm 3 Graph Lifting (Complex)

Input: Unweighted undirected graph $G = (V, E)$, node features $(\mathbf{x}_v)_{v \in V}$, relation types $\mathcal{A} \subseteq$ {boundary, coboundary, lower, upper}, max simplex dimension max_dim
Output: Relational structure $(\mathcal{S}, \{R_r\}_{r \in \mathcal{A}})$ and features $(\mathbf{x}_\sigma)_{\sigma \in \mathcal{S}}$
  1: Initialize $\mathcal{S} = \emptyset$ ▷ Entities
  2: **for** each non-empty clique $\sigma \subseteq V$ with $|\sigma| \leq$ max_dim $+ 1$ **do**
  3:      $\mathcal{S} \leftarrow \mathcal{S} \cup \{\sigma\}$
  4:      $\mathbf{x}_\sigma \leftarrow \frac{1}{|\sigma|} \sum_{v \in \sigma} \mathbf{x}_v$ ▷ Features
  5: **end for**
  6: **for** each relation type $r \in \mathcal{A}$ **do**
  7:      Initialize $R_r = \emptyset$ ▷ Relations
  8:      **for** each entity $\sigma \in \mathcal{S}$ **do**
  9:          **if** $r =$ boundary **then**
10:             **for** each $\tau \in \mathcal{B}(\sigma)$ **do**
11:                $R_r \leftarrow R_r \cup \{(\sigma, \tau)\}$
12:             **end for**
13:          **else if** $r =$ coboundary **then**
14:             **for** each $\tau \in \mathcal{C}(\sigma)$ **do**
15:                $R_r \leftarrow R_r \cup \{(\sigma, \tau)\}$
16:             **end for**
17:          **else if** $r =$ lower **then**
18:             **for** each $\tau \in \mathcal{N}_\downarrow(\sigma)$ **do**
19:                $R_r \leftarrow R_r \cup \{(\sigma, \tau, \sigma \cap \tau)\}$
20:             **end for**
21:          **else if** $r =$ upper **then**
22:             **for** each $\tau \in \mathcal{N}_\uparrow(\sigma)$ **do**
23:                $R_r \leftarrow R_r \cup \{(\sigma, \tau, \sigma \cup \tau)\}$
24:             **end for**
25:          **end if**
26:      **end for**
27: **end for**
28: **return** $(\mathcal{S}, \{R_r\}_{r \in \mathcal{A}})$ and $(\mathbf{x}_\sigma)_{\sigma \in \mathcal{S}}$

F.5 LAYERS

We review the implementation of the layers and update rules used in the models considered in this work. In the following, $\mathcal{B}(\cdot)$, $\mathcal{C}(\cdot)$, $\mathcal{N}_\uparrow(\cdot)$, and $\mathcal{N}_\downarrow(\cdot)$ follow the definitions in Definition 2.3. Given a relational structure $\mathcal{R} = (\mathcal{S}, R_1, \dots, R_k)$, the notation $r \in \mathcal{R}$ refers to iterating over the relation indices $r = 1, 2, \dots, k$. For a binary relation $R \subseteq \mathcal{S} \times \mathcal{S}$, $\mathcal{R}(\cdot)$ denotes the incoming neighborhood of an entity, i.e., for any $\sigma, \tau \in \mathcal{S}$, if $\tau \in \mathcal{N}_R(\sigma)$, then $(\sigma, \tau) \in R$.

SGC   Following Wu et al. (2019), the SGC update rule is:

$$\mathbf{h}_\sigma^{(t+1)} = \sum_{r \in \mathcal{R}} \sum_{\tau \in \mathcal{N}_{R_r}(\sigma)} \mathbf{h}_\tau^{(t)},$$

where no learnable weights nor non-linearities are used, and signals are simply aggregated over neighborhoods.

GCN   Following Kipf & Welling (2017), the GCN update rule is:

$$\mathbf{h}_\sigma^{(t+1)} = \mathbf{f}\left(\mathbf{W}^{(t)}\left(\sum_{r \in \mathcal{R}} \sum_{\tau \in \mathcal{N}_{R_i}(\sigma) \cup \{\sigma\}} \frac{1}{\sqrt{\hat{d}_\sigma \hat{d}_\tau}} \mathbf{h}_\tau^{(t)}\right)\right),$$

where $\mathbf{f}$ is a point-wise non-linearity (e.g., ReLU), $\mathbf{W}^{(t)}$ is a learnable weight matrix, and $\hat{d}_\tau = \deg(\tau) + 1$.

