# Signal Recovery with Non-Expansive Generative Network Priors

**Jorio  Cocola** [*]
Harvard University
jcocola@seas.harvard.edu

## Abstract

We study compressive sensing with a deep generative network prior. Initial theoretical guarantees for efficient recovery from compressed linear measurements have been developed for signals in the range of a ReLU network with Gaussian weights and logarithmic expansivity: that is when each layer is larger than the previous one by a logarithmic factor. It was later shown that constant expansivity is sufficient for recovery. It has remained open whether the expansivity can be relaxed, allowing for networks with contractive layers (as often the case of real generators). In this work we answer this question, proving that a signal in the range of a Gaussian generative network can be recovered from few linear measurements provided that the width of the layers is proportional to the input layer size (up to log factors). This condition allows the generative network to have contractive layers. Our result is based on showing that Gaussian matrices satisfy a matrix concentration inequality which we term *Range Restricted Weight Distribution Condition* (R2WDC) and which weakens the *Weight Distribution Condition* (WDC) upon which previous theoretical guarantees were based. The WDC has also been used to analyze other signal recovery problems with generative network priors. By replacing the WDC with the R2WDC, we are able to extend previous results for signal recovery with expansive generative network priors to non-expansive ones. We discuss these extensions for phase retrieval, denoising, and spiked matrix recovery.

## 1   Introduction

The compressed sensing problem consists in estimating a signal $y_\star \in \mathbb{R}^n$ from (possibly) noisy linear measurements

$$b = Ay_\star + \eta$$

where $A \in \mathbb{R}^{m \times n}$ is the measurements matrix, $m < n$ and $\eta \in \mathbb{R}^m$ is the noise.

To overcome the ill-posedness of the problem, structural priors on the unknown signal $y_\star$ need to be enforced. One now classical approach assumes that the target signal $y_\star$ is sparse with respect to a given basis. In the last 20 years, efficient reconstruction algorithms have been developed that provably estimate $s$-sparse signals in $\mathbb{R}^n$ from $m = \mathcal{O}(s \log n)$ random measurements [5, 12].

Another approach recently put forward, leverages trained generative networks. These networks are trained, in an unsupervised manner, to generate samples from a target distribution of signals. Assuming $y_\star$ belongs to the same distribution used to train a generative network $G : \mathbb{R}^k \to \mathbb{R}^n$ with $k \ll n$, an estimate of $y_\star$ can be found by searching the input $\hat{x}$ ("latent code") of $G$ that minimizes

---

[*]Work done while at Northeastern University

36th Conference on Neural Information Processing Systems (NeurIPS 2022).

the reconstruction error

$$\tilde{x} = \arg\min_{x \in \mathbb{R}^x} f_{\mathrm{cs}}(x) := \frac{1}{2}\|b - AG(x)\|_2^2, \tag{1}$$

$$y_\star \approx G(\tilde{x}).$$

As empirically demonstrated in [3], the minimization problem (1) can be solved efficiently by gradient descent methods. Moreover, solving (1) can effectively regularize the solution of the compressed sensing problem, significantly outperforming sparsity-based algorithms in the low measurements regime[3]. Generative network based inversion algorithms have been subsequently developed for a variety of signal recovery problems, demonstrating their potential to outperform inversion algorithms based on non-learned (hand-crafted) priors [16, 31, 30, 20, 33, 28]. For a recent overview see [32].

The optimization problem (1) is in general non-convex and gradient-based methods could get stuck in local minima. To better understand the empirical success of (1), in [18] the authors established theoretical guarantees for the noiseless compressed sensing problem ($\eta = 0$) where $G : \mathbb{R}^k \to \mathbb{R}^n$ is a $d$-layer ReLU network of the form:

$$G(x) = \mathrm{ReLU}(W_d \cdots \mathrm{ReLU}(W_2 \mathrm{ReLU}(W_1 x))) \tag{2}$$

with $W_i \in \mathbb{R}^{n_i \times n_{i-1}}$, $n_0 = k$, $n_d = n$, and $\mathrm{ReLU}(z) = \max(z, 0)$ is applied entrywise. The authors of [18] used a probabilistic model for the generative network $G$ and measurement matrix $A$. They assumed that each layer $W_i$ has independent Gaussian entries and is *strictly expansive*. Specifically it holds that

$$n_i \geq n_{i-1} \cdot \log n_{i-1} \cdot \mathrm{poly}(d) \qquad \text{for all } i = 1, \ldots, d. \tag{3}$$

Moreover, they considered $A$ to be a Gaussian matrix and $m \geq k \cdot \log n \cdot \mathrm{poly}(d)$. Under this probabilistic model it was shown in [18] that, despite its non-convexity, $f_{\mathrm{cs}}$ has a favorable optimization geometry and no spurious critical points exist apart from $x_\star$ and a negative multiple of it $-\rho_d x_\star$, where $\rho_d$ is a function of the depth $d$ of the network.

The landscape analysis was later extended to recovery guarantees using a gradient based method in [21], under the same probabilistic assumptions of [18]. In particular, [21] has shown that there is an efficient gradient descent method (see Algorithm 1 in Section 3) that given as input $A, G$ and $b$ outputs a latent vector $\tilde{x}$ such that $\|y_\star - G(\tilde{x})\|_2 = O(\|\eta\|_2)$. This result demonstrated that efficient recovery is possible with a number of measurements which is information-theoretic optimal up to $log$-factors in $n$ and polynomials in $d$ ($m = \tilde{\Omega}(k)$).

Generative networks used in practice though, have often contractive layers. For example, the output of the layers near the end of the StyleGAN generators have larger dimension than the generated images [25, 24]. Thus, one major drawback of the theory developed in [18] is constituted by the expansivity condition on the weight matrices (3). Relaxing the condition (3) and accommodating for generative networks with contractive layers was formulated as an open problem in the survey paper[2] [32].

An initial positive result on this problem came from [10]. Using a refined analysis of the concentration of Lipschitz functions, the authors proved that the results of [18, 21] hold true also for weight matrices satisfying $n_i \geq n_{i-1} \cdot \mathrm{poly}(d)$. While not allowing for contractive layers, this condition removed the logarithmic expansivity requirement of (3).

More recently, [22, 23] have studied the denoising and compressive sensing problem with random generative network prior as in [18, 21, 20], and have shown that the expansivity condition can indeed be relaxed. In [23] they have provided an efficient iterative method that given as input $A, b$ and $G$, assuming that up to $log$-factors each layer width satisfies

$$n_i \gtrsim 5^i k, \tag{4}$$

and the number of measurement satisfies

$$m \gtrsim 2^d k, \tag{5}$$

outputs a latent vector $\tilde{x}$ such that for $y_\star = G(x_\star)$ it holds that $\|y_\star - G(\tilde{x})\|_2 = O(2^d \sqrt{\frac{k}{m}} \|\eta\|_2)$ with high probability. Notice that the condition (4) while requiring the width to grow with the depth, can allow for contractive layers $n_i < n_{i-1}$.

---

[2]This open problem was also proposed in the recent talk [11].

## 1.1 Our contributions

It is natural to wonder whether the price to pay to remove the expansivity assumption is indeed the exponential factors in the depth $d$ of the network and the use of less-standard non-gradient based iterative methods, as happens in [22, 23]. In this paper, we answer these questions. Our main result is summarized below and provides guarantees for solving compressed sensing with random generative network priors via a gradient descent method (Algorithm 1 in Section 3).

**Theorem 1.1** (Informal version of Theorem 5.4). *Assume that $A$ has i.i.d. $\mathcal{N}(0, 1/m)$ entries and each $W_i$ has i.i.d. $\mathcal{N}(0, 1/n_i)$ entries. Suppose that $y_\star = G(x_\star)$. Furthermore assume that, up to log-factors,*

*1. $n_i \geq k \cdot poly(d)$;*

*2. $m \geq k \cdot poly(d)$.*

*Suppose that the noise error and the step size $\alpha > 0$ are small enough. Then with high probability, Algorithm 1 with input loss function $f_{cs}$, step size $\alpha$ and number of iterations $T = poly(d)$, outputs an estimate $G(x_T)$ satisfying $\|G(x_T) - y_\star\|_2 = O(\sqrt{\frac{k}{m}} \|\eta\|_2)$.*

Compared to [21] and [10], our result do not require strictly expanding generative networks and allows for contractive layers. Furthermore, we show that the same algorithm proposed in [21] has a denoising effect, leading to a reconstruction of the target signal $y_\star$ of the order $O(\sqrt{\frac{k}{m}} \|\eta\|_2)$ rather than only $O(\|\eta\|_2)$. We show that this holds true even in case of deterministic noise, while [19] discuss only the case of Gaussian noise. Furthermore, the decrease in the reconstruction error with the number of measurements has also been observed for trained generative networks (see for example [3]), and here we give a partial theoretical explanation for this phenomenon.

Compared to the results of [23] we show that it is sufficient for the width of the layers as well as the number of measurements to grow polynomially with the depth rather than exponentially. Similarly, compared to [23], we remove the exponential factor in the depth from the reconstruction error.

The analysis of [18] was based on a deterministic condition on the weight matrices termed *Weight Distribution Condition* (WDC). This condition, together with a deterministic condition on $A$ (see Sec 4 for details), was shown to be sufficient for the absence of spurious local minima in (1) and to be satisfied by expansive Gaussian random generative networks as (2). The WDC was also used in the subsequent [21] to prove convergence of Algorithm 1. Our main technical contribution is to show that the WDC can be replaced by a weaker form of deterministic condition, termed *Range Restricted Weight Distribution Condition* (R2WDC), and still, obtain the absence of spurious local minima and recovery guarantees via Algorithm 1. We will then show that random Gaussian networks satisfying the Assumption 1. of Theorem 1.1 satisfy the R2WDC.

The framework introduced in [18] was used in a number of recent works to analyze other signal recovery problems with generative network priors, from one-bit recovery to blind demodulation [34, 27, 16, 15, 35, 8]. These works considered expansive generative network priors, using the WDC and the results of [18] in their analysis. Replacing the WDC with our R2WDC we can extend the previous results in the literature to more realistic (non-expansive) generative networks. This paper details these extensions for three representative signal recovery problems.

**Theorem 1.2.** *Suppose $G$ is random generative network as in* (2)*, satisfying Assumption 1. of Theorem 1.1. Then Algorithm 1 with appropriate loss functions, step sizes, and number of steps, succeed with high probability for Phase Retrieval, Denoising, and Spiked Matrix Recovery.*

Our result on the denoising problem, implies a similar result on the inversion of a generative network. The problem of inverting a generative neural network has important applications [39, 1, 33], and has been recently analyzed theoretically [26, 22, 2]. Our result shows that a random generative network can be efficiently inverted by gradient descent, even when containing contractive layers. This motivates the empirical use of gradient-based methods for inverting generative networks.

## 1.2 Organization of the paper

This paper is organized as follows. In Section 2 we introduce some notation used in the rest of the paper. In Section 3 we formalize the compressed sensing problem with a generative network prior and describe an algorithm for the recovery. In Section 4 we describe our novel deterministic condition on the weights of the network (R2WDC) and provide theoretical guarantees for solving compressed sensing with a generative network prior satisfying this condition via the algorithm described in Section 3. Then in Section 5 we demonstrate that random non-expansive generative networks satisfy the R2WDC with high probability. The appendix contains the full proof of the results described in the main text. Appendix F contains the extension of the theoretical guarantees for compressed sensing with a generative network prior to other signal recovery problems.

## 2 Preliminaries

We use $I_n$ to denote the $n \times n$ identity matrix. For $j \geq 0$, we define the $j$-th sub-network $G_j : \mathbb{R}^k \to \mathbb{R}^{n_j}$ as $G_j(x) = \mathsf{ReLU}(W_j \cdots \mathsf{ReLU}(W_2 \mathsf{ReLU}(W_1 x)))$, with the convention that $G_0(x) = I_k x = x$. For a matrix $W \in \mathbb{R}^{n \times k}$, let $\mathrm{diag}(Wx > 0)$ be the diagonal matrix with $i$-th diagonal element equal to 1 if $(Wx)_i > 0$ and 0 otherwise, and $W_{+,x} = \mathrm{diag}(Wx > 0)W$. We then define $W_{1,+,x} = (W_1)_{+,x} = \mathrm{diag}(W_1 x > 0)W_1$ and

$$W_{j,+,x} = \mathrm{diag}(W_j W_{j-1,+,x} \cdots W_{2,+,x} W_{1,+,x})W_j.$$

Finally, we let $\Lambda_{0,x} = I_k$ and for $j \geq 1$ $\Lambda_{j,x} = \prod_{\ell=1}^{j} W_{\ell,+,x}$ with $\Lambda_x = \Lambda_{d,x} = \prod_{\ell=1}^{d} W_{\ell,+,x}$. Notice in particular that $G_j(x) = \Lambda_{j,x} x$ and $G(x) = \Lambda_x x$.

For $r, s$ nonzero vectors in $\mathbb{R}^\ell$, we define the matrix

$$Q_{r,s} = \frac{\pi - \theta_{r,s}}{2\pi} I_\ell + \frac{\sin \theta_{r,s}}{2\pi} M_{\hat{r} \leftrightarrow \hat{s}} \tag{6}$$

where $\theta_{r,s} = \angle(r, s)$, $\hat{r} = r/\|r\|_2$, $\hat{s} = s/\|s\|_2$, $I_\ell$ is the $\ell \times \ell$ identity matrix and $M_{\hat{r} \leftrightarrow \hat{s}}$ is the matrix that sends $\hat{r} \mapsto \hat{s}$, $\hat{s} \mapsto \hat{r}$, and with kernel $\mathrm{span}(\{r, s\})^\perp$. If $r$ or $s$ are zero, then we let $Q_{r,s} = 0$. The operator $Q_{r,s}$ is used to define the WDC in the next sections, and allows to control how a random ReLU layer distorts its inputs. Specifically, for very $r, s \in \mathbb{R}^\ell$ we have $\mathbb{E}\left[\mathsf{ReLU}(Wr)^T \mathsf{ReLU}(Ws)\right] = r^T Q_{r,s} s$ when $W \in \mathbb{R}^{n \times \ell}$ has i.i.d. $\mathcal{N}(0, 1/n)$.

## 3 Problem statement and recovery algorithm

Consider a generative network $G : \mathbb{R}^k \to \mathbb{R}^n$ as in (2). The compressive sensing problem with a generative network prior can be formulated as follows.

---

**COMPRESSED SENSING WITH A DEEP GENERATIVE PRIOR**

**Let**: $G : \mathbb{R}^k \to \mathbb{R}^n$ generative network, $A \in \mathbb{R}^{m \times n}$ measurement matrix.
**Let**: $y_\star = G(x_\star)$ for some unknown $x_\star \in \mathbb{R}^k$.

**Given**: $G$ and $A$.
**Given**: Measurements $b = A y_\star + \eta \in \mathbb{R}^m$ with $m \ll n$ and $\eta \in \mathbb{R}^m$ noise.

