Data Augmentation MCMC for Bayesian Inference from Privatized Data

Abstract

Differentially private mechanisms protect privacy by introducing additional randomness into the data. Restricting access to only the privatized data makes it challenging to perform valid statistical inference on parameters underlying the confidential data. Specifically, the likelihood function of the privatized data requires integrating over the large space of confidential databases and is typically intractable. For Bayesian analysis, this results in a posterior distribution that is doubly intractable, rendering traditional MCMC techniques inapplicable. We propose an MCMC framework to perform Bayesian inference from the privatized data, which is applicable to a wide range of statistical models and privacy mechanisms. Our MCMC algorithm augments the model parameters with the unobserved confidential data, and alternately updates each one conditional on the other. For the potentially challenging step of updating the confidential data, we propose a generic approach that exploits the privacy guarantee of the mechanism to ensure efficiency. In particular, we give results on the computational complexity, acceptance rate, and mixing properties of our MCMC. We illustrate the efficacy and applicability of our methods on a naïve-Bayes log-linear model as well as on a linear regression model.

1 Introduction

Motivation. Differential privacy [Dwork et al., 2006] presents a formal mathematical framework to protect the confidentiality of individuals and businesses in aggregate data products. It is the state-of-the-art standard for statistical disclosure limitation (SDL), and has become widely adopted by curators of large-scale scientific, commercial, and official databases. Differentially private data products are produced by probabilistic mechanisms that carry proven privacy guarantees. Generally speaking, these mechanisms work by introducing carefully designed random noise into the query of interest, which is an otherwise deterministic function of the underlying database.

The privatization of data products through noise infusion poses a challenge to statistical analysis in the downstream. Statistical estimators are typically complex functions of the data. If instead of the confidential data, the analyst only has access to a probabilistically processed version of them, how can they maintain the statistical validity of the resulting inference?

A crucial statistical advantage of differentially private mechanisms over traditional SDL counterparts, such as swapping [Dalenius, 1977], is that their probabilistic design is publicly known. This knowledge allows the data analyst to, at least in theory, accurately account for the privatization mechanism and conduct reliable uncertainty quantification. Nevertheless, it remains a substantial computational challenge to incorporate the privacy procedure into the statistical analysis. The

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challenge is a wide-spread and varied one, as the extra layer of privacy protection calls for the revision of a wide range of existing statistical methodologies that previously operate on the original, non-privatized data, most of which are neither low-dimensional nor simply structured. This is the challenge we address in this work, in which we develop a general computational framework for practitioners to obtain valid statistical inference based on privatized data.

Related literature. Current inferential strategies for privatized data fall into two broad categories. One invokes traditional statistical asymptotics to approximate the sampling distribution of a differentially private statistic, on the grounds that the privacy noise is often asymptotically negligible compared to errors due to sampling [e.g. Smith, 2011, Cai et al., 2021]. These approximations are often inaccurate for finite sample sizes [Wang et al., 2018] and call for specific handling to incorporate the privacy mechanism [e.g. Gaboardi et al., 2016, Wang et al., 2015, Gaboardi and Rogers, 2018].

The second category recognizes (as Section 2 will explain) that the marginal likelihood of the model parameters in (3) is central to the problem of inference from privatized data. The marginal likelihood requires a potentially high-dimensional integral over the space of unobserved confidential databases, and one that is analytically tractable only in a few, simple settings [Awan and Slavković, 2018, 2020]. Typically, one must resort to either approximating it or sampling from it using Monte Carlo methods. Markov chain Monte Carlo (MCMC) techniques have been proposed for specific privacy mechanisms and data generating models. Karwa et al. [2017] propose an MCMC procedure for inference on exponential random graph models, and Bernstein and Sheldon [2018, 2019] devise MCMC methods designed to handle the low-dimensional latent sufficient statistics from exponential family models and linear regression. Gong [2019] shows that for certain differentially private statistics, approximate Bayesian computation (ABC) can give samples that are exact with respect to the marginal likelihood and the Bayesian posterior. In addition, when the statistical model for the confidential data is fully parametric, the parametric bootstrap may be used to produce inference accompanied by uncertainty quantification with better accuracy than asymptotic approximation [e.g. Gaboardi et al., 2016, Ferrando et al., 2020]. Variational Bayesian analysis [Karwa et al., 2015] is another alternative which invokes a non-asymptotic approximation to the posterior distribution.

Our contribution. We develop a general-purpose MCMC framework to perform Bayesian inference on the model parameters underlying the privatized data. Our framework allows us to overcome the intractable marginal likelihood resulting from privatization, and is applicable to a wide range of statistical models and privacy mechanisms. The resulting MCMC algorithms are *exact*, in that they target the posterior distribution precisely, without involving any approximation.

