Online Agnostic Multiclass Boosting

Abstract

Boosting is a fundamental approach in machine learning that enjoys both strong theoretical and practical guarantees. At a high-level, boosting algorithms cleverly aggregate weak learners to generate predictions with arbitrarily high accuracy. In this way, boosting algorithms convert weak learners into strong ones. Recently, Brukhim et al. [4] extended boosting to the online agnostic binary classification setting. A key ingredient in their approach is a clean and simple reduction to online convex optimization, one that efficiently converts an arbitrary online convex optimizer to an agnostic online booster. In this work, we extend this reduction to multiclass problems and give the first boosting algorithm for online agnostic mutliclass classification. Our reduction also enables the construction of algorithms for statistical agnostic, online realizable, and statistical realizable multiclass boosting.

1 Introduction

Boosting is a fundamental technique in machine learning that cleverly aggregates the predictions of weak learners to produce a strong learner. Originally studied in the batch (realizable) PAC learning setting for binary classification, boosting has now been extended to a wide variety of settings, including multiclass classification, online learning, and agnostic learning. Recently, Brukhim et al. [4] extended boosting to online agnostic binary classification, marking the completion of boosting algorithms for all four regimes of statistical/online and agnostic/realizable binary classification. However, less can be said about multiclass classification, where boosting algorithms are only well studied under the assumption of realizability. Designing online agnostic multiclass boosting algorithms is important for several reasons. First, realizability is a very strong assumption in real-life: it requires the existence of an expert that *perfectly* labels the data. Second, the vast majority of classification tasks require more than two labels, a prominent example being image classification. Lastly, modern machine learning tasks often require sequential processing of data, making the design and development of online algorithms increasingly relevant. In this work we fill this gap in literature by studying online agnostic boosting for multiclass problems.

1.1 Main Results

We give the first weak learning conditions and algorithms for online agnostic multiclass boosting. A key idea of our algorithms is a reduction from boosting to online convex optimization (OCO), an idea borrowed from Brukhim et al. [4]. As a consequence of this reduction, we are also able to give algorithms for the three other settings of statistical agnostic, online realizable, and statistical realizable multiclass boosting. Finally, we give empirical results showcasing that our OCO-based boosting algorithms are fast and competitive with existing state-of-the-art multiclass boosting algorithms.

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1.2 Related Works

Boosting was first studied for binary classification under the realizable PAC learning setting [30, 13–15, 29] and then later extended to the agnostic PAC learning setting [12, 18, 24, 23, 22, 25, 2]. The success of boosting for binary classification led to significant interest in designing boosting algorithms for multiclass problems. As a result, several multiclass boosting algorithms were proposed for the realizable batch setting [16, 11, 19, 17], culminating in the work by Mukherjee and Schapire [27], who unified the previous approaches under a cost matrix framework. Beyond the batch setting, online boosting algorithms for binary [28, 7, 3] and multiclass classification [8, 21] have been designed assuming realizability (mistake-bound). More recently, Brukhim et al. [4] give the first online *agnostic* (regret-bound) boosting algorithm for binary classification, marking complete all four regimes of statistical/online and agnostic/realizable boosting for *binary classification*.

Moving to agnostic *multiclass* boosting, Brukhim et al. [5] study the resources required for boosting in the statistical/batch setting as the number of labels k grows. However, they consider an alternative model of boosting where the weak learner is a strong agnostic PAC learner for a simple "easy-to-learn" base hypothesis class, and the goal is to learn target concepts outside the base class. Specifically, they assume the target concept can be represented by weighted plurality votes over the base class. In this way, the weakness of a weak learner is manifested in the base hypothesis class. Instead, in our work, we consider the standard boosting model of fixing the base hypothesis class, and ask whether a weak learner's performance can be improved relative to the best fixed hypothesis in that class. Lastly, Abernethy et al. [1] study robust multiclass boosting, however, they consider a setting where the adversary can only perturb the input data, a weaker adversary compared to the agnostic setting.

2 Preliminaries and Notation

We first describe the basic setup for online agnostic multiclass boosting. There are k possible labels $[k] := 1, ..., k$ and k is known to the booster and the weak learners. The booster maintains N copies of a weak learner, $\mathcal{W}$, which themselves are online learning algorithms that sequentially process examples from instance space $\mathcal{X}$ and output predictions in $[k]$. At each iteration $t = 1, ..., T$, an adversary picks a labeled example $(x_t, y_t) \in \mathcal{X} \times [k]$ and reveals x_t to the booster. Once the booster observes the unlabeled data x_t, it gathers the weak learners' predictions and makes a final (possibly randomized) prediction $\hat{y}_t$. After observing the booster's final decision, the adversary reveals the true label y_t, and the booster suffers the loss $\mathbb{1}\{\hat{y}_t \neq y_t\}$ or equivalently, *gains* the reward $2\mathbb{1}\{\hat{y}_t = y_t\} - 1$. Finally, the booster, after observing the true label y_t, updates each weak learner. Throughout this paper, we abuse notation by letting an integer class label $\ell \in [k]$ also denote its corresponding basis vector $e_\ell \in \mathcal{B}_k$, where $\mathcal{B}_k$ is the set of length k basis vectors. In this way, the loss/gain of the booster in round t can be written as $1 - \hat{y}_t \cdot y_t$ and $2\hat{y}_t \cdot y_t - 1$ respectively. We also let Δ_k represent the k-dimensional probability simplex and $\Delta_{\frac{k}{\gamma}}$ represent the $\frac{1}{\gamma}$-scaled k-dimensional probability simplex for $\gamma \in (0, 1)$. Finally, we let $\mathbb{1}_k$ denote the k-dimensional ones vector.