**Estimate**: $y_\star$.

---

To solve the compressed sensing problem with deep generative prior $G$, in [21], the authors propose the gradient descent method described in Algorithm 1 with objective function $f = f_{\mathrm{cs}}$. This algorithm attempts to minimize the objective function $f_{\mathrm{cs}}$ in (1). Because of the ReLU activation function, the loss function $f_{\mathrm{cs}}$ is nonsmooth. Algorithm 1 therefore resorts to the notion of *Clarke subdifferential*. Indeed, being continuous and piecewise smooth, $f_{\mathrm{cs}}$ is differentiable almost everywhere (by Rademacher's theorem) and admits a Clarke subdifferential given by[3]:

$$\partial f_{\mathrm{cs}}(x) = \mathrm{conv}\left\{ \lim_{p \to \infty} \nabla f_{\mathrm{cs}}(x_p) : \ x_p \to x, \ x_p \in \mathrm{dom}(\nabla f_{\mathrm{cs}}) \right\}, \tag{7}$$

---

[3]For details see for example [7].

where with conv($\cdot$) we denote the convex hull and with dom($\nabla f$) the subset of $\mathbb{R}^k$ where $f$ is differentiable. The vectors $v_x \in \partial f_{\text{cs}}(x)$ are called the *subgradients* of $f_{\text{cs}}$ at $x$, and at a point $x$ where $f_{\text{cs}}$ is differentiable it holds that $\partial f_{\text{cs}}(x) = \{\nabla f_{\text{cs}}(x)\}$.

---

**Algorithm 1:** SUBGRADIENT DESCENT [21]

---

**Input:** Objective function $f$, initial point $x_0 \in \mathbb{R}^k \setminus \{0\}$ and step size $\alpha$
**Output:** An estimate of the target signal $y_\star = G(x_\star)$ and latent vector $x_\star$

1 **for** $t = 0, 1, \ldots$ **do**
2     **if** $f(-x_t) < f(x_t)$ **then** $\tilde{x}_t \leftarrow -x_t$
3     **else** $\tilde{x}_t \leftarrow x_t$
4     Compute $v_{\tilde{x}_t} \in \partial f(\tilde{x}_t)$
5     $x_{t+1} \leftarrow \tilde{x}_t - \alpha v_{\tilde{x}_t}$
6 **end**
7 **return** $x_t, G(x_t)$

---

Notice that, as described in line 5, Algorithm 1 corresponds to a subgradient descent method with constant step size $\alpha$. Before taking a step in the direction of the subgradient though, the algorithm checks whether the objective function at the current state $x_t$ has a larger value than the value at its negative $-x_t$, and if so it updates the current state with its negative (line 3-4). This negation step allows the algorithm to escape the spurious critical point in a neighborhood of $-\rho_d x_\star$ where $\rho_d \in (0, 1)$, and it is motivated by the landscape analysis of $f_{\text{cs}}$ under the deterministic and probabilistic assumptions that we describe in the coming sections.

## 4 Recovery guarantees under deterministic conditions

The strategy taken in [18] and [21] to analyze the landscape of the minimization problem (1) and the convergence of Algorithm 1, consists in identifying a set of deterministic conditions on the measurements matrix $A$ and the generative network $G$, that ensure that the objective function $f_{\text{cs}}$ is well behaved and Algorithm 1 converges efficiently to an estimate of $x_\star$ and $y_\star$. These conditions are then shown to hold with high probability under probabilistic models for $A$ and $G$. This is akin to the results on compressed sensing with sparsity where, for example, recovery guarantees were developed under the Restricted Isometry Property [4].

The first condition, introduced in [18], is on the measurement matrix $A$ and ensures that $A^T A$ behaves like an isometry over differences of points in the range of a generative network $G$.

**Definition 4.1** (RRIC [18]). A matrix $A \in \mathbb{R}^{m \times n}$ satisfies the *Restricted Isometry Condition* with respect to $G$ with constant $\epsilon$ if for all $x_1, x_2, x_3, x_4 \in \mathbb{R}^k$, it holds that

$$\left| \left\langle \left( A^T A - I_n \right) \left( G(x_1) - G(x_2) \right), G(x_3) - G(x_4) \right\rangle \right| \leq \epsilon \| G(x_1) - G(x_2) \| \| G(x_3) - G(x_4) \|$$

The second deterministic condition introduced in [18] is on the weight matrices of $G$, ensures that they are approximately distributed like a Gaussian, and allows the control of how the layers of the network distort angles.

**Definition 4.2** (WDC [18]). We say that a generative network $G$ as in (2), satisfies the **Weight Distribution Condition** (WDC) with constant $\epsilon > 0$ if for all $i = 1, \ldots, d$, for all $r, s \in \mathbb{R}^{n_{i-1}}$:

$$\| (W_i)_{+,r}^T (W_i)_{+,s} - Q_{r,s} \|_2 \leq \epsilon, \tag{8}$$

Strictly speaking, in [18] the authors define the WDC as a property of a single weight matrix $W$, and then assume that the WDC is satisfied at each layer $W_i$ of $G$. This is equivalent to the definition above and simplifies the introduction of a novel, weaker, condition on the weight matrices, the R2WDC below.

**Definition 4.3** (R2WDC). We say that a generative network $G$ as in (2), satisfies the **Range Restricted Weight Distribution Condition** (R2WDC) with constant $\epsilon > 0$ if for all $i = 1, \ldots, d$, and for all

$x, y, x_1, x_2, x_3, x_4 \in \mathbb{R}^k$ , it holds that

$$\left| \langle ((W_i)^T_{+,r}(W_i)_{+,s} - Q_{r,s})u, v \rangle \right| \leq \epsilon \|u\| \|v\|,$$
$$\text{where} \ \ r = G_{i-1}(x),$$
$$s = G_{i-1}(y), \tag{9}$$
$$u = G_{i-1}(x_1) - G_{i-1}(x_2),$$
$$\text{and} \ \ v = G_{i-1}(x_3) - G_{i-1}(x_4).$$

Notice that the R2WDC is weaker than the WDC. Indeed, (8) and (9) are equivalent for $i = 1$, but for $i \geq 2$ equation (8) requires $(W_i)^t_{+,r}(W_i)_{+,s}$ to be close to the matrix $Q_{r,s}$ for any vector $r, s \in \mathbb{R}^{n_{i-1}}$ and when acting on any vector $u, v \in \mathbb{R}^{n_{i-1}}$, while equation (9) requires $(W_i)^t_{+,r}(W_i)_{+,s}$ to be close to the matrix $Q_{r,s}$ only for vectors $r, s$ on the range of $G_{i-1}$ and when acting on vectors $u, v \in \mathbb{R}^{n_{i-1}}$ given by the difference of points on the range of $G_{i-1}$. Notice that contrary to (8), defining the R2WDC (9) for layer $i$ requires considering the input/ouput pairs of the layers up to $i - 1$.

Our first technical result provides theoretical guarantees for efficiently estimating a target signal $y_\star$ on the range of a generative network from few linear measurements under the RRIC and the R2WDC .

**Theorem 4.4.** *Suppose $d \geq 2$, and $A$ and $G$ satisfy the RRIC and the R2WDC with constant $\epsilon < K_1/d^{90}$. Assume that $\|\eta\|_2 \leq \frac{K_2 \|x_\star\|_2}{d^{42} 2^{d/2}}$. Let $\{x_t\}$ be the iterates generated by Algorithm 1 with loss function $f_{cs}$, initial point $x_0 \in \mathbb{R}^k \setminus \{0\}$ and step size $\alpha = K_3 \frac{2^d}{d^2}$. Then there exists a number of steps $T$ satisfying $T \leq \frac{K_4 f(x_0) 2^d}{d^4 \epsilon \|x_\star\|_2^2}$ such that*

$$\|x_T - x_\star\|_2 \leq K_5 d^9 \|x_\star\|_2 \sqrt{\epsilon} + K_6 d^6 2^{d/2} \omega \|\eta\|_2.$$

*In addition, for all $t \geq T$, we have*

$$\|x_{t+1} - x_\star\|_2 \leq C^{t+1-T} \|x_T - x_\star\|_2 + K_7 2^{d/2} \|\eta\|_2,$$
$$\|G(x_{t+1}) - y_\star\|_2 \leq \frac{1.2}{2^{d/2}} C^{t+1-T} \|x_T - x_\star\|_2 + 1.2 K_7 \|\eta\|_2,$$

*where $C = 1 - \frac{7}{8} \frac{\alpha}{2^d} \in (0, 1)$. Here, $K_1, \ldots, K_7$ are universal positive constants.*

**Remark 1.** *The exponential factors $2^d$ appearing in the conditions and theses of the theorem are artifacts of the scaling of the weights of the generative network. For example, the output $G(x)$ of the network scales like $\|x\|_2/2^{d/2}$ and the loss function $f_{cs}(x)$ as $\|x\|_2^2/2^d$ (see for example Proposition C.1). Hence, for new constants $K_2', K_4'$ the bounds for $\eta$ and $T$ could be equivalently written as $\|\eta\|_2 \leq K_2' \|y_\star\|_2/d^{42}$ and $T \leq K_4' f(x_0)/(d^4 \epsilon \|y_\star\|_2^2)$. Choosing the weights of the network to be $\{\sqrt{2} W_i\}_{i \in [d]}$ would remove the $2^d$ factors in the above theorem (and scale the definition of R2WDC).*

This theorem shows that, despite the nonconvexity of the minimization problem (1), if the RRIC and the R2WDC hold with constant $\epsilon$, after $T = O(\epsilon^{-1})$ number of iterations the iterates of the subgradient descent method described in Algorithm 1 enter in a region of local convergence around $x_\star$. Moreover, after a large enough number of steps, $G(x_t)$ gives an estimate of the target signal $y_\star$ up to the noise level $O(\|\eta\|)$.

Theorem 3.1 in [21] shows that Theorem 4.4 holds assuming that the RRIC and the WDC hold. Our first technical contribution is to show that the WDC in Theorem 3.1 of [21], can be relaxed into the R2WDC. Relaxing the WDC into the R2WDC, will enable the relaxing of the expansivity assumption needed to show that the WDC holds for Gaussian generative networks as we demonstrate in Section 5.

We next describe the role of these deterministic conditions in the analysis of the problem (1). The full proof of Theorem 4.4 is given in Appendix C.

### 4.1 Global landscape analysis via the R2WDC

The analysis of [18] and [21] follows the approach recent line of works that analyze the global landscape geometry of non-convex optimization problems arising in statistical and signal recovery problems (see for example [36, 37, 14, 13] and [6] for an overview). The analysis roughly consists of two steps:

i) Showing that $f_{cs}(x) \approx f_E(x)$ and $\partial f_{cs}(x) \approx h_x$ uniformly over $x$.

ii) Analyzing the global properties of $f_E(x)$ and $h_x$, and transfer them to $f_{cs}(x)$ and $h_x$ using the first step.

Here $f_E(x)$ and $h_x$ are continuous functions of $x$, corresponding to the expected value of $f_{cs}(x)$ and $\partial f_{cs}(x)$ under Gaussian weights and measurement matrix $A$ (see next section for details) and zero noise. The RRIC and the WDC are used in [18] and [21] to obtain the uniform concentration in the first step, as well as directly proving convexity-like properties of $\partial f_{cs}(x)$ in the vicinity of $x_\star$.

To illustrate how the WDC and the R2WDC come into play, consider for simplicity the noiseless case $\eta = 0$. Then at a point $x \in \mathbb{R}^k$ where $G$ is differentiable, the gradient of $f_{cs}$ is given by

$$\nabla f_{cs}(x) = \Lambda_{d,x}^T A^T (A\Lambda_{d,x} x - A\Lambda_{d,x_\star} x_\star),$$
$$\approx \Lambda_{d,x}^T (\Lambda_{d,x} x - \Lambda_{d,x_\star} x_\star)$$

where $\Lambda_{d,x}$ and $\Lambda_{d,x_\star}$ ar defined in Section 2 and the approximation uses the fact that $A$ satisfies the RRIC with respect to $G$. Then if $G$ satisfies the WDC we have that

$$\nabla f_{cs}(x) \approx \Lambda_{d,x}^T (\Lambda_{d,x} x - \Lambda_{d,x_\star} x_\star)$$
$$= \Lambda_{d-1,x}^T (W_d)_{+,G_{d-1}(x)}^T (W_d)_{+,G_{d-1}(x)} \Lambda_{d-1,x} x - \Lambda_{d-1,x}^T (W_d)_{+,G_{d-1}(x)}^T (W_d)_{+,G_{d-1}(x_\star)} \Lambda_{d-1,x_\star} x_\star$$
$$= \Lambda_{d-1,x}^T \Big[ Q_{G_{d-1}(x),G_{d-1}(x)} + O(\epsilon) \Big] \Lambda_{d-1,x} x - \Lambda_{d-1,x}^T \Big[ Q_{G_{d-1}(x),G_{d-1}(x_\star)} + O(\epsilon) \Big] \Lambda_{d-1,x_\star} x_\star$$

where the last line used the WDC to control the concentration of $(W_d)_{+,G_{d-1}(x)}^T (W_d)_{+,G_{d-1}(x)}$ and $(W_d)_{+,G_{d-1}(x)}^T (W_d)_{+,G_{d-1}(x_\star)}$. The resulting terms are then controlled again applying the WDC to the the other $d-1$ weights of $G$, so that proceeding by induction over $d$ one obtains

$$\nabla f_{cs}(x) \approx h_x := \frac{1}{2^d} x - \frac{1}{2^d} \tilde{h}_{x,x_\star}, \tag{10}$$

where $\tilde{h}$ is a deterministic vector field defined in Appendix C.

In Appendix C we show that the R2WDC can be used to control directly the concentration of the terms

$$\Lambda_{d-1,x}^T (W_d)_{+,G_{d-1}(x)}^T (W_d)_{+,G_{d-1}(x)} \Lambda_{d-1,x} x$$

and

$$\Lambda_{d-1,x}^T (W_d)_{+,G_{d-1}(x)}^T (W_d)_{+,G_{d-1}(x_\star)} \Lambda_{d-1,x_\star} x_\star,$$

around their expectation (with respect to $W_d$) obtaining in this way

$$\nabla f_{cs}(x) \approx \Lambda_{d,x}^T (\Lambda_{d,x} x - \Lambda_{d,x_\star} x_\star)$$
$$= \Lambda_{d-1,x}^T \big[ Q_{G_{d-1}(x),G_{d-1}(x)} \big] \Lambda_{d-1,x} x - \Lambda_{d-1,x}^T \big[ Q_{G_{d-1}(x),G_{d-1}(x_\star)} \big] \Lambda_{d-1,x_\star} x_\star$$
$$+ O(\epsilon \|\Lambda_{d-1,x}\| \|\Lambda_{d-1,x} x\|) + O(\epsilon \|\Lambda_{d-1,x}\| \|\Lambda_{d-1,x_\star} x_\star\|)$$

Then again applying the R2WDC to the other layers of $G$, we can show that (10) still holds. We can then borrow the analysis of $h_x$ from [21] and obtain the same convergence guarantees.

The advantage of using the R2WDC over the original WDC, is that it is satisfied by random generative networks with contractive layers as we demonstrate in the next section.