GIN   Following Xu et al. (2019), the GIN update rule is:

$$\mathbf{h}_\sigma^{(t+1)} = \text{MLP}^{(t)}\left((1 + \epsilon)\mathbf{h}_\sigma^{(t)} + \sum_{r \in \mathcal{R}} \sum_{\tau \in \mathcal{N}_{R_i}(\sigma)} \mathbf{W}^{(t)}\mathbf{h}_\tau^{(t)}\right),$$

where $\mathbf{W}^{(t)}$ is a learnable weight matrix and ϵ is a learnable scalar. Here, MLP consists of two layer fully connected layers, each followed by a point-wise non-linearity (e.g., ReLU) and a BatchNorm layer.

R-GCN   Following Schlichtkrull et al. (2018), the R-GCN update rule is:

$$\mathbf{h}_\sigma^{(t+1)} = \mathbf{f}\left(\mathbf{W}_{\text{root}}^{(t)}\mathbf{h}_\sigma^{(t)} + \sum_{r \in \mathcal{R}} \sum_{j \in \mathcal{N}_r(i)} \frac{1}{|\mathcal{N}_r(i)|} \mathbf{W}_r^{(t)}\mathbf{h}_\sigma^{(t)}\right),$$

where $\mathbf{f}$ is a component-wise non-linearity (e.g., ReLU), and $\mathbf{W}_{\text{root}}^{(t)}$ and $\mathbf{W}_r^{(t)}$ are learnable weight matrices.

R-GIN   Following Xu et al. (2019) and Schlichtkrull et al. (2018), the R-GIN update rule is

$$\mathbf{h}_\sigma^{(t+1)} = \mathbf{W}_{\text{root}}^{(t)}\mathbf{h}_\sigma^{(t)} + \sum_{r \in \mathcal{R}} \text{MLP}_r^{(t)}\left((1 + \epsilon_r)\mathbf{h}_\sigma^{(t)} + \sum_{\tau \in \mathcal{N}_r(\sigma)} \mathbf{h}_\tau^{(t)}\right),$$

where $\mathbf{W}_{\text{root}}^{(t)}$ is a learnable weight matrix and each $\epsilon_\bullet$ is a learnable scalar. Here, each MLP consists of two layer fully connected layers, each followed by a point-wise non-linearity (e.g., ReLU) and a BatchNorm layer. and $\text{MLP}_r^{(t)}$ is a relation-specific MLP.

SIN Following Bodnar et al. (2021b), the SIN update rule is:

$$\mathbf{h}_\sigma^{(t+1)} = \mathrm{MLP}_{U,p}^{(t)} \begin{bmatrix} \mathrm{MLP}_{B,p}\left((1+\epsilon_{\mathcal{B}})\,\mathbf{h}_\sigma^{(t)} + \sum_{\delta \in \mathcal{B}(\sigma)} \mathbf{h}_\delta^{(t)} \right) \\ \mathrm{MLP}_{\uparrow,p}\left((1+\epsilon_\uparrow)\,\mathbf{h}_\sigma^{(t)} + \sum_{\delta \in \mathcal{N}_\uparrow(\sigma)} \mathbf{h}_\delta^{(t)} \right) \\ \textit{additionally, when rewiring:} \\ \boxed{\mathrm{MLP}_{R_{new},p}\left((1+\epsilon_{R_{new}})\,\mathbf{h}_\sigma^{(t)} + \sum_{\delta \in \mathcal{N}_{R_{new}}(\sigma)} \mathbf{h}_\delta^{(t)} \right)} \end{bmatrix},$$

where $[\cdot]$ denotes column-wise concatenation and each $\epsilon_\bullet$ is a learnable scalar. Here, each MLP consists of two layer fully connected layers, each followed by a point-wise non-linearity (e.g., ReLU) and a BatchNorm layer. and $\mathrm{MLP}_r^{(t)}$ is a relation-specific MLP.