Evaluation. Unlike the realizable setting, in the agnostic setting, we place no restrictions on how the stream of examples $x_1, ..., x_T$ are labelled. Thus, for a fixed hypothesis class $\mathcal{H} \subseteq [k]^{\mathcal{X}}$, the goal of the booster is to output predictions $\hat{y}_t$ such that the expected *regret*,

$$\mathbb{E}\left[\max_{h \in \mathcal{H}} \sum_{t=1}^{T} (2h(x_t) \cdot y_t - 1) - \sum_{t=1}^{T} (2\hat{y}_t \cdot y_t - 1)\right],$$

is minimized, where the expectation is over the randomness of the booster and that of the possibly adaptive adversary. Note, this is in contrast to the realizable setting where the stream is labelled by a $h \in \mathcal{H}$ and we wish the booster to minimize the (expected) number of *mistakes* (mistake-bound).

Agnostic Boosting. A key technique in agnostic boosting, first appearing in the work by Kanade and Kalai [24], is to update weak learners by feeding randomly *relabelled* examples. This is in contrast to the realizable setting where we typically update weak learners by passing *reweighted* examples. Accordingly, in order to design a good boosting algorithm, we need to design the booster's strategy for random relabelling while also quantifying the weak learner's ability to maximize cumulative gain, even under relabelled data. The first task will be resolved by allowing the booster to use an Online Convex Optimization (OCO) oracle. In this way, we reduce boosting to OCO, an idea borrowed from Brukhim et al. [4]. For the second task, we give different possible weak learning conditions for the

81 *same* algorithm, all of which capture the ability of a weak learner to maximize cumulative gain with
82 respect to the best fixed competitor in hindsight.

Online Convex Optimization. Our booster will use an OCO oracle to update its weak learners. The
OCO setting is a sequential game between an online player and adversary over N rounds (see [20] for
an in-depth introduction). In each round, the player plays a point x_i in a compact convex set $\mathcal{K} \subset \mathbb{R}^d$,
the adversary reveals a loss function f_i chosen from a family of bounded convex functions over $\mathcal{K}$,
and the player suffers the loss $f_i(x_i)$. The goal of player is to also minimize *regret*, defined as

$$R(N) = \sum_{i=1}^{N} f_i(x_i) - \min_{x \in \mathcal{K}} \sum_{i=1}^{N} f_i(x).$$

83 We will denote an algorithm in this setting as $\mathrm{OCO}(\mathcal{K}, N)$. If $\mathcal{A}$ is a $\mathrm{OCO}(\mathcal{K}, N)$, then we will
84 denote its regret by $R_{\mathcal{A}}(N)$.

85 3 Online Agnostic Boosting

86 In this section, we present an online agnostic boosting algorithm and analyze its performance. We
87 begin in Subsection 3.1 by formally describing our algorithm which reduces boosting to OCO. Then,
88 in Subsection 3.2, we state a weak learning condition and subsequently prove a regret bound for our
89 proposed algorithm. We remind the reader that we let a prediction/label $\ell \in [k]$ also represent its
90 corresponding basis vector e_ℓ.

91 3.1 Algorithm

92 Pseudocode for our online agnostic boosting algorithm is provided in Algorithm 1. The booster
93 maintains oracle access to N copies of a weak learner $\mathcal{W}_1, ..., \mathcal{W}_N$ as well as a $\mathrm{OCO}(\Delta_k, N)$
94 algorithm $\mathcal{A}$. Each weak learner is characterized by some advantage parameter γ (directly proportional
95 to its strength), and we will be precise about what exactly γ quantifies in Subsection 3.2. In round
96 t, the booster uses the weak learners to make a prediction $\hat{y}_t$, observes the true label y_t, and finally
97 simulates a game with the OCO algorithm $\mathcal{A}$ to update each weak learner $\mathcal{W}_i$. Specifically, the
98 booster uses the outputs of $\mathcal{A}$ to feed relabelled examples (x_t, y_t^i) to $\mathcal{W}_i$. A critical component of
99 our algorithm is the L_2 projection operator onto the probability simplex Δ_k, which we denote by $\prod$.
100 This is used by the booster in line 4 to make randomized predictions $\hat{y}_t$.