## 5 Recovery guarantees under probabilistic assumptions

In this section we give probabilistic models for the measurement matrix $A$, generative network $G$, and noise vector $\eta$ that will ensure that the RRIC and the R2WDC are satisfied with high probability and Algorithm 1 efficiently estimate the target signal $y_\star$ up to an error of the order $\tilde{O}(\sqrt{k/m}\|\eta\|)$.

We make the following assumption on the sensing matrix $A \in \mathbb{R}^{m \times n}$.

**Assumptions A.**

*A.1* $A$ is independent from $\{W_i\}_{i=1}^d$.

***A.2*** *A has i.i.d. $\mathcal{N}(0, 1/m)$ entries.*

***A.3*** *There are sufficient number of linear measurements:*

$$m \geq \widehat{C}_\epsilon \cdot k \cdot \log \prod_{j=1}^{d} \frac{e\, n_i}{k}, \tag{11}$$

*where $\widehat{C}_\epsilon$ depends polynomially on $\epsilon^{-1}$.*

Under Assumptions A, the measurement matrix satisfies the RRIC with respect to $G$ with high probability.

**Lemma 5.1** (Consequence of Proposition 6 in [18])**.** *Let Assumptions A be satisfied. Then A satisfies the RRIC with constat $\epsilon > 0$ with respect to $G$, with probability at least*

$$1 - \hat{\gamma} e^{-\hat{c}\epsilon m}$$

*where $\hat{\gamma}$ and $\hat{c}$ are positive universal constants.*

*Proof.* This result is proved in Proposition 6 in [18] for a number of measurements $m$ satisfying $m \geq C'_\epsilon \cdot k \cdot d \cdot \log \prod_{j=1}^{d} n_j$ where $C'_\epsilon$ depends polynomially on $\epsilon$. To imporove the lower bound on $m$ to (11) it is enough to follow the proof of Proposition 6 in [18] and use the sharper upper bound on the number of affine subspaces in the range of a gnerative network given in Lemma D.1. $\square$

We then provide a probabilistic model for a generative network $G : \mathbb{R}^k \to \mathbb{R}^n$ as in (2).

**Assumptions B.**

***B.1*** *Each weight matrix $W_i \in \mathbb{R}^{n_i \times n_{i-1}}$ have i.i.d. $\mathcal{N}(0, 1/n_i)$ entries.*

***B.2*** *The first layer satisfies $n_1 \geq \widetilde{C}_\epsilon \cdot k$, and for any $i = 2, \dots, d$:*

$$n_i \geq \widetilde{C}_\epsilon \cdot k \cdot \log \prod_{j=1}^{i-1} \frac{e\, n_j}{k}, \tag{12}$$

*where $\tilde{C}_\epsilon$ depends polynomially on $\epsilon^{-1}$.*

***B.3*** *The $\{W_j\}_{j=1}^{d}$ are independent.*

Under Assumptions B, the generative network $G$ satisfies the R2WDC .

**Lemma 5.2.** *Fix $0 < \epsilon < 1$. Consider a $d$-layer ReLU network $G$ with weight matrices $\{W_i\}_{i=1}^{d}$. Assume that the $\{W_i\}_{i=1}^{d}$ satisfy Assumptions B. Then $G$ satisfies the R2WDC with constant $\epsilon$ with probability at least*

$$1 - \gamma \left( \frac{en_1}{k} \right)^{2k} e^{-c_\epsilon n_1} - \gamma \sum_{i=2}^{d} \left( \frac{e\, n_i}{k+1} \right)^{4k} e^{-c_\epsilon n_i/2}$$

*where $c_\epsilon$ depends polynomially on $\epsilon^{-1}$ and $\gamma$ is a positive absolute constant.*

We finally conclude with some assumptions on the noise vector $\eta \in \mathbb{R}^m$.

**Assumption C.** *The noise vector $\eta$ is independent from A and the weights $\{W_i\}_{i=1}^{d}$*

The next lemma is used to bound the perturbation of the objective function $f_{\mathrm{cs}}$ and its gradient due to the presence of the noise term $\eta$. These bounds are then used to show that Algorithm 1 leads to a reconstruction of $y_\star$ of the order $O(\sqrt{k/m}\|\eta\|)$.

**Lemma 5.3.** *Suppose $G : \mathbb{R}^k \to \mathbb{R}^n$ satisfies the R2WDC with $\epsilon < 1/(16\pi d^2)^2$ and $d \geq 2$. Let $A \in \mathbb{R}^{m \times n}$ be a matrix with i.i.d. entries $\mathcal{N}(0, 1/m)$ and $\eta \in \mathbb{R}^m$ satisfies Assumption C. Let*

$$\omega := \frac{2}{2^{d/2}} \sqrt{\frac{13}{12}} \sqrt{\frac{k}{m} \log \left( 5 \prod_{j=1}^{d} \frac{e\, n_i}{k} \right)}. \tag{13}$$

*Then with probability at least*

$$1 - e^{-\frac{k}{2}\log(5\prod_{i=1}^{d}\frac{e\,n_i}{k})}$$

*for every $x \in \mathbb{R}^k$ we have that*

$$\langle x, \Lambda_x^T A^T \eta \rangle \leq \omega\|\eta\|\|x\|, \tag{14}$$

*if in addition $G$ is differentiable at $x$ we also have that*

$$\|\Lambda_x^T A^T \eta\| \leq \omega\|\eta\|. \tag{15}$$

Given the previous assumptions, we are now ready to state the main result of this section.

**Theorem 5.4.** *Suppose $d \geq 2$, $\epsilon < K_1/d^{90}$ and $\omega\|\eta\|_2 \leq \frac{K_2\|x_\star\|_2}{d^{42}2^{d/2}}$ where $\omega$ is defined in* (13).
*Assume that $A$, $G$ and $\eta$ satisfy Assumptions A, B and C. Then with probability at least*

$$1 - \gamma\Big(\frac{e\,n_1}{k}\Big)^{2k}e^{-c_\epsilon n_1} - \gamma\sum_{i=2}^{d}\Big(\frac{e\,n_i}{k+1}\Big)^{4k}e^{-c_\epsilon n_i/2} - \hat{\gamma}e^{-\hat{c}\epsilon m} - e^{-\frac{k}{2}\log(5\prod_{i=1}^{d}\frac{e\,n_i}{k})}, \tag{16}$$

*where $\gamma$, $\hat{\gamma}$ and $\hat{c}$ are positive universal constants, the following holds. Let $\{x_t\}$ be the iterates generated by Algorithm 1 with loss function $f_{\mathrm{cs}}$, initial point $x_0 \in \mathbb{R}^k \setminus \{0\}$ and step size $\alpha = K_3\frac{2^d}{d^2}$. There exists a number of steps $T$ satisfying $T \leq \frac{K_4 f(x_0)2^d}{d^4\epsilon\|x_\star\|_2}$ such that*

$$\|x_T - x_\star\|_2 \leq K_5 d^9\|x_\star\|_2\sqrt{\epsilon} + K_6 d^6 2^{d/2}\omega\|\eta\|_2.$$

*In addition, for all $t \geq T$, we have*

$$\|x_{t+1} - x_\star\|_2 \leq C^{t+1-T}\|x_T - x_\star\|_2 + K_7 2^{d/2}\omega\|\eta\|_2,$$

$$\|G(x_{t+1}) - y_\star\|_2 \leq \frac{1.2}{2^{d/2}}C^{t+1-T}\|x_T - x_\star\|_2 + 1.2K_7\omega\|\eta\|_2,$$

*where $C = 1 - \frac{7}{8}\frac{\alpha}{2^d} \in (0,1)$. Here, $K_1,\ldots,K_7$ are universal positive constants.*

*Proof.* Combining Lemma 5.1, Lemma 5.2 and Theorem 4.4 we obtain Theorem 5.4 with $\omega = 1$ and probability at least

$$1 - \gamma\Big(\frac{e\,n_1}{k}\Big)^{2k}e^{-c_\epsilon n_1} - \gamma\sum_{i=2}^{d}\Big(\frac{e\,n_i}{k+1}\Big)^{4k}e^{-c_\epsilon n_i/2} - \hat{\gamma}e^{-\hat{c}\epsilon m}.$$

Inspecting the proof of Theorem 3.1 in [21], it is easy to see that if Lemma 5.3 holds, then the conclusions of Theorem 5.4 hold with $\omega$ given by (13) and probability at least (16). □

**Remark 2.** *As for Theorem 4.4, the exponential factors $2^d$ are artifacts of the scaling of the weights of the network. Had the entries of $W_i$ been drawn from $\mathcal{N}(0, 2/n_i)$ the $2^d$ factors would not be present.*

**Remark 3.** *Notice that $4k\log(en/(k+1)) \leq 4k\log(n)/\log(2)$ for every $n \geq 2$. Thus if for every $i = 1,\ldots,d$, it holds that*

$$\frac{n_i}{\log(n_i)} \geq \frac{16 \cdot k \cdot c_\epsilon^{-1}}{\log(2)} \tag{17}$$

*the conclusions of the theorem hold with nontrivial probability bounds. In Appendix G we provide an example of a generative network $G$ with contractive layers satisfying both* (12) *and* (17).

Theorem 5.4 provides guarantees for the efficient recovery of a signal $y_\star$ in the range of a generative network $G$ from few noisy linear measurements, using a nonconvex (sub)gradient descent method. Notice that the intrinsic dimension of the signal $y_\star$ is $k$ (the dimension of the latent space) and the number of measurements required $m$ is proportional to $k$ and information-theoretically optimal up to log factors in the widths of the network and polynomials in the depth. Notice moreover, that up to these factors, the width $n_i$ of each layer of the network is also required to be linear in $k$. This is necessary to ensure that each subnetwork $G_i : \mathbb{R}^k \to \mathbb{R}^{n_i}$ is invertible, and it is weaker than

the assumptions in the previous works that required $n_i$ to be linear in $n_{i-1}$ in order to ensure the invertibility of every single layer.

In Appendix H we empirically verify the predictions of Theorem 5.4, demonstrating how (a practical variant of) Algorithm 1 recover signals $y_\star$ in the range of non-expansive generative networks from undersampled noisy measurements. We show that the recovery is linear in $k/m$ and that in practice the dependence on the depth $d$ of the networks is milder than that predicted by our theory. We leave for future works the establishing of sharper bounds in the depth $d$.

Limitations of the current and previous works on theoretical guarantees for signal recovery with generative networks are the Gaussian assumption on the weights and the absence of biases. Important directions of future research are the inclusion of biases in the generative network and the departure from the Gaussian weights assumptions for more realistic probabilistic models.

## Acknowledgments and Disclosure of Funding

I would like to thank Paul Hand for comments on an earlier version of this manuscript and Babhru Joshi for helpful discussions. This research was partially supported by the NSF award DMS-1848087.

## References

[1] Rameen Abdal, Yipeng Qin, and Peter Wonka. Image2stylegan: How to embed images into the stylegan latent space? In *Proceedings of the IEEE/CVF International Conference on Computer Vision*, pages 4432–4441, 2019.

[2] Aviad Aberdam, Dror Simon, and Michael Elad. When and how can deep generative models be inverted? *arXiv preprint arXiv:2006.15555*, 2020.

[3] Ashish Bora, Ajil Jalal, Eric Price, and Alexandros G Dimakis. Compressed sensing using generative models. In *Proceedings of the 34th International Conference on Machine Learning-Volume 70*, pages 537–546. JMLR. org, 2017.

[4] Emmanuel J Candès, Justin Romberg, and Terence Tao. Robust uncertainty principles: Exact signal reconstruction from highly incomplete frequency information. *IEEE Transactions on information theory*, 52(2):489–509, 2006.

[5] Emmanuel J Candes, Justin K Romberg, and Terence Tao. Stable signal recovery from incomplete and inaccurate measurements. *Communications on Pure and Applied Mathematics: A Journal Issued by the Courant Institute of Mathematical Sciences*, 59(8):1207–1223, 2006.

[6] Yuejie Chi, Yue M Lu, and Yuxin Chen. Nonconvex optimization meets low-rank matrix factorization: An overview. *IEEE Transactions on Signal Processing*, 67(20):5239–5269, 2019.

[7] Christian Clason. Nonsmooth Analysis and Optimization. *arXiv preprint arXiv:1708.04180*, 2017.

[8] Jorio Cocola, Paul Hand, and Vlad Voroninski. Nonasymptotic guarantees for spiked matrix recovery with generative priors. *Advances in Neural Information Processing Systems*, 33:15185–15197, 2020.

[9] Jorio Cocola, Paul Hand, and Vladislav Voroninski. No statistical-computational gap in spiked matrix models with generative network priors. *Entropy*, 23(1):115, 2021.

[10] Constantinos Daskalakis, Dhruv Rohatgi, and Emmanouil Zampetakis. Constant-expansion suffices for compressed sensing with generative priors. In H. Larochelle, M. Ranzato, R. Hadsell, M. F. Balcan, and H. Lin, editors, *Advances in Neural Information Processing Systems*, volume 33, pages 13917–13926. Curran Associates, Inc., 2020.

[11] Alex Dimakis. Deep generative models and unsupervised methods for inverse problems, `https://youtu.be/OsrR9Ar1tVc?t=2069`, October 2021. In *Algorithmic Advances for Statistical Inference with Combinatorial Structure*.

[12] David L Donoho. For most large underdetermined systems of linear equations the minimal $\ell_1$-norm solution is also the sparsest solution. *Communications on Pure and Applied Mathematics: A Journal Issued by the Courant Institute of Mathematical Sciences*, 59(6):797–829, 2006.

[13] Rong Ge, Chi Jin, and Yi Zheng. No spurious local minima in nonconvex low rank problems: A unified geometric analysis. In *International Conference on Machine Learning*, pages 1233–1242. PMLR, 2017.

[14] Rong Ge, Jason D Lee, and Tengyu Ma. Matrix completion has no spurious local minimum. *Advances in neural information processing systems*, 29, 2016.

[15] Paul Hand and Babhru Joshi. Global guarantees for blind demodulation with generative priors. In *Advances in Neural Information Processing Systems*, pages 11531–11541, 2019.

[16] Paul Hand, Oscar Leong, and Vlad Voroninski. Phase retrieval under a generative prior. In *Advances in Neural Information Processing Systems*, pages 9136–9146, 2018.

[17] Paul Hand, Oscar Leong, and Vladislav Voroninski. Compressive phase retrieval: Optimal sample complexity with deep generative priors. *arXiv preprint arXiv:2008.10579*, 2020.

[18] Paul Hand and Vladislav Voroninski. Global guarantees for enforcing deep generative priors by empirical risk. *IEEE Transactions on Information Theory*, 66(1):401–418, 2019.

[19] Reinhard Heckel and Paul Hand. Deep Decoder: Concise image representations from untrained non-convolutional networks. In *International Conference on Learning Representations*, 2019.

[20] Reinhard Heckel, Wen Huang, Paul Hand, and Vladislav Voroninski. Rate-optimal denoising with deep neural networks. *Information and Inference: A Journal of the IMA*, 10(4):1251–1285, 2021.

