FIDELITY-BASED DEEP ADIABATIC SCHEDULING

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ABSTRACT

Adiabatic quantum computation is a form of computation that acts by slowly interpolating a quantum system between an easy to prepare initial state and a final state that represents a solution to a given computational problem. The choice of the interpolation schedule is critical to the performance: if at a certain time point, the evolution is too rapid, the system has a high probability to transfer to a higher energy state, which does not represent a solution to the problem. On the other hand, an evolution that is too slow leads to a loss of computation time and increases the probability of failure due to decoherence. In this work, we train deep neural models to produce optimal schedules that are conditioned on the problem at hand. We consider two types of problem representation: the Hamiltonian form, and the Quadratic Unconstrained Binary Optimization (QUBO) form. A novel loss function that scores schedules according to their approximated success probability is introduced. We benchmark our approach on random QUBO problems, Grover search, and 3-SAT problems and show that our approach outperforms, by a sizable margin, the linear schedules as well as alternative approaches that were very recently proposed.

1 INTRODUCTION

Many of the algorithms developed for quantum computing employ the quantum circuit model, in which a quantum state involving multiple qubits undergoes a series of invertible transformations. However, an alternative model, called Adiabatic Quantum Computation (AQC) (Farhi et al., 2000; McGeoch, 2014), is used in some of the leading quantum computers, such as those manufactured by D-Wave Systems (Boixo et al., 2014). AQC algorithms can achieve quantum speedups over classical algorithms (Albash & Lidar, 2018), and are polynomially equivalent to the quantum circuit model (Aharonov et al., 2008).

In AQC, given a computational problem Q, e.g., a specific instance of a 3SAT problem, a physical system is slowly evolved until a specific quantum state that represents a proper solution is achieved. Each AQC run involves three components:

1. An initial Hamiltonian $\mathcal{H}_b$, chosen such that its ground state (in matrix terms, the minimal eigenvector of $\mathcal{H}_b$) is easy to prepare and there is a large spectral gap. This is typically independent of the specific instance of Q.
2. A final Hamiltonian $\mathcal{H}_p$ designed such that its ground state corresponds to the solution of the problem instance Q.
3. An adiabatic schedule, which is a strictly increasing function $s(t)$ that maps a point in time $0 \leq t \leq t_f$, where t_f is total computation time, to the entire interval $[0, 1]$ (i.e., $s(0) = 0$, $s(t_f) = 1$, and $s(t_1) < s(t_2)$ iff $t_1 < t_2$ and vice versa).

These three components define a single time-dependent Hamiltonian $\mathcal{H}(t)$, which can be seen as an algorithm for solving Q:

$$\mathcal{H}(t) = (1 - s(t)) \cdot \mathcal{H}_b + s(t) \cdot \mathcal{H}_p \tag{1}$$

At the end of the adiabatic calculation, the quantum state is measured. The square of the overlap between the quantum state and ground state of the final Hamiltonian, is the fidelity, and represents the probability of success in finding the correct solution. An AQC algorithm that is evolved over an insufficient time period (a schedule that is too fast) will have a low fidelity. Finding the optimal

schedule, i.e., the one that would lead to a high fidelity and would keep the time complexity of the algorithm minimal is, therefore, of a great value. However, for most problems, an analytical solution for the optimal schedule does not exist (Albash & Lidar, 2018).

Previous attempts to employ AI for the task of finding the optimal schedule have relied on Reinforcement Learning (Lin et al., 2020; Chen et al., 2020). While these methods were able to find schedules that are better than the linear path, they are limited to either learning one path for a family of problems (without considering the specific instance) or to rerunning the AQC of a specific instance Q multiple times in order to optimize the schedule.

In our work, supervised learning is employed in order to generalize from a training set of problems and their optimal paths to new problem instances. Training is done offline and the schedule our neural model outputs is a function of the specific problem instance. The problem instance is encoded in our model either based on the final Hamiltonian $\mathcal{H}_p$ or directly based on the problem.

The suggested neural models are tested using several different problem types: Grover search problems, 3SAT problems, and randomized QUBO problems. We show that the evolution schedules suggested by our model greatly outperform the naive linear evolution schedule, as well as those schedules provided by the recent RL methods, and allow for much shorter total evolution times.

2 BACKGROUND

The goal of the scheduling task is to find a schedule $s(t)$ that maximizes the probability to get the correct answer for instance Q, using H_b and H_p over an adiabatic quantum computer. The solution to Q is coded as the lowest energy eigenstate of $\mathcal{H}_p$.

In order to achieve the solution state with high probability, the system must be evolved "sufficiently slowly". The adiabatic theorem (Roland & Cerf, 2002; Albash & Lidar, 2018; Rezakhani et al., 2009) is used to analyze how fast could this evolution be. It states that the probability to reach the desired state at the end of the adiabatic calculation is $1-\varepsilon^2$ for $\varepsilon << 1$ if

$$\frac{\left|\langle E_1(t)|\, \frac{d}{dt}\mathcal{H}(t)\, |E_0(t)\rangle\right|}{g^2(t)} \leq \varepsilon \tag{2}$$

where the Dirac notation (Tumulka, 2009) is used[1], $E_0(t)$ ($E_1(t)$) is the ground state (first excited state) of the time dependent Hamiltonian $\mathcal{H}(t)$, i.e., the eigenstates that corresponds to the lowest (2nd lowest) eigenvalue, and $g(t)$ is the time dependent instantaneous spectral gap between the smallest and second smallest eigenvalues of $\mathcal{H}(t)$.