Algorithm 1: Online Agnostic Multiclass Boosting via OCO

Input: Weak Learners $\mathcal{W}_1 ... \mathcal{W}_N$, $\mathrm{OCO}(\Delta_k, N)$ algorithm $\mathcal{A}$

1 **for** $t = 1, ..., T$ **do**
2   Receive example x_t
3   Accumulate weak predictions $h_t = \sum_{i=1}^{N} \mathcal{W}_i(x_t)$
4   Set $\mathcal{D}_t = \prod(\frac{h_t}{\gamma N})$
5   Predict $\hat{y}_t \sim \mathcal{D}_t$
6   Receive true label y_t
7   **for** $i = 1, ..., N$ **do**
8   If $i > 1$, obtain $p_t^i = \mathcal{A}(l_t^1, ..., l_t^{i-1})$. Else, initialize $p_t^1 = \frac{\mathbb{1}_k}{k}$.
9   Reveal loss function: $l_t^i(p) = p \cdot \left(\frac{2\mathcal{W}_i(x_t) - \mathbb{1}_k}{\gamma} - (2y_t - \mathbb{1}_k) \right)$
10   Sample random label $y_t^i \sim p_t^i$, and pass (x_t, y_t^i) to $\mathcal{W}_i$
11   Reset $\mathcal{A}$

---

101 While the framework in Algorithm 1 is inspired from that of Brukhim et al. [4] in the binary setting,
102 several new complications arise in the multiclass setting. Most notably, when $k > 2$, we must figure
103 out how the booster should make predictions, what loss functions the booster should construct and
104 pass to $\mathcal{A}$, and lastly, how the booster should use the output of $\mathcal{A}$ to update each weak learner. The
105 interplay behind these three algorithmic pieces is delicate and they have been carefully designed in
106 Algorithm 1 to enable the analysis in Subsection 3.2. Below, we provide some intuition behind these
107 algorithmic decisions.

Figure 1: γ-scaled votes, p_1 and p_2, and their corresponding L_2 projection onto Δ_2, p_1^* and p_2^*
. $\tilde{p}_1$ and $\tilde{p}_2$ denote the orthogonal projections onto the plane containing Δ_2.

Randomized Prediction. At the start of each round, the booster averages the weak learners votes, scales the average by the parameter γ, projects the scaled vector back into the simplex, and finally samples a random label. When the L_2 projection operator is selected, this approach for randomized prediction achieves a *polarization* effect: as γ gets smaller, the projection concentrates mass on a fewer number of labels, specifically those labels that have achieved the majority of the votes from the weak learners. When γ nears 0, the L_2 projection eventually places all mass on the label with the most votes.

In general, the L_2 projection is not a natural projection operator for Δ_k. However, in the case where we project from $\Delta_{\frac{k}{\gamma}}$ to Δ_k, we find that it is a natural choice from a geometric perspective. Figure 1 provides a visualization of the L_2 projection operator from $\Delta_{\frac{k}{\gamma}}$ to Δ_k for $k = 2$. Because the spaces $\Delta_{\frac{k}{\gamma}}$ and Δ_k are parallel, given any point $p \in \Delta_{\frac{k}{\gamma}}$, one can think of its L_2 projection in the following procedural manner:

1. Compute the *orthogonal* projection, $\tilde{p}$, onto the plane containing Δ_k.

2. If $\tilde{p} \in \Delta_k$, output $p^* = \tilde{p}$.

3. Else, output $p^* = \arg\min_{p \in \Delta_k} ||p - \tilde{p}||_2$.

From this procedural perspective, Figure 1 lends a geometric intuition behind the polarization effect of the L_2 projection. As γ shrinks, the number of points in $\Delta_{\frac{k}{\gamma}}$ that lie orthogonally above Δ_k shrinks. Thus, the orthogonal projection of the vast majority of points in $\Delta_{\frac{k}{\gamma}}$ will lie outside Δ_k, leading to a greater number of sparse projections that lie on the boundary of Δ_k. From the booster's viewpoint, this property of the L_2 projection is desirable. If the weak learners are very weak (small γ) but somehow concentrate votes on a few labels, then it may be likely that the true response is amongst these few labels. In this sense, γ controls how many weak learners need to agree on a particular label to convince the booster to deterministically predict that label.

Updating Weak Learners. Once the true label y_t is revealed in each round, the booster must update each weak learner. As mentioned in Section 2, one strategy for updating weak learners in the agnostic setting is via *random relabelling*. Indeed, in line 10, the booster passes back to the weak learner the example x_t with a random label $y_t^i \sim p_t^i$. Together, the specified loss function in line 9 and random relabelling strategy in line 10 achieve the following effect: if more weak learners make mistakes, the distribution over labels output by the OCO algorithm $\mathcal{A}$ in line 8 concentrates on the true label y_t, increasing the likelihood that y_t is passed to subsequent weak learners. This is desirable as the outputs of the OCO guide each weak learner to correct for the mistakes of preceding learners.

3.2 Regret Analysis

Before we give the regret bound of Algorithm 1, we need to specify the capacity to which the weak learners can make predictions, even under potentially relabelled data. Unfortunately, in the agnostic setting, there is no canonical weak learning condition for multiclass problems. In this paper, we give *several* possible definitions of an Agnostic Weak Online Learner (AWOL) for the multiclass setting, all of which enable a regret analysis of Algorithm 1. We emphasize that we can derive different regret bounds for Algorithm 1 based on what condition we assume our weak learners to satisfy. For this section, we present a weak learning condition based loosely on a *one-vs-one* approach to multiclass classification. In Appendix C, we provide alternative weak learning conditions.