Our approach is general purpose, allowing data analysts to leverage existing inferential tools designed for non-private data. It can be viewed as a flexible, user-friendly wrapper that migrates existing MCMC algorithms for non-private data to the setting of privatized data access, requiring no further algorithm design or tuning. The sampler, formally a Metropolis-within-Gibbs sampler, is presented in Algorithm 1, and only further requires that the analyst can 1) sample from the statistical model for the confidential data and 2) can evaluate the probability density of the noise induced by the privacy mechanism. The algorithm augments the model parameters with the unobserved confidential data, and alternately updates each one conditioned on the other. While the imputation of an entire unobserved database might appear daunting, we demonstrate how knowledge of the privacy mechanism can be exploited to confer performance guarantees to the proposed MCMC algorithm. We provide theoretical results for the computational complexity, Metropolis-Hastings acceptance rate, and mixing properties. In particular, the higher the privacy, the more rapid is our algorithm's exploration of the parameter space. We illustrate the efficacy and applicability of our methods on a privatized naïve-Bayes log-linear model and a linear regression model with clamped and privatized input.

2 Problem Setup

Let $x = (x_1, \ldots, x_n) \in \mathbb{X}^n$ denote the confidential database, containing n records. We assume these records are independent and identically distributed (i.i.d.) draws from a statistical model $f(\cdot \mid \theta)$, though this can be relaxed. The goal of the analyst is to conduct statistical inference on the

unknown model parameter $\theta \in \Theta$. A Bayesian analyst represents *a priori* beliefs about θ with a prior probability distribution $p(\theta)$, and seeks to compute a posterior distribution $p(\theta \mid x) \propto p(\theta) f(x \mid \theta)$ that updates their beliefs in light of the observations x. In many modern applications, this posterior distribution is intractable, and it is common for analysts to represent it using samples drawn via some MCMC algorithm. In this work, we will assume access to such a posterior sampling method:

Assumption 1. *The analyst has available a Markov kernel that targets $p(\theta \mid x) \propto p(\theta) f(x \mid \theta)$, the posterior distribution over the model parameters given the confidential database x.*

Differential privacy. Our work here focuses on the following departure from the usual Bayesian setting: instead of observing the database x, we observe a privatized data product or query, denoted as s_{dp}. The quantity s_{dp} is probabilistically generated based on data x through a *privacy mechanism*, written as $\eta(\cdot \mid x)$. The privacy mechanism η is said to be ϵ-*differentially private* (ϵ-DP) [Dwork et al., 2006] if for all values of s_{dp}, and for all 'neighboring' databases $(x, x') \in \mathbb{X}^n \times \mathbb{X}^n$ differing by one record (denoted by $d(x, x') \leq 1$), the probability ratio is bounded:

$$\frac{\eta(s_{\mathrm{dp}} \mid x)}{\eta(s_{\mathrm{dp}} \mid x')} \leq \exp(\epsilon), \quad \epsilon > 0. \tag{1}$$

The parameter ϵ is called the *privacy loss budget*, and controls how informative s_{dp} is about x. Large values of ϵ guarantee less privacy, while $\epsilon = 0$ corresponds to perfect privacy. A simple and widely used ϵ-differentially private mechanism is the *Laplace mechanism*: for a deterministic query $s : \mathbb{X}^n \to \mathbb{R}^m$, the privatized query is defined as $s_{\mathrm{dp}} = s(x) + u$, where $u = (u_1, \ldots, u_m)$ are i.i.d. Laplace variables. The scale parameter of the Laplace distribution is inversely proportional to ϵ (more privacy requires more noise), and directly proportional to $\Delta(s) = \max_{(x,x') \in \mathbb{X}^n \times \mathbb{X}^n; d(x,x') \leq 1} \|s(x) - s(x')\|_1$, the ℓ_1 *(global) sensitivity* of s (the more sensitive the confidential query is to changes in one record of the database, the more noise we need).

Our methodology requires that the privacy mechanism η is known and can be evaluated. This is true of ϵ- (or *pure*) DP, as well as common variants such as (ϵ, δ)- (or *approximate*) DP, *zero-concentrated* DP (zCDP) [Dwork and Rothblum, 2016, Bun and Steinke, 2016], and *Gaussian*-DP [Dong et al., 2021]. To ensure computational efficiency, we make the following additional assumption.