Let t_f be the total calculation time. let $s(t)$ be an evolution schedule , such that $s(0)=0,\ s(t_f)=1$. Applying the adiabatic condition for $s(t)$, we get

$$\frac{\left|\langle E_1(s(t))|\, \frac{\boldsymbol{ds}}{\boldsymbol{dt}}\frac{d}{ds}\mathcal{H}(s(t))\, |E_0(s(t))\rangle\right|}{g^2(s(t))} \leq \varepsilon \Rightarrow \frac{ds}{dt} \leq \varepsilon \frac{g^2(s)}{\left|\langle E_1(s)|\, \frac{d}{ds}\mathcal{H}(s)\, |E_0(s)\rangle\right|} \tag{3}$$

we could solve for $t(s)$ by integration to get

$$t(s) = \frac{1}{\varepsilon}\int_0^s \frac{\left|\langle E_1(s)|\, \frac{d}{ds}\mathcal{H}(s)\, |E_0(s)\rangle\right|}{g^2(s)} ds \tag{4}$$

and the total required evolution time is

$$t_f = t(s=1) = \frac{1}{\varepsilon}\int_0^1 \frac{\left|\langle E_1(s)|\, \frac{d}{ds}\mathcal{H}(s)\, |E_0(s)\rangle\right|}{g^2(s)} ds \tag{5}$$

We note that finding a numerical solution for eq 4 requires calculating the full eigenvalue decomposition of $\mathcal{H}(x)$.

[1] See appendix A for the conventional matrix notation.

2.1 MOST-RELATED WORK

Two recent contributions use deep learning in order to obtain, for a given t_f, a schedule that outperform the linear schedule. Lin et al. (2020) suggest using deep reinforcement learning in order to find an optimal schedule for each specific class of problems (e.g., 3SAT problems of a certain size). In contrast, we study the problem of finding schedules for generic problem instances. They train and benchmark their performance by simulating an adiabatic quantum computer, and scoring the computation results for randomly chosen problem instances. Their results are generally better than the naive linear schedule, and the solution produced by their neural network is somewhat transferable for larger problem sizes.

Chen et al. (2020) also use RL to construct, given a t_f, a schedule for 3SAT problems. The most successful technique suggested is a Monte Carlo Tree Search (MCTS, Silver et al. (2016)), which produces results that significantly outperform the linear schedule. This technique requires running the adiabatic evolution process many times for each problem, in order to find a successful schedule. An approach inspired by alpha-zero (Silver et al., 2018) is used to adapt the generic MCTS solution to specific problem class, while requiring only a few additional rounds of the adiabatic process for each new instance. In our method, we do not require any run given a new problem instance.

3 METHOD

We consider two types of deep neural models. The first model is designed to get the problem Hamiltonian $\mathcal{H}_p$ as an input. For an n qubit problem, the problem Hamiltonian is generally of size $2^n \times 2^n$. In this work, we consider problem Hamiltonians which are diagonal and can be represented by vector of size 2^n. This scenario covers both the Grover search problem and the 3SAT problem we present in Sec. 4.

The second model is designed to get a quadratic unconstrained binary optimization (QUBO) problem as an input. The QUBO problem has the following form:

$$\bar{x} = argmin_x(x^T Q x)\,, \tag{6}$$

where x is a vector of binary variables and $Q \in \mathbb{R}^{n\times n}$ defines the specific QUBO instance. The QUBO problem is NP-Complete, and many types of common problems can be reduced to QUBO (Glover et al., 2018). The QUBO formulation is of special interest in the context of adiabatic quantum computing, since it allows a relatively easy mapping to real quantum annealing devices that do not possess full qubit connectivity (Cruz-Santos et al., 2019).

A QUBO problem Q can be converted to the Hamiltonian form in the following fashion:

$$\mathcal{H}_p = \sum_{i=1}^{n} Q_{ii}(\frac{I+\sigma_z^i}{2}) + \sum_{i\neq j} Q_{ij}(\frac{I+\sigma_z^i}{2})(\frac{I+\sigma_z^j}{2})\,, \tag{7}$$

where σ_z^i is the Pauli matrix σ_z operating only on qubit i (Liboff, 2003). The resulting $\mathcal{H}_p$ is of size $2^n \times 2^n$ and is diagonal.

The prediction target of our models is the desired normalized schedule $\hat{s}(t)$, which is defined over the range $[0,1]$ as $\hat{s}(t) = s(t/t_f)$. For the purpose of estimation, it is sampled at 100 points in the interval $t = [0,1]$. The representation of this schedule is given as a vector $\boldsymbol{d} \in [0,1]^{99}$, which captures the temporal derivative of the schedule. In other words, $\boldsymbol{d}$ is trained to hold the differences between consecutive points on the path, i.e., element i is given by $\boldsymbol{d}_i = \hat{s}((i+1)/100) - \hat{s}(i/100)$. Note that the sum of $\boldsymbol{d}$ is one.

