Robust Streaming PCA

Abstract

We consider streaming principal component analysis (PCA) when the stochastic data-generating model is subject to perturbations. While existing models assume a fixed covariance, we adopt a robust perspective where the covariance matrix belongs to a temporal uncertainty set. Under this setting, we provide fundamental limits on any algorithm recovering principal components. We analyze the convergence of the noisy power method and Oja's algorithm, both analyzed for the stationary data generating model, and argue that the noisy power method is rate-optimal in our setting. Finally, we demonstrate the validity of our analysis through numerical experiments.

1 Introduction

Principal component analysis (PCA) is one of the most extensively studied methods for obtaining the low-dimensional representation of observed data [21]. However, classical algorithms for PCA store all the observations and use cubic-time complexity, thereby imposing prohibitively large computation-time and space requirements.

Recently, several works on PCA have focused on the design and analysis of single-pass (streaming) algorithms with near-optimal memory and storage complexity [40, 43, 24, 28]. These algorithms assume that all the observations belong to the same low-dimensional space. However, this situation is unlikely when the unknown/unexplored alterations corrupt a system's observations. For instance, it is well known that the typical data attacks on power grids can significantly change the estimated covariance matrix of the data observed from sensors [8, 22, 9]. Similarly, PCA can be used to explain stock returns in terms of the macroeconomic factors [26], product pricing taking into account cross product elasticity and demands [45], etc. In all these scenarios the underlying data generating model changes at every time instant and the decisions are based on identifying the changed model.

Recent work considers perturbations of the data lying in a *fixed* low-dimensional space [49]. They determine the worst-case position of the adversarial data-point to incur the maximum error in the subspace estimated through PCA and measure the distance between the two subspaces using the notion of the principal angle between them. Another line of work considers PCA through the lens of stochastic optimization [51, 52, 42, 4]. Our work differs from these approaches since we assume that the data generating model changes at every time instant. Further, we propose near-optimal algorithms for recovering principal components under this framework.

We assume a system relies on the time-series of p-dimensional vectors sampled from the time-variant model. The available observations are the vectors $\{\mathbf{x}_t\}_{t=1}^T$. We assume that $T \gg p$; justifying the use of streaming (single-pass) algorithms. The noisy observation $\mathbf{x}_t \in \mathbb{R}^p$ is a vector lying in the column space of an unobserved full-rank matrix $\mathbf{A}_t \in \mathbb{R}^{p \times k}$. Precisely, from the standard spiked covariance model [35], we consider the time-dependent environment:

$$\mathbf{x}_t \sim \mathcal{N}(\mathbf{0}_{p \times 1},\, \mathbf{A}_t \mathbf{A}_t^\top + \sigma^2 \mathbf{I}_{p \times p}), \tag{1}$$

where $\mathbf{A}_t$ can vary with time. The parameter $k \ll p$ is the desired number of principal components. Assuming $\mathbf{A}_t$ belongs to a temporal uncertainty set (defined in (2)), our goal is to recover the top-k principal components of the terminal subspace $\mathbf{A}_T$. In financial applications, our model captures the market evolution, in terms of the changing $\mathbf{A}_T$ with the ultimate goal of explaining the market conditions for an appropriately chosen T.

Previous works on the stationary environment assume $\mathbf{A}_t = \mathbf{A}$ for all t and have focused on computing a basis for the column space of the matrix $\mathbf{A}$ using streaming algorithms [43, 28, 38, 32, 20, 36]. The accurate reconstruction of the principal components for the standard streaming PCA problem depends on the magnitude of observation noise σ, the dimension of observations p, the number of principal components k of the matrix $\mathbf{A}$, and the spectral gap δ between k^{th} and $k+1^{\text{th}}$ singular value of $\mathbf{A}$. In a marked departure from previous work, we study the case when the column space of $\mathbf{A}$ varies across time, possibly perturbed by Γ. This paper explores these avenues and proposes a tractable analysis framework for the streaming PCA problem, robust to perturbations in the stochastic data-generating model. Our contributions can be summarized as:

(a) (Lower Bound; Section 4) Our first contribution is establishing a fundamental lower bound for estimating the principal components when the covariance matrix belongs to a temporal uncertainty set. In Theorem 1, we derive the minimax bounds of the expected error for recovering the top singular vectors for any algorithm in the class Φ. A key component of this analysis involves carefully constructed instances to elicit the worst-case behavior of any algorithm in Φ. We show that for $T = \mathcal{O}(\Gamma^{-2/3})$, the minimax estimation error decreases as $\mathcal{O}(p^{1/2}/\delta\sqrt{T})$. On the other hand, for $T = \Omega(\Gamma^{-2/3})$, the estimation error stagnates to $\mathcal{O}(p^{1/3}\Gamma^{1/3}/\delta)$ and does not decrease upon collecting more observations[1]. Since observations can belong to potentially different spaces, the estimation of principal components associated with $\mathbf{A}_T$ should be determined by a subset of samples. This is in sharp contrast to the standard spiked covariance model (where data is generated from the fixed model), observations from the far past become less important when the covariance matrix is time-variant, and it becomes imperative to find the block size B that can be used to recover the principal components.

(b) (Algorithm Analysis; Section 5) We then analyze two algorithms to recover the principal components extensively used in the standard streaming PCA setting; the noisy power method [28] and Oja's algorithm [48]. These algorithms represent two very different design principles for computing principal components from data; processing data in blocks vis-à-vis single observations (However, note that both algorithms perform computations in a streaming fashion, i.e., have $\mathcal{O}(kp)$ storage requirements). We first determine the optimal choice for critical parameters: the block size in the noisy power method on Lemma 1 and the learning rate in Oja's algorithm on Lemma 2. Next, we leverage these results in order to obtain an upper bound on the convergence error for these algorithms in Theorem 2 and Theorem 3 respectively. From these results, we have that: (i) the block size B for the noisy power method and inverse of the learning rate, $1/\zeta$ for Oja's algorithm play a similar role even in the non-stationary environment; (ii) In either case, the optimal parameters scale with $\sim \Gamma^{-2/3}$, where Γ is perturbation budget which will be defined later, and (iii) the derived upper bound on the estimation error for noisy power method matches the minimax error in terms of p, δ, and Γ and becomes rate-optimal if it satisfies mild condition.