Assumption 2 (Record Additivity). *The privacy mechanism can be written in the form $\eta(s_{\mathrm{dp}} \mid x) = g\left(s_{\mathrm{dp}}, \sum_{i=1}^{n} t_i(x_i, s_{\mathrm{dp}})\right)$ for some known and tractable functions $g, t_1, \ldots, t_n$.*

We refer to privacy mechanisms that satisfy Assumption 2 as *record-additive*. An implication of record additivity is that after changing one record in x, we do not have to scan the entire database to reevaluate η. This is satisfied by many commonly used mechanisms, two important examples being: 1) mechanisms that add data-independent noise to a query of the form $s = \sum_{i=1}^{n} s_i(x_i)$, such as the sample mean, sample variance-covariance, and sufficient statistics of an exponential family distribution (see Sections 4 and 5 for examples), and 2) mechanisms designed to optimize empirical risk functions of the form $u(x, s_{dp}) = \sum_{i=1}^{n} u_i(x_i, s_{dp})$, such as the exponential mechanism [McSherry and Talwar, 2007], K-norm gradient mechanism [Reimherr and Awan, 2019], objective perturbation [Chaudhuri et al., 2011, Kifer et al., 2012], and functional mechanism [Zhang et al., 2012].

Doubly intractable Bayesian inference from privatized data. Without access to the confidential database x, and given only the privatized query s_{dp}, the Bayesian analyst is now concerned with the following posterior distribution:

$$p(\theta \mid s_{\mathrm{dp}}) \propto p(\theta) p(s_{\mathrm{dp}} \mid \theta). \tag{2}$$

Here, $p(s_{\mathrm{dp}} \mid \theta)$ is the *marginal likelihood* of θ, integrating over all possible confidential databases:

$$p(s_{\mathrm{dp}} \mid \theta) = \int_{\mathbb{X}^n} \eta(s_{\mathrm{dp}} \mid x) f(x \mid \theta) \, dx. \tag{3}$$

The marginal likelihood contributes all the information that is available in the privatized observation s_{dp} about the parameter θ, and is the foundation to statistical inference using privatized statistics [Williams and McSherry, 2010]. The posterior distribution (2) reveals that the inferential uncertainty

127 about the parameter θ consists of three contributing sources: 1) prior uncertainty as encoded in $p(\theta)$,
128 2) sampling (or modeling) uncertainty of the confidential database as reflected in f, and 3) uncertainty
129 due to privacy as induced by the probabilistic mechanism η.

130 We now come to the core challenge to address in this work: the marginal likelihood in (3) calls for an
131 integral over the entire space of possible input databases $x \in \mathbb{X}^n$. This is usually computationally
132 challenging, especially if the privacy mechanism is not a function of a low-dimensional sufficient
133 statistic. If the integral underlying the marginal likelihood is intractable, then $p(s_{\mathrm{dp}} \mid \theta)$ cannot be
134 analytically evaluated. This makes the corresponding posterior distribution $p(\theta \mid s_{\mathrm{dp}})$ of (2) *doubly*
135 intractable [Murray et al., 2012] in the sense that it cannot be analytically evaluated even up to a
136 normalizing constant. Thus, traditional MCMC techniques are inapplicable and inference strategies
137 devised for privatized statistics must tame this possibly high-dimensional integration problem.

3 Data Augmentation MCMC for Inference from Privatized Data

139 In this paper, we present a simple, efficient, and general *data augmentation* MCMC [Tanner and
140 Wong, 1987, Van Dyk and Meng, 2001] framework, allowing practitioners to perform valid Bayesian
141 inference on a wide-range of data models and privacy mechanisms. Our approach is to augment the
142 MCMC state space with the latent confidential database x, so that the stationary distribution is the
143 *joint* posterior distribution

$$p(\theta, x \mid s_{\mathrm{dp}}) \propto p(\theta) f(x \mid \theta) \eta(s_{\mathrm{dp}} \mid x). \tag{4}$$

144 Marginally, the θ samples produced by such an algorithm follow the posterior $p(\theta \mid s_{\mathrm{dp}})$ in (2). Our
145 sampler is *exact*, targeting the marginal posterior distribution $p(\theta|s_{\mathrm{dp}})$ without any approximation
146 error, despite the fact that the marginal likelihood (3) is intractable.

147 Our approach of imputing the latent confidential database x is motivated by two factors: 1) we wish
148 our algorithm to be *general-purpose*, applicable to a wide range of models and privacy mechanisms,
149 and 2) we wish our algorithm to inherit guarantees on mixing performance from guarantees of the
150 privacy mechanism. Towards these ends, we do not assume any specific form of the underlying
151 model of x and the privacy mechanism beyond Assumptions 1 and 2 respectively. In this light, our
152 contribution can be viewed as a flexible wrapper that allows existing MCMC algorithms for models
153 of the confidential data to be extended to settings where the data is now protected by some privacy
154 mechanism. Though imputing the confidential dataset might appear to present a significant challenge,
155 we show that properties of the mechanism can be exploited to give performance guarantees on our
156 sampling scheme, and show that it has a runtime of the same order as the non-private sampler.