3.1 UNIVERSALITY OF THE OPTIMAL SCHEDULE

The reason that we work with the normalized schedule is that the optimal evolution schedule is not dependent upon the choice of t_f. As shown next, for every time budget t_f, the same normalized schedule would provide the highest fidelity (neglecting decoherence).

Let $s_1(t) : [0, t_f] \to [0,1]$ be a suggested evolution schedule, which outperforms a different suggested schedule $s_2(t)$, for a specific $t_f = \tau_1$, i.e. it achieves a greater fidelity at the end of the

schedule for a specific problem instance Q. Then, Thm. 1 shows that $s_1(t)$ outperforms $s_2(t)$ for every possible choice of t_f for the same problem Q.

Theorem 1. *Let $s_1(t)$ and $s_2(t)$ be two monotonically increasing fully differentiable bijective functions from $[0, t_f = \tau_1]$ to $[0, 1]$. Let Q be an optimization problem, and assume that $s_1(t)$ achieves a greater fidelity than $s_2(t)$ at the end of a quantum adiabatic computation for Q with total evolution time $t_f = \tau_1$. Then, for any other choice $t_f = \tau_2$, the scaled schedule $s_1(\frac{\tau_2}{\tau_1}t)$ will achieve a greater fidelity than $s_2(\frac{\tau_2}{\tau_1}t)$ for an adiabatic computation over the same problem Q with total evolution time $t_f = \tau_2$.*

The proof can be found in appendix B.

3.2 Architecture

The model architectures are straightforward and no substantial effort was done to optimize them. The Hamiltonian as input model has seven fully connected layers, with decreasing sizes: 4096, 2048, 2048, 1024, 512, and finally the output layer, which, as mentioned, is of size 99.

For the QUBO model, in which the input is a matrix, a two part architecture was used. In the first part, five layers of 2D convolution was employed, with kernel size of 3×3, for 64 kernels. The output from the convolution layers was then flattened to a vector of size $64n^2$, and fed to the second part of the network, consisted of five fully connected layers, with decreasing dimensions of 2048, 1024, 1024, 512, and finally the output layer of size 99.

This output layers in both models are normalized to have a sum of one. For both models, the SELU activation function Klambauer et al. (2017) was used for all layers, except the final layer, which used the sigmoid (logistic) function.

3.3 A Fidelity Based Loss Function

Let $|\psi(t)\rangle$ is the state of the quantum system at time $t = st_f$. The fidelity of the QAC is given by (Farhi et al., 2000)

$$p_{success} = |\langle E_0(s=1) \,|\, \psi(t=t_f)\rangle|^2 \ , \tag{8}$$

where $\langle E_\ell(s=1)|$ is the ℓ-th eigenstate of the parameter dependent evolution Hamiltonian $\mathcal{H}(s)$, such that $\langle E_0(s=1)|$ is the ground state of the final Hamiltonian $\mathcal{H}_p$. Finding $\langle E_0(s=1)|$ requires performing eigenvalue decomposition for $\mathcal{H}_p$, which is equivalent to solving the original optimization problem, and is done for the training set.

The quantum state $|\psi(t)\rangle$ is evolving according to the Schrödinger equation

$$i\frac{d}{dt}\,|\psi(t)\rangle = \mathcal{H}(t)\,|\psi(t)\rangle \tag{9}$$

A brute force approach for finding $p_{success}$ is to numerically solve the Schrödinger equation, see appendix C. This full numerical calculation is, however, too intense to be practical. We next develop an approximate method that would be easier to compute and still be physically meaningful. It is based on the adiabatic local evolution speed limit from Eq. 3:

$$\left|\frac{ds}{dt}\right| \leq \varepsilon \frac{g^2(s)}{\left|\langle E_1(s)|\,\frac{d}{ds}\mathcal{H}(s)\,|E_0(s)\rangle\right|} \tag{10}$$

This inequality could be used as a local condition for convergence of any suggested path. We define

$$g_E^2(s) = \frac{g^2(s)}{\left|\langle E_1(s)|\,\frac{d}{ds}\mathcal{H}(s)\,|E_0(s)\rangle\right|} \tag{11}$$

We would like to use the local condition to create a global convergence condition for a full suggested path $s(t), 0 \leq t \leq t_f$. To do so, we integrate both sides of Eq. 10 over the suggested schedule s. This integral represents a mean value of the local adiabatic condition, for every point in the suggested schedule.

$$\varepsilon = \int_0^1 \frac{\frac{ds}{dt}}{g_E^2(s)} ds \tag{12}$$

We note that integrand is always positive (assuming $s(t)$ is monotonically increasing). Recall that the adiabatic theorem ties ε to the fidelity: $\varepsilon = \sqrt{1 - p_{success}}$. By defining the right hand side of Eq.12 as our loss function, we ensure that any training process that minimizes Eq. 12 will maximize the fidelity. Recall that the vector $\boldsymbol{d}$ that the network outputs is a vector of differences, therefore, it approximates the local derivatives of the obtained path. Let $\hat{s}^*$ be the optimal normalized path, which we estimate for each training sample. The loss function is, therefore, defined as:

$$\mathcal{L}(\boldsymbol{d}, \hat{s}^*) = \sum_{i=1}^{99} \frac{\boldsymbol{d}_i^2}{g_E^2(\hat{s}^*(i/100))} \tag{13}$$

The values of g_E are precomputed along the optimal path $\hat{s}^*$ for efficiency. While the denominator is obtained on points that do not correspond to the estimated path (the commutative sum of $\boldsymbol{d}$), the approximation becomes increasingly accurate at the estimated path appraoches the optimal one.