CIN Following Bodnar et al. (2021a), the CIN update rule is:

$$\mathbf{h}_\sigma^{(t+1)} = \mathrm{MLP}_{U,p}^{(t)} \begin{bmatrix} \mathrm{MLP}_{B,p}\left((1+\epsilon_{\mathcal{B}})\,\mathbf{h}_\sigma^{(t)} + \sum_{\tau \in \mathcal{B}(\sigma)} \mathbf{h}_\tau^{(t)} \right) \\ \mathrm{MLP}_{\uparrow,p}\left((1+\epsilon_\uparrow)\,\mathbf{h}_\sigma^{(t)} + \sum_{\substack{\tau \in \mathcal{N}_\uparrow(\sigma) \\ \delta \in C(\sigma,\tau)}} \mathrm{MLP}_{M,p} \begin{bmatrix} \mathbf{h}_\tau^{(t)} \\ \mathbf{h}_\delta^{(t)} \end{bmatrix} \right) \\ \textit{additionally, when rewiring:} \\ \boxed{\mathrm{MLP}_{R_{new},p}\left((1+\epsilon_{R_{new}})\,\mathbf{h}_\sigma^{(t)} + \sum_{\delta \in \mathcal{N}_{R_{new}}(\sigma)} \mathbf{h}_\delta^{(t)} \right)} \end{bmatrix},$$

where $[\cdot]$ denotes column-wise concatenation, each $\epsilon_\bullet$ is a learnable scalar, and $\mathrm{MLP}_{M,p}$ is a single fully connected layer projecting back to the dimension of $\mathbf{h}_\sigma^{(t)}$. The other MLPs consists of two layer fully connected layers, each followed by a point-wise non-linearity (e.g., ReLU) and a BatchNorm layer. When entity unions aren't contained in the relational structure due to dimension constraints, we instead use $\mathbf{h}_\delta^{(t)} := \mathbf{0}$.

CIN++ Following Giusti et al. (2024), the CIN++ update rule is:

$$\mathbf{h}_\sigma^{(t+1)} = \mathrm{MLP}_{U,p}^{(t)} \begin{bmatrix} \mathrm{MLP}_{B,p}\left((1+\epsilon_{\mathcal{B}})\,\mathbf{h}_\sigma^{(t)} + \sum_{\tau \in \mathcal{B}(\sigma)} \mathbf{h}_\tau^{(t)} \right) \\ \mathrm{MLP}_{\uparrow,p}\left((1+\epsilon_\uparrow)\,\mathbf{h}_\sigma^{(t)} + \sum_{\substack{\tau \in \mathcal{N}_\uparrow(\sigma) \\ \delta \in C(\sigma,\tau)}} \mathrm{MLP}_{M,p} \begin{bmatrix} \mathbf{h}_\tau^{(t)} \\ \mathbf{h}_\delta^{(t)} \end{bmatrix} \right) \\ \mathrm{MLP}_{\downarrow,p}\left((1+\epsilon_\downarrow)\,\mathbf{h}_\sigma^{(t)} + \sum_{\substack{\tau \in \mathcal{N}_\downarrow(\sigma) \\ \delta \in B(\sigma,\tau)}} \mathrm{MLP}_{M',p} \begin{bmatrix} \mathbf{h}_\tau^{(t)} \\ \mathbf{h}_\delta^{(t)} \end{bmatrix} \right) \\ \textit{additionally, when rewiring:} \\ \boxed{\mathrm{MLP}_{R_{new},p}\left((1+\epsilon_{R_{new}})\,\mathbf{h}_\sigma^{(t)} + \sum_{\delta \in \mathcal{N}_{R_{new}}(\sigma)} \mathbf{h}_\delta^{(t)} \right)} \end{bmatrix},$$

where $[\cdot]$ denotes column-wise concatenation, each $\epsilon_\bullet$ is a learnable scalar. $\text{MLP}_{M,p}$ and $\text{MLP}_{M',p}$ are single fully connected layers projecting back to the dimension of $\mathbf{h}_\sigma^{(t)}$. The other MLPs consists of two layer fully connected layers, each followed by a point-wise non-linearity (e.g., ReLU) and a BatchNorm layer. When entity unions or intersections aren't contained in the relational structure due to dimension constraints, we instead use $\mathbf{h}_\delta^{(t)} := \mathbf{0}$.