157 In what follows, we outline our proposed MCMC algorithm, derive guarantees on the runtime and
158 acceptance rate of the algorithm, and provide mild conditions for the proposed samplers to be ergodic,
159 as well as additional conditions for our sampler to achieve geometric rates of convergence.

3.1 A Privacy-Aware Metropolis-within-Gibbs Sampler

161 Our approach to sample from the joint posterior distribution $p(\theta, x \mid s_{\mathrm{dp}})$ is through a sequence of
162 alternating Gibbs updates. Let $(x^{(t)}, \theta^{(t)})$ denote the state of the Gibbs sampler at the t-th iterations.
163 Each iteration of the Gibbs sampler entails two steps:
164 (**Step 1**) sample $\theta^{(t+1)}$ from $p(\cdot \mid x^{(t)}, s_{\mathrm{dp}})$, and (**Step 2**) sample $x^{(t+1)}$ from $p(\cdot \mid \theta^{(t+1)}, s_{\mathrm{dp}})$.

165 The conditional distribution in Step 1 simplifies as $p(\theta|x^{(t)}, s_{\mathrm{dp}}) = p(\theta|x^{(t)})$, highlighting why data-
166 augmentation is useful: this conditional distribution is independent of the privacy mechanism, and we
167 can use existing sampling algorithms (Assumption 1) for the confidential data. We note that with the
168 exception of a few models, such as simple models with conjugate priors, it is usually not possible to
169 directly sample from $p(\theta|x)$. Assumption 1 however only requires that we can *conditionally* simulate
170 a new value of θ from a Markov kernel that has $p(\theta|x^{(t)})$ as its stationary distribution. Our overall
171 Gibbs sampler then becomes a Metropolis-within-Gibbs sampler [Gilks et al., 1995], that nevertheless
172 targets the joint posterior $p(\theta, x|s_{\mathrm{dp}})$.

Step 2 is the data-augmentation step, and connects the statistical model and the privacy mechanism on x. Again, we cannot expect to produce a conditionally independent sample of the latent database from $p(x \mid \theta, s_{\mathrm{dp}})$, as this is model and mechanism dependent. Instead, we take the much more tractable approach of cycling through the elements of latent database x, sequentially updating x one element of a time. Writing $x_{-i} = (x_1, \ldots, x_{i-1}, x_{i+1}, \ldots, x_n)$ to denote the vector x excluding the ith element, step 2 then consists of the sequence of updates $p(x_1 \mid \theta, x_{-1}, s_{\mathrm{dp}}), p(x_2 \mid \theta, x_{-2}, s_{\mathrm{dp}}), \ldots, p(x_n \mid \theta, x_{-n}, s_{\mathrm{dp}})$. The complete sweep can be viewed as a dependent update of the latent database x that targets the conditional distribution $p(x \mid \theta, s_{\mathrm{dp}})$.

Before we specify our complete sampler in Algorithm 1, we first address the following questions: (Q1) *What is the performance loss from updating x one element at a time, rather than jointly?*, and (Q2) *How can we efficiently carry out the conditional updates $p(x_i \mid \theta, x_{-i}, s_{\mathrm{dp}})$, $i = 1, \ldots, n$?*

Q1 concerns whether the dependence of x_i given $(x_{-i}, s_{\mathrm{dp}})$ is so strong as to impede efficient exploration of the $\mathbb{X}^n$-space and cause poor mixing. Here we note that the privacy mechanism limits the change in the likelihood $\eta(s_{\mathrm{dp}} \mid x)$ when one element of x is changed, and therefore limits the coupling between x_i and x_{-i}. This suggests a Gibbs sweep through the latent database x will not suffer from poor mixing.

Algorithm 1 One iteration of the privacy-aware Metropolis-within-Gibbs sampler

    1. Conditional update of $p(\theta \mid x)$ using the kernel from Assumption 1.

    2. For each $i = 1, 2, \ldots$, sequentially update $x_i \mid x_{-i}, \theta, s_{\mathrm{dp}}$.

        (a) Propose $x_i^\star \sim f(\cdot \mid \theta)$.

        (b) Update $t(x^\star, s_{dp}) = t(x, s_{dp}) - t_i(x_i, s_{dp}) + t_i(x_i^\star, s_{dp})$ according to Assumption 2.

