An Axiomatic Theory of Provably-Fair Welfare-Centric Machine Learning

Abstract

We address an inherent difficulty in welfare-theoretic fair ML, propose an alternative, and study the resulting computational and statistical learning questions. Welfare metrics quantify *overall wellbeing* across a *population* of *groups*, and welfare-based objectives and constraints have recently been proposed to incentivize *fair ML methods* to satisfy their diverse needs. However, many ML problems are cast as *loss minimization*, rather than *utility maximization* tasks, thus requiring non-trivial modeling to construct *utility functions*. We define a complementary metric, termed *malfare*, measuring overall *societal harm*, with axiomatic justification via the standard axioms of cardinal welfare, and cast fair ML as *malfare minimization* over the *risk values* (expected losses) of each group. Surprisingly, the axioms of cardinal welfare (malfare) dictate that this is not equivalent to simply defining utility as negative loss and maximizing welfare. Building upon these concepts, we define *fair-PAC learning*, where a fair PAC-learner is an algorithm that learns an ε-δ malfare-optimal model, with bounded sample complexity, for *any data distribution* and (axiomatically justified) malfare concept. We show conditions under which many standard PAC-learners may be converted to fair-PAC learners. This places fair-PAC learning on firm theoretical ground, as it yields statistical, and in some cases computational, efficiency guarantees for many well-studied machine-learning models, and is also practically relevant, as it democratizes fair ML by providing concrete training algorithms and rigorous generalization guarantees.

1 Introduction

It is now well-understood that contemporary ML systems exhibit differential accuracy across gender and race for tasks like facial recognition [4, 6, 5], in medical settings [14, 2], and many others. This exacerbates existing inequality, as facial recognition in policing yields disproportionate false-arrest rates, and medical ML yields disproportionate health outcomes. Welfare-centric ML methods encode both *accuracy* and *fairness* in a *welfare function* defined on a *set of groups*, and optimize or constrain welfare to *learn fairly*. This addresses differential accuracy and bias issues across groups by ensuring that (1) each group is *seen* and *considered* during training, and (2) an outcome is incentivized that is desirable overall, ideally according to some mutually-agreed-upon welfare function. Unfortunately, welfare based metrics require a notion of (positive) utility, and we argue that this is not natural to many ML tasks, where we instead *minimize* some *negatively connoted* loss value. We thus define a complementary measure to *welfare*, which we term *malfare*, measuring societal harm (rather than wellbeing). In particular, malfare arises naturally when one applies the standard *axioms of cardinal welfare* to *loss* rather than *utility*. We then cast fair ML as a direct *malfare minimization* problem.

Perhaps surprisingly, minimizing a *malfare function* is *not equivalent* to maximizing some *welfare function*, while taking utility to be negative loss (except in the trivial cases of egalitarian and utilitarian

malfare). This is because nearly every function satisfying the standard axioms of cardinal welfare requires *nonnegative* inputs, and it is in general impossible to contort a loss function into a utility function while satisfying this requirement. For example, while minimizing the 0-1 loss, which simply counts the number of mistakes a classifier makes, is isomorphic to maximizing the 1-0 gain, which counts number of correct classifications, minimizing some *malfare function* defined on 0-1 loss over groups *is not* in general equivalent to maximizing any *welfare function* defined on 1-0 gain.

Building upon these concepts, we develop a notion of generic fair machine learning, termed *fair-PAC* (FPAC) *learning*, where the goal is to learn models for which finite training samples may guarantee high-probability near-optimal malfare, for *any* (axiomatically justified) malfare concept. This definition extends Valiant's [26] classic PAC-learning formalization, and we show that, with appropriate modifications, many (standard) PAC-learners may be converted to FPAC learners. In particular, we show via a constructive polynomial reduction that *realizable FPAC-learning* reduces to *realizable PAC-learning*. Furthermore, we show, non-constructively, that for learning problems where PAC-learnability implies uniform convergence, it is equivalent to FPAC-learnability. We also show that when training is possible via *convex optimization*, then ε-δ training malfare-optimal models, like risk-optimal models, requires polynomial time. We briefly summarize our contributions below.

1. We derive in § 2 the *malfare concept*, extending welfare to measure *negatively-connoted* attributes, and show that *malfare-minimization* naturally generalizes *risk-minimization* to produce *fairness-sensitive* machine-learning objectives that consider multiple protected groups.
2. Section 3 extends PAC-learning to FPAC-learning, where we consider minimization of malfare rather than risk (expected loss) objectives. Both PAC and FPAC learning are parameterized by a *learning task* (model space and loss function), and we explore the resulting learnability hierarchy.
3. We show that for many loss functions, PAC and FPAC learning are *statistically equivalent*, and convexity conditions for *computationally efficient* PAC-learning are also sufficient for FPAC-learning.

1.1 Related Work

Constraint-based notions of algorithmic fairness have risen to prominence in fair ML, with the potential to ensure demographic-parity (e.g., equality of *opportunity* or *outcome*) thus correcting for some forms of data or algorithmic bias. While noble in intent and intuitive by design, fairness by demographic-parity constraints has several prominent flaws: most notably, several popular parity constraints are mutually unsatisfiable [13], and their constraint-based formulation inherently puts *accuracy* and *fairness* at odds, where additional *tolerance parameters* are required to strike a balance between the two. Furthermore, recent works [11, 12] have shown that *welfare* and even *disadvantaged group utility* can be harmed by such fairness constraints, calling into question whether demographic parity constraints even benefit those they purport to aid.

Perhaps in response to these issues, some recent work has trended toward welfare-based fairness-concepts [11, 20], wherein both *accuracy* and *fairness* are encoded in a *welfare function* defined on a *group of subpopulations*. Welfare is then directly optimized [20] or constrained [24, 10] to promote *fair learning* across *all groups*. Perhaps the most similar to our work is a method of [11], wherein they *directly maximize* empirical welfare over linear classifiers; however as with other previous works, an appropriate utility function must be selected, which we avoid by instead using malfare. We argue that *empirical welfare maximization* is an effective strategy when an appropriate and natural measure of *utility* is available, but in machine learning contexts, there is no "correct" or clearly neutral way to convert loss to utility. We avoid this issue by working directly with malfare and loss.

Our most poignant contrast to existing work is the *Seldonian learner* [25] framework, which essentially extends PAC-learning to learning problems with both *constraints* and *arbitrary nonlinear objectives*. We argue that this generality is harmful to the utility of the concept as a mathematical or practical object, as nearly any ML problem can be posed as a constrained nonlinear optimization task. The utility in FPAC learning is that it is sophisticated enough to handle fairness issues, with a particular axiomatically justified objective, but remains simple enough to study as a mathematical object, leading to informative reductions between various PAC and FPAC learnable classes.