[21] Wen Huang, Paul Hand, Reinhard Heckel, and Vladislav Voroninski. A provably convergent scheme for compressive sensing under random generative priors. *Journal of Fourier Analysis and Applications*, 27(2):1–34, 2021.

[22] Babhru Joshi, Xiaowei Li, Yaniv Plan, and Ozgur Yilmaz. Plugin: A simple algorithm for inverting generative models with recovery guarantees. *Advances in Neural Information Processing Systems*, 34, 2021.

[23] Babhru Joshi, Xiaowei Li, Yaniv Plan, and Ozgur Yilmaz. PLUGIn-CS: A simple algorithm for compressive sensing with generative prior. In *NeurIPS 2021 Workshop on Deep Learning and Inverse Problems*, 2021.

[24] Tero Karras, Miika Aittala, Samuli Laine, Erik Härkönen, Janne Hellsten, Jaakko Lehtinen, and Timo Aila. Alias-free generative adversarial networks. *Advances in Neural Information Processing Systems*, 34:852–863, 2021.

[25] Tero Karras, Samuli Laine, Miika Aittala, Janne Hellsten, Jaakko Lehtinen, and Timo Aila. Analyzing and improving the image quality of stylegan. In *Proceedings of the IEEE/CVF conference on computer vision and pattern recognition*, pages 8110–8119, 2020.

[26] Qi Lei, Ajil Jalal, Inderjit S Dhillon, and Alexandros G Dimakis. Inverting deep generative models, one layer at a time. *Advances in neural information processing systems*, 32, 2019.

[27] Fangchang Ma, Ulas Ayaz, and Sertac Karaman. Invertibility of convolutional generative networks from partial measurements. *Advances in Neural Information Processing Systems*, 31, 2018.

[28] Morteza Mardani, Enhao Gong, Joseph Y Cheng, Shreyas S Vasanawala, Greg Zaharchuk, Lei Xing, and John M Pauly. Deep generative adversarial neural networks for compressive sensing mri. *IEEE transactions on medical imaging*, 38(1):167–179, 2018.

[29] Jiri Matousek. *Lectures on discrete geometry*, volume 212. Springer Science & Business Media, 2013.

[30] Sachit Menon, Alexandru Damian, Shijia Hu, Nikhil Ravi, and Cynthia Rudin. Pulse: Self-supervised photo upsampling via latent space exploration of generative models. In *Proceedings of the ieee/cvf conference on computer vision and pattern recognition*, pages 2437–2445, 2020.

[31] Lukas Mosser, Olivier Dubrule, and Martin J Blunt. Stochastic seismic waveform inversion using generative adversarial networks as a geological prior. *Mathematical Geosciences*, 52(1):53–79, 2020.

[32] Gregory Ongie, Ajil Jalal, Christopher A Metzler, Richard G Baraniuk, Alexandros G Dimakis, and Rebecca Willett. Deep learning techniques for inverse problems in imaging. *IEEE Journal on Selected Areas in Information Theory*, 1(1):39–56, 2020.

[33] Xingang Pan, Bo Dai, Ziwei Liu, Chen Change Loy, and Ping Luo. Do 2d {gan}s know 3d shape? unsupervised 3d shape reconstruction from 2d image {gan}s. In *International Conference on Learning Representations*, 2021.

[34] Shuang Qiu, Xiaohan Wei, and Zhuoran Yang. Robust one-bit recovery via relu generative networks: Near-optimal statistical rate and global landscape analysis. In *International Conference on Machine Learning*, pages 7857–7866. PMLR, 2020.

[35] Ganlin Song, Zhou Fan, and John Lafferty. Surfing: Iterative optimization over incrementally trained deep networks. *Advances in Neural Information Processing Systems*, 32, 2019.

[36] Ju Sun, Qing Qu, and John Wright. Complete dictionary recovery over the sphere i: Overview and the geometric picture. *IEEE Transactions on Information Theory*, 63(2):853–884, 2016.

[37] Ju Sun, Qing Qu, and John Wright. A geometric analysis of phase retrieval. *Foundations of Computational Mathematics*, 18(5):1131–1198, 2018.

[38] Roman Vershynin. *High-dimensional probability: An introduction with applications in data science*, volume 47. Cambridge university press, 2018.

[39] Jun-Yan Zhu, Philipp Krähenbühl, Eli Shechtman, and Alexei A Efros. Generative visual manipulation on the natural image manifold. In *European conference on computer vision*, pages 597–613. Springer, 2016.


## A Roadmap

In Appendix C we establish the main consequences of the R2WDC that are then used to prove Theorem 4.4. Then in Appendix D, we prove Lemma 5.2 showing that a Gaussian generative network satisfies the R2WDC with high probability. In Appendix E we analyze the perturbation of the gradient and objective function due to the noise term $\eta$, and provide the proof of Lemma 5.3. Extension of the recovery guarantees for Phase Retrieval, Denosing, and Spiked Matrix Recovery are discussed in Appendices F.1, F.2, F.3 respectively. Finally, in Appendix G we give an example of a network with contractive layers, satisfying the assumptions of our main theorems, and in Appendix H we verify the prediction of our theory with synthetic experiments.

## B Notation

For any vector $x$ we denote with $\|x\|$ its Euclidean norm and for any matrix $A$ we denote with $\|A\|$ its spectral norm and with $\|A\|_F$ its Frobenius norm. The euclidean inner product between two vectors $a$ and $b$ is $\langle a, b \rangle$. For a set $S$ we will write $|S|$ for its cardinality and $S^c$ for its complement. Let $\mathcal{B}(x, r)$ be the Euclidean ball of radius $r$ centered at $x$, and $\mathcal{S}^{k-1}$ be the unit sphere in $\mathbb{R}^k$. We will use $a = b + O_1(\delta)$ when $\|a - b\| \leq \delta$, where the norm is understood to be the absolute value for scalars, the Euclidean norm for vectors and the spectral norm for matrices.

## C Consequences of the R2WDC

Following [18], we define the function $g : [0, \pi] \to \mathbb{R}$ which describes how the operator $x \mapsto W_{+,x}$ distorts angles:

$$g(\theta) := \cos^{-1}\Big(\frac{(\pi - \theta)\cos\theta + \sin\theta}{\pi}\Big). \tag{18}$$

For two nonzero vectors $x, y$ we let $\bar{\theta}_0 = \angle(x, y)$ and define inductively $\bar{\theta}_i := g(\bar{\theta}_{i-1})$. Then we set

$$\tilde{h}_{x,y} := \frac{1}{2^d}\Big[\Big(\prod_{i=0}^{d-1}\frac{\pi - \bar{\theta}_i}{\pi}\Big)y + \sum_{i=1}^{d-1}\frac{\sin\bar{\theta}_i}{\pi}\Big(\prod_{j=i+1}^{d-1}\frac{\pi - \bar{\theta}_j}{\pi}\Big)\|y\|\hat{x}\Big]. \tag{19}$$

**Proposition C.1.** *Fix $\epsilon > 0$ such that $\max(2d\epsilon, 10\epsilon) < 1$. Let $G$ be a generative network as in (2) satisfying the R2WDC with constant $\epsilon$. Then for any $x \in \mathbb{R}^k$ and $j \in [d]$*

$$\|x\|^2\Big(\frac{1}{2} - \epsilon\Big)^j \leq \|G_j(x)\|^2, \leq \Big(\frac{1}{2} + \epsilon\Big)^j\|x\|^2 \tag{20a}$$

$$\|G(x)\|^2 \leq \frac{1 + 4\epsilon d}{2^d}\|x\|^2. \tag{20b}$$

*Moreover, for any $x \neq 0$, $y \neq 0$, $j \in [d]$, the angle $\theta_j = \angle(G_j(x), G_j(y))$ is well-defined and*

$$|\theta_j - g(\theta_{j-1})| \leq 4\sqrt{\epsilon} \tag{21a}$$

$$\langle G(x), G(y) \rangle \geq \frac{1}{4\pi}\frac{1}{2^d}\|x\|\|y\| \tag{21b}$$

$$|\langle G(x), G(y) \rangle - \langle x, \tilde{h}_{x,y} \rangle| \leq 24\frac{d^3\sqrt{\epsilon}}{2^d}\|x\|\|y\| \tag{21c}$$

*where $g$ is given in (18) and $\tilde{h}$ in (19).*

The next result is used to prove concentration of the gradient of the objective function around its expectation.

**Proposition C.2.** *Fix $0 < \epsilon < d^{-4}/(16\pi)^2$ and $d \geq 2$. Suppose that $G$ as in (2) satisfies the R2WDC with constant $\epsilon$. Let $x \in \mathbb{R}^k$ be a point where $G(x)$ is differentiable, and $y \in \mathbb{R}^k \setminus \{0\}$, then*

$$\|\Lambda_{d,x}\|^2 \leq \frac{1 + 4\epsilon d}{2^d} \leq \frac{13}{12}\frac{1}{2^d} \tag{22}$$

$$\|\Lambda_{d,x}^t\Lambda_{d,x} - \frac{1}{2^d}I_k\| \leq \frac{4\epsilon d}{2^d} \tag{23}$$

$$\|\Lambda_{d,x}^t\Lambda_{d,y}y - \tilde{h}_{x,y}\| \leq 24\frac{d^3\sqrt{\epsilon}}{2^d}\|y\| \tag{24}$$

The next proposition uses the R2WDC to bound the local Lipschitz constant of the ReLU-networks $\{G_j\}_{j \in [d]}$.

**Proposition C.3.** *Suppose that $x \in \mathcal{B}(x, d\sqrt{\epsilon}\|y\|)$ and $G$ satisfies the R2WDC with $\epsilon < 1/(200)^4/d^6$. Then for every $i \in [d]$, it holds that*

$$\|G_i(x) - G_i(y)\| \leq \frac{1.2}{2^{i/2}}\|x - y\| \tag{25}$$

The next proposition is used to show that when $x$ is close to $y$, the gradient of the objective function points in a direction that decreases the distance between of $x$ and $y$.

**Proposition C.4.** *Suppose $x \in \mathcal{B}(y, d\sqrt{\epsilon}\|y\|)$ is a differentiable point for $G$, and the R2WDC holds with $\epsilon < 1/(200)^4/d^6$. Then it holds that*

$$\Lambda_x^T(\Lambda_x x - \Lambda_y y) = \frac{1}{2^d}(x - y) + \frac{1}{2^d}\frac{1}{16}\|x - y\|O_1(1) \tag{26}$$

We can now prove Theorem 4.4.

*Proof of Theorem 4.4.* The proof of Theorem 3.1 in [21] only uses the inequality (20a)-(26), which are proved for a network satisfying the WDC. The previous propositions have shown that such inequalities hold under the weaker R2WDC. Therefore from the proof of Theorem 3.1 in [21] combined with the Propositions C.1-C.4, we obtain automatically the proof of Theorem 4.4. □

## C.1 Supplemental Results for Section C

**Proof of Proposition C.1**

*Proof.* For $x, y \in \mathbb{R}^k$ and $j \in [d]$, below we write $x_j := G_j(x)$ and $y_j := G_j(y)$.

*- Proof of* (20a)
Notice that by (9) for $x \in \mathbb{R}^k$

$$\left(\frac{1}{2} - \epsilon\right)\|x_{j-1}\|^2 \leq \|x_j\|^2 \leq \left(\frac{1}{2} + \epsilon\right)\|x_{j-1}\|^2,$$

which proceeding by induction gives (20a).

*- Proof of* (20b)
Next observe that since $\log(1 + z) \leq z$, $e^z \leq 1 + 2z$ for $z < 1$ and $2d\epsilon \leq 1$, from (20a) we have

$$\|G_d(x)\|^2 \leq \frac{(1 + 2\epsilon)^d}{2^d}\|x\|^2 \leq \frac{1}{2^d}e^{d\log(1+2\epsilon)}\|x\| \leq \frac{1 + 4\epsilon d}{2^d}\|x\|,$$

which corresponds to (20b).

*- Proof of* (21a)
Assume that $x, y \in \mathbb{R}^k \setminus \{0\}$. Then, the assumption $2d\epsilon \leq 1$ and the lower bound in (20a) imply that $\theta_j$ are well-defined for all $j \in [d]$. To prove then (21a) notice that it is sufficient to prove that for any $j \in [d]$ it holds that

$$\left|\cos\theta_j - \frac{(\pi - \theta_{j-1})\cos\theta_{j-1} + \sin\theta_{j-1}}{\pi}\right| \leq 5\epsilon$$

By homogeneity of the ReLU activation function, we can assume without loss of generality that $\|x_{j-1}\| = \|y_{j-1}\| = 1$. Let

$$\delta_1 := \langle x_{j-1}, (W_{j,+,x}^T W_{j,+,y} - Q_{x_{j-1},y_{j-1}})y_{j-1}\rangle$$
$$\delta_2 := \langle x_{j-1}, (W_{j,+,x}^T W_{j,+,x} - I_k/2)y_{j-1}\rangle$$
$$\delta_3 := \langle y_{j-1}, (W_{j,+,y}^T W_{j,+,y} - I_k/2)y_{j-1}\rangle$$

and notice that by the R2WDC we have $\max(|\delta_1|, |\delta_2|, |\delta_3|) \leq \epsilon$. Thus,

$$
\begin{aligned}
\cos\theta_j &= \frac{\langle x_j, y_j\rangle}{\|x_j\|\|y_j\|} \\
&= \frac{\langle x_{j-1}, W_{j,+,x}^T W_{j,+,y} y_{j-1}\rangle}{\sqrt{\langle x_{j-1}, W_{j,+,x}^T W_{j,+,x} x_{j-1}\rangle \langle y_{j-1}, W_{j,+,y}^T W_{j,+,y} y_{j-1}\rangle}} \\
&= 2\frac{\langle x_{j-1}, Q_{x_{j-1},y_{j-1}} y_{j-1}\rangle + \delta_1}{\sqrt{(1+2\delta_2)(1+2\delta_3)}}.
\end{aligned}
$$

Finally, notice that $2\langle x_{j-1}, Q_{x_{j-1},y_{j-1}} y_{j-1}\rangle = [(\pi - \theta_{j-1})\cos\theta_{j-1} + \sin\theta_{j-1}]/\pi$ so

$$
\begin{aligned}
|\cos\theta_j - 2\langle x_{j-1}, Q_{x_{j-1},y_{j-1}} y_{j-1}\rangle| &\leq 2|\langle x_{j-1}, Q_{x_{j-1},y_{j-1}} y_{j-1}\rangle| \left|1 - \frac{1}{\sqrt{(1+2\delta_2)(1+2\delta_3)}}\right| \\
&\quad + \frac{2|\delta_1|}{\sqrt{(1+2\delta_2)(1+2\delta_3)}} \\
&\leq \left|1 - \frac{1}{(1-2\epsilon)}\right| + \frac{2\epsilon}{(1-2\epsilon)} \\
&\leq 5\epsilon
\end{aligned}
$$

where the second inequality follows from $|2\langle x_{j-1}, Q_{x_{j-1},y_{j-1}} y_{j-1}\rangle| \leq 1$ and $\max(|\delta_1|, |\delta_2|, |\delta_3|) \leq \epsilon$, and the third inequality from $10\epsilon < 1$.