        (c) Accept the proposed state with probability $\alpha(x_i^\star \mid x_i, x_{-i}, \theta)$ given by:

$$\alpha(x_i^\star \mid x_i, x_{-i}, \theta) = \min\left\{\frac{\eta(s_{dp} \mid x_i^\star, x_{-i})}{\eta(s_{dp} \mid x_i, x_{-i})}, 1\right\} = \min\left\{\frac{g(s_{\mathrm{dp}}, t(x^\star, s_{\mathrm{dp}}))}{g(s_{\mathrm{dp}}, t(x, s_{\mathrm{dp}}))}, 1\right\}. \quad (5)$$

Q2 recognizes that the conditionals $p(x_i \mid \theta, s_{\mathrm{dp}}, x_{-i})$ are model- and mechanism-specific, and simulating from these is challenging in most settings. For this, we take the following simple approach in Algorithm 1: at each step, we propose x_i from the model $f(x \mid \theta)$, and accept it with the appropriate MH probability (5). Observe that our choice of proposal distribution is independent of the privacy mechanism, and ignores the privatized data s_{dp} as well as all other elements x_{-i}. Despite being simple and general purpose, we show in Proposition 3.1 that for $\epsilon-$DP, we can lower-bound the acceptance probability of proposals produced this way by $\exp(-\epsilon)$. This lower bound is key to efficiency: despite the unconstrained nature of the proposal distribution, we can guarantee a minimum acceptance probability. These two facts suggest our sampler will explore the space of databases relatively quickly.

Proposition 3.1. *For a pure ϵ-DP privacy mechanism η, the acceptance probability α from Equation (5) satisfies $\alpha(x_i^\star \mid x_i, x_{-i}, \theta) \geq \exp(-\epsilon)$, for all $\theta, x_{-i}, x_i, x_i^\star$.*

The privacy loss budget ϵ is usually understood to be a small constant, which privacy experts recommend be between .01 and 1 [Dwork, 2011]. When $\epsilon = 1$, Proposition 3.1 ensures that the acceptance rate in Algorithm 1 is no less than 36.7%, and as ϵ approaches zero, the bound on the acceptance rate approaches one. Intuitively, this is because as ϵ decreases, the distribution of the privatized data s_{dp} depends less and less on any individual element of the database.

The simplicity of our approach arises through a decoupling of the data model from the privacy mechanism: the former is used to update θ and propose x_i's, while the latter is used to calculate the acceptance probabilities. The next result formalizes the computational efficiency of our approach. Specifically, for any record-additive mechanism, one iteration of our algorithm requires $O(n)$ operations, where n is the size of the latent database. Essentially, this arises because of Assumption 2, which allows the acceptance probability in (5) to be calculated in $O(1)$ time.

Proposition 3.2. *The Gibbs sampler described in Algorithm 1 requires $O(n)$ number of operations to update the full latent database according to $p(x \mid \theta, s_{\mathrm{dp}})$.*

Note that even without privacy, one round of an MCMC procedure typically takes $O(n)$ time. This is because updating θ given the confidential data requires computing the data likelihood $f(x \mid \theta) = \prod_{i=1}^{n} f(x_i \mid \theta)$, an $O(n)$ operation in general. Thus, as a result of the mild and typical condition that η is record-additive, our MCMC procedure enjoys the *same order* of runtime as the original MCMC algorithm for confidential data.

The previous two results are key to understanding the efficiency of our approach. In the next section we formally establish geometric ergodicity of the sampler in Theorem 3.3 and 3.4.

Computational complexity. The i.i.d. assumption on the records ensures that step 2a of Algorithm 1 take $O(1)$ time, though this assumption can easily be weakened. Assumption 2 allows steps 2b and 2c to also takes $O(1)$ time. Overall, step 2 of our algorithm then takes $O(n)$ (rather than $O(n^2)$) time, as stated in Proposition 3.2. This matches the typical per-iteration cost of samplers for the non-private posterior distribution required in step 1. Thus, the overall cost of an iteration of our MCMC sampler is $O(n)$, which is typical when dealing with datasets of size n.

3.2 Ergodicity of the Privacy-Aware Sampler

Ergodicity ensures the MCMC chain converges to the posterior distribution in total variation distance. [Tierney, 1994], which is essential for an MCMC sampler to consistently estimate functionals of the posterior distribution. In Theorem 3.3, we provide mild and sufficient conditions for our proposed Metropolis-within-Gibbs sampler to be ergodic.

Theorem 3.3. *Under conditions A1 - A3 below, the Metropolis-within-Gibbs sampler of Algorithm 1 on the joint space $(\mathbb{X}^n \times \Theta)$ is ergodic and it admits $p(x, \theta \mid s_{\mathrm{dp}})$ as the unique limiting distribution.*

A1. The prior distribution is proper and $p(\theta) > 0$ for all θ in $\Theta = \{\theta \mid f_\theta(x) > 0 \text{ for some } x\}$.
A2. The model is such that the set $\{x : f(x \mid \theta) > 0\}$ does not depend on θ.
A3. The privacy mechanism satisfies $\eta(s_{\mathrm{dp}} \mid x) > 0$ for all $x \in \mathbb{X}^n$.