2 Quantifying Population-Level Sentiment

A generic *population mean* function $\mathrm{M}(\mathcal{S}; \boldsymbol{w})$ summarizes some *sentiment vector* $\mathcal{S} \in \mathbb{R}_{0+}^g$, weighted by *probability vector* $\boldsymbol{w} \in \mathbb{R}_+^g$ s.t. $\|\boldsymbol{w}\|_1 = 1$, across a population of g groups. When $\mathcal{S}$ measures a

desirable quantity, generally termed *utility*, the population mean is a measure of *cardinal welfare*
[16], and thus quantifies overall *wellbeing*. We also consider the inverse-notion of *illbeing*, termed
malfare, in terms of an *undesirable* $\mathcal{S}$, generally *loss* or *risk*, which naturally extends the concept. We
show an equivalent *axiomatic justification* for malfare, and argue that its use is more natural in many
situations, particularly when considering or optimizing *loss functions*.

Definition 2.1 (Population Means: Welfare and Malfare). A *population mean* function $\mathrm{M}(\mathcal{S}; \boldsymbol{w})$
measures the *overall sentiment* of a population, given *sentiment vector* $\mathcal{S}$ and *probability vector* $\boldsymbol{w}$. If
$\mathcal{S}$ denotes a desirable quantity (e.g., utility), we call $\mathrm{M}(\mathcal{S}; \boldsymbol{w})$ a *welfare function*, written $\mathrm{W}(\mathcal{S}; \boldsymbol{w})$,
and inversely, if it is undesirable (e.g., disutility, loss, or risk), we call $\mathrm{M}(\mathcal{S}; \boldsymbol{w})$ a *malfare function*,
written $\mathrm{\Lambda\!\!\Lambda}(\mathcal{S}; \boldsymbol{w})$.

For now, think of the term *population mean* as signifying that an entire population, with diverse
and subjective desiderata, is considered and summarized, rather than a single perspective (i.e., some
$\mathcal{S}_i$). As we introduce axioms and show consequent properties, the appropriateness of the term shall
become more apparent. We use the term *sentiment* to refer to $\mathcal{S}$ with neutral connotation, but when
discussing welfare or malfare, we often refer to $\mathcal{S}$ as *utility* or *risk*, respectively, as in these cases, $\mathcal{S}$
describes a well-understood preëxisting concept. We shall see that often population means on utilities
and risks are mathematically interchangeable, however, in order to promote fairness, the desirable
axioms of malfare and welfare functions differ slightly. The notation reflects this; $\mathrm{M}(\mathcal{S}; \boldsymbol{w})$ is M
for *mean*, whereas $\mathrm{W}(\mathcal{S}; \boldsymbol{w})$ is W for *welfare*, and $\mathrm{\Lambda\!\!\Lambda}(\mathcal{S}; \boldsymbol{w})$ is $\mathrm{\Lambda\!\!\Lambda}$ (*inverted* W), to emphasize its
inverted nature.

2.1 Axioms of Cardinal Welfare and Malfare

Definition 2.2 (Axioms of Cardinal Welfare and Malfare). We define the *population-mean axioms*
for population-mean function $\mathrm{M}(\mathcal{S}; \boldsymbol{w})$ below. For each item, assume (if necessary) that the axiom
applies $\forall \mathcal{S}, \mathcal{S}' \in \mathbb{R}_{0+}^g$, scalars $\alpha, \beta \in \mathbb{R}_{0+}$, and $\boldsymbol{w} \in \mathbb{R}_+^g$.
1. (Strict) Monotonicity: $\mathcal{S}' \neq \boldsymbol{0} \implies \mathrm{M}(\mathcal{S}; \boldsymbol{w}) < \mathrm{M}(\mathcal{S} + \mathcal{S}'; \boldsymbol{w})$.
2. Symmetry: $\forall$ permutations π over $\{1, \ldots, g\}$: $\mathrm{M}(\mathcal{S}; \boldsymbol{w}) = \mathrm{M}(\pi(\mathcal{S}); \pi(\boldsymbol{w}))$.
3. Continuity: $\{\mathcal{S}' \mid \mathrm{M}(\mathcal{S}'; \boldsymbol{w}) \leq \mathrm{M}(\mathcal{S}; \boldsymbol{w})\}$ and $\{\mathcal{S}' \mid \mathrm{M}(\mathcal{S}'; \boldsymbol{w}) \geq \mathrm{M}(\mathcal{S}; \boldsymbol{w})\}$ are closed sets.
4. Independence of unconcerned agents:
$$\mathrm{M}(\langle \mathcal{S}_{1:g-1}, \alpha \rangle; \boldsymbol{w}) \leq \mathrm{M}(\langle \mathcal{S}'_{1:g-1}, \alpha \rangle; \boldsymbol{w}) \implies \mathrm{M}(\langle \mathcal{S}_{1:g-1}, \beta \rangle; \boldsymbol{w}) \leq \mathrm{M}(\langle \mathcal{S}'_{1:g-1}, \beta \rangle; \boldsymbol{w}) \ .$$
5. Independence of common scale: $\mathrm{M}(\mathcal{S}; \boldsymbol{w}) \leq \mathrm{M}(\mathcal{S}'; \boldsymbol{w}) \implies \mathrm{M}(\alpha \mathcal{S}; \boldsymbol{w}) \leq \mathrm{M}(\alpha \mathcal{S}'; \boldsymbol{w})$.
6. Multiplicative linearity: $\mathrm{M}(\alpha \mathcal{S}; \boldsymbol{w}) = \alpha \mathrm{M}(\mathcal{S}; \boldsymbol{w})$.
7. Unit scale: $\mathrm{M}(\boldsymbol{1}; \boldsymbol{w}) = \mathrm{M}(1, \ldots, 1; \boldsymbol{w}) = 1$.
8. Pigou-Dalton transfer principle: Suppose $\mu = \boldsymbol{w} \cdot \mathcal{S} = \boldsymbol{w} \cdot \mathcal{S}'$, and $\forall i \in \{1, \ldots, g\} : |\mu - \mathcal{S}'_i| \leq |\mu - \mathcal{S}_i|$. Then $\mathrm{W}(\mathcal{S}'; \boldsymbol{w}) \geq \mathrm{W}(\mathcal{S}; \boldsymbol{w})$.
9. Anti-Pigou-Dalton transfer principle: Suppose as in 8, but conclude $\mathrm{\Lambda\!\!\Lambda}(\mathcal{S}'; \boldsymbol{w}) \leq \mathrm{\Lambda\!\!\Lambda}(\mathcal{S}; \boldsymbol{w})$.

We take a moment to comment on each axiom, to preview their purpose and assure the reader of
their necessity. Axioms 1-5 are the standard *axioms of cardinal welfarism* [8, 9, 21, 19]. Together,
they imply (via the Debreu-Gorman theorem) that any population-mean can be decomposed as a
monotonic function of a *sum* (over groups) of *log* or *power* functions. Axiom 6 is a natural and useful
property, and one which enables *dimensional analysis* on mean functions; in particular, the *units*
of population mean functions match those of sentiment values. Note that axiom 6 implies axiom 4,
and is thus a simple strengthening of a traditional cardinal welfare axiom. We will also see that it is
essential to show convenient *statistical* and *learnability* properties. Axiom 7 furthers this theme, as it
ensures that not only do *units* of means match those of $\mathcal{S}$, but scale does as well (making comparisons
like $\mathcal{S}_i$ is *above* the population-welfare meaningful), and also enabling comparison *across populations*
(i.e., comparing *averages* rather than *sums*). Finally, axiom 8 (the *Pigou-Dalton transfer principle*
[18, 7]. is also standard in cardinal welfare theory as it ensures fairness, in the sense that welfare is
higher when utility values are more uniform (i.e., incentivizing *equitable redistribution of "wealth"*
in welfare). Its antithesis, axiom 9, encourages the opposite; in the context of welfare, this perversely
incentivizes inequality, but for malfare, which we generally wish to *minimize*, the opposite occurs,
thus this axiom characterizes *fairness* in the context of *malfare*.