3.4 Training Data and the Training Process

In order to train the QUBO problem model, we produced a training dataset of 10,000 random QUBO instances for each problem size: $n = 6, 8, 10$. The QUBO problems were generated by sampling independently, from the normal distribution, each coefficient of the problem matrix Q. The entire matrix Q was then multiplied by a single random normal variable.

We approximated an optimal evolution schedule for each problem, by calculating the full eigenvalue decomposition of $\mathcal{H}_t$ as described in Sec 2. We also calculated the value of $g(s(t))$ for each problem.

For the model that uses the problem Hamiltonian as input, we used the same prepared QUBO problems, converted to the Hamiltonian form. In addition, we added another 500 cases of randomized Hamiltonians with randomized values around distinct energy levels. For each Hamiltonian, We first randomized an energy level between the following values: 0.5, 1, 1.5 or 2, and then randomized uniformly distributed values around the selected energy level. To each Hamiltonian we added a single ground state with energy 0. This type of Hamiltonian is not commonly created by the random QUBO creation process described above, but is more representative of binary optimization problems, and specifically more closely resembles problem Hamiltonians for the Grover problem and the 3SAT problem, which we later use to benchmark our model performance. We note that the Hamiltonian for these specific problems in our test set are nevertheless different from our randomized problem Hamiltonians, which highlights the generalization capability of our method.

The training was performed using the Adam optimizer (Kingma & Ba, 2014), with batches of size 200. Batch normalization (Ioffe & Szegedy, 2015) was applied during training. A uniform dropout value of 0.1 is employed for all layers during the model training.

4 Results

As a baseline to the loss $\mathcal{L}$ (Eq. 13) we use, we employed the Mean Squared Error (MSE) loss, for which the model output was compared to the known optimal schedule from the dataset, which was calculated in advance.

4.1 Grover Search

The Grover algorithm is a well-known quantum algorithm that finds with high probability the unique input to a black box function that produces a particular output value, using just $\sqrt{N}$ evaluations of the function, where N is size of the search space. For an n qubit space, the search is over the set $\{0, 1, .., 2^n - 1\}$, making $N = 2^n$. It is possible to reproduce the Grover speedup using an adiabatic formulation, with the following problem Hamiltonian:

$$\mathcal{H}_p = I - |m\rangle \langle m| \ , \tag{14}$$

where $|m\rangle$ is the state that represents the value we search. Roland & Cerf (2002) showed that for this problem, a linear schedule does not produce quantum speedup over a classical algorithm, but for a specific initial Hamiltonian $\mathcal{H}_b = I - |\psi_0\rangle \langle \psi_0|$, for ψ_0 as the maximal superposition state (a

sum of the states representing all values from 0 to $N-1$), an optimal schedule could be derived analytically to achieve a quadratic speedup. The optimal path is given by

$$\hat{s}(t) = \frac{1}{2} + \frac{1}{2\sqrt{N-1}} \tan\left[(2s-1)\tan^{-1}\sqrt{N-1}\right] \tag{15}$$

In practice, the proposed $\mathcal{H}_b$ is hard to physically realize, and a simpler initial Hamiltonian is used:

$$\mathcal{H}_b = \frac{1}{2}\sum_{i=1}^{n} I - \sigma_x^i \,, \tag{16}$$

where σ_x^i is the Pauli matrix σ_x operating only on qubit i (Liboff, 2003).

We test our model's performance by using the Grover problem Hamiltonian $\mathcal{H}_p$ as input for several problem sizes. Different Grover problems are completely symmetrical, and are identical after changing variables, so it is sufficient to use a single test case to test our model.

We benchmark our model's performance by simulating AQC for multiple values of t_f, and calculating the fidelity by measuring the overlap between the quantum state at the end of the adiabatic evolution and the solution state.

We also show the convergence pattern for the fidelity (i.e. the overlap with the solution state, measured during the adiabatic evolution) for a single specific t_f. For each problem size, we chose a different t_f, for which a full convergence ($p > 0.95$) is achieved with the evolution schedule suggested by our model. We compare several suggested schedules: the path produced by training our model using our novel loss function, the path produced by training our model using the MSE loss, the linear path, and a numerically calculated optimal path. We also include the results reported by Lin et al. (2020) for the same problem.

The results are reported in Fig. 1 for $n = 6, 10$, see appendix for $n = 8$. It is evident that our model produces paths that are significantly superior to the linear path, and also outperforms Lin et al. (2020). The advantage of the new loss function over the MSE loss is also clear.

4.2 3SAT

In the 3-SAT problem, the logical statement consists of m clauses, C_i, such that each clause contain a disjunction over three variables out of n binary variables. A solution to the 3SAT problem is an assignment for the n variables that satisfies all m clauses. It is possible to construct a problem Hamiltonian for each 3SAT problem, by taking a sum over all clauses

$$\mathcal{H}_p = \frac{1}{2}\sum_{i=1}^{m} I + \sigma_z^{F_i} \,, \tag{17}$$

where $\sigma_z^{F_i}$ is the Pauli matrix σ_z operating only on the state that represents the assignment $|a = \{0,1\}, b = \{0,1\}, c = \{0,1\}\rangle$ which produces False value for clause i. This Hamiltonian counts the number of clauses which are not satisfied by each assignment, and its ground state corresponds to the eigenvalue 0 and represents the solution of the problem, for which all clauses are satisfied.