*- Proof of* (21b)
By (20a) and $\epsilon \leq 1/2$, it follows that $\|x_d\|\|y_d\| \geq \frac{(1-2\epsilon)^d}{2^d}\|x\|\|y\|\frac{1-2d\epsilon}{2^d}$. Moreover, let $\delta := 4\sqrt{\epsilon}$, then by (21a) we have that $\theta_j = g(\theta_{j-1}) + O_1(\delta)$. Thus, $\theta_d = g(g(\cdots g(g(\theta_0) + O_1(\delta)) + O_1(\delta)\cdots) + O_1(\delta))) + O_1(\delta)$ and for $\bar\theta = g^{\circ d}(\theta_0)$, so that, using $g'(\theta) \leq 1$ for all $\theta$, we have

$$
|\theta_d - \bar\theta_d| \leq d\delta. \tag{27}
$$

Then by (27), $\bar\theta_d \leq \cos^{-1}(1/\pi)$ for $d \geq 2$, and $16\pi d\sqrt{\epsilon} < 1$, follows that $\cos\theta_d \geq 3/(4\pi)$.

Finally, if $2d\epsilon \leq 2/3$, we can then conclude that

$$
\langle G(x), G(y)\rangle \geq \cos(\theta_d)\|x_d\|\|y_d\| \geq \frac{1}{4\pi}\frac{1}{2^d}\|x\|\|y\|.
$$

*- Proof of* (21c)
The following result on a recurrence relation will be used in the subsequent analysis

$$
\Gamma_d = s_d \Gamma_{d-1} + r_d, \quad \Gamma_0 = y \implies \Gamma_d = \left(\prod_{i=1}^d s_i\right) y + \sum_{i=1}^d \left(r_i \prod_{j=i+1}^d s_j\right) \tag{28}
$$

Define $\Gamma_d := \langle x_d, y_d\rangle$, then

$$
\begin{aligned}
\Gamma_d &= \langle x_{d-1}, W_{d-1,+,x}^T W_{d-1,+,x} y_{d-1}\rangle \\
&= \langle x_{d-1}, Q_{x_{d-1},y_{d-1}} y_{d-1}\rangle + O_1(\epsilon)\|y_{d-1}\|\|x_{d-1}\|, \\
&= \frac{\pi - \theta_{d-1}}{2\pi}\Gamma_{d-1} + \frac{\sin\theta_{d-1}}{2\pi}\|x_{d-1}\|\|y_{d-1}\| + O_1(\epsilon)\|y_{d-1}\|\|x_{d-1}\|, \\
&= \frac{\pi - \theta_{d-1}}{2\pi}\Gamma_{d-1} + \frac{\sin\theta_{d-1}}{2\pi}\frac{\|x\|\|y\|}{2^{d-1}} + \frac{\epsilon}{2^d}\left(\frac{4\epsilon d}{\pi} + 2(1+4\epsilon d)\right)\|y\|\|x\|O_1(1), \\
&= \frac{\pi - \theta_{d-1}}{2\pi}\Gamma_{d-1} + \frac{\sin\theta_{d-1}}{2\pi}\frac{\|x\|\|y\|}{2^{d-1}} + 11d\epsilon\frac{\|y\|\|x\|}{2^d}O_1(1),
\end{aligned}
$$

Where the second equality follows from the R2WDC , the third from the definition of $Q_{p,q}$. The rest of the proof proceeds as in the proof of Lemma 8 in [18]. □

**Proof of Proposition C.2**

*Proof.*

*- Proof of* (23).
Let $x \in \mathbb{R}^k$ be a point where $G$ is differentiable, and notice that for small enough $z$, by local linearity of $G$, we have $G(x + z) = \Lambda_x z$. Then the R2WDC gives for $j \in [d]$

$$\left| \langle (W_{j,+,x}^T W_{j,+,y} - I_k/2)\Lambda_{j-1,x}z, \Lambda_{j-1,x}z \rangle \right| \leq \epsilon \|\Lambda_{j-1,x}\|^2 \|z\|^2$$

for all $z$, which in turn implies

$$\left\| \Lambda_{j,x}^T \Lambda_{j,x} - \frac{1}{2}\Lambda_{j-1,x}^T \Lambda_{j-1,x} \right\| \leq \epsilon \|\Lambda_{j-1,x}^T \Lambda_{j-1,x}\|. \tag{30}$$

Let now $M_d := \Lambda_{d,x}^T \Lambda_{d,x}$ with $M_0 = I_k$, then

$$M_d = \frac{1}{2}M_{d-1} + \|M_{d-1}\|O_1(\epsilon). \tag{31}$$

We then obtain

$$\|M_d\| \leq \left(\frac{1}{2} + \epsilon\right)\|M_{d-1}\| \leq \frac{(1 + 2\epsilon)^d}{2^d}\|M_0\| \leq \frac{1 + 4\epsilon d}{2^d},$$

where the second inequality, and the third inequality uses $2d\epsilon \leq 1$ and the same reasoning as in the proof of (20b). From (22) and (31) we obtain the following recurrence relation

$$M_d = \frac{1}{2}M_{d-1} + O_1\left(\epsilon\frac{1 + 4\epsilon(d-1)}{2^{d-1}}\right),$$

which, using (28) and $4\epsilon d \leq 1$, gives

$$M_d = \frac{1}{2}I_k + \sum_{i=1}^{d} O_1\left(\epsilon\frac{1 + 4\epsilon(i-1)}{2^{i-1}}\right)\frac{1}{2^{d-i}}$$

$$= \frac{1}{2}I_k + \frac{4\epsilon d}{2^d}O_1(1)$$

*- Proof of* (24).
Notice again that if $x \in \mathbb{R}^k$ is a differentiable point for $G$, the R2WDC gives for any $j \in [d]$

$$\|\Lambda_{j,x}^T \Lambda_{j,y}y - \Lambda_{j-1,x}^T Q_{x_{j-1},y_{j-1}}\Lambda_{j-1,y}y\| \leq \epsilon\|\Lambda_{j-1,x}\|\|G_{j-1}(y)\|. \tag{32}$$

We then let $\Gamma_d := \Lambda_{d,x}^T \Lambda_{d,y}y$ and observe that

$$\Gamma_d = \Lambda_{d-1,x}^T Q_{x_{d-1},y_{d-1}}\Lambda_{d-1,y}y + \|\Lambda_{d-1,x}\|\|G_{d-1}(y)\|O_1(\epsilon)$$

$$= \frac{\pi - \theta_{d-1}}{2\pi}\Gamma_{d-1} + \frac{\sin\theta_{d-1}}{2\pi}\frac{\|y_{d-1}\|}{\|x_{d-1}\|}\Lambda_{d-1,x}^T \Lambda_{d-1,x}x + \epsilon\left(\frac{1 + 4\epsilon d}{2^{d-1}}\right)\|y\|$$

where the first equality is from (32), and the second uses the definition of $Q_{x,y}$, (20b) and (22). The rest of the proof follows as in the proof of Equation (7) in Lemma 8 in [18]. $\square$

**Proof of Proposition C.3**

**Lemma C.5.** *Suppose $G$ satisfies the R2WDC with constant $\epsilon$. Then for any $x, y \in \mathbb{R}^k \setminus \{0\}$ and $i \in [d]$, it holds that*

$$\|G_i(x) - G_i(y)\| \leq \left(\sqrt{\frac{1}{2} + \epsilon} + \sqrt{2(2\epsilon + \theta_{i-1})}\right)\|G_{i-1}(x) - G_{j-1}(y)\|$$

*where $\theta_{i-1} = \angle(G_i(x), G_i(y))$.*

*Proof of Lemma C.5.* We have

$$\|G_j(x) - G_j(y)\| \le \|(W_j)_{+,x_{j-1}}(x_{j-1} - y_{j-1})\| + \|(W_{j,+,x} - W_{j,+,y})y_{j-1}\|. \quad (33)$$

We begin analyzing the first term, noticing that by the R2WDC

$$\|W_{j,+,x}(x_{j-1} - y_{j-1})\|^2 = (x_{j-1} - y_{j-1})^T (W_{j,+,x}^T W_{j,+,x} - \frac{1}{2}I_{n_1})(x_{j-1} - y_{j-1}) + \frac{1}{2}\|x_{j-1} - y_{j-1}\|^2$$

$$\le \left(\frac{1}{2} + \epsilon\right)\|x_{j-1} - y_{j-1}\|^2 \quad (34)$$

We next analyze the second term. Let $W_{j,i} \in \mathbb{R}^{1 \times n_{j-1}}$ be the $i$-th row of $W_j$ then

$$\|(W_{j,+,x} - W_{j,+,y})y_{j-1}\|^2 = \sum_{i=1}^{n} \left(\mathbb{1}_{W_{j,i}x_{j-1}>0} - \mathbb{1}_{W_{j,i}y_{j-1}>0}\right)^2 (W_{j,i}\, y_{j-1})^2$$

$$\le \sum_{i=1}^{n} \left(\mathbb{1}_{W_{j,i}x>0} - \mathbb{1}_{W_{j,i}y>0}\right)(W_{j,i}(x_{j-1} - y_{j-1}))$$

$$= \sum_{i=1}^{n} \mathbb{1}_{W_{j,i}x>0}\mathbb{1}_{W_{j,i}y\le 0}W_{j,i}(x_{j-1} - y_{j-1})$$

$$+ \sum_{i=1}^{n} \mathbb{1}_{W_{j,i}x\le 0}\mathbb{1}_{W_{j,i}y>0}W_{j,i}(x_{j-1} - y_{j-1})$$

$$= (x_{j-1} - y_{j-1})^T (W_j)_{+,x_{j-1}}^T \left((W_j)_{+,x_{j-1}} - (W_j)_{+,y_{j-1}}\right)(x_{j-1} - y_{j-1})$$

$$+ (x_{j-1} - y_{j-1})^T (W_j)_{+,y_{j-1}}^T \left((W_j)_{+,y_{j-1}} - (W_j)_{+,x_{j-1}}\right)(x_{j-1} - y_{j-1}). \quad (35)$$

Observe now that by the R2WDC we have

$$|(x_{j-1} - y_{j-1})^T (W_j)_{+,y_{j-1}}^T \left((W_j)_{+,y_{j-1}} - (W_j)_{+,x_{j-1}}\right)(x_{j-1} - y_{j-1})|$$

$$\le |(x_{j-1} - y_{j-1})^T \left((W_j)_{+,y_{j-1}}^T (W_j)_{+,y_{j-1}} - I_k/2\right)(x_{j-1} - y_{j-1})|$$

$$+ |(x_{j-1} - y_{j-1})^T \left((W_j)_{+,y_{j-1}}^T (W_j)_{+,x_{j-1}} - Q_{x_{j-1},y_{j-1}}\right)(x_{j-1} - y_{j-1})|$$

$$+ |(x_{j-1} - y_{j-1})^T \left(I_{n_{i-1}}/2 - Q_{x_{j-1},y_{j-1}}\right)(x_{j-1} - y_{j-1})|$$

$$\le (2\epsilon + \theta_{j-1})\|x_{j-1} - y_{j-1}\|^2,$$

which together with (35) gives

$$\|(W_{j,+,x} - W_{j,+,y})y_{j-1}\|^2 \le 2(2\epsilon + \theta_{j-1})\|x_{j-1} - y_{j-1}\|^2. \quad (36)$$

We conclude using (34) and (36) in (33). $\qquad \square$

## Proof of Proposition C.4

We next prove the convexity-like property in Proposition C.4.

*Proof of Proposition C.4.* We begin observing that by (27) we have $|\theta_i - \bar{\theta}_i| \le 4i\sqrt{\epsilon} \le 4d\sqrt{\epsilon}$. Furthermore, since $x \in \mathcal{B}(y, d\sqrt{\epsilon}\|y\|)$ it follows that

$$\bar{\theta}_i \le \bar{\theta}_0 \le 2d\sqrt{\epsilon}.$$

Thus by the assumption on $\epsilon$, we have

$$\sqrt{2}\sqrt{\theta_i + 2\epsilon} \le \sqrt{2}\sqrt{\bar{\theta}_i + 4d\sqrt{\epsilon} + 2\epsilon} \le \sqrt{2}\sqrt{2d\sqrt{\epsilon} + 4d\sqrt{\epsilon} + 2\epsilon} \le \frac{1}{30\sqrt{2}d} \quad (37)$$

Let now $\Gamma_d := \Lambda_{d,x}^T(\Lambda_{d,x}x - \Lambda_{d,y}y)$. Then notice that

$$\begin{aligned}
\Gamma_d &= \Lambda_{d-1,x}^T W_{d,+,x}^T (W_{d,+,x}\Lambda_{d-1,x}x - W_{d,+,y}\Lambda_{d-1,y}y)\\
&= \Lambda_{d-1,x}^T W_{d,+,x}^T W_{d,+,x}(\Lambda_{d-1,x}x - \Lambda_{d-1,y}y) + \Lambda_{d,x}^T(W_{d,+,x} - W_{d,+,y})\Lambda_{d-1,y}y\\
&= \frac{1}{2}\Gamma_{d-1} + \epsilon\|\Lambda_{d-1,x}\|\|\Lambda_{d-1,x}x - \Lambda_{d-1,y}y\|O_1(1) + \|\Lambda_{d,x}\|\|(W_{d,+,x} - W_{d,+,y})\Lambda_{d-1,y}y\|O_1(1)\\
&= \frac{1}{2}\Gamma_{d-1} + \left(\epsilon + \sqrt{\frac{1}{2}+\epsilon}\sqrt{2(2\epsilon+\theta_{d-1})}\right)\|\Lambda_{d-1}\|\|\Lambda_{d-1,x}x - \Lambda_{d-1,y}y\|O_1(1)\\
&= \frac{1}{2}\Gamma_{d-1} + \left(\epsilon + \sqrt{\frac{1}{2}+\epsilon}\sqrt{2(2\epsilon+\theta_{d-1})}\right)\frac{1.2\sqrt{1+4\epsilon d}}{2^{d-1}}\|x-y\|O_1(1)\\
&= \frac{1}{2}\Gamma_{d-1} + 2\left(\frac{1}{200^4 d^6} + \frac{1}{30\sqrt{2}d}\right)\frac{\|x-y\|}{2^{d-1}}O_1(1) \qquad (38)
\end{aligned}$$

where the third equality follows from the R2WDC, the fourth the R2WDC and (36), the fifth from (22) and Proposition C.3, and sixth from (37) and the assumption on $\epsilon$. Finally, from (38) and (28) we obtain

$$\Gamma_d = \frac{1}{2^d}\|x-y\| + \frac{1}{16}\frac{\|x-y\|}{2^d}O_1(1)$$

$\square$

# D    Proof of Lemma 5.2

In this section, we prove that a generative network $G$ with random weights satisfies the R2WDC with high-probability (Lemma 5.2). Our proof is inspired by the proof of Proposition 3 in [17].

Notice that because of the piecewise-linear nature of the ReLU activation function, the output of a ReLU network is a subset of a union of affine subspaces. The following lemma from [22] provides an upper bound on the number of such subspaces.