The Power Mean We now define the *p-power mean* $\mathrm{M}_p(\cdot; \cdot)$, for any $p \in \mathbb{R} \cup \pm\infty$, which we
shall use to quantify both malfare and welfare. Power means arise when analyzing population means
obeying the various axioms of def. 2.2, and are thus a particularly important class of population mean.

Definition 2.3 (Power-Mean Welfare and Malfare). Suppose $p \in \mathbb{R} \cup \pm\infty$, sentiment vectors $\mathcal{S} \in \mathbb{R}_{0+}^g$, and probability vector $\boldsymbol{w} \in \mathbb{R}_+^g$. We then define the *p-weighted-power-mean* as

$$\mathrm{M}_p(\mathcal{S}; \boldsymbol{w}) \doteq \sqrt[p]{\sum_{i=1}^g \boldsymbol{w}_i \mathcal{S}_i^p} \ , \ \mathrm{M}_0(\mathcal{S}; \boldsymbol{w}) \doteq \exp\left(\sum_{i=1}^g \boldsymbol{w}_i \ln \mathcal{S}_i\right) \ , \ \mathrm{M}_{\pm\infty}(\mathcal{S}; \boldsymbol{w}) \doteq \pm \max_{i \in 1,\dots,g} \pm\mathcal{S}_g \ ,$$

where $p \in \{-\infty, 0, \infty\}$ attain their *limits*, the *minimum*, *geometric mean*, and *maximum*, respectively.

Theorem 2.4 (Properties of the Power-Mean). Suppose sentiment vectors $\mathcal{S}, \boldsymbol{\varepsilon} \in \mathbb{R}_{0+}^g$, and probability vector $\boldsymbol{w} \in \mathbb{R}_+^g$. The following then hold.
1. Monotonicity: $\mathrm{M}_p(\mathcal{S}; \boldsymbol{w})$ is weakly-monotonically-increasing in p.
2. Subadditivity: $\forall p \geq 1 : \mathrm{M}_p(\mathcal{S} + \boldsymbol{\varepsilon}; \boldsymbol{w}) \leq \mathrm{M}_p(\mathcal{S}; \boldsymbol{w}) + \mathrm{M}_p(\boldsymbol{\varepsilon}; \boldsymbol{w})$.
3. Contraction: $\forall p \geq 1 : \left|\mathrm{M}_p(\mathcal{S}; \boldsymbol{w}) - \mathrm{M}_p(\mathcal{S}'; \boldsymbol{w})\right| \leq \mathrm{M}_p(|\mathcal{S} - \mathcal{S}'|; \boldsymbol{w}) \leq \|\mathcal{S} - \mathcal{S}'\|_\infty$.
4. Curvature: $\mathrm{M}_p(\mathcal{S}; \boldsymbol{w})$ is concave in $\mathcal{S}$ for $p \leq 1$ and convex for $p \geq 1$.

2.2 Properties of Welfare and Malfare Functions

We now derive properties of welfare and malfare from the axioms of def. 2.2.

Theorem 2.5 (Population Mean Properties). Suppose population-mean function $\mathrm{M}(\mathcal{S}; \boldsymbol{w})$, and assume arbitrary sentiment vector $\mathcal{S} \in \mathbb{R}_{0+}^g$ and probability vector $\boldsymbol{w} \in \mathbb{R}_+^g$. If $\mathrm{M}(\cdot; \cdot)$ satisfies (subsets of) the population-mean axioms (see def. 2.2), then $\mathrm{M}(\cdot; \cdot)$ exhibits the following properties.
1. *Identity*: Axioms 6 & 7 imply $\mathrm{M}(\alpha\mathbf{1}; \boldsymbol{w}) = \alpha$.
2. Axioms 1-5 imply $\exists p \in \mathbb{R}$, strictly-monotonically-increasing continuous $F : \mathbb{R} \to \mathbb{R}_{0+}$ s.t.

$$\mathrm{M}(\mathcal{S}; \boldsymbol{w}) = F\left(\sum_{i=1}^g f_p(\mathcal{S}_i)\right) \ , \ \text{with} \begin{cases} p = 0 & f_0(x) \doteq \ln(x) \\ p \neq 0 & f_p(x) \doteq \mathrm{sgn}(p)x^p \end{cases} \ .$$

3. Axioms 1-7 imply $F(x) = f_p^{-1}(x)$, thus $\mathrm{M}(\mathcal{S}; \boldsymbol{w}) = \mathrm{M}_p(\mathcal{S}; \boldsymbol{w})$.
4. Axioms 1-5 and 8 imply $p \in (-\infty, 1]$.
5. Axioms 1-5 and 9 imply $p \in [1, \infty)$.

Thm. 2.5 tells us that the mild conditions of axioms 1-4 (generally assumed for welfare), along with 5 (multiplicative linearity), imply that welfare and utility, or malfare and loss, are measured in the *same units* (e.g., *nats* or of *cross-entropy loss*, or *dollars* of *income utility*), and the power-mean is effectively the only reasonable family of welfare or malfare function; even without multiplicative linearity, axioms 1-5 imply population mean functions are still *monotonic transformations* of power-means. Furthermore, the entirely milquetoast *unit scale* axiom implies that sentiment values and population means have the same *scale*, making comparisons like "the risk of group i is above (or below) the population malfare" meaningful. Finally, we have that $\mathrm{W}_p(\cdot; \boldsymbol{w})$ promotes equitable redistribution of utility for $p \leq 1$, and $\mathcal{M}_p(\cdot; \boldsymbol{w})$ promotes equitable redistribution of loss for $p \geq 1$.

3 Statistical and Computational Learning-Efficiency Guarantees

In this section, we define a formal notion of fair-learnability, termed *FPAC-learning*, where a loss function and hypothesis class are FPAC-learnable essentially if any distribution can be learned to approximate malfare-optimality from a finite sample (w.h.p.). We then construct various FPAC learners, and relate the concept to standard PAC learning [26], with the understanding that this allows the vast breadth of research on PAC-learning algorithms to be applied to FPAC learning. In particular, we show a hierarchy of fair-learnability via generic statistical and computational reductions.

Definition 3.1 (Hypothesis Classes and Class Sequences). A *hypothesis class* is a family of functions mapping domain $\mathcal{X}$ to codomain $\mathcal{Y}$, and a *hypothesis class sequence* $\mathcal{H} = \mathcal{H}_1, \mathcal{H}_2, \dots$ is a *nested sequence* of *hypothesis classes* mapping $\mathcal{X} \to \mathcal{Y}$. In other words, $\mathcal{H}_1 \subseteq \mathcal{H}_2 \subseteq \dots$.