Notation We fix notation throughout the main body of the paper. Matrices are denoted by bold uppercase letters ($\mathbf{M}$) and vectors are denoted by bold lowercase letters ($\mathbf{x}$). Let $\mathbf{M} \in \mathbb{R}^{p \times k}$. Then $\mathbf{M}_i$ denote the ith column of $\mathbf{M}$, $\mathbf{M}_{i:j}$ denote the submatrix consists with $i^{\text{th}} \sim j^{\text{th}}$ column of $\mathbf{M}$ and $\mathbf{M}_{i,j}$ be the (i,j)-element of $\mathbf{M}$. $\|\cdot\|$ denote matrix 2-norm or equivalently operator norm for matrices and standard 2-norm for vectors. We use $\text{St}_k(\mathbb{R}^p)$ for the class of orthogonal matrices in $\mathbb{R}^{p \times k}$. $b(\mathbf{M})$ is the orthogonal matrix where the columns form basis for $\text{ran}(\mathbf{M})$. $\lambda_i(\mathbf{M})$ represents the ith largest eigenvalue of the matrix $\mathbf{M}$ and and $s_i(\mathbf{M})$ be the ith largest singular value. The singular value decomposition of $\mathbf{M}$ is defined as $\text{SVD}(\mathbf{M}) = \mathbf{UDV}^\top$, where $\mathbf{U} \in \text{St}_p(\mathbb{R}^p)$, $\mathbf{V} \in \text{St}_k(\mathbb{R}^k)$, and $\mathbf{D} \in \mathbb{R}^{p \times k}$ is a diagonal matrix whose i^{th} diagonal element equals $s_i(\mathbf{M})$. Therefore, we assume without loss of generality that the singular values and respective singular vectors are ordered from largest to smallest. We denote by $\mathbf{M}_\perp$ the orthogonal projection onto the orthogonal complement of

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[1]For precise explanation for δ and Γ, please refer to the Section 3.

the range of $\mathbf{M}$. Therefore, if $\text{rk}(\mathbf{M}){=}k$ and $\text{SVD}(\mathbf{M}){=}\ \mathbf{U}\mathbf{D}\mathbf{V}^\top$, $\mathbf{M}_\perp$ given by $\mathbf{M}_\perp{=}\mathbf{I} - \mathbf{U}_{1:k}\mathbf{U}_{1:k}^\top$.
Moreover, when $\text{rk}(\mathbf{M}){=}\text{rk}(\tilde{\mathbf{M}})$ and $\text{SVD}(\tilde{\mathbf{M}}){=}\ \tilde{\mathbf{U}}\tilde{\mathbf{D}}\tilde{\mathbf{V}}^\top$, the distance between $\text{ran}(\mathbf{M})$ and $\text{ran}(\tilde{\mathbf{M}})$
is defined by:

$$d(\text{ran}(\mathbf{M}), \text{ran}(\tilde{\mathbf{M}})) = \parallel b(\mathbf{M})b(\mathbf{M})^\top - b(\tilde{\mathbf{M}})b(\tilde{\mathbf{M}})^\top \parallel\ =\ \parallel \mathbf{U}_{1:k}\mathbf{U}_{1:k}^\top - \tilde{\mathbf{U}}_{1:k}\tilde{\mathbf{U}}_{1:k}^\top \parallel .$$

We denote $d(\text{ran}(\mathbf{M}), \text{ran}(\tilde{\mathbf{M}}))$ by $d(\mathbf{M}, \tilde{\mathbf{M}})$ whenever clear from the context. For $\mathcal{A}{=}(\mathbf{A}_1, \ldots, \mathbf{A}_T)$, we denote the expectation of $f(\mathbf{x}_1, \ldots, \mathbf{x}_T)$ over $\mathbf{x}_t \sim \mathcal{N}(\mathbf{0}, \mathbf{A}_t\mathbf{A}_t^\top + \sigma^2\mathbf{I})$ as $\mathbb{E}_{\mathcal{X}\sim\mathcal{A}}[f]$. We denote $\tilde{\mathcal{O}}$ as the big-$\mathcal{O}$ notation with ignore the multiplicative dependency of $\log(p) + 2\log(T)$ and $1 + \Theta(\log(T/\Delta)/k)$.

2   Related Work

Principal component analysis has been extensively studied across operations research, computer science, and other communities. We highlight how our work differs from existing literature. **Robust PCA** Robust PCA deals with the problem of retrieving the principal components robust to the presence of outliers in the data. The cornerstone work in this direction is the Principal Component Pursuit framework wherein they assume that the matrix of observations $\mathbf{M}$ can be decomposed in terms of a low-rank matrix $\mathbf{L}$ and a sparse matrix (with entries of arbitrarily large magnitude) $\mathbf{S}$ [14].

Several works [23, 25, 16, 46] consider the robust PCA problem and propose algorithms in the offline, batch, and online setting. Our work differs from this line of literature in two aspects: assumptions about the data generating model and convex optimization techniques. First, our data generation model is unrelated to those considered in the robust PCA literature. Further, rich theories from convex optimization can be used in the PCA framework to design efficient algorithms, but our problem is not amenable to such techniques.