**Lemma D.1** (Lemma 7 in [22]). *Consider a generative network $G$ as in (2) and assume that $n_i \geq k$ for $i \in [d]$. Then for $i \in [d]$, $range(G_i)$ is contained in a union of affine subspaces. Precisely,*

$$range(G_i) \subseteq \cup_{j\in[\Psi_i]}S_{i,j} \quad where \quad \Psi_i \leq \prod_{j=1}^{i}\left(\frac{en_j}{k}\right)^k.$$

*Here each $S_{i,j}$ is some $k$-dimensional affine subspace (which depends on $\{W_\ell\}_{\ell\in[i]}$) in $\mathbb{R}^{n_i}$.*

We next give the main result upon which the proof of Proposition 5.2 rests.

**Proposition D.2.** *Fix $0 < \epsilon < 1$ and $\ell < n$. Let $W \in \mathbb{R}^{m\times n}$ have i.i.d. $\mathcal{N}(0,1/m)$. Let $R, S$ be $\ell$-dimensional subspaces of $\mathbb{R}^n$, and $T$ be an $\ell'$-dimensional subspaces of $\mathbb{R}^n$ with $l' \geq l$. Then if $m \geq C_\epsilon\ell'$, we have that*

$$|\langle W_{+,r}^T W_{+,s}u, v\rangle - \langle Q_{r,s}u, v\rangle| \leq \epsilon\|u\|_2\|v\|_2 \quad \forall\, u,v \in T, \forall r \in R, \forall s \in S, \qquad (39)$$

*with probability exceeding*

$$1 - \gamma\left(\frac{e\,m}{\ell}\right)^{2\ell}\exp(-c_\epsilon m)$$

*Furthermore, let $U = \bigcup_{i=1}^{N_1} U_i$, $V = \bigcup_{j=1}^{N_2} V_j$ $V = \bigcup_{j=1}^{N_2} V_j$, $R = \bigcup_{p=1}^{N_3} R_p$, and $S = \bigcup_{q=1}^{N_4} S_q$ be union of subspaces of $\mathbb{R}^n$ of dimension at most $\ell$. Then if $m \geq 2C_\epsilon\ell'$*

$$|\langle W_{+,r}^T W_{+,s}u, v\rangle - \langle Q_{r,s}u, v\rangle| \leq \epsilon\|u\|_2\|v\|_2 \quad \forall\, u \in U, v \in V, \forall r \in R, \forall s \in S, \qquad (40)$$

*with probability exceeding*

$$1 - \gamma N_1 N_2 N_3 N_4\left(\frac{e\,m}{\ell}\right)^{2\ell}\exp(-c_\epsilon m).$$

*Here $c_\epsilon$ depends polynomially on $\epsilon$, $C_\epsilon$ depends polynomially on $\epsilon^{-1}$, and $\gamma$ is a positive universal constant.*

With the above two results, we are in a position to prove Lemma 5.2.

*Proof of Lemma 5.2.* We begin establishing the proposition in the $d = 2$ case.

If $n_1 \geq 2C_\epsilon k$ by the second part of Proposition D.2 with $U, V, R, S = \mathbb{R}^k$, $W_1$ satisfies (9) with probability at least

$$1 - \gamma\Big(\frac{en_1}{k}\Big)^{2k} \exp(-c_\epsilon n_1).$$

We next consider the bound (9) for $j = 2$. Fix $W_1$ and observe that, by Lemma D.1, range$(G_1)$ is contained in the union of at most $\Psi_1$ number of $k$-dimensional affine subspaces of $\mathbb{R}^{n_1}$ and $\{G_1(x_1) - G_1(x_2) : x_1, x_2 \in \mathbb{R}^k\}$ is contained in the union of at most $\Psi_1^2$ number of $2k$-dimensional affine subspaces of $\mathbb{R}^{n_1}$. Since then an $\ell$-dimensional affine subspace is also contained in an $\ell + 1$ subspace. We have that range$(G_1) \subset \mathcal{R}_1$ where $\mathcal{R}_1$ is the union of at most $\Psi_1$ number of $k + 1$-dimensional subspaces and $\{G_1(x_1) - G_1(x_2) : x_1, x_2 \in \mathbb{R}^k\} \subset \mathcal{U}_1$ where $\mathcal{U}_1$ is the union of at most $\Psi_1^2$ number of $2k + 1$-dimensional subspaces.

By applying the second part of Proposition D.2 to the sets $\mathcal{U}_1, \mathcal{U}_1, \mathcal{R}_1$ and $\mathcal{R}_1$, we have that for fixed $W_1$,

$$\Big|\Big\langle \Big((W_2)^T_{+,G_1(x)}(W_2)_{+,G_1(y)} - Q_{G_1(x),G_1(y)}\Big)(G_1(x_1) - G_1(x_2)), G_1(x_3) - G_1(x_4)\Big\rangle\Big|$$
$$\leq \epsilon\|G_1(x_1) - G_1(x_2)\|_2\|G_1(x_3) - G_1(x_4)\|_2 \quad (41)$$

with probability at least

$$1 - \gamma\Psi_1^6\Big(\frac{e\, n_2}{k + 1}\Big)^{2k+2} e^{-c_\epsilon n_2} \geq 1 - \gamma\Big(\frac{e\, n_2}{k + 1}\Big)^{4k} e^{-c_\epsilon n_2/2}$$

provided that $n_2 \geq 12c_\epsilon^{-1} \log \Psi_1$ and $n_2 \geq 2C_\epsilon(2k + 1)$. In particular the above holds provided that $n_2 \geq \widetilde{C}_\epsilon k \log(en_1/k)$ where $\widetilde{C}_\epsilon$ depends polynomially on $\epsilon^{-1}$.

Integrating over the probability space of $W_1$, independence of $W_2$ and $W_1$ implies that (41) holds for random $W_1$ with the same probability bound. This allows us to conclude that a two-layer random generative network $G$ satisfies the R2WDC with probability at least

$$1 - \gamma\Big(\frac{en_1}{k}\Big)^{2k} e^{-c_\epsilon n_1} - \gamma\Big(\frac{e\, n_2}{k + 1}\Big)^{4k} e^{-c_\epsilon n_2/2}.$$

The proof of the $d \geq 2$ case follows similarly. In particular, to establish (9) for $W_i$ notice that range$(G_{i-1})$ is contained in the union of at most $\Psi_{i-1}$ number $k + 1$ subspaces, and $\{G_{i-1}(x_1) - G_{i-1}(x_2) : x_1, x_2 \in \mathbb{R}^k\}$ in the union of at most $\Psi_{i-1}^2$ number of $2k + 1$-dimensional subspaces. Applying Proposition D.2 to these subspaces we have that for fixed $\{W_j\}_{j \in [i-1]}$

$$\Big|\Big\langle \Big((W_i)^T_{+,G_{i-1}(x)}(W_i)_{+,G_{i-1}(y)} - Q_{G_{i-1}(x),G_{i-1}(y)}\Big)(G_{i-1}(x_1) - G_{i-1}(x_2)), G_{i-1}(x_3) - G_{i-1}(x_4)\Big\rangle\Big|$$
$$\leq \epsilon\|G_{i-1}(x_1) - G_{i-1}(x_2)\|_2\|G_{i-1}(x_3) - G_{i-1}(x_4)\|_2 \quad (42)$$

with probability at least

$$1 - \gamma\Big(\frac{e\, n_i}{k + 1}\Big)^{4k} e^{-c_\epsilon n_i/2}$$

provided that

$$n_i \geq \widetilde{C}_\epsilon \cdot k \cdot \prod_{j=1}^{i-1} \frac{e\, n_j}{k}.$$

Integrating over the probability space of $\{W_j\}_{j \in [i-1]}$) indpendence of $W_i$ and $(W_1, \ldots, W_{i-1})$ gives that (42) holds with the same probability bound. $\square$

We will devote the following section to the proof of Proposition D.2.

### D.1 Proof of Proposition D.2

We begin by proving a weaker form of Proposition D.2, that characterizes the concentration of $W_{+,r}^T W_{+,s}$ around its mean for fixed $r, s$ and when acting on $\ell$-dimensional subspaces.

**Lemma D.3.** *Fix $0 < \epsilon < 1$ and $k < n$. Let $W \in \mathbb{R}^{m \times n}$ have i.i.d. $\mathcal{N}(0, 1/m)$ entries and fix $r, s \in \mathbb{R}^n$. Let $T$ be a $\ell$-dimensional subspace of $\mathbb{R}^n$. Then if $m \geq \tilde{C}_1 \ell$, we have that with probability exceeding $1 - 2\exp(-\tilde{c}_1 m)$,*

$$|\langle W_{+,r}^T W_{+,s} u, u \rangle - \langle Q_{r,s} u, u \rangle| \leq \epsilon \|u\|_2^2 \quad \forall\, u \in T \tag{43}$$

*and*

$$|\langle W_{+,r}^T W_{+,s} u, v \rangle - \langle Q_{r,s} u, v \rangle| \leq 3\epsilon \|u\|_2 \|v\|_2 \quad \forall\, u, v \in T, \tag{44}$$

*Furthermore, let $U = \bigcup_{i=1}^{N_1} U_i$ and $V = \bigcup_{j=1}^{N_2} V_j$ where $U_i$ and $V_j$ are subspaces of $\mathbb{R}^n$ of dimension at most $\ell$ for all $i \in [N_1]$ and $j \in [N_2]$. Then if $m \geq 2\tilde{C}_1 \ell$*

$$|\langle W_{+,r}^T W_{+,s} u, v \rangle - \langle Q_{r,s} u, v \rangle| \leq 3\epsilon \|u\|_2 \|v\|_2 \quad \forall\, u \in U, \forall\, v \in V, \tag{45}$$

*with probability exceeding $1 - 2N_1 N_2 \exp(-\tilde{c}_1 m)$. Here $\tilde{c}_1$ depends polynomially on $\epsilon$ and $\tilde{C}_1 = \Omega(\epsilon^{-1} \log \epsilon^{-1})$.*

*Proof.* The proof follows the one in Proposition 4 of [17] with minor variations. Set $\Sigma_{r,s} := W_{+,r}^T W_{+,s} - Q_{r,s}$, and notice that for fixed $u \in \mathbb{R}^{n-1}$, $\langle \Sigma_{r,s} u, u \rangle = \sum_{i=1}^m Y_i$ where $Y_i = X_i - \mathbb{E}[X_i]$, $X_i = \mathbb{1}_{\langle w_i, r \rangle > 0} \mathbb{1}_{\langle w_i, s \rangle > 0} \langle w_i, u \rangle^2$ and each $w_i \sim \mathcal{N}(0, I_n/m)$. We then notice that the $Y_i$ are sub-exponential random variables and by standard $\epsilon$-net argument we can show that (43) holds with high-probability. Proposition 5 in [17] can then be adapted to this case as well and used to derive (44) from (43). Finally (45) follows by a union bound over all subspaces of the form $\text{span}(U_i, V_j)$. $\square$

We next observe that the rows of a sufficiently tall random matrix $W$ tessellate the unit sphere in regions of small diameter.

**Lemma D.4.** *Fix $0 < \epsilon < 1$. Let $W \in \mathbb{R}^{m \times n}$ have i.i.d. $\mathcal{N}(0, 1/m)$ entries with rows $\{w_\ell\}_{\ell=1}^m$. Let $Z$ be a $\ell$-dimensional subspace of $\mathbb{R}^n$. Define $E_{Z,W}$ to be the event that there exists a set $Z_0 \subset Z$ with the following properties:*

*i) each $z_0 \in Z_0$ satisfies $\langle w_\ell, z_0 \rangle \neq 0$ for all $\ell \in [m]$,*

*ii) $|Z_0| \leq (\frac{e\,m}{\ell})^\ell$, and*

*iii) for all $z \in Z$ such that $\|z\|_2 = 1$, there exists $z_0 \in Z_0$ such that $\|z - z_0\|_2 \leq \epsilon$.*

*If $m \geq \tilde{C}_2 \ell$, then $\mathbb{P}(E_{Z,W}) \geq 1 - C_2 \exp(-c_2 \epsilon m)$. Here $C_2$ and $c_2$ are positive absolute constants and $\tilde{C}_2$ depends polynomially on $\epsilon^{-1}$.*

*Proof.* The proof of this lemma follows the one in Lemma 24 in [17]. The upper bound $|Z_0| \leq (\frac{e\,m}{\ell})^\ell$ is due to Lemma D.6 in Appendix D.2. $\square$

We are now ready to present the proof of Proposition D.2.

*Proof of Proposition D.2.* Let $E_{R,W}$ be the event defined in Lemma D.4 corresponding to the matrix $W$ and subspace $R$. On the event $E_{R,W}$ there exists a finite set $R_0 \subset R$ satisfying properties i) - iii) of Lemma D.4. Similarly, we can define the event $E_{S,W}$ for the matrix $W$ and subspace $S$, and the finite set $S_0 \subset S$ satisfying properties i) - iii).

We can then define the event $E_{R,S} := E_{R,W} \cap E_{S,W}$ so that if $m \geq \tilde{C}_2 \ell'$ by Lemma D.4 we have

$$\mathbb{P}(E_{R,S}) \geq 1 - 2C_2 \exp(-c_2 \epsilon m).$$

For fixed $r_0 \in R_0$ and $s_0 \in S_0$, Lemma D.3 gives that if $m \geq 2\tilde{C}_1 \ell$ with probability at least $1 - 2\exp(-\tilde{c}_1 m)$

$$|\langle W_{+,r_0}^T W_{+,s_0} u, v \rangle - \langle Q_{r_0, v_0} u, v \rangle| \leq 3\epsilon \|u\|_2 \|v\|_2 \quad \forall u, v \in T.$$

Next, let $E_0$ be the event that

$$|\langle W_{+,r_0}^T W_{+,s_0} u, v\rangle - \langle Q_{r_0,v_0} u, v\rangle| \leq 3\epsilon \|u\|_2 \|v\|_2 \quad \forall u, v \in T, \ r_0 \in R_0, \ s_0 \in S_0.$$

Then, on $E_{R,S}$, a union bound gives

$$\mathbb{P}(E_0) \geq 1 - 2|R_0||S_0| \exp(-\tilde{c}_1 m/2) \geq 1 - 2\left(\frac{e\,m}{\ell}\right)^{2\ell} \exp(-\tilde{c}_1 m/2).$$

We will next work on the event $E_0 \cap E_{R,S}$. Fix nonzero $r \in R$ and $s \in S$, and define the set of indices

$$\Omega_{r,s} := \{j \in [m] : \langle w_j, r\rangle = 0 \text{ or } \langle w_j, s\rangle = 0\}$$

Observe then that by the definition of $W_{+,r}$ and $\Omega_{r,s}$ the following holds

$$
\begin{aligned}
W_{+,r}^T W_{+,s} &= \sum_{j=1}^m \mathbb{1}_{\langle w_j, r\rangle > 0} \mathbb{1}_{\langle w_j, s\rangle > 0} w_j w_j^T \\
&= \sum_{j \in \Omega_{r,s}} \mathbb{1}_{\langle w_j, r\rangle > 0} \mathbb{1}_{\langle w_j, s\rangle > 0} w_j w_j^T + \sum_{j \in \Omega_{r,s}^c} \mathbb{1}_{\langle w_j, r\rangle > 0} \mathbb{1}_{\langle w_j, s\rangle > 0} w_j w_j^T \\
&= \sum_{j \in \Omega_{r,s}^c} \mathbb{1}_{\langle w_j, r\rangle > 0} \mathbb{1}_{\langle w_j, s\rangle > 0} w_j w_j^T
\end{aligned}
$$

On the event $E_{R,S}$, there exist therefore $r_0 \in R_0$ and $s_0 \in S_0$ such that for all $j \in \Omega_{r,s}^c$ it holds that

$$\mathrm{sgn}(\langle w_j, r\rangle) = \mathrm{sgn}(\langle w_j, r_0\rangle) \quad \text{and} \quad \mathrm{sgn}(\langle w_j, s\rangle) = \mathrm{sgn}(\langle w_j, s_0\rangle).$$

In particular, we can write

$$
\begin{aligned}
W_{+,r}^T W_{+,s} &= \sum_{j \in \Omega_{r,s}^c} \mathbb{1}_{\langle w_j, r\rangle > 0} \mathbb{1}_{\langle w_j, s\rangle > 0} w_j w_j^T \\
&= W_{+,r_0}^T W_{+,s_0} - \sum_{j \in \Omega_{r,s}} \mathbb{1}_{\langle w_j, r_0\rangle > 0} \mathbb{1}_{\langle w_j, s_0\rangle > 0} w_j w_j^T \\
&=: W_{+,r_0}^T W_{+,s_0} - \widetilde{W}_{+,r_0}^T \widetilde{W}_{+,s_0}
\end{aligned}
$$

The next lemma shows that the residual $\widetilde{W}_{+,r_0}^T \widetilde{W}_{+,s_0}$ has small norm when acting on $T$.