Define the gain function for input $z \in [k]$,

$$\sigma_{y,\ell}(z) = \mathbb{1}\{z = y\} - \mathbb{1}\{z = \ell\}.$$

For some $y \in [k]$ and $\ell \in [k] \setminus \{y\}$, $\sigma_{y,\ell}(\cdot)$ can be thought of as the binary classification task between labels y and ℓ. Definition 1 requires that for any such sequence of binary classification tasks, an online agnostic weak learner must be able to eventually distinguish between every pair of labels to some non-trivial, but far from optimal, degree.

Definition 1 (Agnostic Weak Online Learning). *Let $\mathcal{H} \subseteq [k]^{\mathcal{X}}$ be a class of experts and let $0 < \gamma \leq 1$ denote the "advantage". An online learning algorithm $\mathcal{W}$ is a (γ, T)-agnostic weak online learner (**AWOL**) for $\mathcal{H}$ if for any adaptively chosen sequence of tuples $(x_t, y_t, \ell_t) \in \mathcal{X} \times [k] \times [k]$ where $\ell_t \neq y_t$, the algorithm outputs $\mathcal{W}(x_t) \in [k]$ at every iteration $t \in [T]$ such that,*

$$\gamma \max_{h \in \mathcal{H}} \mathbb{E}\left[\sum_{t=1}^{T} \sigma_{y_t, \ell_t}(h(x_t))\right] - \mathbb{E}\left[\sum_{t=1}^{T} \sigma_{y_t, \ell_t}(\mathcal{W}(x_t))\right] \leq R_{\mathcal{W}}(T, k),$$

where the expectation is taken w.r.t. the randomness of the weak learner $\mathcal{W}$ and that of the possibly adaptive adversary, and $R_{\mathcal{W}} : \mathbb{N} \times \mathbb{N} \to \mathbb{R}_+$ is the additive regret: a non-decreasing, sub-linear function of T.

The precise dependence of $R_{\mathcal{W}}(T, k)$ on k is explored in more detail in Appendix D where we explicitly construct learners satisfying AWOL. We make few remarks about Definition 1. First, the strength of the weak learner varies directly with γ. Second, for $k = 2$, Definition 1 reduces to the weak learning condition by Brukhim et al. [4] in the binary setting. Finally, we emphasize that Definition 1 holds under an *adaptive* adversary, one that can choose (y_t, ℓ_t) based on $\{\mathcal{W}(x_i)\}_{i=1}^{t-1}$ and its own internal random bits. Importantly, we also allow the adversary to pick ℓ_t even after it has observed $\mathcal{W}(x_t)$ as this only controls how much loss the weak learner suffers when it is incorrect.

Under the assumption that weak learners satisfy Definition 1, Theorem 1 bounds the regret of Algorithm 1 under an *oblivious* adversary. Using a standard reduction, our results can then be generalized to an *adaptive* adversary (see Chapter 4 in [6]).

Theorem 1 (Regret Bound). *The expected regret bound Algorithm 1 satisfies:*

$$\frac{1}{T} \mathbb{E}\left[\max_{h \in \mathcal{H}} \sum_{t=1}^{T} (2h(x_t) \cdot y_t - 1) - \sum_{t=1}^{T} (2\hat{y}_t \cdot y_t - 1)\right] \leq \frac{R_{\mathcal{W}}(T, k)}{\gamma T} + \frac{R_{\mathcal{A}}(N)}{N},$$

where the expectation is over the randomness of the algorithm and weak learners, and $R_{\mathcal{W}}(T, k)$, $R_{\mathcal{A}}(N)$ are the regret terms of the AWOL, OCO algorithms respectively.

If one picks $\mathcal{A}$ to be *Online Gradient Descent* (OGD), then $R_{\mathcal{A}}(N) = O(GD\sqrt{N})$, where D is the diameter of Δ_k and G is the upper-bound on $||\nabla_p l_t^i(p)||$. In our setting, $D = \sqrt{2}$ and $G = O(\frac{1}{\gamma})$ and hence the average regret further simplifies to:

$$\frac{1}{T} \mathbb{E}\left[\max_{h \in \mathcal{H}} \sum_{t=1}^{T} (2h(x_t) \cdot y_t - 1) - \sum_{t=1}^{T} (2\hat{y}_t \cdot y_t - 1)\right] \leq \frac{R_{\mathcal{W}}(T, k)}{\gamma T} + O\left(\frac{1}{\gamma\sqrt{N}}\right).$$

To exemplify the role of N, consider a scenario where $R_{\mathcal{W}}(T, k) = O(k\sqrt{T})$ (see Appendix D for examples). By setting $N = \frac{T}{\gamma^2}$ the overall regret of the booster becomes $O(k\sqrt{T})$. In the next subsection, we give the proof of Theorem 1.