Streaming PCA Streaming algorithms for PCA have been proposed among other works on PCA [43, 34, 28, 59, 2]. Algorithms analyzed in this work, such as the noisy power method [28] and Oja's algorithm [48], are iterative methods for estimating the principal components. These iterative schemes are instances of stochastic approximation-based solutions for the optimization version of the PCA problem [4]. The stochastic approximation is a root-finding framework extensively used for stochastic optimization [37, 10]. Oja's algorithm, originally proposed by [48], was the first such scheme.

This framework is also used to analyze gradient-type and proximal-type incremental methods akin to algorithms for convex optimization. Stochastic gradient descent-based algorithms for the streaming PCA problem, where a single observation is used at every point in time to update the principal components' estimate, have been extensively studied. Along this line, [29] propose GRASTA, an incremental online gradient method for learning over different subspaces. Similarly, [5] proposes GROUSE, based on the idea of gradient updates over the Grassmannian manifold. In these and other related works [53, 15, 60, 30, 30, 1, 58] consider Oja's algorithm for the standard streaming PCA problem. However, all these works are based on a completely different modeling assumption than this paper and do not provide theoretical guarantees for our setting.

Streaming and robust PCA algorithms are used in presence of outliers or data with a lot of missing entries [17, 54]. From this literature, closest to our work is the work on robust subspace tracking [2, 59, 29]. However, the robustness considered their is against erasures or sparse outliers. While these algorithms provide theoretical guarantees in that setting, those guarantees cannot be extended to our model. Considerations of erasures and outliers under the model proposed in this paper is beyond the scope of this work and remains an interesting future direction.

3   Mathematical Framework

Observations from the standard spiked covariance model belong to a fixed k-dimensional column space $\mathbf{A} \in \mathbb{R}^p$. Previous work has focused on reconstructing this space from the observed time series. Under our framework, we assume that the sequence of observations $(\mathbf{x}_t)_{t=1}^T$ follows the time-dependent spiked covariance model (1). We consider the problem of computing top-k singular vectors of $\mathbf{A}_T$. We formulate this model by addressing a minimax optimization as a robust optimization problem.

When the adversary is allowed to select a completely arbitrary sequence of matrices $\mathbf{A}_1, \ldots, \mathbf{A}_T$, it is impossible to accurately recover the column space of $\mathbf{A}_T$ since only the single vector $\mathbf{x}_T$ provides information about this space. Instead, we define temporal uncertainty sets to restrict the power of the adversary.

Definition 1. Let $\Gamma, \delta \geq 0$. We only allow the sequence of matrices $\mathbf{A}_1, \ldots, \mathbf{A}_T$ that lie in an temporal uncertainty set $\text{Tu}(\delta, \Gamma)$ defined as:

$$\text{Tu}(\delta, \Gamma) := \left\{ (\mathbf{A}_t)_{t=1}^{T} \ : \ s_k(\mathbf{A}_t \mathbf{A}_t^\top) \geq \delta \ , \ \max_{2 \leq t \leq T} \|\mathbf{A}_t \mathbf{A}_t^\top - \mathbf{A}_{t-1} \mathbf{A}_{t-1}^\top\| \leq \Gamma \right\} . \tag{2}$$

That is, we constrain the difference between any two consecutive covariance matrices of the underlying (noise-free) process. The assumption $s_k(\mathbf{A}_t \mathbf{A}_t^\top) \geq \delta$ is crucial for establishing bounds of estimation error. It is also justified in any setting where the underlying phenomenology guarantees covariance of rank k; the stated lower bound excludes pathological cases of near-rank $k - 1$ or smaller.

Having described the constraints on the perturbation power, we now present the family of algorithms and the performance metric of interest. In particular, we consider algorithms that are (a) streaming, i.e., make a single pass through the time series in chronological order; only the previous samples are stored at any point in the past. These characteristics model the behavior of an algorithm that receives data in real-time and only stores samples in limited memory. Among those streaming algorithms, we consider the algorithms whose (b) output is a set of k orthonormal vectors in $\mathbb{R}^p$. The first requirement ensures that the storage and computational memory requirements are independent of the number of samples T, and hence, the assumption $T \gg p$ is not prohibitive. Let us denote by Φ the family of algorithms just described. Let $\mathcal{A} = (\mathbf{A}_1, \ldots, \mathbf{A}_T) \in \text{Tu}(\delta, \Gamma)$ be given. We formalize this definition as follows:

Definition 2. $\Sigma(\mathcal{A})$ is the time-dependent spiked covariance model represented by the $\mathcal{A} = (\mathbf{A}_1, \ldots, \mathbf{A}_T)$. Furthermore, $\mathcal{X} := (\mathbf{x}_1, \ldots, \mathbf{x}_T) \sim \Sigma(\mathcal{A})$ is a sequence of observation vectors drawn from $\Sigma(\mathcal{A})$.

The output of algorithm ϕ for a given choice of the observations $\mathbf{x_t}$ sampled using model (2) is a set of k orthonormal vectors which we denote by $\phi_\mathcal{X}$; abusing notation we will also view $\phi_\mathcal{X}$ as a matrix in $\mathbb{R}^{p \times k}$ or the subspace generated by $\phi_\mathcal{X}$.

Definition 3. Let $\phi \in \Phi$, $\mathcal{A} \in \text{Tu}(\delta, \Gamma)$ and $\mathcal{X} \sim \Sigma(\mathcal{A})$.

(a) Given $\mathcal{A}$ and $\mathcal{X}$, the estimation error of ϕ given $\mathcal{X}$ is the distance between the space spanned by $\phi_\mathcal{X}$ and the column space of $\mathbf{A}_T$ ($=d(\text{ran}(\mathbf{A}_T), \phi_\mathcal{X})$). The metric can be easily extended to the cumulative error. Please see the Appendix A for the discussion.