We prove this in the supplementary material by verifying invariance, aperiodicity, and irreducibility. Conditions A1-A3 concern model specification, prior specification, and privacy noise. These mild assumptions are typically true and are easy to verify. While there are some mechanisms, such as the release-one-at-random mechanism, which satisfy approximate-DP but which violate A3 [Barber and Duchi, 2014], most privacy mechanisms of interest satisfy A3. It is easy to verify that if η satisfies ϵ-DP, zero-concentrated DP, or Gaussian-DP, then property A3 is guaranteed.

Next, we establish conditions for Algorithm 1 to be geometrically ergodic [Rosenthal, 1995, Roberts and Rosenthal, 1998]. A chain is said to be *geometrically ergodic* if its total variation distance to the target has a geometrically decaying upper bound. Geometric ergodicity is a desirable property since it provides a rate on convergence to the stationary distribution, guaranteeing central limit theorems, and allowing for the computation of asymptotically valid standard errors.

For simplicity, we focus on the situation where one can directly sample from the conditional posterior $p(\theta \mid x)$. This is an important and common case, relevant when either θ is low-dimensional or where one can place conjugate priors on θ. Both applications we present in this work, a log-linear model in Section 4, and a linear regression model in Section 5, fall under this setting.

Theorem 3.4. *Assume that in step 1 of Algorithm 1, one can directly sample from $p(\theta \mid x)$. Under A1-A3 of Theorem 3.3, the resulting (x, θ) chain, as well as the marginal chains, are geometrically ergodic if there exists $0 < a \leq b < \infty$ such that $a \leq f(x \mid \theta) \leq b \quad \forall \theta, x$.*

To prove Theorem 3.4, we verify the drift and minorization conditions for component-wise Gibbs samplers in Theorem 8 of [Johnson et al., 2013]. See the supplementary material for details. Note that Theorem 3.4 does not depend on the ϵ-DP assumption.

Unsurprisingly, geometric ergodicity requires stronger assumptions than just ergodicity. The first assumption concerning the ability to sample directly from $p(\theta|x)$ can be avoided, but for the sake of clarity we do not try to relax it, since we are mostly concerned with the interface with the privacy mechanism. The second assumption on the boundedness of the likelihood is stronger, but also typical. A common way to achieve bounded likelihoods is to require the sample space $\mathbb{X}$ (and typically also Θ) to be bounded. In many real-world settings, such bounds exist, even if they may be very loose.

4   Naïve Bayes Log-Linear Model

Log-linear models are often used to model categorical data, a popular instance being the naïve Bayes classifier. Following Karwa et al. [2016], we consider the following model: $x = (x_1, \ldots, x_K)$ is the input *feature-vector*, with each x_k taking values in $\{1, 2, \ldots, J_k\}$, and y is the output *class* taking values in $\{1, 2, \ldots, I\}$. Each input-output pair (x, y) forms one record in our confidential database, the entire database consisting of N i.i.d. copies of (x, y). The naïve Bayes classifier assumes that $P(x \mid y) = \prod_{k=1}^{K} P(x_k \mid y)$, the model parameters being $p_{ij}^k = P(x_k = j | y = i)$ and $p_i = P(y = i)$. The sufficient statistics of this model are $n_{ij}^k = \#(y = i, x_k = j)$, which count the number of class-feature co-occurrences. This will form our confidential query s, which we privatize by adding Laplace noise to each of the n_{ij}^k. The resulting quantity s_{dp}, consisting of the noisy counts $m_{ij}^k = n_{ij}^k + L_{ijk}$, is what we release. When $L_{ijk} \overset{\text{i.i.d.}}{\sim} \text{Laplace}(0, 2K/\epsilon)$, the output s_{dp} satisfies ϵ-DP. Placing a $\text{Dirichlet}(2, \ldots, 2)$ on all parameter vectors, our goal is to obtain the marginal posterior distribution of $p, p_{i-}^k \mid s_{dp}$. While Karwa et al. [2016] approximate this private posterior distribution using variational Bayes methods, our MCMC procedure is able to target the exact private posterior distribution.

Simulation setup. We perform several simulation experiments where we apply our MCMC samplers to the log-linear model described above. For the simulation, we set $N = 100$ (number of records), $I = 5$ (number of classes), $K = 5$ (number of features), and $J_K = 3$ (possible values for each feature). We evaluate our sampler for privacy levels corresponding to $\epsilon \in \{.1, .3, 1, 3, 10\}$.