3.2.1 Proof of Theorem 1

We follow a similar procedure to Brukhim et al. [4] by lower and upper bounding the expected sum of losses passed to $\mathcal{A}$ in terms of the regret of the weak learner and the regret of $\mathcal{A}$ respectively. These bounds rely on several lemmas that have been abstracted out and provided in Appendix E. As mentioned previously, we also assume an oblivious adversary.

Let $(x_1, y_1), ..., (x_T, y_T)$ be any sequence of example-label pairs. We start by giving a lower bound on the expected sum of losses passed to $\mathcal{A}$ using Definition 1. Define $h^* = \arg\max_{h \in \mathcal{H}} \sum_{t=1}^{T} (2h(x_t) \cdot y_t - 1)$ as the optimal competitor in hindsight and let

$$
\ell_t^i = \begin{cases} \mathcal{W}_i(x_t), & \text{if } \mathcal{W}_i(x_t) \neq y_t^i \\ \ell \in [k] \setminus \{y_t^i\}, & \text{otherwise} \end{cases} .
$$

Note that the precise choice of ℓ_t^i in the second case is not important. The proof below only uses the fact that $\ell_t^i \in [k] \setminus \{y_t^i\}$ in that case. We then have,

$$
\begin{aligned}
\mathbb{E}\left[\sum_{i=1}^{N}\sum_{t=1}^{T} \ell_t^i(p_t^i)\right] &= \mathbb{E}\left[\sum_{i=1}^{N}\sum_{t=1}^{T} p_t^i \cdot \left(\frac{2\mathcal{W}_i(x_t) - \mathbb{1}_k}{\gamma} - (2y_t - \mathbb{1}_k)\right)\right] \\
&= \frac{1}{\gamma}\sum_{i=1}^{N}\sum_{t=1}^{T} \mathbb{E}\left[p_t^i \cdot 2(\mathcal{W}_i(x_t) - \mathbb{1}_k)\right] - \sum_{i=1}^{N}\sum_{t=1}^{T} \mathbb{E}\left[p_t^i \cdot (2y_t - \mathbb{1}_k)\right] \\
&= \frac{1}{\gamma}\sum_{i=1}^{N}\sum_{t=1}^{T} \mathbb{E}\left[2\mathcal{W}_i(x_t) \cdot y_t^i - 1\right] - \sum_{i=1}^{N}\sum_{t=1}^{T} \mathbb{E}\left[p_t^i \cdot (2y_t - \mathbb{1}_k)\right] \qquad \text{(Lemma 12)} \\
&= \frac{1}{\gamma}\sum_{i=1}^{N}\sum_{t=1}^{T} \mathbb{E}\left[\sigma_{y_t^i, \ell_t^i}(\mathcal{W}_i(x_t))\right] - \sum_{i=1}^{N}\sum_{t=1}^{T} \mathbb{E}\left[p_t^i \cdot (2y_t - \mathbb{1}_k)\right].
\end{aligned}
$$

Using the weak learning condition in Definition 1,

$$
\begin{aligned}
\frac{1}{\gamma}\sum_{i=1}^{N}\sum_{t=1}^{T} \mathbb{E}\left[\sigma_{y_t^i, \ell_t^i}(\mathcal{W}_i(x_t))\right] &\geq \sum_{i=1}^{N} \max_{h \in \mathcal{H}} \sum_{t=1}^{T} \mathbb{E}\left[\sigma_{y_t^i, \ell_t^i}(h(x_t))\right] - \frac{NR_{\mathcal{W}}(T, k)}{\gamma} \qquad \text{(Definition 1)} \\
&\geq \sum_{i=1}^{N}\sum_{t=1}^{T} \mathbb{E}\left[\sigma_{y_t^i, \ell_t^i}(h^*(x_t))\right] - \frac{NR_{\mathcal{W}}(T, k)}{\gamma} \\
&\geq \sum_{i=1}^{N}\sum_{t=1}^{T} \mathbb{E}\left[2h^*(x_t) \cdot y_t^i - 1\right] - \frac{NR_{\mathcal{W}}(T, k)}{\gamma} \\
&= \sum_{i=1}^{N}\sum_{t=1}^{T} \mathbb{E}\left[p_t^i \cdot (2h^*(x_t) - \mathbb{1}_k)\right] - \frac{NR_{\mathcal{W}}(T, k)}{\gamma}. \qquad \text{(Lemma 12)}
\end{aligned}
$$

Putting things together, we find,

$$
\begin{aligned}
\mathbb{E}\left[\sum_{i=1}^{N}\sum_{t=1}^{T} \ell_t^i(p_t^i)\right] &\geq \sum_{i=1}^{N}\sum_{t=1}^{T} \mathbb{E}\left[p_t^i \cdot (2h^*(x_t) - 2y_t)\right] - \frac{NR_{\mathcal{W}}(T, k)}{\gamma} \\
&\geq \sum_{i=1}^{N}\sum_{t=1}^{T} \mathbb{E}\left[2(h^*(x_t) \cdot y_t - 1)\right] - \frac{NR_{\mathcal{W}}(T, k)}{\gamma} \qquad \text{(Lemma 13)} \\
&= N\sum_{t=1}^{T} 2(h^*(x_t) \cdot y_t - 1) - \frac{NR_{\mathcal{W}}(T, k)}{\gamma}.
\end{aligned}
$$