(b) $\mathcal{R}^\phi$ is the maximum expected estimation error of ϕ under $\mathcal{A}$ over all $\mathcal{A} \in \text{Tu}(\delta, \Gamma)$. Precisely,

$$\mathcal{R}^\phi := \sup_{\mathcal{A} \in \text{Tu}(\delta, \Gamma)} \mathbb{E}_{\mathcal{X} \sim \mathcal{A}} \left[d\big(\text{ran}(\mathbf{A}_T), \phi_\mathcal{X}\big) \right] .$$

(c) $\mathcal{R}^*$ is the minimax estimation error defined as the minimum of the largest expected estimation error incurred by $\phi \in \Phi$:

$$\mathcal{R}^* := \inf_\phi \mathcal{R}^\phi = \inf_\phi \sup_{\mathcal{A} \in \text{Tu}(\delta, \Gamma)} \mathbb{E}_{\mathcal{X} \sim \mathcal{A}} \left[d\big(\text{ran}(\mathbf{A}_T), \phi_\mathcal{X}\big) \right] .$$

We note that minimax estimation problems have been widely studied [56, 13, 57].

(d) An algorithm $\phi \in \Phi$ is **rate-optimal** if $\mathcal{R}^\phi \leq C \cdot \mathcal{R}^*$, where constant $C > 0$ is independent of the problem parameters Γ, T, δ, p, k. In this paper we establish the minimax estimation error $\mathcal{R}^*$ and propose rate-optimal sublinear-time, single-pass algorithms for robust streaming PCA.

4 Minimax Lower Bound

When $\mathbf{A}_t$ belongs to the temporal uncertainty set $\text{Tu}(\delta, \Gamma)$ (Definition 1) an algorithm designed to recover the principal components of $\mathbf{A}_T$ from the observations cannot guarantee zero estimation error. Our first goal is to obtain a lower bound on the estimation error in terms of the problem parameters

δ, Γ, p and the number of observations T. We now establish a lower bound on the expected error in recovering the top singular vector of $\mathbf{A}_T$ from observations $\mathcal{X} = (\mathbf{x}_t)_{t=1}^{T}$.

In order to establish the lower bound, we leverage the fundamental limit of hypothesis tests [56]. The crux of the proof lies in constructing the set of worst-case hypotheses and establishing a lower bound on the probability of error in distinguishing between these hypotheses using observed data. The complete proof is provided in Appendix C.

Theorem 1 (Lower Bound). *Assume $0 \leq \Gamma < \delta$ and $T \geq 1$. For any algorithm $\phi \in \Phi$, there exists a sequence $\mathcal{A} \in \text{Tu}(\delta, \Gamma)$ and constant $C > 0$ which is independent with Γ, δ, T, p, k such that:*

$$\mathbb{E}_{\mathcal{X} \sim \mathcal{A}}\big[d\big(\text{ran}(\mathbf{A}_T), \phi_{\mathcal{X}}\big)\big] \geq \min\left\{1, C\bigg(\Big(\tfrac{\Gamma}{\delta}\Big)^{1/3}\Big(\tfrac{p\sigma^2(\sigma^2+\delta)}{\delta^2}\Big)^{1/3} + \tfrac{1}{\sqrt{T}}\Big(\tfrac{p\sigma^2(\sigma^2+\delta)}{\delta^2}\Big)^{1/2}\bigg)\right\} . \quad (3)$$

For the standard streaming PCA problem (Theorem 1 with the case of $\Gamma = 0$), the fundamental limit is $\Theta((\sigma/\delta)(p(\sigma^2+\delta)/T)^{1/2})$ which has the expected $\Theta(1/\sqrt{T})$ dependence [13, 57]. On the other hand, in presence of perturbations ($\Gamma > 0$), the lower bound exhibits a phase transition phenomenon, with the first term representing the effect of model ambiguity. To this end, define the critical time T_c as

$$T_c := \left(\frac{\Gamma}{\delta}\right)^{-2/3}\left(\frac{p\sigma^2(\sigma^2+\delta)}{\delta^2}\right)^{1/3} . \quad (4)$$

For $T = \mathcal{O}(T_c)$, the lower bound decreases with the rate of $1/\sqrt{T}$. However, when $T = \Omega(T_c)$, the first term dominates the second term and $\mathcal{R}^*$ becomes independent to the number of observations T. In this regime, the error stagnates to $\mathcal{O}((\Gamma/\delta)^{1/3}\left(p\sigma^2(\sigma^2+\delta)/\delta^2\right)^{1/3})$. This validates our intuition that in a dynamic environment past information quickly becoming stale.

In Theorem 2 and 3, we will prove that the noisy power method and Oja's algorithm attain a near-optimal bound on the convergence guarantee. Theorem 2 guarantees that if $s_1(\mathbf{A}_t\mathbf{A}_t^\top) = \delta$, the upper bound for estimation error on the noisy power method is of the following order:

$$\tilde{\mathcal{O}}\left(\left(\frac{\Gamma}{\delta}\right)^{1/3}\left(\frac{(p\sigma^2 + k\delta)(\sigma^2+\delta)}{\delta^2}\right)^{1/3}\right) . \quad (5)$$

That is, if $p\sigma^2$ dominates $k\delta$, the noisy power method becomes rate-optimal under the controlled uncertainty with $\text{Tu}(\delta, \Gamma)$. This regime is the case of practical noisy situations, with $\sigma^2 \lll \delta$. Observe also that we have $p > k$ from the setting.

5 Convergence Analysis

In this section, we analyze two algorithms for the robust streaming PCA problem. A generic template for algorithms $\phi \in \Phi$ of interest to us is as follows: (i) ϕ is initialized with a random matrix with orthonormal columns $\hat{\mathbf{U}} \in \mathbb{R}^{p \times k}$; (ii) A running estimate of the principal components is maintained as the columns of $\hat{\mathbf{U}}$; (iii) Observations are projected onto the column space of $\hat{\mathbf{U}}$ to update this estimate.