G HIGHER-ORDER GRAPHS ARE RELATIONAL STRUCTURES

Although we do not focus on higher-order graph neural networks (k-GNNs) (Morris et al., 2019) in this work, we briefly illustrate how they naturally fit into our relational framework. In k-GNNs, graphs are lifted into higher-order graphs following the procedure outlined in Algorithm 4. In our framework, these higher-order graphs can be viewed as relational structures, where each entity represents a set of nodes of size $n = k$. The relations between these entities are of two types: *local* relations, which connect entities differing by one node and where the differing nodes are adjacent in the original graph, and *global* relations, where the differing nodes are not adjacent in the original graph. From this perspective, k-GNN message passing on higher-order graphs is equivalent to R-GCN message passing on the corresponding relational structures. We leave the theoretical and empirical study of these relational structures for future work.

Algorithm 4 Graph Lifting (Higher-Order Graph)

Input: Unweighted undirected graph $G = (V, E)$, node features $(\mathbf{x}_v)_{v \in V}$, entity size n
Output: Relational structure $(\mathcal{S}, \{R_{\texttt{local}}, R_{\texttt{global}}\})$ and features $(\mathbf{x}_\sigma)_{\sigma \in \mathcal{S}}$

1: Initialize $\mathcal{S} = \emptyset$ ▷ Entities
2: **for** each subset $\sigma \subseteq V$ with $|\sigma| = n$ **do**
3:     $\mathcal{S} \leftarrow \mathcal{S} \cup \{\sigma\}$
4:     $\mathbf{x}_\sigma \leftarrow \frac{1}{|\sigma|} \sum_{v \in \sigma} \mathbf{x}_v$ ▷ Features
5: **end for**
6: Initialize $R_{\texttt{local}} = \emptyset, R_{\texttt{global}} = \emptyset$ ▷ Relations
7: **for** each pair $\sigma, \tau \in \mathcal{S}$ **do**
8:     **if** $|\sigma \cap \tau| = n - 1$ **then**
9:         Let $\{s_\sigma\} = \sigma \setminus \tau$ and $\{t_\tau\} = \tau \setminus \sigma$
10:        **if** $\{s_\sigma, t_\tau\} \in E$ **then**
11:            $R_{\texttt{local}} \leftarrow R_{\texttt{local}} \cup \{(\sigma, \tau), (\tau, \sigma)\}$
12:        **else**
13:            $R_{\texttt{global}} \leftarrow R_{\texttt{global}} \cup \{(\sigma, \tau), (\tau, \sigma)\}$
14:        **end if**
15:    **end if**
16: **end for**
17: **return** $(\mathcal{S}, \{R_{\texttt{local}}, R_{\texttt{global}}\})$ and $(\mathbf{x}_\sigma)_{\sigma \in \mathcal{S}}$

H WORKED EXAMPLE: RELATIONS AND OPERATORS FOR A PATH GRAPH

In this section, we illustrate, via a small example, the definitions of a simplicial complex, relational structure, and influence graph.

Simplicial Complex. Consider the simplicial complex $\mathcal{K}$ corresponding to the path graph $i - j - k$. The complex $\mathcal{K}$ consists of

- 0-simplices (vertices): $\{i\}, \{j\}, \{k\}$, and

- 1-simplices (edges): $\{i, j\}, \{j, k\}$.

Relational Structure. The relations on the set of entities $\mathcal{S} = \mathcal{K}$ are as follows:

- Relation R_1 (identity):

$$R_1 = \{(\sigma) \mid \sigma \in \mathcal{K}\}$$
$$= \{(\{i\}), (\{j\}), (\{k\}), (\{i,j\}), (\{j,k\})\}.$$

- Relation R_2 (boundary):

$$R_2 = \{(\sigma, \tau) \mid \sigma \in \mathcal{K}, \tau \in \mathcal{B}(\sigma)\}$$
$$= \left\{ \begin{array}{l} (\{i,j\}, \{i\}), (\{i,j\}, \{j\}), \\ (\{j,k\}, \{j\}), (\{j,k\}, \{k\}) \end{array} \right\}.$$

- Relation R_3 (co-boundary):

$$R_3 = \{(\sigma, \tau) \mid \sigma \in \mathcal{K}, \tau \in \mathcal{C}(\sigma)\}$$
$$= \left\{ \begin{array}{l} (\{i\}, \{i,j\}), (\{j\}, \{i,j\}), \\ (\{j\}, \{j,k\}), (\{k\}, \{j,k\}) \end{array} \right\}.$$