For example, *linear classifiers* naturally form a sequence of families using their *dimension* as $\mathcal{H}_d \doteq \left\{\vec{x} \mapsto \mathrm{sgn}\left(\vec{x} \cdot \langle \vec{w}_{1:d}, \vec{0}\rangle\right) \ \middle| \ \vec{w} \in \mathbb{R}^d\right\}$. The *hypothesis class sequence* concept allows us to distinguish statistically easy problems, like learning hyperplanes in finite-dimensional space, from the challenging problem of learning hyperplanes in $\mathbb{R}^\infty$. It is also used to analyze the *computational efficiency* of learners as d increases.

Figure 1: We minimize malfare on a *weighted hinge-loss* SVM, with $g = 5$ ethnoracial protected groups, listed in the legend with *group weight* $\boldsymbol{w}_i$ (population frequency) and *class bias* $\boldsymbol{b}_i$ (proportion with income $\geq$ \$50k/yr.). Due to existing societal inequity, *class imbalance* varies widely by group, so we weight risk as $\frac{1}{\boldsymbol{b}_i}\hat{\mathrm{R}}(h; \ell_{\text{hinge}}, \boldsymbol{z}_i)$. We report per-group training (dotted) and test (dashed) hinge risk, along with the malfare objective (green) of the EMM solution

$$\hat{h} \doteq \operatorname*{argmin}_{h \in \mathcal{H}} \mathbb{M}_p\left(i \mapsto \frac{1}{\boldsymbol{b}_i}\hat{\mathrm{R}}(h; \ell_{\text{hinge}}, \boldsymbol{z}_i); \boldsymbol{w}\right) \ ,$$

as a function of $p \in [1, 32]$. Detailed description of the experimental setup is provided in app. B.1 .

3.1 Empirical Malfare Minimization and the Rademacher Average

We define the *risk* of hypothesis $h : \mathcal{X} \to \mathcal{Y}$ w.r.t. loss $\ell : \mathcal{Y} \times \mathcal{Y} \to \mathbb{R}_+$ on distribution $\mathcal{D}$ over $(\mathcal{X} \times \mathcal{Y})$, and the *empirical risk* on sample $\boldsymbol{z} \in (\mathcal{X} \times \mathcal{Y})^m$, as

$$\mathrm{R}(h; \ell, \mathcal{D}) \doteq \mathbb{E}_{(x,y)\sim\mathcal{D}}\big[\ell(y, h(x))\big] \quad \& \quad \hat{\mathrm{R}}(h; \ell, \boldsymbol{z}) \doteq \hat{\mathbb{E}}_{(x,y)\in\boldsymbol{z}}\big[\ell(y, h(x))\big] \ ,$$

respectively. The goal in ML is generally to recover the h^* that minimizes (true) risk, and *empirical risk minimization* (ERM) computes $\hat{h} \doteq \operatorname{argmin}_{h\in\mathcal{H}}\hat{\mathrm{R}}(h; \ell, \boldsymbol{z})$ as a proxy for $h^* \doteq \operatorname{argmin}_{h\in\mathcal{H}}\mathrm{R}(h; \ell, \mathcal{D})$. Analogously, for fair machine learning, we define *empirical malfare minimization* (EMM), given $\mathbb{M}(\cdot; \boldsymbol{w})$, $\mathcal{D}_{1:g}$, and $\boldsymbol{z}_{1:g}$, with proxy and ideal models

$$\hat{h} \doteq \operatorname*{argmin}_{h \in \mathcal{H}} \mathbb{M}\left(i \mapsto \hat{\mathrm{R}}(h; \ell, \boldsymbol{z}_i); \boldsymbol{w}\right) \quad \& \quad h^* \doteq \operatorname*{argmin}_{h \in \mathcal{H}} \mathbb{M}\left(i \mapsto \mathrm{R}(h; \ell, \mathcal{D}_i); \boldsymbol{w}\right) \ .$$

Although empirical risk is an *unbiased estimate* of risk, with $\mathbb{M}_p(\cdot; \boldsymbol{w})$ for $p > 1$, *empirical malfare* is a *biased estimator* of malfare. Fortunately, empirical malfare is a *consistent estimator* of malfare, and we show that standard tools for bounding the error of ERM can be applied to EMM. Here it is necessary to consider not the loss function or hypothesis class *in isolation*, but their *composition*

$$\forall h \in \mathcal{H} : \ (\ell \circ h)(x, y) \doteq \ell(y, h(x)) \quad \& \quad \ell \circ \mathcal{H} \doteq \{\ell \circ h \,|\, h \in \mathcal{H}\} \ .$$

The *empirical Rademacher average* is a well-studied [22, 15] data-dependent measurement of the *capacity to overfit* of a model class $\mathcal{H}$ w.r.t. loss function ℓ, defined as

$$\hat{\mathfrak{R}}_m(\ell \circ \mathcal{H}, \boldsymbol{z}) \doteq \mathbb{E}_{\boldsymbol{\sigma}}\left[\sup_{h\in\mathcal{H}}\left|\frac{1}{m}\sum_{i=1}^{m}\boldsymbol{\sigma}_i\ell\big(\boldsymbol{y}_i, h(\boldsymbol{x}_i)\big)\right|\right] \ ,$$

where $\boldsymbol{\sigma}_{1:m}$ are drawn i.i.d. Rademacher (uniform on ± 1), which essentially measures the ability of the hypothesis class to *spuriously correlate* loss values with *noise*.

Theorem 3.2 (Generalization Guarantees for Malfare Estimation)**.** Suppose fair power-mean malfare $\mathbb{M}_p(\cdot; \cdot)$ (i.e., $p \geq 1$), *loss function* $\ell : (\mathcal{Y} \times \mathcal{Y}) \to [0, r]$, samples $\boldsymbol{z}_i \sim \mathcal{D}_i^m$, and hypothesis class $\mathcal{H} \subseteq \mathcal{X} \to \mathcal{Y}$. Then with probability at least $1 - \delta$ over choice of $\boldsymbol{z}$,

$$\sup_{h\in\mathcal{H}}\left|\mathbb{M}_p\big(i \mapsto \mathrm{R}(h; \ell, \mathcal{D}_i); \boldsymbol{w}\big) - \mathbb{M}_p\big(i \mapsto \hat{\mathrm{R}}(h; \ell, \boldsymbol{z}_i); \boldsymbol{w}\big)\right| \leq \mathbb{M}_p\left(i \mapsto 2\hat{\mathfrak{R}}_m(\ell \circ \mathcal{H}, \boldsymbol{z}_i) + 3r\sqrt{\frac{\ln\frac{g}{\delta}}{2m}}; \boldsymbol{w}\right) \ .$$

The *contraction property* (thm. 2.4 item 3) of malfare is key to this result, and this property does not hold for welfare. Strikingly, fair welfare functions $\mathrm{W}_p(\cdot; \mathcal{S})$ for $p \in [0, 1)$ are Lipschitz-discontinuous; e.g., the *Nash social welfare* (geometric welfare) $\mathrm{W}_0(\mathcal{S}; \frac{1}{g}) = \sqrt[g]{\prod_{i=1}^{g}\mathcal{S}_i}$ is unstable to perturbations of each $\mathcal{S}_i$ around 0. Consequently, if utility values are $\textsc{Bernoulli}(p)$-distributed for some group, the sample complexity of $\varepsilon\text{-}\delta$ estimating $\mathrm{W}_0(\mathcal{S}; \boldsymbol{w})$ grows unboundedly as $p \to 0$.