We consider two algorithms: robust version of the noisy power method (Algorithm 1) and Oja's algorithm (Algorithm 2). The key difference between the noisy power method and Oja's algorithm is in the data used to estimate the principal components. In the noisy power method, the estimates are updated after a batch of observations, whereas in Oja's algorithm, the estimates are updated after scaling every observation with the learning rate. Therefore, the parameters determining the performance of these algorithms are the batch size B for the robust power method and the learning rate ζ for Oja's algorithm.

The analysis of these algorithms cannot be readily established with existing techniques when the covariance matrix belongs to a temporal uncertainty set since they rely on showing that the estimates improve every iteration. Further, the application of many concentration results requires random matrices to be bounded which is not the case when the observations are sampled from (1). Therefore, in order to simplify the analysis on both algorithms, we introduce Assumption 1, adapted from [34].

Assumption 1. Let $(\mathbf{A}_t)_{t=1}^{T} \in \text{Tu}(\delta, \Gamma)$. For $\tilde{\delta} \geq \delta$ and $\mathcal{M}, \mathcal{V} > 0$, we consider the observations $\mathbf{x}_t$ for $t \in [T]$ satisfy the following:

(a) $\mathbb{E}[\mathbf{x}_t\mathbf{x}_t^\top] = \mathbf{A}_t\mathbf{A}_t^\top + \sigma^2\mathbf{I}_{p\times p}$ while $\|\mathbf{A}_t\mathbf{A}_t^\top\| \leq \tilde{\delta}$,

(b) $\|\mathbf{x}_t\mathbf{x}_t^\top - (\mathbf{A}_t\mathbf{A}_t^\top + \sigma^2\mathbf{I}_{p\times p})\| \leq \mathcal{M}$ almost surely, and

(c) $\left\|\mathbb{E}\left[(\mathbf{x}_t\mathbf{x}_t^\top - (\mathbf{A}_t\mathbf{A}_t^\top + \sigma^2\mathbf{I}_{p\times p}))^2\right]\right\| \leq \mathcal{V}$.

When the observations follow equation (1), we condition our analysis on event $\mathfrak{E}_\Delta$.

Definition 4. For $\mathbf{z}_t \sim \mathcal{N}(\mathbf{0}_{k\times 1}, \mathbf{I}_{k\times k})$ and $\mathbf{w}_t \in \mathcal{N}(\mathbf{0}_{p\times 1}, \mathbf{I}_{p\times p})$, the event $\mathfrak{E}_\Delta$ is defined as:

$$\mathfrak{E}_\Delta \; : \; \begin{cases} \|\mathbf{z}_t\|^2 \leq k + 2\sqrt{k\log(2T/\Delta)} + 2\log(2T/\Delta)\,, \text{ and} \\[2mm] \|\mathbf{w}_t\|^2 \leq p + 2\sqrt{p\log(2T/\Delta)} + 2\log(2T/\Delta) \end{cases} \quad \forall t \in [T]\,. \tag{6}$$

The observations from model (1) satisfy Assumption 1 with $\mathcal{M} = 2(k\tilde{\delta}+p\sigma^2)(1+\Theta(\log(T/\Delta)/k))$, $\mathcal{V} = 2\mathcal{M}(\tilde{\delta}+\sigma^2)$ with probability $\mathbb{P}[\mathfrak{E}_\Delta] \geq 1 - \Delta$. Although after conditioning $\mathbb{E}[x_tx_t^\top|\mathfrak{E}_\Delta] \neq \mathbf{A}_t\mathbf{A}_t^\top + \sigma^2\mathbf{I}_{p\times p}$, we can use all the results in Section 5.1 and 5.2 with a sufficiently small Δ. Please refer to Appendix H for the details.

5.1 Noisy Power Method

Algorithm 1 Noisy Power Method with block size B [28]

1: **Input:** $\{\mathbf{x}_1\}_{t=1}^T$, Block Size: B, Dimensions: p, k & **Output:** $\hat{\mathbf{U}}(L)$
2: $\hat{\mathbf{U}}(0) \leftarrow p \times k$ random matrix
3: **for** $\ell = 1 : L = T/B$ **do**
4: $\quad \mathbf{Y}(\ell) \leftarrow \mathbf{0} \in \mathbb{R}^{p\times k}$
5: $\quad$ **for** $t = (\ell-1)B+1 : \ell B$ **do**
6: $\quad\quad \mathbf{Y}(\ell) \leftarrow \mathbf{Y}(\ell) + \frac{1}{B}\mathbf{x}_t\mathbf{x}_t^\top\hat{\mathbf{U}}(\ell-1)$
7: $\quad$ **end for**
8: $\quad \hat{\mathbf{U}}(\ell) \leftarrow$ Gram-Schmidt$(\mathbf{Y}(\ell))$
9: **end for**

The noisy power method is an iterative algorithm for computing the top-k principal components of given matrix. Starting from the random matrix $\hat{\mathbf{U}}(0)$ in $\mathbb{R}^{p\times k}$, the algorithm runs for L iterations, each processing B samples. By repeating this procedure, we expect the algorithm to reconstruct the covariance matrix if the observations are derived from the fixed distribution.