- Relation R_4 (lower adjacency):

$$R_4 = \{(\sigma, \tau, \delta) \mid \sigma \in \mathcal{K}, \tau \in \mathcal{N}_\uparrow(\sigma), \delta = \sigma \cap \tau\}$$
$$= \left\{ \begin{array}{l} (\{i,j\}, \{j,k\}, \{j\}), \\ (\{j,k\}, \{i,j\}, \{j\}) \end{array} \right\}.$$

- Relation R_5 (upper adjacency):

$$R_5 = \{(\sigma, \tau, \delta) \mid \sigma \in \mathcal{K}, \tau \in \mathcal{N}_\downarrow(\sigma), \delta = \sigma \cup \tau\}$$
$$= \left\{ \begin{array}{l} (\{i\}, \{j\}, \{i,j\}), (\{j\}, \{i\}, \{i,j\}), \\ (\{j\}, \{k\}, \{j,k\}), (\{k\}, \{j\}, \{j,k\}) \end{array} \right\}.$$

If we assume, for the sake of example, that all relations are involved in message passing, this gives rise to the relational structure

$$\mathcal{R}(\mathcal{K}) = (\mathcal{S}, R_1, R_2, R_3, R_4, R_5).$$

Of course, one could restrict message passing to a smaller subset of the relations $\{R_1, R_2, R_3, R_4, R_5\}$.

Shift Operators. In what follows, we index the entries in matrices as follows:

Index	Simplex
1	$\{i\}$
2	$\{j\}$
3	$\{k\}$
4	$\{i,j\}$
5	$\{j,k\}$

For instance, the 14th entry in a matrix would correspond to the connection $(\{i\}, \{i,j\})$. We assume, for the sake of example, that all shift operators correspond to the adjacency tensors, i.e., the entries with indices corresponding to entities that are related by a relation are equal to 1, and the remaining entries are all 0. Of course, normalized adjacency tensors and others could be similarly considered. As such, we get the following:

- Shift operators $\mathbf{A}^{R_1}$ and $\tilde{\mathbf{A}}^{R_1}$ (identity relation):

$$\mathbf{A}^{R_1} = \tilde{\mathbf{A}}^{R_1} = \begin{pmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{pmatrix}.$$

→ Boundary → Co-boundary → Lower → Upper

(a) Simplicial Complex $\mathcal{K}$ (b) Adjacency Relations on $\mathcal{K}$

(c) Influence Graph $\mathcal{G}(\mathcal{S}, \mathbf{B})$

Figure 12: (a) The simplicial complex $\mathcal{K}$ consisting of nodes i, j, k, and edges $\{i,j\}$, $\{j,k\}$. (b) The adjacency relations on $\mathcal{K}$, showing boundary, co-boundary, lower, and upper relations. (c) The influence graph $\mathcal{G}(\mathcal{S}, \mathbf{B})$ representing the information flow between the entities.

- Shift operators $\mathbf{A}^{R_2}$ and $\tilde{\mathbf{A}}^{R_2}$ (boundary relation):

$$\mathbf{A}^{R_2} = \tilde{\mathbf{A}}^{R_2} = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 \end{pmatrix}.$$

- Shift operators $\mathbf{A}^{R_3}$ and $\tilde{\mathbf{A}}^{R_3}$ (co-boundary relation):

$$\mathbf{A}^{R_3} = \tilde{\mathbf{A}}^{R_3} = \begin{pmatrix} 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}.$$

- Shift operator $\mathbf{A}^{R_4}$ and $\tilde{\mathbf{A}}^{R_4}$ (lower adjacency relation):
  The non-zero entries of $\mathbf{A}^{R_4}$ are

$$\mathbf{A}^{R_4} : \begin{cases} \mathbf{A}^{R_4}_{\{i,j\},\{j,k\},\{j\}} = 1 \\ \mathbf{A}^{R_4}_{\{j,k\},\{i,j\},\{j\}} = 1 \end{cases}.$$

  The aggregated influence for $\tilde{\mathbf{A}}^{R_4}$:

$$\tilde{\mathbf{A}}^{R_4} = \begin{pmatrix} 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 \end{pmatrix}.$$

- Shift Operator $\mathbf{A}^{R_5}$ and $\tilde{\mathbf{A}}^{R_5}$ (upper adjacency relation):

The non-zero entries of $\mathbf{A}^{R_5}$ are

$$\mathbf{A}^{R_5} : \begin{cases} \mathbf{A}^{R_5}_{\{i\},\{j\},\{i,j\}} = 1 \\ \mathbf{A}^{R_5}_{\{j\},\{i\},\{i,j\}} = 1 \\ \mathbf{A}^{R_5}_{\{j\},\{k\},\{j,k\}} = 1 \\ \mathbf{A}^{R_5}_{\{k\},\{j\},\{j,k\}} = 1 \end{cases} .$$

The aggregated influence for $\tilde{\mathbf{A}}^{R_5}$:

$$\tilde{\mathbf{A}}^{R_5} = \begin{pmatrix} 0 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 \end{pmatrix}$$

.