Experimental Validation Figure 1 presents a brief experiment on the lauded `adult` dataset, where the task is to predict whether income is above or below \$50k/yr. We train $\mathbb{M}_p(\cdot; \boldsymbol{w})$-minimizing SVM, and find significant risk-variation between groups; generally low risk for the *white* and *Asian* groups, and high risk for the *native American* and *other* groups. The $p = 1$ model is a standard weighted SVM, with poor performance for small and traditionally marginalized groups, as expected in an 85.43% majority-white population. As p increases (towards egalitarianism), we observe interesting

fairness tradeoffs; training malfare and weighted average risks increase monotonically, and in general (but not monotonically[1]), *white* and *Asian* training risks increase, as the remaining risks decrease. At first, most improvement is in the relatively-large (9.64%), high-risk Black group, but for larger p, the much smaller (0.96%), but higher-risk, native American group sharply improves.

Both training and test performance generally improve for high-risk groups, but significant overfitting occurs in small groups and malfare. This is unsurprising, as although SVM generalization error is well-understood (see Ch. 26 [22]), bounds are generally vaccuous for tiny subpopulations of ≈ 400 individuals. In general, overfitting increases with p, due to higher relative importance of small high-risk groups on $\hat{h}$. This experiment validates EMM as a fair-learning technique, with the capacity to specify tradeoffs between majority and marginalized groups, while demonstrating *overfitting to fairness*, which we formally treat in the sequel. We observe similar fairness tradeoffs in our supplementary experiments (app. B.2), on weighted and unweighted SVM and GLM with race and gender groups.

3.2 Fair Probably Approximately Correct Learning

For context, we first present a generalized notion of PAC-learnability, which we then generalize to FPAC-learnability. Standard presentations consider only classification under 0-1 loss, but we follow the generalized learning setting of [27], which treats generic learning problems.

Definition 3.3 (PAC-Learnability). Suppose *nested hypothesis class sequence* $\mathcal{H}_1 \subseteq \mathcal{H}_2 \subseteq \ldots$, all over $\mathcal{X} \to \mathcal{Y}$, and *loss function* $\ell : \mathcal{Y} \times \mathcal{Y} \to \mathbb{R}_{0+}$. We say $\mathcal{H}$ is *PAC-learnable* w.r.t. ℓ if $\exists$ a (randomized) algorithm $\mathcal{A}$, such that $\forall$:

1. sequence index d;
2. instance distribution $\mathcal{D}$ over $\mathcal{X} \times \mathcal{Y}$;
3. additive approximation error $\varepsilon > 0$; and
4. failure probability $\delta \in (0, 1)$;

$\mathcal{A}$ can identify a hypothesis $\hat{h} \in \mathcal{H}$, i.e., $\hat{h} \leftarrow \mathcal{A}(\mathcal{D}, \varepsilon, \delta, d)$, such that

1. there exists some *sample complexity* function $\mathrm{m}(\varepsilon, \delta, d) : \big(\mathbb{R}_+ \times (0, 1) \times \mathbb{N}\big) \to \mathbb{N}$ s.t. $\mathcal{A}(\mathcal{D}, \varepsilon, \delta, d)$ consumes no more than $\mathrm{m}(\varepsilon, \delta, d)$ samples from $\mathcal{D}$ (i.e., finite sample complexity); and
2. with probability at least $1 - \delta$ (over randomness of $\mathcal{A}$), $\hat{h}$ obeys

$$\mathrm{R}(\hat{h}; \ell, \mathcal{D}) \leq \inf_{h^* \in \mathcal{H}} \mathrm{R}(h^*; \ell, \mathcal{D}) + \varepsilon \ .$$

The class of such learning problems is denoted PAC, with membership denoted $(\mathcal{H}, \ell) \in$ PAC.

Furthermore, if for all d, the space of $\mathcal{D}$ is restricted such that $\exists h \in \mathcal{H}_d$ s.t. $\mathrm{R}(h; \ell, \mathcal{D}) = 0$, then $(\mathcal{H}, \ell)$ is *realizable-PAC-learnable*, written $(\mathcal{H}, \ell) \in$ PAC$^{\mathrm{Rlz}}$.

We now generalize this concept to *fair-PAC learnability*. In particular, we replace the *univariate risk-minimization* task with a *multivariate malfare-minimization* task. Following the theory of § 2.2, we do not commit to any particular objective, but instead require that a FPAC-learner is able to maximize *any* fair malfare function satisfying the standard axioms. Furthermore, here problem instances grow not just in problem complexity d, but also in the number of groups g.

Definition 3.4 (Fair-PAC (FPAC) Learnability). Suppose *nested hypothesis class sequence* $\mathcal{H}_1 \subseteq \mathcal{H}_2 \subseteq \cdots \subseteq \mathcal{X} \to \mathcal{Y}$, and *loss function* $\ell : \mathcal{Y} \times \mathcal{Y} \to \mathbb{R}_{0+}$. We say $\mathcal{H}$ is *fair PAC-learnable* w.r.t. *group count* g and *loss function* ℓ if $\exists$ a (randomized) algorithm $\mathcal{A}$, such that $\forall$:

1. sequence index d;
2. instance distribution $\mathcal{D}_{1:g}$ over $(\mathcal{X} \times \mathcal{Y})^g$;
3. probability vector $\boldsymbol{w} \in \mathbb{R}_+^g$
4. malfare concept $\mathbb{M}$ satisfying axioms 1-7 & 9;
5. additive approximation error $\varepsilon > 0$; and
6. failure probability $\delta \in (0, 1)$;

$\mathcal{A}$ can identify a hypothesis $\hat{h} \in \mathcal{H}$, i.e., $\hat{h} \leftarrow \mathcal{A}(\mathcal{D}_{1:g}, \boldsymbol{w}, \mathbb{M}, \varepsilon, \delta, d)$, such that

1. there exists some *sample complexity* function $\mathrm{m}(\varepsilon, \delta, d, g) : \big(\mathbb{R}_+ \times (0, 1) \times \mathbb{N} \times \mathbb{N}\big) \to \mathbb{N}$ s.t. $\mathcal{A}(\mathcal{D}_{1:g}, \boldsymbol{w}, \mathbb{M}, \varepsilon, \delta, d)$ consumes no more than $\mathrm{m}(\varepsilon, \delta, d, g)$ samples (finite sample complexity); and
2. with probability at least $1 - \delta$ (over randomness of $\mathcal{A}$), $\hat{h}$ obeys

$$\mathbb{M}\big(i \mapsto \mathrm{R}(\hat{h}; \ell, \mathcal{D}_i); \boldsymbol{w}\big) \leq \inf_{h^* \in \mathcal{H}} \mathbb{M}\big(i \mapsto \mathrm{R}(h^*; \ell, \mathcal{D}_i); \boldsymbol{w}\big) + \varepsilon \ .$$

---

[1]Note that for continuous loss functions and $g = 2$ groups, group training risks are monotonic in p, as seen in the supplementary *gender-group* experiments.