When observations are drawn from the model 1 under Assumption 1, the observations belong to different underlying subspace. Unlike the standard case, for any algorithm ϕ, the presence of perturbation prevents the convergence of the columns of $\hat{\mathbf{U}}(\ell)$ to the singular vectors of the final covariance matrix $\mathbf{U}(\ell)$. The main difficulty here is that the columns of $\hat{\mathbf{U}}(\ell)$ do not converge towards a fixed set of vectors but keep tracking the time-varying principal components. Recall that our ultimate objective is to recover the principal components associated with the *terminal* observation. Hence, we decompose the covariance matrix in terms of the last observation and the remaining samples. For ℓ-th block, we denote the covariance matrix for ℓB-th observation as $\mathbf{M}(\ell)$ and have:

$$\frac{1}{B}\sum_{t=(\ell-1)B+1}^{\ell B} \mathbf{x}_t\mathbf{x}_t^\top = \mathbb{E}[\mathbf{x}_{\ell B}\mathbf{x}_{\ell B}^\top] + \mathcal{E}(\ell) = \underbrace{\mathbf{A}_{\ell B}\mathbf{A}_{\ell B}^\top + \sigma^2\mathbf{I}}_{\mathbf{M}(\ell)} + \mathcal{E}(\ell)\,. \tag{7}$$

Due to presence of perturbations in the robust model, $\mathcal{E}(\ell)$ is a non-zero mean random variable. We first decompose $\mathcal{E}(\ell)$ in terms of the contribution due to inherent noise and the perturbations allowed by the robust model. In Lemma 1, we decompose the error $\|\mathcal{E}(\ell)\|$ with respect to the block size B and the allowed perturbations in the robust model. The complete proof is provided in Appendix D.

Lemma 1 (Spectral norm of noise). *Assume that the observations* $(\mathbf{x}_t)_{t=1}^T$ *generated according to the Assumption 1. For probability greater than* $1 - \frac{1}{T}$, *the matrices* $\mathcal{E}(\ell)$ *in the equation* (7) *are bounded by:*

$$\max_{1\leq\ell\leq L} \|\mathcal{E}(\ell)\| \leq \frac{1+3\sqrt{2}}{3}\sqrt{\frac{\mathcal{V}\log(2pT^2)}{B}} + \frac{B\Gamma}{2}\,. \tag{8}$$

252 *when the block size B is larger than $\mathcal{M}^2 \log(2pT^2)/\mathcal{V}$.*

253 Lemma 1 highlights the effect of allowing perturbations in the data generation model. From classical
254 results in statistics, our intuition tells that the effect of noise washes out as the block size increases,
255 i.e., the error decays with the 'inverse square root' of the block size. Hence, barring memory and
256 data issues, a larger block size is better when $\Gamma = 0$. In contrast, covariance perturbations add error
257 proportional to the block size. Consequently, we have a trade-off between two terms in this case, and
258 there exists an optimal block size depending on the Γ.

259 We establish the convergence guarantee of the robust power method in Theorem 2. In the proof of
260 Theorem 2 (Appendix E), we bound the distance, $d\big(\mathbf{U}(L), \hat{\mathbf{U}}(L)\big)$ between the output of Algorithm 1,
261 $\hat{\mathbf{U}}_{1:k}(L)$ and k-orthonormal vectors spanning the column space of $\mathbf{A}_T$, $\mathbf{U}_{1:k}(L)$. We identify the
262 optimal block size B, which is the unique parameter for noisy power method and establish an upper
263 bound on the estimation error of the noisy power method.

264 **Theorem 2** (Robust power method). *Assume that $\delta \geq 0.71\sigma^2$ and $\Gamma = \mathcal{O}(\delta^3/(\mathcal{V}\log(2pT^2)))$. When*
265 *the observations $\{\mathbf{x}_t\}_{t=1}^T$ satisfies the assumption 1, for $B = \Theta(\mathcal{V}^{1/3}\log(2pT^2)^{1/3}/\Gamma^{2/3})$ we have:*

$$\mathcal{R}^{\mathrm{NPM}} = \mathcal{O}\left(\frac{(\mathcal{V}\Gamma\log(2pT^2))^{1/3}}{\delta} + 0.7^{T/B}\right), \tag{9}$$

266 *with probability $1 - 1/T - c^{\Omega(p-k+1)} - \exp(-\Omega(p))$.*

267 When $\Gamma = \Omega(\delta^3/(\mathcal{V}\log(2pT^2)))$ we have that $\frac{(\mathcal{V}\Gamma\log(2pT^2))^{1/3}}{\delta} = \Omega(1)$. Therefore, the condition
268 on Γ in Theorem 2 is necessary to avoid a trivial upper bound (≤ 1) on $\mathcal{R}^{\mathrm{NPM}}$. This condition en-
269 compasses several applications of interest alluded to earlier. For example, in the financial applications
270 alluded to earlier, individual market changes of interest happen on millisecond time-scale and it is of
271 significant interest to terminally detect cumulative market changes. Our results hold for sufficiently
272 large value of cumulative changes (allowing for seconds time-scale) allowing us to study them. They
273 further imply that the noisy power method is rate-optimal for non-trivial values of Γ. When the
274 observations follow model (1), under the event $\mathfrak{E}_\Delta$, Theorem 2 shows that the robust power method
275 can achieve a estimation error of:

$$\tilde{\mathcal{O}}\left(\left(\frac{\Gamma}{\delta}\right)^{1/3}\left(\frac{(p\sigma^2 + k\delta)(\sigma^2 + \delta)}{\delta^2}\right)^{1/3}\right), \tag{10}$$

276 with probability $1 - \Delta - 1/T - c^{\Omega(p-k+1)} - \exp(-\Omega(p))$ if $\tilde{\delta}$ equals to δ. This is order-wise identical
277 to the fundamental limit established in Theorem 1 when $p\sigma^2$ dominates $k\delta$.