190 Now, we compute an upper bound. For any $t \in [T]$ and arbitrary $p^* \in \Delta_k$:

$$
\mathbb{E}\left[\frac{1}{N}\sum_{i=1}^{N} l_t^i(p_t^i)\right] \leq \frac{1}{N}\mathbb{E}\left[\min_{p\in\mathcal{K}}\sum_{i=1}^{N} l_t^i(p)\right] + \frac{R_{\mathcal{A}}(N)}{N} \qquad \text{(OCO Regret)}
$$

$$
\leq \frac{1}{N}\mathbb{E}\left[\sum_{i=1}^{N} l_t^i(p^*)\right] + \frac{R_{\mathcal{A}}(N)}{N}
$$

$$
= \frac{1}{N}\sum_{i=1}^{N}\mathbb{E}\left[p^* \cdot \left(\frac{2\mathcal{W}_i(x_t) - \mathbb{1}_k}{\gamma} - (2y_t - \mathbb{1}_k)\right)\right] + \frac{R_{\mathcal{A}}(N)}{N}
$$

$$
= \frac{1}{\gamma N}\sum_{i=1}^{N}\left(p^* \cdot \mathbb{E}\left[2\mathcal{W}_i(x_t) - \mathbb{1}_k\right]\right) - p^* \cdot (2y_t - \mathbb{1}_k) + \frac{R_{\mathcal{A}}(N)}{N}
$$

$$
= p^* \cdot \left(\frac{1}{\gamma N}\sum_{i=1}^{N}\mathbb{E}\left[2\mathcal{W}_i(x_t) - \mathbb{1}_k\right]\right) - p^* \cdot (2y_t - \mathbb{1}_k) + \frac{R_{\mathcal{A}}(N)}{N}
$$

$$
= \mathbb{E}\left[p^* \cdot \left(\frac{\frac{2}{N}\sum_{i=1}^{N}\mathcal{W}_i(x_t) - \mathbb{1}_k}{\gamma} - (2y_t - \mathbb{1}_k)\right)\right] + \frac{R_{\mathcal{A}}(N)}{N}
$$

$$
\leq \mathbb{E}\left[2\left(\prod\left(\frac{1}{\gamma N}\sum_{i=1}^{N}\mathcal{W}_i(x_t)\right)\cdot y_t - 1\right)\right] + \frac{R_{\mathcal{A}}(N)}{N} \qquad \text{(Lemma 14)}
$$

$$
= 2\left(\mathbb{E}\left[\hat{y}_t\right]\cdot y_t - 1\right) + \frac{R_{\mathcal{A}}(N)}{N}. \qquad \text{(Lemma 16)}
$$

191 Summing over T,

$$
\mathbb{E}\left[\frac{1}{N}\sum_{t=1}^{T}\sum_{i=1}^{N}\ell_t^i(p_t^i)\right] \leq \sum_{t=1}^{T} 2\left(\mathbb{E}\left[\hat{y}_t\right]\cdot y_t - 1\right) + \frac{TR_{\mathcal{A}}(N)}{N}.
$$

192 Combining lower and upper bounds for $\mathbb{E}\left[\frac{1}{NT}\sum_{i=1}^{N}\ell_t^i(p_t^i)\right]$, we get

$$
\frac{1}{T}\mathbb{E}\left[\max_{h\in\mathcal{H}}\sum_{t=1}^{T}(2h(x_t)\cdot y_t - 1) - \sum_{t=1}^{T}(2\hat{y}_t\cdot y_t - 1)\right] \leq \frac{R_{\mathcal{W}}(T,k)}{\gamma T} + \frac{R_{\mathcal{A}}(N)}{N},
$$

193 which completes the proof. Since $2h(x_t)\cdot y_t - 1 = 1 - 2\mathbb{1}\{h(x_t)\neq y_t\}$, we can also write the
194 average regret in terms of the number of *mistakes*,

$$
\frac{1}{T}\mathbb{E}\left[\sum_{t=1}^{T}\mathbb{1}\{\hat{y}_t\neq y_t\} - \min_{h\in\mathcal{H}}\sum_{t=1}^{T}\mathbb{1}\{h(x_t)\neq y_t\}\right] \leq \frac{R_{\mathcal{W}}(T,k)}{2\gamma T} + \frac{R_{\mathcal{A}}(N)}{2N}.
$$

195 4 Beyond Online Agnostic Boosting

196 In this section, we present results of extending our reduction to the three other boosting settings,
197 namely statistical agnostic, online realizable, and statistical realizable learning. The purpose of
198 this section is to showcase the generality of the OCO-based boosting framework and not to achieve
199 state-of-the-art bounds for these settings. Algorithms, discussion of weak learning conditions, and all
200 associated proofs can be found in Appendix A and B.