Posterior mean. We generate one non-private dataset from the model, and hold it fixed. We then create 100 private queries s_{dp} at each ϵ value, and for each s_{dp} we run Algorithm 1 for 10000 iterations. We discard the first 5000 iterations as burn-in. Finally, for each chain, we calculate the posterior mean. Figure 1a plots the 100 different posterior means for each ϵ-value for the parameters $p_i = P(Y = i)$ for $i = 1, \ldots, 5$. In this plot, the solid horizontal lines indicate the non-private posterior means, and the dashed horizonal lines indicate the prior means for each parameter. We see that as ϵ approaches zero, the posterior mean approaches the prior mean, reflecting the intuition that we learn less from the data as the privacy budget gets smaller. On the other hand, as ϵ increases, we see that the private posterior mean approaches the non-private posterior mean, which reflects the fact that as ϵ grows, we learn approximately the same from the data as if there were no privacy mechanism.

Acceptance rate.   Using the same simulation setup as for the posterior mean, we calculate the average and minimum acceptance rate of Step 2 in Algorithm 1. Since the privacy mechanism satisfies ϵ-DP, we know that $\exp(-\epsilon)$ is a lower bound on these acceptance rates. In Figure 1b, we confirm this bound, and see that the average acceptance rate is significantly higher than this lower bound. This suggests that the chain mixes even faster than indicated by Proposition 3.1.

Coverage of credible intervals.   For the next experiment, we sample a set of parameters from the prior, and hold this fixed. Then for each ϵ value, we produce 100 non-private datasets (n_{ij}^k), one private dataset (m_{ij}^k) for each non-private one, and then run a chain for 10,000 iterations discarding the first 5000 iterations for burn-in. From each chain, we produce a 90% credible interval for each $p_i = P(Y = i)$, and calculate the empirical coverage which is reported in Table 1.

At a sample size of only $N = 100$, we do not expect the coverage of the credible intervals to match the nominal level of 90%, but we see in Table 1 that most of the coverage values are above .9. Notable exceptions are the coverage of p_4 when ϵ is small. This may be because when ϵ is small, the private

(a) Posterior means for the log-linear model. The solid horizonal lines indicate the non-private posterior means, and the dashed lines at .2 indicate the prior means.

(b) Observed acceptance rates for the log-linear model. The blue (above) point clouds indicate the average acceptance rate, and the orange (below) points indicate the observed minimum acceptance of each chain. The solid black line is the lower bound of Proposition 3.1.

ϵ	$p_1 = .097$	$p_2 = .148$	$p_3 = .145$	$p_4 = .446$	$p_5 = .163$
.1	1	1	1	**.36**	1
.3	.97	1	1	**.59**	1
1	.94	.99	.97	**.83**	.98
3	.95	.91	.97	.89	.93
10	.92	.88	.94	.92	.9

Table 1: Coverage of $p_i = P(Y = i)$ for the log-linear model at different ϵ. Top row is the true data generating parameter values. Coverage is based on 100 replicates.

posterior is approximately equal to the prior, which is centered at .2; however p_4 is significantly further from .2 than the other parameters, which may explain why the coverage is low in this case.

5 Linear Regression

Next, we consider ordinary linear regression with n subjects and p predictors. We write x_0 for the matrix of predictors excluding the intercept columns, $x = (\underline{1}, x_0)$ for the matrix including the intercept, and y for the vector of outcomes. We model the explanatory variables x_0 as $x_0^i \overset{\text{i.i.d.}}{\sim} \mathcal{N}_p(m, \Sigma)$ for $i = 1, \ldots, n$, with $y|x$ given by $\mathcal{N}_n(x\beta, \sigma^2 I_n)$. Here I_n is the $n \times n$ identity matrix and $\mathcal{N}_n$ denotes the n-dimentional multivariate Normal distribution. The parameters of interest are β, the $(p+1)$−dimension vector of regression coefficients, with σ, m and Σ assumed known. We use independent $\mathcal{N}(0, 2^2)$ priors for the components of β.

To achieve ϵ-DP via the Laplace mechanism, we require a finite global sensitivity. To achieve this, standard practice in the DP literature is to bound each predictor and response variable in a data-independent fashion. The bounds chosen by the privacy expert are $[a_i, b_i]$ for each instance of x_0^i and $[a_y, b_y]$ for the entries of y, and these values are shared with the analyst.

Definition 5.1. For a real value z, and $a \leq b$, define the *clamp function* $[z]_a^b := \min\{\max\{z, a\}, b\}$. If z is a vector of length d, we use the same notation to apply an entry-wise clamp: $[z]_a^b := ([z_1]_a^b, [z_2]_a^b, \ldots, [z_d]_a^b)^\top$.