265 The class of such fair-learning problems is FPAC, with membership denoted $(\mathcal{H}, \ell) \in \text{FPAC}$.

266 Finally, if for all d, the space of $\mathcal{D}$ is restricted such that $\exists h \in \mathcal{H}_d$ s.t. $\max\limits_{i \in 1, \ldots, g} \text{R}(h; \ell, \mathcal{D}_i) = 0$, then

267 $(\mathcal{H}, \ell)$ is *realizable-FPAC-learnable*, written $(\mathcal{H}, \ell) \in \text{FPAC}^{\text{Rlz}}$.

268 **Observation 3.5** (Malfare Functions and Special Cases)**.** By assumption, $\mathbb{M}(\cdot; \cdot)$ must be $\mathbb{M}_p(\cdot; \cdot)$
269 for some $p \in [1, \infty)$. Taking $g = 1$ implies $\boldsymbol{w} = \langle 1 \rangle$, and $\mathbb{M}_p(\mathcal{S}; \boldsymbol{w}) = \mathcal{S}_1$, thus reducing the
270 problem to standard PAC-learning. Similarly, taking $p = 1$ converts the problem to *weighted loss*
271 *minimization* (weights determined by $\boldsymbol{w}$), and $p = \infty$ yields a *minimax optimization problem*, where
272 the max is over groups, as commonly encountered in adversarial and robust learning settings.

273 **On Computational Efficiency** Some authors consider not just the *statistical* but also the *computa-*
274 *tional* performance of learning, generally requiring that $\mathcal{A}$ have *polynomial* time complexity (and
275 thus implicitly sample complexity). In other words, they require that $\mathcal{A}(\mathcal{D}, \varepsilon, \delta, d)$ terminates in
276 $\text{m}(\varepsilon, \delta, d) \in \text{Poly}(\frac{1}{\varepsilon}, \frac{1}{\delta}, d)$ steps. A similar concept of *polynomial-time* FPAC-learnability is equally
277 interesting, where here we assume $\mathcal{A}(\mathcal{D}_{1:g}, \boldsymbol{w}, \mathbb{M}, \varepsilon, \delta, d)$ may be computed by a Turing machine
278 (with access to *sampling* and *entropy* oracles) in $\text{m}(\varepsilon, \delta, d, g) \in \text{Poly}(\frac{1}{\varepsilon}, \frac{1}{\delta}, d, g)$ steps. We denote
279 these concepts $\text{PAC}^{\text{Agn}}_{\text{Poly}}$, $\text{PAC}^{\text{Rlz}}_{\text{Poly}}$, $\text{FPAC}^{\text{Agn}}_{\text{Poly}}$, and $\text{FPAC}^{\text{Rlz}}_{\text{Poly}}$.

280 We first observe (immediately from defs. 3.3 and 3.4) that PAC-learning is a special case of FPAC-
281 learning; in particular, taking $g = 1$ implies $\text{M}_p(\mathcal{S}) = \text{M}_1(\mathcal{S}) = \mathcal{S}_1$, thus *malfare-minimization*
282 coincides with *risk minimization*. The more interesting question, which we seek to answer in the
283 remainder of this document, is *when* and *whether* the *converse* holds. Furthermore, when possible,
284 we would like to show practical *constructive reductions*.

285 We now show that in the *realizable case*, *PAC-learnability* implies *FPAC-learnability* via a construc-
286 tive polynomial-time reduction. Our reduction simply takes a sufficiently number of samples from the
287 *uniform mixture distribution* over all g groups, and PAC-learns on this distribution. As the reduction
288 is constructive (and efficiency-preserving), this gives us generic algorithms for (efficient) realizable
289 FPAC-learning in terms of algorithms for realizable PAC-learning.

290 **Theorem 3.6** (Realizable Reductions)**.** Suppose loss function ℓ and hypothesis class $\mathcal{H}$. Then

291 1. $(\mathcal{H}, \ell) \in \text{PAC}^{\text{Rlz}} \implies (\mathcal{H}, \ell) \in \text{FPAC}^{\text{Rlz}}$; & 292 2. $(\mathcal{H}, \ell) \in \text{PAC}^{\text{Rlz}}_{\text{Poly}} \implies (\mathcal{H}, \ell) \in \text{FPAC}^{\text{Rlz}}_{\text{Poly}}$.

293 We construct a(n) (efficient) FPAC-learner for $(\mathcal{H}, \ell)$ by noting that there exists some $\mathcal{A}'$ with
294 sample-complexity $\text{m}_{\mathcal{A}'}(\varepsilon, \delta, d)$ and time complexity $\text{t}_{\mathcal{A}'}(\varepsilon, \delta, d)$ to PAC-learn $(\mathcal{H}, \ell)$, and tak-
295 ing $\mathcal{A}(\mathcal{D}_{1:g}, \boldsymbol{w}, \mathbb{M}, \varepsilon, \delta, d) \doteq \mathcal{A}'(\text{mix}(\mathcal{D}_{1:g}), \frac{\varepsilon}{g}, \delta, d)$. Then $\mathcal{A}$ FPAC-learns $(\mathcal{H}, \ell)$, with sample-
296 complexity $\text{m}_{\mathcal{A}}(\varepsilon, \delta, d, g) = \text{m}_{\mathcal{A}'}(\frac{\varepsilon}{g}, \delta, d)$, and time-complexity $\text{t}_{\mathcal{A}}(\varepsilon, \delta, d, g) = \text{t}_{\mathcal{A}'}(\frac{\varepsilon}{g}, \delta, d)$.

297 Unfortunately, the argument here strongly depends on the properties of realizablility, and does not
298 extend to the agnostic case. Furthermore, we note that, philosophically speaking, realizable FPAC
299 learning is rather uninteresting, essentially because in a world where all parties may be satisfied
300 completely, the obvious solution is to do so. Thus nontrivial unfairness and bias issues logically only
301 arise in a world of *conflict*, e.g., in zero-sum or resource-constrained settings, which foster *competition*
302 between groups. We henceforth focus our efforts on the more interesting agnostic-learning setting.

3.3 Characterizing Fair Statistical Learnability with FPAC-Learners

304 We first consider only questions of *statistical learning*, i.e., we ignore computation and show that
305 *there exist* FPAC-learning algorithms. In particular, we show a generalization of the *fundamental*
306 *theorem of statistical learning* to fair learning problems. The aforementioned result relates *uniform*
307 *convergence* and *PAC-learnability*, and is generally stated for binary classification only. We define a
308 natural generalization of uniform convergence to arbitrary learning problems within our framework,
309 and then show conditions under which a generalized fundamental theorem of (fair) statistical learning
310 holds. In particular, we show that, neglecting computational concerns, PAC-learnability and FPAC
311 learnability are equivalent for learning problems with a particular *no-free-lunch* guarantee.