**Lemma D.5.** *Fix $0 < \epsilon < 1$ and $\ell < m$. Suppose that $W \in \mathbb{R}^{m \times n}$ has i.i.d. $\mathcal{N}(0, 1/m)$ entries. Let $T \subset \mathbb{R}^n$ be an $\ell$-dimensional subspace and $R_0$ and $S_0$ be subsets of $\mathbb{R}^n$. Let $E_1$ be the event the following inequality holds for all set of indexes $\Omega \subset [m]$ with cardinality $|\Omega| \leq 2\ell$:*

$$|\langle \widetilde{W}_{+,r_0}^T \widetilde{W}_{+,s_0} u, v\rangle| \leq \epsilon \|u\|_2 \|v\|_2 \quad \forall u, v \in T, \ r_0 \in R_0, \ s_0 \in S_0$$

*where*

$$\widetilde{W}_{+,r_0}^T \widetilde{W}_{+,s_0} := \sum_{j \in \Omega} \mathbb{1}_{\langle w_j, r_0\rangle > 0} \mathbb{1}_{\langle w_j, s_0\rangle > 0} w_j w_j^T.$$

*There exists a $\delta_\epsilon > 0$ such that if $m \geq 9\epsilon^{-1}\ell$ an $2\ell \leq \delta_\epsilon m$, then $\mathbb{P}(E_1) \geq 1 - 2m \exp(-\epsilon m/36)$.*

We now consider the event $E := E_1 \cap E_0 \cap E_{R,S}$ where $E_1$ is the event defined in the previous lemma. On $E$ for all $r \in R$, $s \in S$ and $u, v \in T$,

$$
\begin{aligned}
|\langle W_{+,r}^T W_{+,s} u, v\rangle - \langle Q_{r,s} u, v\rangle| &= \left| \langle W_{+,r_0}^T W_{+,s_0} u, v\rangle - \langle \widetilde{W}_{+,r_0}^T \widetilde{W}_{+,s_0} u, v\rangle - \langle Q_{r,s} u, v\rangle \right| \\
&\leq \left| \langle W_{+,r_0}^T W_{+,s_0} u, v\rangle - \langle Q_{r_0,s_0} u, v\rangle \right| \\
&\quad + \left| \langle Q_{r_0,s_0} u, v\rangle - \langle Q_{r_0,s_0} u, v\rangle \right| \\
&\quad + \left| \langle \widetilde{W}_{+,r_0}^T \widetilde{W}_{+,s_0} u, v\rangle \right| \\
&\leq 3\epsilon \|u\|_2 \|v\|_2 + \frac{60}{\pi}\epsilon \|u\|_2 \|v\|_2 + \epsilon \|u\|_2 \|v\|_2 \\
&\leq 24\epsilon \|u\|_2 \|v\|_2, \tag{46}
\end{aligned}
$$

where the first equality used the event $E_{R,S}$ and the definition of $\widetilde{W}_{+,r_0}^T \widetilde{W}_{+,s_0}$. The second inequality used instead the event $E_1 \cap E_0$ and the Lipschitz continuity of $Q_{r,s}$ (Lemma D.7).

In conclusion, there exist $C_\epsilon$ and $c_\epsilon$ such that if $m \geq C_\epsilon \ell'$ then

$$\mathbb{P}(E_1 \cap E_0 \cap E_{R,S}) \geq 1 - 2m \exp(-\epsilon m/36) - 2\Big(\frac{e\,m}{\ell}\Big)^{2\ell} \exp(-\tilde{c}_1 m/2) - 2C_2 \exp(-c_2 \epsilon m)$$

$$\geq 1 - \gamma \Big(\frac{e\,m}{\ell}\Big)^{2\ell} \exp(-c_\epsilon m)$$

Here $C_\epsilon$ depends polynomially on $\epsilon^{-1}$ and $c_\epsilon$ depends polynomially on $\epsilon$, and $\gamma$ is positive absolute constant.

Notice that (46) gives a bound in terms of $24\epsilon\|u\|_2\|v\|_2$. To obtain a bound as in (39) simply rescale $\epsilon$ by $1/24$ in the discussion above, and modify $c_\epsilon$ and $C_\epsilon$ accordingly.

To extend (39) to the union of subspaces, we consider the subspace $T_{i,j} = \text{span}(U_i, V_j)$ with dimension at most $2\ell'$. Then use (39) with subspaces $T_{i,j}$, $R_p$ and $S_q$, and take a union bound. $\quad\square$

## D.2 Supplemental Results for Section D

We begin this section by providing an upper bound on the number of activation patterns of a ReLU layer. This result is used in the proof of Lemma D.4.

**Lemma D.6.** *Let $S$ be an $\ell$-dimensional subspace of $\mathbb{R}^n$ and $m \geq \ell$. Let $W \in \mathbb{R}^{m \times n}$ have i.i.d $\mathcal{N}(0, 1/m)$ entries. Then with probability 1,*

$$|\{diag(Ws > 0)W \mid s \in S\}| \leq \Big(\frac{em}{\ell}\Big)^\ell$$

*Proof.* Observe that by rotational invariance of the Gaussian distribution we may take, without loss of generality, $S$ to be the span of the first $\ell$ standard basis vector, i.e. $S = \text{span}(e_1, \ldots, e_\ell)$. We can then also take $W \in \mathbb{R}^{m \times \ell}$ and $S = \mathbb{R}^\ell$.

Let $\{w_j\}_{j=1}^m$ be the rows of the matrix $W$. Notice that for fixed $W$, $|\{\text{diag}(Ws > 0)W \mid s \in S\}|$ equals the number of binary vectors $(\mathbb{1}_{\langle w_j, v\rangle > 0})_{j \in [n]}$ for $v \in \mathcal{S}^{\ell-1}$. Each $(\mathbb{1}_{\langle w_j, v\rangle > 0})_{j \in [n]}$ uniquely identifies a region of the partitioning of $\mathbb{R}^l$ induced by the set of hyperplanes $\mathcal{H} := \{x : \langle w_j, x\rangle = 0\}$. From the theory of hyperplane arrangements [29] we know that $m \geq \ell$ hyperplanes in $\mathbb{R}^\ell$ partition the space in at most $\sum_{j=0}^\ell \binom{m}{j}$. Thus, with probability 1 we have

$$|\{\text{diag}(Ws > 0)W \mid s \in S\}| \leq \sum_{j=0}^\ell \binom{m}{j}$$

$$\leq \sum_{j=0}^\ell \frac{m^j}{j!} \leq \sum_{j=0}^\ell \frac{\ell^j}{j!}\Big(\frac{m}{\ell}\Big)^j \leq \Big(\frac{m}{\ell}\Big)^\ell \sum_{j=0}^\infty \frac{\ell^j}{j!} = \Big(\frac{em}{\ell}\Big)^\ell$$

$\square$

Next we prove Lemma D.5, providing an upper bound for the random matrix $\widetilde{W}^T \widetilde{W}$ when acting on low-dimensional subspaces.

*Proof of Lemma D.5.* Notice that for any $\Omega \subset [m]$, $u, v \in T$, $r_0 \in R_0$ and $s_0 \in S_0$, it holds that

$$|\langle \widetilde{W}_{+,r_0}^T \widetilde{W}_{+,s_0} u, v\rangle| = |\langle \text{diag}(W_\Omega r_0 > 0) \odot \text{diag}(W_\Omega s_0 > 0) W_\Omega u, W_\Omega v\rangle|$$

$$\leq \|\text{diag}(W_\Omega r_0 > 0) \odot \text{diag}(W_\Omega s_0 > 0)\|\|W_\Omega v\|\|W_\Omega u\|$$

$$\leq \|W_\Omega v\|\|W_\Omega u\|.$$

Therefore, it is sufficient to show that

$$\|W_\Omega u\| \leq \sqrt{\epsilon}\|u\| \qquad \forall u \in T \;\; \forall \Omega \subset [m] \;\; \text{satisfying} \;\; |\Omega| \leq 2\ell \leq \delta_\epsilon m.$$

The rest of the proof follows, *mutatis mutandis*, as in Lemma 26 of [17]. $\quad\square$

We will next show that $Q_{x,y}$ is a Lipschitz function of its arguments.

**Lemma D.7.** *Fix $0 < \epsilon < 1$ and $x, \tilde{x}, y, \tilde{y} \in \mathcal{S}^{n-1}$. If $\|\tilde{x} - x\| \le \epsilon$ and $\|\tilde{y} - y\| \le \epsilon$, then*

$$\|Q_{\tilde{x},\tilde{y}} - Q_{\tilde{x},\tilde{y}}\| \le \left(\frac{2}{\pi} + 2\sqrt{79}\right)\epsilon$$

*Proof.* Recall the following facts:

$$\|x - y\| \ge 2\sin(\angle(x,y)/2), \qquad\qquad \forall x, y \in \mathcal{S}^{n-1} \qquad (47)$$
$$|\angle(x_1, x_2)| \ge |\angle(x_1, y) - \angle(x_2, y)|, \qquad\qquad \forall x_1, x_2, y \in \mathcal{S}^{n-1} \qquad (48)$$
$$\sin(\theta/2) \ge \theta/4, \qquad\qquad \forall \theta \in [0, \pi] \qquad (49)$$

Let $\theta_{\tilde{x},x} = \angle(\tilde{x}, x)$ and $\theta_{\tilde{y},y} = \angle(\tilde{y}, y)$, then

$$\|Q_{x,y} - Q_{\tilde{x},\tilde{y}}\| \le \frac{|\theta_{x,y} - \theta_{\tilde{x},\tilde{y}}|}{2\pi} + \left\|\frac{\sin\theta_{x,y}}{2\pi}M_{x\leftrightarrow y} - \frac{\sin\theta_{\tilde{x},\tilde{y}}}{2\pi}M_{\tilde{x}\leftrightarrow\tilde{y}}\right\|.$$

By (48) it holds that

$$|\theta_{x,y} - \theta_{\tilde{x},\tilde{y}}| \le |\theta_{x,y} - \theta_{\tilde{x},y}| + |\theta_{\tilde{x},y} - \theta_{\tilde{x},\tilde{y}}| \le |\theta_{\tilde{x},x}| + |\theta_{\tilde{y},y}|,$$

while from (47) and (49) it follows that

$$|\theta_{\tilde{x},x}| \le 4\sin(\theta_{\tilde{x},x}/2) \le 2\epsilon,$$
$$|\theta_{\tilde{y},y}| \le 4\sin(\theta_{\tilde{y},y}/2) \le 2\epsilon.$$

Thus $|\theta_{x,y} - \theta_{\tilde{x},\tilde{y}}| \le 4\epsilon$. Lemma B.3 in [10] then proves that

$$\left\|\frac{\sin\theta_{x,y}}{2\pi}M_{x\leftrightarrow y} - \frac{\sin\theta_{\tilde{x},\tilde{y}}}{2\pi}M_{\tilde{x}\leftrightarrow\tilde{y}}\right\| \le 2\sqrt{79}\epsilon,$$

which concludes the proof. $\qquad\square$

# E    Proof of Lemma 5.3

In this section we prove Lemma 5.3 which is used to bound the perturbation of the objective function $f_{\mathrm{cs}}$ and its gradient due to the presence of the noise term $\eta$.

*Proof of Lemma 5.3.* Fix $x, z \in \mathcal{S}^{k-1}$ and notice that by the properties of the Gaussian distribution, for $t \ge 0$ it holds that

$$\mathbb{P}_A\left[\langle z, \Lambda_x^T A^T \eta\rangle \ge \frac{\|\Lambda_x z\|}{\sqrt{m}}\|\eta\|t\right] = \mathbb{P}_{y\sim\mathcal{N}(0,1)}\left[\frac{\|\Lambda_x z\|}{\sqrt{m}}\|\eta\|y \ge \frac{\|\Lambda_x z\|}{\sqrt{m}}\|\eta\|t\right] \le e^{-\frac{t^2}{2}}.$$

If $z = x$ use (20b), while if $z \ne x$ and $G$ differentiable at $x$ use (22), to obtain that

$$\mathbb{P}_A\left[\langle z, \Lambda_x^T A^T \eta\rangle \ge \sqrt{\frac{13}{12}}\frac{\|\eta\|}{2^{d/2}}\frac{t}{\sqrt{m}}\right] \le e^{-\frac{t^2}{2}}$$

Let $\mathcal{N}_{1/2}$ be a $\frac{1}{2}$-net over $\mathcal{S}^{k-1}$ such that $|\mathcal{N}_{1/2}| \le 5^k$ (see for example [38]). Recall that by Lemma D.1 the number of different matrices $\Lambda_x$ is bounded by $\Psi_d$. Thus, a union bound gives

$$\mathbb{P}\left[\langle z, \Lambda_x^T A^T \eta\rangle \ge \sqrt{\frac{13}{12}}\frac{\|\eta\|}{2^{d/2}}\frac{t}{\sqrt{m}}, \;\; \forall x, z \in \mathcal{S}^{k-1}\right] \le |\mathcal{N}_{\frac{1}{2}}|\,\Psi_d\,\mathbb{P}\left[\langle z, \Lambda_x^T A^T \eta\rangle \ge \sqrt{\frac{13}{12}}\frac{\|\eta\|}{2^{d/2}}\frac{t}{\sqrt{m}}\right]$$

$$\le \exp(-\frac{t^2}{2} + \log 5 + \log \Psi_d)$$

Choosing $t = 2\sqrt{k\log(5\prod_{i=1}^d \frac{e\,n_i}{k})}$ we obtain the theses.