278 Establishing bounds on the estimation error when the underlying singular vectors change is intricate
279 since the subspace to which consecutive observations belong are potentially different. Conventional
280 proofs that analyze noisy power methods or Oja's algorithm show that under a variety of assumptions
281 at every iteration ℓ, the distance between the estimated and true subspace, $d(\mathbf{U}(\ell), \hat{\mathbf{U}}(\ell))$ decreases.
282 For instance, the proof in [28] requires $\|\mathcal{E}(\ell)\mathbf{U}(\ell)\| = O(\sqrt{k/p})$ which does not hold under our
283 model since $\|\mathbf{M}(\ell) - \mathbf{M}(\ell-1)\|$ is in general greater than $\sqrt{k/p}$. Similarly, the concentration
284 approach in [32] can be used only when the covariance matrix is time-invariant. We briefly describe
285 our proof technique to establish Theorem 2, deferring details to Appendix E. Let $\mathcal{M}^{(L)}$ denote the
286 product $\prod_{\ell=1}^L \big(\mathbf{M}(\ell) + \mathcal{E}(\ell)\big)$. Then, the output of the algorithm $\hat{\mathbf{U}}_{1:k}(L)$ is an orthonormal basis
287 of $\mathcal{M}^{(L)}\hat{\mathbf{U}}(0)$, which estimates the first k principal components $\mathbf{U}_{1:k}$ of the $\mathbf{M}(L)$. To address
288 this, we construct sequences of k and $(p-k)$-dimensional subspaces of $\mathbb{R}^p$ from observations
289 $\{(\mathbf{x}_t)_{t=(\ell-1)B}^{\ell B}\}_{\ell=1}^L$, denoted by $\{\mathbf{N}^{(\ell)}\}_{\ell=1}^L \in \mathrm{St}_{p-k}(\mathbb{R}^p)$ and $\{\mathbf{W}^{(\ell)}\}_{\ell=1}^L \in \mathrm{St}_k(\mathbb{R}^p)$ respectively,
290 such that for all iterations ℓ:

291 (a) $\mathrm{ran}\big((\mathbf{M}(\ell)+\mathcal{E}(\ell))\mathbf{N}^{(\ell)}\big) \subseteq \mathrm{ran}\big(\mathbf{N}^{(\ell+1)}\big)$ and $\mathrm{ran}\big((\mathbf{M}(\ell)+\mathcal{E}(\ell))\mathbf{W}^{(\ell)}\big) \subseteq \mathrm{ran}\big(\mathbf{W}^{(\ell+1)}\big)$,

293 (b) $\|(\mathbf{M}(\ell) + \mathcal{E}(\ell))\mathbf{N}^{(\ell)}\|\|((\mathbf{M}(\ell) + \mathcal{E}(\ell))\mathbf{W}^{(\ell)})^{-1}\| < 1$,

295 (c) $d(\mathbf{U}_{1:k}(\ell), \mathbf{W}^{(\ell+1)}) = O(\|\mathcal{E}(\ell)\|)$ and $d(\mathbf{U}_{k+1:p}(\ell), \mathbf{N}^{(\ell+1)}) = O(\|\mathcal{E}(\ell)\|)$.

The initial random matrix $\hat{\mathbf{U}}_{1:k}(0)$ consists of both $\mathbf{N}^{(1)}$ and $\mathbf{W}^{(1)}$ with high probability. From the first two properties, at every iteration ℓ, the projection of $\hat{\mathbf{U}}_{1:k}(\ell-1)$ in $\mathbf{W}^{(\ell)}$ is amplified more than that on $\mathbf{N}^{(\ell)}$ and thus $\hat{\mathbf{U}}_{1:k}(L)$ becomes very close to $\mathbf{W}^{(L)}$ after sufficiently large L. From the last property, we can conclude that $\mathbf{W}^{(L)}$ is close to $\mathbf{U}_{1:k}(L)$, where the distance between $\mathbf{W}^{(L)}$ and $\mathbf{U}_{1:k}(L)$ is proportional to $\|\mathcal{E}(L)\|$. Combining these ideas with Lemma 1 establishes the results.

5.2 Oja's Algorithm

Algorithm 2 Oja's Algorithm with learning rate ζ [48]

1: **Input:** $\{\mathbf{x}_1\}_{t=1}^T$, Learning rate: ζ, dimensions: p, k & **Output:** $\hat{\mathbf{U}}(T)$
2: $\hat{\mathbf{U}}(0) \leftarrow p \times k$ random matrix
3: **for** $t = 1 : T$ **do**
4: $\quad \hat{\mathbf{U}}(t) \leftarrow$ Gram-Schmidt$((\mathbf{I} + \zeta \mathbf{x}_t \mathbf{x}_t^\top)\hat{\mathbf{U}}(t-1))$
5: **end for**

We now establish the convergence guarantee for Oja's Algorithm (Algorithm 2) when observations follow (1). Unlike the noisy power method, Oja's Algorithm is multiplicative in its construction of the estimated subspace. We extend the existing analysis for Oja's algorithm [34, 32, 3] by considering a virtual block with $B = B_\zeta = \zeta^{-1}$ observations. Building upon the analysis framework for the noisy power method and intuition from binomial approximation $(1 + \zeta x)^{1/\zeta} \simeq (1 + x)$ we establish the convergence guarantees for Oja's algorithm. Like the previous section, we decompose the block with target $\mathbf{M}(\ell)$ and error matrix $\mathcal{E}'(\ell)$ as:

$$\prod_{t=(\ell-1)B+1}^{\ell B}(\mathbf{I} + \zeta\,\mathbf{x}_t\mathbf{x}_t^\top) = \prod_{t=(\ell-1)B+1}^{\ell B}(\mathbf{I} + \zeta\,\mathbb{E}\,\mathbf{x}_{\ell B}\mathbf{x}_{\ell B}^\top) + \mathcal{E}(\ell) = \mathbf{M}_{\text{Oja}}(\ell) + e^{\tilde{\delta}+\sigma^2}\,\mathcal{E}'(\ell)\,. \tag{11}$$

where $\mathcal{E}'$ is scaled error matrix. In Lemma 2, we provide bound for scaled error matrix with respect to the learning parameter ζ, or the virtual block size $B_\zeta = 1/\zeta$ similar to Lemma 1. The proof is provided in Appendix F.