Lastly, the aggregated influence matrix $\tilde{\mathbf{A}}$ arising from including all relations R_1 to R_5:

$$\begin{aligned} \tilde{\mathbf{A}} &= \tilde{\mathbf{A}}^{R_1} + \tilde{\mathbf{A}}^{R_2} + \tilde{\mathbf{A}}^{R_3} + \tilde{\mathbf{A}}^{R_4} + \tilde{\mathbf{A}}^{R_5} \\ &= \begin{pmatrix} 1 & 1 & 0 & 2 & 0 \\ 1 & 1 & 1 & 2 & 2 \\ 0 & 1 & 1 & 0 & 2 \\ 1 & 2 & 0 & 1 & 1 \\ 0 & 2 & 1 & 1 & 1 \end{pmatrix} . \end{aligned}$$

The maximum row sum of $\tilde{\mathbf{A}}$ is $\gamma = 7$, giving rise to the augmented adjacency matrix

$$\begin{aligned} \mathbf{B} &= \gamma \mathbf{I} + \tilde{\mathbf{A}} \\ &= \begin{pmatrix} 8 & 1 & 0 & 2 & 0 \\ 1 & 8 & 1 & 2 & 2 \\ 0 & 1 & 8 & 0 & 2 \\ 1 & 2 & 0 & 8 & 1 \\ 0 & 2 & 1 & 1 & 8 \end{pmatrix} . \end{aligned}$$

The influence graph is finally constructed from $\mathbf{B}$ as follows:

- Nodes: Each simplex in $\mathcal{K}$.

- Edges: For each (σ, τ), if $\mathbf{B}_{\sigma,\tau} > 0$, there is a directed edge from τ to σ with weight $\mathbf{B}_{\sigma,\tau}$.

I COMPUTATIONAL COMPLEXITY

It is well known in the TDL community that topological methods can require much more computation due to directly modeling higher order structures. This typically incurs combinatorial complexity in the algorithms. A common strategy is to truncate the dimensionality of the objects considered to maintain polynomial complexity. Another option is to use a sparse complex construction method. As noted earlier, more work is needed in understanding graph lifting and which method is preferred in a given setting. In this section, we analyze the running times required for TDL with clique complexes.

I.1 CLIQUE COMPLEX SIZE

The exact computational requirements of performing clique lifting depend on the graph in consideration. However, there are a few cases where we can provide precise estimates.

Clique lifting of dimension 1 In this case, the highest dimensional cell is the edge cell. Consider a graph $G = (V, E)$ with $|V| = n$ nodes, $|E| = m$ edges, and maximum node degree d. The corresponding clique complex $\mathcal{K}$ has:

1. $n + m$ cells

2. $2m$ boundary relations

3. $2m$ co-boundary relations

4. $\sum_{v \in V} 2\binom{\deg(v)}{2} \leq 2md$ lower relations

5. $2m$ upper relations

Note that we need to count the upper and lower relations for each direction of the edge so we get a factor of 2.

Clique lifting of dimension 2 If G has $|F| = k$ faces, then the clique $\mathcal{K}$ contains

1. $n + m + k$ cells

2. $2m + 3k$ boundary relations

3. $2m + 3k$ co-boundary relations

4. $\sum_{v \in V} 2\binom{\deg(v)}{2} + \sum_{\sigma \in F} \sum_{\tau \in F} \chi(\sigma \cap \tau)$ lower relations

5. $2m + 6k$ upper relations

In dense graphs, the number of lower relations can be very large. Depending on the dataset under consideration, one may opt for a model like CIN where lower relations are ignored.