$\qquad\square$

# F  Extensions

## F.1  Compressive Phase Retrieval with a Generative Prior

Consider a generative network $G : \mathbb{R}^k \to \mathbb{R}^n$ as in (2). The compressive phase retrieval problem with a generative network prior can be formulated as follows.

---

**COMPRESSIVE PHASE RETRIEVAL WITH A DEEP GENERATIVE PRIOR**

**Let**:  $G : \mathbb{R}^k \to \mathbb{R}^n$ generative network, $A \in \mathbb{R}^{m \times n}$ measurement matrix.
**Let**:  $y_\star = G(x_\star)$ for some unknown $x_\star \in \mathbb{R}^k$.

**Given**:  $G$ and $A$.
**Given**:  Measurements $b = |A y_\star| + \eta \in \mathbb{R}^m$ with $m \ll n$ and $\eta \in \mathbb{R}^m$ noise.

**Estimate**:  $y_\star$.

---

To estimate $y_\star$, [16] proposes to find the latent code $\hat{x}$ that minimizes the reconstruction error

$$\tilde{x} = \arg \min_{x \in \mathbb{R}^x} f_{\mathrm{pr}}(x) := \frac{1}{2} \| b - |AG(x)| \|_2^2, \tag{50}$$

$$y_\star \approx G(\tilde{x}).$$

In [17] it is shown that Algorithm 1 with inputs $f_{\mathrm{pr}}$, small enough step size and arbitrary initial condition estimates $y_\star$ up to the noise level in polynomial time, provided that the number of phaseless measurements is up-to log-factors $m \geq k \cdot \mathrm{poly}(d)$ and the generative network is logarithmically expansive. The proof uses the WDC and an isometry condition akin to the RRIC. As before, the RWDC can be replaced by the R2WDC and obtain the same convergence guarantees. Moreover, as in the case of compressed sensing, the logarithmic factor in the number of measurements can be improved using Lemma D.1.

## F.2  Denoising with a Generative Prior

Consider a generative network $G : \mathbb{R}^k \to \mathbb{R}^n$ as in (2). The denoising problem with a generative network prior can be formulated as follows.

---

**DENOISING WITH A DEEP GENERATIVE PRIOR**

**Let**:  $G : \mathbb{R}^k \to \mathbb{R}^n$ generative network.
**Let**:  $y_\star = G(x_\star)$ for some unknown $x_\star \in \mathbb{R}^k$.

**Given**:  $G$.
**Given**:  Noisy signal $b = y_\star + \eta \in \mathbb{R}^m$ with $\eta \sim \mathcal{N}(0, \sigma^2 \, \mathrm{I}_n)$ noise.

**Estimate**:  $y_\star$.

---

To estimate $y_\star$, [20] proposes to find the latent code $\hat{x}$ that minimizes the reconstruction error

$$\tilde{x} = \arg \min_{x \in \mathbb{R}^x} f_{\mathrm{den}}(x) := \frac{1}{2} \| b - G(x) \|_2^2, \tag{51}$$

$$y_\star \approx G(\tilde{x}).$$

In [20] recovery guarantees based on this minimization problem are given for an expansive generative network $G$. Specifically, it is shown that Algorithm 1 with input $f_{\mathrm{den}}$, small enough step size $\alpha$ and arbitrary initial point $x_0$, reconstructs the signal $y_\star$ up to an $O(k/n)$ error. The random network $G$ is assumed to be logarithmically expansive in order to satisfy the WDC with high-probability, but inspecting the proof it can be seen that the R2WDC is enough. Using Lemma 5.2 we can extend the result of [20] to the case of a generative network satisfying Assumptions B.

## F.3 Spiked Matrix Recovery with a Generative Prior

Consider a generative network $G : \mathbb{R}^k \to \mathbb{R}^n$ as in (2). The spiked Wishart matrix recovery with a generative prior is formulated as follows.

---

**SPIKED WISHART MATRIX RECOVERY
WITH A DEEP GENERATIVE PRIOR**

**Let**:   $G : \mathbb{R}^k \to \mathbb{R}^n$ generative network.

**Let**:   $y_\star = G(x_\star)$ for some unknown $x_\star \in \mathbb{R}^k$.

**Given**:   $G$.

**Given**:   Noisy matrix $B = u\, y_\star^T + \sigma \mathcal{Z} \in \mathbb{R}^{N \times n}$, with $u \sim \mathcal{N}(0, I_N)$ and $\mathcal{Z}$ with i.i.d. $\mathcal{N}(0, 1)$ entries.

**Estimate**:   $y_\star$.

---

Similarly, the spiked Wigner matrix recovery with a generative prior is formulated as follows.

---

**SPIKED WIGNER MATRIX RECOVERY
WITH A DEEP GENERATIVE PRIOR**

**Let**:   $G : \mathbb{R}^k \to \mathbb{R}^n$ generative network.

**Let**:   $y_\star = G(x_\star)$ for some unknown $x_\star \in \mathbb{R}^k$.

**Given**:   $G$.

**Given**:   Noisy matrix $B = y_\star\, y_\star^T + \sigma \mathcal{H} \in \mathbb{R}^{n \times n}$, with $\mathcal{H}$ from a Gaussian Orthogonal Ensemble.

**Estimate**:   $y_\star$.

---

To estimate $y_\star$, [8] proposes to find the latent code $\hat{x}$ that minimizes the reconstruction error

$$\tilde{x} = \arg\min_{x \in \mathbb{R}^x} f_{\text{spiked}}(x) := \frac{1}{2}\|M - G(x)G(x)^T\|_F^2,$$

$$y_\star \approx G(\tilde{x}),$$

where

- in the spiked Wishart model $M = B^T B / N - \sigma^2 I_n$;
- in the spiked Wigner model $M = B$.

As shown in [9] Algorithm 1 with inputs $f_{\text{spiked}}$, appropriate $\alpha$ and arbitrary initial point $x_0$, estimates in polynomial time the signal $y_\star$ with rate-optimal dependence on the noise level or sample complexity. In particular, this shows that the absence of a computational-statistical gap in spiked matrix recovery with an expansive (random) generative network prior. The proof uses the fact that for $G$ satisfying the WDC the bounds in Proposition C.1, C.2, C.3 and C.4 hold. Since these bounds hold under the weaker R2WDC we can directly extend the results in [9] to non-expansive generative networks $G$ satisfying Assumptions B.

# G   An example of a contractive generative network

In this section we give an example of a generative network as in (2) satisfying the conditions (12) and (17), and with contractive layers.

Let $d \geq 2$ and $\bar{C}_\epsilon := \max(\widetilde{C}_\epsilon, 16c_\epsilon^{-1}/\log(2))$. Then consider a $d$-layer generative network $G$ such that for $i \in [d]$

$$n_i := \bar{C}_\epsilon \cdot k \cdot d(2d - i) \cdot \alpha,$$

where $\alpha \cdot \bar{C}_\epsilon \in \mathbb{N}$ and

$$\alpha \geq \max \left\{ \frac{2 \log \left( \bar{C}_\epsilon \cdot k \right)}{d^2}, \log \left( e^2 \bar{C}_\epsilon \right) \right\}. \tag{52}$$

We now demonstrate that $n_i$ satisfies (12). Notice that

$$
\begin{aligned}
\log \Big( \prod_{j=1}^{i-1} \frac{e n_j}{k} \Big) &= \sum_{j=1}^{i-1} \log \left( \alpha \cdot \bar{C}_\epsilon \cdot d(2d - j) \cdot e \right) \\
&\leq (d-1) \log \left( \alpha \cdot \bar{C}_\epsilon \cdot 2d^2 \cdot e \right) \\
&= (d-1) \big[ \log(e \bar{C}_\epsilon) + 2 \log(d) \big] + (d-1) \log(2\alpha) \\
&\leq (d-1) d \big[ \log(e \bar{C}_\epsilon) + 1 \big] + (d-1)\alpha \\
&\leq (d-1) d \alpha + (d-1)\alpha \\
&= (d^2 - 1)\alpha
\end{aligned}
$$

where in the second inequality we have used $2 \log(x) \leq x$ and $\log(2x) \leq x$ for $x > 0$ and in the third (52). Next since $d(2d - i) \geq (d^2 - 1)$ for every $i \in [d]$, $n_i$ satisfies (12) for every $i \in [d]$.

We now show that $n_i$ satisfies (17). We have

$$
\begin{aligned}
\log(n_i) &= \log(\bar{C}_\epsilon \cdot k \cdot d(2d - i) \cdot \alpha) \\
&= \log(d(2d - i)\alpha) + \log(\bar{C}_\epsilon k) \\
&\leq \frac{d(2d - i)\alpha}{2} + \log(\bar{C}_\epsilon k) \\
&\leq \frac{d(2d - i)\alpha}{2} + \frac{d^2 \alpha}{2} \\
&\leq d(2d - i)\alpha,
\end{aligned}
$$

where in the first inequality we have used $2 \log(x) \leq x$ for $x > 0$, in the second inequality (52) and in the third $2d^2 \leq d(2d - i)$ for every $i \in [d]$. We therefore have

$$\log(n_i) \cdot \frac{16 \cdot k \cdot c_\epsilon^{-1}}{\log(2)} \leq d(2d - i) \cdot \alpha \cdot \frac{16 \cdot k \cdot c_\epsilon^{-1}}{\log(2)} \leq d(2d - i) \cdot \alpha \cdot \bar{C}_\epsilon \cdot k = n_i.$$

## H   Experiments

In this section, we verify the prediction of our theory (Theorem 5.4) on synthetic data. We consider the compressed sensing problem with generative prior formulated in Se ction 3 and show that a practical implementation of Algorithm 1 can estimates the sought signal in the case of generative networks with contractive layers.

At each step Algorithm 1 requires the computation of the subgradient of $f_{\text{cs}}$ at a point $\tilde{x}$. This could be achieved for example by computing $\nabla f_{\text{cs}}(\tilde{x} + \delta)$ where $\delta$ is sufficiently small and $f_{\text{cs}}$ is differentiable at $f_{\text{cs}}(\tilde{x} + \delta)$. As already shown in [21], in practice, one can replace Algorithm 1 with Algorithm 2 below. In the Algorithm 2, $v'_{\tilde{x}_t}$ corresponds to the gradient of $f_{\text{cs}}$ at a point of differentiability. At a point of non-differentiability $v'_{\tilde{x}_t}$ corresponds to the output of autograd applied to $f_{\text{cs}}$ as implemented by the current deep learning libraries. Notice that the set of non-differentiability of $f_{\text{cs}}$ has measure zero and it is extremely unlikely that the iterates of the subgradient descent will hit this set. Hence, in practice Algorithm 1 and Algorithm 2 are equivalent.

We then the performance of Algorithm 2 when solving compressed sensing with generative networks $G$ with two and three layers. The entries of the weights $W_i$ of $G$ are drawn from $\mathcal{N}(0, I/n_i)$, while the entries of the measurements matrix $A$ are drawn from $\mathcal{N}(0, I/m)$. In all the experiments we run Algorithm 2 for 10,000 iterations, or until the norm of the $v'_{\tilde{x}_t}$ is below $10^{-10}$, whichever comes first. We fix the learning rate $\nu$ to 0.7 and randomly initialize $x_0$. The target $x_\star$ is chosen so that $y_\star = G(x_\star)$ has unit norm.

In the first experiment, we consider the noiseless case with $\eta = 0$ and a fixed number of measurements $m = 300$. The 2-layers generative networks have fixed layer widths of dimensions $n_1 = 700$ and

---
**Algorithm 2:** PRACTICAL SUBGRADIENT DESCENT [21]
---
**Input:** Objective function $f$, initial point $x_0 \in \mathbb{R}^k \setminus \{0\}$ and step size $\alpha$
**Output:** An estimate of the target signal $y_\star = G(x_\star)$ and latent vector $x_\star$

1 **for** $t = 0, 1, \dots$ **do**
2     **if** $f(-x_t) < f(x_t)$ **then** $\tilde{x}_t \leftarrow -x_t$
3     **else** $\tilde{x}_t \leftarrow x_t$
4     $v'_{\tilde{x}_t} = \Lambda_{d,\tilde{x}_t}^T A^T (A\Lambda_{d,\tilde{x}_t} \tilde{x}_t - b)$
5     $x_{t+1} \leftarrow \tilde{x}_t - \alpha v'_{\tilde{x}_t}$
6 **end**
7 **return** $x_t, G(x_t)$
---

Figure 1: Empirical probability of successful recovery the latent vector from compressed noiseless random measurements versus the latent dimension. The 2-layers generative networks have fixed layer widths of dimensions $n_1 = 700$ and $n = n_2 = 600$ and the 3-layers generative networks have fixed layer widths of dimensions $n_1 = 150$, $n_2 = 700$ and $n = n_3 = 600$.

$n = n_2 = 650$ with varying latent dimension $n_0 = k$. The 3-layers generative networks have fixed layer widths of dimensions $n_1 = 150$, $n_2 = 700$ and $n = n_3 = 600$ with varying latent dimension $n_0 = k$.

Figure 1 reports the empirical probability of successful recovery from 30 random independent trials (over $A$ and $G$) for of noiseless problems, where a run Algorithm 2 is called successful if the relative reconstruction error $\|x_\star - x_T\|/\|x_\star\|$ is below $10^{-3}$. These experiments show that recovery of $x_\star$ from undersampled measurements can be achieved even in the case of non-expansive generative networks and adding more layers makes the problem harder. Nonetheless, the algorithm succeeds in wider range of parameters as predicted by our theory, in particular with a milder dependence on the depth $d$.

In the second experiment we study the noisy compressive sensing problem with a fixed number of measurements $m = 300$. We consider 2-layers generative networks with fixed layer widths of dimensions $n_1 = 700$ and $n = n_2 = 600$, and varying latent dimension $n_0 = k$ in $\{5, 10, 15, 20, 25, 30, 35, 40, 45\}$. The measurements are taken to be $b = Ay_\star + \eta$ where $\eta = \tau e$ where $e$ is taken uniformly at random over the sphere $\mathcal{S}^{m-1}$ and $\tau$ is chosen so that the signal to noise ratio (SNR) varies in $\{20, 40, 80\}$. The SNR is defined as $10 \log_{10}\left(\|AG(x_\star)\|/\|\eta\|\right)$.

In Figure 2 we plot the average reconstruction error $\|G(x_T) - y_\star\|/\|y_\star\|$ over 30 independent random draws of $A$ and $G$. As predicted by our theory this quantity scales linearly with the ratio $k/m$. Furthermore, the error is proportional to the magnitude of the noise $\|\eta\|$.

Figure 2: Reconstruction error $\|y_\star - G(x)\|/\|y_\star\|$ from noisy compressed linear measurements at varying level of SNR. The number of measurements $m$ is fixed at 300 and latent dimension $k$ is varied. The generative networks have outputs of dimension $n = n_2 = 600$ and hidden layers of dimension $n_1 = 700$.