312 We first define a *generalized notion* of *uniform convergence*. In particular, our definition applies to
313 *any bounded loss function*, thus greatly generalizes the standard notion for binary classification.

314 **Definition 3.7** (Uniform Convergence)**.** Suppose $\ell : \mathcal{Y} \times \mathcal{Y} \to [0, r] \subseteq \mathbb{R}$ and hypothesis class
315 $\mathcal{H} \subseteq \mathcal{X} \to \mathcal{Y}$. We say $(\mathcal{H}, \ell) \in \text{UC}$ if

Figure 2: Implications between membership in PAC and FPAC classes. In particular, for arbitrary fixed ℓ, implication denotes *implication of membership* of some $\mathcal{H}$ (i.e., containment); see thms. 3.6 and 3.9.

$$\lim_{m \to \infty} \sup_{\mathcal{D} \text{ over } \mathcal{X} \times \mathcal{Y}} \mathbb{E}_{\boldsymbol{z} \sim \mathcal{D}^m} \left[\sup_{h \in \mathcal{H}} \left| \hat{R}(h; \ell, \boldsymbol{z}) - R(h; \ell, \mathcal{D}) \right| \right] = 0 \ .$$

We stress that this definition is both uniform over ℓ composed with the *hypothesis class* $\mathcal{H}$ and uniform over *all possible distributions* $\mathcal{D}$. Standard uniform convergence definitions consider only the convergence of *empirical frequencies* of events to their *true frequencies*, whereas we generalize to consider uniform convergence of the *empirical means* of functions to their *expected values*. It is also helpful to consider the *sample complexity* of ε-δ uniform-convergence, where we take

$$m_{\mathrm{UC}}(\ell \circ \mathcal{H}, \varepsilon, \delta) \doteq \operatorname{argmin}\left\{ m \ \middle| \ \sup_{\mathcal{D} \text{ over } \mathcal{X} \times \mathcal{Y}} \mathbb{P}\left(\sup_{h \in \mathcal{H}} \left| \mathbb{E}_{\mathcal{D}}[\ell \circ h] - \hat{\mathbb{E}}_{\boldsymbol{z} \sim \mathcal{D}^m}[\ell \circ h] \right| > \varepsilon \right) \le \delta \right\} \ ,$$

i.e., the *minimum sufficient sample size* to ensure ε-δ uniform-convergence over the *loss family* $\ell \circ \mathcal{H}$. For classification, *uniform convergence* bijectively implies *PAC-learnability* as follows.

Theorem 3.8 (Fundamental Theorem of Statistical Learning (Classification) [22, thm. 6.2]). Suppose ℓ is the 0-1 loss. Then the following are equivalent:

1. $\forall d \in \mathbb{N}$: $\mathcal{H}_d$ has finite VC dimension.
2. $\forall d \in \mathbb{N}$: $(\ell, \mathcal{H}_d)$ has the uniform convergence property.
3. Any ERM rule is a successful agnostic-PAC learner for $\mathcal{H}$, thus $\mathcal{H}$ is agnostic-PAC learnable.
4. Any ERM rule is a successful realizable-PAC learner for $\mathcal{H}$, thus $\mathcal{H}$ is realizable-PAC learnable.

It is rather subtle to generalize this result beyond classification, as there are PAC-learnable problems for which uniform convergence *does not hold*. However, [1] show similar results for various regression problems, and we now consider FPAC learning (summarized in fig. 2).

Theorem 3.9 (Fundamental Theorem of Fair Statistical Learning). Suppose ℓ such that $\forall \mathcal{H}$: $(\mathcal{H}, \ell) \in \mathrm{PAC}^{\mathrm{Rlz}} \implies (\mathcal{H}, \ell) \in \mathrm{UC}$. Then $\forall \mathcal{H}$, the following are equivalent:

1. $\forall d \in \mathbb{N}$: $(\ell, \mathcal{H}_d)$ has the (generalized) uniform convergence property.
2. Any EMM rule is an agnostic-FPAC learner for $(\ell, \mathcal{H})$, thus $(\ell, \mathcal{H})$ is agnostic-FPAC learnable.
3. Any EMM rule is a realizable-FPAC learner for $(\ell, \mathcal{H})$, thus $(\ell, \mathcal{H})$ is realizable-FPAC learnable.

3.4 Characterizing Computational Learnability with Efficient FPAC Learners

In this section, we consider the more granular question of whether FPAC learning is *computationally harder* than PAC learning. In other words, where previously we showed conditions under which $\mathrm{PAC} = \mathrm{FPAC}$, here we focus on the subset of models with polynomial time training efficiency guarantees, i.e., $\mathrm{PAC}_{\mathrm{Poly}} = \mathrm{FPAC}_{\mathrm{Poly}}$. Thm. 3.6 has already characterized the computational complexity of *realizable* FPAC-learning, so we now focus on the agnostic case. Here we show neither a generic reduction nor a non-constructive proof that $\mathrm{PAC}_{\mathrm{Poly}} = \mathrm{FPAC}_{\mathrm{Poly}}$, nor do we show a counterexample, but rather we show that under standard convexity conditions often leveraged as sufficient for efficient PAC-learning, so too is efficient FPAC-learning possible.

Here we show concretely and constructively the existence of FPAC-learners under standard convex optimization assumptions via the *subgradient method* [23]. This result is immediately practical, and can be applied verbatim to problems like *generalized linear models* [17] and many *kernel methods*.

Algorithm 1 Approximate Empirical Malfare Minimization via the Subgradient Method

---
1: **procedure** $\mathcal{A}_{\mathrm{PSG}}(\ell, \mathcal{H}, \theta_0, \mathrm{m}_{\mathrm{UC}}(\cdot, \cdot), \mathcal{D}_{1:g}, \boldsymbol{w}, \mathcal{M}(\cdot; \cdot), \varepsilon, \delta)$
2:     **Input**: λ_ℓ-Lipschitz loss function ℓ, $\lambda_{\mathcal{H}}$-Lipschitz hypothesis class $\mathcal{H}$ with parameter space Θ, initial guess $\theta_0 \in \Theta$, uniform-convergence sample-complexity bound $\mathrm{m}_{\mathrm{UC}}(\cdot, \cdot)$, group distributions $\mathcal{D}_{1:g}$, group weights $\boldsymbol{w} \in \mathbb{R}_+^g$, malfare function $\mathcal{M}(\cdot; \cdot)$, and solution optimality guarantee ε-δ.
3:     **Output**: ε-δ-$\mathcal{M}(\cdot; \cdot)$-optimal $h \in \mathcal{H}$ (under the conditions of thm. 3.10).
4:     $\mathrm{m}_{\mathcal{A}} \leftarrow \mathrm{m}_{\mathrm{UC}}(\frac{\varepsilon}{3}, \frac{\delta}{g}); \boldsymbol{z}_{1:g, 1:\mathrm{m}_{\mathcal{A}}} \sim \mathcal{D}_1^{\mathrm{m}_{\mathcal{A}}} \times \cdots \times \mathcal{D}_g^{\mathrm{m}_{\mathcal{A}}}$ ▷ Draw sufficient sample for each group
5:     $n \leftarrow 1 + \left(\frac{3 \operatorname{Diam}(\Theta) \lambda_\ell \lambda_{\mathcal{H}}}{\varepsilon}\right)^2; \alpha \leftarrow \frac{\operatorname{Diam}(\Theta)}{\lambda_\ell \lambda_{\mathcal{H}} \sqrt{n}}$ ▷ Set iteration count n and learning rate α
6:     $f(\theta): \Theta \mapsto \mathbb{R}_{0+} \doteq \mathcal{M}\big(i \mapsto \hat{\mathrm{R}}(h(\cdot; \theta); \ell, \boldsymbol{z}_i); \boldsymbol{w}\big)$ ▷ Define empirical malfare objective
7:     $\hat{\theta} \leftarrow \textsc{ProjectedSubgradient}(f, \Theta, \theta_0, n, \alpha)$ ▷ Run PSG algorithm on empirical malfare
8:     **return** $h(\cdot; \hat{\theta})$ ▷ Return ε-δ optimal model
9: **end procedure**