201 4.1 Statistical Agnostic Boosting

In the statistical setting, our objective of interest is the *correlation*, which we define below. Let $\mathcal{D}$ be
a distribution over $\mathcal{X} \times \mathcal{Y}$ and let $h : \mathcal{X} \to \mathcal{Y}$ be an hypothesis. Define the *multiclass* correlation of h
with respect to $\mathcal{D}$ as

$$
cor_D(h) = \mathbb{E}_{(x,y)\sim D}\left[h(x)\cdot(2y - \mathbb{1}_k)\right].
$$

Like the online agnostic setting, the boosting algorithm for this setting (provided in Appendix A) can be analyzed under several candidate weak learning conditions. To showcase the dependence on k, we provide a weak learning condition based loosely on a *one-vs-all* approach to multiclass classification.

Definition 2 (Empirical Agnostic Weak Learning). *Let $\mathcal{H} \subseteq [k]^{\mathcal{X}}$ be a hypothesis class and let $0 < \gamma \leq 1$ denote the "advantage". Let $\boldsymbol{x} = (x_1, ..., x_m) \in \mathcal{X}$ denote an unlabeled sample. A learning algorithm $\mathcal{W}$ is a $(\gamma, \epsilon_0, m_0)$-agnostic weak learner (**AWL**) for $\mathcal{H}$ with respect to $\boldsymbol{x}$ if for any labels $\boldsymbol{y} = (y_1, ..., y_m) \in [k]$, and every reference label $\ell \in [k]$,*

$$\mathbb{E}_{S'} \left[\sum_{i:y_i=\ell} \mathcal{W}_{S'}(x_i) \cdot (2y_i - \mathbb{1}_k) \right] \geq \gamma \max_{h \in \mathcal{H}} \sum_{i:y_i=\ell} h(x_i) \cdot (2y_i - \mathbb{1}_k) - m\epsilon_0,$$

where S' is an independent sample size m_0 drawn from the distribution which uniformly assigns to each example (x_i, y_i) probability $1/m$.

Under the assumption that our weak learners satisfy Definition 2, Corollary 2 bounds the expected correlation of a statistical agnostic boosting algorithm.

Corollary 2 (Empirical Agnostic Correlation Bound). *There exists a boosting algorithm whose output after T rounds, denoted by $\bar{h}$, satisfies:*

$$\mathbb{E}\left[cor_S(\bar{h})\right] \geq \max_{h \in \mathcal{H}} \mathbb{E}\left[cor_S(h)\right] - \frac{R_{\mathcal{A}}(T)}{T} - \frac{k\epsilon_0}{\gamma},$$

where S is the distribution which uniformly assigns to each example (x_i, y_i) probability $1/m$.

4.2 Online Realizable Boosting

In the online realizable setting, we are guaranteed that the stream of examples $(x_1, y_1), ..., (x_T, y_T)$ is perfectly labelled by some hypothesis $h \in \mathcal{H}$. Definition 3 then gives an appropriate online realizable weak learning condition.

Definition 3 (Realizable Weak Online Learning). *Let $\mathcal{H} \subseteq [k]^{\mathcal{X}}$ be a class of experts and let $0 < \gamma \leq 1$ denote the "advantage". An online learning algorithm $\mathcal{W}$ is a (γ, T)-realizable weak online learner (**RWOL**) for $\mathcal{H}$ if for any sequence $(x_1, y_1), ..., (x_T, y_T) \in \mathcal{X} \times [k]$ that is realizable by $\mathcal{H}$, at every iteration $t \in [T]$, the algorithm outputs $\mathcal{W}(x_t) \in [k]$ such that,*

$$\mathbb{E}\left[\sum_{t=1}^{T} (2\mathcal{W}(x_t) \cdot y_t - 1) \right] \geq \gamma T - R_{\mathcal{W}}(T, k),$$

where the expectation is taken w.r.t. the randomness of the weak learner $\mathcal{W}$ and $R_{\mathcal{W}} : \mathbb{N} \times \mathbb{N} \to \mathbb{R}_+$ is the additive regret: a non-decreasing, sub-linear function of T.

Under Definition 3, Corollary 3 bounds the expected gain of an online boosting algorithm.

Corollary 3 (Mistake Bound). *There exists a boosting algorithm whose outputs $\hat{y}_1, ..., \hat{y}_T$ satisfy:*

$$\frac{1}{T} \sum_{t=1}^{T} (2\,\mathbb{E}\left[\hat{y}_t\right] \cdot y_t - 1) \geq 1 - \frac{R_{\mathcal{A}}(N)}{N} - \frac{\tilde{R}_{\mathcal{W}}(T, k)}{\gamma T},$$

where $\tilde{R}_{\mathcal{W}}(T, k) = 2R_{\mathcal{W}}(T, k) + \tilde{O}(\sqrt{T})$.

4.3 Statistical Realizable Boosting

In the statistical realizable setting, our metric of interest is again the multiclass correlation. However, under the realizability assumption, $\max_{h^* \in \mathcal{H}} cor_D(h^*) = 1$. Accordingly, Definition 4 gives an appropriate realizable weak learning condition.