Lemma 2 (Spectral norm of noise, Oja's algorithm case)**.** *Assume that the observations* $(\mathbf{x}_t)_{t=1}^T$ *generated according to the Assumption 1. For probability greater than* $1 - \frac{1}{T}$, *the matrices* $\mathcal{E}'(\ell)$ *on the equation* (11) *are bounded by:*

$$\max_{1\le\ell\le L}\|\mathcal{E}'(\ell)\| \le \sqrt{\frac{2e^2\mathcal{M}^2\log(pT^2)}{B_\zeta}} + \frac{B_\zeta\Gamma}{2} + \mathcal{O}(B_\zeta\Gamma)\,, \tag{12}$$

when the virtual block size $B_\zeta = \zeta^{-1}$ *is larger than* $2\mathcal{M}^2\log(pT^2)$.

The difference in parameter $\mathcal{M}$, rather than $\mathcal{V}$ as in the case of Lemma 1 arises due to use of a different concentration inequality. The estimator of the noisy power method averages the outer product of vectors (equation (11)) making it straightforward to use Bernstein's inequality. On the other hand, Oja's algorithm averages the product of random matrices rather than the sum of random matrices. Therefore, we introduce the multiplicative concentration inequality [31], which requires the (probabilistic) norm bound for matrices. Combining multiplicative concentration inequalities with Lemma 2 and the techniques developed for noisy power method, we obtain a convergence guarantees for Oja's algorithm in Theorem 3. The complete proof is provided in Appendix G.

Theorem 3 (Oja's algorithm)**.** *Assume that* $\delta \ge 0.71$ *and* $\Gamma = \mathcal{O}(\delta^3/(e^{3(\tilde{\delta}-\delta)}\mathcal{M}^2\log(pT^2)))$. *When the observations* $\{\mathbf{x}_t\}_{t=1}^T$ *satisfies the assumption 1, for* $1/\zeta = \Theta(\mathcal{M}^{2/3}\log(pT^2)^{1/3}/\Gamma^{2/3})$ *we have:*

$$\mathcal{R}^{\text{Oja}} = \tilde{\mathcal{O}}\left(e^{\tilde{\delta}}\frac{(\mathcal{M}^2\Gamma\log(pT^2))^{1/3}}{\delta} + 0.7^{\zeta T}\right)\,, \tag{13}$$

with probability $1 - 1/T - c^{\Omega(p-k+1)} - \exp(-\Omega(p))$.

328 Unlike the noisy power method the upper bound in Theorem 3 is $O(p^{2/3})$ rather than optimal-
329 dependence of $O(p^{1/3})$ from Theorem 1. It is unclear whether the upper bound can be improved.
330 Sharpening our analysis with a two-phase strategy [32, 39] (wherein the first phase identifies a good
331 initial point and the second phase establishes convergence given an initial point) might be an excellent
332 direction for future investigation.

6 Numerical Results

334 Lastly, in order to verify the established results for both algorithms, we examine the performance of
335 algorithms for a variety of environments . Theorem 2 and 3 show the **(a)** existence of the optimal block
336 size B and the learning rate ζ to obtain the minimum recovery error, and **(b)** $\Gamma^{-3/2}$ dependencies of
337 those optimal B and $1/\zeta$. We demonstrate this phenomenon in Figure 1a and Figure 1b respectively.
Experimental details are described in Appendix I. Our first observation on the Figure 1a is that

(a) Distance between true column space of $\mathbf{A}_T$ ($\mathbf{U}_{1:k}$) (b) Empirical optimal block size B and learning rate
and estimated space at $t = T$ ($\hat{\mathbf{U}}_{1:k}$) as Γ varies. ζ as Γ varies.

Figure 1: Numerical results. We used the setting $(\sigma, \delta, p, k)=(0.15, 1.0, 100, 5)$.

339 convergence error decreases as the block size increases without any covariance perturbation ($\Gamma = 0$).
340 This behavior is expected since in the absence of rotation of underlying space, increasing the number
341 of past information results in better accuracy guarantees for the recovered space. However, if the
342 covariance perturbation exists ($\Gamma > 0$), there exists the optimal learning parameter, and we have the
343 smaller optimal block size as the adversary stronger. This phenomenon is also the expected behavior
344 since an increase in the adversarial budget implies that the past information becomes less relevant.
345 Our observations also corroborate our theoretical results in the Theorem 1. In Figure 1b, we focus
346 on the optimal value of the block size B and inverse of the learning rate $1/\zeta$ and its variation with
347 the perturbation budget Γ. We plot the empirically optimal learning rate for the case of the noisy
348 power method and Oja's algorithm with $\Gamma^{-2/3}$. We observe that the optimal block size and the
349 inverse learning rate is proportional to $\Gamma^{-2/3}$. This experimental dependency of $\Gamma^{-2/3}$ verifies the
350 theoretically prescribed results in Theorem 2 and Theorem 3.

7 Conclusion

352 In this paper, we analyzed the fundamental limit of the minimax error and convergence of the noisy
353 power method and Oja's algorithm on the streaming PCA settings with time-varying covariance.
354 Under this framework, when there exists no perturbation ($\Gamma = 0$), our theoretical result on a lower
355 limit coincides with the order of the traditional outcome. Furthermore, when perturbation exists
356 ($\Gamma > 0$), we have a non-avoidable positive minimax recovery error, although the time horizon
357 becomes arbitrarily long. Next, we found that the noisy power method order-wisely achieves this
358 recovery error and becomes rate-optimal if the $k\delta \ll p\sigma^2$. Finally, we showed that optimal block size
359 and inverse learning rate minimizing recovery error were proportional to $\Gamma^{-3/2}$ both theoretically
360 and experimentally. All our results are based on the spectral gap assumption. However, a recent line
361 of literature considers effective rank [11, 50] as a crucial metric. Establishing guarantees for robust
362 streaming PCA in terms of effective rank is an interesting direction for future work.

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