We test our model's performance, by randomizing 3SAT problems, and converting them to Hamiltonian form. Following Chen et al. (2020), we focus on 3SAT problems with a single solution, and a number of clauses $m = 3n$. This type of 3SAT problems is considered difficult to solve with adiabatic algorithms (Žnidarič, 2005).

We benchmark our model's performance by simulating the adiabatic computation for multiple values of t_f and calculating the fidelity by measuring the overlap between the quantum state at the end of the adiabatic evolution and the solution state.

In addition to the linear path and the paths obtained by training with either $\mathcal{L}$ or MSE, we also include for n=11, the results for the schedules designed by MCTS (Chen et al., 2020). For this purpose, we used the test data obtained by Chen et al. As can be seen in Fig. 2, our method outperform all baselines. Note that the MCTS methdod was optimized, for each problem instance and for each t_f using tens AQC of runs on the specific test problem, while our method does not run on the test data.

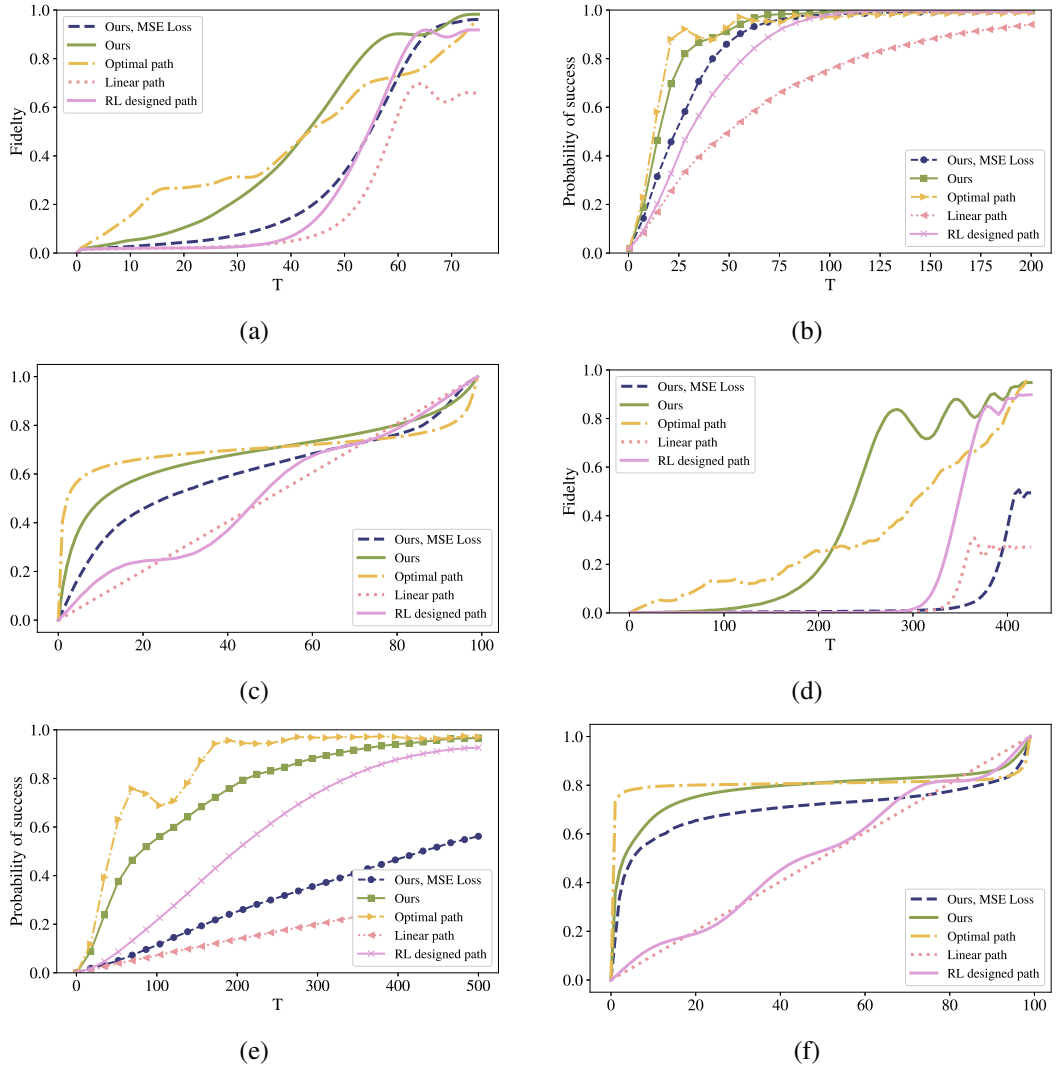


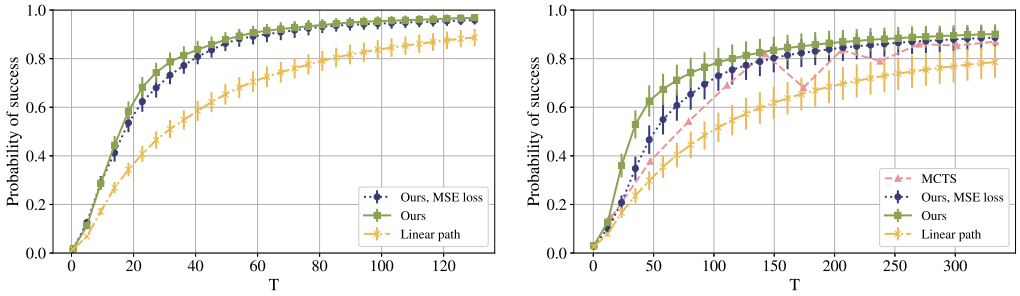
Figure 1: Results for Grover search over (a-c) n=6 or (d-f) n=10 qubits. (a) fidelity for $t_f = 75$, and (d) fidelity for $t_f = 425$. (b,e) fidelity at time t_f for multiple t_f values. (c,f) sample suggested schedules.

4.3 QUBO

To test our models with general QUBO problems, sets of random QUBO test problems of varying difficulty are generated . Since the final energy gap of the corresponding problem Hamiltonian is a critical parameter that determines the difficulty of the problem at hand (problems with a small energy gap require much longer evolution schedules), we generated two sets of test problems. The first has an energy gap of $g \sim 10$ and the second has an energy gap of $g \sim 0.1$. Varying the gap was obtained by multiplying the random Q matrix by the required values of the gap.

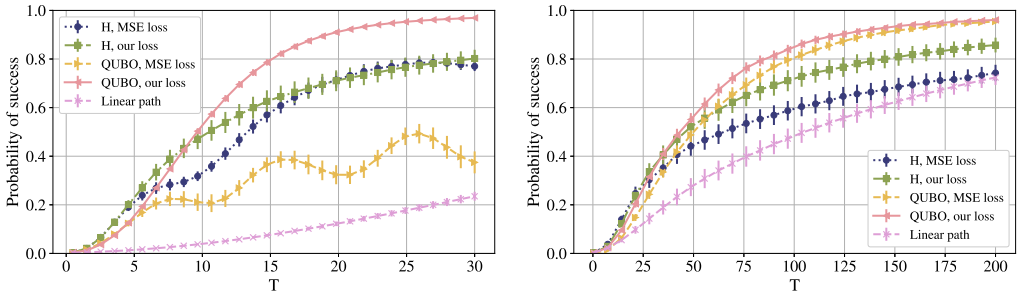
We benchmark the model's performance as in previous problems. However, in this case, we have two alternatives models: the one the receives the matrix Q as input and the one that receives the Hamiltonian $\mathcal{H}_p$.