Before adding noise for privacy, we first clamp the predictors and response, and then normalize them to take values in $[-1, 1]$: $\widetilde{x}_0^i := (b_i - a_i)^{-1} 2([x_0^i]_{a_i}^{b_i} - a_i) - 1$ and $\widetilde{y} := (b_y - a_y)^{-1} 2([y]_{a_y}^{b_y} - a_y) - 1$. Call $\widetilde{x} := [\underline{1}, \widetilde{x}_0^1, \widetilde{x}_0^2, \ldots, \widetilde{x}_0^p]$ and $x := (\widetilde{x}^\top \widetilde{y}, \widetilde{y}^\top \widetilde{y}, \widetilde{x}^\top \widetilde{x})$. The s is the summary statistic to which we will add noise for privacy. The ℓ_1 sensitivity of s (ignoring duplicate entries of $\widetilde{x}^\top \widetilde{x}$, and the constant entry $(\widetilde{x}^\top \widetilde{x})_{1,1}$) is $\Delta = p^2 + 3p + 3$. To satisfy ϵ-DP, we add independent Laplace$(0, \Delta/\epsilon)$ noise to each of the $d = \frac{1}{2}(p+1)(p+2) + (p+1)$ unique entries of x, which gives our final private

330 summary s_{dp}. We notice that s is an additive function and each individual's contribution to s is
331 $t(x_i, y_i) = \left((\tilde{x}^i)^\top \tilde{y}_i, \tilde{y}_i^2, (\tilde{x}_i)^\top \tilde{x}_i\right)$. This mechanism producing s_{dp} is record-additive.

Simulation setup. Our experiments focus on posterior inference about β based on S_{dp}. For simplicity,
333 we fix other parameters σ^2, m, Σ at the true data generating parameters (reported in the supplementary
334 materials). When they are unknown, the posterior distributions of these parameters can be estimated
335 by our Gibbs sampler as well. Confidential predictors and responses are clamped with bounds $b = 10$
336 and $a = -10$. Given a confidential database (x, y), the posterior distribution of β is multivariate
337 Normal and can be sampled directly with a runtime linear in n.

Posterior mean. We generate one confidential dataset (x, y) and hold it fixed. At each ϵ value, we
339 create 100 private outputs s_{dp} and run Gibbs samplers for 10,000 iterations targeting the posterior
340 $\beta \mid s_{\mathrm{dp}}$, discarding the first 5000 iterations. We plot the 100 different posterior means of β in Figure 2.
341 In this plot, the solid horizontal lines indicate posterior means given confidential data (x, y), which
342 we do not expect to fully recover due to clamping. The posterior quantities display the same trend
343 with respect to change in privacy level as observed in Figure 1b. The other experiments from Section
344 4 were also run on this linear regression model, and produced similar results. Simulation details, plots
345 and discussion are in the supplementary materials.

Figure 2: Posterior mean for private linear regression $\beta \mid s_{\mathrm{dp}}$ with fixed confidential data. The solid
horizonal lines indicate the confidential data posterior means, and the dashed lines indicate 0.

6   Discussion

347 We proposed a novel, but simple sampling procedure for parameter inference on $p(\theta \mid s_{\mathrm{dp}})$, which
348 leverage existing samplers for the non-private posterior, as well as the structure of the privacy
349 mechanism. The result is a simple wrapper for practitioners to obtain valid statistical inference from
350 privatized data using the same models for the unobserved confidential data. As a side product, our
351 algorithm also produces multiple copies of the confidential database from the posterior $p(x \mid s_{\mathrm{dp}})$,
352 which could be useful when one is interested in inferring properties of x as well. Although we did not
353 discuss this, our data augmentation scheme can also enable frequentist analysis through the Monte
354 Carlo expectation-maximization algorithm.

355 We acknowledge the limitations of present work. First, we point out that strong assumptions such
356 as bounded parameter space Θ and sample space $\mathbb{X}$ are required to establish geometric ergodicity
357 of the Gibbs sampler in Theorem 3.4. They can likely be relaxed, at the cost of a more complex
358 theorem statement and proof. Second, our current proposal for updating $x_i \mid x_{-i}, \theta, s_{\mathrm{dp}}$ only tailors
359 to the model $f(\cdot \mid \theta)$ and it is not customized for s_{dp} yet. In the future, we might be able to design
360 algorithms that also incorporate the privatized output s_{dp} in these proposals. Third, we point out that
361 strong correlations in the posterior $p(\theta, x \mid s_{\mathrm{dp}})$ can potentially cause poor mixing in practice, despite
362 geometric convergence rate. While in simple problems, this can be fixed by reparameterization, we
363 plan to develop MCMC algorithms for this setting in follow-up studies. Finally, while the proposed
364 algorithm converges so long as the privacy mechanism η is known, for alternative versions of DP
365 (such as zCDP), the acceptance probability results in Proposition 3.1 may no longer hold. Developing
366 alternatives to Proposition 3.1 for different privacy definitions is another goal of future work.

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