Definition 4 (Empirical Realizable Weak Learning). *Let $\mathcal{H} \subseteq [k]^{\mathcal{X}}$ be a hypothesis class and let $0 < \gamma \leq 1$ denote the "advantage". Let $S = \{(x_1, y_1), ..., (x_m, y_m)\} \in \mathcal{X} \times [k]$ be a sample. A learning algorithm $\mathcal{W}$ is a (γ, m_0) - realizable weak learner (**RWL**) for $\mathcal{H}$ with respect to S if for any distribution $\boldsymbol{p} = (p_1, ..., p_m)$ over the examples,*

$$\mathbb{E}_{S'} \left[cor_{\boldsymbol{p}}(\mathcal{W}_{S'}) \right] \geq \gamma,$$

where S' is an independent sample of size m_0 drawn from $\boldsymbol{p}$.

Table 1: Average accuracy and runtime of algorithms on 7 UCI datasets.

Dataset	k	Accuracy(%)				Runtime(s)			
		Agn	Opt	Ada	OCOR	Agn	Opt	Ada	OCOR
Balance	3	**84.4**	79.5	75.2	78.8	11.5	57.6	7.1	4.5
Cars	4	**80.6**	69.9	75.8	69.6	42.5	104.1	20.4	12.1
Landsat	6	67.0	**80.8**	56.9	79.6	1255.6	1814.3	81.8	440.4
Segment	7	75.0	78.9	68.6	**79.3**	647.2	2448.3	51.8	154.7
Mice	8	**86.0**	77.6	71.3	79.6	1025.8	1811.6	90.1	258.2
Yeast	10	42.3	39.8	41.7	**47.6**	348.3	1468.8	25.2	56.6
Abalone	28	22.1	**24.9**	19.2	22.0	4036.7	9775.3	108.0	248.2

Under Definition 4, Corollary 4 bounds the expected correlation of a statistical realizable boosting algorithm.

Corollary 4 (Empirical Correlation Bound). *There exists a boosting algorithm whose output after T rounds, denoted $\bar{h}$, satisfies,*

$$\mathbb{E}\left[cor_S(\bar{h})\right] \geq 1 - \frac{R_{\mathcal{A}}(T)}{T},$$

where S is the distribution which uniformly assigns to each example (x_i, y_i) probability $1/m$.

5 Experiments

We performed experiments with Algorithm 1 on seven UCI datasets [10]. Since, to our knowledge, there are no known online *agnostic* boosting algorithms for multiclass problems, we benchmark performance against three online *realizable* multiclass boosting algorithms: the state-of-the-art optimal (OnlineMBBM) and adaptive (Adaboost.OLM) online boosting algorithms by Jung et al. [21], and the OCO-based online boosting algorithm alluded to in Section 4 (see Appendix B). For weak learners, we used the implementation of the VeryFastDecisionTree from the River package [26] and restricted the maximum depth of the tree to 1. We used Projected OGD [32] for the OCO algorithm and set the number of weak learners, N, to 100 for each boosting algorithm.

Our experiments using Algorithm 1 are performed using *fractional relabeling*. That is, instead of passing just a single example-label pair (x_t, y_t^i) where $y_t^i \sim p_t^i$, we pass all k possible *weighted* example-label tuples $\{(x_t, y_\ell, p_t^i[\ell])\}_{\ell=1}^k$ to weak learner i in round t. Experiments with random relabeling showed that random relabeling runs faster but performs worse than fractional relabeling.

Table 1 summarizes the average accuracy and runtime over five independent shuffles of each dataset for each boosting algorithm. "Agn" refers to Algorithm 1, "Opt" to OnlineMBBM, "Ada" to Adaboost.OLM, and "OCOR" to the OCO-based online realizable algorithm. Standard errors for accuracies are provided in Appendix F. All algorithms **except** AdaBoost.OLM are parameterized by an advantage parameter $\gamma \in (0, 1)$. Thus, γ was tuned separately for each respective cell of Table 1. See Appendix F for more experimental details. All code is available at ⟨URL REDACTED⟩.

Despite having the fastest overall runtime, Adaboost.OLM had the poorest performance. Compared to OnlineMBBM, our OCO-based boosting algorithms achieve comparable if not better performance at a fraction of the runtime. Specifically, for three of our datasets, our online agnostic boosting algorithm achieves the highest accuracy. For the remaining datasets, our OCO-based realizable boosting algorithm achieves comparable accuracy to OnlineMBBM with shorter runtimes.

6 Discussion

We give the first weak learning conditions and algorithm for online agnostic multiclass boosting. Our algorithm relies on a clean and simple reduction from boosting to online convex optimization. This fruitful connection allows us to go beyond the online agnostic setting and design multiclass boosting algorithms for all four regimes of statistical/online and agnostic/realizable learning. As future work, we leave it open to identify the correct weak learning condition, construct adaptive versions of our boosting algorithms, improve regret upper bounds/prove lower bounds, design agnostic boosting algorithms under bandit feedback, and study the impact of the choice of weak learner and OCO algorithm on the empirical performance of our OCO-based boosting algorithms.

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