We noticed that the process of creating samples of varying spectral gaps creates a mismatch in scale with the training set problems. To compensate, we pre-process the inputs to the network models. Specifically, for the model that has Q as input, we normalize the Frobenius norm of each Q such that if it is larger than 60, we scale Q to have a norm of 60. Similarly for the model that accepts



(a) (b)

Figure 2: Fidelity for various t_f values, over random 3SAT instances with m=3n clauses. Keep note of the different time scale for each problem size, as larger problems require longer times to converge. (a) n=8 qubits (b) n=11 qubits. For $n = 11$, we employ the test set of Chen et al. (2020) and directly compare with their MCTS method.



(a) (b)

Figure 3: Fidelity for various t_f values, over random QUBO instances of size n=10. Keep note of the different time scale for each problem, as problems with different energy gaps require different times to converge. The energy gaps are (a) $g \sim$10, (b) $g \sim$0.1

$\mathcal{H}_p$ as input, we clip every value that is larger than 90 to be 90 (Q with high norms translate to Hamiltonians with specific coeefieicents that are high). To clarify, this preprocessing is only applied to the input to the models and does not change the problem we solve.

Our Results are presented at Fig. 3. As can be seen, our dedicated QUBO model (Q as input) constructs successful schedules, outperforming all other models. The Hamiltonian model trained with our loss obtains the second highest results. The advantage of the fidelity-based loss term is evident in all cases.

5 Conclusions

Optimal scheduling of AQC tasks is the main way to reduce the time complexity for an emerging class of quantum computes. While recent work has applied RL for this task, it either provided a generic schedule for each class of problems or required running the exact computation that needs to be solved multiple times. Our solution employs a separate training set, and at test time provides a schedule that is tailored to the specific instance, without performing any runs. Remarkably, although our training was performed for one type of problem (QUBO), it generalizes well to completely different instances: Grover search and 3-SAT. At the heart of our method lies a new type of loss that maximizes the fidelity based on a new approximation of the success probability. Our experiments demonstrate the effectiveness of our method, as well as its advantage over the recent contributions.