Complete Graph We can provide precise estimates when we have a complete graph. This can help us understand performance in the limiting case that the graph is very dense or contains large cliques. Consider a complete graph K_n on n nodes and the cell $\sigma = \{1, 2, ..., m\}$ for $m \geq 1$. This is the boundary of $n - m$ cells and is the co-boundary of m cells. It is upper adjacent to $(n - m)m$ cells. It is lower adjacent to $m(n - m)$ cells. Counting all the cells of dimension $d = m - 1$, we get

1. $n_d = \binom{n}{d+1}$ cells

2. $n_d(n - d - 1)$ boundary relations (if $d > 0$)

3. $n_d(d + 1)$ co-boundary relations

4. $n_d(d + 1)(n - d - 1)$ lower relations (if $d > 0$)

5. $n_d(d + 1)(n - d - 1)$ upper relations

, where the $d > 0$ condition comes from the fact that there isn't a cell $\{\}$ in the complex so there aren't lower or co-boundary relations for zero dimensional cells. Computing the total number of cells gives $2^n - 1$. The total number of edges is

$$-n - n(n-1) + \sum_{d=0}^{n-1} \binom{n}{d+1}(n + (n - d - 1)(d + 1))$$

$$= \frac{n}{2}(2^n n + 2^n - 2) - n^2 = O(n^2 2^n)$$

This bound is prohibitive for large n. If we restrict to two dimensional cells, we get the total number of cells is

$$n + \binom{n}{2} + \binom{n}{3} = \frac{n^3 + 5n}{6} = O(n^3)$$

and the total number of relations is

$$\frac{n}{6}(-24 + 41n - 24n^2 + 7n^3) = O(n^4)$$

I.2 CLIQUE LIFT COMPLEXITY

Our main use case is computing complexes for two dimensional cells. In this section we compute the complexity of clique lifting. Note that computing the 0 and 1 dimensional cells is fast since they are encoded directly in the graph as nodes and edges, respectively. The boundary and coboundary relations encode edge indcidence to a node and are also easily obtained. The upper relations on the 0-cells are just the edges in the graph. The lower relations on the edges are all pairs of edges in the graph sharing a node. This can be efficiently computed by storing the edge incidences in a dictionary in $O(m)$ time and then computing all pairs in the neighborhood of each node. This requires $O\left(n + m + \sum_{v \in V} \binom{\deg(v)}{2}\right)$ computation. For connected graphs, this reduces to $O(md)$.

Computing the two dimensional cells can also be done efficiently. By storing the neighborhood dictionary, we just need to loop over pairs of edges $(u, v), (u, w)$ in the neighborhood of u and check if $(v, w) \in G$. This can be done in $O\left(\sum_{v \in V} \binom{\deg(v)}{2}\right) = O(md)$ time. The boundary, co-boundary, and upper adjacencies can be easily computed in $O(k)$ time. The lower adjacency computation for 2-cells can require more time. For a given edge (u, v), consider the 2-cells σ such that $(u, v) \prec \sigma$. Each pair of these gets a lower adjacency and can be efficiently added. This results in

$$O\left(\sum_{e \in E} \binom{|\mathcal{C}(e)|}{2}\right)$$

computation. We can bound this as

$$\sum_{e \in E} \binom{|\mathcal{C}(e)|}{2} \leq \frac{1}{2} \sum_{e \in E} |\mathcal{C}(e)|^2 \leq \frac{1}{2} \max_e |\mathcal{C}(e)| \sum_{e \in E} |\mathcal{C}(e)| = \frac{3k}{2} \max_e |\mathcal{C}(e)|$$

and we get the complexity bound

$$O\left(k \max_e |\mathcal{C}(e)|\right)$$

This is analogous to the result with lower adjacency of 1-cells that relied on the maximum node degree in G. The maximum co-boundary degree $\max_e |\mathcal{C}(e)|$ is the analogue of node degree for 2-cells. So, the clique complex construction up to dimension 2 has time complexity:

$$O(md + k \max_e |\mathcal{C}(e)|)$$

I.3 RUNNING TIMES

Lift	ENZYMES	IMDB-B	MUTAG	NCI1	PROTEINS	ZINC	MOLHIV
None	2.12	3.51	1.60	4.70	2.99	63.47	12.92
Clique	6.89	45.04	1.22	20.26	16.07	94.70	128.18
Ring	11.81	101.98	1.48	26.62	30.73	106.37	169.45

Table 13: Time to perform graph lifting and download datasets from Torch Geometric.