Theorem 3.10 (Efficient FPAC Learning via Convex Optimization). Suppose each hypothesis space $\mathcal{H}_d \in \mathcal{H}$ is indexed by $\Theta_d \subseteq \mathbb{R}^{\operatorname{Poly}(d)}$, i.e., $\mathcal{H}_d = \{h(\cdot; \theta) \mid \theta \in \Theta_d\}$, s.t. (Euclidean) $\operatorname{Diam}(\Theta_d) \in \operatorname{Poly}(d)$, and $\forall x \in \mathcal{X}, \theta \in \Theta_d, h(x; \theta)$ can be evaluated in $\operatorname{Poly}(d)$ time, and $\tilde{\theta} \in \mathbb{R}^{\operatorname{Poly}(d)}$ can be Euclidean-projected onto Θ_d in $\operatorname{Poly}(d)$ time. Suppose also ℓ such that $\forall x \in \mathcal{X}, y \in \mathcal{Y}: \theta \mapsto \ell(y, h(x; \theta))$ is a *convex function*, and suppose Lipschitz constants $\lambda_\ell, \lambda_{\mathcal{H}} \in \operatorname{Poly}(d)$ and some norm $\|\cdot\|_{\mathcal{Y}}$ over $\mathcal{Y}$ s.t. ℓ is $\lambda_\ell \|\cdot\|_{\mathcal{Y}} \|\cdot\|$-Lipschitz in $\hat{y}$ and each $\mathcal{H}_d$ is $\lambda_{\mathcal{H}} \|\cdot\|_2 \|\cdot\|_{\mathcal{Y}}$-Lipschitz in θ, i.e.,

$$\forall y, \hat{y}, \hat{y}' \in \mathcal{Y}: \left|\ell(y, \hat{y}) - \ell(y, \hat{y}')\right| \leq \lambda_\ell \left\|\hat{y} - \hat{y}'\right\|_{\mathcal{Y}} \ \& \ \forall x \in \mathcal{X}, \theta, \theta' \in \Theta_d: \left\|h(x; \theta) - h(x; \theta')\right\|_{\mathcal{Y}} \leq \lambda_{\mathcal{H}} \left\|\theta - \theta'\right\|_2 \quad .$$

Finally, assume $\ell \circ \mathcal{H}_d$ exhibits ε-δ *uniform convergence* with sample complexity $\mathrm{m}_{\mathrm{UC}}(\varepsilon, \delta, d) \in \operatorname{Poly}(\frac{1}{\varepsilon}, \frac{1}{\delta}, d)$. It then holds that, for arbitrary initial guess $\theta_0 \in \Theta_d$, for any group distributions $\mathcal{D}_{1:g}$, group weights $\boldsymbol{w}$, and fair malfare function $\mathcal{M}(\cdot; \cdot)$, the algorithm (see alg. 1)

$$\mathcal{A}(\mathcal{D}_{1:g}, \boldsymbol{w}, \mathcal{M}(\cdot; \cdot), \varepsilon, \delta, d) \doteq \mathcal{A}_{\mathrm{PSG}}\big(\ell, \mathcal{H}_d, \theta_0, \mathrm{m}_{\mathrm{UC}}(\cdot, \cdot, d), \mathcal{D}_{1:g}, \boldsymbol{w}, \mathcal{M}(\cdot; \cdot), \varepsilon, \delta\big)$$

FPAC-learns $(\mathcal{H}, \ell)$ with sample complexity $\mathrm{m}(\varepsilon, \delta, d, g) = g \cdot \mathrm{m}_{\mathrm{UC}}(\frac{\varepsilon}{3}, \frac{\delta}{g}, d)$, and (training) time-complexity $\in \operatorname{Poly}(\frac{1}{\varepsilon}, \frac{1}{\delta}, d, g)$, thus $(\mathcal{H}, \ell) \in \operatorname{FPAC}_{\operatorname{Poly}}^{\operatorname{Agn}}$.

4   Conclusion

This paper introduces *malfare minimization* as a fair learning task, which we argue is better aligned to address ML tasks cast as loss minimization than is welfare. We then show statistical and computational relationships between malfare and risk minimization. We do not claim that malfare is a better or more useful concept than welfare; rather only that it is *significantly different*, stands on an equal axiomatic footing, and it stands to reason that the right tool (malfare) should be used for the right job (fair risk minimization). We acknowledge that some learning tasks, e.g., bandit problems and reinforcement learning tasks, are more naturally phrased as *maximizing* utility or (discounted) reward, however, most supervised learning problems are naturally cast as minimizing nonnegative *loss functions*.

We are highly interested in exploring a parallel theory of *fair welfare maximization*, however some key malfare properties do not hold for welfare. In particular, non-Lipschitz welfare functions (i.e., $p \in [0, 1)$), create great difficulty in both the *statistical* and *computational* aspects of learning. For this reason, straightforward translation of our FPAC framework into a welfare setting is rather vacuous. In contrast, the *contraction property* (thm. 2.4 item 3) of malfare leads to *uniform sample complexity bounds*, which are crucial to FPAC-learning. It thus seems that similar bounds for welfare would need either to either impose additional assumptions to avoid non-Lipschitz behavior (e.g., artificially limit the permitted range of p), or otherwise provide weaker (non-uniform) learning guarantees.

As FPAC generalizes PAC, known hardness results (e.g., lower-bounds on learning complexity) apply, thus the interesting question is whether *malfare minimization* is harder than *risk minimization*. Thm. 3.9 answers in the negative for *sample complexity*, under appropriate conditions on the loss function, as does thm. 3.6 under realizability. The remaining cases are left open, though § 3.4 shows that convexity conditions sufficient for PAC-learnability are also sufficient for FPAC-learnability. We hope that deeper inquiry into these questions will lead to both a better understanding of what is and is not FPAC-learnable, as well as more practical and efficient reductions and FPAC-learning algorithms.

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