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A CONVENTIONAL MATRIX NOTAION

For those not familiar with the Dirac notation, we repeat the equations in conventional matrix multiplication notation.

$$\frac{\left|E_1(t)^\top(\frac{d}{dt}\mathcal{H})E_0(t)\right|}{g^2(t)} \leq \varepsilon \tag{2}$$

$$\frac{\left|E_1(s(t))^\top \boldsymbol{\frac{ds}{dt}} \frac{d}{ds}\mathcal{H}(s(t))E_0(s(t))\right|}{g^2(s(t))} \leq \varepsilon \Rightarrow \left|\frac{ds}{dt}\right| \leq \varepsilon \frac{g^2(s)}{\left|E_1(s)^\top \frac{d}{ds}\mathcal{H}(s)E_0(s)\right|} \tag{3}$$

$$t(s) = \frac{1}{\varepsilon}\int_0^s \frac{\left|E_1(s)^\top \frac{d}{ds}\mathcal{H}(s)E_0(s)\right|}{g^2(s)} ds \tag{4}$$

$$t_f = t(s=1) = \frac{1}{\varepsilon}\int_0^1 \frac{\left|E_1(s)^\top \frac{d}{ds}\mathcal{H}(s)E_0(s)\right|}{g^2(s)} ds \tag{5}$$

$$p_{success} = \left|E_0(s=1)^\top \psi(t=t_f)\right|^2 \tag{8}$$

$$i\frac{d}{dt}\psi(t) = \mathcal{H}(t)\psi(t) \tag{9}$$

$$\left|\frac{ds}{dt}\right| \leq \varepsilon \frac{g^2(s)}{\left|E_1(s)^\top \frac{d}{ds}\mathcal{H}(s)E_0(s)\right|} \tag{10}$$

$$g_E^2(s) = \frac{g^2(s)}{\left|E_1(s)^\top \frac{d}{ds}\mathcal{H}(s)E_0(s)\right|} \tag{11}$$

$$\mathcal{H}_p = I - m^T m \tag{14}$$

B PROOF OF THM. 1

Theorem 1. *Let $s_1(t)$ and $s_2(t)$ be two monotonically increasing fully differentiable bijective functions from $[0, t_f = \tau_1]$ to $[0, 1]$. Let Q be an optimization problem, and assume that $s_1(t)$ achieves a greater fidelity than $s_2(t)$ at the end of a quantum adiabatic computation for Q with total evolution time $t_f = \tau_1$. Then, for any other choice $t_f = \tau_2$, the scaled schedule $s_1(\frac{\tau_2}{\tau_1}t)$ will achieve a greater fidelity than $s_2(\frac{\tau_2}{\tau_1}t)$ for an adiabatic computation over the same problem Q with total evolution time $t_f = \tau_2$.*

Proof. The adiabatic condition from Eq. 3 defines a local speed limit over the evolution schedule. We define:

$$g_E^2(s) = \frac{g^2(s)}{\left|\langle E_1(s)| \frac{d}{ds}\mathcal{H}(s) |E_0(s)\rangle\right|} \tag{18}$$

Then, for both schedules $s_i(t), i = 1, 2$ the local adiabatic speed is

$$\frac{1}{g_E^2(s_i(t))}\frac{ds_i(t)}{dt} = \varepsilon_i(t), \;\; 0 \leq t \leq \tau_1 \tag{19}$$

We now consider a new $t_f = \tau_2$. We use the same suggested schedules with a scaling factor $a = \frac{\tau_1}{\tau_2}$:

$$s_i^{scaled}(t) = s_i(at) \tag{20}$$

It is clear that $s_1^{scaled}(t=0) = 0$ and $s_1^{scaled}(t=\tau_2) = s_1\left(\frac{\tau_1}{\tau_2}\tau_2\right) = s_1(t=\tau_1) = 1$, and the same is true for $s_2^{scaled}(t)$. We calculate the new derivative

$$\frac{ds_i^{scaled}(t)}{dt} = a\frac{ds_i(at)}{dt} \tag{21}$$

By multiplying Eq. 19 by factor a we can get for the new time axis $0 \leq t \leq \tau_2$

$$a\frac{1}{g_E^2(s_i(at))}\frac{ds_i(at)}{dt} = a \cdot \varepsilon_i(t) \tag{22}$$

then, we can switch to the scaled schedules and finally

$$\frac{1}{g_E^2\left(s_i^{scaled}(t)\right)}\frac{ds_i^{scaled}(t)}{dt} = a \cdot \varepsilon_i(t) = \varepsilon_i^{scaled}(t) \tag{23}$$

We now consider the fidelity for each evolution schedule. According to the adiabatic theorem, the fidelity achieved at the end of the adiabatic evolution for each schedule is dependent solely on the local adiabatic speeds $\varepsilon_i(t)$. The resulting fidelity for the full path is then bounded by some functional $\mathcal{F} : L^2 \to \mathbb{R}$ which transforms all of the local adiabatic speeds to a single number.

$$p_i \geq 1 - \mathcal{F}\left(\varepsilon_i^2(t)\right) \tag{24}$$

Following Roland & Cerf (2002), we assume a global maximum value

$$\mathcal{F}[f(t)] = max(f(t)) \tag{25}$$

$$p_i \geq 1 - max\left(\varepsilon_i^2(t)\right) \tag{26}$$

it is clear that for this choice of $\mathcal{F}$,

$$\mathcal{F}[af(t)] = a\mathcal{F}[f(t)] \tag{27}$$

For any positive scalar a. It follows that the new values for fidelity for the scaled schedules are bounded by

$$p_i^{new} \geq 1 - \mathcal{F}\left(a^2\varepsilon_i^2(t)\right) = 1 - a^2\mathcal{F}\left(\varepsilon_i^2(t)\right) \tag{28}$$

We assumed $p_1 > p_2$, so $\mathcal{F}\left(\varepsilon_1^2(t)\right) < \mathcal{F}\left(\varepsilon_2^2(t)\right)$, and for any a it remains true that

$$a^2\mathcal{F}\left(\varepsilon_1^2(t)\right) < a^2\mathcal{F}\left(\varepsilon_2^2(t)\right) \tag{29}$$

and therefore

$$p_1^{new} \geq p_2^{new} \tag{30}$$

We note that this holds true for many choices for $\mathcal{F}[f(t)]$, as long as

$$\mathcal{F}[af(t)] = q(a)\mathcal{F}[f(t)] \tag{31}$$

for some monotonically increasing function q. □

C SOLVING THE SCHRÖDINGER EQUATION FOR THE ADIABATIC EVOLUTION

It is possibly numerically integrate and solve differential equation in Eq. 9, using the explicit evolution Hamiltonian $\mathcal{H}(s)$ for every $0 < s < 1$, and the boundary condition $|\psi(t=0)\rangle = |E_0(s=0)\rangle$, where $|E_0(s=0)\rangle$ is the known ground state of the initial Hamiltonian $\mathcal{H}_b$. This first order differential equation, could be solved numerically to obtain $|\psi(t=t_f)\rangle$ in the following fashion:

1. Divide the time axis to M slices $1..M$
2. For every time slice, find $\mathcal{H}_m = \mathcal{H}(s(t = \frac{t_f}{M} \cdot m))$
3. Calculate the eigenvalue decomposition of $\mathcal{H}_m$: eigenvectors V_i and eigenvalues E_i

4. Find the projection of the last quantum state $|\psi_{m-1}\rangle$ onto the eigenvectors space

$$|\psi_{m-1}\rangle = \sum_{i=1}^{N} a_i V_i \tag{32}$$

$$a_i = \langle V_i, \psi_{i-1} \rangle \tag{33}$$

5. Evolve the quantum state according to

$$|\psi_m\rangle = \sum_{i=1}^{N} e^{iE_i \cdot \frac{t_f}{M}} \cdot a_i \cdot V_i \tag{34}$$

6. Repeat steps 2-5 until reaching t_f

D ADDITIONAL GROVER SEARCH RESULTS

The results of Grover search for $n = 8$ qubits, for our model, as well as the method of Lin et al. (2020) and other baselines are presented in Fig. 4.

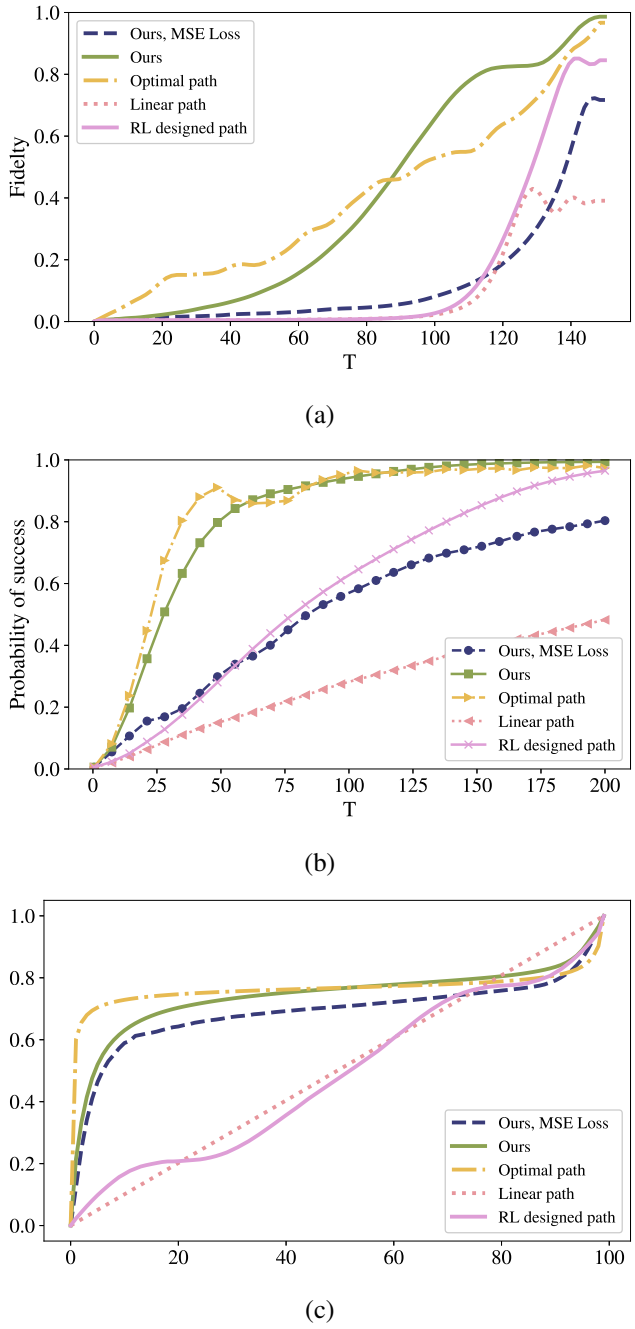

Figure 4: Model Results for Grover search over n=8 qubits. (a) fidelity for $t_f = 150$. (b) fidelity for multiple t_f. (c) suggested schedules for one specific instance.